DEPARTMENT OF ELECTRICAL AND ELECTRONIC ENGINEERING **EXAMINATIONS 2009**

EEE/ISE PART II: MEng, BEng and ACGI

COMMUNICATIONS 2

Corrected Copy

Wednesday, 10 June 2:00 pm

Time allowed: 2:00 hours

There are FOUR questions on this paper.

Q1 is compulsory. Answer Q1 and any two of questions 2-4. Q1 carries 40% of the marks. Questions 2 to 4 carry equal marks (30% each).

Any special instructions for invigilators and information for candidates are on page 1.

Examiners responsible

First Marker(s): K.K. Leung, K.K. Leung

Second Marker(s): J.A. Barria, J.A. Barria

Special Instructions for Invigilator: None

Information for Students: None

Communications II

Page 1 of 5

1. This is a compulsory question.

a.	 Consider two random variables X and Y. i. If X and Y are uncorrelated, express E[XY] in terms of E[X] and E[Y]. ii. Comment on the relationship between the following two statements: "X and Y are independent" and "X and Y are uncorrelated". iii. A property is defined as follows: If X and Y are uncorrelated, then X and Y are independent. Name one form of the joint probability distribution functions (PDFs) for X and Y that possesses such a property. 	[2] [3] [2]
b.	 Channel noise. i. What are the three most important parameters that define the capacity of a communication channel? ii. Name the theoretical principle (theorem) that supports the validity of modeling channel noise as a Gaussian distribution in most cases and why? iii. What is it meant by "white noise"? iv. What are the <i>in-phase</i> and <i>quadrature-phase components</i> (which are denoted by n_c(t) and n_s(t), respectively) of the band-limited white noise n(t) where the centre of the frequency band is ± f_c Hz? v. Compare the average power of the noise waveforms: n(t), n_c(t) and n_s(t). 	[2] [4] [3] [4] [2]
c.	 Digital communications. i. To transmit a signal with a bandwidth of W Hz in a digital communication system, what is the minimum sampling rate so that the original signal can be recovered at the receiver and why? Use a frequency-domain diagram to explain. ii. Let Δ be the uniform gap between two adjacent quantizing levels and assume that the quantization error (noise) is uniformly distributed from -Δ/2 to Δ/2. Write down the probability density function (pdf) for the quantization noise. iii. Using the result in the above part ii, derive the quantization noise power. 	[4] [2] [5]
d.	 Source entropy. i. For a communication system with three-symbol alphabet {A, B, C}, the probabilities of transmitting these symbols, P_A, P_B and P_C, are 0.7, 0.2 and 0.1, respectively. Calculate the source entropy of the system. ii. What is the main objective of the Huffman coding algorithm? 	[4] [3]

2. Noise effects on amplitude modulation.

Consider an amplitude modulation (AM) signal s(t) with the following notation:

- A is the amplitude of the carrier,
- f_c is the carrier frequency,
- m(t) is the message (information) signal,
- n(t) is the noise signal,
- $n_c(t)$ is the in-phase component of the noise signal,
- $n_s(t)$ is the quadrature phase component of the noise signal, and
- N_o is the power spectral density for the noise in the baseband transmission.
- a. Provide an expression for the AM signal s(t). [3]
- Assuming that noise is negligible, state and explain a condition under which an
 envelop detection method can be employed to recover the message signal at the
 receiver.
- c. What is "synchronous detection"? What is the key technical difficulty in using the synchronous detection to receive the AM signal? [3]
- d. Draw a block diagram to depict the AM transmission, the additive-white-Gaussiannoise (AWGN) channel and the receiver structure using the synchronous detection. [3]
- e. Give an expression for the pre-detection signal as a function of the AM signal, $n_c(t)$, and $n_s(t)$ of the noise signal. [3]
- f. Let the signal power at the receiver output be denoted by $P = E\{m^2(t)\}$. Derive the signal-to-noise ratio (SNR) of the receiver output for the synchronous detection. [3]
- g. Derive a relationship between the SNR of the receiver output for the synchronous detection and that for the baseband transmission with the same transmitted power. [5]
- h. Now consider the envelop detection for the AM signal with an assumption that noise is much smaller when compared with the carrier and message signal. By approximating the envelop of the pre-detection signal under the assumption, derive a relationship between the SNR of the receiver output for the envelop detection and that for the baseband transmission with the same transmitted power. [5]
- i. Compare and comment on your results in parts g and h. [2]

3. a. Digital communications.

Consider a digital communication channel that transmits a polar signal of 0 or 1. When a 0 and 1 are sent, the pulse amplitudes are $-A_P$ and A_P , respectively. The channel has additive zero-mean white Gaussian noise with the rms value of the noise denoted by σ_n (i.e., the noise power is σ_n^2). Assume that the detection threshold T is θ at the receiver. That is, if the received signal is equal to or great than θ , the signal is considered to be a 1. Otherwise, a 0 is considered to be sent.

i. Given that a 0 is being sent, derive the error probability $P_E(0)$ in terms of A_P , σ_n and

$$Q(y) = \frac{1}{\sqrt{2\pi}} \int_{y}^{\infty} e^{-x^{2}/2} dx.$$
 [4]

- ii. Similarly, given that a 1 is being sent, obtain the error probability $P_E(1)$ in terms of A_P , σ_n and Q(y).
- iii. Using results in parts i and ii and assuming equal probability for sending 0 and 1, obtain the error probability P_E in terms of A_P , σ_n and Q(y). [3]
- iv. Now, let the probabilities of sending a 0 and 1 be denoted by p_0 and p_1 where $p_1 > p_0$. Discuss whether the new detection threshold T for $p_1 > p_0$ should be greater or less than 0 in order to minimize the overall error probability. Provide a physical interpretation for your answer. [4]

b. Power spectral density.

Consider a linear system with the input, output and impulse-response functions denoted by x(t), y(t) and h(t), respectively. It is known that y(t) = h(t) * x(t) where * represents convolution. The corresponding relationship in Fourier transforms is Y(f) = H(f)X(f). Let $R_x(\tau)$ and $R_y(\tau)$ be the respective autocorrelation function for x(t) and y(t). Furthermore, let $P_x(f)$ and $P_y(f)$ be the power spectral density (PSD) for x(t) and y(t), respectively. Recall that the PSD is the Fourier transform of the autocorrelation function of a given signal.

- i. Using the definition of autocorrelation function, namely $R_y(\tau) = E[y(t)y(t+\tau)]$, and the fact that y(t) = h(t) * x(t), prove that $R_y(\tau) = h(-\tau) * u(-\tau)$ where $u(-\tau) \equiv h(\tau) * R_x(\tau)$.
- ii. Based on the result in part i above, show that $P_{\nu}(f) = |H(f)|^2 P_{\nu}(f)$. [5]

[10]

4. Error control code.

a. For any error control code, express the maximum number of errors the coding scheme can detect in terms of the minimum Hamming distance (denoted by d_{\min}) and explain why.

[3]

b. To continue with part a, express the maximum number of errors t the coding scheme can correct in terms of d_{\min} and explain why.

[4]

c. For the (7,4) Hamming code, the generator matrix is given by

$$G = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{pmatrix}$$

i. Using the generator matrix G, calculate the code words for two source blocks: 1010 and 1101.

[4]

ii. Obtain the Hamming distance between the code words obtained in part i.

[2]

iii. Write the parity check matrix H associated with the above generator matrix G.

[5]

iv. For a received code word y, which may contain errors, describe how the associated syndrome s can be determined.

[2]

v. For the error vector 0000010, find the corresponding syndrome.

[3]

d. For five symbols S_0 , S_1 , S_2 , S_3 , S_4 of the alphabet of a discrete memory-less source, their probabilities are 0.4, 0.2, 0.2, 0.1 and 0.1, respectively. Use the Huffman encoding algorithm to generate the code words for the source.

[7]

1. a. i. Given x, y and uncorrelated, Communications T $E[(X-m_X)(Y-m_Y)] = 0 \qquad \sum_{z \in Y} C_{z} C_{z$

ii. If "x & Y are independent," then "x and Y are uncorrelated."

However, the converse is not necessarily true.

That is, if "x & y are uncorrelated," "x
& Y are not necessarily independent.

vii. One such PDF is normal distribution for X2 X.

1. b. i.

The three parameters are:

bandwidth, signal power and noise power.

ii

The Central Limit Theorem States that the normalized sum of independent R.V.'s has a Gaussian distribution. Since channel noise is originated from many independent sources, thus the total noise behaves like a Gaussian distribution as The Theorem suggests.

iii. Whire noise has the constant power sprectrum dension (PSD) over all frequencies.

iv. The power spectral clarsing of the band-limited white noise n(t) is depicted as follows.

The bandpass represtation of n(t) is

nlt) = nelt) cos (27fct) - ns(t) sin (27fct)

where Nott) = 2 aucoz(27(fre-fo)t + On)

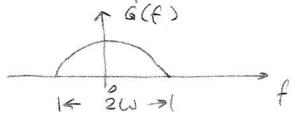
Nott) = 2 Qu Sin (2T (fu-fo) t + Ob).

. b. iv. In words, na(t) is the baseband noise (Cont'd) component that is in-phase with the carrier at frequency fc, while 115th) is that conjunent that is 90° out of phase with the carrier.

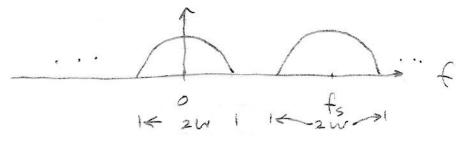
The average of n(t) nc(t) and ns(t) is identical (because nc(t) & ns(t) are 90° our of phase with each other, although $n(t) = n_E(t)\cos(2\pi f_c t) - n_S(t)\sin(2\pi f_c t)$. 1. C.i. The minimum sampling rate is 2W Hz, referred to as Myquist rate.

The physical interpretation of the Myquist rate is given as follows.

The frequency domain representation of the signed



The frequency domain representath of the sampled sopnal:



From the above diagram, it is clear that if the sampling frequency for is greater than 2W to, the replica of the signal do not overlap with each other i.e., with no changes to the signal due to the sampling operations. Therefore, the original signal can be recovered by a low-pass filter.

1.C. ii. Let N denote the noise level and falt) be the paf for the noise.

Then

fult)= { /a for 1/2 st s 1/2 o otherwise

iii. Noise power PN = E((N-E(N))2)

Since E(N) = 0, PN = E(N2)

Thus, $PN = E(N^2) = \int_{-b/2}^{0/2} t^2 \int_{-b}^{2} dt^2$

 $\Rightarrow P_N = \frac{1}{4} \cdot \frac{t^3}{3} \int_{-\Delta/2}^{\Delta/2}$

 $\Rightarrow P_{N} = \frac{1}{\Delta} \cdot \left[\frac{\Delta^{3}}{8} + \frac{8^{3}}{8} \right] \cdot \frac{1}{3}$

 $\Rightarrow P_N = \frac{\Delta^2}{1^2}.$

$$|+(s)| = -\frac{2}{4} P_{4} \log_{2} P_{4}$$

$$\Rightarrow |+(s)| = -\frac{3}{2} P_{4} \log_{2} P_{4}$$

$$= -\left[0.7 \log_{2} 0.7 + 0.2 \log_{2} 0.2 + 0.1\right]$$

$$+ 0.1 \log_{2} 0.1$$

$$\Rightarrow H(s) = 0.7 \times 0.5(5 + 0.2 \times 2.322 + 0.1 \times 3.322 + 0.1 \times$$

ii. The main objective of the Huffman Coding algorithm is to choose codeword lengths for a given tource alphabet so that more probable sequences have shorter codewads. On the contrary, less likely sequences will have longer codewads.

2. a. S(t) = [A + m(t)] Cos (217fot)

b. Let mp = max | m(6) |

With negligible vision a condition under which the envelop detection can be employed to receive the signal is.

 $M = \frac{mp}{A} \le 1 \iff A \ge mp$

That is, the amplitude of the camier is larger than the maximum (peah) amplitude of the message signal. Given Azmp, the envelop of the received signal is the message information.

C. Synchronous eletection is a eletection method where the receiver multiplies the received signal with the carrier with exactly identical phase and frequency of that at the transmitter, in order to recover the transmitted Mynal.

The key technical challenge for synchronous detections is to generate the carrier with precially identical frequency and phase of the carrier at the transmitting side.

e. Pre-detection signal
$$\chi(t) = S(t) + n(t)$$

$$= \left[A + m(t)\right] \cos (2\pi f_0 t) + n(t)$$

$$\Rightarrow \chi(t) = \left[A + m(t) + n_0(t)\right] \cos (2\pi f_0 t)$$

$$-n_s(t) \sin (2\pi f_0 t)$$

The signal power at the receive output P = E[m2(+)]

Noise your: PN = 2 NoW (: used by Am

g. Transmitting power PT = E[S(4)]=E[(A+m(6))] $\Rightarrow P_T = \frac{A^2}{2} + E[m^2(6)]$

Therefore, the SNR for baseband transmissing SNR baseband = PT NoW & Bandwidth is w for baseband transmission.

 \Rightarrow SNR baseband = $\frac{A^2 + P}{2}$. $\frac{1}{N_0 W}$

From part f, SNRAM/sync = P 2 No W

Therefore SNR AM/sync = PA+P · SNR baseband

:. h. The envelop of the signal (xxx)

 $y(t) = \sqrt{(A + m(t) + n_c(t))^2 + n_s(t)^2}$

When noise is much smaller than the carrier and message signal,

| A+ m(t) + no(t) | >> | Ns(t) |

Thus, y(t) = A + m(t) + nc(t)

Based on this approximation.

SUR env = [M(6)] = P 2NOW = 2NOW. Since SNR baseband = A2+P 2NoW

Thus, SNR env = P SNR baseband

Results in parts g & h reveal that

SNRenv = SNR Am/sync.

That is, for low noise, the envelop detection performs very closely to that of synchronom detection.

Since the noise has a zero-mean Gaussian distribution,

$$PE(0) = P_{-}(n > Ap)$$

$$= \frac{1}{\sqrt{2\pi}} \int_{0}^{\infty} e^{-n^{2}/2C_{n}^{2}} dn$$

$$= \frac{1}{\sqrt{2\pi}} \int_{0}^{\infty} e^{-x^{2}/2} dx$$

$$= \frac{1}{\sqrt{2\pi}} \int_{0}^{\infty} e^{-x^{2}/2} dx$$

$$\Rightarrow PE(0) = Q(\frac{Ap}{a})$$

ii. Similarly,

$$P_{E}(1) = P_{r}(A_{p+n} < 0)$$

$$= P_{r}(n < -A_{p})$$

$$= P_{E}(1) = \frac{1}{\sqrt{2\pi} \sqrt{n}} \int_{-\infty}^{-A_{p}} e^{-n^{2}/2\sqrt{n}^{2}} dn$$

Since $e^{-n^2/2\sigma_n^2}$ is symmetrical about The axis of n=0, $PE(1)=\frac{1}{\sqrt{2\pi}\sigma_n}\int_{A_0}^{\infty}e^{-n^2/2\sigma_n^2}dn$

$$P_E = P_E(0) \cdot P_0 + P_E(1) P_1$$

$$= Q(\frac{A_E}{C_n}) \cdot \frac{1}{2} + Q(\frac{A_E}{C_n}) \cdot \frac{1}{2}$$

$$\Rightarrow P_E = Q(\frac{A_E}{C_n}).$$

iv. When PI > Po, The detection thrushold T should be set less than O.

This is so because with P. > No, the received signal is likely to be greater than zero. When a 1 is transmitted, setting T<0 will help reduce the probability of lower. On the other hand, setting T<0 will increase the probability of error for sending a 0. At balance, given P.> p., with T<0, the reduction of the error prob for sending a 1 will be well compensated for the increase of the lover prob for sending a 1 The lover prob for sending a 0. That is, setting T<0 will reduce the ornal error prob, given P. > Po.

3.6. i.
$$R_{y}(z) = E \left[g(t) \ y(t+t) \right]$$

$$= E \left[\int_{-\infty}^{\infty} h(t_{t}) \ x(t-t_{t}) dt_{t} \right] \int_{-\infty}^{\infty} h(t_{t}) x \left(t+\tau-t_{t} \right) dt_{t}$$

$$\Rightarrow R_{y}(\tau) = E \left[\int_{-\infty}^{\infty} \int_{h(t_{t})}^{h(t_{t})} h(t_{t}) x \left(t-t_{t} \right) x \left(t+\tau-t_{t} \right) dt_{t} dt_{t} \right]$$

$$\Rightarrow R_{y}(\tau) = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} h(t_{t}) h(t_{t}) E \left[x(t-t_{t}) x \left(t+\tau-t_{t} \right) dt_{t} dt_{t} \right]$$

$$\Rightarrow R_{y}(\tau) = \int_{-\infty}^{\infty} h(t_{t}) h(t_{t}) f(t_{t}) f(t_{t}) f(t_{t}) dt_{t} dt_{t}$$

$$\Rightarrow R_{y}(\tau) = \int_{-\infty}^{\infty} h(t_{t}) h(\tau+t_{t}) f(\tau+t_{t}) dt_{t} dt_{t}$$

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$$\Rightarrow R_{y}(\tau) = \int_{-\infty}^{\infty} h(t_{t}) f(\tau+\tau+t_{t}) dt_{t} dt$$

3.6. ii.

Ry(-t) = h(-t) * h(t) * Rx(t)

Take the Fourier transferms on hoth sides of the above. We obtain

 $P_y(f) = \mathcal{F}[h(-\tau)] \cdot H(f) \cdot P_x(f)$

By the property of Fourier transform. F[h(-t)] = H*(f) Where * is the complex conjugate

Substitute this into the expression for Py (f),

Py(f) = H*(f) H(f) · Px(f)

 $\Rightarrow P_{y}(f) = \left| H(f) \right|^{2} P_{x}(f).$

4. a. The number of errors that can be detected is donin -1.

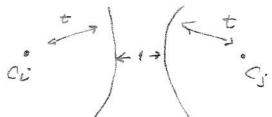
This is so because then must be at least dain errors that can turn an input code into another different but valid code and, escaping the detection mechanism.

For any input codeword with less than down errors, the visultant victor cannot be a valid code ward, thus enabling detection.

6. The number of errors to that can be corrected

is $t = \lfloor \frac{d_{min}-1}{z} \rfloor$ where LXI is the largest
integer that is smaller or equal to X.

Consider the coderads Ci and G.



when there are at most t emoss for both Ci and Cj, the resultant codewords still deffer from each other by one Hammi's distance. Therefore, the enroneous codewords can be corrected (or treated) as the "newest" to the valid codewords.

4. C. i. The output codeword

$$x = u \cdot G$$
 where u is the source

code block

$$\begin{pmatrix}
1010
\end{pmatrix}
\begin{pmatrix}
1 & 0 & 0 & 0 & 0 & 1 & 1 \\
0 & 1 & 0 & 0 & 0 & 1 & 1 \\
0 & 0 & 0 & 0 & 1 & 1 & 1
\end{pmatrix}$$

$$(1101)$$
 $\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 \end{pmatrix}$

ii. The Itamming distance between the obtained codewards

Hamming distance = 4.

iii. The parity check matrix H

$$H = \begin{pmatrix} 0 & 1 & 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 \end{pmatrix}$$

iv. The syndrome $s = y \cdot H^T$ where H is the parity check matrix.

4.0. v. Syndrome
$$S = eH^T$$

So, $S = (0000010) \begin{pmatrix} 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix}$

$$\Rightarrow S = (010)$$

4.d.

Symbol	Stage 1	2	3	4	Code
So Si Si Si Si Si	0.4 0.2 0.2 0.1 0.1	0.4		and a second	J 00 10 11 010

Answerp.