## DISCRETE MATHEMATICS AND COMPUTATIONAL COMPLEXITY

## 1. [Compulsory]

Let  $\mathcal{R}$  be the set of all relations on a set A.

Express the predicate P(x), meaning that  $x \in \mathcal{R}$  is a transitive relation, in terms of appropriate symbolic logic.

[2]

b) Prove that the proposition p given by  $\forall R (P(R) \leftrightarrow (\forall n \in Z^+ R^n \subseteq R))$  is true, where the universe of discourse is  $\mathscr{R}$ .

[ 12 ]

c) Prove that the proposition q given by  $\forall R (P(R) \rightarrow (\forall n \in Z^+ ((n \ge 2) \rightarrow P(R^n))))$  is true, where the universe of discourse is  $\mathcal{R}$ .

[ 10 ]

d) Replacing the implication in the definition of q by its converse yields another proposition r. Prove that r is false.

[8]

e) A relation R' is said to be the transitive closure of R when R' is the smallest transitive relation containing R. Define the connectivity relation  $R^*$  and prove that  $R^* = R'$ .

[8]

lere the convers of discourse is A. 1. YR P(R) (Ynezt Rns R) First prove  $\forall R[P(R) \rightarrow (\forall n \in 2^+ R^n \subseteq R)]$ for 1=1, we have  $\forall R[P(R) \rightarrow R \subseteq R]$  which is time as the RHS is always time. Assure true for n & use induct for n>1. Let (2,5) & R'N+1 = R. RN. =) ]a (a, x) e R n (2,5) e R^] line R'ER, (a,6) & R Daniel Rntisk D record years YR [ (An & Z+ R^SR) > P(R)] To be jule, we would need R s.t. ( $\forall n \in \mathbb{Z}^+ \ R^n \subseteq R$ )  $\wedge \neg P(R)$ . YNEZ RUSR => RZSR let (a, b) & R , (b, c) & R. Then (a, c) & K2 CR ⇒ (a,c) FR ⇒ P(R) D

YR (AR) → (Yn EZ\* ((n>2) → P(R^n))) the will more ARP(R) > YNEZ+ P(R)], Arried for the p(R) = P(R), while is time.

Arrived for the p(R2n) and p(R2n+1) are time. First P(R2n) Let  $(x,b) \in \mathbb{R}^{2n} = \mathbb{E} \mathbb{R}^n \cdot \mathbb{R}^n$   $(b,c) \in \mathbb{R}^{2n} = \mathbb{R}^n \cdot \mathbb{R}^n$ Then  $\exists x \exists y \lceil (x,x) \in \mathbb{R}^n \mid (x,b) \in \mathbb{R}^n \mid (b,y) \in \mathbb{R}^n$  $\Lambda$   $(\gamma, c) \in \mathbb{R}^{N}$ rince Kn is truité, (a, b) ERN , (b, c) ERN  $\Rightarrow (a,c) \in \mathbb{R}^n \cdot \mathbb{R}^n = \mathbb{R}^{2n}$ xcond, P(R2n+1) Let  $(a,b) \in \mathbb{R}^{2n+1} = \mathbb{R} \cdot \mathbb{R}^n \cdot \mathbb{R}^n$  at  $(b,c) \in \mathbb{R}^{2n+1} = \mathbb{R}^n \mathbb{R}^n \mathbb{R}^n$ .

Let  $(a,b) \in \mathbb{R}^{2n+1} = \mathbb{R} \cdot \mathbb{R}^n \cdot \mathbb{R}^n$  at  $(b,c) \in \mathbb{R}^n = \mathbb{R}^n \mathbb{R}^n \mathbb{R}^n$ . 1 (5,9) ER^ 1 (7,9) ER^ 1 (9,0) ER^ me  $R^n$  is familie,  $(a, p) \in R^n$   $(y, c) \in R^n$ .  $= (A,C) \in \mathbb{R}^n \cdot \mathbb{R} \cdot \mathbb{R}^n = \mathbb{R}^{2n+1} \square$ 

∀R (Hn € 2 + ((n > 2) > P(R))) > P(R)) o-sider R = {(a, b), (b, c)}. R= /(ac)?  $R'' = \phi$  pr n > 2. p(R) is the 5t  $p(R^2)$ ,  $p(R^3)$ , etc. are the.  $\forall n \in \mathbb{Z}^+$   $(n \ni 2) \Rightarrow p(R^n)$  is  $\forall R \cup E$ . ~ 3R ( 4n ∈2+ ((n ?, 2) → P(R))) N ¬P(R)) INDE = original proposite or is take. R\* = RUR2 UR3 U... ried to moe R'=R\*

RIR\* directly:
Need to show R'=S wherever RES, Strinte ty (a, b) ∈ R° , (b, c) ∈ R° it pllors for (+) Ht (n c) ER". Gt RES. Pu. P(S") , S" S S. RES => R\* SS\*. R'SS'SS D.

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