

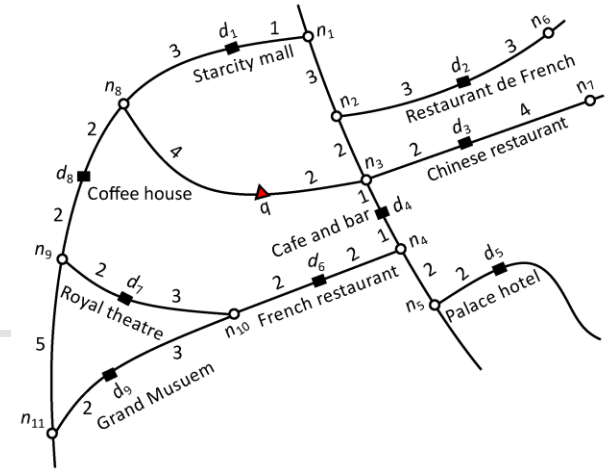


COMP 1002/COMP 1001

Lecture 9

Problem Solving III

Graph



- **Graphs** are so common.
- How can we represent a graph?
 - There are two parts:
 - **Nodes**: can be represented as a list or a set.
 - **Edges**: can be represented as a list or a set.
 - For each edge, we would need a pair or a tuple expressing the two nodes.
 - Some graphs need more information on the edges:
 - **Weight** of an edge.
 - **Direction** of an edge.

Graph

- In computer science, a graph G is often written like:

- $G = (V, E)$

- V is a set of vertices or nodes.
- E is a set of edges.

- Example

- Nodes

- $\{a, b, c, d, e, f\}$ (6 nodes)

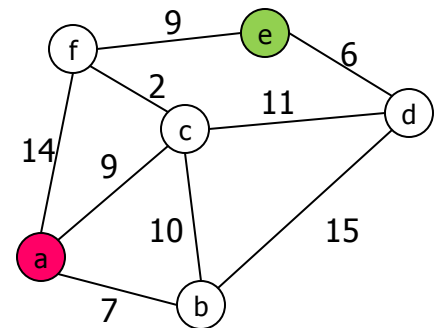
- Edges

- $\{ab, ac, af, bc, bd, cd, cf, de, ef\}$ (9 edges)

- With **weights** on edges:

- $\{ab=7, ac=9, af=14, bc=10, bd=15, cd=11, cf=2, de=6, ef=9\}$

- This graph is **undirected** (no direction on edges).

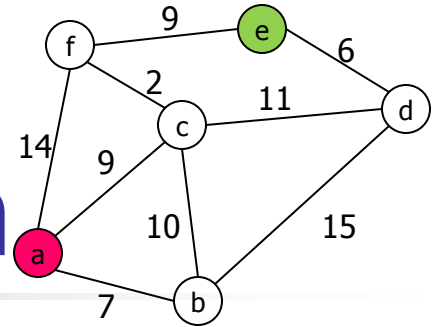


Graph

- We need to build the **data model** for the **graph** before we can use the computer to process it.
 - It is easy to represent the **nodes**, but perhaps harder with the **edges**.
 - We may represent nodes and edges separately as two different groups.
 - It is more natural to represent **edges as linked to nodes**. There are two common representations.
 - **Adjacency list**
 - **Adjacency matrix**
 - Note that an **edge** in an **undirected graph** is equivalent to a **pair of edges** in a **directed graph**.

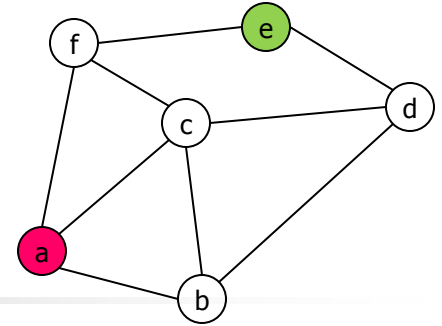


Graph Representation



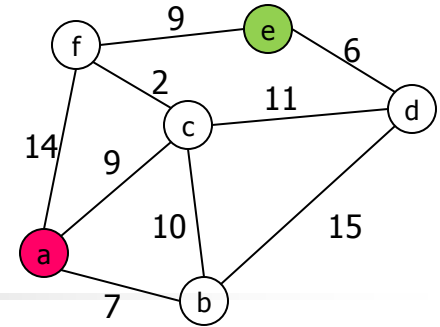
- Simply $G = (V, E)$.
 - Set of nodes:
 - $V = \{a, b, c, d, e, f\}$
 - Set of edges:
 - $E = \{ab, ac, af, bc, bd, cd, cf, de, ef\}$
 - $E = \{(a,b), (a,c), (a,f), (b,c), (b,d), (c,d), (c,f), (d,e), (e,f)\}$
 - Edges storing the weights:
 - $E = \{ab=7, ac=9, af=14, bc=10, bd=15, cd=11, cf=2, de=6, ef=9\}$
 - $E = \{(a,b,7), (a,c,9), (a,f,14), (b,c,10), (b,d,15), (c,d,11), (c,f,2), (d,e,6), (e,f,9)\}$
 - Should we try to store nodes and edges together instead of separately?

Adjacency List



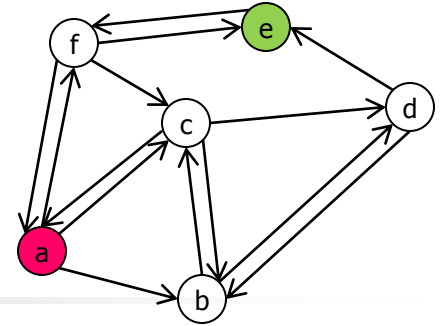
- For each node, put together the edges for each node.
- An **adjacency list** is a **list for each node**, showing the neighbors of that node, i.e. edges.
- Example
 - **Nodes**, a set
 - $\{a,b,c,d,e,f\}$
 - **Edges** in 6 adjacency lists: $D(\text{node})$, each being a set
 - $D(a) = \{b,c,f\}$
 - $D(b) = \{a,c,d\}$
 - $D(c) = \{a,b,d,f\}$
 - $D(d) = \{b,c,e\}$
 - $D(e) = \{d,f\}$
 - $D(f) = \{a,c,e\}$

Adjacency List



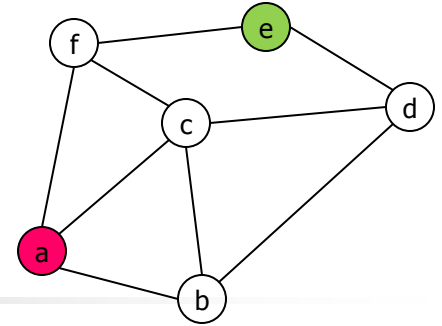
- For edges with **weights**, an adjacency list is a list for each node, showing the **edges and the weights**.
- Example
 - **Nodes**, a set
 - $\{a,b,c,d,e,f\}$
 - **Edges** in adjacency lists: $D(\text{node})$, each being a set of tuples
 - $D(a) = \{(b,7),(c,9),(f,14)\}$
 - $D(b) = \{(a,7),(c,10),(d,15)\}$
 - $D(c) = \{(a,9),(b,10),(d,11),(f,2)\}$
 - $D(d) = \{(b,15),(c,11),(e,6)\}$
 - $D(e) = \{(d,6),(f,9)\}$
 - $D(f) = \{(a,14),(c,2),(e,9)\}$

Adjacency List



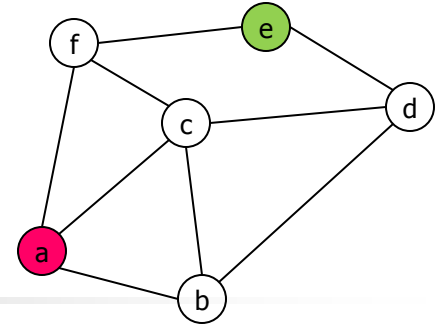
- For a **directed graph** (graph with **direction** on edges), an adjacency list is a list for each node, showing the **next reachable node** and perhaps also the **weights**.
- Example
 - **Nodes**, a set
 - $\{a, b, c, d, e, f\}$
 - **Edges** in adjacency lists: $D(\text{node})$, each being a set
 - $D(a) = \{b, c, f\}$
 - $D(b) = \{c, d\}$
 - $D(c) = \{a, b, d\}$
 - $D(d) = \{b, e\}$
 - $D(e) = \{f\}$
 - $D(f) = \{a, c, e\}$

Adjacency List



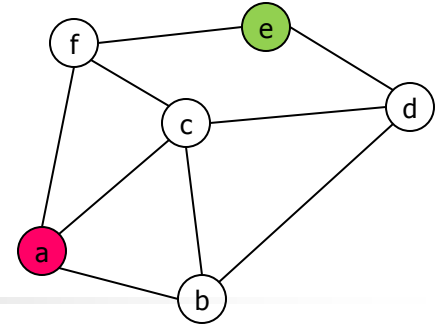
- Modeling in Python
 - Nodes can be represented as a **list**.
 - Edges can be represented as a **list of lists**.
- Example
 - Nodes as a **list**:
 - `N = ["a", "b", "c", "d", "e", "f"]`
 - Edges as a **list of lists**:
 - `D = [["b", "c", "f"],
 ["a", "c", "d"],
 ["a", "b", "d", "f"],
 ["b", "c", "e"],
 ["d", "f"],
 ["a", "c", "e"]]`

Adjacency List



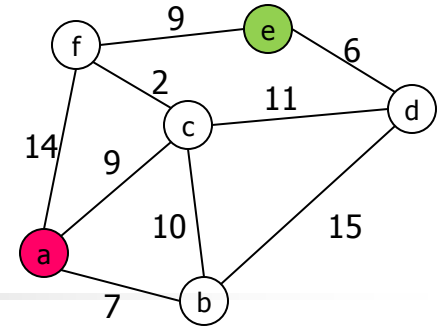
- Modeling in Python
 - Nodes can be represented as a list.
 - Edges can be represented as a list of sets.
- Example
 - Nodes as a list:
 - `N = ["a", "b", "c", "d", "e", "f"]`
 - Edges as a list of sets:
 - `D = [{"b", "c", "f"},
 {"a", "c", "d"},
 {"a", "b", "d", "f"},
 {"b", "c", "e"},
 {"d", "f"},
 {"a", "c", "e"}]`

Adjacency List



- Modeling in Python
 - Nodes can be represented as dictionary keys.
 - Edges can be represented as dictionary values in list/set.
- Example
 - Node and edges as a dictionary:
 - $D = \{ \text{"a": } \{ \text{"b"}, \text{"c"}, \text{"f"} \},$
 $\text{"b": } \{ \text{"a"}, \text{"c"}, \text{"d"} \},$
 $\text{"c": } \{ \text{"a"}, \text{"b"}, \text{"d"}, \text{"f"} \},$
 $\text{"d": } \{ \text{"b"}, \text{"c"}, \text{"e"} \},$
 $\text{"e": } \{ \text{"d"}, \text{"f"} \},$
 $\text{"f": } \{ \text{"a"}, \text{"c"}, \text{"e"} \} \}$

Adjacency List



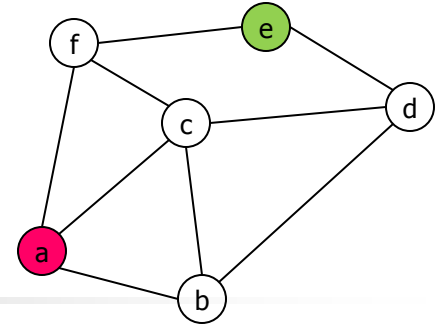
- Modeling in Python

- Nodes can be represented as a **list**.
- Edges with weights can be represented as a **list of lists of tuples**.

- Example

- Nodes as a **list**:
 - `N = ["a", "b", "c", "d", "e", "f"]`
- Edges as a **list of lists of tuples**:
 - `D = [[("b",7), ("c",9), ("f",14)],
[("a",7), ("c",10), ("d",15)],
[("a",9), ("b",10), ("d",11), ("f",2)],
[("b",15), ("c",11), ("e",6)],
[("d",6), ("f",9)],
[("a",14), ("c",2), ("e",9)]]`

Adjacency Matrix

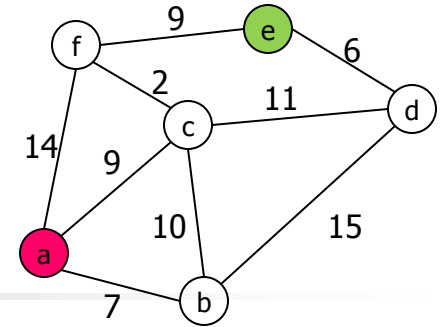


- For **each pair of nodes**, maintain a **matrix** to show the neighborhood between the pair.
 - An **adjacency matrix** has each node in a row (source) and in a column (destination).
 - **1 / True** means an edge and **0 / False** means no edge.
- Example
 - Matrix $D[6,6]$ for a graph with 6 nodes

■ $D =$

	a	b	c	d	e	f
a	0	1	1	0	0	1
b	1	0	1	1	0	0
c	1	1	0	1	0	1
d	0	1	1	0	1	0
e	0	0	0	1	0	1
f	1	0	1	0	1	0

Adjacency Matrix

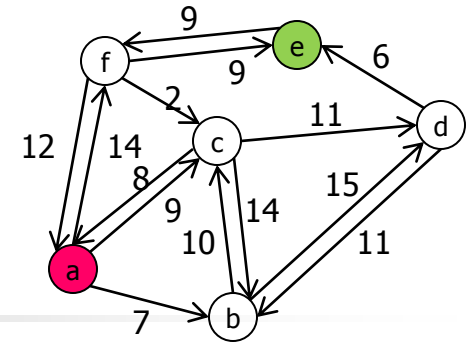


- An **adjacency matrix** for a graph with **weight** have matrix elements showing neighborhood and **weight**.
- This is often called a **distance matrix** when the weight is the distance.
- Example
 - Matrix $D[6,6]$ for a graph with 6 nodes

■ $D =$

	a	b	c	d	e	f
a	0	7	9	∞	∞	14
b	7	0	10	15	∞	∞
c	9	10	0	11	∞	2
d	∞	15	11	0	6	∞
e	∞	∞	∞	6	0	9
f	14	∞	2	∞	9	0

Adjacency Matrix

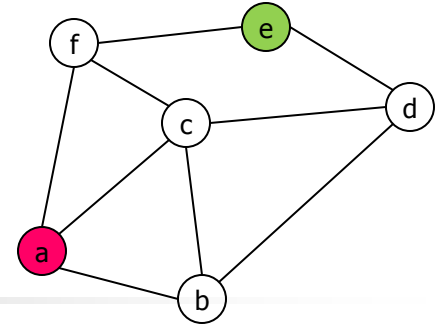


- An **distance matrix** for a **directed graph** with **weight** have matrix elements showing the **weight/distance** from node in row i to node in column j .
- Example
 - Matrix $D[6,6]$ for a graph with 6 nodes

■ $D =$

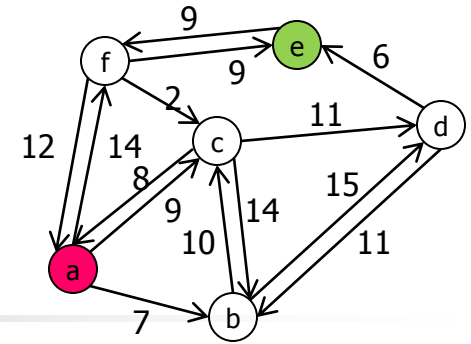
	a	b	c	d	e	f
a	0	7	9	∞	∞	14
b	∞	0	10	15	∞	∞
c	8	14	0	11	∞	∞
d	∞	11	∞	0	6	∞
e	∞	∞	∞	∞	0	9
f	12	∞	2	∞	9	0

Adjacency Matrix



- Modeling in Python
 - The 2-D matrix is normally represented as a list of lists as a logical 2-D array.
- Example
 - Nodes as a list:
 - `N = ["a", "b", "c", "d", "e", "f"]`
 - Matrix as a list of lists:
 - `D = [[0, 1, 1, 0, 0, 1],
 [1, 0, 1, 1, 0, 0],
 [1, 1, 0, 1, 0, 1],
 [0, 1, 1, 0, 1, 0],
 [0, 0, 0, 1, 0, 1],
 [1, 0, 1, 0, 1, 0]]`

Adjacency Matrix



- Modeling in Python

- The 2-D matrix can also be represented as a dictionary of dictionary.

- Example

- Nodes as a list:

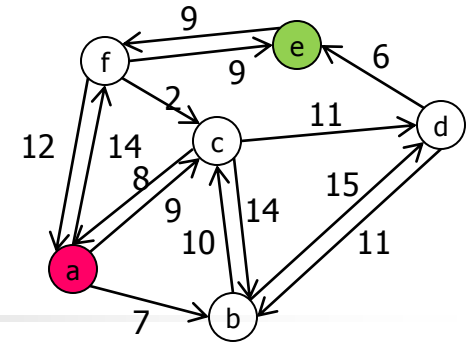
- `N = ["a", "b", "c", "d", "e", "f"]` Not really needed with d

- Matrix as a dictionary of dictionary.

- `D = { "a": {"a":0,"b":7,"c":9,"d":inf,"e":inf,"f":14},
"b": {"a":inf,"b":0,"c":10,"d":15,"e":inf,"f":inf},
"c": {"a":8,"b":14,"c":0,"d":11,"e":inf,"f":inf},
"d": {"a":inf,"b":11,"c":inf,"d":0,"e":6,"f":inf},
"e": {"a":inf,"b":inf,"c":inf,"d":inf,"e":0,"f":9},
"f": {"a":12,"b":inf,"c":2,"d":inf,"e":9,"f":0} }`

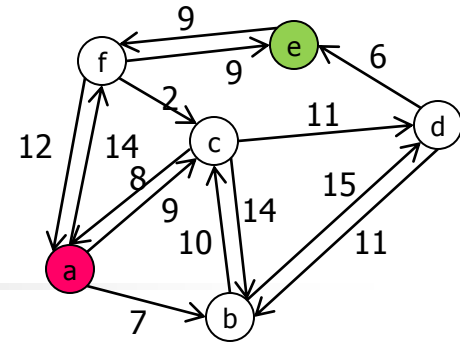
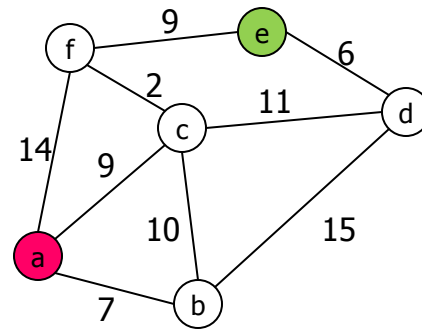
- Here, `inf = float("infinity")` in Python, meaning infinity (∞).

Adjacency Matrix



- Modeling in Python
 - The 2-D matrix can also be represented as a dictionary of dictionary.
- Example
 - Nodes as a list: Not really needed with dictionary D
 - `N = ["a", "b", "c", "d", "e", "f"]`
 - Matrix as a dictionary of dictionary.
 - `D = { "a": {"a":0,"b":7,"c":9,"d":inf,"e":inf,"f":14},
"b": {"a":inf,"b":0,"c":10,"d":15,"e":inf,"f":inf},
"c": {"a":8,"b":14,"c":0,"d":11,"e":inf,"f":inf},
"d": {"a":inf,"b":11,"c":inf,"d":0,"e":6,"f":inf},
"e": {"a":inf,"b":inf,"c":inf,"d":inf,"e":0,"f":9},
"f": {"a":12,"b":inf,"c":2,"d":inf,"e":9,"f":0} }`
 - Here, `inf = float("infinity")` in Python, meaning infinity (∞).

Shortest Path



- A highly common application on a graph is to find the **shortest path** from one node to another.

- The **starting** node is often called **source** node in graph theory.

The **target** node is often called **destination** node.

The graphs may contain **weights**, or without.

- We call them **weighted graphs** and **unweighted graphs** respectively.

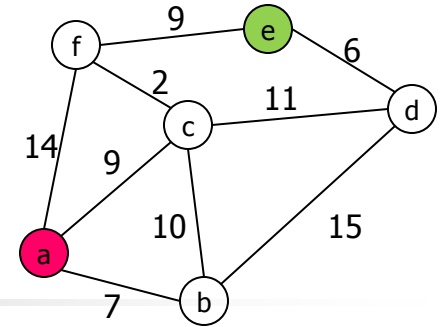
The graphs may contain edges with or without **directions**.

- We call these two types **directed graphs** and **undirected graphs** respectively.

- There are other applications on a graph, e.g. **breadth first search**, **depth first search**, **topological sort**.

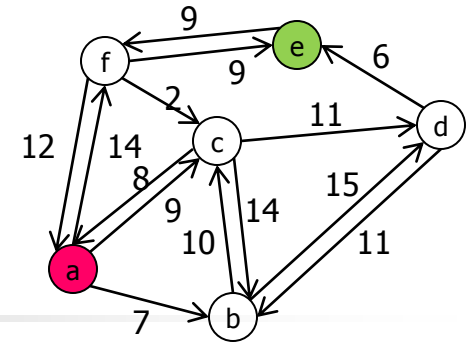


Shortest Path



- Can you find the **shortest path** from **a** to **e**?
 - Algorithm 1: Layman approach to find **all possible paths** first.
 - $a \rightarrow f \rightarrow e$ ■ 23
 - $a \rightarrow c \rightarrow f \rightarrow e$ ■ 20
 - $a \rightarrow c \rightarrow d \rightarrow e$ ■ 26
 - $a \rightarrow b \rightarrow c \rightarrow d \rightarrow e$ ■ 34
 - $a \rightarrow b \rightarrow d \rightarrow e$ ■ 28
 - $a \rightarrow c \rightarrow b \rightarrow d \rightarrow e$ ■ 40
 - Any more?
 - $a \rightarrow b \rightarrow c \rightarrow f \rightarrow e$ ■ 28
 - $a \rightarrow b \rightarrow d \rightarrow c \rightarrow f \rightarrow e$ ■ 44
 - Yet more?
 - $a \rightarrow f \rightarrow c \rightarrow d \rightarrow e$ ■ 33
 - $a \rightarrow f \rightarrow c \rightarrow b \rightarrow d \rightarrow e$ ■ 47
 - **Mission impossible** for larger graph!

Shortest Path

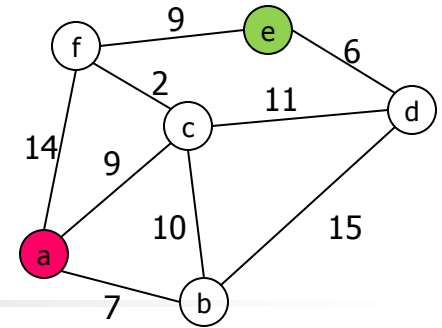


- Can you find the **shortest path** from **a** to **e**?
 - This is a **directed graph**.
 - There are **fewer** possible paths than previous one.

■ $a \rightarrow f \rightarrow e$	■ 23
■ $a \rightarrow c \rightarrow d \rightarrow e$	■ 26
■ $a \rightarrow b \rightarrow c \rightarrow d \rightarrow e$	■ 34
■ $a \rightarrow b \rightarrow d \rightarrow e$	■ 28
■ $a \rightarrow c \rightarrow b \rightarrow d \rightarrow e$	■ 44
■ $a \rightarrow f \rightarrow c \rightarrow d \rightarrow e$	■ 33
■ $a \rightarrow f \rightarrow c \rightarrow b \rightarrow d \rightarrow e$	■ 51
 - Still **mission impossible** for larger graph!
 - What is the **maximum number of possible paths** from a to e in this small graph with 6 nodes?

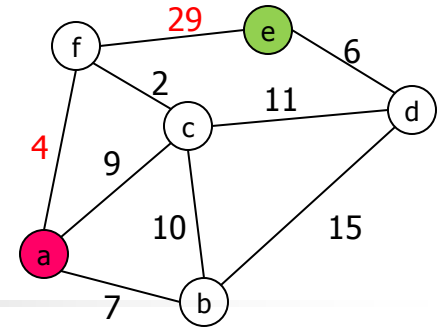
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Improvement



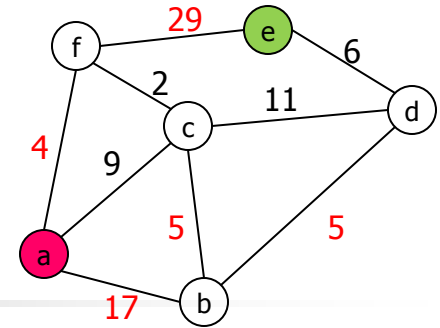
- There are so many possible paths.
- Algorithm 2: search for paths **starting with fewer steps.**
 - $a \rightarrow f \rightarrow e$ ■ 23
 - $a \rightarrow b \rightarrow d \rightarrow e$ ■ 28
 - $a \rightarrow c \rightarrow d \rightarrow e$ ■ 26
 - $a \rightarrow c \rightarrow f \rightarrow e$ ■ 20
 - Stop here?
 - $a \rightarrow b \rightarrow c \rightarrow d \rightarrow e$ ■ 34
 - $a \rightarrow b \rightarrow c \rightarrow f \rightarrow e$ ■ 28
 - $a \rightarrow c \rightarrow b \rightarrow d \rightarrow e$ ■ 40
 - $a \rightarrow f \rightarrow c \rightarrow d \rightarrow e$ ■ 33
 - Stop here?
 - $a \rightarrow b \rightarrow d \rightarrow c \rightarrow f \rightarrow e$ ■ 44
 - $a \rightarrow f \rightarrow c \rightarrow b \rightarrow d \rightarrow e$ ■ 47

Improvement



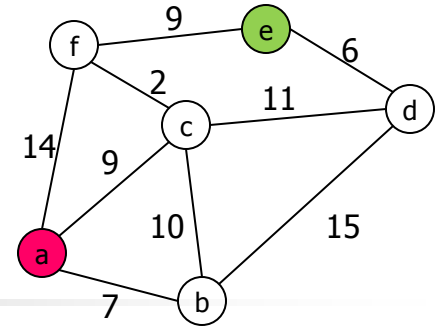
- Consider this graph.
- Search for paths **starting with fewer steps**.
 - $a \rightarrow f \rightarrow e$ ■ 33
 - $a \rightarrow b \rightarrow d \rightarrow e$ ■ 28
 - $a \rightarrow c \rightarrow d \rightarrow e$ ■ 26
 - $a \rightarrow c \rightarrow f \rightarrow e$ ■ 40
 - Stop here?
 - $a \rightarrow b \rightarrow c \rightarrow d \rightarrow e$ ■ 34
 - $a \rightarrow b \rightarrow c \rightarrow f \rightarrow e$ ■ 48
 - $a \rightarrow c \rightarrow b \rightarrow d \rightarrow e$ ■ 40
 - $a \rightarrow f \rightarrow c \rightarrow d \rightarrow e$ ■ 23
 - Stop here?
 - $a \rightarrow b \rightarrow d \rightarrow c \rightarrow f \rightarrow e$ ■ 64
 - $a \rightarrow f \rightarrow c \rightarrow b \rightarrow d \rightarrow e$ ■ 37

Improvement

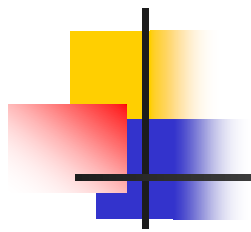


- Consider this graph.
- Search for paths **starting with fewer steps**.
 - $a \rightarrow f \rightarrow e$ ■ 33
 - $a \rightarrow b \rightarrow d \rightarrow e$ ■ 28
 - $a \rightarrow c \rightarrow d \rightarrow e$ ■ 26
 - $a \rightarrow c \rightarrow f \rightarrow e$ ■ 40
 - Stop here?
 - $a \rightarrow b \rightarrow c \rightarrow d \rightarrow e$ ■ 39
 - $a \rightarrow b \rightarrow c \rightarrow f \rightarrow e$ ■ 53
 - $a \rightarrow c \rightarrow b \rightarrow d \rightarrow e$ ■ 25
 - $a \rightarrow f \rightarrow c \rightarrow d \rightarrow e$ ■ 23
 - Stop here?
 - $a \rightarrow b \rightarrow d \rightarrow c \rightarrow f \rightarrow e$ ■ 64
 - $a \rightarrow f \rightarrow c \rightarrow b \rightarrow d \rightarrow e$ ■ 22

Improvement



- We may miss the correct answer if not trying out all possible paths.
 - We **cannot afford** to try all possible paths in larger graph.
 - We need a **more clever and systematic** approach.
- Possible approach:
 - Look at **edges** and use them to **improve** on existing paths.
 - An edge is said to lead to improvement, if **passing** through it would lead to a better path.
 - Example:
 - Going from **a to c** directly, the distance is 9.
 - Going from **a to f** directly, the distance is 14.
 - Going from **a to f via c**, the distance improves to **11** < 14.
 - We will try to implement this idea.



End of Lecture 9