

Paper Number(s): **E4.05**
SO7

IMPERIAL COLLEGE OF SCIENCE, TECHNOLOGY AND MEDICINE
UNIVERSITY OF LONDON

DEPARTMENT OF ELECTRICAL AND ELECTRONIC ENGINEERING
EXAMINATIONS 2002

MSc and EEE PART IV: M.Eng. and ACGI

TRAFFIC THEORY & QUEUEING SYSTEMS

Wednesday, 24 April 10:00 am

There are FIVE questions on this paper.

Answer FOUR questions.

Time allowed: 3:00 hours

Examiners responsible:

First Marker(s): Barria, J.A.

Second Marker(s): De Wilde, P.

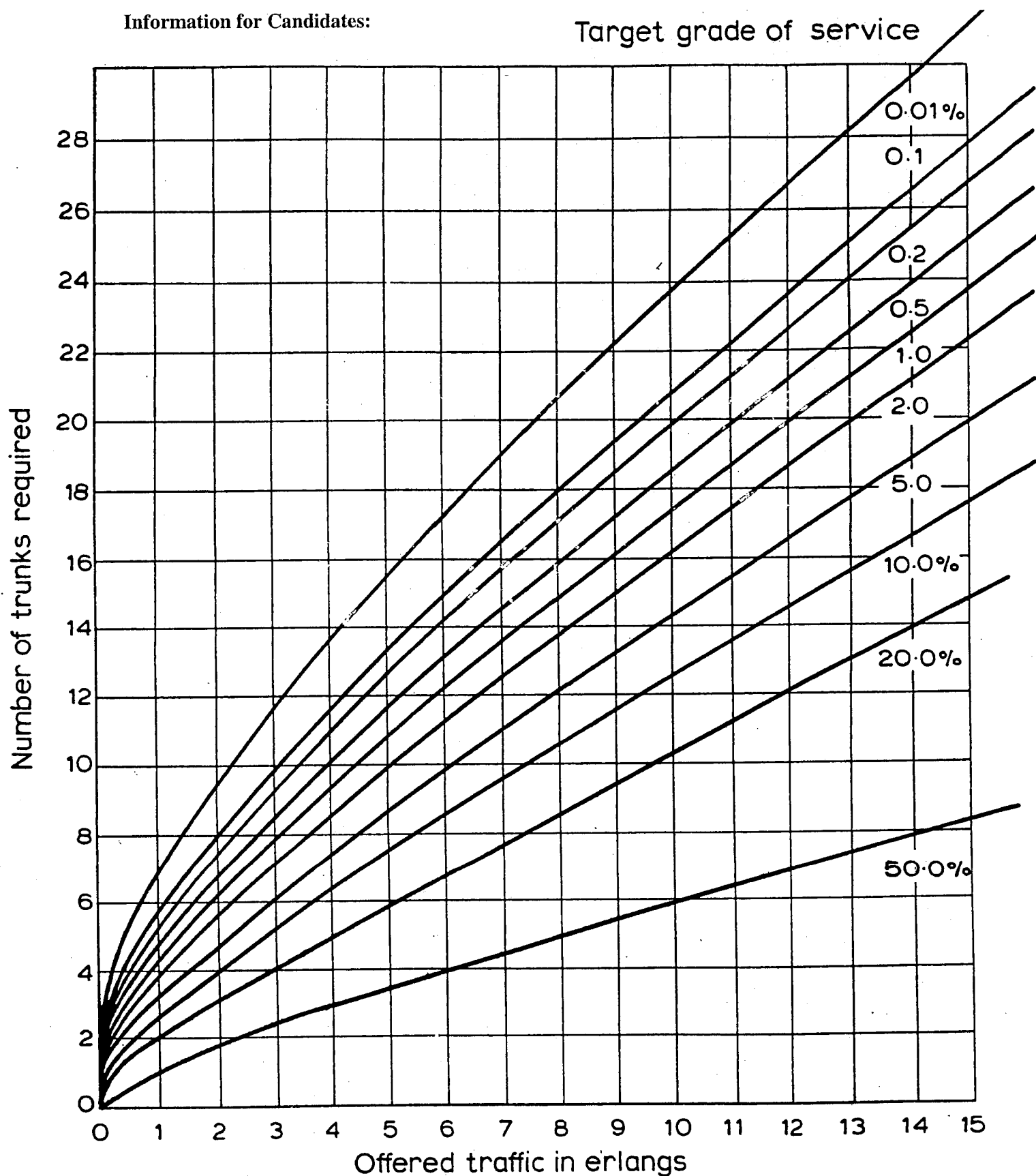
Corrected Copy

Special Information for Invigilators:

NIL

Information for Candidates:

Target grade of service



*Traffic capacity on basis of Erlang B.
formula.*

1. (a) For the Engset model:

- (i) Describe and list the basic assumptions underlying the model. [2]
- (ii) Derive the death and birth coefficients. [2]
- (iii) Derive the equilibrium traffic distribution if $N < M$. Where M is the number of independent active Poisson sources that is offered to an N -channel communications link. [6]

(b) Two telephone exchanges are connected by a multi-channel link operating with a loss probability of 0.005. If the aggregated calling rate in the two directions is 1200 calls/hour, and the average call holding time is 180 seconds:

- (i) Determine the total offered traffic and the total carried traffic for the link. [5]
- (ii) Estimate the size (i.e. number of channels) of the link. [5]

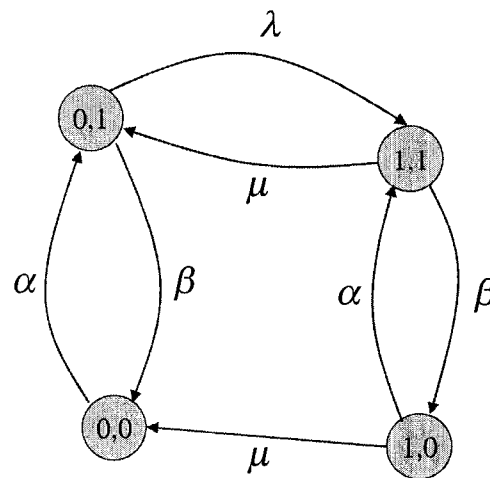
Hint: you can assume that for large values of the offered traffic, the B-curves are approximately linear.

2. (a) For the $M/M/K/N$ system, derive the Queue-length distribution. [10]

(b) The two dimensional Birth and Death process represented in *Figure 2.1* represents the traffic being offered to an overflow link. If we define the state of the system by (N_t, Y_t) where N_t is the number of busy channels on the overflow link, and Y_t is an ON/OFF switching process:

(i) define and explain the meaning of all parameters, [5]

(ii) calculate the blocking probability for the overflow link. [5]



$$\alpha = 0.145$$

$$\beta = 7.08$$

$$\lambda = 3$$

$$\mu = 0.33$$

$$\pi_0 = P(N_t = 0, Y_t = 0)$$

$$\pi_1 = P(N_t = 1, Y_t = 0)$$

$$\sigma_0 = P(N_t = 0, Y_t = 1)$$

$$\sigma_1 = P(N_t = 1, Y_t = 1)$$

Figure 2.1

3. (a) (i) For queuing systems define and explain two types of priority schemes known to you. [4]
- (ii) Derive the expression for the mean waiting time of class k traffic in a non-pre-emptive priority scheme. Explain all variables involved in the derivation and clearly state and discuss any assumptions made. [6]
- (b) A Poisson stream of messages with rate λ is fed into a single-channel communication link via a large input buffer. The message stream consists of a random mixture of 1-packet messages and 2-packet messages, each packet being of length B bits. If $\lambda = 1800$ messages/minute, $B = 60$ bits, and 60 % of the messages are single-packet messages:
- determine the overall mean message transit time across the link (i.e. waiting time plus transmission time) when the queue discipline is FIFO and the channel transmission rate is 64 Kbits/seconds. [10]

4. (a) (i) Define the stochastic Knapsack problem. [4]

(ii) *Figure 4.1* represents a Knapsack problem. The resources required by n_1 type traffic is $B1= 1$ [unit] and by n_2 type traffic is $B2= 2$ [units].

What is the capacity of the system ? [2]

Suppose the probabilities associated with the state space (n_1, n_2) are given by:

$$P(i,0) = 0.05 \text{ for } i = 0, \dots, 8$$

$$P(i,1) = 0.05 \text{ for } i = 0, \dots, 6$$

$$P(i,2) = 0.03 \text{ for } i = 0, \dots, 4$$

$$P(i,3) = 0.01 \text{ for } i = 0, \dots, 2$$

$$P(i,4) = 0.02 \text{ for } i = 0.$$

Derive the blocking probability of n_2 type traffic. [4]

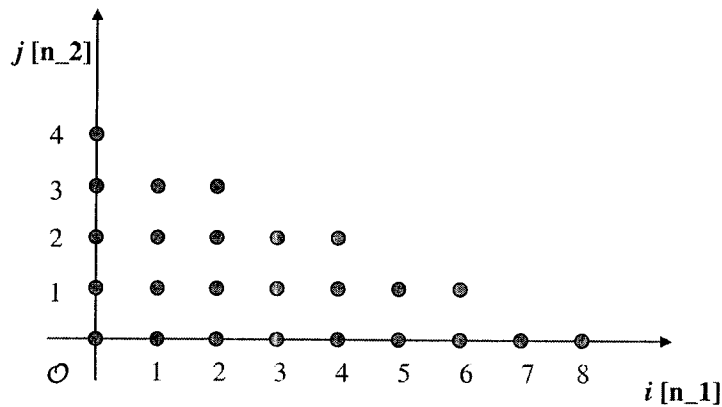


Figure 4.1

(b) In broadband networks, admission and access control schemes play important roles. Discuss briefly the importance of admission and access control in broadband networks. Give examples as necessary. [10]

5. (a) For the N - sources stochastic fluid model represented in *Figure 5.1*:

- (i) explain the meaning of the random variable x , [3]
- (ii) explain the relationship between the parameter αC and the system capacity VC cells/seconds, [3]
- (iii) derive the stationary probability that the buffer occupancy is less than or equal to x given that i sources are in talkspurt. [6]

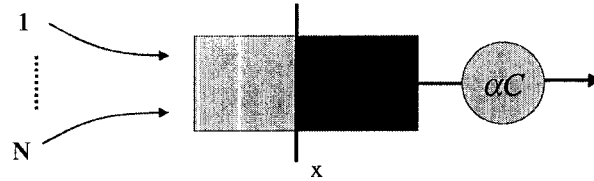


Figure 5.1

(b) A well known measurement of equivalent capacity is given by the following expression:

$$C_l = \min[C_{ls}, C_{lf}]$$

where

$$C_{ls} = mR_p + R_p\sigma\sqrt{-\ln(2\pi) - 2\ln \epsilon}$$

$$C_{lf} = R_p N \left(\frac{1-k}{2} \right) + R_p N \sqrt{\left(\frac{1-k}{2} \right)^2 + kp}.$$

- (i) Explain the meaning of all parameters of C_{ls} . [4]
- (ii) Explain the meaning of all parameters of C_{lf} . [4]

Traffic Theory & Queueing Systems

Question Number etc. in left margin

Mark allocation in right margin

1a
(i)

Engset model.



since there is no buffering j channels busy $\Leftrightarrow j$ sources busy
 \Rightarrow total arrival fall as N increases

Assumptions

(1) Each idle source is a Poisson source :

 $P[\text{source generates new demand in } (t, t+\Delta t) | \text{source is idle}] = \lambda \Delta t$
(2) channel holding times are exponential with mean $(1/\mu)$

(3) full-availability access

(ii)

Birth coefficients

no. of idle sources in state i is $(M-i)$:

$$\begin{aligned} \lambda_i &= (M-i)\lambda \\ &= (1-i/M) M\lambda \\ &= (1-i/M) \lambda_0 \end{aligned} \quad i < i_{\max}$$

Death coefficients (as in Erlang model)

$$\mu_i = i\mu \quad ; \quad i > 0$$

(iii)

$$N < M$$

$$\pi_i = \left(\frac{\lambda_{i-1}}{\mu_i} \right) \pi_{i-1} = \left[\frac{M - (i-1)}{i} \right] \pi_{i-1}$$

$$\pi_i = \binom{M}{i} \alpha^i \pi_0 \quad , \quad i = 1, 2, \dots, N \quad \text{and} \quad \alpha = \left(\frac{\lambda}{\mu} \right)$$

$$\pi_0 = \left[\sum_{j=0}^N \binom{M}{j} \alpha^j \right]^{-1} \quad \alpha = \frac{\rho}{1-\rho}$$

$$\pi_i = \left[\frac{\binom{M}{i} \rho^i (1-\rho)^{M-i}}{\sum_{j=0}^N \binom{M}{j} \rho^j (1-\rho)^{M-j}} \right] \quad i = 0, 1, \dots, N$$

Question Number etc. in left margin

Mark allocation in right margin

- 1b Total callup rate = 1200 calls/hour = 20 calls/min
(i) mean call duration = 180 sec = 3 min
 \Rightarrow offered traffic = $20 \times 3 = 60$ Erlangs
loss probability = $B_L = 0.005$
 \Rightarrow carried traffic = $60(1 - 0.005) = 59.7$ Erlangs
- (ii) From Erlang chart, N is approximately linear in ρ for large ρ .
For $B_L = 0.005$ $N \approx 1.33\rho + 5$
and so for $\rho = 60$ Erlangs:
 $N \approx 85$ channels

Question Number etc. in left margin

Mark allocation in right margin

2(a) Queue length distribution for M/M/K/N system
Local balance equation give

$$\pi_i = \left(\frac{A^i}{i!} \right) \pi_0 \quad 0 \leq i \leq K$$

$$= \left(\frac{A^k}{k!} \right) \rho^{i-k} \pi_0 \quad K \leq i \leq K+B$$

$$\pi_0 = S^{-1} ; S = \left(\frac{A^k}{k!} \left[E_K(A) + \frac{\rho(1-\rho^B)}{1-\rho} \right] \right)$$

$$\pi_0 = \frac{1}{\left(\frac{A^k}{k!} \right) \left[\frac{(1-\rho) E_K(A)}{(1-\rho) + \rho(1-\rho^B) E_K(A)} \right]}$$

$$P[\text{Delay}] = P[\text{all } K \text{ server busy} / \text{buffer not full}]$$

$$= P[K \leq N_t < K+B]$$

$$= \pi_K \left[\frac{1-\rho^B}{1-\rho} \right]$$

$$P[\text{loss}] = P[\text{buffer full}]$$

$$= P[N_t = K+B]$$

$$= \pi_K \rho^B$$

$$\pi_K = \left(\frac{A^k}{k!} \right) \pi_0$$

$$P[Q_t = i | \text{Delay}] = P[Q_t = i | K \leq N_t < K+B]$$

$$= \frac{P[N_t = K+i]}{P[\text{Delay}]} = \frac{\pi_K \rho^i}{\pi_K \left[\frac{1-\rho^B}{1-\rho} \right]} \quad i=0,1,\dots,B-1$$

$$= \rho^i \left[\frac{1-\rho}{1-\rho^B} \right] \quad i=0,1,\dots,B-1$$

- The queue length seen by rejected arrivals is

$Q_t = B$ (i.e. buffer full)

- $P[Q_t = i | \text{Delay}]$: Queue length seen by arrivals which are accepted but delayed.

4.05

MODEL ANSWER and MARKING SCHEME

First Examiner	Paper Code
Second Examiner	Question Page out of
Question labels in left margin	Marks allocations in right margin
<p>24 (i)</p> <p>This B/D process represent an IPP model for a single channel overflow link</p> <p>N_t = no. of busy channel on the overflow link</p> <p>$Y_t = 0$ arrival stream is off</p> <p>$Y_t = 1$ arrival stream is on</p> <p>λ = mean overflow arrival rate</p> <p>μ = holdup time of a call.</p> <p>(ii)</p> <p>$P(\text{overflow traffic is ON}) = \frac{\alpha}{\alpha + \beta} = 0.02$</p> <p>$P(\text{overflow traffic is OFF}) = \frac{\beta}{\alpha + \beta} = 0.98$</p> <p>$P(\text{overflow channel busy}) = \pi_1 + \sigma_1 = 0.115$</p> <p>$P(\text{overflow channel idle}) = \pi_0 + \sigma_0 = 0.885$</p> <p><u>global balance equations</u></p> <p>$\alpha \pi_0 = \mu \sigma_0 + \mu \pi_1$</p> <p>$(\mu + \lambda) \sigma_0 = \alpha \pi_0 + \mu \sigma_1$</p> <p>$(\alpha + \mu) \pi_1 = \mu \sigma_1$</p> <p>$(\mu + \mu) \sigma_1 = \alpha \pi_1 + \lambda \sigma_0$</p> <p>$\pi_0 + \sigma_0 = \left[\frac{\alpha + \mu(1+\delta)}{\alpha(1+\beta) + \mu(1+\delta)} \right]$</p> <p>$\delta = \frac{\lambda}{\mu + \alpha + \mu}$</p>	

Question Number etc. in left margin

Mark allocation in right margin

3a
(i)

Two main type of priority in common use:

- (1) Non-pre-emptive priority: service of an item may not be interrupted by a high-priority arrival
 (2) Pre-emptive priority: service is interrupted by any higher-priority arrival

In this case there are two possibilities

(2a): after the period of interruption, the interrupted service is resumed at the point where it was interrupted - Pre-emptive resume

(2b): after the period of interruption, service starts again from the beginning - Pre-emptive re-start

(ii)

$$\text{class 1} \Rightarrow E(W_1) = E(R) + E(Q_1)E(S_1)$$

$$E(W_1) = \left[\frac{E(R)}{1-\rho_1} \right]$$

$$\text{class 2} \Rightarrow E(W_2) = E(R) + [E(Q_1)E(S_1) + E(Q_2)E(S_2)] + [\lambda_1 E(W_2)]E(S_1)$$

$$E(W_2) = E(R) + [\rho_1 E(W_1) + \rho_2 E(W_2)] + \rho_1 E(W_2)$$

$$E(W_2) = \frac{E(R) + \rho_1 E(W_1)}{1-\rho_1-\rho_2} = \frac{E(W_1)}{1-\rho_1-\rho_2}$$

$$E(W_2) = \frac{E(R)}{(1-\rho_1)(1-\rho_2-\rho_1)}$$

class k

$$(\text{define } \sigma_{k-1} = \sum_{i=1}^{k-1} \rho_i)$$

$$E(W_k) = \frac{E(R)}{(1-\sigma_{k-1})(1-\sigma_k)}$$

Question Number etc. in left margin

Mark allocation in right margin

3b This is an M/G/1 system

$$\text{mean message length} = B \frac{6}{10} + 2 \frac{4}{10} B = 1.2 B$$

$$\text{mean square message length} = \frac{6}{10} B^2 + \frac{4}{10} (2B)^2 = 2.2 B^2$$

$$E(S) = \frac{1.2 B}{64} = \frac{1.2 \times 60}{64} = 1.125 \text{ ms}$$

$$E(S^2) = \frac{6}{10} \frac{60^2}{64^2} + \frac{4}{10} \frac{4 \times 60^2}{64^2} = \left(\frac{60}{64} \right)^2 = 0.8789 \text{ msec}^2$$

$$E(W) = \left[\frac{\lambda E(S^2)}{2(1-\rho)} \right] = \frac{300 \times 0.8789}{2(0.6625)} = 0.19899 \text{ msec}$$

$$E(T) = E(W) + E(S)$$

$$\rho = \lambda E(S) = 1.125 \times 10^{-3} \times 300 = 0.3375$$

Question Number etc. in left margin

Mark allocation in right margin

4a. Stochastic Knapsack

(i) - C resource units

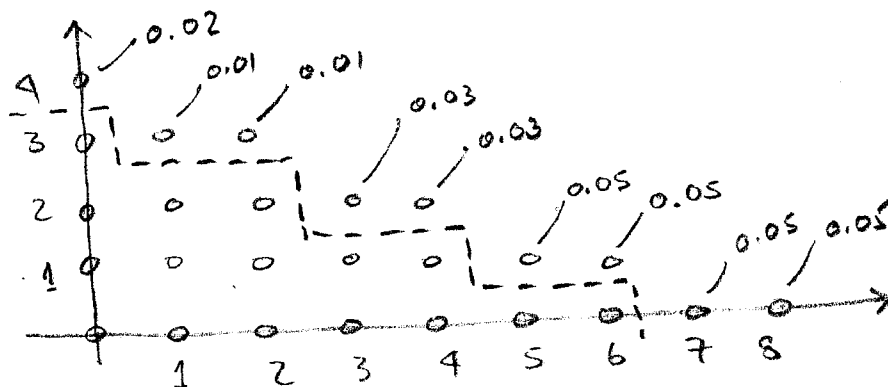
- Arrivals from k classes- Poisson arrival λ_k - Exponential holding time $1/\mu_k$ - hold b_k resources units

- pure loss system

 $n = (n_1, \dots, n_k)$ state of the system $b = (b_1, \dots, b_k)$ - admit class- k arrival iff $b_k < C - b \cdot n$ Dynamic Knapsack problem

$$S = \{n \in \mathbb{I}^k : b \cdot n \leq C\}$$

$$X(t) = (X_1(t), \dots, X_k(t)) \text{ state at } t$$

(ii) $C=8, b_1=1, b_2=2$ (capacity of system = 8)

$$S = \text{all } 0$$

$$S_k = \{n \in S : b \cdot n \leq C - b_k\}$$

$$S_2 = \{n \in S : b \cdot n \leq C - 2\}$$

$$B_k = 1 - \sum_{n \in S_k} \pi(n)$$

$$B_2 = 1 - 0.3 = 0.7$$

Question Number etc. in left margin

Mark allocation in right margin

46 Admission and access control aim at answering the following question: How much traffic can it handle if a prescribed QoS for each traffic class is to be maintained while the network utilisation (i.e. throughput) is to meet some minimum goal?

- Admission control in the context of B-ISDN: given VPs set up in a network, how many virtual connections (i.e. calls with specified QoS) can it handle? given this number, does one admit a new call with specified QoS?

- Discussion on admission policies

Access control: traffic can only be described statistically and therefore congestion may develop despite good Admission policy.

To prevent congestion from occurring, control at the access point as well as within the network must be exerted (also called "policing function")

(ATM refers to user parameter control UPC)

UPC is used to ensure that user do not violate their traffic contract, negotiated during admission control.

- Discussion on e.g. leaky bucket scheme.

Question Number etc. in left margin

Mark allocation in right margin

5a
(i)

- one voice source will generate cell at a rate of $V \frac{\text{cells}}{\text{sec}}$ during a talk spurt of average length $\frac{1}{\alpha} \text{ sec}$
- Increment x by $\frac{V}{\alpha}$ cells during a talk spurt
 $\xrightarrow{\text{unit of information}}$

(ii)

- System with capacity $VC \frac{\text{cell}}{\text{sec}}$ will have an equivalent capacity of $\frac{VC}{\frac{V}{\alpha}} = \alpha C$ "unit of information" per second



- i sources are "pumping" $i\alpha$ "unit of information" to the system $\therefore i\alpha \left[\frac{V}{\alpha} \right] = iV = i \left[\frac{\text{cells}}{\text{sec}} \right]$
- The buffer at the same time empties at the rate αC

(iii)

$$F_i(t+\Delta t, x) = [N-(i-1)] \Delta t F_{i-1}(t, x) + (i+1) \alpha \Delta t F_{i+1}(t, x) + \{1 - [(N-i)\Delta t + i\alpha]\Delta t\} F_i[t, x - (i-C)\alpha \Delta t] + o(\Delta t)$$

- Expand $F_i(t+\Delta t, x)$ and $F_i(t, x - \Delta x)$ ($\Delta x \equiv (i-C)\alpha \Delta t$) in their respective Taylor series and $\Delta t \rightarrow 0$

$$\frac{\partial F_i(x, t)}{\partial t} = [N-(i-1)] F_{i-1}(t, x) + (i+1) \alpha F_{i+1}(t, x) - [(N-i)\Delta t + i\alpha] F_i(t, x) - (i-C) \alpha \frac{\partial F_i}{\partial x}(t, x)$$

$$\dots$$

$$(i-C) \alpha \frac{dF_i(x)}{dx} = [N-(i-1)] F_{i-1}(x) - [(N-i)\Delta t + i\alpha] F_i(x) + (i+1) \alpha F_{i+1}(x)$$

...

$$\frac{dF_i(x)}{dx} D = F_i(x) M$$

Question Number etc. in left margin

Mark allocation in right margin

5b ON-OFF traffic sources

(i) R_p = peak rate $m = Np$ (ON-OFF sources model)

$$\sigma^2 = Np(1-p) = m(1-p)$$

$$E = \int_0^{\infty} \frac{e^{-(x-m)^2/2\sigma^2}}{\sqrt{2\pi}\sigma^2} dx$$

$$\text{if } C-m > 3\sqrt{2}\sigma$$

$$E = \frac{\sigma e^{-(C-m)^2/2\sigma^2}}{\sqrt{2\pi}(C-m)}$$

$$C_L = mR_p + \sigma \sqrt{-\ln(2\pi) - 2\ln E} R_p$$

plus explanation

(ii)

$$P_L = e^{-\beta R_p x / R_p}$$

$$r = (1-p) \left(1 + \frac{\alpha}{\beta}\right) / \left(1 - \frac{C_L}{N R_p}\right)$$

$$r \equiv \frac{\beta x}{R_p} (1-p) \ln \left(\frac{1}{P_L}\right)$$

$$C_{LF} = \left(\frac{1-h}{2}\right) R_p N + \sqrt{\left(\frac{1-h}{2}\right)^2 + h p} R_p N$$