Consensus Algorithms in Wireless Blockchain System

1 Consensus Algorithm in Each Round

Algorithm 1 The SWIB Protocol

Input: List of consensus nodes $\{1, 2, \cdots, N\}$ with active time $\{T_1, T_2, \cdots, T_N\}$; Transactions $Txs = \{tx_1, tx_2, \cdots, \}$; Target contention success probability ς; Target transmission success probability ξ; channel compete probability p Transmission parameters $\{P_t, \alpha, \beta\}$

Output: Blockchain BC

//** Achieving consensus on block round-by-round,r round

▷ Initialization:Slots 1:

Get N nodes sorted based on public key value;

Compute elected probability of all nodes according to stability;

Compute required contention time slots x_{comp};

Compute required transmission time slots x_{trans}

▷ Block Proposer Election:Slots 2:

$Rdm^r = GenerateRandomValue(r, B_H^{r-1}, sig_{full}^{r-1})$

ID_{BP} = Block Proposer Election$(NodeList, Rdm^r)$

▷ Block Generation:

for $x_{comp} + T_B \cdot x_{trans}$ Slots **do**

  if $BP_{ID} == Node_{ID}$ **then**

    B^r = Generate Block(B_H^{r-1}, Txs)

    Broadcast the block B^r to other nodes

  else

    Listen on the channel to receive the block

  end if

end for

▷ Block Verification and Finalization

for $N \cdot x_{comp} + T_{sig} \cdot x_{trans}$ Slots **do**

  if $Node_{ID} != BP_{ID}$ **then**

    if $isLegal(BPNode_{ID})$ and $isValid(B^r)$ **then**

      sig^r = Generate Signature(B^r, sk)

    end if

    //**Broadcast signatures

    if v decided to send a transaction based on $\hat{p}_v$ **then**

      $broadcast(MSGs)$ with power P_t

    else

      if channel is idle **then**

        //** Count idle slots within T_v

        $e_v = e_v + 1$

      else

        Receive a message from others

      end if

    end if

    //**maintain the estimate of adversary time window

    $count_v = count_v + 1$

    if $count_v > T_v$ **then**

      $count_v = 1$

      if $e_v == 0$ **then**

        $\hat{p}_v = \hat{p}_v/(1+\gamma)$

        $\hat{T}_v = \hat{T}_v + 2$

      else if $e_v >= 1$ **then**

        $\hat{p}_v = \hat{p}_v * (1+\gamma)$

        $\hat{T}_v = \hat{T}_v - 1$

46: **end if**
47: **end if**
48: **end if**
49: **end for**
50: **while** $!finalized$ **do**
51: //** Broadcast signatures
52: $BroadcastMSG()$
53: $B^r, proof, sig_u^r, sig_{full}^r = RcvMSG()$
54: //**Check the Finalization of new block
55: **if** $isValid(sig_{full}^r)$ **then**
56: $AddSig(B^r, sig_{full}^r)$
57: $Append(BC, B^r)$
58: $finalized = True$
59: **else if** $Count(Sigs^r) \geq \lceil \frac{N+1}{2} \rceil$ **then**
60: $sig_{full}^r =$ Recover Full Signature$(Sigs^r)$
61: $broadcast(sig_{full}^r)$ with probability p_{max} and power P_{max}
62: $AddSig(B^r, sig_{full}^r)$
63: $Append(BC, B^r)$
64: $finalized = True$
65: **else if** $sig_u^r \notin Signs^r$ **then**
66: Append Signature$(Sigs^r, sig_u^r)$
67: **else**
68: //**Check the validation of new block
69: **if** $isValid(B^r, pk_{BP}, proof, Rdm^r)$ **then**
70: $sig_v^r =$ Generate Signature(B_H^r, sk_v)
71: **end if**
72: **end if**
73: **end while**

Algorithm 2 BroadcastMSG subroutine

if v decided to send a transaction based on $\hat{p}_v$ **then**
  $broadcast(MSGs)$ with power P_t
else
  if channel is idle **then**
    //** Count idle slots within T_v
    $e_v = e_v + 1$
  else
    Receive a message from others
  end if
end if
//**maintain the estimate of adversary time window
$count_v = count_v + 1$
if $count_v > T_v$ **then**
  $count_v = 1$
  if $e_v == 0$ **then**
    $\hat{p}_v = \hat{p}_v/(1+\gamma)$
    $\hat{T}_v = \hat{T}_v + 2$
  else if $e_v >= 1$ **then**
    $\hat{p}_v = \hat{p}_v * (1+\gamma)$
    $\hat{T}_v = \hat{T}_v - 1$
  end if
end if

Algorithm 3 The SWIB Blockchain Consensus Protocol for each node v

while true **do**
  //** Iteration for round r
  ▷ Initialization:
  for $j < K$ slots **do**
    $BroadcastMSG()$
    $j = j + 1$
  end for
  $Rds^r = GenerateRandomValue(r, B_H^{r-1}, sig_{full}^{r-1})$
  ▷ Consensus Process:
  Block Proposer Election();
  Block Verification();
  Block Finalization();
  $r = r + 1$
end while

Algorithm 4 Block Verification for each node v

$B^r, proof = RcvMSG()$
//**Check the validation of new block
$result_v$ = Verify Block Proposer($pk_{BP}, proof, Rdm^r$)
if $result_v == True$ **then**
    if $H^r_{pre} == B^{r-1}_H$ **then**
        if isvalid(Txs) **then**
            sig^r_v = Generate Signature(B^r_H, sk_v)
        end if
    end if
end if

Algorithm 5 Block Finalization for each node v

while $!finalized$ **do**
  $BroadcastMSG()$
  $sig_u^r, sig_{full}^r = RcvMSG()$
  //**Check the Finalization of new block
  if $isValid(sig_{full}^r)$ **then**
    $AddSig(B^r, sig_{full}^r)$
    $Append(BC, B^r)$
    $finalized = True$
  else if $Count(Sigs^r) \geq \lceil \frac{N+1}{2} \rceil$ **then**
    $sig_{full}^r =$ Recover Full Signature$(Sigs^r)$
    $broadcast(sig_{full}^r)$ with probability p_{max} and power P_{max}
    $AddSig(B^r, sig_{full}^r)$
    $Append(BC, B^r)$
    $finalized = True$
  else if $sig_u^r \notin Signs^r$ **then**
    Append Signature$(Sigs^r, sig_u^r)$
  end if
end while

Algorithm 6 Stable Wireless Blockchain Protocol

▷ Initialization:
$Sortition(PKs^r, S^r)$
$Rds^r = GenerateRandomness(r, B^{r-1}_{hash}, sig^r_{final})$
▷ Leader Election and Block Proposal:
$result = BlockProposerSelection(sk, Rds^r)$
if $result == True$ **then** ▷ As Block Proposer
  $B^r = GenerateBlock(B^{r-1}, Txs)$
  $sig^r_{partial} = Sign(B^r_{hash})$
  $broadcast(B^r, sig^r_{partial})$ with probability p
else ▷ As Ordinary Nodes
  Waiting to receive new Block
end if
▷ Block Verification and Finalization:
while $!finalized$ **do** ▷ All Consensus Nodes
  $(B^r, Signs^r, sig^r_{full}, Tx) = RcvMSG()$
  //**Check the validation of new block
  if $isValid(B^r)$ and $VerifyBlockProposer(pk_{BP}, Rds^r)$ **then**
    $sig^r_v = GenerateSignature(B^r_{hash}, sk_v)$
  end if
  if $isValid(sig^r_{full})$ **then**
    $\sigma^r_F = sig^r_{full}$
    $broadcast(\sigma^r_F)$ with probability p
    $Append(B^r, \sigma^r_F)$
    $finalized = True$
  else if $Count(Signs^r) >= \lceil \frac{N}{2} \rceil$ **then**
    $\sigma^r_F = RecoverFullSignature(Signs^r)$
    $broadcast(\sigma^r_F)$ with probability p
    $Append(B^r, \sigma^r_F)$
    $finalized = True$
  else if $sig^r_u \notin Signs^r$ **then**
    $Signs^r = AppendSignature(sig^r_u)$
  else if v did not broadcast its partial signature **then**
    $broadcast(sig^r_v)$ with probability p
  else
    $broadcast(Tx)$ with probability p
  end if
  $count = count + 1$
  if $count > T$ **then**
    $count = 1$
    if Received T consecutive transactions in the past T rounds **then**
      $p = p * (1 + \delta)^{-1}$
      $T = T + 2$
    end if
  end if
end while

46: **function** RecNewBlock(m_B, σ_v)
47: **if** $\sigma_v \notin sigShares$ **then**
48: $sigShares = AppendSignature(\sigma_v)$
49: **end if**
50: **if** $Count(sigShares) > K$ **then**
51: $FinalSig = RecoverFinalSig(sigShares)$
52: **else**
53: $FinalSig = null$
54: **end if**
55: **return** $sigShares, FinalSig, B_v^{new}$
56: **end function**
57: **function** AppendSignature(σ_v)
58: **if** $\sigma_v \notin sigShares$ **then**
59: $sigShares \leftarrow sigShares + \sigma_v)$
60: **end if**
61: **return** $sigShares$
62: **end function**