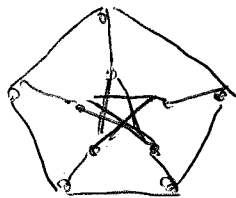


Discrete Math

counting
relations

graph theory



pentagram
inside
pentagon



not graphs
like ~~A~~

group project: 1. mathematical & latex

2. Four programs
lecture 4/4

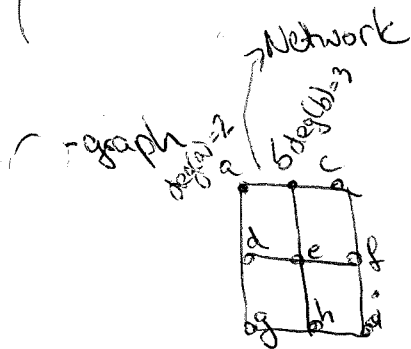
3. Project problems

Presentations 17/5

- { 1. Recurrence
- { 2. Relations

Game of Life

- { 3. Cellular Automata
- { 4.



vertex
 $V = \{a, b, c, \dots, i\}$

find all the pairs

$E = \{(a, b), (a, d), \dots, (h, i)\}$

vertices with degree ~~2~~ → which are attached to it

2: a, c, g, i $\deg(a) = 2$

3: b, d, f, h

4: e

Handshake theorem: $2e = \sum_{v \in V} \deg(v) = 4 \cdot 2 + 4 \cdot 3 + 4 = 24$

$e = \# \text{ edges}$

$e = 12$

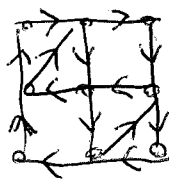
we have 12 edges

vertices with large degrees are nowadays called
hub

The postman wants to go over an edge exactly once!

Euler circuit

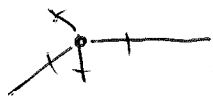
sufficient and necessary conditions



Theorem: The graph has an Euler circuit

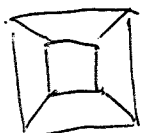
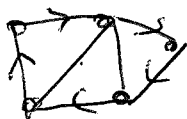
proof by induction \leftarrow sufficient condition \Rightarrow \Leftarrow
all degrees of the vertices are even numbers

Proof:



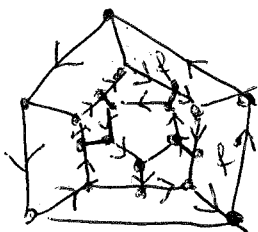
\rightarrow we come in and we go out

Travelling salesman wants to visit each town exactly once



$$\begin{aligned} f &= 6 \\ e &= 12 \\ v &= 8 \end{aligned}$$

$$v + f - e = 2$$



$$v = 20 \rightarrow \text{vertices}$$

$$e = 30 \rightarrow \text{edges}$$

$$f = 12 \rightarrow \text{faces}$$

$$e - v + f = 2$$

He wants to make a

Hamilton cycle

Sufficient condition for a Hamilton cycle

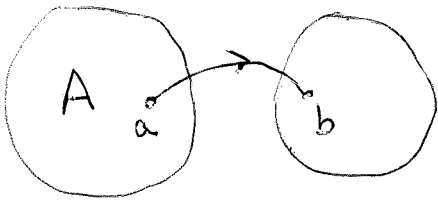
$\deg(u) + \deg(v) \geq v$ where u and v are
two non-adjacent vertices.

Necessary condition

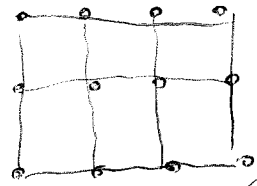
6 pentagons inside

6 pentagons outside

Bijection - $\begin{cases} \text{one-to-one} \\ \text{on to} \end{cases}$



Graph



$V = 12$

vertices in G

Edges in HC

Ex.) $f: \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z} \rightarrow$ Is this function on to?

$$f(m, n) = m^2 - n^2 = (m-n)(m+n)$$

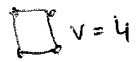
Can you find pairs (m, n)

such that $f(m, n) = 0$ $f(2, 2) = 0$

$f(m, n) = 1$ $f(1, 0) = 1$

$f(m, n) = 2$

If there is a HC in a graph with V vertices then the HC has 66... edges



$V = 4$

$$2 \cdot 1 = (m+n)(m-n)$$

$$m+n=2$$

$$2m=3$$

$$m-n=1$$

$$m = \frac{3}{2}$$

Answer: we can't get 2 with such a function!

splitting up: $60 = 2^2 \cdot 15 = 2^2 \cdot 3 \cdot 5$ (into primes)

Arithmetic sum: $S_n = 1 + 2 + 3 + 4 + \dots + n$

$$= \frac{n \cdot (n+1)}{2}$$

$$S_1 = 1$$

$$S_2 = 1 + 2 = 3$$

$$S_3 = 1 + 2 + 3 = 6$$

Proof 1: $S_n = 1 + 2 + \dots + n$

$$S_n = n + (n-1) + \dots + 1$$

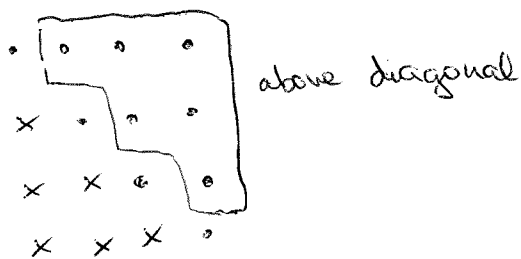
Add them $2S_n = (n+1) + (n+1) + \dots + (n+1)$

$$= n \cdot (n+1)$$

$$S_n = \frac{n \cdot (n+1)}{2}$$

Proof 2:

$n=4$



28.3.2017

(2)

$$S_4 = \underbrace{\frac{(4^2 - 4)}{2}}_{\text{# points above diagonal}} + \underbrace{4}_{\text{on the diagonal}} = \frac{4^2 + 4}{2} = \frac{4(4+1)}{2}$$

points
above diagonal

$$S_n = \frac{(n^2 - n)}{2} + n = \frac{n^2 + n}{2} = \frac{n(n+1)}{2}$$

Proof 3: by induction

well-ordering
principle

Every subset of the positive
integers has a smallest element



$$S_n = \frac{n(n+1)}{2} = 1+2+\dots+n$$

$$S_1 = 1 = \frac{1 \cdot (1+1)}{2} = 1 \quad \text{OK}$$

$$S_2 = 1+2 = 3 = \frac{2 \cdot (2+1)}{2} = 3 \quad \text{OK}$$

Assume $S_k = \frac{k(k+1)}{2}$ Induction Assumption

$$S_{k+1} = \underbrace{1+2+3+\dots+k}_{\frac{k(k+1)}{2}} + (k+1) \stackrel{\text{Induction Assumption}}{=} \frac{k(k+1)}{2} + (k+1)$$

$$= (k+1) \left[\frac{k}{2} + 1 \right] = \frac{(k+1)(k+2)}{2} \quad \text{yes!}$$

$$= \frac{(k+1)(k+1+1)}{2} = \frac{(k+1)(k+2)}{2} \quad \text{← dreamy result}$$



exclude \rightarrow 2×2

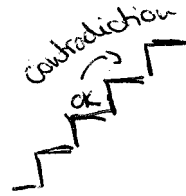
exclude \rightarrow 4×4

8×8

16×16

\vdots

Prove by Induction, that it is always possible to tile $2^m \times 2^m$ boards with L-shaped figures.

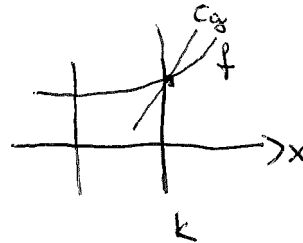


Big-oh notation

$$\frac{n^2}{2} < S_n = \frac{n^2 + n}{2} < n^2$$

So S_n is $O(n^2)$

f is $O(g)$ if $f \leq Cg$ if $x > k$
alias function
2 constants



$$f(n) = \frac{n(n+1)}{2}$$

$$g(n) = n^2$$

$$C = 1$$

$$k = 1$$

For large n :

$$n! > 2^n > n^2 > n \log n > \log n > 1$$

Bit

A bit operation takes 10^{-9} s. 2^{100} bit operations takes ≈ 1000 years

Multiplication

$$(1101)_2 = 13 = 1 \cdot 2^3 + 1 \cdot 2^2 + 0 \cdot 2^1 + 1 \cdot 2^0$$

$$(1010)_2 = 10$$

$$13 \cdot 10 = 130$$

$T(n)$ = the time for multiplication of two n digit numbers.

$$T(n) \text{ is } O(?)$$

$$\uparrow$$

$$n^2$$

$$\begin{array}{r} 1101 \\ 1010 \\ \hline 0000 \\ 1101 \\ \hline 0000 \\ 1101 \\ \hline 1000010 \end{array}$$

$$1000010 = 2^7 + 2^1 = 128 + 2 = 130 \quad T(n) \text{ is } O(n^2)$$

$$1101 = 2^2 \cdot 11 + 01$$

$\underbrace{\quad\quad}_{A_1} \quad \underbrace{\quad\quad}_{A_0}$

$$1010 = 2^2 \cdot 10 + 10$$

$$A_1 = (11)_2 = 2^1 + 2^0 = (3)_{10}$$

$$T(2n) = 3T(n) + C_1 \cdot n$$

$$\vdots$$

$$T(n) = O(n^{\log_2 3}) \approx 1.6$$

28.3.2017

(4)

$$(2^2 A_1 + A_0) \cdot (2^2 B_1 + B_0) = (2^4 A_1 B_1 + 2^2 A_1 B_0 + 2^2 (A_1 - A_0)(B_0 - B_1) + (2^2 + 1) A_0 B_0)$$

Fermat's little theorem

$a^p \equiv a \pmod{p}$ if p is a prime
p > a

$$35 \equiv 25 \pmod{5}$$

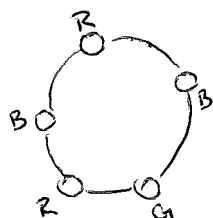
$$\text{since } \frac{35-25}{5} = 2$$

$$\text{Is } 36 \equiv 20 \pmod{17}?$$

No!

$$a=3 \quad 3^5 \equiv 3 \pmod{5}$$

$$p=5 \quad \frac{3^5 - 3}{5} = \frac{243 - 3}{5} = 48$$



5 bricks red blue green
3 colors R, B, G

Total number of colorings
= 3^5

Exclude the 3 uni-colored necklaces

Left is $3^5 - 3$ necklaces

There are only $\frac{3^5 - 3}{5}$ unique necklaces

→ because rotated would not be unique

48

does it work if it is not with a prime number?

$6 = 3 \cdot 2$
think of it!





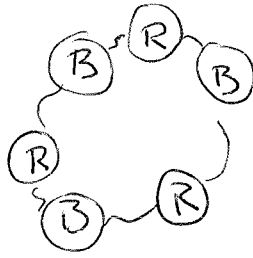
ABC → 123

$$a^p \equiv a \pmod{p}$$

$$\frac{3^5 - 3}{5} = 48$$

$$\frac{5^3 - 5}{3} = 40$$

$$\frac{2^6 - 2}{6} = \frac{62}{6}$$



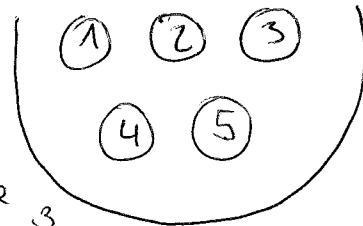
In how many ways ...?

Combinatorics

* Product rule

* Sum rule

Balls in a bowl



Take 3 of them

with respect to order?

Yes

5 5 1



No

$$\binom{7}{3}$$

New!

Here we are only interested of how many of each sort
 two fives & one three
 553 355

So our problem is really about the number of integer solutions to the equation

$$x_1 + x_2 + x_3 + x_4 + x_5 = 3$$

$$x_i \in \mathbb{N}$$

$$0 + 0 + 1 + 0 + 2 = 3$$

$x_1 = \#$ of 1s. (Number of Balls with 1 on it)
 \vdots
 $x_5 = \#$ of 5s.

Back to the stone age! + + + + + + +

+ + + + + + +
 7 positions

$$1 + 1 + 0 + 1 + 0$$

$$1 + 1 + + 1 + \Rightarrow 7 \text{ pos}$$

$$0 + 0 + 3 + 0 + 0 \Leftrightarrow + + ||| + + \Rightarrow 7 \text{ pos}$$

Allow for repetitions, replacing the ball back in

Yes

using the product rule

542, 245

another outcome

No

$$5 \cdot 4 \cdot 3 =$$

$$\frac{5!}{2!} = P(5,3)$$

$$\frac{5 \cdot 4 \cdot 3}{1 \cdot 2 \cdot 3} =$$

$$\frac{5!}{2!3!} = \binom{5}{3}$$

$$= C(5,3) =$$

$$10 = \binom{5}{2}$$

In how many ways can we put 3 sticks on 7 positions?

$$\frac{7 \cdot 6 \cdot 5}{1 \cdot 2 \cdot 3} = \binom{7}{3}$$

13 balls, 6 repetitions
we don't care about order.

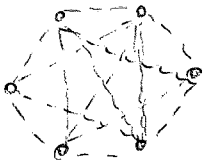
In how many ways? (sticks, ...)

$$\binom{12+6}{6} = \binom{18}{6} = \binom{18}{12}$$

Ex: 6 people shake hands with each other

How many handshakes?

Ex: K_6

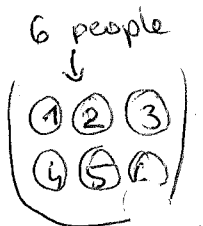


coloring of the edges
with red or blue.

How many colorings?

$$\binom{2}{15} = 2^{15}$$

$$n=6, r=2, \binom{6}{2}=15$$



we take out 2

combinatorics → no repetition
no respect to order

$$n=2, r=15$$

$$\frac{6!}{2!4!} = \frac{6 \cdot 5}{2} = 15$$

$$= 5+4+3+2+1$$

Sum rule

How many bit strings of length 8 have
at least 3 zeros?

0011110
0111111

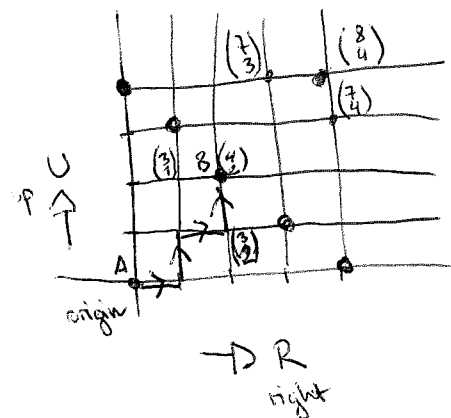
"universe" in total we have 2^8 bit strings

0 zero	At least 3 zeros
1 zero	
2 zeros	

□□□□□□□□

$$N(\text{at least 3 0s}) = 2^8 - \underbrace{N(\text{0 0s})}_1 - \underbrace{N(\text{1 zero})}_{8 = \binom{8}{1}} - \underbrace{N(\text{2 zeros})}_{\binom{8}{2} = \frac{8 \cdot 7}{2} = 28}$$

$$= 256 - 1 - 8 - 28 = 256 - 37 = 219$$



ways to go from
A to B = $\binom{4}{2} = 6$

RURU (right, up, right, up) $\frac{4 \cdot 3}{2} = 6$

come from low come from left

$$\binom{4}{2} = \binom{3}{2} + \binom{3}{1}$$

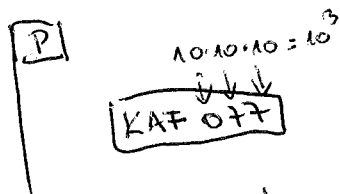
$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$$

$$\binom{8}{4} = \binom{7}{3} + \binom{7}{4}$$

4 ways

$$\binom{4}{0} + \binom{4}{1} + \binom{4}{2} + \binom{4}{3} + \binom{4}{4} = 2^4 = 16$$

Box-principle



1001 swedish crowns

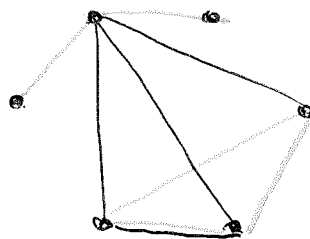
At least two cars have the same number combination!

understand this!

← Parking lot

Ramsey theory

$$5 = 4 + 1 = 3 + 2 = 2 + 3$$



Edge coloring of K_6 with 2 colors, green and black

A green or a black triangle is unavoidable

in a group of 6 people you can find 3 who are complete strangers

$$\leftarrow R(3,3) = 6$$

$$R(4,4) = 18$$

$$43 \leq R(5,5) \leq 49 \rightarrow \text{impossible}$$

$$102 \leq R(6,6) \leq 165 \rightarrow \text{to compute}$$

I + moves!

Geometric Series

$$S = 1 + k + k^2 + k^3 = \frac{k^4 - 1}{k - 1}$$

$$a^3 - b^3 = (a - b)(a^2 + ab + b^2)$$

$$a^4 - b^4 = (a - b)(a^3 + a^2b + ab^2 + b^3)$$

⋮
⋮

① Linear RR

external influence

memory	NO <small>homogeneous</small>	YES <small>non-homogeneous</small>
1	$a_0 = 100 \text{ IV}$ $a_{n+1} = 1.03 a_n$ $a_n = 100 \cdot 1.03^n$ $n = 0, 1, 2, 3, \dots$	$a_{n+1} = 1.03 a_n + 100$
2	$a_n = a_{n-1} + 6 a_{n-2}$ $a_0 = 2$ $a_1 = 5$ $a_2 = 17$	$a_n = a_{n-1} + a_{n-2} + 2^n$

How to solve the RR?

1.) R Solve

2.) Find a pattern. Prove your formula by induction

3.) Matrices

4.) Two solutions is also a solution!

Add them and fix the constants by initial values.

Ex.) $a_{n+1} = 2a_n + 2^n$
 $a_0 = 1$

Use Method ② → find pattern

Ex.) $a_{n+1} = 3a_n + 2^n$
 $a_0 = 1$

$$a_1 = 3 \cdot 1 + 2^0 = 4$$

$$a_2 = 3 \cdot 4 + 2^1 = 14$$

$$a_2 = 3(3 \cdot 1 + 2^0) + 2^1$$

$$= 3^2 + 3 \cdot 2^0 + 2^1 = 14$$

$$a_3 = 3(3^2 + 3 \cdot 2^0 + 2^1) + 2^2$$

$$= 3^3 + 3^2 \cdot 2^0 + 3 \cdot 2^1 + 2^2 \Rightarrow$$

$$3^3 + \frac{3^3 - 2^3}{3 - 2} = 2 \cdot 3^3 - 2^3$$

$$54 - 8 = 46$$

Guery:

$$a_n = 2 \cdot 3^n - 2^{n+1}$$

④, Fix the particular solution

$$a_n = C \cdot 2^n$$

plug it in

$$C \cdot 2^{n+1} = 3C \cdot 2^n + 2^n$$

$$(2c - 3c)2^n = 2^n$$

$$C = -1$$

$$a_n = -2^n$$

ii) $a_{n+1} = 3a_n$ $a_n = c_1 3^n$

$$a_n = c_1 3^n$$

iii) put them together

$$a_n = c_1 3^n - 2^n$$

Determine C_1 by IV:

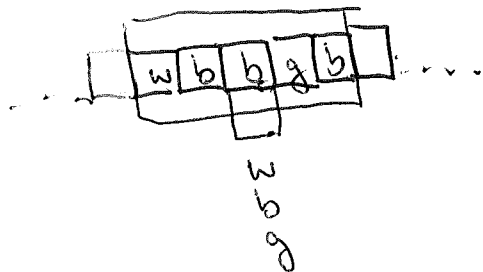
$$a_0 = 1 = C_1 3^0 - 2^0$$

 \Rightarrow

$$1 = C_1 - 1$$

$$c_1 = 2$$

The Feigenbaum Constant - YouTube Numberphile



$$3 \cdot 3 \cdot 3 \cdot 3 \cdot 3 = 243$$

3²⁴³

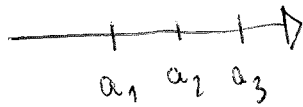
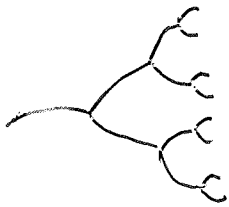
$3^{243} \rightarrow$ so many roles

Game of life!

YouTube Numberphile

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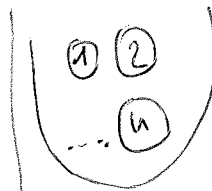
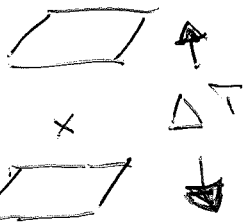
(1)



$$\frac{a_3 - a_2}{a_2 - a_1} \approx 4.66...$$

for all humps!

Turbulence



pick r

put them back

care about order?

	Y	N
Y	n^r	$\binom{n+r-1}{r}$
N	$\frac{n!}{(n-r)!}$	$\binom{n}{r}$

Sum Rule

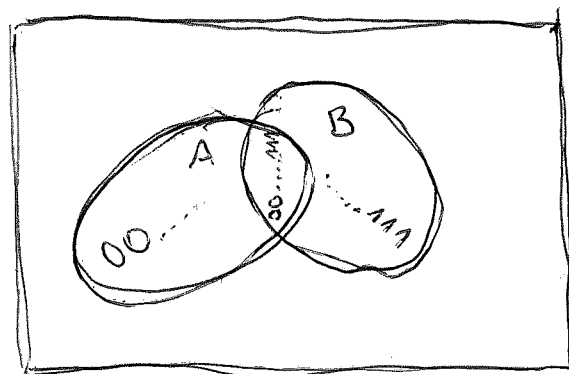
Product Rule

Today Inclusion-Exclusion +
vertex colorings of graphs

bit strings of length 10 that start
with 00.....

or end with 111 = ?

All bit strings of length 10 = 2^{10}



$$|A \cup B| =$$

$$|A| + |B| - |A \cap B|$$

$$= 2^8 + 2^7 - 2^5 =$$

$$2^5(8+4-1)$$

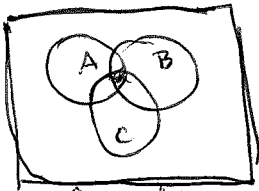
$$= 12 \cdot 32$$

$$= 384$$

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3-sets

(2)



$$|A \cup B \cup C| = |A| + |B| + |C| -$$

$$|A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$$

consider an element in $A \cap B \cap C$:

$$1 + 1 + 1 - 1 - 1 - 1$$

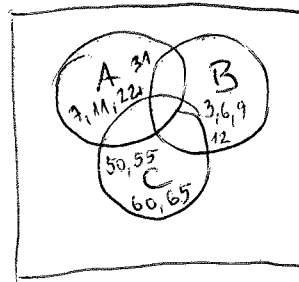
8.5.19 Pick 4 numbers from one to 100

A - all 4 numbers odd

B - multiples of 3

C - multiples of 5

$$|A \cup B \cup C| = ?$$



Beware! Maybe wrong figure!!!

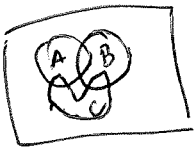
odd number = 50

multiples of 3 = 33

multiples of 5 = 20

15, 30, 45, 60, 75, 90
wrong figure because we can't pick 4 -> only 3

correct figure



$$|A \cup B \cup C| = \frac{50 \cdot 49 \cdot 48 \cdot 47}{1 \cdot 2 \cdot 3 \cdot 4} + \binom{3}{4} + \binom{20}{4} = \binom{6}{4} + \binom{10}{4}$$

onto functions

$$\{a, b, c, d, e\} \rightarrow \{1, 2, 3\}$$



$$\# \text{ functions} = 3^5$$

But here all boxes are not empty!

First without Inclusion-Exclusion

$$\text{sum } 3 + 1 + 1$$

$$\frac{|abc|}{|c|} \cdot \frac{|d|}{|d|}$$

$$\binom{5}{3} \cdot 3 \cdot 2$$

$$2 + 2 + 1$$

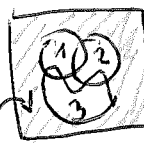
$$\frac{|bdc|}{|dc|} \cdot \frac{|ae|}{|ae|}$$

$$+ \frac{\binom{5}{2} \cdot \binom{3}{2} \cdot 3!}{2} = 60 + \frac{10 \cdot 3 \cdot 6}{2}$$

$$= 60 + 90 = 150$$

with I-E

$$3^5 = 243$$



here are the onto functions

1 - the set where container 1 is empty

2 - container 2 empty

3 - all 3 empty

6.4.2017

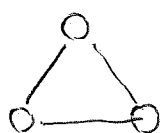
(3)

→ containers continued

$$3^5 - 3 \cdot 2^5 + 3 \cdot 1^5 = 243 - 96 + 3 = 246 - 96 = 150$$

We want to keep things apart

⇒ vertex-coloring of graphs



with 2 colors we can do 2^3 colorings of the Δ

3 colors

3^3

x^3

x colors

Proper colorings (PC): Vertices with an edge in common must have different colors

The smallest number of colors needed for a PC is called the chromatic number χ

#PC of Δ using 2 colors = 0

3 colors = $3 \cdot 2 \cdot 1 = 6$

$x = x \cdot (x-1) \cdot (x-2) = P_{\Delta}(x)$

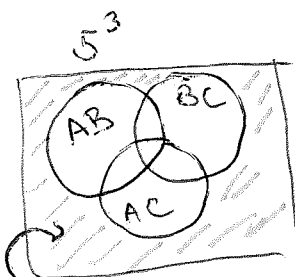
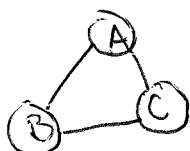
$P(5) = 5 \cdot 4 \cdot 3 = 60$

ways of coloring

chromatic polynomial

by I-E

Δ and 5 colors



here are the PC

AB - vertex A and B

have the same color

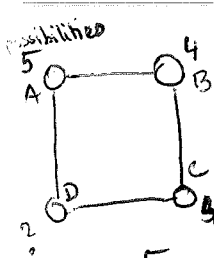
intersections all uncolored Δ

$$5^3 - 3 \cdot 5^2 + 3 \cdot 5^1 - 5^0 = 60$$

$$= 125 - 75 + 15 - 5 = 60$$

so in set AB vertex A and B have same color





PC with same color on A and C = $5 \cdot 4 \cdot 4 = 80$ (4)

PC with different colors on A and C = $5 \cdot 4 \cdot 3^2 = 180$

260

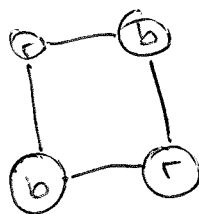
$X = 5$

Colors

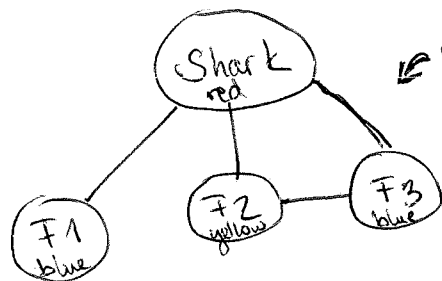
How many PCs?

$P_4(5) = 260$

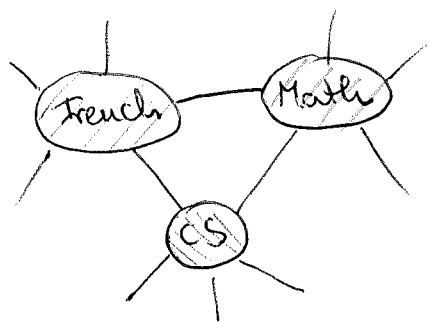
$\chi_4 = 2$ ← minimum number needed to make a proper coloring



Do it with I-E!



← must be apart! not together in aquarium $\Rightarrow \chi_G = 3$



χ_G gives the minimum number of exam days

Generating Functions (GF) an alternative to I-E.

11.4.2017

①

x - very small

$$S_3 = 1 + x + x^2 + x^3 \leftarrow \text{polynomial with coefficient 1}$$

we don't care about x, only about the number in front of it

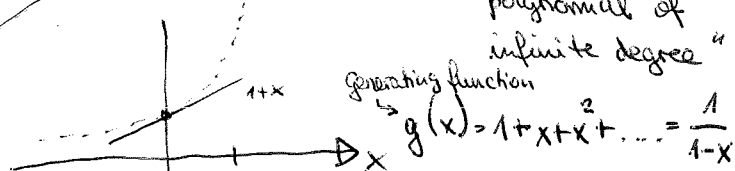
$$x \cdot S_3 = x + x^2 + x^3 + x^4$$

$$S_3 - x \cdot S_3 = 1 + x - x + x^2 - x^2 + x^3 - x^3 + x^4 = 1 - x^4$$

$$S_3 = \frac{1 - x^4}{1 - x} \quad x \neq 1$$

$$S_n = \frac{1 - x^{n+1}}{1 - x} \quad \text{if } n \rightarrow \infty \text{ goes to infinity} \quad S_\infty = \frac{1}{1 - x}$$

$= 1 + x + x^2 + x^3 + \dots$
"polynomial of infinite degree"



$$g(x) = 1 + x + x^2 + \dots = \frac{1}{1 - x}$$

$$g'(x) = \frac{1}{(1-x)^2} = 1 + 2x + 3x^2 + 4x^3 + \dots$$

$$g''(x) = \frac{2}{(1-x)^3} = 2 + 6x + 12x^2 + \dots$$

$$x \cdot x = \frac{1}{2}$$

$$\frac{1}{1 - \frac{1}{2}} = 2 = 1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \frac{1}{16} + \dots$$

$$\frac{1}{(1-x)^3} = 1 + 3x + 6x^2 + \dots$$

GF	Number sequence
$\frac{1 - x^4}{1 - x}$	1, 1, 1, 1, 0, 0, ...
$\frac{1}{1 - x}$	1, 1, 1, 1, 1, ...
$\frac{1}{(1 - x)^2}$	1, 2, 3, 4, 5, ...
$\frac{1}{(1 - x)^3}$	1, 3, 6, ..., $\binom{3+n-1}{n}$, ...

Question ①: Generating function for $\binom{3}{0}, \binom{3}{1}, \binom{3}{2}, \binom{3}{3}, \dots, 0, \dots, 0, \dots$

Question ②: Number sequence for $\frac{1}{1 - 2x}$ is

Question ③: Generating function for 1, 1, 0, 1, 1, 1, 1, ... is

$$\textcircled{1}: \binom{3}{0} + \binom{3}{1}x + \binom{3}{2}x^2 + \binom{3}{3}x^3 = (1 + x)^3$$

②: 1, 2, 4, 8, 16, 32, ... \rightarrow is the answer!
think about it

$$\textcircled{3}: \frac{1}{1-x} - \frac{2}{x} = \frac{1 - x^2 + x^3}{1 - x}$$

$$(a+b)^n = \sum_{k=0}^n \binom{n}{k} a^k b^{n-k}$$

Binomial theorem

reminder: power laws

$$\text{Ex: } x^4 \cdot x^6 \cdot x^7 = x^{4+6+7} = x^{17}$$

(2)

$$x^5 \cdot x^6 \cdot x^3 = x^{5+6+3} = x^{14}$$

The GF for this problem is:

$$G(x) = \underbrace{(1+x+x^2+x^3+\dots)}_{\text{child 1}} \underbrace{(1+x+x^2+x^3+\dots)}_{\text{child 2}} \underbrace{(1+x+x^2+x^3+\dots)}_{\text{child 3}}$$

What is the coefficient in front of x^{17} ?

$$G(x) = \frac{1}{(1-x)^3} = 1 + 3x + 6x^2 + \dots + \binom{19}{17} x^{17}$$

In how many ways can we distribute 17 cakes to 3 kids.

This is the number of integer solutions to $e_1 + e_2 + e_3 = 17$

$$\text{Ans. } \binom{19}{17} = \binom{19}{2} = \frac{19 \cdot 18}{2} = 171$$

$$\frac{19!}{2! \cdot 17!} \quad \text{cancel out the common factorial factors}$$

$$\binom{3+n-1}{n} \quad n=17$$

$$\binom{19}{17} = \frac{x^2(1+x+x^2+x^3) \cdot x^3(1+x+x^2+x^3) \cdot x^3(1+x+x^2+x^3)}{x^2(1+x+x^2+x^3) \cdot x^3(1+x+x^2+x^3) \cdot x^3(1+x+x^2+x^3)}$$

$$GF: G(x) = \underbrace{(x^2+x^3+x^4+x^5)}_{\text{child 1}} \underbrace{(x^3+x^4+x^5+x^6)}_{\text{child 2}} \underbrace{(x^4+x^5+x^6+x^7)}_{\text{child 3}} =$$

what is the coefficient in front of x^{17} ?

$$= x^2 \cdot x^3 \cdot x^4 \cdot (1+x+x^2+x^3)^3 =$$

$$= x^9 \frac{(1-x^4)^3}{(1-x)^3} = x^9 (1-3x^4+3x^8-x^{12}) \cdot (1+3x+6x^2+\dots)$$

$$\left(1+3x+6x^2+\dots+\binom{n-1}{n}x^{n-1}\right)$$

How can we get x^{17} ?

$$i) x^9 \cdot 1 \cdot \binom{10}{8} x^8$$

$$ii) x^9 \cdot (-3x^4) \cdot \binom{6}{4} x^4$$

$$iii) 3 \cdot x^9 \cdot x^8$$

$$\text{In total: } \boxed{\binom{10}{8} - 3\binom{6}{4} + 3} = 45 - 45 + 3 = 3$$

$$\text{Ex: } e_1 + e_2 + e_3 = 17$$

$$2 \leq e_1 \leq 5$$

$$3 \leq e_2 \leq 6 \quad \text{sum of maximum values: 18}$$

$$4 \leq e_3 \leq 7$$

Q. In how many ways can we distribute these 17 cakes with the restrictions above?

Ans: 3

Previous example

with I-E instead

$$e_1 + e_2 + e_3 = 17$$

$$2 \leq e_1 \leq 5$$

$$3 \leq e_2 \leq 6$$

$$4 \leq e_3 \leq 7$$

$$e_1 = 2 + y_1$$

$$e_2 = 3 + y_2$$

$$e_3 = 4 + y_3$$

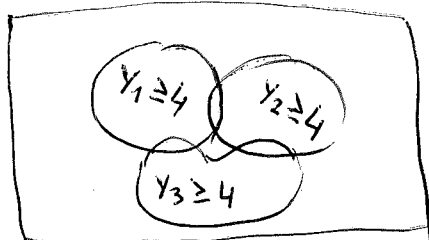
$$2 + y_1 + 3 + y_2 + 4 + y_3 = 17$$

$$0 \leq y_i \leq 3$$

$$y_1 + y_2 + y_3 = 8$$

$$0 \leq y_i \leq 3$$

All solutions without upper restrictions = $\binom{10}{2} = \binom{10}{8}$



$$I) y_1 \geq 4$$

$$y_1 = 4 + z_1$$

$$4 + z_1 + y_2 + y_3 = 8$$

$$z_1 + y_2 + y_3 = 4$$

$$z_1 \geq 0$$

$$y_2 \geq 0$$

$$y_3 \geq 0$$

$$\text{Ans.: } \binom{6}{2} = \binom{6}{4}$$

Throw

E: 5 dice



$$S = x_1 + x_2 + \dots + x_5$$

$$5 \leq S \leq 30$$

in how many ways can we get 17?

in Function + I-E!

Ex.) → not important! will not come to exam!!

$$a_{n+1} = 2a_n + 2^n$$

$$a_0 = 1$$

$$a_n = (n+2) \cdot 2^{n-1}$$

$$a_0 = 1, a_1 = 3, a_2 = 8, \dots$$

$$\sum_{n=0}^{\infty} a_{n+1} \cdot x^{n+1} = \sum_{n=0}^{\infty} 2a_n \cdot x^{n+1} + \sum_{n=0}^{\infty} 2^n \cdot x^{n+1}$$

$$(G(x) - 1) = 2xG(x) + \frac{x}{1-2x}$$

$$G(x) \cdot (1-2x) = 1 + \frac{x}{1-2x} = \frac{1-x}{1-2x}$$

$$G(x) = \frac{1-x}{(1-2x)^2}$$

ways to get sum 17
throwing 5 dice =

$$\binom{16}{12} - 5 \binom{10}{6} + \binom{5}{2} \cdot 1$$

Block chains "A new kind
of trust"

10 likes \geq colleague

70 likes \geq friend

300 likes \geq partner!

Cambridge Analytica

Relations (This week)

1) Reflexive

2) Symmetric

3) Anti-symmetric

4) Transitive

Posets: 1, 3, 4 \geq \geq

Have diagram, lattice

Equivalence relation 1, 2, 4

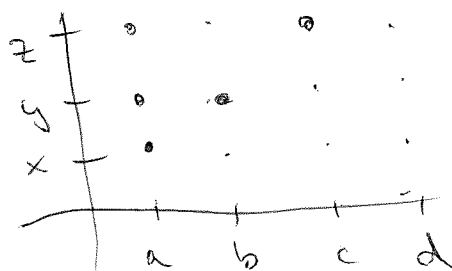
"vote on the same party"



A relation from A to B is a subset
of $A \times B$

$$A = \{a, b, c, d\} \quad |A \times B| = |A| \cdot |B| = 4 \cdot 3 = 12$$

$$B = \{x, y, z\}$$



$$R = \{(a, x), (a, y), (a, z), (b, y), (c, z)\}$$

On $A \times B$ we have

2^{12} relations

$\hookrightarrow 4096$

A relation on A is a subset of $A \times A$.

(2)

1.) A relation is reflexive if $(a,a) \in R$ for every $a \in A$.

2.) A relation R is symmetric if $(a,b) \in R \rightarrow (b,a) \in R$ for all $a,b \in A$

3.) Anti-symmetric relation: $aRb \wedge bRa \rightarrow a=b$
(or $(a,b) \in R$)

4.) A relation is transitive if whenever
 $(a,b) \in R \wedge (b,c) \in R$
 \rightarrow
 $(a,c) \in R$

$$A = \{a, b, c\}$$

$$R = \{(a,a), (b,b), (c,c), (a,c)\}$$

c	.	.	.
b	.	.	.
a	.	.	.
	a	b	c

Ex 1.) A - webpages
www....

$(a,b) \in R$ Relations	Reflexive?	Symmetric?	Anti-sym?	Transitive
\Leftrightarrow Everyone who has visited webpage a has also visited webpage b	Y	N	N	Y
There are no common links on a and b	N	Y	N	N
$x \cdot y \geq 1$	No $0 \cdot 0 < 1$	Y	N	Y
$x \equiv y \pmod{7}$				

none of them = 0

They have to have the same sign

$$x \cdot y \geq 1 \text{ is then } y \cdot z \geq 1 \text{ } x \cdot z \geq 1$$

$$(x \cdot y)(y \cdot z) \geq 1$$

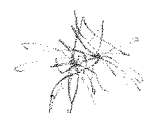
$$x \cdot z \geq 1$$

$$x \cdot z \geq 1$$

A relation on $\mathbb{Z} \rightarrow \mathbb{Z} (a,b) \in R \Leftrightarrow$

$$\frac{17-10}{7} = 1$$

$$10 R 17 \quad 17 R 51$$



18.4.2017

(3)

$$x \equiv y \pmod{7}$$

1.) Reflexive $x - x = 0 \cdot 7$ Yes

2.) Symmetric $x - y = k \cdot 7 \quad k \in \mathbb{Z}$

$$y - x = (-k) \cdot 7 \quad \text{Yes}$$

$$y - x = l \cdot 7$$

$$\frac{24-17}{7} = 1 \quad \begin{array}{l} 24-17 = 1 \cdot 7 \\ 17-24 = (-1) \cdot 7 \end{array}$$

$$4.) \quad \begin{array}{l} x \quad y \quad k \\ 77-49 = 4 \cdot 7 \quad (1) \end{array}$$

$$\begin{array}{l} l \\ 49-35 = 2 \cdot 7 \quad (2) \end{array}$$

$$(1) + (2) \quad 77-35 = (4+2) \cdot 7 = 6 \cdot 7$$

$$x - z = (k+l) \cdot 7 \quad x R z$$

Grand-parents,

Train trips with one stop

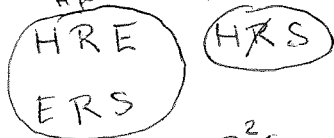
From R we can form the composite relation.

$$R \circ R = R^2$$

composite Relation

$(a,b) \in R^2$ if there is a c such that $(a,c) \in R$ and $(c,b) \in R$

Ex.) R - parent relation
H parent to E, E is parent to S



This means HR^2S
grand parent

→ grand-grand-parent

Theorem: R is transitive $\Leftrightarrow R^n \subseteq R$

$n = 1, 2, 3, 4, 5, \dots$

Ex.) VR^2S

$$R^3 = R^2 \cdot R$$

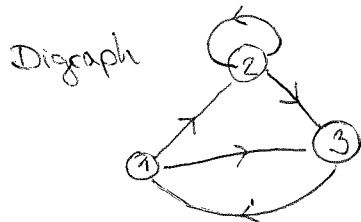
⋮

Representing relations by

1.) 0-1 matrices

2.) Directed graphs (Digraphs)

Ex.) $A = \{1, 2, 3\}$
 $R = \{(1, 2), (1, 3), (2, 2), (2, 3), (3, 1)\}$



0-1 matrix $M_R = \begin{bmatrix} 0 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \end{bmatrix}$

Is R reflexive? No

Is R symmetric? No
 You can see it in diagonals

Is R anti-sym. ? No

Is R transitive? No

$$\begin{bmatrix} 0 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \end{bmatrix} \quad \begin{bmatrix} 0 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

$$M_{R^2} = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

Ex.) $A = \{1, 2, 3, 8, 12, 21, 24\}$

$a R b \Leftrightarrow a \mid b$

$a \mid b$ means $b = k \cdot a$ for some integer k .

$$M_R = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 8 & 12 & 21 & 24 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 8 \\ 12 \\ 21 \\ 24 \end{matrix} & \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$

symmetric: yes $1, 1, 1 \rightarrow$ diagonal

anti-sym: $b = k \cdot a$
 $a = l \cdot b = k \cdot l \cdot a$
 $k \cdot l = 1$
 $k = l = 1$

Next week: graph theory until 11/5
 16/5, 18/5 old exams repetition
 17/5 presentation Deadline 12/5

Ex.) $A = \{x, y, z\}$

A relation on a set A is a subset of $A \times A$

$$M_R = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

the diagonal

$$M_{R^2} = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix}$$

$$R^2 = R \cdot R$$

on A we have $2^9 = 512$ relations

• How many of them are reflexive? $\rightarrow 2^6 = 64$

• How many of them are symmetric? If $aRb \rightarrow bRa$ $\rightarrow 2^6 = 64$

diagonal would have 1's everywhere on the diagonal

• How many are anti-symmetric? $2^3 \cdot 3^3 = 216$

• How many are reflexive and anti-symmetric? $\rightarrow 3^3$

Too difficult: How many transitive relations on A

$$\begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix}$$

← anti-symm.
Think of divides relation $a|b$ and $b|a \Rightarrow a=b$

~~Theorem~~ Theorem: The relation R on A is transitive $\Leftrightarrow R^n \subseteq R$
 $n=1, 2, 3, 4, \dots$

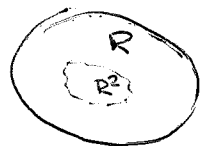
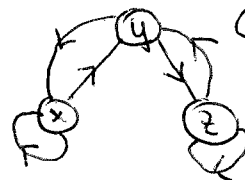


Diagram:



(= sufficient +
 \Rightarrow necessary
 Do it with induction

it is not possible to go from x to z without stopping in y

$$R^1 = R \subseteq R$$

Induction Assumption (IA):

Assume $R^k \subseteq R$

Take an element (a, c) in R^{k+1} . Is it also an element in R ? $(a, c) \in R$?

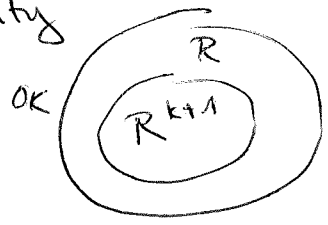
$$R^{k+1} = R^k \cdot R$$

There is, by definition, an element b such that $aR^k b$ and bRc

But $R^k \subseteq R$ due to IA
 so aRb

continued

→ we are allowed to use the transitivity
 $aRb \wedge bRc \Rightarrow (a,c) \in R$



Equivalence relations are

- 1, reflexive
- 2, symmetric
- 3, transitive

ex: Vote on some party

S	MP	PD
M	L	
V	C	FT

EX) $A = \mathbb{Z}$

$xRy \Leftrightarrow x \equiv y \pmod{4}$

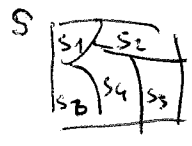
$5R9$ since $\frac{9-5}{4} = 1$

$13 \not R 18$ $\frac{18-13}{4} = \frac{5}{4}$

-8	-7	-6	-5
-4	-3	-2	-1
0	1	2	3
4	5	6	7
...
16

Partition of a set

A partition: 5 subsets



every element in S is in a subset S_i

$S_i \cap S_j = \emptyset \quad i \neq j$

Equivalence classes:

↪ is the same set, are related

$[0] = \{ \dots -8, -4, 0, 4, 8, \dots \} = [20] = [100]$

$[1] = \{ \dots -3, 1, 5, \dots \} = [17]$

$[2] = \{ \dots -2, 2, 6, \dots \} = [38]$

$[3] = \{ \dots -1, 3, 7, \dots \} = [51]$

↑
all elements related to three

Equivalence rel. \Rightarrow Partition of the set

↪ from Partition we get E. relation

Theorem 3 equivalent statements
about an equiv. rel R on A

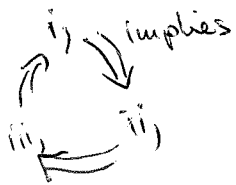
i, $a R b$

i \Rightarrow ii

ii, $[a] = [b]$

iii, $[a] \cap [b] \neq \emptyset$

if we can prove $[a] \subseteq [b]$
AND $[b] \subseteq [a]$
 $\Rightarrow [a] = [b]$



Take an element c in $[a]$.

$c R a$ and also $a R c$ \rightarrow here we use symmetry

Remember $a R b$

Don't forget transitivity

$c R a \wedge a R b \Rightarrow c R b$

That is $c \in [b]$

$[a] \subseteq [b]$

Ex.) $A = \{x, y, z\}$

How many equiv. rel. on A ?

Remember partitions!

How many partitions do we have on A ?

Symmetric and reflexive 0-1 matrices

$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$

They are 8.

1.) $\{x, y, z\}$

2.) $\{x\}, \{y\}, \{z\}$

3.) $\{x, y\}, \{z\}$

4.) $\{y, z\}, \{x\}$

5.) $\{x, z\}, \{y\}$

Partial orders (posets)

- 1.) reflexive
- 2.) anti-symmetric
- 3.) Transitivity

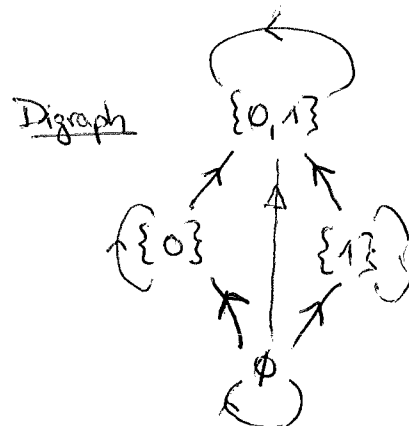
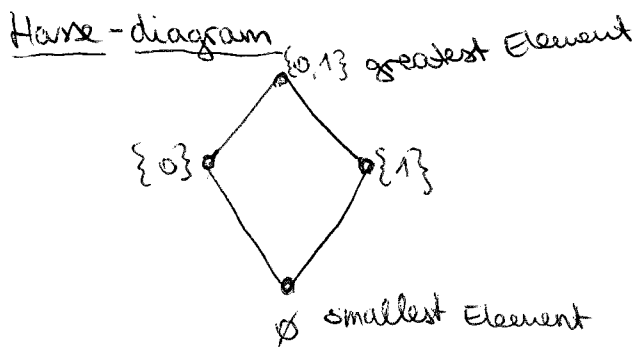
Examples \leq
 \subseteq

Ex.) ~~on the subsets of $\{0,1\}$~~

\subseteq on the subsets of $\{0,1\} = A$

$$\mathcal{P}(A) = \{ \emptyset, \{0\}, \{1\}, \{0,1\} \}$$

↑
empty set



Ex.) Draw the Hasse-diagram for the divides relation
 on $A = \{1,2,4,8\}$

$$M = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$



→ total order



25.4.2017

①

A relation on a set A is a subset of $A \times A$

Partial order:

$$a \leq b$$

like $\subseteq, \leq, 1$

reflexive anti-symmetric transitive

Equivalence relations

$$a \sim b$$

gives a partition of the set



reflexive symmetric transitive

Ex. 1) Divides relation on the positive integers
Is it a partial order?

$$a R b \Leftrightarrow a | b$$

i) reflexive?

$a | a$ since $a = 1 \cdot a$ Yes!

ii) anti-symm?

$$6 | 30 \quad 30 \nmid 6$$

$$\left. \begin{array}{l} a | b \\ b | a \end{array} \right\} \begin{array}{l} b = k \cdot a \\ a = l \cdot b \end{array} \Rightarrow a = l \cdot k \cdot a$$

$$l \cdot k = 1$$

$$k, l \in \mathbb{Z}^+ \leftarrow \text{positive integers} \quad \begin{array}{l} l = 1 \\ k = 1 \end{array}$$

iii) Transitive?

$$\begin{array}{l} x | y \\ 6 | 30 \end{array} \quad \begin{array}{l} y = k \cdot x \quad k \in \mathbb{Z}^+ \\ \text{since } 30 = 5 \cdot 6 \\ z = l \cdot y \quad l \in \mathbb{Z}^+ \end{array}$$

$$\begin{array}{l} y | z \\ 30 | 90 \end{array} \quad \text{since } 90 = 3 \cdot 30 = 3 \cdot 5 \cdot 6 = 15 \cdot 6$$

$$\Rightarrow 6 | 90$$

$$z = l \cdot k \cdot x$$

$$z = (l \cdot k) x$$

$$\begin{array}{l} x | y \\ y | z \end{array} \Rightarrow x | z$$

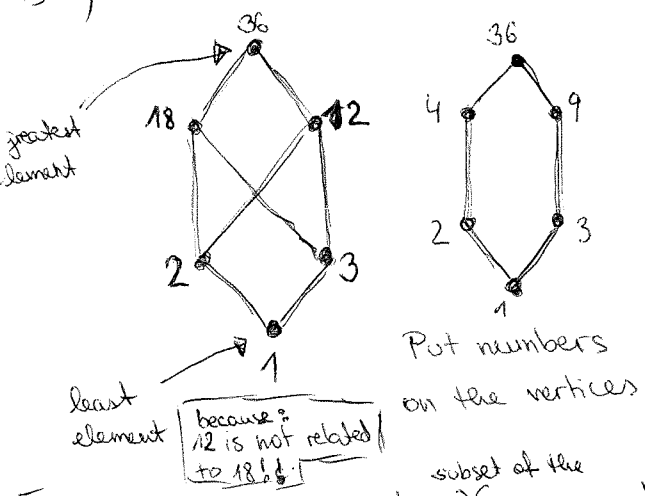
Total order on $A = \{x, y, z\}$



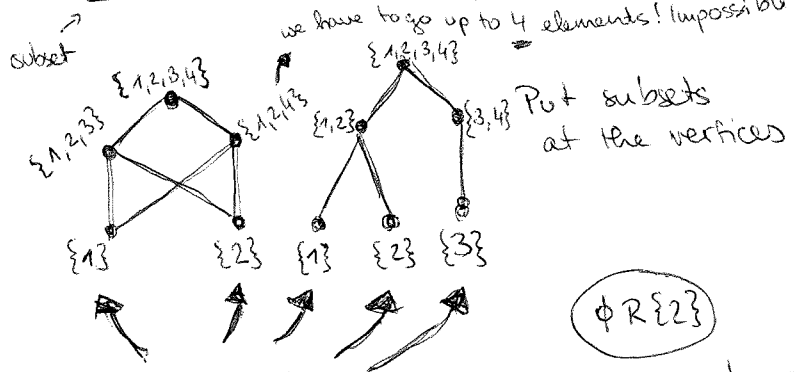
Q: How many total orders on $\{x, y, z\}$?
 Ans: $3! = 6$

$x R y$ $x \leq y$
 $y R z$ $y \leq z$

Ex) Divides relation on a subset of the positive integer



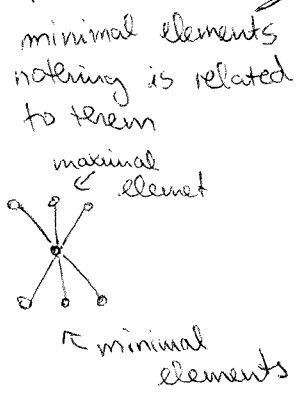
Ex) \subseteq relation on the powerset of some set $\mathcal{P}(\{1, 2, 3\}) = \{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$



→ both not LATTICE

$\emptyset R \{2\}$

we can not use the empty set in this Hasse-Diagrams because of this Relation



LCM - least common multiple $\rightarrow l \cup b \{a, b\}$ least upper bound for a and b
 GCD - greatest common divisor $\rightarrow g \cup b \{a, b\}$ greatest lower bound for a and b

LCM (12, 21) :

$$12 = 2^2 \cdot 3^1 \cdot 7^0$$

$$21 = 2^0 \cdot 3^1 \cdot 7^1$$

$$\text{LCM}(12, 21) = 2^2 \cdot 3^1 \cdot 7^1 = 84$$

maximum
of each power

$$\text{GCD}(30, 12) = 6$$

$$30 = 2 \cdot 12 + 6$$

$$12 = 2 \cdot 6$$

$$30 = 2 \cdot 3 \cdot 5$$

$$12 = 2^2 \cdot 3 \cdot 5^0$$

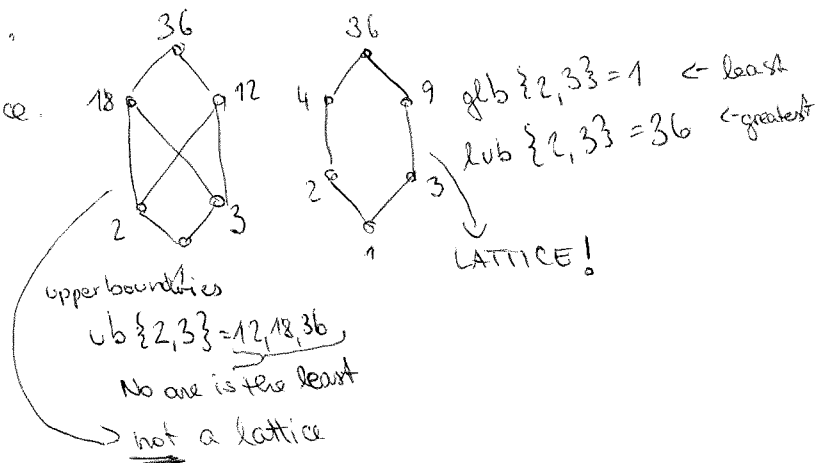
$$\text{GCD}(30, 12) = 2^1 \cdot 3^1 \cdot 5^0 = 6$$

we take ~~minimum~~ minimum
of each power

LATTICE : $l \cup b \{a, b\}$ and $g \cup b \{a, b\}$
 exist for all pair of
 elements in the poset.

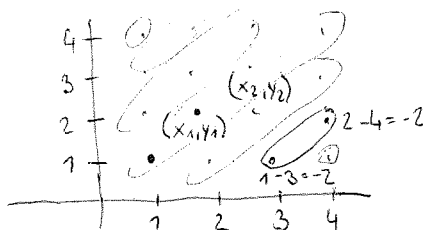
Ex: Divides relation on \mathbb{Z}^+ is a lattice.

$g \cup b$ - greatest lower bound
 $l \cup b$ - least upper bound



Ex: $S = \{1, 2, 3, 4\}$

A relation on $S \times S$ that is a subset of $(S \times S) \times (S \times S)$



$$(x_1, y_1) R (x_2, y_2) \Leftrightarrow y_1 - x_1 = y_2 - x_2$$

Show that this is an equivalence relation!
 Give the equivalence classes

i, Reflexive? Yes

$$y_1 - x_1 = y_1 - x_1$$

ii, symmetric? $(x_1, y_1) R (x_2, y_2) \Rightarrow y_1 - x_1 = y_2 - x_2$
 is then

$$(x_2, y_2) R (x_1, y_1) \Rightarrow y_2 - x_2 = y_1 - x_1$$

yes!

iii, $(x_1, y_1) R (x_2, y_2)$
 $(x_2, y_2) R (x_3, y_3) \Rightarrow (x_1, y_1) R (x_3, y_3)$?
 Yes!



Ex.)

$$S = \{ (a, b) \mid \underset{\substack{\text{such} \\ \text{that}}}{a, b \in \mathbb{Z} \wedge b \neq 0} \}$$

$$(a, b) R (c, d) \Leftrightarrow \frac{a}{b} = \frac{c}{d}$$

$$\rightarrow \text{ex: } (1, 3) R (2, 6) \text{ since } \frac{1}{3} = \frac{2}{6}$$

Check that it is an equivalence relation on S .

$$[(2, 3)] = \{ (4, 6), (6, 9), (8, 12), \dots \}$$

(answer is on Hans page)

$$\frac{2}{3} \neq \frac{3}{4}$$

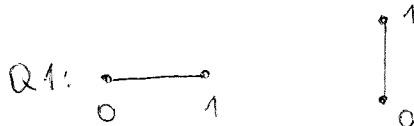
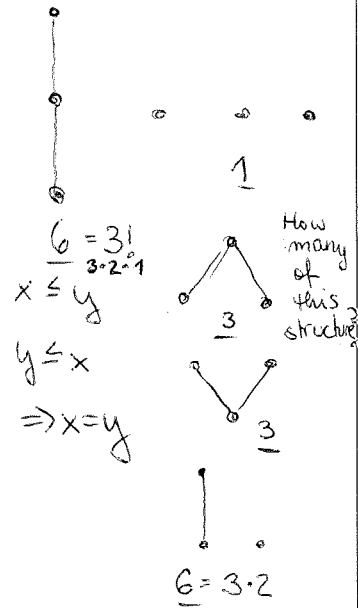
$$\frac{2}{3} = \frac{4}{6} = \frac{6}{9} = \frac{8}{12} = \frac{10}{15}$$

think of Rational number as equivalence class

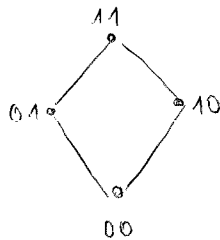
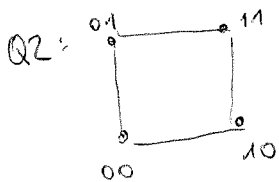
Graphs

Almost everything but Not
 10.6 shortest paths...
 10.3 matrices
 just know the concept connectivity

19 partial orders on
 $A = \{x, y, z\}$
 (Binary relations -
 wikipedia)

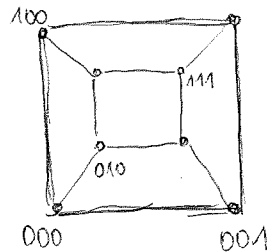
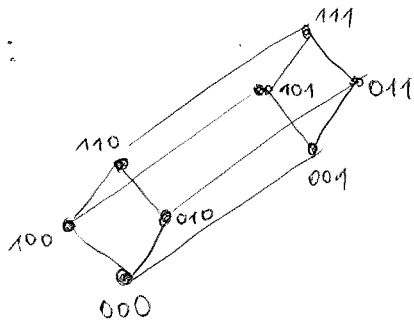


(Hasse-diagram)
 wikipedia

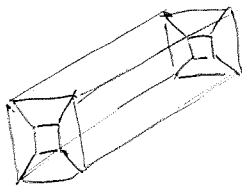


$xyRzw$
 \Leftrightarrow
 $x \leq z$ and
 $y \leq w$

Q3:



Cubes



Q4:

Graph Vertex Edge

Def: $G = (V, E)$

$E \subseteq V \times V$

Handshake-theorem:

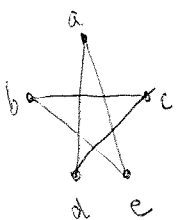
$$2e = \sum_{v \in V} \deg(v)$$

$v=5$

$v=5$

$e=5$

$|E|=5$

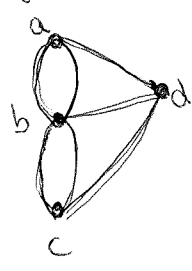


$\deg(a)=2$

$2e = 5 \cdot 2$

$e=5$

Multigraphs more than one edge between 2 vertices

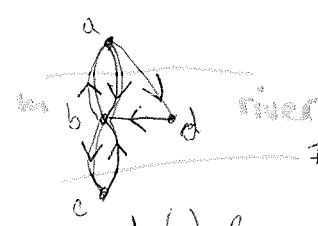


$\deg(a) = 3$
 $\deg(b) = 5$
 $\deg(c) = 3$
 $\deg(d) = 3$

$2e = 14$
 $e = 7$

the people in Königsberg wanted to find an Euler circuit (go over all edges once) closed

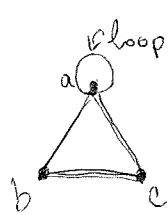
path not repeating edges



Find an Euler path

$\deg(a) = 3$
 $\deg(b) = 5$
 $\deg(c) = 2$
 $\deg(d) = 2$

$a \rightarrow d \rightarrow b \rightarrow c \rightarrow b$
 $\rightarrow a \rightarrow b$



$\deg(a) = 4$
 $\deg(b) = 2$
 $\deg(c) = 2$

$2e = 8$
 $e = 4$

Directed graphs

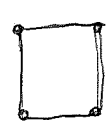


www (1999)

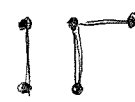
$|V| = 3 \cdot 10^9$

70% was connected?

A graph is connected if there is a path between all pairs of vertices



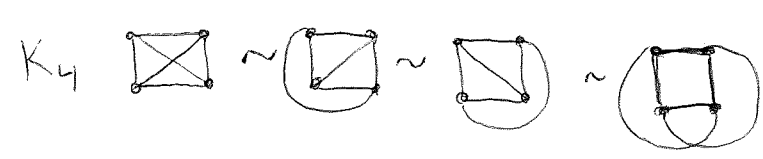
connected



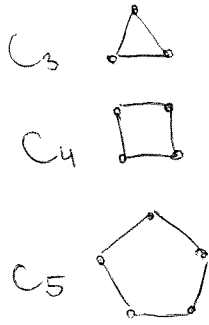
not connected

we consider mainly simple graphs (not multigraphs) without loops and undirected.

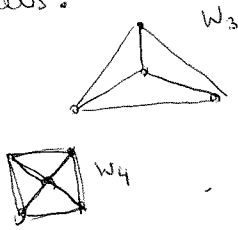
complete graphs



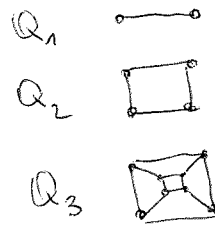
C-cycles no vertices repeated



Wheels:



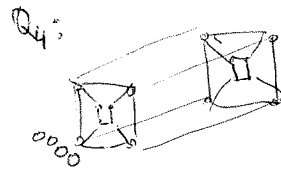
Hypercubes



e
1
4
12

can we see a pattern?
How many edges in Hypercube

How many edges in Q_n ?



$12 + 12 + 8 = 32$

$|V| = 2^4 = 16$

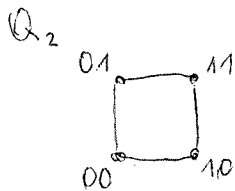
$\deg(v) = 4$ ← Handshaking Theorem

$2e = 16 \cdot 4$

$e = \frac{16 \cdot 4}{2} = 32$

Ex.) # edges in Q_n

$2e = \sum_{v \in V} \deg(v)$



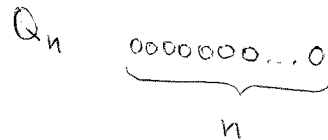
bit strings

$|V| = 2^2 = 4$

$\deg(v) = 2 \quad \forall v \in V$

$2e = 4 \cdot 2$

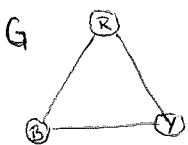
$e = 4$



$|V| = 2^n$

$\deg(v) = n$

Proper coloring of a graph G

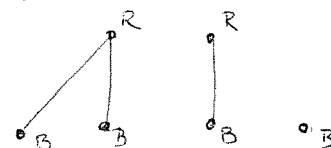


chi

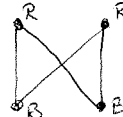
$\chi(G) = 3$

the chromatic number

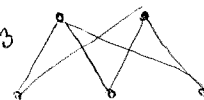
Bipartite graphs



$K_{2,2}$



$K_{2,3}$



~~Handwritten scribbles~~

$2e = 3 + 3 + 2 + 2 + 2$

$e =$

edges in $K_{m,n}$

$m = 1, 2, 3, \dots$

$n = 1, 2, 3, \dots$

$1 \ 2 \ 3 \ \dots \ m$

$1 \ 2 \ 3 \ \dots \ n$

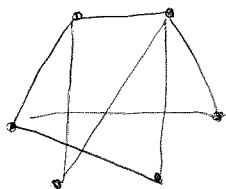
$2e = m \cdot n + m \cdot n \rightarrow e = m \cdot n$

Theorem: A graph G is bipartite

(\Rightarrow)

$\chi(G) = 2$

Ex.) G

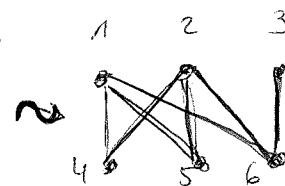
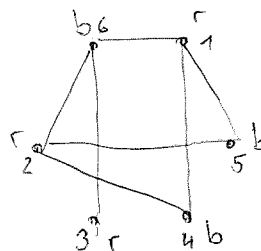


Is G bipartite?

Answer:
No!

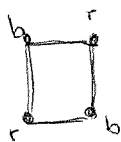
would be if 6

(4)



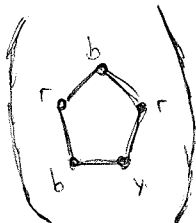
bipartite: because we can do the proper coloring with 2 colors

Ex.) for which n are C_n bipartite



C_4 ?

Answer: even n



Ex.) Draw $\overline{Q_2}$ complementary graph to Q_2

The graph has the same number of vertices as Q_2

Two vertices are adjacent in $\overline{Q_2}$

\Leftrightarrow They are not adjacent in Q_2

Answer:



$\overline{Q_2}$?



~



Q_2 :  $\overline{Q_2}$: ?

Ex.) Try draw $K_{3,3}$ without crossings

Answer:

\rightarrow impossible here



* path (walk) vandrings

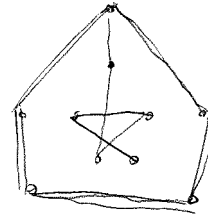
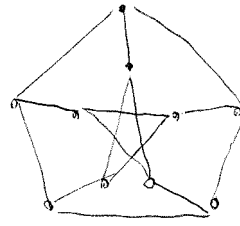
* simple path-(trail)
no repetitions of edges (våg)

* circuit (closed walk)

path that starts and ends (krets)
at the same point

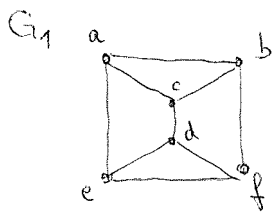
* cycle - simple closed path
without repetitions
of vertices

G



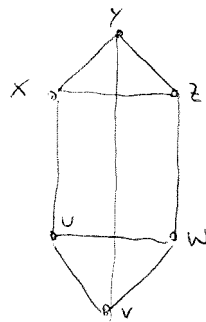
a subgraph
of G

Isomorphism
(same structure)



$$V_1 = 6$$

$$e_1 = 9$$

G₂

$$V_2 = 6$$

$$e_2 = 9$$

Def: 2 graphs are isomorphic if
there is a bijection* f

$f: V_1 \rightarrow V_2$ such that

* onto and one-to-one



$\{a, b\}$ an edge in $G_1 \Leftrightarrow \{f(a), f(b)\}$ is
an edge in G_2

$$f(a) = x$$

$$f(b) = z$$

$$f(c) = y$$

$$f(d) = v$$

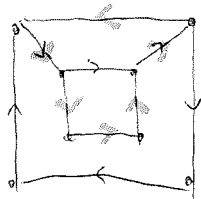
$$f(e) = u$$

$$f(f) = w$$

Q: Draw all non-isomorphic graphs with 3 vertices.

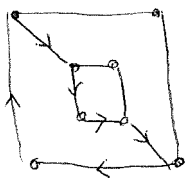
opposite \rightarrow would be isomorphic

operation: G_1



$v=8$

$e=10$



$v=8$

$e=10$

vertices with degree 2 = 4

= 4

in \square

cycles of $G = 2$

= 4

this two graphs are not the same, they differ in cycles.

Steps: check for vertices

check for cycles \rightarrow we found that they differ here \rightarrow we don't have to continue

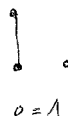
(isomorphic: same number of vertices + same # edges)

1,



$e=0$

2,



$e=1$

3,



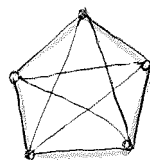
$e=2$

4,



$e=3$

Q: C_5



Is it isomorphic to its complement

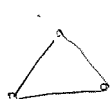
$\overline{C_5}$



$C_5 \sim \overline{C_5}$ isomorphic?

Yes!

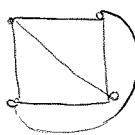
K_3



$\binom{3}{2} = 3$ edges

K_n has $\binom{n}{2} = \frac{n(n-1)}{2}$ edges

K_4



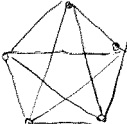
$\binom{4}{2} = 6$

If a graph G is isomorphic

to its complement ~~then~~ we must have

$n \equiv 0, 1 \pmod{4}$

K_5



$\binom{5}{2} = 10$

To be isomorphic G and \overline{G} must have

$\frac{n \cdot (n-1)}{2 \cdot 2}$ edges each

$\frac{n(n-1)}{4}$

(For example $n=5$
 $\frac{5 \cdot 4}{4} = 5$)

Handshake Theorem: $2e = \sum_{v \in V} \deg(v)$

$n=5$

$2e = 5 \cdot 4$

$e = 10$

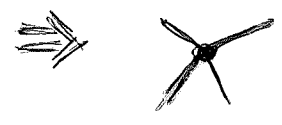
Euler circuit: go over all edges once (allowed to repeat vertices)

3

(EC)

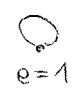
EC exists $\Leftrightarrow \deg(v_i)$ is even for all vertices $v_i \in V$

deg(v)=2

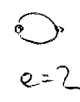


2+2+... will always be an even number

Proof by induction. Over the number of edges



e=1

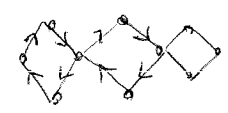


e=2



Assume it is checked up to $e=k$

Consider a graph with $e=k+1$



Since all degrees are even there must be cycles in the Graph! This needs an explanation

A tree is a graph without cycles.

$V = e + 1$ for trees.

On the other hand

The Handshake-theorem tells us

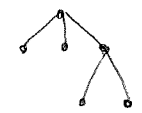
$$2e = \sum_{v_i \in V} \deg(v_i) \geq 2 \cdot V$$

so $V \leq e$

so they are not trees!

\hookrightarrow There are cycles in the graph!

Example: tree



$V=6$
 $e=5$

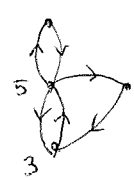
$V = 5 + 1$

For trees

$V = e + 1$

if and only if

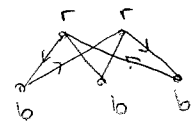
Euler path exists \Leftrightarrow Exactly two vertices are of odd degree



Hamilton cycle/path

Q: For which m, n has $K_{m,n}$ a Hamilton cycle

$K_{2,3}$

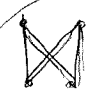


$\chi(K_{2,3}) = 2$

Answer: $m=n \geq 2$

≥ 2 because in \downarrow we are not allowed to repeat edges

\hookrightarrow no HC here!



Yes



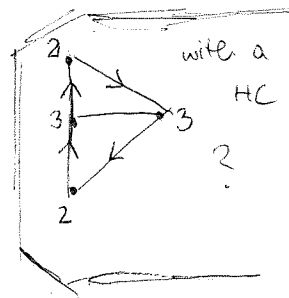
Yes

See pattern so Answer

Q: Construct a graph with 4 vertices without a HC
 where $\deg(x) + \deg(y) \geq 3$ for all non-adjacent vertices

Ore's theorem: $\deg(x) + \deg(y) \geq V$ then there is a HC

④



* Q: Show that $e \leq \frac{V^2}{4}$

for bipartite graphs

Ex: $K_{2,2}$



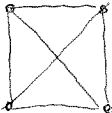
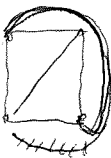
$$e = 4 \leq \frac{4^2}{4} \text{ OK}$$

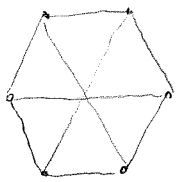
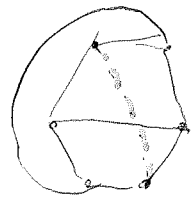
$K_{2,3}$

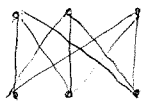


$$e = 6 \leq \frac{5^2}{4} = 6 \frac{1}{4}$$

Planar graphs

Is  planar? Yes! \sim isomorphic 

Is  planar? \neq  No! non-planar

 $\sim K_{3,3}$

For a given V we can not have too many edges if G should be planar.

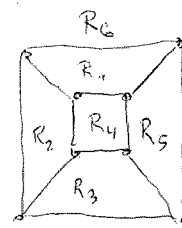
Kuratowski's theorem:

In all non-planar graphs $K_{3,3}$ or K_5 are "hidden"



Euler's formula for planar graphs

Q_3



Regions

$$r = 6$$

$$e = 12$$

$$v = 8$$

$$12 - 8 + 2 = 6$$

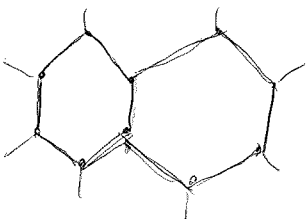
$$r = e - v + 2$$

always for planar Graphs

Assume the formula holds for K regions (IA)

Consider a graph with $r = k+1$ regions

Table 1989



$$r \rightarrow r-1$$

$$e \rightarrow e-1$$

Using IA:

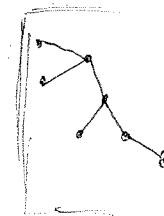
$$(r-1) = (e-1) - v + 2$$

$$r = e - v + 2$$

Proof: Induction proof over the number of regions.

Ex: $r=1$

$$1 = e - (e+1) + 2$$

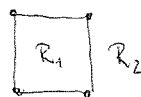


Tree!

$$v = e + 1$$

Ex: $r=2$ $v = e?$

$$v = e$$



$$2 = e - e + 2$$

OK!

Handshake - formula: $2e = \sum_{v_i \in V} \deg(v_i)$

For planar graphs: $2e = \sum_{i=1}^r \deg(R_i)$

Ex.) $v=6$

$\deg(v)=3$

for all vertices.

It is planar, How many regions?

Draw it.

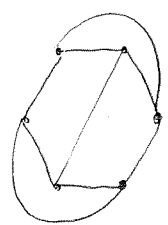
1.) Handshake theorem

2.) E. formula $r=e-v+2$

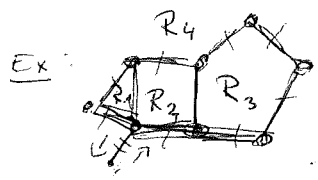
$2e = 6 \cdot 3$

$e = 9$

$r = 9 - 6 + 2 = 5$



$r=5$
 $v=6$
 $e=9$



Ex:

$\deg(R_1)=3 \leftarrow \# \text{ of edges}$

$\deg(R_2)=4$

$\deg(R_3)=5$

$\deg(R_4)=10$

$10+5+4+3=22$

$2e = 22$

$e = 11$

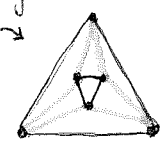
instead $\deg(v)=4$

$2e = 6 \cdot 4$

$e = 12$

$r = 12 - 6 + 2$
 $= 8$

2D drawing of the octahedron



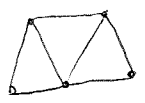
$v=6$
 $e=12$



Not too many edges...

$2e = \sum \deg(R_i) \geq 3+3+3+\dots = 3 \cdot r$

$r = e - v + 2$



$\deg(R_i) \geq 3$

$2e \geq 3r$

$r \leq \frac{2e}{3}$

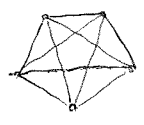
$e - v + 2 \leq \frac{2e}{3}$

$\frac{e}{3} \leq v - 2$

$e \leq 3v - 6$

Ex: K_5 $v=5$

$e = 10 = \binom{5}{2}$



$= 4+3+2+1$

Non-planar! Is $10 \leq 3 \cdot 5 - 6 = 9$?
No!

$$K_{3,3} \quad v=6$$

$$e = 3 \cdot 3 = 9$$

$$\text{Is } 9 \leq 3 \cdot 6 - 6 = 12 ?$$

Yes! It is not a sufficient condition

↳ bipartite graph

Conclusion: $e \leq 3v - 6$ is necessary but not sufficient condition for planarity.

In bipartite graphs the minimum cycle length is 4 ~

$$2e \geq 4 + 4 + 4 + \dots + 4 = 4 \cdot r$$

$$v + 2 = r \leq \frac{e}{2}$$

$$\frac{e}{2} \leq v - 2$$

$$e \leq 2v - 4$$

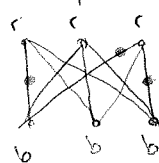
$$K_{3,3} \quad e=9$$

$$v=6$$

$$\text{Is } 9 \leq 2 \cdot 6 - 4 = 8 ?$$

No! $K_{3,3}$ is non-planar

K_5 or $K_{3,3}$ is hiding in every non planar graph

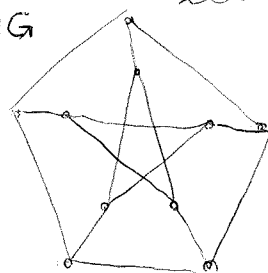


3 new vertices are all deg 2

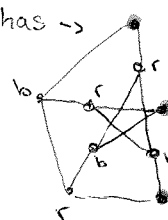
→ is homeomorphic to $K_{3,3}$

Theorem by Kuratowski

Find subgraph with structure of $K_{3,3}$ or K_5



has →



→ as a subgraph

check that this subgraph is homeomorphic to $K_{3,3}$

5 platonic solids

(degree of every region)

$$\deg(R_i) = m \quad i = 1, 2, \dots, r$$

$$\deg(v_i) = n \quad i = 1, \dots, |v|$$



$$r = e - v + 2$$

$$2e = \sum \deg(R_i) = r \cdot m \quad (1)$$

$$2e = \sum \deg(v_i) = v \cdot n \quad (2)$$

$$2 = r - e + v \geq 0 \quad (3)$$

(1) and (2) into (3)

$$\frac{2e}{m} - e + \frac{2e}{n} \geq 0$$

$$e \left[\frac{2n - mn + 2m}{mn} \right] \geq 0$$

$$2n - mn + 2m \geq 0$$

m	n	name
3	3	tetrahedron (4 faces)
3	4	octahedron
4	3	cube
3	5	
5	3	

symmetry

$$\text{Ex: } m=3, n=4$$

$$r = e - v + 2$$

$$\frac{2e}{3} = e - \frac{2e}{4} + 2$$

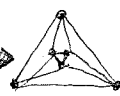
$$\frac{2e}{3} + \frac{e}{2} - e = 2$$

$$\frac{4e + 3e - 6e}{6} = 2$$

$$e=12$$

$$r=8$$

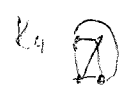
$$v=6$$



Vertex colorings of graphs

$$\chi(\triangle) = 3$$

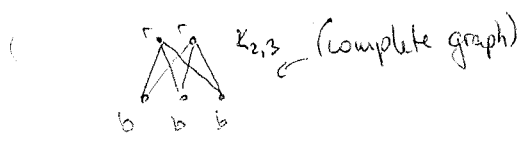
$$\chi(K_n) = n$$



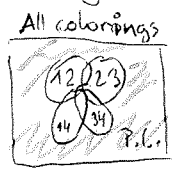
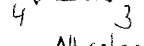
Proper colorings (P.C.)

two vertices with edge in common must have different colors. The minimum number of colors needed for a p.c. of G is called the chromatic number $\chi(G)$.

The G bipartite $\Rightarrow \chi(G) = 2$



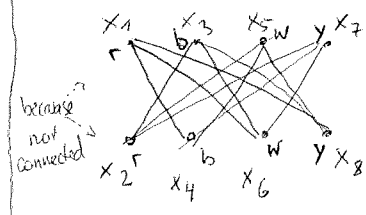
Ex.) How many proper colorings of C_4 with 3 colors?



\rightarrow try with I.E.

$$3 \cdot 2 \cdot 2 + 3 \cdot 2 \cdot 1^2 = 12 + 6 = 18$$

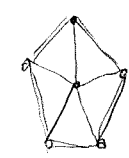
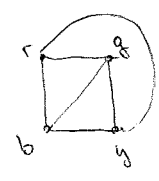
Greedy algorithm



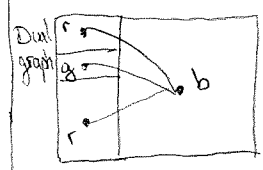
r
b
w
y

4 colors needed

Q: Find two different planar graphs with chromatic number = 4



4-color theorem: 1976 with computer



Note: The dual graph is planar.

Theorem: $\chi \leq 6$

Q: Can all vertices in a planar graph have degree 6 or more?

$$2e = \sum \deg(v_i) \geq 6 + 6 + 6 + \dots = 6v$$

$$e \geq 3v$$

\rightarrow Answer: No!

There are vertices with smaller degree, like 5

$\chi(\text{The dual graph}) = 4$
at most, (No 5th needed)

red	blue	one point is not enough!
blue	red	

But from last week: $2e = \sum \deg(v_i) \geq 3 + 3 + 3 + \dots$
 $= 3r$

$$\text{so } r \leq \frac{2e}{3}$$

$$\text{Euler's formula: } r = e - v + 2 \leq \frac{2e}{3}$$

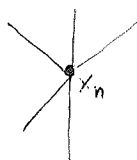
$$\frac{e}{3} \leq v - 2$$

$$e \leq 3v - 6$$

$$|V| = n$$

$G \quad x_n$

(2)



$$G - \{x_n\}^* \quad x_{n-1}$$

$$G - \{x_n, x_{n-1}\} \quad x_{n-2}$$

G without x_n and

*The subgraph of all edges attached to it

In this way

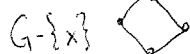
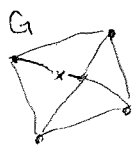
we get a list of the

vertices $x_1, x_2, x_3, x_4, x_5, x_6, x_7, \dots, x_n$

x_7 is connected with at most 5 of the vertices $x_1, x_2, x_3, x_4, x_5, x_6$

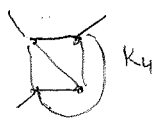
Let's do the greedy algorithm

I can reuse a color for vertex x_7



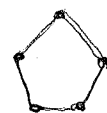
Q: Show that a graph G has at least $\binom{\chi(G)}{2}$ edges.

$\chi(G)$	$e \geq$
2	$\binom{2}{2} = 1$
3	$\binom{3}{2} = 3$
4	$\binom{4}{2} = 6$
5	$\binom{5}{2} = 10$
6	$\binom{6}{2} = 15$



K_n has $\binom{n}{2}$ edges

$$\chi(K_n) = n$$



$$\chi(G) = 3$$

$$5 > \binom{3}{2}$$

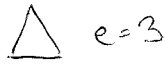
Today trees

11.5.2017

①

5 Tuesday: K-building Problems
1-3

$\chi = 3$



$e = 3$

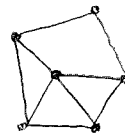


$e = 5$

$\chi = 4$



$e = 6$



$e = 10$

1/5 Thursday: The rest

2/3/5 Last minute questions

Extremal problems

Given a v

$e \leq \lfloor \frac{v^2}{4} \rfloor$ if we want to avoid \triangle



and the graph is $K_{\lfloor \frac{v}{2} \rfloor, \lceil \frac{v}{2} \rceil}$

$v = 7$ floor function

$$\lfloor \frac{7^2}{4} \rfloor = \lfloor 12 \frac{1}{4} \rfloor = 12$$

$$\lfloor \frac{v}{2} \rfloor = 3$$

ceiling function

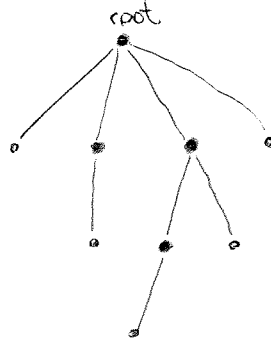
$$\lceil \frac{v}{2} \rceil = 4$$



$e = 12$

Trees - no cycles

$$v = e + 1$$



$l = 5$

$i = 4$

$v = 9$

$e = 8$

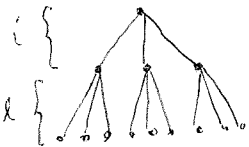
l = leaves

i = internal vertices

$$v = l + i$$

root = internal

A full 3-ary



root children grandchildren

$$1 + 3 + 3^2 = e$$

$$1 + 3 \cdot (1 + 3) = 1 + 3i$$

$$v = l + i = 1 + 3i$$

only for full

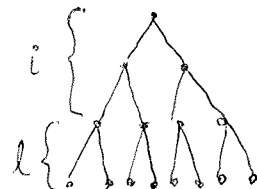
$$\frac{l-1}{3-1} = i$$

General formula for m-ary

$$\frac{l-1}{m-1} = i$$

exp: $m = 2 \quad i = l - 1$

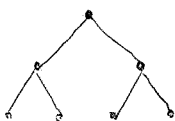
full binary tree



$l = 8$

$i = 7$

Tournaments

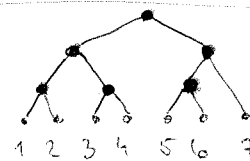


How many games will be played during the whole tournament?

$$\# \text{ games} = i = l - 1$$

$$\# \text{ games} = 64 - 1 = 63$$

1 2 3 ... 64



← not a full tree

How many games will be played?

$$\# \text{ games} = 7 - 1 = 6$$

Everyone is a loser exactly once except the winner !!!



There is a bijection between the set of losers and the set of gamers.

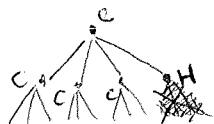
Q: How many edges in a full binary tree with ~~1023~~ 1023 internal vertices?

$$l = i + 1 = 1024 \quad v = 2047$$

$$e = 2046$$

Q: C_4H_{10} is it a tree?

Q: Which complete bipartite graphs are trees?



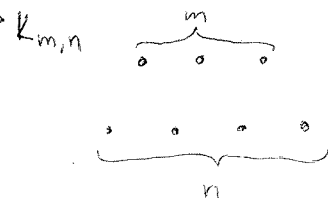
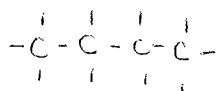
$$v = 11$$

$$2e = 26$$

$$e = 13$$

$$v = e + 1$$

Yes!



$$v = e + 1$$

$$v = m + n$$

$$e = m \cdot n$$

$$m + n = m \cdot n + 1$$

$$m = 1$$

$$\text{and/or } n = 1$$

Ex: 25 bagels

$$x_3, x_4, x_1, x_2$$

$$3 \leq x_i \quad i = 2, 3, 4, \dots$$

$$3 \leq x_i = 6$$

$$GF: G(x) = (x^3 + x^4 + x^5 + \dots)^3 \cdot (x^3 + x^4 + x^5 + x^6)$$

$$= \dots + 2x^{25} + x^9 \cdot \frac{1}{(1-x)^3} \cdot \frac{x^3 \cdot (1-x)^4}{1-x}$$

$$= \frac{x^{12} - x^{16}}{(1-x)^4}$$

$$\begin{array}{l|l} y_1 \geq 0 & z_4 + 4 = y_4 \\ y_2 \geq 0 & y_1 + y_2 + y_3 + y_4 + z_4 = 13 \\ y_3 \geq 0 & y_1 + y_2 + y_3 + z_4 = 9 \\ z_4 \geq 0 & \end{array}$$

Chemistry

2

isomers - non isomorphic graphs

Ex:



$$\deg(C) = 4$$

$$\deg(H) = 1$$

$$v = 3 + 8 = 11$$

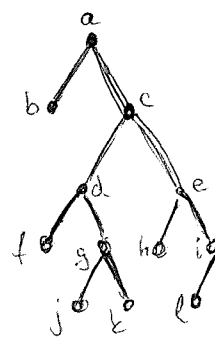
$$\text{Is } v = e + 1?$$

$$\text{if so } e = 10$$

let's check with handshake - theorem

$$2e = \sum_{v \in V} \deg(v) = 3 \cdot 4 + 8 \cdot 1 = 20$$

$$e = 10$$



center = the vertex with the smallest

eccentricity of

a vertex is the length of the longest path that

can be found starting at the vertex.

vertex	ecc.
a	4
b	5
c	3
d	4
e	4
f	5
g	5
h	5
i	5
j	6
k	6
l	6

center (the smallest number)

Q: with two centers?
think of it

$$G(x) = \left(\frac{x^{12} - x^{16}}{1-x} \right) \left(1 + 4x^9 + 6x^{13} \right) = \dots + 2x^{25}$$

$$\text{Ans. } b - a =$$

$$\begin{array}{l} y_1, y_2, y_3, y_4 = 13 \\ \text{in total (restrictions)} \end{array} \cdot \binom{16}{13} = \binom{16}{13}$$

Inclusion-Exclusion

$$x_1 = 3 + y_1$$

$$x_2 = 3 + y_2$$

$$x_3 = 3 + y_3$$

$$x_4 = 3 + y_4$$

$$y_i \geq 0 \quad 0 \leq y_4 \leq 3$$

Today

Combinatorics
&
Relations

Combinatorics

* 4 cases

* sum rule

* multiplication rule

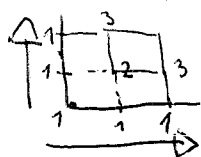
Thursday: Graphs & Trees

3 exams: 150829

160526

160609

1.) 9/6 2016

a) sum of the n th row in Pascal's Δ 

$$\begin{array}{ccccccc} & & & & 1 & & \\ & & & 1 & & 1 & \\ & & 1 & & 2 & & 1 \\ & 1 & & 3 & & 3 & & 1 \\ 2^4 = 16 & 1 & 4 & 6 & 4 & 1 & \end{array}$$

$$\text{Ans} = 2^n$$

$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^k y^{n-k}$$

b) How many six-letter words can be formed with 2U, 2R and 2D

R, G, C, A

D U R D U R

$$(x+y)^4 = \sum_{k=0}^4 \binom{4}{k} x^k y^{4-k}$$

$$(1+1)^4 = \sum_{k=0}^4 \binom{4}{k} = 2^4$$

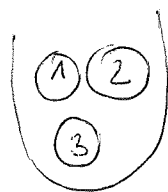
$$1.) \frac{6!}{2!2!2!}$$

2.) First decide the positions for the U's

D U R U D R

$$\begin{aligned} \binom{6}{2} \text{ way to do this} & \cdot \binom{4}{2} \cdot \binom{2}{2} = \frac{6!}{2!4!} \cdot \frac{4!}{2!2!} \cdot \frac{2!}{2!0!} \\ & = \frac{6!}{2!2!2!} \quad \downarrow \quad 0! = 1 \end{aligned}$$

c.)



Pick them 13 times
with repetition
Don't care about
order

$$x_1 + x_2 + x_3 = 13$$

How many integer
solutions if $x_i \geq 0$?

↓
non-negative

$$\binom{15}{13} = \binom{15}{2} =$$

$$6 + 7 + 0 \rightarrow \underbrace{||||| + ||||| + }_{15 \text{ positions} \rightarrow (1+...)} +$$

15 positions $\rightarrow (1+...)$
 $\rightarrow 13$ sticks
 $\rightarrow 2 +$

d.) 1, 1, 1, 1, 1, 1, ... (infinite string of ones)

$$G(x) = 1 + 1 \cdot x + 1 \cdot x^2 + 1 \cdot x^3 + 1 \cdot x^4 + \dots$$

$$= \frac{1}{1-x} \quad |x| < 1$$

$$x \cdot S = (1 + x + x^2 + x^3) \cdot x$$

$$S \geq 1 + \dots + x^3$$

$$S_3 - x \cdot S_3 = 1 - x^4$$

$$S_3 = \frac{1-x^4}{1-x}$$

$$\frac{S_3(1-x)}{1-x} = \frac{1-x^4}{(1-x)}$$

$$S_n = \frac{1-x^{n+1}}{1-x}$$

$$S_\infty = \frac{1}{1-x} \quad n \rightarrow \infty$$

not exam:

$$x_1 + x_2 + x_3 = 13$$

integer
solutions?

$$x_i \geq -4$$

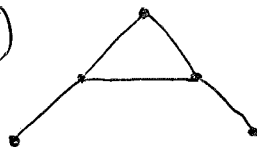
$$i = 1, 2, 3$$

$$x_i = -4 + y_i \quad y_i \geq 0$$

$$y_1 + y_2 + y_3 = 25 \quad y_i \geq 0$$

$$\text{Ans } \binom{27}{25}$$

e.)



• 5 colors
available

ways to
do a proper
coloring with
5 colors

1.) Go from left
to right

$$\rightarrow 5 \cdot 4 \cdot 4 \cdot 3 \cdot 4 = 5 \cdot 4^3 \cdot 3$$

2.) Start with triangle

$$5 \cdot 4 \cdot 3 \cdot 4 \cdot 4$$

We can make — colorings of the Δ

$$\binom{5}{3} \cdot 3! = \frac{5!}{3! \cdot 2!} \cdot 3!$$

$$= 5 \cdot 4 \cdot 3$$

$$\text{Totally } 5 \cdot 4 \cdot 3 \cdot 4 \cdot 4 = 5 \cdot 4^3 \cdot 3$$

Relations

$$A = \{1, 2, 3, 4\}$$

A relation on A is a subset of $A \times A$

4	•	•	•	•
3	•	•	•	•
2	•	•	•	•
1	•	•	•	•
	1	2	3	4

- A. reflexive
- B. symmetric
- C. anti-symmetric
- D. transitive

A, B, D: Equivalence relations $[x] = \{x, \dots\}$

A, C, D: Partial order

Partition



Lattice

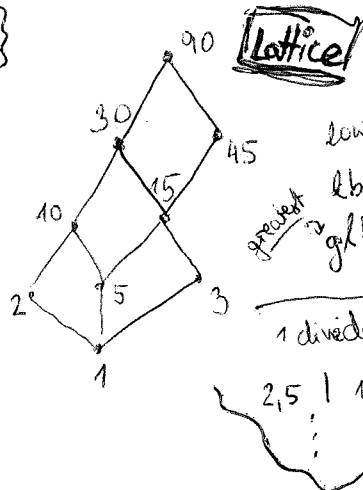
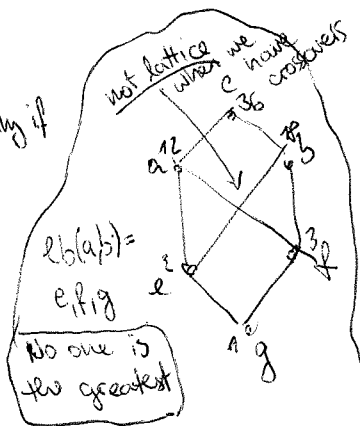
Ex: $GCD(a, b)$ $glb(a, b)$
 $LCM(a, b)$ $lub(a, b)$
 least common multiple

3. $26/5 = 16$

$$\{1, 2, 3, 5, 10, 15, 30, 45, 90\}$$

aRb

\Rightarrow if and only if $a|b$



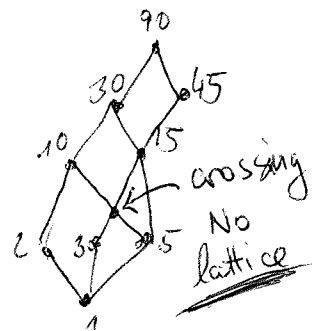
lower bounds

$$lb(10, 15) = 1, 5$$

$$glb(10, 15) = 5$$

1 divides 2, 5, 3

$$2, 5 \mid 10$$



$$M_2 = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 1 & 1 & 0 & 0 & 1 \\ 2 & 1 & 1 & 0 & 0 & 1 \\ 3 & 0 & 0 & 1 & 0 & 0 \\ 4 & 0 & 0 & 1 & 1 & 0 \\ 5 & 1 & 1 & 0 & 0 & 1 \end{pmatrix}$$

equivalence Relation

$$[1] = \{1, 2, 5\}$$

$$[4] = \{3, 4\}$$

$$\begin{pmatrix} 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ \hline 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \end{pmatrix}$$

4. The *odd graph* O_k , k is an integer ≥ 2 , is defined in the following way:
The vertices represent the subsets with $k - 1$ elements that can be obtained from a set with $2k - 1$ elements. Two vertices are joined with an edge if and only if the corresponding subsets are disjoint.
 - a) Draw O_2 and O_3 (3p)
 - b) Show that O_k is k -regular for all $k \geq 2$, that is, show that all vertices in O_k has degree k . (3p)
5.
 - a) Draw a planar graph with six vertices which is 3-regular (i.e. all vertices have degree 3). (2p)
 - b) Let G be a graph with 11 vertices and let \bar{G} be its complement. Show that not both G and \bar{G} can be planar. (4p)

6. The *girth* of a graph G is the length, i.e. the number of edges, of the shortest cycle it is possible to find in G and it is denoted $g(G)$. For graphs with n vertices and $n+1$ edges, $n \geq 4$, it is possible to show that

$$g(G) \leq \left\lfloor \frac{2(n+1)}{3} \right\rfloor$$

- a) Draw a graph with maximal girth for $n = 7$. (1p)
 b) To show the inequality above start by first showing that a graph with n vertices and $n+1$ edges always has at least two cycles. Consider then separately the two cases where the two cycles are edge disjoint and when they have edges in common. What is the maximal girth if the cycles are edge disjoint? Note that when the two cycles have edges in common there must be three different paths between two vertices in the graph. (5p)
7. Suppose that a connected planar simple graph with e edges and v vertices contains no simple cycles of length 4 or less. Show that

$$e \leq \frac{5v-10}{3} \quad (1)$$

8. Draw all non-isomorphic trees with five vertices. Note trees are simple connected graphs. (4p)
9. Let G be a simple graph with $|V| = n \geq 3$. If $\deg(x) + \deg(y) \geq n$ for all nonadjacent $x, y \in V$, then G contains a Hamilton cycle. This is a sufficient condition for the existence of a Hamilton cycle proved by Ore 1960.
- a) Draw a graph for which the condition holds and find the Hamilton cycle. Show also that it is not a necessary condition, that is, find a graph with a Hamilton cycle but for which the condition is not fulfilled. (3p)
 b) A group of 12 people meet for dinner at a big circular table. In this group everyone knows at least 6 others. Can they be seated around the table in such a way that each person knows the person to the left and to the right? (2p)
10. a) Consider connected simple graphs, G , with 11 vertices. Prove that either G or its complement \overline{G} must be nonplanar. (3p)
 b) This theorem does not hold for eight vertices. Find a counterexample to part a above, that is, find a planar G for which also the complement is planar. (2p)

Graphs

- * Handshake theorem
- * Euler circuit & Hamilton Cycles
- * isomorphisms
- * Trees, $v = e + 1$, no cycles, $v = l + i$
- * Planar graphs

$$2e = \sum_i \deg(R_i)$$



Euler's formula $\rightarrow r = e - v + 2$

$r = 1$ for trees

$r = 6$ for Q_3



$$3 = 7$$

$$5b = 10a$$

$$4b *$$

$$5b *$$

$$6b ***$$

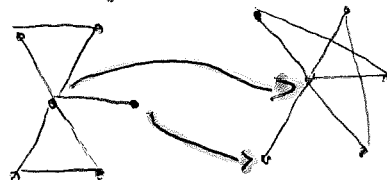
2a.) $5, 2, 2, 2, 2, 1$

$$v = 6$$

$$2e = \sum_{v_i \in V} \deg(v_i) = 14$$

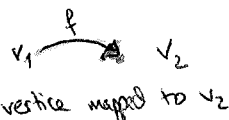
$$e = 7$$

planar G \hookrightarrow no crossings or

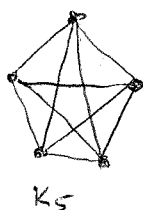
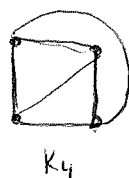
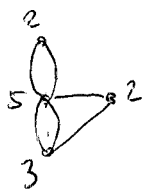


(Regions is something different)

both isomorphic \rightarrow same structure to each other

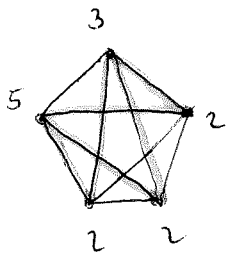


b.) For which n have K_n an Euler circuit?
(go over all edges, didn't want to repeat edges)



In K_n all vertices have degree $n-1$
(if it's 4 it's 3
if it's 3 it's 2)
So $n-1$ must be even
 n must be odd, $n \geq 3$

2)



→ handshake swedes with chinese

→ black is the handshake between swedes themselves and chines among them

$K_{2,3}$ is a subgraph of K_5

$\overline{K_{2,3}}$

3) Necess. cond. for planar graphs

$$e \leq 3v - 6$$

$$e \leq 2v - 4$$

$$e \leq \frac{5v - 10}{3}$$

Shortest cycle in the graph

3

→ we do this one!

4

5

→ Pentagon

$$r = e - v + 2$$

$$2e = \sum \deg(R_i) \geq 3 + 3 + 3 + \dots + 3 = 3r$$

Regions

$$e - v + 2 = r \leq \frac{2e}{3}$$

$$\text{or: } \begin{cases} 2e \geq 3r \\ \frac{2e}{3} \geq r \end{cases}$$

$$\frac{e}{3} \leq v - 2$$

$$e \leq 3v - 6$$

Do 4 or 5 on your own! Same way!

$$v = 7$$

$$e \leq \frac{7 \cdot 5 - 10}{3} = 8 \frac{1}{3}$$

edges = 8

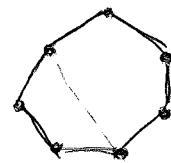
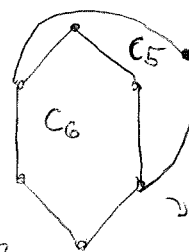
vertices = 7

avoid squares or triangles

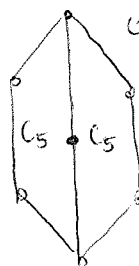
↓
No C_3
No C_4

$$v = 7$$

$$e = 8$$



→ something wrong here



Cor both work!

is it now possible to avoid C_3 & C_4

$$2e \geq 4 + 4 + 4 + \dots + 4 = 4r$$

$$2e \geq 4r$$

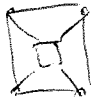
$$e \geq 2r$$

$$e - v + 2 = r \leq \frac{e}{2}$$

$$e - v + 2 \leq \frac{e}{2}$$

$$\frac{e}{2} \leq v - 2$$

$$\boxed{e \leq 2v - 4}$$

Q₂or degrees
are 4

$$v = 8$$

$$e = 12$$

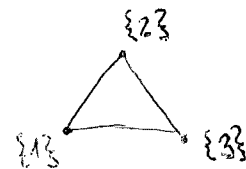
4) Odd graph

$$O_k \quad k = 2, 3, \dots$$

$k=2$ subsets with $2-1$ elements
that can be obtained from
a set with $2 \cdot 2 - 1 = 3$
elements

$$\{1, 2, 3\}$$

So O_2 has \dots vertices.



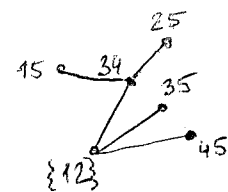
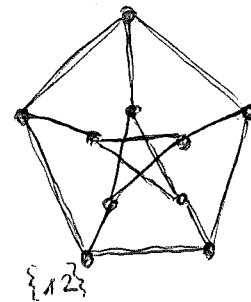
O_2 isomorphic
to K_3 or C_3

$k=3$ subsets with $3-1=2$ elements
..... from a set with $2 \cdot 3 - 1 = 5$ elements
so O_3 has $\binom{5}{2} = 10$ vertices.

so O_3 is 3-regular that is

degree = 3 for all vertices

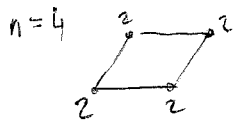
O_3 has \dots edges $\rightarrow \frac{3 \cdot 10}{2} \rightarrow$ handshake theorem

Q₃

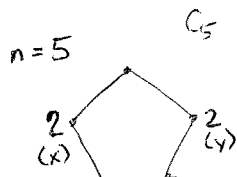
9. $|V| = n \geq 3$

If there are many edges we should find a HC

$\deg(x) + \deg(y) \geq n$



we find it directly. Every vertex has a degree 2



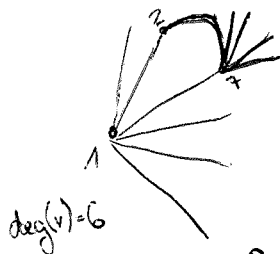
non-adjacent

$2+2 < 5$

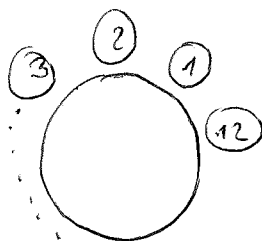
→ condition not fulfilled but cycle is there!

A graph with 66 vertices has a Hamilton cycle
This cycle has how many
...66... edges.

b.) 12 people



$6+6 \geq 12$ Yes!
There is a HC.



Put the HC around the table

6.) girth = smallest
length = # edges

$V = n$

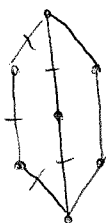
$E = n+1$

There are cycles in the Graph!

$g(G) \leq \left\lceil \frac{2(n+1)}{3} \right\rceil \quad n \geq 4$

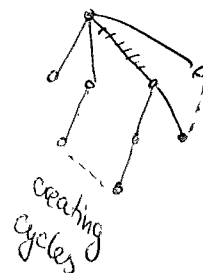
$n=7 \quad g \leq \left\lceil \frac{16}{3} \right\rceil = \underline{\underline{5}}$

$n=7$
 $e=8$



girth = shortest cycle = 5

$v=n=9$
 $e=8$



29/8-2 Bagels

I-E

$$x_p + x_t + x_b = 25$$

$$11 \geq x_p \geq 5$$

$$11 \geq x_t \geq 5$$

$$11 \geq x_b \geq 5$$

$$x_p = 5 + y_p$$

$$x_t = 5 + y_t$$

$$x_b = 5 + y_b$$

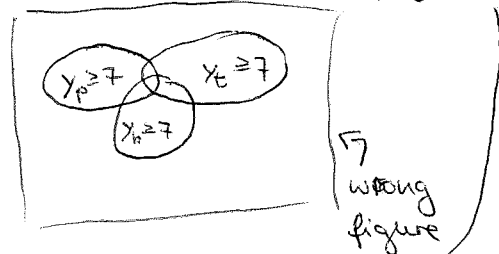
$$y_p + y_t + y_b = 10$$

$$0 \leq y_p \leq 6$$

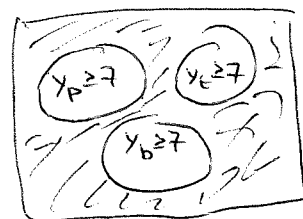
$$0 \leq y_t \leq 6$$

$$0 \leq y_b \leq 6$$

All poss. no upper restriction = $\binom{12}{10} = 66$



wrong figure



correct!

we are interested what is outside

$$y_p \geq 7:$$

$$y_p = 7 + z_p$$

$$7 + z_p + y_t + y_b = 10$$

$$z_p + y_t + y_b = 3$$

$$\binom{5}{2} = \binom{5}{3}$$

$$\text{Answer} = \binom{12}{10} - 3 \cdot \binom{5}{3} = 66 - 3 \cdot 10 = 36$$

With GF (generating function)

$$G(x) = \left(\frac{x^5 + x^6 + \dots + x^{11}}{x} \right) \cdot \left(\frac{x^5 + x^6 + \dots + x^{11}}{x} \right) \cdot \left(\frac{x^5 + x^6 + \dots + x^{11}}{x} \right)$$

$$= x^{15} + \dots + x^{25} + x^{33}$$

what is in front of x^{25}

$$G(x) = x^5 \cdot x^5 \cdot x^5 (1 + x + \dots + x^6)^3$$

$$= x^{15} \cdot \left(\frac{1-x^7}{1-x} \right)^3 = x^{15} \cdot (1 - 3x^7 + 3x^{14} - x^{21}) \cdot \frac{1}{(1-x)^3}$$

table formulas on the exam

$$= \left(x^{15} - 3x^{22} + 3x^{29} - x^{36} \right) \cdot \left(1 + 3x + \binom{4}{2}x^2 + \binom{5}{3}x^3 + \dots + \binom{12}{10}x^{10} \right)$$

$$\text{Ans } \binom{12}{10} - 3\binom{5}{3}$$

