

MSc and EEE/EIE PART IV: MEng and ACGI

## INFORMATION THEORY

**Time allowed: 3:00 hours**

**Answer ALL questions.**

*All questions carry equal marks*

**Examiners responsible**

<b>First Marker(s) :</b>	C. Ling
<b>Second Marker(s) :</b>	D. Gunduz

## The Questions

### I. Basics of information theory.

- a)  $X$  and  $Y$  are correlated binary random variables with  $p(X \neq Y) = 0$  and all other joint probabilities equal to  $1/3$ . Calculate  $H(X)$ ,  $H(Y)$ ,  $H(X|Y)$ ,  $H(Y|X)$ ,  $H(X,Y)$ ,  $I(X;Y)$ .

[6]

- b) Suppose  $x_1$  and  $x_2$  are i.i.d. Bernoulli random variables taking values of 0 and 1 with equal probabilities ( $p = 0.5$ ). Let  $y_1 = x_2$ ,  $y_2 = x_1$ , and  $y_3 = x_1 \oplus x_2$ . Compute the following mutual information:

- i)  $I(x_1; y_1)$
- ii)  $I(x_2; y_2)$
- iii)  $I(x_{1,2}; y_{1,2})$
- iv)  $I(x_1; x_2 | y_3)$

[8]

- c) Consider a Markov process with two states, 0 and 1, and transition matrix

$$T = \begin{pmatrix} 1-p & p \\ q & 1-q \end{pmatrix}.$$

- i) Determine the stationary distribution.
- ii) Calculate the entropy rate,  $H(X)$ .
- iii) Find the values of  $p$  and  $q$  that maximize  $H(X)$ .

[11]

- b) Upper bound on the rate-distortion function. For the case of a continuous random variable  $X$  with mean zero and variance  $\sigma^2$  and squared-error distortion, show that the Gaussian distribution has the largest rate-distortion function, i.e., the rate-distortion function for  $X$  is bounded as follows:

$$R(D) \leq \frac{1}{2} \log \frac{\sigma^2}{D}.$$

Hint: use the following joint distribution of  $X$  and  $\hat{X}$  in Fig. 2.2.

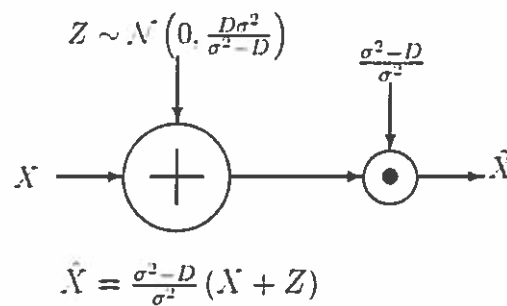


Fig. 2.2. Joint distribution of  $X$  and  $\hat{X}$ .  $X$  and  $Z$  are independent.

[10]

4. Network information theory.

- a) Consider the inference channel in Fig. 4.1. There are two senders with equal power  $P$ , two receivers, with crosstalk coefficient  $a$ . The noise is Gaussian with zero mean and variance  $N$ . Show that the capacity under very strong interference (i.e.,  $a^2 \geq 1 + P/N$ ) is equal to the capacity under no interference at all.

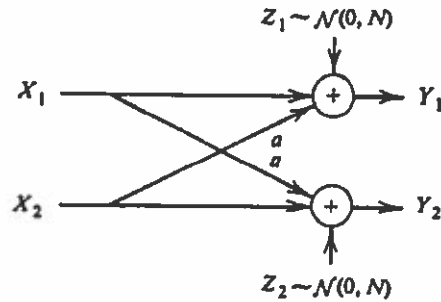


Fig. 4.1. Interference channel.

[10]

- b) Slepian-Wolf coding. Two senders know random variables  $U_1$  and  $U_2$  respectively. Let the random variables  $(U_1, U_2)$  have the following joint distribution:

$U_1 \backslash U_2$	0	1	2	...	$m-1$
0	$\alpha$	$\frac{\beta}{m-1}$	$\frac{\beta}{m-1}$	...	$\frac{\beta}{m-1}$
1	$\frac{\gamma}{m-1}$	0	0	...	0
2	$\frac{\gamma}{m-1}$	0	0	...	0
$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\ddots$	$\vdots$
$m-1$	$\frac{\gamma}{m-1}$	0	0	...	0

where  $\alpha + \beta + \gamma = 1$ . Find the region of rates  $(R_1, R_2)$  that allow a common receiver to decode both random variables reliably.

[15]