

LEMPER ZIV (LZ) ALGORITHMS

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Dictionary Coding

- Observation: Correlations between parts of the data (patterns)
- Idea: Replace recurring patterns with references to a dictionary
- Static, semi-adaptive, adaptive

Static Dictionary

- Static dictionary technique is most appropriate when considerable prior knowledge about the source is available.
- Similar to the concept of fixed length coding.

Example

- Encode the sequence ***abracadabra***

Code	Entry	Code	Entry
000	a	100	r
001	b	101	ab
010	c	110	ac
011	d	111	ad

Result: 101100110111101100000.

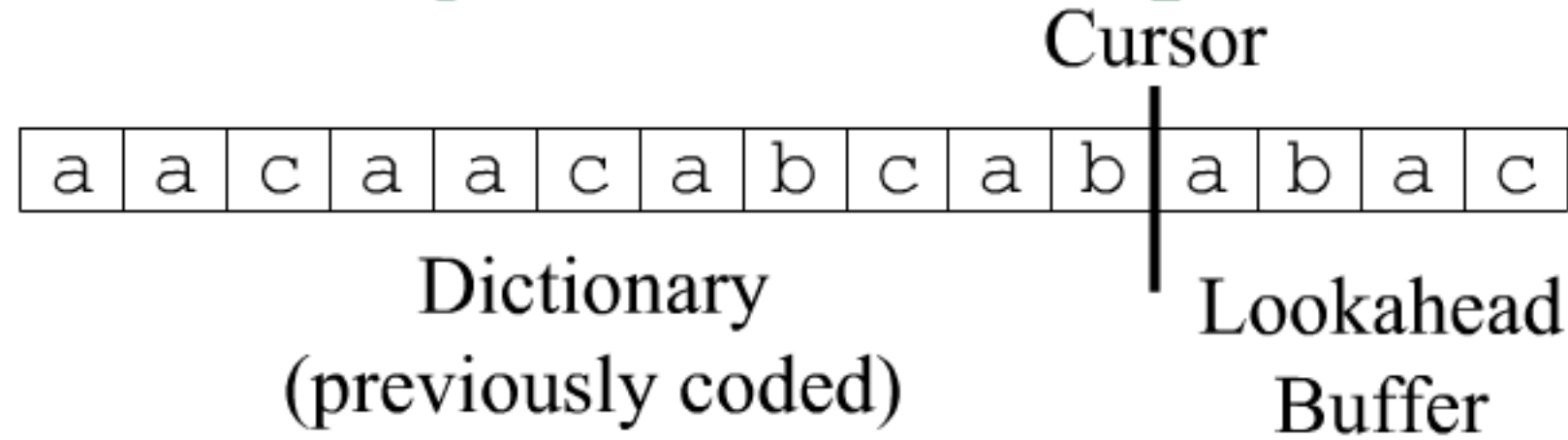
Adaptive dictionary

- LZ algorithms use this approach
- Coding scheme is universal
- No need to transmit/store dictionary
- Single-pass (dictionary creation “on-the-fly”)
- LZ77 and LZ78 are two lossless data compression algorithms developed by Abraham Lempel and Jacob Ziv in 1977 and 1978 respectively.

LZ77

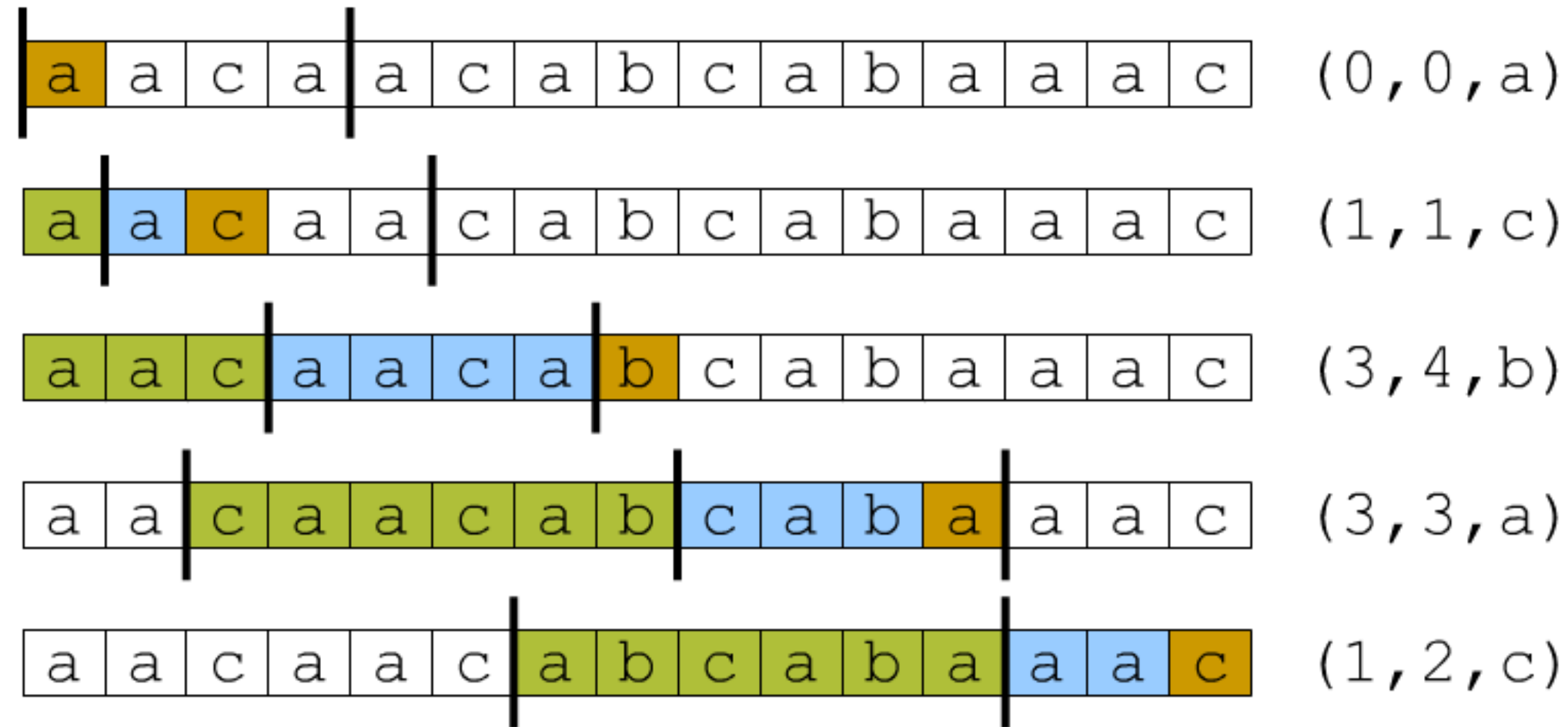
- LZ77 coding uses the dictionary which is a portion of the previously encoded sequence.
- The input sequence is encoded through a sliding windows which consists of a search buffer and a look-ahead buffer.
- The encoder tries to find the match of patterns in the windows and encodes it with a triple $\langle o, l, c \rangle$
o: offset, l: length of the match, c: next character following the match.

LZ77: Sliding Window Lempel-Ziv



- Dictionary and buffer “windows” are fixed length and slide with the cursor
- On each step:
Output (o,l,c)
o = relative position of the longest match in the dictionary
l = length of longest match
c = next char in buffer beyond longest match
- Advance window by $l + 1$

LZ77: Example



Dictionary (size = 6)

Longest match

Next character

LZ77 Decoding

- Decoder keeps same dictionary window as encoder.
- For each message it looks it up in the dictionary and inserts a copy

LZ77 limitations

- The LZ77 implicitly assumes that the like pattern will occur closely.
- Sliding Window LZ is Asymptotically Optimal [Wyner-Ziv,94]
- Will compress long enough strings to the source entropy as the window size goes to infinity.

LZ78: Dictionary Lempel-Ziv

Basic algorithm:

- Keep dictionary of words with integer *id* for each entry (e.g. keep it as a trie).
- Coding loop
 - find the longest match *S* in the dictionary
 - Output the entry *id* of the match and the next character past the match from the input (*id,c*)
 - Add the string *Sc* to the dictionary
- Decoding keeps same dictionary and looks up *ids*

LZ78: Coding Example

	Output	Dict.
a a b a a c a b c a b c b	(0, a)	1 = a
a a b a a c a b c a b c b	(1, b)	2 = ab
a a b a a c a b c a b c b	(1, a)	3 = aa
a a b a a c a b c a b c b	(0, c)	4 = c
a a b a a c a b c a b c b	(2, c)	5 = abc
a a b a a c a b c a b c b	(5, b)	6 = abcb

LZ78: Decoding Example

Input

(0, a)

a												
---	--	--	--	--	--	--	--	--	--	--	--	--

(1, b)

a	a	b										
---	---	---	--	--	--	--	--	--	--	--	--	--

(1, a)

a	a	b	a	a								
---	---	---	---	---	--	--	--	--	--	--	--	--

(0, c)

a	a	b	a	a	c							
---	---	---	---	---	---	--	--	--	--	--	--	--

(2, c)

a	a	b	a	a	c	a	b	c				
---	---	---	---	---	---	---	---	---	--	--	--	--

(5, b)

a	a	b	a	a	c	a	b	c	a	b	c	b
---	---	---	---	---	---	---	---	---	---	---	---	---

Dict.

1 = a

2 = ab

3 = aa

4 = c

5 = abc

6 = abcb

LZ78 Weaknesses

- Dictionary grows without bound
- Long phrases appear late
- Inclusion of first non-matching symbol may prevent a good match
- Few substrings of the processed input are entered into the dictionary

LZW (Lempel-Ziv-Welch)

- Don't send extra character c , but still add Sc to the dictionary.
- The dictionary is initialized with byte values being the first 256 entries (e.g. $a=112$, $ascii$), otherwise there is no way to start it up.
- The decoder is one step behind the coder since it does not know c

LZW: Encoding Example

	Output	Dict.
a a b a a c a b c a b c b	112	256=aa
a a b a a c a b c a b c b	112	257=ab
a a b a a c a b c a b c b	113	258=ba
a a b a a c a b c a b c b	256	259=aac
a a b a a c a b c a b c b	114	260=ca
a a b a a c a b c a b c b	257	261=abc
a a b a a c a b c a b c b	260	262=cab

LZW: Decoding Example

Input

Dict

112	a a b a a c a b c a b c b	
112	a a b a a c a b c a b c b	256=aa
113	a a b a a c a b c a b c b	257=ab
256	a a b a a c a b c a b c b	258=ba
114	a a b a a c a b c a b c b	259=aac
257	a a b a a c a b c a b c b	260=ca
260	a a b a a c a b c a b c b	261=abc

LZ78 and LZW issues

What happens when the dictionary gets too large?

- Throw the dictionary away when it reaches a certain size (used in GIF)
- Throw the dictionary away when it is no longer effective at compressing (used in unix compress)
- Throw the least-recently-used (LRU) entry away when it reaches a certain size (used in BTLZ, the British Telecom standard)

LZ Advantages

- The LZ algorithms are popular because they run in a single pass,
- Provide good compression, are easy to code, and run quickly
- Used in popular compression utilities such as compress, gzip, and WinZip

Lempel-Ziv Summary

LZ77 (Sliding Window)

- Variants: LZSS (Lempel-Ziv-Storer-Szymanski)
- Applications: gzip, Squeeze, LHA, PKZIP, ZOO

LZ78 (Dictionary Based)

- Variants: LZW (Lempel-Ziv-Welch), LZC (Lempel-Ziv-Compress)
- Applications: compress, GIF, CCITT (modems), ARC, PAK

REFERENCES

- <http://en.wikipedia.org>
- Data Compression: The Complete Reference, 3rd Edition by David Saloman, published by Springer (2004).
- ieeexplore.ieee.org/

■ THANK YOU