

## Exercise Solution: P and NP

1. Consider the following problems, give a certificate that helps proving the NP membership.

- (a)  $\text{PARTITION} = \{\langle S \rangle \mid S = \{x_1, x_2, \dots, x_k\} \text{ and there is a partition of } S, T = \{x_{i_1}, \dots, x_{i_t}\} \subseteq S \text{ such that the sum in } T \text{ is equal to the sum in } S \setminus T\}$

We can use a subset of  $S$ , which is one of the two parts, as a certificate.

- (b)  $\text{HAMPATH} = \{\langle G, s, t \rangle \mid G \text{ is a directed graph with a Hamiltonian path from } s \text{ to } t\}$

We can use a permutation of vertices in  $G$ , which in order forms a Hamiltonian path from  $s$  to  $t$ , as a certificate.

2. Show that  $\text{PARTITION}$  is in NP.

First, for any yes-instance (that is, a set of numbers with a equal-sum bipartition) we define a certificate  $C$  as a subset of elements in one of the two parties.

Next, we design a polynomial-time verifier for  $\text{PARTITION}$ .

$V = \text{"On input } \langle \langle S \rangle, C \rangle :$

1. Check if  $C$  is a subset of  $S$ . If it is not true, *reject*.
2. Check if the sum of elements in  $C$  is equal to the sum of elements in  $S \setminus C$ .  
If it is not true, *reject*
3. If both Steps 1 and 2 pass, *accept*."

Finally, we show that  $V$  can verify  $\text{PARTITION}$  in polynomial time. For Step 1,  $V$  scans the input string for at most  $\min\{|C|, |S|\}$  times. For Step 2,  $V$  scans the input string twice. Hence, the total running time of  $V$  is polynomial in the size of input string.

3. Show that  $3\text{SAT}$  is in NP.