

Span VM: Multi-Hypervisor Virtual Machines Enabling An Ecosystem of Hypervisor-Level Services In Cloud



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Problem: Support for 3rd-Party Hypervisor-level Services

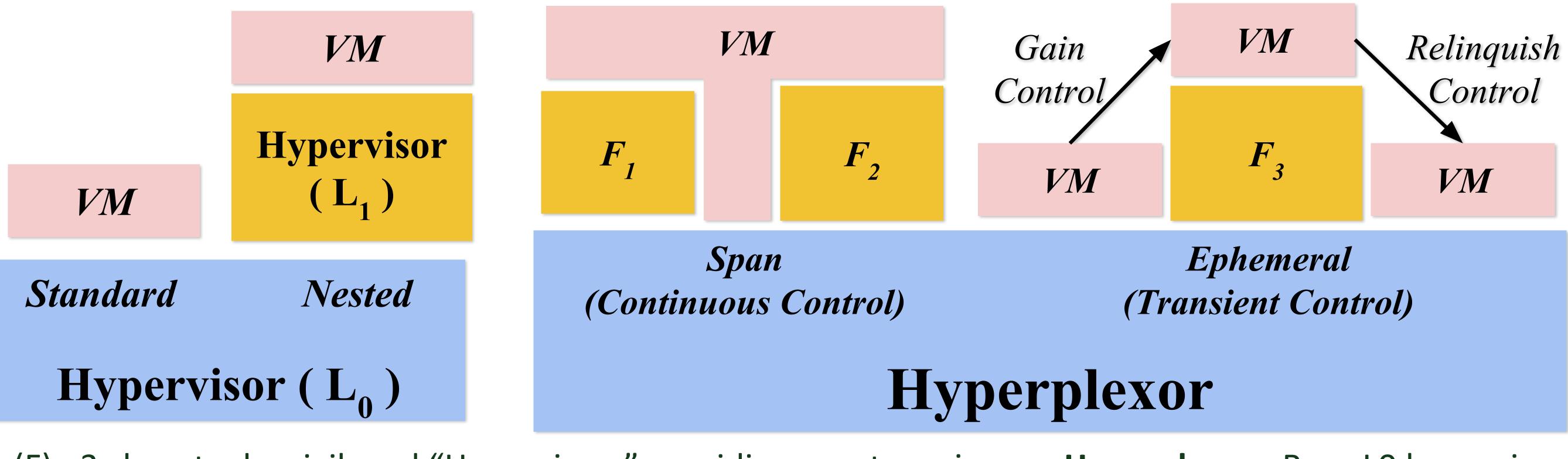
1. Growing Number of Hypervisor-

level Services: VM Introspection, Intrusion Detection, High Availability, Live Migration, Live Patching, etc.

2. Guests Cannot Simultaneously Use

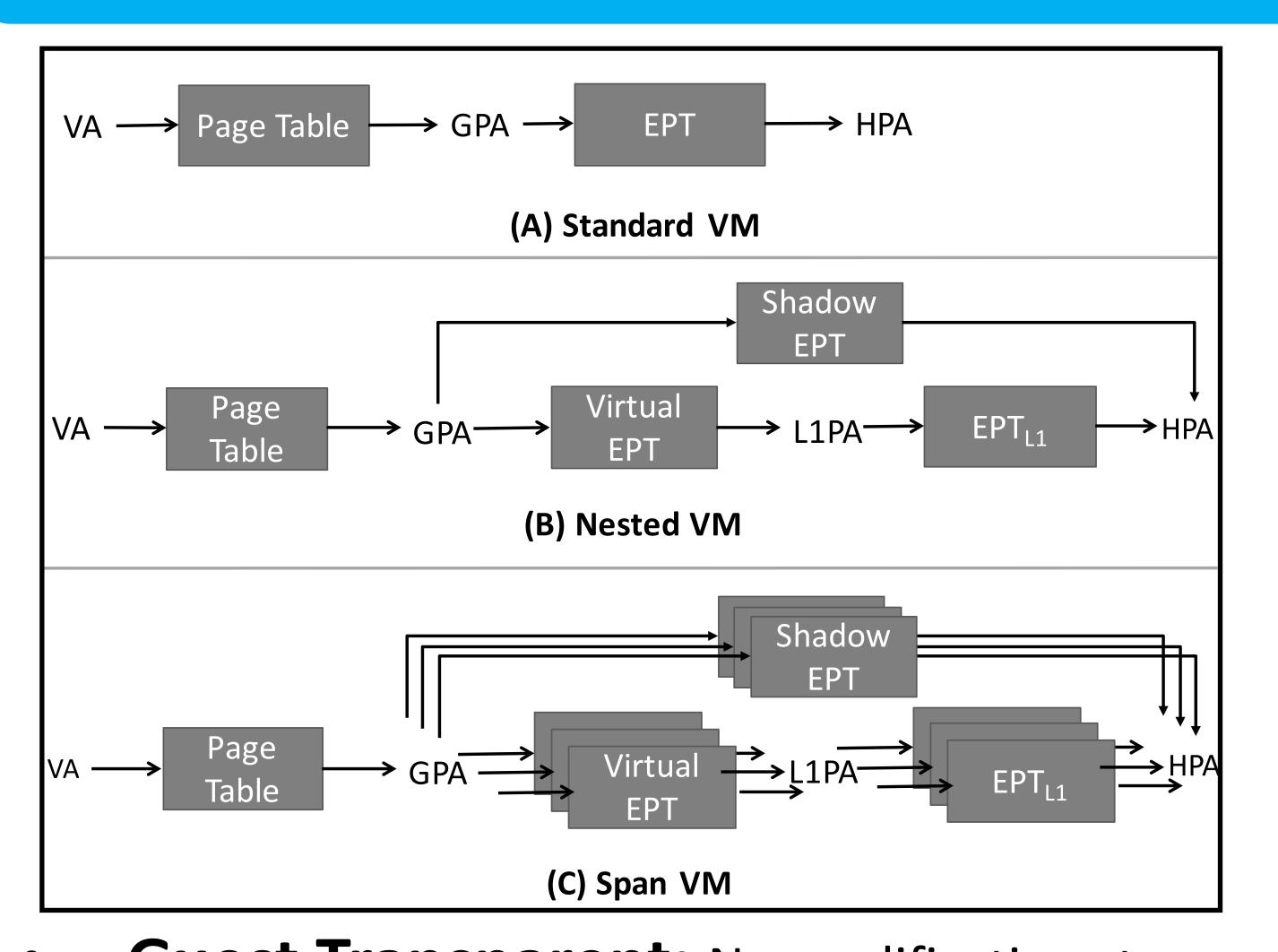
Multiple 3rd-party Services: E.g. Cross-cloud migration, Customized guest security, Attestation, etc.

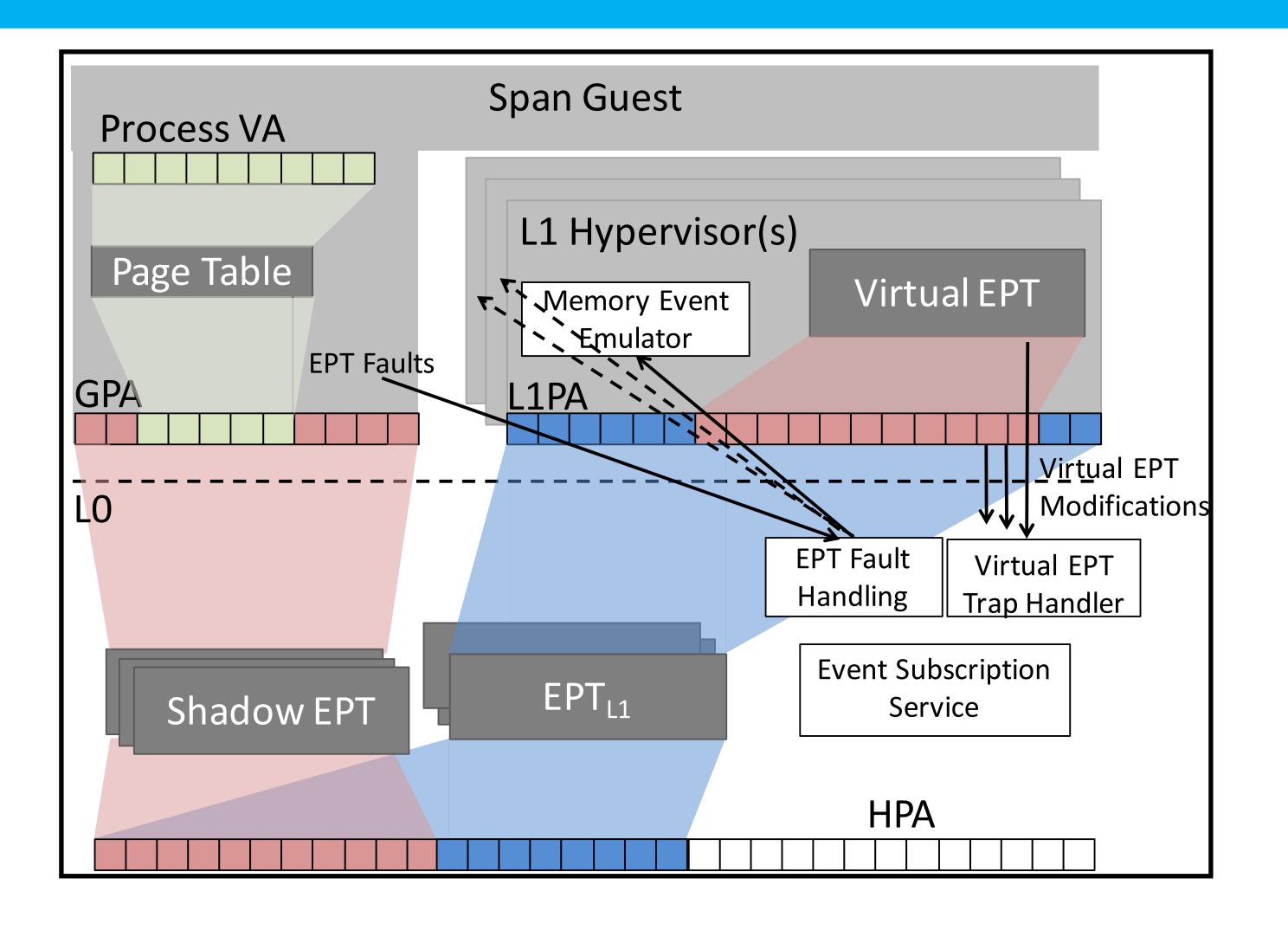
Solution: Compartmentalize Services & Share Guest Control

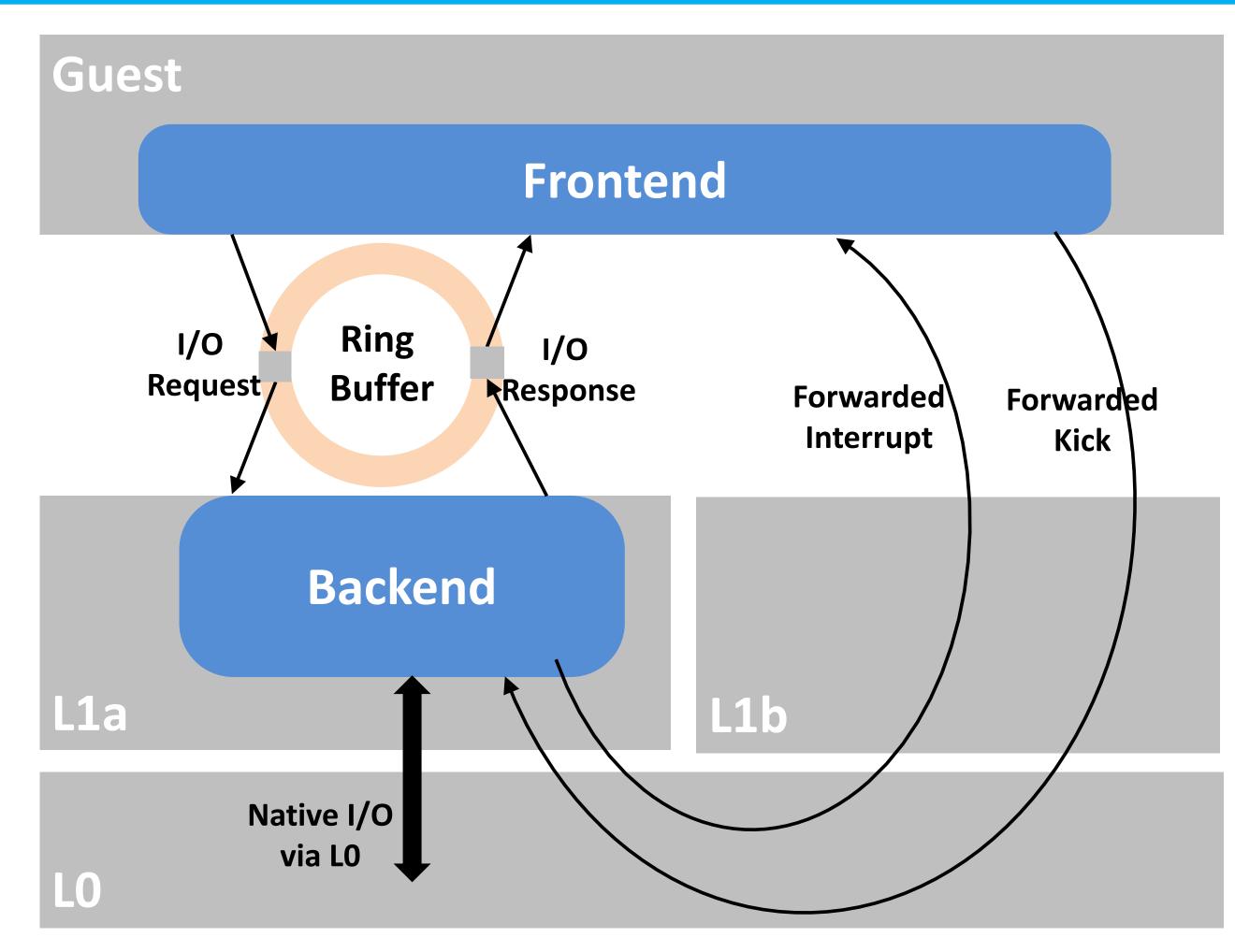


Featurevisors (F): 3rd-party deprivileged "Hypervisors" providing guest services. Hyperplexor: Base LO hypervisor.

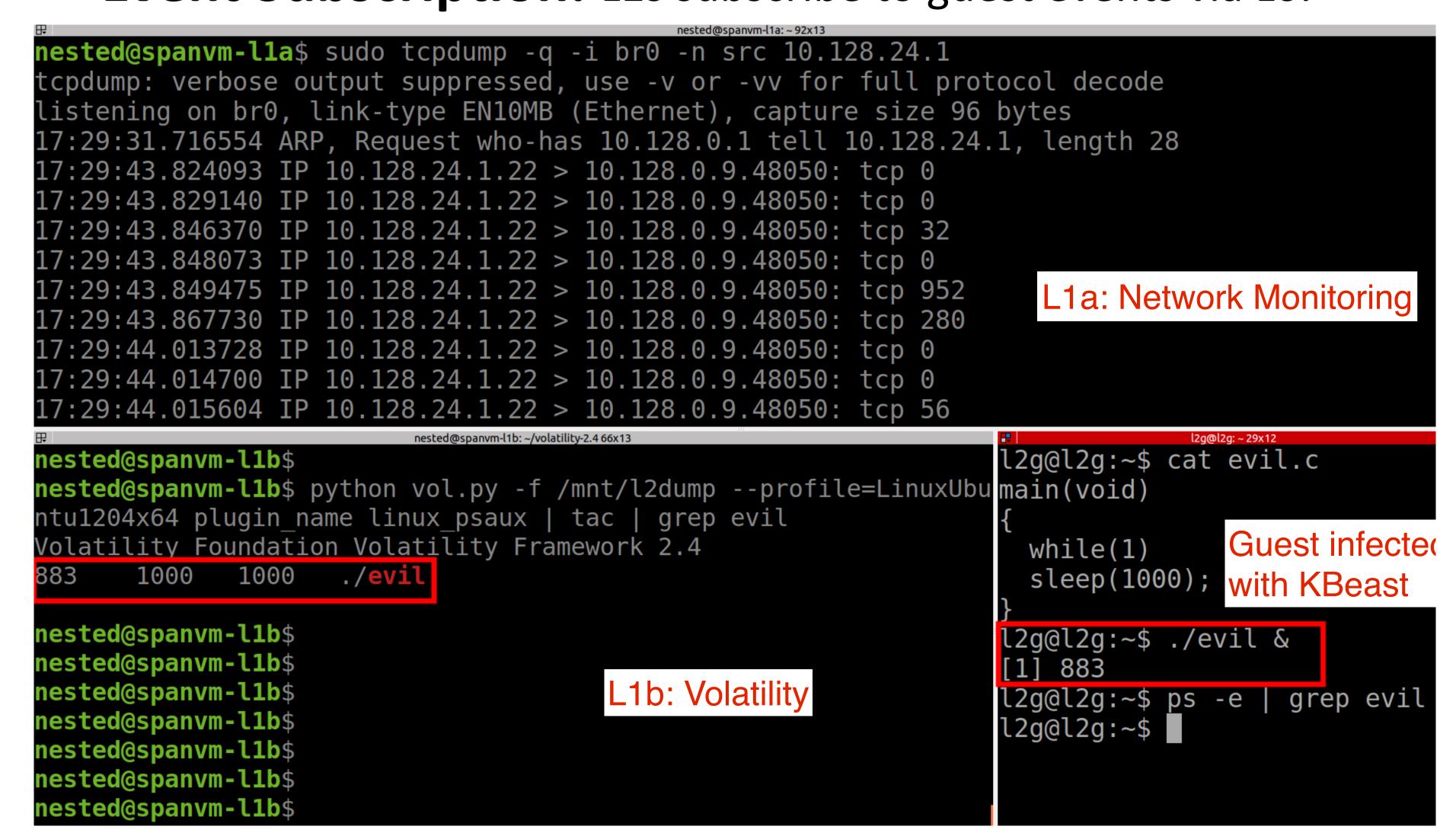
Approach: Transparent and Simultaneous Control of Guest by Multiple L1 Hypervisors







- Guest Transparent: No modifications to guest.
- Attach/Detach L1s to/from guest at runtime: Partial/full control over guest memory, VCPUs, and I/O devices.
- Event Subscription: L1s subscribe to guest events via L0.



Status, Results, and Future Work

Key Publications:

- 1. Multi-hypervisor Virtual Machines: Enabling an Eco-system of Hypervisor-level Services, Accepted in USENIX ATC, 2017
- 2. Enabling Hypervisor-as-a-service Clouds with Ephemeral Virtualization, **VEE 2016.**

Prototype on KVM/QEMU Platform

- 0—15% overhead on benchmarks: Kernbench, iperf, quicksort.
- Ephemeral virtualization: 80ms average switching times
- Page fault servicing: 3.6—4.2us; Event Redirection: 13-41us.

Ongoing/Future Work:

- Supporting unmodified L1 hypervisors.
- Live hypervisor patching.
- Support on public clouds.