

Q 14 P.1 a) solⁿ

The Boolean expression implemented by the digital circuit is:

$$Y = (\neg A \uparrow \neg B) \uparrow ((A \uparrow B) \uparrow C)$$

We can also write the above exp. as:

$$\neg \left((\neg(A \wedge \neg B)) \wedge \neg(\neg(A \wedge B) \wedge C) \right)$$

$$= \neg \left((\neg \neg A \vee \neg \neg B) \wedge (\neg \neg(A \wedge B) \vee \neg C) \right)$$

$$= \neg \left((A \vee B) \vee \neg((A \wedge B) \vee \neg C) \right) \quad [\text{De Morgan's law}]$$

$$= \neg \left((A \vee B) \vee \neg((A \wedge B) \vee \neg C) \right) \quad [\text{double negation}]$$

$$= ((\neg A \wedge \neg B) \vee \neg(\neg(A \wedge B) \wedge \neg \neg C))$$

[De Morgan's law]

$$= (\neg A \wedge \neg B) \vee ((\neg A \vee \neg B) \wedge C)$$

[double negation]

$$= (\neg A \wedge \neg B) \vee (\neg A \wedge C) \vee (\neg B \wedge C)$$

[Distributivity]

$$= (\neg A \wedge \neg B) \vee (\neg A \wedge C) \vee (\neg B \wedge C)$$

The above eq. can also be written as,

$$\begin{aligned} & (\neg A \wedge \neg B \wedge C) \vee (\neg A \wedge \neg B \wedge \neg C) \vee \\ & (\neg A \wedge B \wedge C) \vee (\neg A \wedge \neg B \wedge C) \vee \\ & (A \wedge \neg B \wedge C) \vee (\neg A \wedge \neg B \wedge C) \end{aligned}$$

So, in minterm form we can write,

$$\begin{aligned} & = m_1 + m_0 + m_3 + m_1 + m_5 + m_1 \\ & = m_0 + m_1 + m_3 + m_5 \quad [\text{idempotency}] \end{aligned}$$

Q. N 8(2)

$$S = (A \vee B) \vee C_{in}$$

$$C_{out} = (A \wedge B) \vee (C_{in} \wedge (A \vee B))$$

for S:

The S can also be written as

$$\neg((\neg A \wedge B) \vee (A \wedge \neg B)) \wedge C_{in}$$

$$\vee ((\neg A \wedge B) \vee (A \wedge \neg B)) \wedge \neg C_{in}$$

$$= (\neg((\neg A \wedge B) \wedge C_{in}) \vee ((\neg A \wedge B) \wedge C_{in})) \vee ((A \wedge \neg B) \wedge C_{in}) \vee ((A \wedge \neg B) \wedge \neg C_{in})$$

[distributivity]

0	0	0	0
0	1	0	0
1	0	0	0
1	1	0	0

0	0	0	0
0	1	0	0
1	0	0	0
1	1	0	0

~~Q. 11~~
Q. 11 N 8.2) Soln

$$S = (A \vee B) \vee C_{in}$$

$$C_{out} = (A \wedge B) \vee (C_{in} \wedge (A \vee B))$$

for S and C_{out}

As the truth table is easier to follow than the algebraic computation
so,

A	B	C _{in}	S	C _{out}
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1
X	X	X	X	X

so, sum for S is

$$(\neg A \wedge \neg B \wedge C_{in}) \vee (\neg A \wedge B \wedge \neg C_{in}) \vee (A \wedge \neg B \wedge \neg C_{in}) \vee (A \wedge B \wedge C_{in})$$

The product for the S is:

$$(\neg A \vee \neg B \vee \neg C_{in}) \wedge (\neg A \vee B \vee C_{in}) \wedge \\ (A \vee \neg B \vee C_{in}) \wedge (A \vee B \vee \neg C_{in})$$

The sum for the Cout is:

$$\cancel{(\neg A \vee \neg B \vee C_{in}) \wedge} \\ (A \wedge B \wedge \neg C_{in}) \vee (\neg A \wedge B \wedge C_{in}) \vee \\ (A \wedge \neg B \wedge C_{in}) \vee (A \wedge B \wedge C_{in})$$

The product for Cout is:

$$(\neg A \vee \neg B \vee \neg C_{in}) \wedge (\neg A \vee \neg B \vee C_{in}) \wedge \\ (\neg A \vee B \vee \neg C_{in}) \wedge (A \vee \neg B \vee \neg C_{in})$$

C)

$$\begin{aligned} S &= A \oplus B \oplus C_{in} \\ &= (A \oplus B) \oplus C_{in} \\ &= ((A \vee B) \wedge (\neg A \vee \neg B)) \oplus C_{in} \\ &= (\neg(\neg A \wedge \neg B) \wedge (\neg(A \wedge B))) \oplus C_{in} \end{aligned}$$

[De Morgan's law?]

$$= \cancel{(\neg A \uparrow \neg B) \wedge (A \uparrow B)) \uparrow \neg A} \\ = \cancel{(\neg A \wedge A)}$$

c) let's consider for x, y

$$x \downarrow y = ? \text{ (in terms of } \neg \text{ and } \uparrow \text{)}$$

	x	y	$x \uparrow y$	$\neg(x \uparrow y)$	$(x \uparrow y) \uparrow \neg(x \uparrow y)$	$x \downarrow y$
	0	0	1	0	0	0
	0	1	1	1	1	1
	1	0	1	1	1	1
	1	1	0	1	0	0

As we can see $x \downarrow y$ is identical to
 $(x \uparrow y) \uparrow \neg(x \uparrow y)$

So,

$$x \downarrow y = (x \uparrow y) \uparrow \neg(x \uparrow y)$$

$$\text{So } S = A \downarrow B \downarrow C \text{ in}$$

$$= (A \downarrow B) \downarrow C \text{ in}$$

$$= ((A \uparrow B) \uparrow (\neg A \uparrow \neg B)) \vee C_{in}$$

$$= (((A \uparrow B) \uparrow (\neg A \uparrow \neg B)) \uparrow C_{in}) \uparrow a$$

$$(\neg((A \uparrow B) \uparrow (\neg A \uparrow \neg B)) \uparrow \neg C_{in})$$

which is the required form
consisting only \uparrow and \neg .

for C_{out} :

$$C_{out} = (A \wedge B) \vee (C_{in} \wedge (A \vee B))$$

$$= \neg(\neg A \uparrow \neg B) \vee \neg(\neg C_{in} \uparrow \neg(A \vee B))$$

[obvious]

$$= \neg((\neg A \uparrow \neg B) \wedge (\neg C_{in} \uparrow \neg(A \vee B)))$$

[De Morgan's law]

$$= (\neg A \uparrow \neg B) \uparrow (\neg C_{in} \uparrow \neg(A \vee B))$$

$$= (\neg A \uparrow \neg B) \uparrow (C_{in} \uparrow \neg((A \uparrow B) \uparrow (\neg A \uparrow \neg B)))$$

which is the required form
consisting only \uparrow and \neg .

d)

for S:

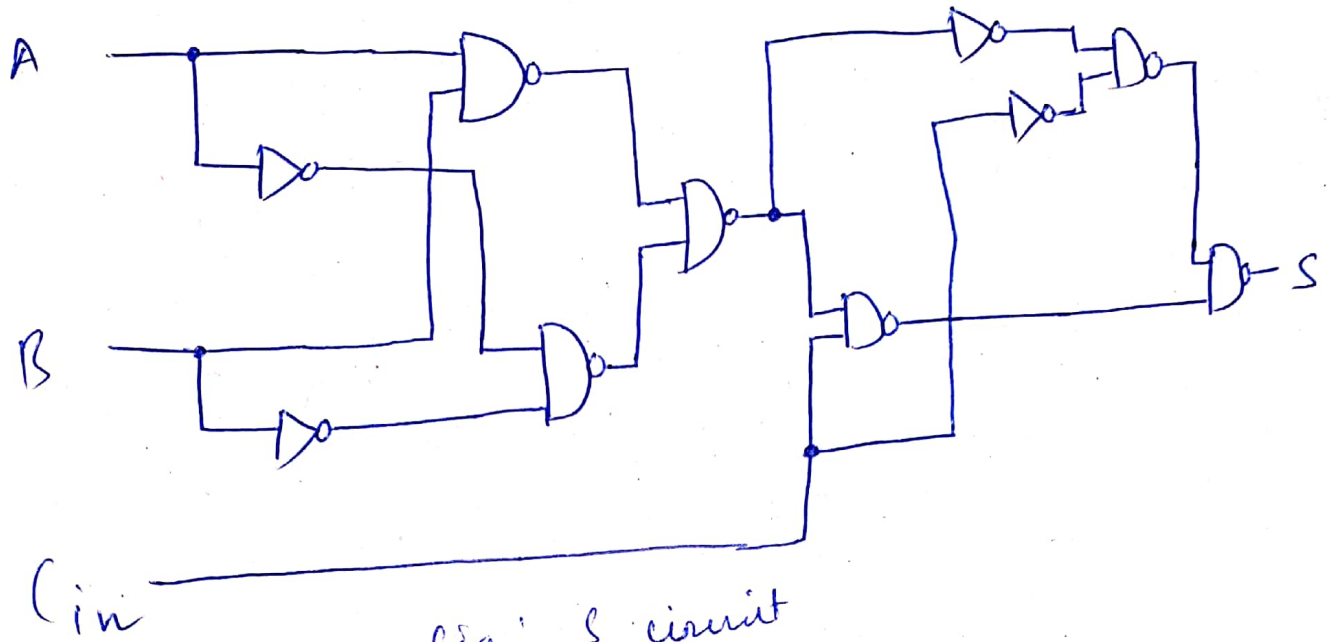


fig: S circuit

for Cout:

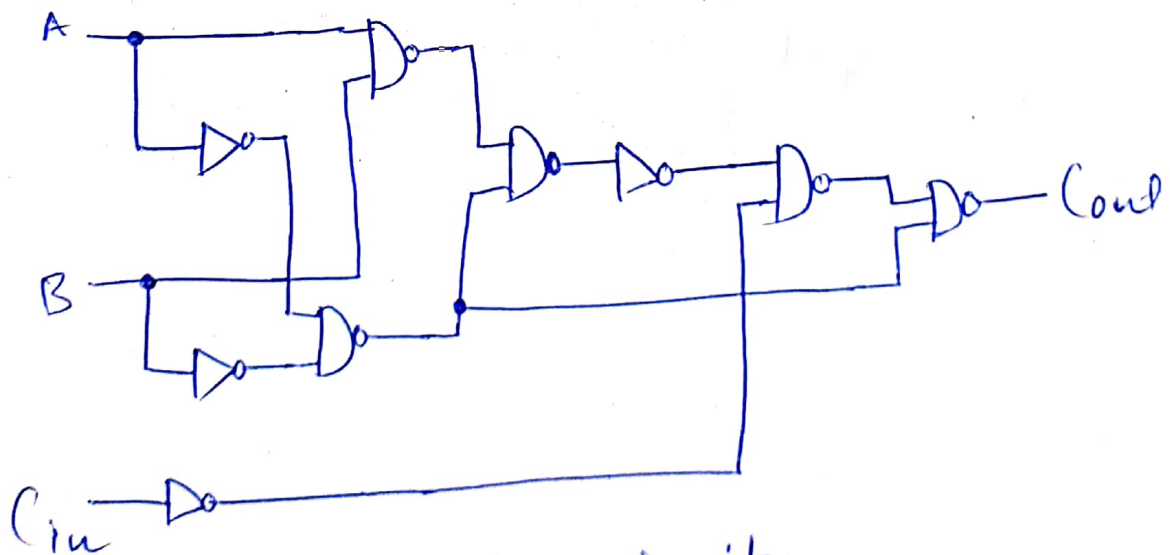


fig: Cout circuit