

Lab9 Solution

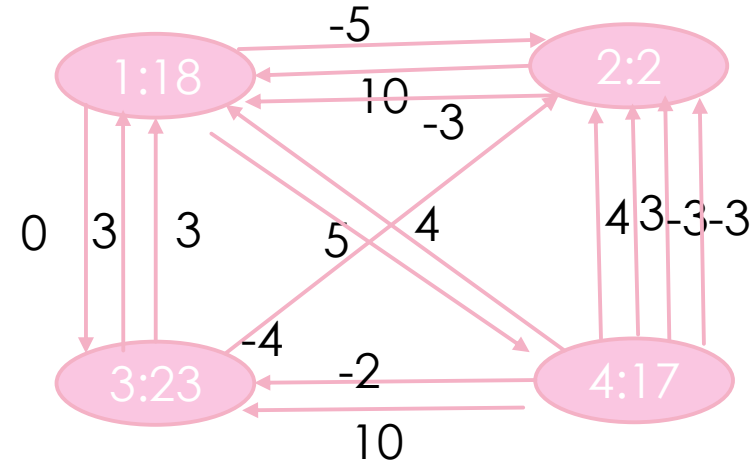
YAO ZHAO

Lab9.A: Voidwalker

- ▶ Once there was a voidwalker in the VOID who was eager for power.
- ▶ The VOID consists of N spots and M one-way tunnels. The voidwalker can walk through tunnel i by spending w_i energy (a positive w_i means losing w_i energy, a negative w_i means gaining $|w_i|$ energy, while $w_i = 0$ means no effect).
- ▶ The characteristic of energy differs in different spots. In spot i , the voidwalker can compose an energy core using a_i energy or decompose an energy core into a_i energy. The voidwalker can carry at most one energy core with him while walking through a tunnel, since there will be a disaster otherwise.
- ▶ The voidwalker begins its journey at spot S with initial power 0. It wonders the maximum energy it can reach with no more than $2K$ operations (an operation is either a compose or a decompose). It is valid for the energy to fall to negative.
- ▶ If the voidwalker can gain infinite energy, output **INVINCIBLE**.

Sample 2 Input

4 15 2 1
18 2 23 17
4 1 4
1 4 5
3 1 3
4 3 10
2 1 10
4 2 3
3 2 -4
3 1 3
1 2 -5
1 3 0
4 2 4
4 2 -3
2 1 -3
4 3 -2
4 2 -3



start from $S=1$ initial energy $E = 0$

Easy find there are some cycles which total $w_i < 0$:

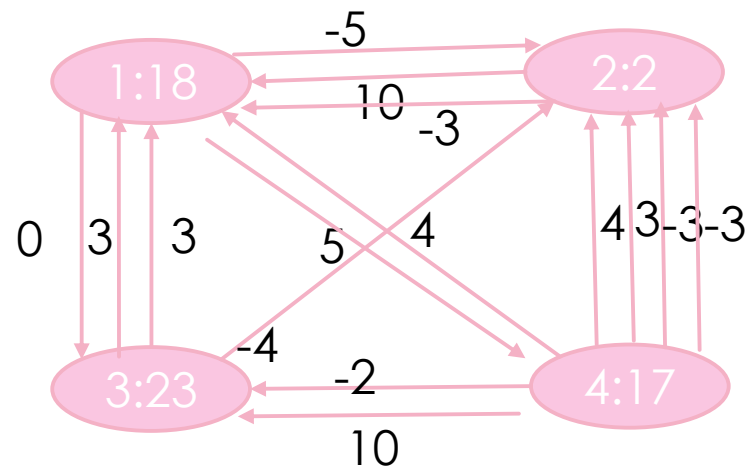
$1 \rightarrow 2 \rightarrow 1: -8$

$4 \rightarrow 2 \rightarrow 1 \rightarrow 4: -1$

$3 \rightarrow 2 \rightarrow 1 \rightarrow 3: -7$



INVINCIBLE



1	2	3	4
0	∞	∞	∞



bellmanFord iteration 1

1	2	3	4
0	-5	0	5



bellmanFord iteration 2

1	2	3	4
-8	-5	0	5



bellmanFord iteration 3

1	2	3	4
-8	-13	-8	-3



bellmanFord iteration 4

1	2	3	4
-16	-13	-8	-3

repeat n **iterations**

If some values are still updated in the n^{th} iteration, output “invincible” , stop.
If no values are updated in iteration $i \leq n$, You can return.

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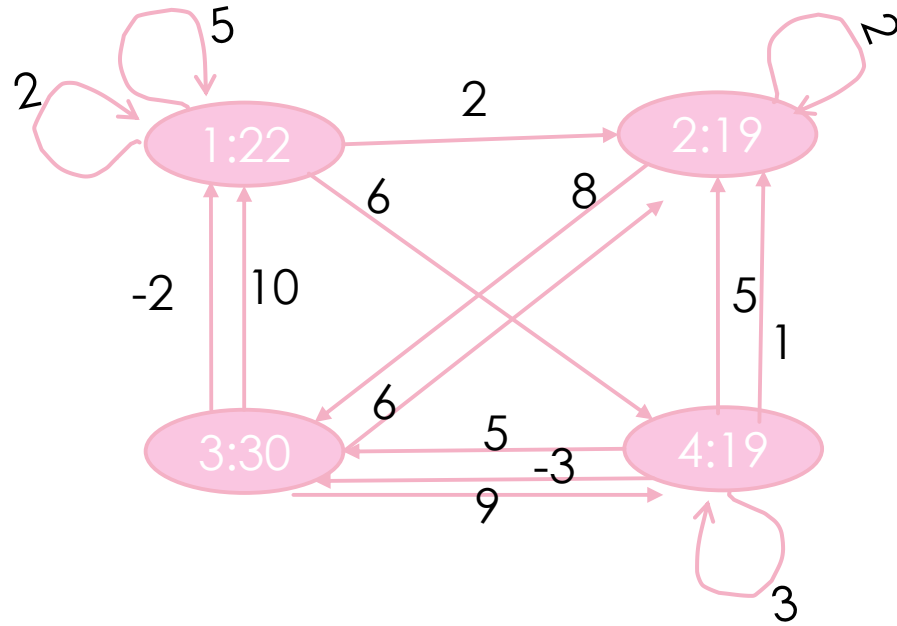
Push-Based-Shortest-Path( $G, s, t$ ) {
  foreach node  $v \in V$  {
     $M[v] \leftarrow \infty$ 
     $\text{successor}[v] \leftarrow \phi$ 
  }

   $M[t] = 0$ 
  for  $i = 1$  to  $n$  {
    foreach node  $w \in V$  {
      if ( $M[w]$  has been updated in previous iteration) {
        foreach node  $v$  such that  $(v, w) \in E$  {
          if ( $M[v] > M[w] + c_{vw}$ ) {
             $M[v] \leftarrow M[w] + c_{vw}$ 
             $\text{successor}[v] \leftarrow w$ 
          }
        }
      }
    }
    If no  $M[w]$  value changed in iteration  $i$ , stop.
    If  $i == n$  and still have  $M[w]$  value changed, output "INVINCIBLE"
  }
}

```

Sample 3 Input

4 15 3 3
 22 19 30 19
 2 3 8
 3 2 6
 1 2 2
 3 1 10
 4 3 5
 3 4 9
 1 1 2
 4 2 5
 4 4 3
 1 1 5
 4 3 -3
 4 2 1
 2 2 2
 3 1 -2
 1 4 6



using $f[0][n]$ to record the minimum distance from each node to s

1st compose an energy core

let $f[1][i] = f[0][i] + a[i]$

then do a bellman Ford using $f[1][n]$

$f[1][n]$ records the minimum cost from s to i and compose an energy core in the midway.

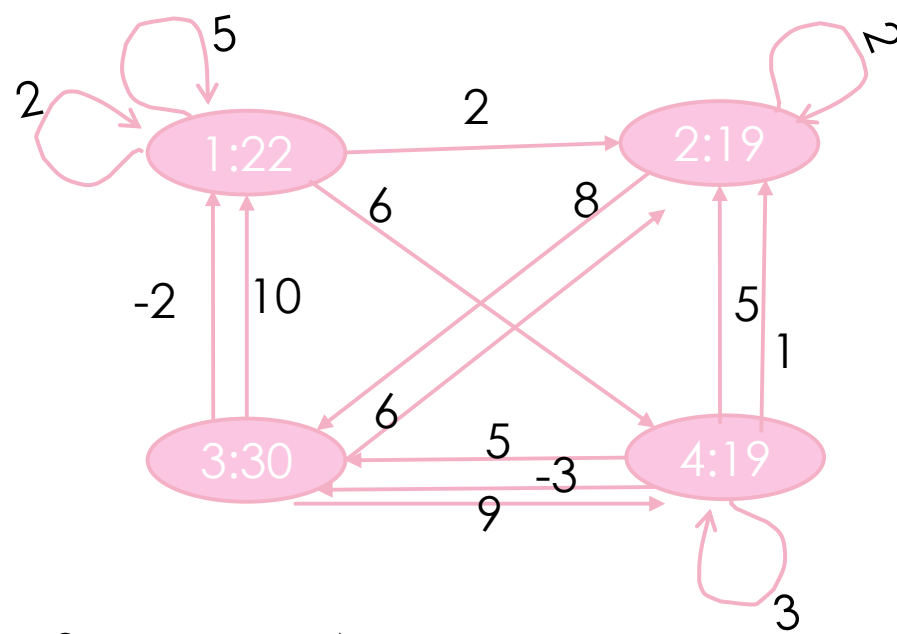
let $f[2][i] = f[1][i] - a[i]$

then do a bellman Ford using $f[2][n]$

$f[2][n]$ records the minimum cost from s to i and compose an energy core and decompose an energy in the midway.

repeat until spend all $2K$ operations.

output the minimum value in $f[0][n] \sim f[2k][n]$



f[0]: initial

1	2	3	4
∞	∞	0	∞

f[0]: final

1	2	3	4
-2	0	0	4

Bellman Ford

f[1]: initial

1	2	3	4
20	19	30	23

f[1]: final

1	2	3	4
18	19	20	23

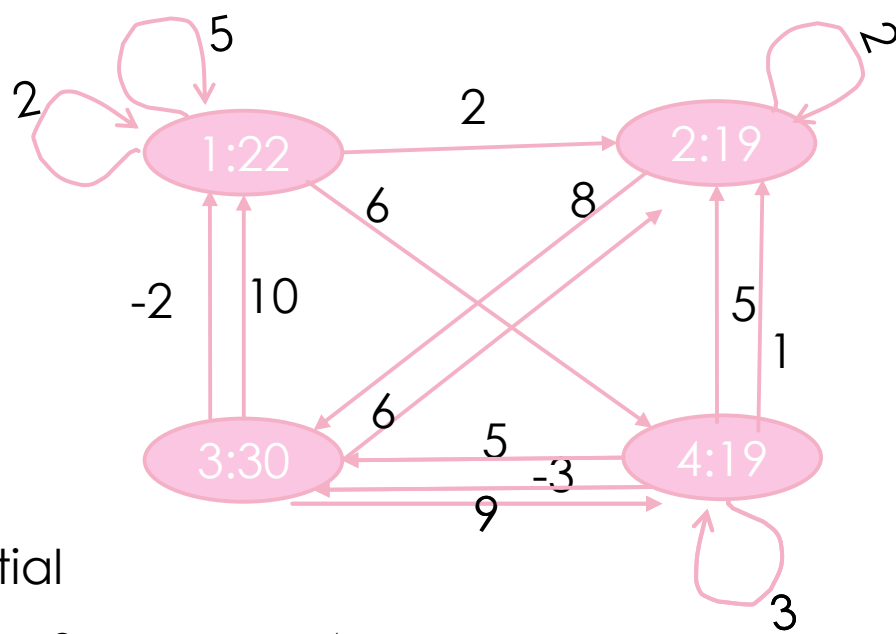
Bellman Ford

f[2]: initial

1	2	3	4
-4	0	-10	4

-a[i]

+a[i]



f[2]:initial

1	2	3	4
-4	0	-10	4

f[2]:final

Bellman Ford

1	2	3	4
-12	-10	-10	-6

f[3]:initial

1	2	3	4
10	9	20	13

f[3]:final

Bellman Ford

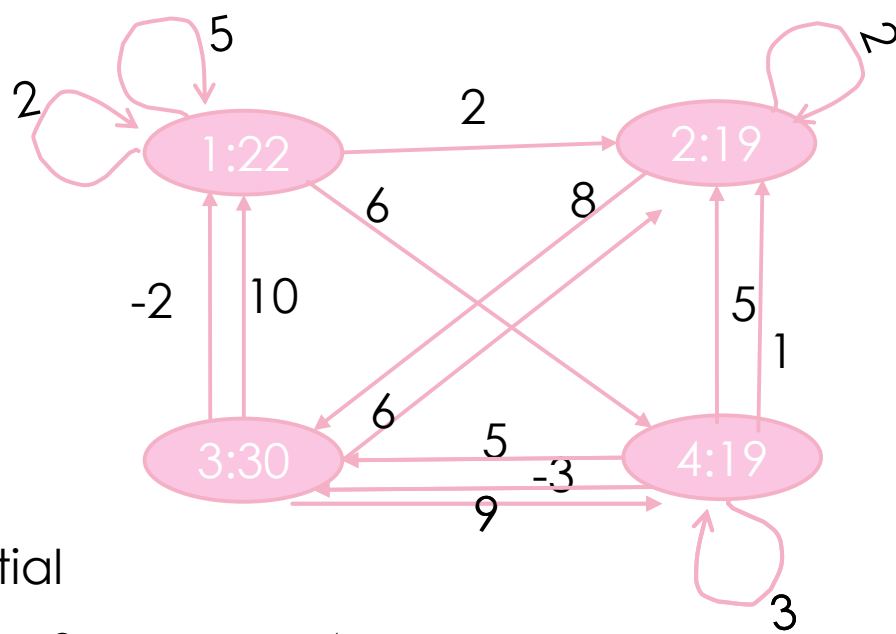
1	2	3	4
8	9	10	13

+a[i]

f[4]: initial

-a[i]

1	2	3	4
-14	-10	-20	-6



f[4]:initial

1	2	3	4
-14	-10	-20	-6

f[4]:final



Bellman Ford

1	2	3	4
-22	-20	-20	-16

f[5]:initial

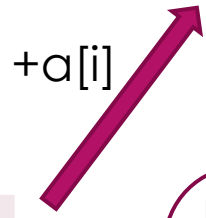
1	2	3	4
0	-1	10	3

f[5]:final



Bellman Ford

1	2	3	4
-2	0	0	3

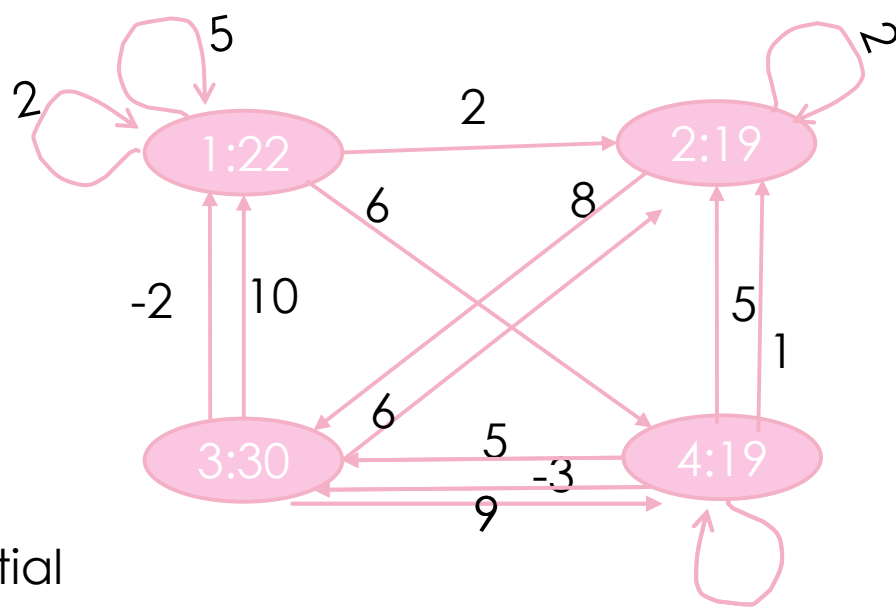


f[6]:initial

-a[i]



1	2	3	4
-24	-19	-30	-16



f[6]:initial

1	2	3	4
-24	-19	-30	-16

f[6]:final



Bellman Ford

1	2	3	4
-32	-30	-30	-26

Sample 3 Output

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Lab9.B: Scream Out Loud

- ▶ Lida Pu has long suffered from a compulsion to obtain symmetric things, for example, palindrome strings.
- ▶ One day, Lida Pu received a secret mail, in which he saw a string template. The template contains lowercase letters, symbol '?' corresponding to an arbitrary letter and symbol '*' corresponding to a zero or more arbitrary letters.
- ▶ Please tell Lida Pu the minimum length of the palindrome string which can be obtained from the given template. If he cannot get a palindrome string anyhow, just tell him to face the reality.
- ▶ Note that in Lida Pu's mind, an empty string is also a palindrome string.

symbol '?' corresponding to an arbitrary letter

symbol '*' corresponding to a zero or more arbitrary letters

Sample 1 Input

***ac?ba**

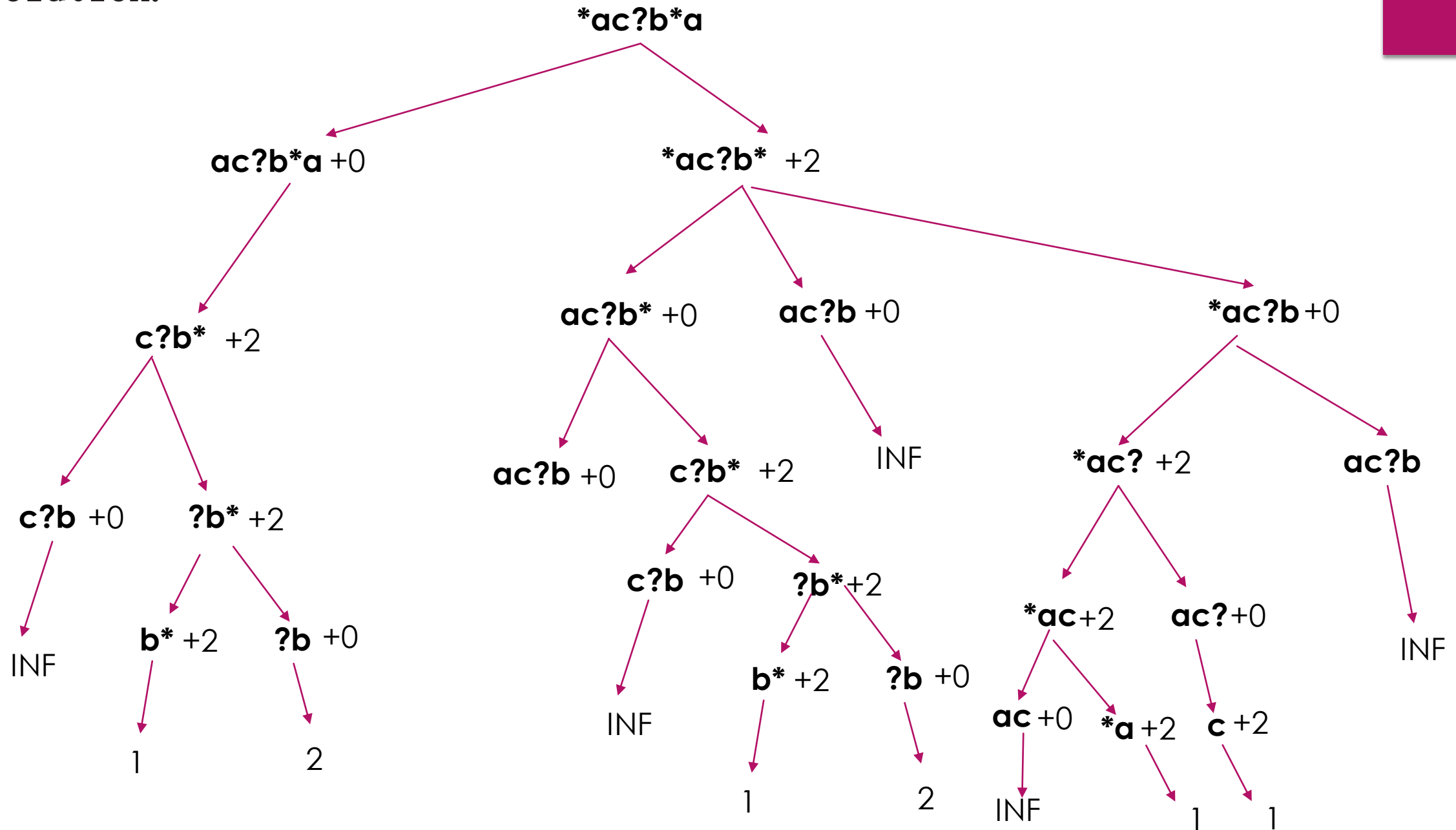
try length = 5 → **ac?ba** **fail**

try length = 6 $\xrightarrow{\text{let } * = a}$ **aac?ba** **fail**

try length = 7 $\xrightarrow{\text{let } * = ab}$ **abac?ba** $\xrightarrow{\text{let } ? = a}$ **abacaba** **success!**

↓
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1. Characterize the structure of an optimal solution.



2. Recursively define the value of an optimal solution.

let $|s| = n$

Input String: $c[1]c[2]\dots c[n]$

Let $\text{opt}[i][j]$ is the **minimum** length of the palindrome string $c[i]c[i+1]\dots c[j]$, then:

if $c[i] == '*'$

$\text{opt}[i][j] = \min\{\text{opt}[i+1][k]+0, \text{opt}[i][j-1]+2\}$

if $c[j] == '*'$

if $c[i] == '*' \ \&\& \ c[j] == '*'$

if $c[i]$ can match $c[j]$

$\text{opt}[i][j] = \min\{\text{opt}[i+1][j-1]\}$

3. Compute the value of an optimal solution, typically in a bottom-up fashion.

***ac?b*a**

	1	2	3	4	5	6	7
1	0	1	3	3	5	5	6
2		1	∞	3	∞	6	6
3			1	2	∞	4	∞
4				1	2	2	3
5					1	1	∞
6						0	1
7							1

acbbca