

Design Analysis and Algorithm – Lab Work

Week 4

Question 1: Write a program to Illustrate Merge Sort to sort the array

{157,110,147,122,111,149,151,141,123,112,117,133}.

Code:

```
1 //CH.SC.U4CSE24122
2 #include <stdio.h>
3 #define MAX 50
4 void merge(int arr[], int l, int m, int r){
5     int i, j, k;
6     int n1 = m - 1 + 1;
7     int n2 = r - m;
8     int L[MAX], R[MAX];
9     for(i = 0; i < n1; i++)
10        L[i] = arr[l + i];
11     for(j = 0; j < n2; j++)
12        R[j] = arr[m + 1 + j];
13     i = 0;
14     j = 0;
15     k = l;
16    while(i < n1 && j < n2){
17        if(L[i] <= R[j])
18            arr[k++] = L[i++];
19        else
20            arr[k++] = R[j++];
21    }
22    while(i < n1)
23        arr[k++] = L[i++];
24    while(j < n2)
25        arr[k++] = R[j++];
```

```

27  void mergeSort(int arr[], int l, int r){
28      int m;
29      if(l < r){
30          m = l + (r - 1) / 2;
31          mergeSort(arr, l, m);
32          mergeSort(arr, m + 1, r);
33          merge(arr, l, m, r);
34      }
35  }
36  int main(){
37      printf("CH.SC.U4CSE24122\n");
38      int i;
39      int arr[] = {157,110,147,122,111,149,151,141,123,112,117,133};
40      int n = sizeof(arr) / sizeof(arr[0]);
41      mergeSort(arr, 0, n - 1);
42      printf("Sorted array: ");
43      for(i = 0; i < n; i++)
44          printf("%d ", arr[i]);
45  }

```

Output:

```

C:\Users\kpriy\Downloads\uh  X  +  ▾
CH.SC.U4CSE24122
Sorted array: 110 111 112 117 122 123 133 141 147 149 151 157
-----
Process exited after 0.07928 seconds with return value 0
Press any key to continue . . .

```

Time Complexity:

The program recursively divides the array into two halves until each subarray has one element. At each level of recursion, the merge step requires n comparisons and assignments. Since the array is divided $\log n$ times, the total number of operations is proportional to $n \log n$. Hence, the **Time Complexity of Merge Sort in the best case, average case, and worst case is $O(n \log n)$.**

Space Complexity:

The space occupied will be proportional to the size of the input array. For an array of size n , additional temporary arrays $L[]$ and $R[]$ are created during each merge step, together holding n elements. Apart from these, a few integer variables (i, j, k, l, m, r) are used for indexing and recursion. The recursion stack depth is $\log n$. Hence, the total extra memory used grows linearly with the input size.

Therefore, the **Space Complexity is $O(n)$.**

Question 2: Write a program to Illustrate Quick Sort the array

{157,110,147,122,111,149,151,141,123,112,117,133}.

Code:

```

1 //CH.SC.U4CSE24122
2 #include<stdio.h>
3 int partition(int arr[],int l,int h){
4     int pivot=arr[h];
5     int i=l-1,j;
6     for(j=l;j<h;j++){
7         if(arr[j]<pivot){
8             i++;
9             int temp=arr[i];
10            arr[i]=arr[j];
11            arr[j]=temp;
12        }
13    }
14    int temp=arr[i+1];
15    arr[i+1]=arr[h];
16    arr[h]=temp;
17    return (i+1);
18 }
19 void quicksort(int arr[],int l,int h){
20     if(l<h){
21         int pi=partition(arr,l,h);
22         quicksort(arr,l,pi-1);
23         quicksort(arr,pi+1,h);
24     }
25 }
26 int main(){

```

```

27     printf("CH.SC.U4CSE24122\n");
28     int arr[]={157,110,147,122,111,149,151,141,123,112,117,133};
29     int h=sizeof(arr)/sizeof(arr[0]);
30     quicksort(arr,0,h-1);
31     printf("Sorted Array: ");
32     int i;
33     for(i=0;i<h;i++){
34         printf("%d ",arr[i]);
35     }
36     return 0;
37 }
38

```

Output:

```

CH.SC.U4CSE24122
Sorted Array: 110 111 112 117 122 123 133 141 147 149 151 157
-----
Process exited after 0.06539 seconds with return value 0
Press any key to continue . . .

```

Space Complexity:

The space occupied will be proportional to the recursion depth. For an array of size n , the program uses a few integer variables (pivot, i, j, l, h, pi, temp) for indexing and swapping. The recursion stack depth depends on how the pivot divides the array: in the best and average case, the depth is $\log n$, while in the worst case it can go as deep as n . Hence, the total extra memory used depends on recursion depth.

Therefore, the **Space Complexity is $O(\log n)$ on average and $O(n)$ in the worst case.**

Time Complexity:

The program partitions the array around a pivot and recursively sorts the two halves. At each level of recursion, the partition step requires n comparisons and swaps. If the pivot divides the array evenly, the array is split $\log n$ times, giving a total of $n \log n$ operations. If the pivot divides poorly (always smallest or largest element), one side has size $n-1$ and the other has size 0, leading to n^2 operations.

Hence, the **Time Complexity of Quick Sort is $O(n \log n)$ in the best case and average case, and $O(n^2)$ in the worst case.**