## CPE 301 - Embedded C Programming hand-out from textbook

Example 2.58 Suppose the following variables are allocated in the specified order. Let r be a pointer that points to 1, q be a pointer that points to s, and p be a pointer that points to c. Also, let s be a 16-bit short, 1 be a 32-bit long, and c be an 8-bit char. Note that in this hypothetical example, pointers are 32-bit variables meaning the processor has 32-bit addressable space.

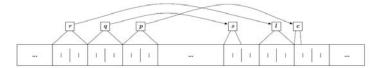


Figure 2.6: A schematic view of memory. The smallest division represents one 8-bit memory location. The height of the division line represents the byte-, short-, and long-address boundaries.

As indicated by Ex. 2.58, the amount of space reserved for any pointer is the same, no matter to what type the pointer is pointing. In the example, 32-bits are assumed for the addressable space; this is based on the architecture of the CPU. One way to make the pointer point to some variable is via the unary operator &, which gives the address of the label on its right-hand-side.

#### Example 2.59

```
p = &c;
q = &s;
r = &1;
```

Note that &p is the address of the pointer variable.

One way to access the contents of a variable using a pointer is via the unary operator \*, which is called the *dereferencing* operator.

Example 2.60 At the end of these three statements, variable c is loaded with the value 10 via the dereference of its address in pointer p.

```
p = &c;
c = 0;
*p = 10;
/* now it is true that (c == 10) */
```

To declare a pointer, just add the \* symbol to the left of the variable name.

### Example 2.61

```
char *p;
short *q;
long *r;
```

This syntax is intended as a mnemonic. Using Ex. 2.61, the notation implies that the expression \*p is a char, \*q is a short and \*r is a long. Note that as seen before, the space allocated to hold p, q and r is all the same (usually 32-bits on modern microprocessors), but what they point to is different. This matters to the compiler when pointer indexing is used.

The unary operators \* and & have a very high precedence. However, the unary operators ++ and -- have the same level of operator precedence. When the compiler parses a line of source code, it resolves operators with the same precedence from right-to-left. Thus, the statement \*p++; will have a very different result compared to (\*p) ++;. The former case will increment the address stored in p first, then dereference the result. The latter case will read the dereferenced address first and increment the resulting value without changing the address stored in p. Several examples of using pointer indexing are listed in Table 2.13, assuming that char c = 5;, char \*p;, and p = &c;. Notice that all but the final statement are equivalent.

	Be	fore		After				
Instruction	&cc = 100	101	Р	&c = 100	101	p	*p	
c = *p + 1;	5	0	100	6	0	100	6	
*p += 1;	5	0	100	6	0	100	6	
++*p;	5	0	100	6	0	100	6	
(*p)++;	5	0	100	6	0	100	6	
*p++;	5	0	100	5	0	101	0	

	Table 2.9: Bitwise Operators
Operator	Operation
&	AND (boolean intersection)
- 1	OR (boolean union)
^	XOR (boolean exclusive-or)
<<	left shift
>>	right shift
~	NOT (boolean negation, i.e., ones' complement)

Statement		mask	d	Embedded usefulness
d = (c & mask);	0x55	0x0F	0x05	Clear bits that are 0 in the mask
d = (c   mask);	0x55	0x0F	0x5F	Set bits that are 1 in the mask
d = (c ^ mask);	0x55	0x0F	0x5A	Invert bits that are 1 in the mask
d = (c << 3);	0x55		0xA8	Multiply by a power of 2
d = (c >> 2);	0x55		0x15	Divide by a power of 2
d = "c;	0x55		0xAA	Invert all bits

Tab	le 2.10: Assig	nment Operators
Operator	Syntax	Equivalent Operation
+=	i += j;	i = (i + j);
-=	i -= j;	i = (i - j);
*=	i *= j;	i = (i * j);
/=	i /= j;	i = (i / j);
%=	i %= j;	i = (i % j);
&=	i &= j;	i = (i & j);
=	i  = j;	i = (i   j);
^=	i ^= j;	i = (i ^ j);
<<=	i <<= j;	i = (i << j);
>>=	i >>= j;	i = (i >> j);

Bit	7	6	5	4	3	2	1	0
0x25	PORTB7	PORTB6	PORTB5	PORTB4	PORTB3	PORTB2	PORTB1	PORTBO
Read/Write	R/W							
Default	0	0	0	0	0	0	0	0

• PORTB7-0: GPIO data value stored in bit n.

Bit	7	6	5	4	3	2	1	0
0x24	DDRB7	DDRB6	DDRB5	DDRB4	DDRB3	DDRB2	DDRB1	DDRBO
Read/Write	R/W							
Default	0	0	0	0	0	0	0	0

DDRB7-0: selects the direction of pin n. If DDRBn is written '1', then PORTBn is configured
as an output pin. If DDRBn is written '0', then PORTBn is configured as an input pin.

Bit	7	6	5	4	3	2	1	0
0x23	PINB7	PINB6	PINB5	PINB4	PINB3	PINB2	PINB1	PINBO
Read/Write	R	R	R	R	R	R	R	R
Default	-	-	-	-	-	-	-	-

• PINB7-0: logic value present on external pin n.

Bit	7	6	5	4	3	2	1	0
0x28	-	PORTC6	PORTC5	PORTC4	PORTC3	PORTC2	PORTC1	PORTCO
Read/Write	R	R/W						
Default	0	0	0	0	0	0	0	0

• PORTC6-0: GPIO data value stored in bit n.

Bit	7	6	5	4	3	2	1	0
0x27	-	DDRC6	DDRC5	DDRC4	DDRC3	DDRC2	DDRC1	DDRCO
Read/Write	R	R/W						
Default	0	0	0	0	0	0	0	0

• DDRC6-0: selects the direction of pin n. If DDRCn is written '1', then PORTCn is configured as an output pin. If DDRCn is written '0', then PORTCn is configured as an input pin.

Bit	7	6	5	4	3	2	1	0
0x26	-	PINC6	PINCS	PINC4	PINC3	PINC2	PINC1	PINCO
Read/Write	R	R	R	R	R	R	R	R
Default	0	-	-	-	-	-	-	-

PINC6-0: logic value present on external pin n.

Bit	7	6	5	4	3	2	1	0
0x2B	PORTD7	PORTD6	PORTD5	PORTD4	PORTD3	PORTD2	PORTD1	PORTDO
Read/Write	R/W							
Default	0	0	0	0	0	0	0	0

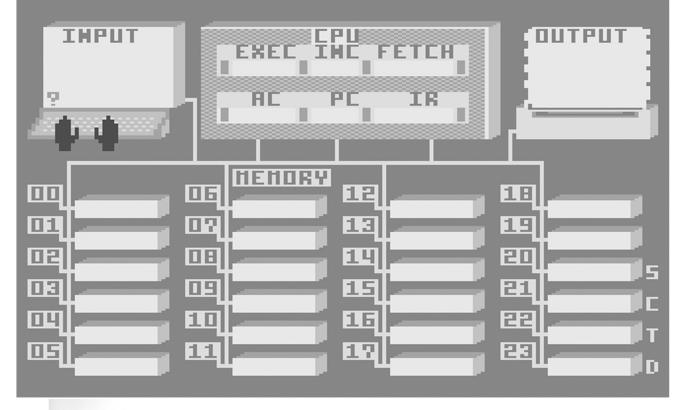
• PORTD7-0: GPIO data value stored in bit n.

Bit	7	6	5	4	3	2	1	0
0x2A	DDRD7	DDRD6	DDRD5	DDRD4	DDRD3	DDRD2	DDRD1	DDRDO
Read/Write	R/W							
Default	0	0	0	0	0	0	0	0

• DDRD7-0: selects the direction of pin n. If DDRDn is written '1', then PORTDn is configured as an output pin. If DDRDn is written '0', then PORTDn is configured as an input pin.

Bit	7	6	5	4	3	2	1	0
0x29	PIND7	PIND6	PIND5	PIND4	PIND3	PIND2	PIND1	PINDO
Read/Write	R	R	R	R	R	R	R	R
Default	-	-	-	-	-	-	-	-

• PIND7-0: logic value present on external pin n.



## ADDENIDIX II ·

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PROGRAM	INSTRUCTIONS

Operation Code	Mnemonic	Explanation
1	LDAxx	Load the accumulator with the value in location xx.
2	STAxx	Store the value of the accumulator into location xx.
3	ADDxx	Add the value in location xx to the accumulator.
4	SUBxx	Subtract the value in location xx from the accumulator.
5	MULxx	Multiply the accumulator by the value in location xx.
6	DIVxx	Divide the accumulator by the value in location xx. This rounds off the
		answer down to the nearest one.
7 8 9 000 001	INPxx OUTxx JMPxx STP SKP01	Input to location xx. Output from location xx. Jump to location xx. Stop. Skip the next instruction if the accumulator is less than 0.
002	SKP02	Skip the next instruction if the accumulator is greater than 0.
003	SKP03	Skip the next instruction if the accumulator is 0.
004	SKP04	Skip the next instruction if the accumulator is either less than or equal to 0.
005	SKP05	Skip the next instruction if the accumulator is either greater than or equal to 0.
006	SKP06	Skip the next instruction if the accumulator is not

# APPENDIX III : COMMANDS

1	YT			V	L	•	111	٠	,
$\mathcal{C}$	nm	m.	and	Ev	nla	n	atio	n	

RUN Run the program starting in location

RUNxx Run the program starting in location xx.

RUNSPEDx Run the program using a speed of x. x can be a value from 0 to 9, with 0 being the slowest and 9 the fastest.

Break. The computer program will stop

running.

CONT Continue at preset speed. CONSPEDX Continue at speed x. CONSTP Continue, one step at a time.

NEW Clears memory.

Get ready to load information begin-LOADxx

ning at location xx.

## JOYSTICK CONTROL

The joystick can be used to perform the following

**SPEED** 

STEP

В

functions: BREAK Press the red button to stop a program.

After a RUN command, move the joystick forward to increase run speed,

and backward to decrease it. Push the joystick forward after a

RUNSTEP command. This does the

same thing as pressing the spacebar.

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