

IMPERIAL COLLEGE LONDON

DEPARTMENT OF ELECTRICAL AND ELECTRONIC ENGINEERING  
EXAMINATIONS 2016

MSc and EEE/EIE PART IV: MEng and ACGI

**DISCRETE-EVENT SYSTEMS**

**Corrected copy**

Monday, 9 May 10:00 am

Time allowed: 3:00 hours

**There are FOUR questions on this paper.**

**Answer ALL questions.**

*All questions carry equal marks*

**Any special instructions for invigilators and information for candidates are on page 1.**

Examiners responsible      First Marker(s) :      D. Angeli  
Second Marker(s) :      E.C. Kerrigan



1. A buffer has a maximum capacity of 6 units. Packets arrive in blocks of 2 or 4 units (events  $a_2$  and  $a_4$ , respectively) and depart in blocks of 3 units (event  $d_3$ ).
  - a) Build a finite deterministic automaton  $G$  that models the arrival and departure of packets as well as buffer occupancy, assuming an empty buffer at initial time. [ 4 ]
  - b) Assume that events  $a_2$  and  $a_4$  are partially observable, viz. only an  $a$  event is detected but no information on packet size is available. Model this as a non-deterministic automaton  $G_N$ . [ 2 ]
  - c) Denote by an even interval (or an odd interval) any set of integers of type  $\{n, n+2, \dots, n+2k\}$ , for some  $k$  non-negative and integer and  $n$  even (respectively odd). Let  $n$  and  $n+2k$  be in the range  $\{0, \dots, 6\}$  so as to make physical sense as occupancy levels of the buffer. Show that the state transition map of  $f_N(\cdot, a)$  of  $G_N$  maps even intervals to even intervals and odd intervals to odd intervals. Instead,  $f_N(\cdot, d_3)$  maps even intervals to odd intervals and vice versa. (*Hint: write explicit expressions for the map*). [ 4 ]
  - d) Build an observer automaton  $G_O$  for  $G_N$ . [ 7 ]
  - e) Remark from the graph of  $G_O$  that, regardless of the arrival and departure sequences, one can never be certain to have exactly four units of space taken in the buffer. Can you explain by alternative types of considerations why this is the case ? [ 3 ]

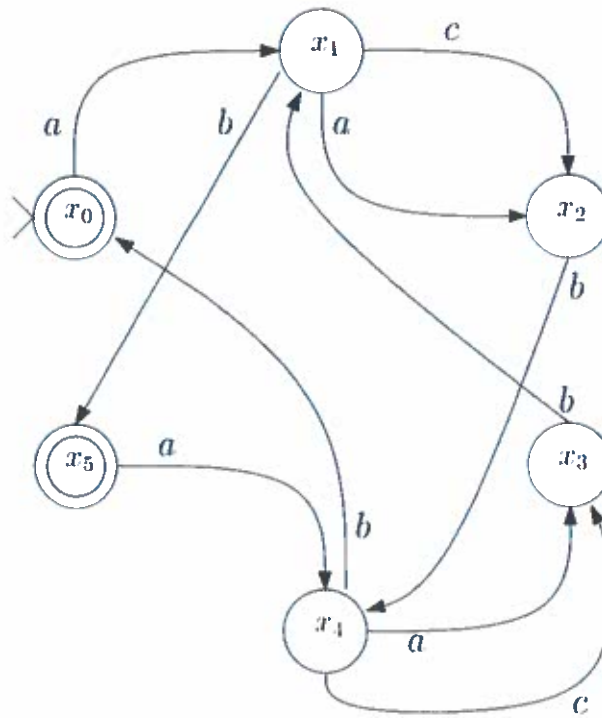


Figure 2.1 The transition diagram of  $G$

2. Consider the Automaton  $G$  shown in Fig. 2.1, with state space  $X = \{x_0, x_1, x_2, x_3, x_4, x_5\}$ , and event set  $E = \{a, b, c\}$ .
  - a) Find in  $X$  the equivalence classes of  $\sim$ , where  $x_a \sim x_b$  denotes  $\mathcal{L}_m(G/x_a) = \mathcal{L}_m(G/x_b)$  and  $\mathcal{L}(G/x_a) = \mathcal{L}(G/x_b)$ . [ 5 ]
  - b) Find a minimal automaton  $G_{\min}$  equivalent to  $G$ , in the sense that  $\mathcal{L}_m(G) = \mathcal{L}_m(G_{\min})$  and  $\mathcal{L}(G) = \mathcal{L}(G_{\min})$  while the cardinality of  $X_{\min}$  (the state space of  $G_{\min}$ ) is as small as possible. [ 5 ]
  - c) Build an automaton generating the language  $K \subset \{a, c\}^*$  corresponding to the following verbal specification: “between two  $a$  events there should always be at least one  $c$  event, and between two  $c$  events there should always be at least one  $a$  event”. [ 3 ]
  - d) Assume that  $E_{uc} = \{c\}$ . Is the specification modeled by  $K$  controllable with respect to  $\mathcal{L}(G)$ ? [ 2 ]
  - e) Design an automaton  $G_S$  that can represent a supervisor  $S$  so that  $\mathcal{L}(S/G_{\min})$  is the supremal controllable sublanguage of  $K$  with respect to  $\mathcal{L}(G)$ . [ 5 ]

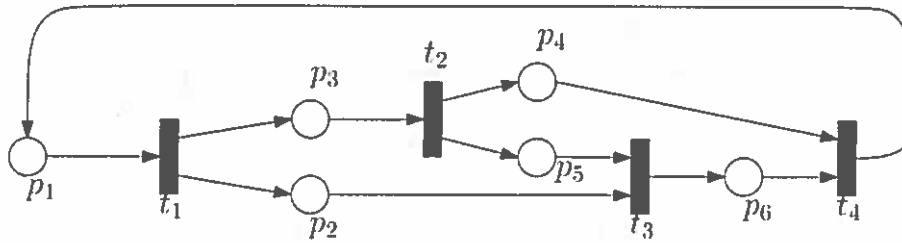


Figure 3.1 The Petri Net  $N$

3. Consider the Petri Net  $N$  shown in Fig. 3.1.
  - a) Compute the incidence matrix  $C$ , when ordering places and transitions according to the numbering provided in Fig. 3.1. [ 3 ]
  - b) Compute all the minimal  $P$ -semiflows of the net. [ 5 ]
  - c) Is the network structurally bounded? (Justify your answer). [ 3 ]
  - d) Consider the initial marking  $M_0 = [1, 0, 0, 0, 0, 0]^T$ . Find the reachable set  $\mathcal{R}(M_0)$  and sketch the transition diagram of the associated automaton  $G_N$ . [ 3 ]
  - e) Consider two automata  $G_a$  and  $G_b$  that have the same structure of  $G_N$  and associated event sets  $E_a = \{t_1^a, t_2^a, t_3^a, t_4^a\}$ ,  $E_b = \{t_1^b, t_2^b, t_3^b, t_4^b\}$  (notice that  $E_a \cap E_b = \emptyset$ ). Compute their parallel composition  $G_a || G_b$  and draw its transition diagram. [ 3 ]
  - f) Do you expect the previous graph to be isomorphic to the one arising by taking an initial marking  $M_0 = [2, 0, 0, 0, 0, 0]^T$ ? (Justify your answer) [ 3 ]

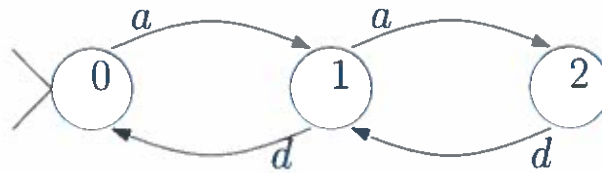


Figure 4.1 Queue automaton  $G$

4. Consider the automaton  $G$  shown in Fig. 4.1. Assume both events  $a$  and  $d$  (arrivals and departures) happen at random times with exponentially distributed clocks of rate  $\lambda$ .
  - a) Model the situation in which the initial state of the queue is unknown by means of a non-deterministic automaton  $G_N$  (Hint: add an extra state playing the role of initial state and connect it to the remaining states in a suitable way). [ 4 ]
  - b) Build a deterministic automaton that, by observing  $a$  and  $d$  events alone and assuming an arbitrary initial state, keeps track of an estimate of all possible current states of the queue. [ 4 ]
  - c) Considering the observer automaton designed at the previous step and the occurrence rates of  $a$  and  $d$  events sketch the transition diagram and transition matrix  $Q$  of the associated Markov Chain. [ 3 ]
  - d) Argue that with probability 1 the state of the system will be correctly estimated (without uncertainty) after a finite time. [ 3 ]
  - e) Compute the average time it takes to have a correct state estimate. [ 3 ]
  - f) Would you expect this to happen for any deterministic automaton  $G$ ? (Justify your answer. Hint: Would minimality of the automaton play a role?) [ 3 ]



