## Hadoop & HDFS

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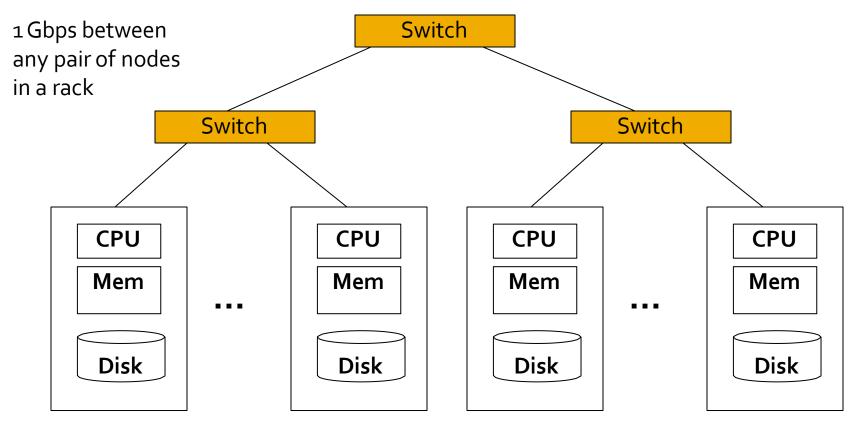
## Hadoop

 A large-scale distributed & parallel batchprocessing infrastructure

- Large-scale:
  - Handle a large amount of data and computation
- Distributed:
  - Distribute data & computation over multiple machines
- Batch processing
  - Process a series of jobs without human intervention

#### **Cluster Architecture**

#### 2-10 Gbps backbone between racks



Each rack contains 16-64 nodes

In 2011 it was guestimated that Google had 1M machines, <a href="http://bit.ly/ShhoRO">http://bit.ly/ShhoRO</a>



### History

- 1<sup>st</sup> version released by Yahoo! in 2006
  - named after an elephant toy

- Originated from Google's work
  - GFS: Google File System (2003)
  - MapReduce (2004)



## Roadmap

Hadoop architecture



- HDFS
- MapReduce

Installing Hadoop & HDFS

#### Key components

- HDFS (Hadoop distributed file system)
  - Distributed data storage with high reliability

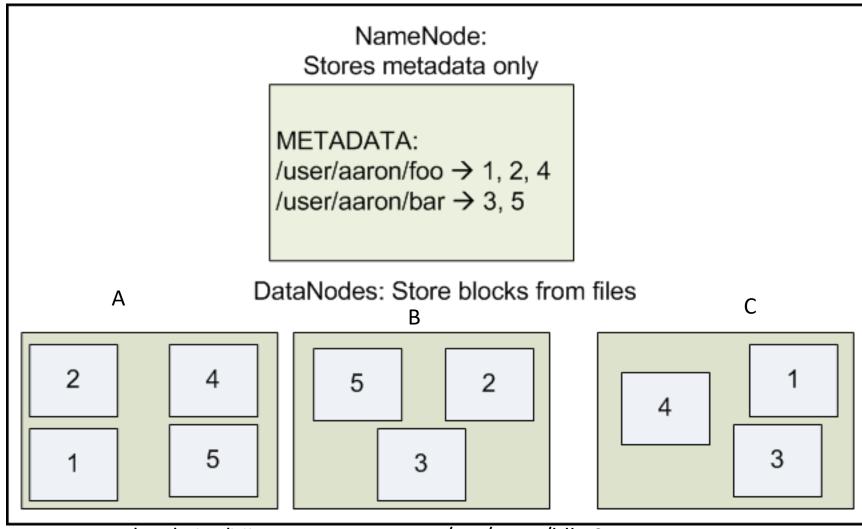
- MapReduce
  - A parallel, distributed computational paradigm
  - With a simplified programming model

#### **HDFS**

- Data are distributed among multiple data nodes
  - Data nodes may be added on demand for more storage space

- Data are replicated to cope with node failure
  - Typically replication factor: 2 or 3
- Requests can go to any replica
  - Removing the bottleneck (as in single file server)

#### HDFS architecture



/usr/john/blk\_5\_1.csv

/usr/mary/blk\_3\_1.csv

writing a block in hdfs is done in pipeline. it might start from B blocks can be broken down into packets. so instead of large block as 256MB, hdfs would break it into small packets of like 64KB and send it to the pipeline.

B gets it and write it to the storage, and forward to machine C

#### HDFS has ...

- A single NameNode, storing meta data:
  - A hierarchy of directories and files (name space)
  - Attributes of directories and files (in inodes), e.g., permission, access/modification times, etc.
  - Mapping of files to blocks on data nodes

- A number of DataNodes:
  - Storing contents/blocks of files

#### Compute nodes

Data nodes are compute nodes too

这什么意思啊

- Advantage:
  - Allow schedule computation close to data

#### HDFS also has ...

- A SecondaryNameNode
  - Maintaining checkpoints/images of NameNode
  - For recovery
  - not a fail over node

- In a single-machine setup
  - all nodes correspond to the same machine

#### Metadata in NameNode

同p10

NameNode has an inode for each file and dir

- Record attributes of file/dir such as
  - Permission
  - Access time
  - Modification time

Also record mapping of files to blocks

#### Mapping information in NameNode

E.g., file /user/aaron/foo consists of blocks 1,
2, and 4

- Block 1 is stored on data nodes 1 and 3
- Block 2 is stored on data nodes 1 and 2

• ...

#### Block size

- HDFS: 128 MB (version 2 & above)
  - Much larger than disk block size (4KB)
  - A: 128MB; B: 4KB
  - -128MB/4KB = 32K
  - A: 1GB/128MB = 8; B: 1GB/4KB = 2^30/2^12 = 2^18 = 2^8K = 256K
- Why larger size in HDFS?
  - Reduce metadata required per file
  - Fast streaming <u>read</u> of data (since larger amount of data are sequentially laid out on disk)

#### **HDFS**

HDFS exposes the concept of blocks to client

- Reading and writing are done in two phases
  - Phase 1: client asks NameNode for block locations
    - By calling (sending request) getBlockLocations(), if reading
    - Or calling <u>addBlock()</u> for allocating new blocks (one at a time), if <u>writing</u> (need to call <u>create()/append()</u> first)
  - Phase 2: client talks to DataNode for data transfer
    - Reading blocks via <u>readBlock()</u> or writing blocks via <u>writeBlock()</u>

getBlockLocations(<file\_path>, <offset>, <size of data to be read>)
具体见p20

```
readBlock()
```

if we read the 1st 100MB of data. only first block will be visited.

if we read the 2nd 100MB of data, both the first and second block will be visited:

```
block 1: offset = 100MB, length = 28MB
```

block 2: offset = 0, length = 72MB

```
e.g. foo has 200MB, block 1 (128MB), block 2 (72MB)
```

client reads 2nd 100MB

- 1. talk to NameNode, getBlockLocations(...) block 1, block 2
- 2. talk to DataNode, readBlock(...) readBlock(block1, 100MB, 28MB) readBlock(block2, 0, 72MB)

#### Client and Namenode communication

- Source code (version 2.8.1)
  - Definition of protocol
    - ClientNamenodeProtocol.proto
    - <hadoop-src-dir>\hadoop-hdfs-project\hadoop-hdfsclient\src\main\proto
  - Implementation
    - ClientProtocol.java
    - <hadoop-src-dir>\hadoop-hdfs-project\hadoop-hdfsclient\src\main\java\org\apache\hadoop\hdfs\protocol

#### Key operations

of Phase 1: client asks NameNode for block locations

- Reading:
  - getBlockLocations()

- Writing
  - create()
  - append()
  - addBlock()

### getBlockLocations

Before reading, client needs to first obtain locations of blocks

```
message GetEditsFromTxidResponseProto {
  required EventsListProto eventsList = 1;
service ClientNamenodeProtocol {
  rpc getBlockLocations(GetBlockLocationsRequestProto)
      returns (GetBlockLocationsResponseProto);
  rpc getServerDefaults(GetServerDefaultsRequestProto)
      returns (GetServerDefaultsResponseProto);
  rpc create (CreateRequestProto) returns (CreateResponseProto);
  rpc append (AppendRequestProto) returns (AppendResponseProto);
  rpc setReplication(SetReplicationRequestProto)
      returns (SetReplicationResponseProto);
  rpc setStoragePolicy(SetStoragePolicyRequestProto)
      returns (SetStoragePolicyResponseProto);
  rpc getStoragePolicies(GetStoragePoliciesRequestProto)
      returns (GetStoragePoliciesResponseProto);
  rpc setPermission(SetPermissionRequestProto)
      returns (SetPermissionResponseProto);
```

## getBlockLocations

#### Input:

- File name
- Offset (to start reading)
- Length (how much data to be read)

#### Output:

Located blocks (data nodes + offsets)

```
// File contents
/**
 * Get locations of the blocks of the specified file
 * within the specified range.
 * DataNode locations for each block are sorted by
 * the proximity to the client.
 * 
 * Return {@link LocatedBlocks} which contains
 * file length, blocks and their locations.
 * DataNode locations for each block are sorted by
 * the distance to the client's address.
 * 
 * The client will then have to contact.
 * one of the indicated DataNodes to obtain the actual data.
 * @param src file name
 * @param offset range start offset
 * @param length range length
 * @return file length and array of blocks with their locations
 * @throws org.apache.hadoop.security.AccessControlException If access is
            denied
 * @throws java.io.FileNotFoundException If file <code>src</code> does not
            exist.
 * @throws org.apache.hadoop.fs.UnresolvedLinkException If <code>src</code>
            contains a symlink
 * @throws IOException If an I/O error occurred
 */
@Idempotent
LocatedBlocks getBlockLocations (String src, long offset, long length)
   throws IOException;
```

```
../java/...hdfs/protocol/LocatedBlocks.java
```

```
public class LocatedBlocks {
  private final long fileLength;
  // array of blocks with prioritized locations
  private final List<LocatedBlock> blocks;
  private final boolean underConstruction;
                                                                      Block
  private final LocatedBlock lastLocatedBlock;
                                                                      Offset of this block
  private final boolean isLastBlockComplete;
  private final FileEncryptionInfo fileEncryptionInfo;
                                                                        in the entire file
                                                                      Data nodes with
          public class LocatedBlock {
  public
                                                                         replicas of block
    fileI
           private final ExtendedBlock b;
           private long offset; // offset of the first byte of the block in the file
    block
           private final DatanodeInfoWithStorage[] locs; '
    under
           /** Cached storage ID for each replica */
    lastI
           private final String[] storageIDs;
    isLas
           /** Cached storage type for each replica, if reported. */
    fileE
           private final StorageType[] storageTypes;
            // corrupt flag is true if all of the replicas of a block are corrupt.
            // else false. If block has few corrupt replicas, they are filtered and
            // their locations are not part of this object
            private boolean corrupt;
            private Token<BlockTokenIdentifier> blockToken = new Token<BlockTokenIdentifie
</pre>
          r>();
            /**
             * List of cached datamode locations
             */
           private DatanodeInfo[] cachedLocs;
            // Used when there are no locations
           private static final DatanodeInfoWithStorage[] EMPTY LOCS =
               new DatanodeInfoWithStorage[0];
```

## Create/append a file

```
message GetEditsFromTxidResponseProto {
                                                This opens the file for
  required EventsListProto eventsList = 1;
                                                   create/append
service ClientNamenodeProtocol {
  rpc getBlockLocations(GetBlockLocationsRequestProto)
      returns (GetBlockLocationsResponseProto);
  rpc getServerDefaults(GetServerDefaultsReguestProto)
      returns (GetServerDefaultsResponseProto);
  rpc create (CreateRequestProto) returns (CreateResponseProto);
  rpc append(AppendRequestProto) returns(AppendResponseProto);
  rpc setReplication(SetReplicationRequestProto)
      returns (SetReplicationResponseProto);
  rpc setStoragePolicy(SetStoragePolicyRequestProto)
      returns (SetStoragePolicyResponseProto);
  rpc getStoragePolicies(GetStoragePoliciesRequestProto)
      returns (GetStoragePoliciesResponseProto);
  rpc setPermission(SetPermissionRequestProto)
      returns (SetPermissionResponseProto);
```

## Creating a file

- Needs to specify:
  - Path to the file to be created, e.g., /foo/bar
  - Permission mask
  - Client name
  - Flag on whether to overwrite (entire file!) if already exists
  - How many replicas
  - Block size

```
/**
                                                   A hierarchy of files and directories
 * Create a new file entry in the namespace.
 * This will create an empty file specified by the source path.
 * The path should reflect a full path originated at the root.
 * The name-node does not have a notion of "current" directory for a client.
 * 
 * Once created, the file is visible and available for read to other clients.
 * Although, other clients cannot {@link #delete(String, boolean)}, re-create
 * or {@link #rename(String, String)} it until the file is completed
 * or explicitly as a result of lease expiration.
 * 
 * Blocks have a maximum size. Clients that intend to create
 * multi-block files must also use
 * {@link #addBlock}
 * @param src path of the file being created.
 * @param masked masked permission.
 * @param clientName name of the current client.
 * @param flag indicates whether the file should be
 * overwritten if it already exists or create if it does not exist or append.
 * @param createParent create missing parent directory if true
 * @param replication block replication factor.
 * @param blockSize maximum block size.
 * @param supportedVersions CryptoProtocolVersions supported by the client
                            Creating a new file
@AtMostOnce
HdfsFileStatus create(String src, FsPermission masked,
    String clientName, EnumSetWritable<CreateFlag> flag,
   boolean createParent, short replication, long blockSize,
    CryptoProtocolVersion[] supportedVersions)
   throws IOException;
```

## Allocating new blocks for writing

Asking NameNode to allocate a new block + data nodes holding its replicas

```
rpc setPermission(SetPermissionRequestProto)
    returns (SetPermissionResponseProto);
rpc setOwner(SetOwnerRequestProto) returns(SetOwnerResponseProto);
rpc abandonBlock(AbandonBlockRequestProto) returns(AbandonBlockResponseProto);
rpc addBlock(AddBlockRequestProto) returns(AddBlockResponseProto);
rpc getAdditionalDatanode(GetAdditionalDatanodeReguestProto)
    returns (GetAdditionalDatanodeResponseProto);
rpc complete (CompleteRequestProto) returns (CompleteResponseProto);
rpc reportBadBlocks(ReportBadBlocksRequestProto)
    returns (ReportBadBlocksResponseProto);
rpc concat(ConcatRequestProto) returns(ConcatResponseProto);
rpc truncate(TruncateRequestProto) returns(TruncateResponseProto);
rpc rename (RenameRequestProto) returns (RenameResponseProto);
rpc rename2(Rename2RequestProto) returns(Rename2ResponseProto);
rpc delete(DeleteRequestProto) returns(DeleteResponseProto);
rpc mkdirs(MkdirsRequestProto) returns(MkdirsResponseProto);
rpc getListing(GetListingRequestProto) returns(GetListingResponseProto);
rpc renewLease(RenewLeaseRequestProto) returns(RenewLeaseResponseProto);
```

```
/**
 * A client that wants to write an additional block to the
 * indicated filename (which must currently be open for writing)
 * should call addBlock().
  addBlock() allocates a new block and datanodes the block data
  should be replicated to.
 * addBlock() also commits the previous block by reporting
 * to the name-node the actual generation stamp and the length
 * of the block that the client has transmitted to data-nodes.
 * @param src the file being created
 * @param clientName the name of the client that adds the block
 * @param previous previous block
 * @param excludeNodes a list of nodes that should not be
 * allocated for the current block
 * @param fileId the id uniquely identifying a file
 * @param favoredNodes the list of nodes where the client wants the blocks.
            Nodes are identified by either host name or address.
 * @param addBlockFlags flags to advise the behavior of allocating and placing
                        a new block.
 * @return LocatedBlock allocated block information.
 * ...
@Idempotent
LocatedBlock addBlock String src, String clientName,
    ExtendedBlock previous, DatanodeInfo[] excludeNodes, long fileId,
    String[] favoredNodes, EnumSet<AddBlockFlag> addBlockFlags)
    throws IOException;
```

#### Client and Datanode communication

- Source code (version 2.8.1)
  - Definition of protocol
    - datatransfer.proto
    - Located at: <hadoop-src-dir>\hadoop-hdfsproject\hadoop-hdfs-client\src\main\proto
  - Implementation
    - DataTransferProtocol.java
    - <hadoop-src-dir>\hadoop-hdfs-project\hadoop-hdfsclient\src\main\java\org\apache\hadoop\hdfs\protocol \datatransfer

#### Operations

of Phase 2: client talks to DataNode for data transfer

readBlock()

writeBlock()

copyBlock() – for load balancing

- replaceBlock() for load balancing
  - Move a block from one DataNode to another

## Reading a file

- Client first contacts NameNode which informs the client of the closest DataNodes storing blocks of the file
  - This is done by making which RPC call? getBlockLocations?
- 2. Client contacts the DataNodes directly for reading the blocks

  NameNode:
  Stores metadata only
  - Calling readBlock()

DataNodes: Store blocks from files

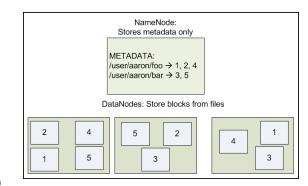
## datatransfer.proto

```
message OpReadBlockProto {
  required ClientOperationHeaderProto header = 1;
                                                       Block, offset, length
  required uint64 offset = 2; _
  required uint64 len = 3; ←
  optional bool sendChecksums = 4 [default = true];
  optional CachingStrategyProto cachingStrategy = 5;
message ChecksumProto {
  required ChecksumTypeProto type = 1;
  required uint32 bytesPerChecksum = 2;
message OpWriteBlockProto {
  required ClientOperationHeaderProto header = 1;
  repeated DatanodeInfoProto targets = 2;
  optional DatanodeInfoProto source = 3;
  enum BlockConstructionStage {
    PIPELINE SETUP APPEND = 0;
    // pipeline set up for failed PIPELINE SETUP APPEND recovery
    PIPELINE SETUP APPEND RECOVERY = 1;
    // data atmosmina
```

## DataTransferProtocol.java

```
/**
 * Read a block.
 * @param blk the block being read.
 * @param blockToken security token for accessing the block.
 * @param clientName client's name.
 * @param blockOffset offset of the block.
 * @param length maximum number of bytes for this read.
 * @param sendChecksum if false, the DN should skip reading and sending
          checksums
 * @param cachingStrategy The caching strategy to use.
 */
public void readBlock(final ExtendedBlock blk, 
    final Token<BlockTokenIdentifier> blockToken Block, offset, length
    final String clientName,
    final long blockOffset,
    final long length, <
    final boolean sendChecksum,
    final CachingStrategy cachingStrategy) throws IOException;
 * Write a block to a datanode pipeline.
 * The receiver datanode of this call is the next datanode in the pipeline.
 * The other downstream datanodes are specified by the targets parameter.
 * Note that the receiver {@link DatanodeInfo} is not required in the
 * parameter list since the receiver datanode knows its info. However, the
 * {@link StorageType} for storing the replica in the receiver datanode is a
 * parameter since the receiver datanode may support multiple storage types.
```

## Writing a file



- Blocks are written one at a time
  - In a pipelined fashion through the data nodes

For each block:

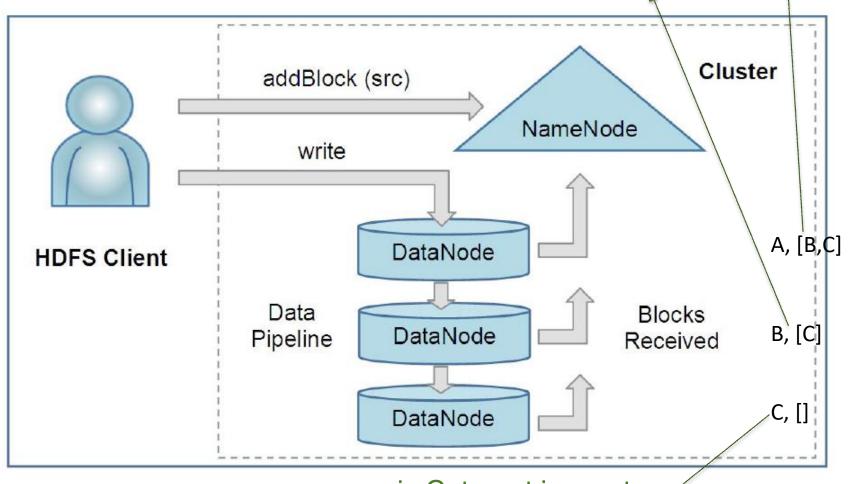
高亮这句会考到, 背诵

- Client asks NameNode to select DataNodes for holding its replica (using which rpc call?) addBlock()
  - e.g., DataNodes 1 and 3 for the first block of /user/aaron/foo
- It then forms the pipeline to send the block

#### source is A, target is B and C

# Writing a file

source is B, target is C



source is C, target is empty. In now every replica are up-to-date

```
* Write a block to a datanode pipeline.
 * The receiver datanode of this call is the next datanode in the pipeline.
 * The other downstream datanodes are specified by the targets parameter.
 * Note that the receiver {@link DatanodeInfo} is not required in the
 * parameter list since the receiver datanode knows its info. However, the
 * {@link StorageType} for storing the replica in the receiver datanode is a
 * parameter since the receiver datanode may support multiple storage types.
 * @param blk the block being written.
 * @param storageType for storing the replica in the receiver datanode.
 * @param blockToken security token for accessing the block.
 * @param clientName client's name.
 * @param targets other downstream datanodes in the pipeline.
 * @param targetStorageTypes target {@link StorageType}s corresponding
                             to the target datanodes.
 * @param source source datanode.
 * @param stage pipeline stage.
 * @param pipelineSize the size of the pipeline.
 * @param minBytesRcvd minimum number of bytes received.
 * @param maxBytesRcvd maximum number of bytes received.
 * @param latestGenerationStamp the latest generation stamp of the block.
 * @param pinning whether to pin the block, so Balancer won't move it.
 * @param targetPinnings whether to pin the block on target datanode
 */
void writeBlock (final ExtendedBlock blk, Block to be written
   final StorageType storageType,
   final Token<BlockTokenIdentifier> blockToken, Rest of data nodes
    final String clientName,
    final DatanodeInfo[] targets,
    final StorageType[] targetStorageTypes,
                                             Current data node in the pipeline
    final DatanodeInfo source,
    final BlockConstructionStage stage,
    final int pipelineSize,
    final long minBytesRcvd,
   final long maxBytesRcvd,
```

# Data pipelining

Consider a block X to be written to DataNode
 A, B, and C (replication factor = 3)

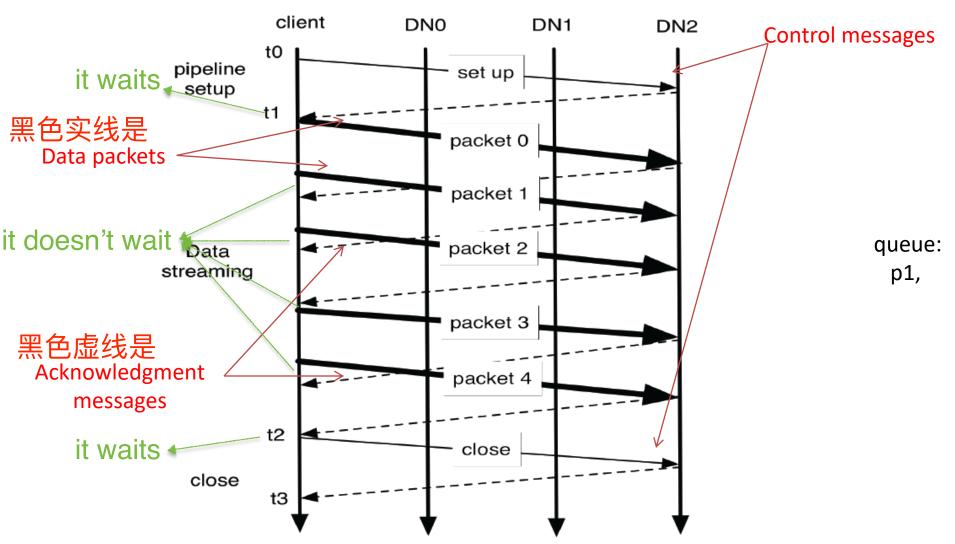
- 1. X is broken down into packets (typically 64KB/packet)
  - 128MB/64KB = 2048
- 2. Client sends the packet to DataNode A
- 3. A sends it further to B & B further to C

#### Acknowledgement

- Client maintains an ack (acknowledgment) queue
- Packet removed from ack queue once received by all data nodes

- When all packets were written, client notifies NameNode
  - NameNode will update the metadata for the file
  - Reflecting that a new block has been added to the file

# Data pipelining for writing blocks



At the time t1, all three dataNodes are ready to be streamed.

before the client receives the ack of packet 0, the client already started sending packet 1.

The client is not waiting. This is asynchronous. This is called non-blocking(?)

Why no waiting? because 他没说。。

How does the client make sure all the packets are received? the client need to maintain some data structure to store the status of all the acknowledgements (probably called ack queue).

## Acknowledgement

 Client does not wait for the acknowledgement of previous packet before sending next one

Is this synchronous or asynchronous?

Advantage?

他没说。只说因此需要ack queue来记录哪些顺利收到了

#### Roadmap

- Hadoop architecture
  - HDFS
  - MapReduce

Installing Hadoop & HDFS



## Hadoop & HDFS installation

 Refer to the installation note posted on course web site on how to install Hadoop and setup HDFS

- Setting up home directory in hdfs
  - hdfs dfs -mkdir /user
  - hdfs dfs -mkdir /user/ec2-user
     (ec2-user is user name of your EC2 account)
- Create a directory "input" under home
  - hdfs dfs -mkdir /user/ec2-user/input
  - Or simply:
  - hdfs dfs -mkdir input

- Copy data from local file system
  - hdfs dfs -put etc/hadoop/\*.xml /user/ec2user/input
  - Ignore error if you see one like this: "WARN hdfs.
     DataStreamer: Caught exception..."

- List the content of directory
  - hdfs dfs -ls /user/ec2-user/input

- Copy data from hdfs
  - hdfs dfs -get /user/ec2-user/input input1
  - If input1 does not exist, it will create one
  - If it does, it will create another one under it

- Examine the content of file in hdfs
  - hdfs dfs -cat /user/ec2-user/input/core-site.xml

- Remove files
  - hdfs dfs -rm /user/ec2-user/input/core-site.xml
  - hdfs dfs -rm /user/ec2-user/input/\*

- Remove directory
  - hdfs dfs -rmdir /user/ec2-user/input
  - Directory "input" needs to be empty first

#### Where is hdfs located?

/tmp/hadoop-ec2-user/dfs/

#### References

 K. Shvachko, H. Kuang, S. Radia, and R. Chansler, "The hadoop distributed file system," in Mass Storage Systems and Technologies (MSST), 2010 IEEE 26th Symposium on, 2010, pp. 1-10.

- HDFS File System Shell Guide:
  - https://hadoop.apache.org/docs/current/hadoopproject-dist/hadoop-common/FileSystemShell.html