

# COM S 352 Homework 6

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## Question 1

## Question 2

## Question 3

```
int serve, int order, int eat;
void init() {
    serve = 0;
    order = 1;
    eat = 0;
}

void chef() {
    while (true) {
        wait(serve);
        signal(eat);
    }
}

void customer() {
    wait(seats);
    wait(order);
    signal(serve);
    wait(eat);
    signal(order);
    signal(seat);
}
```

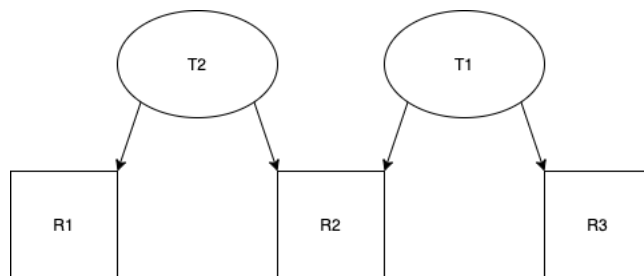
## Question 4

a.

- *Mutual Exclusion*  
Only a single street of cars can use the resource they are on
- *Hold & Wait*  
Each line of cars is *holding* a resource (the street) and *waiting* for the perpendicular street to open up
- *No Preemption*  
A resource cannot be released until the entire line of cars have passed
- *Circular Wait* If we label each set of vehicles, v1, v2, v3, v4, then, v1 waits for v2, v2 waits for v3, etc... resulting in a circular wait

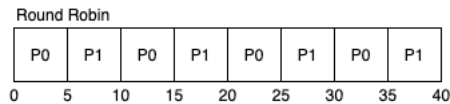
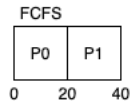
b. A rule to ensure no deadlocking could be to remove any sort of pre-emption. This would force each process to wait until its next resources are available. If a process makes a request but the request is not satisfied it will need to make the request again once the resources are available.

## Question 5



This diagram shows that there is no way to have a cycle which, means that this system is deadlock free.

## Question 6



The Round Robin algorithm could result in deadlock. If  $P_0$  uses S and  $P_1$  uses Q then they will both be waiting which means  $P_0$  won't be able to use Q and  $P_1$  can't use S. The FCFS algorithm is sequential however, which results in no deadlock.

## Question 7

- **a.**  
No deadlock because no cycle is present. A possible execution plan could be:  
 $T_2 \rightarrow T_3 \rightarrow T_1$
- **b.** This graph is deadlocked. The cycle is:  
 $T_1 \rightarrow R_3 \rightarrow T_3 \rightarrow R_1 \rightarrow T_1$
- **c.** No deadlock because no cycle is present. A possible execution plan could be:  
 $T_3 \rightarrow T_1 \rightarrow T_2$
- **d.** This graph is deadlocked. The cycle is:  
 $T_1 \rightarrow R_2 \rightarrow T_3 \rightarrow R_1 \rightarrow T_1$
- **e.** No deadlock because there is no hold and wait condition. A possible execution plan could be:  $T_2 \rightarrow T_1 \rightarrow T_3 \rightarrow T_4$
- **f.** No deadlock, R2, has 3 available instances and 3 connecting threads. A possible execution plan could be:  $T_2 \rightarrow T_4 \rightarrow T_1 \rightarrow T_3$