Database System Principles

Test One

Class	No	Name	
1. ((10points) Fill in blan	ıks		
•		latabase at a particular moment is cal	led an
<u>instance</u> of the database.		1	
<u> </u>	ovides users with th	ree levels of data abstraction, the _v	iew
level of abstraction describ	es only part of the en	ntire database.	
(3) Database design involv	es the following pha	ses: requirements analysis,	
Conceptual schema de	esign, logical design	and physical design.	
(4) <u>Data model</u> is	a collection of con	nceptual tools for describing data,	, data
relationships, data ser	nantics, and data cor	nstraints.	
(5) As human-machine into	erfaces, the database	language consists of two parts, i.e th	ie data
definition language (I	DDL) and <u>DML (da</u>	ata manipulation language)	
(6) With respect to integrit	y mechanisms in DB	SS, the <u>trigger</u> defines actions to b	e
executed automatically satisfied.	when some events oc	ecur and corresponding conditions are	e
(7) An entity set that does	not have a primary k	ey is referred to as a <u>weak entity se</u>	<u>et</u> .
(8) For a relation r(R), am	ong schema, the attr	ributes the values of which can be u	sed to
uniquely identify the tuple	s in $r(R)$ is called th	e key/superkey/ primarykey/candida	te key
of R.			
(9) The six fundamental o	perations in the relat	tional algebra are select, project, u	ınion ,
set difference,	, and renar	ne.	
(10) Let $r_1(R_1)$ and $r_2(R_2)$) be relations with j	primary keys K_1 and K_2 respectively	ly, the
subset α of R_2 is called the	foreign key	referencing K_1 in relation r_1 , if for ex-	very t ₂
in r_2 there must be a tuple	t_1 in r_1 such that $t_1[K]$	$[X_1] = t_2[\alpha]$	

2. (5 points) Given a table Employees and some SQL queries on it, why are these queries wrong?

Employees(employee-id, employee-name, company-id, employee-city, age, salary)

It is assumed that each employee has an unique id and name.

```
(1) create table Employees

( employee-id char(20),
        employee-name char(20),
        company-id char(20),
        employee-city char(20),
        age integer,
        salary integer,
        primary key (employee-id),
        primary key (employee-name),
        check (age >0)

)
```

(2) select employeeid, sum(salary) from *Employees* group by *company-id* having avg(salary)>1000

不能有2个主键

employeeid 不是分组属性,不能出现在 select 子句中

- 3. (10 points) 给出下列关系代数操作对应的 SQL 语句
 - (1) $\sigma_{p}(\mathbf{r})$
- (2) $\Pi_{A1, A2, ..., Am}(r)$
- $(3) \quad r \infty_S$

, 假设 r(A, B, C), s(C, E, F)

Answers:

- (1) select * from r where P
- (2) select A1, A2, ..., Am from r
- (3) select * from r natual join s 或者:

select * from r where r.C = s.C

- 4. (10 points) 给出下列 SQL 语句对应的关系代数表达式
- (1) select branch-name, max (salary)

from *pt-works*

group by branch-name

假设 pt-works(employee-name, branch-name, salary)

(2) insert into r

select $A_1, A_2, ..., A_m$

from $r_1, r_2, ..., r_n$

where P

(3) update loan

set amount = amount *1.2

where amount > 1000

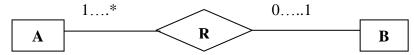
Answers:

- (1) branch_name G max(salary) (pt-works)
- (2) $\mathbf{r} \leftarrow \mathbf{r} \cup \pi_{A1,A2,...Am}(\sigma_{p}(\mathbf{r}_{1} \times \mathbf{r}_{2} \times ... \times \mathbf{r}_{n}))$
- (3) $T1 \leftarrow \Pi_{loan-number, branch_name, amount*1.2} \sigma_{amount > 1000} (loan)$

T2 $\leftarrow \sigma_{amount \leq 1000}$ (loan)

 $loan \leftarrow T1 \cup T2$

- 5. (5 points) For the entity sets A and B and the relationship set R among them in the following figure,
- (1) point out the participation constraints of **A** and **B** in **R**
- (2) what is the mapping cardinality form A to B



Answers:

- (1) A: total B: partial
- (2) one-to-many
- **6.** (5points) Convert the entity set "学生", of which the attribute "老乡" is a multivalued attribute, in Fig.4 into relational **tables**

student-id	籍贯	老乡	性别	年龄
07494	北京	07596,	男	20
		07611		
07498	河北	07320,	女	19
		07321		

Fig.4

Answers:

student-id	籍贯	性别	年龄
07494	北京	男	20
07498	河北	女	19

student-id	<u>老乡</u>
07494	07596
07494	07611

07498	07320
07498	07321

7. (20 points) There are four relations in a *Student* database.

Student(studentID, studentname, sex, age, birthday, schoolID)

Teacher(teacherID, teachername, sex, birthday, schoolID)

Course(courseID, coursename, teacherID)

Score(studentID, courseID, grade)

School(SchoolID, schoolname, dean, location)

Give SQL statements for the following queries:

(1) (5 points) Use a SQL statement to define the relational table *Student*, in which {studentID} is the primary key, and {studentname} is the candidate key and is not permitted to be null; there also exists the referential integrity between the table *Student* and *School*. It is required that the student's age is larger than 10.

create table student

(studentID integer primary key /*也可以放在后面用单独 primary key 语句定义

```
studentname varch(50) /*也可以采用其它长度的 varch、char 类型
sex
           varch(50)
age
           int,
birthday
           date
                        /*或: varch(50),
schoolID
           integer
primary key (studentID) /*也可以在前面 studentID 处定义主键
unique (studentname), not null,
foreign key (schoolID) references School,
check (age>10)
)
主键、候选键、外键、check 定义,每个0.5分。其它1分。
```

(2) (5points) Find the student who takes a course and the grade he gets on this course is higher than the average grade of this course. List the student's name, the name of the course he takes and gets higher grade, and his grade on this course.

```
Select studentname, coursename, grade
From Score as A, Student, Course

Where A.studentID=Student.studentID and A.courseID=Course.courseID

A.grade >

(Select avg(grade)
From Score as B

Where A.courseID=B.courseID
```

关键部分答出,2分。

(3) (5points) Find the courses that are taught by the teachers in School of Computer Science (its SchoolID is SCS) and are taken by at least 5 students. For these courses, list their <u>courseIDs</u>, coursenames, average grades maximal grades and minimal grades, in descending order of the average grades.

(4) (5 points) Find all students who take the course "Database System Principles", and increase their scores on this course by 5 points.

```
Update Score as A
Set grade=grade+5
From Course as B
Where A.CourseID=B.CourseID and B.coursename= Database System Principles
或者:
Update Score as A
Set grade=grade+5
Where CourseID in
```

```
{ Select courseID
From Course
Where coursename= Database System Principles
}
```

8. (5 points) Given *R* (A, B, C, D, E, F, G, H, I), and *F* = {A→F, B→E, BE→F, E→C, A→G, G→CD, I→E} holding on *R*, find out all candidate keys of *R* (利用求候选键算法,给出计算过程)

Answers:

参见讲义附录部分的解法

- 9. (10 points) The relation schema R(U,F),U=(A,B,C,D,E,F,G), The functional dependencies set $F=\{E\rightarrow G,AC\rightarrow D,FG\rightarrow E,AFG\rightarrow B\}$:
 - (1) Compute (ACF)⁺ (2 points)

$$(ACF)^{+}=\{A, C, F, D\}$$
 — $\uparrow 0.5$

(2) Consider the decomposition $\rho = \{R1(A,C,D), R2(A,B,C,E,F,G)\}$, is this decomposition lossless or lossy? Why? (4points)

是 lossless。 $(2 \, \beta)$ 因为 $R1 \cap R2 \rightarrow R1$ $(2 \, \beta)$

- **10.** (20points) The functional dependency set $F = \{A \rightarrow C, C \rightarrow A, B \rightarrow A, D \rightarrow AC, B \rightarrow E\}$ holds on the relation schema R = (A, B, C, D, E),
- (1) Compute (AB)⁺

(3points)

(2) List all the candidate keys of R.

(5points)

- (3) What is the highest normal form of R, and why? (2 points)
- (4) Compute the canonical cover F_c

(5points)

(5) Give a lossless and dependency-preserving decomposition of R into 3NF. (5 points)

Answers:

- $(1) (AB)^{+} = ABCE$
- (2) 唯一候选键 BD
- (3) R 所属最高范式为 1NF (1 point)

因为:

存在非主属性 A 对候选键 BD 的部分依赖 $B \rightarrow A$ (或 $D \rightarrow A$),

或者: 非主属性 E 对键 BD 的部分依赖 $B\rightarrow E$; 非主属性 C 对键 BD 的部分依赖 $D\rightarrow C$

(4) $Fc=\{A\rightarrow C, C\rightarrow A, B\rightarrow AE, D\rightarrow A\}$ 或者: $Fc=\{A\rightarrow C, C\rightarrow A, B\rightarrow AE, D\rightarrow C\}$

答对1个即可,Fc不完全正确,扣1分

(5) 分解 1: 根据 Fc={ A→C, C→A, B→AE, D→A } R1(AC), R2(ABE), R3(AD), R4(BD)

分解 2: 根据 $Fc=\{A\rightarrow C, C\rightarrow A, B\rightarrow AE, D\rightarrow C\}$ R1(AC), R2(ABE), R3(CD), R4(BD)