# ${\bf IUT\ SuperSonic,\ Islamic\ University\ of\ Technology}$

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#### Sublime Build

```
"cmd" : ["g++ -std=c++14 $file_name -o
    $file_base_name && timeout 4s ./
    $file_base_name<inputf.in>outputf.in"],
"selector" : "source.cpp",
"file_regex": "^(..[^:]*):([0-9]+):?([0-9]+)
    ?:? (.*)$",
"shell": true,
"working_dir" : "$file_path"
```

### All Macros

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack
    :200000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,ssse3,
    sse4,popcnt,abm,mmx,avx,tune=native")
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
   //find_by_order(k) --> returns iterator
        to the kth largest element counting
         from 0
    //order_of_key(val) --> returns the
        number of items in a set that are
        strictly smaller than our item
template <typename DT>
using ordered_set = tree <DT, null_type,</pre>
    less<DT>, rb_tree_tag,
    tree_order_statistics_node_update>;
int kx[] =
    \{-2,-2,-1,+1,+2,+2,+1,-1\};
int ky[] =
    \{-1,+1,+2,+2,+1,-1,-2,-2\};
#define fastio
                        ios_base::
    sync_with_stdio(0);cin.tie(0);
#define Make(x,p)
                      (x | (1 << p))
#define DeMake(x,p)
                      (x & ~(1<<p))
                      (x & (1<<p))
#define Check(x,p)
#define DEBUG(x)
                      cerr << #x << " = " <<
     x << endl
template<size_t N>
bitset<N> rotl( std::bitset<N> const& bits,
    unsigned count ) {
   count %= N; // Limit count to range [0,N
   return bits << count | bits >> (N -
        count);
```

#### 2 DP

### Convex Hull Trick

```
struct line {
 11 m, c;
 line() {}
 line(ll m, ll c) : m(m), c(c) {}
};
struct convex_hull_trick {
 vector<line>lines;
 int ptr = 0;
 convex_hull_trick() {}
 bool bad(line a, line b, line c) {
   return 1.0 * (c.c - a.c) * (a.m - b.m)
        < 1.0 * (b.c - a.c) * (a.m - c.m);
 void add(line L) {
   int sz = lines.size();
   while (sz >= 2 && bad(lines[sz - 2],
        lines[sz - 1], L)) {
     lines.pop_back(); sz--;
   lines.pb(L);
```

```
ll get(int idx, int x) {
   return (111 * lines[idx].m * x + lines[
        idx].c);
  11 query(int x) {
    if (lines.empty()) return 0;
    if (ptr >= lines.size()) ptr = lines.
        size() - 1;
    while (ptr < lines.size() - 1 && get(</pre>
        ptr, x) > get(ptr + 1, x)) ptr++;
   return get(ptr, x);
};
11 sum[MAX];
11 dp[MAX];
int arr[MAX];
int main() {
  fastio;
  int t;
  cin >> t;
  while (t--) {
   int n, a, b, c;
    cin >> n >> a >> b >> c;
   for (int i = 1; i <= n; i++) cin >> sum
         [i];
    for (int i = 1; i <= n; i++) dp[i] = 0,</pre>
         sum[i] += sum[i - 1];
    convex_hull_trick cht;
    cht.add( line(0, 0) );
    for (int pos = 1; pos <= n; pos++) {</pre>
     dp[pos] = cht.query(sum[pos]) - 1ll *
           a * sqr(sum[pos]) - c;
      cht.add( line(211 * a * sum[pos], dp[
          pos] - a * sqr(sum[pos])) );
   11 \text{ ans } = (-111 * dp[n]);
    ans += (111 * sum[n] * b);
    cout << ans << "\n";
}
```

#### Divide and Conquer DP 2.2

```
inline void compute(int cur, int L, int R,
    int best_L, int best_R) {
 if (L > R) return;
 int mid = (L + R) >> 1;
 pair<11, int>best = {inf, -1};
 for (int k = best_L; k <= min(best_R, mid)</pre>
      ; k++) {
   best = min(best, {dp[cur ^ 1][k - 1] +}
        getCost(k, mid), k});
 dp[cur][mid] = best.ff;
 int best_id = best.ss;
 compute(cur, L, mid - 1, best_L, best_id);
 compute(cur, mid + 1, R, best_id, best_R);
// in main
int cur = 0;
for (int i = 1; i <= n; i++) dp[1][i] = inf;</pre>
for (int guard = 1; guard <= g; guard++) {</pre>
 compute(cur, 1, n, 1, n); cur ^= 1;
ll ans = dp[cur ^ 1][n];
```

```
2.3 Knuth Iterative
for (int i = 1; i <= n; i++) {
 path[i][i] = i;
 dp[i][i] = 0;
for (int len = 2; len <= n; len++) {</pre>
 for (int st = 1; st + len - 1 <= n; st++)</pre>
   int ed = st + len - 1;
   int L = max(st, path[st][ed - 1]);
   int R = min(ed - 1, path[st + 1][ed]);
   dp[st][ed] = INT_MAX;
   for (int i = L; i <= R; i++) {</pre>
     int cur = dp[st][i] + dp[i + 1][ed] +
          arr[ed] - arr[st - 1];
     if (dp[st][ed] > cur) {
       dp[st][ed] = cur;
```

```
path[st][ed] = i;
 }
cout << dp[1][n] << "\n";
```

# 2.4 Knuth Optimization

```
11 solve(int st, int ed) { ///recursive
  if (st == ed) {
    path[st][ed] = st;
    return 0:
  11 &ret = dp[st][ed];
  if (ret != -1) return ret;
  solve(st, ed - 1); solve(st + 1, ed);
  int L = max(st, path[st][ed - 1]);
  int R = min(ed - 1, path[st + 1][ed]);
  ret = inf:
  for (int i = L; i <= R; i++) {</pre>
    ll cur = solve(st, i) + solve(i + 1, ed)
    cur += (arr[ed] - arr[st - 1]);
    if (cur < ret) ret = cur; path[st][ed] =</pre>
  return ret;
///knuth for divide and conquer
int solve(int group, int pos) {
  if (!pos) return dp[group][pos] = 0;
  if (!group) return dp[group][pos] =
      INT_MAX;
  int &ret = dp[group][pos];
  if (ret != -1) return ret;
  int L = 1, R = pos;
  if (pos - 1 > 0) {
    solve(group, pos - 1);
    L = max(L, path[group][pos - 1]);
  if (group + 1 <= m) {</pre>
    solve(group + 1, pos);
    R = min(R, path[group + 1][pos]);
  ret = INT MAX:
  for (int i = L; i <= R; i++) {</pre>
    int cur = solve(group - 1, i - 1) + 111
        * (arr[pos] - arr[i]) * (arr[pos] -
         arr[i]);
    if (cur < ret) {</pre>
      ret = cur:
     path[group][pos] = i;
  return ret;
```

```
Li Chao Tree
2.5
struct line {
 11 m, c;
 line(ll m = 0, ll c = 0) : m(m), c(c) {}
ll calc(line L, ll x) {
 return 111 * L.m * x + L.c:
struct node {
 11 m, c;
 line L;
 node *lft, *rt;
 node(11 m = 0, 11 c = 0, node *lft = NULL,
       node *rt = NULL) : L(line(m, c)),
      lft(lft), rt(rt) {}
};
struct LiChao {
 node *root;
 LiChao() {
   root = new node();
 void update(node *now, int L, int R, line
     newline) {
   int mid = L + (R - L) / 2;
   line lo = now->L, hi = newline;
```

```
if (calc(lo, L) > calc(hi, L)) swap(lo,
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi;
     return;
   if (calc(lo, mid) < calc(hi, mid)) {</pre>
     now->L = hi;
     if (now->rt == NULL) now->rt = new
         node();
     update(now->rt, mid + 1, R, lo);
     now->L = lo;
     if (now->lft == NULL) now->lft = new
         node();
     update(now->lft, L, mid, hi);
   }
 11 query(node *now, int L, int R, 11 x) {
   if (now == NULL) return -inf;
   int mid = L + (R - L) / 2;
   if (x <= mid) return max( calc(now->L, x
        ), query(now->lft, L, mid, x) );
   else return max( calc(now->L, x), query(
        now->rt, mid + 1, R, x) );
 }
};
```

#### **Number Permutation** 2.6

```
11 dp[2][3005]; 11 sum[2][3005];
int dir[3005]:
int arr[MAX];
int main() {
 fastio;
 int n;
 string s:
 cin >> n >> s;
 s = '#' + s;
 s.pb('<'); ///last element less than the</pre>
       element placed after it
 sum[1][0] = 1;
 int cur = 0;
 for (int baki = 1; baki <= n; baki++) {</pre>
   if (s[baki] == '<') dp[cur][0] = 0;</pre>
   else dp[cur][0] = sum[cur ^ 1][baki -
   for (int small = 1; small <= baki; small</pre>
        ++) {
     if (s[baki] == '<') dp[cur][small] =</pre>
          sum[cur ^ 1][small - 1];
       int big = baki - small;
       dp[cur][small] = sum[cur ^ 1][small
            + big - 1];
       dp[cur][small] -= sum[cur ^ 1][small
             - 1];
       if (dp[cur][small] < 0) dp[cur][</pre>
            small] += MOD;
     }
   sum[cur][0] = dp[cur][0];
   for (int small = 1; small <= baki; small</pre>
        ++) {
     sum[cur][small] = (sum[cur][small - 1]
           + dp[cur][small]);
     if (sum[cur][small] >= MOD) sum[cur][
          small] -= MOD;
   }
   cur ^= 1;
 11 ans = dp[cur ^ 1][n];
  cout << ans << "\n";
```

### Same Color Group

```
int prv[21]; 11 cost[21][21];
11 dp[1 << 21]; int m, n;</pre>
bool ok[1 << 21];</pre>
11 solve(ll mask) {
 if (mask == (1 << m) - 1) return Oll;</pre>
 11 &ret = dp[mask];
 if (ok[mask]) return ret;
 ok[mask] = true; ret = inf;
```

```
for (int i = 0; i < m; i++) {</pre>
   if (!(mask & (1 << i) )) {</pre>
      11 c = 0;
      for (int j = 0; j < m; j++) {
        if ((mask & (1 << j)))</pre>
          c += cost[i][j];
      ret = min(ret, c + solve((mask | (1 <<
            i))));
 7
 return ret:
int arr[MAX];
int main() {
 for (int i = 0; i < n; i++) {</pre>
   int val = arr[i];
   val--; prv[val]++;
   for (int j = 0; j < m; j++) {
  if (val == j) continue;</pre>
      cost[val][j] += prv[j];
 11 ans = solve(0);
```

#### Sum of Subsets

```
//submask == all i such that mask&i == i ||
      mask&i == mask (all i such that all 0
      in mask are fixed and the 1's change)
//sos dp memory optimized
for (int i = 0; i < (1 << N); ++i) F[i] = A</pre>
     [i];
for (int i = 0; i < N; ++i) {</pre>
  for (int mask = 0; mask < (1 << N); ++</pre>
      mask) {
    if (mask & (1 << i)) F[mask] += F[mask</pre>
          (1 << i)]; /// doing -= can work
          like inclusion-exclusion on unset
 }
}
```

#### 2.9Triangulation DP

```
bool valid[205][205];
11 dp[205][205];
11 solve(int L, int R) {
 if (L + 1 == R) return 1;
 if (dp[L][R] != -1) return dp[L][R];
 ll ret = 0;
 for (int mid = L + 1; mid < R; mid++) {</pre>
   if (valid[L][mid] && valid[mid][R]) {
     ///selecting triangle(P[L], P[mid], P[
         R1)
     11 temp = ( solve(L, mid) * solve(mid,
           R) ) % MOD;
     ret = (ret + temp) % MOD;
 return dp[L][R] = ret;
```

# **Data Structures**

#### DSU on Tree

```
///Query: Number of distinct names among all
     the k'th son of a node.
const int N = 100005;
string name[N];
vector<int>G[N];
vector<pii>Q[N];
int L[N],ans[N];
void dfs(int v,int d){
   L[v]=d;
   for(int i:G[v]) dfs(i,d+1);
   return:
void dsu(int v,map<int,set<string>>&mp){
   for(int i:G[v]){
       map<int,set<string>>s;
```

```
dsu(i,s);
       if(s.size()>mp.size()) swap(mp,s);
       for(auto it:s) mp[it.ff].insert(all(
            it.ss));
   if(v!=0) mp[L[v]].insert(name[v]); //
        Here zero is not a actual node
    for(pii p:Q[v]) ans[p.ss] = mp[p.ff].
        size();
    return:
int main(){
   int n;
   cin >> n:
   FOR(i.1.n){
       int u;
       cin >> name[i] >> u;
       G[u].pb(i);
   dfs(0,0);
   int q;
   cin >>q;
   FOR(i,1,q){
       int v,k;
       cin >> v >> k;
       Q[v].pb(pii(k+L[v],i)); //Actual
   map<int,set<string>>mp;
    dsu(0,mp);
   FOR(i,1,q) \ cout << \ ans[i] << \ '\n';
   return 0;
```

#### 3.2Dominator Tree

```
struct dominator {
 int n. d t:
 vector<vector<int>>> g, rg, tree, bucket;
 vector<int> sdom, dom, par, dsu, label,
      val, rev;
 dominator() {}
 dominator(int n) :
   n(n), d_t(0), g(n + 1), rg(n + 1),
   tree(n + 1), bucket(n + 1), sdom(n + 1),
   dom(n + 1), par(n + 1), dsu(n + 1),
   label(n + 1), val(n + 1), rev(n + 1)
 { for (int i = 1; i <= n; i++) sdom[i] =
      dom[i] = dsu[i] = label[i] = i; }
 void add_edge(int u, int v) { g[u].pb(v);
 int dfs(int u) {
   d t++:
   val[u] = d_t, rev[d_t] = u;
   label[d_t] = sdom[d_t] = dom[d_t] = d_t;
   for (int v : g[u]) {
     if (!val[v]) {
       dfs(v);
      par[val[v]] = val[u];
     rg[val[v]].pb(val[u]);
   }
 int findpar(int u, int x = 0) {
   if (dsu[u] == u) return x ? -1 : u;
   int v = findpar(dsu[u], x + 1);
   if (v < 0) return u;</pre>
   if (sdom[label[dsu[u]]] < sdom[label[u</pre>
        ]]) label[u] = label[dsu[u]];
   dsu[u] = v;
   return x ? v : label[u];
 void join(int u, int v) { dsu[v] = u; }
 vector<vector<int>> build(int s) {
   dfs(s);
   for (int i = n; i >= 1; i--) {
     for (int j = 0; j < rg[i].size(); j++)</pre>
       sdom[i] = min(sdom[i], sdom[ findpar
            (rg[i][j]) ]);
     if (i > 1) bucket[sdom[i]].pb(i);
```

```
for (int w : bucket[i]) {
      int v = findpar(w);
      if (sdom[v] == sdom[w]) dom[w] =
          sdom[w];
     else dom[w] = v;
   if (i > 1) join(par[i], i);
  for (int i = 2; i <= n; i++) {</pre>
    if (dom[i] != sdom[i]) dom[i] = dom[
         dom[i]];
    tree[rev[i]].pb(rev[dom[i]]);
   tree[rev[dom[i]]].pb(rev[i]);
  return tree;
}
```

# Hopcroft Karp

#include<bits/stdc++.h>

```
using namespace std;
const int N = 3e5 + 9;
struct HopcroftKarp {
  static const int inf = 1e9;
  int n:
  vector<int> 1, r, d;
  vector<vector<int>> g;
  HopcroftKarp(int _n, int _m) {
   n = _n;
   int p = _n + _m + 1;
   g.resize(p);
   1.resize(p, 0);
   r.resize(p, 0);
   d.resize(p, 0);
  void add_edge(int u, int v) {
   g[u].push_back(v + n); //right id is
        increased by n, so is l[u]
  }
  bool bfs() {
   queue<int> q;
   for (int u = 1; u <= n; u++) {</pre>
     if (!1[u]) d[u] = 0, q.push(u);
     else d[u] = inf;
   d[0] = inf;
   while (!q.empty()) {
     int u = q.front();
     q.pop();
     for (auto v : g[u]) {
       if (d[r[v]] == inf) {
         d[r[v]] = d[u] + 1;
         q.push(r[v]);
   return d[0] != inf;
  }
  bool dfs(int u) {
   if (!u) return true:
   for (auto v : g[u]) {
       1 [ii] = v:
       r[v] = u;
       return true;
     }
   d[u] = inf:
   return false;
  int maximum_matching() {
   int ans = 0;
   while (bfs()) {
     for(int u = 1; u <= n; u++) if (!1[u]</pre>
          && dfs(u)) ans++;
   return ans;
 }
}:
int32_t main() {
  ios_base::sync_with_stdio(0);
```

```
cin.tie(0);
int n, m, q;
cin >> n >> m >> q;
HopcroftKarp M(n, m);
while (q--) {
  int u, v;
  cin >> u >> v;
 M.add_edge(u, v);
cout << M.maximum_matching() << '\n';</pre>
return 0;
```

#### Implicit Segment Tree 3.4

```
struct node {
  int val;
 node *lft, *rt;
 node() {}
 node(int val = 0) : val(val), lft(NULL),
      rt(NULL) {}
struct implicit_segtree {
 node *root;
  implicit_segtree() {}
  implicit_segtree(int n) {
   root = new node(n);
  void update(node *now, int L, int R, int
      idx, int val) {
   if (L == R) {
     now -> val += val;
     return;
   int mid = L + (R - L) / 2;
   if (now->lft == NULL) now->lft = new
        node(mid - L + 1);
    if (now->rt == NULL) now->rt = new node(
        R - mid);
    if (idx <= mid) update(now->lft, L, mid,
         idx, val);
   else update(now->rt, mid + 1, R, idx,
        val);
   now->val = (now->lft)->val + (now->rt)->
        val;
  int query(node *now, int L, int R, int k)
   if (L == R) return L;
   int mid = L + (R - L) / 2;
   if (now->lft == NULL) now->lft = new
        node(mid - L + 1);
   if (now->rt == NULL) now->rt = new node(
        R - mid);
   if (k <= (now->lft)->val) return query(
        now->lft, L, mid, k);
   else return query(now->rt, mid + 1, R, k
         - (now->lft)->val);
 }
};
```

#### Implicit Treap 3.5

```
if(d[r[v]] == d[u] + 1 && dfs(r[v])) { | mt19937 rnd(chrono::steady_clock::now().
                                             time_since_epoch().count());
                                        typedef struct node* pnode;
                                        struct node {
                                         int prior, sz;
                                          11 val, sum, lazy;
                                          bool rev:
                                          node *lft, *rt;
                                          node(int val = 0, node *lft = NULL, node *
                                              rt = NULL) : lft(lft), rt(rt), prior(
                                              rnd()), sz(1), val(val), rev(false),
                                              sum(0), lazy(0) {}
                                        struct implicit_treap {
                                         pnode root:
                                          implicit_treap() {
                                           root = NULL;
                                          int get_sz(pnode now) {
                                           return now ? now->sz : 0;
```

```
void update_sz(pnode now) {
 if (!now) return;
 now->sz = 1 + get_sz(now->lft) + get_sz(
      now->rt);
// lazy sum
void push(pnode now) {
 if (!now || !now->lazy) return;
 now->val += now->lazy;
 now->sum += get_sz(now) * now->lazy;
  if (now->lft) now->lft->lazy += now->
      lazy;
  if (now->rt) now->rt->lazy += now->lazy;
 now->lazy = 0;
void combine(pnode now) {
 if (!now) return;
 now->sum = now->val; // reset the node
  push(now->lft), push(now->rt); // update
       lft and rt
  now->sum += (now->lft ? now->lft->sum :
      0) + (now->rt ? now->rt->sum : 0);
// reverse substring
void push(pnode now) {
 if (!now || !now->rev) return;
 now->rev = false;
 swap(now->lft, now->rt);
 if (now->lft) now->lft->rev ^= true;
 if (now->rt) now->rt->rev ^= true;
sort ascending or descending
void push(pnode now) {
 if (!now || !now->sort_kor) return;
  if (now->sort_kor == -1) swap(now->lft,
      now->rt);
  int cnt[26];
  for (int i = 0; i < 26; i++) cnt[i] =</pre>
      now->cnt[i]:
  int idx = 0:
  if (now->lft) {
   memset(now->lft->cnt, 0, sizeof now->
        lft->cnt);
   int lft_sz = get_sz(now->lft);
   while (idx < 26 && lft_sz) {</pre>
     int mn = min(cnt[idx], lft_sz);
     now->lft->cnt[idx] = mn;
     cnt[idx] -= mn; lft_sz -= mn;
     if (!cnt[idx]) idx++;
   now->lft->sort_kor = now->sort_kor;
  while (!cnt[idx]) idx++;
 now->val = idx, cnt[idx]--;
  if (!cnt[idx]) idx++:
  if (now->rt) {
   memset(now->rt->cnt, 0, sizeof now->rt
        ->cnt);
   int rt_sz = get_sz(now->rt);
   while (idx < 26 && rt_sz) {</pre>
     int mn = min(cnt[idx], rt_sz);
     now->rt->cnt[idx] = mn;
     cnt[idx] -= mn; rt_sz -= mn;
     if (!cnt[idx]) idx++;
   now->rt->sort_kor = now->sort_kor;
  if (now->sort_kor == -1) swap(now->lft,
      now->rt):
 now->sort_kor = 0;
void combine(pnode now) {
 if (!now) return;
 memset(now->cnt, 0, sizeof now->cnt);
  for (int i = 0; i < 26; i++) {</pre>
   now->cnt[i] = (now->lft ? now->lft->
        cnt[i] : 0) + (now->rt ? now->rt
        ->cnt[i] : 0);
 now->cnt[now->val]++;
///first pos ta elements go to left,
    others go to right
```

```
void split(pnode now, pnode &lft, pnode &
    rt, int pos, int add = 0) {
 if (!now) return void(lft = rt = NULL);
 push(now);
 int cur = add + get_sz(now->lft);
 if (cur < pos) split(now->rt, now->rt,
      rt, pos, cur + 1), lft = now;
 else split(now->lft, lft, now->lft, pos,
       add), rt = now;
 update_sz(now); combine(now);
}
void merge(pnode &now, pnode lft, pnode rt
    ) {
 push(lft);
 push(rt);
 if (!lft || !rt) now = lft ? lft : rt;
 else if (lft->prior > rt->prior) merge(
      lft->rt, lft->rt, rt), now = lft;
 update_sz(now); combine(now);
void insert(int pos, ll val) {
 if (!root) return void(root = new node(
 pnode lft, rt;
 split(root, lft, rt, pos - 1);
 pnode notun = new node(val);
 merge(root, lft, notun);
 merge(root, root, rt);
void erase(int pos) {
 pnode lft, rt, temp;
 split(root, lft, rt, pos);
 split(lft, lft, temp, pos - 1);
 merge(root, lft, rt);
 delete(temp);
void reverse(int 1, int r) {
 pnode lft, rt, mid;
 split(root, lft, mid, l - 1);
 split(mid, mid, rt, r - l + 1);
 mid->rev ^= true;
 merge(root, lft, mid);
 merge(root, root, rt);
}
void right_shift(int 1, int r) {
 pnode lft, rt, mid, last;
 split(root, lft, mid, l - 1);
 split(mid, mid, rt, r - l + 1);
 split(mid, mid, last, r - 1);
 merge(mid, last, mid);
 merge(root, lft, mid);
 merge(root, root, rt);
void output(pnode now, vector<int>&v) {
 if (!now) return;
 push(now);
 output(now->lft, v);
 v.pb(now->val);
 output(now->rt, v);
vector<int>get_arr() {
 vector<int>ret:
 output(root, ret);
 return ret;
```

#### LCA 3.6

```
/*Hey, What's up?*/
#include<bits/stdc++.h>
using namespace std;
#define pi acos(-1.0)
#define fastio ios_base::sync_with_stdio(
    false);cin.tie(NULL);cout.tie(NULL)
vector<long long>v[100005],vc;
long long x[200005][40],mp[100005],ml
     [100005],nd,pos[100005];
void build(long long n)
   long long a,i,j,k,b,c;
```

```
for(i=0; i<n; i++)</pre>
                                                   x[i][0]=vc[i];
                                               b=1:
                                                while(a<n)</pre>
                                                {
                                                   for(i=0; i<n-a; i++)</pre>
                                                       x[i][b]=min(x[i][b-1],x[i+a][b
                                                   a*=2:
                                                   b++;
                                                return;
else merge(rt->lft, lft, rt->lft), now = long long query(long long a, long long b)
                                                long long c,d,e,f;
                                                //if(a==b)return x[a][0];
                                                c=log2(1.0*(b-a+1));
                                                //cout<<c<' ';
                                                f=powl(1.0*2,1.0*c);
                                                d=x[a][c];
                                                e=x[b-f+1][c];
                                                //cout<<b-f+1<<' ';
                                                return min(d,e);
                                            void tour_de_euler(long long p, long long q)
                                                vc.push_back(mp[p]);
                                                //nd++:
                                                if(!pos[mp[p]])pos[mp[p]]=nd;
                                               nd++;
                                                for(int i=0;i<v[p].size();i++){</pre>
                                                   if(v[p][i]==q)continue;
                                                   tour_de_euler(v[p][i],p);
                                                   vc.push_back(mp[p]); nd++;
                                               return;
                                            void dfs(long long p, long long q)
                                               mp[p]=nd;
                                               ml[nd]=p;
                                               nd++;
                                                for(int i=0;i<v[p].size();i++){</pre>
                                                    if(v[p][i]==q)continue;
                                                   dfs(v[p][i],p);
                                               return:
                                            long long lca(long long a, long long b)
                                                a=pos[mp[a]];
                                                b=pos[mp[b]];
                                                   swap(a,b);
                                               long long c=query(a,b);
                                                return ml[c];
                                            int main()
                                                long long a=0,b=0,c,d,e,f=0,1,g,m,n,k,i,
                                                     j,t,p,q;
                                                cin>>n;
                                                for(i=1; i<n; i++)</pre>
                                                   cin>>a>>b;
                                                   v[a].push_back(b);
                                                   v[b].push_back(a);
                                               }
                                               nd=1:
                                                dfs(1,-1);
                                                vc.push_back(696969696969);
```

```
tour_de_euler(1,-1);
l=vc.size();
build(1+2);
cin>>q;
while (q--) {
    cin>>a>>b;
    cout<<lca(a,b)<<endl;</pre>
return 0;
```

```
Link Cut Tree
struct SplayTree {
 struct node {
   int ch[2] = \{0, 0\}, p = 0;
   11 self = 0, path = 0;
   11 \text{ sub} = 0, \text{ extra} = 0;
   bool rev = false;
 };
 vector<node> T;
 SplayTree(int n) : T(n + 1) {}
 void push(int x) {
   if (!x) return;
   int 1 = T[x].ch[0], r = T[x].ch[1];
   if (T[x].rev) {
     T[1].rev ^= true, T[r].rev ^= true;
     swap(T[x].ch[0], T[x].ch[1]);
     T[x].rev = false;
 void pull(int x) {
   int 1 = T[x].ch[0], r = T[x].ch[1];
   push(1), push(r);
   T[x].path = T[x].self + T[l].path + T[r]
        ].path;
   T[x].sub = T[x].self + T[x].extra + T[1]
        ].sub + T[r].sub;
 void set(int parent, int child, int d) {
   T[parent].ch[d] = child;
   T[child].p = parent;
   pull(parent);
 int dir(int x) {
   int parent = T[x].p;
   if (!parent) return -1;
   return (T[parent].ch[0] == x) ? 0 : (T[
        parent].ch[1] == x) ? 1 : -1;
 void rotate(int x) {
   int parent = T[x].p, gparent = T[parent
        ].p;
   int dx = dir(x), dp = dir(parent);
   set(parent, T[x].ch[!dx], dx);
   set(x, parent, !dx);
   if (~dp) set(gparent, x, dp);
   T[x].p = gparent;
 void splay(int x) {
   push(x);
   while (~dir(x)) {
     int parent = T[x].p;
     int gparent = T[parent].p;
     push(gparent), push(parent), push(x);
     int dx = dir(x), dp = dir(parent);
     if (~dp) rotate(dx != dp ? x : parent)
     rotate(x);
struct LinkCut : SplayTree {
 LinkCut(int n) : SplayTree(n) {}
 void cut_right(int x) {
   splay(x);
   int r = T[x].ch[1];
   T[x].extra += T[r].sub;
   T[x].ch[1] = 0, pull(x);
 int access(int x) {
   int u = x, v = 0;
   for (; u; v = u, u = T[u].p) {
```

```
cut_right(u);
     T[u].extra -= T[v].sub;
     T[u].ch[1] = v, pull(u);
   return splay(x), v;
  void make_root(int x) {
   access(x);
T[x].rev ^= true, push(x);
  void link(int u, int v) {
   make_root(v), access(u);
   T[u].extra += T[v].sub;
   T[v].p = u, pull(u);
  void cut(int u) {
   access(u);
   T[u].ch[0] = T[T[u].ch[0]].p = 0;
   pull(u);
  void cut(int u, int v) {
   make_root(u), access(v);
   T[v].ch[0] = T[u].p = 0, pull(v);
  int find_root(int u) {
   access(u), push(u);
   while (T[u].ch[0]) {
     u = T[u].ch[0], push(u);
   return splay(u), u;
  int lca(int u, int v) {
   if (u == v) return u;
   access(u);
   int ret = access(v);
   return T[u].p ? ret : 0;
  // subtree query of {\tt u} if {\tt v} is the root
  11 subtree(int u, int v) {
   make_root(v), access(u);
   return T[u].self + T[u].extra;
 11 path(int u, int v) {
   make_root(u), access(v);
   return T[v].path;
 }
  // point update
  void update(int u, ll val) {
   access(u);
   T[u].self = val, pull(u);
 }
};
```

```
Merge Sort Tree
3.8
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template<typename T> using ordered_set =
    tree<T, null_type, less<T>, rb_tree_tag
    , tree_order_statistics_node_update>;
ordered_set<pii> bst[MAX << 2];
void init(int n) {
 for (int i = 0; i <= 4 * n; i++) bst[i].</pre>
void build(int now, int L, int R) {
 if (L == R) {
   bst[now].insert({arr[L], L});
 }
 for (int i = L; i <= R; i++) bst[now].</pre>
      insert({arr[i], i});
  int mid = (L + R) / 2;
 build(now << 1, L, mid);</pre>
 build((now << 1) | 1, mid + 1, R);
void update(int now, int L, int R, int idx,
    int ager_val, int val) {
```

```
if (L == R) {
   bst[now].erase(bst[now].find({ager_val,
        idx})):
   bst[now].insert({val, idx});
   return;
 int mid = (L + R) / 2;
 if (idx <= mid) update(now << 1, L, mid,</pre>
      idx, ager_val, val);
 else update((now << 1) | 1, mid + 1, R,</pre>
      idx, ager_val, val);
 bst[now].erase(bst[now].find({ager_val,
      idx}));
 bst[now].insert({val, idx});
ll query(int now, int L, int R, int i, int j
     int val) {
 if (R < i || L > j) return 0;
 if (L >= i && R <= j) {</pre>
   int ret = bst[now].order_of_key({val,
        INT_MAX});
   return ret;
 int mid = (L + R) / 2;
 return query(now << 1, L, mid, i, j, val)</pre>
      + query((now << 1) | 1, mid + 1, R, i
      , j, val);
```

# 3.9 Mo on Tree

```
#include <bits/stdc++.h>
using namespace std:
const int N = 3e5 + 9;
//unique elements on the path from u to v
vector<int> g[N];
int st[N], en[N], T, par[N][20], dep[N], id[
    N * 21:
void dfs(int u, int p = 0) {
 st[u] = ++T;
 id[T] = u;
 dep[u] = dep[p] + 1;
 par[u][0] = p;
  for(int k = 1; k < 20; k++) par[u][k] =</pre>
      par[par[u][k - 1]][k - 1];
 for(auto v : g[u]) if(v != p) dfs(v, u);
 en[u] = ++T;
 id[T] = u;
int lca(int u, int v) {
 if(dep[u] < dep[v]) swap(u, v);</pre>
 for(int k = 19; k \ge 0; k--) if(dep[par[u]
     ][k]] >= dep[v]) u = par[u][k];
 if(u == v) return u;
  for(int k = 19; k >= 0; k--) if(par[u][k]
      != par[v][k]) u = par[u][k], v = par[
      v][k];
 return par[u][0];
int cnt[N], a[N], ans;
inline void add(int u) {
 int x = a[u];
 if(cnt[x]++ == 0) ans++;
inline void rem(int u) {
 int x = a[u]:
 if(--cnt[x] == 0) ans--;
bool vis[N]:
inline void yo(int u) {
 if(!vis[u]) add(u);
 else rem(u);
 vis[u] ^= 1;
const int B = 320;
struct query {
 int 1, r, id;
 bool operator < (const query &x) const {</pre>
```

```
if(1 / B == x.1 / B) return r < x.r;</pre>
   return 1 / B < x.1 / B;</pre>
} Q[N];
int res[N];
int main() {
 ios_base::sync_with_stdio(0);
  cin.tie(0);
  int n, q;
  while(cin >> n >> q) {
    for(int i = 1; i <= n; i++) cin >> a[i];
    map<int, int> mp;
    for(int i = 1; i <= n; i++) {</pre>
      if(mp.find(a[i]) == mp.end()) mp[a[i
           ]], mp[a[i]] = mp.size();
     a[i] = mp[a[i]];
    for(int i = 1; i < n; i++) {</pre>
      int u, v;
      cin >> u >> v;
      g[u].push_back(v);
      g[v].push_back(u);
    T = 0:
    dfs(1);
    for(int i = 1; i <= q; i++) {</pre>
     int u, v;
      cin >> u >> v;
      if(st[u] > st[v]) swap(u, v);
      int lc = lca(u, v);
      if(lc == u) Q[i].1 = st[u], Q[i].r =
          st[v];
      else Q[i].1 = en[u], Q[i].r = st[v];
      Q[i].id = i;
    sort(Q + 1, Q + q + 1);
    ans = 0:
    int 1 = 1, r = 0;
    for(int i = 1; i <= q; i++) {</pre>
      int L = Q[i].1, R = Q[i].r;
      if(R < 1) {
        while (1 > L) yo(id[--1]);
        while (1 < L) yo(id[1++]);</pre>
        while (r < R) yo(id[++r]);
        while (r > R) yo(id[r--]);
      } else {
        while (r < R) yo(id[++r]);
        while (r > R) yo(id[r--]);
        while (1 > L) yo(id[--1]);
        while (1 < L) yo(id[1++]);</pre>
      int u = id[1], v = id[r], lc = lca(u)
          v):
      if(lc != u && lc != v) yo(lc); //take
           care of the lca separately
      res[Q[i].id] = ans;
      if(lc != u && lc != v) yo(lc);
    for(int i = 1; i <= q; i++) cout << res[</pre>
        i] << '\n';
    for(int i = 0; i <= n; i++) {</pre>
      g[i].clear();
      vis[i] = cnt[i] = 0;
      for(int k = 0; k < 20; k++) par[i][k]</pre>
           = 0:
 return 0:
```

### 3.10 Mo's Algorithm

```
vector<pair<long long,pair<long long,</pre>
    long long> > x[d+2];
//map<pair<long long,long long> ,long
     long>mp;
v.push_back(-37);
for(i=0;i<n;i++){</pre>
   cin>>a;
   v.push_back(a);
}
cin>>q;
for(i=0;i<q;i++){</pre>
    cin>>a>>b:
   e=a/d;
   x[e].push_back({b,{a,i}});
for(i=0;i<=d+1;i++){</pre>
    sort(x[i].begin(),x[i].end());
for(i=0;i<=d;i++){</pre>
   memset(vis,0,sizeof(vis));
   l=i*d+1;
   r=i*d:
   p=x[i].size();
   f=0:
    for(j=0;j<p;j++){</pre>
       b=x[i][j].first;
       a=x[i][j].second.first;
       while(r<b){</pre>
           r++:
           vis[v[r]]++;
           if(vis[v[r]]==1)f++;
       //cout<<l<' '<<r<'='<<f<<endl;
       if(1<a){
       while(1<a){
           vis[v[1]]--;
           if(vis[v[1]]==0)f--;
           1++;
       else if(1>a){
           while(1>a){
               vis[v[1]]++;
               if(vis[v[1]]==1)f++;
       }
       //cout<<l<' '<<r<'='<<f<<endl;
       y[x[i][j].second.second]=f;
        //cout<<a<<' '<<b<<'='<<f<<endl:
for(i=0;i<q;i++){</pre>
   cout<<y[i]<<'\n';
return 0:
```

### 3.11 Persistent Segment Tree

```
struct node {
  int val, lft, rt;
 node(int val = 0, int lft = 0, int rt = 0)
       : val(val), lft(lft), rt(rt) {}
node nodes[30 * MAX]; ///take at least 2*n*
    log(n) nodes
int root[MAX], sz;
inline int update(int &now, int L, int R,
    int idx, int val) {
  if (L > idx || R < idx) return now;</pre>
 if (L == R) {
   nodes[sz] = nodes[now];
   nodes[sz].val += val;
   return sz;
 }
 int mid = (L + R) >> 1;
 int ret = ++sz;
 if (idx <= mid) {</pre>
   if (!nodes[now].lft) nodes[now].lft = ++
   nodes[ret].lft = update(nodes[now].lft,
        L, mid, idx, val);
```

```
nodes[ret].rt = nodes[now].rt;
 } else {
   if (!nodes[now].rt) nodes[now].rt = ++sz
   nodes[ret].rt = update(nodes[now].rt,
        mid + 1, R, idx, val);
   nodes[ret].lft = nodes[now].lft;
 nodes[ret].val = nodes[ nodes[ret].lft ].
      val + nodes[ nodes[ret].rt ].val;
 return ret;
inline int query(int &now, int L, int R, int | };
     i, int j) {
 if (L > j || R < i) return 0;</pre>
 if (L \ge i \&\& R \le j) return nodes [now].
      val;
 int mid = (L + R) \gg 1;
 return query(nodes[now].lft, L, mid, i, j)
       + query(nodes[now].rt, mid + 1, R, i
      , j);
/// in main(make segtree for every prefix)
root[0] = 0;
for (int i = 1; i <= n; i++) root[i] =</pre>
    update(root[i - 1], 1, n, p[i], 1);
3.12
         Sparse Table
const int maxn = (1 << 20) + 5;
int logs[maxn] = {0};
```

```
void compute_logs(){
   logs[1] = 0;
   for(int i=2;i<(1<<20);i++){</pre>
       logs[i] = logs[i/2]+1;
class Sparse_Table
   public:
       vector <vector<LL>> table;
       function < LL(LL,LL) > func;
       LL identity;
   Sparse_Table(vector <LL> &v, function <</pre>
        LL(LL,LL)> _func, LL id){
       if(logs[2] != 1) compute_logs();
       int sz = v.size();
       table.assign(sz,vector <LL>(logs[sz
            1+1) ):
       func = _func, identity = id;
       for(int j=0;j<=logs[sz];j++){</pre>
           for(int i=0;i+(1<<j)<=sz;i++){</pre>
              if(j==0) table[i][j] = func(v
                    [i],id); // base case,
                    when only 1 element in
                   range
               else table[i][j] = func(table
                    [i][j-1], table[i +
                    (1<<(j-1))][j-1] );
           }
   // when intersection of two ranges wont
        be a problem like min, gcd, max
   LL query(int 1, int r){
       assert(r>=1);
       int pow = logs[r-l+1];
       return func(table[1][pow], table[r-
            (1<<pow) + 1][pow]);
   // other cases like sum
   LL Query(int 1,int r){
       if(l>r) return identity; // handle
       int pow = logs[r - 1 + 1];
       return func(table[1][pow], Query(1
            +(1<<pow), r));
```

### 3.13 Treap

```
mt19937 rnd(chrono::steady_clock::now().
    time_since_epoch().count());
typedef struct node* pnode;
struct node {
 int prior, val, sz;
 ll sum;
 node *lft, *rt;
 node(int val = 0, node *lft = NULL, node *
      rt = NULL)
   lft(lft), rt(rt), prior(rnd()), val(val)
        , sz(1), sum(0) {}
struct treap {
 pnode root;
 treap() {
   root = NULL:
 int get_sz(pnode now) {
   return now ? now->sz : 0;
 void update_sz(pnode now) {
   if (!now) return;
   now->sz = 1 + get_sz(now->lft) + get_sz(
        now->rt):
 11 get(pnode now) {
   return now ? now->sum : 0;
 void push(pnode now) {}
 void combine(pnode now) {
   if (!now) return;
   now->sum = now->val + get(now->lft) +
        get(now->rt);
 pnode unite(pnode lft, pnode rt) {
   if (!lft || !rt) return lft ? lft : rt;
   // push(lft), push(rt); this not tested
   if (lft->prior < rt->prior) swap(lft, rt
       );
   pnode 1, r;
   split(rt, 1, r, lft->val);
   lft->lft = unite(lft->lft, 1), update_sz
        (lft):
   lft->rt = unite(lft->rt, r), update_sz(
        lft);
   // combine(lft); this not tested
   return lft;
 ///value < val goes to left, value >= val
      goes to right
 void split(pnode now, pnode &lft, pnode &
      rt, int val, int add = 0) {
   push(now);
   if (!now) return void(lft = rt = NULL);
   if (now->val < val) split(now->rt, now->
       rt, rt, val), lft = now;
   else split(now->lft, lft, now->lft, val)
        , rt = now;
   update_sz(now), combine(now);
 void merge(pnode &now, pnode lft, pnode rt
      ) {
   push(lft), push(rt);
   if (!lft || !rt) now = lft ? lft : rt;
   else if (lft->prior > rt->prior) merge(
        lft->rt, lft->rt, rt), now = lft;
   else merge(rt->lft, lft, rt->lft), now =
   update_sz(now), combine(now);
 void insert(pnode &now, pnode notun) {
   if (!now) return void(now = notun);
   push(now);
   if (notun->prior > now->prior) split(now
        , notun->lft, notun->rt, notun->val
        ), now = notun;
   else insert(notun->val < now->val ? now
        ->lft : now->rt, notun);
   update_sz(now), combine(now);
 void erase(pnode &now, int val) {
   push(now);
```

```
if (now->val == val) {
   pnode temp = now;
   merge(now, now->lft, now->rt);
   delete(temp);
 } else erase(val < now->val ? now->lft :
       now->rt, val);
 update_sz(now), combine(now);
int get_idx(pnode &now, int val) {
 if (!now) return INT_MIN;
 else if (now->val == val) return 1 +
      get_sz(now->lft);
 else if (val < now->val) return get_idx( |4
      now->lft, val);
 else return (1 + get_sz(now->lft) +
      get_idx(now->rt, val));
int find_kth(pnode &now, int k) {
 if (k < 1 || k > get_sz(now)) return -1;
 if (get_sz(now->lft) + 1 == k) return
      now->val;
 if (k <= get_sz(now->lft)) return
      find_kth(now->lft, k);
 return find_kth(now->rt, k - get_sz(now
      ->lft) - 1);
ll prefix_sum(pnode &now, int k) {
 if (k < 1 || k > get_sz(now)) return -
      inf;
 if (get_sz(now->lft) + 1 == k) return
      get(now->lft) + now->val;
 if (k <= get_sz(now->lft)) return
      prefix_sum(now->lft, k);
 return get(now->lft) + now->val +
      prefix_sum(now->rt, k - get_sz(now
       ->lft) - 1);
pnode get_rng(int 1, int r) { ///gets all
    1 <= values <= r
 pnode lft, rt, mid;
 split(root, lft, mid, 1);
 split(mid, mid, rt, r + 1);
 merge(root, lft, rt);
 return mid;
void output(pnode now, vector<int>&v) {
 if (!now) return;
 output(now->lft, v);
 v.pb(now->val);
 output(now->rt, v);
vector<int>get_arr() {
 vector<int>ret;
 output(root, ret);
 return ret;
```

# 3.14 Trie

```
int trie[30 * 100000 + 5][2];
int mark[30 * 100000 + 5];
int node = 1;
void add(int n) {
  int now = 1:
 for (int i = 27; i >= 0; i--) {
    int d = (bool)(n & (1 << i));</pre>
   if (!trie[now][d]) trie[now][d] = ++node Point rotate(Point a, Tf rad) {
   now = trie[now][d];
   mark[now]++;
void del(int n) {
 int now = 1;
 deque<int>v;
 for (int i = 27; i >= 0; i--) {
   int d = (bool)(n & (1 << i));</pre>
   if (trie[now][d]) {
     v.push_front(now);
     now = trie[now][d];
     mark[now]--;
```

```
v.push_front(now);
for (int i = 1; i < v.size(); i++) {</pre>
  if (!mark[v[i - 1]]) {
    if (trie[v[i]][0] == v[i - 1]) trie[v[ Point scale(Point a, Tf s) {
        i]][0] = 0;
    if (trie[v[i]][1] == v[i - 1]) trie[v[
        i]][1] = 0;
 }
}
```

/// use long long for

# Geometry

# 4.1 2D Point typedef double Tf;

exactness

 $x<0 ? -1 : 1);}$ 

const Tf PI = acos(-1), EPS = 1e-9;

typedef Tf Ti;

```
struct Point {
   Ti x, y;
    Point(Ti x = 0, Ti y = 0) : x(x), y(y) {}
    Point operator + (const Point& u) const {
              return Point(x + u.x, y + u.y); }
    Point operator - (const Point& u) const {
               return Point(x - u.x, y - u.y); }
    Point operator * (const long long u) const
                  { return Point(x * u, y * u); }
    Point operator * (const Tf u) const {
              return Point(x * u, y * u); }
    Point operator / (const Tf u) const {
               return Point(x / u, y / u); }
    bool operator == (const Point& u) const {
               return dcmp(x - u.x) == 0 \&\& dcmp(y -
                  u.y) == 0; }
    bool operator != (const Point& u) const {
               return !(*this == u); }
    bool operator < (const Point& u) const {</pre>
               return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x)) \mid
                 u.x) == 0 && dcmp(y - u.y) < 0); }
Ti dot(Point a, Point b) { return a.x * b.x
          + a.y * b.y; }
Ti cross(Point a, Point b) { return a.x * b.
           y - a.y * b.x; }
Tf length(Point a) { return sqrt(dot(a, a));
Ti sqLength(Point a) { return dot(a, a); }
Tf distance(Point a, Point b) {return length
           (a-b);
Tf angle(Point u) { return atan2(u.y, u.x);
// returns angle between oa, ob in (-PI, PI]
Tf angleBetween(Point a, Point b) {
   double ans = angle(b) - angle(a);
   return ans <= -PI ? ans + 2*PI : (ans > PI
                  ? ans - 2*PI : ans);
// Rotate a ccw by rad radians
   static_assert(is_same<Tf, Ti>::value);
    return Point(a.x * cos(rad) - a.y * sin(
               rad), a.x * sin(rad) + a.y * cos(rad)
// rotate a ccw by angle th with cos(th) =
         co && sin(th) = si
Point rotatePrecise(Point a, Tf co, Tf si) {
    static_assert(is_same<Tf, Ti>::value);
   return Point(a.x * co - a.y * si, a.y * co };
                  + a.x * si);
Point rotate90(Point a) { return Point(-a.y, |4.2|
              a.x); }
```

```
// scales vector a by s such that length of
                                                  a becomes s
                                               static_assert(is_same<Tf, Ti>::value);
                                               return a / length(a) * s;
                                              // returns an unit vector perpendicular to
                                                  vector a
                                              Point normal(Point a) {
                                                static_assert(is_same<Tf, Ti>::value);
                                               Tf 1 = length(a);
                                               return Point(-a.y / 1, a.x / 1);
                                              // returns 1 if c is left of ab, 0 if on ab
                                                   && -1 if right of ab
                                              int orient(Point a, Point b, Point c) {
                                               return dcmp(cross(b - a, c - a));
int dcmp(Tf x) { return abs(x) < EPS ? 0 : (</pre>
                                              class polarComp {
                                                 point 0, dir;
                                                  bool half(point p) {
                                                     return dcmp(dir & p) < 0 || (dcmp(</pre>
                                                          dir & p) == 0 && dcmp(dir ^ p) >
                                                  public:
                                                  polarComp(point 0 = point(0, 0), point
                                                      dir = point(1, \bar{0})
                                                     : 0(0), dir(dir) {}
                                                  bool operator() (point p, point q) {
                                                     return make_tuple(half(p), 0) <</pre>
                                                          make_tuple(half(q), (p & q));
                                              }; // given a pivot point and an initial
                                                   direction, sorts by Angle with the
                                                   given direction
                                              struct Segment {
                                                Point a, b;
                                                Segment(Point aa, Point bb) : a(aa), b(bb)
                                                     {}
                                              typedef Segment Line;
                                              struct Circle {
                                                 Point o;
                                                 Tf r;
                                                 Circle(Point o = Point(0, 0), Tf r = 0)
                                                      : o(o), r(r) {}
                                                  // returns true if point p is in || on
                                                      the circle
                                                  bool contains(Point p) {
                                                   return dcmp(sqLength(p - o) - r * r)
                                                        <= 0:
                                                  // returns a point on the circle rad
                                                      radians away from +X CCW
                                                 Point point(Tf rad) {
                                                   static_assert(is_same<Tf, Ti>::value);
                                                   return Point(o.x + cos(rad) * r, o.y +
                                                         sin(rad) * r);
                                                  // area of a circular sector with
                                                      central angle rad
                                                  Tf area(Tf rad = PI + PI) { return rad *
                                                       r * r / 2; }
                                                  // area of the circular sector cut by a
                                                      chord with central angle alpha
                                                  Tf sector(Tf alpha) { return r * r * 0.5
                                                       * (alpha - sin(alpha)); }
```

#### Circle Algorithms

```
// Extremely inaccurate for finding near
    touches
// compute intersection of line 1 with
    circle c
// The intersections are given in order of
    the ray (1.a, 1.b)
vector<Point> circleLineIntersection(Circle
    c. Line 1) {
   static_assert(is_same<Tf, Ti>::value);
   vector<Point> ret;
   Point b = 1.b - 1.a, a = 1.a - c.o;
   Tf A = dot(b, b), B = dot(a, b);
   Tf C = dot(a, a) - c.r * c.r, D = B * B
        - A * C:
   if (D < -EPS) return ret;</pre>
   ret.push_back(l.a + b * (-B - sqrt(D +
        EPS)) / A);
   if (D > EPS) ret.push_back(1.a + b * (-B
         + sqrt(D)) / A);
   return ret;
// signed area of intersection of circle(c.o
    , c.r) &&
// triangle(c.o, s.a, s.b) [cross(a-o, b-o)
Tf circleTriangleIntersectionArea(Circle c,
    Segment s) {
   Tf OA = length(c.o - s.a);
   Tf OB = length(c.o - s.b);
   // sector
   if (dcmp(distancePointSegment(c.o, s) -
        c.r) >= 0)
       return angleBetween(s.a - c.o, s.b -
             c.o) * (c.r * c.r) / 2.0;
   // triangle
   if (dcmp(OA - c.r) \le 0 \&\& dcmp(OB - c.r)
        ) <= 0)
       return cross(c.o - s.b, s.a - s.b) /
             2.0;
   // three part: (A, a) (a, b) (b, B)
   vector<Point> Sect =
        circleLineIntersection(c, s);
   return circleTriangleIntersectionArea(c,
         Segment(s.a, Sect[0])) +
         circleTriangleIntersectionArea(c,
              Segment(Sect[0], Sect[1])) +
          circleTriangleIntersectionArea(c,
              Segment(Sect[1], s.b));
// area of intersecion of circle(c.o, c.r)
    && simple polyson(p[])
// Tested : https://codeforces.com/gym
    /100204/problem/F - Little Mammoth
Tf circlePolyIntersectionArea(Circle c,
    Polygon p) {
   Tf res = 0;
   int n = p.size();
   for (int i = 0; i < n; ++i)</pre>
            \verb|circleTriangleIntersectionArea| (c
            , Segment(p[i], p[(i + 1) % n]))
   return abs(res);
// locates circle c2 relative to c1
// interior
            (d < R - r) ----> -2
// interior tangents (d = R - r) ----> -1
// concentric (d = 0)
               (R - r < d < R + r) \longrightarrow 0
// secants
// exterior tangents (d = R + r) ----> 1
                (d > R + r) ----> 2
// exterior
int circleCirclePosition(Circle c1, Circle
    c2) {
   Tf d = length(c1.o - c2.o);
   int in = dcmp(d - abs(c1.r - c2.r)), ex
        = dcmp(d - (c1.r + c2.r));
```

```
return in < 0 ? -2 : in == 0 ? -1 : ex
        == 0 ? 1 : ex > 0 ? 2 : 0;
// compute the intersection points between
    two circles c1 && c2
vector<Point> circleCircleIntersection(
    Circle c1, Circle c2) {
   static_assert(is_same<Tf, Ti>::value);
   vector<Point> ret;
   Tf d = length(c1.o - c2.o);
   if (dcmp(d) == 0) return ret;
   if (dcmp(c1.r + c2.r - d) < 0) return
        ret:
   if (dcmp(abs(c1.r - c2.r) - d) > 0)
        return ret;
   Point v = c2.o - c1.o;
   Tf co = (c1.r * c1.r + sqLength(v) - c2.
        r * c2.r) / (2 * c1.r * length(v));
   Tf si = sqrt(abs(1.0 - co * co));
   Point p1 = scale(rotatePrecise(v, co, -
        si), c1.r) + c1.o;
   Point p2 = scale(rotatePrecise(v, co, si
        ), c1.r) + c1.o;
   ret.push_back(p1);
   if (p1 != p2) ret.push_back(p2);
   return ret;
// intersection area between two circles c1,
Tf circleCircleIntersectionArea(Circle c1,
    Circle c2) {
   Point AB = c2.o - c1.o;
   Tf d = length(AB);
   if (d >= c1.r + c2.r) return 0;
   if (d + c1.r <= c2.r) return PI * c1.r *</pre>
         c1.r:
   if (d + c2.r <= c1.r) return PI * c2.r *</pre>
         c2.r:
   Tf alpha1 = acos((c1.r * c1.r + d * d -
        c2.r * c2.r) / (2.0 * c1.r * d));
   Tf alpha2 = acos((c2.r * c2.r + d * d -
        c1.r * c1.r) / (2.0 * c2.r * d));
   return c1.sector(2 * alpha1) + c2.sector
        (2 * alpha2);
// returns tangents from a point p to circle
vector<Point> pointCircleTangents(Point p,
    Circle c) {
   static_assert(is_same<Tf, Ti>::value);
   vector<Point> ret;
   Point u = c.o - p;
   Tf d = length(u);
   if (d < c.r)
   else if (dcmp(d - c.r) == 0) {
      ret = {rotate(u, PI / 2)};
   } else {
       Tf ang = asin(c.r / d);
       ret = {rotate(u, -ang), rotate(u,
           ang)};
   return ret;
// returns the points on tangents that
    touches the circle
vector<Point> pointCircleTangencyPoints(
    Point p, Circle c) {
   static_assert(is_same<Tf, Ti>::value);
   Point u = p - c.o;
```

Tf d = length(u);

return {};

else if (dcmp(d - c.r) == 0)

return {c.o + u};

if (d < c.r)

```
else {
       Tf ang = acos(c.r / d);
       u = u / length(u) * c.r;
       return {c.o + rotate(u, -ang), c.o +
             rotate(u, ang)};
// for two circles c1 && c2, returns two
    list of points a && b
// such that a[i] is on c1 && b[i] is c2 &&
    for every i
// Line(a[i], b[i]) is a tangent to both
    circles
// CAUTION: a[i] = b[i] in case they touch |
     -1 for c1 = c2
int circleCircleTangencyPoints(Circle c1,
    Circle c2, vector<Point> &a,
                            vector<Point> &b
   a.clear(), b.clear();
   int cnt = 0;
   if (dcmp(c1.r - c2.r) < 0) {
       swap(c1, c2);
       swap(a, b);
   Tf d2 = sqLength(c1.o - c2.o);
   Tf rdif = c1.r - c2.r, rsum = c1.r + c2.
   if (dcmp(d2 - rdif * rdif) < 0) return</pre>
   if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r)
        == 0) return -1;
   Tf base = angle(c2.o - c1.o);
   if (dcmp(d2 - rdif * rdif) == 0) {
       a.push_back(c1.point(base));
       b.push_back(c2.point(base));
       cnt++:
       return cnt;
   Tf ang = acos((c1.r - c2.r) / sqrt(d2));
   a.push_back(c1.point(base + ang));
   b.push_back(c2.point(base + ang));
   cnt++:
   a.push_back(c1.point(base - ang));
   b.push_back(c2.point(base - ang));
   if (dcmp(d2 - rsum * rsum) == 0) {
       a.push_back(c1.point(base));
       b.push_back(c2.point(PI + base));
       cnt++;
   } else if (dcmp(d2 - rsum * rsum) > 0) {
       Tf ang = acos((c1.r + c2.r) / sqrt(
            d2)):
       a.push_back(c1.point(base + ang));
       b.push_back(c2.point(PI + base + ang
           ));
       a.push_back(c1.point(base - ang));
       b.push_back(c2.point(PI + base - ang
           )):
       cnt++;
   return cnt;
```

#### 4.3 Closest Pair of Points

```
}
   auto it1 = s.lower_bound({pts[i].ss
        - d, pts[i].ff});
   auto it2 = s.upper_bound({pts[i].ss
        + d, pts[i].ff});
   for (auto it = it1; it != it2; ++it)
       int dx = pts[i].ff - it->ss;
       int dy = pts[i].ss - it->ff;
       best_dist = min(best_dist, 1LL *
            dx * dx + 1LL * dy * dy;
   s.insert({pts[i].ss, pts[i].ff});
return best_dist;
```

### 4.4 Linear

```
// returns true if point p is on segment s
bool onSegment(Point p, Segment s) {
   return dcmp(cross(s.a - p, s.b - p)) ==
        0 &&
          dcmp(dot(s.a - p, s.b - p)) <= 0;
// returns true if segment p && q touch or
    intersect
bool segmentsIntersect(Segment p, Segment q)
     {
   if (onSegment(p.a, q) || onSegment(p.b,
        q)) return true;
   if (onSegment(q.a, p) || onSegment(q.b,
        p)) return true;
   Ti c1 = cross(p.b - p.a, q.a - p.a);
   Ti c2 = cross(p.b - p.a, q.b - p.a);
   Ti c3 = cross(q.b - q.a, p.a - q.a);
   Ti c4 = cross(q.b - q.a, p.b - q.a);
   return dcmp(c1) * dcmp(c2) < 0 && dcmp(</pre>
        c3) * dcmp(c4) < 0;
bool linesParallel(Line p, Line q) {
   return dcmp(cross(p.b - p.a, q.b - q.a))
// lines are represented as a ray from a
    point: (point, vector)
// returns false if two lines (p, v) && (q,
    w) are parallel or collinear
// true otherwise, intersection point is
    stored at o via reference
bool lineLineIntersection(Point p, Point v,
    Point q, Point w, Point& o) {
   static_assert(is_same<Tf, Ti>::value);
   if (dcmp(cross(v, w)) == 0) return false
   Point u = p - q;
   o = p + v * (cross(w, u) / cross(v, w)); Tf yvalSegment(const Line &s, Tf x) {
   return true;
// returns false if two lines p && q are
    parallel or collinear
// true otherwise, intersection point is
    stored at o via reference
bool lineLineIntersection(Line p, Line q,
   return lineLineIntersection(p.a, p.b - p
        .a, q.a, q.b - q.a, o);
// returns the distance from point a to line
Tf distancePointLine(Point p, Line 1) {
   return abs(cross(1.b - 1.a, p - 1.a) /
        length(1.b - 1.a));
```

```
// returns the shortest distance from point | iter prev(iter it) { return it == st.begin()
    a to segment s
Tf distancePointSegment(Point p, Segment s)
   if (s.a == s.b) return length(p - s.a);
   Point v1 = s.b - s.a, v2 = p - s.a, v3 =
         p - s.b;
   if (dcmp(dot(v1, v2)) < 0)</pre>
       return length(v2);
   else if (dcmp(dot(v1, v3)) > 0)
       return length(v3);
       return abs(cross(v1, v2) / length(v1
           ));
// returns the shortest distance from
    segment p to segment q
Tf distanceSegmentSegment(Segment p, Segment
   if (segmentsIntersect(p, q)) return 0;
   Tf ans = distancePointSegment(p.a, q);
   ans = min(ans, distancePointSegment(p.b,
         q));
   ans = min(ans, distancePointSegment(q.a,
         p));
   ans = min(ans, distancePointSegment(q.b,
        p));
   return ans;
// returns the projection of point p on line
Point projectPointLine(Point p, Line 1) {
   static_assert(is_same<Tf, Ti>::value);
   Point v = 1.b - 1.a;
   return 1.a + v * ((Tf)dot(v, p - 1.a) /
        dot(v, v));
```

#### 4.5Pair of Intersecting segments using Line Sweep

Checking for the intersection of two

intersect ()

iterator iter;

segments is carried out by the

function, using an algorithm based on the

```
oriented area of the triangle.
The queue of segments is the global variable
     s, a set<event>. Iterators that
specify the position of each segment in the
    queue (for convenient removal of
segments from the queue) are stored in the
    global array where.
Two auxiliary functions prev() and next()
    are also introduced, which return
iterators to the previous and next elements
    (or end(), if one does not exist).
   if (dcmp(s.a.x - s.b.x) == 0) return s.a
   return s.a.y + (s.b.y - s.a.y) * (x - s.
        a.x) / (s.b.x - s.a.x);
struct SegCompare {
   bool operator()(const Segment &p, const
        Segment &q) const {
       Tf x = max(min(p.a.x, p.b.x), min(q.
           a.x, q.b.x));
       return dcmp(yvalSegment(p, x) -
           yvalSegment(q, x)) < 0;
multiset<Segment, SegCompare> st;
typedef multiset<Segment, SegCompare>::
```

```
? st.end() : --it; }
iter next(iter it) { return it == st.end() ?
      st.end() : ++it; }
struct Event {
   Tf x;
    int tp, id;
    Event(Ti x, int tp, int id) : x(x), tp(
        tp), id(id) {}
    bool operator<(const Event &p) const {</pre>
       if (dcmp(x - p.x)) return x < p.x;</pre>
       return tp > p.tp;
};
bool anyIntersection(const vector<Segment> &
    using Linear::segmentsIntersect;
    vector<Event> ev;
    for (int i = 0; i < (int)v.size(); ++i)</pre>
       ev.push_back(Event(min(v[i].a.x, v[i
            ].b.x), +1, i));
       ev.push_back(Event(max(v[i].a.x, v[i
            ].b.x), -1, i));
   sort(ev.begin(), ev.end());
    st.clear();
    vector<iter> where(v.size());
    for (auto &cur : ev) {
       int id = cur.id;
       if (cur.tp == 1) {
           iter nxt = st.lower_bound(v[id]);
           iter pre = prev(nxt);
if (pre != st.end() &&
                segmentsIntersect(*pre, v[id
                ])) return true;
           if (nxt != st.end() &&
                segmentsIntersect(*nxt, v[id
                ])) return true;
           where[id] = st.insert(nxt, v[id])
       } else {
           iter nxt = next(where[id]);
           iter pre = prev(where[id]);
           if (pre != st.end() && nxt != st.
                end() &&
               segmentsIntersect(*pre, *nxt)
                   )
               return true;
           st.erase(where[id]);
       }
   return false;
```

```
4.6 Polygons
// returns the signed area of polygon p of n
      vertices
using Polygon = vector<Point>;
Tf signedPolygonArea(Polygon p) {
   Tf ret = 0;
   for (int i = 0; i < (int)p.size() - 1; i</pre>
        ++)
       ret += cross(p[i] - p[0], p[i + 1] -
             p[0]);
   return ret / 2;
// given a polygon p of n vertices,
    generates the convex hull in ch
   in CCW && returns the number of vertices
    in the convex hull
int convexHull(Polygon p, Polygon &ch) {
   sort(p.begin(), p.end());
   int n = p.size();
   ch.resize(n + n);
   int m = 0; // preparing lower hull
   for (int i = 0; i < n; i++) {</pre>
       while (m > 1 &&
```

```
dcmp(cross(ch[m - 1] - ch[m -
                   2], p[i] - ch[m - 1])) <=
       ch[m++] = p[i];
   int k = m; // preparing upper hull
   for (int i = n - 2; i >= 0; i--) {
       while (m > k &&
             dcmp(cross(ch[m - 1] - ch[m -
                  2], p[i] - ch[m - 2])) <=
           m--;
       ch[m++] = p[i];
   if (n > 1) m--;
   ch.resize(m);
   return m;
// for a point o and polygon p returns:
// -1 if o is strictly inside p
// 0 if o is on a segment of p
// 1 if o is strictly outside p
// computes via winding numbers
int pointInPolygon(Point o, Polygon p) {
   int wn = 0, n = p.size();
   for (int i = 0; i < n; i++) {</pre>
       int j = (i + 1) \% n;
       if (onSegment(o, Segment(p[i], p[j])
            ) || o == p[i]) return 0;
       int k = dcmp(cross(p[j] - p[i], o -
            p[i]));
       int d1 = dcmp(p[i].y - o.y);
int d2 = dcmp(p[j].y - o.y);
       if (k > 0 && d1 <= 0 && d2 > 0) wn
       if (k < 0 && d2 <= 0 && d1 > 0) wn
   return wn ? -1 : 1;
// returns the longest line segment of 1
    that is inside or on the
// simply polygon p. O(n lg n). TESTED:
    TIMUS 1955
Tf longestSegInPoly(Line 1, const Polygon &p
    ) {
   int n = p.size();
   vector<pair<Tf, int>> ev;
   for (int i = 0; i < n; ++i) {</pre>
       Point a = p[i], b = p[(i + 1) \% n],
            z = p[(i - 1 + n) \% n];
       int ora = orient(1.a, 1.b, a), orb =
             orient(1.a, 1.b, b),
           orz = orient(l.a, l.b, z);
       if (!ora) {
           Tf d = dot(a - 1.a, 1.b - 1.a);
           if (orz && orb) {
              if (orz != orb) ev.
                   emplace_back(d, 0);
           } else if (orz)
              ev.emplace_back(d, orz);
           else if (orb)
              ev.emplace_back(d, orb);
       } else if (ora == -orb) {
           Point ins;
           lineLineIntersection(1, Line(a, b
               ). ins):
           ev.emplace_back(dot(ins - 1.a, 1.
               b - 1.a). 0:
       }
   }
   sort(ev.begin(), ev.end());
   Tf ret = 0, cur = 0, pre = 0;
   bool active = false;
   int sign = 0;
   for (auto &qq : ev) {
       int tp = qq.second;
       Tf d = qq.first;
       if (sign) {
           cur += d - pre;
```

```
ret = max(ret, cur);
           if (tp != sign) active = !active;
           sign = 0;
       } else {
           if (active) cur += d - pre, ret =
                max(ret, cur);
           if (tp == 0)
              active = !active;
           else
              sign = tp;
       pre = d;
       if (!active) cur = 0;
   ret /= length(l.b - l.a);
   return ret;
// EVERYTHING AFTER THIS ONLY FOR CONVEX
    POLYGON
/// Tested on Kattis::fenceortho
void rotatingCalipersGetRectangle(Point *p,
    int n, Tf &area, Tf &perimeter) {
   static_assert(is_same<Tf, Ti>::value);
   p[n] = p[0];
   int 1 = 1, r = 1, j = 1;
   area = perimeter = 1e100;
   for (int i = 0; i < n; i++) {</pre>
       Point v = (p[i + 1] - p[i]) / length
            (p[i + 1] - p[i]);
       while (dcmp(dot(v, p[r % n] - p[i])
            - dot(v, p[(r + 1) \% n] - p[i]))
             <
             0)
           r++;
       while (j < r || dcmp(cross(v, p[j % 
            n] - p[i]) -
                           cross(v, p[(j +
1) % n] - p[i
                                ])) < 0)
           j++;
       while (1 < j \mid \mid dcmp(dot(v, p[1 % n]
             - p[i]) -
                           dot(v, p[(1 + 1)
                               % n] - p[i]))
                                > 0)
       Tf w = dot(v, p[r \% n] - p[i]) - dot
            (v, p[1 % n] - p[i]);
       Tf h = distancePointLine(p[j % n],
           Line(p[i], p[i + 1]));
       area = min(area, w * h);
       perimeter = min(perimeter, 2 * w + 2)
             * h);
   }
// returns the left side of polygon u after
    cutting it by ray a->b
Polygon cutPolygon(Polygon u, Point a, Point
   Polygon ret;
   int n = u.size();
   for (int i = 0; i < n; i++) {</pre>
       Point c = u[i], d = u[(i + 1) \% n];
       if (dcmp(cross(b - a, c - a)) >= 0)
            ret.push_back(c);
       if (dcmp(cross(b - a, d - c)) != 0)
            {
           Point t:
           lineLineIntersection(a, b - a, c,
                 d - c, t);
           if (onSegment(t, Segment(c, d)))
                ret.push_back(t);
       }
   }
   return ret;
// returns true if point p is in or on
    triangle abc
```

```
bool pointInTriangle(Point a, Point b, Point
      c, Point p) {
    return dcmp(cross(b - a, p - a)) >= 0 &&
         dcmp(cross(c - b, p - b)) >= 0 &&
          dcmp(cross(a - c, p - c)) >= 0;
// Tested : https://www.spoj.com/problems/
    INOROUT
// pt must be in ccw order with no three
     collinear points
// returns inside = -1, on = 0, outside = 1
int pointInConvexPolygon(const Polygon &pt,
    Point p) {
    int n = pt.size();
   assert(n >= 3);
    int lo = 1, hi = n - 1;
    while (hi - lo > 1) {
       int mid = (lo + hi) / 2;
       if (dcmp(cross(pt[mid] - pt[0], p -
            pt[0])) > 0)
           lo = mid;
       else
           hi = mid;
    bool in = pointInTriangle(pt[0], pt[lo],
         pt[hi], p);
    if (!in) return 1;
    if (dcmp(cross(pt[lo] - pt[lo - 1], p -
        pt[lo - 1])) == 0) return 0;
    if (dcmp(cross(pt[hi] - pt[lo], p - pt[
        lo])) == 0) return 0;
    if (dcmp(cross(pt[hi] - pt[(hi + 1) % n
        ], p - pt[(hi + 1) \% n])) == 0)
       return 0:
    return -1:
// Extreme Point for a direction is the
     farthest point in that direction
// O'Rourke, page 270, http://crtl-i.com/PDF
     /comp_c.pdf
// also https://codeforces.com/blog/entry
     /48868
// poly is a convex polygon, sorted in CCW,
     doesn't contain redundant points
// u is the direction for extremeness
int extremePoint(const Polygon &poly, Point
     u = Point(0, 1)) {
    int n = (int)poly.size();
    int a = 0, b = n;
    while (b - a > 1) {
       int c = (a + b) / 2;
       if (dcmp(dot(poly[c] - poly[(c + 1)
            % n], u)) >= 0 &&
           dcmp(dot(poly[c] - poly[(c - 1 +
               n) % n], u)) >= 0) {
           return c;
       bool a_up = dcmp(dot(poly[(a + 1) %
            n] - poly[a], u)) >= 0;
       bool c_up = dcmp(dot(poly[(c + 1) %
           n] - poly[c], u)) >= 0;
       bool a_above_c = dcmp(dot(poly[a] -
            poly[c], u)) > 0;
       if (a_up && !c_up)
           b = c:
       else if (!a_up && c_up)
           a = c;
       else if (a_up && c_up) {
           if (a_above_c)
              b = c;
           else
              a = c;
       } else {
           if (!a_above_c)
           else
               a = c;
```

```
}
   if (dcmp(dot(poly[a] - poly[(a + 1) % n
        ], u)) > 0 &&
       dcmp(dot(poly[a] - poly[(a - 1 + n)
           % n], u)) > 0)
       return a;
   return b % n;
// For a convex polygon p and a line 1,
    returns a list of segments
// of p that are touch or intersect line 1.
// the i'th segment is considered (p[i], p[(
    i + 1) modulo |p|])
// #1 If a segment is collinear with the
    line, only that is returned
// #2 Else if 1 goes through i'th point, the
     i'th segment is added
// If there are 2 or more such collinear
    segments for #1,
// any of them (only one, not all) should be
     returned (not tested)
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(const
     Polygon &p, Line 1) {
   assert((int)p.size() >= 3);
   assert(1.a != 1.b);
   int n = p.size();
   vector<int> ret;
   Point v = 1.b - 1.a;
   int lf = extremePoint(p, rotate90(v));
   int rt = extremePoint(p, rotate90(v) *
        Ti(-1));
   int olf = orient(l.a, l.b, p[lf]);
   int ort = orient(1.a, 1.b, p[rt]);
   if (!olf || !ort) {
       int idx = (!olf ? lf : rt);
       if (orient(l.a, l.b, p[(idx - 1 + n)
             n == 0)
           ret.push_back((idx - 1 + n) % n);
       else
          ret.push_back(idx);
       return ret;
   if (olf == ort) return ret;
   for (int i = 0; i < 2; ++i) {</pre>
       int lo = i ? rt : lf;
       int hi = i ? lf : rt;
       int olo = i ? ort : olf;
       while (true) {
           int gap = (hi - lo + n) \% n;
           if (gap < 2) break;</pre>
           int mid = (lo + gap / 2) % n;
           int omid = orient(1.a, 1.b, p[mid
               1):
           if (!omid) {
              lo = mid;
              break;
           if (omid == olo)
              lo = mid;
           else
              hi = mid;
       ret.push_back(lo);
   return ret;
// Tested : https://toph.co/p/cover-the-
// Calculate [ACW, CW] tangent pair from an
    external point
constexpr int CW = -1, ACW = 1;
bool isGood(Point u, Point v, Point Q, int
    dir) {
                                              double d; //ax + by + cz = d
```

```
return orient(Q, u, v) != -dir;
Point better(Point u, Point v, Point Q, int
   return orient(Q, u, v) == dir ? u : v;
Point pointPolyTangent(const Polygon &pt,
    Point Q, int dir, int lo, int hi) {
   while (hi - lo > 1) {
       int mid = (lo + hi) / 2;
       bool pvs = isGood(pt[mid], pt[mid -
           1], Q, dir);
       bool nxt = isGood(pt[mid], pt[mid +
           1], Q, dir);
       if (pvs && nxt) return pt[mid];
       if (!(pvs || nxt)) {
          Point p1 = pointPolyTangent(pt, Q
               , dir, mid + 1, hi);
          Point p2 = pointPolyTangent(pt, Q
              , dir, lo, mid - 1);
          return better(p1, p2, Q, dir);
       if (!pvs) {
          if (orient(Q, pt[mid], pt[lo]) ==
                dir)
              hi = mid - 1;
          else if (better(pt[lo], pt[hi], Q
               , dir) == pt[lo])
              hi = mid - 1;
          else
              lo = mid + 1;
       if (!nxt) {
          if (orient(Q, pt[mid], pt[lo]) ==
               dir)
              lo = mid + 1;
          else if (better(pt[lo], pt[hi], Q
              , dir) == pt[lo])
hi = mid - 1;
              lo = mid + 1;
   Point ret = pt[lo];
   for (int i = lo + 1; i <= hi; i++) ret = };
         better(ret, pt[i], Q, dir);
   return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(const
    Polygon &pt, Point Q) {
    int n = pt.size();
   Point cw_tan = pointPolyTangent(pt, Q,
        CW, 0, n - 1);
   return make_pair(acw_tan, cw_tan);
```

### Three Dimensional

```
point get_perp(point p){ // returns a random
     perpendicular line to the vector p
   assert(sgn(norm(p)));
   point ret = point(-p.y, p.x, 0);
   if(sgn(norm(ret))) return ret;
   ret = point(0, -p.z, p.y);
   if(sgn(norm(ret))) return ret;
   assert(false)
struct plane{ // Caution: directed plane,
    directed on the direction of (p2 \times p3)
point n; // {a, b, c}
```

```
// d = n . p [ where p is any point on the
     plane ]
plane(){;}
 plane(point _n, double _d){
       n = _n;
       d = _d;
 plane(point p1, point p2, point p3){
 n = crsp(p2 - p1, p3 - p1);
  if(norm(n) < eps) {assert(false);} //</pre>
      doesn't define a plance
  d = dotp(p1, n);
    //Preserves the direction
 point get_p1(){ return univ(n) * d / norm(n
 point get_p2(){ return get_p1() + get_perp(
 point get_p3(){ return crsp(n, get_p2() -
     get_p1()) + get_p1();}
 int get_side(point p){ return sgn(dotp(n, p)
      ) - d);} ///OK
double sgn_dist(point p) {return (dotp(n, p
     ) - d) / norm(n);}
 double dist(point p) {return fabs(sgn_dist(
     p));}
 point project(point p){ return p - sgn_dist
      (p) * univ(n);}
 point reflect(point p){ return p - 2 *
     sgn_dist(p) * univ(n);}
    ///OK
   point get_coords(point p){ // "2-D"-fies
         the plane. All points on this
        plane have z = 0
       point 0 = get_p1();
       point ox = univ(get_p2() - 0);
       point oy = univ(get_p3() - 0);
       point oz = univ(n);
       p = p - 0;
       return {dotp(p, ox), dotp(p, oy),
            dotp(p, oz)};
plane translate(plane p, point t) {return {p
     .n, p.d + dotp(p.n, t)};}
plane shiftUp(plane p, double d) {return {p.
    n, p.d + d * norm(p.n)};}
point projection(point p, point st, point ed
    ) { return dotp(ed - st, p - st) / norm
     (ed - st) * univ(ed - st) + st;} //OK
point extend(point st, point ed, double len)
      { return ed + univ(ed-st) * len;} //OK
point rtt(point axis, point p, double theta)
   axis = univ(axis);
   return p * cos(theta) + sin(theta) *
        crsp(axis, p) + axis * (1-cos(theta))
        )) * dotp(axis, p);
} //OK
point segmentProjection(point p, point st,
     point ed)
    double d = dotp(p - st, ed - st) / norm(
        ed - st);
    if(d < 0) return st;</pre>
    if(d > norm(ed - st) + eps) return ed;
   return st + univ(ed - st) * d;
} //OK
double distPointSegment(point p, point st,
     point ed) {return norm(p -
     segmentProjection(p, st, ed)); } //OK
double distPointLine( point P, point st,
     point ed) { return norm( projection(P,
```

```
st, ed) - P); } //OK
double pointPlanedist(plane P, point q){
    return fabs(dotp(P.n, q) - P.d) / norm(
    P.n);}
double pointPlanedist(point p1, point p2,
    point p3, point q){ return
    pointPlanedist(plane(p1,p2,p3), q); }
point reflection(point p, point st, point ed
   point proj = projection(p, st, ed);
   if(p != proj) return extend(p, proj,
       norm(p - proj));
   return proj;
} //OK
bool coplanar(point p1, point p2, point p3,
   p2 = p2-p1, p3 = p3-p1, q = q-p1;
   if( fabs( dotp(q, crsp(p2, p3)) ) < eps</pre>
        ) return true;
   return false;
int linePlaneIntersection(point u, point v,
    point 1, point m, point r, point &x){
       -> 1, m, r defines the plane
       -> u. v defines the line
       -> returns 0 when does not intersect
       -> returns 1 when there exists one
           unique common point
       -> returns -1 when there exists
            infinite number of common point
   \mathtt{assert(1 != m \&\& m != r \&\& 1 != r \&\& u}
        ! = v):
        , r, v)) return -1;
   1 = 1 - m;
   r = r - m;
   u = u - m;
   v = v - m;
   point C = crsp(1, r);
   double denom = dotp(v - u, C);
   if(fabs(denom) < eps) return 0;</pre>
   double alpha = -dotp(C, u) / denom;
   x = u + (v - u) * alpha + m;
   return 1;
double angle(point u, point v) { return acos |}
    ( max(-1.0, min(1.0, dotp(u, v) / (norm
    (u) * norm(v)))));}
struct line3d{ //directed
   point d, o; // dir = direction, o =
        online point
   line3d(point p, point q){
       d = q - p;
       o = p;
       assert(sgn(norm(d)));
   line3d(plane p1, plane p2){
       d = crsp(p1.n, p2.n);
       o = (crsp(p2.n*p1.d - p1.n*p2.d, d))
            /sq(d);
   point get_p1(){return o;}
   point get_p2(){return o + d;}
   double dist(point p){ return norm(crsp(d
        , p - o)) / norm(d);};
```

```
point project(point p){ return
                                                  projection(p, o, o + d); }
                                             point reflect(point p) {return
                                                  reflection(p, o, o + d);}
                                          line3d perpThrough(plane p, point o){return
                                              line3d(o, o + p.n);
                                          plane perpThrough(line3d 1, point o){return
                                              plane(l.d, dotp(l.d, o));}
                                          double dist(line3d l1, line3d l2) {
                                             point n = crsp(11.d, 12.d);
                                              if (!sgn(norm(n))) return 11.dist(12.o);
                                             return abs(dotp(12.o-11.o, n))/norm(n);
                                          point closestOnL1(line3d 11, line3d 12) {
                                             point n2 = crsp(12.d, crsp(11.d, 12.d));
                                             return 11.0 + (11.d * (dotp(12.o-11.o,
                                                  n2))) / dotp(11.d,n2);
                                          double angle(plane p1, plane p2){return
                                              angle(p1.n, p2.n);}
                                          bool isparallel(plane p1, plane p2){return !
                                              sgn(norm(crsp(p1.n, p2.n)));}
                                          bool isperp(plane p1, plane p2) {return !sgn
                                               (dotp(p1.n, p2.n));}
                                          double angle(line3d 11, line3d 12){return
                                              angle(11.d, 12.d);}
                                          bool isparallel(line3d 11, line3d 12){return
                                                !sgn(norm(crsp(11.d, 12.d)));}
                                          bool isperp(line3d 11, line3d 12) {return !
                                              sgn(dotp(11.d, 12.d));}
                                          double angle(plane p, line3d 1) {return pi/2
                                                - angle(p.n, l.d);}
if(coplanar(1, m, r, u) && coplanar(1, m | bool isParallel(plane p, line3d 1) {return !
                                              sgn(dotp(p.n, 1.d));}
                                          bool isPerpendicular(plane p, line3d 1) {
                                              return !sgn(norm(crsp(p.n, 1.d)));}
                                          point vector_area2(vector <point> &poly){
                                             point S = \{0, 0, 0\};
                                             for(int i = 0; i < (int) poly.size(); i</pre>
                                                 S = S + crsp(poly[i], poly[(i + 1)
                                                     % poly.size()]);
                                             return S;
                                          double area(vector < point > &poly){ // All
                                              points must be co-planer
                                              return norm(vector_area2(poly)) * 0.5;
                                             Polyhedrons
                                          bool operator <(point p, point q) { ///OK</pre>
                                             return tie(p.x, p.y, p.z) < tie(q.x, q.y</pre>
                                                  , q.z);
                                          struct edge {
                                             int v:
                                             bool same; // = is the common edge in
                                                  the same order?
                                          // Given a series of faces (lists of points) int main()
                                              , reverse some of them
                                          // so that their orientations are consistent
                                                [ every face then will point in the
                                               same direction, inside / outside ]
                                          void reorient(vector< vector<point> > &fs) { //
                                             int n = fs.size();
                                              // Find the common edges and create the
```

resulting graph

```
vector< vector<edge> > g(n);
   map<pair<point,point>, int> es;
   for (int u = 0; u < n; u++) {</pre>
       for (int i = 0, m = fs[u].size(); i
           < m; i++) {
           point a = fs[u][i], b = fs[u][(i
               +1)%m];
           // Let look at edge [AB]
           if (es.count({a,b})) { // seen in
                same order
              int v = es[{a,b}];
              g[u].push_back({v,true});
              g[v].push_back({u,true});
           else if (es.count({b,a})) { //
               seen in different order
                  int v = es[{b,a}];
                  g[u].push_back({v,false});
                  g[v].push_back({u,false});
           else es[{a,b}] = u;
   vector<bool> vis(n,false), flip(n);
   flip[0] = false;
   queue<int> q;
   q.push(0);
   while (!q.empty()) {
       int u = q.front();
       q.pop();
       for (edge e : g[u]) {
           if (!vis[e.v]) {
              vis[e.v] = true;
              // If the edge was in the
                   same order,
              // exactly one of the two
                   should be flipped
              flip[e.v] = (flip[u] ^ e.same
                   ):
              q.push(e.v);
       }
   for (int u = 0; u < n; u++)
       if (flip[u])
          reverse(fs[u].begin(), fs[u].end
               ());
double volume(vector< vector<point> > fs) {
   double vol6 = 0.0;
   for (vector<point> f : fs)
       vol6 += dotp(vector_area2(f), f[0]);
   return abs(vol6) / 6.0;
   Spherical Co-ordinate System
point sph(double r, double lat, double lon)
    { // lat, lon in degrees
   lat *= pi/180, lon *= pi/180;
   return {r*cos(lat)*cos(lon), r*cos(lat)*
        sin(lon), r*sin(lat)};
double greatCircleDist(point o, double r,
    point a, point b) {
   return r * angle(a-o, b-o);
   plane p = {point(0, 0, 10), point(0, 1,
       10), point(1, 0, 10)};
     plane q = \{p.get_p1(), p.get_p2(), p.
    get_p3()};
   double d = dotp(p.get_p2() - p.get_p1(),
         p.get_p3() - p.get_p1());
```

```
D(eq(d, 0))
  return 0;
}
```

# 5 Graph

### 5.1 2 Satisfiability

```
2-Sat Note: Assign true or false values to
      n variables in order to satisfy
 a system of constraints on pairs of
      variables.
 E.g: (x1 or !x2) and (x2 or x3) and (!x3
      or !x3)
 x1 = true
 x2 = true
 x3 = false
 is a solution to make the above formula
 MAX must be equal to the maximum number of
       variables.
 n passed in init() is the number of
      variables.
 !a is represented as neg(a).
 example xor:
  lalbl
  |0|0| x or(a,b)
  10111
 11101
  |1|1| x or(!a, !b)
 do OR of negation of values of variables
      for each undesired situation
 to make it impossible.
struct two_sat {
 int n. id:
 vector<int> g[2 * MAX], rg[2 * MAX], order
       . st:
 bool state[2 * MAX], vis[2 * MAX];
 int scc[2 * MAX];
 void init(int _n) {
   n = _n;
   for (int i = 0; i <= 2 * n; i++) {</pre>
     g[i].clear(), rg[i].clear();
     state[i] = vis[i] = false;
     scc[i] = -1;
   st.clear(), order.clear();
 }
 void add_edge(int u, int v) {
   g[u].pb(v);
   rg[v].pb(u);
 void OR(int u, int v) {
   add_edge(neg(u), v);
   add_edge(neg(v), u);
 void XOR(int u, int v) {
   OR(u, v);
   OR(neg(u), neg(v));
 void ForceTrue(int u) {
   add_edge(neg(u), u);
 void ForceFalse(int u) {
   add_edge(u, neg(u));
 void imply(int u, int v) {
   OR(neg(u), v);
```

```
int neg(int u) {
   if (u <= n) return u + n;</pre>
   return u - n;
 void dfs(int u, vii g[], bool topsort) {
   vis[u] = true;
   for (int v : g[u]) {
     if (!vis[v]) dfs(v, g, topsort);
   if (topsort) st.pb(u);
   else scc[u] = id, order.pb(u);
 void build_scc() {
   for (int i = 1; i <= 2 * n; i++) {</pre>
     if (!vis[i]) dfs(i, g, true);
   reverse(st.begin(), st.end());
   fill(vis, vis + 2 * n + 1, false);
   for (int u : st) {
     if (!vis[u]) id++, dfs(u, rg, false);
 bool solve() {
   build_scc();
   for (int i = 1; i <= n; i++) {</pre>
     if (scc[i] == scc[i + n]) return false
   for (int i = (int)order.size() - 1; i >=
         0; i--) {
     int u = order[i];
     if (state[neg(u)] == false) state[u] =
   return true;
} solver;
```

#### 5.2 Block Cut Tree

```
bool ap[MAX];
int id[MAX], koyta[MAX];
int d[MAX], low[MAX];
bool vis[MAX];
vii g[MAX], tree[MAX];
int d_t;
stack<int>st;
vector<vector<int>>comp;
void articulation(int u, int p) {
 vis[u] = true;
  d[u] = low[u] = ++d_t;
  int child = 0; st.push(u);
 for (int v : g[u]) {
   if (v == p) continue;
   if (!vis[v]) {
     child++;
     articulation(v, u);
     low[u] = min(low[u], low[v]);
     if (p == -1 && child > 1) ap[u] = true
     if (low[v] >= d[u]) {
       if (p != -1) ap[u] = true;
       comp.pb({u}); int top;
       do {
         top = st.top(); st.pop();
         comp.back().pb(top);
       } while (top != v);
   } else low[u] = min(low[u], d[v]);
 }
int node = 0;
void make_tree(int n) {
 for (int i = 1; i <= n; i++) {</pre>
   if (ap[i]) id[i] = ++node;
 for (int i = 0; i < comp.size(); i++) {</pre>
   ++node:
   int cnt = 0;
   for (int u : comp[i]) {
```

```
if (ap[u]) tree[node].pb(id[u]), tree[
          id[u]].pb(node), koyta[id[u]] =
     1;
    else id[u] = node, cnt++;
}
    koyta[node] = cnt;
}
```

### 5.3 Bridge Tree

```
vector<int> tree[MAX];
bool vis[MAX]:
int d[MAX], low[MAX];
int id[MAX];
int d_t;
struct edge {
 int v, rev;
  edge() {}
 edge(int v, int rev) : v(v), rev(rev) {}
vector<edge>g[MAX];
vector<bool>is_bridge[MAX];
queue<int>q[MAX];
int comp = 1;
void add_edge(int u, int v) {
 edge _u = edge(v, g[v].size());
  edge _v = edge(u, g[u].size());
 g[u].pb(_u);
 g[v].pb(_v);
  is_bridge[u].pb(false);
 is_bridge[v].pb(false);
void bridge(int u, int p) {
 vis[u] = true;
 d[u] = low[u] = ++d_t;
  for (int i = 0; i < g[u].size(); i++) {</pre>
   edge e = g[u][i]; int v = e.v;
   if (v == p) continue;
   if (!vis[v]) {
     bridge(v, u);
     low[u] = min(low[v], low[u]);
     if (low[v] > d[u]) {
       is_bridge[u][i] = true;
       is_bridge[v][e.rev] = true;
   } else low[u] = min(low[u], d[v]);
void make_tree(int node) {
 int cur = comp; q[cur].push(node);
 vis[node] = true; id[node] = cur;
 while (!q[cur].empty()) {
   int u = q[cur].front(); q[cur].pop();
   for (int i = 0; i < g[u].size(); i++) {</pre>
     edge e = g[u][i]; int v = e.v;
     if (vis[v]) continue;
     if (is_bridge[u][i]) {
       comp++;
       tree[cur].pb(comp);
       tree[comp].pb(cur);
       make_tree(v);
     } else {
       q[cur].push(v);
       vis[v] = true; id[v] = cur;
 }
```

### 5.4 Centroid Decomposition

```
// problem: calculate the sum of number of
    distinct colors in the path between any
    two pair of nodes
//centroid decomposition (res[i] contains
    the sum of numbers of distinct colors
    in all paths from i)
set<int>g[MAX];
int col[MAX], child[MAX], used[18][MAX];
```

```
11 ans[MAX], res[MAX];
int sz = 0, uniq = 0, n;
bool vis[MAX];
void dfs(int u, int p) {
 sz++; child[u] = 1;
 for (auto v : g[u]) {
   if (v != p) {
     dfs(v, u);
     child[u] += child[v];
 }
int get_centroid(int u, int p) {
 for (auto v : g[u]) {
   if (v != p && child[v] > sz / 2) return
        get_centroid(v, u);
 }
 return u;
void add(int u, int p, int depth, int
    centroid) {
 bool check = false; child[u] = 1;
 if (!vis[col[u]]) {
   uniq++; check = true;
   vis[col[u]] = true;
 ans[centroid] += uniq;
 for (auto v : g[u]) {
   if (v != p) {
     add(v, u, depth, centroid);
     child[u] += child[v];
 }
 if (check) {
   used[depth][col[u]] += child[u];
   vis[col[u]] = false;
void del(int u, int p, int depth, int
    centroid) {
  bool check = false;
 if (!vis[col[u]]) {
   used[depth][col[u]] -= child[u];
   vis[col[u]] = true; check = true;
 ans[centroid] -= uniq;
 for (auto v : g[u]) {
   if (v != p) del(v, u, depth, centroid);
 child[u] = 0;
 if (check) uniq--; vis[col[u]] = false;
void solve(int u, int p, int depth, int
    centroid) {
 ans[u] += (ans[p] + child[centroid] - used
      [depth][col[u]]);
 res[u] += ans[u];
  int temp = used[depth][col[u]];
 used[depth][col[u]] = child[centroid];
 for (auto v : g[u]) {
   if (v != p) solve(v, u, depth, centroid)
 }
 ans[u] = 0;
 used[depth][col[u]] = temp;
void reset_col(int u, int p, int depth) {
  used[depth][col[u]] = 0;
 for (auto v : g[u]) {
   if (v != p) reset_col(v, u, depth);
void decompose(int u, int depth) {
 sz = 0;
 uniq = 0;
 dfs(u, -1);
 int centroid = get_centroid(u, -1);
 reset_col(centroid, -1, depth);
 add(centroid, -1, depth, centroid); ///get
       ans for centroid and get the number
      of paths where each color is used
 res[centroid] += ans[centroid];
```

```
uniq++;
 vis[col[centroid]] = true;
 for (auto v : g[centroid]) {
   child[centroid] -= child[v];
   ///remove all contribution of the
        subtree of v
   del(v, centroid, depth, centroid);
   used[depth][col[centroid]] = child[
        centroid];
   solve(v, centroid, depth, centroid);
   ///add back the contribution of the
        subtree of v
   add(v, centroid, depth, centroid);
   child[centroid] += child[v];
  vis[col[centroid]] = false;
 for (auto it = g[centroid].begin(); it !=
      g[centroid].end(); it++) {
   g[*it].erase(centroid);
   decompose(*it, depth + 1);
int arr[MAX];
int main() {
 fastio;
 cin >> n:
 for (int i = 1; i <= n; i++) cin >> col[i
      ];
 for (int i = 0; i < n - 1; i++) {</pre>
   int u, v;
   cin >> u >> v:
   g[u].insert(v); g[v].insert(u);
 decompose(1, 0);
 for (int i = 1; i <= n; i++) cout << res[i //create new src and sink
      ] << "\n";
```

### 5.5 Dinic Max-Flow

```
int src, sink;
int dis[MAX], q[MAX], work[MAX];
int n, m, nodes;
struct edge {
 int v, rev, cap, flow;
 edge() {}
 edge(int v, int rev, int cap) : v(v), rev(|int tree[4 * MAX];
      rev), cap(cap), flow(0) {}
vector<edge>g[MAX];
void add_edge(int u, int v, int cap, int rev
     = 0) {
 edge _u = edge(v, g[v].size(), cap);
 edge _v = edge(u, g[u].size(), rev);
 g[u].pb(_u);
 g[v].pb(_v);
bool dinic bfs(int s) {
 fill(dis, dis + nodes, -1);
 dis[s] = 0;
 int id = 0:
 q[id++] = s;
 for (int i = 0; i < id; i++) {</pre>
   int u = q[i];
   for (int j = 0; j < g[u].size(); j++) {</pre>
     edge &e = g[u][j];
     if (dis[e.v] == -1 \&\& e.flow < e.cap)
       dis[e.v] = dis[u] + 1;
       q[id++] = e.v;
 return dis[sink] >= 0;
int dinic_dfs(int u, int f) {
 if (u == sink) return f;
```

```
for (int &i = work[u]; i < g[u].size(); i</pre>
      ++) {
   edge &e = g[u][i];
   if (e.cap <= e.flow) continue;</pre>
   if (dis[e.v] == dis[u] + 1) {
     int flow = dinic_dfs(e.v, min(f, e.cap
           - e.flow));
     if (flow) {
       e.flow += flow;
       g[e.v][e.rev].flow -= flow;
       return flow;
 }
 return 0;
int max_flow(int _src, int _sink) {
 src = _src;
 sink = _sink;
 int ret = 0;
 while (dinic_bfs(src)) {
   fill(work, work + nodes, 0);
   while (int flow = dinic_dfs(src, INT_MAX
        )) {
     ret += flow;
 return ret;
```

### 5.6 Flow with Lower Bound

```
///flow with demand(lower bound) only for
    DAG
//add_edge(new src, u, sum(in_demand[u]))
//add_edge(u, new sink, sum(out_demand[u]))
//add_edge(old sink, old src, inf)
// if (sum of lower bound == flow) then
    demand satisfied
//flow in every edge i = demand[i] + e.flow
```

# Heavy Light Decomposi-

```
int arr[MAX], n;
vector<int> parent, depth, heavy, head, pos;
int cur_pos, sub[MAX];
vii g[MAX];
void update(int now, int L, int R, int idx,
     int val) {
  if (L == R) {
   tree[now] = val;
   return;
  int mid = (L + R) / 2:
  if (idx <= mid) update(now << 1, L, mid,</pre>
      idx, val);
  else update( (now << 1) | 1, mid + 1, R,</pre>
      idx, val);
  tree[now] = tree[now << 1] + tree[(now <<</pre>
      1) | 1];
11 segtree_query(int now, int L, int R, int
     i, int j) {
  if (R < i || L > j) return 0;
 if (L >= i && R <= j) return tree[now];</pre>
  int mid = (L + R) / 2;
  return segtree_query(now << 1, L, mid, i,</pre>
       j) + segtree_query((now << 1) | 1,</pre>
      mid + 1, R, i, j);
int dfs(int u) {
 sub[u] = 1;
  int mx_size = 0;
  for (int v : g[u]) {
   if (v != parent[u]) {
     parent[v] = u, depth[v] = depth[u] +
          1:
      int v_size = dfs(v);
      sub[u] += v_size;
      if (v_size > mx_size) {
```

```
mx_size = v_size;
       heavy[u] = v;
   }
 }
 return sub[u];
void decompose(int u, int h) {
  head[u] = h, pos[u] = cur_pos++;
  if (heavy[u] != -1) decompose(heavy[u], h) //cycle decomposition of the k-th root of p
  for (int v : g[u]) {
   if (v != parent[u] && v != heavy[u])
        decompose(v, v);
void init(int n) {
 parent = vector<int>(n, -1);
  depth = vector<int>(n);
  heavy = vector<int>(n, -1);
 head = vector<int>(n);
  pos = vector<int>(n);
  cur_pos = 1;
  dfs(1); decompose(1, 1);
11 query(int a, int b) {
  11 \text{ res} = 0;
  for (; head[a] != head[b]; b = parent[head
      [b]]) {
   if (depth[head[a]] > depth[head[b]])
        swap(a, b);
   11 cur_heavy_path_res = segtree_query(1,
         1, n, pos[head[b]], pos[b]);
   res += cur_heavy_path_res;
  if (depth[a] > depth[b]) swap(a, b);
  11 last_heavy_path_res = segtree_query(1,
      1, n, pos[a], pos[b]);
  res += last_heavy_path_res;
  return res;
```

#### K-th Root of a Permuta-5.8 tion

```
vector<vector<int>> decompose(vector<int> &p
    ) {
  int n = p.size();
 vector<vector<int>> cycles;
 vector<bool> vis(n, 0);
 for (int i = 0; i < n; i++) {</pre>
   if (!vis[i]) {
     vector<int> v:
     while (!vis[i]) {
       v.push_back(i);
       vis[i] = 1;
       i = p[i];
     cycles.push_back(v);
 }
 return cycles;
vector<int> restore(int n, vector<vector<int</pre>
    >> &cycles) {
 vector<int> p(n);
 for (auto v : cycles) {
   int m = v.size();
   for (int i = 0; i < m; i++) p[v[i]] = v</pre>
        [(i + 1) % m];
 }
//cycle decomposition of the k-th power of p
vector<vector<int>> power(vector<int> &p,
    int k) {
 int n = p.size();
 auto cycles = decompose(p);
 vector<vector<int>> ans;
 for (auto v : cycles) {
   int len = v.size(), g = __gcd(k, len);
        //g cycles of len / g length
   for (int i = 0; i < g; i++) {</pre>
     vector<int> w;
```

```
for (int j = i, cnt = 0; cnt < len / g
          ; cnt++, j = (j + k) \% len) {
       w.push_back(v[j]);
     ans.push_back(w);
 return ans;
    with minimum disjoint cycles
//if toggle = 1, then the parity of number
    of cycles will be different from the
    other(toggle = 0) if possible
//returns empty vector if no solution exists
vector<vector<int>> root(vector<int> &p, int
     k, int toggle = 0) {
  int n = p.size();
  vector<vector<int>> cycles[n + 1];
  auto d = decompose(p);
  for (auto v : d) cycles[(int)v.size()].
      push_back(v);
  vector<vector<int>> ans;
  for (int len = 1; len <= n; len++) {</pre>
   if (cycles[len].empty()) continue;
   int tmp = k, t = 1, x = __gcd(len, tmp);
   while (x != 1) {
     t *= x;
     tmp /= x;
     x = \_gcd(len, tmp);
   if ((int)cycles[len].size() % t != 0) {
     ans.clear();
     return ans; //no solution exists
   int id = 0;
   //we can merge t * z cycles iff tmp \% z
    if (toggle && tmp % 2 == 0 && (int)
        cycles[len].size() >= 2 * t) {
     int m = 2 * t * len;
     vector<int> cycle(m);
     for (int x = 0; x < 2 * t; x++) { //
          merging 2t cycles to perform the
          toggle
       for (int y = 0; y < len; y++) {</pre>
         cycle[(x + 1LL * y * k) % m] =
              cycles[len][x][v];
     ans.push_back(cycle);
     id = 2 * t;
     toggle = 0;
   for (int i = id; i < (int)cycles[len].</pre>
        size(); i += t) {
     int m = t * len:
     vector<int> cycle(m);
     for (int x = 0; x < t; x++) { //
          merging t cycles
       for (int y = 0; y < len; y++) {
         cycle[(x + 1LL * y * k) % m] =
              cycles[len][i + x][y];
     }
     ans.push_back(cycle);
   }
 return ans;
//minimum swaps to obtain this perm from
    unit perm
vector<pair<int, int>> transpositions(vector
    <vector<int>> &cycles) {
  vector<pair<int, int>> ans;
  for (auto v : cycles) {
   int m = v.size();
   for (int i = m - 1; i > 0; i--) ans.
        push_back({v[0], v[i]});
 return ans;
int32_t main() {
  ios_base::sync_with_stdio(0);
```

```
cin.tie(0);
int n, 1, k;
cin >> n >> 1 >> k;
vector<int> p(n);
for (auto &x : p) cin >> x, --x;
auto a = root(p, k);
if (a.empty()) return cout << "no solution</pre>
     n'', 0;
auto t = transpositions(a);
if (t.size() % 2 != 1 % 2) {
 a = root(p, k, 1);
  t = transpositions(a);
if (t.size() % 2 != 1 % 2 || t.size() > 1)
     return cout << "no solution\n", 0;</pre>
auto z = restore(n, a);
auto w = power(z, k);
auto x = restore(n, w);
assert(p == x);
for (auto x : t) cout << x.first + 1 << '</pre>
      << x.second + 1 << '\n';</pre>
for (int i = t.size(); i < 1; i++) cout <<</pre>
     1 << ' ' << 2 << '\n';
return 0;
```

#### 5.9Lowest Common Ancestor

```
int lv1[MAX], P[MAX][25];
void dfs(int u, int par, int d) {
 lvl[u] = d:
 P[u][0] = par;
 for (int v : g[u]) {
   if (v == par) continue;
   dfs(v, u, d + 1);
void init() {
 for (int j = 0; j < 25; j++) {</pre>
   for (int i = 0; i <= n; i++) P[i][j] =</pre>
        -1:
 dfs(1, -1, 0);
 for (j = 1; j < 25; j++) {</pre>
   for (int i = 1; i <= n; i++) {</pre>
     if (P[i][j - 1] != -1) {
       P[i][j] = P[P[i][j-1]][j-1];
       ///to find max weight between two
       // weight[i][j] = max(weight[i][j
            -1], weight[p[i][j-1]][j-1]);
     }
int lca(int u, int v) {
 int i, lg;
 if (lvl[u] < lvl[v]) swap(u, v);</pre>
 for (lg = 0; (1 << lg) <= lvl[u]; lg++);</pre>
 lg--;
 for (i = lg; i >= 0; i--) {
   if (lvl[u] - (1 << i) >= lvl[v]) {
     u = P[u][i];
 if (u == v) return u;
 for (i = lg; i >= 0; i--) {
   if (P[u][i] != -1 && P[u][i] != P[v][i])
     u = P[u][i], v = P[v][i];
     // ret = max(ret, weight[u][i]);
     // ret = max(ret, weight[v][i]);
```

}

// ret = max(ret, weight[u][0]);

```
return P[u][0];
//Get the ancestor of node "u"
//which is "dis" distance above.
int getAncestor(int u, int dis) {
 dis = lvl[u] - dis;
 int i, lg = 0;
 for (; (1 << lg) <= lvl[u]; lg++) continue</pre>
 for (i = lg; i >= 0; i--) {
   if (lvl[u] - (1 << i) >= dis) {
     u = P[u][i];
 }
 return u;
//returns the distance between
//two nodes "u" and "v".
int dis(int u, int v) {
 if (lv1[u] < lv1[v]) swap(u, v);</pre>
 int p = lca(u, v);
 return lvl[u] + lvl[v] - 2 * lvl[p];
```

#### 5.10Min Cost Max Flow

```
mt19937 rnd(chrono::steady_clock::now().
    time_since_epoch().count());
struct edge {
 int v, rev;
 ll cap, cost, flow;
 edge() {}
 edge(int v, int rev, ll cap, ll cost) : v(
      v), rev(rev), cap(cap), cost(cost),
      flow(0) {}
struct mcmf {
 int src, sink, nodes;
 vector<int> par, idx, Q;
 vector<bool> inq;
 vector<ll> dis;
 vector<vector<edge>> g;
 mcmf() {}
 mcmf(int src, int sink, int nodes) : src(
      src), sink(sink), nodes(nodes),
   par(nodes), idx(nodes), inq(nodes),
   dis(nodes), g(nodes), Q(10000005) \{\} //
        use Q(nodes) if not using random
  void add_edge(int u, int v, ll cap, ll
      cost, bool directed = true) {
    edge _u = edge(v, g[v].size(), cap, cost
   edge _v = edge(u, g[u].size(), 0, -cost)
   g[u].pb(_u);
   g[v].pb(_v);
   if (!directed) add_edge(v, u, cap, cost,
         true):
  bool spfa() {
   for (int i = 0; i < nodes; i++) {</pre>
     dis[i] = inf, inq[i] = false;
   int f = 0, 1 = 0;
   dis[src] = 0, par[src] = -1, Q[1++] =
        src, inq[src] = true;
   while (f < 1) {
     int u = Q[f++];
     for (int i = 0; i < g[u].size(); i++)</pre>
       edge &e = g[u][i];
       if (e.cap <= e.flow) continue;</pre>
       if (dis[e.v] > dis[u] + e.cost) {
         dis[e.v] = dis[u] + e.cost;
         par[e.v] = u, idx[e.v] = i;
         if(!inq[e.v]) inq[e.v] = true, Q[1
             ++] = e.v;
         // if (!inq[e.v]) {
         // inq[e.v] = true;
```

```
// if (f && rnd() & 7) Q[--f] = e
             .v;
            else Q[1++] = e.v;
        // }
     inq[u] = false;
   return (dis[sink] != inf);
 pair<11, 11> solve() {
   11 mincost = 0, maxflow = 0;
   while (spfa()) {
     11 bottleneck = inf;
     for (int u = par[sink], v = idx[sink];
          u != -1; v = idx[u], u = par[u])
       edge &e = g[u][v];
       bottleneck = min(bottleneck, e.cap -
            e.flow):
     for (int u = par[sink], v = idx[sink];
          u != -1; v = idx[u], u = par[u])
       edge &e = g[u][v];
       e.flow += bottleneck;
      g[e.v][e.rev].flow -= bottleneck;
     mincost += bottleneck * dis[sink],
         maxflow += bottleneck;
   return make_pair(mincost, maxflow);
 }
// want to minimize cost and don't care
    about flow
// add edge from sink to dummy sink (cap =
    inf, cost = 0)
// add edge from source to sink (cap = inf,
    cost = 0
// run mcmf, cost returned is the minimum
```

#### 5.11Tree Isomorphism

```
mp["01"] = 1;
ind = 1;
int dfs(int u, int p) {
 int cnt = 0:
 vector<int>vs;
 for (auto v : g1[u]) {
   if (v != p) {
     int got = dfs(v, u);
     vs.pb(got):
     cnt++;
 if (!cnt) return 1;
 sort(vs.begin(), vs.end());
 string s = "0";
 for (auto i : vs) s += to_string(i);
 vs.clear();
 s.pb('1');
 if (mp.find(s) == mp.end()) mp[s] = ++ind;
 int ret = mp[s];
 return ret;
```

#### 6 Math

#### 6.1Berlekamp Massey

```
struct berlekamp_massey { // for linear
    recursion
 typedef long long LL;
 static const int SZ = 2e5 + 5;
 static const int MOD = 1e9 + 7; /// mod
      must be a prime
 LL m , a[SZ] , h[SZ] , t_[SZ] , s[SZ] , t[  
     SZ];
 // bigmod goes here
 inline vector <LL> BM( vector <LL> &x ) {
   LL lf , ld;
   vector <LL> ls , cur;
```

```
for ( int i = 0; i < int(x.size()); ++i</pre>
      ) {
    LL t = 0;
    for ( int j = 0; j < int(cur.size());</pre>
        ++j) t = (t + x[i - j - 1] * cur
        [j]) % MOD;
    if ( (t - x[i]) \% MOD == 0 ) continue;
   if ( !cur.size() ) {
     cur.resize( i + 1 );
     lf = i; ld = (t - x[i]) % MOD;
     continue;
   LL k = -(x[i] - t) * bigmod(ld, MOD)
        - 2 , MOD ) \% MOD;
    vector <LL> c(i - lf - 1);
    c.push_back( k );
    for ( int j = 0; j < int(ls.size());</pre>
         ++j ) c.push_back(-ls[j] * k %
        MOD);
    if ( c.size() < cur.size() ) c.resize(</pre>
         cur.size() );
    for ( int j = 0; j < int(cur.size());</pre>
        ++j ) c[j] = (c[j] + cur[j]) %
        MOD;
    if (i - lf + (int)ls.size() >= (int)
        cur.size() ) ls = cur, lf = i, ld
         = (t - x[i]) \% MOD;
  for ( int i = 0; i < int(cur.size()); ++</pre>
       i ) cur[i] = (cur[i] % MOD + MOD) %
       MOD:
  return cur;
inline void mull( LL *p , LL *q ) {
  for ( int i = 0; i < m + m; ++i ) t_[i]</pre>
      = 0;
  for ( int i = 0; i < m; ++i ) if ( p[i]</pre>
     for ( int j = 0; j < m; ++j ) t_[i +
           j] = (t_[i + j] + p[i] * q[j])
          % MOD;
  for ( int i = m + m - 1; i >= m; --i )
      if ( t_[i] )
     for ( int j = m - 1; ~j; --j ) t_[i
          -j-1] = (t_[i-j-1] + t_[i
          ] * h[j]) % MOD;
  for ( int i = 0; i < m; ++i ) p[i] = t_[</pre>
      i];
inline LL calc( LL K ) {
 for ( int i = m; ~i; --i ) s[i] = t[i] =
       0;
  s[0] = 1; if (m!=1) t[1] = 1; else t
      [0] = h[0]:
  while (K) {
   if (K&1) mull(s,t);
   mull( t , t ); K >>= 1;
  LL su = 0;
  for ( int i = 0; i < m; ++i ) su = (su +</pre>
       s[i] * a[i]) % MOD;
  return (su % MOD + MOD) % MOD;
/// already calculated upto k , now
    calculate upto n.
inline vector <LL> process( vector <LL> &x
     , int n , int k ) {
  auto re = BM( x );
  x.resize(n + 1);
  for ( int i = k + 1; i <= n; i++ ) {</pre>
   for ( int j = 0; j < re.size(); j++ )</pre>
     x[i] += 1LL * x[i - j - 1] % MOD *
          re[j] % MOD; x[i] %= MOD;
   }
 return x;
inline LL work( vector <LL> &x , LL n ) {
  if ( n < int(x.size()) ) return x[n] %</pre>
  vector <LL> v = BM( x ); m = v.size();
       if (!m) return 0;
```

```
for ( int i = 0; i < m; ++i ) h[i] = v[i</pre>
        ], a[i] = x[i];
   return calc( n ) % MOD;
 }
} rec;
vector <LL> v;
void solve() {
 int n;
 cin >> n;
 cout << rec.work(v, n - 1) << endl;</pre>
```

### 6.2Chinese Remainder Theo-

```
// given a, b will find solutions for
// ax + by = 1
tuple <LL,LL,LL> EGCD(LL a, LL b){
   if(b == 0) return {1, 0, a};
   elsef
       auto [x,y,g] = EGCD(b, a\%b);
       return {y, x - a/b*y,g};
// given modulo equations, will apply CRT
PLL CRT(vector <PLL> &v){
   LL V = 0, M = 1;
   for(auto &[v, m]:v){
       auto [x, y, g] = EGCD(M, m);
       if((v - V) \% g != 0)
          return {-1, 0};
       V += x * (v - V) / g % (m / g) * M,
           M *= m / g;
       V = (V \% M + M) \% M;
   return make_pair(V, M);
```

#### Derangements 6.3

```
array <int, N + 1> Drng;
void init(){
   Drng[0] = 1, Drng[1] = 0;
   for(int i = 2; i <= N; i++)</pre>
       Drng[i] = (LL) (i - 1) * (Drng[i -
            1] + Drng[i - 2]) % mod;
int D(int n) {
   return n < 0 ? 0 : Drng[n];</pre>
```

#### 6.4 FFT in Mod

```
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a, const
     vector<LL> &b) {
    int n = 1;
   while (n < a.size() + b.size()) n <<= 1;</pre>
   vector<CD> al(n), ar(n), bl(n), br(n);
   for (int i = 0; i < a.size(); i++) al[i]</pre>
         = a[i] % M / B, ar[i] = a[i] % M %
         В;
   for (int i = 0; i < b.size(); i++) bl[i]</pre>
         = b[i] \% M / B, br[i] = b[i] % M %
   pairfft(al, ar); pairfft(bl, br);
        fft(al); fft(ar); fft(bl); fft(br);
   for (int i = 0; i < n; i++) {</pre>
       CD ll = (al[i] * bl[i]), lr = (al[i]
             * br[i]);
       CD rl = (ar[i] * bl[i]), rr = (ar[i]
             * br[i]);
       al[i] = ll; ar[i] = lr;
       bl[i] = rl; br[i] = rr;
   pairfft(al, ar, true); pairfft(bl, br,
         fft(al, true); fft(ar, true); fft(
    bl, true); fft(br, true);
   vector<LL> ans(n);
```

```
for (int i = 0; i < n; i++) {</pre>
   LL right = round(br[i].real()), left
         = round(al[i].real());;
   LL mid = round(round(bl[i].real()) +
         round(ar[i].real()));
   ans[i] = ((left \% M) * B * B + (mid)
        % M) * B + right) % M;
}
return ans;
```

#### 6.5Fast Fourier Transformation

```
1. Whenever possible remove leading zeros.
2. Custom Complex class may slightly improve
     performance.
3. Use pairfft to do two ffts of real
    vectors at once, slightly less accurate
than doing two ffts, but faster by about
4. FFT accuracy depends on answer. x \le 5e14
     (double), x <= 1e18(long double)</pre>
   where x = max(ans[i]) for FFT, and x = N*
       mod for anymod **/
const double PI = acos(-1.0L);
int N:
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
   assert((n & (n - 1)) == 0);
   N = n;
   perm = vector<int> (N, 0);
   for (int k = 1; k < N; k <<= 1)</pre>
       for (int i = 0; i < k; i++)</pre>
           perm[i] \ll 1, perm[i + k] = 1 + |6.6|
               perm[i];
   wp[0] = wp[1] = vector < CD > (N);
   for (int i = 0; i < N; i++) {</pre>
       wp[0][i] = CD(cos(2 * PI * i / N),
            sin(2 * PI * i / N));
       wp[1][i] = CD(cos(2 * PI * i / N), -
            sin(2 * PI * i / N));
   }
void fft(vector<CD> &v, bool invert = false)
   if (v.size() != perm.size())
        precalculate(v.size());
   for (int i = 0; i < N; i++)</pre>
       if (i < perm[i])</pre>
           swap(v[i], v[perm[i]]);
   for (int len = 2; len <= N; len *= 2) {</pre>
       for (int i = 0, d = N / len; i < N;</pre>
            i += len) {
           for (int j = 0, idx = 0; j < len
               / 2; j++, idx += d) {
               CD x = v[i + j];
               CD y = wp[invert][idx] * v[i
                   + j + len / 2];
               v[i + j] = x + y;
               v[i + j + len / 2] = x - y;
           }
       }
   }
   if (invert) {
       for (int i = 0; i < N; i++) v[i] /=</pre>
            N;
void pairfft(vector<CD> &a, vector<CD> &b,
    bool invert = false) {
    int N = a.size();
   vector<CD> p(N);
   for (int i = 0; i < N; i++) p[i] = a[i]</pre>
        + b[i] * CD(0, 1);
```

```
fft(p, invert);
   p.push_back(p[0]);
   for (int i = 0; i < N; i++) {</pre>
       if (invert) {
           a[i] = CD(p[i].real(), 0);
           b[i] = CD(p[i].imag(), 0);
       } else {
           a[i] = (p[i] + conj(p[N - i])) *
                CD(0.5, 0);
           b[i] = (p[i] - conj(p[N - i])) *
                CD(0, -0.5);
       }
   }
vector<LL> multiply(const vector<LL> &a,
    const vector<LL> &b) {
   int n = 1;
   while (n < a.size() + b.size()) n <<= 1;</pre>
   vector<CD> fa(a.begin(), a.end()), fb(b.
        begin(), b.end());
   fa.resize(n); fb.resize(n);
         fft(fa); fft(fb);
   pairfft(fa, fb);
   for (int i = 0; i < n; i++) fa[i] = fa[i</pre>
        ] * fb[i];
   fft(fa, true);
   vector<LL> ans(n);
   for (int i = 0; i < n; i++) ans[i] =</pre>
        round(fa[i].real());
   return ans;
```

#### Walsh Fast Hadamord

```
Transformation
#define bitwiseXOR 1
//#define bitwiseAND 2
//#define bitwiseOR 3
void FWHT(vector< LL >&p, bool inverse){
    LL n = p.size();
    assert((n&(n-1))==0);
    for (LL len = 1; 2*len <= n; len <<= 1)</pre>
        for (LL i = 0; i < n; i += len+len)</pre>
            {
           for (LL j = 0; j < len; j++) {
               LL u = p[i+j];
               LL v = p[i+len+j];
               #ifdef bitwiseXOR
               p[i+j] = u+v;
               p[i+len+j] = u-v;
               #endif // bitwiseXOR
               #ifdef bitwiseAND
               if (!inverse) {
                   p[i+j] = v \% MOD;
                   p[i+len+j] = (u+v) % MOD;
               } else {
                   p[i+j] = (-u+v) \% MOD;
                   p[i+len+j] = u % MOD;
               #endif // bitwiseAND
               #ifdef bitwiseOR
               if (!inverse) {
                   p[i+j] = u+v;
                   p[i+len+j] = u;
               } else {
                   p[i+j] = v;
                   p[i+len+j] = u-v;
               #endif // bitwiseOR
           }
       }
```

}

```
#ifdef bitwiseXOR
   if (inverse) {
       //LL val=BigMod(n,MOD-2); //Option
            2: Exclude
       for (LL i = 0; i < n; i++) {</pre>
           //assert(p[i]%n==0); //Option 2:
           //p[i] = (p[i]*val)%MOD; //Option
                 2: p[i]/=n;
           p[i]/=n;
       }
   #endif // bitwiseXOR
vector<pair<int,int> >V[100005];
int dis[100005];
void dfs(int s,int pr){
   for(auto p:V[s]){
       if(p.first==pr) continue;
       dis[p.first]=dis[s]^p.second;
       dfs(p.first,s);
int main(){
    int t;
   cin >> t;
    const int len=(1<<16);</pre>
    for(int tc=1;tc<=t;tc++){</pre>
       LL n:
       cin >> n;
       for(int i=1;i<=n-1;i++){</pre>
           int u,v,w;
           cin >> u >> v >> w;
           V[u].push_back({v,w});
           V[v].push_back({u,w});
       dfs(1,0);
       vector<LL>a(len,0);
       for(int i=1;i<=n;i++) a[dis[i]]++;</pre>
       FWHT(a,false);
       for(int i=0;i<len;i++) a[i]*=a[i];</pre>
       FWHT(a,true);
       a[0]-=n:
       cout << "Case " << tc << ":\n";
       for(int i=0;i<len;i++) cout << a[i</pre>
            ]/2 << '\n';
       memset(dis,0,sizeof(dis));
       for(int i=1;i<=n;i++) V[i].clear();</pre>
   }
```

### Gaussian Elimination

```
const double EPS = 1e-9;
const int INF = 2; // it doesn't actually
     have to be infinity or a big number
template <typename DT> int gauss (vector <</pre>
     vector<DT> > a, vector<DT> & ans) {
    int n = (int) a.size();
   int m = (int) a[0].size() - 1;
   vector<int> where (m. -1):
   for (int col=0, row=0; col<m && row<n;</pre>
         ++col) {
       int sel = row;
       for (int i=row; i<n; ++i)</pre>
           if (abs (a[i][col]) > abs (a[sel
                ][col]))
               sel = i;
       if (abs (a[sel][col]) < EPS)</pre>
            continue;
       for (int i=col; i<=m; ++i)</pre>
           swap (a[sel][i], a[row][i]);
       where[col] = row;
       for (int i=0; i<n; ++i)</pre>
           if (i != row) {
               DT c = a[i][col] / a[row][col]
                    ];
               for (int j=col; j<=m; ++j)</pre>
```

```
a[i][j] -= a[row][j] * c;
           }
       ++row;
   ans.assign (m, 0);
   for (int i=0; i<m; ++i)</pre>
       if (where[i] != -1)
           ans[i] = a[where[i]][m] / a[where int XOR[40 * MAX][26];
                [i]][i];
   for (int i=0; i<n; ++i) {</pre>
       DT sum = 0:
       for (int j=0; j<m; ++j)</pre>
           sum += ans[j] * a[i][j];
       if (abs (sum - a[i][m]) > EPS)
           return 0;
   for (int i=0; i<m; ++i)</pre>
       if (where[i] == -1)
           return INF;
   return 1;
int compute_rank(vector<vector<double>> A) {
   int n = A.size();
   int m = A[0].size();
   int rank = 0;
   vector<bool> row_selected(n, false);
   for (int i = 0; i < m; ++i) {</pre>
       int j;
       for (j = 0; j < n; ++j) {
           if (!row_selected[j] && abs(A[j][
                i]) > EPS)
               break;
       if (j != n) {
           ++rank:
           row_selected[j] = true;
           for (int p = i + 1; p < m; ++p)
               A[j][p] /= A[j][i];
           for (int k = 0; k < n; ++k) {</pre>
               if (k != j && abs(A[k][i]) >
                    EPS) {
                   for (int p = i + 1; p < m; | scanf("%d %d",&r,&c);</pre>
                         (\alpha++
                       A[k][p] -= A[j][p] * A FOR(i,1,r){
                            [k][i];
               }
           }
   }
```

#### 6.8Gradient Descent

return rank;

```
Given n 3D point. Find a point from where
     distance to farthest of those n point
     is minimum.
int main(){
    int n:
    cin >> n;
    double x[n],y[n],z[n];
    \label{eq:rep} \texttt{REP(i,n)} \ \texttt{cin} >> \texttt{x[i]} >> \texttt{y[i]} >> \texttt{z[i]};
    double px=0.0,py=0.0,pz=0.0,alpha=1.0;
    REP(loop,100000){
        int idx=0;
        double dis=-1.0;
        REP(i,n){
            double tmp=SQ(x[i]-px)+SQ(y[i]-py
                 )+SQ(z[i]-pz);
            if(tmp>dis) {
                dis=tmp;
                idx=i:
            }
        px+=alpha*(x[idx]-px);
        py+=alpha*(y[idx]-py);
        pz+=alpha*(z[idx]-pz);
```

```
alpha*=0.999;
cout << px << ' ' << py << ' ' << pz;
```

#### Green Hackenbush on Trie

```
int trie[40 * MAX][26];
int valu[40 * MAX];
int node = 1:
int add(string s) {
 int now = 1;
  stack<int>st;
 for (int i = 0; i < s.size(); i++) {</pre>
   int c = s[i] - 'a';
   if (!trie[now][c]) trie[now][c] = ++node
   st.push(now);
   now = trie[now][c];
  int nxt = now:
  int nxt_val = 0;
  for (int i = 0; i < 26; i++) nxt_val ^=</pre>
      XOR[now][i]:
  while (!st.empty()) {
   now = st.top();
   st.pop();
   int val = 0;
   for (int i = 0; i < 26; i++) {</pre>
     if (trie[now][i] == nxt) {
       XOR[now][i] = nxt_val + 1;
     val ^= XOR[now][i];
   nxt_val = val;
   nxt = now;
 return nxt_val;
```

#### Grid Nim 6.10

```
int r,c;
int nim=0:
   FOR(j,1,c){
       int tmp;
       scanf("%d",&tmp);
       if(((r-i)+(c-j))%2){
           nim^=tmp;
if(nim) printf("Case %d: win\n",tc);
else printf("Case %d: lose\n",tc);
```

#### Linear Seive with Multi-6.11plicative Functions

```
const int maxn = 1e7;
vector <int> primes;
int spf[maxn+5], phi[maxn+5], NOD[maxn+5],
    cnt[maxn+5], POW[maxn+5], SOD[maxn+5];
bool prime[maxn+5];
void sieve(){
   fill(prime+2, prime+maxn+1, 1);
   SOD[1] = NOD[1] = phi[1] = spf[1] = 1;
   for(11 i=2;i<=maxn;i++){</pre>
       if(prime[i]) {
          primes.push_back(i), spf[i] = i;
          phi[i] = i-1;
           NOD[i] = 2, cnt[i] = 1;
           SOD[i] = i+1, POW[i] = i;
       for(auto p:primes){
           if(p*i>maxn or p > spf[i]) break;
           prime[p*i] = false, spf[p*i] = p;
           if(i\%p == 0){
              phi[p*i]=p*phi[i];
```

# 6.12 Matrix Exponentiation

```
struct matrix {
  vector<vector<ll>> mat:
  int n, m;
 matrix() {}
 matrix(int n, int m) : n(n), m(m), mat(n)
   for (int i = 0; i < n; i++) mat[i] =</pre>
         vector<ll>(m);
  }
  void identity() { for (int i = 0; i < n; i</pre>
      ++) mat[i][i] = 1; }
  void print() {
   for (int i = 0; i < n; i++) {</pre>
     for (int j = 0; j < n; j++) cout <<</pre>
          mat[i][j] << " ";
      cout << "\n":
  }
  vector<ll> &operator[](int i) {
   return mat[i];
 }
}:
// make sure a.m == b.n
matrix operator * (matrix &a, matrix &b) {
  int n = a.n, m = b.m;
 matrix ret(n, m);
  for (int i = 0; i < n; i++) {</pre>
   for (int j = 0; j < m; j++) {
     for (int k = 0; k < a.m; k++) {
       ll val = (1ll * a[i][k] * b[k][j]) %
       ret[i][j] = (ret[i][j] + val) % MOD;
   }
 }
  return ret;
matrix mat_exp(matrix &mat, 11 p) {
 int n = mat.n, m = mat.m;
  matrix ret(n, m);
 ret.identity();
 matrix x = mat;
  while (p) {
   if (p & 1) ret = ret * x;
   x = x * x;
   p = p >> 1;
  }
 return ret;
```

#### 6.13 Mobius Function

```
const int N=1000001;
int mu[N];
void mobius(){
    MEM(mu,-1);
    mu[1]=1;
    for(int i = 2; i<N; i++){
        if(mu[i]){</pre>
```

# 6.14 Modular Binomial Coefficients

```
const int N = 2e5+5;
const int mod = 1e9+7;
array <int, N+1> fact, inv, inv_fact;
void init(){
   fact[0] = inv_fact[0] = 1;
   for(int i=1; i<=N; i++){</pre>
       inv[i] = i == 1 ? 1 : (LL) inv[i -
            mod\%i] * (mod/i + 1) \% mod;
       fact[i] = (LL) fact[i-1] * i % mod;
       inv_fact[i] = (LL) inv_fact[i-1] *
            inv[i] % mod;
   }
int C(int n,int r){
   if(fact[0] != 1) init();
   return (r < 0 or r > n) ? 0 : (LL) fact[
        n]*inv_fact[r] % mod * inv_fact[n-r
        ] % mod;
```

# 6.15 Number Theoretic Transformation

```
#define pii pair<LL,LL>
const LL N= 1<<18;
const LL MOD=786433;
vector<LL>P[N];
LL rev[N], w[N|1], a[N], b[N], inv_n, g;
LL Pow(LL b, LL p){
   LL ret=1;
   while(p){
       if(p & 1) ret=(ret*b)%MOD;
       b=(b*b)\%MOD;
       p>>=1;
   return ret;
LL primitive_root(LL p){
   vector<LL>factor;
   LL phi = p-1,n=phi;
   for(LL i=2;i*i<=n;i++){</pre>
       if(n%i) continue;
       factor.emplace_back(i);
       while(n%i==0) n/=i;
   if(n>1) factor.emplace_back(n);
   for(LL res=2;res<=p;res++){</pre>
       bool ok=true;
       for(LL i=0;i<factor.size() && ok;i</pre>
            ++) ok &= Pow(res,phi/factor[i])
             != 1;
       if(ok) return res;
   }
   return -1;
void prepare(LL n){
   LL sz=abs(31-__builtin_clz(n));
   LL r=Pow(g,(MOD-1)/n);
   inv_n=Pow(n,MOD-2);
   w[0]=w[n]=1;
   for(LL i=1;i<n;i++) w[i]= (w[i-1]*r)%MOD</pre>
   for(LL i=1;i<n;i++) rev[i]=(rev[i</pre>
        >>1]>>1) | ((i & 1)<<(sz-1));
```

void NTT(LL \*a,LL n,LL dir=0){

```
for(LL i=1;i<n-1;i++) if(i<rev[i]) swap(</pre>
         a[i],a[rev[i]]);
    for(LL m=2;m<=n;m <<= 1) {</pre>
        for(LL i=0;i<n;i+=m){</pre>
           for(LL j=0;j< (m>>1);j++){
               LL &u=a[i+j], &v=a[i+j+(m>>1)
                    ];
               LL t=v*w[dir ? n-n/m*j:n/m*j
                    ]%MOD;
                v=u-t<0?u-t+MOD:u-t;
               u=u+t>=MOD?u+t-MOD:u+t;
           }
       }
    }
    if(dir) for(LL i=0;i<n;i++) a[i]=(inv_n*</pre>
         a[i])%MOD;
vector<LL> mul(vector<LL>p,vector<LL>q){
    LL n=p.size(),m=q.size();
    LL t=n+m-1, sz=1;
    while(sz<t) sz <<= 1;</pre>
    prepare(sz);
    for(LL i=0;i<n;i++) a[i]=p[i];</pre>
    for(LL i=0;i<m;i++) b[i]=q[i];</pre>
    for(LL i=n;i<sz;i++) a[i]=0;</pre>
    for(LL i=m;i<sz;i++) b[i]=0;</pre>
    NTT(a,sz);
    NTT(b.sz):
    for(LL i=0;i<sz;i++) a[i]=(a[i]*b[i])%</pre>
    NTT(a,sz,1);
    vector<LL> c(a,a+sz);
    while(c.size() && c.back()==0) c.
        pop_back();
    return c;
N different number box
Number of ways to make a number by picking
     any number from any of the boxes
vector<LL> solve(LL 1,LL r){
    if(l==r) return P[1];
    LL m=(1+r)/2:
    return mul(solve(1,m),solve(m+1,r));
int main(){
    cin >> m:
    for(LL i=1;i<=m;i++){</pre>
       LL num;
        cin >> num;
        vector<pii>v;
        LL mx=0:
        while(num--){
           LL typ, cnt;
            cin >> typ >> cnt;
            v.emplace_back(typ,cnt);
           mx=max(mx,typ);
        P[i].resize(mx+1);
        for(pii p:v) P[i][p.first]=p.second;
    g=primitive_root(MOD);
    vector<LL>c=solve(1,m);
    for(LL i=0;i<c.size();i++){</pre>
        if(c[i]){
            cout << i << ' ' << c[i] << '\n';
   }
```

# 6.16 Pollard Rho and Factorization

// fast factorize

```
map <ull,int> fast_factorize(ull n){
   map <ull,int> ans;
   for(;n>1;n/=spf[n])
       ans[spf[n]]++;
   return ans;
inline ULL mul(ULL a,ULL b,ULL mod){
   LL ans = a * b - mod * (ULL) (1.L / mod)
        * a * b);
   return ans + \mod * (ans < 0) - \mod * (
        ans >= (LL) mod);
inline ULL bigmod(ULL num,ULL pow,ULL mod){
   ULL ans = 1;
   for( ; pow > 0; pow >>= 1, num = mul(num
          num, mod))
       if(pow&1) ans = mul(ans,num,mod);
   return ans;
inline bool is_prime(ULL n){
   if(n < 2 or n % 6 % 4 != 1)
       return (n|1) == 3;
   ULL a[] = \{2, 325, 9375, 28178, 450775,
        9780504, 1795265022};
   ULL s = __builtin_ctzll(n-1), d = n >> s
   for(ULL x: a){
       ULL p = bigmod(x \% n, d, n), i = s;
       for(; p != 1 \text{ and } p != n-1 \text{ and } x \% n
             and i--; p = mul(p, p, n));
       if(p != n-1 and i != s)
           return false;
   }
   return true;
ULL get_factor(ULL n) {
   auto f = [&](LL x) { return mul(x, x, n)
         + 1; };
   ULL x = 0, y = 0, t = 0, prod = 2, i = 0
        2, q;
   for( ; t++ %40 or gcd(prod, n) == 1; x =
         f(x), y = f(f(y)) (
       (x == y) ? x = i++, y = f(x) : 0;
       prod = (q = mul(prod, max(x,y) - min)
            (x,y), n)) ? q: prod;
   return gcd(prod, n);
map <ULL, int> factorize(ULL n){
   map <ULL, int> res;
    if(n < 2) return res;</pre>
   ULL small_primes[] = {2, 3, 5, 7, 11,
        13, 17, 19, 23, 29, 31, 37, 41, 43,
         47, 53, 59, 61, 67, 71, 73, 79,
        83, 89, 97 };
   for (ULL p: small_primes)
       for( ; n % p == 0; n /= p, res[p]++)
   auto _factor = [&](ULL n, auto &_factor)
       if(n == 1) return;
       if(is_prime(n))
          res[n]++:
       else {
           ULL x = get_factor(n);
           _factor(x, _factor);
           _factor(n / x, _factor);
    _factor(n, _factor);
   return res;
```

### 6.17 Prime Counting Function

```
int len = 0; // total number of primes
     generated by sieve
int primes[MAX_PRIMES];
                     // pref[i] --> number
int pref[MAXN];
     of primes <= i
int dp[PHI_N][PHI_K]; // precal of yo(n,k)
bitset<MAXN> f;
void sieve(int n) {
    f[1] = true;
    for (int i = 4; i <= n; i += 2) f[i] =
        true;
    for (int i = 3; i * i <= n; i += 2) {</pre>
       if (!f[i]) {
           for (int j = i * i; j <= n; j +=</pre>
                i << 1) f[j] = 1;
    }
    for (int i = 1; i <= n; i++) {</pre>
       if (!f[i]) primes[len++] = i;
       pref[i] = len;
void init() {
    sieve(MAXN - 1);
    // precalculation of phi upto size (
        PHI_N,PHI_K)
    for (int n = 0; n < PHI_N; n++) dp[n][0]
          = n;
    for (int k = 1; k < PHI_K; k++) {</pre>
       for (int n = 0; n < PHI_N; n++) {</pre>
           dp[n][k] = dp[n][k - 1] - dp[n /
                primes[k - 1]][k - 1];
       }
// returns the number of integers less or
    equal n which are
primes
// recurrence \rightarrow yo(n, k) = yo(n, k-1) - yo
     (n / p_k, k-1)
 // for sum of primes yo(n, k) = yo(n, k-1)
      p_k * yo(n / p_k , k-1)
long long yo(long long n, int k) {
    if (n < PHI_N && k < PHI_K) return dp[n</pre>
        ][k];
    if (k == 1) return ((++n) >> 1);
    if (primes[k - 1] >= n) return 1;
    return yo(n, k - 1) - yo(n / primes[k -
        1], k - 1);
// complexity: n^(2/3).(log n^(1/3))
long long Legendre(long long n) {
    if (n < MAXN) return pref[n];</pre>
    int lim = sqrt(n) + 1;
    int k = upper_bound(primes, primes + len
        , lim) - primes;
    return yo(n, k) + (k - 1);
// runs under 0.2s for n = 1e12
long long Lehmer(long long n) {
    if (n < MAXN) return pref[n];</pre>
    long long w, res = 0;
    int b = sqrt(n), c = Lehmer(cbrt(n)), a
        = Lehmer(sqrt(b));
    b = Lehmer(b);
    res = yo(n, a) + ((1LL * (b + a - 2) * (
        b - a + 1)) >> 1);
    for (int i = a; i < b; i++) {</pre>
       w = n / primes[i];
       int lim = Lehmer(sqrt(w));
       res -= Lehmer(w):
       if (i <= c) {</pre>
           for (int j = i; j < lim; j++) {
               res += j;
               res -= Lehmer(w / primes[j]);
       }
    }
    return res;
```

### 6.18 Seive Upto 1e9

```
// credit: min_25
// takes 0.5s for n = 1e9
vector<int> sieve(const int N, const int Q =
      17, const int L = 1 << 15) {
  static const int rs[] = {1, 7, 11, 13, 17,
        19, 23, 29};
  struct P {
   P(int p) : p(p) {}
    int p; int pos[8];
  auto approx_prime_count = [] (const int N)
    return N > 60184 ? N / (log(N) - 1.1)
                    : \max(1., N / (\log(N) -
                         1.11)) + 1;
  const int v = sqrt(N), vv = sqrt(v);
  vector<bool> isp(v + 1, true);
  for (int i = 2; i <= vv; ++i) if (isp[i])</pre>
    for (int j = i * i; j <= v; j += i) isp[</pre>
         j] = false;
  const int rsize = approx_prime_count(N +
      30);
  vector<int> primes = {2, 3, 5}; int psize
      = 3;
  primes.resize(rsize);
  vector<P> sprimes; size_t pbeg = 0;
  int prod = 1;
  for (int p = 7; p <= v; ++p) {</pre>
    if (!isp[p]) continue;
    if (p <= Q) prod *= p, ++pbeg, primes[</pre>
        psize++] = p;
    auto pp = P(p);
    for (int t = 0; t < 8; ++t) {</pre>
      int j = (p \le Q) ? p : p * p;
      while (j % 30 != rs[t]) j += p << 1;</pre>
     pp.pos[t] = j / 30;
    sprimes.push_back(pp);
  vector<unsigned char> pre(prod, 0xFF);
  for (size_t pi = 0; pi < pbeg; ++pi) {</pre>
    auto pp = sprimes[pi]; const int p = pp.
        p;
    for (int t = 0; t < 8; ++t) {</pre>
      const unsigned char m = ~(1 << t);</pre>
      for (int i = pp.pos[t]; i < prod; i +=</pre>
            p) pre[i] &= m;
  const int block_size = (L + prod - 1) /
      prod * prod;
  vector<unsigned char> block(block_size);
      unsigned char* pblock = block.data();
  const int M = (N + 29) / 30;
  for (int beg = 0; beg < M; beg +=</pre>
       block_size, pblock -= block_size) {
    int end = min(M, beg + block_size);
    for (int i = beg; i < end; i += prod) {</pre>
      copy(pre.begin(), pre.end(), pblock +
          i);
    if (beg == 0) pblock[0] &= 0xFE;
    for (size_t pi = pbeg; pi < sprimes.size</pre>
        (); ++pi) {
      auto& pp = sprimes[pi];
      const int p = pp.p;
      for (int t = 0; t < 8; ++t) {</pre>
        int i = pp.pos[t]; const unsigned
            char m = ~(1 << t);</pre>
        for (; i < end; i += p) pblock[i] &=</pre>
       pp.pos[t] = i;
    for (int i = beg; i < end; ++i) {</pre>
```

```
for (int m = pblock[i]; m > 0; m &= m
           - 1) {
       primes[psize++] = i * 30 + rs[
            __builtin_ctz(m)];
 assert(psize <= rsize);</pre>
  while (psize > 0 && primes[psize - 1] > N) |}
        --psize;
 primes.resize(psize);
 return primes;
int32_t main() {
 ios_base::sync_with_stdio(0);
  cin.tie(0);
 int n, a, b; cin >> n >> a >> b;
 auto primes = sieve(n);
 vector<int> ans;
 for (int i = b; i < primes.size() &&</pre>
       primes[i] <= n; i += a) ans.push_back</pre>
       (primes[i]);
  cout << primes.size() << ' ' << ans.size()</pre>
        << '\n';
  for (auto x: ans) cout << x << ' '; cout</pre>
      << '\n';
 return 0;
```

# 6.19 Simpson Integration

```
For finding the length of an arc in a
        range
   L = integrate(ds) from start to end of
        range
   where ds = sqrt(1+(d/dy(x))^2)dy
const double SIMPSON_TERMINAL_EPS = 1e-12;
/// Function whose integration is to be
    calculated
double F(double x);
double simpson(double minx, double maxx)
   return (maxx - minx) / 6 * (F(minx) + 4)
        * F((minx + maxx) / 2.) + F(maxx));
double adaptive_simpson(double minx, double
    maxx, double c, double EPS)
     if(maxx - minx < SIMPSON_TERMINAL_EPS)</pre>
     return 0;
   double midx = (minx + maxx) / 2;
   double a = simpson(minx, midx);
   double b = simpson(midx, maxx);
   if(fabs(a + b - c) < 15 * EPS) return a</pre>
        + b + (a + b - c) / 15.0;
   return adaptive_simpson(minx, midx, a,
        EPS / 2.) + adaptive_simpson(midx,
        maxx, b, EPS / 2.);
double adaptive_simpson(double minx, double
    maxx, double EPS)
   return adaptive_simpson(minx, maxx,
        simpson(minx, maxx, i), EPS);
```

### 6.20 Stirling Numbers

vis[marble][box] = cs;

```
//stirling number 2nd kind variation(number of ways to place n marbles in k boxes so that each box has at least x marbles )

11 solve(int marble, int box) {
   if (marble < 111 * box * x) return 0;
   if (box == 1 && marble >= x) return 1;
   if (vis[marble][box] == cs) return dp[
        marble][box];
```

```
ll a = ( 111 * box * solve(marble - 1, box
      ) ) % MOD;
 11 b = ( 111 * box * ncr(marble - 1, x -
      1) ) % MOD;
 b = (b * solve(marble - x, box - 1)) % MOD
 11 ret = (a + b) % MOD;
 return dp[marble][box] = ret;
//number of ways to place n marbles in k
    boxes so that no box is empty
11 stir(ll n, ll k) {
 ll ret = 0;
 for (int i = 0; i <= k; i++) {</pre>
   11 v = ncr(k, i) * bigmod(i, n) % MOD;
   if ( (k - i) % 2 == 0 ) ret = (ret + v)
        % MOD;
   else ret = (ret - v + MOD) % MOD;
 return ret:
```

#### 6.21 Subset Convolution

```
array<array<int, N>, b> Fh = {0}, Gh = {0},
    H = \{0\};
for (int mask = 0; mask < N; mask++)</pre>
   Fh[__builtin_popcount(mask)][mask] = F[
        maskl.
   Gh[__builtin_popcount(mask)][mask] = G[
        mask1:
for (int i = 0; i < b; i++)</pre>
   for (int j = 0; j < b; j++)
       for (int mask = 0; mask < N; mask++)</pre>
           if ((mask & (1 << j)) != 0)</pre>
               Fh[i][mask] += Fh[i][mask 1
                     (1 << j)],
                   Gh[i][mask] += Gh[i][mask
                         (1 << j)];
for (int mask = 0; mask < N; mask++)</pre>
   for (int i = 0; i < b; i++)</pre>
       for (int j = 0; j \le i; j++)
           H[i][mask] += Fh[j][mask] * Gh[i]
                 - j][mask];
for (int i = 0; i < b; i++)</pre>
   for (int j = 0; j < b; j++)
       for (int mask = 0; mask < N; mask++)</pre>
           if ((mask & (1 << j)) != 0) H[i][</pre>
                mask] -= H[i][mask ^ (1 << j
                )];
for (int mask = 0; mask < N; mask++)</pre>
   Ans[mask] = H[__builtin_popcount(mask)][
        mask];
```

# 7 String Algorithms

### 7.1 Aho Corasick

```
struct state {
 int len, par, link, next_lif;
 11 val:
  int next[26];
 char p_char;
 bool lif;
  state(int par = -1, char p_char = '$', int
       len = 0) : par(par), p_char(p_char),
       len(len) {
   lif = false;
   link = 0:
   next_lif = 0;
   val = 0:
   memset(next, 0, sizeof next);
vector<state>aho;
inline void add_str(const string &s, ll val)
     {
  int now = 0;
 for (int i = 0; i < s.size(); i++) {</pre>
```

```
int c = s[i] - 'a';
   if (!aho[now].next[c]) {
     aho[now].next[c] = (int)aho.size();
     aho.emplace_back(now, s[i], aho[now].
          len + 1);
   now = aho[now].next[c];
 aho[now].lif = true;
  aho[now].val = val;
inline void push_link() {
  queue<int>q;
  q.push(0);
  while (!q.empty()) {
   int cur = q.front();
   int link = aho[cur].link;
   q.pop();
   aho[cur].next_lif = aho[link].lif ? link
         : aho[link].next_lif;
   for (int c = 0; c < 26; c++) {</pre>
     if (aho[cur].next[c]) {
       aho[ aho[cur].next[c] ].link = cur ?
             aho[link].next[c] : 0;
       q.push( aho[cur].next[c] );
     } else aho[cur].next[c] = aho[link].
          next[c];
   aho[cur].val += aho[link].val;
inline int count(string &s) {
 int now = 0, ret = 0;
 for (int i = 0; i < s.size(); i++) {</pre>
   now = aho[now].next[s[i] - 'a'];
   ret += aho[now].val;
  return ret;
struct dynamic_aho {
 aho_corasick ac[20];
  vector<string> dict[20];
  dynamic_aho() {
   for (int i = 0; i < 20; i++) {</pre>
     ac[i].aho.clear();
     dict[i].clear();
  void add_str(string &s) {
   int idx = 0;
   for (; idx < 20 && !ac[idx].aho.empty();</pre>
         idx++) {}
   ac[idx] = aho_corasick();
   ac[idx].add_str(s, 1), dict[idx].pb(s);
   for (int i = 0; i < idx; i++) {</pre>
     for (string x : dict[i]) {
       ac[idx].add_str(x, 1);
       dict[idx].pb(x);
     ac[i].aho.clear(), dict[i].clear();
   ac[idx].push_link();
 inline int count(string &s) {
   int ret = 0;
   for (int i = 0; i < 20; i++) {</pre>
     if (!ac[i].aho.empty()) ret += ac[i].
          count(s);
   return ret:
 }
int arr[MAX];
int main() {
 fastio;
 dynamic_aho add, del;
 int m:
  cin >> m;
 while (m--) {
   int type;
   string s;
   cin >> type >> s;
   if (type == 1) add.add_str(s);
```

```
else if (type == 2) del.add_str(s);
   else cout << add.count(s) - del.count(s)</pre>
         << "\n" << flush;
}
```

#### 7.2 Double Hash

```
ostream& operator << (ostream& os, pll hash)
 return os << "(" << hash.ff << ", " <<
      hash.ss << ")";
pll operator + (pll a, ll x) {return pll(a.
    ff + x, a.ss + x);
pll operator - (pll a, ll x) {return pll(a.
    ff - x, a.ss - x);}
pll operator * (pll a, ll x) {return pll(a.
    ff * x, a.ss * x);}
pll operator + (pll a, pll x) {return pll(a.
    ff + x.ff, a.ss + x.ss);}
pll operator - (pll a, pll x) {return pll(a.
    ff - x.ff, a.ss - x.ss);}
pll operator * (pll a, pll x) {return pll(a.
    ff * x.ff, a.ss * x.ss);}
pll operator % (pll a, pll m) {return pll(a.
    ff % m.ff, a.ss % m.ss);}
pll base(1949313259, 1997293877);
pll mod(2091573227, 2117566807);
pll power (pll a, ll p) {
 if (!p) return pll(1, 1);
 pll ans = power(a, p / 2);
 ans = (ans * ans) % mod;
 if (p % 2) ans = (ans * a) % mod;
 return ans;
pll inverse(pll a) {
 return power(a, (mod.ff - 1) * (mod.ss -
      1) - 1);
pll inv_base = inverse(base);
pll val;
vector<pll> P;
void hash_init(int n) {
 P.resize(n + 1);
 P[0] = pll(1, 1);
 for (int i = 1; i <= n; i++) P[i] = (P[i -</pre>
       1] * base) % mod;
///appends c to string
pll append(pll cur, char c) {
 return (cur * base + c) % mod;
///prepends c to string with size k
pll prepend(pll cur, int k, char c) {
 return (P[k] * c + cur) % mod;
///replaces the i-th (0-indexed) character
    from right from a to b;
pll replace(pll cur, int i, char a, char b)
 cur = (cur + P[i] * (b - a)) % mod;
 return (cur + mod) % mod;
///Erases c from the back of the string
pll pop_back(pll hash, char c) {
 return (((hash - c) * inv_base) % mod +
      mod) % mod:
///Erases c from front of the string with
    size len
pll pop_front(pll hash, int len, char c) {
```

```
return ((hash - P[len - 1] * c) % mod +
      mod) % mod;
///concatenates two strings where length of
    the right is k
pll concat(pll left, pll right, int k) {
 return (left * P[k] + right) % mod;
///Calculates hash of string with size len
    repeated cnt times
///This is O(\log n). For O(1), pre-calculate
     inverses
pll repeat(pll hash, int len, ll cnt) {
 pll mul = (P[len * cnt] - 1) * inverse(P[
      len] - 1);
 mul = (mul % mod + mod) % mod;
 pll ret = (hash * mul) % mod;
 if (P[len].ff == 1) ret.ff = hash.ff * cnt |};
 if (P[len].ss == 1) ret.ss = hash.ss * cnt unordered_map<11, int, custom_hash> mp;
 return ret;
ll get(pll hash) {
 return ( (hash.ff << 32) ^ hash.ss );</pre>
struct hashlist {
 int len:
 vector<pll> H, R;
 hashlist() {}
 hashlist(string &s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, pll(0, 0)), R.resize(
        len + 2, pll(0, 0));
   for (int i = 1; i <= len; i++) H[i] =</pre>
        append(H[i - 1], s[i - 1]);
   for (int i = len; i >= 1; i--) R[i] =
        append(R[i + 1], s[i - 1]);
 /// 1-indexed
 inline pll range_hash(int 1, int r) {
   int len = r - 1 + 1;
   return ((H[r] - H[l - 1] * P[len]) % mod
         + mod) % mod;
 inline pll reverse_hash(int 1, int r) {
   int len = r - 1 + 1:
   return ((R[1] - R[r + 1] * P[len]) % mod vector<int> d2(n); ///even palindromes
         + mod) % mod;
 }
 inline pll concat_range_hash(int 11, int
      r1, int 12, int r2) {
   int len_2 = r2 - 12 + 1;
   return concat(range_hash(l1, r1),
        range_hash(12, r2), len_2);
  inline pll concat_reverse_hash(int 11, int
       r1, int 12, int r2) {
   int len_1 = r1 - l1 + 1;
   return concat(reverse_hash(12, r2),
        reverse_hash(l1, r1), len_1);
 }
```

#### Faster Hash Table

```
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
struct custom_hash {
 const 11 rnd = chrono::
      high_resolution_clock::now().
      time_since_epoch().count();
```

```
static unsigned long long hash_f(unsigned
     long long x) {
   x += 0x9e3779b97f4a7c15;
   x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9
   x = (x ^(x >> 27)) * 0x94d049bb133111eb
   return x ^ (x >> 31);
 11 operator() (11 x) const { return hash_f
      (x) ^ rnd; }
 // static int combine_hash(pii hash) {
     return hash.ff * 31 + hash.ss; }
 // static ll combine_hash(pll hash) {
     return (hash.ff << 32) ^ hash.ss; }
 // ll operator() (pll x) const {
     x.ff = hash_f(x.ff) ^ rnd; x.ss =
      hash_f(x.ss) ^ rnd;
     return combine_hash(x);
gp_hash_table<11, int, custom_hash> mp;
```

#### 7.4 Knuth Morris Pratt

```
vector<int> prefix_function(string s) {
 int n = (int)s.length();
 vector<int> pi(n);
 for (int i = 1; i < n; i++) {</pre>
   int j = pi[i - 1];
   while (j > 0 \&\& s[i] != s[j]) j = pi[j -
        1];
   if (s[i] == s[j]) j++;
  pi[i] = j;
 return pi;
```

#### 7.5Manacher

```
vector<int> d1(n); ///odd palindromes
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 1 : min(d1[l + r - i], r
       - i + 1);
 while (0 \le i - k \&\& i + k \le n \&\& s[i - k]
       == s[i + k]) {
 d1[i] = k--;
 if (i + k > r) {
   1 = i - k;
   r = i + k;
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 0 : min(d2[1 + r - i +
      1], r - i + 1);
 while (0 <= i - k - 1 && i + k < n && s[i
      -k-1] == s[i+k]) {
   k++;
 d2[i] = k--;
 if (i + k > r) {
   l = i - k - 1;
   r = i + k;
 }
```

#### Palindromic Tree 7.6

```
struct state {
 int len. link:
 map<char, int> next;
state st[MAX];
int id, last;
string s;
11 ans[MAX];
void init() {
 for (int i = 0; i <= id; i++) {</pre>
    st[i].len = 0; st[i].link = 0;
    st[i].next.clear(); ans[i] = 0;
```

```
st[1].len = -1; st[1].link = 1;
 st[2].len = 0; st[2].link = 1;
 id = 2; last = 2;
void extend(int pos) {
 while (s[pos - st[last].len - 1] != s[pos
      ]) last = st[last].link;
 int newlink = st[last].link;
 char c = s[pos];
 while (s[pos - st[newlink].len - 1] != s[
      pos]) newlink = st[newlink].link;
 if (!st[last].next.count(c)) {
   st[last].next[c] = ++id;
   st[id].len = st[last].len + 2;
   st[id].link = (st[id].len == 1 ? 2 : st[
        newlink].next[c]);
   ans[id] += ans[st[id].link];
   if (st[id].len > 2) {
     int 1 = st[id].len / 2 + (st[id].len %
          2 ? 1 : 0);
     if (h.range_hash(pos - st[id].len + 1,
          pos - st[id].len + 1) == h.
          reverse_hash(pos - st[id].len +
          1, pos - st[id].len + 1)) ans[id
          ]++;
   }
 last = st[last].next[c];
```

### String Match FFT

```
//find occurrences of t in s where '?'s are
    automatically matched with any
     character
//res[i + m - 1] = sum_j = 0 to m - 1_{s[i + j]}
    ] * t[j] * (s[i + j] - t[j])
vector<int> string_matching(string &s,
    string &t) {
 int n = s.size(), m = t.size();
  vector<int> s1(n), s2(n), s3(n);
 for(int i = 0; i < n; i++) s1[i] = s[i] ==</pre>
        '?' ? 0 : s[i] - 'a' + 1; //assign
      any non zero number for non '?'s
 for(int i = 0; i < n; i++) s2[i] = s1[i] *</pre>
       s1[i];
 for(int i = 0; i < n; i++) s3[i] = s1[i] *</pre>
       s2[i];
 vector<int> t1(m), t2(m), t3(m);
 for(int i = 0; i < m; i++) t1[i] = t[i] ==</pre>
        '?' ? 0 : t[i] - 'a' + 1;
  for(int i = 0; i < m; i++) t2[i] = t1[i] *</pre>
       t1[i];
 for(int i = 0; i < m; i++) t3[i] = t1[i] *</pre>
       t2[i];
 reverse(t1.begin(), t1.end());
 reverse(t2.begin(), t2.end());
 reverse(t3.begin(), t3.end());
 vector<int> s1t3 = multiply(s1, t3);
 vector<int> s2t2 = multiply(s2, t2);
 vector<int> s3t1 = multiply(s3, t1);
 vector<int> res(n);
 for(int i = 0; i < n; i++) res[i] = s1t3[i]
      ] - s2t2[i] * 2 + s3t1[i];
  vector<int> oc;
 for(int i = m - 1; i < n; i++) if(res[i]</pre>
      == 0) oc.push_back(i - m + 1);
 return oc;
```

#### 7.8Suffix Array

```
vector<vector<int> >c:
vector<int>sort_cyclic_shifts(string const&
    s) {
 int n = s.size();
 const int alphabet = 256;
 vector<int> p(n), cnt(alphabet, 0);
 c.clear(); c.emplace_back(); c[0].resize(n
 for (int i = 0; i < n; i++) cnt[s[i]]++;</pre>
 for (int i = 1; i < alphabet; i++) cnt[i]</pre>
      += cnt[i - 1];
```

return lcp;

```
for (int i = 0; i < n; i++) p[--cnt[s[i]]] }</pre>
       = i;
 c[0][p[0]] = 0;
  int classes = 1;
 for (int i = 1; i < n; i++) {</pre>
   if (s[p[i]] != s[p[i - 1]]) classes++;
   c[0][p[i]] = classes - 1;
 vector<int> pn(n), cn(n); cnt.resize(n);
 for (int h = 0; (1 << h) < n; h++) {
   for (int i = 0; i < n; i++) {</pre>
     pn[i] = p[i] - (1 << h);
     if (pn[i] < 0) pn[i] += n;</pre>
   fill(cnt.begin(), cnt.end(), 0);
   /// radix sort
   for (int i = 0; i < n; i++) cnt[c[h][pn[</pre>
        i]]]++;
   for (int i = 1; i < classes; i++) cnt[i]</pre>
         += cnt[i - 1];
   for (int i = n - 1; i >= 0; i--) p[--cnt
        [c[h][pn[i]]] = pn[i];
   cn[p[0]] = 0; classes = 1;
   for (int i = 1; i < n; i++) {</pre>
     pii cur = {c[h][p[i]], c[h][(p[i] + (1 | int main() {
           << h)) % n]};
     pii prev = {c[h][p[i - 1]], c[h][(p[i
           - 1] + (1 << h)) % n]};
     if (cur != prev) ++classes;
     cn[p[i]] = classes - 1;
   c.push_back(cn);
 7
 return p;
vector<int> suffix_array_construction(string
     s) {
 s += "!":
 vector<int> sorted_shifts =
      sort_cyclic_shifts(s);
  sorted_shifts.erase(sorted_shifts.begin())
 return sorted_shifts;
/// compare two suffixes starting at i and j
      with length 2<sup>k</sup>
int compare(int i, int j, int n, int k) {
 pii a = \{c[k][i], c[k][(i + 1 - (1 << k))
      % n]};
 pii b = \{c[k][j], c[k][(j + 1 - (1 << k))\}
      % n]};
 return a == b ? 0 : a < b ? -1 : 1;
int lcp(int i, int j) {
 int log_n = c.size() - 1;
 int ans = 0;
 for (int k = log_n; k >= 0; k--) {
   if (c[k][i] == c[k][j]) {
     ans += 1 << k;
     i += 1 << k; j += 1 << k;
 }
 return ans:
vector<int> lcp_construction(string const& s
    , vector<int> const& p) {
  int n = s.size();
 vector<int> rank(n, 0);
 for (int i = 0; i < n; i++) rank[p[i]] = i</pre>
  int k = 0;
 vector<int> lcp(n, 0);
 for (int i = 0; i < n; i++) {</pre>
   if (rank[i] == n - 1) {
     k = 0;
     continue;
   int j = p[rank[i] + 1];
   while (i + k < n \&\& j + k < n \&\& s[i + k]
        ] == s[j + k]) k++;
   lcp[rank[i]] = k;
   if (k) k--;
```

```
//kth lexicographically smallest substring (
    considering duplicates)
int arr[MAX], lg[MAX], mn[MAX][25];
vector<int>p, lcp;
string s;
int k;
// sparse table for min in lcp goes here
///find the rightmost position where get(1,r
int khoj(int 1, int r, int val) {
 int lo = 1 + 1, hi = r, ret = -1;
 while (lo <= hi) {</pre>
   int mid = lo + (hi - lo) / 2;
   if (get(1, mid - 1) > val) {
     ret = mid; lo = mid + 1;
   } else hi = mid - 1;
 return ret;
int done[MAX];
int arr[MAX];
 fastio;
 cin >> s >> k;
 p = suffix_array_construction(s);
 lcp = lcp_construction(s, p);
 build();
  int n = s.size();
 int milaisi = 0;
 for (int i = 0; i < n; i++) {</pre>
   milaisi += done[i];
   int len = n - p[i];
   int cur = milaisi;
   /// cur = i ? lcp[i-1] : 0; this can
        replace all the milaisi and done
        parts
   while (cur < len) {</pre>
     int r = khoj(i, n - 1, cur);
     int koyta, milabo;
     if (r == -1) {
       milabo = len - cur;
       kovta = 1;
     } else {
       milabo = get(i, r - 1) - cur;
       koyta = r - i + 1;
     if (koyta * milabo < k) k -= (koyta *</pre>
          milabo);
     else {
       int d = k / koyta; int m = k % koyta
       if (!m) {
         cout << s.substr(p[i], cur + d) <<</pre>
              "\n":
         return 0;
       } else {
         cout << s.substr(p[i], cur + d + 1)
               << "\n";
         return 0;
     if (r == -1) break;
     done[r + 1] -= milabo;
     cur = get(i, r - 1);
     milaisi += milabo;
 cout << "No such line.\n";</pre>
```

#### 7.9 Suffix Automaton

```
struct state {
 int len, link;
 map<char, int> next;
 bool is_clone;
 int first_pos;
 vector<int>inv link:
}:
state st[2 * MAX];
int mn[2 * MAX], mx[2 * MAX];
```

```
int sz, last;
void sa_init() {
  st[0].len = 0;
  st[0].link = -1;
  st[0].next.clear();
 sz = 1;
 last = 0;
}
void sa_extend(char c) {
 int cur = sz++;
  st[cur].len = st[last].len + 1;
  st[cur].first_pos = st[cur].len - 1;
  st[cur].is_clone = false;
  st[cur].next.clear();
  ///for lcs of n strings
  // mn[cur] = st[cur].len;
  int p = last;
  while (p != -1 && !st[p].next.count(c)) {
   st[p].next[c] = cur;
   p = st[p].link;
  if (p == -1) {
   st[cur].link = 0;
  } else {
    int q = st[p].next[c];
   if (st[p].len + 1 == st[q].len) {
     st[cur].link = q;
   } else {
     int clone = sz++;
     st[clone].len = st[p].len + 1;
     st[clone].next = st[q].next;
     st[clone].link = st[q].link;
     st[clone].first_pos = st[q].first_pos;
     st[clone].is_clone = true;
     ///for lcs of n strings
     // mn[clone] = st[clone].len;
     while (p != -1 && st[p].next[c] == q)
       st[p].next[c] = clone;
       p = st[p].link;
     st[q].link = st[cur].link = clone;
   }
 last = cur;
}
void radix_sort() {
  for (int i = 0; i < sz; i++) cnt[st[i].len</pre>
      ]++;
  for (int i = 1; i < sz; i++) cnt[i] += cnt</pre>
       [i - 1];
  for (int i = 0; i < sz; i++) order[--cnt[</pre>
      st[i].len]] = i;
// after constructing the automaton
for (int v = 1; v < sz; v++) {</pre>
  st[st[v].link].inv_link.push_back(v);
// output all positions of occurrences
void output_all_occurrences(int v, int
    P_length) {
  if (!st[v].is_clone)
    cout << st[v].first_pos - P_length + 1</pre>
        << endl:
  for (int u : st[v].inv_link)
   output_all_occurrences(u, P_length);
//lcs of two strings
string lcs (string S, string T) {
  sa init():
  for (int i = 0; i < S.size(); i++)</pre>
      sa extend(S[i]):
  int v = 0, 1 = 0, best = 0, bestpos = 0;
  for (int i = 0; i < T.size(); i++) {</pre>
   while (v && !st[v].next.count(T[i])) {
     v = st[v].link; l = st[v].len;
   if (st[v].next.count(T[i])) {
     v = st[v].next[T[i]]; 1++;
   if (1 > best) best = 1; bestpos = i;
 return T.substr(bestpos - best + 1, best);
```

```
//lcs of n strings
void add_str(string s) {
 for (int i = 0; i < sz; i++) mx[i] = 0;</pre>
 int v = 0, 1 = 0;
 for (int i = 0; i < s.size(); i++) {</pre>
   while (v && !st[v].next.count(s[i])) {
     v = st[v].link; l = st[v].len;
   if (st[v].next.count(s[i])) {
     v = st [v].next[s[i]]; 1++;
   mx[v] = max(mx[v], 1);
 for (int i = sz - 1; i > 0; i--) mx[st[i].
      link] = max(mx[st[i].link], mx[i]);
  for (int i = 0; i < sz; i++) mn[i] = min(</pre>
      mn[i], mx[i]);
int lcs() {
 int ret = 0;
 for (int i = 0; i < sz; i++) ret = max(ret</pre>
      , mn[i]);
 return ret;
string s[15];
int arr[MAX];
int main() {
 fastio;
 int n = 0;
 while (cin \gg s[n]) n++;
 sa_init();
 for (int i = 0; i < s[0].size(); i++)</pre>
      sa_extend(s[0][i]);
 for (int i = 1; i < n; i++) add_str(s[i]);</pre>
 cout << lcs() << "\n";
```

## |7.10 Z Algorithm

```
vector<int> calcz(string s) {
 int n = s.size();
 vector<int> z(n);
 int 1, r; 1 = r = 0;
 for (int i = 1; i < n; i++) {</pre>
   if (i > r) {
     1 = r = i;
     while (r < n \&\& s[r] == s[r - 1]) r++;
     z[i] = r - 1; r--;
     int k = i - 1:
     if (z[k] < r - i + 1) z[i] = z[k];
     else {
       1 = i;
       while (r < n \&\& s[r] == s[r - 1]) r
       z[i] = r - 1; r--;
     }
 return z;
```

# 8 Tricks

### 8.1 Array Compression

### 8.2 Different Cumulative Sum

```
///cost = sum of (i * a[i]) where i starts
    from the beginning for every range
///sum[] = prefix sum of value[i]
///isum[] = prefix sum of i*value[i]
```

```
11 cost(int i, int j) {
    ll ret = isum[j];
    if (i) ret -= isum[i - 1];
    ll baad = sum[j];
    if (i) baad -= sum[i - 1];
    return ret - i * baad;
}
```

### 8.3 Fractional Binary Search

```
Given a function f and n. finds the smallest
     fraction p / q in [0, 1] or [0,n]
such that f(p / q) is true, and p, q \le n.
Time: O(log(n))
struct frac { long long p, q; };
bool f(frac x) {
return 6 + 8 * x.p >= 17 * x.q + 12;
frac fracBS(long long n) {
 bool dir = 1, A = 1, B = 1;
 frac lo{0, 1}, hi{1, 0}; // Set hi to 1/0
      to search within [0, n] and {1, 1} to
       search within [0, 1]
  if (f(lo)) return lo;
  assert(f(hi)); //checking if any solution
      exists or not
  while (A || B) {
   long long adv = 0, step = 1; // move hi
        if dir, else lo
    for (int si = 0; step; (step *= 2) >>=
        si) {
     adv += step;
     frac mid{lo.p * adv + hi.p, lo.q * adv
           + hi.q};
     if (abs(mid.p) > n \mid\mid mid.q > n \mid\mid dir
           == !f(mid)) {
       adv -= step; si = 2;
     }
   hi.p += lo.p * adv;
   hi.q += lo.q * adv;
   dir = !dir;
    swap(lo, hi);
   A = B; B = !!adv;
 return dir ? hi : lo;
```

#### 8.4 Time Count

# Equations and Formulas

### Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n}$$
  $C_0 = 1, C_1 = 1$  and  $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$ 

The number of ways to completely parenthesize n+1 factors.

The number of triangulations of a convex polygon with n+2 sides (i.e. the number of partitions of polygon into disjoint triangles) If  $P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$ , then, (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

The number of ways to connect the 2n points on a circle to form n disjoint i.e. non-intersecting chords.

The number of rooted full binary trees with n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex 9.5 has either two children or no children.

Number of permutations of  $1, \ldots, n$  that avoid the pattern 123 (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n=3, these permutations are 132, 213, 231, 312 and 321.

### Stirling Numbers First Kind

The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), where, S(0,0) =$$

$$1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = n!$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$[n - k] = \sum_{0 \le i_1 \le i_2 \le i_k \le n} i_1 i_2 \dots i_k.$$

### Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

 $S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1), \text{ where } S(0,0) =$ 1, S(n,0) = S(0,n) = 0  $S(n,2) = 2^{n-1} - 1$   $S(n,k) \cdot k! = \text{number}$ of ways to color n nodes using colors from 1 to k such that each color is used at least once.

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets, with each subset containing at least r elements. It is denoted by  $S_r(n,k)$  and obeys the recurrence relation.  $S_r(n+1,k) =$ 

$$kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$$

Denote the  $\hat{n}$  objects to partition by the integers  $1, 2, \ldots, n$ . Define the reduced Stirling numbers of the second kind, denoted  $S^d(n,k)$ , to be the number of ways to partition the integers 1, 2, ..., n into k nonempty subsets such that all ele-  $\sum_{k=1}^{n} \frac{n}{\gcd(k, n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k, n)} - 1$ , for n > 1is, for any integers i and j in a given subset, it is required that  $|i-j| \ge d$ . It has been shown that these numbers satisfy,  $\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{i=1}^{n} \mu(d) \lfloor \frac{n}{d} \rfloor^2$  $S^{d}(n,k) = S(n-d+1, k-d+1), n \ge k \ge d$ 

# 9.4 Other Combinatorial Identities

If 
$$P(n) = \sum_{k=0}^{n} \binom{n}{k} \cdot Q(k)$$
, then,
$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$
If  $P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$ , then,
$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

### Different Math Formulas

$$(1-x)(1-x^2)(1-x^3)\cdots = 1-x-x^2+x^5+x^7-x^{12}-x^{15}+x^{22}+x^{26}-\cdots$$

The exponents  $1,2,5,7,12,\cdots$  on the right hand side are given by the formula  $g_k = \frac{k(3k-1)}{2}$  for  $k=1,-1,2,-2,3,\cdots$  and are called (generalized) pentagonal numbers. It is useful to find the partition number in  $O(n\sqrt{n})$ 

Let a and b be coprime positive integers, and find integers a'and b' such that  $aa' \equiv 1 \mod b$  and  $bb' \equiv 1 \mod a$ . Then the number of representations of a positive integers (n) as a non negative linear combination of a and b is

$$\frac{n}{ab}$$
 -  $\left\{\frac{b'n}{a}\right\}$  -  $\left\{\frac{a'n}{b}\right\}$  + 1

#### 9.6GCD and LCM

if m is any integer, then  $gcd(a + m \cdot b, b) = gcd(a, b)$ 

The gcd is a multiplicative function in the following sense: if  $a_1$  and  $a_2$  are relatively prime, then  $gcd(a_1 \cdot a_2, b) = gcd(a_1, b)$ .  $\gcd(a_2,b).$ 

 $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c)).$ 

lcm(a, gcd(b, c)) = gcd(lcm(a, b), lcm(a, c)).

For non-negative integers a and b, where a and b are not both zero,  $gcd(n^a - 1, n^b - 1) = n^{gcd(a,b)} - 1$ 

$$\gcd(a,b) = \sum_{k|a \text{ and } k|b} \phi(k)$$

$$\begin{split} &\sum_{i=1}^{n} [\gcd(i,n) = k] = \phi\left(\frac{n}{k}\right) \\ &\sum_{k=1}^{n} \gcd(k,n) = \sum_{d|n} d \cdot \phi\left(\frac{n}{d}\right) \\ &\sum_{k=1}^{n} x^{\gcd(k,n)} = \sum_{d|n} x^{d} \cdot \phi\left(\frac{n}{d}\right) \\ &\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d) \\ &\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d) \\ &\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for } n > 1 \\ &\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \lfloor \frac{n}{d} \rfloor^{2} \end{split}$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \lfloor \frac{n}{d} \rfloor$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \lfloor \frac{n}{d} \rfloor^{2}$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[\gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i)i^{2}$$

$$F(n) = \sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left( \frac{\left(1 + \left\lfloor \frac{n}{l} \right\rfloor\right) \left(\left\lfloor \frac{n}{l} \right\rfloor\right)}{2} \right)^{2} \sum_{d|l} \mu(d) l d$$