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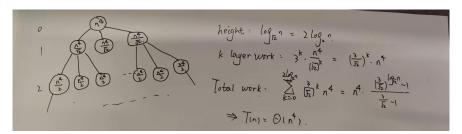
# Problem 2-1.

(a)  $T(n) = \Theta(n^2)$ 

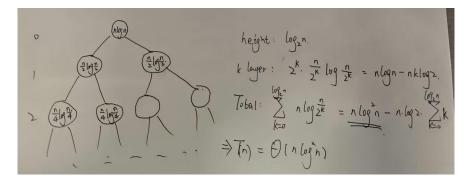
total height.  $\Rightarrow \log_2 n$ The k layer:  $4^k \cdot \frac{n}{2^k} = n \cdot 2^k$ .

Total work:  $\lim_{k \to 0} n \cdot 2^k = n \cdot 2^k = n$   $\Rightarrow T(n) = \theta(n^2)$ .

(b)  $T(n) = \Theta(n^4)$ 



(c)  $T(n) = \Theta(n \log^2 n)$ 



(d) 
$$T(n) = \Theta(n^2)$$

#### Problem 2-2.

### (a) Selection Sort

For each single item, to get it in the proper place, select sort needs a swap operation, which is 2 time set\_at operation.

But for insert sort, each single item needs n times set\_at operation, which is worse. Also same for merge sort.

## (b) Merge Sort

For select sort and insert sort, both need  $O(n^2)$  times to compare, take  $O(n^2 \log n)$ . But for Merge Sort, only takes  $O(n \log^2 n)$ .

#### (c) Insertion Sort

As most items in A is already sorted( $\log \log n$  is much smaller than n).

#### Problem 2-3.

• Binary Search is quite easy to think of, I'm not sure whether I should consider that the end of the island is submerging??

#### Problem 2-4.

Store a viewer as a struct, includes the viewer's ID, as well as an array store pointers that point to all messages the viewer had said. Store all viewers in a sorted order with merge sort which is  $O(n \log n)$ , with a comparison of each viewer's ID. Then store all messages as a linked list, a single node represents a viewer's message.

#### For those operations:

- build(V): Sort all viewers by their ID using merge sort cost  $O(n \log n)$ .
- send(v, m): Find viewer v with binary search cost  $O(\log n)$ , then sends a message, change the state of the struct v and the linked list both cost O(1).
- recent(k): Here we can assume that each time we insert a new message, we insert it in the front rather than from the end. Then we can return the k most recent messages in O(k) time.
- ban(v): Find viewer v cost  $O(\log n)$  and delete all the message from the linked list which is stored in struct v cost  $O(n_v)$ . So totally cost  $O(n_v + \log n)$ .

## Problem 2-5.

- (a) Firstly, try to find out how many time scope we need to take. We can do this by checking every tuple in  $B_1$  and  $B_2$ , apparently this can be done in O(n) time. Then, check every time scope to see how many rooms  $B_1$  and  $B_2$  have reserved and then B should take the amount of  $B_1 + B_2$ .
- (b) Take (a), seems like merge sort
- (c) Submit your implementation to alg.mit.edu.