

UNIVERSITY OF LONDON
IMPERIAL COLLEGE OF SCIENCE, TECHNOLOGY AND MEDICINE

EXAMINATIONS 2004

BEng Honours Degree in Computing Part I
MEng Honours Degrees in Computing Part I
for Internal Students of the Imperial College of Science, Technology and Medicine

*This paper is also taken for the relevant examinations for the
Associateship of the City and Guilds of London Institute*

PAPER C142

DISCRETE MATHEMATICS

Wednesday 12 May 2004, 10:00

Duration: 90 minutes
(Reading time 5 minutes)

Answer THREE questions

Paper contains 4 questions
Calculators not required

- 1 Let A be a set. Recall the definition of the powerset $\mathcal{P}(A)$ of A given by

$$\mathcal{P}(A) = \{B : B \subseteq A\}.$$

- a Let $A = \{a, \{b\}\}$ and $B = \{a, b\}$. Determine the following sets:

$$A \cup B, \quad A \cap B, \quad A \times B, \quad \mathcal{P}(A), \quad \mathcal{P}(\emptyset), \quad \mathcal{P}(\{\emptyset\}).$$

- b
- i) Give the definition of a binary relation between sets A and B .
 - ii) Let A and B be finite sets. Determine how many possible binary relations there are between A and B , in terms of the cardinality of A and B .
- c Let A and B be arbitrary sets.
- i) Prove that $\mathcal{P}(A) \cap \mathcal{P}(B) = \mathcal{P}(A \cap B)$.
 - ii) Show that $\mathcal{P}(A) \cup \mathcal{P}(B)$ and $\mathcal{P}(A \cup B)$ are not equal by giving a simple counter-example. In fact, one set is contained in the other. Prove this.
 - iii) Show that $\mathcal{P}(A) = \mathcal{P}(B)$ implies $A = B$.

The three parts carry, respectively, 30%, 15%, and 55% of the marks.

- 2a
- i) Define what it means for a binary relation R on set A to be reflexive, symmetric and transitive.
 - ii) Let R be an equivalence relation on A . Define the equivalence class $[a]$ of element $a \in A$ with respect to R .
 - iii) Let \mathbb{Z} denote the set of integers. For $z \in \mathbb{Z}$, define $\lfloor z/2 \rfloor$ to be the integer part of $z/2$: for example, $\lfloor -3/2 \rfloor = -1$. Define a binary relation S on the integers by

$$\forall z_1, z_2 \in \mathbb{Z}. z_1 S z_2 \text{ if and only if } \lfloor z_1/2 \rfloor = \lfloor z_2/2 \rfloor.$$

Show that S is an equivalence relation, and describe the equivalence classes of \mathbb{Z} with respect to S .

- b
- i) Let R be a binary relation between sets A and B . State the conditions required for R to be a function.
 - ii) Let $f : A \rightarrow B$ be a function. Define a binary relation R_f on A by

$$(a_1, a_2) \in R_f \text{ if and only if } f(a_1) = f(a_2).$$

Show that R_f is an equivalence relation.

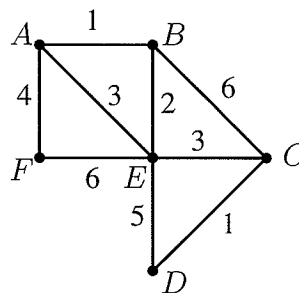
- iii) Let $f : A \rightarrow B$ be an onto function, and R_f be the equivalence relation defined in part 2b(ii). Let $[a]$ denote the equivalence class of a with respect to R_f , and let A/R_f denote the set of equivalence classes generated by the elements of A with respect to R_f .

Define a binary relation S between sets B and A/R_f by

$$b S [a] \text{ if and only if } f(a) = b.$$

Prove that S is a function.

- 3 a
- i) For a graph G with m arcs, what is the relationship between m and the sum S of the degrees of the nodes of G ? Explain your answer briefly.
 - ii) What does it mean for a graph to be a *tree*?
 - iii) State without proof the number of arcs in a tree with n nodes.
 - iv) Let T be a tree. Suppose that $d \geq 2$ is the maximum value of the degrees of nodes of T . Show that T has at least d nodes of degree one.
- b Use Prim's algorithm (starting at node A) to obtain a minimum spanning tree for the following graph:



Give your answer in the form of a diagram. Also state the order in which nodes are added to the spanning tree.

- c Let $G = (G, W)$ be a connected weighted graph, and let T be a minimum spanning tree for G . Let $nodes(G)$ be partitioned into two nonempty sets X and Y . Call an arc a *crossing arc* if it joins a node in X to a node in Y .
- i) Explain why T must contain a crossing arc.
 - ii) Show that at least one crossing arc of minimum weight must belong to T .

The three parts carry, respectively, 45%, 20%, and 35% of the marks.

- 4 a
- i) What is the worst-case number of comparisons for Insertion Sort applied to a list of length n ? Give a brief explanation.
 - ii) Suppose that Insertion Sort is applied to a list L of length $2n$ ($n \geq 1$), which is composed of n distinct numbers in ascending order followed by the same n numbers in the same order. An example for $n = 4$ is $[1, 2, 3, 4, 1, 2, 3, 4]$. How many comparisons does Insertion Sort take when applied to L ? Your answer will be a formula which works for all values of $n \geq 1$.
- b
- The sorting algorithm MaxSort finds the maximum element x in the unsorted portion of the list (initially the whole list), and then swaps x with the last element in the unsorted portion, at which point x is in its correct position and joins the sorted portion. This procedure is repeated until the list is sorted. Thus the unsorted portion lies to the left of the sorted portion.
- i) Write out MaxSort in pseudocode.
 - ii) Use your answer to (i) to calculate the worst-case number of comparisons for MaxSort.
- From the point of view of time complexity measured by comparisons, should we prefer Insertion Sort or MaxSort?
- c
- You are given five coins of identical appearance. They all weigh the same, except for one which is counterfeit and which weighs a different amount from the others. You are given scales of the balance type, allowing you to compare the weights of the coins against each other. Each weighing has two outcomes: the weights are either equal or unequal.
- i) Give an algorithm in the form of a decision tree to determine which is the counterfeit coin. The leaves of the tree will represent outcomes, and the internal nodes will represent weighings. Your algorithm should be *optimal*, in that it uses no more weighings in worst case than necessary.
 - ii) What is the worst-case number of comparisons for your algorithm?
Explain why your algorithm is optimal in worst-case.

The three parts carry, respectively, 35%, 35%, and 30% of the marks.