3. Other tokens are **ID** and **NUM**, defined by the following regular expressions:

```
ID = letter letter*

NUM = digit digit*

letter = a|..|z|A|..|Z

digit = 0|..|9
```

Lower- and uppercase letters are distinct.

- 4. White space consists of blanks, newlines, and tabs. White space is ignored except that it must separate **ID**'s, **NUM**'s, and keywords.
- 5. Comments are surrounded by the usual C notations /*..*/. Comments can be placed anywhere white space can appear (that is, comments cannot be placed within tokens) and may include more than one line. Comments may not be nested.

A2 SYNTAX AND SEMANTICS OF C-

A BNF grammar for C- is as follows:

```
1. program \rightarrow declaration-list
 2. declaration-list → declaration-list declaration | declaration
 3. declaration \rightarrow var-declaration | fun-declaration
 4. var-declaration → type-specifier ID; type-specifier ID [ NUM ];
 5. type-specifier \rightarrow int | void
 6. fun-declaration → type-specifier ID ( params ) compound-stmt
 7. params \rightarrow param-list \mid void
 8. param-list \rightarrow param-list _{\bullet} param | param
 9. param → type-specifier ID type-specifier ID [ ]
10. compound-stmt \rightarrow { local-declarations statement-list }
11. local-declarations \rightarrow local-declarations var-declaration | empty
12. statement-list → statement-list statement | empty
13. statement → expression-stmt | compound-stmt | selection-stmt
                   | iteration-stmt | return-stmt
14. expession-stmt \rightarrow expression ; | ;
15. selection-stmt \rightarrow if (expression) statement
                       if (expression) statement else statement
16. iteration-stmt \rightarrow while (expression) statement
17. return-stmt → return ; return expression;
18. expression \rightarrow var = expression \mid simple-expression
19. var \rightarrow ID \mid ID \mid expression
20. simple-expression \rightarrow additive-expression relop additive-expression
                               additive-expression
21. relop \rightarrow \langle = | \langle | \rangle | \rangle = | = | 1 =
22. additive-expression \rightarrow additive-expression addop term | term
23. addop \rightarrow + |-
24. term \rightarrow term mulop factor | factor |
25. mulop \rightarrow * /
26. factor \rightarrow (expression) | var | call | NUM
```

```
27. call \rightarrow ID ( args )
```

- **28.** $args \rightarrow arg$ -list | empty
- **29.** arg-list $\rightarrow arg$ -list $\downarrow expression$ | expression

For each of these grammar rules we give a short explanation of the associated semantics.

- 1. program → declaration-list
- 2. declaration-list → declaration-list declaration | declaration
- 3. declaration → var-declaration | fun-declaration

A program consists of a list (or sequence) of declarations, which may be function or variable declarations, in any order. There must be at least one declaration. Semantic restrictions are as follows (these do not occur in C). All variables and functions must be declared before they are used (this avoids backpatching references). The last declaration in a program must be a function declaration of the form **void main(void)**. Note that C—lacks prototypes, so that no distinction is made between declarations and definitions (as in C).

```
4. var-declaration → type-specifier ID; | type-specifier ID [ NUM ];
5. type-specifier → int | void
```

A variable declaration declares either a simple variable of integer type or an array variable whose base type is integer, and whose indices range from $0 ... \mathtt{NUM} - 1$. Note that in C- the only basic types are integer and void. In a variable declaration, only the type specifier **int** can be used. **Void** is for function declarations (see below). Note, also, that only one variable can be declared per declaration.

- 6. fun-declaration → type-specifier ID (params) compound-stmt
- 7. $params \rightarrow param-list \mid \mathbf{void}$
- 8. $param-list \rightarrow param-list$, param param
- **9.** param \rightarrow type-specifier **ID** | type-specifier **ID** []

A function declaration consists of a return type specifier, an identifier, and a comma-separated list of parameters inside parentheses, followed by a compound statement with the code for the function. If the return type of the function is **void**, then the function returns no value (i.e., is a procedure). Parameters of a function are either **void** (i.e., there are no parameters) or a list representing the function's parameters. Parameters followed by brackets are array parameters whose size can vary. Simple integer parameters are passed by value. Array parameters are passed by reference (i.e., as pointers) and must be matched by an array variable during a call. Note that there are no parameters of type "function." The parameters of a function have scope equal to the compound statement of the function declaration, and each invocation of a function has a separate set of parameters. Functions may be recursive (to the extent that declaration before use allows).

10. compound-stmt \rightarrow { local-declarations statement-list }

A compound statement consists of curly brackets surrounding a set of declarations and statements. A compound statement is executed by executing the statement

sequence in the order given. The local declarations have scope equal to the statement list of the compound statement and supersede any global declarations.

- 11. local-declarations → local-declarations var-declaration | empty
- **12.** statement-list \rightarrow statement-list statement | empty

Note that both declarations and statement lists may be empty. (The nonterminal *empty* stands for the empty string, sometimes written as ε .)

```
13. statement → expression-stmt | compound-stmt | selection-stmt | iteration-stmt | return-stmt | return-stmt | compound-stmt → expression ; | ;
```

An expression statement has an optional expression followed by a semicolon. Such expressions are usually evaluated for their side effects. Thus, this statement is used for assignments and function calls.

```
15. selection-stmt → if (expression) statement | if (expression) statement else statement
```

The if-statement has the usual semantics: the expression is evaluated; a nonzero value causes execution of the first statement; a zero value causes execution of the second statement, if it exists. This rule results in the classical dangling else ambiguity, which is resolved in the standard way: the else part is always parsed immediately as a substructure of the current if (the "most closely nested" disambiguating rule).

```
16. iteration-stmt \rightarrow while (expression) statement
```

The while-statement is the only iteration statement in C-. It is executed by repeatedly evaluating the expression and then executing the statement if the expression evaluates to a nonzero value, ending when the expression evaluates to 0.

```
17. return-stmt → return ; | return expression ;
```

A return statement may either return a value or not. Functions not declared **void** must return values. Functions declared **void** must not return values. A return causes transfer of control back to the caller (or termination of the program if it is inside **main**).

```
18. expression \rightarrow var = expression \mid simple-expression
19. var \rightarrow ID \mid ID \mid expression
```

An expression is a variable reference followed by an assignment symbol (equal sign) and an expression, or just a simple expression. The assignment has the usual storage semantics: the location of the variable represented by var is found, then the subexpression to the right of the assignment is evaluated, and the value of the subexpression is stored at the given location. This value is also returned as the value of the entire expression. A var is either a simple (integer) variable or a subscripted array variable. A negative subscript causes the program to halt (unlike C). However, upper bounds of subscripts are not checked.

Vars represent a further restriction of C- from C. In C the target of an assignment must be an **l-value**, and **l-values** are addresses that can be obtained by many operations. In C- the only **l-values** are those given by the *var* syntax, and so this category is checked syntactically, instead of during type checking as in C. Thus, pointer arithmetic is forbidden in C-.

```
    20. simple-expression → additive-expression relop additive-expression | additive-expression
    21. relop → <= | < | > | >= | !=
```

A simple expression consists of relational operators that do not associate (that is, an unparenthesized expression can only have one relational operator). The value of a simple expression is either the value of its additive expression if it contains no relational operators, or 1 if the relational operator evaluates to true, or 0 if it evaluates to false.

```
22. additive-expression \rightarrow additive-expression addop term | term
```

- 23. $addop \rightarrow + | -$
- 24. term → term mulop factor | factor
- 25. $mulop \rightarrow * /$

Additive expressions and terms represent the typical associativity and precedence of the arithmetic operators. The / symbol represents integer division; that is, any remainder is truncated.

A factor is an expression enclosed in parentheses, a variable, which evaluates to the value of its variable; a call of a function, which evaluates to the returned value of the function; or a NUM, whose value is computed by the scanner. An array variable must be subscripted, except in the case of an expression consisting of a single ID and used in a function call with an array parameter (see below).

```
27. call → ID (args)
28. args → arg-list | empty
29. arg-list → arg-list , expression | expression
```

A function call consists of an **ID** (the name of the function), followed by parentheses enclosing its arguments. Arguments are either empty or consist of a commaseparated list of expressions, representing the values to be assigned to parameters during a call. Functions must be declared before they are called, and the number of parameters in a declaration must equal the number of arguments in a call. An array parameter in a function declaration must be matched with an expression consisting of a single identifier representing an array variable.

Finally, the above rules give no input or output statement. We must include such functions in the definition of C-, since unlike C, C- has no separate compilation or linking facilities. We, therefore, consider two functions to be predefined in the global environment, as though they had the indicated declarations:

```
int input(void) { . . . }
void output(int x) { . . . }
```

The **input** function has no parameters and returns an integer value from the standard input device (usually the keyboard). The **output** function takes one integer parameter, whose value it prints to the standard output (usually the screen), together with a newline.

A3 SAMPLE PROGRAMS IN C-

The following is a program that inputs two integers, computes their greatest common divisor, and prints it:

```
/* A program to perform Euclid's
   Algorithm to compute gcd. */
int gcd (int u, int v)
{ if (v == 0) return u;
   else return gcd(v,u-u/v*v);
   /* u-u/v*v == u mod v */
}

void main(void)
{ int x; int y;
   x = input(); y = input();
   output(gcd(x,y));
}
```

The following is a program that inputs a list of 10 integers, sorts them by selection sort, and outputs them again:

```
/* A program to perform selection sort on a 10
  element array. */
int x[10]:
int minloc ( int a[], int low, int high )
{ int i; int x; int k;
  k = low;
  x = a[low];
  i = low + 1;
  while (i < high)
    { if (a[i] < x)
         {\mathbf x = \mathbf a[i]}
           k = i; }
      i = i + 1
    }
  return k;
}
```

```
void sort( int a[], int low, int high)
{ int i; int k;
  i = low;
  while (i < high-1)
    { int t;
      k = minloc(a,i,high);
      t = a[k];
      a[k] = a[i];
      a[i] = t;
      i = i + 1:
    }
}
void main(void)
{ int i;
  i = 0;
  while (i < 10)
    {x[i] = input();}
       i = i + 1;
  sort(x,0,10);
  i = 0;
  while (i < 10)
     { output(x[i]);
       i = i + 1;
}
```

A4 A TINY MACHINE RUNTIME ENVIRONMENT FOR THE C— LANGUAGE

The following description assumes a knowledge of the Tiny Machine as given in Section 8.7 and an understanding of stack-based runtime environments from Chapter 7. Since C- (unlike TINY) has recursive procedures, the runtime environment must be stack based. The environment consists of a global area at the top of dMem, and the stack just below it, growing downward toward 0. Since C- contains no pointers or dynamic allocation, there is no need for a heap. The basic organization of each activation record (or stack frame) in C- is

