## Esercizio 1

$$T(n) = \begin{cases} 1, \text{ se } n \leq 2 \\ 3, T(\sqrt{n}) + 2, T(\sqrt{n}) + \log(n), \text{ altriment:} \end{cases}$$

				1 1 2	1 2	1	4					
		1 4	1 4	188	18	118	1 16	16				
3 -	1		2 . :									

=> 9.1 + 12.1 + 4.1

Livello	Nodi per livello	in in put	Contributo pur nodo	Contributo totale per eiveleo
0	1	n	log(n)	log(n)
1	5	12 0 n	1/2 log(n) 0 1/4 log(n)	2 log(n)
2	25	n o n o n 46	1 lag(n) 0 1 lag(n) 0 1 lag(n)	Alag(n)
i	5 <sup>i</sup>	,	/	ei log(n)

$$\frac{4}{n^{4\min}} \leq 2 \quad \stackrel{\text{$<$}}{=} \quad \frac{1}{4^{\min}} \quad \log(n) \leq 1 \quad \stackrel{\text{$<$}}{=} \quad 2^{2\min} \quad 2 \log(n) \quad \stackrel{\text{$<$}}{=} \quad 1 \log(n) \leq 1 \log(n)$$

$$\frac{4}{2^{\text{lmax}}} \leq 2 \quad < \approx > \quad \underbrace{4}_{2^{\text{lmax}}} \log(n) \leq 4 < \approx > 2^{\text{lmax}} > \log(n) < \approx > \text{lmax} > \log(\log(n))$$

Calcolo della sommatoria

$$\frac{\sum_{i=0}^{n} 2^{i} \log(n)}{\sum_{i=0}^{n} 2^{i} \log(n)} \leq T(n) \leq \sum_{i=0}^{n} 2^{i} \log(n)$$

$$lag(n) \stackrel{hmin}{\underset{i=0}{\not=}} 2^{i} \leq T(n) \leq lag(n) \stackrel{hmax}{\underset{i=0}{\not=}} 2^{i}$$

$$lag(n)$$
  $\frac{2^{n+n+4}}{2-1} \le T(n) \le lag(n)$   $\frac{2^{n+n+4}}{2-1}$ 

$$\log \binom{n}{2} \binom{\frac{1}{2} \log \binom{\log (n)+1}{2}}{-1} \leq T \binom{n}{2} \leq \log \binom{n}{2} \binom{\log (\log (n))+1}{-1}$$

$$\log\left(n\right)\left(2\cdot2^{\frac{4}{2}\log(\log(n))}-1\right)\leq\tau(n)\leq\log\left(n\right)\left(2\cdot2^{\log(\log(n))}-1\right)$$

$$log(n)(2.\frac{1}{2}log(n)-1) \le T(n) \le log(n)(2log(n)-1)$$

$$\log(n)^2 - \log(n) \le \tau(n) \le 2\log(n)^2 - \log(n) \Longrightarrow \Re(\log(n)^2)$$

