Constraints

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Kinds of Constraints

- Keys
- Foreign-key, or referential-integrity
- Value-based constraints
 - Constrain values of a particular attribute
- Tuple-based constraints
 - Constrain relationship among attributes
- Assertions: any SQL boolean expression
 - Very expressive

Keys

Specified using "primary key" or "unique"

Foreign Keys

- Consider Relation Sells(bar, beer, price).
- We might expect that a beer value is a real beer --- something appearing in Beers.name.
- A constraint that requires a beer in Sells to be a beer in Beers is called a *foreign -key* constraint.

Expressing Foreign Keys

- Use the keyword REFERENCES, either:
 - 1. Within the declaration of an attribute, when only one attribute is involved, or
 - 2. As an element of the schema, as:

```
FOREIGN KEY ( < list of attributes > )
REFERENCES < relation > ( < attributes > )
```

 Note MySQL seems to enforce FK only when defined as an element

Example: Express FK with Attribute

```
CREATE TABLE Beers (
        CHAR (20) PRIMARY KEY,
 name
 manf CHAR(20);
CREATE TABLE Sells (
        CHAR (20),
 bar
 beer CHAR(20) REFERENCES Beers(name),
 price REAL );
```

Example: Express FK as Element

```
CREATE TABLE Beers (
        CHAR (20) PRIMARY KEY,
 name
 manf CHAR(20);
CREATE TABLE Sells (
 bar CHAR(20),
 beer CHAR(20),
                     Note parentheses are necessary!
 price REAL,
 FOREIGN KEY (beer) REFERENCES
    Beers (name));
```

Primary Key vs. Unique

- Referenced attributes must be declared as PRIMARY KEY or UNIQUE.
 - Otherwise, MySQL does not allow creation of the table
 - Note that primary key can not be null, but unique attribute can

- Null values can be inserted into attribute of foreign key
 - Even though it refers to primary key in referenced table

Example of FKs with Unique Attributes

- create table R (a int primary key);
 - insert into R values (1);
 - Select * from R;

Or "a int unique"

- create table S(b int, foreign key (b) references
 R(a));
 - insert into S values (1);
 - insert into S values (null); // this works even though"a" is primary key in R
 - select * from S;

Enforcing Foreign-Key Constraints

- If there is a foreign-key constraint from attributes of relation *S* to the primary key (or unique attribute) of relation *R*, two violations are possible:
 - 1. An insert or update to *S* introduces values not found in *R*.
 - 2. A deletion or update to R causes some tuples of S to "dangle."

Actions Taken

- Suppose R = Beers, S = Sells.
- An insert or update to Sells that introduces a nonexistent beer must be rejected.
- A deletion or update to Beers that removes a beer value found in some tuples of Sells can be handled in three ways.

 Beers

Actions Taken (Cont'd)

- The three possible ways to handle beers that suddenly cease to exist are:
 - 1. Default: Reject the modification.
 - 2. Cascade: Make the same changes in Sells.
 - ☐ Deleted beer: delete Sells tuple.
 - ☐ Updated beer: change value in Sells.
 - 3. Set NULL: Change the beers in Sells to NULL.

Example: Cascade

- Suppose we delete the Bud tuple from Beers.
 - Then delete all tuples from Sells that have beer = 'Bud'.
- Suppose we update the Bud tuple by changing 'Bud' to 'Budweiser'.
 - Then change all Sells tuples with beer = 'Bud' so that beer = 'Budweiser'.

Example: Set NULL

- Suppose we delete the Bud tuple from Beers.
 - Change all tuples of Sells that have beer = 'Bud' to have beer = NULL.
- Suppose we update the Bud tuple by changing 'Bud' to 'Budweiser'.
 - Same change.

Choosing a Policy

- When we declare a foreign key, we may choose policies SET NULL or CASCADE independently for deletions and updates.
- Follow the foreign-key declaration by:
- ON [UPDATE, DELETE][SET NULL/CASCADE]
- Two such clauses may be used.
- Otherwise, the default (reject) is used.

Example

```
CREATE TABLE Sells (
 bar CHAR(20),
 beer CHAR (20),
 price REAL,
 FOREIGN KEY (beer)
   REFERENCES Beers (name)
   ON DELETE SET NULL
   ON UPDATE CASCADE );
```

Multi-attribute keys (unique, PK, FK)

```
create table R(a int, b int, unique(a, b));
insert into R values(1, null);
                                                          Are
create table P(a int, b int, primary key(a, b));
                                                       they ok?
insert into P values(1, 2);
insert into R values(2, null);
create table F(a int, b int, foreign key (a, b)
        references P (a, b));
insert into F values(1, null);
create table F1 (a int, foreign key (a) references P(a));
insert into F1 values(2);
```

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- Constrain values of a particular attribute.
- Tuple-based constraints.
 - Relationship among components.
- Assertions: any SQL boolean expression.

Attribute-Based Checks

- Put a constraint on the value of a particular attribute.
- CHECK(<condition>) must be added to the declaration for the attribute.
- The condition may use the name of the attribute, but any other relation or attribute name must be in a subquery.
- Note: MySQL does not seem to support this
 - Accept definition, but does not enforce it
- Other DBMS, e.g., PostgreSQL, may support it

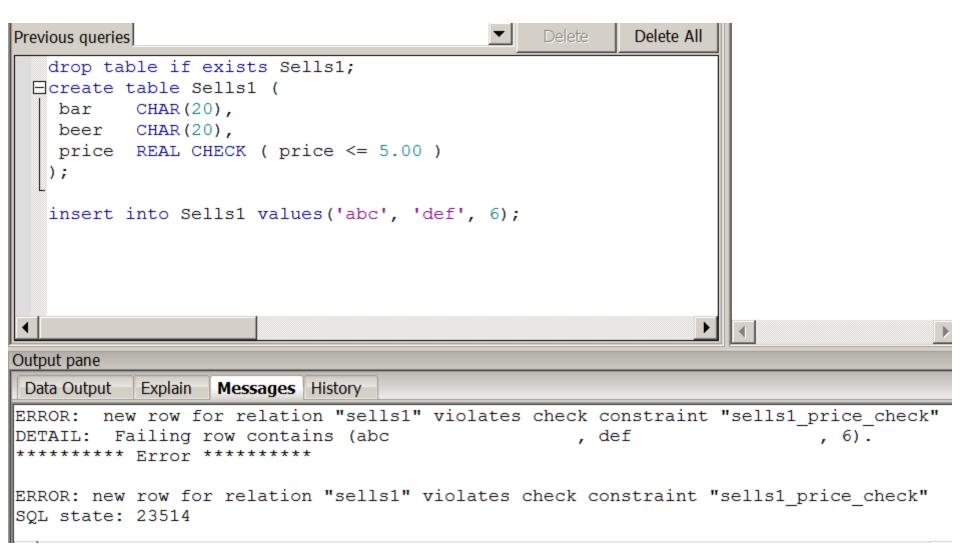
Example

```
eg.
CREATE TABLE Employees (
  EmployeeID INT PRIMARY KEY,
  Name VARCHAR(50),
  Age INT CHECK (Age >= 18 AND Age <= 65),
  Department VARCHAR(30)
);
eg.
CREATE TABLE Customers (
  CustomerID INT PRIMARY KEY,
  Name VARCHAR(50),
  Email VARCHAR(50) CHECK (Email LIKE '%@%.__%'),
  Phone VARCHAR(20) CHECK (Phone LIKE '+__%')
);
```

Timing of Checks

- An attribute-based check is checked only when a value for that attribute is inserted or updated.
 - Example: CHECK (price <= 5.00) checks every new price and rejects it if it is more than \$5.
 - Example: CHECK (beer IN (SELECT name FROM Beers)) not checked if a beer is deleted from Beers (unlike foreign-keys).

PostgreSQL example



Tuple-Based Checks

- CHECK (<condition>) may be added as another element of a schema definition.
- The condition may refer to any attribute of the relation, but any other attributes or relations require a subquery.
- Checked on insert or update only.

Example: Tuple-Based Check

• Only Joe's Bar can sell beer for more than \$5:

```
CREATE TABLE Sells (
bar CHAR(20),
beer CHAR(20),
price REAL,
CHECK (bar = 'Joe' OR
price <= 5.00)
);
```

Example (work in PostgreSQL)

- insert into sells values('Joe', 'bud', 8);
 - This insert is ok

- update sells set bar = 'joe1'
 - This update is not ok

Assertions

- These are database-schema elements, like relations or views.
- Defined by:

CREATE ASSERTION < name > CHECK (< condition >);

• Condition may refer to any relation or attribute in the database schema.

- Very expensive to enforce
 - Neither PostgreSQL nor MySQL supports this

Example: Assertion

• In Sells(bar, beer, price), no bar may charge an average of more than \$5.

CREATE ASSERTION NoRipoffBars CHECK (
NOT EXISTS (

SELECT bar FROM Sells
GROUP BY bar
HAVING 5.00 < AVG(price)

Bars with an average price above \$5

Example: Assertion

• In Drinkers(name, addr, phone) and Bars(name, addr, license), there cannot be more bars than drinkers.

```
CREATE ASSERTION FewBar CHECK (
   (SELECT COUNT(*) FROM Bars) <=
   (SELECT COUNT(*) FROM Drinkers)
);</pre>
```

Timing of Assertion Checks

- In principle, we must check every assertion after every modification to any relation of the database.
- A clever system can observe that only certain changes could cause a given assertion to be violated.
 - Example: No change to Beers can affect FewBar.
 Neither can an insertion to Drinkers.