
Stack ADT

**IT5003: Data Structures and Algorithms
(AY2019/20 Semester 1)**

Lecture Overview

■ Stack

- ❑ Introduction
- ❑ Specification
- ❑ Implementations
 - Python List (dynamic array)
 - Linked List
- ❑ Applications
 - Tower of Hanoi
 - Bracket Matching
 - Maze Exploration

Stack: A specialized list

- List ADT allows user to manipulate (insert/retrieve/remove) item at **any position within the sequence of items**
- There are cases where we only want to consider a few specific positions only
 - E.g. only the first/last position
 - Can be considered as special cases of list
- **Stack** is one such example:
 - Only manipulation at the **last position** is allowed
- **Queue** (coming next) is another example

Stack: Overview

- Real life example:
 - A stack of books, A stack of plates, Etc
- It is easier to add/remove item to/from the **top of the stack**
- The latest item added is the first item you can get out from the stack
 - Known as **Last In First Out (LIFO)** order
- Major Operations:
 - **Push** : Place item on top of the stack
 - **Pop** : Remove item from the top of the stack
- It is also common to provide:
 - **Top** : Take a look at the topmost item without removing it

Stack : Illustration

Top of stack
(accessible)

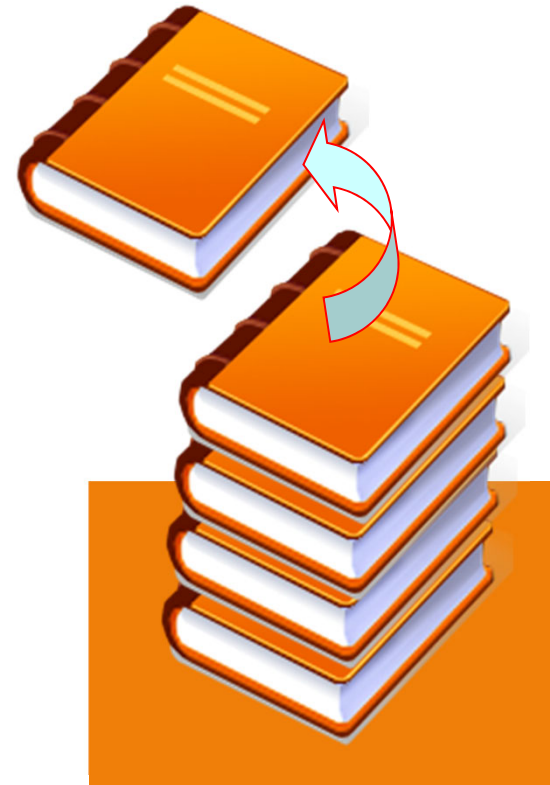


Bottom of
stack
(inaccessible)

A **stack** of
four books



Push a new
book on top



Pop a book
from top

Stack ADT: Python Specification

```
#imports not shown
class StackBase(ABC):
    @abstractmethod
    def getTop(self):
        pass

    @abstractmethod
    def push(self, newItem):
        pass

    @abstractmethod
    def pop(self):
        pass
```

```
@abstractmethod
def size(self):
    pass

@abstractmethod
def isEmpty(self):
    pass
```

Question:

It is not meaningful to provide a *toString()* method in stack, why?

Stack ADT: Sample User Program

- Generate **Fibonacci Sequence** { 1, 1, 2, 3, 5, 8, ... }
 - Basic idea: $Fib(N) = Fib(N-1) + Fib(N-2)$

```
def main():
```

```
    mystack =
```

To be replaced by actual implementations
of the `StackBase` class

```
    mystack.push( 1 )
```

```
    mystack.push( 1 )
```

We start off with a stack with 2 items,
"1" and "1" on top

```
    for i in range(10):
```

```
        prev1 = mystack.getTop( )
```

```
        mystack.pop( )
```

```
        prev2 = mystack.getTop( )
```

```
        cur = prev1 + prev2;
```

```
        mystack.push( prev1 );
```

```
        mystack.push( cur );
```

```
        print(cur)
```

Stack ADT : Implementations

- Many ways to implement Stack ADT, we will cover:
 - ❑ **Python List**
 - Make use of the built-in Python List
 - Essentially a **dynamic array**
 - ❑ **Linked List implementation:**
 - Study the best way to make use of linked list
- Learn how to weight the **pros** and **cons** for each implementations

Stack ADT : Design Consideration

- How to choose appropriate implementation?
 - ❑ Concentrate on the major operations in ADT
 - ❑ Match with data structures you have learned
 - Pick one to be the internal data structure of an ADT
 - Can the internal data structure support what you need?
 - Is the internal data structure efficient in those operations?
- Internal data structure like array, linked list etc are usually very flexible:
 - ❑ Make sure you use them in the best possible way

Python List is quite versatile

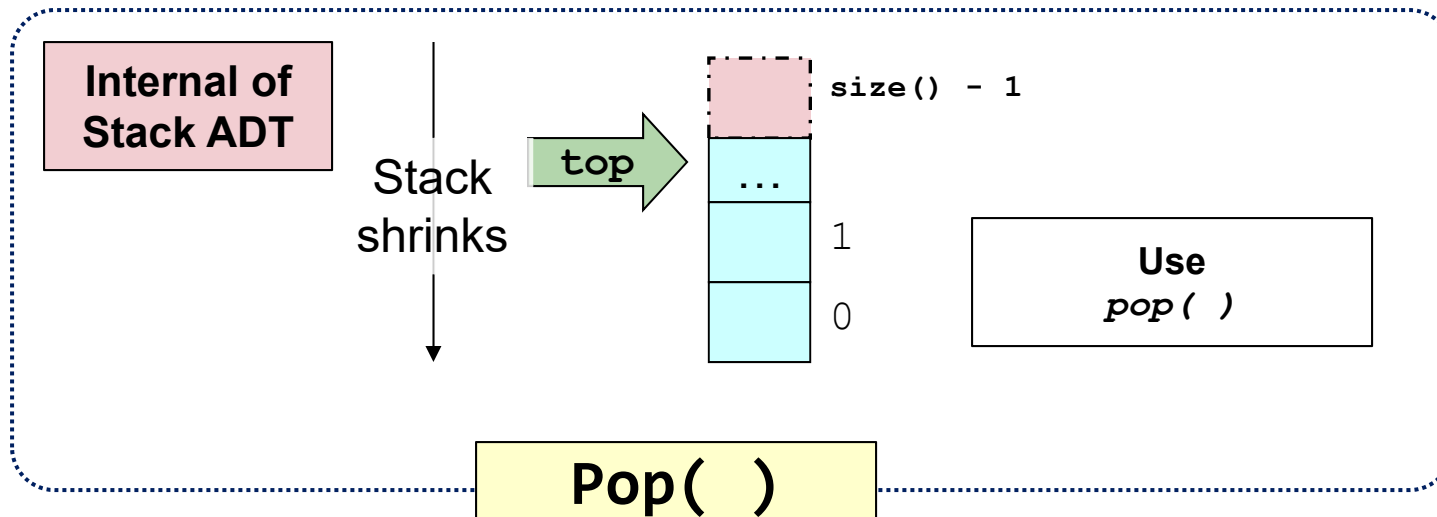
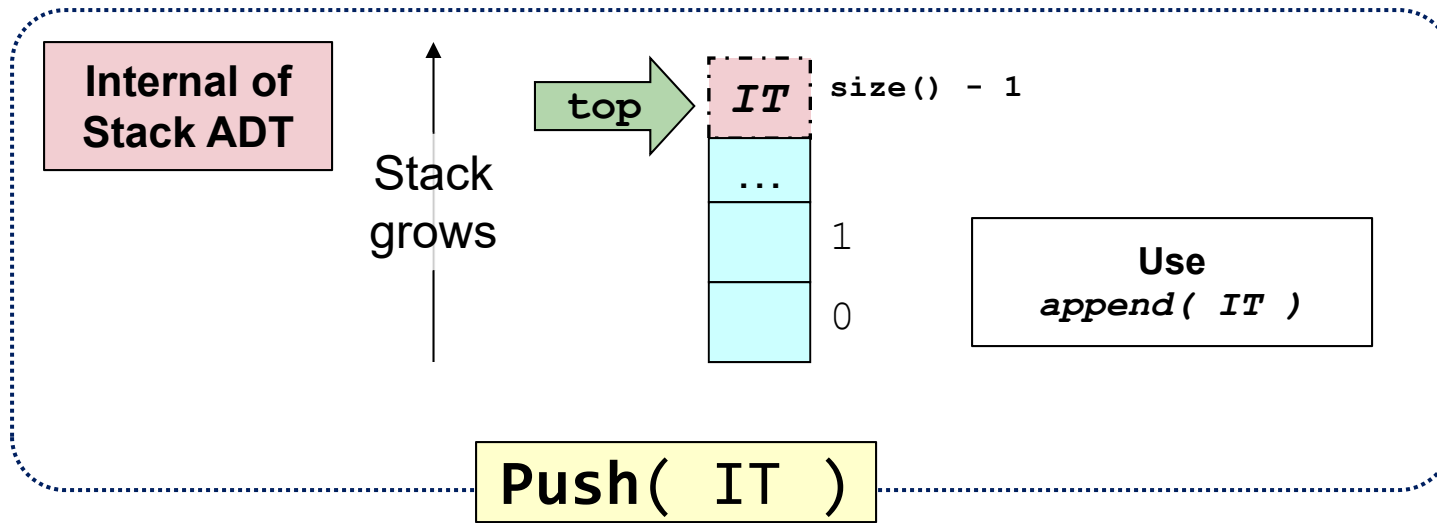
STACK ADT

USING PYTHON LIST

Stack ADT: Using **P**ython **L**ist

- Python List has the following capabilities:
 - ❑ Add **the last item**
 - *append()* [Very efficient]
 - ❑ Remove item from **any location**
 - Many ways: *pop(index)*, *remove(item)*
 - Only *pop()* last item is efficient (why?)
- What Stack ADT needs:
 - ❑ Add/Remove from **top of stack**
 - **No manipulation of other locations**
 - ❑ Hence, to make the best use of Python List:
 - Use the **back of list** as the **top of stack**

Using **P**ython **L**ist as Stack: Illustration



StackList: Implementation

```
class StackList(StackBase):

    def __init__(self):
        self._items = []

    def getTop(self):
        if not self.isEmpty():
            return self._items[-1]
            #negative index = count from the end of list
        else:
            return None

    def push(self, newItem):
        self._items.append(newItem)
```

StackList: Implementation

```
def pop(self):
    if not self.isEmpty():
        self._items.pop() #last item is removed
        return True
    else:
        return False

def size(self):
    return len(self._items)

def isEmpty(self):
    return len(self._items) == 0
```

- Very straightforward as **Python List** provides all the required operation with good time complexity

Python List: Comments

- Python List is designed with versatility in mind:
 - Especially to "double up" as simple **stack** / **queue**
- We provide the ***Fibonacci*** program written directly with Python List (but used as a "stack") to highlight the differences
- Contrast with Stack ADT to understand how ADT wraps around and "protect" the property of stack

Fibonacci: Using Python List

```
def main():
    mystack = []
    #Initialize the first two fibonacci numbers
    mystack.append( 1 )
    mystack.append( 1 )

    for i in range(10):
        prev1 = mystack[-1]
        mystack.pop()
        prev2 = mystack[-1]

        cur = prev1 + prev2
        mystack.append( prev1 )
        mystack.append( cur )
        print(cur)
```

Question:

What are the potential issues of using list directly as stack?

Not linked list again!

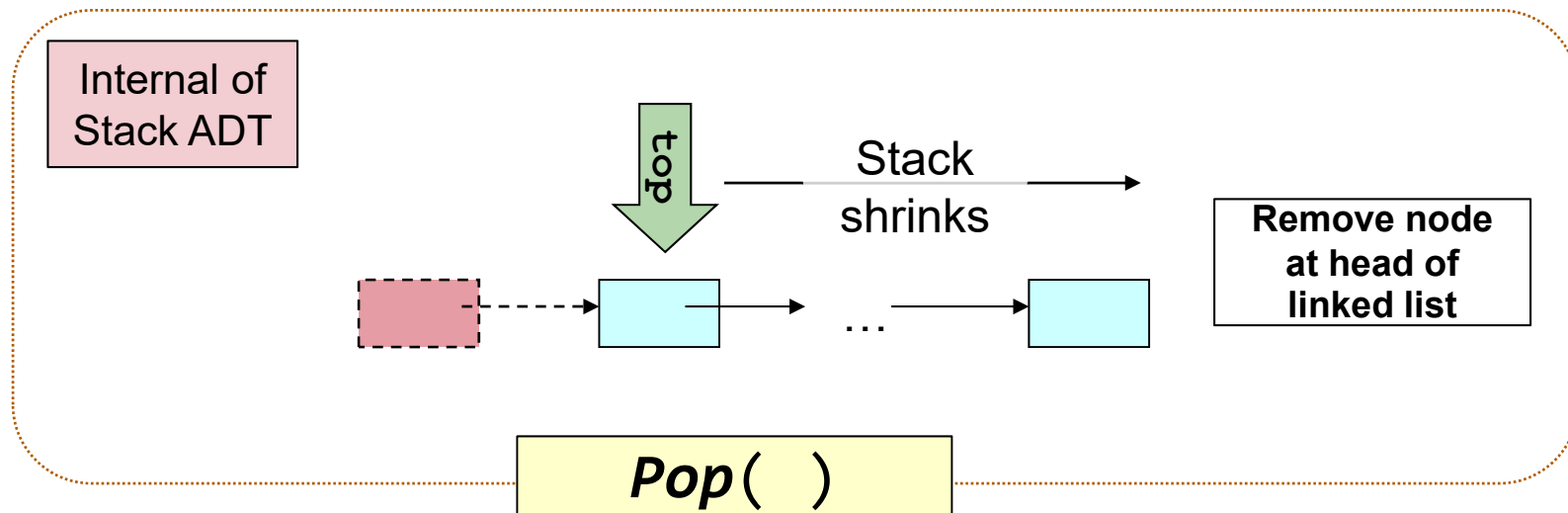
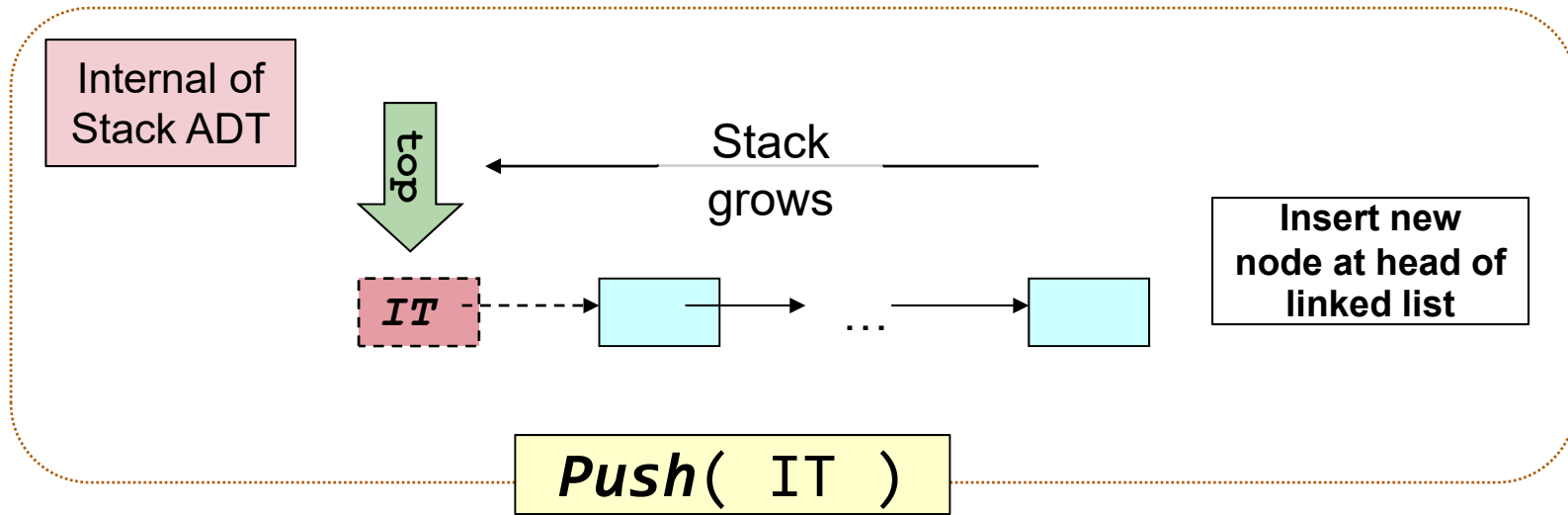
STACK ADT

USING LINKED LIST

Stack ADT: Using **Linked List**

- Characteristics of singly linked list:
 - Efficient manipulation of 1st Node:
 - Has a **head** pointer directly pointing to it
 - No need to traverse the list
 - Manipulation of other locations is possible:
 - Need to first traverse the list, less efficient
- Hence, best way to use singly linked list:
 - Use 1st Node as the top of stack
- Question:
 - How would you use **other variations of linked list?**

Stack ADT: Using Linked List (Illustration)



StackLinkedList: Implementation

```
class StackLinkedList(StackBase):  
  
    def __init__(self):  
        self._head = None  
        self._size = 0  
  
    def getTop(self):  
        if not self.isEmpty():  
            return self._head.item  
        else:  
            return None  
  
    def push(self, newItem):  
        newPtr = SinglyNode(newItem)  
  
        newPtr.next = self._head  
        self._head = newPtr  
  
        self._size += 1  
        return True
```

Reuse the **SinglyNode** linked list node from the List ADT

As we only insert at head position. General insertion code not needed.

StackLinkedList: Implementation

```
def pop(self):
    if not self.isEmpty():
        self._head = self._head.next
        self._size -= 1
        return True
    else:
        return False

def size(self):
    return self._size

def isEmpty(self):
    return self._size == 0
```

Similarly, only need to consider removal from head position in `pop()`

LIFO? Is it good for anything?

STACK APPLICATIONS

Stack Applications

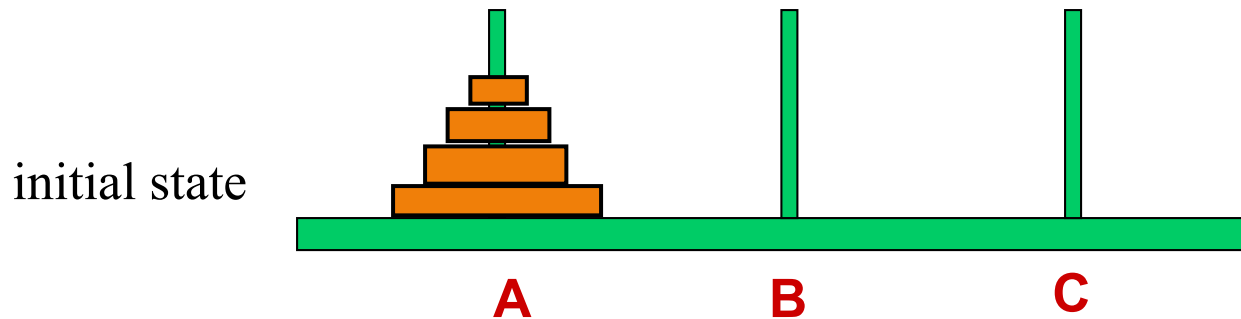
- Many useful applications for stack, we cover:
 - a. **Tower of Hanoi**
 - b. **Bracket Matching**
 - c. **Maze Exploration**

- More “Computer Science” inclined examples:
 - ❑ **Base-N number conversion**
 - ❑ **Postfix evaluation**
 - ❑ **Infix to postfix conversion**

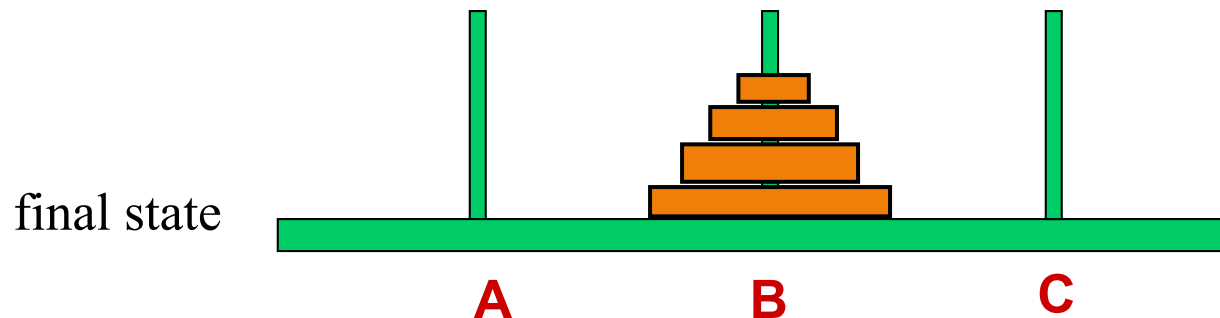
Stack Application Example One

TOWER OF HANOI

Tower of Hanoi : **D**escription



- How do we move all the disks from pole “**A**” to pole “**B**”, using pole “**C**” as temporary storage
 - ❑ One disk at a time
 - ❑ Disk must rest on top of larger disk



Each pole is a stack...

- We are not writing a program to ***solve the puzzle automatically***
 - ❑ Coming soon.... 😊
- Just a simple program to let user **play the puzzle:**
 - ❑ Keep track of the discs
 - ❑ Check movement
 - ❑ Display the current state
 - ❑ etc
- Since we can only
 - ❑ Remove the topmost disc from a pole, then
 - ❑ Place the disc on top of other pole
- Clearly, **each pole is a stack**

Tower of Hanoi : **I**mplementation

- Use a **List of 3 Stacks** to represent the 3 poles

```
class TowerOfHanoi():  
  
    def __init__(self, nDiscs):  
        self._poles = [StackList() for i in range(3)]  
        self._nDiscs = nDiscs  
  
        for i in range(nDiscs,0,-1):  
            self._poles[0].push(i)
```

Initialize the first pole to have a number of discs. The number represent the disc diameter.

Tower of Hanoi : Implementation

```
def move(self, src, dst):
    if src == dst:
        return True #pretend we have moved

    if src < 0 or src > 2 or dst < 0 or dst > 2:
        return False #illegal pole number

    if self._poles[src].isEmpty():
        return False #no disc at the source pole

    srcDisc = self._poles[src].getTop()

    if (not self._poles[dst].isEmpty()) and \
        (srcDisc > self._poles[dst].getTop()) :
        return False #destination pole has a larger disc

    #all checks passed, we can move safely
    self._poles[dst].push( srcDisc )
    self._poles[src].pop()
```

A simple function to check and perform a disc movement

Tower of Hanoi : Implementation

```
def _displayPole(self, pole):
    copy = StackList()
    while not pole.isEmpty():
        disc = pole.getTop()
        print(disc)
        pole.pop()
        copy.push(disc)

    while not copy.isEmpty():
        pole.push(copy.getTop())
        copy.pop()
```

A straightforward implementation
to show the discs on each pole.
Functional but not pretty 😊

```
def display(self):
    for i in range(3):
        self._displayPole(self._poles[i])
        print("Pole " + chr(ord('A')+i))
        print("-----")
```

Try to write a better looking
display function?

Tower of Hanoi : Main Function

```
def main():
    toh = TowerOfHanoi(3)

    toh.display()

    srcPole = int(input("src pole [0,1,2; -1 dst exit:"))
    dstPole = int(input("dst pole [0,1,2; -1 dst exit:"))
    while srcPole >= 0 and dstPole >= 0:
        if toh.move(srcPole, dstPole):
            print("Move Ok!")
        else:
            print("Illegal Move!")

    toh.display()
    srcPole = int(input("src pole [0,1,2; -1 dst exit:"))
    dstPole = int(input("dst pole [0,1,2; -1 dst exit:"))
```

A simple program to allow a human player to play the puzzle

Stack Application Example Two

BRACKET MATCHING



Bracket Matching : Description

- Mathematical expression can get quite convoluted:
 - E.g. $\{ [x+2(i-4!)]^e+4\pi/5*(\phi-7.28) \dots \}$
- We are interested in checking whether all brackets are matched correctly (with), [with] and { with }
- Bracket matching is equally useful for checking programming code

Bracket Matching : Pseudo-Code

1. Go through the input string character by character
 - Non-bracket characters
 - Ignore
 - Open bracket { , [or (
 - Push into stack
 - Close bracket },] or)
 - Pop from stack and check
 - If the stack top bracket does not agree with the closing bracket, complain and exit
 - Else continue
2. If the stack is not empty after we read through the whole string
 - The input is wrong also

Bracket Matching : Implementation (1)

```
def matchBracket( input ):  
    openBrackets = ['(', '[', '{']  
    closeBrackets = [')', ']', '}']  
    bracketStack = StackList()  
  
    for cur in input:  
        if cur in openBrackets: #cur is an open bracket  
            idx = openBrackets.index(cur)  
            bracketStack.push(closeBrackets[idx])  
  
        elif cur in closeBrackets:  
            if bracketStack.isEmpty() or \  
                bracketStack.getTop() != cur:  
                return False  
            else:  
                bracketStack.pop()  
  
    return bracketStack.isEmpty()
```

Why do we push the opposite bracket into the stack?

What is the meaning of these conditions?

Shouldn't we just return **True**?

Bracket Matching : Implementation (2)

```
def main():  
    userInput = input("Enter a math expression:")  
  
    if matchBracket(userInput):  
        print("Bracket matched!")  
    else:  
        print("Bracket NOT matched!")
```

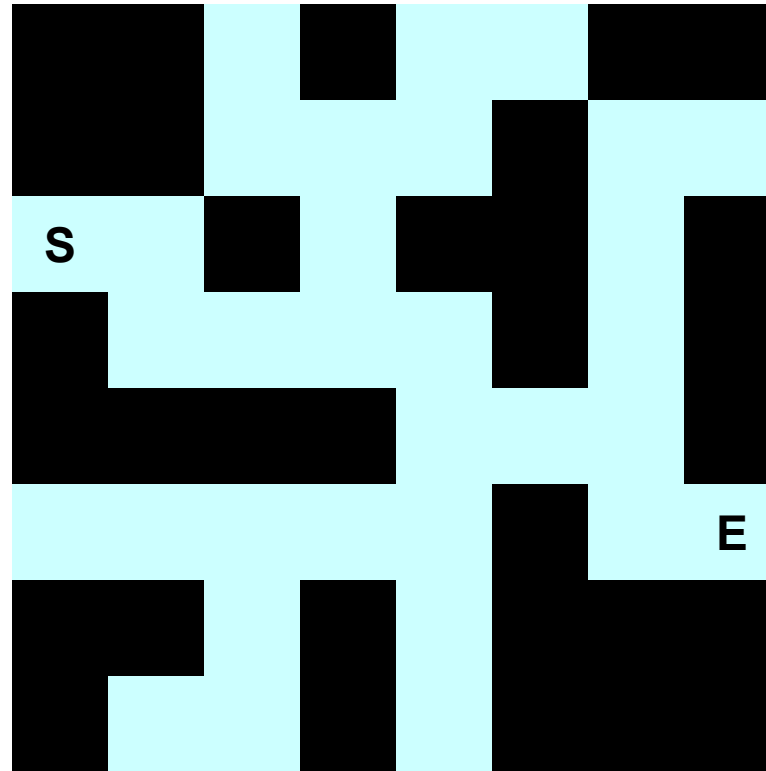
- A simple main program to test the matchBracket() function

Stack Application Example Three

MAZE EXPLORATION



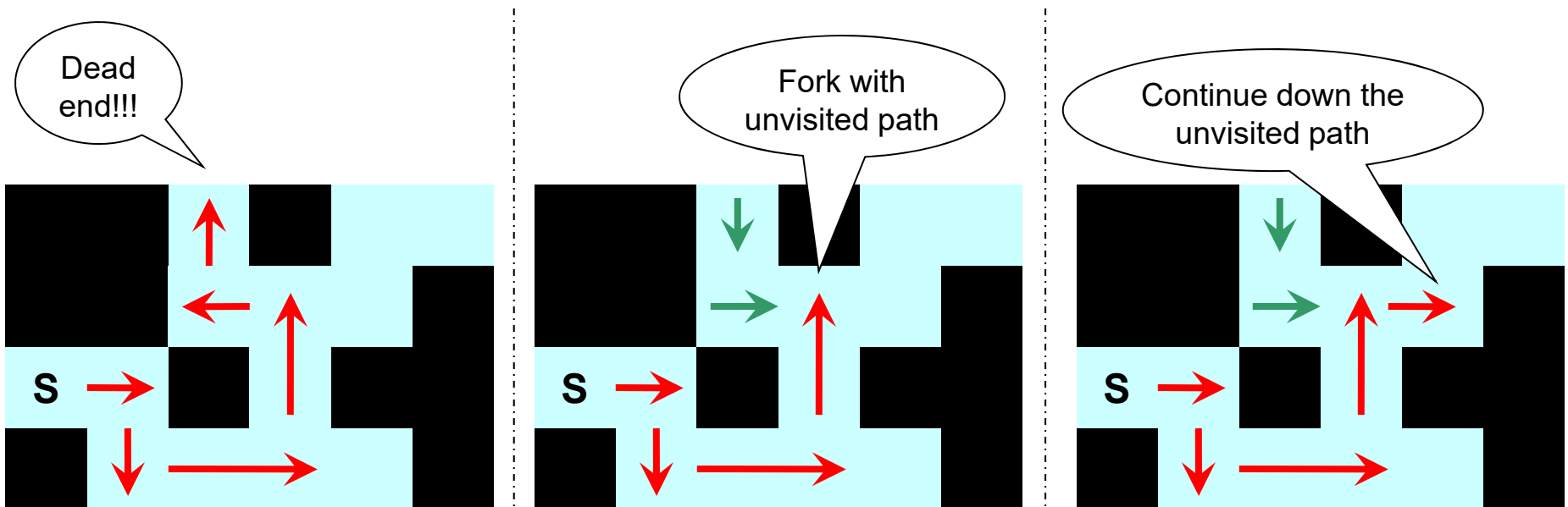
Exploring Maze: **D**escription



- How to define an algorithm that **always** get you from **S** to **E** (as long as there is a path)?
 - What should you do when you reached a dead end?

Exploring Maze: **B**asic **I**deas

- When we reached a dead end
 - ❑ Always restart from **S** is usually not a good idea
- Instead, we retrace our steps until:
 - ❑ the most recent fork which has an unvisited path
 - ❑ take the unvisited path and continue exploration



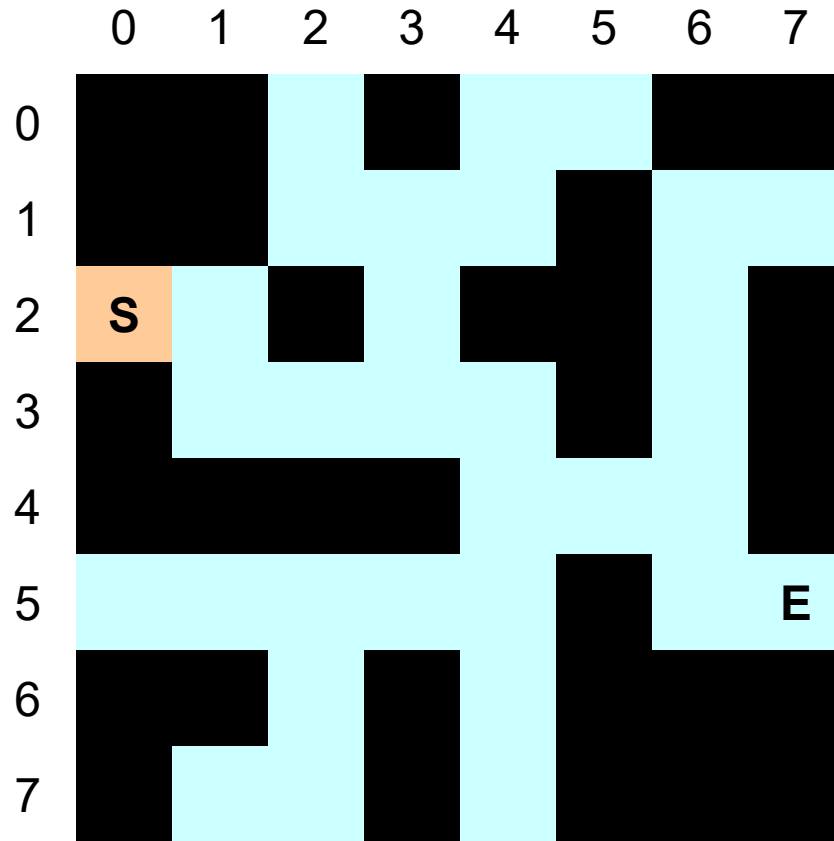
Exploring Maze: Some design issues

- The maze is represented as a **N x N 2D array**
 - ❑ Each square has a unique (`row`, `column`) coordinate
- There are 4 directions of movement from a square:
 - ❑ {`Up`, `Left`, `Down`, `Right`}
- Each square will know about:
 - ❑ Which direction is ***unvisited***
 - ❑ Assume a method `getUnvisitedDir()` is implemented for this purpose
- When a square has multiple unvisited exits:
 - ❑ We visit them in the order of the description above
- The path traveled is kept as a **stack of coordinates**

Exploring Maze: **P**seudo **C**ode

1. **Path** = empty
2. **done** = false //are we are the end yet?
3. **Path.push**(coordinate of **S**)
4. While (**Path** is not empty && not **done**)
 - i. **CurSq** = **Path.top**() //where are we now?
 - ii. **NewDir** = **CurSq.getUnvisitedDir**()
 - iii. If (**NewDir** == None) //dead end!
 Path.pop() //move back one square
 - iv. Else //there is an exit
 - a) **NewSq** = **CurSq.move**(**NewDir**)
 - b) **Path.push**(coordinate of **NewSq**)
 - c) If (**NewSq** == **E**) //Yes! We reached the end!
 done = true

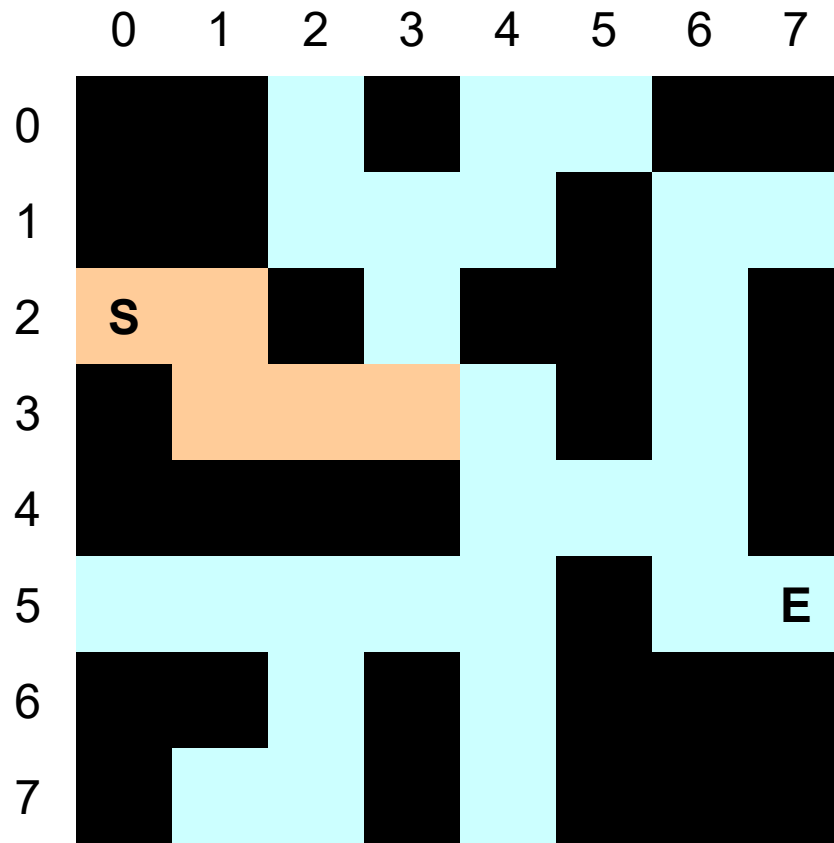
Exploring Maze: Test Run (1)



(2, 0) Path

- Just started at (2, 0)

Exploring Maze: Test Run(2)

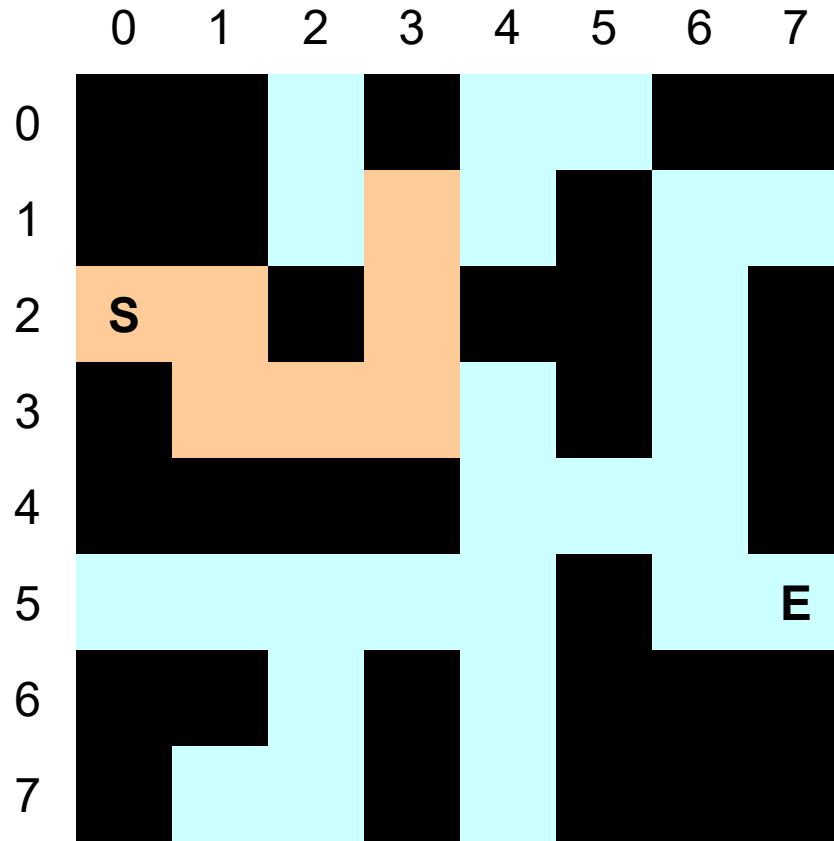


(3, 3)
(3, 2)
(3, 1)
(2, 1)
(2, 0)

Path

- (3, 3) is the first square with multiple exits
 - As stated, we will first try to go Up

Exploring Maze: Test Run(3)

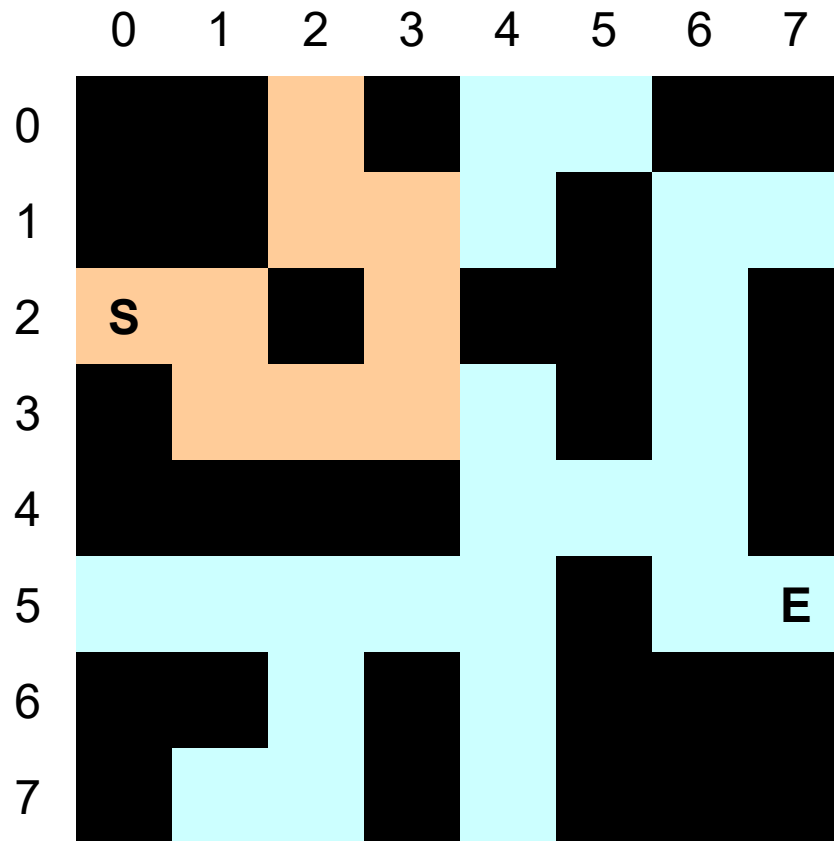


(1, 3)
(2, 3)
(3, 3)
(3, 2)
(3, 1)
(2, 1)
(2, 0)

Path

- Multiple exits at (1, 3)
 - go Up is impossible, so go Left is the 2nd choice

Exploring Maze: Test Run(4)

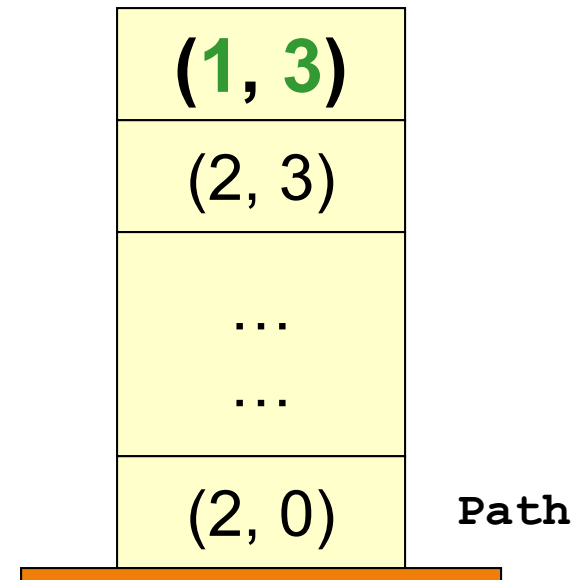
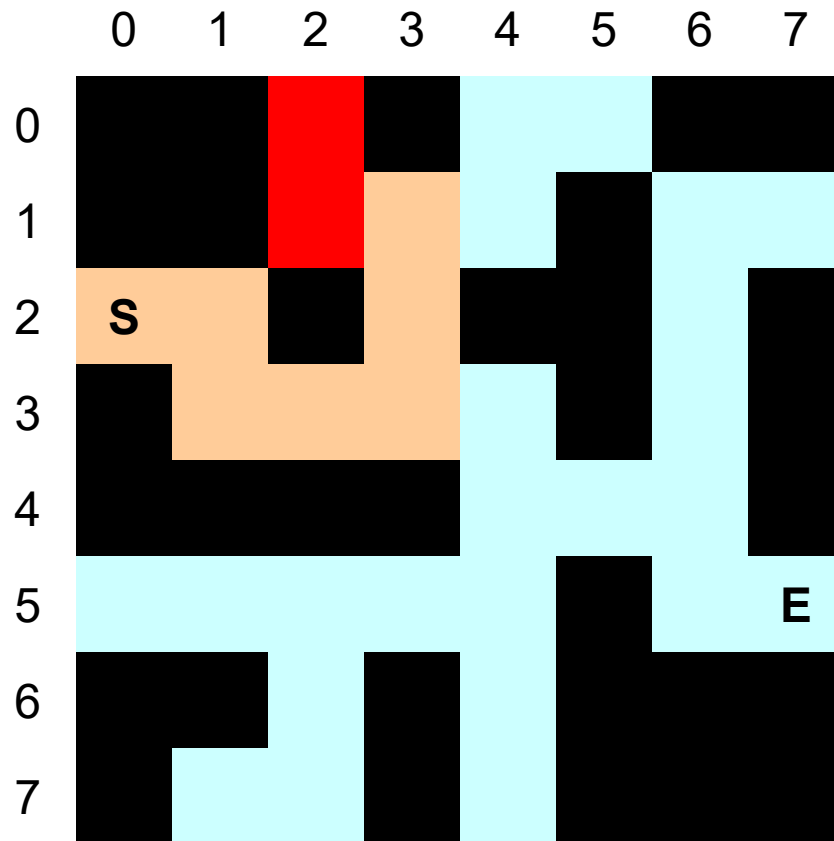


(0, 2)
(1, 2)
(1, 3)
(2, 3)
...
...
(2, 0)

Path

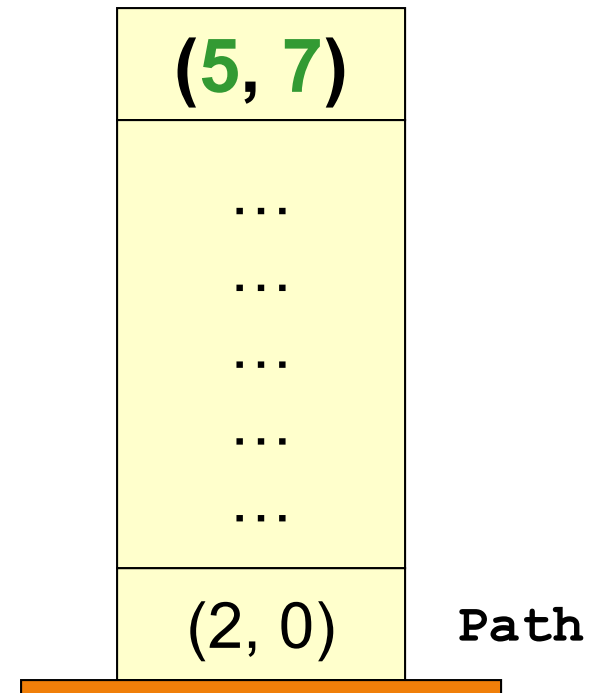
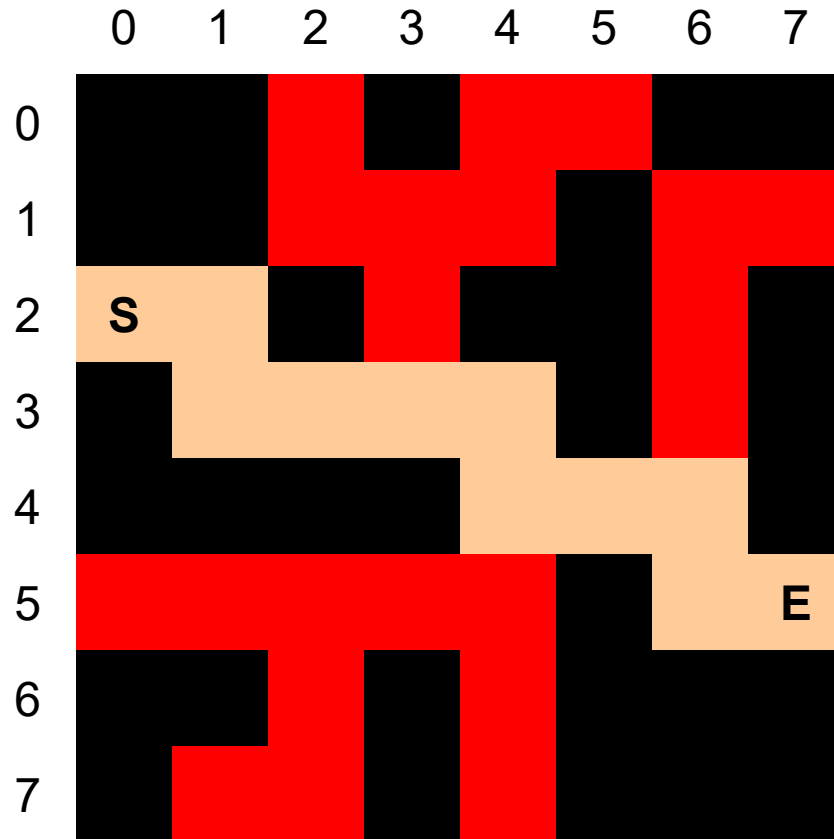
- No exits from (0, 2)
 - Back trace: pop until a square with unvisited exits

Exploring Maze: Test Run(5)



- Back to (1,3) after several pops
 - Up, Left, Down all impossible, going Right to (1,4)

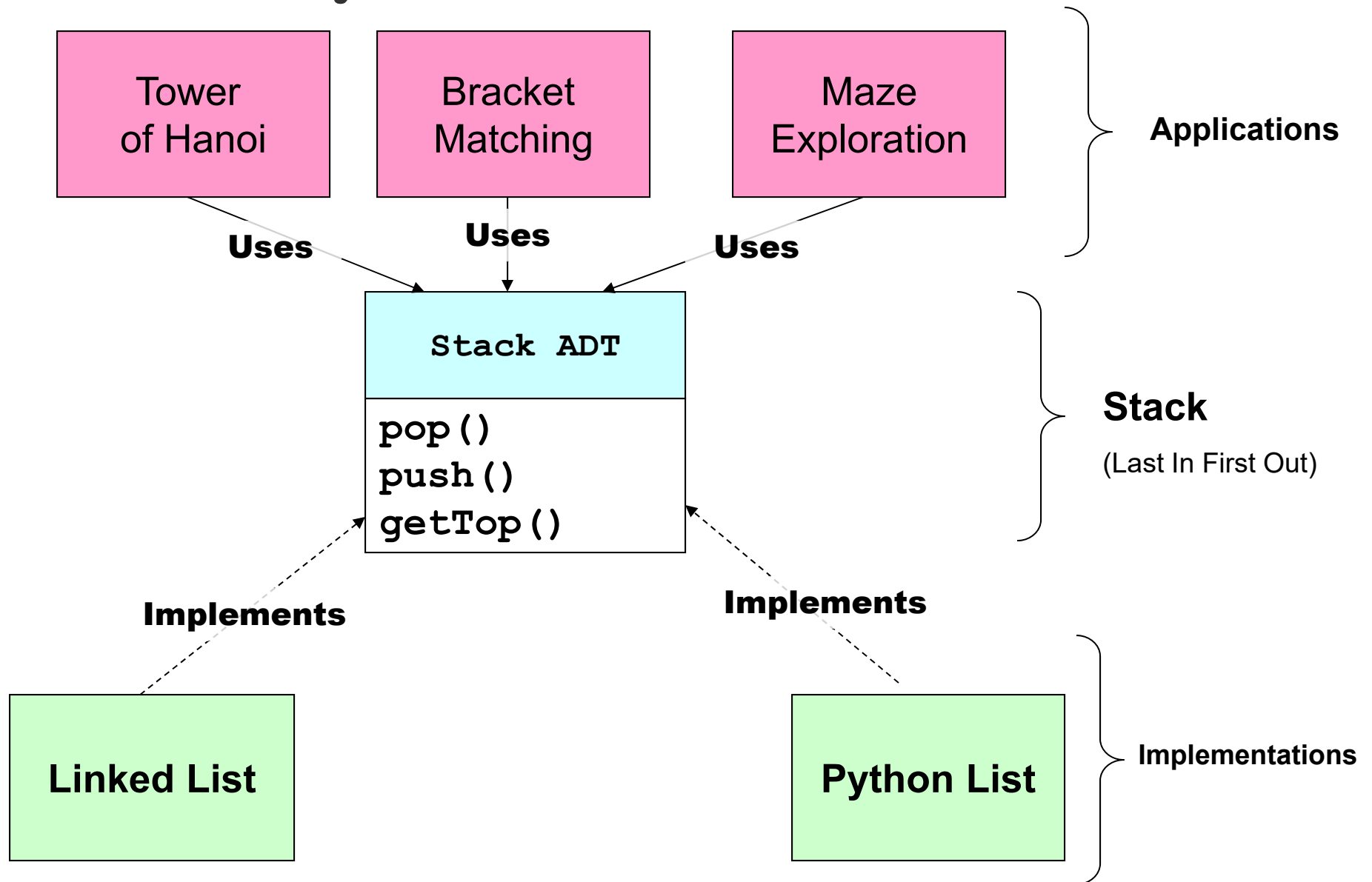
Exploring Maze: Test Run (much later)



■ Poor Guy/Girl ☹

- ❑ Traveled the whole maze to find the exit....
- ❑ Can we refine the algorithm that **always perform better?**

Summary





END