

Internship defense M2  
Parisian Master of Research in Computer Science  
Élimination non-singleton :  
Avec les faisceaux, c'est pas faux

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# Table of contents

1 Topoi

2 Sheaves in Type theory

3 Bibliography

# The starting point

Cartesian closed categories	$\Leftrightarrow$	Simply typed $\lambda$ -calculus	$\Leftrightarrow$	Negative fragment of intuitionistic propositional logic
$\text{Hom}(A \times B, C)$ $\simeq$ $\text{Hom}(A, B \rightarrow C)$		$\lambda x : A. t : A \rightarrow B$		$\frac{\Gamma, A \vdash B}{\Gamma \vdash A \rightarrow B}$

# Subobject classifier

$$\text{Sub}(A) \simeq \text{Hom}(A, \Omega)$$

$\text{Sub}(A) :=$  equivalence classes of monomorphisms  $X \rightarrowtail A$ .

Introduces a type of proposition in the  $\lambda$ -calculus, as well as subobjects

- All finite limits as well. Finite colimits are recovered[1].

# Formulas

$\wedge, \vee, \forall, \exists, \neg, \perp, \top, \Rightarrow$  can all be expressed e.g.

$$\begin{array}{ccc} X \wedge Y & \longrightarrow & Y \\ \downarrow & & \downarrow \\ X & \longrightarrow & A \end{array}$$

So can  $=$  and  $\in$ . Equality is the predicate  $A \times A \rightarrow \Omega$  given by the diagonal  $A \rightarrowtail A \times A$ .

Objects similar to the set theoretic  $\{x \in A \mid \phi(x)\}$  can be constructed.

# Internal logic of $\Omega$

- $\wedge : \text{Sub}(A) \times \text{Sub}(A) \rightarrow \text{Sub}(A)$
- $\wedge : \text{Hom}(A, \Omega) \times \text{Hom}(A, \Omega) \rightarrow \text{Hom}(A, \Omega)$
- $\wedge : \text{Hom}(A, \Omega \times \Omega) \rightarrow \text{Hom}(A, \Omega)$
- $\wedge : \Omega \times \Omega \rightarrow \Omega$

Similarly for  $\vee, \perp, \top, \Rightarrow$ .

# Topoi of presheaves

For a small category  $\mathbf{C}$ , its category of presheaves is noted  $\widehat{\mathbf{C}} := \mathbf{Set}^{\mathbf{C}^{op}}$ .

$$(P \times Q)(C) = P(C) \times Q(C)$$

$$P^Q(C) \simeq \mathbf{Nat}(yC, P^Q) \simeq \mathbf{Nat}(yC \times Q, P)$$

$\Omega(C)$  is the set of subpresheaves of  $yC$  (also called sieves on  $C$ ).  
(A subpresheaf is pointwise a subset, with function restricted)

Small limits are pointwise.

# Sheaves

A sheaf = A presheaf with a compatibility property.

A Grothendieck topology introduces a notion of cover.

- If  $f : D \rightarrow C$  is in  $S(D)$ , then it is covered by  $S$ .
- If  $f$  is covered then  $f \circ g$  is covered.
- If  $S$  covers  $f$  and  $R$  covers any map of  $S(D)$  for all  $D$ , then  $R$  covers  $f$

We note the sieve of all maps covered by  $S$  as  $\overline{S}$ .

A sheaf is a presheaf where the inclusion of  $\text{Nat}(\overline{S}, F)$  in  $\text{Nat}(S, F)$  is an equivalence.



# Example : Sheaves on $\mathbb{R}^2$

$\mathbf{C} := \mathcal{O}(\mathbb{R}^2)$  (poset : at most one map between object).

We fix  $C := \mathbb{R}^2$ .

$S$  covers  $U := (U \subseteq \bigcup S)$  :

- if  $U \in S$  then  $U \subseteq \bigcup S$
- if  $U \subseteq \bigcup S$ , and  $V \subset U$ , then  $V \subseteq \bigcup S$
- if  $U \subseteq \bigcup S$ , and for all  $V \in S$ ,  $V \subseteq \bigcup R$ , then  $U = \bigcup R$

# The topos of sheaves

Finite limits of sheaves as presheaves are sheaves and are finite limits for sheaves.

$F^P$  is a sheaf whenever  $F$  is.

$\Omega$  contains the sieves which are also sheaves, that is which only cover themselves.

# Lawvere-Tierney topology

Sheaves are defined through an operator  $j : \Omega \rightarrow \Omega$  for which

- $id \leq j$ , since  $S \subseteq \overline{S}$
- $\Rightarrow \circ id \times j \leq \Rightarrow \circ j \times j : \Omega \times \Omega \rightarrow \Omega$ , since  $\overline{S} \subseteq \overline{\overline{R}}$  whenever  $S \subseteq \overline{R}$

This can be defined in any topos.

Or, rather, any type theory with a type of (strict) propositions.

# Table of contents

1 Topoi

2 Sheaves in Type theory

3 Bibliography

# $J$ -Sheaves

Given  $J : \text{Prop} \rightarrow \text{Prop}$  with

- $\text{ret}_J : \prod P, P \rightarrow JP$
- $\text{bind}_J : \prod QP, JP \rightarrow (P \rightarrow JQ) \rightarrow JQ$

A sheaf is then a type  $F$  together with

- A map  $\text{ask} : \prod P, JP \rightarrow (P \rightarrow F) \rightarrow F$
- Such that  $\prod P j, \text{ask } P j (\lambda \_, f) = f$

# $I$ O-sheaves

Noting  $I := \Sigma P. JP : \text{Type}$  and  $O(P, j) := P : I \rightarrow \text{Prop}$ ,

A sheaf is equivalently a type  $F$  where:

- $\text{ask} : \Pi i : I, (O\ i \rightarrow F) \rightarrow F$
- $\Pi i : I, \text{ask}\ i\ (\lambda \_, f) = f$

# Table of contents

1 Topoi

2 Sheaves in Type theory

**3 Bibliography**

# Bibliography I

- [1] Saunders MacLane and Ieke Moerdijk. *Sheaves in geometry and logic: A first introduction to topos theory*. Springer Science & Business Media, 2012.