

RSA Algorithm for Cryptography

Introduction To Discrete Mathematics
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Class: Data Science & Engineering(DSE)



Plain text



Encryption



Encrypted text



Decryption



Plain text

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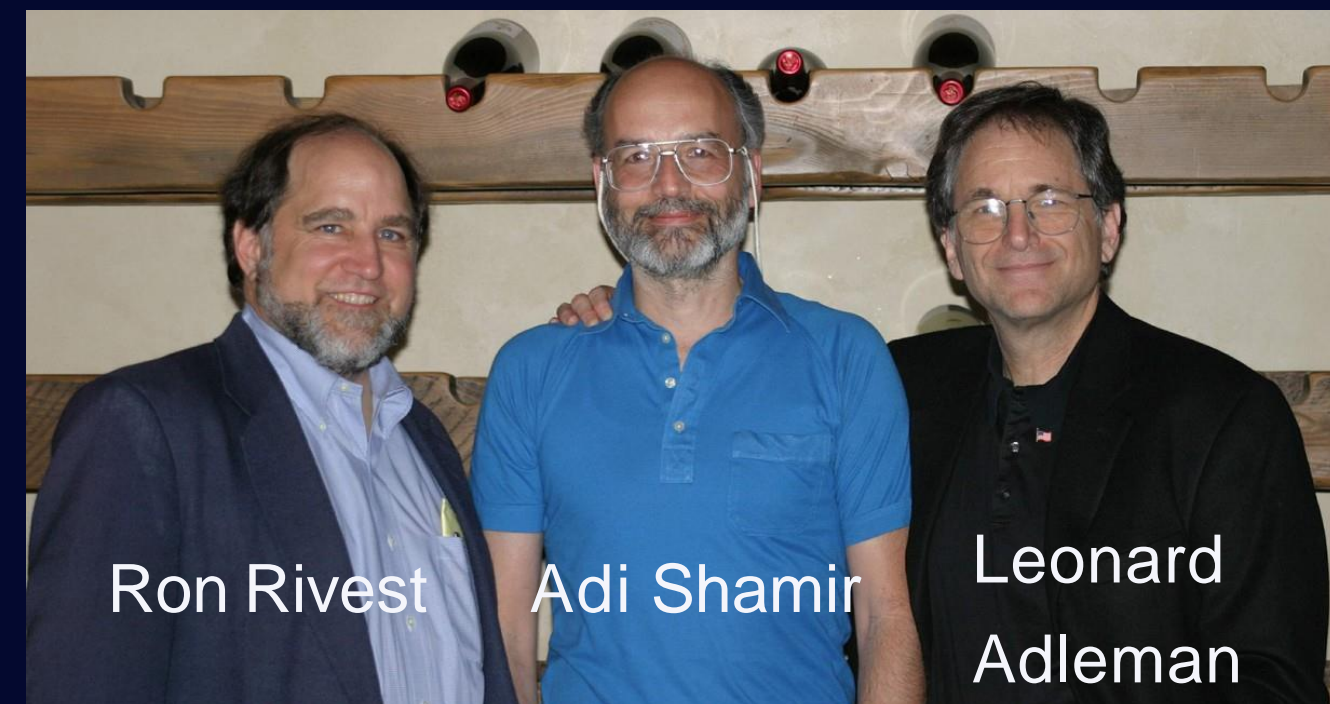
Example how to
Decrypt a Message



Invented by Ron Rivest, Adi Shamir,
and Leonard Adleman in 1977

RSA is a popular algorithm used to securely transmit data over the internet. It works by using a public key to encrypt messages, which can only be decrypted using a private key

Inventors



Ron Rivest

Adi Shamir

Leonard
Adleman

Key Generation



Key Generation

01



Public Key

02



Private Key

03



Key Generation

- Choose two large prime numbers (p) and (q).
- Calculate their product ($n = p * q$). And the value of (n) is part of the public key.
- Calculate the Euler's totient function of (n), $\phi(n)$ or $T = (p-1) * (q-1)$.
- Choose an integer (e), such that $1 < e < T$ and $[\gcd(e, T)=1]$. The value of (e) is part of a public key.
- Calculate the modular multiplicative inverse of (e) modulo T , which is another integer (d). The value of d is a private key.

Public Key (n,e)

- n : The modulus, which is the product of the two large prime numbers.
- e : The public exponent, used for encryption.

Private Key (n,d)

- n : The modulus, The same as in the public key.
- d : The private exponent, used for decryption.



Encryption

To encrypt a message (M)
represented as an integer into a
cipher text: (C)
 $C = M^e \pmod{n}$



01

Choose a Message to Send

The message can be any text
or data that you want to
transmit securely



02

Convert Message to Number

Each character in the message
is converted to a number,
based on its ASCII value



03

Apply Public Key

Using the public key, each
number in the message is
raised to the power of the
public key e , and then taken
modulo the public key n



Decryption

To decrypt the cipher-text (C)
back to the original message (M)
 $M = C^d \pmod{n}$



01

Apply Private Key

Using the private key, each number in the encrypted message is raised to the power of the private key, and then taken modulo the public key n



02

Convert Number to Message

The resulting numbers are converted back to characters, based on their ASCII values, to reveal the original message



Example of Key Generation

✓ Key generation:

- Choose two prime numbers $(p)=2$ and $(q)=7$
- Calculate $(n: p*q) = 2 * 7 = 14$
- Calculate $T = (p-1) * (q-1) = 1*6 = 6$
- Choose $\{(e, d) \bmod 6 = 1\}$ [e: encryption, d: decryption]; [Note: e and T should be relatively prime, and d is the inverse of e in mod T]
- Choose $e = 5, d = 11$

So, the public key is (14, 5) And The private key is (14, 11)



How to encrypt a message

➤ The published key is (14, 5)

So, let's send a one-letter secret message "B", Instead of (B) Let's use the number 2.

Suppose: A=1, B=2, C=3 and etc.

We get Encryption value = $2^5 \bmod 14$ is 4.

Therefore, the encrypted value is 4. The encrypted message of 4 will be translated to letter 'D' if we were to translate directly.

How to decrypt a message

➤ The private key is (14, 11)

So, we know that the secret value is 4

Supposed: A=1, B=2, C=3.... Etc.

We get the decrypted value = $4^{11} \bmod 14$

→ $4^{11} = 4,194,304 \bmod 14$

→ $4,194,304 / 14$ We get remainder 2

Therefore: 2 is a decrypted message which is letter "B"



```
import random

#function to the gcd of 2 numbers
def gcd(a, b):
    remainder = 0
    while(1):
        remainder = a % b
        if (remainder == 0):
            return b
        a = b
        b = remainder

#function to do extended euclidean for gcd, x and y which are coefficients for linear combination
def extended_euclidean(a, b):
    if a == 0:
        return b, 0, 1
    else:
        gcd, x1, y1 = extended_euclidean(b % a, a)
        x = y1 - (b // a) * x1
        y = x1
        return gcd, x, y

#function to find the modular inverse
def mod_inverse(e, t):
    gcd, x, y = extended_euclidean(e, t)
    if gcd != 1:
        raise ValueError("Modular inverse does not exist since 'e' and 't' are not coprime.")
    return x % t
```



```

if __name__ == '__main__':
    #initialize values for prime numbers p, q, calculate n and t and set initial value for e
    p = 124070563
    q = 334003807
    n = p * q
    e = 2
    t = (p - 1) * (q - 1)

    #find e which is co-prime with t
    while(1):
        if gcd(e, t) == 1:
            break
        else:
            e = random.randint(2, t-1)

    #calculate d as the modular inverse of e modulo t
    d = mod_inverse(e, t)

    #message to be encrypted
    message = 12
    print("")
    print("Original message data = ", message, end="\n\n")

    #encrypt the message
    ciphertext = pow(message, e, n)
    print("+ Encryption:")
    print("Encrypted message data (ciphertext) = ", ciphertext)
    print(f"Using: n = {n}, e = {e}", end = "\n\n")

    #decrypt the message
    decrypted_message = pow(ciphertext, d, n)
    print("+ Decryption:")
    print("Decrypted message data (original) = ", decrypted_message)
    print(f"Using: n = {n}, d = {d}", end = "\n\n")

```

Output

```
Original message data = 12
```

```
+ Encryption:
```

```
Encrypted message data (ciphertext) = 14352020436792370
```

```
Using: n = 41440040378633341, e = 4811942912821487
```

```
+ Decryption:
```

```
Decrypted message data (original) = 12
```

```
Using: n = 41440040378633341, d = 24038365170580919
```

Thank You for Paying Attention!