# Rock Lee do Pagode Namora D+

# University of Brasilia

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```
set ts=4 sw=4 sta nu rnu sc stl+=%F cindent
set bg=dark ruler clipboard=unnamed,unnamedplus
  timeoutlen=100
imap {<CR> {<CR>}}<Esc>0
nmap <F2> 0V$%d
nmap <C-down> :m+1<CR>
nmap <C-up> :m-2<CR>
vmap < C-c > "+y
nmap <C-a> ggVG
syntax on
alias cmp='g++ -Wall -Wformat=2 -Wshadow -Wconversion -
  fsanitize=address -fsanitize=undefined -fno-sanitize-
  recover -std=c++14'
```

# **Data Structures**

# **Merge Sort Tree**

```
struct MergeTree{
   int n:
   vector<vector<int>>> st;
   void build(int p, int L, int R, const int v[]){
          st[p].push_back(v[L]);
          return:
      int mid = (L+R)/2;
      build(2*p, L, mid, v);
      build(2*p+1, mid+1, R, v);
       st[p].resize(R-L+1);
      merge(st[2*p].begin(), st[2*p].end(),
              st[2*p+1].begin(), st[2*p+1].end(),
              st[p].begin());
   }
   int query(int p, int L, int R, int i, int j, int x)
     const{
       if(L > j \mid \mid R < i) return 0;
      if(L >= i && R <= j){
          int id = lower_bound(st[p].begin(), st[p].end
             (), x) - st[p].begin();
          return int(st[p].size()) - id;
      int mid = (L+R)/2;
      return query(2*p, L, mid, i, j, x) +
          query(2*p+1, mid+1, R, i, j, x);
   }
public:
   MergeTree(int sz, const int v[]): n(sz), st(4*sz){
      build(1, 1, n, v);
   //number of elements >= x on segment [i, j]
   int query(int i, int j, int x) const{
      if(i > j) swap(i, j);
      return query(1, 1, n, i, j, x);
   }
};
```

## **Wavelet Tree**

```
template<typename T>
class wavelet{
   T L, R;
   vector<int> 1;
   vector<T> sum; // <<</pre>
```

```
wavelet *lef, *rig;
   int r(int i) const{ return i - l[i]; }
public:
   template<typename ITER>
   wavelet(ITER bg, ITER en){
       lef = rig = nullptr;
       L = *bg, R = *bg;
       for(auto it = bg; it != en; it++)
          L = min(L, *it), R = max(R, *it);
       if(L == R) return;
       T mid = L + (R - L)/2;
       1.reserve(std::distance(bg, en) + 1);
       sum.reserve(std::distance(bg, en) + 1);
       1.push_back(0), sum.push_back(0);
       for(auto it = bg; it != en; it++)
          1.push_back(1.back() + (*it <= mid)),</pre>
          sum.push_back(sum.back() + *it);
       auto tmp = stable_partition(bg, en, [mid](T x){
          return x <= mid;</pre>
       });
       if(bg != tmp) lef = new wavelet(bg, tmp);
       if(tmp != en) rig = new wavelet(tmp, en);
    ~wavelet(){
       delete lef;
       delete rig;
    // 1 index, first is 1st
   T kth(int i, int j, int k) const{
       if(L >= R) return L;
       int c = 1[j] - 1[i-1];
       if(c \ge k) return lef \ge kth(l[i-1]+1, l[j], k);
       else return rig->kth(r(i-1)+1, r(j), k - c);
   }
   // # elements > x on [i, j]
   int cnt(int i, int j, T x) const{
       if(L > x) return j - i + 1;
       if(R <= x || L == R) return 0;
       int ans = 0;
       if(lef) ans += lef->cnt(l[i-1]+1, l[j], x);
       if(rig) ans += rig->cnt(r(i-1)+1, r(j), x);
       return ans;
   // sum of elements <= k on [i, j]</pre>
   T sumk(int i, int j, T k){
       if(L == R) return R <= k ? L * (j - i + 1) : 0;
       if(R <= k) return sum[j] - sum[i-1];</pre>
       int ans = 0;
       if(lef) ans += lef->sumk(l[i-1]+1, l[j], k);
       if(rig) ans += rig->sumk(r(i-1)+1, r(j), k);
      return ans;
   // swap (i, i+1) just need to update "array" l[i]
}:
Order Set
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/detail/standard_policies.hpp>
```

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```
using namespace __gnu_pbds; // or pb_ds;
                                                                   int mid = (L+R)/2;
                                                                   if(cht[mid].eval(x) >= cht[mid+1].eval(x)) // <<<</pre>
template<typename T, typename B = null_type>
                                                                       L = mid + 1:
using oset = tree<T, B, less<T>, rb_tree_tag,
                                                                   else bans = mid, R = mid - 1;
  tree_order_statistics_node_update>;
// find_by_order / order_of_key
                                                               return cht[bans].eval(x);
Hash table
                                                            Convex Hull Trick
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
                                                            * Author: Simon Lindholm
                                                             * source: https://github.com/kth-competitive-
struct custom_hash {
                                                               programming/kactl/blob/master/content/data-structures
   static uint64_t splitmix64(uint64_t x) {
       // http://xorshift.di.unimi.it/splitmix64.c
                                                               /LineContainer.h
      x += 0x9e3779b97f4a7c15;
                                                             * License: CC0
                                                            */
      x = (x \hat{ } (x >> 30)) * 0xbf58476d1ce4e5b9;
       x = (x \hat{ } (x >> 27)) * 0x94d049bb133111eb;
      return x \hat{ } (x \gg 31);
                                                            struct Line {
                                                               mutable 11 m, b, p;
   }
                                                               bool operator<(const Line& o) const { return m < o.m</pre>
   size_t operator()(uint64_t x) const {
      static const uint64_t FIXED_RANDOM = chrono::
                                                               bool operator<(11 x) const { return p < x; }</pre>
         steady_clock::now().time_since_epoch().count();
                                                            };
      return splitmix64(x + FIXED_RANDOM);
   }
                                                            struct LineContainer : multiset<Line, less<>> { // CPP14
};
                                                               // (for doubles, use inf = 1/.0, div(a,b) = a/b)
gp_hash_table<long long, int, custom_hash> table;
                                                               const 11 inf = LLONG_MAX;
unordered_map<long long, int, custom_hash> uhash;
                                                               11 div(ll a, ll b) { // floored division
                                                                   return a / b - ((a ^ b) < 0 && a % b); }
uhash.reserve(1 << 15);</pre>
uhash.max_load_factor(0.25);
                                                               bool isect(iterator x, iterator y) {
                                                                   if (y == end()) { x->p = inf; return false; }
Convex Hull Trick Simple
                                                                   if (x->m == y->m) x->p = x->b > y->b? inf : -inf
struct Line{
                                                                   else x->p = div(y->b - x->b, x->m - y->m);
                                                                   return x->p >= y->p;
   inline 11 eval(11 x) const{
      return x * m + b;
                                                               void add(ll m, ll b) {
   }
                                                                   auto z = insert(\{m, b, 0\}), y = z++, x = y;
};
                                                                   while (isect(y, z)) z = erase(z);
                                                                   if (x != begin() \&\& isect(--x, y)) isect(x, y =
// min => cht.back().m >= L.m
                                                                     erase(y));
// max => cht.back().m <= L.m
                                                                   while ((y = x) != begin() && (--x)->p >= y->p)
void push_line(vector<Line> &cht, Line L){
                                                                      isect(x, erase(y));
 while((int)cht.size() >= 2){
                                                               }
   int sz = (int)cht.size();
                                                               11 query(ll x) {
   if((long double)(L.b-cht[sz-1].b)*(cht[sz-2].m-L.m)
                                                                   assert(!empty());
  <= (long double)(L.b-cht[sz-2].b)*(cht[sz-1].m-L.m)){</pre>
                                                                   auto 1 = *lower_bound(x);
     cht.pop_back();
                                                                   return 1.m * x + 1.b;
   }
   else break;
                                                            };
 3
                                                            Min queue
 cht.push_back(L);
                                                            template<typename T>
// x increasing; pos = 0 in first call
                                                            class minQ{
11 linear_search(const vector<Line> &cht,ll x,int &pos){
                                                               deque<tuple<T, int, int> > p;
   while(pos+1 < (int)cht.size()){</pre>
                                                               T delta;
/*>>*/ if(cht[pos].eval(x) >= cht[pos+1].eval(x)) pos++;
                                                               int sz;
      else break;
                                                            public:
                                                               minQ() : delta(0), sz(0) {}
                                                               inline int size() const{ return sz; }
   return cht[pos].eval(x);
}
                                                               inline void add(T x){ delta += x; }
                                                               inline void push(T x, int id){
11 binary_search(const vector<Line> &cht, ll x){
                                                                   x -= delta, sz++;
   int L = 0, R = (int)cht.size()-2;
                                                                   int t = 1;
   int bans = (int)cht.size()-1;
                                                                   while(p.size() > 0 \& get<0>(p.back()) >= x)
   while(L <= R){</pre>
                                                                       t += get<1>(p.back()), p.pop_back();
```

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```
p.emplace_back(x, t, id);
                                                               unlaze(u);
   }
                                                               if(!u) u = it;
                                                               else if(Y[it] > Y[u]) split_val(u, X[it], L[it], R[
   inline void pop(){
       get<1>(p.front())--, sz--;
                                                                  it]), u = it;
       if(!get<1>(p.front())) p.pop_front();
                                                                else insert(X[it] < X[u] ? L[u] : R[u], it);
                                                                calc(u):
   T getmin() const{ return get<0>(p.front())+delta; }
   int getid() const{ return get<2>(p.front()); }
                                                            void erase(int &u, num key){
                                                               unlaze(u);
                                                                if(!u) return;
Sparse Table
                                                               if(X[u] == key) u = merge(L[u], R[u]);
                                                               else erase(key < X[u] ? L[u] : R[u], key);</pre>
int fn(int i, int j){
                                                                calc(u);
   if(j == 0) return v[i];
   if(~dn[i][j]) return dn[i][j];
                                                            int create_node(num key){
   return dn[i][j] = min(fn(i, j-1), fn(i + (1 << (j-1)))
                                                               X[en] = key;
     ), i-1));
                                                               sz[en] = 1;
                                                               L[en] = R[en] = 0;
                                                               return en++;
int getmn(int 1, int r){ // [1, r]
   int 1z = 1g(r - 1 + 1);
                                                            int query(int u, int 1, int r){//0 index
   return min(fn(1, 1z), fn(r - (1 << 1z) + 1, 1z));
                                                               unlaze(u):
                                                                if(u! or r < 0 or 1 >= sz[u]) return
Treap
                                                                 identity_element;
                                                               if(1 \le 0 \text{ and } r \ge sz[u] - 1) \text{ return } subt\_data[u];
// source: https://github.com/victorsenam/caderno/blob/
                                                               int ans = query(L[u], 1, r);
  master/code/treap.cpp
                                                                if(1 \le sz[L[u]] and sz[L[u]] \le r)
                                                                   ans = max(ans, st[u]);
//const int N = ; typedef int num;
num X[N]; int en = 1, Y[N], sz[N], L[N], R[N];
                                                                ans = max(ans, query(R[u], l-sz[L[u]]-1, r-sz[L[u]])
void calc (int u) { // update node given children info
                                                                 11-1)):
   if(!u) return;
                                                               return ans;
   sz[u] = sz[L[u]] + 1 + sz[R[u]];
   // code here, no recursion
                                                            ColorUpdate
void unlaze (int u) {
                                                            // source: https://github.com/tfg50/Competitive-
   if(!u) return;
   // code here, no recursion
                                                              Programming/tree/master/Biblioteca/Data%20Structures
void split_val(int u, num x, int &l, int &r) { // l gets
                                                            #include <set>
                                                            #include <vector>
   <= x, r gets > x
   unlaze(u); if(!u) return (void) (1 = r = 0);
                                                            template <class Info = int>
   if(X[u] \le x) \{ split_val(R[u], x, 1, r); R[u] = 1;
                                                            class ColorUpdate {
     1 = u; }
   else { split_val(L[u], x, 1, r); L[u] = r; r = u; }
                                                            public:
   calc(u);
                                                                struct Range {
                                                                   Range(int 1 = 0) { this->1 = 1; }
void split_sz(int u, int s, int &l, int &r) { // l gets
                                                                   Range(int 1, int r, Info v) {
  first s, r gets remaining
                                                                      this->1 = 1:
   unlaze(u); if(!u) return (void) (1 = r = 0);
                                                                       this -> r = r;
   if(sz[L[u]] < s)  { split_sz(R[u], s - sz[L[u]] - 1,
                                                                      this->v = v;
     1, r); R[u] = 1; 1 = u; }
   else { split_sz(L[u], s, l, r); L[u] = r; r = u; }
                                                                   int 1, r;
   calc(u);
                                                                   Info v;
int merge(int 1, int r) { // els on 1 <= els on r</pre>
                                                                   bool operator < (const Range &b) const { return 1</pre>
   unlaze(l); unlaze(r); if(!l || !r) return l + r; int
                                                                      < b.1; }
                                                               };
   if(Y[1] > Y[r]) { R[1] = merge(R[1], r); u = 1; }
   else { L[r] = merge(1, L[r]); u = r; }
                                                                std::vector<Range> upd(int 1, int r, Info v) {
   calc(u); return u;
                                                                   std::vector<Range> ans;
                                                                   if(1 >= r) return ans;
void init(int n=N-1) { // XXX call before using other
                                                                   auto it = ranges.lower_bound(1);
                                                                   if(it != ranges.begin()) {
  funcs
   for(int i = en = 1; i \le n; i++) { Y[i] = i; sz[i] =
                                                                      it--;
      1; L[i] = R[i] = 0; }
                                                                       if(it->r>1) {
   random\_shuffle(Y + 1, Y + n + 1);
                                                                          auto cur = *it;
                                                                          ranges.erase(it);
void insert(int &u, int it){
                                                                          ranges.insert(Range(cur.1, 1, cur.v));
```

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```
ranges.insert(Range(1, cur.r, cur.v));
          }
                                                           void update(int p, T val){
      }
                                                               st[p += n] = val;
      it = ranges.lower_bound(r);
                                                               while(p >>= 1) st[p] = merge(st[2*p], st[2*p+1]);
      if(it != ranges.begin()) {
          it--;
                                                           Recursive Segtree + lazy
          if(it->r>r) {
              auto cur = *it;
                                                           class SegTree{
              ranges.erase(it);
                                                               vi st:
              ranges.insert(Range(cur.1, r, cur.v));
                                                               vi lazv:
              ranges.insert(Range(r, cur.r, cur.v));
                                                               int size;
                                                               int el_neutro = -oo;
       for(it = ranges.lower_bound(1); it != ranges.end
         () && it->l < r; it++) {
                                                               int f(int a, int b){
          ans.push_back(*it);
                                                                  return max(a,b);
      ranges.erase(ranges.lower\_bound(1), \ ranges.
                                                               void propagate(int sti, int stl, int str){
         lower_bound(r));
                                                                  if(lazy[sti]){
       ranges.insert(Range(l, r, v));
                                                                      st[sti] += lazy[sti];
       return ans:
                                                                      if(stl != str)
   }
private:
                                                                          lazy[sti*2 + 1] += lazy[sti];
   std::set<Range> ranges;
                                                                          lazy[sti*2 + 2] += lazy[sti];
Heavy Light Decomposition
                                                                      lazy[sti] = 0;
                                                                  }
void dfs_sz(int u){
                                                               }
   sz[u] = 1;
                                                               int query(int sti, int stl, int str, int l, int r){
   for(auto &v : g[u]) if(v == p[u]){
                                                                  propagate(sti, stl, str);
       swap(v, g[u].back()); g[u].pop_back();
      break;
                                                                  if(str < 1 \mid \mid r < stl)
                                                                      return el_neutro;
   for(auto &v : g[u]){
      p[v] = u; dfs_sz(v); sz[u] += sz[v];
      if(sz[v] > sz[g[u][0]])
                                                                  if(stl >= 1 and str <= r)
          swap(v, g[u][0]);
                                                                      return st[sti];
   }
                                                                  int mid = (str+st1)/2;
// nxt[u] = start of path with u
// set nxt[root] = root beforehand
                                                                  return f(query(sti*2+1,stl,mid,l,r),query(sti
void dfs_hld(int u){
                                                                     *2+2,mid+1,str,l,r));
   in[u] = t++;
   rin[in[u]] = u;
                                                               void update_range(int sti, int stl, int str, int l,
   for(auto v : g[u]){
                                                                 int r, int amm){
      nxt[v] = (v == g[u][0] ? nxt[u] : v); dfs_hld(v);
                                                                  propagate(sti, stl, str);
                                                                  if(stl >= 1 and str <= r){
   out[u] = t;
                                                                      lazy[sti] = amm;
                                                                      propagate(sti, stl, str);
// subtree of u => [ in[u], out[u] )
                                                                      return:
// path from nxt[u] to u => [ in[ nxt[u] ], in[u] ]
                                                                  if(stl > r or str < l)</pre>
Iterative Segtree
                                                                      return:
T query(int 1, int r){ // [1, r]
                                                                  int mid = (stl + str)/2;
   T rl, rr;
                                                                  update_range(sti*2+1,stl,mid,l,r,amm);
   for(1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1){
                                                                  update_range(sti*2+2,mid+1,str,l,r,amm);
      if(l & 1) rl = merge(rl, st[l++]);
                                                                  st[sti] = f(st[sti*2+1], st[sti*2+2]);
       if(r & 1) rr = merge(st[--r], rr);
                                                               void update(int sti, int stl, int str, int i, int
   return merge(rl, rr);
                                                                 amm){
                                                                  propagate(sti, stl, str);
                                                                  if(stl == i and str == i){
// initially save v[i] in st[n+i] for all i in [0, n)
                                                                      st[sti] = amm;
void build(){
                                                                      return:
   for(int p = n-1; p > 0; p--)
                                                                  if(stl > i or str < i)</pre>
       st[p] = merge(st[2*p], st[2*p+1]);
}
                                                                      return:
```

```
int mid = (stl + str)/2;
                                                           void add_letter(char c)
      update(sti*2+1,stl,mid,i,amm);
                                                           {
      update(sti*2+2,mid+1,str,i,amm);
                                                               s[n++] = c -= 'a';
       st[sti] = f(st[sti*2+1], st[sti*2+2]);
                                                               last = get_link(last);
                                                               if(!to[last][c])
   public:
                                                               {
       SegTree(int n): st(4*n,0), lazy(4*n,0){size = n;}
                                                                   len[sz] = len[last] + 2;
       int query(int 1, int r){return query(0,0,size-1,1
                                                                   link[sz] = to[get_link(link[last])][c];
                                                                   diff[sz] = len[sz] - len[link[sz]];
       void update_range(int 1, int r, int amm){
                                                                   if(diff[sz] == diff[link[sz]])
         update_range(0,0,size-1,1,r,amm);}
                                                                      slink[sz] = slink[link[sz]];
      void update(int i, int amm){update(0,0,size-1,i,
                                                                   else
         amm);}
                                                                      slink[sz] = link[sz];
};
                                                                   to[last][c] = sz++;
LiChao's Segtree
                                                               last = to[last][c];
                                                           }
void add_line(line nw, int v = 1, int l = 0, int r = 0
  maxn) { // [1, r)}
                                                           int main()
   int m = (1 + r) / 2;
                                                           {
   bool lef = nw.eval(1) < st[v].eval(1);</pre>
                                                               ios::sync_with_stdio(0);
   bool mid = nw.eval(m) < st[v].eval(m);</pre>
                                                               cin.tie(0);
   if(mid) swap(st[v], nw);
                                                               init();
   if(r - 1 == 1) {
                                                               string s;
      return;
                                                               cin >> s;
   } else if(lef != mid) {
                                                               int n = s.size();
       add_line(nw, 2 * v, 1, m);
                                                               int ans[n + 1];
   } else {
                                                               memset(ans, 63, sizeof(ans));
       add_line(nw, 2 * v + 1, m, r);
                                                               ans[0] = 0;
                                                               for(int i = 1; i <= n; i++)
}
                                                                   add_letter(s[i - 1]);
int get(int x, int v = 1, int l = 0, int r = maxn) {
                                                                   for(int v = last; len[v] > 0; v = slink[v])
   int m = (1 + r) / 2;
   if(r - 1 == 1) {
                                                                      series_ans[v] = ans[i - (len[slink[v]] + diff
      return st[v].eval(x);
                                                                        [v])];
   } else if(x < m) {
                                                                      if(diff[v] == diff[link[v]])
      return min(st[v].eval(x), get(x, 2*v, 1, m));
                                                                          series_ans[v] = min(series_ans[v],
                                                                            series_ans[link[v]]);
      return min(st[v].eval(x), get(x, 2*v+1, m, r));
                                                                      ans[i] = min(ans[i], series_ans[v] + 1);
                                                                  }
                                                                   cout << ans[i] << "\n";</pre>
Palindromic tree
                                                               }
                                                               return 0;
#include <bits/stdc++.h>
                                                           Math
using namespace std;
                                                           Extended Euclidean Algorithm
const int maxn = 3e5 + 1, sigma = 26;
int len[maxn], link[maxn], to[maxn][sigma];
                                                           // a*x + b*y = gcd(a, b), < gcd, x, y>
int slink[maxn], diff[maxn], series_ans[maxn];
                                                           tuple<int, int, int> gcd(int a, int b) {
int sz, last, n;
                                                               if(b == 0) return make_tuple(a, 1, 0);
char s[maxn];
                                                               int q, w, e;
                                                               tie(q, w, e) = gcd(b, a % b);
void init()
                                                               return make_tuple(q, e, w - e * (a / b));
   s[n++] = -1;
   link[0] = 1;
                                                           Chinese Remainder Theorem
   len[1] = -1;
   sz = 2;
                                                            // x = vet[i].first (mod vet[i].second)
}
                                                           11 crt(vector<pair<ll, ll>> vet){
                                                               11 ans = vet[0].first, lcm = vet[0].second;
int get_link(int v)
                                                               ll a, b, g, x, y;
                                                               for(int i = 1; i < (int)vet.size(); i++){</pre>
   while(s[n - len[v] - 2] != s[n - 1]) v = link[v];
                                                                   tie(a, b) = vet[i];
   return v;
                                                                   tie(g, x, y) = gcd(lcm, b);
                                                                   ans = ans + x * (a - ans) / g % (b / g) * lcm;
}
                                                                   lcm = lcm * b / g;
```

```
ans = (ans \% lcm + lcm) \% lcm;
                                                                          if(x == n-1) break;
   }
                                                                      if(x != n-1) return 0;
   return ans:
}
                                                                  }
Preffix inverse
                                                               return 1;
inv[1] = 1;
for(int i = 2; i < p; i++)
                                                           Totiente
   inv[i] = (p - (p/i) * inv[p%i] % p) % p;
Pollard Rho
                                                           11 totiente(ll n){
                                                               11 \text{ ans} = n;
11 rho(11 n){
                                                               for(ll i = 2; i*i <= n; i++){
   if(n % 2 == 0) return 2;
                                                                   if(n \% i == 0){
   11 d, c, x, y;
                                                                      ans = ans / i * (i - 1);
   do{
                                                                      while(n \% i == 0) n /= i;
       c = 11rand() % n, x = 11rand() % n, y = x;
       do{
                                                               }
          x = add(mul(x, x, n), c, n);
          y = add(mul(y, y, n), c, n);
                                                               if(n > 1) ans = ans / n * (n - 1);
          y = add(mul(y, y, n), c, n);
                                                               return ans;
          d = \_gcd(abs(x - y), n);
       }while(d == 1);
                                                           Primitive root
   }while(d == n);
   return d;
                                                           // a primitive root modulo n is any number g such that
11 pollard_rho(ll n){
                                                             any c coprime to n is congruent to a power of g modulo
   11 x, c, y, d, k;
   int i;
                                                           bool exists_root(ll n){
   do{
                                                               if(n == 1 || n == 2 || n == 4) return true;
       i = 1;
                                                               if(n \% 2 == 0) n /= 2;
       x = 11rand() % n, c = 11rand() % n;
                                                               if(n % 2 == 0) return false;
       y = x, k = 4;
                                                               // test if n is a power of only one prime
       do{
                                                               for(11 i = 3; i * i <= n; i += 2) if(n % i == 0){
          if(++i == k) y = x, k *= 2;
                                                                  while(n \% i == 0) n /= i;
          x = add(mul(x, x, n), c, n);
                                                                  return n == 1;
          d = \_gcd(abs(x - y), n);
       }while(d == 1);
                                                               return true;
   }while(d == n);
   return d:
                                                           ll primitive_root(ll n){
                                                               if(n == 1 || n == 2 || n == 4) return n - 1;
void factorize(ll val, map<ll, int> &fac){
                                                               if(not exists_root(n)) return -1;
   if(rabin(val)) fac[ val ]++;
                                                               11 x = phi(n);
   else{
                                                               auto pr = factorize(x);
       11 d = pollard_rho(val);
                                                               auto check = [x, n, pr](11 m){
       factorize(d, fac);
                                                                   for(ll p : pr) if(fexp(m, x / p, n) == 1)
       factorize(val / d, fac);
                                                                      return false;
   }
                                                                  return true;
}
map<ll, int> factor(ll val){
                                                               for(11 m = 2; ; m++) if(\_gcd(m, n) == 1)
   map<ll, int> fac;
                                                                   if(check(m)) return m;
   if(val > 1) factorize(val, fac);
                                                           }
   return fac;
                                                           // Let's denote R(n) as the set of primitive roots
                                                             modulo n, p is prime
Miller Rabin
                                                           // g \ln R(p) => (pow(g, p-1, p * p) == 1 ? g+p : g) \ln r
                                                              R(pow(p, k)), for all k > 1
bool rabin(ll n){
                                                           // g \text{ in } R(pow(p, k)) \Rightarrow (g \% 2 == 1 ? g : g + pow(p, k))
   if(n \ll 1) return 0;
                                                               if(n <= 3) return 1;
   11 s = 0, d = n - 1;
                                                           Mobius Function
   while(d % 2 == 0) d /= 2, s++;
   for(int k = 0; k < 64; k++){
                                                           memset(mu, 0, sizeof mu);
       11 a = (11rand() \% (n - 3)) + 2;
                                                           mu[1] = 1;
       11 x = fexp(a, d, n);
                                                           for(int i = 1; i < N; i++)
       if(x != 1 \&\& x != n-1){
                                                               for(int j = i + i; j < N; j += i)
          for(int r = 1; r < s; r++){
                                                                  mu[j] -= mu[i];
              x = mul(x, x, n);
                                                           // g(n) = sum{f(d)} => f(n) = sum{mu(d)*g(n/d)}
              if(x == 1) return 0;
```

# Mulmod TOP

```
constexpr uint64_t mod = (1ull<<61) - 1;</pre>
uint64_t modmul(uint64_t a, uint64_t b){
   uint64_t 11 = (uint32_t)a, h1 = a>>32, 12 = (
     uint32_t)b, h2 = b >> 32;
   uint64_t l = 11*12, m = 11*h2 + 12*h1, h = h1*h2;
   uint64_t ret = (1\&mod) + (1>>61) + (h << 3) + (m >>
     29) + (m << 35 >> 3) + 1;
   ret = (ret & mod) + (ret>>61);
   ret = (ret & mod) + (ret>>61);
   return ret-1:
Matrix Determinant
int n;
long double a[n][n];
long double gauss(){
   long double det = 1;
   for(int i = 0; i < n; i++){
       int q = i;
       for(int j = i+1; j < n; j++){
          if(abs(a[j][i]) > abs(a[q][i]))
             q = j;
      if(abs(a[q][i]) < EPS){
          det = 0;
          break;
      if(i != q){
          for(int w = 0; w < n; w++)
              swap(a[i][w], a[q][w]);
          det = -det;
      det *= a[i][i];
       for(int j = i+1; j < n; j++) a[i][j] /= a[i][i];</pre>
       for(int j = 0; j < n; j++) if(j != i){
          if(abs(a[j][i]) > EPS)
             for(int k = i+1; k < n; k++)
                 a[j][k] -= a[i][k] * a[j][i];
      }
   }
   return det;
Simplex Method
typedef long double dbl;
const dbl eps = 1e-6;
const int N = , M = ;
mt19937 rng(chrono::steady_clock::now().time_since_epoch
  ().count());
struct simplex {
   int X[N], Y[M];
   dbl A[M][N], b[M], c[N];
   dbl ans:
   int n, m;
   dbl sol[N];
   void pivot(int x, int y){
       swap(X[y], Y[x]);
      b[x] /= A[x][y];
       for(int i = 0; i < n; i++)
          if(i != y)
```

```
A[x][i] /= A[x][y];
       A[x][y] = 1. / A[x][y];
       for(int i = 0; i < m; i++)
          if(i != x \& abs(A[i][y]) > eps) {
              b[i] -= A[i][y] * b[x];
              for(int j = 0; j < n; j++) if(j != y)
                 A[i][j] -= A[i][y] * A[x][j];
              A[i][y] = -A[i][y] * A[x][y];
       ans += c[y] * b[x];
       for(int i = 0; i < n; i++)
          if(i != y)
              c[i] -= c[y] * A[x][i];
       c[y] = -c[y] * A[x][y];
   }
   // maximiza sum(x[i] * c[i])
   // sujeito a
   // sum(a[i][j] * x[j]) <= b[i] para 0 <= i < m (Ax)
   // x[i] >= 0 para 0 <= i < n (x >= 0)
   // (n variaveis, m restricoes)
   // guarda a resposta em ans e retorna o valor otimo
   dbl solve(int _n, int _m) {
       this->n = _n; this->m = _m;
       for(int i = 1; i < m; i++){</pre>
          int id = uniform_int_distribution<int>(0, i)(
            rng);
          swap(b[i], b[id]);
          for(int j = 0; j < n; j++)
              swap(A[i][j], A[id][j]);
       }
       ans = 0.;
       for(int i = 0; i < n; i++) X[i] = i;
       for(int i = 0; i < m; i++) Y[i] = i + n;
       while(true) {
          int x = min_element(b, b + m) - b;
          if(b[x] >= -eps)
              break:
          int y = find_if(A[x], A[x] + n, [](dbl d) {
            return d < -eps; }) - A[x];</pre>
          if(y == n) throw 1; // no solution
          pivot(x, y);
       while(true) {
          int y = max_{element(c, c + n)} - c;
          if(c[y] <= eps) break;</pre>
          int x = -1;
          dbl mn = 1. / 0.;
          for(int i = 0; i < m; i++)
              if(A[i][y] > eps \&\& b[i] / A[i][y] < mn)
                 mn = b[i] / A[i][y], x = i;
          if(x == -1) throw 2; // unbounded
          pivot(x, y);
      memset(sol, 0, sizeof(dbl) * n);
       for(int i = 0; i < m; i++)</pre>
          if(Y[i] < n)
              sol[Y[i]] = b[i];
       return ans;
   }
};
FFT
void fft(vector<base> &a, bool inv){
```

```
int n = (int)a.size();
                                                                  base++;
   for(int i = 1, j = 0; i < n; i++){
                                                                 }
       int bit = n \gg 1;
                                                               }
       for(; j >= bit; bit >>= 1) j -= bit;
       i += bit;
                                                               void fft(vector<num> &a, int n = -1) {
       if(i < j) swap(a[i], a[j]);</pre>
                                                                 if(n == -1) {
                                                                  n = a.size();
   for(int sz = 2; sz <= n; sz <<= 1) {
                                                                 assert((n \& (n-1)) == 0);
       double ang = 2 * PI / sz * (inv ? -1 : 1);
                                                                 int zeros = __builtin_ctz(n);
       base wlen(cos(ang), sin(ang));
                                                                 ensure_base(zeros);
       for(int i = 0; i < n; i += sz){
                                                                 int shift = base - zeros;
                                                                 for(int i = 0; i < n; i++) {
          base w(1, 0);
          for(int j = 0; j < sz / 2; j++){
                                                                  if(i < (rev[i] >> shift)) {
              base u = a[i+j], v = a[i+j + sz/2] * w;
                                                                    swap(a[i], a[rev[i] >> shift]);
              a[i+j] = u + v;
              a[i+j+sz/2] = u - v;
                                                                 3
              w *= wlen;
                                                                 for(int k = 1; k < n; k <<= 1) {
                                                                  for(int i = 0; i < n; i += 2 * k) {
          }
                                                                    for(int j = 0; j < k; j++) {
                                                                      num z = a[i+j+k] * roots[j+k];
   if(inv) for(int i = 0; i < n; i++) a[i] /= 1.0 * n;
                                                                      a[i+j+k] = a[i+j] - z;
                                                                      a[i+j] = a[i+j] + z;
FFT Tourist
                                                                }
namespace fft {
                                                               }
 typedef double dbl;
                                                               vector<num> fa. fb:
 struct num {
                                                               vector<int> multiply(vector<int> &a, vector<int> &b) {
   dbl x, y;
                                                                 int need = a.size() + b.size() - 1;
   num() \{ x = y = 0; \}
                                                                 int nbase = 0;
   num(dbl x, dbl y) : x(x), y(y) \{ \}
                                                                 while((1 << nbase) < need) nbase++;</pre>
 }:
                                                                 ensure_base(nbase);
                                                                 int sz = 1 << nbase;</pre>
 inline num operator+(num a, num b) { return num(a.x +
                                                                 if(sz > (int) fa.size()) {
   b.x, a.y + b.y); }
                                                                  fa.resize(sz);
 inline num operator-(num a, num b) { return num(a.x -
    b.x, a.y - b.y); }
                                                                 for(int i = 0; i < sz; i++) {
 inline num operator*(num a, num b) { return num(a.x *
                                                                  int x = (i < (int) a.size() ? a[i] : 0);
    b.x - a.y * b.y, a.x * b.y + a.y * b.x); }
                                                                  int y = (i < (int) b.size() ? b[i] : 0);</pre>
 inline num conj(num a) { return num(a.x, -a.y); }
                                                                   fa[i] = num(x, y);
 int base = 1;
                                                                 fft(fa, sz);
 vector<num> roots = \{\{0, 0\}, \{1, 0\}\};
                                                                num r(0, -0.25 / sz);
 vector<int> rev = \{0, 1\};
                                                                 for(int i = 0; i \le (sz >> 1); i++) {
                                                                  int j = (sz - i) & (sz - 1);
 const dbl PI = acosl(-1.0);
                                                                  num z = (fa[j] * fa[j] - conj(fa[i] * fa[i])) * r;
                                                                  if(i != j) {
 void ensure_base(int nbase) {
                                                                    fa[j] = (fa[i] * fa[i] - conj(fa[j] * fa[j])) * r
   if(nbase <= base) return;</pre>
   rev.resize(1 << nbase);</pre>
                                                                  fa[i] = z;
   for(int i = 0; i < (1 << nbase); i++) {
     rev[i] = (rev[i >> 1] >> 1) + ((i & 1) << (nbase -
                                                                 fft(fa, sz);
       1));
                                                                 vector<int> res(need);
                                                                 for(int i = 0; i < need; i++) {</pre>
   roots.resize(1 << nbase);</pre>
                                                                  res[i] = fa[i].x + 0.5;
   while(base < nbase) {</pre>
                                                                return res:
     dbl \ angle = 2*PI / (1 << (base + 1));
     for(int i = 1 << (base - 1); i < (1 << base); i++)</pre>
                                                               vector<int> multiply_mod(vector<int> &a, vector<int> &
       roots[i << 1] = roots[i];</pre>
                                                                 b, int m, int eq = 0) {
       dbl \ angle_i = angle * (2 * i + 1 - (1 << base));
                                                                 int need = a.size() + b.size() - 1;
       roots[(i \ll 1) + 1] = num(cos(angle_i), sin(
                                                                 int nbase = 0;
         angle_i));
```

```
while ((1 << nbase) < need) nbase++;</pre>
 ensure_base(nbase);
 int sz = 1 << nbase;</pre>
 if (sz > (int) fa.size()) {
   fa.resize(sz);
 for (int i = 0; i < (int) a.size(); i++) {</pre>
   int x = (a[i] \% m + m) \% m;
   fa[i] = num(x & ((1 << 15) - 1), x >> 15);
 fill(fa.begin() + a.size(), fa.begin() + sz, num {0,
     0});
 fft(fa, sz);
 if (sz > (int) fb.size()) {
   fb.resize(sz);
 if (eq) {
   copy(fa.begin(), fa.begin() + sz, fb.begin());
 } else {
   for (int i = 0; i < (int) b.size(); i++) {</pre>
     int x = (b[i] \% m + m) \% m;
     fb[i] = num(x & ((1 << 15) - 1), x >> 15);
   fill(fb.begin() + b.size(), fb.begin() + sz, num
     {0, 0});
   fft(fb, sz);
 dbl ratio = 0.25 / sz;
 num r2(0, -1);
 num r3(ratio, 0);
 num r4(0, -ratio);
 num r5(0, 1);
 for (int i = 0; i \le (sz >> 1); i++) {
   int j = (sz - i) & (sz - 1);
   num a1 = (fa[i] + conj(fa[j]));
   num a2 = (fa[i] - conj(fa[j])) * r2;
   num b1 = (fb[i] + conj(fb[j])) * r3;
   num b2 = (fb[i] - conj(fb[j])) * r4;
   if (i != j) {
     num c1 = (fa[j] + conj(fa[i]));
     num c2 = (fa[j] - conj(fa[i])) * r2;
     num d1 = (fb[j] + conj(fb[i])) * r3;
     num d2 = (fb[j] - conj(fb[i])) * r4;
     fa[i] = c1 * d1 + c2 * d2 * r5;
     fb[i] = c1 * d2 + c2 * d1;
   fa[j] = a1 * b1 + a2 * b2 * r5;
   fb[j] = a1 * b2 + a2 * b1;
 fft(fa, sz);
 fft(fb, sz);
 vector<int> res(need);
 for (int i = 0; i < need; i++) {
   long long aa = fa[i].x + 0.5;
   long long bb = fb[i].x + 0.5;
   long long cc = fa[i].y + 0.5;
   res[i] = (aa + ((bb \% m) << 15) + ((cc \% m) << 30))
      % m;
 }
 return res;
}
vector<int> square_mod(vector<int> &a, int m) {
 return multiply_mod(a, a, m, 1);
```

}

# NTT

```
const int mod = 7340033;
const int root = 5;
const int root_1 = 4404020;
const int root_pw = 1<<20;</pre>
void fft (vector<int> & a, bool invert) {
   int n = (int) a.size();
   for (int i=1, j=0; i<n; ++i) {
       int bit = n \gg 1;
       for (; j>=bit; bit>>=1)
          j -= bit;
       j += bit;
       if (i < j)
           swap (a[i], a[j]);
   }
   for (int len=2; len<=n; len<<=1) {</pre>
       int wlen = invert ? root_1 : root;
       for (int i=len; i<root_pw; i<<=1)</pre>
          wlen = int (wlen * 111 * wlen % mod);
       for (int i=0; i<n; i+=len) {</pre>
           int w = 1;
           for (int j=0; j<len/2; ++j) {</pre>
              int u = a[i+j], v = int (a[i+j+len/2] * 1
                ll * w % mod);
              a[i+j] = u+v < mod ? u+v : u+v-mod;
              a[i+j+len/2] = u-v >= 0 ? u-v : u-v+mod;
              w = int (w * 111 * wlen % mod);
           }
       }
   if (invert) {
       int nrev = reverse (n, mod);
       for (int i=0; i<n; ++i)
          a[i] = int (a[i] * 111 * nrev % mod);
   }
}
Gauss
// Solves systems of linear equations.
// To use, build a matrix of coefficients and call run(
  mat, R, C). If the i-th variable is free, row[i] will
  be -1, otherwise it's value will be ans[i].
namespace Gauss {
 const int MAXC = 1001;
 int row[MAXC];
 double ans[MAXC];
 void run(double mat[][MAXC], int R, int C) {
   REP(i, C) row[i] = -1;
   int r = 0;
   REP(c, C) {
     int k = r;
     FOR(i, r, R) if(fabs(mat[i][c]) > fabs(mat[k][c]))
       k = i:
     if(fabs(mat[k][c]) < eps) continue;</pre>
     REP(j, C+1) swap(mat[r][j], mat[k][j]);
     REP(i, R) if (i != r) {
       double w = mat[i][c] / mat[r][c];
       REP(j, C+1) mat[i][j] -= mat[r][j] * w;
```

```
row[c] = r++;
                                                                int from, to; ll flow, cap;
   }
                                                            } edge[E];
                                                            int lvl[N], vis[N], pass, start = N-2, target = N-1;
   REP(i, C) {
                                                            int qu[N], qt, px[N];
     int r = row[i];
     ans[i] = r == -1 ? 0 : mat[r][C] / mat[r][i];
                                                            ll run(int s, int sink, ll minE){
                                                                if(s == sink) return minE;
 }
                                                                11 \text{ ans} = 0;
Gauss Xor
                                                                for(; px[s] < (int)g[s].size(); px[s]++){</pre>
                                                                   int e = g[s][ px[s] ];
const 11 \text{ MAX} = 1e9;
                                                                   auto &v = edge[e], &rev = edge[e^1];
const int LOG_MAX = 64 - __builtin_clzll((11)MAX);
                                                                   if(lvl[v.to] != lvl[s]+1 || v.flow >= v.cap)
                                                                       continue; // v.cap - v.flow < lim</pre>
struct Gauss {
                                                                   11 tmp = run(v.to, sink,min(minE, v.cap-v.flow));
   array<11, LOG_MAX> vet;
                                                                   v.flow += tmp, rev.flow -= tmp;
   int size:
                                                                    ans += tmp, minE -= tmp;
   Gauss() size(0) {}
                                                                    if(minE == 0) break;
   Gauss(vector<ll> vals) size(0) {
                                                                }
       for(ll val : vals) add(val);
                                                                return ans;
   bool add(ll val) {
                                                            bool bfs(int source, int sink){
       for(int i = 0; i < LOG_MAX; i++) if(val & (1LL <<</pre>
                                                                qt = 0;
          i)) {
                                                                qu[qt++] = source;
          if(vet[i] == 0) {
                                                                lvl[source] = 1;
              vet[i] = val;
                                                                vis[source] = ++pass;
              size++;
                                                                for(int i = 0; i < qt; i++){
              return true;
                                                                    int u = qu[i];
                                                                   px[u] = 0;
          val ^= vet[i];
                                                                    if(u == sink) return true;
                                                                    for(auto& ed : g[u]) {
      return false;
                                                                       auto v = edge[ed];
   }
                                                                       if(v.flow >= v.cap || vis[v.to] == pass)
};
                                                                           continue; // v.cap - v.flow < lim</pre>
Simpson
                                                                       vis[v.to] = pass;
                                                                       lvl[v.to] = lvl[u]+1;
inline double simpson(double fl,double fr,double fmid,
                                                                       qu[qt++] = v.to;
  double 1,double r) {
                                                                   }
   return (fl + fr + 4.0 * fmid) * (r - 1) / 6.0;
                                                                }
                                                                return false;
double rsimpson(double slr, double fl, double fr, double
  fmid,double 1,double r) {
                                                            11 flow(int source = start, int sink = target){
   double mid = (1+r)*0.5;
                                                                11 \text{ ans} = 0;
   double fml = f((1+mid)*0.5), fmr = f((mid+r)*0.5);
                                                                //for(lim = (1LL << 62); lim >= 1; lim /= 2)
   double slm = simpson(fl, fmid, fml, l, mid);
                                                                while(bfs(source, sink))
   double smr = simpson(fmid, fr, fmr, mid, r);
                                                                    ans += run(source, sink, oo);
   if(fabs(slr-slm-smr) < eps and r - 1 < delta) return
                                                                return ans;
      slr;
   return rsimpson(slm,fl,fmid,fml,l,mid) + rsimpson(
                                                            void addEdge(int u, int v, ll c = 1, ll rc = \emptyset){
     smr,fmid,fr,fmr,mid,r);
                                                                edge[ne] = \{u, v, 0, c\};
                                                                g[u].push_back(ne++);
double integrate(double 1,double r) {
                                                                edge[ne] = {v, u, 0, rc};
   double mid = (1+r)*0.5;
                                                                g[v].push_back(ne++);
   double fl = f(1), fr = f(r), fmid = f(mid);
   return rsimpson(simpson(fl,fr,fmid,l,r),fl,fr,fmid,l
                                                            void reset_flow(){
                                                                for(int i = 0; i < ne; i++)
                                                                    edge[i].flow = 0;
                                                            }
Graphs
                                                            Push relabel
Dinic
                                                            // Push relabel in O(V^2 E^0.5) with gap heuristic
const int N = 100005:
                                                             // It's quite fast
const int E = 2000006;
                                                            template<typename flow_t = long long>
vector<int> g[N];
                                                            struct PushRelabel {
int ne;
                                                                struct Edge { int to, rev; flow_t f, c; };
struct Edge{
                                                                vector<vector<Edge> > g;
```

```
vector<flow_t> ec;
                                                            11 d[N];
   vector<Edge*> cur;
                                                            bool spfa(int source, int sink){
                                                               for(int i = 0; i < N; i++) d[i] = oo;</pre>
   vector<vector<int> > hs;
   vector<int> H:
                                                               inqueue[i] = 0;
   PushRelabel(int n) : g(n), ec(n), cur(n), hs(2*n), H
                                                               d[source] = 0; queue<int> q; q.push(source);
                                                               inqueue[source] = 1;
   void add_edge(int s, int t, flow_t cap, flow_t rcap
     =0) {
      if (s == t) return;
                                                               while(!q.empty()){
      Edge a = \{t, (int)g[t].size(), 0, cap\};
                                                                   int u = q.front(); q.pop();
      Edge b = \{s, (int)g[s].size(), 0, rcap\};
                                                                   inqueue[u] = 0;
       g[s].push_back(a);
                                                                   for(int e : g[u]){
       g[t].push_back(b);
                                                                      auto v = edge[e];
                                                                      if(v.cap > 0 \text{ and } d[u] + v.cost < d[v.to]){
   void add_flow(Edge& e, flow_t f) {
                                                                          d[v.to] = d[u] + v.cost; p[v.to] = e;
       Edge &back = g[e.to][e.rev];
                                                                          if(!inqueue[v.to]){
      if (!ec[e.to] && f)
                                                                              q.push(v.to); inqueue[v.to] = 1;
          hs[H[e.to]].push_back(e.to);
                                                                          }
       e.f += f, ec[e.to] += f;
                                                                      }
       back.f -= f, ec[back.to] -= f;
                                                                   }
   flow_t max_flow(int s, int t) {
                                                               return d[sink] != oo;
      int v = g.size();
                                                            // <max flow, min cost>
      H[s] = v; ec[t] = 1;
                                                            pair<11, 11> mincost(int source = start, int sink =
      vector<int> co(2 * v);
                                                              target){
       co[0] = v-1;
                                                               ll ans = 0, mf = 0;
       for(int i = 0; i < v; ++i) cur[i] = g[i].data();</pre>
       for(auto &e : g[s]) add_flow(e, e.c);
                                                               while(spfa(source, sink)){
                                                                   11 f = oo;
      if(hs[0].size())
                                                                   for(int u = sink; u != source; u = edge[ p[u] ].
       for (int hi = 0; hi >= 0;) {
          int u = hs[hi].back();
                                                                      f = min(f, edge[ p[u] ].cap);
                                                                   for(int u = sink; u != source; u = edge[ p[u] ].
          hs[hi].pop_back();
          while (ec[u] > 0) // discharge u
                                                                     from) {
              if (cur[u] == g[u].data() + g[u].size()) {
                                                                      edge[p[u]].cap -= f;
                                                                      edge[ p[u] ^ 1 ].cap += f;
                 H[u] = 1e9;
                 for(auto &e:g[u])
                                                                   }
                     if (e.c - e.f && H[u] > H[e.to]+1)
                                                                   mf += f;
                                                                   ans += f * d[sink];
                        H[u] = H[e.to]+1, cur[u] = &e;
                 if (++co[H[u]], !--co[hi] && hi < v)</pre>
                     for(int i = 0; i < v; ++i)
                                                               return {mf, ans};
                        if (hi < H[i] && H[i] < v){</pre>
                            --co[H[i]];
                                                            void addEdge(int u, int v, ll c, ll cost){
                            H[i] = v + 1;
                                                               edge[ne] = \{u, v, c, cost\};
                                                               g[u].push_back(ne++);
                 hi = H[u];
                                                               edge[ne] = \{v, u, 0, -cost\};
              } else if (cur[u]->c - cur[u]->f && H[u]
                                                               g[v].push_back(ne++);
                == H[cur[u]->to]+1)
                 add_flow(*cur[u], min(ec[u], cur[u]->c
                                                            Blossom Algorithm for General Matching
                    - cur[u]->f));
              else ++cur[u];
                                                            const int MAXN = 2020 + 1;
          while (hi >= 0 && hs[hi].empty()) --hi;
                                                            // 1-based Vertex index
                                                            int vis[MAXN], par[MAXN], orig[MAXN], match[MAXN], aux[
       return -ec[s];
                                                              MAXN], t, N;
                                                            vector<int> conn[MAXN];
};
                                                            queue<int> Q;
Min Cost Max Flow
                                                            void addEdge(int u, int v) {
                                                               conn[u].push_back(v); conn[v].push_back(u);
const 11 oo = 1e18;
const int N = 222, E = 2 * 1000006;
                                                            void init(int n) {
                                                               N = n; t = 0;
                                                               for(int i=0; i<=n; ++i)
vector<int> g[N];
int ne;
                                                                   conn[i].clear(), match[i] = aux[i] = par[i] = 0;
struct Edge{
                                                            }
   int from, to; 11 cap, cost;
                                                            void augment(int u, int v) {
} edge[E];
                                                               int pv = v, nv;
int start = N-1, target = N-2, p[N]; int inqueue[N];
```

```
pv = par[v]; nv = match[pv];
                                                               for(int v : g[u]) if(v != p && v != big)
      match[v] = pv; match[pv] = v;
                                                                  add(v, u);
      v = nv:
                                                           }
   } while(u != pv);
                                                           void dfs(int u, int p, int keep){
                                                               int big = -1, mmx = -1;
int lca(int v, int w) {
                                                               for(int v : g[u]) if(v != p \&\& sz[v] > mmx)
                                                                  mmx = sz[v], big = v;
   ++t;
   while(true) {
                                                               for(int v : g[u]) if(v != p && v != big)
       if(v) {
                                                                   dfs(v, u, 0);
          if(aux[v] == t) return v; aux[v] = t;
                                                               if(big != -1) dfs(big, u, 1);
          v = orig[par[match[v]]];
                                                               add(u, p, big);
                                                               for(auto x : q[u]){
       swap(v, w);
                                                                   // answer all queries for this vx
   }
                                                               if(!keep){ /*Remove data from this subtree*/ }
void blossom(int v, int w, int a) {
   while(orig[v] != a) {
                                                           Centroid Decomposition
      par[v] = w; w = match[v];
      if(vis[w] == 1) Q.push(w), vis[w] = 0;
                                                           void decomp(int v, int p){
       orig[v] = orig[w] = a; v = par[w];
                                                               int treesize = calc_sz(v, v);
                                                               if(treesize < k) return;</pre>
                                                               int cent = centroid(v, v, treesize);
bool bfs(int u) {
                                                               erased[cent] = 1;
   fill(vis+1, vis+1+N, -1); iota(orig + 1, orig + N +
     1, 1);
                                                               for(int i = 1; i <= treesize; i++) dist[i] = 1e18;</pre>
   Q = queue < int > (); Q.push(u); vis[u] = 0;
   while(!Q.empty()) {
                                                               for(pair<int,int> x : G[cent]) if(!erased[x.ff]){
      int v = Q.front(); Q.pop();
                                                                  procurar_ans(x.ff, cent, 1, x.ss); // linear
       for(int x: conn[v]) {
                                                                  atualiza_dist(x.ff, cent, 1, x.ss); // linear
          if(vis[x] == -1) {
                                                               }
             par[x] = v; vis[x] = 1;
              if(!match[x]) return augment(u, x), true;
                                                               for(pair<int, int> x : G[cent]) if(!erased[x.ff])
              Q.push(match[x]); vis[match[x]] = 0;
                                                                   decomp(x.ff, cent);
          else if(vis[x] == 0 && orig[v] != orig[x]) {
                                                           Kosaraju
              int a = lca(orig[v], orig[x]);
             blossom(x, v, a); blossom(v, x, a);
                                                           vector<int> g[N], gt[N], S; int vis[N], cor[N];
          }
                                                           void dfs(int u){
       }
                                                               vis[u] = 1; for(int v : g[u]) if(!vis[v]) dfs(v);
   }
                                                               S.push_back(u);
   return false;
                                                           void dfst(int u, int e){
int Match() {
                                                               cor[u] = e;
   int ans = 0;
                                                               for(int v : gt[u]) if(!cor[v]) dfst(v, e);
   // find random matching (not necessary, constant
     improvement)
                                                           void kosaraju(){
   vector<int> V(N-1); iota(V.begin(), V.end(), 1);
                                                               for(int i = 1; i <= n; i++) if(!vis[i]) dfs(i);</pre>
   shuffle(V.begin(), V.end(), mt19937(0x94949));
                                                               for(int i = 1; i \le n; i++) for(int j : g[i])
   for(auto x: V) if(!match[x]){
                                                                  gt[j].push_back(i);
       for(auto y: conn[x]) if(!match[y]) {
                                                               int e = 0; reverse(S.begin(), S.end());
          match[x] = y, match[y] = x;
                                                               for(int u : S) if(!cor[u]) dfst(u, ++e);
          ++ans; break;
                                                           }
                                                           Tarjan
   for(int i=1; i<=N; ++i) if(!match[i] && bfs(i)) ++</pre>
                                                           int cnt = 0, root;
                                                           void dfs(int u, int p = -1){
   return ans;
                                                               low[u] = num[u] = ++t;
                                                               for(int v : g[u]){
Small to Large
                                                                   if(!num[v]){
                                                                      dfs(v, u);
void cnt_sz(int u, int p = -1){
                                                                         if(u == root) cnt++;
   sz[u] = 1;
                                                                      if(low[v] >= num[u]) u PONTO DE ARTICULACAO;
                                                                      if(low[v] > num[u]) ARESTA u->v PONTE;
   for(int v : g[u]) if(v != p)
       cnt_sz(v, u), sz[u] += sz[v];
                                                                      low[u] = min(low[u], low[v]);
void add(int u, int p, int big = -1){
                                                                   else if(v != p) low[u] = min(low[u], num[v]);
   // Update info about this vx in global answer
```

```
}
                                                            vector<int> bucket[N];
root PONTO DE ARTICULAÇÃO <=> cnt > 1
                                                            vector<int> down[N];
void tarjanSCC(int u){
                                                            void prep(int u){
   low[u] = num[u] = ++cnt;
                                                               S.push_back(u);
   vis[u] = 1;
                                                               id[u] = ++dfs_time;
   S.push_back(u);
                                                               label[u] = sdom[u] = dsu[u] = u;
   for(int v : g[u]){
                                                               for(int v : g[u]){
       if(!num[v]) tarjanSCC(v);
                                                                   if(!id[v])
       if(vis[v]) low[u] = min(low[u], low[v]);
                                                                       prep(v), down[u].push_back(v);
   if(low[u] == num[u]){
                                                                   gt[v].push_back(u);
       ssc[u] = ++ssc_cnt; int v;
                                                               }
                                                            }
       do{
          v = S.back(); S.pop_back(); vis[v] = 0;
          ssc[v] = ssc_cnt;
                                                            int fnd(int u, int flag = 0){
       }while(u != v);
                                                               if(u == dsu[u]) return u;
   }
                                                               int v = fnd(dsu[u], 1), b = label[ dsu[u] ];
                                                               if(id[ sdom[b] ] < id[ sdom[ label[u] ] ])</pre>
                                                                   label[u] = b;
Max Clique
                                                               dsu[u] = v;
                                                               return flag ? v : label[u];
long long adj[N], dp[N];
                                                            }
for(int i = 0; i < n; i++){
                                                            void build_dominator_tree(int root, int sz){
   for(int j = 0; j < n; j++){
                                                               // memset(id, 0, sizeof(int) * (sz + 1));
       int x;
                                                               // for(int i = 0; i <= sz; i++) T[i].clear();
       scanf("%d",&x);
                                                               prep(root);
       if(x \mid \mid i == j)
                                                               reverse(S.begin(), S.end());
          adj[i] |= 1LL << j;
   }
                                                               int w;
}
                                                               for(int u : S){
                                                                   for(int v : gt[u]){
int resto = n - n/2;
                                                                      w = fnd(v);
int C = n/2;
                                                                      if(id[ sdom[w] ] < id[ sdom[u] ])
for(int i = 1; i < (1 << resto); i++){</pre>
                                                                          sdom[u] = sdom[w];
   int x = i;
                                                                   }
   for(int j = 0; j < resto; j++)
                                                                   gt[u].clear();
      if(i & (1 << j))
          x \&= adj[j + C] >> C;
                                                                   if(u != root) bucket[ sdom[u] ].push_back(u);
   if(x == i){
      dp[i] = __builtin_popcount(i);
                                                                   for(int v : bucket[u]){
   }
                                                                      w = fnd(v);
}
                                                                      if(sdom[w] == sdom[v]) idom[v] = sdom[v];
                                                                       else idom[v] = w;
for(int i = 1; i < (1 << resto); i++)</pre>
   for(int j = 0; j < resto; j++)
                                                                   bucket[u].clear();
       if(i & (1 << j))
          dp[i] = max(dp[i], dp[i ^ (1 << j)]);
                                                                   for(int v : down[u]) dsu[v] = u;
                                                                   down[u].clear();
int maxCliq = 0;
for(int i = 0; i < (1 << C); i++){
   int x = i, y = (1 << resto) - 1;
                                                               reverse(S.begin(), S.end());
   for(int j = 0; j < C; j++)
                                                               for(int u : S) if(u != root){
      if(i & (1 << j))
                                                                   if(idom[u] != sdom[u]) idom[u] = idom[ idom[u] ];
          x \&= adj[j] \& ((1 << C) - 1), y \&= adj[j] >>
                                                                   T[ idom[u] ].push_back(u);
            C:
   if(x != i) continue;
                                                               S.clear();
   maxCliq = max(maxCliq, __builtin_popcount(i) + dp[y
     ]);
}
                                                            Min Cost Matching
Dominator Tree
                                                            // Min cost matching
vector<int> g[N], gt[N], T[N];
                                                            // O(n^2 * m)
vector<int> S;
                                                           // n == nro de linhas
int dsu[N], label[N];
                                                            // m == nro de colunas
int sdom[N], idom[N], dfs_time, id[N];
                                                            // n <= m | flow == n
```

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```
// a[i][j] = custo pra conectar i a j
vector<int> u(n + 1), v(m + 1), p(m + 1), way(m + 1);
for(int i = 1; i \le n; ++i){
   p[0] = i;
   int j0 = 0;
   vector<int> minv(m + 1 , oo);
   vector<char> used(m + 1 , false);
   do{
       used[j0] = true;
       int i0 = p[j0], delta = oo, j1;
       for(int j = 1; j \le m; ++j)
          if(! used[j]){
              int cur = a[i0][j] - u[i0] - v[j];
              if(cur < minv[j])</pre>
                  minv[j] = cur, way[j] = j0;
              if(minv[j] < delta)</pre>
                  delta = minv[j] , j1 = j;
       for(int j = 0; j \le m; ++j)
          if(used[j])
              u[p[j]] += delta, v[j] -= delta;
          else
              minv[j] -= delta;
       j0 = j1;
   {\bf while}(p[j0] != 0);
   do{
       int j1 = way[j0];
       p[j0] = p[j1];
       j0 = j1;
   }while(j0);
// match[i] = coluna escolhida para linha i
vector<int> match(n + 1);
for(int j = 1; j \le m; ++j)
   match[p[j]] = j;
int cost = -v[0];
```

# **Strings**

# **Aho Corasick**

```
int to[N][A];
int ne = 2, fail[N], term[N];
void add_string(const char *str, int id){
   int p = 1;
   for(int i = 0; str[i]; i++){
      int ch = str[i] - 'a';
      if(!to[p][ch]) to[p][ch] = ne++;
      p = to[p][ch];
   }
   term[p]++;
void init(){
   for(int i = 0; i < ne; i++) fail[i] = 1;</pre>
   queue<int> q; q.push(1);
   int u, v; char c;
   while(!q.empty()){
      u = q.front(); q.pop();
       for(int i = 0; i < A; i++){
          if(to[u][i]){
              v = to[u][i]; q.push(v);
              if(u != 1){
                 fail[v] = to[ fail[u] ][i];
                 term[v] += term[ fail[v] ];
```

```
else if(u != 1) to[u][i] = to[ fail[u] ][i];
          else to[u][i] = 1;
       }
   }
}
void clean() {
   memset(to, 0, ne * sizeof(to[0]));
   memset(fail, 0, ne * sizeof(fail[0]));
   memset(term, 0, ne * sizeof(term[0]));
   memset(to, 0, ne * sizeof(to[0]));
   ne = 2;
Suffix Array
int lcp[N], c[N];
// Caractere final da string '\0' esta sendo considerado
   parte da string s
void build_sa(char s[], int n, int a[]){
   const int A = 300; // Tamanho do alfabeto
   int c1[n], a1[n], h[n + A];
   memset(h, 0, sizeof h);
   for(int i = 0; i < n; i++) {
       c[i] = s[i];
      h[c[i] + 1]++;
   }
   partial_sum(h, h + A, h);
   for(int i = 0; i < n; i++)
       a[h[c[i]]++] = i;
   for(int i = 0; i < n; i++)
      h[c[i]]--;
   for(int L = 1; L < n; L <<= 1) {
       for(int i = 0; i < n; i++) {
          int j = (a[i] - L + n) \% n;
          a1[h[c[j]]++] = j;
       }
       int cc = -1;
       for(int i = 0; i < n; i++) {
          if(i == 0 || c[a1[i]] != c[a1[i-1]] || c[(a1[
            i] + L) % n] != c[(a1[i-1] + L) % n])
              h[++cc] = i;
          c1[a1[i]] = cc;
       }
       memcpy(a, a1, sizeof a1);
       memcpy(c, c1, sizeof c1);
       if(cc == n-1) break;
   }
}
void build_lcp(char s[], int n, int a[]){ // lcp[i] =
  lcp(s[:i], s[:i+1])
   int k = 0;
   //memset(lcp, 0, sizeof lcp);
   for(int i = 0; i < n; i++){
       if(c[i] == n-1) continue;
       int j = a[c[i]+1];
       while(i+k < n \&\& j+k < n \&\& s[i+k] == s[j+k]) k
       lcp[c[i]] = k;
```

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```
if(k) k--;
                                                           bool is_leaf(int u) { return len[u] > n; }
   }
}
                                                           int get_len(int u) {
                                                             if (!u) return 0;
int comp_lcp(int i, int j){
                                                             if (is_leaf(u)) return n - fpos[u];
                                                             return len[u];
   if(i == j) return n - i;
   if(c[i] > c[j]) swap(i, j);
   return min(lcp[k]  for k in [c[i], c[j]-1]);
                                                           int leafs[maxn];
                                                           int calc_leafs(int u = 0) {
                                                             leafs[u] = is_leaf(u);
Adamant Suffix Tree
                                                             for (const auto &c : to[u]) leafs[u] += calc_leafs(c.
                                                               second);
namespace sf {
                                                             return leafs[u];
const int inf = 1e9;
                                                           }; // namespace sf
const int maxn = 200005;
char s[maxn];
                                                           int main() { sf::len[0] = sf::inf; }
map<int, int> to[maxn];
                                                           Z Algorithm
int len[maxn], fpos[maxn], link[maxn];
int node, pos;
                                                           vector<int> z_algo(const string &s) {
int sz = 1, n = 0;
                                                               int n = s.size(), L = 0, R = 0;
                                                               vector<int> z(n, 0);
int make_node(int _pos, int _len) {
                                                               for(int i = 1; i < n; i++){
 fpos[sz] = _pos;
                                                                  if(i \le R) z[i] = min(z[i-L], R - i + 1);
 len[sz] = _len;
                                                                  while(z[i]+i < n \& s[z[i]+i] == s[z[i]])
 return sz++;
                                                                      z[i]++;
                                                                  if(i+z[i]-1 > R) L = i, R = i + z[i] - 1;
void go_edge() {
 while (pos > len[to[node][s[n - pos]]]) {
                                                               return z;
   node = to[node][s[n - pos]];
   pos -= len[node];
                                                           Prefix function/KMP
}
                                                           vector<int> preffix_function(const string &s){
void add_letter(int c) {
                                                               int n = s.size(); vector<int> b(n+1);
 s[n++] = (char)c;
                                                               b[0] = -1; int i = 0, j = -1;
 pos++;
                                                               while(i < n){
 int last = 0;
                                                                  while(j \ge 0 \& s[i] != s[j]) j = b[j];
 while (pos > 0) {
                                                                  b[++i] = ++j;
   go_edge();
                                                               }
   int edge = s[n - pos];
                                                               return b;
   int &v = to[node][edge];
   int t = s[fpos[v] + pos - 1];
                                                           void kmp(const string &t, const string &p){
   if (v == 0) {
                                                               vector<int> b = preffix_function(p);
     v = make_node(n - pos, inf);
                                                               int n = t.size(), m = p.size();
     link[last] = node;
                                                               int j = 0;
     last = 0;
                                                               for(int i = 0; i < n; i++){
   } else if (t == c) {
                                                                  while(j \ge 0 \& t[i] != p[j]) j = b[j];
     link[last] = node;
                                                                  j++;
     return;
                                                                  if(j == m){
   } else {
                                                                      //patern of p found on t
     int u = make_node(fpos[v], pos - 1);
                                                                      j = b[j];
     to[u][c] = make\_node(n - 1, inf);
                                                                  }
     to[u][t] = v;
                                                               }
     fpos[v] += pos - 1;
     len[v] -= pos - 1;
     v = u;
                                                           Min rotation
     link[last] = u;
                                                           int min_rotation(int *s, int N) {
     last = u:
                                                             REP(i, N) s[N+i] = s[i];
   if (node == 0)
     pos--;
                                                             int a = 0;
                                                             REP(b, N) REP(i, N) {
   else
                                                               if (a+i == b \mid \mid s[a+i] < s[b+i]) { b += max(0, i-1);}
     node = link[node];
                                                                  break; }
 }
                                                               if (s[a+i] > s[b+i]) { a = b; break; }
void add_string(char *str) {
 for (int i = 0; str[i]; i++) add_letter(str[i]);
                                                             return a;
 add_letter('$');
```

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# Manacher

```
void manacher(char *s, int N, int *rad) {
 static char t[2*MAX];
 int m = 2*N - 1;
 REP(i, m) t[i] = -1;
 REP(i, N) t[2*i] = s[i];
 int x = 0;
 FOR(i, 1, m) {
   int &r = rad[i] = 0;
   if (i <= x+rad[x]) r = min(rad[x+x-i], x+rad[x]-i);</pre>
   while (i-r-1 >= 0 \& i+r+1 < m \& t[i-r-1] == t[i+r]
     +1]) ++r:
   if (i+r >= x+rad[x]) x = i;
 REP(i, m) if (i-rad[i] == 0 || i+rad[i] == m-1) ++rad[
 REP(i, m) rad[i] \neq 2;
Suffix Automaton
map<char, int> to[2*N];
int link[2*N], len[2*N], last = 0, sz = 1;
void add_letter(char c){
   int p = last;
   last = sz++;
   len[last] = len[p] + 1;
   for(; !to[p][c]; p = link[p]) to[p][c] = last;
   if(to[p][c] == last){
       link[last] = 0;
       return:
   int u = to[p][c];
   if(len[u] == len[p]+1){
       link[last] = u;
       return;
   int c1 = sz++;
   to[c1] = to[u];
   link[c1] = link[u];
   len[c1] = len[p]+1;
   link[last] = link[u] = c1;
   for(; to[p][c] == u; p = link[p]) to[p][c] = c1;
Suffix Tree
namespace sf {
// const int NS = ; const int N = * 2;
int cn, cd, ns, en = 1, lst;
string S[NS]; int si = -1;
vector<int> sufn[N]; // sufn[si][i] no do sufixo S[si][i
  . . . ]
struct node {
   int 1, r, si, p, suf;
   map<char, int> adj;
   node() : 1(0), r(-1), suf(0), p(0) {}
   node(int L, int R, int S, int P) : l(L), r(R), si(S)
     , p(P) \{ \}
   inline int len() { return r - l + 1; }
   inline int operator[](int i) { return S[si][l + i];
   inline int& operator()(char c) { return adj[c]; }
} t[N];
```

```
inline int new_node(int L, int R, int S, int P) { t[en]
  = node(L, R, S, P); return en++; }
void add_string(string s) {
   s += '$'; S[++si] = s; sufn[si].resize(s.size() + 1)
     ; cn = cd = 0;
   int i = 0; const int n = s.size();
   for(int j = 0; j < n; j++)
       for(; i <= j; i++) {</pre>
          if(cd == t[cn].len() \& t[cn](s[j])) { cn = t}
             [cn](s[j]); cd = 0; }
          if(cd < t[cn].len() \& t[cn][cd] == s[j]) {
              cd++;
              if(j < s.size() - 1) break;
              else {
                 if(i) t[lst].suf = cn;
                 for(; i <= j; i++) { sufn[si][i] = cn;</pre>
                    cn = t[cn].suf; }
              }
          } else if(cd == t[cn].len()) {
              sufn[si][i] = en;
              if(i) t[lst].suf = en; lst = en;
              t[cn](s[j]) = new_node(j, n - 1, si, cn);
              cn = t[cn].suf; cd = t[cn].len();
          } else {
              int mid = new_node(t[cn].1, t[cn].1 + cd -
                 1, t[cn].si, t[cn].p);
              t[t[cn].p](t[cn][0]) = mid;
              if(ns) t[ns].suf = mid;
              if(i) t[lst].suf = en; lst = en;
              sufn[si][i] = en;
              t[mid](s[j]) = new_node(j, n - 1, si, mid)
              t[mid](t[cn][cd]) = cn;
              t[cn].p = mid; t[cn].l += cd; cn = t[mid].
              int g = cn? j - cd : i + 1; cn = t[cn].suf
              while(g < j \&\& g + t[t[cn](S[si][g])].len
                () <= i) {
                 cn = t[cn](S[si][g]); g += t[cn].len();
              if(g == j) \{ ns = 0; t[mid].suf = cn; cd =
                 t[cn].len(); }
              else { ns = mid; cn = t[cn](S[si][g]); cd
                = j - g; }
          }
      }
   }
};
2D basics
```

# Geometry

```
typedef double cod;
double eps = 1e-7;
bool eq(cod a, cod b){ return abs(a - b) <= eps; }</pre>
struct vec{
   cod x, y; int id;
   vec(cod a = 0, cod b = 0) : x(a), y(b) {}
   vec operator+(const vec &o) const{
       return \{x + o.x, y + o.y\};
   vec operator-(const vec &o) const{
       return \{x - o.x, y - o.y\};
   vec operator*(cod t) const{
```

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```
return {x * t, y * t};
}
vec operator/(cod t) const{
   return {x / t, y / t};
cod operator*(const vec &o) const{ // cos
   return x * o.x + y * o.y;
cod operator^(const vec &o) const{ // sin
   return x * o.y - y * o.x;
bool operator==(const vec &o) const{
   return eq(x, o.x) \& eq(y, o.y);
bool operator<(const vec &o) const{</pre>
   if(!eq(x, o.x)) return x < o.x;
   return y < o.y;</pre>
}
cod cross(const vec &a, const vec &b) const{
   return (a-(*this)) ^ (b-(*this));
int ccw(const vec &a, const vec &b) const{
   cod tmp = cross(a, b);
   return (tmp > eps) - (tmp < -eps);</pre>
}
cod dot(const vec &a, const vec &b) const{
   return (a-(*this)) * (b-(*this));
cod len() const{
   return sqrt(x * x + y * y); // <
double angle(const vec &a, const vec &b) const{
   return atan2(cross(a, b), dot(a, b));
double tan(const vec &a, const vec &b) const{
   return cross(a, b) / dot(a, b);
}
vec unit() const{
   return operator/(len());
int quad() const{
   if(x > 0 \& y >= 0) return 0;
   if(x \le 0 \&\& y > 0) return 1;
   if(x < 0 \& y <=0) return 2;
   return 3;
bool comp(const vec &a, const vec &b) const{
   return (a - *this).comp(b - *this);
bool comp(vec b){
   if(quad() != b.quad()) return quad() < b.quad();</pre>
   if(!eq(operator^(b), 0)) return operator^(b) > 0;
   return (*this) * (*this) < b * b;
template<class T>
void sort_by_angle(T first, T last) const{
   std::sort(first, last, [=](const vec &a, const
     vec &b){
      return comp(a, b);
   });
vec rot90() const{ return {-y, x}; }
vec rot(double a) const{
   return \{\cos(a)*x - \sin(a)*y, \sin(a)*x + \cos(a)*y\};
vec proj(const vec &b) const{ // proj of *this onto
  b
```

```
cod k = operator*(b) / (b * b);
       return b * k;
   }
   // proj of (*this) onto the plane orthogonal to b
   vec rejection(vec b) const{
       return (*this) - proj(b);
};
struct line{
   cod a, b, c; vec n;
   line(vec q, vec w){ // q.cross(w, (x, y)) = 0
       a = -(w.y-q.y);
      b = w.x-q.x;
      c = -(a * q.x + b * q.y);
      n = \{a, b\};
   cod dist(const vec &o) const{
       return abs(eval(o)) / n.len();
   bool contains(const vec &o) const{
       return eq(a * o.x + b * o.y + c, \emptyset);
   cod dist(const line &o) const{
       if(!parallel(o)) return 0;
       if(!eq(o.a * b, o.b * a)) return 0;
       if(!eq(a, 0))
          return abs(c - o.c * a / o.a) / n.len();
       if(!eq(b, 0))
          return abs(c - o.c * b / o.b) / n.len();
       return abs(c - o.c);
   bool parallel(const line &o) const{
       return eq(n ^ o.n, 0);
   bool operator==(const line &o) const{
       if(!eq(a*o.b, b*o.a)) return false;
       if(!eq(a*o.c, c*o.a)) return false;
       if(!eq(c*o.b, b*o.c)) return false;
       return true;
   bool intersect(const line &o) const{
       return !parallel(o) || *this == o;
   vec inter(const line &o) const{
       if(parallel(o)){
          if(*this == o){ }
          else{ /* dont intersect */ }
       }
       auto tmp = n \cdot o.n;
       return {(o.c*b -c*o.b)/tmp, (o.a*c -a*o.c)/tmp};
   vec at_x(cod x) const{
       return \{x, (-c-a*x)/b\};
   vec at_y(cod y) const{
       return {(-c-b*y)/a, y};
   cod eval(const vec &o) const{
      return a * o.x + b * o.y + c;
};
struct segment{
   segment(vec a = vec(), vec b = vec()): p(a), q(b) {}
```

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```
bool onstrip(const vec &o) const{ // onstrip strip
      return p.dot(o, q) >= -eps && q.dot(o, p) >= -eps
   }
   cod len() const{
      return (p-q).len();
   cod dist(const vec &o) const{
      if(onstrip(o)) return line(p, q).dist(o);
      return min((o-q).len(), (o-p).len());
   bool contains(const vec &o) const{
      return eq(p.cross(q, o), 0) && onstrip(o);
   bool intersect(const segment &o) const{
      if(contains(o.p)) return true;
      if(contains(o.q)) return true;
      if(o.contains(q)) return true;
      if(o.contains(p)) return true;
      return p.ccw(q, o.p) * p.ccw(q, o.q) == -1
      && o.p.ccw(o.q, q) * o.p.ccw(o.q, p) == -1;
   bool intersect(const line &o) const{
      return o.eval(p) * o.eval(q) <= 0;</pre>
   }
   cod dist(const segment &o) const{
      if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || onstrip(o.q)
          || o.onstrip(p) || o.onstrip(q))
             return line(p, q).dist(line(o.p, o.q));
      else if(intersect(o)) return 0;
      return min(min(dist(o.p), dist(o.q)),
                min(o.dist(p), o.dist(q)));
   cod dist(const line &o) const{
      if(line(p, q).parallel(o))
          return line(p, q).dist(o);
      else if(intersect(o)) return 0;
      return min(o.dist(p), o.dist(q));
};
struct hray{
   vec p, q;
   hray(vec a = vec(), vec b = vec()): p(a), q(b){}
   bool onstrip(const vec &o) const{ // onstrip strip
      return p.dot(q, o) >= -eps;
   cod dist(const vec &o) const{
      if(onstrip(o)) return line(p, q).dist(o);
      return (o-p).len();
   bool intersect(const segment &o) const{
      if(!o.intersect(line(p,q))) return false;
      if(line(o.p, o.q).parallel(line(p,q)))
          return contains(o.p) || contains(o.q);
      return contains(line(p,q).inter(line(o.p,o.q)));
   bool contains(const vec &o) const{
      return eq(line(p, q).eval(o), 0) && onstrip(o);
   cod dist(const segment &o) const{
      if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || onstrip(o.q))
             return line(p, q).dist(line(o.p, o.q));
          return o.dist(p);
```

```
else if(intersect(o)) return 0;
       return min(min(dist(o.p), dist(o.q)),
                o.dist(p));
   bool intersect(const hray &o) const{
       if(!line(p, q).parallel(line(o.p, o.q)))
          return false;
       auto pt = line(p, q).inter(line(o.p, o.q));
       return contains(pt) && o.contains(pt); // <<</pre>
   bool intersect(const line &o) const{
       if(line(p, q).parallel(o)) return line(p, q)== o;
       if(o.contains(p) || o.contains(q)) return true;
       return (o.eval(p) >= -eps)^(o.eval(p)<o.eval(q));</pre>
       return contains(o.inter(line(p, q)));
   cod dist(const line &o) const{
       if(line(p,q).parallel(o))
          return line(p,q).dist(o);
       else if(intersect(o)) return 0;
       return o.dist(p);
   cod dist(const hray &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || o.onstrip(p))
              return line(p,q).dist(line(o.p, o.q));
          return (p-o.p).len();
       else if(intersect(o)) return 0;
       return min(dist(o.p), o.dist(p));
   }
};
double heron(cod a, cod b, cod c){
   cod s = (a + b + c) / 2;
   return sqrt(s * (s - a) * (s - b) * (s - c));
}
line mediatrix(const vec &a, const vec &b) {
   auto tmp = (b - a) * 2;
   return line(tmp.x, tmp.y, a * a - b * b);
}
struct circle {
   vec c; cod r;
   circle() : c(0, 0), r(0) {}
   circle(const vec o) : c(o), r(0) {}
   circle(const vec &a, const vec &b) {
       c = (a + b) * 0.5; r = (a - c).len();
   circle(const vec &a, const vec &b, const vec &cc) {
       c = mediatrix(a, b).inter(mediatrix(b, cc));
       r = (a - c).len();
   bool inside(const vec &a) const {
       return (a - c).len() \ll r;
circle min_circle_cover(vector<vec> v) {
   random_shuffle(v.begin(), v.end());
   circle ans:
   int n = (int)v.size();
   for(int i = 0; i < n; i++) if(!ans.inside(v[i])) {
       ans = circle(v[i]);
       for(int j = 0; j < i; j++) if(!ans.inside(v[j])){
          ans = circle(v[i], v[j]);
          for(int k=0; k< j; k++)if(!ans.inside(v[k])){
              ans = circle(v[i], v[j], v[k]);
```

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```
}
   return ans;
Circle line intersection
// intersection of line a * x + b * y + c = 0
// and circle centered at the origin with radius r
double r, a, b, c; // given as input
double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);
if(c*c > r*r*(a*a+b*b)+EPS)
   puts("no points");
else if(abs(c*c - r*r*(a*a+b*b)) < EPS){
   puts("1 point");
   cout << x0 << ' ' << y0 << ' 'n';
else {
   double d = r*r - c*c/(a*a+b*b);
   double mult = sqrt (d / (a*a+b*b));
   double ax, ay, bx, by;
   ax = x0 + b * mult;
   bx = x0 - b * mult;
   ay = y0 - a * mult;
   by = y0 + a * mult;
   puts ("2 points");
   cout<<ax<<' '<<ay<<'\n'<<bx<<' '<<by<<'\n';
Half plane intersection
const double eps = 1e-8;
typedef pair<long double, long double> pi;
bool z(long double x){ return fabs(x) < eps; }</pre>
struct line{
   long double a, b, c;
   bool operator<(const line &l)const{</pre>
       bool flag1 = pi(a, b) > pi(0, 0);
       bool flag2 = pi(1.a, 1.b) > pi(0, 0);
       if(flag1 != flag2) return flag1 > flag2;
       long double t = ccw(pi(0, 0), pi(a, b), pi(1.a, 1)
         .b)):
       return z(t) ? c * hypot(1.a, 1.b) < 1.c * hypot(a
         , b) : t > 0;
   pi slope(){ return pi(a, b); }
};
pi cross(line a, line b){
   long double det = a.a * b.b - b.a * a.b;
   return pi((a.c * b.b - a.b * b.c) / det, (a.a * b.c
     - a.c * b.a) / det);
bool bad(line a, line b, line c){
   if(ccw(pi(0, 0), a.slope(), b.slope()) <= 0) return</pre>
     false;
   pi crs = cross(a, b);
   return crs.first * c.a + crs.second * c.b >= c.c;
bool solve(vector<line> v, vector<pi> &solution){ // ax
  + by <= c;
   sort(v.begin(), v.end());
   deque<line> dq;
   for(auto &i : v){
       if(!dq.empty() \&\& z(ccw(pi(0, 0), dq.back().slope))
         (), i.slope()))) continue;
       while(dq.size() >= 2 && bad(dq[dq.size()-2], dq.
         back(), i)) dq.pop_back();
       while(dq.size() \geq 2 && bad(i, dq[0], dq[1])) dq.
```

```
pop_front();
       dq.push_back(i);
   }
   while(dq.size() > 2 && bad(dq[dq.size()-2], dq.back
      (), dq[0])) dq.pop_back();
   while(dq.size() > 2 && bad(dq.back(), dq[0], dq[1]))
      dq.pop_front();
   vector<pi> tmp;
   for(int i=0; i<dq.size(); i++){</pre>
       line cur = dq[i], nxt = dq[(i+1)%dq.size()];
       if(ccw(pi(0, 0), cur.slope(), nxt.slope()) <= eps</pre>
         ) return false;
       tmp.push_back(cross(cur, nxt));
   solution = tmp;
   return true:
Detect empty Half plane intersection
// abs(point a) = absolute value of a
// ccw(a, b, c) = a.ccw(b, c)
pair<bool, point> half_inter(vector<pair<point,point> >
  &vet){
   random_shuffle(all(vet));
   rep(i,0,sz(vet)) if(ccw(vet[i].x,vet[i].y,p) != 1){
       point dir = (vet[i].y - vet[i].x) / abs(vet[i].y
         - vet[i].x);
       point l = vet[i].x - dir*1e15;
       point r = vet[i].x + dir*1e15;
       if(r < 1) swap(1, r);
       rep(j, 0, i){
          if(ccw(point(), vet[i].x-vet[i].y, vet[j].x-
            vet[j].y) == 0){
              if(ccw(vet[j].x, vet[j].y, p) == 1)
                 continue;
              return mp(false, point());
          if(ccw(vet[j].x, vet[j].y, 1) != 1)
              l = max(l, line_intersect(vet[i].x,vet[i].
                y,vet[j].x,vet[j].y));
          if(ccw(vet[j].x, vet[j].y, r) != 1)
              r = min(r, line_intersect(vet[i].x,vet[i].
                y,vet[j].x,vet[j].y));
          if(!(1 < r)) return mp(false, point());</pre>
       }
      p = r;
   }
   return mp(true, p);
```

# **Circle Circle intersection**

Assume that the first circle is centered at the origin and second at (x2, y2). Find circle line intersection of first circle and line Ax + By + C = 0, where  $A = -2x_2$ ,  $B = -2y_2$ ,  $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$ .

Be aware of corner case with two circles centered at the same point.

# Tangents of two circles

```
// solve first for same circle(and infinitely many
  tangents)
// Find up to four tangents of two circles
void tangents(pt c, double r1, double r2, vector<line> &
  ans){
```

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```
double r = r2 - r1;
   double z = c.x * c.x + c.y * c.y;
   double d = z - r * r;
   if(d < -EPS) return;</pre>
   d = sqrt(abs(d));
   line 1;
   1.a = (c.x * r + c.y * d) / z;
   1.b = (c.y * r - c.x * d) / z;
   1.c = r1;
   ans.push_back (1);
vector<line> tangents(circle a, circle b){
   vector<line> ans:
   pt aux = a.center - b.center;
   for(int i = -1; i \le 1; i += 2)
       for(int j = -1; j <= 1; j += 2)
          tangents(aux, a.r * i, b.r * j, ans);
   for(size_t i = 0; i < ans.size(); ++i)</pre>
       ans[i].c = ans[i].a * a.x + ans[i].b * a.y;
   return ans:
Convex Hull
vector<vec> monotone_chain_ch(vector<vec> P){
   sort(P.begin(), P.end());
                                                               }
   vector<vec> L, U;
   for(auto p : P){
       while(L.size() >= 2 && L[L.size() - 2].cross(L.
         back(), p) < 0)
          L.pop_back();
       L.push_back(p);
   }
   reverse(P.begin(), P.end());
   for(auto p : P){
       while(U.size() >= 2 && U[U.size() - 2].cross(U.
         back(), p) < 0)
          U.pop_back();
       U.push_back(p);
   }
   L.pop_back(), U.pop_back();
   L.reserve(L.size() + U.size());
   L.insert(L.end(), U.begin(), U.end());
   return L;
Check point inside polygon
bool below(const vector<vec> &vet, vec p){
   auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
   if(it == vet.begin()) return *it == p;
   return prev(it)->cross(*it, p) <= 0;</pre>
                                                            Center of mass
}
bool above(const vector<vec> &vet, vec p){
   auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
   if(it == vet.begin()) return *it == p;
   return prev(it)->cross(*it, p) >= 0;
```

```
// lowerhull, upperhull and point, borders included
bool inside_poly(const vector<vec> &lo, const vector<vec</pre>
  > &hi, vec p){
   return below(hi, p) && above(lo, p);
Check point inside polygon without lower/upper
// borders included
// must not have 3 colinear consecutive points
bool inside_poly(const vector<vec> &v, vec p){
   if(v[0].ccw(v[1], p) < 0) return false;
   if(v[0].ccw(v.back(), p) > 0) return 0;
   if(v[0].ccw(v.back(), p) == 0)
      return v[0].dot(p, v.back()) >= 0
          && v.back().dot(p, v[0]) >= 0;
   int L = 1, R = (int)v.size() - 1, ans = 1;
   while(L <= R){</pre>
      int mid = (L+R)/2;
      if(v[0].ccw(v[mid], p) >= 0) ans = mid, L = mid
         +1:
       else R = mid-1;
   return v[ans].ccw(v[(ans+1)%v.size()], p) >= 0;
Minkowski sum
vector<vec> mk(const vector<vec>&a,const vector<vec>&b){
   int i = 0, j = 0;
   for(int k = 0; k < (int)a.size(); k++)if(a[k] < a[i])
   for(int k = 0; k < (int)b.size(); k++)if(b[k] < b[j])
       j = k;
   vector<vec> c:
   c.reserve(a.size() + b.size());
   for(int k = 0; k < int(a.size()+b.size()); k++){
       vec pt{a[i] + b[j]};
       if((int)c.size() >= 2
       && c[c.size()-2].ccw(c.back(), pt) == 0)
          c.pop_back();
       c.push_back(pt);
      int q = i+1, w = j+1;
      if(q == int(a.size())) q = 0;
      if(w == int(b.size())) w = 0;
      if(c.back().ccw(a[i]+b[w], a[q]+b[j]) < 0) i = q;
      else j = w;
   c.shrink_to_fit();
   return c;
Geo Notes
```

**System of points(2D/3D):** Mass weighted average of points. Frame(2D/3D): Get middle point of each segment solve as previously.

**Triangle:** Average of vertices.

**2D Polygon:** Compute **signed** area and center of mass of triangle  $((0,0), p_i, p_{i+1})$ . Then solve as system of points.

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**Polyhedron surface:** Solve each face as a 2D polygon(be aware of (0, 0)) then replace each face with its center of mass and solve as system of points.

**Tetrahedron(Triangular pyramid):** As triangles, its the average of points.

**Polyhedron:** Can be done as 2D polygon, but with tetrahedralization intead of triangulation.

### Pick's Theorem

Given a polygon without self-intersections and all its vertices on integer coordinates in some 2D grid. Let A be its area, I the number of points with interger coordinates stricly inside the polygon and B the number of points with interger coordinates in the border of the polygon. The following formula holds:  $A = I + \frac{B}{2} - 1$ .

# Miscellaneous

# LIS

```
multiset<int> S;
for(int i = 0; i < n; i++){
   auto it = S.upper_bound(a[i]); // low for inc
   if(it != S.end()) S.erase(it);
   S.insert(a[i]);
}
ans = S.size();
DSU rollback
struct DSU{
   vector<int> sz, p, change;
   vector<tuple<int, int, int>> modifications;
   vector<size_t> saves;
   bool bipartite;
   DSU(int n): sz(n+1, 1), p(n+1), change(n+1),
     bipartite(true){
      iota(p.begin(), p.end(), 0);
   void add_edge(int u, int v){
      if(!bipartite) return;
      int must_change = get_colour(u) == get_colour(v);
      int a = rep(u), b = rep(v);
      if(sz[a] < sz[b]) swap(a, b);
      if(a != b){
          p[b] = a;
          modifications.emplace_back(b, change[b],
            bipartite);
          change[b] ^= must_change;
          sz[a] += sz[b];
      else if(must_change){
          modifications.emplace_back(0, change[0],
            bipartite);
          bipartite = false;
      }
   }
   int rep(int u){
      return p[u] == u ? u : rep(p[u]);
   }
   int get_colour(int u){
      if(p[u] == u) return change[u];
```

```
return change[u] ^ get_colour(p[u]);
   }
   void reset(){
       modifications.clear();
       saves.clear();
       iota(p.begin(), p.end(), 0);
       fill(sz.begin(), sz.end(), 1);
       fill(change.begin(), change.end(), 0);
       bipartite = true;
   }
   void rollback(){
       int u = get<0>(modifications.back());
       tie(ignore, change[u], bipartite) = modifications
         .back();
       sz[ p[u] ] -= sz[u];
      p[u] = u;
       modifications.pop_back();
   }
   void reload(){
       while(modifications.size() > saves.back())
          rollback();
       saves.pop_back();
   }
   void save(){
       saves.push_back(modifications.size());
};
Buildings
// count the number of circular arrays of size m, with
  elements on range [1, c**(n*n)]
int n, m, c; cin >> n >> m >> c;
int x = f_{exp}(c, n * n); int ans = f_{exp}(x, m);
for(int i = 1; i \le m; i++) if(m \% i == 0) {
 int y = f_exp(x, i);
  for(int j = 1; j < i; j++) if(i \% j == 0)
     y = sub(y, mult(j, dp[j]));
 dp[i] = mult(y, inv(i));
 ans = sub(ans, mult(i - 1, dp[i]));
cout \ll ans \ll '\n';
Rand
#include <random>
#include <chrono>
cout << RAND_MAX << endl;</pre>
mt19937 rng(chrono::steady_clock::now().time_since_epoch
  ().count());
vector<int> permutation(N);
iota(permutation.begin(), permutation.end(), 0);
shuffle(permutation.begin(), permutation.end(), rng);
iota(permutation.begin(), permutation.end(), 0);
for(int i = 1; i < N; i++){
   swap(permutation[i], permutation[
     uniform_int_distribution<int>(0, i)(rng)]);
}
Klondike
// minimum number of moves to make
// all elements equal
// move: change a segment of equal value
// elements to any value
```

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```
Slope Trick
int v[305], dp[305][305], rec[305][305];
                                                           ///By wogja125, contest: Codeforces Round #371 (Div. 1),
int f(int 1, int r){
                                                              problem: (C) Sonya and Problem Wihtout a Legend,
 if(r == 1) return 1;
                                                             Accepted. #
 if(r < 1) return 0;</pre>
                                                           int main() {
 if(dp[l][r] != -1) return dp[l][r];
                                                              int n, t; long long ans = 0; priority_queue<int> Q;
 int ans = f(1+1, r) + 1;
                                                              scanf("%d%d", &n, &t); Q.push(t);
 for(int i = l+1; i \le r; i++)
                                                              for(int i = 1; i < n; i++) {
   if(v[i] == v[1])
                                                                  scanf("%d", &t); t -= i; Q.push(t);
     ans = min(ans, f(1, i - 1) + f(i+1, r));
                                                                  if(Q.top() > t) {
 return dp[l][r] = ans;
                                                                     ans += Q.top() - t; Q.pop(); Q.push(t);
Hilbert Order
                                                              printf("%lld", ans);
// maybe use B = n / sqrt(q)
inline int64_t hilbertOrder(int x, int y, int pow = 21,
                                                           Knapsack Bounded with Cost
  int rotate = 0) {
   if(pow == 0) return 0;
   int hpow = 1 \ll (pow-1);
                                                           // menor custo para conseguir peso ate M usando N tipos
                                                             diferentes de elementos, sendo que o i-esimo elemento
   int seg = (x < hpow) ? (
                                                             pode ser usado b[i] vezes, tem peso w[i] e custo c[i]
       (y < hpow) ? 0 : 3
                                                           // O(N * M)
   ):(
       (y < hpow) ? 1 : 2
                                                           int b[N], w[N], c[N];
                                                           MinOueue O[M]
   seg = (seg + rotate) & 3;
                                                           int d[M] //d[i] = custo minimo para conseguir peso i
   const int rotateDelta[4] = {3, 0, 0, 1};
   int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
   int nrot = (rotate + rotateDelta[seg]) & 3;
                                                           for(int i = 0; i \le M; i++) d[i] = i ? oo : 0;
                                                           for(int i = 0; i < N; i++){
   int64_t subSquareSize = int64_t(1) << (2*pow - 2);</pre>
                                                              for(int j = 0; j < w[i]; j++)
   int64_t ans = seg * subSquareSize;
                                                                  Q[j].clear();
   int64_t add = hilbertOrder(nx, ny, pow-1, nrot);
   ans += (seg == 1 || seg == 2) ? add : (subSquareSize
                                                              for(int j = 0; j \le M; j++){
                                                                  q = Q[j \% w[i]];
      - add - 1):
   return ans;
                                                                  if(q.size() >= q) q.pop();
                                                                  q.add(c[i]);
                                                                  q.push(d[j]);
Modular Factorial
                                                                  d[j] = q.getmin();
// Compute (1*2*...*(p-1)*1*(p+1)*(p+2)*..*n) % p
// in O(p*lg(n))
                                                           LCA < O(nlgn), O(1) >
int factmod(int n, int p){
   int ans = 1;
                                                           int start[N], dfs_time;
   while (n > 1) {
                                                           int tour[2*N], id[2*N];
      for(int i = 2; i \le n \% p; i++)
          ans = (ans * i) % p;
                                                           void dfs(int u){
      n /= p;
      if(n \% 2) ans = p - ans;
                                                              start[u] = dfs_time;
   }
                                                              id[dfs_time] = u;
                                                              tour[dfs_time++] = start[u];
   return ans % p;
                                                              for(int v : g[u]){
int fac_pow(int n, int p){
                                                                  dfs(v);
   int ans = 0:
                                                                  id[dfs_time] = u;
   while(n) n \neq p, ans += n;
                                                                  tour[dfs_time++] = start[u];
   return ans;
                                                              }
                                                           }
int C(int n, int k, int p){
   if(fac_pow(n, p) > fac_pow(n-k, p) + fac_pow(k, p))
                                                           int LCA(int u, int v){
                                                              if(start[u] > start[v]) swap(u, v);
      return 0;
   int tmp = factmod(k, p) * factmod(n-k, p) % p;
                                                              return id[min(tour[k]for k in [start[u],start[v]])];
   return (f_{exp}(tmp, p - 2, p) * factmod(n, p)) % p;
                                                           Buffered reader
Enumeration all submasks of a bitmask
                                                           // source: https://github.com/ngthanhtrung23/
// loop through all submask of a given bitmask
                                                             ACM_Notebook_new/blob/master/buffered_reader.h
// it does not include mask 0
                                                           int INP.AM.REACHEOF;
for(int sub = mask; sub; sub = (sub-1)&mask){
                                                           #define BUFSIZE (1<<12)</pre>
                                                           char BUF[BUFSIZE+1], *inp=BUF;
                                                           #define GETCHAR(INP) { \
```

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```
if(!*inp && !REACHEOF) { \
       memset(BUF,0,sizeof BUF);\
       int inpzzz = fread(BUF,1,BUFSIZE,stdin);\
       if (inpzzz != BUFSIZE) REACHEOF = true;\
       inp=BUF; \
   } \
   INP=*inp++; \
#define DIG(a) (((a)>='0')&&((a)<='9'))
#define GN(j) { \
   AM=0; \
   GETCHAR(INP); while(!DIG(INP) && INP!='-') GETCHAR(
   if (INP=='-') {AM=1;GETCHAR(INP);} \
   j=INP-'0'; GETCHAR(INP); \
   while(DIG(INP)){j=10*j+(INP-'0');GETCHAR(INP);} \
   if (AM) j=-j;\
Modular summation
//calcula (sum(0 \le i \le n) P(i)) % mod,
//onde P(i) eh uma PA modular (com outro modulo)
namespace sum_pa_mod{
   11 calc(11 a, 11 b, 11 n, 11 mod){
       assert(a&&b);
       if(a >= b){
          ll ret = ((n*(n+1)/2) \mod)*(a/b);
          if(a\%b) ret = (ret + calc(a%b,b,n,mod))%mod;
          else ret = (ret+n+1) mod;
          return ret;
       return ((n+1)*(((n*a)/b+1)%mod) - calc(b,a,(n*a)/b+1)%mod)
         b, mod) + mod + n/b + 1)%mod;
   }
   //P(i) = a*i \mod m
   11 solve(l1 a, l1 n, l1 m, l1 mod){
       a = (a\%m + m)\%m;
       if(!a) return 0;
       11 \text{ ret} = (n*(n+1)/2) \% mod;
       ret = (ret*a)%mod;
       ll g = \_gcd(a,m);
       ret -= m*(calc(a/g,m/g,n,mod)-n-1);
       return (ret%mod + mod)%mod;
   }
   //P(i) = a + r*i \mod m
   11 solve(l1 a, l1 r, l1 n, l1 m, l1 mod){
       a = (a\%m + m)\%m;
       r = (r\%m + m)\%m;
       if(!r) return (a*(n+1))%mod;
       if(!a) return solve(r, n, m, mod);
       11 g, x, y;
       g = gcdExtended(r, m, x, y);
       x = (x\%m + m)\%m;
       11 d = a - (a/g)*g;
       a -= d;
       x = (x*(a/g))%m;
       return (solve(r, n+x, m, mod) - solve(r, x-1, m,
         mod) + mod + d*(n+1))%mod;
   }
Edge coloring CPP
const int MX = 300;
int C[MX][MX] = {}, G[MX][MX] = {};
```

```
void solve(vector<pii> &E, int N){
   int X[MX] = \{\}, a, b;
   auto update = [&](int u){ for(X[u] = 1; C[u][X[u]];
     X[u]++); };
   auto color = [&](int u, int v, int c){
       int p = G[u][v];
       G[u][v] = G[v][u] = c;
       C[u][c] = v; C[v][c] = u;
       C[u][p] = C[v][p] = 0;
       if(p) X[u] = X[v] = p;
       else update(u), update(v);
       return p; };
   auto flip = [&](int u, int c1, int c2){
       int p = C[u][c1], q = C[u][c2];
       swap(C[u][c1], C[u][c2]);
       if(p) G[u][p] = G[p][u] = c2;
       if( !C[u][c1] ) X[u] = c1;
       if( !C[u][c2] ) X[u] = c2;
       return p; };
   for(int i = 1; i <= N; i++) X[i] = 1;</pre>
   for(int t = 0; t < E.size(); t++){</pre>
       int u = E[t].first, v0 = E[t].second, v = v0, c0
         = X[u], c = c0, d;
       vector<pii> L;
       int vst[MX] = {};
       while(!G[u][v0]){
          L.emplace_back(v, d = X[v]);
          if(!C[v][c]) for(a = (int)L.size()-1; a >= 0;
             a--) c = color(u, L[a].first, c);
          else if(!C[u][d])for(a=(int)L.size()-1;a>=0;a
             --)color(u,L[a].first,L[a].second);
          else if( vst[d] ) break;
          else vst[d] = 1, v = C[u][d];
       if( !G[u][v0] ){
          for(;v; v = flip(v, c, d), swap(c, d));
          if(C[u][c0]){
              for(a = (int)L.size()-2; a >= 0 \&\& L[a].
                second != c; a--);
              for(; a >= 0; a--) color(u, L[a].first, L[
                a].second);
          } else t--;
      }
   }
Burnside's Lemma
```

Let  $(G, \oplus)$  be a finite group that acts on a set X. It should hold that  $e_g * x = x$  and  $g_1 * (g_2 * x) = (g_1 \oplus g_2) * x$ ,  $\forall x \in X, g_1, g_2 \in G$ . For each  $g \in G$  let  $X^g = \{x \in X \mid g * x = x\}$ . The number of orbits its given by:

 $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$ 

# Wilson's Theorem

 $(n-1)! = -1 \mod n \iff n \text{ is prime}$ 

## Fibonacci

- $F_{n-1}F_{n+1} F_n^2 = (-1)^n$
- $\bullet \ F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$
- $GCD(F_n, F_m) = F_{GCD(n,m)}$

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• 
$$F_n = \frac{(\frac{1+\sqrt{5}}{2})^n - (\frac{1-\sqrt{5}}{2})^n}{\sqrt{5}}$$

# Lucas's Theorem

For non-negative integers m and n and a prime p, the following congruence holds:

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p}$$

where  $m_i$  is the i-th digit of m in base p.  $\binom{a}{b} = 0$  if a < b.

# Kirchhoff's Theorem

Laplacian matrix is L = D - A, where D is a diagonal matrix with vertex degrees on the diagonals and A is adjacency matrix.

The number of spanning trees is any cofactor of L. i-th cofactor is determinant of the matrix gotten by removing i-th row and column of L.

# Multigraphs

In D[i][i] all loops are excluded. A[i][j] = number of edges from i to j.

# Directed multigraphs

D[i][i] = indegree of i minus the number of loops at i. A[i][j] = number of edges from i to j.

The number of oriented spanning trees rooted at a vertex i is the determinant of the matrix gotten by removing the ith row and column of L.

### Matroid

Let *X* set of objects,  $I \subseteq 2^X$  set of independents sets such that:

- 1.  $\emptyset \in I$
- 2.  $A \in I, B \subseteq A \implies B \in I$
- 3. Exchange axiom,  $A \in I, B \in I, |B| > |A| \implies \exists x \in B \setminus A : A \cup \{x\} \in I$
- 4.  $A \subseteq X$  and I and I' are maximal independent subsets of A then |I| = |I'|

Then (X, I) is a matroid. The combinatorial optimization problem associated with it is: Given a weight  $w(e) \ge 0 \ \forall e \in X$ , find an independet subset that has the largest possible total weight.

# Matroid intersection

```
// Input two matroids (X, I_a) and (X, I_b)
// output set I of maximum size, I \in I_a and I \in I_b
set<> I;
while(1){
    for(e_i : X \ I)
        if(I + e_i \in I_a and I + e_i \in I_b)
        I = I + e_i;
    set<> A, T; queue<> Q;
    for(x : X) label[x] = MARK1;
    for(e_i : X \ I){
        if(I + e_i \in I_a)
```

```
Q.push(e_i), label[e_i] = MARK2;
   else{
       for(x such that I - x + e_i \in I_a)
           A[x].push(e_i);
   if(I + e_i \setminus in I_b)
       T = T + \{e_i\}
   else{
       for(x such that I - x + e_i \in I_b)
           A[e_i].push(x);
   }
if(T.empty()) break;
bool found = false;
while(!Q.empty() and !found){
   auto e = Q.front(); Q.pop();
   for(x : A[e]) if(label[x] == MARK1){
       label[x] = e; Q.push(x);
       if(x \setminus in T){
           found = true; put = 1;
           while(label[x] != MARK2){
              I = put ? (I + x) : (I - x);
              put = 1 - put;
           I = I + x;
           break:
       }
   }
if(!found) break;
```

return I;

Where path(e) = [e] if label[e] = MARK2, path(label[e]) + [e] otherwise.

# Matroid Union

Given k matroids over the same set of objects  $(X, I_1)$ ,  $(X, I_2)$ , ...,  $(X, I_k)$  find  $A_1 \in I_1$ ,  $A_2 \in I_2$ , ...,  $A_k \in I_k$  such that  $i \neq j, A_i \cap A_j = \emptyset$  and  $|\bigcup_{i=1}^k A_i|$  is maximum. Matroid union can be reduced to matroid intersection as follows.

Let  $X' = X \times \{1, 2, ..., k\}$ , ie, k copies of each element of X with different colors. M1 = (X', Q) where  $B \in Q \iff \forall 1 \le i \le k$ ,  $\{x \mid (x, i) \in B\} \in I_i$ , ie, for each color, B is independent. M2 = (X', W) where  $B \in W \iff i \ne j \implies \neg((x, i) \in B) \land (x, j) \in B$ ), ie, each element is picked by at most one color.

Intersection of *M*1 and *M*2 is the answer for the combinatorial problem of matroid union.

### **Notes**

When we repeat something and each time we have probability p to succeed then the expected number or tries is  $\frac{1}{p}$ , till we succeed.

# Small to large

**Trick in statement** If k sets are given you should note that the amount of different set sizes is  $O(\sqrt{s})$  where s is total size of those sets. And no more than  $\sqrt{s}$  sets have size greater than  $\sqrt{s}$ . For example, a path to the root in Aho-Corasick through suffix links will have at most  $O(\sqrt{s})$  vertices.

**gcd on subsegment**, we have at most  $log(a_i)$  different values in  $\{gcd(a_i, a_{i+1}, ..., a_i) \text{ for } j < i\}$ .

**From static set to expandable**. To insert, create a new set with the new element. While there are two sets with same size, merge them. There will be at most  $\log(n)$  disjoints sets.

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**Matrix exponentiation optimization**. Save binary power of  $A_{nxn}$  and answer q queries  $b = A^m x$  in  $O((n^3 + qn^2)log(m))$ .

**Ternary search on integers into binary search**, comparing f(mid) and f(mid+1), binary search on derivative

**Dynamic offline set** For each element we will wind segment of time [a,b] such that element is present in the set during this whole segment. Now we can come up with recursive procedure which handles [l,r] time segment considering that all elements such that  $[l,r] \subset [a,b]$  are already included into

the set. Now, keeping this invariant we recursively go into [l, m] and [m + 1, r] subsegments. Finally when we come into segment of length 1.

$$a > b \implies a \mod b < \frac{a}{2}$$

**Convex Hull**. The expected number of points in the convex hull of a random set of points is O(log(n)). The number of points in a convex hull with points coordinates limited by L is  $O(L^{2/3})$ .