

Consensus Algorithms in Wireless Blockchain System

1 Consensus Algorithm in Each Round

Algorithm 1 The SWIB Protocol

Input: List of consensus nodes $\{1, 2, \dots, N\}$ with active time $\{T_1, T_2, \dots, T_N\}$; Transaction $Txs = \{tx_1, tx_2, \dots\}$; Target contention success probability ς ; Target transmission success probability ξ ; channel compete probability p Transmission parameters $\{P_t, \alpha, \beta\}$

Output: Blockchain BC

```
1: /** Achieving consensus on block round-by-round,  $r$  round
2: ▷ Initialization: Slots 1:
3: Get  $N$  nodes sorted based on public key value;
4: Compute elected probability of all nodes according to stability;
5: Compute required contention time slots  $x_{comp}$ ;
6: Compute required transmission time slots  $x_{trans}$ 
7: ▷ Block Proposer Election: Slots 2:
8:  $Rdm^r = \text{GenerateRandomValue}(r, B_H^{r-1}, sig_{full}^{r-1})$ 
9:  $ID_{BP} = \text{Block Proposer Election}(NodeList, Rdm^r)$ 
10: ▷ Block Generation:
11: if  $BP_{ID} == Node_{ID}$  then
12:    $B^r = \text{Generate Block}(B_H^{r-1}, Txs)$ 
13:   Broadcast the block  $B^r$  to other nodes with transmit power  $P_t$ 
14: else
15:   Listen on the channel to receive the block
16: end if
17: ▷ Block Verification
18: for Any node in  $\{1, \dots, N\}$  do
19:   if  $Node_{ID} \neq BP_{ID}$  then
20:     if  $isLegal(BP_{Node_{ID}})$  and  $isValid(B^r)$  then
21:        $sig^r = \text{Generate Signature}(B^r, sk)$ 
22:       Broadcast  $sig^r$  to other nodes
23:     else
24:       Discard the new block and generate an empty block  $B_{empty}^r$ 
25:        $sig^r = \text{Generate Signature}(B_{empty}^r, sk)$ 
26:       Broadcast  $sig^r$  to other nodes
27:     end if
28:   end if
29:   if received enough partial signatures  $Count(Sigs^r) > threshold$  then
30:      $sig_{full}^r = \text{Recover Full Signature}(Sigs^r)$ 
31:     Broadcast  $sig^r$  to other nodes
32:   end if
33: /** Broadcast signatures
34: Listen on the channel
35: if Node  $v$  decides to send a signature with transmit probability  $p_v$  then
36:   Broadcast( $sig_v^r$ ) with transmit power  $P_t$ 
37: else
38:   if channel is idle then
39:      $e_v = e_v + 1$  /** Count idle slots within  $T_v$ 
40:   else
41:     Receive a message from others
42:   end if
43: end if
44: /** maintain the estimate of adversary time window
45: for Any nodes  $v \in \{1, \dots, N\}$  do
```

```

46:      $count_v = count_v + 1$ 
47:     if  $count_v > T_v$  then
48:          $count_v = 1$ 
49:         if  $e_v == 0$  then /** No idle round in the past  $T_v$  slots
50:              $\hat{p}_v = p_v / (1 + \frac{1}{T_v})$ 
51:              $\hat{T}_v = \hat{T}_v + 2$ 
52:         else if  $e_v >= 1$  then
53:              $\hat{p}_v = \hat{p}_v * (1 + \frac{e_v}{T_v})$ 
54:              $\hat{T}_v = \hat{T}_v - 1$ 
55:         end if
56:     end if
57: end for
58: end for
59: ▷ Block Finalization Slots 2:
60: for Any node  $\{1, \dots, N\}$  do
61:     if  $r$  then received or generated  $sig_{full}^r$ 
62:          $AddSig(B^r, sig_{full}^r)$ 
63:          $Append(BC, B^r)$ 
64:     end if
65: end for

```

Algorithm 1 The SWIB Protocol

Input: List of consensus nodes $\{1, 2, \dots, N\}$ with active time $\{T_1, T_2, \dots, T_N\}$; Transactions Txs ; Target contention success probability ς ; Target transmission success probability ξ

Output: Blockchain BC

```
1: /** Achieving consensus on block round-by-round,  $r$  round
2: ▷ Initialization:
3:  $Finalized = False$ 
4: ▷ Block Proposer Election:
5:  $Rdm^r = GenerateRandomValue(r, B_H^{r-1}, sig_{full}^{r-1})$ 
6: Compute elected probability of all nodes according to stability
7: Get sorted node list  $NodeList$  based on public key value
8:  $BP_{ID} = \text{Block Proposer Election}(NodeList, Rdm^r)$ 
9: ▷ Block Generation:
10: if  $Node_{ID} == BP_{ID}$  then
11:   Listen on the channel to collect transactions
12: else
13:   run MroadcastMessage( $Txs$ )
14: end if
15: if  $Node_{ID} == BP_{ID}$  then
16:    $B^r = \text{Generate Block}(B_H^{r-1}, Txs)$ 
17:   Broadcast the block  $B^r$  to other nodes
18: else
19:   Listen on the channel to receive the block
20: end if
21: ▷ Block Verification
22: if  $Node_{ID} \neq BP_{ID}$  then
23:   if  $isLegal(BP_{Node_{ID}})$  and  $isValid(B^r)$  then
24:      $sig^r = \text{Generate Signature}(B^r, sk)$ 
25:   end if
26: end if
27: run BroadcastMessage( $sig$ )
28: if collect enough partial signatures  $Count(Sigs^r) > threshold$  then
29:    $sig_{full}^r = \text{Recover Full Signature}(Sigs^r)$ 
30:   Broadcast the full signature  $sig_{full}^r$  to other nodes
31: end if
32: ▷ Block Finalization:
33: if Nodes have received or generated the  $sig_{full}^r$  then
34:    $AddSig(B^r, sig_{full}^r)$ 
35:    $Append(BC, B^r)$ 
36:    $Update(NodeList)$ 
37:    $Finalized = True$ 
38: end if
```

Algorithm 3 Synchronization Mechanism

Input: Latest blockchain length $Length$; Neighbor list $Neighbors$

```
1: for node  $v$  in  $Neighbors$  do
2:   if  $Len(BC_v) == Length$  and  $Stability_v > threshold$  then
3:     add the node to Candidates
4:   end if
5: end for
6: for node  $i$  in Candidates do
```

```

7:   Request  $m$  blocks from node  $i$ 
8:   if all blocks are valid then
9:       Add missing block to local blockchain
10:  else
11:      request these blocks from other nodes in Candidates
12:  end if
13: end for

```

Algorithm 2 BroadcastMessage Subroutine

```

1: if Node  $v$  decides to send a message with transmit probability  $p_v$  then
2:    $Broadcast(msg)$  with transmit power  $P_t$ 
3: else
4:   if channel is idle then
5:        $e_v = e_v + 1$  /** Count idle slots within  $T_v$ 
6:   else
7:       Receive a message from other nodes
8:   end if
9: end if
10: /**maintain the estimate of adversary time window
11:  $count_v = count_v + 1$ 
12: if  $count_v > \hat{T}_v$  then
13:    $count_v = 1$ 
14:   if  $e_v == 0$  then /** No idle slot in the past  $\hat{T}_v$  slots
15:        $\hat{p}_v = p_v / (1 + \frac{1}{\hat{T}_v})$ 
16:        $\hat{T}_v = \hat{T}_v + 2$ 
17:   else if  $e_v \geq 1$  then
18:        $\hat{p}_v = \hat{p}_v * (1 + \frac{e_v}{\hat{T}_v})$ 
19:        $\hat{T}_v = \hat{T}_v - 1$ 
20:   end if
21: end if

```

Algorithm 3 The SWIB Blockchain Consensus Protocol for each node v

```

1: while true do
2:   /** Iteration for round  $r$ 
3:    $\triangleright$  Initialization:
4:   for  $j < K$  slots do
5:        $BroadcastMSG()$ 
6:        $j = j + 1$ 
7:   end for
8:    $Rds^r = GenerateRandomValue(r, B_H^{r-1}, sig_{full}^{r-1})$ 
9:    $\triangleright$  Consensus Process:
10:  Block Proposer Election();
11:  Block Verification();
12:  Block Finalization();
13:   $r = r + 1$ 
14: end while

```

Algorithm 4 Block Verification for each node v

```
1:  $B^r, proof = RcvMSG()$ 
2: /**Check the validation of new block
3:  $result_v = \text{Verify Block Proposer}(pk_{BP}, proof, Rdm^r)$ 
4: if  $result_v == True$  then
5:   if  $H_{pre}^r == B_H^{r-1}$  then
6:     if  $isvalid(Txs)$  then
7:        $sig_v^r = \text{Generate Signature}(B_H^r, sk_v)$ 
8:     end if
9:   end if
10: end if
```

Algorithm 5 Block Finalization for each node v

```
1: while !finalized do
2:   BroadcastMSG()
3:    $sig_u^r, sig_{full}^r = RcvMSG()$ 
4:   /**Check the Finalization of new block
5:   if isValid( $sig_{full}^r$ ) then
6:     AddSig( $B^r, sig_{full}^r$ )
7:     Append(BC,  $B^r$ )
8:     finalized = True
9:   else if Count( $Sigs^r$ )  $\geq \lceil \frac{N+1}{2} \rceil$  then
10:     $sig_{full}^r = \text{Recover Full Signature}(Sigs^r)$ 
11:    broadcast( $sig_{full}^r$ ) with probability  $p_{max}$  and power  $P_{max}$ 
12:    AddSig( $B^r, sig_{full}^r$ )
13:    Append(BC,  $B^r$ )
14:    finalized = True
15:   else if  $sig_u^r \notin Sigs^r$  then
16:     Append Signature( $Sigs^r, sig_u^r$ )
17:   end if
18: end while
```

Algorithm 6 Stable Wireless Blockchain Protocol

```
1:  $\triangleright$  Initialization:
2:  $Sortition(PKs^r, S^r)$ 
3:  $Rds^r = GenerateRandomness(r, B_{hash}^{r-1}, sig_{final}^r)$ 
4:  $\triangleright$  Leader Election and Block Proposal:
5:  $result = BlockProposerSelection(sk, Rds^r)$ 
6: if  $result == True$  then  $\triangleright$  As Block Proposer
7:    $B^r = GenerateBlock(B^{r-1}, Tx)$ 
8:    $sig_{partial}^r = Sign(B_{hash}^r)$ 
9:    $broadcast(B^r, sig_{partial}^r)$  with probability  $p$ 
10: else  $\triangleright$  As Ordinary Nodes
11:   Waiting to receive new Block
12: end if
13:  $\triangleright$  Block Verification and Finalization:
14: while  $!finalized$  do  $\triangleright$  All Consensus Nodes
15:    $(B^r, Signs^r, sig_{full}^r, Tx) = RcvMSG()$ 
16:   /**Check the validation of new block
17:   if  $isValid(B^r)$  and  $VerifyBlockProposer(pk_{BP}, Rds^r)$  then
18:      $sig_v^r = GenerateSignature(B_{hash}^r, sk_v)$ 
19:   end if
20:   if  $isValid(sig_{full}^r)$  then
21:      $\sigma_F^r = sig_{full}^r$ 
22:      $broadcast(\sigma_F^r)$  with probability  $p$ 
23:      $Append(B^r, \sigma_F^r)$ 
24:      $finalized = True$ 
25:   else if  $Count(Signs^r) \geq \lceil \frac{N}{2} \rceil$  then
26:      $\sigma_F^r = RecoverFullSignature(Signs^r)$ 
27:      $broadcast(\sigma_F^r)$  with probability  $p$ 
28:      $Append(B^r, \sigma_F^r)$ 
29:      $finalized = True$ 
30:   else if  $sig_u^r \notin Signs^r$  then
31:      $Signs^r = AppendSignature(sig_u^r)$ 
32:   else if  $v$  did not broadcast its partial signature then
33:      $broadcast(sig_v^r)$  with probability  $p$ 
34:   else
35:      $broadcast(Tx)$  with probability  $p$ 
36:   end if
37:    $count = count + 1$ 
38:   if  $count > T$  then
39:      $count = 1$ 
40:     if Received  $T$  consecutive transactions in the past  $T$  rounds then
41:        $p = p * (1 + \delta)^{-1}$ 
42:        $T = T + 2$ 
43:     end if
44:   end if
45: end while
```



```

46: function RECNEWBLOCK( $m_B, \sigma_v$ )
47:   if  $\sigma_v \notin sigShares$  then
48:      $sigShares = AppendSignature(\sigma_v)$ 
49:   end if
50:   if  $Count(sigShares) > K$  then
51:      $FinalSig = RecoverFinalSig(sigShares)$ 
52:   else
53:      $FinalSig = null$ 
54:   end if
55:   return  $sigShares, FinalSig, B_v^{new}$ 
56: end function
57: function APPENDSIGNATURE( $\sigma_v$ )
58:   if  $\sigma_v \notin sigShares$  then
59:      $sigShares \leftarrow sigShares + \sigma_v$ 
60:   end if
61:   return  $sigShares$ 
62: end function

```
