

QN 3.3) Solution:

We have,

```
cnt :: Eq a => a -> [a] -> Int
cnt x [] = 0
cnt x (y:ys)
    | x == y = 1 + (cnt x ys)
    | otherwise = cnt x ys
con :: [a] -> [a] -> [a]
con [] ys = ys
con (x:xs) ys = x : (con xs ys)
```

The cnt function is for counting the number of appearance of value of x in a list. And, the con function is for joining two lists into one single list containing the values of those two list.

To Prove:

$$\text{cnt } x \text{ (con } s \text{ } t) == (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t)$$

Using the mathematical induction,

Basic Steps: Let's consider the empty list.

Taking $s = []$

So our expression would be,

$$\Rightarrow \text{cnt } x \text{ (con } [] \text{ } t) == (\text{cnt } x \text{ } []) + (\text{cnt } x \text{ } t)$$
$$\Rightarrow \text{cnt } x \text{ } t == 0 + (\text{cnt } x \text{ } t)$$
$$\Rightarrow \text{cnt } x \text{ } t == \text{cnt } x \text{ } t \quad (\text{True})$$

[As cnt returns Int, so numeric operations validate]

Inductive Steps:

Assumption: $\text{cnt } x \text{ (con } s \text{ } t) == (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t)$

Let's take a list k such that the $\text{length}(k) = (\text{length}(s)+1)$, and the extra value that k contains be c.

To prove:

$$\text{cnt } x \text{ (con } k \text{ } t) == (\text{cnt } x \text{ } k) + (\text{cnt } x \text{ } t)$$

So, our expression would be,

$$\Rightarrow \text{cnt } x \text{ (con } k \text{ } t) == (\text{cnt } x \text{ } k) + (\text{cnt } x \text{ } t)$$

$$\Rightarrow \text{cnt } x \text{ (} k ++ t \text{)} == (\text{cnt } x \text{ } k) + (\text{cnt } x \text{ } t) \quad [\text{As, (con } k \text{ } t) = k ++ t]$$

From our program,

If $x == c$, the above expression is equal to

$$\Rightarrow 1 + \text{cnt } x \text{ (} s ++ t \text{)} == (1 + (\text{cnt } x \text{ } s)) + (\text{cnt } x \text{ } t)$$

$$\Rightarrow 1 + \text{cnt } x \text{ (con } s \text{ } t) == 1 + (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t)$$

$$[\text{As, } s ++ t = (\text{con } s \text{ } t)]$$

$$\Rightarrow 1 + (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t) == 1 + (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t) \quad \textbf{(True)}$$

[From our assumption,

$$\text{cnt } x \text{ (con } s \text{ } t) == (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t)]$$

Otherwise,

$$\Rightarrow 0 + \text{cnt } x \text{ (} s ++ t \text{)} == (0 + (\text{cnt } x \text{ } s)) + (\text{cnt } x \text{ } t)$$

$$\Rightarrow 0 + \text{cnt } x \text{ (con } s \text{ } t) == 0 + (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t)$$

$$[\text{As, } s ++ t = (\text{con } s \text{ } t)]$$

$$\Rightarrow 0 + (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t) == 0 + (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t) \quad \textbf{(True)}$$

[From our assumption,

$$\text{cnt } x \text{ (con } s \text{ } t) == (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t)]$$

Hence, we can conclude inductively that,

$$\text{cnt } x \text{ (con } s \text{ } t) == (\text{cnt } x \text{ } s) + (\text{cnt } x \text{ } t)$$

QN 3.4a) solution:

```
rotate :: Int -> [a] -> [a]
rotate _ [] = []
rotate x xs
    | (x `mod` length(xs)) == 0 = xs
rotate x ys = rotate ((x-1) `mod` length(ys)) (tail ys ++ [head ys])
```

QN 3.4b) solution:

--please note that you must load the rotate function from question 3.4a in order to run this program successfully.

```
circle :: [a] -> [[a]]
circle [] = []
circle xs = [(rotate x xs) | x<-[0..(length(xs)-1)]]
```