

## Design Analysis and Algorithm – Lab Work

## Week 4

Question 1: Write a program to Illustrate Merge Sort to sort the array

{157,110,147,122,111,149,151,141,123,112,117,133}.

Code:

```
1 //CH.SC.U4CSE24122
2 #include <stdio.h>
3 #define MAX 50
4 void merge(int arr[], int l, int m, int r){
5     int i, j, k;
6     int n1 = m - l + 1;
7     int n2 = r - m;
8     int L[MAX], R[MAX];
9     for(i = 0; i < n1; i++)
10         L[i] = arr[l + i];
11     for(j = 0; j < n2; j++)
12         R[j] = arr[m + 1 + j];
13     i = 0;
14     j = 0;
15     k = l;
16     while(i < n1 && j < n2){
17         if(L[i] <= R[j])
18             arr[k++] = L[i++];
19         else
20             arr[k++] = R[j++];
21     }
22     while(i < n1)
23         arr[k++] = L[i++];
24     while(j < n2)
25         arr[k++] = R[j++];
26 }
```

```

27 void mergeSort(int arr[], int l, int r){
28     int m;
29     if(l < r){
30         m = l + (r - l) / 2;
31         mergeSort(arr, l, m);
32         mergeSort(arr, m + 1, r);
33         merge(arr, l, m, r);
34     }
35 }
36 int main(){
37     printf("CH.SC.U4CSE24122\n");
38     int i;
39     int arr[] = {157,110,147,122,111,149,151,141,123,112,117,133};
40     int n = sizeof(arr) / sizeof(arr[0]);
41     mergeSort(arr, 0, n - 1);
42     printf("Sorted array: ");
43     for(i = 0; i < n; i++)
44         printf("%d ", arr[i]);
45     return 0;
46 }

```

Output:

```

C:\Users\kpriy\Downloads\uh
CH.SC.U4CSE24122
Sorted array: 110 111 112 117 122 123 133 141 147 149 151 157
-----
Process exited after 0.07928 seconds with return value 0
Press any key to continue . . . |

```

### Time Complexity:

The program recursively divides the array into two halves until each subarray has one element. At each level of recursion, the merge step requires  $n$  comparisons and assignments. Since the array is divided  $\log n$  times, the total number of operations is proportional to  $n \log n$ . Hence, the **Time Complexity of Merge Sort in the best case, average case, and worst case is  $O(n \log n)$ .**

**Space Complexity:**

The space occupied will be proportional to the size of the input array. For an array of size  $n$ , additional temporary arrays  $L[]$  and  $R[]$  are created during each merge step, together holding  $n$  elements. Apart from these, a few integer variables ( $i, j, k, l, m, r$ ) are used for indexing and recursion. The recursion stack depth is  $\log n$ . Hence, the total extra memory used grows linearly with the input size.

Therefore, the **Space Complexity is  $O(n)$** .

**Question 2: Write a program to Illustrate Quick Sort the array {157,110,147,122,111,149,151,141,123,112,117,133}.**

Code:

```

1  //CH.SC.U4CSE24122
2  #include<stdio.h>
3  int partition(int arr[],int l,int h){
4      int pivot=arr[h];
5      int i=l-1,j;
6      for(j=l;j<h;j++){
7          if(arr[j]<pivot){
8              i++;
9              int temp=arr[i];
10             arr[i]=arr[j];
11             arr[j]=temp;
12         }
13     }
14     int temp=arr[i+1];
15     arr[i+1]=arr[h];
16     arr[h]=temp;
17     return (i+1);
18 }
19 void quicksort(int arr[],int l,int h){
20     if(l<h){
21         int pi=partition(arr,l,h);
22         quicksort(arr,l,pi-1);
23         quicksort(arr,pi+1,h);
24     }
25 }
26 int main(){

```

```

27     printf("CH.SC.U4CSE24122\n");
28     int arr[]={157,110,147,122,111,149,151,141,123,112,117,133};
29     int h=sizeof(arr)/sizeof(arr[0]);
30     quicksort(arr,0,h-1);
31     printf("Sorted Array: ");
32     int i;
33     for(i=0;i<h;i++){
34         printf("%d ",arr[i]);
35     }
36     return 0;
37 }
38

```

Output:

```

CH.SC.U4CSE24122
Sorted Array: 110 111 112 117 122 123 133 141 147 149 151 157
-----
Process exited after 0.06539 seconds with return value 0
Press any key to continue . . . |

```

### Space Complexity:

The space occupied will be proportional to the recursion depth. For an array of size  $n$ , the program uses a few integer variables (pivot,  $i$ ,  $j$ ,  $l$ ,  $h$ ,  $pi$ ,  $temp$ ) for indexing and swapping. The recursion stack depth depends on how the pivot divides the array: in the best and average case, the depth is  $\log n$ , while in the worst case it can go as deep as  $n$ . Hence, the total extra memory used depends on recursion depth.

Therefore, the **Space Complexity is  $O(\log n)$  on average and  $O(n)$  in the worst case.**

### Time Complexity:

The program partitions the array around a pivot and recursively sorts the two halves. At each level of recursion, the partition step requires  $n$  comparisons and swaps. If the pivot divides the array evenly, the array is split  $\log n$  times, giving a total of  $n \log n$  operations. If the pivot divides poorly (always smallest or largest element), one side has size  $n-1$  and the other has size 0, leading to  $n^2$  operations.

Hence, the **Time Complexity of Quick Sort is  $O(n \log n)$  in the best case and average case, and  $O(n^2)$  in the worst case.**