

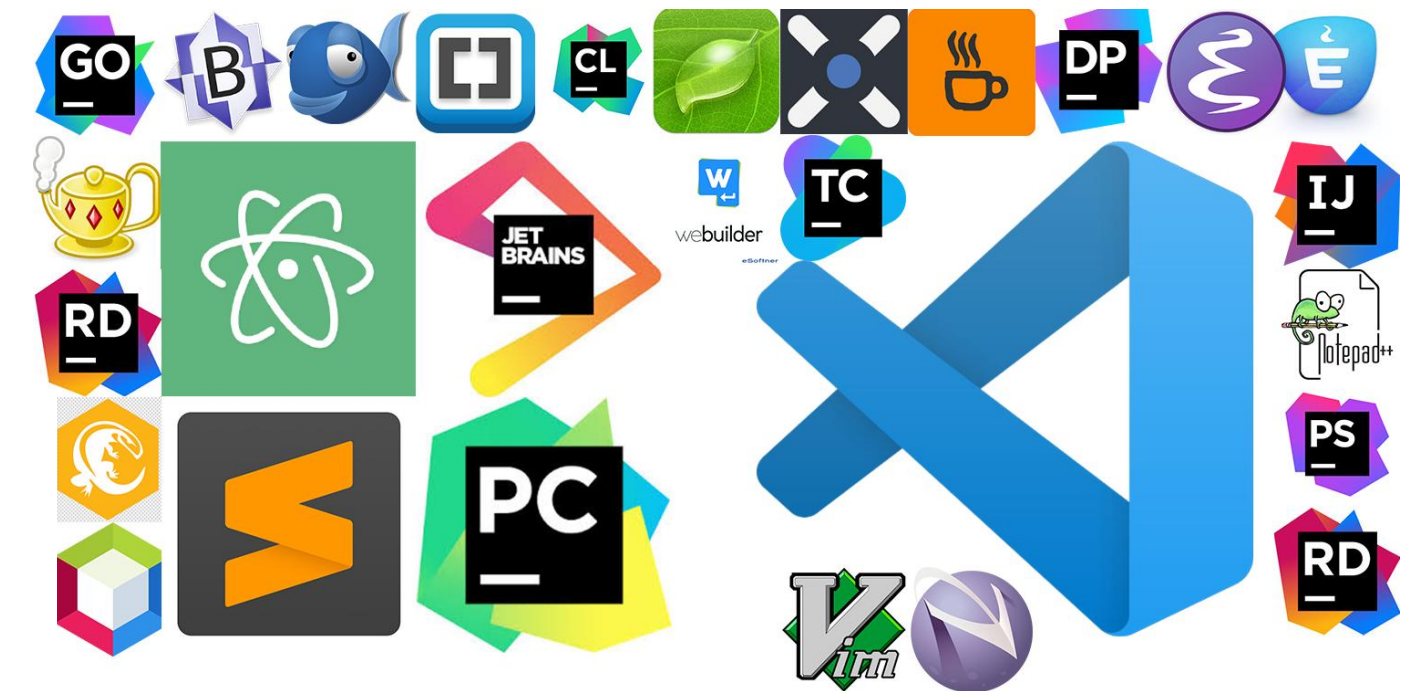
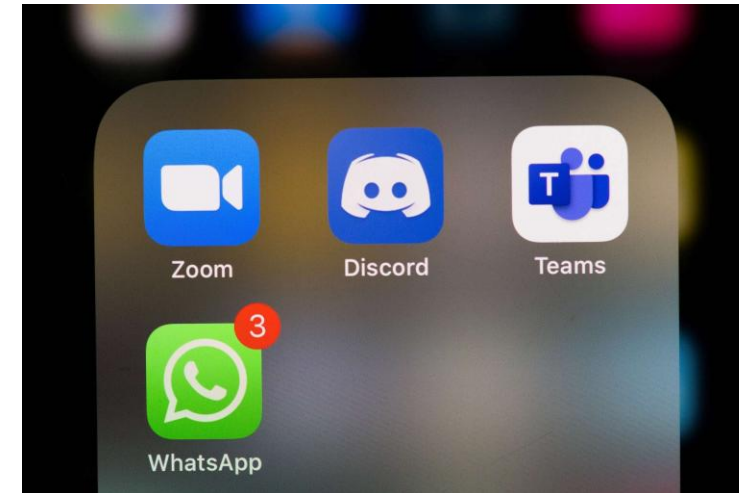


Operating System Concepts: Lecture 3

programs/processes and threads : from bios to application startup

Processes

Processes: Building blocks of what we do on a computer



Job vs Process

- Job: slightly older term (used in batch systems) but still used in literature (“job scheduling”)
- Process: Actually the same thing
- View processes :
 - Windows
 - Linux : ps command

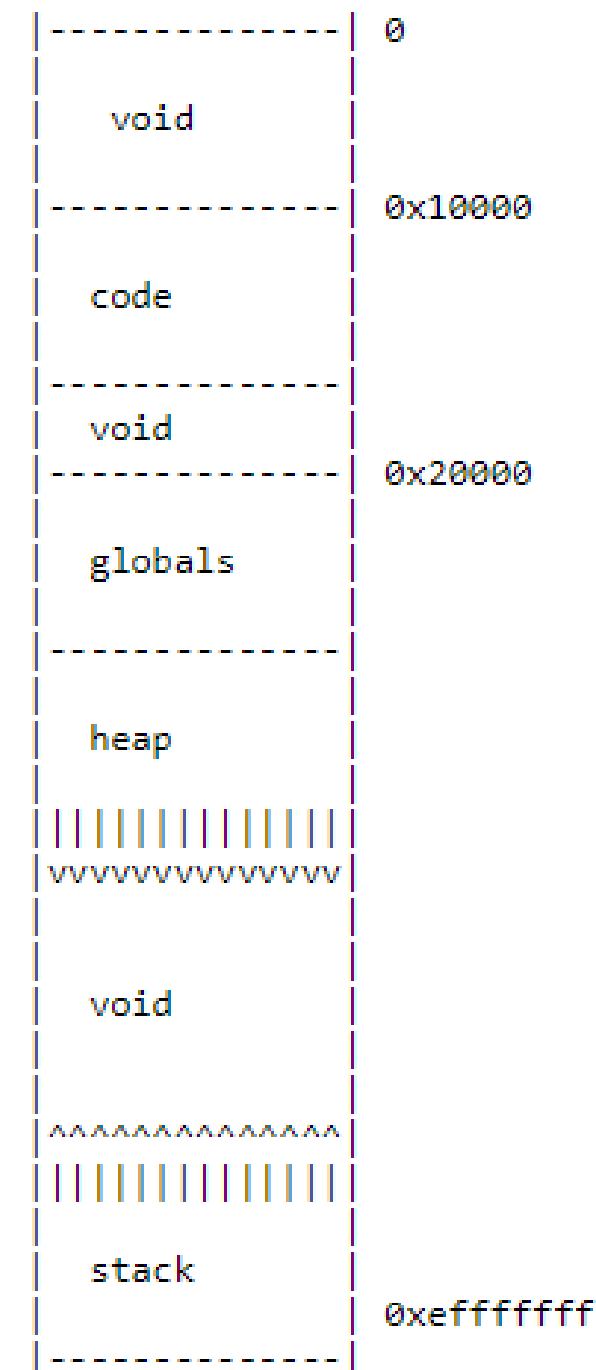
```
sagar@itsFOSS: ~  
sagar@itsFOSS:~$ ps -aux  
USER      PID %CPU %MEM    VSZ   RSS TTY      STAT START   TIME COMMAND  
root         1   0.0   0.0 166872 11456 ?        Ss   15:31   0:02 /sbin/init sp  
root         2   0.0   0.0      0     0 ?        S    15:31   0:00 [kthreadd]  
root         3   0.0   0.0      0     0 ?        S    15:31   0:00 [pool_workque  
root         4   0.0   0.0      0     0 ?        I<   15:31   0:00 [kworker/R-rc  
root         5   0.0   0.0      0     0 ?        I<   15:31   0:00 [kworker/R-rc  
root         6   0.0   0.0      0     0 ?        I<   15:31   0:00 [kworker/R-sl  
root         7   0.0   0.0      0     0 ?        I<   15:31   0:00 [kworker/R-ne  
root         9   0.0   0.0      0     0 ?        I<   15:31   0:00 [kworker/0:0H  
root        12   0.0   0.0      0     0 ?        I<   15:31   0:00 [kworker/R-mm  
root        13   0.0   0.0      0     0 ?        I    15:31   0:00 [rcu_tasks_kt
```

Process	CPU	Private Bytes	Working Set	PID	Description	Company Name	Threads	VirusTotal
Registry		3.952 K	40.036 K	120			4	
System Idle Process	97.19	56 K	8 K	0			8	
System	0.18	196 K	1.248 K	4			232	
csrss.exe	< 0.01	1.984 K	3.508 K	636			13	
wininit.exe		1.736 K	4.552 K	720			2	
GoogleCrashHandler.exe		1.688 K	1.208 K	5020			4	
GoogleCrashHandler64.exe		1.696 K	888 K	5360			4	
csrss.exe	0.08	2.964 K	5.400 K	12644			15	
winlogon.exe		2.600 K	8.264 K	16044			5	
igfxEM.exe		4.220 K	12.688 K	240	igfxEM Module	Intel Corporation	8	
explorer.exe	0.03	149.908 K	166.220 K	6536	Windows Verkenner	Microsoft Corporation	98	
SecurityHealthSystray.exe		1.880 K	7.620 K	16172	Windows Security notification ...	Microsoft Corporation	3	
RtkNGUI64.exe		4.716 K	11.408 K	9952	Realtek HD Audio configuratie	Realtek Semiconductor	10	
RAVBg64.exe	< 0.01	4.756 K	9.880 K	2172	HD Audio Background Proce...	Realtek Semiconductor	7	
WavesSvc64.exe	< 0.01	22.420 K	13.956 K	4668	Waves MaxxAudio Service A...	Waves Audio Ltd.	12	
OneDrive.exe	0.15	338.120 K	307.448 K	4532	Microsoft OneDrive	Microsoft Corporation	37	
ZoomIt.exe		1.544 K	6.356 K	11888	Sysinternals Screen Magnifier	Sysinternals - www.sysinter...	3	
openvpn-gui.exe		2.808 K	7.624 K	5472			3	
CCXProcess.exe		776 K	2.588 K	13636	CCXProcess	Adobe Systems Incorporat...	2	
ONENOTEM.EXE		2.908 K	1.912 K	14308	Send to OneNote Tool	Microsoft Corporation	3	
notepad.exe		3.052 K	12.540 K	14744	Kladblok	Microsoft Corporation	3	
chrome.exe	0.04	115.444 K	158.872 K	15680	Google Chrome	Google LLC	32	
OUTLOOK.EXE	< 0.01	188.040 K	244.900 K	10748	Microsoft Outlook	Microsoft Corporation	64	
vmware.exe	0.01	73.332 K	73.376 K	10276	VMware Workstation	VMware, Inc.	16	
vmware-unity-helper.exe		7.052 K	15.440 K	428	VMware Unity Helper	VMware, Inc.	3	
POWERPNT.EXE	< 0.01	271.200 K	315.496 K	11072	Microsoft PowerPoint	Microsoft Corporation	40	
AcroRd32.exe		11.232 K	17.380 K	9504	Adobe Acrobat Reader DC	Adobe Systems Incorporat...	11	
AcroRd32.exe	< 0.01	84.636 K	80.424 K	10160	Adobe Acrobat Reader DC	Adobe Systems Incorporat...	20	
RdrCEF.exe		17.532 K	29.300 K	6520	Adobe RdrCEF	Adobe Systems Incorporat...	22	
RdrCEF.exe		50.252 K	32.764 K	6544	Adobe RdrCEF	Adobe Systems Incorporat...	13	
procexp64.exe	0.90	88.144 K	104.408 K	16736	Sysinternals Process Explorer	Sysinternals - www.sysinter...	7	
ApntEx.exe		1.996 K	6.764 K	15896	Alps Pointing-device Driver f...	Alps Electric Co., Ltd.	4	
vpnui.exe	0.04	6.572 K	19.692 K	12956	Cisco AnyConnect User Interf...	Cisco Systems, Inc.	7	

Process vs Program

- Program/Binary/Executable (/application/app):
 - passive
 - has no state
 - program code + initialized data
 - (= is not running!)
- Process:
 - active
 - has a state
 - program code + program counter + stack + data section + heap + ...

If we view memory as a big array, the regions (or ``segments'') look as follows:



Process states

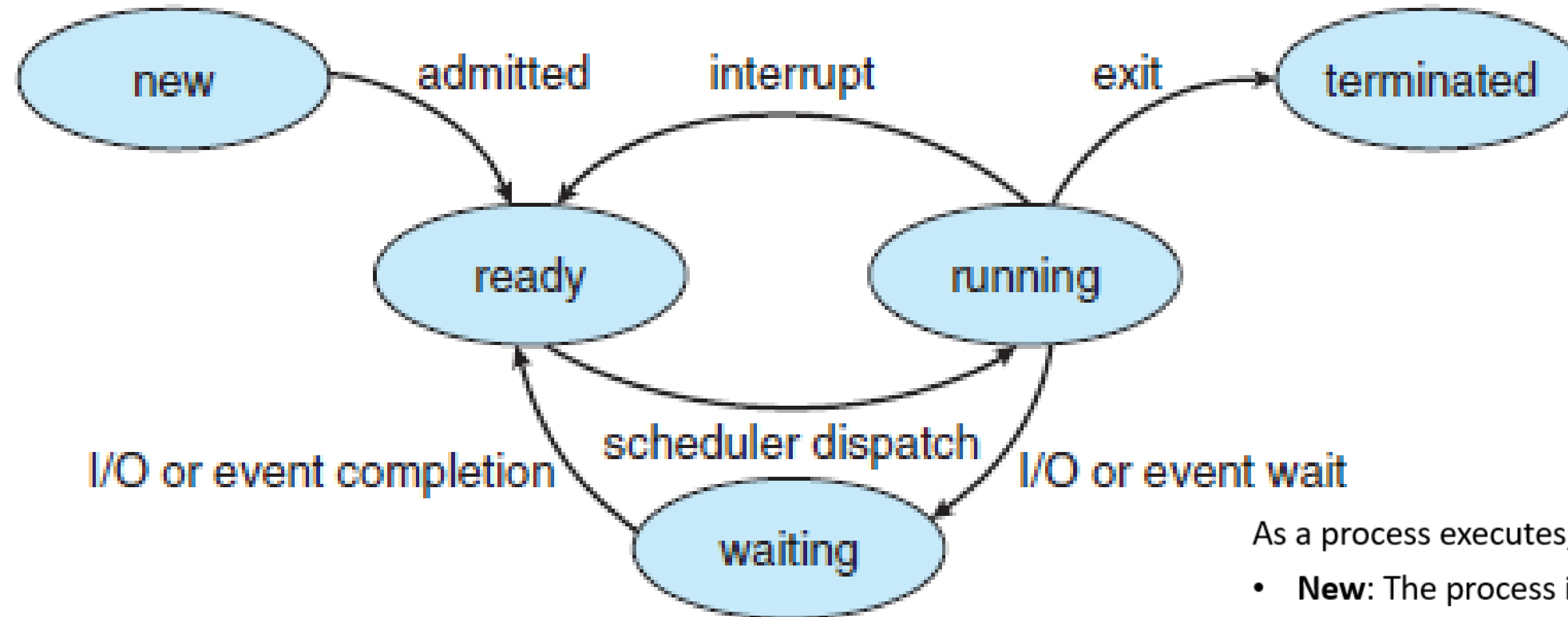


Figure 3.2 Diagram of process state.

As a process executes, the **state** changes:

- **New**: The process is being created
- **Running**: Instructions are being executed
- **Waiting**: The process is waiting for some event to occur
- **Ready**: The process is waiting to be assigned to a processor
- **Terminated**: The process has finished execution

Context Switch

= Saving/storing the state of a process (or thread)

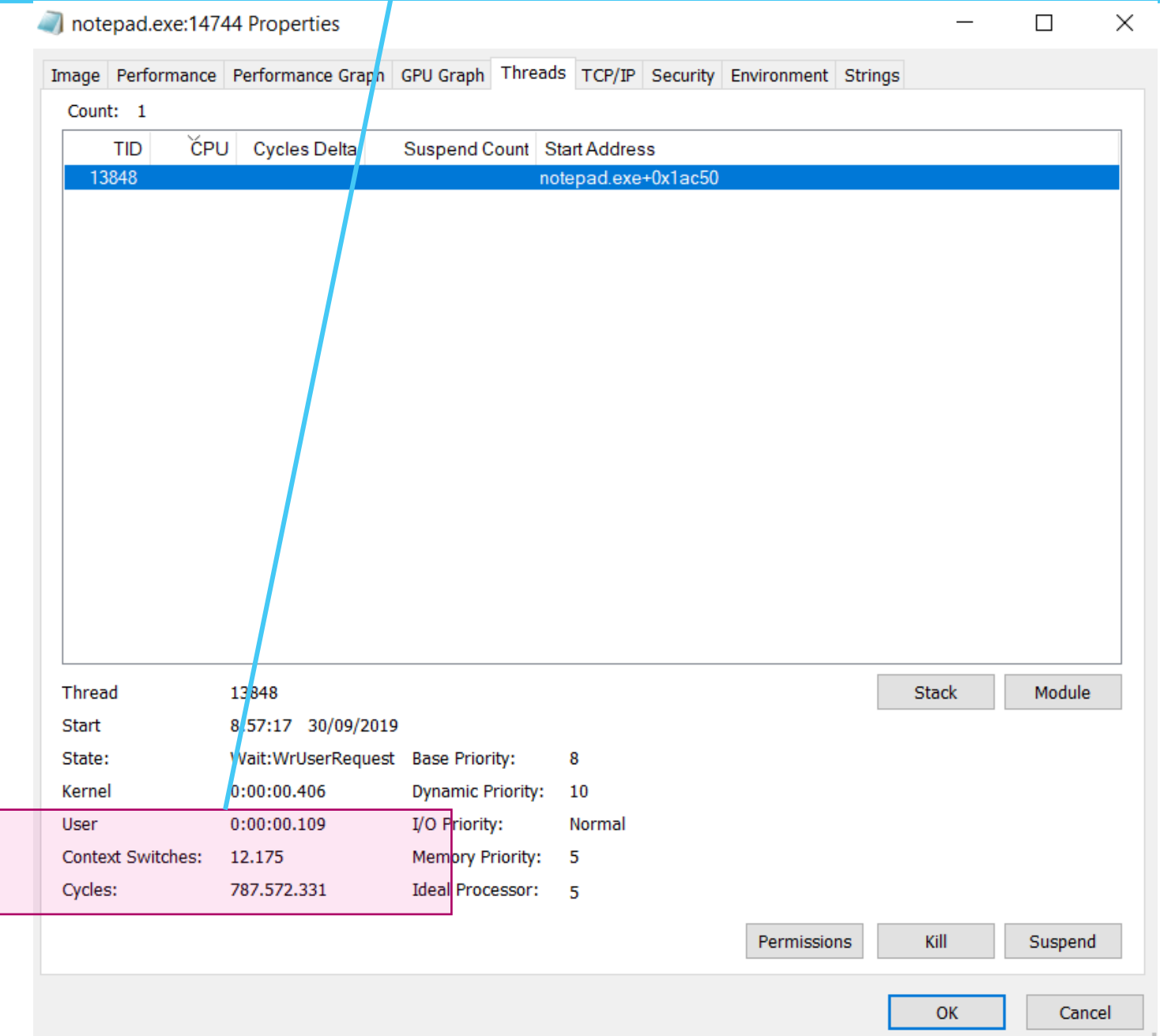
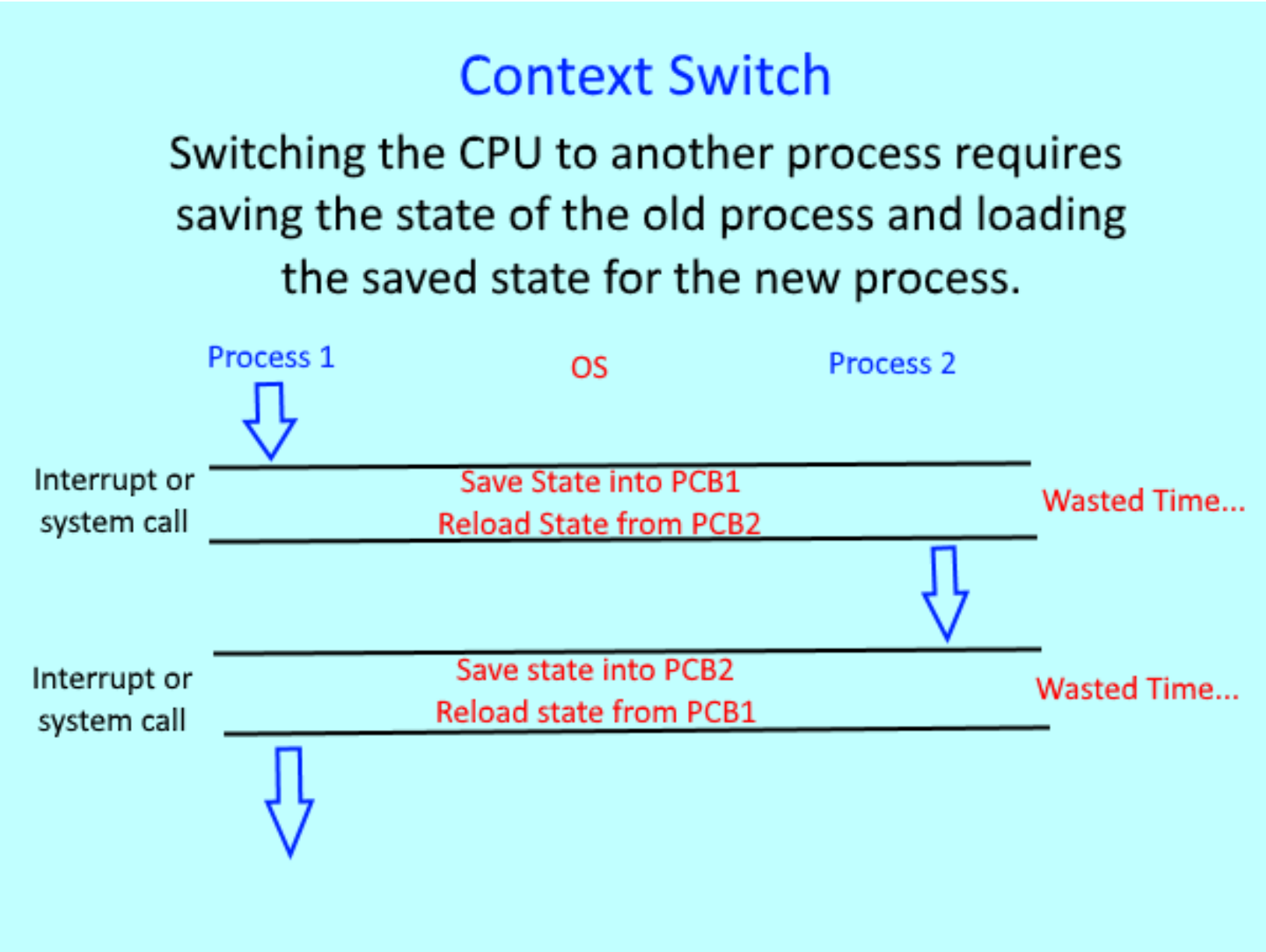
As a result a **single CPU** can **allow multiple processes** -> multitasking operating system!

What actions the context switch performs depends on multiple things:

- Type of OS
- One or more register set
- Sometimes it is needed to go from user mode to kernel mode
- Sometimes as a result of interrupts
- ...

Context Switch

Context Switches: 12 175



<https://stackoverflow.com/questions/7439608/steps-in-context-switching>

(OS) Stack vs Heap

Stack vs. Heap

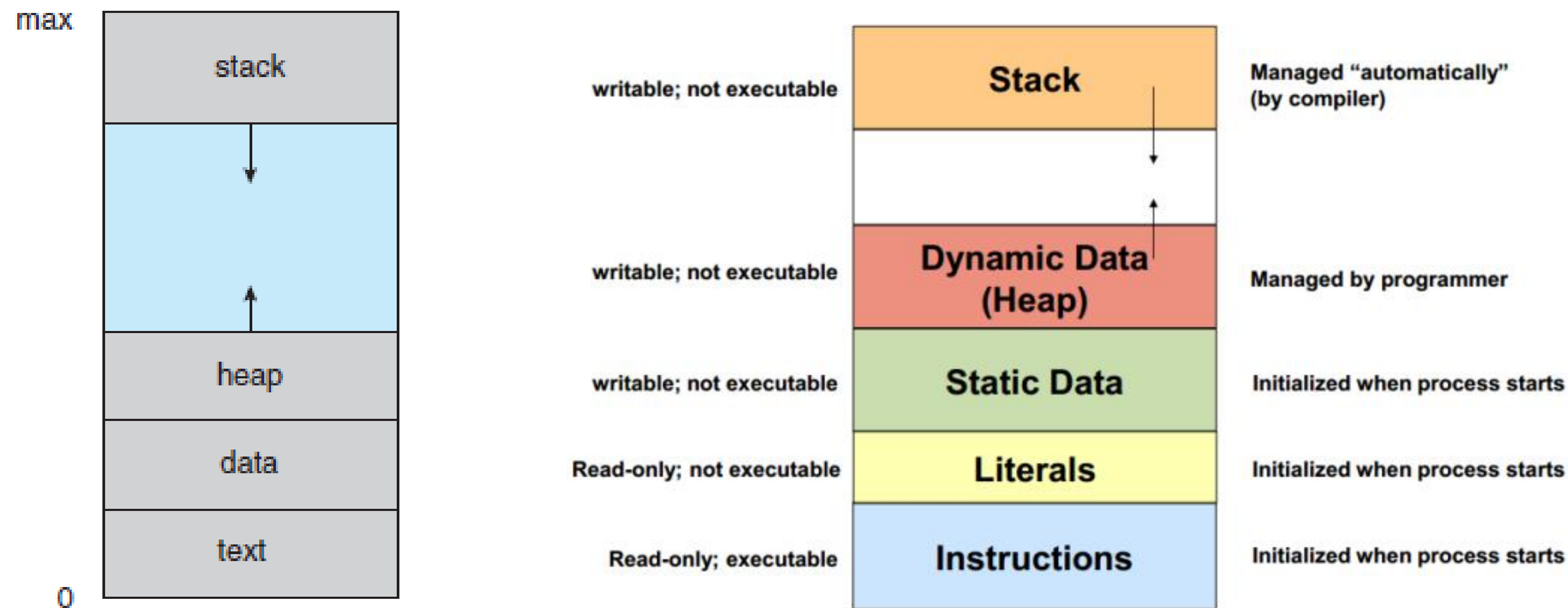
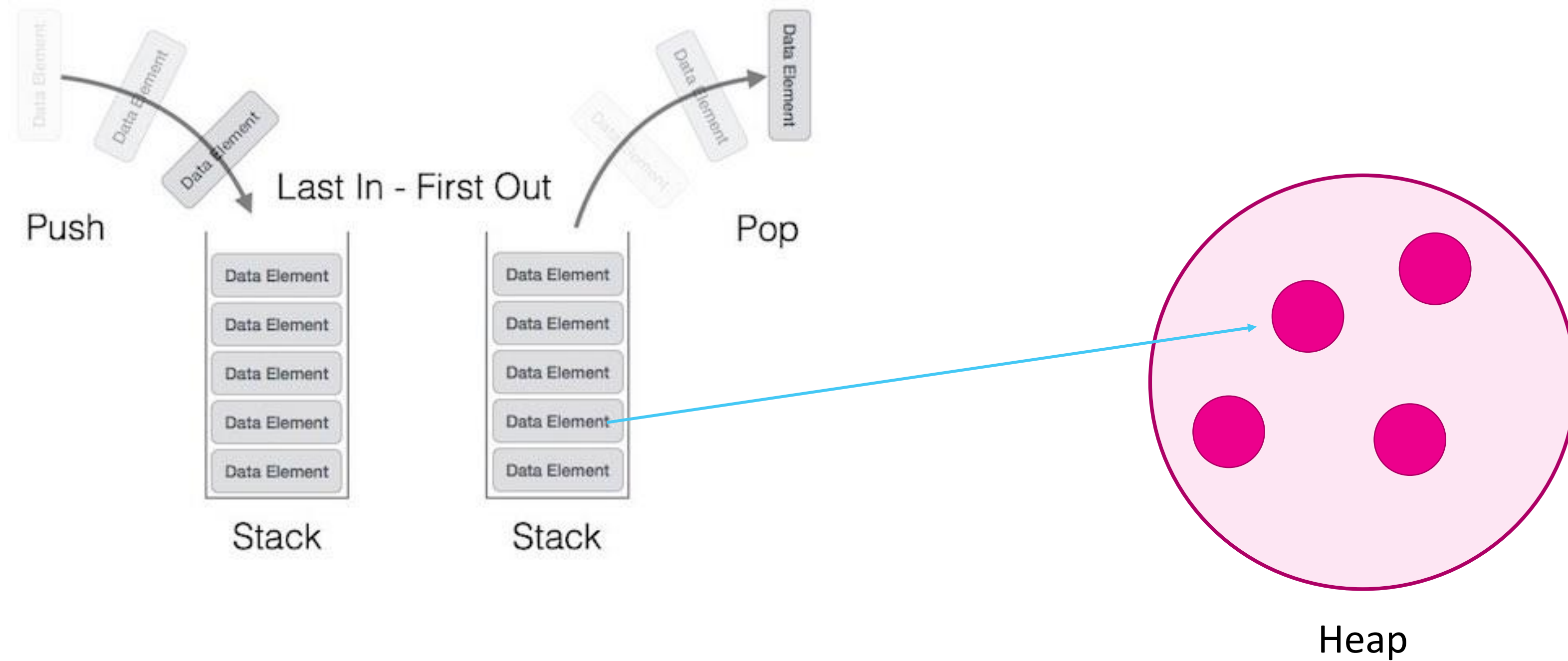
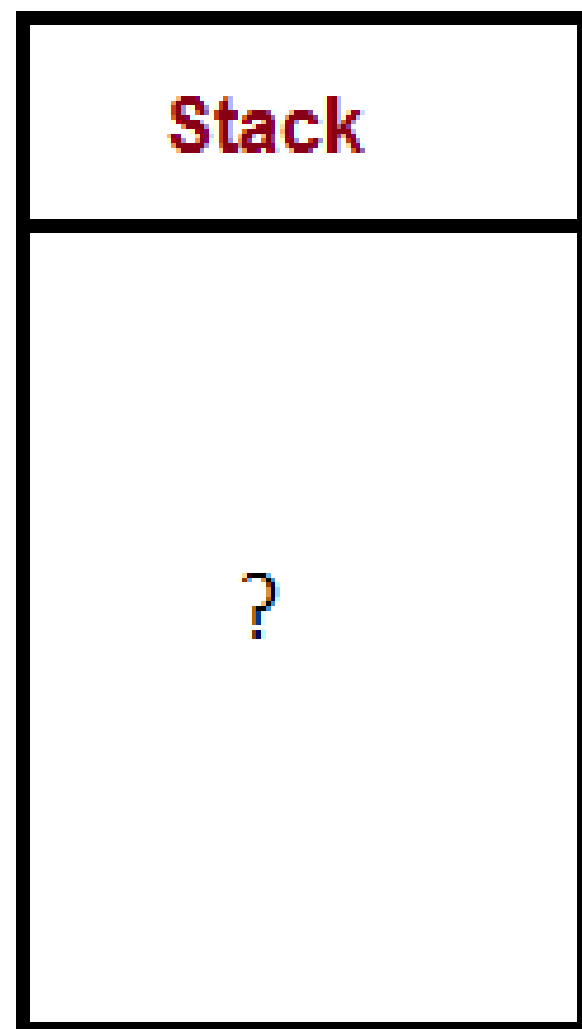


Figure 3.1 Process in memory.

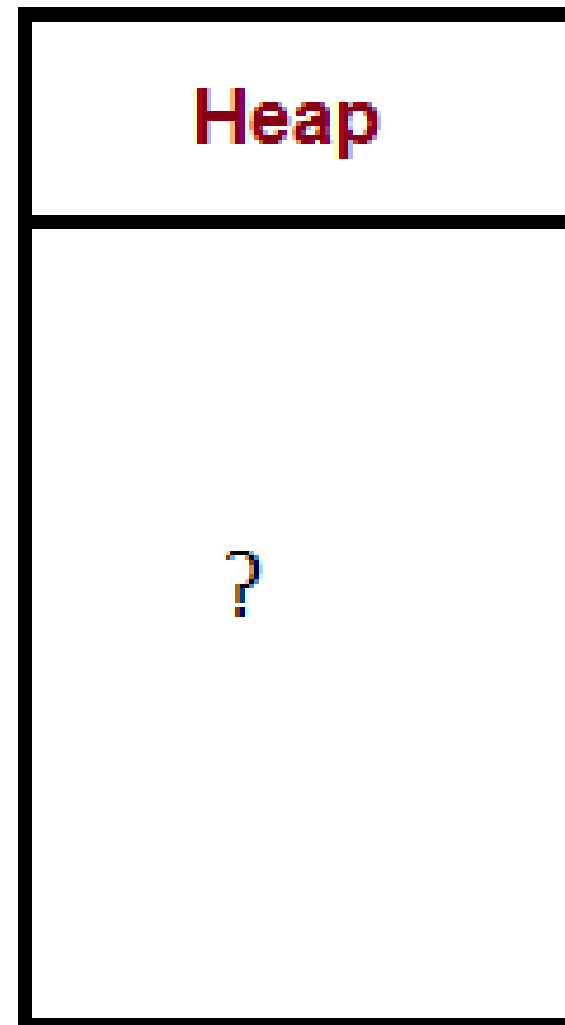
Stack vs. Heap



Stack vs. Heap



```
int x=1  
int y = 2  
  
Form1 frm = new Form1()
```



Stack or heap?

- X?
- Y?
- Frm-object?
- Pointer/reference to Frm-object?

Memzoom

<https://justine.lol/memzoom/index.html>

Linear Memory Order

Renders memory bytes using traditional left-to-right ordering.

```
??ZuP 6x!+94\δ      Pδ      :Qy.i1δ      :\\4/-i:ptj@BRég5X(pU_N  K45DHj:_!!J  K4  K4)qo      e,j  GJ, k %f "Éd 10008002b400
10008002bc00
10008002c400
10008002cc00
10008002d400
10008002dc00
10008002e400
10008002ec00
10008002f400
10008002fc00
100080030400
100080030c00
100080031400
100080031c00
100080032400
100080032c00
100080033400
100080033c00
100080034400
100080034c00
100080035400
100080035c00
100080036400
100080036c00
100080037400
100080037c00
100080038400
100080038c00
100080039400
100080039c00
10008003a400
10008003ac00
memzoom 16x
```

“Low” level programming

Introduction to C



- Pretty old language (1972-1973)
- Was created to make UNIX !
- Interesting in this module because it is low-level
- Needs a Compiler (gcc for linux)
- Benefits from a make/builder (gmake)

Structure of C Program

<i>Header</i>	<code>#include <stdio.h></code>
<i>main()</i>	<code>int main()</code> <code>{</code>
<i>Variable declaration</i>	<code>int a = 10;</code>
<i>Body</i>	<code>printf("%d ", a);</code>
<i>Return</i>	<code>return 0;</code> <code>}</code>

Introduction to C

To **Build** a c program from source you can either use:

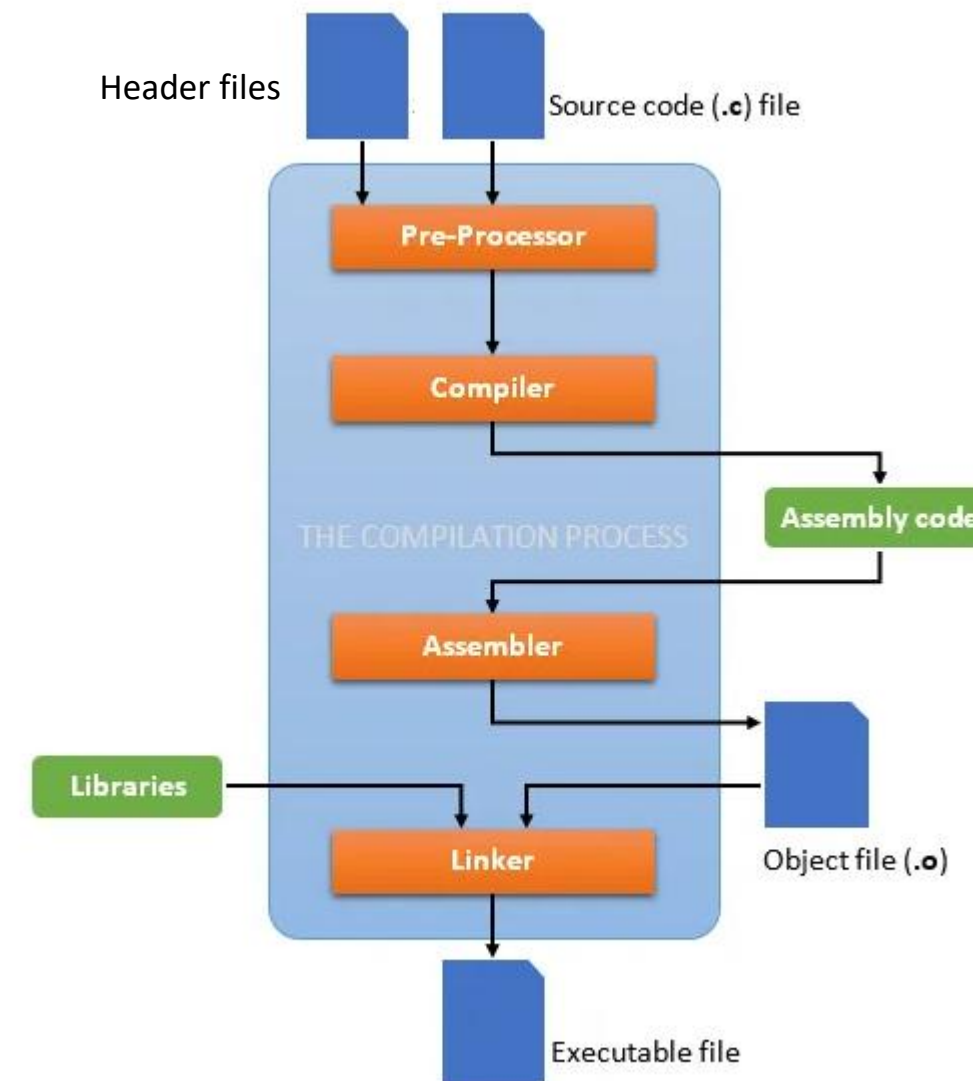
`gcc hello.c -o hello`

Or

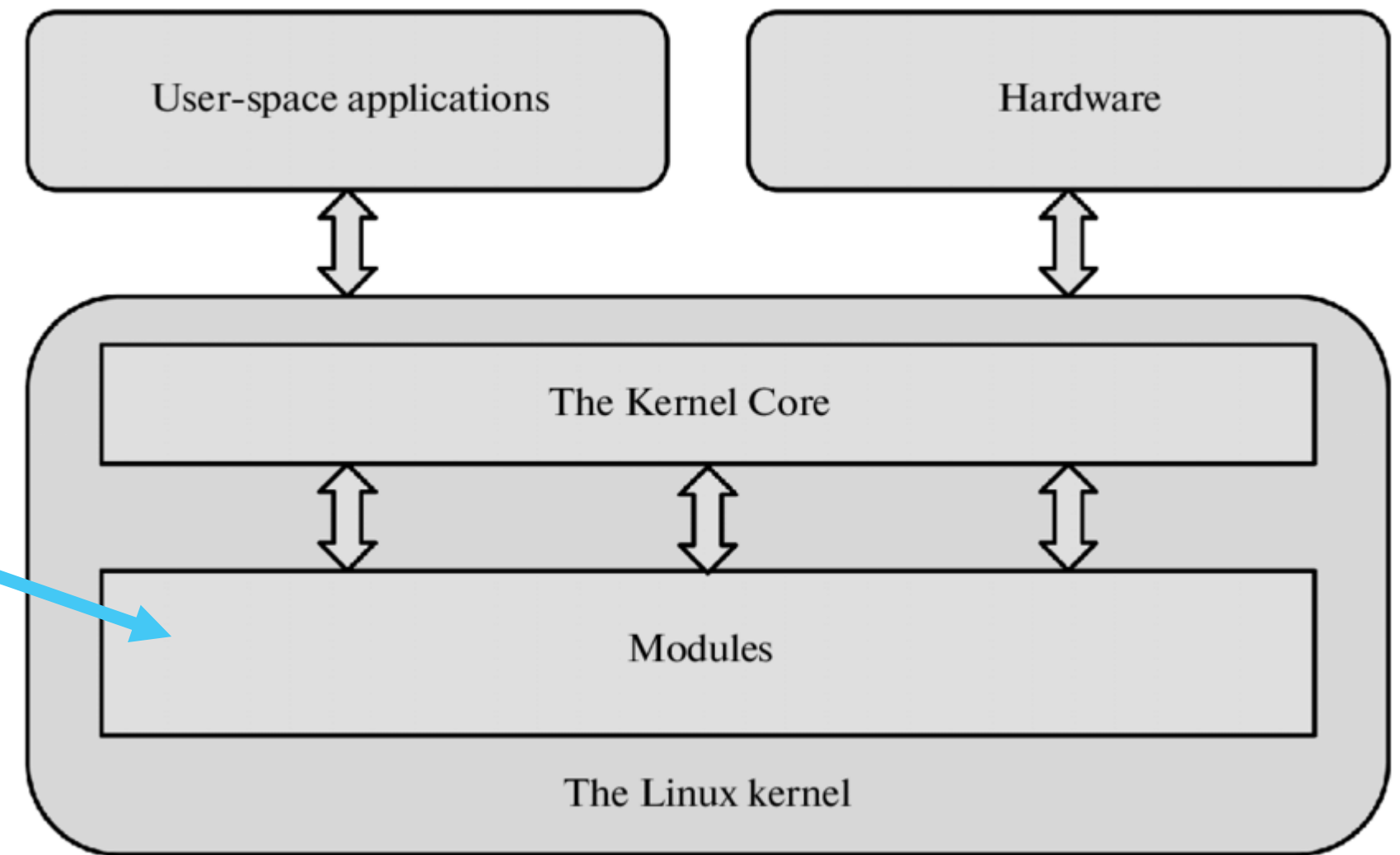
`make hello`

This will generate an executable.

Makefiles exist to make the life of a C programmer easier.



Kernel module



C can do low level things if run in kernel mode ! (i.e. wreck your system)

Pass by value vs. Pass by reference intermezzo

Nice visual representation: <https://www.mathwarehouse.com/programming/passing-by-value-vs-by-reference-visual-explanation.php>

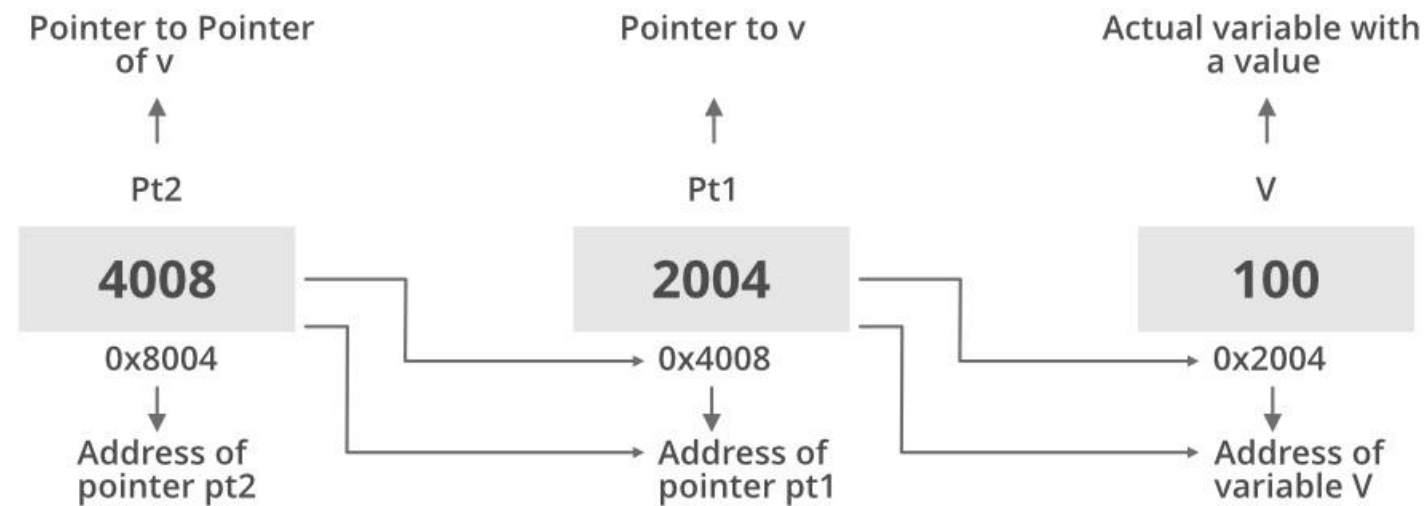
Authentic definition:

- When a parameter is **passed by reference**, the caller and the called **use the same variable** for the parameter. If the called modifies the parameter variable, the effect is visible to the caller's variable.
- When a parameter is **passed by value**, the caller and called have **two independent variables** with the same value. If the called modifies the parameter variable, the effect is not visible to the caller.

Pointers in low level programming languages

<https://www.geeksforgeeks.org/go-pointer-to-pointer-double-pointer/>

Pointer to Pointer



How Pointers works in Go



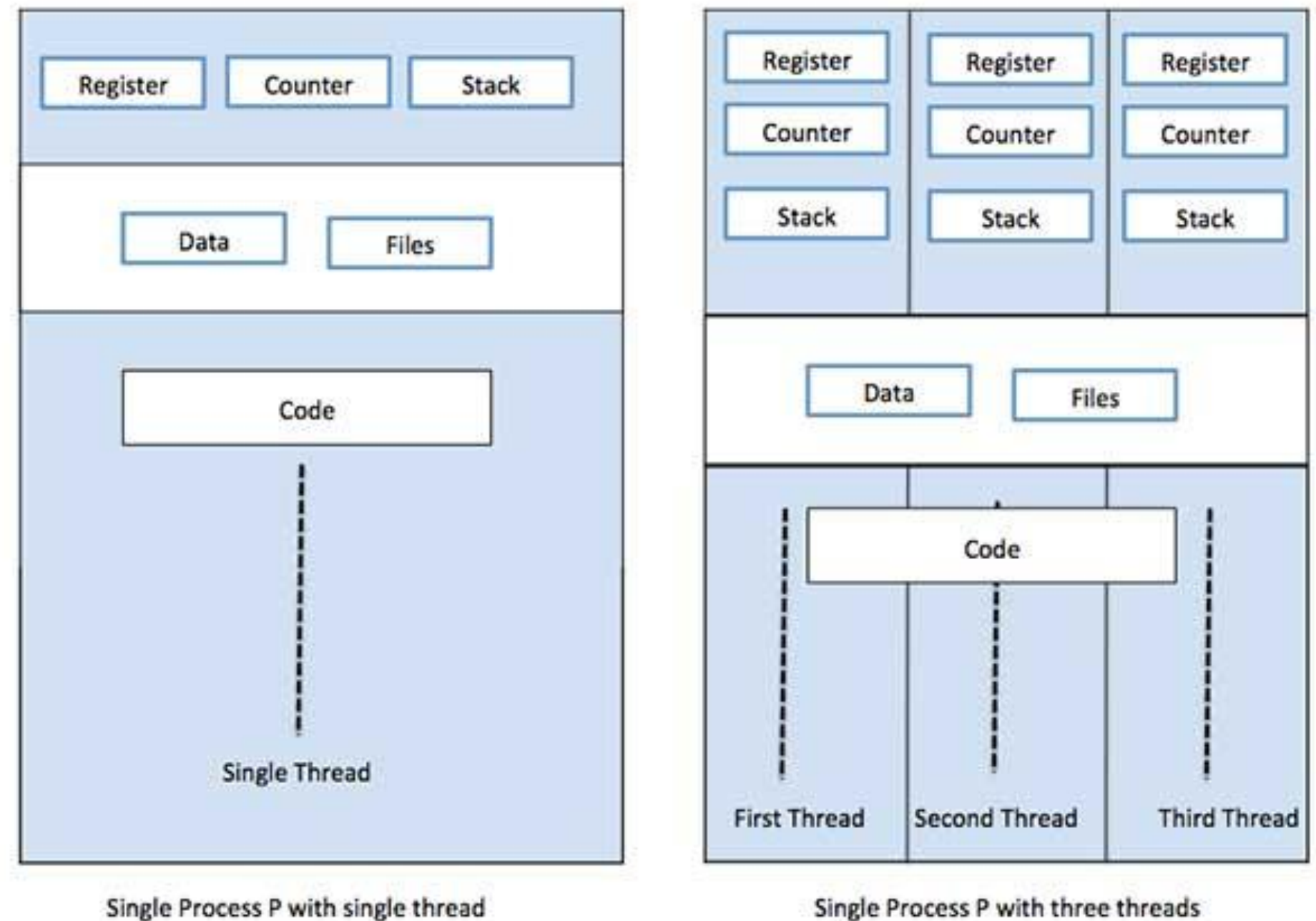
Threading

What is a “Thread”

Often called a “**lightweight process**”

- Minimize **context switching** time
- A “blocking” thread does not block other threads

Question: We know that switching between processes requires interaction from the OS. Is this also the case for threading (thread switching)?



Img src: https://www.tutorialspoint.com/operating_system/os_multi_threading.htm

Threads vs Processes

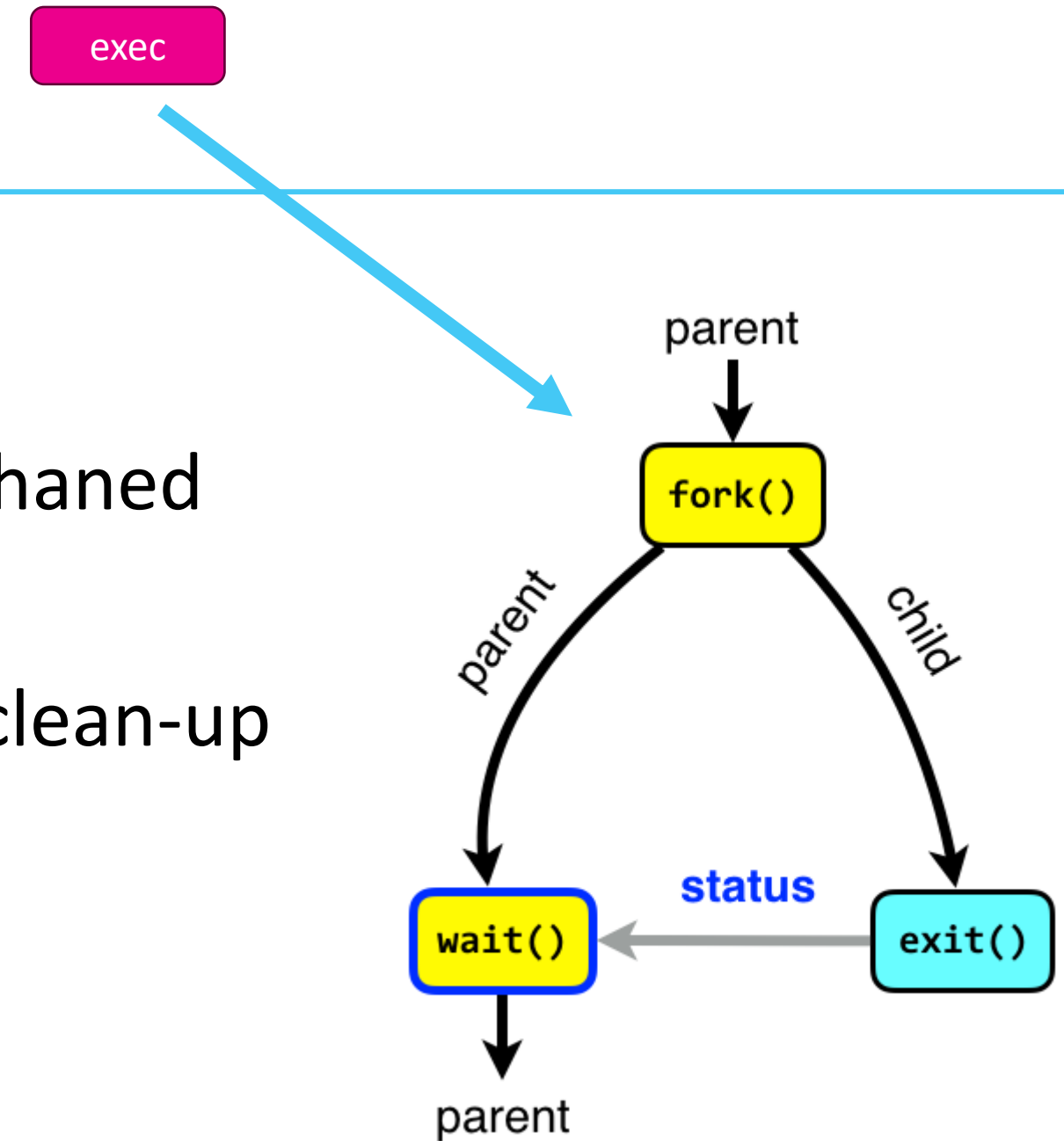
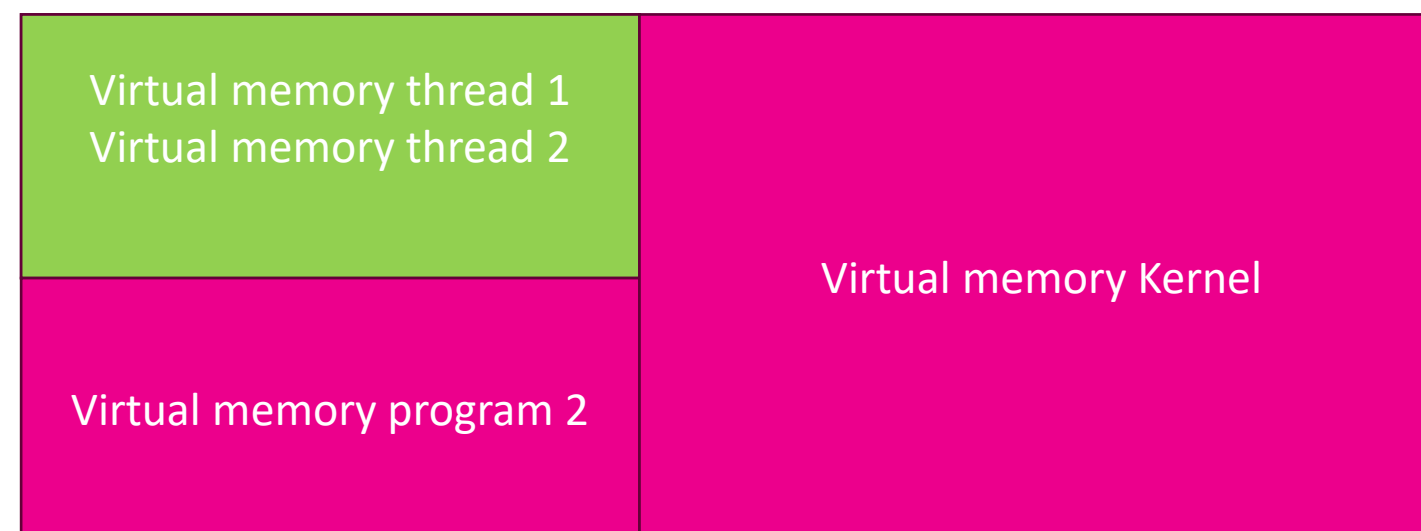
- Processes & scheduling
 - Responsibility with the OS -> CPU scheduling
- Threads & scheduling
 - Responsibility with the application / program designer

```
0110011001111000101101100011000110110
0110001101100110001101100110001101101
1001100011011011001100011011011001100
0110011001111000101101100011000110110
0110001101100110001101100110001101101
0110011001111000101101100011000110110
0110001101100110001101100110001101101
0110001101100110001101100110001101101
0110001101100110001101100110001101101
1001100011011011001100011011011001100
0110011001111000101101100011000110110
0110001101100110001101100110001101101
```

First, solve the problem.
Then, write the code.
~ John Johnson

Threads vs forks (in c)

- Fork -> System call that creates a new child **process**
 - If parent ends before child, the child becomes orphaned
- Pthread_create -> All “part of” the same process
 - There is no exit as virtual memory does not need clean-up



Fork example

How many times will “hello” be printed?

```
File: hello.c
1  #include <stdio.h>
2  #include <sys/types.h>
3  int main()
4  {
5      fork();
6      fork();
7      fork();
8      printf("hello\n");
9      return 0;
10 }
```

```
fork ();    // Line 1
fork ();    // Line 2
fork ();    // Line 3
          L1      // There will be 1 child process
          /  \    // created by line 1.
        L2    L2  // There will be 2 child processes
       / \   / \  // created by line 2
      L3 L3 L3 L3 // There will be 4 child processes
                   // created by line 3
```


Child processes vs Parent processes

Why is this important?

- What happens with the parent process if a child process gets killed?
- **What happens with the child process if a parent process gets killed?**
 - → Actually OS dependend

Multiple options are possible:

- Child process gets killed as soon as parent gets killed (behind the scenes “kill” signals are send to the child process)
- Child becomes an “orphan”
- Child gets to live on without any issues
- Etc.

Attention: Sometimes it gets tricky!

Example 1:

→ Typically (on Unix/Linux)

if a parent dies before the child process. The child becomes an orphan and get's a new parent process (init/systemd (= one of the first processes that starts with the OS)).

Example 2:

→ If the parent process receives proper exit signals – and it is able to send signals to its children – child behaviour depends on what the programmer implemented.

→ Example: SSH into a linux machine, issue a ping and in another session kill the SSH connection. Is your ping still going?

Thread example

```
File: printthread.c
1  #include <stdio.h>
2  #include <stdlib.h>
3  #include <unistd.h> //Header file for sleep(). man 3 sleep for details.
4  #include <pthread.h>
5
6  // A normal C function that is executed as a thread
7  // when its name is specified in pthread_create()
8  void *myThreadFun(void *vargp)
9  {
10     sleep(1);
11     printf("Printing GeeksQuiz from Thread \n");
12     sleep(200);
13     return NULL;
14 }
15
16 int main()
17 {
18     pthread_t thread_id;
19     printf("Before Thread\n");
20     pthread_create(&thread_id, NULL, myThreadFun, NULL);
21     pthread_join(thread_id, NULL);
22     printf("After Thread\n");
23
24     exit(0);
25 }
```

936	943	1	943	debian	20	0	8836	5576	3820	S	0.0	0.1	0:00.14	-bash
943	1379	2	1379	debian	20	0	10648	540	464	S	0.0	0.0	0:00.00	./printthread
943	1379	2	1380	debian	20	0	10648	540	464	S	0.0	0.0	0:00.00	printthread

```
debian@debian:~/threads$ gcc printthread.c -lpthread -o printthread
```

Java example

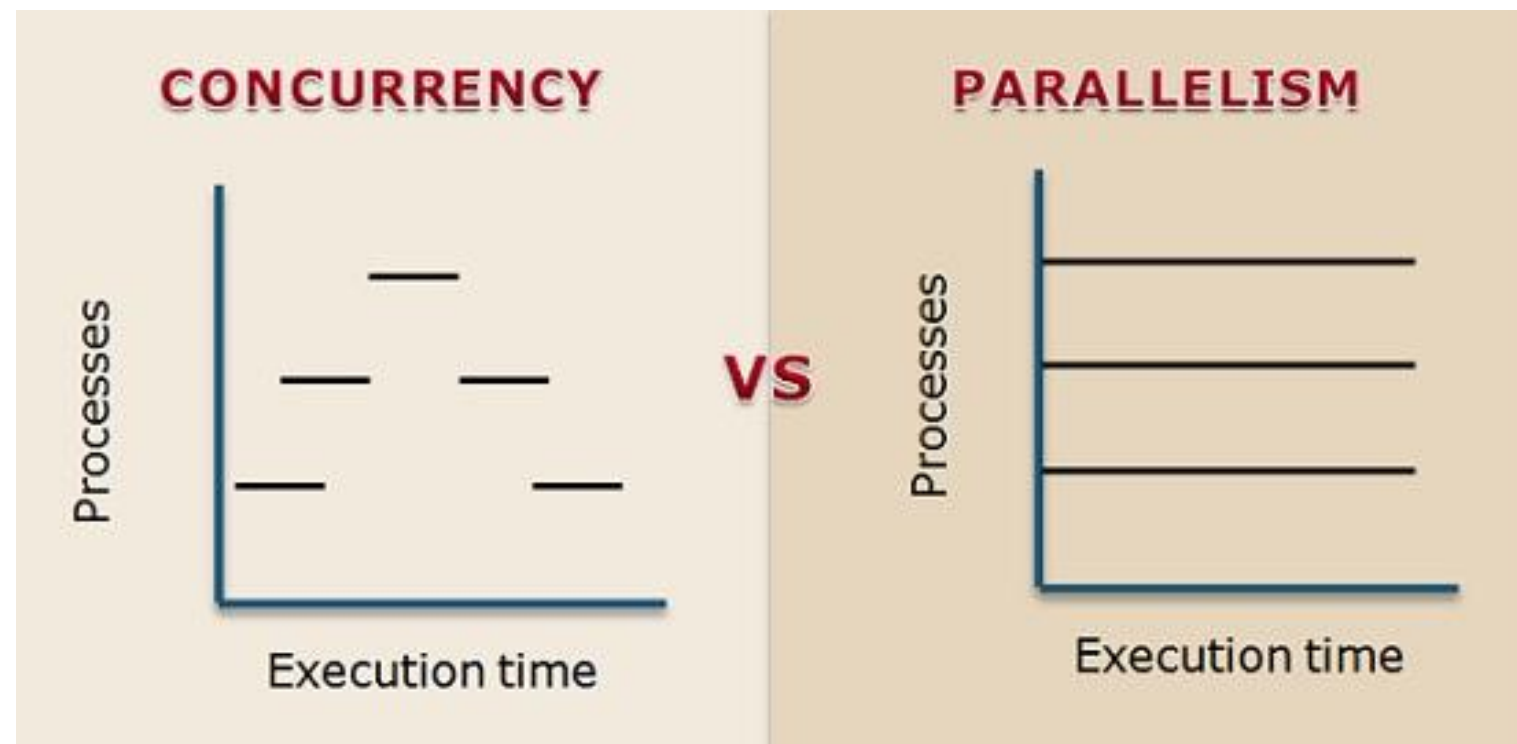
```
File: HelloWorld.java
1  import java.util.*;
2
3  public class HelloWorld {
4      public static void main(String[] args) {
5          System.out.println("Hello World");
6          Scanner sc= new Scanner(System.in);
7          System.out.print("Enter a string: ");
8          String str= sc.nextLine();
9      }
10 }
```

How many threads ?

936	943	1	943	debian	20	0	8820	5800	3884	S	0.0	0.1	0:00.18	-bash
943	1486	12	1486	debian	20	0	2957M	40664	25396	S	0.7	1.0	0:00.13	└─ java HelloWorld
943	1486	12	1487	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.04	└─ java
943	1486	12	1488	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.00	└─ VM Thread
943	1486	12	1489	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.00	└─ Reference Handl
943	1486	12	1490	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.00	└─ Finalizer
943	1486	12	1491	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.00	└─ Signal Dispatch
943	1486	12	1492	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.00	└─ Service Thread
943	1486	12	1493	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.01	└─ C2 CompilerThre
943	1486	12	1494	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.02	└─ C1 CompilerThre
943	1486	12	1495	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.00	└─ Sweeper thread
943	1486	12	1496	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.03	└─ VM Periodic Tas
943	1486	12	1497	debian	20	0	2957M	40664	25396	S	0.0	1.0	0:00.00	└─ Common-Cleaner

Example : where do we use threads?

- “Everywhere”?!
- “We don’t want clients to wait!”
 - Parallelism vs concurrency?
 - Race conditions



Process vs. Thread

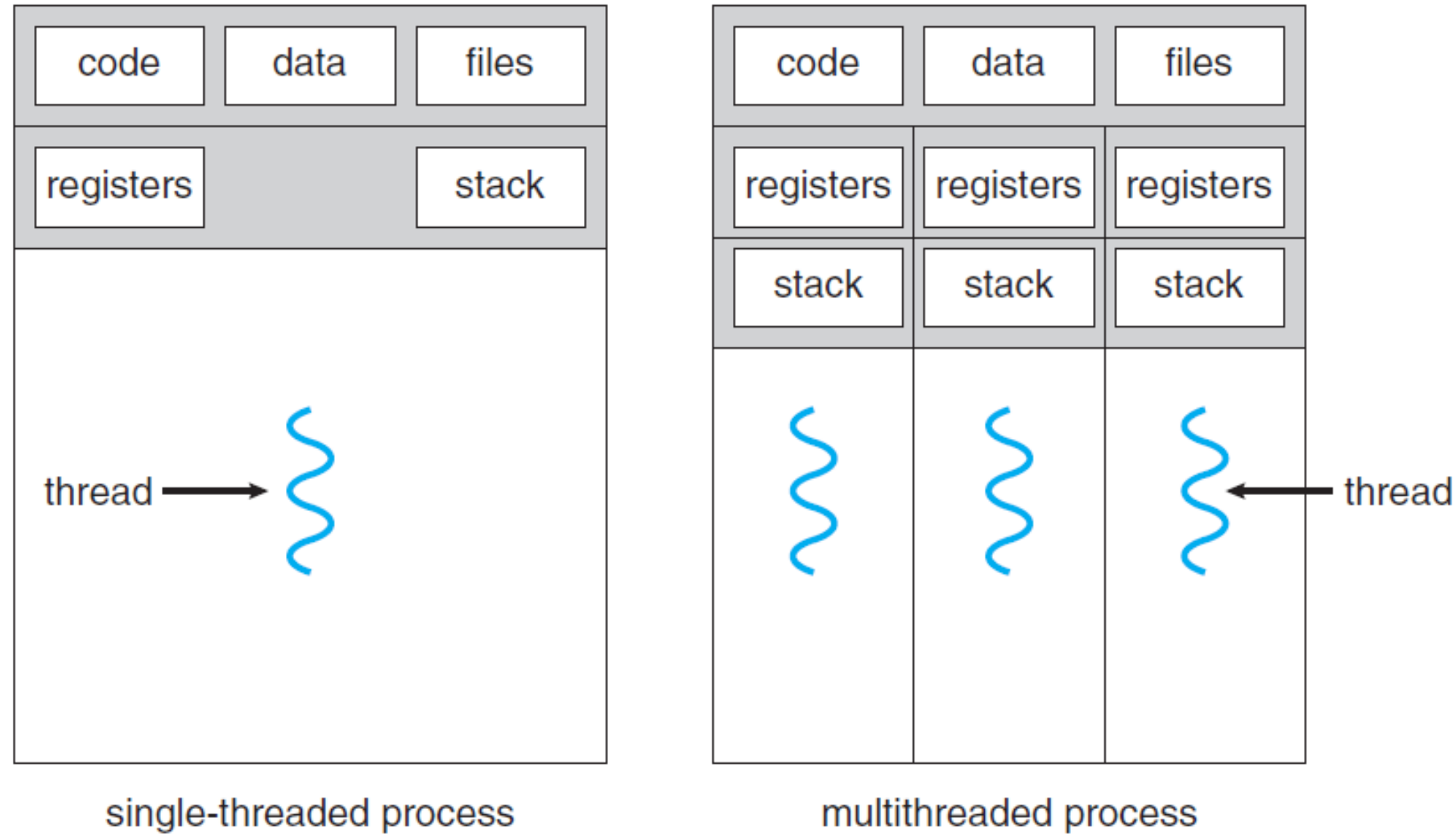


Figure 4.1 Single-threaded and multithreaded processes.

Kernel vs user-level threads

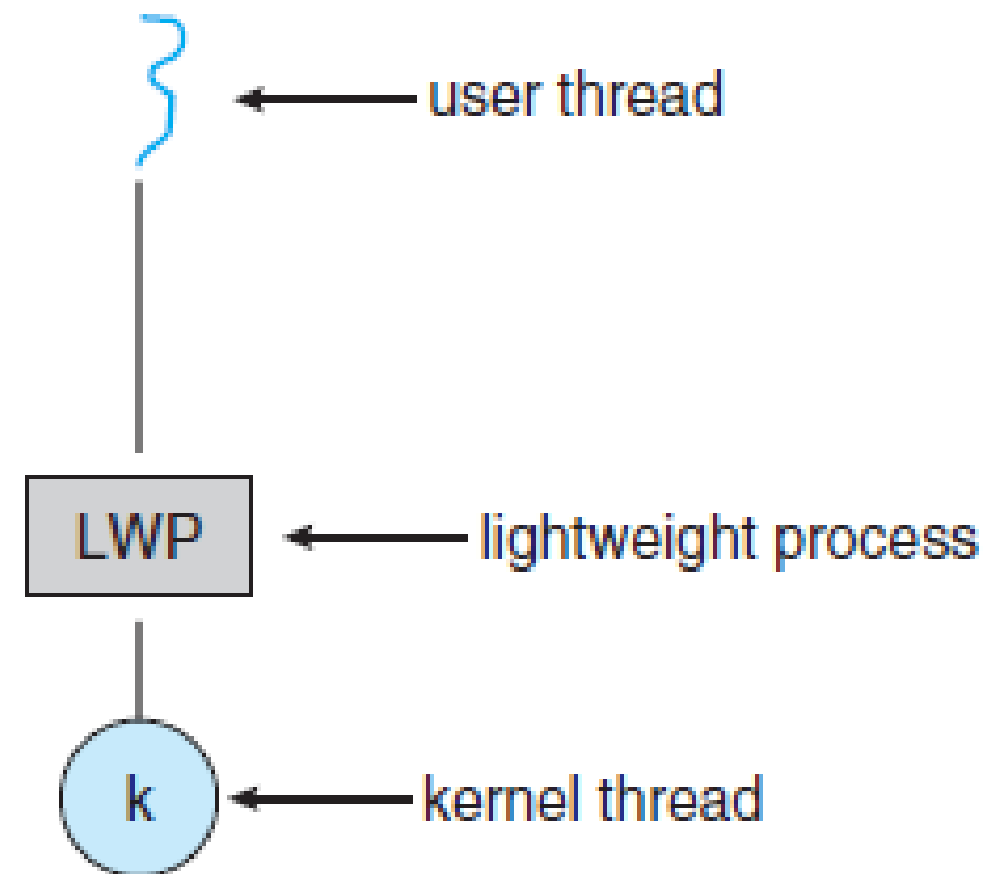


Figure 4.13 Lightweight process (LWP).

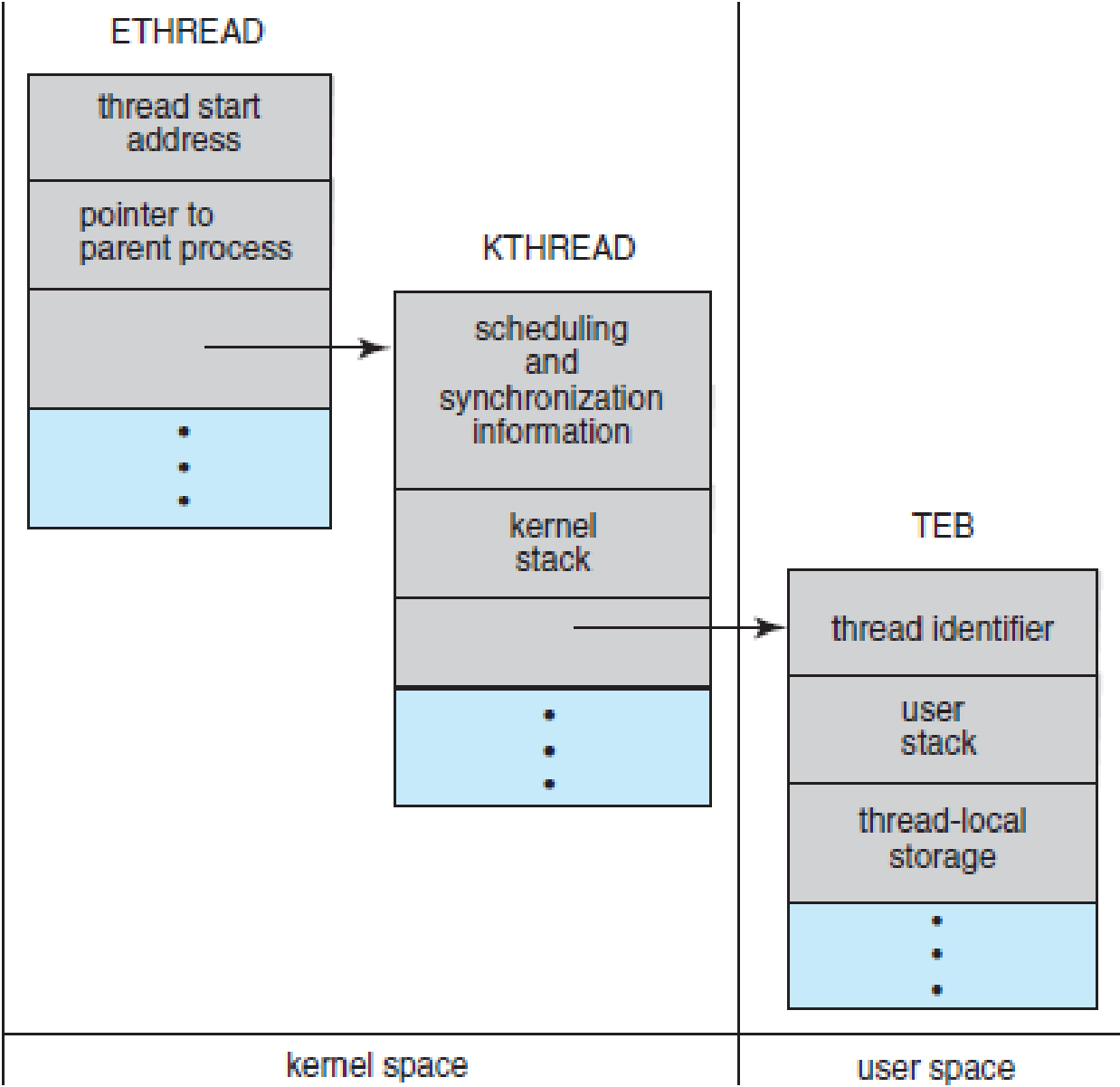


Figure 4.14 Data structures of a Windows thread.

Kernel vs user-level threads & multithreading models

User-level threads

All **code** and **data structures** for the library exist in user space.

Invoking a function in the API results in a **local function call** in user **space** and not a system call.

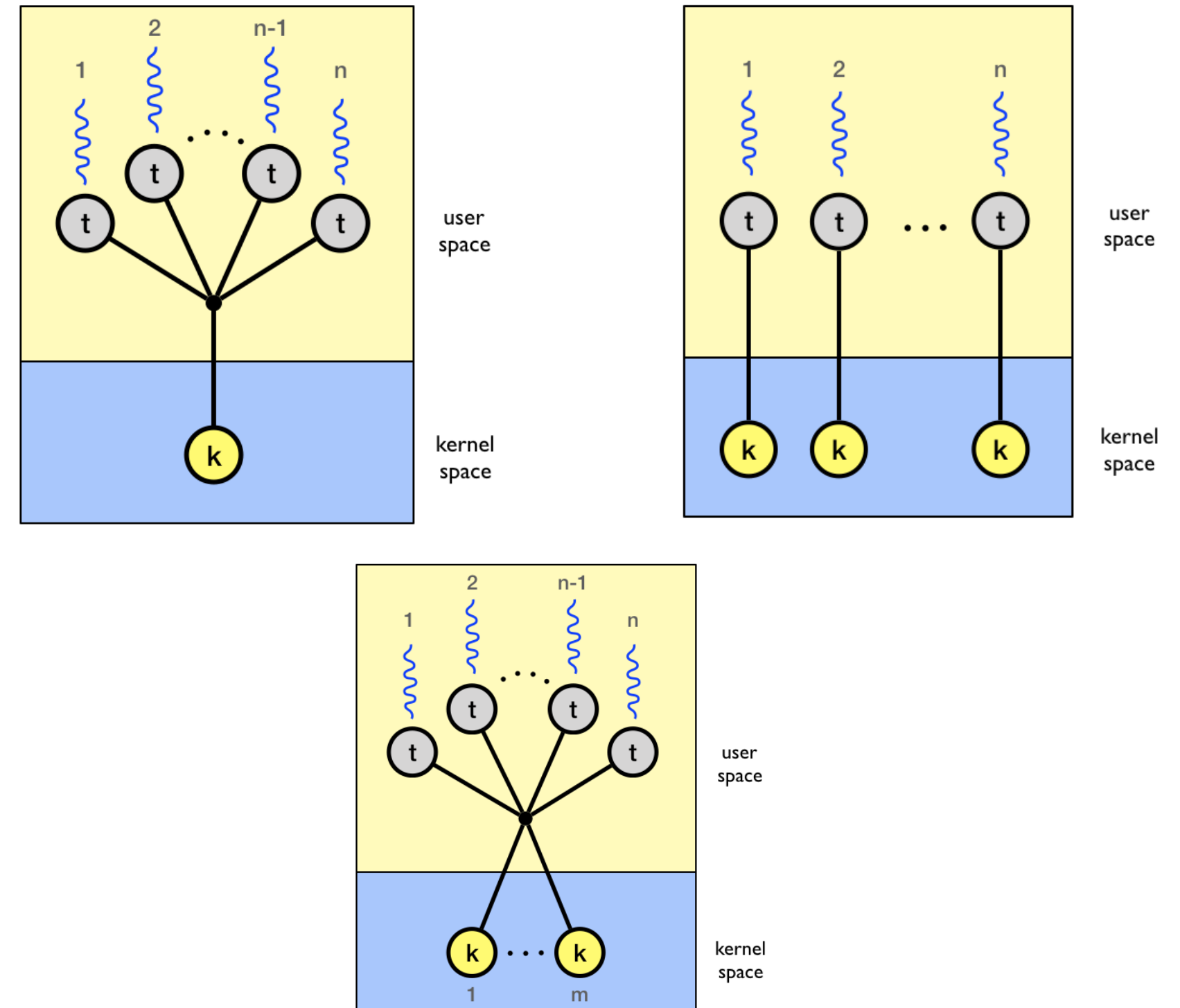
user mode

Kernel-level threads

All **code** and **data structures** for the library exist in **kernel space**.

Invoking a function in the API typically results in a **system call** to the kernel.

kernel mode



Kernel vs user-level threads

Threads

A thread is a basic unit of CPU utilization.

It comprises

- A thread ID
- A program counter
- A register set and
- A stack

It shares with other threads belonging to the same process its code section, data section, and other operating-system resources, such as open files and signals.

A traditional / heavyweight process has a **single thread** of control.

If a process has **multiple threads** of control, it can perform **more than one task at a time**.

The benefits of multithreaded programming can be broken down into four major categories:

- Responsiveness**
Multithreading an interactive application may allow a program to continue running even if part of it is blocked or is performing a lengthy operation, thereby increasing responsiveness to the user.
- Resource sharing**
By default, threads share the memory and the resources of the process to which they belong. The benefit of sharing code and data is that it allows an application to have several different threads of activity within the same address space.
- Economy**
Allocating memory and resources for process creation is costly. Because threads share resources of the process to which they belong, it is more economical to create and context-switch threads.
- Utilization of multiprocessor architectures**
The benefits of multithreading can be greatly increased in a multiprocessor architecture, where threads may be running in parallel on different processors. A single-threaded process can only run on one CPU, no matter how many are available. Multithreading on a multi-CPU machine increases concurrency.

Src: <https://www.youtube.com/watch?v=LOfGJcVnvAk>



On Linux

Some interesting things:

- htop
 - nlwp via f2 setup, add tgid
- ps -e -T
- /proc/<nr>/status
- pstree -p

Main		I/O												Command	
TGID	PID	USER	PRI	NI	VIRT	RES	SHR	S	CPU%	MEM%	TIME+				
1	1	root	20	0	163M	5624	3604	S	0.0	0.0	0:06.43				/sbin/init
730	730	lightdm	20	0	19268	368	368	S	0.0	0.0	0:00.18				/lib/systemd/systemd --user
731	731	lightdm	20	0	164M	128	0	S	0.0	0.0	0:00.00				(sd-pam)
781	781	lightdm	20	0	9128	0	0	S	0.0	0.0	0:00.01				/usr/bin/dbus-daemon --session --addr
773	773	lightdm	9	-11	313M	112	0	S	0.0	0.0	0:00.31				/usr/bin/pulseaudio --daemonize=no --
773	783	lightdm	-6	0	313M	112	0	S	0.0	0.0	0:00.00				/usr/bin/pulseaudio --daemonize=no
787	787	lightdm	20	0	303M	4	4	S	0.0	0.0	0:00.00				/usr/libexec/at-spi-bus-launcher
793	793	lightdm	20	0	8996	0	0	S	0.0	0.0	0:00.00				/usr/bin/dbus-daemon --config-file
787	788	lightdm	20	0	303M	4	4	S	0.0	0.0	0:00.00				/usr/libexec/at-spi-bus-launcher
787	789	lightdm	20	0	303M	4	4	S	0.0	0.0	0:00.00				/usr/libexec/at-spi-bus-launcher
787	791	lightdm	20	0	303M	4	4	S	0.0	0.0	0:00.00				/usr/libexec/at-spi-bus-launcher
821	821	lightdm	20	0	160M	0	0	S	0.0	0.0	0:00.00				/usr/libexec/at-spi2-registr

Process vs. Thread (generally speaking)

- A **thread** is a basic unit of CPU utilization; each thread has a:
 - Thread ID
 - Program Counter
 - Register set
 - Stack
- It **shares** with other threads (belonging to the **same process**):
 - Code section
 - Data section
 - Other OS-resources such as “Open-files”

Why would we want to use threads? - Motivation

- Imagine a browser tab “being stuck” while downloading a file?
 - Other examples?
- **Advantages of multithreaded programming**
 - Responsiveness
 - Resource sharing
 - Economy
(it is less costly to manage threads than processes)
 - Scalability

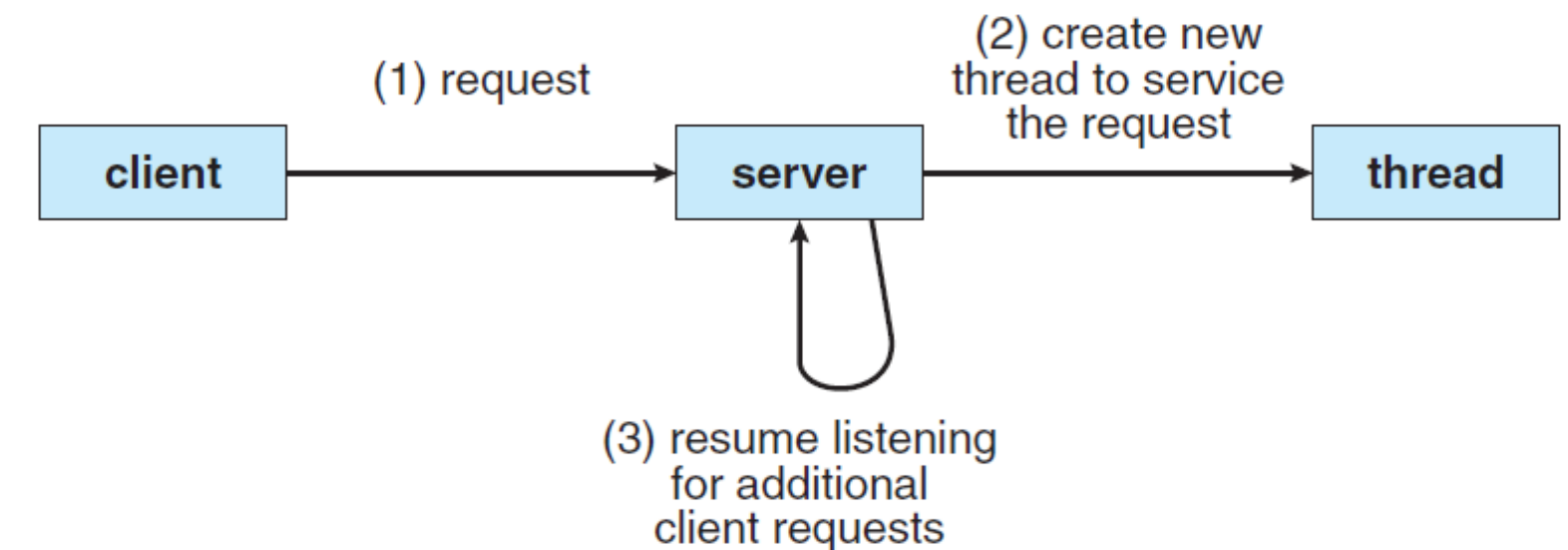
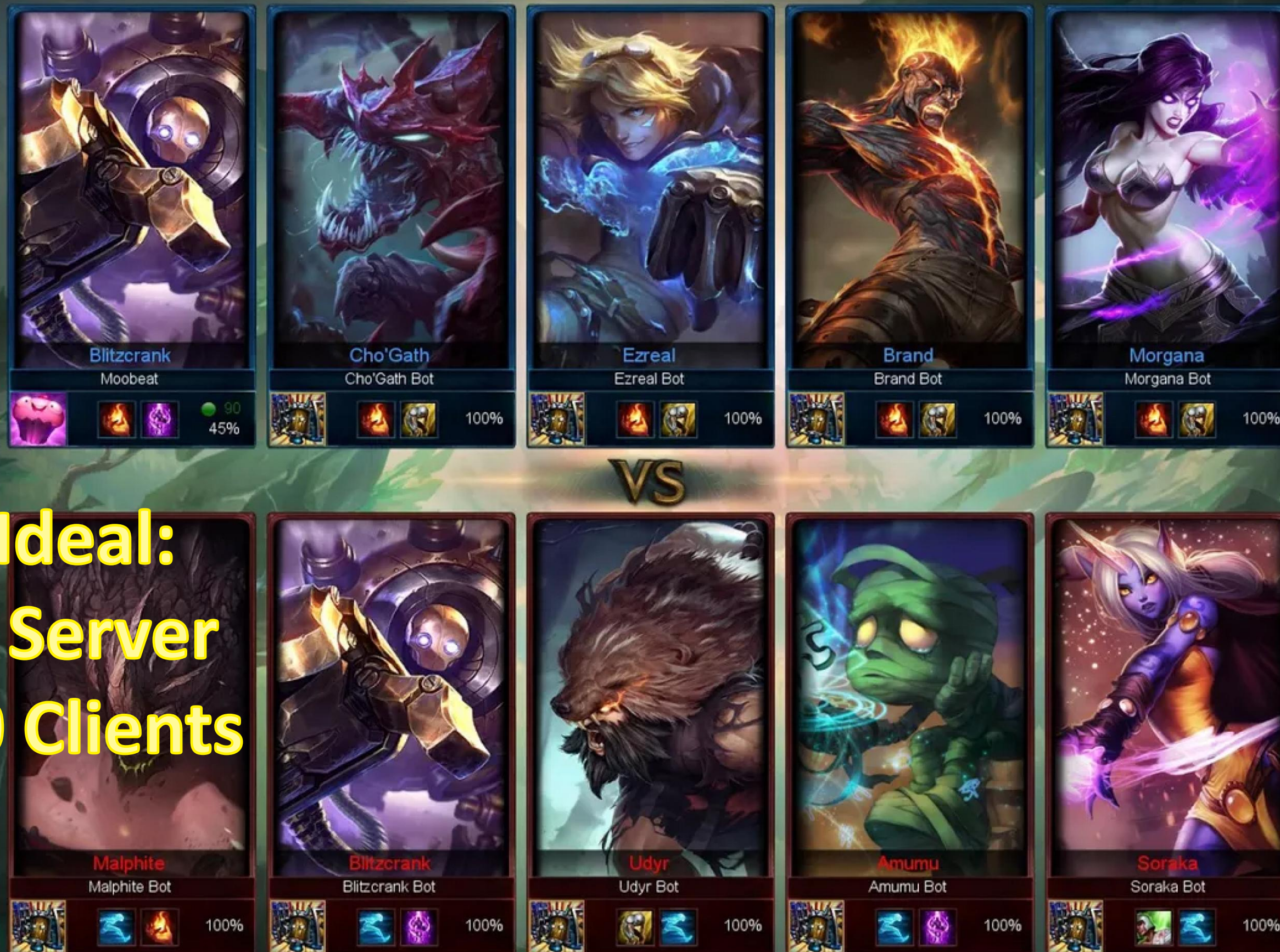


Figure 4.2 Multithreaded server architecture.

**Ideal:
1 Server
10 Clients**



Without threads

- leagueserver.exe
- league_accept_client.exe (for every client)
- leagueclient.exe

With threads

- leagueserver.exe
- leagueclient.exe

Server creates a thread for each client

Disadvantages?

- **Programming** (harder to write, harder to debug, harder to manage)
- **Concurrency issues with shared resources**
 - Similar problems with transactions in databases!
 - Dirty read
 - Loss update
 - Phantom read
- It is never set in stone that a “Java thread” is always mapped to an “OS thread”.
 - You have to trust the libraries/frameworks or write everything yourself if you want that.
 - Luckily we rarely have to worry about that as a programmer 😊
 - -> User and kernel threads

Disadvantages? – Race Conditions

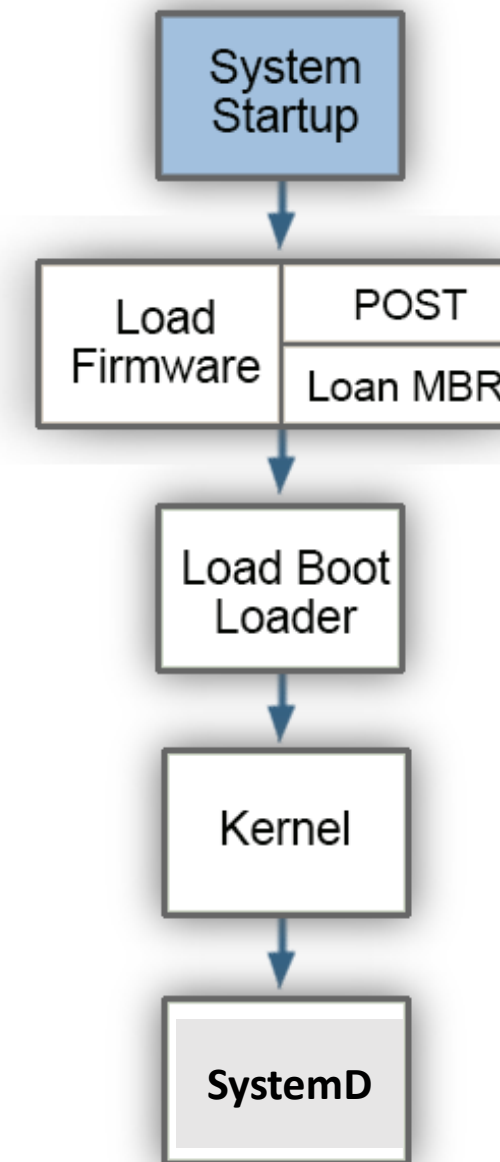
Thread 1	Thread 2		Integer value
			0
read value		←	0
increase value			0
write back		→	1
	read value	←	1
	increase value		1
	write back	→	2

Thread 1	Thread 2		Integer value
			0
read value		←	0
	read value	←	0
increase value			0
	increase value		0
write back		→	1
	write back	→	1

The boot process

Boot Process Overview

- Four main stages
- Some stages can be modified by administrators



Firmware Stage

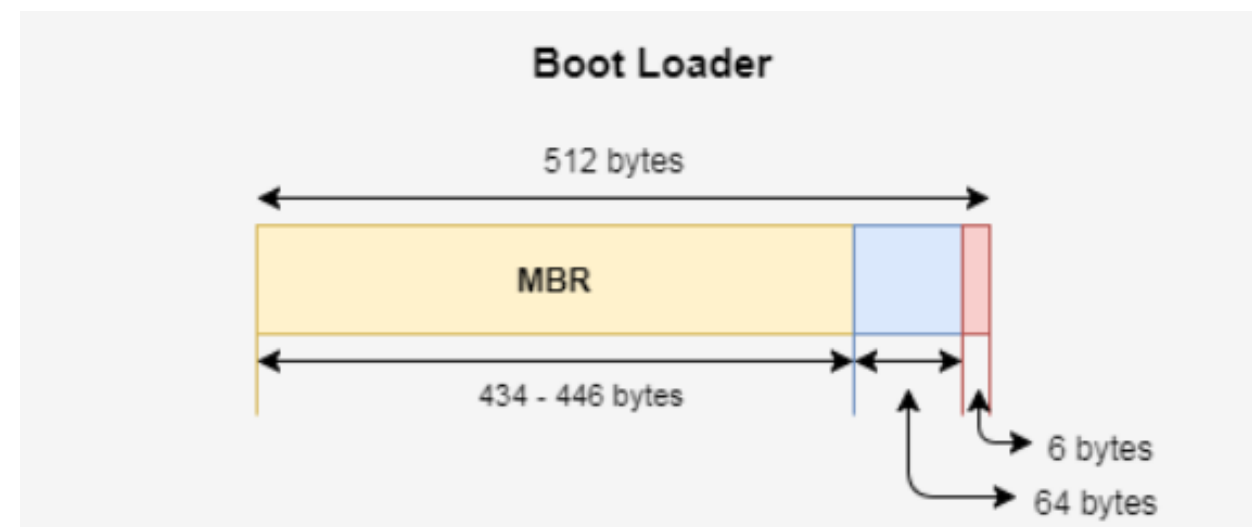
- Firmware is referred to as the BIOS (Basic Input Output System)
- UEFI (Unified Extensible Firmware Interface) has replaced the BIOS on most systems but typically still referred to as “BIOS”

Firmware Stage

- Two primary jobs in this stage:
 - 1) Power-On Self Test (POST) – ensures system hardware (CPU, RAM, peripherals, etc.) is functioning properly
 - 2) Load Master Boot Record (MBR) – contains drive partition table and loads *first stage bootloader* whose purpose is to load the *second stage bootloader* (next stage)

First Stage Boot Loader

BIOS looks for Master Boot Record (MBR) on first found hdd or equivalent – finds partition table and loads *first stage bootloader*



- *Max 446 bytes!*
 - *64 bytes Partition Table, 6 bytes crc*
 - *That bootloader will load the second stage bootloader, who lies elsewhere*
- * (exceptionally this First stage bootloader loads kernel directly)

Bootloaders

- Most common bootloaders
 - LILO - Linux Loader
 - *Supports systems with BIOS
 - *ELILO – Efi Linux Loader supports systems with UEFI
 - GRUB - Grand Unified Bootloader
 - *Supports systems with UEFI
 - *Supports Kernel Flavor choice

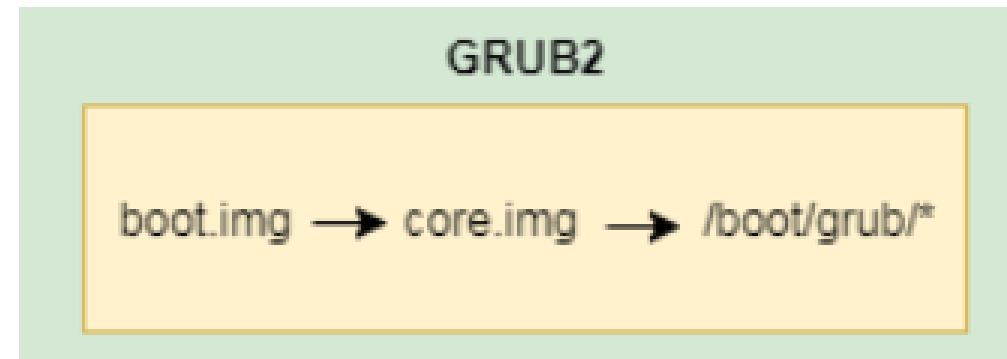
Bootloaders

- Other bootloaders
 - SILO – SPARC Improved bootLoader
 - *Supports Linux on Sun SPARC hardware
 - YABOOT – Yet Another Bootloader
 - Supports PowerPC hardware
- Network booting
 - PXE– Preboot Execution Environment
 - *For hardware that supports TFTP used to download the bootloader from a server

Bootloaders

- Windows
 - BOOTMGR (can be chainloaded from GRUB)
- MacOS
 - IBOOT for INTEL
 - LLB for ARM (Low-Level Bootloader)
 - Can be used alongside GRUB

GRUB2



- Boot.img (1st stage bootloader)
 - <512 bytes on MBR
 - Points to location of core.img
- Core.img (2nd stage | 1,5th stage bootloader)
 - Chosen Kernel image + generic modules to
 - Access files in /boot/grub (the REAL 2nd stage ?)
 - Loads other (kernel) modules at runtime

Kernel phase

- GRUB2 loads the chosen kernel in memory and passes control to it using the files in
- the /boot directory :
 - An initial RAM disk image (initrd / initramfs)
 - temporary / as tmpfs ←!!!
 - the kernel files (vmlinuz) with basic devices
 - Kernel start up : SystemD becomes active

SystemD phase

- Set kernel options using `/etc/sysctl.conf`
- Starts `udev` daemon to detect all (other) devices
- Imports network configuration
- file system check (`fsck`) the root file system if necessary

SystemD phase

- Decrypts filesystems if encrypted.
- Mount filesystems according to /etc/fstab
- the REAL (root) / is mounted
- Enable swap devices
- Eventually
 - The system boots in a specific target mode
- SystemD is managed by systemctl

dmesg Command

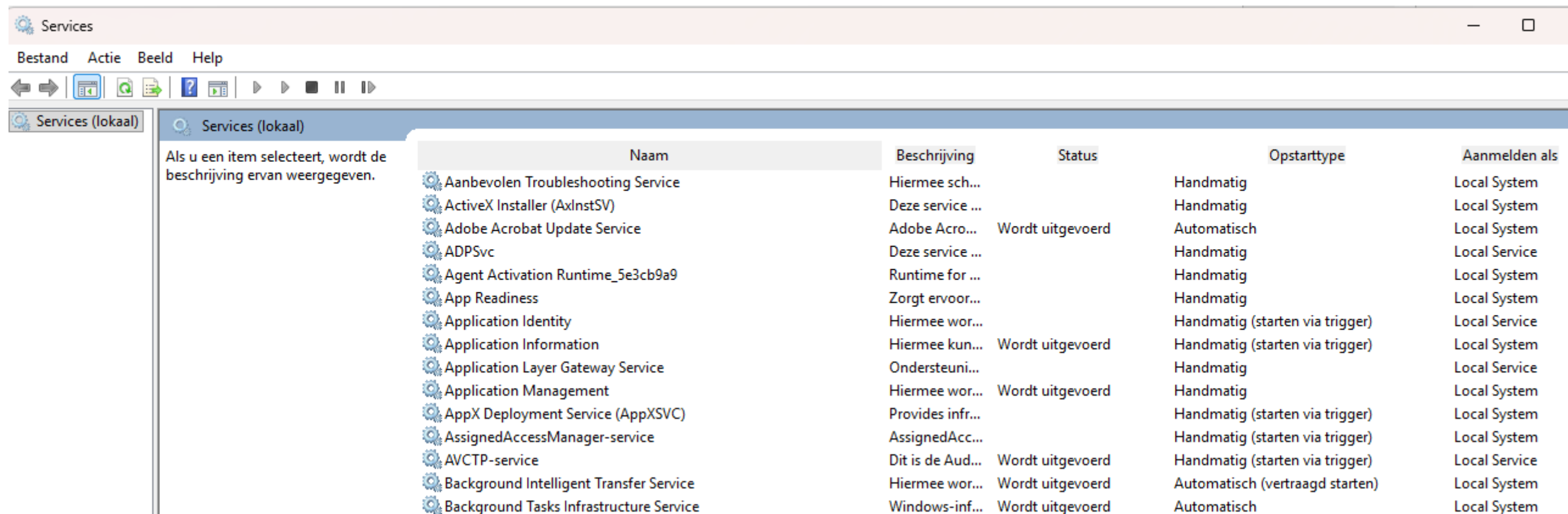
- Executed after boot to view messages generated by the kernel during the boot process
 - Useful for troubleshooting boot issues
- Also executed upon connecting a new device to see device pathname

Daemons

- The administrator can control which services will be provided by the various *daemons*
- A daemon is a running program that provides a service

No Daemons in Windows

- A daemon in windows is called a service

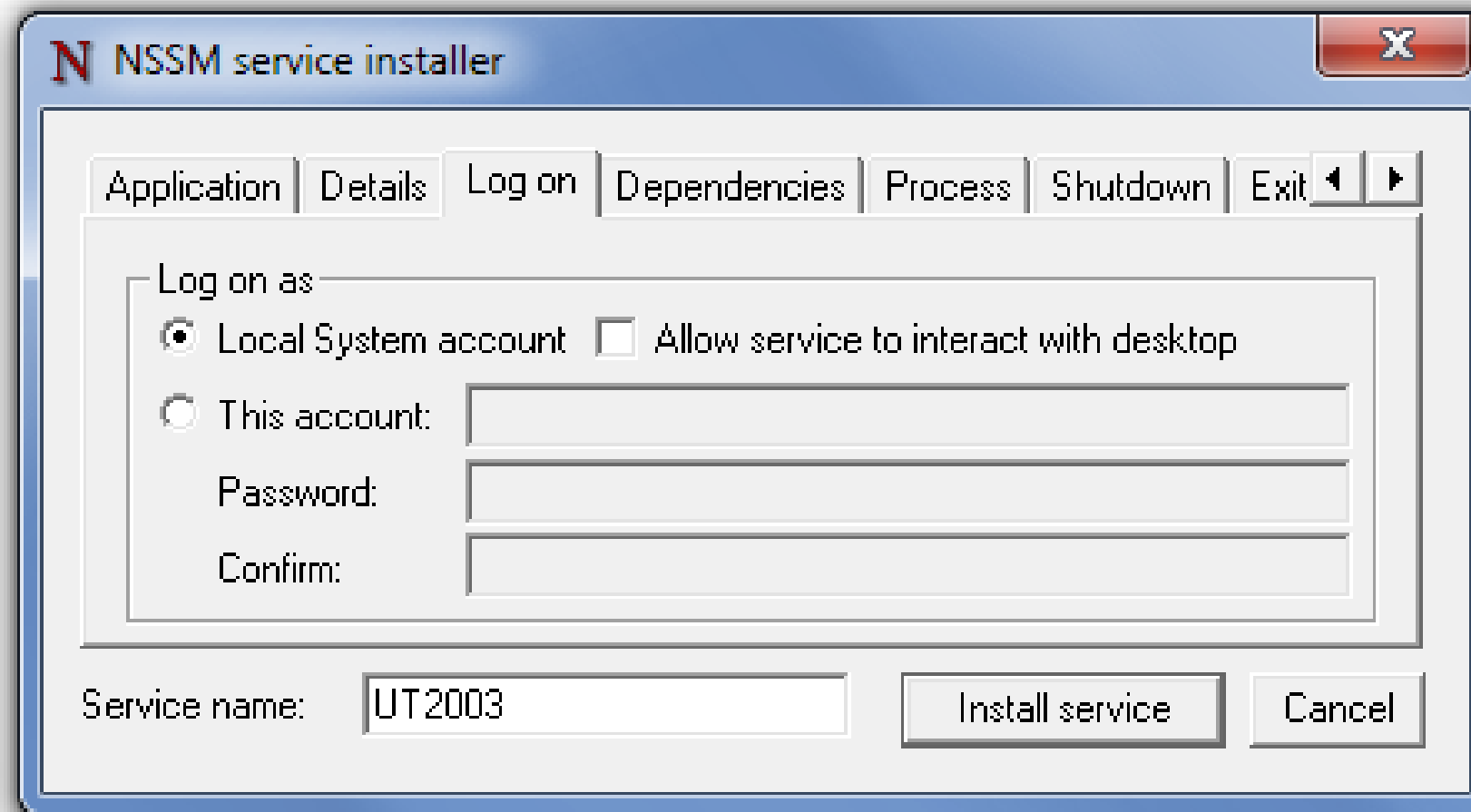


No Daemons in Windows

- Limited wrapping of executables into services is possible
- Need Admin powers
- E.g.
 - `sc create ServiceName binPath= "C:\Path\To\YourExecutable.exe"`
 - `sc start ServiceName`
 - `New-Service -Name "ServiceName" -BinaryPathName "C:\Path\To\YourExecutable.exe" -DisplayName "My Service" -StartupType Automatic`
 - `Start-Service -Name "ServiceName"`

No Daemons in Windows

- NSSM (third party open source tool)



Systemctl controls the daemons

- The `systemctl` command is used in systems that have Systemd
- Services are found in
 - `/etc/systemd/system`
 - They are formal configuration files
- To start a service:
`#systemctl start httpd.service`
- To stop a service:
`#systemctl stop httpd.service`

The systemctl command

- To start check the state of a service:
`#systemctl status httpd.service`
- To view all running services:
`#systemctl -a`
- To configure a service to start automatically:
`#systemctl enable httpd.service`

Targets

A Target is a state where a consistent number of services are running

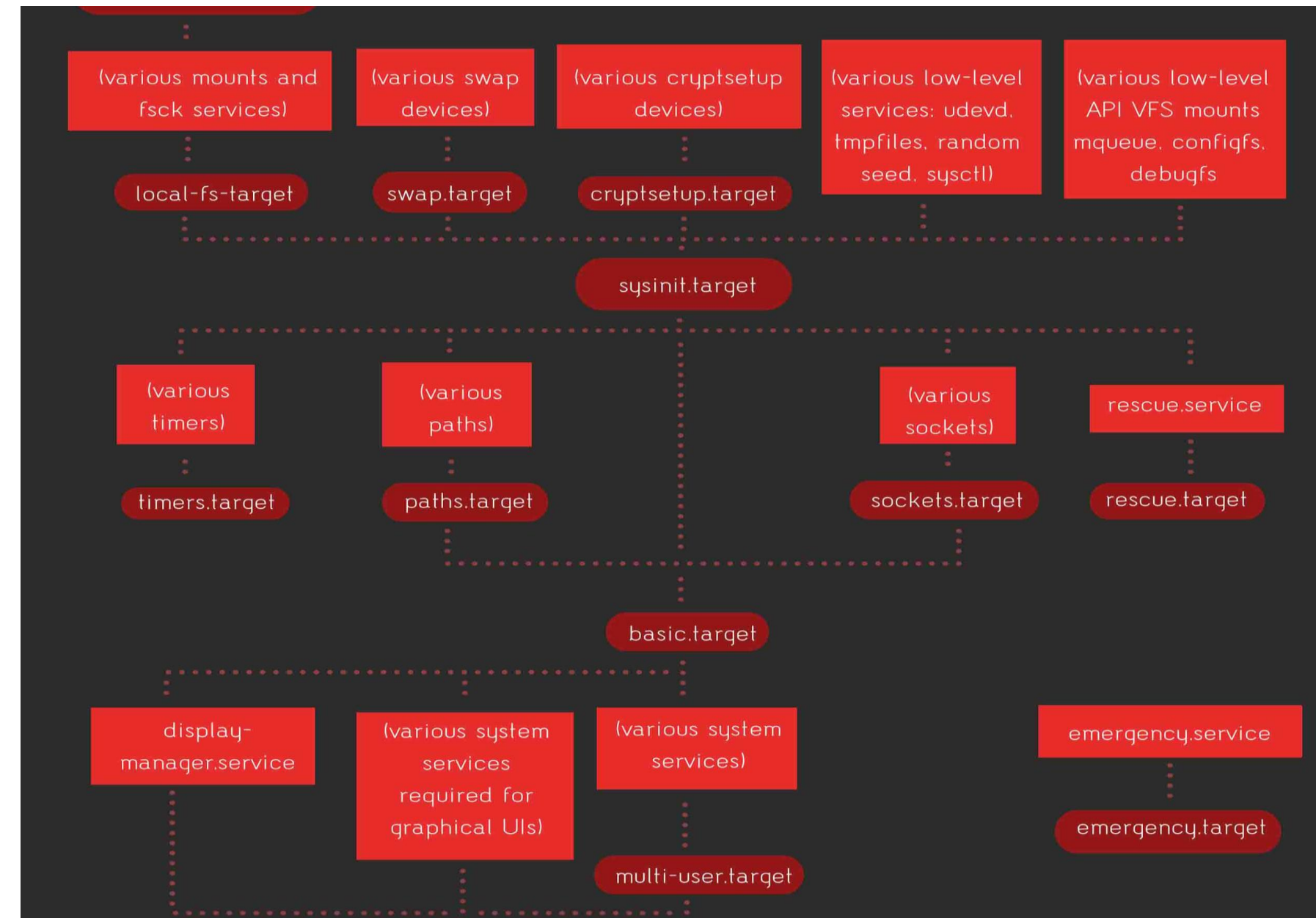
A target has

- Required dependencies

- Parallel dependencies

- Conflicting dependencies

- Sequential dependencies (Before / After)



The systemctl command

- To change the current target state
`#systemctl isolate graphical-target`
- To change the default target state
`#systemctl set-default multi-user.target`

SYSTEMD

In Linux, SYSTEMD is in control

SystemD is in Control

It determines which daemons are running
Consult it using systemctl command (no args)


```
vaneekhouthoutguy@debian-vaneekhouthout-guy-prime: ~  
run-user-1000.mount loaded active mounted /run/user/1000  
sys-fs-fuse-connections.mount loaded active mounted FUSE Control File System  
sys-kernel-config.mount loaded active mounted Kernel Configuration File System  
sys-kernel-debug.mount loaded active mounted Kernel Debug File System  
sys-kernel-tracing.mount loaded active mounted Kernel Trace File System  
systemd-ask-password-plymouth.path loaded active waiting Forward Password Requests to Plymouth Directory Watch  
systemd-ask-password-wall.path loaded active waiting Forward Password Requests to Wall Directory Watch  
init.scope loaded active running System and Service Manager  
session-2.scope loaded active running Session 2 of user vaneekhouthoutguy  
session-5.scope loaded active running Session 5 of user vaneekhouthoutguy  
apache2.service loaded active running The Apache HTTP Server  
apparmor.service loaded active exited Load AppArmor profiles  
atd.service loaded active running Deferred execution scheduler  
bluetooth.service loaded active running Bluetooth service  
connman-wait-online.service loaded active exited Wait for network to be configured by ConnMan  
connman.service loaded active running Connection service  
console-setup.service loaded active exited Set console font and keymap  
cron.service loaded active running Regular background program processing daemon  
dbus.service loaded active running D-Bus System Message Bus  
dnsmasq.service loaded active running dnsmasq - A lightweight DHCP and caching DNS server  
dundee.service loaded active running DUN service  
exim4.service loaded active running LSB: exim Mail Transport Agent  
getty@tty1.service loaded active running Getty on tty1  
ifup@ens33.service loaded active exited ifup for ens33  
ifupdown-pre.service loaded active exited Helper to synchronize boot up for ifupdown  
keyboard-setup.service loaded active exited Set the console keyboard layout
```

.service ?

.service

Configured in a .service file !

Constructed so SystemD knows what (you want it) to do

 Selecteren vaneckhoutguy@debian-vaneckhout-guy-prime: ~

```
vaneckhoutguy@debian-vaneckhout-guy-prime:~$ cat apache.service
[Unit]
Description=Apache web server
After=network.target
Before=nextcloud-web.service
[Service]
ExecStart=/usr/local/apache2/bin/httpd -D FOREGROUND -k start
ExecReload=/usr/local/apache2/bin/httpd -k graceful
Type=notify
Restart=always
[Install]
WantedBy=default.target
RequiredBy=network.target

vaneckhoutguy@debian-vaneckhout-guy-prime:~$
```

The Unit Section

DESCRIPTION :- Human-readable title of the systemd service.

AFTER :- Set dependency on a service. (I want the server to start after the network is online. This typically includes targets or other services.)

BEFORE :- Start current service before specified service. (I want Apache web server running before the service for Nextcloud is started).

```
Selecteren vaneekhoutguy@debian-vaneekhout-guy-prime: ~  
vaneekhoutguy@debian-vaneekhout-guy-prime:~$ cat apache.service  
[Unit]  
Description=Apache web server  
After=network.target  
Before=nextcloud-web.service  
[Service]  
ExecStart=/usr/local/apache2/bin/httpd -D FOREGROUND -k start  
ExecReload=/usr/local/apache2/bin/httpd -k graceful  
Type=notify  
Restart=always  
[Install]  
WantedBy=default.target  
RequiredBy=network.target  
  
vaneekhoutguy@debian-vaneekhout-guy-prime:~$
```

The Service Section

ExecStart:- The command that needs to be executed when the service starts

ExecReload:- (optional).

how a service is restarted. Use this field in case you wish to have a specific restart mechanism.

Type:- start-up type of a process for a given systemd service. Options are simple, exec, forking, oneshot, dbus, notify and idle.

Restart:- (optional)

specifies if/when a service should be restarted or not.

options are no, on-success, on-failure, on-abnormal, on-watchdog, on-abort and always.

The Install Section

is used when you run either `systemctl enable` and `systemctl disable` command for enabling or disabling a service.

WantedBy:

similar to the After and Before fields

used to specify systemd-equivalent "runlevels".

The default.target is when all the system initialization is complete (when the user is asked to log in. user-facing services (like Apache, cron, GNOME-stuff, etc.) use this target.)

RequiredBy:

similar to WantedBy

specifies hard dependencies.

(if a dependency, this service will fail).

Take control

Create/start/Stop/Manage your own service

Creating a root service

Start with a (root) script **update-on-boot.sh**

```
#!/usr/bin/env bash
```

```
if [ ${EUID} -ne 0 ]
```

```
then
```

```
    exit 1 # this is meant to be run as root
```

```
fi
```

```
apt-get update 1>/dev/null 2>>/root/sys-update.log
```

Put it in /root/ directory

Make sure root can execute this script

Test it !

```
./update-on-boot.sh
```

Turning your script into a service

cd to /etc/systemd/system

*create a file **update-on-boot.service***

[Unit]

Description=Keeping my sources minty fresh

After=multi-user.target

[Service]

ExecStart=/usr/bin/bash /root/update-on-boot.sh

Type=simple

[Install]

WantedBy=multi-user.target

Enabling the service

sudo systemctl daemon-reload

makes the systemd daemon aware of the new service

sudo systemctl enable update-on-boot.service

```
root@debian-vaneeckhout-guy-prime:/etc/systemd/system# sudo systemctl enable update-on-boot.service
Created symlink /etc/systemd/system/multi-user.target.wants/update-on-boot.service → /etc/systemd/system/update-on-boot.service.
root@debian-vaneeckhout-guy-prime:/etc/systemd/system#
```

This creates a symlink in the multi-user.target.wants (remember the targets ?)

How to check if it is enabled ?

```
root@debian-vaneeckhout-guy-prime:/etc/systemd/system# sudo systemctl is-enabled update-on-boot.service
enabled
```

So every time a system enters this target, the script will be executed !

Creating a USER script

Create a script keep-uptime.sh to keep the time your system is up

```
#!/usr/bin/env bash  
uptime | tee -a ${HOME}/uptime.log
```

Make it executable

Test it !

Creating a service for a regular user

Make another .service file 😊

cd to /etc/systemd/system ?

~/.config/systemd/user/ !

Directory not present by default

mkdir -p

cd to ~/.config/systemd/user/

Creating a service for a regular user

Create a file keep-uptime.service in ~/.config/systemd/user/

Content :

[Unit]

Description=Log uptime in scoreboard

DefaultDependencies=no

Before=shutdown.target

[Service]

Type=oneshot

ExecStart=/usr/bin/bash %h/keep-uptime.sh

TimeoutStartSec=0

[Install]

WantedBy=shutdown.target

Enabling your user service

makes the systemd daemon aware of the new service

No sudo necessary !

systemctl --user daemon-reload

systemctl --user enable keep-uptime.service

```
vaneckhoutguy@debian-vaneckhout-guy-prime:~$ systemctl --user enable keep-uptime.service
Created symlink /home/vaneckhoutguy/.config/systemd/user/shutdown.target.wants/keep-uptime.service → /home/vaneckhoutguy/.config/systemd/user/keep-uptime.service.
vaneckhoutguy@debian-vaneckhout-guy-prime:~$
```

Start & Check your user service

Start it !

systemctl --user start keep-uptime

Check what happened (log)

systemctl --user status keep-uptime

```
vaneckhoutguy@debian-vaneckhout-guy-prime:~$ systemctl --user status keep-uptime
• keep-uptime.service - Log uptime in scoreboard
   Loaded: loaded (/home/vaneckhoutguy/.config/systemd/user/keep-uptime.service; enabled; vendor preset: enabled)
   Active: inactive (dead) since Wed 2022-11-16 13:41:28 CET; 4s ago
   Process: 3108 ExecStart=/usr/bin/bash /home/vaneckhoutguy/keep-uptime.sh (code=exited, status=0/SUCCESS)
   Main PID: 3108 (code=exited, status=0/SUCCESS)
      CPU: 9ms

Nov 16 13:41:28 debian-vaneckhout-guy-prime systemd[1206]: Starting Log uptime in scoreboard...
Nov 16 13:41:28 debian-vaneckhout-guy-prime bash[3110]: 13:41:28 up 3:18, 1 user, load average: 0.13, 0.15, 0.12
Nov 16 13:41:28 debian-vaneckhout-guy-prime systemd[1206]: keep-uptime.service: Succeeded.
Nov 16 13:41:28 debian-vaneckhout-guy-prime systemd[1206]: Finished Log uptime in scoreboard.
vaneckhoutguy@debian-vaneckhout-guy-prime:~$
```



Want more

Let's build a true server service

Chatbot 1.0

Dumb chatbot.php that answers anything you type in gibberish

```
<?php
$sock = socket_create(AF_INET, SOCK_DGRAM, SOL_UDP);
socket_bind($sock, '0.0.0.0', 10000);
for (;;) {
    socket_recvfrom($sock, $message, 1024, 0, $ip, $port);
    $reply = str_rot13($message);
    socket_sendto($sock, $reply, strlen($reply), 0, $ip, $port);
}
```

Run it !

Php chatbot.php

Test it in another terminal

nc -u 127.0.0.1 10000

Turning Chatbot 1.0 into a service

Let's create a chatbot.service file for our service

[Unit]

Description=Dumb chat service

After=network.target

StartLimitIntervalSec=0

[Service]

Type=simple

Restart=always



RestartSec=1



ExecStart=/usr/bin/env php %h/chatbot.php

[Install]

WantedBy=default.target

What steps do you need to get it running ?

Examine your service process

systemctl --user status chatbot

```
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ systemctl --user status chatbot
• chatbot.service - Dumb chat service
   Loaded: loaded (/home/vaneckhoutguy/.config/systemd/user/chatbot.service; enabled; vendor preset: enabled)
   Active: active (running) since Wed 2022-11-16 14:27:27 CET; 3s ago
 Main PID: 3494 (php)
    Tasks: 1 (limit: 2284)
   Memory: 5.2M
      CPU: 20ms
   CGroup: /user.slice/user-1000.slice/user@1000.service/app.slice/chatbot.service
           └─3494 php /home/vaneckhoutguy/chatbot.php

Nov 16 14:27:27 debian-vaneckhout-guy-prime systemd[1206]: Started Dumb chat service.
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ nc -u 127.0.0.1 10000
Hi Daemon
Jv Qnrzba
^C
```

Examine your service process

ps -ef | grep php


```
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ ps -ef | grep php
vaneck+    3494    1206  0 14:27 ?        00:00:00 php /home/vaneckhoutguy/chatbot.php
vaneck+    3505    3222  0 14:28 pts/1    00:00:00 grep php
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ systemctl --user stop chatbot
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ ps -ef | grep php
vaneck+    3508    3222  0 14:28 pts/1    00:00:00 grep php
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ systemctl --user start chatbot
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ ps -ef | grep php
vaneck+    3510    1206  0 14:28 ?        00:00:00 php /home/vaneckhoutguy/chatbot.php
vaneck+    3512    3222  0 14:28 pts/1    00:00:00 grep php
```

Examine your service process

Try and kill your server process !

Kill pid_of_chatbot

```
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ ps -ef | grep php
vaneck+    3494    1206  0 14:27 ?        00:00:00 php /home/vaneckhoutguy/chatbot.php
vaneck+    3505    3222  0 14:28 pts/1    00:00:00 grep php
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ systemctl --user stop chatbot
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ ps -ef | grep php
vaneck+    3508    3222  0 14:28 pts/1    00:00:00 grep php
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ systemctl --user start chatbot
vaneckhoutguy@debian-vaneckhout-guy-prime:~/.config/systemd/user$ ps -ef | grep php
vaneck+    3510    1206  0 14:28 ?        00:00:00 php /home/vaneckhoutguy/chatbot.php
vaneck+    3512    3222  0 14:28 pts/1    00:00:00 grep php
```



Examine your service process

Your chatbot gets respawned automatically !

```
vaneekhoutguy@debian-vaneekhout-guy-prime:~/.config/systemd/user$ ps -ef | grep chat
vaneck+  3530    1206  0 14:30 ?        00:00:00 php /home/vaneekhoutguy/chatbot.php
vaneck+  3532    3222  0 14:30 pts/1    00:00:00 grep chat
vaneekhoutguy@debian-vaneekhout-guy-prime:~/.config/systemd/user$ kill 3530
vaneekhoutguy@debian-vaneekhout-guy-prime:~/.config/systemd/user$ ps -ef | grep chat
vaneck+  3533    1206  0 14:30 ?        00:00:00 php /home/vaneekhoutguy/chatbot.php
vaneck+  3535    3222  0 14:30 pts/1    00:00:00 grep chat
vaneekhoutguy@debian-vaneekhout-guy-prime:~/.config/systemd/user$ systemctl --user status chatbot
• chatbot.service - Dumb chat service
   Loaded: loaded (/home/vaneekhoutguy/.config/systemd/user/chatbot.service; enabled; vendor preset: enabled)
   Active: active (running) since Wed 2022-11-16 14:30:39 CET; 20s ago
 Main PID: 3533 (php)
   Tasks: 1 (limit: 2284)
  Memory: 5.2M
    CPU: 22ms
   CGroup: /user.slice/user-1000.slice/user@1000.service/app.slice/chatbot.service
           └─3533 php /home/vaneekhoutguy/chatbot.php

Nov 16 14:30:39 debian-vaneekhout-guy-prime systemd[1206]: chatbot.service: Scheduled restart job, restart counter is at 1.
Nov 16 14:30:39 debian-vaneekhout-guy-prime systemd[1206]: Stopped Dumb chat service.
Nov 16 14:30:39 debian-vaneekhout-guy-prime systemd[1206]: Started Dumb chat service.
vaneekhoutguy@debian-vaneekhout-guy-prime:~/.config/systemd/user$
```