

0/9 Questions Answered

Week 2 Wednesday Review Quiz

Student Name

Search students by name or email...



Q1 Iterated transition function

1 Point

For some proofs it is useful to define the iterated transition function of a finite automaton. Given a finite automaton $(Q, \Sigma, \delta, q_0, F)$ the iterated transition function δ^* has domain $Q \times \Sigma^*$ and codomain Q and, on input (q, w) , outputs the state that the machine is in when starting in state q and processing the entire string w character-by-character according to the transition function δ .

Since the string w is built up recursively, we will define the iterated transition function recursively. Choose the correct option to fill in this definition:

The basis case is for $w = \varepsilon$ and we define

$$\delta^*((q, \varepsilon)) =$$

ε

Σ

q

q_0

F

The recursive step is when $w = ua$ for $u \in \Sigma^*$ and $a \in \Sigma$, and we define

$$\delta^*((q, w)) =$$

$\delta((\delta^*((q, u))), a))$

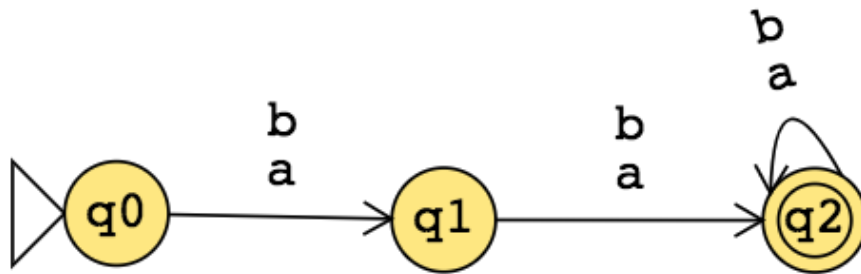
$\delta((\delta((q, u))), a))$

Save Answer

Q2 Iterated transition function example

2 Points

Let's consider the finite automaton whose state diagram is below. Notice that we can deduce the alphabet from the diagram: the arrow labels are the symbols in the alphabet so $\Sigma = \{a, b\}$. Let's refer to the transition function of this automaton by δ .



Q2.1 (a)

1 Point

What is $\delta((q_0, a))$?

q_0

q_1

q_2

a

b

Undefined

What is $\delta((q_0, abb))$?

q_0

q_1

q_2

a

b

Undefined

Save Answer

Q2.2 (c)

1 Point

What is $\delta^*((q_0, a))$?

q_0

q_1

q_2

a

b

Undefined

What is $\delta^*((q_0, abb))$?

q_0

q_1

q_2

a

b

Undefined

Save Answer

Q3 General constructions

3 Points

Fix the alphabet $\Sigma = \{a, b\}$. For each positive integer n , define L_n to be the language over Σ given by $L_n = \{w \in \Sigma^* \mid |w| \text{ is an integer multiple of } n\}$

Q3.1 Example

1 Point

Select all and only the languages below in which $abbaab$ is an element.

☐ L_1

☐ L_2

☐ L_3

☐ L_4

☐ L_5

☐ L_6

☐ L_7

☐ L_8

☐ L_9

☐ L_{10}

☐ L_{11}

☐ L_{12}

Save Answer

Q3.2 Automata and L_n

1 Point

True or false: for each n , there is some finite automaton that recognizes L_n .

True

False

Save Answer

Q3.3 A related language

1 Point

True or false: The set of all strings over $\{a, b\}$ with odd length is regular.

True

False

Save Answer

Q4 DFA construction

2 Points

Consider an arbitrary finite automaton $M = (Q, \{a, b\}, \delta, q_0, F)$ and let's call the language recognized by this finite automaton L .

We can define a new finite automaton which recognizes the collection of strings that result from taking each string in L and replacing each a in the string with 0 and each b in the string with 1. For example, if $L = \{a, aab\}$, then this process would produce the new language $\{0, 001\}$.

Informally: the construction is to keep the states and arrows more or less the same, but change the labels so that the label a on an arrow is replaced by the label 0 and the label b on an arrow is replaced by the label 1.

Fill in the formal definition below;

The new machine is $M' = (Q', \Sigma', \delta', q', F')$ where

$Q' =$

Q

\overline{Q} , aka Q^c , aka the complement of Q

$Q \times Q$

$\Sigma' =$

$\{a, b\}$

$\{0, 1\}$

$\{0, 1, a, b\}$

$\delta' : Q' \times \Sigma' \rightarrow Q'$ is defined by $\delta'((q, 0)) = \delta((q, a))$ and $\delta'((q, 1)) = \delta((q, b))$ for each $q \in Q$.

$q' =$

0

1

a

b

q

q_0

$F' =$

F

\overline{F} , aka F^c , aka the complement of F

$F \times F$

Save Answer

Q5 Closure

2 Points

A set X is said to be **closed** under an operation OP if, for any elements in X , applying OP to them gives an element in X . For example, the set of integers is closed under multiplication because if we take any two integers, their product is also an integer.

For each of the sentences below, (1) first determine if it is a closure claim, and, if it is, then (2) determine if the sentence is true or false.

Concatenating two strings over the alphabet Σ gives a string over the alphabet Σ

Not a closure claim.

Is a closure claim, but false.

Is a closure claim, and true.

The intersection of two infinite sets of integers is an infinite set of integers.

Not a closure claim.

Is a closure claim, but false.

Is a closure claim, and true.

The complement of a regular language is also regular

Not a closure claim.

Is a closure claim, but false.

Is a closure claim, and true.

Save Answer

Q6 Feedback

0 Points

Any feedback or questions about today's material or comments you'd like to share? (Optional; not for credit)

Save Answer

Save All Answers

Submit & View Submission >