

Deterministic Optimization

Linear Optimization Modeling
Network Flow Problems

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Optimal Transportation Problem

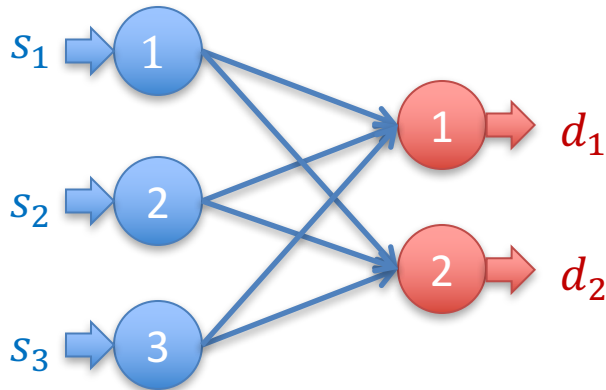
Modeling using Linear Programs

Learning Objectives

- Discover an interesting example of LP called the **transportation problem**, and discuss when two LP models are equivalent.
- Use CVX and read solver output.

Transportation Problem

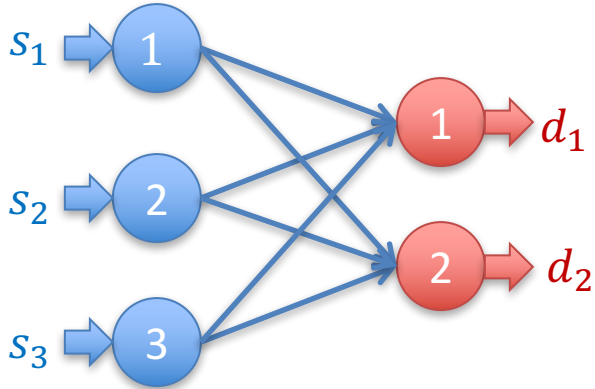
There are m suppliers, n customers. Supplier i can supply up to s_i units of supply, and customer j has d_j units of demand. It costs c_{ij} to transport a unit of product from supplier i to customer j . We want to find a transportation schedule to satisfy all the demand within minimum transportation cost.



Transportation Problem

Decision variables: x_{ij} for $i = 1, \dots, m, j = 1, \dots, n$

x_{ij} : the amount of product transported from supply i to warehouse j .

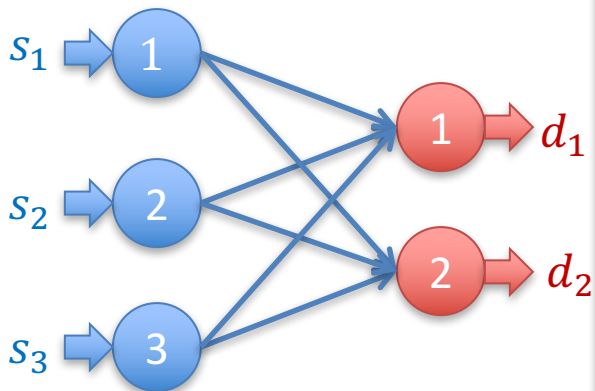


$$\begin{aligned} \min \quad & \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} && \text{(M1)} \\ \text{s.t.} \quad & \sum_{i=1}^m x_{ij} = d_j, \quad j = 1, \dots, n && \text{Demand is exactly satisfied.} \\ & \sum_{j=1}^n x_{ij} \leq s_i, \quad i = 1, \dots, m \\ & x_{ij} \geq 0, \quad i = 1, \dots, m, \quad j = 1, \dots, n. \end{aligned}$$

Another Formulation

Decision variables: x_{ij} for $i = 1, \dots, m, j = 1, \dots, n$

x_{ij} : the amount of product transported from supply i to warehouse j .



$$\min \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} \quad (\text{M2})$$

$$\text{s.t.} \quad \sum_{i=1}^m x_{ij} \geq d_j, \quad j = 1, \dots, n$$

$$\sum_{j=1}^n x_{ij} \leq s_i, \quad i = 1, \dots, m$$

$$x_{ij} \geq 0, \quad i = 1, \dots, m, \quad j = 1, \dots, n.$$

Demand is satisfied or exceeded.

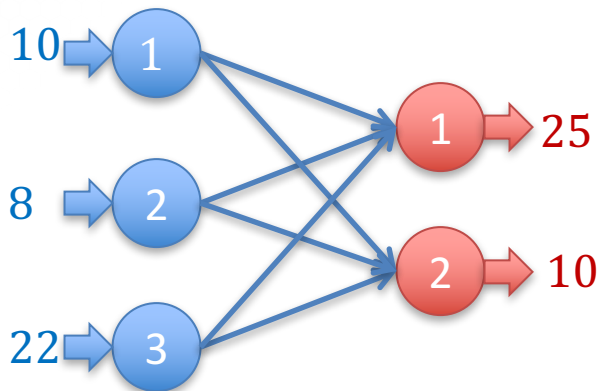
Are Two Formulations Equivalent?

$$\begin{aligned} \min \quad & \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} \\ \text{s.t.} \quad & \sum_{i=1}^m x_{ij} \stackrel{\text{red}}{=} d_j, \quad \forall j \\ & \sum_{j=1}^n x_{ij} \stackrel{\text{blue}}{\leq} s_i, \quad \forall i \\ & x_{ij} \geq 0, \quad \forall i, j. \end{aligned}$$

$$\begin{aligned} \min \quad & \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} \\ \text{s.t.} \quad & \sum_{i=1}^m x_{ij} \stackrel{\text{red}}{\geq} d_j, \quad \forall j \\ & \sum_{j=1}^n x_{ij} \stackrel{\text{blue}}{\leq} s_i, \quad \forall i \\ & x_{ij} \geq 0, \quad \forall i, j. \end{aligned}$$

If $c_{ij} \geq 0$ for all i, j , then we can claim that the \geq inequality in the second model will be satisfied as $=$ at optimal solution, thus, the two formulations are equivalent.

A Concrete Example



CVX code for (M2):

```
m = 3;
n = 2;
c = [8, 2;
     4, 10;
     6, 7];
s = [10; 8; 22];
d = [25; 10];
cvx_begin
variable x(3,2) nonnegative;
minimize sum(sum(c.*x));
subject to
sum(x,1) >= d';
sum(x,2) <= s;
cvx_end
```

Results

How to read the output log of Gurobi?

- Preprocessing
- Algorithm iterations
- Status
- Optimal solution

Calling Gurobi 6.00: 11 variables, 5 equality constraints

Gurobi optimizer, licensed to CVX for CVX

Optimize a model with 5 rows, 11 columns and 17 nonzeros

Coefficient statistics:

Matrix range [1e+00, 1e+00]

Objective range [2e+00, 1e+01]

Bounds range [0e+00, 0e+00]

RHS range [8e+00, 3e+01]

Presolve removed 0 rows and 5 columns

Presolve time: 0.00s

Presolved: 5 rows, 6 columns, 12 nonzeros

Iteration	Objective	Primal Inf.	Dual Inf.	Time
0	0.0000000e+00	3.500000e+01	0.000000e+00	0s
3	1.5400000e+02	0.000000e+00	0.000000e+00	0s

Solved in 3 iterations and 0.01 seconds

Optimal objective 1.540000000e+02

Status: Solved

Optimal value (cvx_optval): +154

>> x

x =

0 10
8 0
17 0

Summary

- We have learned the transportation problem modeled as linear programs.
- Optimization problems need analysis: Different problems may be equivalent.
- CVX + Gurobi is a powerful tool for solving LPs.