$\underline{\hbox{IUT SuperSonic, Islamic University of Technology}}$

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Sublime Build

1 All Macros

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack
    :200000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,ssse3,
    sse4,popcnt,abm,mmx,avx,tune=native")
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
   //find_by_order(k) --> returns iterator
        to the kth largest element counting
         from 0
   //order_of_key(val) --> returns the
        number of items in a set that are
        strictly smaller than our item
template <typename DT>
using ordered_set = tree <DT, null_type,</pre>
    less<DT>, rb_tree_tag,
    tree_order_statistics_node_update>;
    \{-2,-2,-1,+1,+2,+2,+1,-1\};
int ky[] =
    \{-1,+1,+2,+2,+1,-1,-2,-2\};
#define fastio
                        ios_base::
    sync_with_stdio(0);cin.tie(0);
#define Make(x,p)
                      (x | (1<<p))
                      (x & ~(1<<p))
#define DeMake(x,p)
#define Check(x,p)
                      (x & (1<<p))
                      cerr << #x << " = " <<
#define DEBUG(x)
     x << endl
```

2 DP

2.1 Sum of Subsets

```
//submask == all i such that mask&i == i ||
    mask&i == mask (all i such that all 0
    in mask are fixed and the 1's change)
//sos dp memory optimized
for (int i = 0; i < (1 << N); ++i) F[i] = A
    [i];
for (int i = 0; i < N; ++i) {
    for (int mask = 0; mask < (1 << N); ++
        mask) {
      if (mask & (1 << i)) F[mask] += F[mask
        ^ (1 << i)]; /// doing -= can work
        like inclusion-exclusion on unset
        bits
    }
}</pre>
```

2.2 Divide and Conquer DP

```
inline void compute(int cur, int L, int R,
    int best_L, int best_R) {
    if (L > R) return;
    int mid = (L + R) >> 1;
    pair<ll, int>best = {inf, -1};
    for (int k = best_L; k <= min(best_R, mid)
        ; k++) {
        best = min(best, {dp[cur ^ 1][k - 1] +
            getCost(k, mid), k});
    }
    dp[cur][mid] = best.ff;
    int best_id = best.ss;</pre>
```

```
compute(cur, L, mid - 1, best_L, best_id);
compute(cur, mid + 1, R, best_id, best_R);
}
// in main
int cur = 0;
for (int i = 1; i <= n; i++) dp[1][i] = inf;
for (int guard = 1; guard <= g; guard++) {
   compute(cur, 1, n, 1, n); cur ^= 1;
}
ll ans = dp[cur ^ 1][n];</pre>
```

2.3 Li Chao Tree

```
struct line {
 11 m. c:
 line(ll m = 0, ll c = 0) : m(m), c(c) {}
ll calc(line L, ll x) {
 return 111 * L.m * x + L.c;
struct node {
 11 m, c;
 node *lft, *rt;
 node(11 m = 0, 11 c = 0, node *lft = NULL,
       node *rt = NULL) : L(line(m, c)),
      lft(lft), rt(rt) {}
struct LiChao {
 node *root;
 LiChao() {
   root = new node();
 void update(node *now, int L, int R, line
      newline) {
   int mid = L + (R - L) / 2;
   line lo = now->L, hi = newline;
   if (calc(lo, L) > calc(hi, L)) swap(lo,
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi;
     return:
   if (calc(lo, mid) < calc(hi, mid)) {</pre>
     now->L = hi;
     if (now->rt == NULL) now->rt = new
          node():
     update(now->rt, mid + 1, R, lo);
     now->L = lo:
     if (now->lft == NULL) now->lft = new
          node();
     update(now->lft, L, mid, hi);
 11 query(node *now, int L, int R, 11 x) {
   if (now == NULL) return -inf;
   int mid = L + (R - L) / 2;
   if (x <= mid) return max( calc(now->L, x
        ), query(now->lft, L, mid, x));
   else return max( calc(now->L, x), query(
        now->rt, mid + 1, R, x) );
 }
};
```

2.4 Knuth Iterative

```
for (int i = 1; i <= n; i++) {
   path[i][i] = i;
   dp[i][i] = 0;
}

for (int len = 2; len <= n; len++) {
   for (int st = 1; st + len - 1 <= n; st++)
        {
      int ed = st + len - 1;
      int L = max(st, path[st][ed - 1]);
      int R = min(ed - 1, path[st + 1][ed]);
      dp[st][ed] = INT_MAX;
   for (int i = L; i <= R; i++) {
      int cur = dp[st][i] + dp[i + 1][ed] +
            arr[ed] - arr[st - 1];
      if (dp[st][ed] > cur) {
        dp[st][ed] = cur;
      path[st][ed] = i;
    }
}
```

```
}
}
}
cout << dp[1][n] << "\n";
```

2.5 Number Permutation

```
11 dp[2][3005]; 11 sum[2][3005];
int dir[3005];
int arr[MAX];
int main() {
 fastio:
 int n;
 string s;
 cin >> n >> s:
 s = '#' + s;
 s.pb('<'); ///last element less than the</pre>
      element placed after it
 sum[1][0] = 1;
 int cur = 0:
 for (int baki = 1; baki <= n; baki++) {</pre>
   if (s[baki] == '<') dp[cur][0] = 0;</pre>
   else dp[cur][0] = sum[cur ^ 1][baki -
        1];
   for (int small = 1; small <= baki; small</pre>
        ++) {
     if (s[baki] == '<') dp[cur][small] =</pre>
          sum[cur ^ 1][small - 1];
       int big = baki - small;
       dp[cur][small] = sum[cur ^ 1][small
            + big - 1];
       dp[cur][small] -= sum[cur ^ 1][small
             - 1];
       if (dp[cur][small] < 0) dp[cur][</pre>
            small] += MOD;
     }
   sum[cur][0] = dp[cur][0];
   for (int small = 1; small <= baki; small</pre>
     sum[cur][small] = (sum[cur][small - 1]
           + dp[cur][small]);
     if (sum[cur][small] >= MOD) sum[cur][
          small] -= MOD;
   }
   cur ^= 1;
 11 ans = dp[cur ^ 1][n];
 cout << ans << "\n";
```

2.6 Knuth Optimization

```
11 solve(int st, int ed) { ///recursive
 if (st == ed) {
   path[st][ed] = st;
   return 0;
 11 &ret = dp[st][ed];
  if (ret != -1) return ret;
  solve(st, ed - 1); solve(st + 1, ed);
 int L = max(st, path[st][ed - 1]);
  int R = min(ed - 1, path[st + 1][ed]);
 ret = inf:
 for (int i = L; i <= R; i++) {</pre>
   ll cur = solve(st, i) + solve(i + 1, ed)
   cur += (arr[ed] - arr[st - 1]);
   if (cur < ret) ret = cur; path[st][ed] =</pre>
         i:
 return ret;
///knuth for divide and conquer
int solve(int group, int pos) {
 if (!pos) return dp[group][pos] = 0;
  if (!group) return dp[group][pos] =
      INT MAX:
  int &ret = dp[group][pos];
  if (ret != -1) return ret;
 int L = 1, R = pos;
 if (pos - 1 > 0) {
   solve(group, pos - 1);
```

2.7 Same Color Group

```
int prv[21]; 11 cost[21][21];
11 dp[1 << 21]; int m, n;</pre>
bool ok[1 << 21];</pre>
11 solve(ll mask) {
 if (mask == (1 << m) - 1) return Oll;</pre>
 11 &ret = dp[mask];
 if (ok[mask]) return ret;
 ok[mask] = true; ret = inf;
 for (int i = 0; i < m; i++) {</pre>
   if (!(mask & (1 << i) )) {</pre>
     11 c = 0;
     for (int j = 0; j < m; j++) {
       if ((mask & (1 << j)))</pre>
          c += cost[i][j];
     ret = min(ret, c + solve((mask | (1 <<
   }
 }
 return ret;
}
int arr[MAX];
int main() {
 for (int i = 0; i < n; i++) {</pre>
   int val = arr[i];
   val--; prv[val]++;
   for (int j = 0; j < m; j++) {
     if (val == j) continue;
     cost[val][j] += prv[j];
   }
 11 \text{ ans} = \text{solve}(0);
```

2.8 Triangulation DP

```
bool valid[205][205];
ll dp[205][205];
ll solve(int L, int R) {
   if (L + 1 == R) return 1;
   if (dp[L][R] != -1) return dp[L][R];
   ll ret = 0;
   for (int mid = L + 1; mid < R; mid++) {
      if (valid[L][mid] && valid[mid][R]) {
            ///selecting triangle(P[L], P[mid], P[
            R])
      ll temp = ( solve(L, mid) * solve(mid,
            R) ) % MOD;
      ret = (ret + temp) % MOD;
   }
   return dp[L][R] = ret;
}</pre>
```

2.9 Convex Hull Trick

```
struct line {
    ll m, c;
    line() {}
    line(ll m, ll c) : m(m), c(c) {}
};
struct convex_hull_trick {
    vector<line>lines;
    int ptr = 0;
```

```
convex_hull_trick() {}
  bool bad(line a, line b, line c) {
   return 1.0 * (c.c - a.c) * (a.m - b.m)
        < 1.0 * (b.c - a.c) * (a.m - c.m);
  void add(line L) {
   int sz = lines.size();
   while (sz >= 2 && bad(lines[sz - 2],
        lines[sz - 1], L)) {
     lines.pop_back(); sz--;
   lines.pb(L);
 11 get(int idx, int x) {
   return (111 * lines[idx].m * x + lines[
        idx].c);
 11 query(int x) {
   if (lines.empty()) return 0;
   if (ptr >= lines.size()) ptr = lines.
        size() - 1;
   while (ptr < lines.size() - 1 && get(</pre>
        ptr, x) > get(ptr + 1, x)) ptr++;
   return get(ptr, x);
};
11 sum[MAX];
11 dp[MAX];
int arr[MAX];
int main() {
 fastio:
 int t:
  cin >> t;
 while (t--) {
   int n, a, b, c;
   cin >> n >> a >> b >> c;
   for (int i = 1; i <= n; i++) cin >> sum }
        [i];
   for (int i = 1; i \le n; i++) dp[i] = 0, for (int v = 1; v \le sz; v++) {
         sum[i] += sum[i - 1];
   convex_hull_trick cht;
   cht.add( line(0, 0) );
   for (int pos = 1; pos <= n; pos++) {</pre>
     dp[pos] = cht.query(sum[pos]) - 111 *
          a * sqr(sum[pos]) - c;
     pos] - a * sqr(sum[pos])) );
   ll ans = (-111 * dp[n]);
   ans += (111 * sum[n] * b);
   cout << ans << "\n";
 }
}
```

3 String Algorithms

3.1 Suffix Automaton

```
struct state {
 int len, link;
 map<char, int> next;
 bool is_clone;
 int first_pos;
 vector<int>inv_link;
state st[2 * MAX]:
int mn[2 * MAX], mx[2 * MAX];
int sz, last;
void sa_init() {
 st[0].len = 0;
 st[0].link = -1;
 st[0].next.clear();
 sz = 1;
 last = 0:
void sa_extend(char c) {
 int cur = sz++;
 st[cur].len = st[last].len + 1;
 st[cur].first_pos = st[cur].len - 1;
 st[cur].is_clone = false;
 st[cur].next.clear();
 ///for lcs of n strings
  // mn[cur] = st[cur].len;
 int p = last;
```

```
while (p != -1 && !st[p].next.count(c)) {
   st[p].next[c] = cur;
   p = st[p].link;
 if (p == -1) {
   st[cur].link = 0;
  } else {
    int q = st[p].next[c];
    if (st[p].len + 1 == st[q].len) {
     st[cur].link = q;
   } else {
     int clone = sz++:
     st[clone].len = st[p].len + 1;
     st[clone].next = st[q].next;
     st[clone].link = st[q].link;
     st[clone].first_pos = st[q].first_pos;
     st[clone].is_clone = true;
     ///for lcs of n strings
     // mn[clone] = st[clone].len;
     while (p != -1 && st[p].next[c] == q)
       st[p].next[c] = clone;
       p = st[p].link;
     st[q].link = st[cur].link = clone;
 last = cur;
void radix_sort() {
 for (int i = 0; i < sz; i++) cnt[st[i].len</pre>
      ]++;
  for (int i = 1; i < sz; i++) cnt[i] += cnt</pre>
       [i - 1];
  for (int i = 0; i < sz; i++) order[--cnt[</pre>
       st[i].len]] = i;
// after constructing the automaton
 st[st[v].link].inv_link.push_back(v);
// output all positions of occurrences
void output_all_occurrences(int v, int
     P_length) {
  if (!st[v].is_clone)
   cout << st[v].first_pos - P_length + 1</pre>
        << endl:
  for (int u : st[v].inv_link)
   output_all_occurrences(u, P_length);
//lcs of two strings
string lcs (string S, string T) {
 sa_init();
  for (int i = 0; i < S.size(); i++)</pre>
      sa_extend(S[i]);
  int v = 0, 1 = 0, best = 0, bestpos = 0;
  for (int i = 0; i < T.size(); i++) {</pre>
    while (v && !st[v].next.count(T[i])) {
     v = st[v].link; l = st[v].len;
   if (st[v].next.count(T[i])) {
     v = st[v].next[T[i]]; l++;
   if (1 > best) best = 1; bestpos = i;
 }
 return T.substr(bestpos - best + 1, best);
//lcs of n strings
void add_str(string s) {
 for (int i = 0; i < sz; i++) mx[i] = 0;</pre>
  int v = 0, 1 = 0;
 for (int i = 0; i < s.size(); i++) {</pre>
   while (v && !st[v].next.count(s[i])) {
     v = st[v].link; l = st[v].len;
    if (st[v].next.count(s[i])) {
     v = st [v].next[s[i]]; 1++;
   mx[v] = max(mx[v], 1);
  for (int i = sz - 1; i > 0; i--) mx[st[i].
       link] = max(mx[st[i].link], mx[i]);
  for (int i = 0; i < sz; i++) mn[i] = min(</pre>
       mn[i], mx[i]);
```

```
int lcs() {
 int ret = 0;
 for (int i = 0; i < sz; i++) ret = max(ret</pre>
      , mn[i]);
 return ret;
string s[15];
int arr[MAX];
int main() {
 fastio;
 int n = 0;
 while (cin >> s[n]) n++;
 sa_init();
 for (int i = 0; i < s[0].size(); i++)</pre>
      sa_extend(s[0][i]);
 for (int i = 1; i < n; i++) add_str(s[i]);</pre>
 cout << lcs() << "\n";
```

3.2 Palindromic Tree

```
struct state {
 int len, link;
 map<char, int> next;
state st[MAX];
int id, last;
string s;
11 ans[MAX]:
void init() {
 for (int i = 0; i <= id; i++) {</pre>
   st[i].len = 0; st[i].link = 0;
   st[i].next.clear(); ans[i] = 0;
 st[1].len = -1; st[1].link = 1;
 st[2].len = 0; st[2].link = 1;
 id = 2; last = 2;
void extend(int pos) {
 while (s[pos - st[last].len - 1] != s[pos
      ]) last = st[last].link;
  int newlink = st[last].link;
  char c = s[pos];
  while (s[pos - st[newlink].len - 1] != s[
      pos]) newlink = st[newlink].link;
  if (!st[last].next.count(c)) {
   st[last].next[c] = ++id;
   st[id].len = st[last].len + 2;
   st[id].link = (st[id].len == 1 ? 2 : st[
        newlink].next[c]);
   ans[id] += ans[st[id].link];
   if (st[id].len > 2) {
     int 1 = st[id].len / 2 + (st[id].len %
           2 ? 1 : 0);
     if (h.range_hash(pos - st[id].len + 1,
           pos - st[id].len + 1) == h.
          reverse_hash(pos - st[id].len +
          1, pos - st[id].len + 1)) ans[id
          ]++;
   }
 }
 last = st[last].next[c];
```

3.3 Manacher

3.4 Knuth Morris Pratt

3.5 Aho Corasick

struct state {

```
int len, par, link, next_lif;
 ll val;
 int next[26];
 char p_char;
 bool lif;
 state(int par = -1, char p_char = '$', int
       len = 0) : par(par), p_char(p_char),
       len(len) {
   lif = false:
   link = 0;
   next_lif = 0;
   val = 0:
   memset(next, 0, sizeof next);
vector<state>aho;
inline void add_str(const string &s, ll val)
  int now = 0;
 for (int i = 0; i < s.size(); i++) {</pre>
   int c = s[i] - 'a';
   if (!aho[now].next[c]) {
     aho[now].next[c] = (int)aho.size();
     aho.emplace_back(now, s[i], aho[now].
          len + 1):
   now = aho[now].next[c];
 aho[now].lif = true;
 aho[now].val = val;
inline void push_link() {
 queue<int>q;
  q.push(0);
  while (!q.empty()) {
   int cur = q.front();
   int link = aho[cur].link;
   q.pop();
   aho[cur].next_lif = aho[link].lif ? link
         : aho[link].next lif:
   for (int c = 0; c < 26; c++) {
     if (aho[cur].next[c]) {
       aho[ aho[cur].next[c] ].link = cur ?
             aho[link].next[c] : 0;
       q.push( aho[cur].next[c] );
     } else aho[cur].next[c] = aho[link].
          next[c];
   aho[cur].val += aho[link].val;
 }
inline int count(string &s) {
```

```
for (int i = 0; i < s.size(); i++) {</pre>
   now = aho[now].next[s[i] - 'a'];
   ret += aho[now].val;
 return ret;
struct dynamic_aho {
 aho_corasick ac[20];
 vector<string> dict[20];
 dynamic_aho() {
   for (int i = 0; i < 20; i++) {</pre>
     ac[i].aho.clear();
     dict[i].clear();
   }
 void add_str(string &s) {
   int idx = 0;
   for (; idx < 20 && !ac[idx].aho.empty();</pre>
         idx++) {}
   ac[idx] = aho_corasick();
   ac[idx].add_str(s, 1), dict[idx].pb(s);
   for (int i = 0; i < idx; i++) {</pre>
     for (string x : dict[i]) {
       ac[idx].add_str(x, 1);
       dict[idx].pb(x);
     ac[i].aho.clear(), dict[i].clear();
   ac[idx].push_link();
 }
 inline int count(string &s) {
   int ret = 0:
   for (int i = 0; i < 20; i++) {</pre>
     if (!ac[i].aho.empty()) ret += ac[i].
          count(s);
   return ret;
int arr[MAX];
int main() {
 fastio:
 dynamic_aho add, del;
  int m;
 cin >> m;
 while (m--) {
   int type;
   string s;
   cin >> type >> s;
   if (type == 1) add.add_str(s);
   else if (type == 2) del.add_str(s);
   else cout << add.count(s) - del.count(s)</pre>
         << "\n" << flush;
```

3.6 Z Algorithm

```
vector<int> calcz(string s) {
 int n = s.size();
 vector<int> z(n);
 int 1, r; 1 = r = 0;
 for (int i = 1; i < n; i++) {</pre>
   if (i > r) {
     1 = r = i;
     while (r < n \&\& s[r] == s[r - 1]) r++;
     z[i] = r - 1; r--;
   } else {
     int k = i - 1;
     if (z[k] < r - i + 1) z[i] = z[k];
       1 = i;
       while (r < n \&\& s[r] == s[r - 1]) r
       z[i] = r - 1; r--;
   }
 return z;
```

3.7 String Match FFT

//find occurrences of t in s where '?'s are
 automatically matched with any

s) {

```
character
                                               vector<int> suffix_array_construction(string
//res[i + m - 1] = sum_j = 0 to m - 1_{s[i + j]}
   ] * t[j] * (s[i + j] - t[j])
vector<int> string_matching(string &s,
    string &t) {
 int n = s.size(), m = t.size();
 vector<int> s1(n), s2(n), s3(n);
 for(int i = 0; i < n; i++) s1[i] = s[i] ==</pre>
       '?' ? 0 : s[i] - 'a' + 1; //assign
      any non zero number for non '?'s
 for(int i = 0; i < n; i++) s2[i] = s1[i] *</pre>
       s1[i];
 for(int i = 0; i < n; i++) s3[i] = s1[i] *
       s2[i];
 vector<int> t1(m), t2(m), t3(m);
 for(int i = 0; i < m; i++) t1[i] = t[i] ==</pre>
        '?' ? 0 : t[i] - 'a' + 1;
 for(int i = 0; i < m; i++) t2[i] = t1[i] * }
       t1[i];
  for(int i = 0; i < m; i++) t3[i] = t1[i] *</pre>
       t2[i];
 reverse(t1.begin(), t1.end());
 reverse(t2.begin(), t2.end());
 reverse(t3.begin(), t3.end());
 vector<int> s1t3 = multiply(s1, t3);
 vector<int> s2t2 = multiply(s2, t2);
 vector<int> s3t1 = multiply(s3, t1);
  vector<int> res(n);
 for(int i = 0; i < n; i++) res[i] = s1t3[i]}</pre>
      ] - s2t2[i] * 2 + s3t1[i];
 vector<int> oc;
 for(int i = m - 1; i < n; i++) if(res[i]</pre>
      == 0) oc.push_back(i - m + 1);
 return oc;
```

3.8 Suffix Array

```
vector<vector<int> >c;
vector<int>sort_cyclic_shifts(string const&
    s) {
 int n = s.size();
 const int alphabet = 256;
 vector<int> p(n), cnt(alphabet, 0);
 c.clear(); c.emplace_back(); c[0].resize(n
 for (int i = 0; i < n; i++) cnt[s[i]]++;</pre>
 for (int i = 1; i < alphabet; i++) cnt[i]</pre>
       += cnt[i - 1];
  for (int i = 0; i < n; i++) p[--cnt[s[i]]]</pre>
        = i;
 c[0][p[0]] = 0;
 int classes = 1;
 for (int i = 1; i < n; i++) {</pre>
   if (s[p[i]] != s[p[i - 1]]) classes++;
c[0][p[i]] = classes - 1;
 vector<int> pn(n), cn(n); cnt.resize(n);
 for (int h = 0; (1 << h) < n; h++) {
   for (int i = 0; i < n; i++) {</pre>
     pn[i] = p[i] - (1 << h);
     if (pn[i] < 0) pn[i] += n;</pre>
   fill(cnt.begin(), cnt.end(), 0);
   /// radix sort
   for (int i = 0; i < n; i++) cnt[c[h][pn[</pre>
        i]]]++;
   for (int i = 1; i < classes; i++) cnt[i]</pre>
         += cnt[i - 1];
   for (int i = n - 1; i >= 0; i--) p[--cnt
         [c[h][pn[i]]] = pn[i];
    cn[p[0]] = 0; classes = 1;
   for (int i = 1; i < n; i++) {</pre>
     pii cur = {c[h][p[i]], c[h][(p[i] + (1
           << h)) % n]};
     pii prev = {c[h][p[i - 1]], c[h][(p[i
          - 1] + (1 << h)) % n]};
     if (cur != prev) ++classes;
     cn[p[i]] = classes - 1;
   c.push_back(cn);
 return p;
```

```
s += "!";
  vector<int> sorted_shifts =
      sort_cyclic_shifts(s);
  sorted_shifts.erase(sorted_shifts.begin())
 return sorted_shifts;
/// compare two suffixes starting at i and j
     with length 2<sup>k</sup>
int compare(int i, int j, int n, int k) {
 pii a = \{c[k][i], c[k][(i + 1 - (1 << k))
      % n]};
 pii b = \{c[k][j], c[k][(j + 1 - (1 << k))
     % n]};
 return a == b ? 0 : a < b ? -1 : 1;
int lcp(int i, int j) {
  int log_n = c.size() - 1;
  int ans = 0;
  for (int k = log_n; k >= 0; k--) {
   if (c[k][i] == c[k][j]) {
     ans += 1 << k;
     i += 1 << k; j += 1 << k;
 }
 return ans;
vector<int> lcp_construction(string const& s
    , vector<int> const& p) {
  int n = s.size();
  vector<int> rank(n, 0);
  for (int i = 0; i < n; i++) rank[p[i]] = i</pre>
  int k = 0;
  vector<int> lcp(n, 0);
  for (int i = 0; i < n; i++) {</pre>
   if (rank[i] == n - 1) {
     k = 0:
     continue;
   int j = p[rank[i] + 1];
   while (i + k < n \&\& j + k < n \&\& s[i + k]
        ] == s[j + k]) k++;
   lcp[rank[i]] = k;
   if (k) k--;
 return lcp;
//kth lexicographically smallest substring (
    considering duplicates)
int arr[MAX], lg[MAX], mn[MAX][25];
vector<int>p, lcp;
string s;
// sparse table for min in lcp goes here
///find the rightmost position where get(1,r
    ) > val
int khoj(int 1, int r, int val) {
 int lo = 1 + 1, hi = r, ret = -1;
 while (lo <= hi) {</pre>
   int mid = lo + (hi - lo) / 2;
    if (get(1, mid - 1) > val) {
     ret = mid; lo = mid + 1;
   } else hi = mid - 1;
 return ret;
int done[MAX]:
int arr[MAX];
int main() {
 fastio;
 cin >> s >> k;
 p = suffix_array_construction(s);
 lcp = lcp_construction(s, p);
 build();
 int n = s.size();
  int milaisi = 0;
  for (int i = 0; i < n; i++) {</pre>
   milaisi += done[i];
   int len = n - p[i];
```

```
int cur = milaisi;
  /// cur = i ? lcp[i-1] : 0; this can
      replace all the milaisi and done
  while (cur < len) {</pre>
   int r = khoj(i, n - 1, cur);
   int koyta, milabo;
   if (r == -1) {
     milabo = len - cur;
     koyta = 1;
   } else {
     milabo = get(i, r - 1) - cur;
     koyta = r - i + 1;
   if (koyta * milabo < k) k -= (koyta *</pre>
        milabo);
   else {
     int d = k / koyta; int m = k % koyta
     if (!m) {
       cout << s.substr(p[i], cur + d) <<</pre>
            "\n";
       return 0:
     } else {
       cout \ll s.substr(p[i], cur + d + 1)
             << "\n":
       return 0:
   if (r == -1) break;
   done[r + 1] -= milabo;
   cur = get(i, r - 1);
   milaisi += milabo:
cout << "No such line.\n";</pre>
```

Faster Hash Table

```
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
struct custom_hash {
 const ll rnd = chrono::
      high_resolution_clock::now().
      time_since_epoch().count();
 static unsigned long long hash_f(unsigned
     long long x) {
   x += 0x9e3779b97f4a7c15;
   x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9
   x = (x ^ (x >> 27)) * 0x94d049bb133111eb
   return x ^ (x >> 31);
 11 operator() (11 x) const { return hash_f
      (x) ^ rnd; }
 // static int combine_hash(pii hash) {
      return hash.ff * 31 + hash.ss; }
 // static ll combine_hash(pll hash) {
      return (hash.ff << 32) ^ hash.ss; }</pre>
 // 11 operator() (pll x) const {
 // x.ff = hash_f(x.ff) ^ rnd; x.ss =
      hash_f(x.ss) ^ rnd;
     return combine_hash(x);
 // }
};
gp_hash_table<11, int, custom_hash> mp;
unordered_map<11, int, custom_hash> mp;
```

3.10Double Hash

 $ff * x, a.ss * x);}$

```
ostream& operator << (ostream& os, pll hash)
 return os << "(" << hash.ff << ", " <<
      hash.ss << ")";
pll operator + (pll a, ll x) {return pll(a.
    ff + x, a.ss + x);
pll operator - (pll a, ll x) {return pll(a.
    ff - x, a.ss - x);}
pll operator * (pll a, ll x) {return pll(a.
```

```
pll operator + (pll a, pll x) {return pll(a.
    ff + x.ff, a.ss + x.ss);}
pll operator - (pll a, pll x) {return pll(a.
    ff - x.ff, a.ss - x.ss);}
pll operator * (pll a, pll x) {return pll(a.
    ff * x.ff, a.ss * x.ss);}
pll operator % (pll a, pll m) {return pll(a.
    ff % m.ff, a.ss % m.ss);}
pll base(1949313259, 1997293877);
pll mod(2091573227, 2117566807);
pll power (pll a, ll p) {
 if (!p) return pll(1, 1);
 pll ans = power(a, p / 2);
 ans = (ans * ans) % mod;
 if (p % 2) ans = (ans * a) % mod;
 return ans;
pll inverse(pll a) {
 return power(a, (mod.ff - 1) * (mod.ss -
      1) - 1);
pll inv_base = inverse(base);
pll val;
vector<pll> P;
void hash_init(int n) {
 P.resize(n + 1):
 P[0] = pll(1, 1);
 for (int i = 1; i <= n; i++) P[i] = (P[i -
       1] * base) % mod;
///appends c to string
pll append(pll cur, char c) {
 return (cur * base + c) % mod;
///prepends c to string with size k
pll prepend(pll cur, int k, char c) {
 return (P[k] * c + cur) % mod;
///replaces the i-th (0-indexed) character
    from right from a to b;
pll replace(pll cur, int i, char a, char b)
  cur = (cur + P[i] * (b - a)) \% mod;
 return (cur + mod) % mod;
///Erases c from the back of the string
pll pop_back(pll hash, char c) {
 return (((hash - c) * inv_base) % mod +
      mod) % mod:
///Erases c from front of the string with
pll pop_front(pll hash, int len, char c) {
 return ((hash - P[len - 1] * c) % mod +
      mod) % mod;
///concatenates two strings where length of
    the right is k
pll concat(pll left, pll right, int k) {
 return (left * P[k] + right) % mod;
///Calculates hash of string with size len
    repeated cnt times
///This is O(\log n). For O(1), pre-calculate
     inverses
pll repeat(pll hash, int len, ll cnt) {
 pll mul = (P[len * cnt] - 1) * inverse(P[
      len] - 1);
 mul = (mul % mod + mod) % mod;
 pll ret = (hash * mul) % mod;
```

```
if (P[len].ff == 1) ret.ff = hash.ff * cnt } else {
 if (P[len].ss == 1) ret.ss = hash.ss * cnt
 return ret;
ll get(pll hash) {
 return ( (hash.ff << 32) ^ hash.ss );</pre>
struct hashlist {
 int len;
 vector<pll> H, R;
 hashlist() {}
 hashlist(string &s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, pll(0, 0)), R.resize(
       len + 2, pl1(0, 0));
   for (int i = 1; i <= len; i++) H[i] =</pre>
        append(H[i - 1], s[i - 1]);
   for (int i = len; i >= 1; i--) R[i] =
        append(R[i + 1], s[i - 1]);
 /// 1-indexed
 inline pll range_hash(int 1, int r) {
   int len = r - l + 1;
   return ((H[r] - H[l - 1] * P[len]) % mod
         + mod) % mod;
 inline pll reverse_hash(int 1, int r) {
   int len = r - l + 1;
   return ((R[1] - R[r + 1] * P[len]) % mod
         + mod) % mod;
 inline pll concat_range_hash(int l1, int
      r1, int 12, int r2) {
   int len_2 = r2 - 12 + 1;
   return concat(range_hash(l1, r1),
        range_hash(12, r2), len_2);
 inline pll concat_reverse_hash(int 11, int
       r1, int 12, int r2) {
   int len_1 = r1 - l1 + 1;
   return concat(reverse_hash(12, r2),
        reverse_hash(l1, r1), len_1);
 }
```

Data Structures $\mathbf{4}$

4.1Persistent Segment Tree

```
struct node {
 int val, lft, rt;
 node(int val = 0, int lft = 0, int rt = 0)
       : val(val), lft(lft), rt(rt) {}
node nodes[30 * MAX]; ///take at least 2*n*
    log(n) nodes
int root[MAX], sz;
inline int update(int &now, int L, int R,
    int idx, int val) {
 if (L > idx || R < idx) return now;</pre>
 if (L == R) {
   nodes[sz] = nodes[now];
   nodes[sz].val += val:
   return sz;
 int mid = (L + R) >> 1;
 int ret = ++sz;
 if (idx <= mid) {</pre>
   if (!nodes[now].lft) nodes[now].lft = ++
   nodes[ret].lft = update(nodes[now].lft,
        L, mid, idx, val);
   nodes[ret].rt = nodes[now].rt;
```

```
if (!nodes[now].rt) nodes[now].rt = ++sz
    nodes[ret].rt = update(nodes[now].rt,
         mid + 1, R, idx, val);
    nodes[ret].lft = nodes[now].lft;
 nodes[ret].val = nodes[ nodes[ret].lft ].
       val + nodes[ nodes[ret].rt ].val;
 return ret;
inline int query(int &now, int L, int R, int
      i, int j) {
  if (L > j \mid \mid R < i) return 0;
  if (L >= i && R <= j) return nodes[now].</pre>
       val:
  int mid = (L + R) >> 1;
  return query(nodes[now].lft, L, mid, i, j)
        + query(nodes[now].rt, mid + 1, R, i
       , j);
/// in main(make segtree for every prefix)
root[0] = 0;
for (int i = 1; i <= n; i++) root[i] =
    update(root[i - 1], 1, n, p[i], 1);</pre>
```

4.2LCA

```
/*Hey, What's up?*/
#include<bits/stdc++.h>
using namespace std;
#define pi acos(-1.0)
#define fastio ios_base::sync_with_stdio(
     false);cin.tie(NULL);cout.tie(NULL)
vector<long long>v[100005],vc;
long long x[200005][40],mp[100005],ml
     [100005],nd,pos[100005];
void build(long long n)
   long long a,i,j,k,b,c;
   a=1:
   for(i=0; i<n; i++)</pre>
       x[i][0]=vc[i];
   b=1:
   while(a<n)</pre>
       for(i=0; i<n-a; i++)</pre>
           x[i][b]=min(x[i][b-1],x[i+a][b
                -1]);
       a*=2:
       b++;
   }
   return:
long long query(long long a, long long b)
   long long c,d,e,f;
   //if(a==b)return x[a][0];
   c=log2(1.0*(b-a+1));
   //cout<<c<' ';
   f=powl(1.0*2,1.0*c);
   d=x[a][c];
   e=x[b-f+1][c];
   //cout<<b-f+1<<' ';
   return min(d,e);
void tour_de_euler(long long p, long long q)
   vc.push_back(mp[p]);
   //nd++:
   if(!pos[mp[p]])pos[mp[p]]=nd;
   for(int i=0;i<v[p].size();i++){</pre>
       if(v[p][i]==q)continue;
       tour_de_euler(v[p][i],p);
       vc.push_back(mp[p]); nd++;
   return;
```

```
void dfs(long long p, long long q)
   mp[p]=nd;
   ml[nd]=p;
   nd++:
   for(int i=0;i<v[p].size();i++){</pre>
       if(v[p][i]==q)continue;
       dfs(v[p][i],p);
   return:
long long lca(long long a, long long b)
   a=pos[mp[a]];
   b=pos[mp[b]];
   if(a>b)
       swap(a,b);
   long long c=query(a,b);
   return ml[c];
int main()
{
   long long a=0,b=0,c,d,e,f=0,1,g,m,n,k,i,
        j,t,p,q;
   cin>>n;
   for(i=1; i<n; i++)</pre>
       cin>>a>>b;
       v[a].push_back(b);
       v[b].push_back(a);
   nd=1:
   dfs(1,-1):
   vc.push_back(696969696969);
   nd=1:
   tour_de_euler(1,-1);
   l=vc.size();
   build(1+2);
   cin>>q;
   while(q--){
       cin>>a>>b:
       cout << lca(a,b) << endl;
   return 0;
```

Hopcroft Karp

```
#include<bits/stdc++.h>
using namespace std;
const int N = 3e5 + 9;
struct HopcroftKarp {
 static const int inf = 1e9;
 int n:
 vector<int> 1, r, d;
  vector<vector<int>> g;
 HopcroftKarp(int _n, int _m) {
   n = _n;
   int p = _n + _m + 1;
   g.resize(p);
   1.resize(p, 0);
   r.resize(p, 0);
   d.resize(p, 0);
 void add_edge(int u, int v) {
   g[u].push_back(v + n); //right id is
        increased by n, so is l[u]
 }
 bool bfs() {
   queue<int> q;
   for (int u = 1; u <= n; u++) {</pre>
     if (!1[u]) d[u] = 0, q.push(u);
     else d[u] = inf;
   d[0] = inf;
```

```
while (!q.empty()) {
     int u = q.front();
     q.pop();
     for (auto v : g[u]) {
       if (d[r[v]] == inf) {
         d[r[v]] = d[u] + 1;
         q.push(r[v]);
   return d[0] != inf;
 bool dfs(int u) {
   if (!u) return true;
   for (auto v : g[u]) {
     if(d[r[v]] == d[u] + 1 && dfs(r[v])) {
       l[u] = v;
       r[v] = u;
       return true;
     }
   d[u] = inf;
   return false;
 int maximum_matching() {
   int ans = 0;
   while (bfs()) {
     for(int u = 1; u <= n; u++) if (!1[u]</pre>
          && dfs(u)) ans++;
   return ans;
 }
int32_t main() {
 ios_base::sync_with_stdio(0);
 cin.tie(0);
 int n, m, q;
 cin >> n >> m >> q;
 HopcroftKarp M(n, m);
 while (q--) {
   int u, v;
   cin >> u >> v;
   M.add_edge(u, v);
 cout << M.maximum_matching() << '\n';</pre>
 return 0;
```

Merge Sort Tree 4.4

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template<typename T> using ordered_set =
    tree<T, null_type, less<T>, rb_tree_tag
     , tree_order_statistics_node_update>;
ordered_set<pii> bst[MAX << 2];
void init(int n) {
 for (int i = 0; i <= 4 * n; i++) bst[i].</pre>
void build(int now, int L, int R) {
 if (L == R) {
   bst[now].insert({arr[L], L});
   return;
 for (int i = L; i <= R; i++) bst[now].</pre>
      insert({arr[i], i});
  int mid = (L + R) / 2;
 build(now << 1, L, mid);</pre>
 build((now << 1) | 1, mid + 1, R);
void update(int now, int L, int R, int idx,
    int ager_val, int val) {
 if (L == R) {
   bst[now].erase(bst[now].find({ager_val,
        idx})):
   bst[now].insert({val, idx});
```

return:

```
int mid = (L + R) / 2;
  if (idx <= mid) update(now << 1, L, mid,</pre>
       idx, ager_val, val);
  else update((now << 1) | 1, mid + 1, R,</pre>
      idx, ager_val, val);
  bst[now].erase(bst[now].find({ager_val,
       idx}));
  bst[now].insert({val, idx});
ll query(int now, int L, int R, int i, int j
     , int val) {
  if (R < i || L > j) return 0;
 if (L >= i && R <= j) {
    int ret = bst[now].order_of_key({val,
        INT_MAX});
   return ret;
  int mid = (L + R) / 2;
  return query(now << 1, L, mid, i, j, val)</pre>
       + query((now << 1) | 1, mid + 1, R, i
       , j, val);
4.5
       Trie
```

```
int trie[30 * 100000 + 5][2];
int mark[30 * 100000 + 5];
int node = 1:
void add(int n) {
 int now = 1;
 for (int i = 27; i >= 0; i--) {
   int d = (bool)(n & (1 << i));</pre>
   if (!trie[now][d]) trie[now][d] = ++node
   now = trie[now][d];
   mark[now]++:
 }
void del(int n) {
 int now = 1;
 deque<int>v;
 for (int i = 27; i >= 0; i--) {
   int d = (bool)(n & (1 << i));</pre>
   if (trie[now][d]) {
     v.push_front(now);
     now = trie[now][d];
     mark[now] --;
   }
 v.push_front(now);
 for (int i = 1; i < v.size(); i++) {</pre>
   if (!mark[v[i - 1]]) {
     if (trie[v[i]][0] == v[i - 1]) trie[v[
          i]][0] = 0;
     if (trie[v[i]][1] == v[i - 1]) trie[v[
          i]][1] = 0;
 }
```

Implicit Treap 4.6

```
mt19937 rnd(chrono::steady_clock::now().
    time_since_epoch().count());
typedef struct node* pnode;
struct node {
 int prior, sz;
 11 val, sum, lazy;
 bool rev;
 node *lft. *rt:
 node(int val = 0, node *lft = NULL, node *
      rt = NULL) : lft(lft), rt(rt), prior(
      rnd()), sz(1), val(val), rev(false),
      sum(0), lazy(0) {}
struct implicit_treap {
 pnode root;
 implicit_treap() {
   root = NULL:
 int get_sz(pnode now) {
```

```
return now ? now->sz : 0;
void update_sz(pnode now) {
 if (!now) return;
 now->sz = 1 + get_sz(now->lft) + get_sz(
      now->rt);
// lazy sum
void push(pnode now) {
 if (!now || !now->lazy) return;
 now->val += now->lazy;
 now->sum += get_sz(now) * now->lazy;
 if (now->lft) now->lft->lazy += now->
      lazy;
 if (now->rt) now->rt->lazy += now->lazy;
 now->lazy = 0;
}
void combine(pnode now) {
 if (!now) return;
 now->sum = now->val; // reset the node
 push(now->lft), push(now->rt); // update
       lft and rt
 now->sum += (now->lft ? now->lft->sum :
      0) + (now->rt ? now->rt->sum : 0);
// reverse substring
// void push(pnode now) {
   if (!now || !now->rev) return;
// now->rev = false;
    swap(now->lft, now->rt);
    if (now->lft) now->lft->rev ^= true;
    if (now->rt) now->rt->rev ^= true;
// }
// sort ascending or descending
// void push(pnode now) {
   if (!now || !now->sort_kor) return;
    if (now->sort_kor == -1) swap(now->
    lft, now->rt);
    int cnt[26];
    for (int i = 0; i < 26; i++) cnt[i] =
     now->cnt[i]:
    int idx = 0;
//
    if (now->lft) {
      memset(now->lft->cnt, 0, sizeof now
    ->lft->cnt);
      int lft_sz = get_sz(now->lft);
      while (idx < 26 && lft_sz) {
11
        int mn = min(cnt[idx], lft_sz);
//
        now->lft->cnt[idx] = mn;
        cnt[idx] -= mn; lft_sz -= mn;
        if (!cnt[idx]) idx++;
//
//
     now->lft->sort_kor = now->sort_kor;
//
//
    while (!cnt[idx]) idx++;
11
    now->val = idx, cnt[idx]--;
    if (!cnt[idx]) idx++;
    if (now->rt) {
//
      memset(now->rt->cnt, 0, size of now
    ->rt->cnt);
      int rt_sz = get_sz(now->rt);
//
      while (idx < 26 && rt_sz) {
        int mn = min(cnt[idx], rt_sz);
        now->rt->cnt[idx] = mn;
        cnt[idx] -= mn; rt_sz -= mn;
//
        if (!cnt[idx]) idx++;
11
//
     now->rt->sort_kor = now->sort_kor;
//
   if (now->sort_kor == -1) swap(now->
    lft, now->rt);
11
   now->sort_kor = 0;
// }
// void combine(pnode now) {
// if (!now) return;
    memset(now->cnt, 0, sizeof now->cnt);
    for (int i = 0; i < 26; i++) {
      now->cnt[i] = (now->lft ? now->lft
     ->cnt[i] : 0) + (now->rt ? now->rt->
    cnt[i] : 0);
    now->cnt[now->val]++;
// }
```

```
///first pos ta elements go to left,
    others go to right
void split(pnode now, pnode &lft, pnode &
    rt, int pos, int add = 0) {
 if (!now) return void(lft = rt = NULL);
 push(now);
 int cur = add + get_sz(now->lft);
 if (cur < pos) split(now->rt, now->rt,
      rt, pos, cur + 1), lft = now;
 else split(now->lft, lft, now->lft, pos,
       add), rt = now;
 update_sz(now); combine(now);
void merge(pnode &now, pnode lft, pnode rt
 push(lft);
 push(rt);
 if (!lft || !rt) now = lft ? lft : rt;
 else if (lft->prior > rt->prior) merge(
      lft->rt, lft->rt, rt), now = lft;
 else merge(rt->lft, lft, rt->lft), now =
       rt:
 update_sz(now); combine(now);
void insert(int pos, ll val) {
 if (!root) return void(root = new node(
      val));
 pnode lft, rt;
 split(root, lft, rt, pos - 1);
 pnode notun = new node(val);
 merge(root, lft, notun);
 merge(root, root, rt);
void erase(int pos) {
 pnode lft, rt, temp;
 split(root, lft, rt, pos);
 split(lft, lft, temp, pos - 1);
 merge(root, lft, rt);
 delete(temp);
void reverse(int 1, int r) {
 pnode lft, rt, mid;
 split(root, lft, mid, l - 1);
 split(mid, mid, rt, r - l + 1);
 mid->rev ^= true;
 merge(root, lft, mid);
 merge(root, root, rt);
void right_shift(int 1, int r) {
 pnode lft, rt, mid, last;
 split(root, lft, mid, l - 1);
 split(mid, mid, rt, r - l + 1);
 split(mid, mid, last, r - 1);
 merge(mid, last, mid);
 merge(root, lft, mid);
 merge(root, root, rt);
void output(pnode now, vector<int>&v) {
 if (!now) return;
 push(now);
 output(now->lft, v);
 v.pb(now->val);
 output(now->rt, v);
vector<int>get_arr() {
 vector<int>ret;
 output(root, ret);
 return ret;
```

4.7 Mo's Algorithm

```
vector<pair<long long,pair<long long,</pre>
    long long> > x[d+2];
//map<pair<long long,long long> ,long
     long>mp;
v.push_back(-37);
for(i=0;i<n;i++){</pre>
    cin>>a:
    v.push_back(a);
cin>>q;
for(i=0;i<q;i++){</pre>
    cin>>a>>b:
    e=a/d;
    x[e].push_back({b,{a,i}});
for(i=0;i<=d+1;i++){</pre>
    sort(x[i].begin(),x[i].end());
for(i=0;i<=d;i++){</pre>
    memset(vis,0,sizeof(vis));
    l=i*d+1;
    r=i*d:
    p=x[i].size();
    f=0:
    for(j=0;j<p;j++){</pre>
        b=x[i][j].first;
        a=x[i][j].second.first;
        while(r<b){</pre>
           r++;
            vis[v[r]]++;
            if(vis[v[r]]==1)f++;
        //cout<<l<< ' '<<r<<'='<<f<<endl;
        if(1<a){</pre>
        while(l<a){</pre>
            vis[v[1]]--;
            if(vis[v[1]]==0)f--;
            1++;
        else if(l>a){
            while(l>a){
               1--:
                vis[v[1]]++;
                if(vis[v[1]]==1)f++;
        //cout<<l<' '<<r<'='<<f<<endl;
       y[x[i][j].second.second]=f;
        //cout<<a<<' '<<b<<'='<<f<<endl:
for(i=0;i<q;i++){</pre>
    cout << y[i] << '\n';
return 0;
```

4.8 Link Cut Tree

```
struct SplayTree {
 struct node {
   int ch[2] = {0, 0}, p = 0;
   11 self = 0, path = 0;
   11 \text{ sub} = 0, \text{ extra} = 0;
   bool rev = false;
 };
 vector<node> T;
 SplayTree(int n) : T(n + 1) {}
 void push(int x) {
   if (!x) return;
   int 1 = T[x].ch[0], r = T[x].ch[1];
   if (T[x].rev) {
     T[1].rev ^= true, T[r].rev ^= true;
     swap(T[x].ch[0], T[x].ch[1]);
     T[x].rev = false;
 void pull(int x) {
   int 1 = T[x].ch[0], r = T[x].ch[1];
   push(1), push(r);
   T[x].path = T[x].self + T[1].path + T[r]
        ].path;
```

```
T[x].sub = T[x].self + T[x].extra + T[1]
        ].sub + T[r].sub;
  void set(int parent, int child, int d) {
   T[parent].ch[d] = child;
   T[child].p = parent;
   pull(parent);
 }
 int dir(int x) {
   int parent = T[x].p;
   if (!parent) return -1;
   return (T[parent].ch[0] == x) ? 0 : (T[
        parent].ch[1] == x) ? 1 : -1;
 void rotate(int x) {
   int parent = T[x].p, gparent = T[parent
        ].p;
   int dx = dir(x), dp = dir(parent);
   set(parent, T[x].ch[!dx], dx);
   set(x, parent, !dx);
   if (~dp) set(gparent, x, dp);
   T[x].p = gparent;
 void splay(int x) {
   push(x);
   while (~dir(x)) {
     int parent = T[x].p;
     int gparent = T[parent].p;
     push(gparent), push(parent), push(x);
     int dx = dir(x), dp = dir(parent);
     if (~dp) rotate(dx != dp ? x : parent)
     rotate(x);
 }
struct LinkCut : SplayTree {
 LinkCut(int n) : SplayTree(n) {}
 void cut_right(int x) {
   splay(x);
   int r = T[x].ch[1];
   T[x].extra += T[r].sub;
   T[x].ch[1] = 0, pull(x);
 int access(int x) {
   int u = x, v = 0;
   for (; u; v = u, u = T[u].p) {
     cut_right(u);
     T[u].extra -= T[v].sub;
     T[u].ch[1] = v, pull(u);
   return splay(x), v;
 void make_root(int x) {
   access(x);
   T[x].rev ^= true, push(x);
 void link(int u, int v) {
   make_root(v), access(u);
   T[u].extra += T[v].sub;
   T[v].p = u, pull(u);
 void cut(int u) {
   access(u):
   T[u].ch[0] = T[T[u].ch[0]].p = 0;
   pull(u);
 void cut(int u, int v) {
   make_root(u), access(v);
T[v].ch[0] = T[u].p = 0, pull(v);
 int find_root(int u) {
   access(u), push(u);
   while (T[u].ch[0]) {
     u = T[u].ch[0], push(u);
   return splay(u), u;
 }
 int lca(int u, int v) {
   if (u == v) return u;
   access(u);
   int ret = access(v);
   return T[u].p ? ret : 0;
```

```
// subtree query of u if v is the root
 11 subtree(int u, int v) {
   make_root(v), access(u);
   return T[u].self + T[u].extra;
 ll path(int u, int v) {
   make_root(u), access(v);
   return T[v].path;
 // point update
 void update(int u, ll val) {
   T[u].self = val, pull(u);
 }
};
```

Sparse Table 4.9

```
const int maxn = (1 << 20) + 5;
int logs[maxn] = {0};
void compute_logs(){
   logs[1] = 0;
   for(int i=2;i<(1<<20);i++){</pre>
       logs[i] = logs[i/2]+1;
class Sparse_Table
       vector <vector<LL>> table:
       function < LL(LL,LL) > func;
       LL identity;
   Sparse_Table(vector <LL> &v, function <
        LL(LL,LL)> _func, LL id){
       if(logs[2] != 1) compute_logs();
       int sz = v.size();
       table.assign(sz,vector <LL>(logs[sz
            ]+1));
       func = _func, identity = id;
       for(int j=0;j<=logs[sz];j++){</pre>
           for(int i=0;i+(1<<j)<=sz;i++){</pre>
              if(j==0) table[i][j] = func(v
                   [i],id); // base case,
                   when only 1 element in
              else table[i][j] = func(table
                   [i][j-1], table[i +
                    (1<<(j-1))][j-1]);
          }
       }
   // when intersection of two ranges wont
        be a problem like min, gcd, max
   LL query(int 1, int r){
       assert(r>=1);
       int pow = logs[r-l+1];
       return func(table[1][pow], table[r-
            (1<<pow) + 1] [pow]);
   // other cases like sum
   LL Query(int 1,int r){
       if(l>r) return identity; // handle
            basecase
       int pow = logs[r - l + 1];
       return func(table[1][pow], Query(1
            +(1<<pow), r));
   }
};
```

4.10DSU on Tree

```
///Query: Number of distinct names among all
     the k'th son of a node.
const int N = 100005;
string name[N];
vector<int>G[N]:
vector<pii>Q[N];
int L[N],ans[N];
void dfs(int v,int d){
   L[v]=d;
```

```
for(int i:G[v]) dfs(i,d+1);
   return;
void dsu(int v,map<int,set<string>>&mp){
    for(int i:G[v]){
       map<int,set<string>>s;
       dsu(i,s);
       if(s.size()>mp.size()) swap(mp,s);
       for(auto it:s) mp[it.ff].insert(all(
            it.ss));
   if(v!=0) mp[L[v]].insert(name[v]); //
        Here zero is not a actual node
    for(pii p:Q[v]) ans[p.ss] = mp[p.ff].
        size();
   return;
int main(){
   int n;
   cin >> n:
    FOR(i,1,n){
       int u;
       cin >> name[i] >> u;
       G[u].pb(i);
   dfs(0,0);
   int q;
   cin >>q;
    FOR(i,1,q){
       int v,k;
       cin >> v >> k;
       Q[v].pb(pii(k+L[v],i)); //Actual
            level
   map<int,set<string>>mp;
    dsu(0,mp);
    FOR(i,1,q) cout << ans[i] << '\n';
    return 0;
```

4.11**Dominator Tree**

```
struct dominator {
 int n, d_t;
 vector<vector<int>> g, rg, tree, bucket;
 vector<int> sdom, dom, par, dsu, label,
      val, rev;
 dominator() {}
 dominator(int n) :
   n(n), d_t(0), g(n + 1), rg(n + 1),
   tree(n + 1), bucket(n + 1), sdom(n + 1),
   dom(n + 1), par(n + 1), dsu(n + 1),
   label(n + 1), val(n + 1), rev(n + 1)
 { for (int i = 1; i <= n; i++) sdom[i] =
      dom[i] = dsu[i] = label[i] = i; }
  void add_edge(int u, int v) { g[u].pb(v);
 int dfs(int u) {
   d_t++;
   val[u] = d_t, rev[d_t] = u;
   label[d_t] = sdom[d_t] = dom[d_t] = d_t;
   for (int v : g[u]) {
     if (!val[v]) {
       dfs(v);
      par[val[v]] = val[u];
     rg[val[v]].pb(val[u]);
 int findpar(int u, int x = 0) {
   if (dsu[u] == u) return x ? -1 : u;
   int v = findpar(dsu[u], x + 1);
   if (v < 0) return u;</pre>
   if (sdom[label[dsu[u]]] < sdom[label[u</pre>
        ]]) label[u] = label[dsu[u]];
   dsu[u] = v;
   return x ? v : label[u];
 void join(int u, int v) { dsu[v] = u; }
```

vector<vector<int>> build(int s) {

dfs(s);

```
for (int i = n; i >= 1; i--) {
     for (int j = 0; j < rg[i].size(); j++)</pre>
       sdom[i] = min(sdom[i], sdom[ findpar
            (rg[i][j]) ]);
     if (i > 1) bucket[sdom[i]].pb(i);
     for (int w : bucket[i]) {
       int v = findpar(w);
       if (sdom[v] == sdom[w]) dom[w] =
            sdom[w];
       else dom[w] = v;
     if (i > 1) join(par[i], i);
   for (int i = 2; i <= n; i++) {</pre>
     if (dom[i] != sdom[i]) dom[i] = dom[
          dom[i]];
     tree[rev[i]].pb(rev[dom[i]]);
     tree[rev[dom[i]]].pb(rev[i]);
   return tree;
};
```

```
Treap
4.12
mt19937 rnd(chrono::steady_clock::now().
    time_since_epoch().count());
typedef struct node* pnode;
struct node {
 int prior, val, sz;
 ll sum:
 node *lft, *rt;
 node(int val = 0, node *lft = NULL, node *
      rt = NULL) :
   lft(lft), rt(rt), prior(rnd()), val(val)
        , sz(1), sum(0) {}
struct treap {
 pnode root;
 treap() {
   root = NULL;
 int get_sz(pnode now) {
   return now ? now->sz : 0;
 }
 void update_sz(pnode now) {
   if (!now) return;
   now->sz = 1 + get_sz(now->lft) + get_sz(
        now->rt);
 ll get(pnode now) {
   return now ? now->sum : 0;
 void push(pnode now) {}
 void combine(pnode now) {
   if (!now) return;
   now->sum = now->val + get(now->lft) +
        get(now->rt);
 pnode unite(pnode lft, pnode rt) {
   if (!lft || !rt) return lft ? lft : rt;
   // push(lft), push(rt); this not tested
   if (lft->prior < rt->prior) swap(lft, rt
   pnode 1, r;
   split(rt, 1, r, lft->val);
   lft->lft = unite(lft->lft, 1), update_sz
        (lft):
   lft->rt = unite(lft->rt, r), update_sz(
        lft):
   // combine(lft); this not tested
   return lft;
 ///value < val goes to left, value >= val
      goes to right
 void split(pnode now, pnode &lft, pnode &
      rt, int val, int add = 0) {
   push(now):
   if (!now) return void(lft = rt = NULL);
   if (now->val < val) split(now->rt, now->
        rt, rt, val), lft = now;
```

```
else split(now->lft, lft, now->lft, val) 4.13 Implicit Segment Tree
      , rt = now;
 update_sz(now), combine(now);
void merge(pnode &now, pnode lft, pnode rt
 push(lft), push(rt);
 if (!lft || !rt) now = lft ? lft : rt;
 else if (lft->prior > rt->prior) merge(
      lft->rt, lft->rt, rt), now = lft;
  else merge(rt->lft, lft, rt->lft), now =
 update_sz(now), combine(now);
void insert(pnode &now, pnode notun) {
 if (!now) return void(now = notun);
 push(now);
 if (notun->prior > now->prior) split(now
      , notun->lft, notun->rt, notun->val
      ), now = notun;
 else insert(notun->val < now->val ? now
      ->lft : now->rt, notun);
 update_sz(now), combine(now);
void erase(pnode &now, int val) {
 push(now);
 if (now->val == val) {
   pnode temp = now;
   merge(now, now->lft, now->rt);
   delete(temp);
 } else erase(val < now->val ? now->lft :
       now->rt, val);
 update_sz(now), combine(now);
int get_idx(pnode &now, int val) {
 if (!now) return INT_MIN;
 else if (now->val == val) return 1 +
      get_sz(now->lft);
 else if (val < now->val) return get_idx(
      now->lft, val);
 else return (1 + get_sz(now->lft) +
      get_idx(now->rt, val));
int find_kth(pnode &now, int k) {
 if (k < 1 | | k > get_sz(now)) return -1;
 if (get_sz(now->lft) + 1 == k) return
      now->val;
  if (k <= get_sz(now->lft)) return
      find_kth(now->lft, k);
 return find_kth(now->rt, k - get_sz(now
      ->lft) - 1);
11 prefix_sum(pnode &now, int k) {
 if (k < 1 || k > get_sz(now)) return -
  if (get_sz(now->lft) + 1 == k) return
      get(now->lft) + now->val;
  if (k <= get_sz(now->lft)) return
      prefix_sum(now->lft, k);
 return get(now->lft) + now->val +
      prefix_sum(now->rt, k - get_sz(now
      ->lft) - 1);
pnode get_rng(int 1, int r) { ///gets all
    1 <= values <= r
 pnode lft, rt, mid;
 split(root, lft, mid, 1);
 split(mid, mid, rt, r + 1);
 merge(root, lft, rt);
 return mid;
void output(pnode now, vector<int>&v) {
 if (!now) return;
 output(now->lft, v);
 v.pb(now->val);
 output(now->rt, v);
vector<int>get_arr() {
 vector<int>ret;
 output(root, ret);
 return ret;
```

```
struct node {
 int val;
 node *lft, *rt;
 node() {}
 node(int val = 0) : val(val), lft(NULL),
      rt(NULL) {}
struct implicit_segtree {
 node *root;
 implicit_segtree() {}
 implicit_segtree(int n) {
   root = new node(n);
 void update(node *now, int L, int R, int
      idx, int val) {
   if (L == R) {
     now -> val += val;
     return;
   int mid = L + (R - L) / 2;
   if (now->lft == NULL) now->lft = new
       node(mid - L + 1);
   if (now->rt == NULL) now->rt = new node(
        R - mid);
   if (idx <= mid) update(now->lft, L, mid,
        idx, val);
   else update(now->rt, mid + 1, R, idx,
       val):
   now->val = (now->lft)->val + (now->rt)->
        val;
 int query(node *now, int L, int R, int k)
   if (L == R) return L;
   int mid = L + (R - L) / 2;
   if (now->lft == NULL) now->lft = new
       node(mid - L + 1);
   if (now->rt == NULL) now->rt = new node(
       R - mid);
   if (k <= (now->lft)->val) return query(
       now->lft, L, mid, k);
   else return query(now->rt, mid + 1, R, k
         - (now->lft)->val);
 }
```

SumOfDivisors 4.14

```
#include<bits/stdc++.h>
using namespace std;
#define hlw ios_base::sync_with_stdio(false)
     ;cin.tie(NULL);cout.tie(NULL)
long long sod[10000007];
bool pm[10000007];
vector<long long>prm;
int main()
   hlw;
   long long i,j,n,t,d,p;
   sod[1]=1:
   prm.push_back(2);
    for(i=3; i<10000007; i+=2)</pre>
       if(!pm[i])
           prm.push_back(i);
           for(j=i*i; j<10000007; j+=2*i)</pre>
               pm[j]=1;
       }
   for(i=1; i<=10000007;i++)</pre>
       sod[i]=1;
   for(i=0; iprm.size(); i++)
       p=prm[i];
       for(j=p; j<10000007; j+=p)</pre>
```

d=p;

```
while(j%d==0)
           d*=p;
       d/=(p-1);
       sod[j]*=d;
   }
return 0;
```

Geometry

2D Point

```
const double PI = acos(-1), EPS = 1e-10;
template <typename DT> DT sq(DT x) {return x
     * x; }
template <typename DT> int dcmp(DT x) {
    return fabs(x) < EPS ? 0 : (x<0 ? -1 :
template <typename DT>
class point{
   public:
       DT x,y;
   point() = default;
   point(DT x, DT y): x(x), y(y) {};
       : x(p.x), y(p.y) {};
   //opeartions on complex numbers
   point operator * (point rhs) const {
        return point(x * rhs.x - y * rhs.y,
         x * rhs.y + y * rhs.x);}
   point operator / (point rhs) const {
        ) / ~(rhs);}
   bool operator < (point rhs) const {</pre>
        return x < rhs.x or (x == rhs.x and
         y < rhs.y); }
   DT operator & (point rhs) const {
        return x * rhs.y - y * rhs.x; } //
        cross product
   DT operator ^ (point rhs) const {
        return x * rhs.x + y * rhs.y; } //
        dot product
   DT operator ~()
                     const {return sq(x) +
        sq(y); }
                                   //square
   friend istream& operator >> (istream &is
        , point &p) { return is >> p.x >> p
        .y; }
   friend DT DisSq(point a, point b){
        return sq(a.x-b.x) + sq(a.y-b.y); }
   friend DT TriArea(point a, point b,
        point c) { return (b-a) & (c-a); }
   friend DT UTriArea(point a, point b,
        point c) { return abs(TriArea(a, b,
   friend bool Collinear(point a, point b,
        point c) { return UTriArea(a, b, c)
         < EPS; }
   friend double Angle(point u) { return
        atan2(u.y, u.x); }
   friend double Angle(point a, point b) {
       double ans = Angle(b) - Angle(a);
       return ans <= -PI ? ans + 2*PI : (
            ans > PI ? ans - 2*PI : ans);
   friend point Perp(point a){
       return point(-a.y, a.x);
   friend operator Orientation(point a,
        point b, point c) {return dcmp(
        TriArea(a, b, c));}
template <typename DT> using polygon =
    vector <point <DT>>;
template <typename DT>
class polarComp {
   point <DT> 0, dir;
```

```
bool half(point <DT> p) {
       return dcmp(dir & p) < 0 || (dcmp(</pre>
            dir & p) == 0 && dcmp(dir ^ p) >
   public:
   polarComp(point <DT> 0 = point(0, 0),
        point <DT> dir = point(1, 0))
        : 0(0), dir(dir) {}
   bool operator() (point <DT> p, point <DT</pre>
       return make_tuple(half(p), 0) <</pre>
            make_tuple(half(q), (p & q));
}; // given a pivot point and an initial
    direction, sorts by Angle with the
    given direction
```

Pair of Intersecting segments using Line Sweep

Checking for the intersection of two

```
segments is carried out by the
                                               intersect () function, using an
                                               algorithm based on the oriented area of
                                                the triangle.
template <typename X> point(point <X> p) | The queue of segments is the global variable
                                                s, a set<event>. Iterators that
                                               specify the position of each segment in
                                                the queue (for convenient removal of
                                               segments from the queue) are stored in
                                               the global array where.
    return *this * point(rhs.x, - rhs.y | Two auxiliary functions prev() and next()
                                               are also introduced, which return
                                               iterators to the previous and next
                                               elements (or end(), if one does not
                                          set<seg> s;
                                          vector<set<seg>::iterator> where;
                                          set<seg>::iterator prev(set<seg>::iterator
                                              it) {
                                              return it == s.begin() ? s.end() : --it;
                                          set<seg>::iterator next(set<seg>::iterator
                                              it) {
                                              return ++it;
                                          pair<int, int> solve(const vector<seg>& a) { | Point projectPointLine(Point p, Line 1) {
                                              int n = (int)a.size();
                                              vector<event> e;
                                              for (int i = 0; i < n; ++i) {</pre>
                                                 e.push_back(event(min(a[i].p.x, a[i
                                                      ].q.x), +1, i));
                                                 e.push_back(event(max(a[i].p.x, a[i
                                                      ].q.x), -1, i));
                                              sort(e.begin(), e.end());
                                              s.clear();
                                              where.resize(a.size());
                                              for (size_t i = 0; i < e.size(); ++i) {</pre>
                                                 int id = e[i].id;
                                                 if (e[i].tp == +1) {
                                                     set<seg>::iterator nxt = s.
                                                          lower_bound(a[id]), prv =
                                                          prev(nxt);
                                                     if (nxt != s.end() && intersect(*
                                                         nxt, a[id]))
                                                         return make_pair(nxt->id, id)
                                                     if (prv != s.end() && intersect(*
                                                         prv, a[id]))
                                                         return make_pair(prv->id, id)
                                                     where[id] = s.insert(nxt, a[id]);
                                                 } else {
```

```
set<seg>::iterator nxt = next(
            where[id]), prv = prev(where
            [id]);
       if (nxt != s.end() && prv != s.
            end() && intersect(*nxt, *
            prv))
           return make_pair(prv->id, nxt
               ->id):
       s.erase(where[id]);
return make_pair(-1, -1);
```

Some More 2D Geo 5.3

```
Tf distancePointLine(Point p, Line 1) {
   return abs(cross(l.b - l.a, p - l.a) /
        length(1.b - 1.a));
// returns the shortest distance from point
    a to segment s
Tf distancePointSegment(Point p, Segment s)
   if (s.a == s.b) return length(p - s.a);
   Point v1 = s.b - s.a, v2 = p - s.a, v3 =
         p - s.b;
   if (dcmp(dot(v1, v2)) < 0)
       return length(v2);
   else if (dcmp(dot(v1, v3)) > 0)
       return length(v3);
   else
       return abs(cross(v1, v2) / length(v1
// returns the shortest distance from
    segment p to segment q
Tf distanceSegmentSegment(Segment p, Segment
     q) {
   if (segmentsIntersect(p, q)) return 0;
   Tf ans = distancePointSegment(p.a, q);
   ans = min(ans, distancePointSegment(p.b,
         q));
   ans = min(ans, distancePointSegment(q.a,
         p));
   ans = min(ans, distancePointSegment(q.b,
         p));
   return ans;
// returns the projection of point p on line
   static_assert(is_same<Tf, Ti>::value);
   Point v = 1.b - 1.a;
   return 1.a + v * ((Tf)dot(v, p - 1.a) /
        dot(v, v));
// returns the left side of polygon u after
    cutting it by ray a->b
Polygon cutPolygon(Polygon u, Point a, Point
   using Linear::lineLineIntersection,
        Linear::onSegment;
   Polygon ret;
   int n = u.size();
   for (int i = 0; i < n; i++) {</pre>
       Point c = u[i], d = u[(i + 1) \% n];
       if (dcmp(cross(b - a, c - a)) >= 0)
           ret.push_back(c);
       if (dcmp(cross(b - a, d - c)) != 0)
           {
          Point t;
           lineLineIntersection(a, b - a, c,
                d - c, t):
           if (onSegment(t, Segment(c, d)))
               ret.push_back(t);
       }
   }
   return ret;
```

```
// returns false if points are collinear,
    true otherwise
// circle p touch each arm of triangle abc
bool inscribedCircle(Point a, Point b, Point
     c, Circle &p) {
   using Linear::distancePointLine;
   static_assert(is_same<Tf, Ti>::value);
   if (orient(a, b, c) == 0) return false;
   Tf u = length(b - c);
   Tf v = length(c - a);
   Tf w = length(a - b);
   p.o = (a * u + b * v + c * w) / (u + v + c * w)
        w);
   p.r = distancePointLine(p.o, Line(a, b))
   return true;
// set of points A(x, y) such that PA : QA =
     rp : rq
Circle apolloniusCircle(Point P, Point Q, Tf
     rp, Tf rq) {
   static_assert(is_same<Tf, Ti>::value);
   rq *= rq;
   rp *= rp;
   Tf a = rq - rp;
   assert(dcmp(a));
   Tf g = (rq * P.x - rp * Q.x) / a, h = (
        rq * P.y - rp * Q.y) / a;
   Tf c =
       (rq * P.x * P.x - rp * Q.x * Q.x +
            rq * P.y * P.y - rp * Q.y * Q.y
             / a:
   Point o(g, h);
   Tf R = sqrt(g * g + h * h - c);
   return Circle(o, R);
// returns false if points are collinear
bool circumscribedCircle(Point a, Point b,
    Point c, Circle &p) {
   using Linear::lineLineIntersection;
   if (orient(a, b, c) == 0) return false;
   Point d = (a + b) / 2, e = (a + c) / 2;
   Point vd = rotate90(b - a), ve =
        rotate90(a - c);
   bool f = lineLineIntersection(d, vd, e,
        ve, p.o);
   if (f) p.r = length(a - p.o);
   return f;
// Following three methods implement Welzl's
     algorithm for
// the smallest enclosing circle problem:
    Given a set of
// points, find out the minimal circle that
    covers them all.
// boundary(p) determines (if possible) a
    circle that goes
// through the points in p. Ideally |p| <=
    3.
// welzl() is a recursive helper function
    doing the most part
// of Welzl's algorithm. Call minidisk with
    the set of points
// Randomized Complexity: O(CN) with C^{-}10 (
    practically lesser)
// TESTED: ICPC Dhaka 2019 - I (CodeMarshal)
Circle boundary(const vector<Point> &p) {
   Circle ret;
   int sz = p.size();
   if (sz == 0)
       return Circle({0, 0}, 0);
   else if (sz == 1)
   else if (sz == 2)
       ret.o = (p[0] + p[1]) / 2, ret.r =
           length(p[0] - p[1]) / 2;
   else if (!circumscribedCircle(p[0], p
        [1], p[2], ret))
       ret.r = 0:
   return ret;
```

```
Circle welzl(const vector<Point> &p, int fr,
      vector<Point> &b) {
    if (fr >= (int)p.size() || b.size() ==
        3) return boundary(b);
    Circle c = welzl(p, fr + 1, b);
    if (!c.contains(p[fr])) {
       b.push_back(p[fr]);
       c = welzl(p, fr + 1, b);
       b.pop_back();
    return c;
Circle minidisk(vector<Point> p) {
    random_shuffle(p.begin(), p.end());
    vector<Point> q;
    return welzl(p, 0, q);
       Closest Pair of Points
```

```
LL ClosestPair(vector<pii> pts) {
   int n = pts.size();
    sort(all(pts));
   set<pii> s;
   LL best_dist = 1e18;
   int j = 0;
   for (int i = 0; i < n; ++i) {</pre>
       int d = ceil(sqrt(best_dist));
       while (pts[i].ff - pts[j].ff >=
            best_dist) {
           s.erase({pts[j].ss, pts[j].ff});
           j += 1;
       auto it1 = s.lower_bound({pts[i].ss
             - d, pts[i].ff});
       auto it2 = s.upper_bound({pts[i].ss
            + d, pts[i].ff});
       for (auto it = it1; it != it2; ++it)
           int dx = pts[i].ff - it->ss;
int dy = pts[i].ss - it->ff;
           best_dist = min(best_dist, 1LL *
                dx * dx + 1LL * dy * dy);
       s.insert({pts[i].ss, pts[i].ff});
   return best_dist;
```

3D Geo Templates 5.5

```
point get_perp(point p){ // returns a random
     perpendicular line to the vector p
   assert(sgn(norm(p)));
   point ret = point(-p.y, p.x, 0);
   if(sgn(norm(ret))) return ret;
   ret = point(0, -p.z, p.y);
   if(sgn(norm(ret))) return ret;
   assert(false)
struct plane{ // Caution: directed plane,
    directed on the direction of (p2 \times p3)
point n; // {a, b, c}
double d; //ax + by + cz = d
 // d = n . p [ where p is any point on the
     plane ]
plane(){;}
plane(point _n, double _d){
       n = _n;
       d = _d;
plane(point p1, point p2, point p3){
 n = crsp(p2 - p1, p3 - p1);
```

```
if(norm(n) < eps) {assert(false);} //</pre>
      doesn't define a plance
 d = dotp(p1, n);
   //Preserves the direction
 point get_p1(){ return univ(n) * d / norm(n
     );}
 point get_p2(){ return get_p1() + get_perp(
     n);}
 point get_p3(){ return crsp(n, get_p2() -
     get_p1()) + get_p1();}
 int get_side(point p){ return sgn(dotp(n, p
     ) - d);} ///OK
 double sgn_dist(point p) {return (dotp(n, p
     ) - d) / norm(n);}
 double dist(point p) {return fabs(sgn_dist(
     p));}
 point project(point p){ return p - sgn_dist
     (p) * univ(n);}
 point reflect(point p){ return p - 2 *
     sgn_dist(p) * univ(n);}
   point get_coords(point p){ // "2-D"-fies
         the plane. All points on this
        plane have z = 0
       point 0 = get_p1();
       point ox = univ(get_p2() - 0);
       point oy = univ(get_p3() - 0);
       point oz = univ(n);
       p = p - 0;
       return {dotp(p, ox), dotp(p, oy),
            dotp(p, oz));
};
plane translate(plane p, point t) {return {p
     .n, p.d + dotp(p.n, t));}
plane shiftUp(plane p, double d) {return {p.
    n, p.d + d * norm(p.n)};}
point projection(point p, point st, point ed
    ) { return dotp(ed - st, p - st) / norm
     (ed - st) * univ(ed - st) + st;} //OK
point extend(point st, point ed, double len)
      { return ed + univ(ed-st) * len;} //OK
point rtt(point axis, point p, double theta)
   axis = univ(axis);
   return p * cos(theta) + sin(theta) *
        crsp(axis, p) + axis * (1-cos(theta))
        )) * dotp(axis, p);
point segmentProjection(point p, point st,
    point ed)
   double d = dotp(p - st, ed - st) / norm(
       ed - st);
   if(d < 0) return st;</pre>
   if(d > norm(ed - st) + eps) return ed;
   return st + univ(ed - st) * d;
} //OK
double distPointSegment(point p, point st,
    point ed) {return norm(p
     segmentProjection(p, st, ed)); } //OK
double distPointLine( point P, point st,
    point ed) { return norm( projection(P,
     st, ed) - P ); } //OK
double pointPlanedist(plane P, point q){
    return fabs(dotp(P.n, q) - P.d) / norm(
double pointPlanedist(point p1, point p2,
    point p3, point q){ return
    pointPlanedist(plane(p1,p2,p3), q); }
```

```
point reflection(point p, point st, point ed double dist(line3d 11, line3d 12) {
   point proj = projection(p, st, ed);
   if(p != proj) return extend(p, proj,
        norm(p - proj));
   return proj;
} //OK
bool coplanar(point p1, point p2, point p3,
    point q)
   p2 = p2-p1, p3 = p3-p1, q = q-p1;
   if( fabs( dotp(q, crsp(p2, p3)) ) < eps</pre>
        ) return true;
   return false:
int linePlaneIntersection(point u, point v,
    point 1, point m, point r, point &x){
       -> 1, m, r defines the plane
       -> u, v defines the line
       -> returns 0 when does not intersect
       -> returns 1 when there exists one
            unique common point
       -> returns -1 when there exists
            infinite number of common point
   assert(1 != m && m != r && 1 != r && u
   if(coplanar(1, m, r, u) && coplanar(1, m
        , r, v)) return -1;
   1 = 1 - m;
   r = r - m;
   u = u - m;
   v = v - m;
   point C = crsp(l, r);
   double denom = dotp(v - u, C);
   if(fabs(denom) < eps) return 0;</pre>
   double alpha = -dotp(C, u) / denom;
   x = u + (v - u) * alpha + m;
   return 1:
double angle(point u, point v) { return acos
     ( max(-1.0, min(1.0, dotp(u, v) / (norm
     (u) * norm(v)))));}
struct line3d{ //directed
   point d, o; // dir = direction, o =
        online point
   line3d(point p, point q){
       d = q - p;
       o = p;
       assert(sgn(norm(d)));
   line3d(plane p1, plane p2){
       d = crsp(p1.n, p2.n);
       o = (crsp(p2.n*p1.d - p1.n*p2.d, d))
   point get_p1(){return o;}
   point get_p2(){return o + d;}
   double dist(point p){ return norm(crsp(d
        , p - o)) / norm(d);};
   point project(point p){ return
        projection(p, o, o + d); }
   point reflect(point p) {return
        reflection(p, o, o + d);}
}:
line3d perpThrough(plane p, point o){return
    line3d(o, o + p.n);}
plane perpThrough(line3d 1, point o){return
    plane(1.d, dotp(1.d, o));}
```

```
point n = crsp(11.d, 12.d);
   if (!sgn(norm(n))) return l1.dist(l2.o);
   return abs(dotp(12.o-11.o, n))/norm(n);
point closestOnL1(line3d 11, line3d 12) {
   point n2 = crsp(12.d, crsp(11.d, 12.d));
   return 11.0 + (11.d * (dotp(12.o-11.o,
        n2))) / dotp(11.d,n2);
double angle(plane p1, plane p2){return
    angle(p1.n, p2.n);}
bool isparallel(plane p1, plane p2){return !
    sgn(norm(crsp(p1.n, p2.n)));}
bool isperp(plane p1, plane p2) {return !sgn
     (dotp(p1.n, p2.n));}
double angle(line3d 11, line3d 12){return
    angle(11.d, 12.d);}
bool isparallel(line3d 11, line3d 12){return
     !sgn(norm(crsp(11.d, 12.d)));}
bool isperp(line3d 11, line3d 12) {return !
    sgn(dotp(11.d, 12.d));}
double angle(plane p, line3d 1) {return pi/2
     - angle(p.n, 1.d);}
bool isParallel(plane p, line3d l) {return !
    sgn(dotp(p.n, 1.d));}
bool isPerpendicular(plane p, line3d l) {
    return !sgn(norm(crsp(p.n, 1.d)));}
point vector_area2(vector <point> &poly){
   point S = \{0, 0, 0\};
   for(int i = 0; i < (int) poly.size(); i</pre>
       S = S + crsp(poly[i], poly[(i + 1)
           % poly.size()]);
   return S:
double area(vector < point > &poly){ // All
    points must be co-planer
   return norm(vector_area2(poly)) * 0.5;
   Polyhedrons
bool operator <(point p, point q) { ///OK</pre>
   return tie(p.x, p.y, p.z) < tie(q.x, q.y
        , q.z);
struct edge {
   bool same; // = is the common edge in
        the same order?
// Given a series of faces (lists of points) int main()
    , reverse some of them
// so that their orientations are consistent
     [ every face then will point in the
     same direction, inside / outside ]
void reorient(vector< vector<point> > &fs) { |//
   int n = fs.size();
   // Find the common edges and create the
        resulting graph
   vector< vector<edge> > g(n);
   map<pair<point,point>, int> es;
   for (int u = 0; u < n; u++) {</pre>
       for (int i = 0, m = fs[u].size(); i
            < m; i++) {
           point a = fs[u][i], b = fs[u][(i
               +1)%m];
           // Let look at edge [AB]
           if (es.count({a,b})) { // seen in}
                same order
```

 $int v = es[{a,b}];$

```
g[u].push_back({v,true});
              g[v].push_back({u,true});
           else if (es.count({b,a})) { //
                seen in different order
                  int v = es[{b,a}];
                  g[u].push_back({v,false});
                  g[v].push_back({u,false});
           else es[{a,b}] = u;
       }
   }
   vector<bool> vis(n,false), flip(n);
   flip[0] = false;
   queue<int> q;
   q.push(0);
   while (!q.empty()) {
       int u = q.front();
       q.pop();
       for (edge e : g[u]) {
          if (!vis[e.v]) {
              vis[e.v] = true;
              // If the edge was in the
                   same order.
              // exactly one of the two
                   should be flipped
              flip[e.v] = (flip[u] ^ e.same
                   );
              q.push(e.v);
       }
   }
   for (int u = 0; u < n; u++)</pre>
       if (flip[u])
          reverse(fs[u].begin(), fs[u].end
double volume(vector< vector<point> > fs) {
   double vol6 = 0.0;
   for (vector<point> f : fs)
       vol6 += dotp(vector_area2(f), f[0]);
   return abs(vol6) / 6.0;
   Spherical Co-ordinate System
point sph(double r, double lat, double lon)
    { // lat, lon in degrees
   lat *= pi/180, lon *= pi/180;
   return {r*cos(lat)*cos(lon), r*cos(lat)*
        sin(lon), r*sin(lat)};
double greatCircleDist(point o, double r,
    point a, point b) {
   return r * angle(a-o, b-o);
   plane p = \{point(0, 0, 10), point(0, 1,
        10), point(1, 0, 10)};
     plane q = \{p.get_p1(), p.get_p2(), p.
    get_p3()};
   double d = dotp(p.get_p2() - p.get_p1(),
         p.get_p3() - p.get_p1());
   D(eq(d, 0))
   return 0;
```

5.6 Point in Polygon

```
template <typename DT> DT
    FarthestPairOfPoints(polygon <DT> p){
   p = ConvexHull(p);
```

```
int n = p.size();
   DT ans = -1e9;
   for(int i = 0, j = 1; i < n; i++)</pre>
       for( ; UTriArea(p[i], p[(i + 1) % n
            ], p[(j + 1) % n]) > UTriArea(p[
            i], p[(i + 1) % n], p[j]) ; j =
            (j + 1) % n );
       ans = max(ans, DisSq(p[i], p[j]));
       ans = max(ans, DisSq(p[(i + 1) % n],
             p[j]));
   return ans; // will return square of the
         answer.
template <typename DT> int
    PointInConvexPolygon(polygon <int> ::
    iterator b, polygon <int> :: iterator e
     , const point <DT> &0){
   polygon <int> :: iterator lo = b + 2, hi
         = e - 1, ans = e;
   while(lo <= hi) {</pre>
       auto mid = lo + (hi - lo) / 2;
       if(TriArea(*b, 0, *mid) >= 0) ans =
            mid, hi = mid - 1;
       else lo = mid + 1;
   if (ans == e or abs(UTriArea(*b, *(ans -
        1), *ans) - UTriArea(*b, *(ans -
        1), 0) - UTriArea(*b, *ans, 0) -
        UTriArea(*(ans - 1), *ans, 0)) >
       return 0;
   else return (Collinear(*b, *(b + 1), 0)
        or Collinear(*(e - 1), *b, 0) or
        Collinear(*(ans), *(ans - 1), 0)) ?
         -1 : 1;
} // 0 for outside, -1 for on border, 1 for
template <typename DT> int PointInPolygon(
    polygon <DT> &P, point <DT> pt) {
   int n = P.size();
   int cnt = 0;
   for(int i = 0, j = 1; i < n; i++, j = (j</pre>
         + 1) % n) {
       if(TriArea(pt, P[i], P[j]) == 0 and
            min(P[i], P[j]) <= pt and pt <=
            max(P[i], P[j]))
          return -1;
       cnt += ((P[j].y \ge pt.y) - (P[i].y
            >= pt.y)) * TriArea(pt, P[i], P[
            i]) > 0;
   return cnt & 1;
// returns 1e9 if the point is on the
    polygon
int winding_number(vector<PT> &p, const PT&
    z) { // O(n)}
   if (is_point_on_polygon(p, z)) return 1
   int n = p.size(), ans = 0;
   for (int i = 0; i < n; ++i) {</pre>
       int j = (i + 1) \% n;
       bool below = p[i].y < z.y;
       if (below != (p[j].y < z.y)) {</pre>
           auto orient = orientation(z, p[j
               ], p[i]);
           if (orient == 0) return 0;
           if (below == (orient > 0)) ans +=
                below ? 1 : -1;
   }
   return ans;
// -1 if strictly inside, 0 if on the
    polygon, 1 if strictly outside
int is_point_in_polygon(vector<PT> &p, const
     PT& z) { // O(n)
    int k = winding_number(p, z);
   return k == 1e9 ? 0 : k == 0 ? 1 : -1;
```

5.7 Rotating Calipers

```
template <typename DT> polygon <DT>
    ConvexHull(polygon <DT> &PT){
   sort(PT.begin(), PT.end());
   int m = 0, n = PT.size();
   polygon \langle DT \rangle hull(n + n + 2);
   for(int i = 0; i < n; i++){</pre>
       for( ; m > 1 and TriArea(hull[m-2],
           hull[m-1], PT[i]) <= 0; m-- );
       hull[m++] = PT[i];
   for(int i = n - 2, k = m; i \ge 0; i--){
       for( ; m > k and TriArea(hull[m -
            2], hull[m - 1], PT[i]) <= 0; m
       hull[m++] = PT[i];
   if(n > 1)
   while(hull.size() > m)
       hull.pop_back();
   return hull;
template <typename DT> double
    MinimumBoundingBox(polygon <DT> P){
   auto p = ConvexHull(P);
   int n = p.size();
   double area = 1e20 + 5;
   for(int i = 0, 1 = 1, r = 1, u = 1; i <
         n ; i++){
       point <DT> edge = (p[(i+1)%n]-p[i])
            /sqrt(DisSq(p[i], p[(i+1)%n]));
       for( ; (edge p[r_n]-p[i]) < (edge
              p[(r+1)%n] - p[i]); r++);
       for( ; u<r || (edge & p[u%n] - p[i])</pre>
             < (edge & p[(u+1)%n] - p[i]); u
            ++);
       for( ; 1<u || (edge ^ p[1%n] - p[i])</pre>
            > (edge ^ p[(1+1)%n] - p[i]); 1
       double w = (edge ^p[r_n]-p[i]) - (
            edge ^ p[l%n] - p[i]);
       double h = UTriArea(p[u%n], p[i], p
            [(i+1)%n])/sqrt(DisSq(p[i], p[(i
            +1)%n]));
       area = min(area,w*h);
   if(area>1e19)
       area = 0;
   return area;
```

5.8 Line and Circles

```
template <typename DT>
class line{
   public:
       point <DT> dir, 0; // direction of
           vector and starting point
   line(point <DT> p,point <DT> q): dir(q -
         p), O(p) {};
   bool Contains(const point <double> &p){
       return fabs(p - 0 & dir ) < EPS;</pre>
   } // checks whether the line Contains {\tt a}
        certain point
   template <typename XT> point <XT> At(XT
       return point <XT> (dir) * t + 0;
   } // inserts value of t in the vector
        representation, finds the point
        which is 0 + Dir*t
   double AtInv(const point <double> &p){
       return abs(dir.x) > 0 ? (p - 0).x /
            dir.x : (p - 0).y / dir.y;
   } // if the line Contains a point, gives
         the value t such that, p = 0+Dir*t
   line Perp(point <DT> p){
       return line(p, p + (-dir.y,dir.x));
```

```
point <DT> ProjOfPoint(const point <DT>
        &P) {
       return 0 + dir * ((P - 0) ^ dir) /
            (~dir):
   double DisOfPoint(const point <DT> &P) {
       return fabs(dir & (P - 0))/sqrt(~(
   friend bool Parallel(line& L, line& R){
       return fabs(R.dir & L.dir) < EPS;</pre>
   friend int Intersects(line& L, line& R){
       return Parallel(L, R) ? R.Contains(L
            .0) ? -1 : 0 : 1;
   friend pair <double, double>
        IntersectionAt(line &L, line &R){
       double r = double((L.0 - R.0) & L.
            dir)/(R.dir & L.dir);
       double 1 = double((R.0 - L.0) & R.
            dir)/(L.dir & R.dir);
       return {1, r};
   friend pair <int, point<double>>
        IntersectionPoint(line L, line R,
        int _L = 0, int _R = 0){
       // _L and _R can be 0 to 3, 0 is a
            normal line, 3 is a segment, 1
            and 2 are rays (considered
       int ok = Intersects(L, R);
       if(ok == 0)
           return {0, {0, 0}};
       if(ok == 1){
           auto [1,r] = IntersectionAt(L, R)
           if(1 < (0-EPS) and _L & 2)
              return {0, {0, 0}};
           if(1 > (1+EPS) and _L & 1)
              return {0, {0, 0}};
           if(r < (0-EPS) and _R & 2)
               return {0, {0, 0}};
           if(r > (1+EPS) and _R & 1)
              return {0, {0, 0}};
           return {1, L.At(1)};
       return {-1, {0,0}}; // they are the
            same line
template <typename DT>
class circle {
   public:
       point <DT> 0;
       DT R:
   circle(const point <DT> &O = \{0, 0\}, DT
        R = 0) : O(0), R(R) {}
   // the next two make sense only on
        circle <double>
   circle(const point <DT> &A, const point
        <DT> &B, const point <DT> &C){
       point \langle DT \rangle X = (A + B) / 2, Y = (B +
            C) / 2, d1 = Perp(A - B), d2 =
            Perp(B - C);
       0 = IntersectionPoint(line(X, d1),
           line(Y, d2)).second;
       R = \operatorname{sqrt}(^{\sim}(0 - A));
   circle(const point <DT> &A, const point
        <DT> &B, DT R){
       point \langle DT \rangle X = (A + B) / 2, d = Perp
           (A - B);
       d = d * (R / sqrt(~(d)));
       0 = X + d;
       R = \operatorname{sqrt}(^{\sim}(0 - A));
   double SectorArea(double ang) {
       // Area of a sector of cicle
       return ang* R * R * .5;
   double SectorArea(const point <DT> &a,
        const point <DT> &b) {
```

```
return SectorArea(Angle(a - 0, b - 0
            ));
   7
   double ChordArea(const point <DT> &a,
        const point <DT> &b) {
       // Area between sector and its chord
       return SectorArea(a, b) - 0.5 *
            TriArea(0, a, b);
   int Contains(const point <DT> &p){
       // 0 for outside, 1 for inside, -1
            for on the circle
       DT d = DisSq(0, p);
       return d > R * R ? O : (d == R * R ?
             -1 : 1);
   friend tuple <int, point <DT>, point <DT</pre>
        >> IntersectionPoint(const circle &
        a,const circle &b) {
       if(a.R == b.R \text{ and } a.0 == b.0)
          return {-1, {0, 0}, {0, 0}};
       double d = sqrt(DisSq(a.0, b.0));
       if(d > a.R + b.R or d < fabs(a.R - b)
            .R.))
           return {0, {0, 0}, {0, 0}};
       double z = (sq(a.R) + sq(d) - sq(b.R)
           )) / (2 * d);
       double y = sqrt(sq(a.R) - sq(z));
       point \langle DT \rangle 0 = b.0 - a.0, h = Perp(0
            ) * (y / sqrt(~0));
       0 = a.0 + 0 * (z / sqrt(~0));
       return make_tuple(1 + (~(h) > EPS),
            0 - h, 0 + h);
   friend tuple <int, point <DT>, point <DT</pre>
        >> IntersectionPoint(const circle &
        C, line <DT> L) {
       point <DT> P = L.ProjOfPoint(C.0);
       double D = DisSq(C.O, P);
       if(D > C.R * C.R)
           return {0, {0, 0}, {0, 0}};
       double x = sqrt(C.R * C.R - D);
       point <DT> h = L.dir * (x / sqrt(~L.
            dir));
       return \{1 + (x > EPS), P - h, P + h\}
            };
   double SegmentedArea(point <DT> &a,
        point <DT> &b) {
       // signed area of the intersection
            between the circle and triangle
            NAB
       double ans = SectorArea(a, b);
       line <DT> L(a, b);
       auto [cnt, p1, p2] =
            IntersectionPoint(*this, L);
       if(cnt < 2)
          return ans;
       double t1 = L.AtInv(p1), t2 = L.
            AtInv(p2);
       if(t2 < 0 or t1 > 1)
           return ans:
       if(t1 < 0)
          p1 = a;
       if(t2 > 1)
          p2 = b;
       return ans - ChordArea(p1, p2);
   }
};
```

6 Math

Stirling Numbers 6.1

```
//stirling number 2nd kind variation(number
    of ways to place n marbles in k boxes
    so that each box has at least x marbles
11 solve(int marble, int box) {
 if (marble < 111 * box * x) return 0;</pre>
  if (box == 1 && marble >= x) return 1;
 if (vis[marble][box] == cs) return dp[
      marble][box];
 vis[marble][box] = cs;
```

```
11 a = ( 111 * box * solve(marble - 1, box double adaptive_simpson(double minx, double
      ) ) % MOD;
 11 b = ( 111 * box * ncr(marble - 1, x -
      1) ) % MOD;
 b = (b * solve(marble - x, box - 1)) % MOD
 ll ret = (a + b) \% MOD;
 return dp[marble][box] = ret;
//number of ways to place n marbles in k
    boxes so that no box is empty
11 stir(ll n, ll k) {
 ll ret = 0;
 for (int i = 0; i <= k; i++) {</pre>
   11 v = ncr(k, i) * bigmod(i, n) % MOD;
   if ( (k - i) % 2 == 0 ) ret = (ret + v)
        % MOD;
   else ret = (ret - v + MOD) % MOD;
 return ret:
```

FFT in Mod 6.2

```
// need to add modulo to res[i] in Mul
vector<ll> Mul_mod(vector<ll>& a, vector<ll</pre>
    >& b, 11 mod) {
 11 sqrt_mod = (11)sqrtl(mod);
 vector<11> a0(a.size()), a1(a.size());
 vector<ll> b0(b.size()), b1(b.size());
 for (int i = 0; i < a.size(); i++) {</pre>
   a0[i] = a[i] % sqrt_mod;
   a1[i] = a[i] / sqrt_mod;
 for (int i = 0; i < b.size(); i++) {</pre>
   b0[i] = b[i] % sqrt_mod;
   b1[i] = b[i] / sqrt_mod;
 vector<ll> a01(a.size()), b01(b.size());
 for (int i = 0; i < a.size(); i++) {</pre>
   a01[i] = a0[i] + a1[i];
   if (a01[i] >= mod) a01[i] -= mod;
 for (int i = 0; i < b.size(); i++) {</pre>
   b01[i] = b0[i] + b1[i];
   if (b01[i] >= mod) b01[i] -= mod;
 vector<ll> mid = Mul(a01, b01);
 vector<11> a0b0 = Mul(a0, b0);
  vector<ll> a1b1 = Mul(a1, b1);
 for (int i = 0; i < mid.size(); i++) {</pre>
   mid[i] = (mid[i] - a0b0[i] + mod) % mod;
   mid[i] = (mid[i] - a1b1[i] + mod) % mod;
 vector<11> res = a0b0;
 for (int i = 0; i < res.size(); i++) {</pre>
   res[i] += (sqrt_mod * mid[i]) % mod;
   if (res[i] >= mod) res[i] -= mod;
 sqrt_mod = (sqrt_mod * sqrt_mod) % mod;
 for (int i = 0; i < res.size(); i++) {</pre>
   res[i] += (sqrt_mod * a1b1[i]) % mod;
   if (res[i] >= mod) res[i] -= mod;
 return res;
```

Simpson Integration 6.3

```
For finding the length of an arc in a
        range
   L = integrate(ds) from start to end of
        range
   where ds = sqrt(1+(d/dy(x))^2)dy
const double SIMPSON_TERMINAL_EPS = 1e-12;
\ensuremath{///} Function whose integration is to be
     calculated
double F(double x):
double simpson(double minx, double maxx)
   return (maxx - minx) / 6 * (F(minx) + 4)
        * F((minx + maxx) / 2.) + F(maxx));
```

```
maxx, double c, double EPS)
     if(maxx - minx < SIMPSON_TERMINAL_EPS)</pre>
     return 0;
   double midx = (minx + maxx) / 2;
   double a = simpson(minx, midx);
   double b = simpson(midx, maxx);
   if(fabs(a + b - c) < 15 * EPS) return a
        + b + (a + b - c) / 15.0;
   return adaptive_simpson(minx, midx, a,
        EPS / 2.) + adaptive_simpson(midx,
        maxx, b, EPS / 2.);
double adaptive_simpson(double minx, double
    maxx, double EPS)
   return adaptive_simpson(minx, maxx,
        simpson(minx, maxx, i), EPS);
```

6.4Matrix Exponentiation

```
struct matrix {
 vector<vector<11>> mat;
  int n, m;
 matrix() {}
 matrix(int n, int m) : n(n), m(m), mat(n)
    for (int i = 0; i < n; i++) mat[i] =</pre>
         vector<ll>(m);
  void identity() { for (int i = 0; i < n; i</pre>
      ++) mat[i][i] = 1; }
  void print() {
   for (int i = 0; i < n; i++) {</pre>
     for (int j = 0; j < n; j++) cout <<
    mat[i][j] << " ";</pre>
     cout << "\n";
   }
 vector<ll> &operator[](int i) {
   return mat[i];
// make sure a.m == b.n
matrix operator * (matrix &a, matrix &b) {
 int n = a.n, m = b.m;
 matrix ret(n, m);
 for (int i = 0; i < n; i++) {</pre>
   for (int j = 0; j < m; j++) {
     for (int k = 0; k < a.m; k++) {</pre>
       ll val = (1ll * a[i][k] * b[k][j]) %
             MOD;
       ret[i][j] = (ret[i][j] + val) % MOD;
 }
 return ret;
matrix mat_exp(matrix &mat, ll p) {
 int n = mat.n, m = mat.m;
 matrix ret(n, m);
 ret.identity();
 matrix x = mat;
  while (p) {
   if (p & 1) ret = ret * x;
   x = x * x;
   p = p >> 1;
 return ret;
```

Fast Fourier Transforma-6.5tion

```
typedef complex<double> base;
void fft(vector<base> & a, bool invert) {
    int n = (int)a.size();
   for (int i = 1, j = 0; i<n; ++i) {</pre>
       int bit = n >> 1;
       for (; j >= bit; bit >>= 1)j -= bit;
       j += bit;
       if (i < j)swap(a[i], a[j]);</pre>
   for (int len = 2; len <= n; len <<= 1) {</pre>
       double ang = 2 * PI / len * (invert
            ? -1 : 1);
       base wlen(cos(ang), sin(ang));
       for (int i = 0; i<n; i += len) {</pre>
           base w(1);
           for (int j = 0; j<len / 2; ++j) {</pre>
               base u = a[i + j], v = a[i +
                   j + len / 2] * w;
               a[i + j] = u + v;
               a[i + j + len / 2] = u - v;
               w *= wlen;
           }
       }
   }
   if (invert)for (int i = 0; i<n; ++i)a[i]</pre>
          /= n;
vector<LL> Mul(vector<LL>& a, vector<LL>& b)
    vector < base > fa(a.begin(), a.end()), fb(
        b.begin(), b.end());
    int n = 1;
   while (n < max(a.size(), b.size())) n</pre>
        <<= 1:
   fa.resize(n), fb.resize(n);
   fft(fa, false), fft(fb, false);
   for (int i = 0; i<n; ++i)fa[i] *= fb[i];</pre>
   fft(fa, true);
   vector<LL> res;
   res.resize(n):
   for (int i = 0; i<n; ++i)res[i] = round(</pre>
        fa[i].real());
   return res;
}
n degree Polynomial division: (one of the
     vector)
coefficient of x^i is replaced with
     coefficient of x^(n-i).
x^{(i-j)} starts from n+1 ends at 2*n-1
All possible sum of 3 different index.
vector<ll>v=Mul(v1,v2);
v=Mul(v,v3);
vector<ll>dbl(v.size()),tri(v.size());
for(ll i=0;i<v1.size();i++){</pre>
    dbl[i+i]=v1[i]*v2[i]; // All (i,i) Pairs
    tri[i+i+i]=v1[i]*v2[i]*v3[i]; // All (i,
        i,i) Triplets
dbl=Mul(dbl,v3); // All (i,i,j) Triplets
for(ll i=0;i<v.size();i++){</pre>
    v[i] = v[i] - (3 * dbl[i] - 2 * tri[i]);
         // 3 (i,i,j) Triplets have 3 (i,i,
        i) triplets. So Remove 2 (i,i,i)
         triplets.
   v[i]/=6; // (i,j,k) can be oriented in
         3! way
    if(v[i])cout << i - 60000 << " : " << v[
        i] << "\n"; // Handle negative
         value
}
*/
int main()
   int t;
```

```
cin >> t;
while(t--){
    string a,b;
    cin >> a >> b;
    vector<LL>v1,v2;
    int sign = 0;
    if(a[0] == '-'){
        sign = 1 - sign;
        a.erase(a.begin());
    if(b[0] == '-'){
       sign = 1 - sign;
        b.erase(b.begin());
    for(int i = 0;i < a.size();i++){</pre>
        int d = a[i] - '0';
        v1.push_back(d);
    for(int i = 0;i < b.size();i++){</pre>
        int d = b[i] - '0';
        v2.push_back(d);
    reverse(all(v1)),reverse(all(v2));
         //Reverse needed if v1 is in x^n
         +x^n-1+....+x^1+1 form
    vector<LL>v = Mul(v1,v2);
    int carry = 0;
    vector<int>answer;
    for(int i = 0;i < v.size();i++){</pre>
        int temp = v[i];
        temp += carry;
        answer.push_back(temp % 10);
        carry = temp/10;
    while(answer.size() > 1 and answer.
        back() == 0)answer.pop_back();
    reverse(all(answer));
    for(int i : answer)cout << i;</pre>
    cout << "\n";
}
```

Linear Seive with Multi-6.6 plicative Functions

```
const int maxn = 1e7;
vector <int> primes;
int spf[maxn+5], phi[maxn+5], NOD[maxn+5],
    cnt[maxn+5], POW[maxn+5], SOD[maxn+5];
bool prime[maxn+5];
void sieve(){
   fill(prime+2, prime+maxn+1, 1);
   SOD[1] = NOD[1] = phi[1] = spf[1] = 1;
   for(11 i=2;i<=maxn;i++){</pre>
       if(prime[i]) {
          primes.push_back(i), spf[i] = i;
          phi[i] = i-1;
          NOD[i] = 2, cnt[i] = 1;
          SOD[i] = i+1, POW[i] = i;
       for(auto p:primes){
          if(p*i>maxn or p > spf[i]) break;
          prime[p*i] = false, spf[p*i] = p;
          if(i\%p == 0){
              phi[p*i]=p*phi[i];
              NOD[p*i]=NOD[i]/(cnt[i]+1)*(
                   cnt[i]+2), cnt[p*i]=cnt[
                   i]+1;
              SOD[p*i]=SOD[i]/SOD[POW[i]]*(
                   SOD[POW[i]]+p*POW[i]),
                   POW[p*i]=p*POW[i];
              break;
          } else {
              phi[p*i]=phi[p]*phi[i];
              NOD[p*i]=NOD[p]*NOD[i], cnt[p
```

*i]=1;

```
SOD[p*i]=SOD[p]*SOD[i], POW[p]
                *i]=p;
       }
   }
}
```

Gaussian Elimination 6.7

```
const double EPS = 1e-9;
const int INF = 2; // it doesn't actually
    have to be infinity or a big number
template <typename DT> int gauss (vector <
    vector<DT> > a, vector<DT> & ans) {
    int n = (int) a.size();
   int m = (int) a[0].size() - 1;
   vector<int> where (m, -1);
   for (int col=0, row=0; col<m && row<n;</pre>
        ++col) {
       int sel = row;
       for (int i=row; i<n; ++i)</pre>
           if (abs (a[i][col]) > abs (a[sel
                ][col]))
               sel = i;
       if (abs (a[sel][col]) < EPS)</pre>
           continue;
       for (int i=col; i<=m; ++i)</pre>
           swap (a[sel][i], a[row][i]);
       where[col] = row;
       for (int i=0; i<n; ++i)</pre>
           if (i != row) {
               DT c = a[i][col] / a[row][col]
               for (int j=col; j<=m; ++j)</pre>
                   a[i][j] -= a[row][j] * c;
           }
       ++row;
   ans.assign (m, 0);
   for (int i=0; i<m; ++i)</pre>
       if (where[i] != -1)
           ans[i] = a[where[i]][m] / a[where
                [i]][i];
   for (int i=0; i<n; ++i) {</pre>
       DT sum = 0;
       for (int j=0; j<m; ++j)</pre>
           sum += ans[j] * a[i][j];
       if (abs (sum - a[i][m]) > EPS)
           return 0;
   for (int i=0; i<m; ++i)</pre>
       if (where[i] == -1)
           return INF;
   return 1;
int compute rank(vector<vector<double>> A) {
   int n = A.size();
   int m = A[0].size();
   int rank = 0;
   vector<bool> row_selected(n, false);
   for (int i = 0; i < m; ++i) {</pre>
       int j;
       for (j = 0; j < n; ++j) {
           if (!row_selected[j] && abs(A[j][
                i]) > EPS)
               break:
       if (j != n) {
           row_selected[j] = true;
           for (int p = i + 1; p < m; ++p)</pre>
               A[j][p] /= A[j][i];
           for (int k = 0; k < n; ++k) {
               if (k != j && abs(A[k][i]) >
                    EPS) {
```

```
for (int p = i + 1; p < m;</pre>
                      ++p)
                    A[k][p] -= A[j][p] * A
                         [k][i];
            }
       }
}
return rank;
```

Seive Upto 1e9

```
// credit: min_25
// takes 0.5s for n = 1e9
vector<int> sieve(const int N, const int Q =
     17, const int L = 1 << 15) {
  static const int rs[] = {1, 7, 11, 13, 17,
       19, 23, 29};
 struct P {
   P(int p) : p(p) {}
   int p; int pos[8];
 }:
 auto approx_prime_count = [] (const int N)
       -> int {
   return N > 60184 ? N / (log(N) - 1.1)
                   : max(1., N / (log(N) -
                        1.11)) + 1;
 };
  const int v = sqrt(N), vv = sqrt(v);
  vector<bool> isp(v + 1, true);
 for (int i = 2; i <= vv; ++i) if (isp[i])</pre>
   for (int j = i * i; j <= v; j += i) isp[</pre>
        j] = false;
  const int rsize = approx_prime_count(N +
  vector<int> primes = {2, 3, 5}; int psize
      = 3:
 primes.resize(rsize);
 vector<P> sprimes; size_t pbeg = 0;
  int prod = 1;
 for (int p = 7; p <= v; ++p) {</pre>
   if (!isp[p]) continue;
   if (p <= Q) prod *= p, ++pbeg, primes[</pre>
        psize++] = p;
   auto pp = P(p);
   for (int t = 0; t < 8; ++t) {
     int j = (p \le Q) ? p : p * p;
     while (j % 30 != rs[t]) j += p << 1;</pre>
     pp.pos[t] = j / 30;
   sprimes.push_back(pp);
 vector<unsigned char> pre(prod, 0xFF);
 for (size_t pi = 0; pi < pbeg; ++pi) {</pre>
   auto pp = sprimes[pi]; const int p = pp.
        p;
   for (int t = 0; t < 8; ++t) {</pre>
     const unsigned char m = ~(1 << t);</pre>
     for (int i = pp.pos[t]; i < prod; i += }</pre>
           p) pre[i] &= m;
   }
 }
 const int block_size = (L + prod - 1) /
      prod * prod;
  vector<unsigned char> block(block_size);
      unsigned char* pblock = block.data();
  const int M = (N + 29) / 30;
 for (int beg = 0; beg < M; beg +=</pre>
      block_size, pblock -= block_size) {
   int end = min(M, beg + block_size);
   for (int i = beg; i < end; i += prod) {</pre>
     copy(pre.begin(), pre.end(), pblock +
          i):
   if (beg == 0) pblock[0] &= 0xFE;
```

```
for (size_t pi = pbeg; pi < sprimes.size</pre>
        (); ++pi) {
     auto& pp = sprimes[pi];
     const int p = pp.p;
     for (int t = 0; t < 8; ++t) {</pre>
       int i = pp.pos[t]; const unsigned
            char m = ~(1 << t);</pre>
       for (; i < end; i += p) pblock[i] &=</pre>
       pp.pos[t] = i;
   for (int i = beg; i < end; ++i) {</pre>
     for (int m = pblock[i]; m > 0; m &= m
          - 1) {
       primes[psize++] = i * 30 + rs[
            __builtin_ctz(m)];
 assert(psize <= rsize);
 while (psize > 0 && primes[psize - 1] > N)
        --psize;
 primes.resize(psize);
 return primes;
int32_t main() {
 ios_base::sync_with_stdio(0);
 cin.tie(0);
 int n, a, b; cin >> n >> a >> b;
 auto primes = sieve(n);
 vector<int> ans;
 for (int i = b; i < primes.size() &&</pre>
      primes[i] <= n; i += a) ans.push_back</pre>
      (primes[i]);
 cout << primes.size() << ' ' << ans.size()</pre>
       << '\n';
 for (auto x: ans) cout << x << ' '; cout</pre>
      << '\n';
 return 0;
```

Derangements

```
array <int, N + 1> Drng;
void init(){
   Drng[0] = 1, Drng[1] = 0;
   for(int i = 2; i <= N; i++)</pre>
       Drng[i] = (LL) (i - 1) * (Drng[i -
            1] + Drng[i - 2]) % mod;
int D(int n) {
   return n < 0 ? 0 : Drng[n];</pre>
```

6.10Pollard Rho and Factor-const int M = 1e9+7; ization

```
// fast factorize
map <ull,int> fast_factorize(ull n){
   map <ull,int> ans;
   for(;n>1;n/=spf[n])
       ans[spf[n]]++;
   return ans:
inline ULL mul(ULL a,ULL b,ULL mod){
   LL ans = a * b - mod * (ULL) (1.L / mod)
       * a * b);
   return ans + mod * (ans < 0) - mod * (</pre>
        ans >= (LL) mod);
inline ULL bigmod(ULL num,ULL pow,ULL mod){
   ULL ans = 1:
   for( ; pow > 0; pow >>= 1, num = mul(num
         num, mod))
       if(pow&1) ans = mul(ans,num,mod);
   return ans;
inline bool is_prime(ULL n){
   if(n < 2 or n % 6 % 4 != 1)</pre>
      return (n|1) == 3;
   ULL a[] = \{2, 325, 9375, 28178, 450775,
```

9780504, 1795265022};

```
ULL s = \_builtin_ctzll(n-1), d = n >> s
    for(ULL x: a){
       ULL p = bigmod(x % n, d, n), i = s;
       for( ; p != 1 and p != n-1 and x % n
             and i--; p = mul(p, p, n));
       if(p != n-1 and i != s)
           return false;
   }
   return true;
ULL get_factor(ULL n) {
   auto f = [&](LL x) { return mul(x, x, n)
         + 1; };
    ULL x = 0, y = 0, t = 0, prod = 2, i =
        2, q;
    for( ; t++ %40 or gcd(prod, n) == 1; x =
         f(x), y = f(f(y))  ){
       (x == y) ? x = i++, y = f(x) : 0;
       prod = (q = mul(prod, max(x,y) - min
            (x,y), n)) ? q : prod;
   return gcd(prod, n);
map <ULL, int> factorize(ULL n){
   map <ULL, int> res;
    if(n < 2) return res;</pre>
    ULL small_primes[] = {2, 3, 5, 7, 11,
        13, 17, 19, 23, 29, 31, 37, 41, 43,
         47, 53, 59, 61, 67, 71, 73, 79,
        83, 89, 97 };
    for (ULL p: small_primes)
       for( ; n % p == 0; n /= p, res[p]++)
    auto _factor = [&](ULL n, auto &_factor)
       if(n == 1) return;
       if(is_prime(n))
          res[n]++:
       else {
           ULL x = get_factor(n);
           _factor(x, _factor);
           _factor(n / x, _factor);
   };
    _factor(n, _factor);
    return res;
```

Sqrt Field 6.11

```
///Author: anachor
#include<bits/stdc++.h>
using namespace std;
typedef long long LL;
LL mod(LL x) {
   LL ans = x\%M;
   if (ans < 0) ans += M;</pre>
   return ans:
template <LL X>
struct SqrtField {
   LL a, b, c; /// (a + b*sqrt(X))/c;
   {\tt SqrtField(LL\ A=0,\ LL\ B=0,\ LL\ C=1)\ :\ a(A)}
        , b(B), c(C) {}
   SqrtField operator+(const SqrtField &y)
       return SqrtField(mod(a*y.c + y.a*c),
            mod(b*y.c + y.b*c), mod(c*y.c))
   SqrtField operator-(const SqrtField &y)
       return SqrtField(mod(a*y.c - y.a*c),
             mod(b*y.c - y.b*c), mod(c*y.c))
```

```
}
   SqrtField operator*(const SqrtField &y)
       return SqrtField(mod(a*y.a + X*y.b*b
            ));
   }
   SqrtField operator/(const SqrtField &y)
       LL A = mod(a*y.a - X*y.b*b);
       LL B = mod(b*y.a - a*y.b);
       LL C = mod(y.a*y.a - X*y.b*y.b);
       A = mod(A*y.c);
       B = mod(B*y.c);
       C = mod(C*c);
       return SqrtField(A,B,C);
template<LL X>
ostream& operator<<(ostream &os, const
    SqrtField<X> &x) {
   return os<<"("<<x.a<<"+"<<x.b<<"V"<<X<<" }
        )/"<<x.c:
int main() {
   SqrtField<2> c(5);
                             ///5
                             ///3 sqrt(2)
   SqrtField<2> b(0, 1);
                            ///(3+7*sqrt(2)
   SqrtField<2> a(3,7,2);
        )/2
   cout<<a+b<<" "<<a-b<<" "<<a*b<<" "<<a/c
   cout<<a*2<<" "<<a/2<<" "<<a+2<<" "<<a
        -1<<endl;
```

6.12 Prime Counting Function

```
// initialize once by calling init()
#define MAXN 20000010
                        // initial sieve
    limit
#define MAX_PRIMES 2000010 // max size of
    the prime array for sieve
#define PHI_N 100000
#define PHI_K 100
int len = 0; // total number of primes
    generated by sieve
int primes[MAX_PRIMES];
int pref[MAXN];
                     // pref[i] --> number
    of primes <= i
int dp[PHI_N][PHI_K]; // precal of yo(n,k)
bitset<MAXN> f;
void sieve(int n) {
   f[1] = true;
   for (int i = 4; i <= n; i += 2) f[i] =
        true;
   for (int i = 3; i * i <= n; i += 2) {
       if (!f[i]) {
           for (int j = i * i; j <= n; j +=</pre>
                i << 1) f[j] = 1;
       }
   for (int i = 1; i <= n; i++) {</pre>
       if (!f[i]) primes[len++] = i;
       pref[i] = len;
void init() {
   sieve(MAXN - 1);
   // precalculation of phi upto size (
        PHI_N,PHI_K)
   for (int n = 0; n < PHI_N; n++) dp[n][0]</pre>
         = n;
   for (int k = 1; k < PHI_K; k++) {</pre>
       for (int n = 0; n < PHI_N; n++) {</pre>
           dp[n][k] = dp[n][k - 1] - dp[n /
                primes[k - 1]][k - 1];
       }
```

```
// returns the number of integers less or
                                       equal n which are
                                   // not divisible by any of the first k
                                       primes
), mod(a*y.b + b*y.a), mod(c*y.c) // recurrence --> yo(n, k) = yo(n, k-1) - yo
                                       (n / p_k, k-1)
                                   // for sum of primes yo(n, k) = yo(n, k-1)
                                        p_k * yo(n / p_k , k-1)
                                  long long yo(long long n, int k) {
                                      if (n < PHI_N && k < PHI_K) return dp[n</pre>
                                           l[k]:
                                      if (k == 1) return ((++n) >> 1);
                                      if (primes[k - 1] >= n) return 1;
                                      return yo(n, k - 1) - yo(n / primes[k -
                                           1], k - 1);
                                  // complexity: n^(2/3).(\log n^(1/3))
                                  long long Legendre(long long n) {
                                      if (n < MAXN) return pref[n];</pre>
                                      int lim = sqrt(n) + 1;
                                      int k = upper_bound(primes, primes + len
                                          , lim) - primes;
                                      return yo(n, k) + (k - 1);
                                  // runs under 0.2s for n = 1e12
                                  long long Lehmer(long long n) {
                                      if (n < MAXN) return pref[n];</pre>
                                      long long w, res = 0;
                                      int b = sqrt(n), c = Lehmer(cbrt(n)), a
                                           = Lehmer(sqrt(b));
                                      b = Lehmer(b):
                                      res = yo(n, a) + ((1LL * (b + a - 2) * (
                                           b - a + 1)) >> 1);
                                      for (int i = a; i < b; i++) {
                                          w = n / primes[i];
                                          int lim = Lehmer(sqrt(w));
                                          res -= Lehmer(w);
                                          if (i <= c) {</pre>
                                              for (int j = i; j < lim; j++) {</pre>
                                                 res += j;
                                                 res -= Lehmer(w / primes[j]);
                                              }
                                      return res;
```

Walsh Hadamord Fast 6.13Transformation

```
#define bitwiseXOR 1
//#define bitwiseAND 2
//#define bitwiseOR 3
void FWHT(vector< LL >&p, bool inverse){
   LL n = p.size();
   assert((n&(n-1))==0);
   for (LL len = 1; 2*len <= n; len <<= 1)</pre>
       for (LL i = 0; i < n; i += len+len)</pre>
           for (LL j = 0; j < len; j++) {</pre>
               LL u = p[i+j];
               LL v = p[i+len+j];
               #ifdef bitwiseXOR
               p[i+j] = u+v;
               p[i+len+j] = u-v;
               #endif // bitwiseXOR
               #ifdef bitwiseAND
               if (!inverse) {
                  p[i+j] = v \% MOD;
                  p[i+len+j] = (u+v) % MOD;
               } else {
                  p[i+j] = (-u+v) \% MOD;
                  p[i+len+j] = u \% MOD;
               #endif // bitwiseAND
               #ifdef bitwiseOR
               if (!inverse) {
```

```
p[i+j] = u+v;
                   p[i+len+j] = u;
               } else {
                   p[i+j] = v;
                   p[i+len+j] = u-v;
               #endif // bitwiseOR
       }
    #ifdef bitwiseXOR
    if (inverse) {
        //LL val=BigMod(n,MOD-2); //Option
            2: Exclude
        for (LL i = 0; i < n; i++) {</pre>
           //assert(p[i]%n==0); //Option 2:
            //p[i] = (p[i]*val)%MOD; //Option
                  2: p[i]/=n;
           p[i]/=n;
       }
    #endif // bitwiseXOR
vector<pair<int,int> >V[100005];
int dis[100005];
void dfs(int s,int pr){
    for(auto p:V[s]){
        if(p.first==pr) continue;
        dis[p.first] = dis[s]^p.second;
        dfs(p.first,s);
int main(){
    int t:
    cin >> t:
    const int len=(1<<16);</pre>
    for(int tc=1;tc<=t;tc++){</pre>
       LL n;
        cin >> n;
        for(int i=1;i<=n-1;i++){</pre>
           int u,v,w;
            cin >> u >> v >> w;
            V[u].push_back({v,w});
            V[v].push_back({u,w});
        dfs(1,0);
        vector<LL>a(len,0);
        for(int i=1;i<=n;i++) a[dis[i]]++;</pre>
       FWHT(a,false);
        for(int i=0;i<len;i++) a[i]*=a[i];</pre>
        FWHT(a,true);
        a[0]-=n:
        cout << "Case " << tc << ":\n";</pre>
        for(int i=0;i<len;i++) cout << a[i</pre>
            ]/2 << '\n';
        memset(dis,0,sizeof(dis));
        for(int i=1;i<=n;i++) V[i].clear();</pre>
```

6.14 Green Hackenbush on Trie

```
int trie[40 * MAX][26];
int XOR[40 * MAX][26];
int valu[40 * MAX];
int node = 1;
int add(string s) {
 int now = 1;
  stack<int>st;
 for (int i = 0; i < s.size(); i++) {</pre>
   int c = s[i] - 'a';
   if (!trie[now][c]) trie[now][c] = ++node
   st.push(now);
   now = trie[now][c];
```

```
int nxt = now;
int nxt_val = 0;
for (int i = 0; i < 26; i++) nxt_val ^=</pre>
     XOR[now][i];
while (!st.empty()) {
 now = st.top();
  st.pop();
  int val = 0;
  for (int i = 0; i < 26; i++) {</pre>
   if (trie[now][i] == nxt) {
     XOR[now][i] = nxt_val + 1;
   val ^= XOR[now][i];
 nxt_val = val;
 nxt = now;
}
return nxt_val;
```

6.15 Berlekamp Massey

```
struct berlekamp_massey { // for linear
    recursion
  typedef long long LL;
 static const int SZ = 2e5 + 5;
 static const int MOD = 1e9 + 7; /// mod
      must be a prime
 LL m , a[SZ] , h[SZ] , t_[SZ] , s[SZ] , t[  
      SZ];
  // bigmod goes here
  inline vector <LL> BM( vector <LL> &x ) {
   LL lf , ld;
   vector <LL> ls , cur;
   for ( int i = 0; i < int(x.size()); ++i</pre>
        ) {
     LL t = 0;
     for ( int j = 0; j < int(cur.size());</pre>
          ++j) t = (t + x[i - j - 1] * cur
          [j]) % MOD;
     if ( (t - x[i]) \% MOD == 0 ) continue;
     if ( !cur.size() ) {
       cur.resize( i + 1 );
       lf = i; ld = (t - x[i]) % MOD;
       continue;
     LL k = -(x[i] - t) * bigmod(ld, MOD)
          - 2 , MOD ) \% MOD;
     vector <LL> c(i - lf - 1);
     c.push_back( k );
     for ( int j = 0; j < int(ls.size());</pre>
          ++j ) c.push_back(-ls[j] * k %
     if ( c.size() < cur.size() ) c.resize(</pre>
           cur.size() );
     for ( int j = 0; j < int(cur.size());</pre>
          ++j ) c[j] = (c[j] + cur[j]) %
          MOD;
     if (i - lf + (int)ls.size() >= (int)
          cur.size() ) ls = cur, lf = i, ld
           = (t - x[i]) \% MOD;
     cur = c;
   for ( int i = 0; i < int(cur.size()); ++</pre>
        i ) cur[i] = (cur[i] % MOD + MOD) %
         MOD:
   return cur;
 inline void mull( LL *p , LL *q ) {
   for ( int i = 0; i < m + m; ++i ) t_[i]</pre>
        = 0:
   for ( int i = 0; i < m; ++i ) if ( p[i]</pre>
       for ( int j = 0; j < m; ++j ) t_{i} = 0
             j] = (t_[i + j] + p[i] * q[j])
   for ( int i = m + m - 1; i >= m; --i )
        if ( t_[i] )
       for ( int j = m - 1; ~j; --j ) t_[i
            -j-1] = (t_[i - j - 1] + t_[i
            ] * h[j]) % MOD;
   for ( int i = 0; i < m; ++i ) p[i] = t_[</pre>
        i];
```

```
inline LL calc( LL K ) {
   for ( int i = m; ~i; --i ) s[i] = t[i] =
   s[0] = 1; if ( m != 1 ) t[1] = 1; else t
        [0] = h[0];
   while ( K ) {
     if ( K & 1 ) mull( s , t );
     mull( t , t ); K >>= 1;
   LL su = 0;
   for ( int i = 0; i < m; ++i ) su = (su +
         s[i] * a[i]) % MOD;
   return (su % MOD + MOD) % MOD;
 }
 \///\ already calculated upto k , now
      calculate upto n.
  inline vector <LL> process( vector <LL> &x |}
       , int n , int k ) {
   auto re = BM( x );
   x.resize(n+1);
   for ( int i = k + 1; i <= n; i++ ) {</pre>
     for ( int j = 0; j < re.size(); j++ )</pre>
       x[i] += 1LL * x[i - j - 1] % MOD *
            re[j] % MOD; x[i] %= MOD;
   }
   return x;
 inline LL work( vector <LL> &x , LL n ) {
   if ( n < int(x.size()) ) return x[n] %</pre>
   vector \langle LL \rangle v = BM( x ); m = v.size();
        if (!m) return 0;
   for ( int i = 0; i < m; ++i ) h[i] = v[i</pre>
        ], a[i] = x[i];
   return calc( n ) % MOD;
 }
} rec;
vector <LL> v;
void solve() {
 int n;
 cin >> n;
 cout << rec.work(v, n - 1) << endl;</pre>
```

6.16 Modular Binomial Coefficients

```
const int N = 2e5+5;
const int mod = 1e9+7;
array <int, N+1> fact, inv, inv_fact;
void init(){
   fact[0] = inv_fact[0] = 1;
   for(int i=1; i<=N; i++){</pre>
       inv[i] = i == 1 ? 1 : (LL) inv[i -
            mod\%i] * (mod/i + 1) % mod;
       fact[i] = (LL) fact[i-1] * i % mod;
       inv_fact[i] = (LL) inv_fact[i-1] *
            inv[i] % mod;
int C(int n,int r){
   if(fact[0] != 1) init();
   return (r < 0 or r > n) ? 0 : (LL) fact[
        n]*inv_fact[r] % mod * inv_fact[n-r
        ] % mod;
```

6.17 Gradient Descent

6.18 Grid Nim

```
int r,c;
scanf("%d %d",&r,&c);
int nim=0;
FOR(i,1,r){
    FOR(j,1,c){
        int tmp;
        scanf("%d",&tmp);
        if(((r-i)+(c-j))%2){
            nim^=tmp;
        }
    }
}
if(nim) printf("Case %d: win\n",tc);
else printf("Case %d: lose\n",tc);
```

6.19 Mobius Function

6.20 Number Theoretic Transformation

```
pair<LL,LL>
#define pii
const LL N= 1<<18;
const LL MOD=786433;
vector<LL>P[N]:
LL rev[N],w[N|1],a[N],b[N],inv_n,g;
LL Pow(LL b,LL p){
   LL ret=1:
   while(p){
       if(p & 1) ret=(ret*b)%MOD;
       b=(b*b)%MOD;
       p>>=1;
   return ret;
LL primitive_root(LL p){
   vector<LL>factor;
   LL phi = p-1,n=phi;
   for(LL i=2;i*i<=n;i++){</pre>
       if(n%i) continue;
       factor.emplace_back(i);
       while(n%i==0) n/=i;
   if(n>1) factor.emplace back(n):
   for(LL res=2;res<=p;res++){</pre>
       bool ok=true;
```

```
for(LL i=0;i<factor.size() && ok;i</pre>
            ++) ok &= Pow(res,phi/factor[i])
             != 1;
       if(ok) return res;
   return -1;
void prepare(LL n){
   LL sz=abs(31-__builtin_clz(n));
   LL r=Pow(g,(MOD-1)/n);
   inv_n=Pow(n,MOD-2);
   w[0]=w[n]=1;
   for(LL i=1;i<n;i++) w[i]= (w[i-1]*r)%MOD</pre>
    for(LL i=1;i<n;i++) rev[i]=(rev[i</pre>
        >>1]>>1) | ((i & 1)<<(sz-1));
void NTT(LL *a,LL n,LL dir=0){
   for(LL i=1;i<n-1;i++) if(i<rev[i]) swap(</pre>
        a[i],a[rev[i]]);
   for(LL m=2;m<=n;m <<= 1) {</pre>
       for(LL i=0;i<n;i+=m){</pre>
           for(LL j=0;j< (m>>1);j++){
               LL &u=a[i+j],&v=a[i+j+(m>>1)
                    ];
               LL t=v*w[dir ? n-n/m*j:n/m*j
                    ]%MOD;
               v=u-t<0?u-t+MOD:u-t;
               u=u+t>=MOD?u+t-MOD:u+t;
       }
    if(dir) for(LL i=0;i<n;i++) a[i]=(inv_n*</pre>
        a[i])%MOD;
vector<LL> mul(vector<LL>p,vector<LL>q){
   LL n=p.size(),m=q.size();
   LL t=n+m-1,sz=1;
   while(sz<t) sz <<= 1;</pre>
   prepare(sz);
   for(LL i=0;i<n;i++) a[i]=p[i];</pre>
   for(LL i=0;i<m;i++) b[i]=q[i];</pre>
   for(LL i=n;i<sz;i++) a[i]=0;</pre>
   for(LL i=m;i<sz;i++) b[i]=0;</pre>
   NTT(a,sz);
   NTT(b,sz);
   for(LL i=0;i<sz;i++) a[i]=(a[i]*b[i])%</pre>
        MOD;
   NTT(a.sz.1):
   vector<LL> c(a,a+sz);
   while(c.size() && c.back()==0) c.
        pop_back();
    return c;
N different number box
Number of ways to make a number by picking
     any number from any of the boxes
vector<LL> solve(LL 1,LL r){
   if(l==r) return P[1];
   LL m=(1+r)/2;
   return mul(solve(1,m),solve(m+1,r));
int main(){
   LL m:
    cin >> m;
   for(LL i=1;i<=m;i++){</pre>
       LL num;
       cin >> num;
       vector<pii>v;
       LL mx=0;
       while(num--){
           LL typ, cnt;
```

```
cin >> typ >> cnt;
    v.emplace_back(typ,cnt);
    mx=max(mx,typ);
}
P[i].resize(mx+1);
for(pii p:v) P[i][p.first]=p.second;
}
g=primitive_root(MOD);
vector<LL>c=solve(1,m);
for(LL i=0;i<c.size();i++){
    if(c[i]){
        cout << i << ' ' << c[i] << '\n';
    }
}</pre>
```

6.21 Chinese Remainder Theorem

```
// given a, b will find solutions for
// ax + by = 1
tuple <LL,LL,LL> EGCD(LL a, LL b){
    if(b == 0) return {1, 0, a};
    else{
        auto [x,y,g] = EGCD(b, a\%b);
        return {y, x - a/b*y,g};
// given modulo equations, will apply CRT
PLL CRT(vector <PLL> &v){
    LL V = 0, M = 1;
    for(auto &[v, m]:v){
        auto [x, y, g] = EGCD(M, m);
if((v - V) % g != 0)
           return {-1, 0};
        V += x * (v - V) / g % (m / g) * M,
            M *= m / g;
        V = (V \% M + M) \% M;
    return make_pair(V, M);
```

7 Tricks

7.1 Fractional Binary Search

```
Given a function f and n, finds the smallest
     fraction p / q in [0, 1] or [0,n]
such that f(p / q) is true, and p, q <= n.
Time: O(log(n))
**/
struct frac { long long p, q; };
bool f(frac x) {
return 6 + 8 * x.p >= 17 * x.q + 12;
frac fracBS(long long n) {
 bool dir = 1, A = 1, B = 1;
 frac lo{0, 1}, hi{1, 0}; // Set hi to 1/0
      to search within [0, n] and {1, 1} to
       search within [0, 1]
 if (f(lo)) return lo;
 assert(f(hi)); //checking if any solution
      exists or not
 while (A || B) {
   long long adv = 0, step = 1; // move hi
       if dir, else lo
   for (int si = 0; step; (step *= 2) >>=
       si) {
     adv += step;
     frac mid{lo.p * adv + hi.p, lo.q * adv
           + hi.q};
     if (abs(mid.p) > n || mid.q > n || dir
           == !f(mid)) {
       adv -= step; si = 2;
   hi.p += lo.p * adv;
   hi.q += lo.q * adv;
   dir = !dir;
   swap(lo, hi);
   A = B; B = !!adv;
 return dir ? hi : lo;
```

7.2 Different Cumulative Sum

```
///cost = sum of (i * a[i]) where i starts
    from the beginning for every range
///sum[] = prefix sum of value[i]
///isum[] = prefix sum of i*value[i]
ll cost(int i, int j) {
    ll ret = isum[j];
    if (i) ret -= isum[i - 1];
    ll baad = sum[j];
    if (i) baad -= sum[i - 1];
    return ret - i * baad;
}
```

7.3 Array Compression

7.4 Factoradic Permutation Trick

```
vector<int> to_factoradic(ll n, int sz = 0)
  vector<int> ret;
 int base = 1;
  while (n) {
   ret.pb(n % base);
   n /= base; base++;
 while ((int)ret.size() < sz) ret.pb(0);</pre>
 reverse(ret.begin(), ret.end());
 return ret;
11 to_decimal(vector<int> &v) {
  int n = (int)v.size();
 11 \text{ ret} = 0, \text{ base} = n;
  for (int i = 0; i < n; i++, base--) {</pre>
   ret = ( ( (ret * base) % MOD ) + v[i] )
        % MOD;
 return ret;
// returns permutation of size n from given
    factoradic number
vector<int> to_permutation(vector<int> &v) {
 vector<int> ret;
  ordered_set<int> st;
  int n = (int)v.size();
  for (int i = 0; i < n; i++) st.insert(i);</pre>
 for (int x : v) {
   int val = *st.find_by_order(x);
   st.erase(val); ret.pb(val);
 return ret:
// returns lexicographical index of
    permutation in factoradic system
vector<int> order_of_permutation(vector<int>
      } (a%
 vector<int> ret; ordered_set<int> st;
  int n = (int)p.size();
 for (int i = 0; i < n; i++) st.insert(i);</pre>
 for (int x : p) {
   int idx = st.order_of_key(x);
   st.erase(x); ret.pb(idx);
 return ret:
// returns sum of indices a and b in
    factoradic system
```

vector<int> add_order(vector<int> &a, vector

<int> &b) {

```
int n = (int)a.size();
 vector<int> ret(n); int carry = 0;
 for (int i = n - 1, base = 1; i >= 0; i--,
       base++) {
   ret[i] = a[i] + b[i] + carry;
   carry = ret[i] / (base); ret[i] %= base;
 return ret;
// returns kth lexicographically smallest
    permutation of size n
// Oth permutation is 0 1 2 ... n-1
vector<int> kth_permutation(int k, int n) {
 // need to handle k \ge n! if necessary
 vector<int> k_factoradic = to_factoradic(k
      , n);
 return to_permutation(k_factoradic);
```

2 Satisfiability

```
2-Sat Note: Assign true or false values to
       n variables in order to satisfy
 a system of constraints on pairs of
      variables.
 E.g: (x1 \text{ or } !x2) and (x2 \text{ or } x3) and (!x3
      or !x3)
 x1 = true
 x2 = true
 x3 = false
  is a solution to make the above formula
 MAX must be equal to the maximum number of
       variables.
 n passed in init() is the number of
      variables.
 O(V+E)
  !a is represented as neg(a).
 example xor:
  |a|b|
  |0|0| \times or(a,b)
  10111
  |1|1| x or(!a, !b)
 do OR of negation of values of variables
      for each undesired situation
 to make it impossible.
struct two_sat {
 int n, id;
 vector<int> g[2 * MAX], rg[2 * MAX], order
  bool state[2 * MAX], vis[2 * MAX];
 int scc[2 * MAX]:
 void init(int _n) {
   n = _n;
   for (int i = 0; i <= 2 * n; i++) {</pre>
     g[i].clear(), rg[i].clear();
     state[i] = vis[i] = false;
     scc[i] = -1:
   st.clear(), order.clear();
 void add_edge(int u, int v) {
   g[u].pb(v);
   rg[v].pb(u);
 void OR(int u, int v) {
   add_edge(neg(u), v);
   add_edge(neg(v), u);
 void XOR(int u, int v) {
   OR(u, v);
   OR(neg(u), neg(v));
```

```
void ForceTrue(int u) {
   add_edge(neg(u), u);
 void ForceFalse(int u) {
   add_edge(u, neg(u));
  void imply(int u, int v) {
   OR(neg(u), v);
 int neg(int u) {
   if (u <= n) return u + n;</pre>
   return u - n;
 void dfs(int u, vii g[], bool topsort) {
   vis[u] = true:
   for (int v : g[u]) {
     if (!vis[v]) dfs(v, g, topsort);
   if (topsort) st.pb(u);
   else scc[u] = id, order.pb(u);
 void build_scc() {
   for (int i = 1; i <= 2 * n; i++) {</pre>
     if (!vis[i]) dfs(i, g, true);
   reverse(st.begin(), st.end());
   fill(vis, vis + 2 * n + 1, false);
   for (int u : st) {
     if (!vis[u]) id++, dfs(u, rg, false);
 bool solve() {
   build scc():
   for (int i = 1; i <= n; i++) {</pre>
     if (scc[i] == scc[i + n]) return false
   for (int i = (int)order.size() - 1; i >= |clock_t start,finish;
         0; i--) {
     int u = order[i];
     if (state[neg(u)] == false) state[u] =
   return true;
} solver;
```

Bitset With Range Opera-

```
tions
struct Bitset {
   const static int B = 6, K = 64, X = 63:
    ///returns mask with bits 1 to r set,
        and others reset
   static inline ULL getmask(int 1, int r)
       if (r==X) return -(1ULL<<1):</pre>
       return (1ULL<<(r+1)) - (1ULL<<1);</pre>
   vector<ULL> bs:
   int N;
   Bitset(int n) {
       N = n/K+1;
       bs.resize(N);
   void assign(ULL x) {
       fill(bs.begin()+1, bs.end(), 0);
       bs[0] = x:
   bool get(int i) {
       return bs[i>>B] & (1ULL<<(i&X));</pre>
   void set(int i) {
       bs[i>>B] |= (1ULL<<(i&X));
   void reset(int i) {
       bs[i>>B] &= ~(1ULL<<(i&X));
```

```
void flip(int i) {
       bs[i>>B] ^= (1ULL<<(i&X));
   void set(int 1, int r) {
       int idl = 1>>B, idr = r>>B;
       int posl = 1&X, posr = r&X;
       if (idl == idr) {
           bs[idl] |= getmask(posl, posr);
       bs[idl] |= getmask(posl, X);
       bs[idr] |= getmask(0, posr);
       for (int id = idl+1; id < idr; id++)</pre>
             bs[id] = -1;
   void reset(int 1, int r) {
       int idl = 1>>B, idr = r>>B;
       int posl = 1&X, posr = r&X;
       if (idl == idr) {
           bs[idl] &= ~getmask(posl, posr);
           return;
       bs[idl] &= ~getmask(posl, X);
       bs[idr] &= ~getmask(0, posr);
       for (int id = idl+1; id < idr; id++)</pre>
             bs[id] = 0;
   void flip(int 1, int r) {
       int idl = 1>>B, idr = r>>B;
       int posl = 1&X, posr = r&X;
       if (idl == idr) {
           bs[idl] ^= getmask(posl, posr);
       bs[idl] ^= getmask(posl, X);
       bs[idr] ^= getmask(0, posr);
       for (int id = idl+1; id < idr; id++)</pre>
             bs[id] = "bs[id];
};
```

7.7Time Count

```
double timespent;
start=clock():
//Here Comes your Code/Function
finish=clock();
timespent=(double)(finish-start)/
     CLOCKS_PER_SEC;
cout << timespent << '\n';</pre>
```

8 Graph

8.1 Lowest Common Ancestor

```
int lvl[MAX], P[MAX][25];
void dfs(int u, int par, int d) {
 lvl[u] = d:
 P[u][0] = par;
 for (int v : g[u]) {
   if (v == par) continue;
   dfs(v, u, d + 1);
void init() {
 for (int j = 0; j < 25; j++) {</pre>
   for (int i = 0; i <= n; i++) P[i][j] =</pre>
        -1;
 dfs(1, -1, 0);
 for (j = 1; j < 25; j++) {</pre>
   for (int i = 1; i <= n; i++) {
     if (P[i][j - 1] != -1) {
       P[i][j] = P[P[i][j-1]][j-1];
       ///to find max weight between two
            edges
       // weight[i][j] = max(weight[i][j
            -1], weight[p[i][j-1]][j-1]);
     }
   }
```

```
int lca(int u, int v) {
 int i, lg;
 if (lvl[u] < lvl[v]) swap(u, v);</pre>
 for (lg = 0; (1 << lg) <= lvl[u]; lg++);</pre>
 for (i = lg; i >= 0; i--) {
   if (lvl[u] - (1 << i) >= lvl[v]) {
     u = P[u][i];
   }
 if (u == v) return u;
 for (i = lg; i >= 0; i--) {
   if (P[u][i] != -1 && P[u][i] != P[v][i])
     u = P[u][i], v = P[v][i];
     // ret = max(ret, weight[u][i]);
     // ret = max(ret, weight[v][i]);
 // ret = max(ret, weight[u][0]);
 return P[u][0];
//Get the ancestor of node "u"
//which is "dis" distance above.
int getAncestor(int u, int dis) {
 dis = lvl[u] - dis;
 int i, lg = 0;
 for (; (1 << lg) <= lvl[u]; lg++) continue</pre>
 lg--;
 for (i = lg; i >= 0; i--) {
  if (lvl[u] - (1 << i) >= dis) {
     u = P[u][i];
   }
 return u;
//returns the distance between
//two nodes "u" and "v".
int dis(int u, int v) {
 if (lvl[u] < lvl[v]) swap(u, v);</pre>
 int p = lca(u, v);
 return lvl[u] + lvl[v] - 2 * lvl[p];
```

Flow with Lower Bound

```
///flow with demand(lower bound) only for
   DAG
//create new src and sink
//add_edge(new src, u, sum(in_demand[u]))
//add_edge(u, new sink, sum(out_demand[u]))
//add_edge(old sink, old src, inf)
// if (sum of lower bound == flow) then
    demand satisfied
//flow in every edge i = demand[i] + e.flow
```

Steiner Tree 8.3

```
/*
Find the minimum cost connected tree where
    at least the important nodes are
    connected
dp(x,i) = minimum cost of a tree rooted at i
     connecting the important node in
    bitmask x.
Complexity: 0(3^k * n + 2^k * m \log m)
int n. k. m:
vector<int> imp;//k important nodes
vector<pair<int, long long>> g[N];
long long d[32][N]; //[2^k][edge count]
const long long inf = LLONG_MAX / 3;
bool vis[N];
```

```
long long MST() {
  for(int i = 0; i < (1 << k); i++) fill(d[i</pre>
      ], d[i] + N, inf);
  for(int i = 0; i < k; ++i) {</pre>
   d[1 << i][imp[i]] = 0;
 priority_queue<pair<long long, int>> q;
  for(int mask = 1; mask < (1 << k); ++mask)</pre>
    for(int a = 0; a < mask; ++a) { //you</pre>
         can still fasten this loop to \ensuremath{\operatorname{get}}
         exact 3<sup>k</sup> complexity
     if((a | mask) != mask) continue; //we
          only need the subsets
     int b = mask ^ a;
     if(b > a) continue;
     for(int v = 0; v < n; ++v) {</pre>
       d[mask][v] = min(d[mask][v], d[a][v]
              + d[b][v]):
   memset(vis, 0, sizeof vis);
    for(int v = 0; v < n; ++v) {
     if(d[mask][v] == inf) continue;
     q.emplace(-d[mask][v], v);
    while(!q.empty()) {
     long long cost = -q.top().first;
     int v = q.top().second;
     q.pop();
     if(vis[v]) continue;
     vis[v] = true;
     for(auto edge : g[v]) {
       long long ec = cost + edge.second;
        if(ec < d[mask][edge.first]) {</pre>
          d[mask][edge.first] = ec;
          q.emplace(-ec, edge.first);
   }
  long long res = inf;
 for(int v = 0; v < n; ++v) {
   res = min(res, d[(1 << k) - 1][v]);
 return res;
```

8.4 Max Clique

```
const int N = 42;
int g[N][N];
int res;
long long edges[N];
//3 ^ (n / 3)
void BronKerbosch(int n, long long R, long
    long P, long long X) {
 if (P == OLL && X == OLL) { //here we will
       find all possible maximal cliques (
      not maximum) i.e. there is no node
      which can be included in this set
   int t = __builtin_popcountl1(R);
   res = max(res, t);
 int u = 0;
 while (!((1LL << u) & (P | X))) u ++;</pre>
 for (int v = 0; v < n; v++) {</pre>
   if (((1LL << v) & P) && !((1LL << v) &</pre>
        edges[u])) {
     BronKerbosch(n, R | (1LL << v), P &
         edges[v], X & edges[v]);
     P -= (1LL << v);
     X = (1LL << v);
int max_clique (int n) {
 for (int i = 1; i <= n; i++) {</pre>
```

```
edges[i - 1] = 0;
  for (int j = 1; j \le n; j++) if (g[i][j
      ]) edges[i - 1] |= ( 1LL << (j - 1)
BronKerbosch(n, 0, (1LL \ll n) - 1, 0);
return res:
```

```
Centroid Decomposition
// problem: calculate the sum of number of
     distinct colors in the path between any
      two pair of nodes
//centroid decomposition (res[i] contains
     the sum of numbers of distinct colors
     in all paths from i)
set<int>g[MAX];
int col[MAX], child[MAX], used[18][MAX];
11 ans[MAX], res[MAX];
int sz = 0, uniq = 0, n;
bool vis[MAX];
void dfs(int u, int p) {
 sz++; child[u] = 1;
 for (auto v : g[u]) {
   if (v != p) {
     dfs(v, u);
     child[u] += child[v];
 }
int get_centroid(int u, int p) {
 for (auto v : g[u]) {
   if (v != p && child[v] > sz / 2) return
        get_centroid(v, u);
 return u;
void add(int u, int p, int depth, int
     centroid) {
  bool check = false; child[u] = 1;
  if (!vis[col[u]]) {
   uniq++; check = true;
   vis[col[u]] = true;
 ans[centroid] += uniq;
 for (auto v : g[u]) {
   if (v != p) {
     add(v, u, depth, centroid);
     child[u] += child[v];
   }
 if (check) {
   used[depth][col[u]] += child[u];
   vis[col[u]] = false;
void del(int u, int p, int depth, int
    centroid) {
 bool check = false;
  if (!vis[col[u]]) {
   unia++:
   used[depth][col[u]] -= child[u];
   vis[col[u]] = true; check = true;
  ans[centroid] -= uniq;
 for (auto v : g[u]) {
   if (v != p) del(v, u, depth, centroid);
 child[u] = 0:
 if (check) uniq--; vis[col[u]] = false;
void solve(int u, int p, int depth, int
     centroid) {
  ans[u] += (ans[p] + child[centroid] - used
      [depth][col[u]]);
  res[u] += ans[u];
  int temp = used[depth][col[u]];
 used[depth][col[u]] = child[centroid];
 for (auto v : g[u]) {
   if (v != p) solve(v, u, depth, centroid)
```

```
ans[u] = 0;
 used[depth][col[u]] = temp;
void reset_col(int u, int p, int depth) {
 used[depth][col[u]] = 0;
 for (auto v : g[u]) {
   if (v != p) reset_col(v, u, depth);
 }
void decompose(int u, int depth) {
 sz = 0:
 uniq = 0;
 dfs(u, -1);
 int centroid = get_centroid(u, -1);
 reset_col(centroid, -1, depth);
 add(centroid, -1, depth, centroid); ///get
       ans for centroid and get the number
      of paths where each color is used
 res[centroid] += ans[centroid];
 vis[col[centroid]] = true;
 for (auto v : g[centroid]) {
   child[centroid] -= child[v];
   ///remove all contribution of the
        subtree of v
   del(v, centroid, depth, centroid);
   used[depth][col[centroid]] = child[
        centroid];
   solve(v, centroid, depth, centroid);
   /\!//add back the contribution of the
        subtree of v
   add(v, centroid, depth, centroid);
   child[centroid] += child[v];
 }
  vis[col[centroid]] = false;
 for (auto it = g[centroid].begin(); it !=
      g[centroid].end(); it++) {
   g[*it].erase(centroid);
   decompose(*it, depth + 1);
 }
int arr[MAX];
int main() {
 fastio;
 cin >> n:
 for (int i = 1; i <= n; i++) cin >> col[i
      1:
 for (int i = 0; i < n - 1; i++) {</pre>
   int u, v;
   cin >> u >> v;
   g[u].insert(v); g[v].insert(u);
 decompose(1, 0);
 for (int i = 1; i <= n; i++) cout << res[i</pre>
      ] << "\n";
```

```
8.6 Dinic Max-Flow
int src, sink;
int dis[MAX], q[MAX], work[MAX];
int n, m, nodes;
struct edge {
 int v, rev, cap, flow;
 edge() {}
 edge(int v, int rev, int cap) : v(v), rev(
      rev), cap(cap), flow(0) {}
vector<edge>g[MAX];
void add_edge(int u, int v, int cap, int rev
    = 0) {
  edge _u = edge(v, g[v].size(), cap);
 edge _v = edge(u, g[u].size(), rev);
 g[u].pb(_u);
 g[v].pb(_v);
bool dinic_bfs(int s) {
 fill(dis, dis + nodes, -1);
 dis[s] = 0;
 int id = 0;
```

```
q[id++] = s;
 for (int i = 0; i < id; i++) {</pre>
   int u = q[i];
   for (int j = 0; j < g[u].size(); j++) {</pre>
     edge &e = g[u][j];
     if (dis[e.v] == -1 && e.flow < e.cap)</pre>
       dis[e.v] = dis[u] + 1;
       q[id++] = e.v;
 return dis[sink] >= 0;
int dinic_dfs(int u, int f) {
 if (u == sink) return f;
 for (int &i = work[u]; i < g[u].size(); i</pre>
      ++) {
   edge &e = g[u][i];
   if (e.cap <= e.flow) continue;</pre>
   if (dis[e.v] == dis[u] + 1) {
     int flow = dinic_dfs(e.v, min(f, e.cap
           - e.flow));
     if (flow) {
       e.flow += flow;
       g[e.v][e.rev].flow -= flow;
       return flow;
 return 0;
int max_flow(int _src, int _sink) {
 src = _src;
sink = _sink;
 int ret = 0;
 while (dinic_bfs(src)) {
   fill(work, work + nodes, 0);
   while (int flow = dinic_dfs(src, INT_MAX
        )) {
     ret += flow;
 return ret;
```

Min Cost Max Flow

```
mt19937 rnd(chrono::steady_clock::now().
    time_since_epoch().count());
struct edge {
 int v, rev;
 11 cap, cost, flow;
 edge() {}
 edge(int v, int rev, ll cap, ll cost) : v(
      v), rev(rev), cap(cap), cost(cost),
      flow(0) {}
struct mcmf {
 int src, sink, nodes;
 vector<int> par, idx, Q;
 vector<bool> inq;
 vector<ll> dis:
 vector<vector<edge>> g;
 mcmf() {}
 mcmf(int src, int sink, int nodes) : src(
      src), sink(sink), nodes(nodes),
   par(nodes), idx(nodes), inq(nodes),
   dis(nodes), g(nodes), Q(10000005) {} //
        use Q(nodes) if not using random
 void add_edge(int u, int v, ll cap, ll
      cost, bool directed = true) {
   edge _u = edge(v, g[v].size(), cap, cost
       );
   edge _v = edge(u, g[u].size(), 0, -cost)
   g[u].pb(_u);
   g[v].pb(_v);
   if (!directed) add_edge(v, u, cap, cost,
```

```
bool spfa() {
   for (int i = 0; i < nodes; i++) {</pre>
     dis[i] = inf, inq[i] = false;
   int f = 0, 1 = 0;
   dis[src] = 0, par[src] = -1, Q[1++] =
       src, inq[src] = true;
   while (f < 1) {
     int u = Q[f++];
     for (int i = 0; i < g[u].size(); i++)</pre>
       edge &e = g[u][i];
       if (e.cap <= e.flow) continue;</pre>
       if (dis[e.v] > dis[u] + e.cost) {
         dis[e.v] = dis[u] + e.cost;
         par[e.v] = u, idx[e.v] = i;
         if(!inq[e.v]) inq[e.v] = true, Q[1
              ++] = e.v;
         // if (!inq[e.v]) {
         // inq[e.v] = true;
         // if (f && rnd() & 7) Q[--f] = e
         // else Q[1++] = e.v;
        // }
     inq[u] = false;
   return (dis[sink] != inf);
 pair<ll, ll> solve() {
   11 mincost = 0, maxflow = 0;
   while (spfa()) {
     11 bottleneck = inf;
     for (int u = par[sink], v = idx[sink];
           u != -1; v = idx[u], u = par[u])
       edge &e = g[u][v];
       bottleneck = min(bottleneck, e.cap -
             e.flow):
     for (int u = par[sink], v = idx[sink];
           u != -1; v = idx[u], u = par[u])
       edge &e = g[u][v];
       e.flow += bottleneck;
       g[e.v][e.rev].flow -= bottleneck;
     mincost += bottleneck * dis[sink].
          maxflow += bottleneck;
   return make_pair(mincost, maxflow);
// want to minimize cost and don't care
    about flow
// add edge from sink to dummy sink (cap =
    inf, cost = 0)
// add edge from source to sink (cap = inf,
    cost = 0)
// run mcmf, cost returned is the minimum
```

K-th Root of a Permuta-8.8 tion

```
vector<vector<int>> decompose(vector<int> &p
    ) {
 int n = p.size();
 vector<vector<int>> cycles;
 vector<bool> vis(n, 0);
 for (int i = 0; i < n; i++) {</pre>
   if (!vis[i]) {
     vector<int> v:
     while (!vis[i]) {
       v.push_back(i);
       vis[i] = 1;
       i = p[i];
     cycles.push_back(v);
   }
 return cycles;
```

vector<int> restore(int n, vector<vector<int</pre>

```
>> &cycles) {
  vector<int> p(n);
 for (auto v : cycles) {
   int m = v.size();
   for (int i = 0; i < m; i++) p[v[i]] = v</pre>
        [(i + 1) % m];
//cycle decomposition of the k-th power of p
vector<vector<int>> power(vector<int> &p,
    int k) {
 int n = p.size();
 auto cycles = decompose(p);
 vector<vector<int>> ans;
 for (auto v : cycles) {
   int len = v.size(), g = __gcd(k, len);
    //g cycles of len / g length
   for (int i = 0; i < g; i++) {</pre>
     vector<int> w;
     for (int j = i, cnt = 0; cnt < len / g
          ; cnt++, j = (j + k) \% len) {
       w.push_back(v[j]);
     ans.push_back(w);
   }
 }
 return ans;
}
//cycle decomposition of the k-th root of p
    with minimum disjoint cycles
//if toggle = 1, then the parity of number
    of cycles will be different from the
     other(toggle = 0) if possible
//returns empty vector if no solution exists
vector<vector<int>> root(vector<int> &p, int
     k, int toggle = 0) {
  int n = p.size();
 vector<vector<int>> cycles[n + 1];
  auto d = decompose(p);
 for (auto v : d) cycles[(int)v.size()].
      push_back(v);
  vector<vector<int>> ans;
 for (int len = 1; len <= n; len++) {</pre>
   if (cycles[len].empty()) continue;
   int tmp = k, t = 1, x = \_gcd(len, tmp);
   while (x != 1) {
     t *= x;
     tmp /= x;
     x = \_gcd(len, tmp);
   }
   if ((int)cycles[len].size() % t != 0) {
     ans.clear();
     return ans; //no solution exists
   int id = 0:
   //we can merge t * z cycles iff tmp % z
        === 0
    if (toggle && tmp % 2 == 0 && (int)
        cycles[len].size() >= 2 * t) {
     int m = 2 * t * len;
     vector<int> cycle(m);
     for (int x = 0; x < 2 * t; x++) { //
          merging 2t cycles to perform the
          toggle
       for (int y = 0; y < len; y++) {</pre>
         cycle[(x + 1LL * y * k) % m] =
              cycles[len][x][y];
     }
     ans.push_back(cycle);
     id = 2 * t;
     toggle = 0;
   for (int i = id; i < (int)cycles[len].</pre>
        size(); i += t) {
     int m = t * len;
     vector<int> cycle(m);
     for (int x = 0; x < t; x++) { //</pre>
          merging t cycles
       for (int y = 0; y < len; y++) {</pre>
         cycle[(x + 1LL * y * k) % m] =
              cycles[len][i + x][y];
```

```
ans.push_back(cycle);
   }
 return ans;
//minimum swaps to obtain this perm from
vector<pair<int, int>> transpositions(vector
    <vector<int>> &cycles) {
  vector<pair<int, int>> ans;
 for (auto v : cycles) {
   int m = v.size();
   for (int i = m - 1; i > 0; i--) ans.
        push_back({v[0], v[i]});
 return ans;
int32_t main() {
 ios_base::sync_with_stdio(0);
 cin.tie(0);
 int n, 1, k;
 cin >> n >> 1 >> k;
  vector<int> p(n);
 for (auto &x : p) cin >> x, --x;
 auto a = root(p, k);
  if (a.empty()) return cout << "no solution vector<edge>g[MAX];
      \n", 0;
 auto t = transpositions(a);
 if (t.size() % 2 != 1 % 2) {
   a = root(p, k, 1);
   t = transpositions(a);
 if (t.size() % 2 != 1 % 2 || t.size() > 1)
       return cout << "no solution\n", 0;</pre>
 auto z = restore(n, a);
 auto w = power(z, k);
 auto x = restore(n, w);
 assert(p == x);
  for (auto x : t) cout << x.first + 1 << '</pre>
       ' << x.second + 1 << '\n';
  for (int i = t.size(); i < 1; i++) cout <<</pre>
       1 << ' ' << 2 << '\n';
 return 0;
```

8.9Block Cut Tree

```
bool ap[MAX];
int id[MAX], koyta[MAX];
int d[MAX], low[MAX];
bool vis[MAX];
vii g[MAX], tree[MAX];
int d_t;
stack<int>st;
vector<vector<int>>comp;
void articulation(int u, int p) {
 vis[u] = true;
 d[u] = low[u] = ++d_t;
 int child = 0; st.push(u);
 for (int v : g[u]) {
   if (v == p) continue;
   if (!vis[v]) {
     articulation(v, u);
     low[u] = min(low[u], low[v]);
     if (p == -1 \&\& child > 1) ap[u] = true
     if (low[v] >= d[u]) {
       if (p != -1) ap[u] = true;
       comp.pb({u}); int top;
         top = st.top(); st.pop();
         comp.back().pb(top);
       } while (top != v);
   } else low[u] = min(low[u], d[v]);
int node = 0:
void make_tree(int n) {
 for (int i = 1; i <= n; i++) {</pre>
```

if (ap[i]) id[i] = ++node;

```
for (int i = 0; i < comp.size(); i++) {</pre>
  ++node;
  int cnt = 0:
  for (int u : comp[i]) {
    if (ap[u]) tree[node].pb(id[u]), tree[
        id[u]].pb(node), koyta[id[u]] =
    else id[u] = node, cnt++;
  koyta[node] = cnt;
}
```

8.10Bridge Tree

```
vector<int> tree[MAX];
bool vis[MAX];
int d[MAX], low[MAX];
int id[MAX];
int d_t;
struct edge {
 int v, rev;
 edge() {}
 edge(int v, int rev) : v(v), rev(rev) {}
vector<bool>is_bridge[MAX];
queue<int>q[MAX];
int comp = 1;
void add_edge(int u, int v) {
 edge _u = edge(v, g[v].size());
  edge _v = edge(u, g[u].size());
 g[u].pb(_u);
 g[v].pb(_v);
 is_bridge[u].pb(false);
  is_bridge[v].pb(false);
void bridge(int u, int p) {
 vis[u] = true;
 d[u] = low[u] = ++d_t;
 for (int i = 0; i < g[u].size(); i++) {</pre>
   edge e = g[u][i]; int v = e.v;
   if (v == p) continue;
   if (!vis[v]) {
     bridge(v, u);
     low[u] = min(low[v], low[u]);
     if (low[v] > d[u]) {
       is_bridge[u][i] = true;
       is_bridge[v][e.rev] = true;
   } else low[u] = min(low[u], d[v]);
void make_tree(int node) {
 int cur = comp; q[cur].push(node);
 vis[node] = true; id[node] = cur;
  while (!q[cur].empty()) {
   int u = q[cur].front(); q[cur].pop();
   for (int i = 0; i < g[u].size(); i++) {</pre>
     edge e = g[u][i]; int v = e.v;
     if (vis[v]) continue;
     if (is_bridge[u][i]) {
       comp++;
       tree[cur].pb(comp);
       tree[comp].pb(cur);
       make_tree(v);
     } else {
       q[cur].push(v);
       vis[v] = true; id[v] = cur;
 }
```

Heavy Light Decomposi-8.11tion

```
int arr[MAX], n;
```

```
vector<int> parent, depth, heavy, head, pos;
int cur_pos, sub[MAX];
int tree[4 * MAX];
vii g[MAX];
void update(int now, int L, int R, int idx,
    int val) {
  if (L == R) {
   tree[now] = val;
 int mid = (L + R) / 2;
 if (idx <= mid) update(now << 1, L, mid,</pre>
      idx, val);
  else update( (now << 1) | 1, mid + 1, R,</pre>
      idx, val);
 tree[now] = tree[now << 1] + tree[(now <<</pre>
      1) | 1];
11 segtree_query(int now, int L, int R, int
    i, int j) {
 if (R < i || L > j) return 0;
 if (L >= i && R <= j) return tree[now];</pre>
 int mid = (L + R) / 2;
 return segtree_query(now << 1, L, mid, i,</pre>
      j) + segtree_query((now << 1) | 1,</pre>
      mid + 1, R, i, j);
int dfs(int u) {
 sub[u] = 1;
 int mx_size = 0;
 for (int v : g[u]) {
   if (v != parent[u]) {
     parent[v] = u, depth[v] = depth[u] +
         1;
     int v_size = dfs(v);
     sub[u] += v_size;
     if (v_size > mx_size) {
       mx_size = v_size;
       heavy[u] = v;
   }
 }
 return sub[u];
void decompose(int u, int h) {
 head[u] = h, pos[u] = cur_pos++;
 if (heavy[u] != -1) decompose(heavy[u], h)
 for (int v : g[u]) {
   if (v != parent[u] && v != heavy[u])
        decompose(v, v);
void init(int n) {
 parent = vector<int>(n, -1);
 depth = vector<int>(n);
 heavy = vector<int>(n, -1);
 head = vector<int>(n);
 pos = vector<int>(n);
 cur_pos = 1;
 dfs(1); decompose(1, 1);
11 query(int a, int b) {
 11 \text{ res} = 0:
 for (; head[a] != head[b]; b = parent[head
      [b]]) {
   if (depth[head[a]] > depth[head[b]])
        swap(a, b);
   11 cur_heavy_path_res = segtree_query(1,
         1, n, pos[head[b]], pos[b]);
   res += cur_heavy_path_res;
 if (depth[a] > depth[b]) swap(a, b);
 11 last_heavy_path_res = segtree_query(1,
      1, n, pos[a], pos[b]);
 res += last_heavy_path_res;
 return res;
}
```

8.12 Tree Isomorphism

```
mp["01"] = 1;
ind = 1;
int dfs(int u, int p) {
```

```
int cnt = 0;
vector<int>vs;
for (auto v : g1[u]) {
 if (v != p) {
   int got = dfs(v, u);
   vs.pb(got);
   cnt++;
 }
}
if (!cnt) return 1;
sort(vs.begin(), vs.end());
string s = "0";
for (auto i : vs) s += to_string(i);
vs.clear();
s.pb('1');
if (mp.find(s) == mp.end()) mp[s] = ++ind;
int ret = mp[s];
return ret;
```

Equations and Formulas

Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n}$$
 $C_0 = 1, C_1 = 1$ and $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$

The number of ways to completely parenthesize n+1 factors.

The number of triangulations of a convex polygon with n+2 sides (i.e. the number of partitions of polygon into disjoint triangles) If $P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$, then, (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

The number of ways to connect the 2n points on a circle to form n disjoint i.e. non-intersecting chords.

The number of rooted full binary trees with n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex 9.5 has either two children or no children.

Number of permutations of $1, \ldots, n$ that avoid the pattern 123 (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n=3, these permutations are 132, 213, 231, 312 and 321.

Stirling Numbers First Kind

The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), where, S(0,0) =$$

$$1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = n!$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$[n - k] = \sum_{0 \le i_1 \le i_2 \le i_k \le n} i_1 i_2 \dots i_k.$$

Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

 $S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1), \text{ where } S(0,0) =$ 1, S(n,0) = S(0,n) = 0 $S(n,2) = 2^{n-1} - 1$ $S(n,k) \cdot k! = \text{number}$ of ways to color n nodes using colors from 1 to k such that each color is used at least once.

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets, with each subset containing at least r elements. It is denoted by $S_r(n,k)$ and obeys the recurrence relation. $S_r(n+1,k) =$

$$kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$$

Denote the \hat{n} objects to partition by the integers $1, 2, \ldots, n$. Define the reduced Stirling numbers of the second kind, denoted $S^d(n,k)$, to be the number of ways to partition the integers 1, 2, ..., n into k nonempty subsets such that all ele- $\sum_{k=1}^{n} \frac{n}{\gcd(k, n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k, n)} - 1$, for n > 1is, for any integers i and j in a given subset, it is required that $|i-j| \ge d$. It has been shown that these numbers satisfy, $\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{i=1}^{n} \mu(d) \lfloor \frac{n}{d} \rfloor^2$ $S^{d}(n,k) = S(n-d+1, k-d+1), n \ge k \ge d$

9.4 Other Combinatorial Identities

$$\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$$

$$\sum_{i=0}^{k} \binom{n+i}{i} = \sum_{i=0}^{k} \binom{n+i}{n} = \binom{n+k+1}{k}$$

$$n, r \in N, n > r, \sum_{i=r}^{n} \binom{i}{r} = \binom{n+1}{r+1}$$

If
$$P(n) = \sum_{k=0}^{n} \binom{n}{k} \cdot Q(k)$$
, then,
$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$
If $P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$, then,
$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

Different Math Formulas

$$(1-x)(1-x^2)(1-x^3)\cdots = 1-x-x^2+x^5+x^7-x^{12}-x^{15}+x^{22}+x^{26}-\cdots$$

The exponents $1,2,5,7,12,\cdots$ on the right hand side are given by the formula $g_k = \frac{k(3k-1)}{2}$ for $k=1,-1,2,-2,3,\cdots$ and are called (generalized) pentagonal numbers. It is useful to find the partition number in $O(n\sqrt{n})$

Let a and b be coprime positive integers, and find integers a'and b' such that $aa' \equiv 1 \mod b$ and $bb' \equiv 1 \mod a$. Then the number of representations of a positive integers (n) as a non negative linear combination of a and b is

$$\frac{n}{ab}$$
 - $\left\{\frac{b'n}{a}\right\}$ - $\left\{\frac{a'n}{b}\right\}$ + 1

9.6GCD and LCM

if m is any integer, then $gcd(a + m \cdot b, b) = gcd(a, b)$

The gcd is a multiplicative function in the following sense: if a_1 and a_2 are relatively prime, then $gcd(a_1 \cdot a_2, b) = gcd(a_1, b)$. $\gcd(a_2,b).$

 $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c)).$

lcm(a, gcd(b, c)) = gcd(lcm(a, b), lcm(a, c)).

For non-negative integers a and b, where a and b are not both zero, $gcd(n^a - 1, n^b - 1) = n^{gcd(a,b)} - 1$

$$\gcd(a,b) = \sum_{k|a \text{ and } k|b} \phi(k)$$

$$\sum_{i=1}^{n} [\gcd(i,n) = k] = \phi\left(\frac{n}{k}\right)$$

$$\sum_{k=1}^{n} \gcd(k,n) = \sum_{d|n} d \cdot \phi\left(\frac{n}{d}\right)$$

$$\sum_{k=1}^{n} x^{\gcd(k,n)} = \sum_{d|n} x^{d} \cdot \phi\left(\frac{n}{d}\right)$$

$$\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$$

$$\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$$

$$\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{d|n} \frac{k}{d} \cdot \frac{n}{d} = 1, \text{ for } n > 1$$

$$\sum_{k=1}^{n} \frac{\gcd(k,n)}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for } n > 1$$

$$\sum_{k=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \lfloor \frac{n}{d} \rfloor^{2}$$

$$\sum_{k=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{k=1}^{n} \mu(d) \lfloor \frac{n}{d} \rfloor^{2}$$

$$\left| \sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \left\lfloor \frac{n}{d} \right\rfloor^{2} \right|$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[\gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i)i^{2}$$

$$F(n) = \sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left(\frac{\left(1 + \left\lfloor \frac{n}{l} \right\rfloor\right) \left(\left\lfloor \frac{n}{l} \right\rfloor\right)}{2} \right)^{2} \sum_{d|l} \mu(d) l d$$