Week 1 Friday Review Quiz

Student Name

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Q1 Combining languages

5 Points

For this question, consider an arbitrary alphabet Σ and, whenever L_1, L_2 are sets of strings over Σ , we can use the following rules to define associated sets of strings:

$$\mathrm{SUBSTRING}(L_1) := \{ w \in \Sigma^* \mid \mathrm{there} \; \mathrm{exist} \; a, b \in \Sigma^* \; \mathrm{such} \; \mathrm{that} \; awb \in L_1 \}$$

and

$$L_1 \circ L_2 := \{ w \in \Sigma^* \mid w = uv \text{ for some strings } u \in L_1 \text{ and } v \in L_2 \}$$

For the statements below, let $\Sigma = \{0,1\}$ be the alphabet.

Select all and only the true statements.

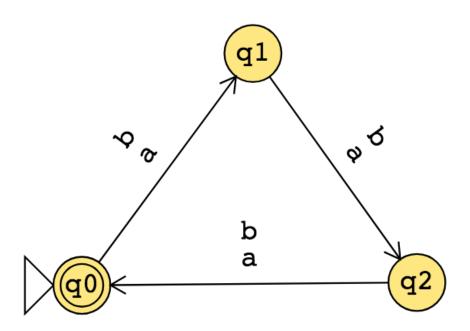
$$\ \, \square \ \, SUBSTRING(\{0\}) = SUBSTRING(\{1\}) = \{\varepsilon\}$$

$$\square$$
 $SUBSTRING(\emptyset) = \emptyset$

$$\square$$
 $SUBSTRING(\Sigma^*) = \Sigma^*$

Q2 Strings in a language recognized by a DFA 1 Point

Select all (and only) the strings below that are accepted by the DFA.

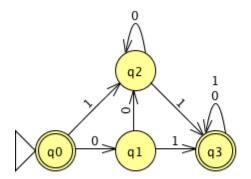


☐ The empty string
а
b
abab
ab
☐ bbb
☐ bba

Q3 Describing DFA

4 Points

Consider the DFA, ${\cal M}$, given by the state diagram:



The author of this DFA claims that its formal definition is: $M = (\{q0,\,q1,\,q2,\,q3\},\,\{0,\,1,\,2,\,3\},\,\delta,\,q0,\,q3) \text{ with } \delta \text{ given by the table below:}$

	0	1
$\overline{q0}$	q1	$\overline{q2}$
q1	q2	q3
q2	q2	q3
q3	q3	q3

Select all and only the components of the formal definition that are correct.

☐ Set of states
☐ Input alphabet
☐ Transition function
☐ Start state
☐ Set of accept states

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Q3.2 (b)
1 Point
True or False: The empty string is accepted by this DFA.
   True
   False
  Save Answer
Q3.3 (c)
1 Point
True or False: L(M) is infinite.
   True
   False
  Save Answer
Q3.4 (d)
1 Point
True or False: If x \in L(M), the string obtained by flipping each
bit in x (changing 0 to 1 and 1 to 0) is also in L(M).
   True
   False
```

	7		
Save Answer			

Q4 Feedback

0 Points