Two-level Implementation

CS207 Chapter 4

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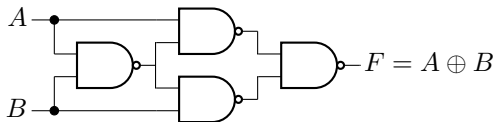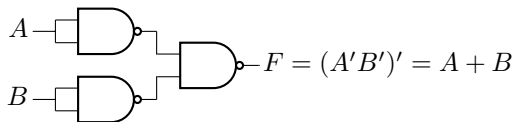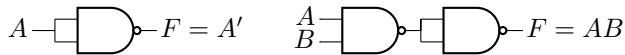
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NAND and NOR implementation

- Digital circuits are frequently constructed with NAND or NOR gates rather than with AND and OR gates.
  - NAND and NOR gates are easier to fabricate with electronic components.
  - They are the basic gates used in all IC digital logic families.
- Rules and procedures have been developed for the conversion from Boolean functions given in terms of AND, OR, and NOT into equivalent NAND and NOR logic diagrams.

NAND circuits

- The NAND gate is said to be a universal gate.
- A convenient way to implement a Boolean function with NAND gates:
  - Obtain the simplified Boolean function in terms of Boolean operators;
  - Convert the function to NAND logic.
- Recall that:

$$A \longrightarrow F = A' \qquad \begin{array}{c} A \\ B \end{array} \longrightarrow \longrightarrow F = AB$$

$$\begin{array}{c} A \\ B \end{array} \longrightarrow F = (A'B')' = A + B$$

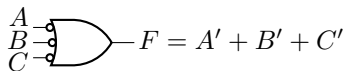$$\begin{array}{c} A \\ B \end{array} \longrightarrow F = A \oplus B$$

NAND circuits

- To facilitate the conversion to NAND logic, it is convenient to define an alternative graphic symbol for the gate.
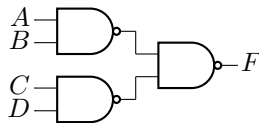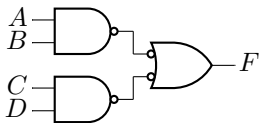
AND-invert:

$$\begin{matrix} A \\ B \\ C \end{matrix}\! \rangle\!\!\circ - F = (ABC)' = A' + B' + C'$$

Invert-OR:

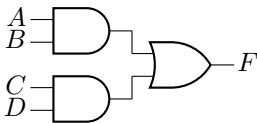$$\begin{matrix} A \\ B \\ C \end{matrix}\! \rangle - F = A' + B' + C'$$

NAND circuits

- The implementation of Boolean functions with NAND gates requires that the functions be in sum-of-products form.
- Take $F = AB + CD$ as an example:

- $F = AB + CD = ((AB)'(CD)')'$ according to DeMorgan property.

NAND circuits

- Example: Implement the following Boolean function with NAND gates:
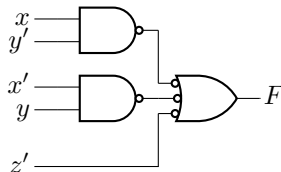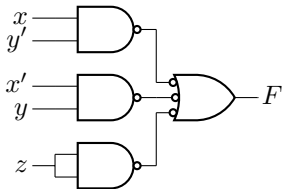 $F(x, y, z) = \sum(1, 2, 3, 4, 5, 7)$.

- $F = xy' + x'y + z$.

NAND circuits

- Example: Implement the following Boolean function with NAND gates: $F = xy' + x'y + z$.

NAND circuits

- A Boolean function can be implemented with two levels of NAND gates.
  1. Simplify the function and express it in **sum-of-products form**.
  2. Draw **a NAND gate for each product term** of the expression that has at least two literals. The inputs to each NAND gate are the literals of the term. This procedure produces a group of first-level gates.
  3. Draw **a single gate** using the AND-invert or the invert-OR graphic symbol **in the second level**, with inputs coming from outputs of first-level gates.
  4. A term with **a single literal requires an inverter** in the first level. However, if the single literal is complemented, it can be connected directly to an input of the second level NAND gate.

Multilevel NAND circuits

- The standard form of expressing Boolean functions results in a two-level implementation.
  - There are occasions when the design of digital systems results in gating structures with three or more levels.
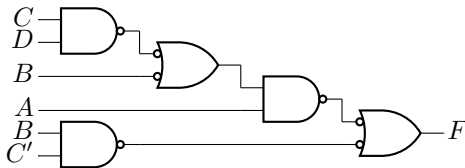 - Example: $F = A(CD + B) + BC'$.
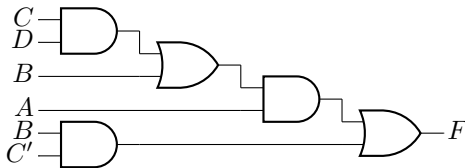
Multilevel NAND circuits

- The standard form of expressing Boolean functions results in a two-level implementation.
    - There are occasions when the design of digital systems results in gating structures with three or more levels.
    - Example: $F = A(CD + B) + BC'$.

Multilevel NAND circuits

- The general procedure for converting a multilevel AND–OR diagram into an all-NAND diagram using mixed notation is as follows:
  1. Convert all AND gates to NAND gates with AND-invert graphic symbols.
  2. Convert all OR gates to NAND gates with invert-OR graphic symbols.
  3. Check all the bubbles in the diagram. For every bubble that is not compensated by another small circle along the same line, insert an inverter (a one-input NAND gate) or complement the input literal.

NOR circuits

- The NOR operation is the dual of the NAND operation.
- All procedures and rules for NOR logic are the duals of the corresponding procedures and rules developed for NAND logic.
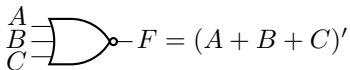
$$A \;—\!\!\Large\supset\!\circ— F = A'$$

$$\begin{array}{c}A\\B\end{array}\!\Large\supset\!\circ—\!\!\Large\supset\!\circ— F = A + B$$

$$\begin{array}{c}A\\[2em]B\end{array} \quad F = (A' + B')' = AB$$

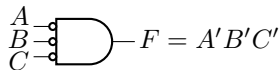## NOR circuits

- To facilitate the conversion to NOR logic, it is convenient to define an alternative graphic symbol for the gate.

OR-invert:

$$\begin{matrix} A \\ B \\ C \end{matrix} \!\!\!\!\!\!\Rightarrow\!\!\!\circ - F = (A + B + C)'$$

Invert-AND:

$$\begin{matrix} A \\ B \\ C \end{matrix} \!\!\!\!\!\!\Rightarrow\!\!\!- F = A'B'C'$$

NOR circuits

- A two-level implementation with NOR gates requires that the function be simplified into product-of-sums form.
- Change the OR gates to NOR gates with OR-invert graphic symbols and the AND gate to a NOR gate with an invert-AND graphic symbol.

Non-degenerate forms

- It will be instructive from a theoretical point of view to find out how many two-level combinations of gates are possible.
- We consider four types of gates: AND, OR, NAND, and NOR.
  - There are 16 possible combinations of two-level forms.
- Eight of these combinations are said to be *degenerate* forms.
  - They degenerate to a single operation.
  - Example: AND in the first level and second level degenerates to an AND of all inputs.
- The remaining eight nondegenerate forms produce an implementation in sum-of-products form or product-of-sums form.
  1) AND-OR 2) OR-AND 3) NAND-NAND 4) NOR-NOR
  5) NOR-OR 6) NAND-AND 7) OR-NAND 8) AND-NOR

AND-OR-INVERT implementation

- The two forms, NAND-AND and AND-NOR, are equivalent.
  - Both perform the AND-OR-INVERT function.
  - Example: $F = (AB + CD + E)'$.
- An AND-OR implementation requires an expression in sum-of-products form.
- The AND-OR-INVERT implementation is similar, except for the inversion.
  - If the complement of the function is simplified into sum-of-products form (by combining the 1's in the map), it will be possible to implement F with the AND-OR part of the function.

OR-AND-INVERT implementation

- The two forms, OR-NAND and NOR-OR, are equivalent.
  - Both perform the OR-AND-INVERT function.
  - Example: $F = [(A + B)(C + D)E]'$.
- The AND-OR-INVERT implementation requires an expression in product-of-sums form.
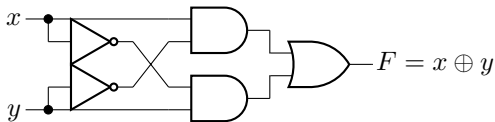
Exclusive-OR function

- Exclusive-OR, *XOR*: $x \oplus y = xy' + x'y$.
- Exclusive-NOR, *XNOR* or *equivalency*: $(x \oplus y)' = xy + x'y'$.
- The following identities apply to the XOR operation:
  - $x \oplus 0 = x$.
  - $x \oplus 1 = x'$.
  - $x \oplus x = 0$.
  - $x \oplus x' = 1$.
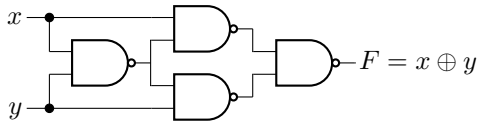 - $x \oplus y' = x' \oplus y = (x \oplus y)'$.

XOR

- XOR is hard to fabricate, so it is typically constructed by other gates.
  - $(x' + y')x + (x' + y')y = xy' + xy' = x \oplus y$.

$$F = x \oplus y$$

- Or use NAND gates:

$$F = x \oplus y$$

  - The first NAND gate performs the operation $(xy)' = x' + y'$.
  - The other two-level NAND circuit produces the sum of products.

XOR

- Only a limited number of Boolean functions can be expressed in terms of XOR operations.
- Particularly useful in arithmetic operations and error detection/correction circuits.
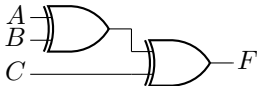
Odd function

- The XOR operation with three or more variables can be converted into an ordinary Boolean function by replacing the $\oplus$ with its equivalent Boolean expression.

$$\begin{aligned} A \oplus B \oplus C &= (AB' + A'B)C' + (AB + A'B')C \\ &= AB'C' + A'BC' + ABC + A'B'C \\ &= \sum(1, 2, 4, 7) \end{aligned}$$
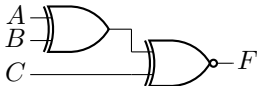
- The three-variable XOR function is equal to 1 if only one variable is equal to 1 or if all three variables are equal to 1.
    - An **odd** number of variables are equal to 1.
    - *Odd function*.

Odd function

- The three-input odd function is implemented by means of two-input XOR gates.

- *Even function* can also be implemented:

Parity generation and checking

- XOR functions are very useful in systems requiring error detection and correction codes.
  - A parity bit is an extra bit included with a binary message to make the number of 1's either odd or even.
- The circuit that generates the parity bit in the transmitter is called a *parity generator*.
- The circuit that checks the parity in the receiver is called a *parity checker*.

Parity generation and checking

- Consider a three-bit message to be transmitted together with an even-parity bit.

x	y	z	Parity bit p
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1
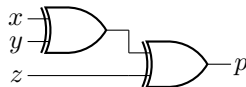
- p constitutes an odd function.

Parity generation and checking

- The three bits in the message, together with the parity bit, are transmitted to their destination.
- The four bits received must have an even number of 1's with even parity.