
Queue ADT

**IT5003: Data Structures and Algorithms
(AY2019/20 Semester 1)**

Lecture Overview

- Queue
 - ❑ Introduction
 - ❑ Specification
 - ❑ Implementations
 - Linked List Based
 - Array Based
 - ❑ Application
 - Palindrome checking

What is a Queue

- Real life example:
 - ❑ A queue for movie tickets, Airline reservation queue, etc
- First item added will be the first item to be removed
 - ❑ Has the **First In First Out (FIFO)** property
- Major Operations:
 - ❑ **Enqueue:** Items are added to the **back of the queue**
 - ❑ **Dequeue:** Items are removed from the **front of the queue**
 - ❑ **Get Front:** Take a look at the first item

Queue: **I**llustration



A **queue** of
3 persons



Enqueue a new
person to the **back**
of the **queue**



Dequeue a person from
the **front of the queue**

Queue ADT: **P**ython **S**pecification

```
class QueueBase(ABC):  
    @abstractmethod  
    def getFront(self):  
        pass  
  
    @abstractmethod  
    def enqueue(self, newItem):  
        pass  
  
    @abstractmethod  
    def dequeue(self):  
        pass  
  
    @abstractmethod  
    def size(self):  
        pass  
  
    @abstractmethod  
    def isEmpty(self):  
        pass
```

**Major
Operations
for Queue**

Design Considerations

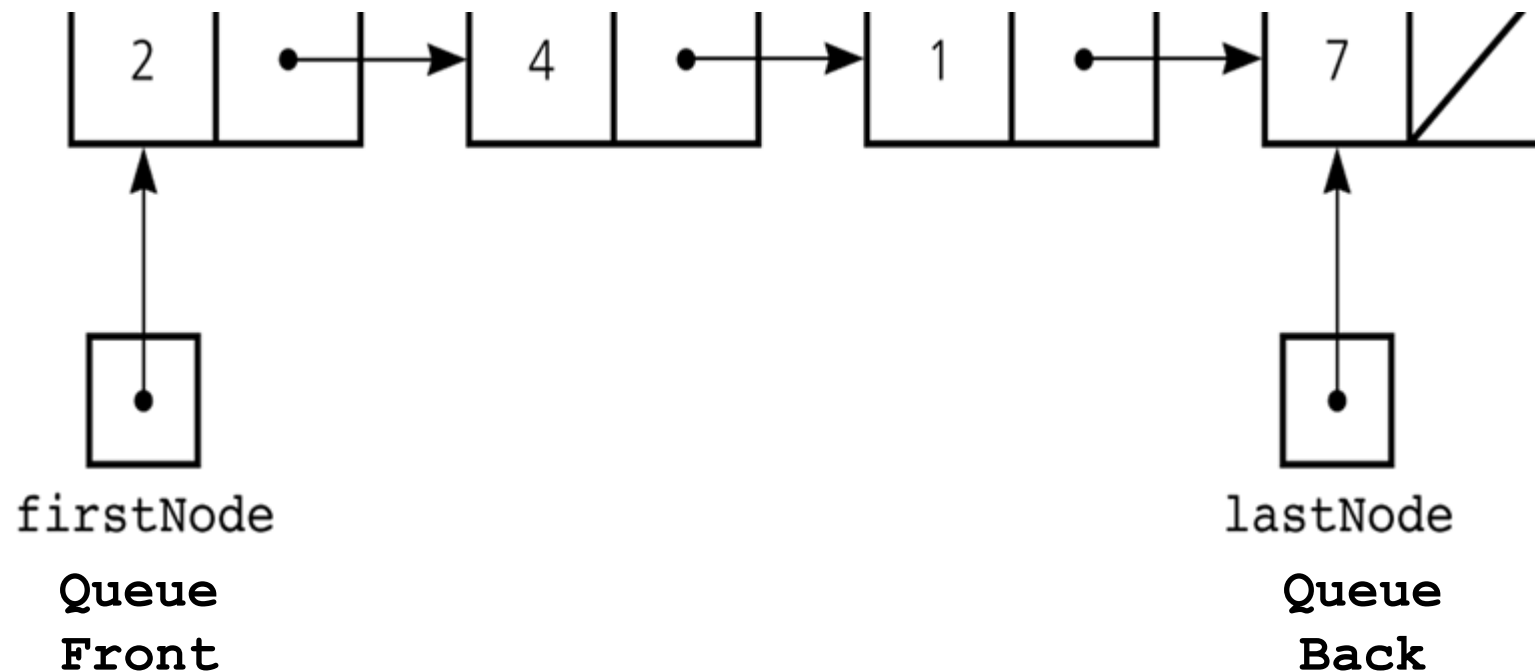
- How about the common coding choices?
 - Efficiency of **singly linked list implementation**:
 - 👍 Removing item at the head is the best case
 - 👎 Adding item at the back is the worst case
 - Efficiency of **array based implementation**:
 - 👎 Removing item at the head is the worst case
 - 👍 Adding item at the back is the best case
- Is it possible to have both efficient **enqueue()** and **dequeue()** operations?

QUEUE ADT USING LINKED LIST VARIANT

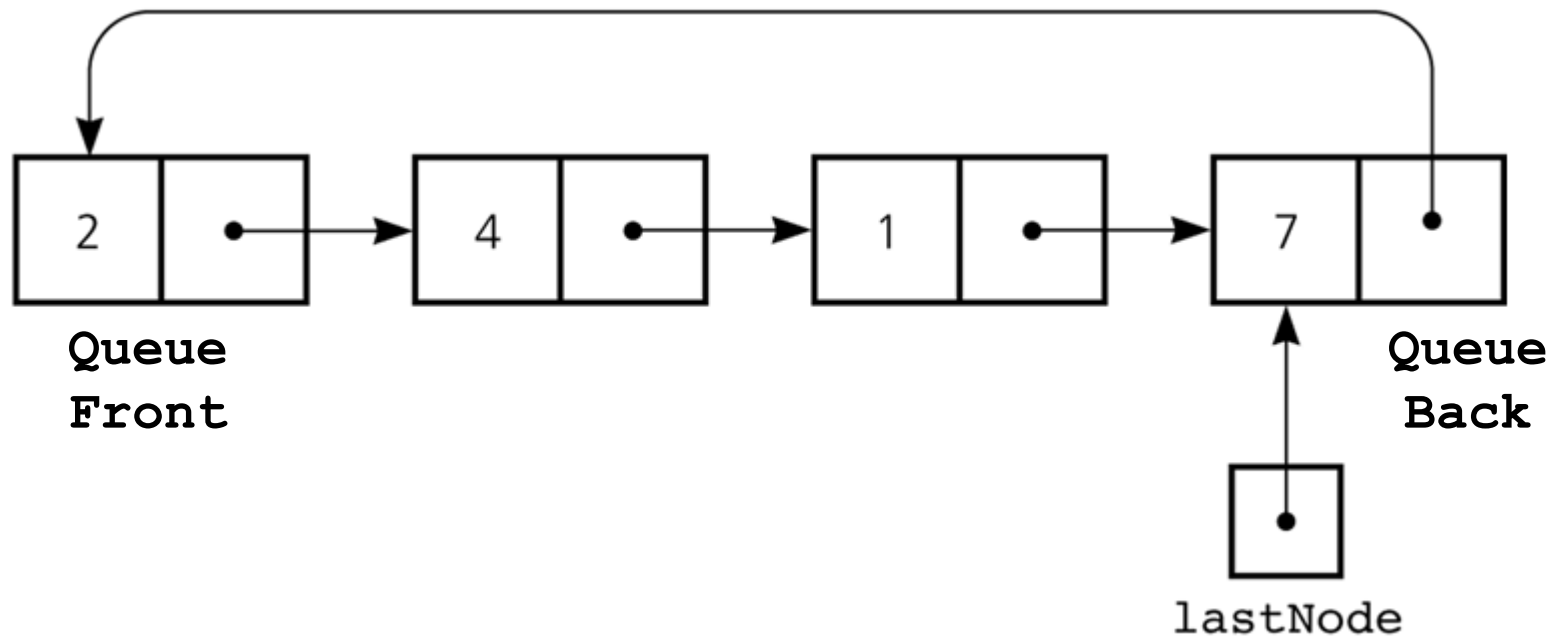
Improving the Singly Linked List

- Singly linked list performs badly for **enqueue()**
 - ❑ Need to traverse all the way to the last node
 - ❑ Takes longer time as the queue grows
- How to avoid the traversal to the last node?
 - ❑ Just need to “know” where is the last node all the time.....
- **Solutions:**
 - ❑ Keep an additional reference to the last node, OR
 - ❑ Circular linked list with a tail reference

Linked List : **head** and **tail**

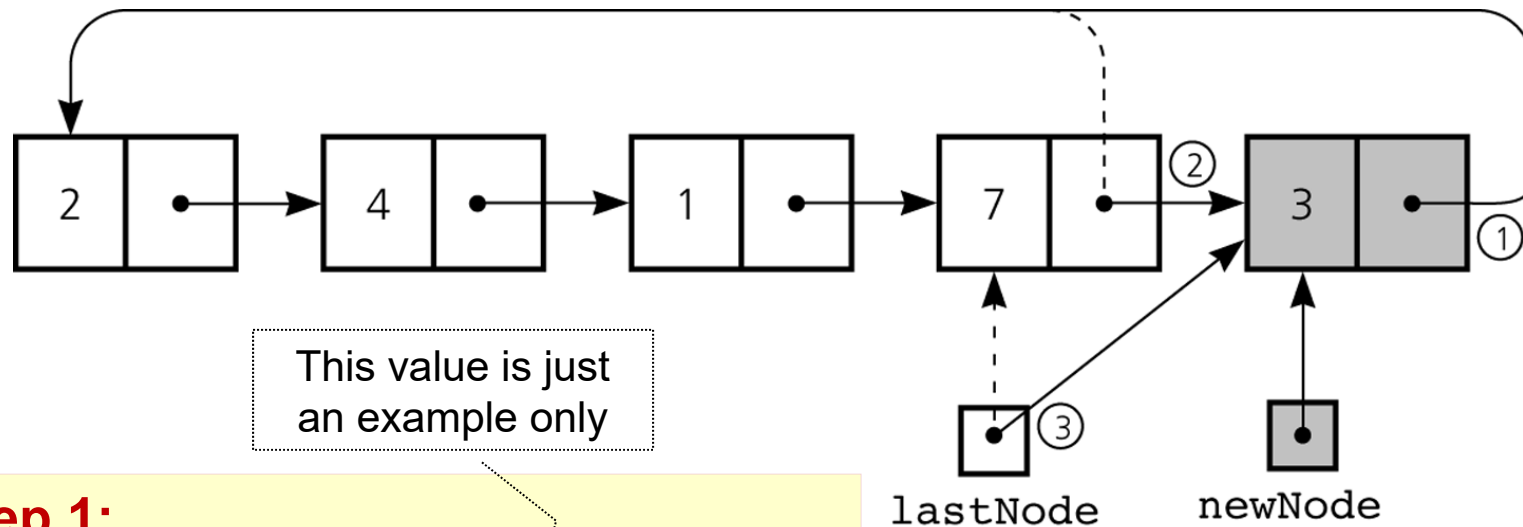


Circular Linked List



- Only keep tracks of `lastNode` reference
 - ❑ `firstNode` reference can be set when needed:
 - `firstNode = lastNode.next`
- We use circular linked list in our implementation

Enqueue: Non-Empty Queue



Step 1:

```
newNode = SinglyNode(3)  
newNode.next = _lastNode.next
```

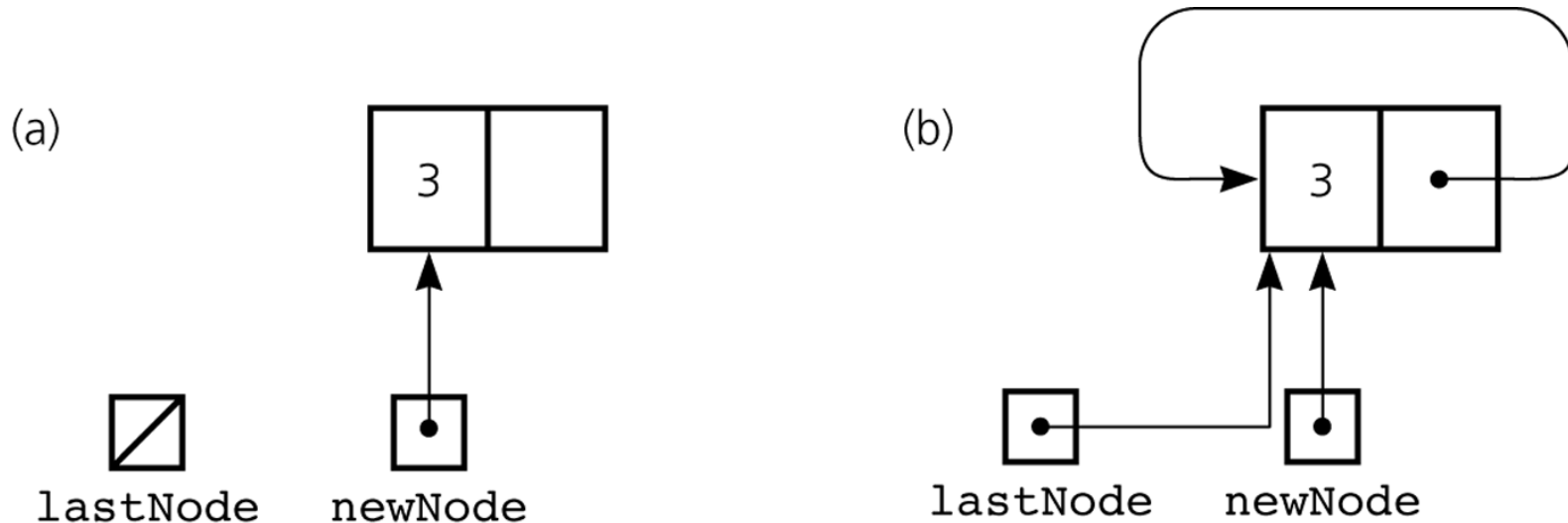
Step 2:

```
self._lastNode.next = newNode
```

Step 3:

```
_lastNode = newNode
```

Enqueue: Empty Queue



Step (a):

`newNode = SinglyNode(3)`

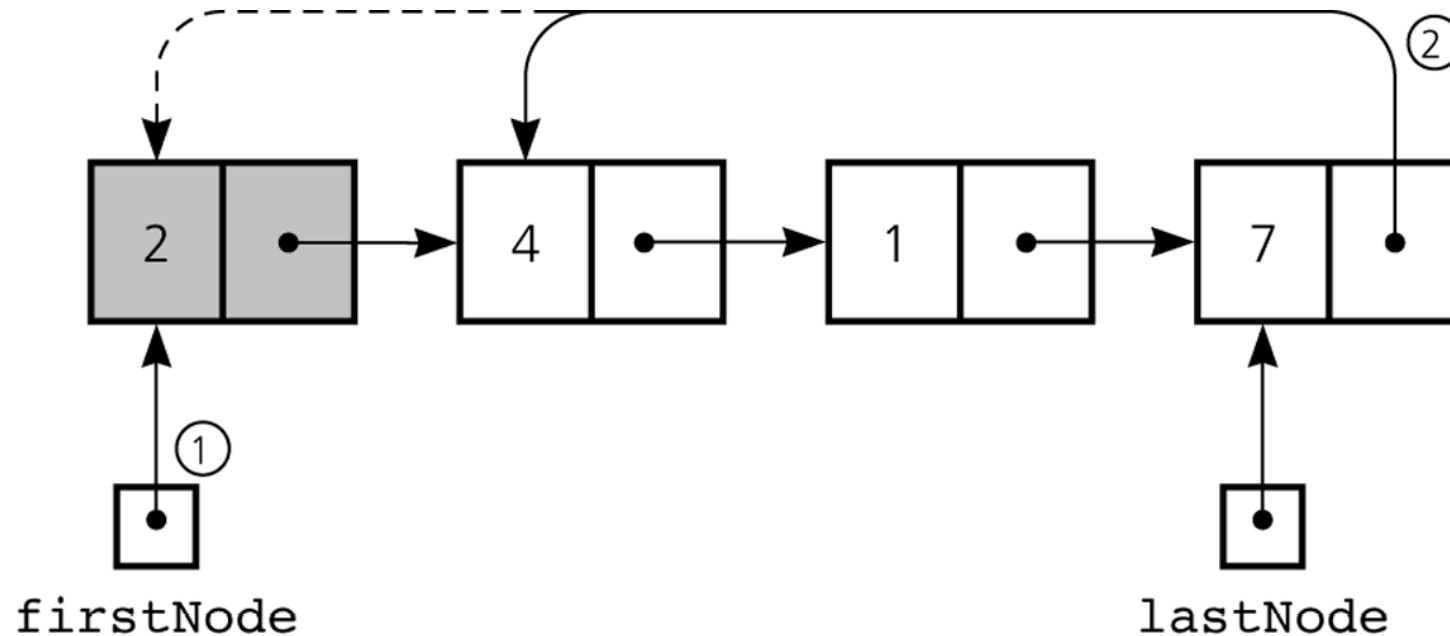
Step (b):

`newNode.next = newNode`

`self._lastNode = newNode`

Set up the "loop"

Deque: Queue with > 1 item



Step 1:

`firstNode = lastNode.next`

Step 2:

`lastNode.next = firstNode.next`

QueueLinkedList: Implementation

```
class QueueLinkedList(QueueBase):  
  
    def __init__(self):  
        self._lastNode = None  
        self._size = 0  
  
    def getFront(self):  
        if not self.isEmpty():  
            return self._lastNode.next.item  
        else:  
            return None
```

QueueLinkedList: Implementation

```
def enqueue(self, newItem):  
    newNode = SinglyNode(newItem)
```

```
    if self._lastNode == None:  
        newNode.next = newNode  
        self._lastNode = newNode
```

Refer to the algorithm
discussion slides

```
    else:  
        newNode.next = self._lastNode.next  
        self._lastNode.next = newNode  
        self._lastNode = newNode
```

```
    self._size += 1  
    return True
```

QueueLinkedList: Implementation

```
def dequeue(self):
    if not self.isEmpty():
        firstNode = self._lastNode.next
        self._lastNode.next = firstNode.next

        if self._size == 1:
            self._lastNode = None

        self._size -= 1
        return True
    else:
        return False

def size(self):
    return self._size

def isEmpty(self):
    return self._size == 0
```

Refer to the algorithm
discussion slides

Take note of the 1
node special case

What?! Array can be circular?

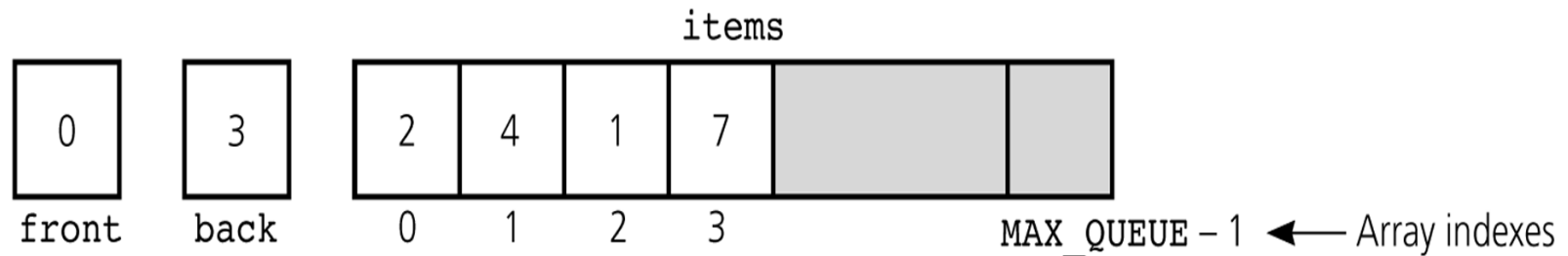
QUEUE ADT USING CIRCULAR ARRAY

Array Implementation Issues

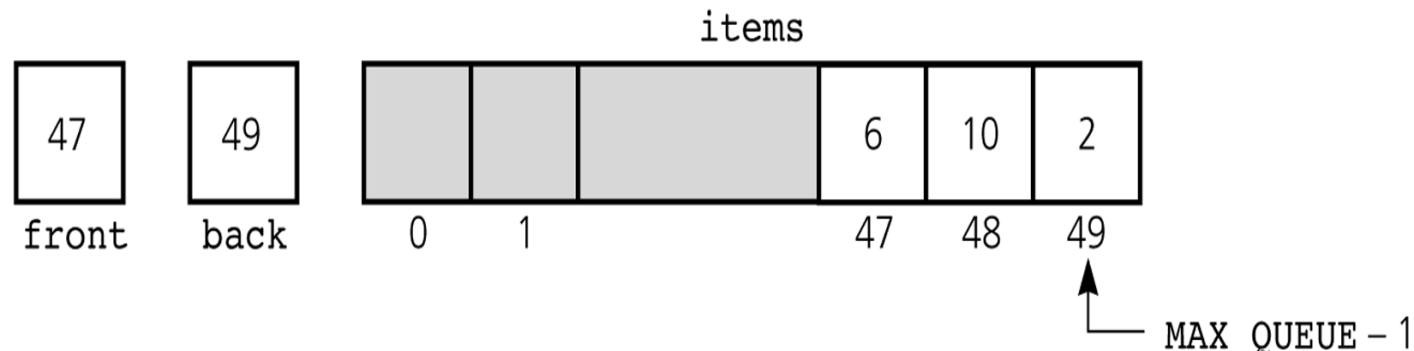
- Removing item from the front is inefficient
 - ❑ **Shifting items is too expensive**
- **Basic Idea:**
 - ❑ The reason for shifting is:
 - Queue's **front** is assumed to be at **index 0**
 - ❑ Instead of shifting items:
 - Why don't we just **shift the front index**
- So, we have two indices:
 - ❑ **Front:** index of the queue front
 - ❑ **Back:** index of the queue back

Incorrect Implementation

- At the beginning, with 4 items queued

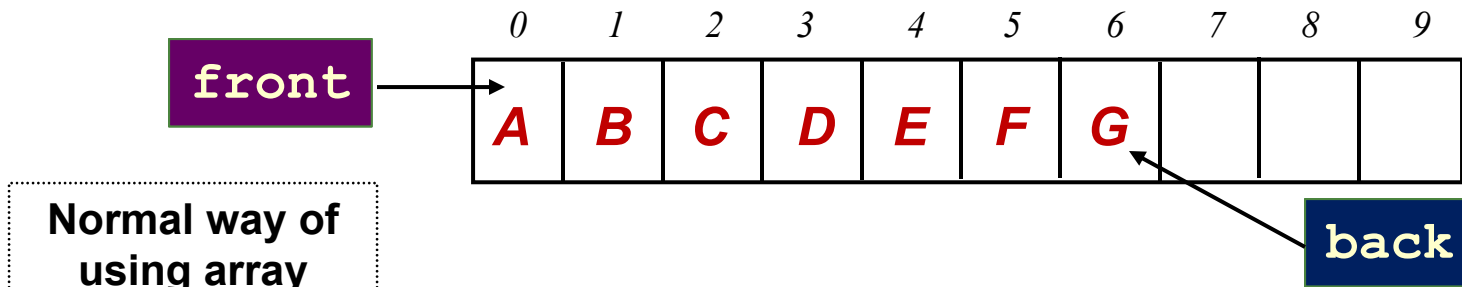


- After many queue operations



- The front index will drift to the right:
 - Most array locations empty and unusable

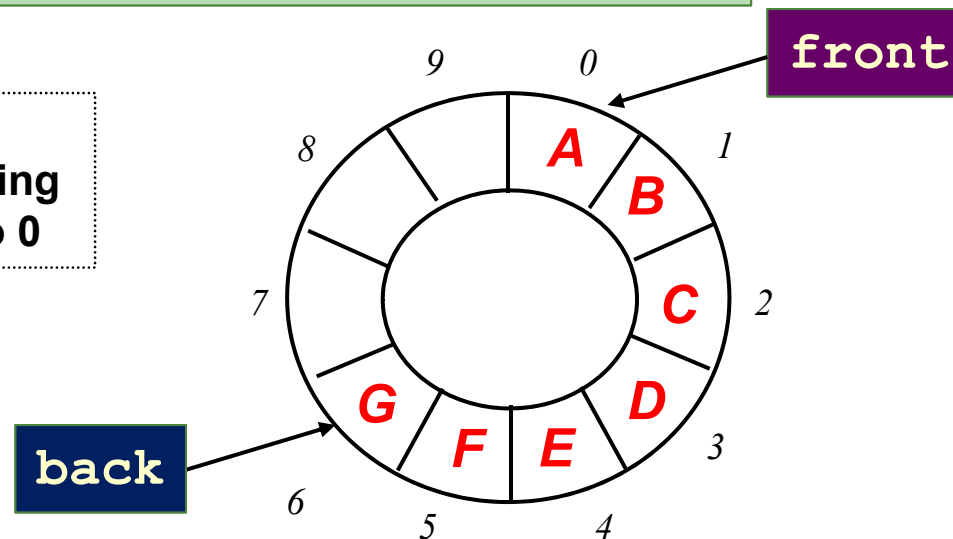
Circular Array: Illustration



The “magic” ingredients

$$\text{front} = (\text{front} + 1) \% \text{maxsize}$$
$$\text{back} = (\text{back} + 1) \% \text{maxsize}$$

Make the array
“circular” by wrapping
the indices back to 0



QueueCircularArray: Implementation

```
class QueueCircularArray(QueueBase):  
    MAXSIZE = 50 #arbitrarily chosen as example  
  
    def __init__(self):  
        self._items \  
        = (QueueCircularArray.MAXSIZE * ctypes.py_object)()  
  
        self._size = 0  
  
        self._front = 0  
        self._back = QueueCircularArray.MAXSIZE - 1
```

MAXSIZE is a class attribute

**Use static array to reduce "syntax burden".
Can be extended to dynamic array**

**_front : array index of the queue front
_back : array index of the queue back**

QueueCircularArray: Implementation

```
def enqueue(self, newItem):
    if self._size == QueueCircularArray.MAXSIZE:
        return False
    else:
        self._back \
            = (self._back + 1) % QueueCircularArray.MAXSIZE
        self._items[self._back] = newItem
        self._size += 1

    return True

def dequeue(self):
    if not self.isEmpty():
        self._front \
            = ( self._front+1 ) % QueueCircularArray.MAXSIZE
        self._size -= 1
        return True
    else:
        return False
```

Focus on the update of the
_front and _back indices

Only selected methods
implementation are shown. As
you should be able to code the
rest easily by now 😊.

Queue has a standard implementation

PYTHON BUILT-IN QUEUE

Python Collections: deque

```
from collections import deque
```

```
def main():  
    dq = deque()
```

#Items can be added both ends

```
dq.append(111) #added to the "right", the tail end
```

```
dq.append(122)
```

```
dq.appendleft(999) #added to the "left", the front end
```

```
dq.appendleft(988)
```

```
print(dq)
```

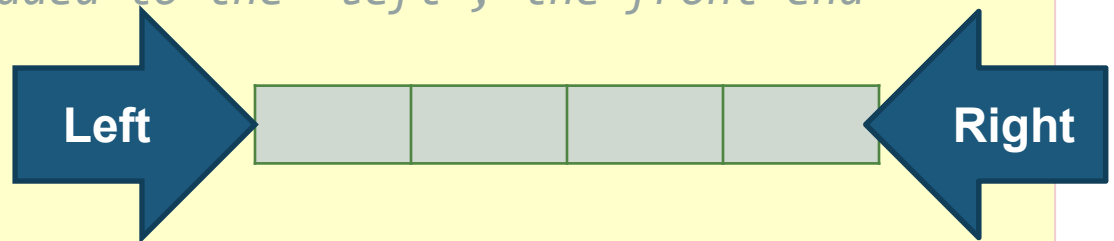
#Items can be removed from both ends

```
dq.pop() #removed from the right end
```

```
dq.popleft() #removed from the left end
```

```
print(dq)
```

Essentially a container that can be accessed from both ends: Left and Right. So, can be used as stack or queue!



"Queue" Application

PALINDROME **C**CHECKING

Palindrome : Problem Description

- **Palindrome** is a string which reads the same either *left to right*, or *right to left*
 - ❑ **Palindromes:** “r a d a r” and “d e e d”
 - ❑ **Counter Example:** “d a t a”
- Many solutions! We chose to:
 - a. Use *stack* to reverse the input
 - b. Use *queue* to preserve the input
 - ❑ The two sequence should be the same for palindrome
 - ❑ Also demonstrate the LIFO and FIFO property

Palindrome: Implementation

```
def isPalindrome( input ):
```

```
    s = StackList()
```

```
    q = QueueLinkedList()
```

Any Stack / Queue is fine!

```
    for ch in input:
```

```
        s.push(ch)
```

```
        q.enqueue(ch)
```

Push the same character
into both queue and stack

```
    while not s.isEmpty():
```

```
        if s.getTop() != q.getFront():
```

```
            return False
```

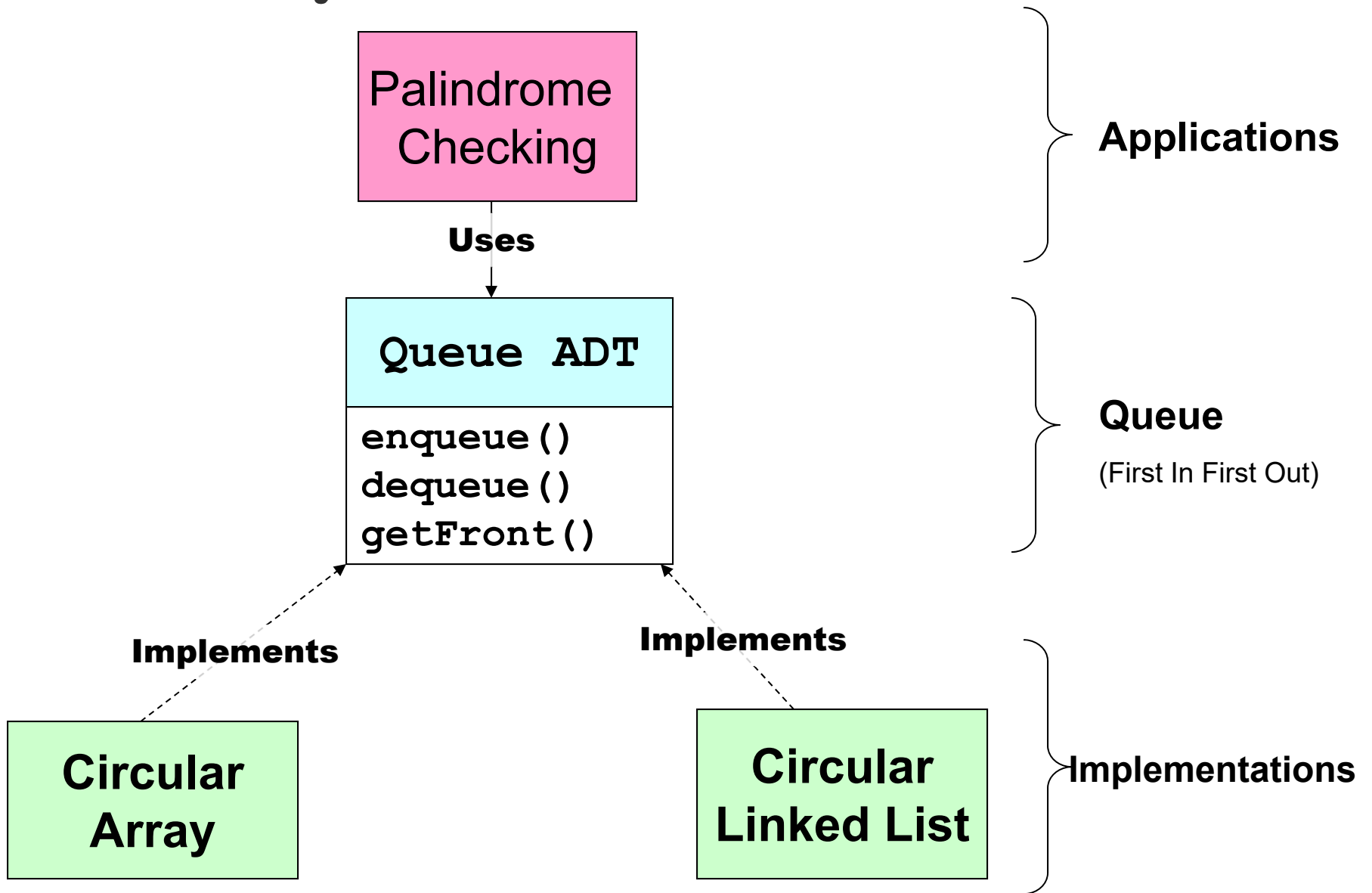
```
        s.pop()
```

```
        q.dequeue()
```

Queue has the original
sequence, Stack has the
reversed. Compare to
make sure they are the
same

```
    return True
```

Summary





END