

UNIVERSITY OF LONDON
IMPERIAL COLLEGE OF SCIENCE, TECHNOLOGY AND MEDICINE

EXAMINATIONS 1999

BEng Honours Degree in Computing Part I
MEng Honours Degrees in Computing Part I
for Internal Students of the Imperial College of Science, Technology and Medicine

*This paper is also taken for the relevant examinations for the
Associateship of the City and Guilds of London Institute*

PAPER 1.3

MATHEMATICAL REASONING – DISCRETE MATHEMATICS

Monday, May 10th 1999, 4.00 – 5.30

Answer THREE questions

For admin. only:
paper contains 4 questions

- 1a i) Define the *Cartesian product* $A \times B$ of two sets A, B . Also, given $R \subseteq A \times B$, $S \subseteq B \times C$, define the *relational composition* $R \circ S$.
- ii) Prove the following law of set algebra:

$$(A \times B) \cap (A' \times B') = (A \cap A') \times (B \cap B').$$

- iii) State with reasons whether each of the following “laws” holds:

$$(A \times B) \cup (A' \times B') = (A \cup A') \times (B \cup B')$$

$$(R \cup S) \circ T = (R \circ T) \cup (S \circ T)$$

- b i) Define the terms *reflexive*, *transitive*, and *anti-symmetric* as applied to a binary relation R on a set A .
- ii) Give an example of a set A with three elements, and a relation R on A which is reflexive but not transitive. Can such an example still be given if the set A is required to have fewer than three elements? (Explain.)
- iii) Determine (by a proof or counterexample) whether the following is necessarily true:

If a relation R on a set B is reflexive, symmetric and anti-symmetric, then R is the relation $=$ (that is, the identity relation on B).

- 2a
- i) Let P be a set with a partial order \leq defined on it. What is meant by saying that an element of P is (a) *least*, (b) *minimal*, with respect to \leq ? Give a simple example to illustrate the difference between the two.
 - ii) Briefly explain the principle of the “topological sort” algorithm for arranging a finite partially ordered set in linear order.
 - iii) List the partial orders on the set $\{a, b\}$ (that is, the distinct relations on this set which are partial orders) which have b as a minimal element.
- b Let a function $f: A \rightarrow B$ be given.
- i) Explain what is meant by the *image* of a set $X \subseteq A$, and by the *pre-image* of a set $Y \subseteq B$, under f .
 - ii) Suppose that $|B| = n$, and that $|f^{-1}(b)| \leq 1$ for all b in B . What can we conclude about $|A|$?
 - iii) State the Pigeonhole Principle, with regard to the function f . How does this relate to your conclusion in part ii) ?
 - iv) Suppose instead that $|A| > 2|B|$. Show that there is an element of B such that at least three distinct elements of A are mapped to it.
 - v) Suppose that, in a network of k computers, each machine is directly linked to at least one, but not more than five, other machines. Show that if $k \geq 11$ there are (at least) three computers with exactly the same number of direct links to other computers.

The two parts carry, respectively, 40%, 60% of the marks.

Turn over ...

- 3a Define the following as applied to (undirected) graphs:
- i) *connected*
 - ii) *simple*
 - iii) *Euler circuit*
 - iv) *degree* of a node
- b For the purposes of this question, let an *RB graph* be defined as a graph where each arc (edge) is coloured either red or blue. An RB Euler circuit is an Euler circuit where successive arcs have different colours.
- i) Draw two simple connected RB graphs G_1 and G_2 , satisfying the following:
 - G_1 has an Euler circuit but no RB Euler circuit.
 - G_2 has an RB Euler circuit.
- Label each graph clearly. For convenience, you can represent red arcs by solid lines and blue arcs by dotted lines.
- ii) Suggest a condition on RB graphs which is necessary and sufficient for a connected RB graph to have an RB Euler circuit.
 - iii) Explain why your condition in (ii) is necessary.
 - iv) Give an informal description of an algorithm which, given a connected RB graph satisfying your condition in (ii), constructs an RB Euler circuit.

The two parts carry, respectively, 25%, 75% of the marks.

- 4a i) Describe the algorithm Insertion Sort.
- ii) What is the worst case number of comparisons made by Insertion Sort when sorting a list of length n ?
- b i) Briefly describe the algorithm Quicksort.
- ii) Give an example of a list for which Insertion Sort performs better than Quicksort, explaining your answer briefly.
- iii) State a recurrence relation for the average number of comparisons made by Quicksort when sorting a list of length n . Do not solve the recurrence relation.
- c A sorted list of length $2n$ is disarranged to create a new list L by placing the elements in the odd positions (still in order) before the elements in even positions (still in order). Thus e.g. the list $[1,2,3,4,5,6]$ would become $[1,3,5,2,4,6]$.
- i) How many comparisons does Insertion Sort take when applied to L ? Explain your working.
- ii) How many comparisons does Quicksort take when applied to L ? Assume that Quicksort always splits around the first element. Explain your working.

The three parts carry, respectively, 25%, 35%, 40% of the marks.

End of paper