

Assignment 5 CECS 451

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1. True or False? You don't need to explain your answers.
 - (a) (2 points) $h(n) = 0$ is an admissible heuristic for the 8-queens problem.
 - i. True, because when $h(n)$ is at 0 then that would mean there are no queens in attacking range of each other.
 - (b) (2 points) Assume that a rook can move on a chessboard one square at a time in vertically or horizontally, but cannot jump over other pieces. Manhattan distance is an admissible heuristic for the problem of moving the rook from square A to square B in the smallest number of moves.
 - i. True
2. (6 points) The heuristic path algorithm is a best-first search in which the evaluation function is $f(n) = (2 - w)g(n) + wh(n)$. What kind of search does this perform for $w = 0, w = 1$, and $w = 2$?
 - (a) If $w = 0$ then $f(n) = (2 - 0)g(n) + 0h(n) = f(n) = 2g(n)$ the algorithm would be a Breadth-First search or a Depth first Search.
 - (b) If $w = 1$ then $f(n) = (2 - 1)g(n) + 1h(n) = f(n) = g(n) + h(n)$ would be a A* search algorithm.
 - (c) If $w = 2$ then $f(n) = (2 - 2)g(n) + 2h(n) = f(n) = 2h(n)$ would be a Greedy Algorithm because it expands with minimal $h(n)$.
3. Give the name of the algorithm that results from each of the following cases:
 - (a) (2 points) Local beam search with $k = 1$.
 - i. It would be Hill Climb algorithm.
 - (b) (2 points) Simulated annealing with $T = \infty$ at all times.
 - i. It would be the Random-restart Hill Climbing algorithm.

4. Imagine that, one of the friends wants to avoid the other. The problem then becomes a two-player pursuit–evasion game. We assume now that the players take turns moving. The game ends only when the players are on the same node; the terminal payoff to the pursuer is minus the total move taken. An example is shown in Figure 1.
 - (a) (2 points) What is the terminal payoff at the node (1)?
 - i. The terminal payoff at the node (1) would be -4.
 - (b) (2 points) What are the positions of the two players at the node (2) and (2)'s children?
 - i. The positions of the two players at the node (2) and (2)'s children would be (be) and (bd).
 - (c) (3 points) Can we assume the terminal payoff at the node (2) is less than < -4 ? Answer yes or no, then explain your answers.
 - i. Yes, we can assume the terminal payoff at the node (2) is less than -4 because as we go down the tree to whatever minimum there is, the payoff would keep subtracting 1 so it would be soon be less than -4.
 - (d) (3 points) Assume the terminal payoff at the node (4) is less than -4. Do we need expand the node (5) and (6)? Answer yes or no, then explain your answers.
 - i. If the terminal payoff at the node (4) is less than -4 then no we should not expand the nodes (5) and (6) because we would offset the balance of the tree since the left side of the tree is at a terminal payoff of -4 at node (1).

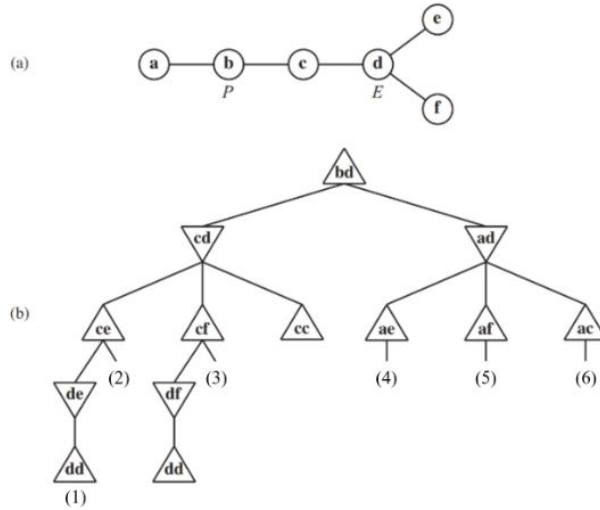


Figure 1: (a) A map where the cost of every edge is 1. Initially the pursuer P is at node b and the evader E is at node d. (b) A partial game tree for this map. Each node is labeled with the P, E positions. P moves first.

5. True or False? You don't need to explain your answers.

(a) (2 points) $(A \wedge B) \models (A \Leftrightarrow B)$

i. True

(b) (2 points) $(C \vee (\neg A \wedge \neg B)) \equiv ((A \Rightarrow C) \wedge (B \Rightarrow C))$

i. True

(c) (2 points) $(A \vee B) \wedge (\neg C \vee \neg D \vee E) \equiv (A \vee B)$

i. True

(d) (2 points) $(A \vee B) \wedge \neg(A \Rightarrow B)$ is satisfiable.

i. True

6. (4 points) Prove using Venn diagram, or find a counterexample to the following assertion:

$$\alpha \models (\beta \wedge \gamma) \text{ then } \alpha \models \beta \text{ and } \alpha \models \gamma$$

