

# 15-210 Assignment BignumLab 1

Roy Sung  
roysung@andrew.cmu.edu  
Section E  
2/12/14

---

## 4: Addition

---

Task 4.1	GEN	GEN	STOP
	GEN	PROP	STOP
	GEN	STOP	STOP

---

## 5: Recurrences

---

### Task 5.1

1. So we just use the tree method. We begin by saying that  $T(n) \leq T(n-1) + k_1 \log(n)$ . By definition of Big theta. Then So if we go down the the tree it only expands into a single branch and at each level we get that  $\log(n-i)$  where  $i$  is the level of the tree. So then we if take the summation of all of the above we get that the final summation is  $\sum_{i=0}^{n-1} \log(n-i)$  This summation reduces down  $\log(n!)$ . Thus we get that  $T(n) \in O(\log(n!))$ . To get big omega we just flip the inequality and go the other direction. With this we get that  $T(n) \in \Omega(\log(n!))$  Thus we have that  $T(n) \in \Theta(\log(n!))$ .
2. So we use the tree method again. We know that there are  $\log(\log(n))$  levels. (We get this from recitation). At each level we have  $n$  work. To see this, we look to the first split. At the split there are  $\sqrt{n}$  nodes and each node doing  $\sqrt{n}$  work. Thus we have that at this level there is  $\sqrt{n} * \sqrt{n} = n$ . So each level we have  $n$  work so the summation is  $\sum_{i=1}^{\log \log(n)} n$ , which will simplify down to  $n \log \log(n)$ . Thus we since we have calculated the exact work we can say that  $T(n) \in \Omega(n \log \log(n))$ .
3. So we start with the base case and this which from the definition of Big Theta we have that  $T(1) \leq c$  where  $c$  is some constant. Thus the base case holds. We assume that for any  $n'$  that is  $1 \leq n' \leq n$  that  $T(n) \leq k_1 n^{\frac{3}{2}}$ . So we start off with the recurrence that  $T(n) = 4T(\frac{n}{4}) + c\sqrt{n}$ . We apply the induction hypothesis on  $T(\frac{n}{4})$  so that we get  $T(n) \leq 4k_1 \frac{n^{\frac{3}{2}}}{4} + c\sqrt{n}$ . We can simplify this down (all the constants will just make another bigger constant  $k_3$  to  $k_3 n^{\frac{3}{2}} + c\sqrt{n} = k_3 \sqrt{n}(n + c)$  We can then lose  $c$  and and we get that  $k_3 n^{\frac{3}{2}}$ . Thus the induction step holds. Since the induction step and the base case holds the we have that  $T(n) \leq kn^{\frac{3}{2}}$ . Thus we have that  $T(n) \in O(n^{\frac{3}{2}})$ . We can then follow the same steps but just switch around the inequality sign to get the proof for the Big omega. Thus we have that  $T(n) \in \Theta(n^{\frac{3}{2}})$ .