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Contents			4	4 Strings		
				4.1		14
1	Data Structures	2		4.2	Suffix Array	
1		- 1		4.3	Z Algorithm	
	O	$\frac{2}{2}$		4.4	Prefix function/KMP	
		2		4.5		15
		2		4.6	Manacher	15
		3		4.7		16
	1	3		4.8	Suffix Tree	16
	1.6 Convex Hull Trick	3	_	C		10
	1.7 Min queue	3	5		metry 2D basis	16
	1.8 Sparse Table	4		5.1		16
	1.9 Treap	4		5.2	Circle line intersection	18
		4		5.3	Half plane intersection	18
	<u> </u>	5		5.4	Detect empty Half plane intersection	19
		5		5.5	Circle Circle intersection	19
	· ·	5		5.6	Tangents of two circles	19
	<u> </u>	5		5.7	Charles a intrinside a classes	19
	1.14 Tamidionne dee	٦		5.8	Check point inside polygon	20
2	Math	6		5.9	Check point inside polygon without lower/up-	20
_		6		E 10	per hull	20
	0					20 20
		6		3.11	Geo Notes	
		6			5.11.1 Center of mass	20
		6			5.11.2 Fick's Theorem	20
		6	6	Miso	cellaneous	20
		7	Ů	6.1	LIS	
		7		6.2	DSU rollback	
		7		6.3	Buildings	21
	2.9 Mulmod TOP	7		6.4	Rand	21
	2.10 Matrix Determinant	7		6.5	Klondike	21
	2.11 Simplex Method	7		6.6	Hilbert Order	21
	2.12 FFT	8		6.7	Modular Factorial	22
	2.13 FFT Tourist	8		6.8	Enumeration all submasks of a bitmask	22
		9		6.9	Slope Trick	
	2.15 Gauss	`. I			Knapsack Bounded with Cost	
	2.16 Gauss Xor	- 1			LCA <o(nlgn), o(1)=""></o(nlgn),>	
	2.17 Simpson	- 1			Buffered reader	22
	2.17 Shirpson			6.13	Modular summation	22
3	Graphs 10	ا ۵		6.14	Edge coloring CPP	23
J	3.1 Dinic				Burnside's Lemma	23
	3.2 Push relabel	· [6.16	Wilson's Theorem	23
		- 1			Fibonacci	23
	3.3 Min Cost Max Flow	- 1		6.18	Kirchhoff's Theorem	23
	3.4 Blossom Algorithm for General Matching 12	- 1			6.18.1 Multigraphs	23
	3.5 Small to Large	- 1			6.18.2 Directed multigraphs	23
	3.6 Centroid Decomposition	- 1		6.19	Matroid	24
	3.7 Kosaraju	3			6.19.1 Matroid intersection	24
	3.8 Tarjan	3			6.19.2 Matroid Union	24
	3.9 Max Clique	3		6.20	Notes	24
	3.10 Dominator Tree	3				
	3.11 Min Cost Matching	4				

University of Brasilia Data Structures, 2

```
set ts=4 sw=4 sta nu rnu sc stl+=%F cindent
set bg=dark ruler clipboard=unnamed,unnamedplus
  timeoutlen=100
imap {<CR> {<CR>}}<Esc>0
nmap <F2> 0V$%d
nmap <C-down> :m+1<CR>
nmap <C-up> :m-2<CR>
vmap < C-c > "+y
nmap <C-a> ggVG
syntax on
alias cmp='g++ -Wall -Wformat=2 -Wshadow -Wconversion -
  fsanitize=address -fsanitize=undefined -fno-sanitize-
  recover -std=c++14'
```

Data Structures

Merge Sort Tree

```
struct MergeTree{
   int n:
   vector<vector<int>>> st;
   void build(int p, int L, int R, const int v[]){
          st[p].push_back(v[L]);
          return:
      int mid = (L+R)/2;
      build(2*p, L, mid, v);
      build(2*p+1, mid+1, R, v);
       st[p].resize(R-L+1);
      merge(st[2*p].begin(), st[2*p].end(),
              st[2*p+1].begin(), st[2*p+1].end(),
              st[p].begin());
   }
   int query(int p, int L, int R, int i, int j, int x)
     const{
       if(L > j \mid \mid R < i) return 0;
      if(L >= i && R <= j){
          int id = lower_bound(st[p].begin(), st[p].end
             (), x) - st[p].begin();
          return int(st[p].size()) - id;
      int mid = (L+R)/2;
      return query(2*p, L, mid, i, j, x) +
          query(2*p+1, mid+1, R, i, j, x);
   }
public:
   MergeTree(int sz, const int v[]): n(sz), st(4*sz){
      build(1, 1, n, v);
   //number of elements >= x on segment [i, j]
   int query(int i, int j, int x) const{
      if(i > j) swap(i, j);
      return query(1, 1, n, i, j, x);
   }
};
```

Wavelet Tree

```
template<typename T>
class wavelet{
   T L, R;
   vector<int> 1;
   vector<T> sum; // <<</pre>
```

```
wavelet *lef, *rig;
   int r(int i) const{ return i - l[i]; }
public:
   template<typename ITER>
   wavelet(ITER bg, ITER en){
       lef = rig = nullptr;
       L = *bg, R = *bg;
       for(auto it = bg; it != en; it++)
          L = min(L, *it), R = max(R, *it);
       if(L == R) return;
       T mid = L + (R - L)/2;
       1.reserve(std::distance(bg, en) + 1);
       sum.reserve(std::distance(bg, en) + 1);
       1.push_back(0), sum.push_back(0);
       for(auto it = bg; it != en; it++)
          1.push_back(1.back() + (*it <= mid)),</pre>
          sum.push_back(sum.back() + *it);
       auto tmp = stable_partition(bg, en, [mid](T x){
          return x <= mid;</pre>
       });
       if(bg != tmp) lef = new wavelet(bg, tmp);
       if(tmp != en) rig = new wavelet(tmp, en);
    ~wavelet(){
       delete lef;
       delete rig;
    // 1 index, first is 1st
   T kth(int i, int j, int k) const{
       if(L >= R) return L;
       int c = 1[j] - 1[i-1];
       if(c \ge k) return lef \ge kth(l[i-1]+1, l[j], k);
       else return rig->kth(r(i-1)+1, r(j), k - c);
   }
   // # elements > x on [i, j]
   int cnt(int i, int j, T x) const{
       if(L > x) return j - i + 1;
       if(R <= x || L == R) return 0;
       int ans = 0;
       if(lef) ans += lef->cnt(l[i-1]+1, l[j], x);
       if(rig) ans += rig->cnt(r(i-1)+1, r(j), x);
       return ans;
   // sum of elements <= k on [i, j]</pre>
   T sumk(int i, int j, T k){
       if(L == R) return R <= k ? L * (j - i + 1) : 0;
       if(R <= k) return sum[j] - sum[i-1];</pre>
       int ans = 0;
       if(lef) ans += lef->sumk(l[i-1]+1, l[j], k);
       if(rig) ans += rig->sumk(r(i-1)+1, r(j), k);
      return ans;
   // swap (i, i+1) just need to update "array" l[i]
}:
Order Set
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/detail/standard_policies.hpp>
```

University of Brasilia Data Structures, 3

```
using namespace __gnu_pbds; // or pb_ds;
                                                                   int mid = (L+R)/2;
                                                                   if(cht[mid].eval(x) >= cht[mid+1].eval(x)) // <<<</pre>
template<typename T, typename B = null_type>
                                                                       L = mid + 1:
using oset = tree<T, B, less<T>, rb_tree_tag,
                                                                   else bans = mid, R = mid - 1;
  tree_order_statistics_node_update>;
// find_by_order / order_of_key
                                                               return cht[bans].eval(x);
Hash table
                                                            Convex Hull Trick
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
                                                             * Author: Simon Lindholm
                                                             * source: https://github.com/kth-competitive-
struct custom_hash {
                                                               programming/kactl/blob/master/content/data-structures
   static uint64_t splitmix64(uint64_t x) {
       // http://xorshift.di.unimi.it/splitmix64.c
                                                               /LineContainer.h
      x += 0x9e3779b97f4a7c15;
                                                             * License: CC0
                                                             */
      x = (x \hat{ } (x >> 30)) * 0xbf58476d1ce4e5b9;
       x = (x \hat{ } (x >> 27)) * 0x94d049bb133111eb;
      return x \hat{ } (x \gg 31);
                                                            struct Line {
                                                               mutable 11 m, b, p;
   }
                                                               bool operator<(const Line& o) const { return m < o.m</pre>
   size_t operator()(uint64_t x) const {
      static const uint64_t FIXED_RANDOM = chrono::
                                                               bool operator<(11 x) const { return p < x; }</pre>
         steady_clock::now().time_since_epoch().count();
                                                            };
      return splitmix64(x + FIXED_RANDOM);
   }
                                                            struct LineContainer : multiset<Line, less<>> { // CPP14
};
                                                                // (for doubles, use inf = 1/.0, div(a,b) = a/b)
gp_hash_table<long long, int, custom_hash> table;
                                                                const ll inf = LLONG_MAX;
unordered_map<long long, int, custom_hash> uhash;
                                                                11 div(ll a, ll b) { // floored division
                                                                   return a / b - ((a ^ b) < 0 && a % b); }
uhash.reserve(1 << 15);</pre>
uhash.max_load_factor(0.25);
                                                               bool isect(iterator x, iterator y) {
                                                                   if (y == end()) { x->p = inf; return false; }
Convex Hull Trick Simple
                                                                   if (x->m == y->m) x->p = x->b > y->b? inf : -inf
struct Line{
                                                                   else x->p = div(y->b - x->b, x->m - y->m);
                                                                   return x->p >= y->p;
   inline 11 eval(11 x) const{
      return x * m + b;
                                                               void add(ll m, ll b) {
   }
                                                                   auto z = insert(\{m, b, 0\}), y = z++, x = y;
};
                                                                   while (isect(y, z)) z = erase(z);
                                                                   if (x != begin() \&\& isect(--x, y)) isect(x, y =
// min => cht.back().m >= L.m
                                                                     erase(y));
// max => cht.back().m <= L.m
                                                                   while ((y = x) != begin() && (--x)->p >= y->p)
void push_line(vector<Line> &cht, Line L){
                                                                       isect(x, erase(y));
 while((int)cht.size() >= 2){
                                                               }
   int sz = (int)cht.size();
                                                               11 query(ll x) {
   if((long double)(L.b-cht[sz-1].b)*(cht[sz-2].m-L.m)
                                                                   assert(!empty());
  <= (long double)(L.b-cht[sz-2].b)*(cht[sz-1].m-L.m)){</pre>
                                                                   auto 1 = *lower_bound(x);
     cht.pop_back();
                                                                   return 1.m * x + 1.b;
   }
   else break;
                                                            };
 3
                                                            Min queue
 cht.push_back(L);
                                                            template<typename T>
// x increasing; pos = 0 in first call
                                                            class minQ{
11 linear_search(const vector<Line> &cht,ll x,int &pos){
                                                                deque<tuple<T, int, int> > p;
   while(pos+1 < (int)cht.size()){</pre>
                                                               T delta;
/*>>*/ if(cht[pos].eval(x) >= cht[pos+1].eval(x)) pos++;
                                                               int sz;
      else break;
                                                            public:
                                                               minQ() : delta(0), sz(0) {}
                                                               inline int size() const{ return sz; }
   return cht[pos].eval(x);
}
                                                               inline void add(T x){ delta += x; }
                                                               inline void push(T x, int id){
11 binary_search(const vector<Line> &cht, ll x){
                                                                   x \rightarrow delta, sz++;
   int L = 0, R = (int)cht.size()-2;
                                                                   int t = 1;
   int bans = (int)cht.size()-1;
                                                                   while(p.size() > 0 \& get<0>(p.back()) >= x)
   while(L <= R){</pre>
                                                                       t += get<1>(p.back()), p.pop_back();
```

University of Brasilia Data Structures, 4

```
p.emplace_back(x, t, id);
                                                               unlaze(u);
   }
                                                               if(!u) u = it;
                                                               else if(Y[it] > Y[u]) split_val(u, X[it], L[it], R[
   inline void pop(){
       get<1>(p.front())--, sz--;
                                                                  it]), u = it;
       if(!get<1>(p.front())) p.pop_front();
                                                                else insert(X[it] < X[u] ? L[u] : R[u], it);
                                                                calc(u):
   T getmin() const{ return get<0>(p.front())+delta; }
   int getid() const{ return get<2>(p.front()); }
                                                            void erase(int &u, num key){
                                                               unlaze(u);
                                                                if(!u) return;
Sparse Table
                                                               if(X[u] == key) u = merge(L[u], R[u]);
                                                               else erase(key < X[u] ? L[u] : R[u], key);</pre>
int fn(int i, int j){
                                                                calc(u);
   if(j == 0) return v[i];
   if(~dn[i][j]) return dn[i][j];
                                                            int create_node(num key){
   return dn[i][j] = min(fn(i, j-1), fn(i + (1 << (j-1)))
                                                               X[en] = key;
     ), i-1));
                                                               sz[en] = 1;
                                                               L[en] = R[en] = 0;
                                                               return en++;
int getmn(int 1, int r){ // [1, r]
   int 1z = 1g(r - 1 + 1);
                                                            int query(int u, int 1, int r){//0 index
   return min(fn(1, 1z), fn(r - (1 << 1z) + 1, 1z));
                                                               unlaze(u):
                                                                if(u! or r < 0 or 1 >= sz[u]) return
Treap
                                                                 identity_element;
                                                               if(1 \le 0 \text{ and } r \ge sz[u] - 1) \text{ return } subt\_data[u];
// source: https://github.com/victorsenam/caderno/blob/
                                                               int ans = query(L[u], 1, r);
  master/code/treap.cpp
                                                                if(1 \le sz[L[u]] and sz[L[u]] \le r)
                                                                   ans = max(ans, st[u]);
//const int N = ; typedef int num;
num X[N]; int en = 1, Y[N], sz[N], L[N], R[N];
                                                                ans = max(ans, query(R[u], l-sz[L[u]]-1, r-sz[L[u]])
void calc (int u) { // update node given children info
                                                                 11-1)):
   if(!u) return;
                                                               return ans;
   sz[u] = sz[L[u]] + 1 + sz[R[u]];
   // code here, no recursion
                                                            ColorUpdate
void unlaze (int u) {
                                                            // source: https://github.com/tfg50/Competitive-
   if(!u) return;
   // code here, no recursion
                                                              Programming/tree/master/Biblioteca/Data%20Structures
void split_val(int u, num x, int &l, int &r) { // l gets
                                                            #include <set>
                                                            #include <vector>
   <= x, r gets > x
   unlaze(u); if(!u) return (void) (1 = r = 0);
                                                            template <class Info = int>
   if(X[u] \le x) \{ split_val(R[u], x, 1, r); R[u] = 1;
                                                            class ColorUpdate {
     1 = u; }
   else { split_val(L[u], x, 1, r); L[u] = r; r = u; }
                                                            public:
   calc(u);
                                                                struct Range {
                                                                   Range(int 1 = 0) { this->1 = 1; }
void split_sz(int u, int s, int &l, int &r) { // l gets
                                                                   Range(int 1, int r, Info v) {
  first s, r gets remaining
                                                                      this->1 = 1:
   unlaze(u); if(!u) return (void) (1 = r = 0);
                                                                       this -> r = r;
   if(sz[L[u]] < s)  { split_sz(R[u], s - sz[L[u]] - 1,
                                                                      this->v = v;
     1, r); R[u] = 1; 1 = u; }
   else { split_sz(L[u], s, l, r); L[u] = r; r = u; }
                                                                   int 1, r;
   calc(u);
                                                                   Info v;
int merge(int 1, int r) { // els on 1 <= els on r</pre>
                                                                   bool operator < (const Range &b) const { return 1</pre>
   unlaze(l); unlaze(r); if(!l || !r) return l + r; int
                                                                      < b.1; }
                                                               };
   if(Y[1] > Y[r]) { R[1] = merge(R[1], r); u = 1; }
   else { L[r] = merge(1, L[r]); u = r; }
                                                                std::vector<Range> upd(int 1, int r, Info v) {
   calc(u); return u;
                                                                   std::vector<Range> ans;
                                                                   if(1 >= r) return ans;
void init(int n=N-1) { // XXX call before using other
                                                                   auto it = ranges.lower_bound(1);
                                                                   if(it != ranges.begin()) {
  funcs
   for(int i = en = 1; i \le n; i++) { Y[i] = i; sz[i] =
                                                                      it--;
      1; L[i] = R[i] = 0; }
                                                                       if(it->r>1) {
   random\_shuffle(Y + 1, Y + n + 1);
                                                                          auto cur = *it;
                                                                          ranges.erase(it);
void insert(int &u, int it){
                                                                          ranges.insert(Range(cur.1, 1, cur.v));
```

University of Brasilia Data Structures, 5

```
ranges.insert(Range(1, cur.r, cur.v));
          }
                                                           void update(int p, T val){
      }
                                                               st[p += n] = val;
      it = ranges.lower_bound(r);
                                                               while(p >>= 1) st[p] = merge(st[2*p], st[2*p+1]);
      if(it != ranges.begin()) {
          it--;
                                                           LiChao's Segtree
          if(it->r>r) {
              auto cur = *it;
                                                           void add_line(line nw, int v = 1, int l = 0, int r = 1
              ranges.erase(it);
                                                             maxn) { // [1, r)}
              ranges.insert(Range(cur.1, r, cur.v));
                                                               int m = (1 + r) / 2;
              ranges.insert(Range(r, cur.r, cur.v));
                                                               bool lef = nw.eval(1) < st[v].eval(1);</pre>
                                                               bool mid = nw.eval(m) < st[v].eval(m);</pre>
                                                               if(mid) swap(st[v], nw);
       for(it = ranges.lower_bound(1); it != ranges.end
                                                               if(r - 1 == 1) {
         () && it->l < r; it++) {
                                                                  return:
          ans.push_back(*it);
                                                               } else if(lef != mid) {
                                                                  add_line(nw, 2 * v, 1, m);
      ranges.erase(ranges.lower\_bound(1), \ ranges.
                                                               } else {
        lower_bound(r));
                                                                  add_line(nw, 2 * v + 1, m, r);
       ranges.insert(Range(1, r, v));
                                                               }
       return ans:
                                                           }
   }
private:
                                                           int get(int x, int v = 1, int l = 0, int r = maxn) {
   std::set<Range> ranges;
                                                               int m = (1 + r) / 2;
                                                               if(r - 1 == 1) {
Heavy Light Decomposition
                                                                  return st[v].eval(x);
                                                               } else if(x < m) {
void dfs_sz(int u){
                                                                  return min(st[v].eval(x), get(x, 2*v, 1, m));
   sz[u] = 1;
   for(auto &v : g[u]) if(v == p[u]){
                                                                  return min(st[v].eval(x), get(x, 2*v+1, m, r));
       swap(v, g[u].back()); g[u].pop_back();
      break;
                                                           Palindromic tree
   for(auto &v : g[u]){
      p[v] = u; dfs_sz(v); sz[u] += sz[v];
                                                           #include <bits/stdc++.h>
      if(sz[v] > sz[g[u][0]])
          swap(v, g[u][0]);
                                                           using namespace std;
   }
                                                           const int maxn = 3e5 + 1, sigma = 26;
// nxt[u] = start of path with u
                                                           int len[maxn], link[maxn], to[maxn][sigma];
// set nxt[root] = root beforehand
                                                           int slink[maxn], diff[maxn], series_ans[maxn];
void dfs_hld(int u){
                                                           int sz, last, n;
   in[u] = t++;
                                                           char s[maxn];
   rin[in[u]] = u;
   for(auto v : g[u]){
                                                           void init()
      nxt[v] = (v == g[u][0] ? nxt[u] : v); dfs_hld(v);
                                                               s[n++] = -1;
   out[u] = t;
                                                               link[0] = 1;
                                                               len[1] = -1;
// subtree of u => [ in[u], out[u] )
                                                               sz = 2;
// path from nxt[u] to u \Rightarrow [in[nxt[u]], in[u]]
                                                           }
Iterative Segtree
                                                           int get_link(int v)
T query(int 1, int r){ // [1, r]
                                                           {
                                                               while(s[n - len[v] - 2] != s[n - 1]) v = link[v];
   T rl, rr;
   for(1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1){
                                                               return v;
      if(l & 1) rl = merge(rl, st[l++]);
      if(r & 1) rr = merge(st[--r], rr);
                                                           void add_letter(char c)
   return merge(rl, rr);
                                                               s[n++] = c -= 'a';
                                                               last = get_link(last);
                                                               if(!to[last][c])
// initially save v[i] in st[n+i] for all i in [0, n)
void build(){
   for(int p = n-1; p > 0; p--)
                                                                  len[sz] = len[last] + 2;
       st[p] = merge(st[2*p], st[2*p+1]);
                                                                  link[sz] = to[get_link(link[last])][c];
                                                                  diff[sz] = len[sz] - len[link[sz]];
}
```

```
if(diff[sz] == diff[link[sz]])
          slink[sz] = slink[link[sz]];
       else
          slink[sz] = link[sz];
       to[last][c] = sz++;
   last = to[last][c];
}
int main()
   ios::sync_with_stdio(0);
   cin.tie(0);
   init();
   string s;
   cin >> s;
   int n = s.size();
   int ans[n + 1];
   memset(ans, 63, sizeof(ans));
   ans[0] = 0;
   for(int i = 1; i \le n; i++)
      add_letter(s[i - 1]);
      for(int v = last; len[v] > 0; v = slink[v])
          series_ans[v] = ans[i - (len[slink[v]] + diff
            [v])];
          if(diff[v] == diff[link[v]])
              series_ans[v] = min(series_ans[v],
                series_ans[link[v]]);
          ans[i] = min(ans[i], series_ans[v] + 1);
       cout << ans[i] << "\n";</pre>
   return 0;
```

Math

inv[1] = 1;

for(int i = 2; i < p; i++)

Extended Euclidean Algorithm

```
// a*x + b*y = gcd(a, b), <gcd, x, y>
tuple<int, int, int> gcd(int a, int b) {
   if(b == 0) return make_tuple(a, 1, 0);
   int q, w, e;
   tie(q, w, e) = gcd(b, a % b);
   return make_tuple(q, e, w - e * (a / b));
}
```

Chinese Remainder Theorem

```
// x = vet[i].first (mod vet[i].second)
ll crt(vector<pair<ll, ll>> vet){
    ll ans = vet[0].first, lcm = vet[0].second;
    ll a, b, g, x, y;
    for(int i = 1; i < (int)vet.size(); i++){
        tie(a, b) = vet[i];
        tie(g, x, y) = gcd(lcm, b);
        ans = ans + x * (a - ans) / g % (b / g) * lcm;
        lcm = lcm * b / g;
        ans = (ans % lcm + lcm) % lcm;
    }
    return ans;
}
Preffix inverse</pre>
```

```
inv[i] = (p - (p/i) * inv[p%i] % p) % p;
Pollard Rho
11 rho(11 n){
   if(n \% 2 == 0) return 2;
   11 d, c, x, y;
       c = llrand() % n, x = llrand() % n, y = x;
          x = add(mul(x, x, n), c, n);
          y = add(mul(y, y, n), c, n);
          y = add(mul(y, y, n), c, n);
          d = \_gcd(abs(x - y), n);
       \}while(d == 1);
   }while(d == n);
   return d;
ll pollard_rho(ll n){
   11 x, c, y, d, k;
   int i;
   do{
      i = 1;
      x = 11rand() % n, c = 11rand() % n;
       y = x, k = 4;
       do{
          if(++i == k) y = x, k *= 2;
          x = add(mul(x, x, n), c, n);
          d = \underline{gcd(abs(x - y), n)};
       \}while(d == 1);
   }while(d == n);
   return d;
void factorize(ll val, map<ll, int> &fac){
   if(rabin(val)) fac[ val ]++;
       11 d = pollard_rho(val);
       factorize(d, fac);
       factorize(val / d, fac);
map<ll, int> factor(ll val){
   map<ll, int> fac;
   if(val > 1) factorize(val, fac);
   return fac;
Miller Rabin
bool rabin(ll n){
   if(n <= 1) return 0;
   if(n <= 3) return 1;
   11 s = 0, d = n - 1;
   while(d % 2 == 0) d /= 2, s++;
   for(int k = 0; k < 64; k++){
       11 a = (11rand() \% (n - 3)) + 2;
       11 x = fexp(a, d, n);
       if(x != 1 \&\& x != n-1){
          for(int r = 1; r < s; r++){
              x = mul(x, x, n);
              if(x == 1) return 0;
              if(x == n-1) break;
          if(x != n-1) return 0;
       }
   }
   return 1;
```

ret = (ret & mod) + (ret>>61);

```
Totiente
11 totiente(ll n){
   11 \text{ ans} = n;
   for(ll i = 2; i*i <= n; i++){
       if(n \% i == 0){
          ans = ans / i * (i - 1);
          while(n % i == 0) n /= i;
   }
   if(n > 1) ans = ans / n * (n - 1);
   return ans:
Primitive root
// a primitive root modulo n is any number g such that
  any c coprime to n is congruent to a power of g modulo
  n.
bool exists_root(ll n){
   if(n == 1 || n == 2 || n == 4) return true;
   if(n \% 2 == 0) n /= 2;
   if(n % 2 == 0) return false;
   // test if n is a power of only one prime
   for(11 i = 3; i * i <= n; i += 2) if(n % i == 0){
      while(n % i == 0) n /= i;
      return n == 1;
   return true;
11 primitive_root(l1 n){
   if(n == 1 || n == 2 || n == 4) return n - 1;
   if(not exists_root(n)) return -1;
   11 x = phi(n);
   auto pr = factorize(x);
   auto check = [x, n, pr](ll m){
       for(11 p : pr) if(fexp(m, x / p, n) == 1)
          return false;
      return true;
   };
   for(11 m = 2; ; m++) if(__gcd(m, n) == 1)
       if(check(m)) return m;
}
// Let's denote R(n) as the set of primitive roots
 modulo n, p is prime
// g \in R(p) \Rightarrow (pow(g, p-1, p * p) == 1 ? g+p : g) \in n
   R(pow(p, k)), for all k > 1
// g \text{ in } R(pow(p, k)) \Rightarrow (g \% 2 == 1 ? g : g + pow(p, k))
   Mobius Function
memset(mu, 0, sizeof mu);
mu[1] = 1;
for(int i = 1; i < N; i++)
   for(int j = i + i; j < N; j += i)
      mu[j] -= mu[i];
// g(n) = sum{f(d)} => f(n) = sum{mu(d)*g(n/d)}
Mulmod TOP
constexpr uint64_t mod = (1ul1 << 61) - 1;
uint64_t modmul(uint64_t a, uint64_t b){
   uint64_t 11 = (uint32_t)a, h1 = a>>32, 12 = (
     uint32_t)b, h2 = b >> 32;
   uint64_t 1 = 11*12, m = 11*h2 + 12*h1, h = h1*h2;
   uint64_t ret = (1&mod) + (1>>61) + (h << 3) + (m >>
     29) + (m << 35 >> 3) + 1;
```

```
ret = (ret & mod) + (ret>>61);
   return ret-1:
Matrix Determinant
long double a[n][n];
long double gauss(){
   long double det = 1;
   for(int i = 0; i < n; i++){</pre>
       int q = i;
       for(int j = i+1; j < n; j++){
          if(abs(a[j][i]) > abs(a[q][i]))
              q = j;
       if(abs(a[q][i]) < EPS){
          det = 0;
          break;
       if(i != q){
          for(int w = 0; w < n; w++)
              swap(a[i][w], a[q][w]);
          det = -det;
       det *= a[i][i];
       for(int j = i+1; j < n; j++) a[i][j] /= a[i][i];</pre>
       for(int j = 0; j < n; j++) if(j != i){
          if(abs(a[j][i]) > EPS)
              for(int k = i+1; k < n; k++)
                 a[j][k] -= a[i][k] * a[j][i];
       }
   }
   return det;
Simplex Method
typedef long double dbl;
const dbl eps = 1e-6;
const int N = , M = ;
mt19937 rng(chrono::steady_clock::now().time_since_epoch
  ().count());
struct simplex {
   int X[N], Y[M];
   dbl A[M][N], b[M], c[N];
   dbl ans:
   int n, m;
   dbl sol[N];
   void pivot(int x, int y){
       swap(X[y], Y[x]);
       b[x] /= A[x][y];
       for(int i = 0; i < n; i++)
          if(i != y)
              A[x][i] /= A[x][y];
       A[x][y] = 1. / A[x][y];
       for(int i = 0; i < m; i++)
          if(i != x \&\& abs(A[i][y]) > eps) {
              b[i] -= A[i][y] * b[x];
              for(int j = 0; j < n; j++) if(j != y)
                 A[i][j] -= A[i][y] * A[x][j];
              A[i][y] = -A[i][y] * A[x][y];
          }
```

```
ans += c[y] * b[x];
       for(int i = 0; i < n; i++)
          if(i != y)
              c[i] -= c[y] * A[x][i];
       c[y] = -c[y] * A[x][y];
   }
   // maximiza sum(x[i] * c[i])
   // sujeito a
   // sum(a[i][j] * x[j]) <= b[i] para 0 <= i < m (Ax)
   // x[i] >= 0 para 0 <= i < n (x >= 0)
   // (n variaveis, m restricoes)
   // guarda a resposta em ans e retorna o valor otimo
   dbl solve(int _n, int _m) {
      this->n = _n; this->m = _m;
       for(int i = 1; i < m; i++){
          int id = uniform_int_distribution<int>(0, i)(
            rng);
          swap(b[i], b[id]);
          for(int j = 0; j < n; j++)
              swap(A[i][j], A[id][j]);
      }
      ans = 0.;
       for(int i = 0; i < n; i++) X[i] = i;</pre>
       for(int i = 0; i < m; i++) Y[i] = i + n;</pre>
       while(true) {
          int x = min_element(b, b + m) - b;
          if(b[x] >= -eps)
              break;
          int y = find_if(A[x], A[x] + n, [](dbl d) {
            return d < -eps; }) - A[x];
          if(y == n) throw 1; // no solution
          pivot(x, y);
      while(true) {
          int y = max_element(c, c + n) - c;
          if(c[y] <= eps) break;</pre>
          int x = -1;
          dbl mn = 1. / 0.;
          for(int i = 0; i < m; i++)
              if(A[i][y] > eps \&\& b[i] / A[i][y] < mn)
                 mn = b[i] / A[i][y], x = i;
          if(x == -1) throw 2; // unbounded
          pivot(x, y);
      memset(sol, 0, sizeof(dbl) * n);
       for(int i = 0; i < m; i++)</pre>
          if(Y[i] < n)
              sol[Y[i]] = b[i];
      return ans;
};
FFT
void fft(vector<base> &a, bool inv){
   int n = (int)a.size();
   for(int i = 1, j = 0; i < n; i++){
       int bit = n \gg 1;
       for(; j >= bit; bit >>= 1) j -= bit;
       j += bit;
       if(i < j) swap(a[i], a[j]);
   }
```

```
for(int sz = 2; sz <= n; sz <<= 1) {</pre>
       double ang = 2 * PI / sz * (inv ? -1 : 1);
       base wlen(cos(ang), sin(ang));
       for(int i = 0; i < n; i += sz){
           base w(1, 0);
           for(int j = 0; j < sz / 2; j++){
              base u = a[i+j], v = a[i+j + sz/2] * w;
              a[i+j] = u + v;
              a[i+j+sz/2] = u - v;
              w *= wlen;
       }
   if(inv) for(int i = 0; i < n; i++) a[i] /= 1.0 * n;
FFT Tourist
namespace fft {
 typedef double dbl;
 struct num {
   dbl x, y;
   num() \{ x = y = 0; \}
   num(dbl x, dbl y) : x(x), y(y) {}
 inline num operator+(num a, num b) { return num(a.x +
    b.x, a.y + b.y); }
 inline num operator-(num a, num b) { return num(a.x -
    b.x, a.y - b.y); }
 inline num operator*(num a, num b) { return num(a.x *
    b.x - a.y * b.y, a.x * b.y + a.y * b.x); }
 inline num conj(num a) { return num(a.x, -a.y); }
 int base = 1;
 vector<num> roots = \{\{0, 0\}, \{1, 0\}\};
 vector < int > rev = \{0, 1\};
 const dbl PI = acosl(-1.0);
 void ensure_base(int nbase) {
   if(nbase <= base) return;</pre>
   rev.resize(1 << nbase);</pre>
   for(int i = 0; i < (1 << nbase); i++) {
     rev[i] = (rev[i >> 1] >> 1) + ((i \& 1) << (nbase -
   roots.resize(1 << nbase);</pre>
   while(base < nbase) {</pre>
     dbl \ angle = 2*PI / (1 << (base + 1));
     for(int i = 1 << (base - 1); i < (1 << base); i++)</pre>
       {
       roots[i << 1] = roots[i];</pre>
       dbl \ angle_i = angle * (2 * i + 1 - (1 << base));
       roots[(i << 1) + 1] = num(cos(angle_i), sin(
         angle_i));
     base++;
   }
 }
 void fft(vector<num> &a, int n = -1) {
   if(n == -1) {
     n = a.size();
```

```
assert((n & (n-1)) == 0);
 int zeros = __builtin_ctz(n);
 ensure_base(zeros);
 int shift = base - zeros;
  for(int i = 0; i < n; i++) {
   if(i < (rev[i] >> shift)) {
     swap(a[i], a[rev[i] >> shift]);
  for(int k = 1; k < n; k <<= 1) {
   for(int i = 0; i < n; i += 2 * k) {
     for(int j = 0; j < k; j++) {
       num z = a[i+j+k] * roots[j+k];
       a[i+j+k] = a[i+j] - z;
       a[i+j] = a[i+j] + z;
     }
   }
 }
}
vector<num> fa, fb;
vector<int> multiply(vector<int> &a, vector<int> &b) {
 int need = a.size() + b.size() - 1;
 int nbase = 0;
 while((1 << nbase) < need) nbase++;</pre>
  ensure_base(nbase);
 int sz = 1 << nbase;</pre>
 if(sz > (int) fa.size()) {
   fa.resize(sz);
 for(int i = 0; i < sz; i++) {
   int x = (i < (int) a.size() ? a[i] : 0);</pre>
   int y = (i < (int) b.size() ? b[i] : 0);</pre>
   fa[i] = num(x, y);
 fft(fa, sz);
 num r(0, -0.25 / sz);
  for(int i = 0; i \le (sz >> 1); i++) {
   int j = (sz - i) & (sz - 1);
   num z = (fa[j] * fa[j] - conj(fa[i] * fa[i])) * r;
   if(i != j) {
     fa[j] = (fa[i] * fa[i] - conj(fa[j] * fa[j])) * r
   }
   fa[i] = z;
 }
 fft(fa, sz);
 vector<int> res(need);
  for(int i = 0; i < need; i++) {</pre>
   res[i] = fa[i].x + 0.5;
 return res;
vector<int> multiply_mod(vector<int> &a, vector<int> &
  b, int m, int eq = 0) {
 int need = a.size() + b.size() - 1;
 int nbase = 0;
 while ((1 << nbase) < need) nbase++;</pre>
 ensure base(nbase):
 int sz = 1 << nbase;</pre>
 if (sz > (int) fa.size()) {
   fa.resize(sz);
  for (int i = 0; i < (int) a.size(); i++) {</pre>
   int x = (a[i] \% m + m) \% m;
   fa[i] = num(x & ((1 << 15) - 1), x >> 15);
```

```
fill(fa.begin() + a.size(), fa.begin() + sz, num {0,
      0}):
   fft(fa, sz);
   if (sz > (int) fb.size()) {
     fb.resize(sz);
   if (eq) {
     copy(fa.begin(), fa.begin() + sz, fb.begin());
   } else {
     for (int i = 0; i < (int) b.size(); i++) {</pre>
       int x = (b[i] \% m + m) \% m;
       fb[i] = num(x & ((1 << 15) - 1), x >> 15);
     fill(fb.begin() + b.size(), fb.begin() + sz, num
       {0, 0});
     fft(fb, sz);
   dbl ratio = 0.25 / sz;
   num r2(0, -1);
   num r3(ratio, 0);
   num r4(0, -ratio);
   num r5(0, 1);
   for (int i = 0; i \le (sz >> 1); i++) {
     int j = (sz - i) & (sz - 1);
     num a1 = (fa[i] + conj(fa[j]));
     num a2 = (fa[i] - conj(fa[j])) * r2;
     num b1 = (fb[i] + conj(fb[j])) * r3;
     num b2 = (fb[i] - conj(fb[j])) * r4;
     if (i != j) {
      num c1 = (fa[j] + conj(fa[i]));
      num c2 = (fa[j] - conj(fa[i])) * r2;
      num d1 = (fb[j] + conj(fb[i])) * r3;
      num d2 = (fb[j] - conj(fb[i])) * r4;
       fa[i] = c1 * d1 + c2 * d2 * r5;
       fb[i] = c1 * d2 + c2 * d1;
     fa[j] = a1 * b1 + a2 * b2 * r5;
     fb[j] = a1 * b2 + a2 * b1;
   fft(fa, sz);
   fft(fb, sz);
   vector<int> res(need);
   for (int i = 0; i < need; i++) {
     long long aa = fa[i].x + 0.5;
     long long bb = fb[i].x + 0.5;
     long long cc = fa[i].y + 0.5;
     res[i] = (aa + ((bb \% m) << 15) + ((cc \% m) << 30))
        % m;
   return res;
 vector<int> square_mod(vector<int> &a, int m) {
   return multiply_mod(a, a, m, 1);
}
NTT
const int mod = 7340033;
const int root = 5;
const int root_1 = 4404020;
const int root_pw = 1<<20;</pre>
void fft (vector<int> & a, bool invert) {
   int n = (int) a.size();
```

```
Gauss Xor
   for (int i=1, j=0; i<n; ++i) {
       int bit = n \gg 1;
                                                            const ll MAX = 1e9;
       for (; j>=bit; bit>>=1)
                                                            const int LOG_MAX = 64 - __builtin_clzll((11)MAX);
          j -= bit;
       j += bit;
       if (i < j)
                                                            struct Gauss {
                                                                array<11, LOG_MAX> vet;
          swap (a[i], a[j]);
                                                                int size;
   }
                                                                Gauss() size(0) {}
                                                                Gauss(vector<ll> vals) size(0) {
   for (int len=2; len<=n; len<<=1) {</pre>
                                                                    for(ll val : vals) add(val);
       int wlen = invert ? root_1 : root;
       for (int i=len; i<root_pw; i<<=1)</pre>
          wlen = int (wlen * 111 * wlen % mod);
                                                                bool add(ll val) {
       for (int i=0; i<n; i+=len) {</pre>
                                                                    for(int i = 0; i < LOG_MAX; i++) if(val & (1LL <<</pre>
                                                                       i)) {
          int w = 1:
          for (int j=0; j<len/2; ++j) {</pre>
                                                                       if(vet[i] == 0) {
              int u = a[i+j], v = int (a[i+j+len/2] * 1
                                                                           vet[i] = val;
                11 * w % mod);
                                                                           size++:
                                                                           return true;
              a[i+j] = u+v < mod ? u+v : u+v-mod;
              a[i+j+len/2] = u-v >= 0 ? u-v : u-v+mod;
                                                                       }
                                                                       val ^= vet[i];
              w = int (w * 111 * wlen % mod);
                                                                   }
                                                                   return false;
      }
                                                                }
   if (invert) {
                                                            };
      int nrev = reverse (n, mod);
                                                            Simpson
       for (int i=0; i<n; ++i)</pre>
          a[i] = int (a[i] * 111 * nrev % mod);
                                                            inline double simpson(double fl, double fr, double fmid,
   }
                                                               double 1,double r) { return (fl+fr+4.0*fmid)*(r-1)
}
                                                            double rsimpson(double slr, double fl, double fr, double
Gauss
                                                               fmid,double 1,double r)
                                                                double mid = (1+r)*0.5;
// Solves systems of linear equations.
                                                                double fml = f((1+mid)*0.5);
// To use, build a matrix of coefficients and call run(
                                                                double fmr = f((mid+r)*0.5);
 mat, R, C). If the i-th variable is free, row[i] will
                                                                double slm = simpson(fl,fmid,fml,l,mid);
 be -1, otherwise it's value will be ans[i].
                                                                double smr = simpson(fmid, fr, fmr, mid, r);
                                                                if(fabs(slr-slm-smr) < eps) return slm+smr;</pre>
namespace Gauss {
                                                                return rsimpson(slm,fl,fmid,fml,l,mid)+rsimpson(smr,
 const int MAXC = 1001;
                                                                  fmid, fr, fmr, mid, r);
 int row[MAXC];
 double ans[MAXC];
                                                            double integrate(double 1,double r) {
                                                                double mid = (1+r)*0.5;
 void run(double mat[][MAXC], int R, int C) {
                                                                double fl = f(1);
   REP(i, C) row[i] = -1;
                                                                double fr = f(r);
                                                                double fmid = f(mid);
   int r = 0;
                                                                return rsimpson(simpson(fl,fr,fmid,l,r),fl,fr,fmid,l
   REP(c, C) {
     int k = r;
     FOR(i, r, R) if(fabs(mat[i][c]) > fabs(mat[k][c]))
                                                            Graphs
     if(fabs(mat[k][c]) < eps) continue;</pre>
                                                            Dinic
     REP(j, C+1) swap(mat[r][j], mat[k][j]);
     REP(i, R) if (i != r) {
                                                            const int N = 100005;
      double w = mat[i][c] / mat[r][c];
                                                            const int E = 2000006;
      REP(j, C+1) mat[i][j] -= mat[r][j] * w;
                                                            vector<int> g[N];
                                                            int ne;
     row[c] = r++;
                                                            struct Edge{
   }
                                                                int from, to; ll flow, cap;
                                                            } edge[E];
   REP(i, C) {
                                                            int lvl[N], vis[N], pass, start = N-2, target = N-1;
     int r = row[i];
                                                            int qu[N], qt, px[N];
     ans[i] = r == -1 ? 0 : mat[r][C] / mat[r][i];
   }
                                                            11 run(int s, int sink, ll minE){
 }
                                                                if(s == sink) return minE;
}
```

```
11 \text{ ans} = 0;
                                                                   if (s == t) return;
                                                                   Edge a = \{t, (int)g[t].size(), 0, cap\};
   for(; px[s] < (int)g[s].size(); px[s]++){</pre>
                                                                   Edge b = \{s, (int)g[s].size(), 0, rcap\};
       int e = g[s][ px[s] ];
                                                                    g[s].push_back(a);
       auto &v = edge[e], &rev = edge[e^1];
                                                                    g[t].push_back(b);
       if(lvl[v.to] != lvl[s]+1 || v.flow >= v.cap)
          continue; // v.cap - v.flow < lim</pre>
                                                                void add_flow(Edge& e, flow_t f) {
       11 tmp = run(v.to, sink,min(minE, v.cap-v.flow));
                                                                    Edge &back = g[e.to][e.rev];
                                                                    if (!ec[e.to] && f)
       v.flow += tmp, rev.flow -= tmp;
       ans += tmp, minE -= tmp;
                                                                       hs[H[e.to]].push_back(e.to);
       if(minE == 0) break;
                                                                    e.f += f, ec[e.to] += f;
                                                                   back.f -= f, ec[back.to] -= f;
   return ans;
                                                                flow_t max_flow(int s, int t) {
bool bfs(int source, int sink){
                                                                    int v = g.size();
                                                                   H[s] = v; ec[t] = 1;
   qt = 0;
   qu[qt++] = source;
                                                                   vector<int> co(2 * v);
   lvl[source] = 1;
                                                                    co[0] = v-1;
   vis[source] = ++pass;
                                                                    for(int i = 0; i < v; ++i) cur[i] = g[i].data();
   for(int i = 0; i < qt; i++){
                                                                    for(auto &e : g[s]) add_flow(e, e.c);
       int u = qu[i];
       px[u] = 0;
                                                                   if(hs[0].size())
       if(u == sink) return true;
                                                                    for (int hi = 0; hi >= 0;) {
       for(auto& ed : g[u]) {
                                                                       int u = hs[hi].back();
          auto v = edge[ed];
                                                                       hs[hi].pop_back();
          if(v.flow >= v.cap || vis[v.to] == pass)
                                                                       while (ec[u] > 0) // discharge u
                                                                           if (cur[u] == g[u].data() + g[u].size()) {
              continue; // v.cap - v.flow < lim</pre>
                                                                              H[u] = 1e9;
          vis[v.to] = pass;
          lvl[v.to] = lvl[u]+1;
                                                                              for(auto &e:g[u])
          qu[qt++] = v.to;
                                                                                  if (e.c - e.f && H[u] > H[e.to]+1)
                                                                                     H[u] = H[e.to]+1, cur[u] = &e;
                                                                              if (++co[H[u]], !--co[hi] && hi < v)</pre>
                                                                                  for(int i = 0; i < v; ++i)
   return false;
                                                                                     if (hi < H[i] && H[i] < v){</pre>
11 flow(int source = start, int sink = target){
                                                                                         --co[H[i]];
   11 \text{ ans} = 0;
                                                                                         H[i] = v + 1;
   //for(lim = (1LL << 62); lim >= 1; lim /= 2)
                                                                                     }
   while(bfs(source, sink))
                                                                              hi = H[u];
       ans += run(source, sink, oo);
                                                                           } else if (cur[u]->c - cur[u]->f \&\& H[u]
   return ans;
                                                                             == H[cur[u]->to]+1)
                                                                              add_flow(*cur[u], min(ec[u], cur[u]->c
void addEdge(int u, int v, ll c = 1, ll rc = 0){
                                                                                 - cur[u]->f));
   edge[ne] = \{u, v, 0, c\};
                                                                           else ++cur[u];
   g[u].push_back(ne++);
                                                                       while (hi \geq 0 && hs[hi].empty()) --hi;
   edge[ne] = {v, u, 0, rc};
   g[v].push_back(ne++);
                                                                   return -ec[s];
                                                                }
                                                            };
void reset_flow(){
   for(int i = 0; i < ne; i++)
                                                             Min Cost Max Flow
       edge[i].flow = 0;
                                                             const 11 oo = 1e18;
Push relabel
                                                             const int N = 222, E = 2 * 1000006;
// Push relabel in O(V^2 E^0.5) with gap heuristic
                                                            vector<int> g[N];
// It's quite fast
                                                             int ne;
template<typename flow_t = long long>
                                                             struct Edge{
struct PushRelabel {
                                                                int from, to; 11 cap, cost;
   struct Edge { int to, rev; flow_t f, c; };
                                                             } edge[E];
   vector<vector<Edge> > g;
                                                            int start = N-1, target = N-2, p[N]; int inqueue[N];
   vector<flow_t> ec;
                                                            11 d[N];
   vector<Edge*> cur;
                                                            bool spfa(int source, int sink){
   vector<vector<int> > hs;
                                                                for(int i = 0; i < N; i++) d[i] = oo;</pre>
   vector<int> H:
                                                                inqueue[i] = 0;
   PushRelabel(int n) : g(n), ec(n), cur(n), hs(2*n), H
      (n) {}
                                                                d[source] = 0; queue<int> q; q.push(source);
   void add_edge(int s, int t, flow_t cap, flow_t rcap
                                                                inqueue[source] = 1;
     =0) {
```

```
while(!q.empty()){
       int u = q.front(); q.pop();
       inqueue[u] = 0;
       for(int e : g[u]){
          auto v = edge[e];
          if(v.cap > 0 \text{ and } d[u] + v.cost < d[v.to]){
              d[v.to] = d[u] + v.cost; p[v.to] = e;
              if(!inqueue[v.to]){
                 q.push(v.to); inqueue[v.to] = 1;
          }
       }
   }
   return d[sink] != oo;
// <max flow, min cost>
pair<11, 11> mincost(int source = start, int sink =
  target){
   11 ans = 0, mf = 0;
   while(spfa(source, sink)){
       11 f = oo;
       for(int u = sink; u != source; u = edge[ p[u] ].
          f = min(f, edge[p[u]].cap);
       for(int u = sink; u != source; u = edge[ p[u] ].
         from) {
          edge[p[u]].cap -= f;
          edge[p[u] ^ 1].cap += f;
      mf += f;
       ans += f * d[sink];
   return {mf, ans};
void addEdge(int u, int v, ll c, ll cost){
   edge[ne] = \{u, v, c, cost\};
   g[u].push_back(ne++);
   edge[ne] = \{v, u, 0, -cost\};
   g[v].push_back(ne++);
Blossom Algorithm for General Matching
const int MAXN = 2020 + 1;
```

```
// 1-based Vertex index
int vis[MAXN], par[MAXN], orig[MAXN], match[MAXN], aux[
  MAXN], t, N;
vector<int> conn[MAXN];
queue<int> Q;
void addEdge(int u, int v) {
   conn[u].push_back(v); conn[v].push_back(u);
void init(int n) {
   N = n; t = 0;
   for(int i=0; i<=n; ++i)
       conn[i].clear(), match[i] = aux[i] = par[i] = 0;
void augment(int u, int v) {
   int pv = v, nv;
   do {
      pv = par[v]; nv = match[pv];
      match[v] = pv; match[pv] = v;
       v = nv:
   } while(u != pv);
int lca(int v, int w) {
   ++t;
   while(true) {
```

```
if(v) {
          if(aux[v] == t) return v; aux[v] = t;
          v = orig[par[match[v]]];
       swap(v, w);
   }
void blossom(int v, int w, int a) {
   while(orig[v] != a) {
       par[v] = w; w = match[v];
       if(vis[w] == 1) Q.push(w), vis[w] = 0;
       orig[v] = orig[w] = a; v = par[w];
bool bfs(int u) {
   fill(vis+1, vis+1+N, -1); iota(orig + 1, orig + N +
   Q = queue < int > (); Q.push(u); vis[u] = 0;
   while(!Q.empty()) {
       int v = Q.front(); Q.pop();
       for(int x: conn[v]) {
          if(vis[x] == -1) {
              par[x] = v; vis[x] = 1;
              if(!match[x]) return augment(u, x), true;
              Q.push(match[x]); vis[match[x]] = 0;
          else if(vis[x] == 0 && orig[v] != orig[x]) {
              int a = lca(orig[v], orig[x]);
              blossom(x, v, a); blossom(v, x, a);
          }
       }
   return false;
int Match() {
   int ans = 0;
   // find random matching (not necessary, constant
     improvement)
   vector<int> V(N-1); iota(V.begin(), V.end(), 1);
   shuffle(V.begin(), V.end(), mt19937(0x94949));
   for(auto x: V) if(!match[x]){
       for(auto y: conn[x]) if(!match[y]) {
          match[x] = y, match[y] = x;
          ++ans; break;
       }
   for(int i=1; i<=N; ++i) if(!match[i] && bfs(i)) ++</pre>
     ans:
   return ans;
Small to Large
void cnt_sz(int u, int p = -1){
   sz[u] = 1;
   for(int v : g[u]) if(v != p)
       cnt_sz(v, u), sz[u] += sz[v];
void add(int u, int p, int big = -1){
    // Update info about this vx in global answer
   for(int v : g[u]) if(v != p && v != big)
       add(v, u);
void dfs(int u, int p, int keep){
   int big = -1, mmx = -1;
   for(int v : g[u]) if(v != p \&\& sz[v] > mmx)
       mmx = sz[v], big = v;
   for(int v : g[u]) if(v != p && v != big)
```

```
dfs(v, u, 0);
                                                               for(int v : g[u]){
   if(big != -1) dfs(big, u, 1);
                                                                   if(!num[v]) tarjanSCC(v);
   add(u, p, big);
                                                                   if(vis[v]) low[u] = min(low[u], low[v]);
   for(auto x : q[u]){
       // answer all queries for this vx
                                                               if(low[u] == num[u]){
                                                                   ssc[u] = ++ssc_cnt; int v;
   if(!keep){ /*Remove data from this subtree*/ }
                                                                   do{
                                                                      v = S.back(); S.pop_back(); vis[v] = 0;
                                                                      ssc[v] = ssc_cnt;
Centroid Decomposition
                                                                   }while(u != v);
void decomp(int v, int p){
   int treesize = calc_sz(v, v);
   if(treesize < k) return;</pre>
                                                            Max Clique
   int cent = centroid(v, v, treesize);
   erased[cent] = 1;
                                                           long long adj[N], dp[N];
   for(int i = 1; i <= treesize; i++) dist[i] = 1e18;</pre>
                                                            for(int i = 0; i < n; i++){
                                                               for(int j = 0; j < n; j++){
   for(pair<int,int> x : G[cent]) if(!erased[x.ff]){
                                                                   int x;
      procurar_ans(x.ff, cent, 1, x.ss); // linear
                                                                   scanf("%d",&x);
       atualiza_dist(x.ff, cent, 1, x.ss); // linear
                                                                   if(x \mid | i == j)
                                                                      adj[i] |= 1LL << j;
                                                               }
   for(pair<int,int> x : G[cent]) if(!erased[x.ff])
                                                           }
       decomp(x.ff, cent);
                                                            int resto = n - n/2;
                                                            int C = n/2;
Kosaraju
                                                            for(int i = 1; i < (1 << resto); i++){}
                                                               int x = i;
vector<int> g[N], gt[N], S; int vis[N], cor[N];
                                                               for(int j = 0; j < resto; j++)
void dfs(int u){
                                                                   if(i & (1 << j))
   vis[u] = 1; for(int v : g[u]) if(!vis[v]) dfs(v);
                                                                      x \&= adj[j + C] >> C;
   S.push_back(u);
                                                               if(x == i){
                                                                   dp[i] = __builtin_popcount(i);
void dfst(int u, int e){
   cor[u] = e;
                                                           }
   for(int v : gt[u]) if(!cor[v]) dfst(v, e);
                                                            for(int i = 1; i < (1 << resto); i++)</pre>
void kosaraju(){
                                                               for(int j = 0; j < resto; j++)
   for(int i = 1; i \le n; i++) if(!vis[i]) dfs(i);
                                                                   if(i & (1 << j))
   for(int i = 1; i <= n; i++) for(int j : g[i])</pre>
                                                                      dp[i] = max(dp[i], dp[i ^ (1 << j)]);
       gt[j].push_back(i);
   int e = 0; reverse(S.begin(), S.end());
                                                           int maxCliq = 0;
   for(int u : S) if(!cor[u]) dfst(u, ++e);
                                                            for(int i = 0; i < (1 << C); i++){}
                                                               int x = i, y = (1 << resto) - 1;
Tarjan
                                                               for(int j = 0; j < C; j++)
                                                                   if(i & (1 << j))
int cnt = 0, root;
                                                                      x \&= adj[j] \& ((1 << C) - 1), y \&= adj[j] >>
void dfs(int u, int p = -1){
                                                                        C:
   low[u] = num[u] = ++t;
                                                               if(x != i) continue;
   for(int v : g[u]){
                                                               maxCliq = max(maxCliq, __builtin_popcount(i) + dp[y
      if(!num[v]){
                                                                 ]);
          dfs(v, u);
                                                           }
              if(u == root) cnt++;
                                                            Dominator Tree
          if(low[v] >= num[u]) u PONTO DE ARTICULAÇÃO;
          if(low[v] > num[u]) ARESTA u->v PONTE;
                                                            vector<int> g[N], gt[N], T[N];
          low[u] = min(low[u], low[v]);
                                                            vector<int> S;
      else if(v != p) low[u] = min(low[u], num[v]);
                                                            int dsu[N], label[N];
   }
                                                           int sdom[N], idom[N], dfs_time, id[N];
}
                                                            vector<int> bucket[N];
root PONTO DE ARTICULAÇÃO <=> cnt > 1
                                                            vector<int> down[N];
void tarjanSCC(int u){
                                                           void prep(int u){
   low[u] = num[u] = ++cnt;
                                                               S.push_back(u);
   vis[u] = 1;
                                                               id[u] = ++dfs_time;
   S.push_back(u);
                                                               label[u] = sdom[u] = dsu[u] = u;
```

University of Brasilia Strings, 14

```
used[j0] = true;
   for(int v : g[u]){
                                                                   int i0 = p[j0] , delta = oo, j1;
                                                                   for(int j = 1; j \le m; ++j)
      if(!id[v])
          prep(v), down[u].push_back(v);
                                                                      if(! used[j]){
       gt[v].push_back(u);
                                                                          int cur = a[i0][j] - u[i0] - v[j];
   }
                                                                          if(cur < minv[j])</pre>
}
                                                                              minv[j] = cur, way[j] = j0;
                                                                          if(minv[j] < delta)</pre>
int fnd(int u, int flag = 0){
                                                                              delta = minv[j], j1 = j;
   if(u == dsu[u]) return u;
   int v = fnd(dsu[u], 1), b = label[ dsu[u] ];
                                                                   for(int j = 0; j \le m; ++j)
   if(id[ sdom[b] ] < id[ sdom[ label[u] ] ])
                                                                      if(used[j])
       label[u] = b;
                                                                          u[p[j]] += delta, v[j] -= delta;
   dsu[u] = v;
   return flag ? v : label[u];
                                                                          minv[j] -= delta;
}
                                                                   j0 = j1;
                                                               }while(p[j0] != 0);
void build_dominator_tree(int root, int sz){
   // memset(id, 0, sizeof(int) * (sz + 1));
                                                               do{
   // for(int i = 0; i <= sz; i++) T[i].clear();
                                                                   int j1 = way[j0];
   prep(root);
                                                                   p[j0] = p[j1];
   reverse(S.begin(), S.end());
                                                                   j0 = j1;
                                                               }while(j0);
                                                            }
   int w;
   for(int u : S){
       for(int v : gt[u]){
                                                            // match[i] = coluna escolhida para linha i
          w = fnd(v);
                                                            vector<int> match(n + 1);
                                                            for(int j = 1; j \leftarrow m; ++j)
          if(id[ sdom[w] ] < id[ sdom[u] ])
              sdom[u] = sdom[w];
                                                               match[p[j]] = j;
      gt[u].clear();
                                                            int cost = -v[0];
                                                            Strings
      if(u != root) bucket[ sdom[u] ].push_back(u);
                                                            Aho Corasick
       for(int v : bucket[u]){
          w = fnd(v);
                                                            int to[N][A];
          if(sdom[w] == sdom[v]) idom[v] = sdom[v];
                                                            int ne = 2, fail[N], term[N];
          else idom[v] = w;
                                                            void add_string(const char *str, int id){
                                                               int p = 1;
      bucket[u].clear();
                                                               for(int i = 0; str[i]; i++){
                                                                   int ch = str[i] - 'a';
       for(int v : down[u]) dsu[v] = u;
                                                                   if(!to[p][ch]) to[p][ch] = ne++;
       down[u].clear();
                                                                   p = to[p][ch];
   }
                                                               term[p]++;
   reverse(S.begin(), S.end());
   for(int u : S) if(u != root){
                                                            void init(){
       if(idom[u] != sdom[u]) idom[u] = idom[ idom[u] ];
                                                               for(int i = 0; i < ne; i++) fail[i] = 1;</pre>
      T[ idom[u] ].push_back(u);
                                                               queue<int> q; q.push(1);
                                                               int u, v; char c;
   S.clear();
                                                               while(!q.empty()){
                                                                   u = q.front(); q.pop();
Min Cost Matching
                                                                   for(int i = 0; i < A; i++){
                                                                      if(to[u][i]){
// Min cost matching
                                                                          v = to[u][i]; q.push(v);
// O(n^2 * m)
                                                                          if(u != 1){
// n == nro de linhas
                                                                              fail[v] = to[ fail[u] ][i];
// m == nro de colunas
                                                                              term[v] += term[ fail[v] ];
// n <= m | flow == n
// a[i][j] = custo pra conectar i a j
vector<int> u(n + 1), v(m + 1), p(m + 1), way(m + 1);
                                                                      else if(u != 1) to[u][i] = to[ fail[u] ][i];
for(int i = 1; i \le n; ++i){
                                                                      else to[u][i] = 1;
   p[0] = i;
                                                                   }
   int j0 = 0;
                                                               }
   vector<int> minv(m + 1 , oo);
                                                            }
   vector<char> used(m + 1 , false);
                                                            void clean() {
                                                               memset(to, 0, ne * sizeof(to[0]));
```

University of Brasilia Strings, 15

```
memset(fail, 0, ne * sizeof(fail[0]));
   memset(term, 0, ne * sizeof(term[0]));
   memset(to, 0, ne * sizeof(to[0]));
   ne = 2:
Suffix Array
int lcp[N], c[N];
// Caractere final da string '\0' esta sendo considerado
   parte da string s
void build_sa(char s[], int n, int a[]){
   const int A = 300; // Tamanho do alfabeto
   int c1[n], a1[n], h[n + A];
   memset(h, 0, sizeof h);
   for(int i = 0; i < n; i++) {
       c[i] = s[i];
       h[c[i] + 1]++;
   }
   partial_sum(h, h + A, h);
   for(int i = 0; i < n; i++)
       a[h[c[i]]++] = i;
   for(int i = 0; i < n; i++)
       h[c[i]]--;
   for(int L = 1; L < n; L <<= 1) {
       for(int i = 0; i < n; i++) {
          int j = (a[i] - L + n) \% n;
          a1[h[c[j]]++] = j;
       int cc = -1;
       for(int i = 0; i < n; i++) {
                                                                  }
          if(i == 0 || c[a1[i]] != c[a1[i-1]] || c[(a1[
            i] + L) % n] != c[(a1[i-1] + L) % n])
              h[++cc] = i;
          c1[a1[i]] = cc;
       }
       memcpy(a, a1, sizeof a1);
       memcpy(c, c1, sizeof c1);
       if(cc == n-1) break;
   }
void build_lcp(char s[], int n, int a[]){ // lcp[i] =
  lcp(s[:i], s[:i+1])
   int k = 0;
   //memset(lcp, 0, sizeof lcp);
   for(int i = 0; i < n; i++){
       if(c[i] == n-1) continue;
       int j = a[c[i]+1];
       while(i+k < n \& j+k < n \& s[i+k] == s[j+k]) k
       lcp[c[i]] = k;
       if(k) k--;
   }
}
int comp_lcp(int i, int j){
   if(i == j) return n - i;
   if(c[i] > c[j]) swap(i, j);
   return min(lcp[k] for k in [c[i], c[j]-1]);
```

```
Z Algorithm
vector<int> z_algo(const string &s) {
   int n = s.size(), L = 0, R = 0;
   vector<int> z(n, 0);
   for(int i = 1; i < n; i++){
       if(i \le R) z[i] = min(z[i-L], R - i + 1);
       while(z[i]+i < n \& s[z[i]+i] == s[z[i]])
       if(i+z[i]-1 > R) L = i, R = i + z[i] - 1;
   return z;
Prefix function/KMP
vector<int> preffix_function(const string &s){
   int n = s.size(); vector<int> b(n+1);
   b[0] = -1; int i = 0, j = -1;
   while(i < n){
       while(j \ge 0 \& s[i] != s[j]) j = b[j];
      b[++i] = ++j;
   return b;
void kmp(const string &t, const string &p){
   vector<int> b = preffix_function(p);
   int n = t.size(), m = p.size();
   int j = 0;
   for(int i = 0; i < n; i++){
       while(j \ge 0 \& t[i] != p[j]) j = b[j];
       j++;
       if(j == m){
          //patern of p found on t
          j = b[j];
Min rotation
int min_rotation(int *s, int N) {
 REP(i, N) s[N+i] = s[i];
 int a = 0;
 REP(b, N) REP(i, N) {
   if (a+i == b \mid \mid s[a+i] < s[b+i]) { b += max(0, i-1);}
      break: }
   if (s[a+i] > s[b+i]) \{ a = b; break; \}
 return a;
Manacher
void manacher(char *s, int N, int *rad) {
 static char t[2*MAX];
 int m = 2*N - 1;
 REP(i, m) t[i] = -1;
 REP(i, N) t[2*i] = s[i];
 int x = 0;
 FOR(i, 1, m) {
   int &r = rad[i] = 0;
   if (i <= x+rad[x]) r = min(rad[x+x-i], x+rad[x]-i);</pre>
   while (i-r-1 >= 0 \& i+r+1 < m \& t[i-r-1] == t[i+r]
     +1]) ++r;
   if (i+r >= x+rad[x]) x = i;
```

```
}
                                                                               cn = t[cn].suf; }
                                                                          }
 REP(i, m) if (i-rad[i] == 0 || i+rad[i] == m-1) ++rad[
                                                                      } else if(cd == t[cn].len()) {
                                                                          sufn[si][i] = en;
 REP(i, m) rad[i] /= 2;
                                                                          if(i) t[lst].suf = en; lst = en;
                                                                          t[cn](s[j]) = new_node(j, n - 1, si, cn);
                                                                          cn = t[cn].suf; cd = t[cn].len();
Suffix Automaton
                                                                      } else {
                                                                          int mid = new_node(t[cn].1, t[cn].1 + cd -
map<char, int> to[2*N];
                                                                             1, t[cn].si, t[cn].p);
int link[2*N], len[2*N], last = 0, sz = 1;
                                                                          t[t[cn].p](t[cn][0]) = mid;
                                                                          if(ns) t[ns].suf = mid;
void add_letter(char c){
                                                                          if(i) t[lst].suf = en; lst = en;
   int p = last;
                                                                          sufn[si][i] = en;
   last = sz++;
                                                                          t[mid](s[j]) = new_node(j, n - 1, si, mid)
   len[last] = len[p] + 1;
   for(; !to[p][c]; p = link[p]) to[p][c] = last;
                                                                          t[mid](t[cn][cd]) = cn;
   if(to[p][c] == last){
                                                                          t[cn].p = mid; t[cn].l += cd; cn = t[mid].
       link[last] = 0;
                                                                            p:
       return;
                                                                          int g = cn? j - cd : i + 1; cn = t[cn].suf
   }
   int u = to[p][c];
                                                                          while(g < j \&\& g + t[t[cn](S[si][g])].len
   if(len[u] == len[p]+1){
                                                                            () <= i) {
       link[last] = u;
                                                                             cn = t[cn](S[si][g]); g += t[cn].len();
       return;
                                                                          if(g == j) \{ ns = 0; t[mid].suf = cn; cd =
   int c1 = sz++;
                                                                             t[cn].len(); }
   to[c1] = to[u];
                                                                          else { ns = mid; cn = t[cn](S[si][g]); cd
   link[c1] = link[u];
                                                                            = j - g; }
   len[c1] = len[p]+1;
                                                                      }
   link[last] = link[u] = c1;
                                                                   }
   for(; to[p][c] == u; p = link[p]) to[p][c] = c1;
                                                               }
Suffix Tree
                                                            Geometry
namespace sf {
                                                            2D basics
// const int NS = ; const int N = * 2;
int cn, cd, ns, en = 1, lst;
string S[NS]; int si = -1;
                                                            typedef double cod;
vector<int> sufn[N]; // sufn[si][i] no do sufixo S[si][i
                                                           double eps = 1e-7;
  . . . ]
                                                           bool eq(cod a, cod b){ return abs(a - b) <= eps; }</pre>
struct node {
   int 1, r, si, p, suf;
                                                           struct vec{
   map<char, int> adj;
                                                               cod x, y; int id;
   node() : 1(0), r(-1), suf(0), p(0) {}
                                                               vec(cod a = 0, cod b = 0) : x(a), y(b) {}
   node(int L, int R, int S, int P) : 1(L), r(R), si(S)
                                                               vec operator+(const vec &o) const{
      , p(P) \{ \}
                                                                   return \{x + o.x, y + o.y\};
   inline int len() { return r - l + 1; }
   inline int operator[](int i) { return S[si][l + i];
                                                               vec operator-(const vec &o) const{
                                                                   return \{x - o.x, y - o.y\};
   inline int& operator()(char c) { return adj[c]; }
} t[N];
                                                               vec operator*(cod t) const{
inline int new_node(int L, int R, int S, int P) { t[en]
                                                                  return {x * t, y * t};
  = node(L, R, S, P); return en++; }
void add_string(string s) {
                                                               vec operator/(cod t) const{
   s += '; S[++si] = s; sufn[si].resize(s.size() + 1)
                                                                  return {x / t, y / t};
      ; cn = cd = 0;
   int i = 0; const int n = s.size();
                                                               cod operator*(const vec &o) const{ // cos
   for(int j = 0; j < n; j++)
                                                                  return x * o.x + y * o.y;
       for(; i <= j; i++) {
          if(cd == t[cn].len() \&\& t[cn](s[j])) { cn = t}
                                                               cod operator^(const vec &o) const{ // sin
             [cn](s[j]); cd = 0; 
                                                                  return x * o.y - y * o.x;
          if(cd < t[cn].len() \& t[cn][cd] == s[j]) {
              cd++;
                                                               bool operator==(const vec &o) const{
              if(j < s.size() - 1) break;
                                                                   return eq(x, o.x) \& eq(y, o.y);
              else {
                 if(i) t[lst].suf = cn;
                                                               bool operator<(const vec &o) const{</pre>
                 for(; i <= j; i++) { sufn[si][i] = cn;</pre>
                                                                   if(!eq(x, o.x)) return x < o.x;
```

```
return y < o.y;</pre>
   }
   cod cross(const vec &a, const vec &b) const{
      return (a-(*this)) ^ (b-(*this));
   int ccw(const vec &a, const vec &b) const{
       cod tmp = cross(a, b);
       return (tmp > eps) - (tmp < -eps);</pre>
   cod dot(const vec &a, const vec &b) const{
       return (a-(*this)) * (b-(*this));
   cod len() const{
      return sqrt(x * x + y * y); // <
   double angle(const vec &a, const vec &b) const{
      return atan2(cross(a, b), dot(a, b));
   double tan(const vec &a, const vec &b) const{
      return cross(a, b) / dot(a, b);
   vec unit() const{
      return operator/(len());
   int quad() const{
      if(x > 0 \& y >= 0) return 0;
      if(x \le 0 \& y > 0) return 1;
       if(x < 0 \& y <=0) return 2;
       return 3;
   bool comp(const vec &a, const vec &b) const{
       return (a - *this).comp(b - *this);
   bool comp(vec b){
      if(quad() != b.quad()) return quad() < b.quad();</pre>
      if(!eq(operator^(b), 0)) return operator^(b) > 0;
      return (*this) * (*this) < b * b;
   template<class T>
   void sort_by_angle(T first, T last) const{
       std::sort(first, last, [=](const vec &a, const
         vec &b){
          return comp(a, b);
   }
   vec rot90() const{ return {-y, x}; }
   vec rot(double a) const{
      return \{\cos(a)*x - \sin(a)*y, \sin(a)*x + \cos(a)*y\};
   vec proj(const vec &b) const{ // proj of *this onto
       cod k = operator*(b) / (b * b);
      return b * k;
   // proj of (*this) onto the plane orthogonal to b
   vec rejection(vec b) const{
      return (*this) - proj(b);
};
struct line{
   cod a, b, c; vec n;
   line(vec q, vec w){ // q.cross(w, (x, y)) = 0
      a = -(w.y-q.y);
      b = w.x-q.x;
      c = -(a * q.x + b * q.y);
      n = \{a, b\};
```

```
cod dist(const vec &o) const{
      return abs(eval(o)) / n.len();
   bool contains(const vec &o) const{
       return eq(a * o.x + b * o.y + c, \emptyset);
   cod dist(const line &o) const{
       if(!parallel(o)) return 0;
       if(!eq(o.a * b, o.b * a)) return 0;
       if(!eq(a, 0))
          return abs(c - o.c * a / o.a) / n.len();
       if(!eq(b, 0))
          return abs(c - o.c * b / o.b) / n.len();
       return abs(c - o.c);
   bool parallel(const line &o) const{
      return eq(n ^ o.n, 0);
   bool operator==(const line &o) const{
       if(!eq(a*o.b, b*o.a)) return false;
       if(!eq(a*o.c, c*o.a)) return false;
       if(!eq(c*o.b, b*o.c)) return false;
      return true;
   bool intersect(const line &o) const{
       return !parallel(o) || *this == o;
   vec inter(const line &o) const{
       if(parallel(o)){
          if(*this == o){ }
          else{ /* dont intersect */ }
       auto tmp = n \hat{o}.n;
       return {(o.c*b -c*o.b)/tmp, (o.a*c -a*o.c)/tmp};
   }
   vec at_x(cod x) const{
      return \{x, (-c-a*x)/b\};
   vec at_y(cod y) const{
      return \{(-c-b*y)/a, y\};
   cod eval(const vec &o) const{
      return a * o.x + b * o.y + c;
};
struct segment{
   vec p. a:
   segment(vec a = vec(), vec b = vec()): p(a), q(b) {}
   bool onstrip(const vec &o) const{ // onstrip strip
       return p.dot(o, q) >= -eps && q.dot(o, p) >= -eps
   cod len() const{
      return (p-q).len();
   cod dist(const vec &o) const{
       if(onstrip(o)) return line(p, q).dist(o);
       return min((o-q).len(), (o-p).len());
   bool contains(const vec &o) const{
       return eq(p.cross(q, o), 0) && onstrip(o);
   bool intersect(const segment &o) const{
       if(contains(o.p)) return true;
```

```
if(contains(o.q)) return true;
       if(o.contains(q)) return true;
      if(o.contains(p)) return true;
      return p.ccw(q, o.p) * p.ccw(q, o.q) == -1
       && o.p.ccw(o.q, q) * o.p.ccw(o.q, p) == -1;
   bool intersect(const line &o) const{
      return o.eval(p) * o.eval(q) <= 0;</pre>
   cod dist(const segment &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || onstrip(o.q)
          || o.onstrip(p) || o.onstrip(q))
             return line(p, q).dist(line(o.p, o.q));
      else if(intersect(o)) return 0;
      return min(min(dist(o.p), dist(o.q)),
                min(o.dist(p), o.dist(q)));
   cod dist(const line &o) const{
      if(line(p, q).parallel(o))
          return line(p, q).dist(o);
       else if(intersect(o)) return 0;
      return min(o.dist(p), o.dist(q));
   }
};
struct hray{
   vec p. a:
   hray(vec a = vec(), vec b = vec()): p(a), q(b){}
   bool onstrip(const vec &o) const{ // onstrip strip
      return p.dot(q, o) >= -eps;
   cod dist(const vec &o) const{
      if(onstrip(o)) return line(p, q).dist(o);
      return (o-p).len();
   bool intersect(const segment &o) const{
      if(!o.intersect(line(p,q))) return false;
       if(line(o.p, o.q).parallel(line(p,q)))
          return contains(o.p) || contains(o.q);
       return contains(line(p,q).inter(line(o.p,o.q)));
   bool contains(const vec &o) const{
      return eq(line(p, q).eval(o), 0) && onstrip(o);
   cod dist(const segment &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || onstrip(o.q))
              return line(p, q).dist(line(o.p, o.q));
          return o.dist(p);
       else if(intersect(o)) return 0;
      return min(min(dist(o.p), dist(o.q)),
                o.dist(p));
   bool intersect(const hray &o) const{
      if(!line(p, q).parallel(line(o.p, o.q)))
          return false;
       auto pt = line(p, q).inter(line(o.p, o.q));
       return contains(pt) && o.contains(pt); // <<</pre>
   bool intersect(const line &o) const{
      if(line(p, q).parallel(o)) return line(p, q)== o;
       if(o.contains(p) || o.contains(q)) return true;
      return (o.eval(p) >= -eps)^(o.eval(p)<o.eval(q));</pre>
      return contains(o.inter(line(p, q)));
```

```
cod dist(const line &o) const{
      if(line(p,q).parallel(o))
          return line(p,q).dist(o);
       else if(intersect(o)) return 0;
      return o.dist(p);
   cod dist(const hray &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || o.onstrip(p))
              return line(p,q).dist(line(o.p, o.q));
          return (p-o.p).len();
       else if(intersect(o)) return 0;
      return min(dist(o.p), o.dist(p));
};
double heron(cod a, cod b, cod c){
   cod s = (a + b + c) / 2;
   return sqrt(s * (s - a) * (s - b) * (s - c));
}
struct circle {
   vec c; cod r;
   circle() : c(0, 0), r(0) {}
   circle(const vec o) : c(o), r(0) {}
   circle(const vec &a, const vec &b) {
      c = (a + b) * 0.5;
      r = (a - c).len();;
   circle(const vec &a, const vec &b, const vec &cc) {
       auto tmp = (b - a) * 2;
       auto q = line(tmp.x, tmp.y, a * a - b * b);
       tmp = (cc - a) * 2;
       auto w = line(tmp.x, tmp.y, a * a - cc * cc);
      c = q.inter(w);
      r = (a - c).len();
   bool inside(const vec &a) const {
      return (a - c).len() \ll r;
};
circle min_circle(vector<vec> v) {
   random_shuffle(v.begin(), v.end());
   circle ans;
   int n = (int)v.size();
   for(int i = 0; i < n; i++) if(!ans.inside(v[i])) {
       ans = circle(v[i]);
       for(int j = 0; j < i; j++) if(!ans.inside(v[j])){
          ans = circle(v[i], v[j]);
          for(int k=0; k<j; k++)if(!ans.inside(v[k])){
              ans = circle(v[i], v[j], v[k]);
      }
   return ans;
Circle line intersection
// intersection of line a * x + b * y + c = 0
// and circle centered at the origin with radius r
```

double r, a, b, c; // given as input

double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);

```
puts("no points");
else if(abs(c*c - r*r*(a*a+b*b)) < EPS){
       puts("1 point");
       cout << x0 << ' ' << y0 << '\n';
else {
       double d = r*r - c*c/(a*a+b*b);
       double mult = sqrt(d/(a*a+b*b));
       double ax, ay, bx, by;
       ax = x0 + b * mult;
       bx = x0 - b * mult;
       ay = y0 - a * mult;
       by = y0 + a * mult;
       puts ("2 points");
       cout<<ax<<' '<<ay<<'\n'<<bx<<' '<<by<<'\n';
Half plane intersection
const double eps = 1e-8;
typedef pair<long double, long double> pi;
bool z(long double x){ return fabs(x) < eps; }</pre>
struct line{
       long double a, b, c;
       bool operator<(const line &l)const{</pre>
              bool flag1 = pi(a, b) > pi(0, 0);
              bool flag2 = pi(1.a, 1.b) > pi(0, 0);
              if(flag1 != flag2) return flag1 > flag2;
              long double t = ccw(pi(0, 0), pi(a, b), pi(1.a, 1)
                   .b));
              return z(t) ? c * hypot(l.a, l.b) < l.c * hypot(a
                   , b) : t > 0;
       }
       pi slope(){ return pi(a, b); }
pi cross(line a, line b){
       long double det = a.a * b.b - b.a * a.b;
       return pi((a.c * b.b - a.b * b.c) / det, (a.a * b.c
           - a.c * b.a) / det);
bool bad(line a, line b, line c){
       if(ccw(pi(0, 0), a.slope(), b.slope()) <= 0) return</pre>
            false:
       pi crs = cross(a, b);
       return crs.first * c.a + crs.second * c.b >= c.c;
bool solve(vector<line> v, vector<pi> &solution){ // ax
    + by <= c;
       sort(v.begin(), v.end());
       deque<line> dq;
       for(auto &i : v){
              if(!dq.empty() \&\& z(ccw(pi(0, 0), dq.back().slope))
                    (), i.slope()))) continue;
              while(dq.size() >= 2 && bad(dq[dq.size()-2], dq.
                   back(), i)) dq.pop_back();
              while(dq.size() \geq 2 && bad(i, dq[0], dq[1])) dq.
                   pop_front();
              dq.push_back(i);
       while(dq.size() > 2 && bad(dq[dq.size()-2], dq.back
            (), dq[0])) dq.pop_back();
       \label{eq:while} \begin{subarray}{ll} \begin{suba
              dq.pop_front();
       vector<pi> tmp;
       for(int i=0; i<dq.size(); i++){</pre>
              line cur = dq[i], nxt = dq[(i+1)%dq.size()];
              if(ccw(pi(0, 0), cur.slope(), nxt.slope()) <= eps | }
```

if(c*c > r*r*(a*a+b*b)+EPS)

```
) return false;
  tmp.push_back(cross(cur, nxt));
}
solution = tmp;
return true;
```

Detect empty Half plane intersection

```
// abs(point a) = absolute value of a
// ccw(a, b, c) = a.ccw(b, c)
pair<bool, point> half_inter(vector<pair<point,point> >
  &vet){
   random_shuffle(all(vet));
   point p;
   rep(i,0,sz(vet)) if(ccw(vet[i].x,vet[i].y,p) != 1){
       point dir = (vet[i].y - vet[i].x) / abs(vet[i].y
         - vet[i].x);
       point l = vet[i].x - dir*1e15;
       point r = vet[i].x + dir*1e15;
       if(r < 1) swap(1, r);
       rep(j, 0, i){
          if(ccw(point(), vet[i].x-vet[i].y, vet[j].x-
            vet[j].y) == 0){
              if(ccw(vet[j].x, vet[j].y, p) == 1)
                 continue;
              return mp(false, point());
          if(ccw(vet[j].x, vet[j].y, 1) != 1)
              1 = max(1, line_intersect(vet[i].x,vet[i].
                y,vet[j].x,vet[j].y));
          if(ccw(vet[j].x, vet[j].y, r) != 1)
              r = min(r, line_intersect(vet[i].x,vet[i].
                y,vet[j].x,vet[j].y));
          if(!(1 < r)) return mp(false, point());</pre>
       }
      p = r;
   }
   return mp(true, p);
```

Circle Circle intersection

Assume that the first circle is centered at the origin and second at (x2, y2). Find circle line intersection of first circle and line Ax + By + C = 0, where $A = -2x_2$, $B = -2y_2$, $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$.

Be aware of corner case with two circles centered at the same point.

Tangents of two circles

```
// solve first for same circle(and infinitely many
  tangents)
// Find up to four tangents of two circles
void tangents(pt c, double r1, double r2, vector<line> &
  ans){
  double r = r2 - r1;
  double z = c.x * c.x + c.y * c.y;
  double d = z - r * r;
  if(d < -EPS) return;
  d = sqrt(abs(d));
  line l;
  l.a = (c.x * r + c.y * d) / z;
  l.b = (c.y * r - c.x * d) / z;
  l.c = r1;
  ans.push_back (l);</pre>
```

```
vector<line> tangents(circle a, circle b){
   vector<line> ans;
   pt aux = a.center - b.center;
   for(int i = -1; i \le 1; i += 2)
       for(int j = -1; j \le 1; j += 2)
          tangents(aux, a.r * i, b.r * j, ans);
   for(size_t i = 0; i < ans.size(); ++i)</pre>
       ans[i].c = ans[i].a * a.x + ans[i].b * a.y;
   return ans;
Convex Hull
vector<vec> monotone_chain_ch(vector<vec> P){
   sort(P.begin(), P.end());
   vector<vec> L, U;
   for(auto p : P){
       while(L.size() >= 2 && L[L.size() - 2].cross(L.
         back(), p) < 0)
          L.pop_back();
       L.push_back(p);
   }
   reverse(P.begin(), P.end());
   for(auto p : P){
       while(U.size() >= 2 && U[U.size() - 2].cross(U.
         back(), p) < 0)
          U.pop_back();
       U.push_back(p);
   L.pop_back(), U.pop_back();
   L.reserve(L.size() + U.size());
   L.insert(L.end(), U.begin(), U.end());
   return L;
Check point inside polygon
bool below(const vector<vec> &vet, vec p){
   auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
   if(it == vet.begin()) return *it == p;
   return prev(it)->cross(*it, p) <= 0;</pre>
}
bool above(const vector<vec> &vet, vec p){
   auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
   if(it == vet.begin()) return *it == p;
   return prev(it)->cross(*it, p) >= 0;
}
// lowerhull, upperhull and point, borders included
bool inside_poly(const vector<vec> &lo, const vector<vec</pre>
  > &hi, vec p){
   return below(hi, p) && above(lo, p);
```

Check point inside polygon without lower/upper hull

```
// borders included
```

```
// must not have 3 colinear consecutive points
bool inside_poly(const vector<vec> &v, vec p){
   if(v[0].ccw(v[1], p) < 0) return false;
   if(v[0].ccw(v.back(), p) > 0) return 0;
   if(v[0].ccw(v.back(), p) == 0)
      return v[0].dot(p, v.back()) >= 0
          && v.back().dot(p, v[0]) >= 0;
   int L = 1, R = (int)v.size() - 1, ans = 1;
   while(L <= R){</pre>
       int mid = (L+R)/2;
       if(v[0].ccw(v[mid], p) >= 0) ans = mid, L = mid
         +1:
       else R = mid-1;
   }
   return v[ans].ccw(v[(ans+1)%v.size()], p) >= 0;
Minkowski sum
vector<vec> mk(const vector<vec>&a,const vector<vec>&b){
   int i = 0, j = 0;
```

```
for(int k = 0; k < (int)a.size(); k++)if(a[k] < a[i])
for(int k = 0; k < (int)b.size(); k++)if(b[k] < b[j])
   j = k;
vector<vec> c;
c.reserve(a.size() + b.size());
for(int k = 0; k < int(a.size()+b.size()); k++){
   vec pt{a[i] + b[j]};
   if((int)c.size() >= 2
    && c[c.size()-2].ccw(c.back(), pt) == 0)
       c.pop_back();
   c.push_back(pt);
   int q = i+1, w = j+1;
   if(q == int(a.size())) q = 0;
   if(w == int(b.size())) w = 0;
   if(c.back().ccw(a[i]+b[w], a[q]+b[j]) < 0) i = q;
   else j = w;
c.shrink_to_fit();
return c:
```

Geo Notes

Center of mass

System of points(2D/3D): Mass weighted average of points. **Frame(2D/3D):** Get middle point of each segment solve as previously.

Triangle: Average of vertices.

2D Polygon: Compute **signed** area and center of mass of triangle $((0,0), p_i, p_{i+1})$. Then solve as system of points.

Polyhedron surface: Solve each face as a 2D polygon(be aware of (0, 0)) then replace each face with its center of mass and solve as system of points.

Tetrahedron(Triangular pyramid): As triangles, its the average of points.

Polyhedron: Can be done as 2D polygon, but with tetrahedralization intead of triangulation.

Pick's Theorem

Given a polygon without self-intersections and all its vertices on integer coordinates in some 2D grid. Let A be its area, I the number of points with interger coordinates stricly inside the polygon and B the number of points with interger coordinates in the border of the polygon. The following formula holds: $A = I + \frac{B}{2} - 1$.

Miscellaneous

LIS

```
multiset<int> S;
for(int i = 0; i < n; i++){
   auto it = S.upper_bound(a[i]); // low for inc
   if(it != S.end()) S.erase(it);
   S.insert(a[i]);
}
ans = S.size();
DSU rollback
struct DSU{
   vector<int> sz, p, change;
   vector<tuple<int, int, int>> modifications;
   vector<size_t> saves;
   bool bipartite;
   DSU(int n): sz(n+1, 1), p(n+1), change(n+1),
     bipartite(true){
      iota(p.begin(), p.end(), 0);
   }
   void add_edge(int u, int v){
      if(!bipartite) return;
      int must_change = get_colour(u) == get_colour(v);
      int a = rep(u), b = rep(v);
      if(sz[a] < sz[b]) swap(a, b);
      if(a != b){
          p[b] = a;
          modifications.emplace_back(b, change[b],
            bipartite);
          change[b] ^= must_change;
          sz[a] += sz[b];
      else if(must_change){
          modifications.emplace_back(0, change[0],
            bipartite);
          bipartite = false;
   }
   int rep(int u){
      return p[u] == u ? u : rep(p[u]);
   int get_colour(int u){
      if(p[u] == u) return change[u];
      return change[u] ^ get_colour(p[u]);
   }
   void reset(){
      modifications.clear();
      saves.clear();
      iota(p.begin(), p.end(), 0);
       fill(sz.begin(), sz.end(), 1);
       fill(change.begin(), change.end(), 0);
```

```
bipartite = true;
   }
   void rollback(){
       int u = get<0>(modifications.back());
       tie(ignore, change[u], bipartite) = modifications
         .back();
       sz[p[u]] = sz[u];
      p[u] = u;
       modifications.pop_back();
   void reload(){
       while(modifications.size() > saves.back())
          rollback();
       saves.pop_back();
   }
   void save(){
       saves.push_back(modifications.size());
};
Buildings
// count the number of circular arrays of size m, with
  elements on range [1, c**(n*n)]
int n, m, c; cin >> n >> m >> c;
int x = f_exp(c, n * n); int ans = f_exp(x, m);
for(int i = 1; i \le m; i++) if(m \% i == 0) {
 int y = f_{exp}(x, i);
 for(int j = 1; j < i; j++) if(i % j == 0)
     y = sub(y, mult(j, dp[j]));
 dp[i] = mult(y, inv(i));
 ans = sub(ans, mult(i - 1, dp[i]));
cout << ans << '\n';
Rand
#include <random>
#include <chrono>
cout << RAND_MAX << endl;</pre>
mt19937 rng(chrono::steady_clock::now().time_since_epoch
  ().count());
vector<int> permutation(N);
iota(permutation.begin(), permutation.end(), 0);
shuffle(permutation.begin(), permutation.end(), rng);
iota(permutation.begin(), permutation.end(), 0);
for(int i = 1; i < N; i++){
   swap(permutation[i], permutation[
     uniform_int_distribution<int>(0, i)(rng)]);
}
Klondike
// minimum number of moves to make
// all elements equal
// move: change a segment of equal value
// elements to any value
int v[305], dp[305][305], rec[305][305];
int f(int 1, int r){
 if(r == 1) return 1;
 if(r < 1) return 0;
 if(dp[l][r] != -1) return dp[l][r];
 int ans = f(1+1, r) + 1;
 for(int i = 1+1; i <= r; i++)</pre>
   if(v[i] == v[1])
```

```
ans = min(ans, f(1, i - 1) + f(i+1, r));
                                                                  scanf("%d", &t); t -= i; Q.push(t);
 return dp[l][r] = ans;
                                                                  if(Q.top() > t)  {
                                                                      ans += Q.top() - t; Q.pop(); Q.push(t);
Hilbert Order
                                                              printf("%lld", ans);
// maybe use B = n / sqrt(q)
inline int64_t hilbertOrder(int x, int y, int pow = 21,
  int rotate = 0) {
                                                           Knapsack Bounded with Cost
   if(pow == 0) return 0;
   int hpow = 1 \ll (pow-1);
                                                           // menor custo para conseguir peso ate M usando N tipos
   int seg = (x < hpow) ? (
                                                             diferentes de elementos, sendo que o i-esimo elemento
       (y < hpow) ? 0 : 3
                                                             pode ser usado b[i] vezes, tem peso w[i] e custo c[i]
   ):(
                                                           // O(N * M)
       (y < hpow) ? 1 : 2
   ):
                                                           int b[N], w[N], c[N];
   seg = (seg + rotate) & 3;
                                                           MinQueue Q[M]
   const int rotateDelta[4] = {3, 0, 0, 1};
                                                           int d[M] //d[i] = custo minimo para conseguir peso i
   int nx = x & (x \hat{phow}), ny = y & (y \hat{phow});
   int nrot = (rotate + rotateDelta[seg]) & 3;
                                                           for(int i = 0; i \le M; i++) d[i] = i ? oo : 0;
   int64_t subSquareSize = int64_t(1) << (2*pow - 2);</pre>
                                                           for(int i = 0; i < N; i++){
   int64_t ans = seg * subSquareSize;
                                                               for(int j = 0; j < w[i]; j++)
   int64_t add = hilbertOrder(nx, ny, pow-1, nrot);
                                                                  Q[j].clear();
   ans += (seg == 1 || seg == 2) ? add : (subSquareSize
                                                               for(int j = 0; j \le M; j++){
      - add - 1);
                                                                  q = Q[j \% w[i]];
   return ans:
                                                                  if(q.size() >= q) q.pop();
}
                                                                  q.add(c[i]);
Modular Factorial
                                                                  q.push(d[j]);
                                                                  d[j] = q.getmin();
// Compute (1*2*...*(p-1)*1*(p+1)*(p+2)*..*n) % p
// in O(p*lg(n))
int factmod(int n, int p){
                                                           LCA < O(nlgn), O(1) >
   int ans = 1;
   while(n > 1){
                                                           int start[N], dfs_time;
       for(int i = 2; i <= n % p; i++)</pre>
                                                           int tour[2*N], id[2*N];
          ans = (ans * i) % p;
      n /= p;
                                                           void dfs(int u){
      if(n \% 2) ans = p - ans;
                                                               start[u] = dfs_time;
                                                              id[dfs_time] = u;
   return ans % p;
                                                              tour[dfs_time++] = start[u];
                                                               for(int v : g[u]){
int fac_pow(int n, int p){
                                                                  dfs(v);
   int ans = 0;
                                                                  id[dfs_time] = u;
   while(n) n \neq p, ans += n;
                                                                  tour[dfs_time++] = start[u];
   return ans;
                                                              }
                                                           }
int C(int n, int k, int p){
   if(fac_pow(n, p) > fac_pow(n-k, p) + fac_pow(k, p))
                                                           int LCA(int u, int v){
      return 0;
                                                              if(start[u] > start[v]) swap(u, v);
   int tmp = factmod(k, p) * factmod(n-k, p) % p;
                                                              return id[min(tour[k]for k in [start[u],start[v]])];
   return (f_exp(tmp, p - 2, p) * factmod(n, p)) % p;
                                                           }
                                                           Buffered reader
Enumeration all submasks of a bitmask
                                                           // source: https://github.com/ngthanhtrung23/
// loop through all submask of a given bitmask
                                                             ACM_Notebook_new/blob/master/buffered_reader.h
// it does not include mask 0
                                                           int INP,AM,REACHEOF;
for(int sub = mask; sub; sub = (sub-1)&mask){
                                                           #define BUFSIZE (1<<12)</pre>
                                                           char BUF[BUFSIZE+1], *inp=BUF;
                                                           #define GETCHAR(INP) { \
Slope Trick
                                                               if(!*inp && !REACHEOF) { \
                                                                  memset(BUF,0,sizeof BUF);\
                                                                  int inpzzz = fread(BUF,1,BUFSIZE,stdin);\
///By woqja125, contest: Codeforces Round #371 (Div. 1),
  problem: (C) Sonya and Problem Wihtout a Legend,
                                                                  if (inpzzz != BUFSIZE) REACHEOF = true;\
  Accepted. #
                                                                  inp=BUF; \
int main() {
                                                              } \
                                                               INP=*inp++; \
   int n, t; long long ans = 0; priority_queue<int> Q;
   scanf("%d%d", &n, &t); Q.push(t);
   for(int i = 1; i < n; i++) {
                                                           #define DIG(a) (((a)>='0')\&\&((a)<='9'))
```

```
#define GN(j) { \
   AM=0; \
   GETCHAR(INP); while(!DIG(INP) && INP!='-') GETCHAR(
   if (INP=='-') {AM=1;GETCHAR(INP);} \
   j=INP-'0'; GETCHAR(INP); \
   while(DIG(INP)){j=10*j+(INP-'0');GETCHAR(INP);} \
   if (AM) j=-j;\
Modular summation
//calcula (sum(0 <= i <= n) P(i)) % mod,
//onde P(i) eh uma PA modular (com outro modulo)
namespace sum_pa_mod{
   11 calc(l1 a, l1 b, l1 n, l1 mod){
       assert(a&&b);
       if(a >= b){
          11 ret = ((n*(n+1)/2) \mod)*(a/b);
          if(a%b) ret = (ret + calc(a%b,b,n,mod))%mod;
          else ret = (ret+n+1)%mod;
          return ret;
       }
       return ((n+1)*(((n*a)/b+1)*mod) - calc(b,a,(n*a)/b+1)*mod)
         b, mod) + mod + n/b + 1)%mod;
   }
   //P(i) = a*i \mod m
   11 solve(ll a, ll n, ll m, ll mod){
       a = (a\%m + m)\%m;
       if(!a) return 0;
       11 ret = (n*(n+1)/2)%mod;
       ret = (ret*a)%mod;
       11 g = \underline{gcd(a,m)};
       ret -= m*(calc(a/g,m/g,n,mod)-n-1);
       return (ret%mod + mod)%mod;
   //P(i) = a + r*i \mod m
   11 solve(l1 a, l1 r, l1 n, l1 m, l1 mod){
       a = (a\%m + m)\%m;
       r = (r\%m + m)\%m;
       if(!r) return (a*(n+1))%mod;
       if(!a) return solve(r, n, m, mod);
       11 g, x, y;
       g = gcdExtended(r, m, x, y);
       x = (x\%m + m)\%m;
       11 d = a - (a/g)*g;
       a -= d;
       x = (x*(a/g))%m;
       return (solve(r, n+x, m, mod) - solve(r, x-1, m,
         mod) + mod + d*(n+1))%mod;
   }
};
Edge coloring CPP
const int MX = 300;
int C[MX][MX] = {}, G[MX][MX] = {};
void solve(vector<pii> &E, int N){
   int X[MX] = \{\}, a, b;
   auto update = [&](int u){ for(X[u] = 1; C[u][X[u]];
     X[u]++); };
   auto color = [&](int u, int v, int c){
       int p = G[u][v];
       G[u][v] = G[v][u] = c;
```

C[u][c] = v; C[v][c] = u;

```
C[u][p] = C[v][p] = 0;
   if(p) X[u] = X[v] = p;
   else update(u), update(v);
   return p; };
auto flip = [\&] (int u, int c1, int c2){
   int p = C[u][c1], q = C[u][c2];
   swap(C[u][c1], C[u][c2]);
   if(p) G[u][p] = G[p][u] = c2;
   if( !C[u][c1] ) X[u] = c1;
   if( !C[u][c2] ) X[u] = c2;
   return p; };
for(int i = 1; i <= N; i++) X[i] = 1;</pre>
for(int t = 0; t < E.size(); t++){</pre>
   int u = E[t].first, v0 = E[t].second, v = v0, c0
     = X[u], c = c0, d;
   vector<pii> L;
   int vst[MX] = {};
   while(!G[u][v0]){
       L.emplace_back(v, d = X[v]);
       if(!C[v][c]) for(a = (int)L.size()-1; a >= 0;
          a--) c = color(u, L[a].first, c);
       else if(!C[u][d])for(a=(int)L.size()-1;a>=0;a
         --)color(u,L[a].first,L[a].second);
       else if( vst[d] ) break;
       else vst[d] = 1, v = C[u][d];
   if( !G[u][v0] ){
       for(;v; v = flip(v, c, d), swap(c, d));
       if(C[u][c0]){
          for(a = (int)L.size()-2; a >= 0 \&\& L[a].
             second != c; a--);
          for(; a >= 0; a--) color(u, L[a].first, L[
            a].second);
       } else t--;
   }
}
```

Burnside's Lemma

Let (G, \oplus) be a finite group that acts on a set X. It should hold that $e_g * x = x$ and $g_1 * (g_2 * x) = (g_1 \oplus g_2) * x$, $\forall x \in X, g_1, g_2 \in G$. For each $g \in G$ let $X^g = \{x \in X \mid g * x = x\}$. The number of orbits its given by:

 $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$

Wilson's Theorem

 $(n-1)! = -1 \mod n \iff n \text{ is prime}$

Fibonacci

- $F_{n-1}F_{n+1} F_n^2 = (-1)^n$
- $F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$
- $GCD(F_n, F_m) = F_{GCD(n,m)}$
- $F_n = \frac{(\frac{1+\sqrt{5}}{2})^n (\frac{1-\sqrt{5}}{2})^n}{\sqrt{5}}$

Kirchhoff's Theorem

Laplacian matrix is L = D - A, where D is a diagonal matrix with vertex degrees on the diagonals and A is adjacency matrix.

The number of spanning trees is any cofactor of L. i-th cofactor is determinant of the matrix gotten by removing i-th row and column of L.

Multigraphs

In D[i][i] all loops are excluded. A[i][j] = number of edges from i to j.

Directed multigraphs

D[i][i] = indegree of i minus the number of loops at i. A[i][j] = number of edges from i to j.

The number of oriented spanning trees rooted at a vertex i is the determinant of the matrix gotten by removing the ith row and column of L.

Matroid

Let *X* set of objects, $I \subseteq 2^X$ set of independents sets such that:

- 1. $\emptyset \in I$
- 2. $A \in I, B \subseteq A \implies B \in I$
- 3. Exchange axiom, $A \in I, B \in I, |B| > |A| \implies \exists x \in B \setminus A : A \cup \{x\} \in I$
- 4. $A \subseteq X$ and I and I' are maximal independent subsets of A then |I| = |I'|

Then (X, I) is a matroid. The combinatorial optimization problem associated with it is: Given a weight $w(e) \ge 0 \ \forall e \in X$, find an independet subset that has the largest possible total weight.

Matroid intersection

```
// Input two matroids (X, I_a) and (X, I_b)
// output set I of maximum size, I \in I_a and I \in I_b
set<> I;
while(1){
    for(e_i : X \setminus I)
       if(I + e_i \in I_a \text{ and } I + e_i \in I_b)
           I = I + e_i;
    set<> A, T; queue<> Q;
    for(x : X) label[x] = MARK1;
   for(e_i : X \setminus I){
       if(I + e_i \setminus in I_a)
           Q.push(e_i), label[e_i] = MARK2;
       else{
           for (x \text{ such that } I - x + e_i \in I_a)
               A[x].push(e_i);
       if(I + e_i \setminus in I_b)
           T = T + \{e_i\}
           for(x such that I - x + e_i \setminus in I_b)
               A[e_i].push(x);
   if(T.empty()) break;
   bool found = false;
   while(!Q.empty() and !found){
       auto e = Q.front(); Q.pop();
```

```
for(x : A[e]) if(label[x] == MARK1){
    label[x] = e; Q.push(x);
    if(x \in T){
        found = true; put = 1;
        while(label[x] != MARK2){
            I = put ? (I + x) : (I - x);
            put = 1 - put;
        }
        I = I + x;
        break;
    }
}
if(!found) break;
}
return I;
```

Where path(e) = [e] if label[e] = MARK2, path(label[e]) + [e] otherwise.

Matroid Union

Given k matroids over the same set of objects (X, I_1) , (X, I_2) , ..., (X, I_k) find $A_1 \in I_1$, $A_2 \in I_2$, ..., $A_k \in I_k$ such that $i \neq j, A_i \cap A_j = \emptyset$ and $|\bigcup_{i=1}^k A_i|$ is maximum. Matroid union can be reduced to matroid intersection as follows.

Let $X' = X \times \{1, 2, ..., k\}$, ie, k copies of each element of X with different colors. M1 = (X', Q) where $B \in Q \iff \forall 1 \le i \le k$, $\{x \mid (x, i) \in B\} \in I_i$, ie, for each color, B is independent. M2 = (X', W) where $B \in W \iff i \ne j \implies \neg((x, i) \in B \land (x, j) \in B)$, ie, each element is picked by at most one color.

Intersection of *M*1 and *M*2 is the answer for the combinatorial problem of matroid union.

Notes

When we repeat something and each time we have probability p to succeed then the expected number or tries is $\frac{1}{p}$, till we succeed.

Small to large

Trick in statement If k sets are given you should note that the amount of different set sizes is $O(\sqrt{s})$ where s is total size of those sets. And no more than \sqrt{s} sets have size greater than \sqrt{s} . For example, a path to the root in Aho-Corasick through suffix links will have at most $O(\sqrt{s})$ vertices.

gcd on subsegment, we have at most $log(a_i)$ different values in $\{gcd(a_j, a_{j+1}, ..., a_i) \text{ for } j < i\}$.

From static set to expandable. To insert, create a new set with the new element. While there are two sets with same size, merge them. There will be at most log(n) disjoints sets.

Matrix exponentiation optimization. Save binary power of A_{nxn} and answer q queries $b = A^m x$ in $O((n^3 + qn^2)log(m))$.

Ternary search on integers into binary search, comparing f(mid) and f(mid+1), binary search on derivative

Dynamic offline set For each element we will wind segment of time [a,b] such that element is present in the set during this whole segment. Now we can come up with recursive procedure which handles [l,r] time segment considering that all elements such that $[l,r] \subset [a,b]$ are already included into the set. Now, keeping this invariant we recursively go into [l,m] and [m+1,r] subsegments. Finally when we come into segment of length 1.

```
a > b \implies a \mod b < \frac{a}{2}
```