# Homework 3

May 18, 2021

# 0.1 Algorithmic Design Homework 3

## 0.1.1 Exercise 1.

Implement the binary heap-based version of the Dijkstra's algorithm.

#### 0.1.2 Solution

In the cell below we implement the heap-based version of the Dijkstra's algorithm on the Graph data structure written in the file Graph.py and using the binary heap defined in Heap.py.

The heap implements all the methods we saw during the lectures and some more which we found useful during the implementation of the Dijkstra's algorithm, in particular it implements a way to decrease a value without needing the index of it in the heap, obviously to accomplish this it needs more space because it needs to store a dictionary that maps a value to a specific index, the key of this mapping is defined by the user. We used this feature in the Dijkstra's algorithm to set as a key the name of the vertex, which is unique in the graph

Designing the graph was more complicated because it can be implemented in different ways, from a high level perspective we chosed to define the graph as a set of vertexes and a vertex as a set of edges, in the implementation however we don't use a set but a dictionary because it allows finding a specific vertex or edge of the graph in  $\Theta(1)$  time on average. Defining the edge as an object is very useful if we want to store additional information about the connection, we didn't use this feature in the code below, however we used the fact that the vertex is an object to add the members v.d, v.prev and v.importance when they were needed. Obviously this implementation requires more space than simply storing a dictionary of {source, destination} pairs, however it allows for a very high generalizability

```
[1]: from numpy import inf
  from src import *
  from warnings import warn
  from random import randint
  from typing import TypeVar, Union
  from copy import deepcopy

n = TypeVar("n")
```

```
def init_sssp(G: Graph, s: n):
   """Initializes Dijkstra algorithm
   Arqs:
        G (Graph): Graph to initialize
        s (n): source
   Raises:
        RuntimeWarning: If s is not present in the graph
   u = G.get_vertex(s)
   if u is None:
       message = f"The source: {s} is not present in the graph, exiting"
       warn(message, RuntimeWarning)
       return
   for v in G.V():
       v.d = inf
       v.pred = None
   u.d = 0.
def relax(Q: Binary_heap, u: Vertex, v: Vertex):
    """Relaxes a node
   Args:
        Q (Heap): Queue implemented by a binary heap
       u (Vertex): Vertex already relaxed
        v (Vertex): Vertex to be relaxed
   w = u.get_weight(v.name)
   if v.d > u.d + w:
       v.d = u.d + w
       v.pred = u
       Q.decrease_value((v.d, v.name), (v.d, v.name))
def dijkstra(G: Graph, s: n):
    """Implements the dijkstra algorithm for finding the single source shortest
   path from the source s
   Args:
        G (Graph): Graph to complute the distance
        s (n): Name of the source in the graph
   Raises:
```

```
RuntimeWarning: If s is not present in the graph
         11 II II
         init_sssp(G, s)
         # Initialize Q with tuples (v.d, v.name), this is much lighter than storing
      \rightarrow v itself
         Q = Binary_heap([(v.d, v.name) for v in G.V()], total_order=lambda x, y:_{\sqcup}
      \rightarrow x[0] < y[0], dict_key=lambda x: x[1])
         while not Q.is_empty():
             u = Q.remove_min()
             u = G.get_vertex(u[1])
             # Relax all the childs of v
             for v in G.get childs(u.name):
                 relax(Q, u, v)
     # Create a simple graph for testing
     G = Graph(directed=True)
     G.add_edge("A", "B", 1000)
     G.add_edge("A", "C", 1)
     G.add_edge("C", "D", 5)
     G.add_edge("D", "B", 8)
     G.add_vertex("E")
     dijkstra(G, "A")
     [(v.name, v.d) for v in G.V()]
[1]: [('B', 14.0), ('A', 0.0), ('C', 1.0), ('D', 6.0), ('E', inf)]
[2]: # Creates the graph from the slides for testing
     G_slides = Graph()
     edges = [(1,6,1), (1,5,1), (2,3,2), (3,2,1), (3,4,3), (4,7,1),
              (4,8,3), (5,6,1), (5,1,3), (6,8,1), (7,8,1), (8,1,1), (8,7,1)
     [G_slides.add_edge(s, d, w) for s, d, w in edges]
     [setattr(v, "importance", v.name) for v in G_slides.V()]
     dijkstra(G slides, 1)
     [(v.name, v.d) for v in G_slides.V()]
[2]: [(6, 1.0),
      (1, 0.0),
      (5, 1.0),
      (3, inf),
      (2, inf),
      (4, inf),
      (7, 3.0),
```

(8, 2.0)

#### 0.1.3 Exercise 2.

#### 0.1.4 Solution

The implementation of the function to add the shortcuts to a graph can be found in the file Graph.py at line 90, below we will implement the bidirectional version of the Dijkstra algorithm

```
[3]: def dijkstra_iteration(G: Graph, Q: Binary_heap) -> n:
         """Performs a single iteration of the bidirectional dijkstra algorithm
         Arqs:
             G (Graph): Graph to iterate
             Q (Binary_heap): Corresponding queue
         Returns:
             n: name of the vertex that has been completed
         u = Q.remove_min()
         u = G.get_vertex(u[1])
         for v in G.Adj(u.name):
             v = G.get_vertex(v)
             # Only relax if the importance is greater
             if v.importance < u.importance:</pre>
                 continue
             Q.insert((v.d, v.name))
             relax(Q, u, v)
         return u.name
     def get_path(G: Graph, v: n, d: n) -> "tuple(list[n], float)":
         """Returns the path to be travelled to reach d from v and its associated_{\sqcup}
      \hookrightarrow weight
         Args:
             v (n): Source of the path
             d (n): Destination of the path
         Returns:
             tuple(list[n], float): the first element is the path to travel to reach
             d from v, the second is it's associated weight
         v = deepcopy(v)
         path = [v]
```

```
w = 0.
   while v != d:
        v = G.get_vertex(v) # Get the vertex corresponding to v
        edge = v.pred[v.name] # Get the edge from v.pred to v
       path.extend(edge.decompose()[::-1]) # Add to the path the vertexes the
→edge passes by in reverse order
       w += edge.weight # Update the overall weight
       v = v.pred.name # Update v
   return path, w
def remove_duplicates(a: list) -> list:
    """Simple function that removes subsequent duplicate values from a list
   Args:
        a (list): list to remove the duplicates
   Returns:
        list: list with the duplicates removed
   return [a[i] for i in range(0, len(a)) if a[i-1]!=a[i]]
def Bidirectional_dijkstra(G: Graph, s: n, d: n) -> "tuple(list[n], float)":
    """Implements the bidirectional Dijkstra algorithm to find the
    shortest path from s to d of G, note that if the node importance
    is not set properly this algorithm can find sub-optimal paths
   Args:
        G (Graph): Graph to compute the shortest path
        s (n): Source or starting point of the path
        d (n): Destination or ending of the path
   Returns:
        tuple(list[n], float): the first element is the path to travel to reach
        d from s, the second is it's associated weight
   Raises:
        RuntimeWarning:
            If a path between s and d does not exist, in this case returns None
    # Create G_up and G_down
   G_down = G.T()
   G_{up} = G
```

```
# Add the shortcuts
   G_up.add_shortcuts()
   G_down.add_shortcuts()
   # Initialize the search
   init_sssp(G_up, s)
   init_sssp(G_down, d)
   # Create the binary heaps
   u = G up.get vertex(s)
   Q_{up} = Binary_{heap}([(u.d, u.name)], total_order=lambda x, y: x[0] < y[0],_u
→dict key=lambda x: x[1])
   u = G_up.get_vertex(d)
   Q_{\text{down}} = Binary_{\text{heap}}([(u.d, u.name)], total_{\text{order}}=lambda x, y: x[0] < y[0],_u
→dict_key=lambda x: x[1])
   # Create two sets to keep track of the vertexes encountered so far
   vertices up = set()
   vertices_down = set()
   common vertex = ""
   # I iterate on G up and G down iteratively until I find a common vertex
   # or until both queues are empty, if this is the case then it means that
   # I haven't found a common vertex and so there is no path from s to d
   while not Q_up.is_empty() or not Q_down.is_empty():
       if not Q_up.is_empty():
           up = dijkstra_iteration(G_up, Q_up)
           if up in vertices_down:
               common vertex = up
               break
       if not Q_down.is_empty():
           down = dijkstra_iteration(G_down, Q_down)
           if down in vertices_up:
               common_vertex = down
               break
       vertices_up.add(up)
       vertices_down.add(down)
   if common_vertex == "":
       warn(f"A path between {s} and {d} does not exist", RuntimeWarning)
       return
   # Get the path to be travelled to reach s from d and its weight
   p1, w1 = get_path(G_up, common_vertex, s)
   p2, w2 = get_path(G_down, common_vertex, d)
   p1.reverse()
   return remove_duplicates([*p1, *p2]), w1+w2
```

```
[4]: # Test on the simple graph defined above
for v in G.V():
    v.importance=randint(0, 10)

Bidirectional_dijkstra(G, "A", "B")

[4]: (['A', 'C', 'D', 'B'], 14.0)

[5]: # Test on the graph on the slides
    Bidirectional_dijkstra(G_slides, 1, 8), Bidirectional_dijkstra(G_slides, 3, 6)
[5]: (([1, 6, 8], 2.0), ([3, 4, 7, 8, 1, 6], 7.0))
```

## 0.2 Conclusion

The code of the bidirectional Dijkstra algorithm can be improved by creating two ad-hoc procedures to build G\_up and G\_down, in particular in these graphs the connections that go from a vertex with higher importance to one with lower importance are never used, so we could create G\_up and G\_down with these connections removed by default, this saves both time (we don't have to check if v.importance < u.importance) and space (because we store less connections). We sticked to this implementation because the homework says that the bidirectional Dijkstra has to operate on the graphs decorated by the algorithm at Point 2a, for completeness, below I implement the two procedures to create G\_up and G\_down like we described above

```
[]: def remove useless edges(G: Graph) -> Graph:
         """Removes connections that go from a higher importance to a lower_{\sqcup}
      \hookrightarrow importance,
         be aware that the removal is done in place!
         Args:
              G (Graph): Graph to remove the connections
         Returns:
              Graph: The graph with the connections removed
         Raises:
              RuntimeWarning: If a vertex does not have an importance
         try:
              for edge in G.E():
                  s = edge.get_source()
                  d = edge.get_destination()
                  if G.get_vertex(s).importance > G.get_vertex(d).importance:
                      G.remove_edge(s, d)
         except AttributeError:
```

```
warn(f"The graph contains vertexes without an importance attribute, ⊔
 ⇔exiting", RuntimeWarning)
        return
    return G
def get_G_up(G: Graph) -> Graph:
    """Computes G_up, useful for the bidirectional Dijkstra algorithm
    Args:
        G (Graph): Graph to compute G_up
    Returns:
        Graph: The G_up graph
    G_{up} = deepcopy(G)
    G_up.add_shortcuts()
    return remove_useless_edges(G_up)
def get_G_down(G: Graph) -> Graph:
    """Computes G_down, useful for the bidirectional Dijkstra algorithm
    Arqs:
        G (Graph): Graph to compute G\_down
    Returns:
       Graph: The G_down graph
    G_down = G.T()
    G_down.add_shortcuts()
    return remove_useless_edges(G_down)
```

Now we can use get\_G\_up and get\_G\_down in the Bidirectional\_dijkstra function, we can also update dijkstra\_iteration to avoid the checking of the importances