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**Problem 1**

Solve the following recurrences:

1.  $T(n) = 4T(n/2) + n^2$
2.  $T(n) = T(n/2) + 2^n$
3.  $T(n) = 16T(n/4) + n$
4.  $T(n) = 2T(n/2) + n \log n$
5.  $T(n) = \sqrt{2}T(n/2) + \log n$
6.  $T(n) = 3T(n/2) + n$
7.  $T(n) = 3T(n/3) + \sqrt{n}$

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**Problem 2**

Consider the following algorithm.

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**Algorithm 1** REVERSE

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1: function REVERSE( $A$ )
2:   if  $|A| = 1$  then
3:     return  $A$ 
4:   else
5:     Let  $B$  and  $C$  be the first and second half of  $A$ 
6:     return concatenate REVERSE( $C$ ) and REVERSE( $B$ )
7:   end if
8: end function
```

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Let  $T(n)$  be the running time of the algorithm on a instance of size  $n$ . Write down the recurrence relation for  $T(n)$  and solve it by unrolling it.

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**Problem 3**

The product of two  $n \times n$  matrices  $X$  and  $Y$  is a third  $n \times n$  matrix  $Z = XY$ , where the  $(i, j)$  entry of  $Z$  is  $Z_{ij} = \sum_{k=1}^n X_{ik}Y_{kj}$ . Suppose that  $X$  and  $Y$  are divided into four  $n/2 \times n/2$  blocks each:

$$X = \begin{bmatrix} A & B \\ C & D \end{bmatrix} \text{ and } Y = \begin{bmatrix} E & F \\ G & H \end{bmatrix}.$$

Using this block notation we can express the product of  $X$  and  $Y$  as follows

$$XY = \begin{bmatrix} AE + BG & AF + BH \\ CE + DG & CF + DH \end{bmatrix}.$$

In this way, one multiplication of  $n \times n$  matrices can be expressed in terms of 8 multiplications and 4 additions that involve  $n/2 \times n/2$  matrices. Let  $T(n)$  be the time complexity of multiplying two  $n \times n$  matrices using this recursive algorithm.

1. Derive the recurrence for  $T(n)$ . (Assume adding two  $k \times k$  matrices takes  $O(k^2)$  time.)
2. Solve the recurrence by unrolling it.

#### Problem 4

Your friend Alex is very excited because he has discovered a novel algorithm for sorting an array of  $n$  numbers. The algorithm makes three recursive calls on arrays of size  $\frac{2n}{3}$  and spends only  $O(1)$  time per call.

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#### Algorithm 2 NEW-SORT

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1: procedure NEW-SORT( $A$ )
2:   if  $|A| < 3$  then
3:     sort  $A$  directly
4:   else
5:     NEW-SORT( $A[0], \dots, A[\frac{2n}{3}]$ )
6:     NEW-SORT( $A[\frac{n}{3}], \dots, A[n]$ )
7:     NEW-SORT( $A[0], \dots, A[\frac{2n}{3}]$ )
8:   end if
9: end procedure

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Alex thinks his breakthrough sorting algorithm is very fast but has no idea how to analyze its complexity or prove its correctness. Your task is to help Alex:

1. Find the time complexity of NEW-SORT
2. Prove that the algorithm actually sorts the input array

#### Problem 5

Given an array  $A$  holding  $n$  objects, we want to test whether there is a *majority* element; that is, we want to know whether there is an object that appears in more than  $n/2$  positions of  $A$ .

Assume we can test equality of two objects in  $O(1)$  time, but we cannot use a dictionary indexed by the objects. Your task is to design an  $O(n \log n)$  time algorithm for solving the majority problem.

1. Show that if  $x$  is a majority element in the array then  $x$  is a majority element in the first half of the array or the second half of the array
2. Show how to check in  $O(n)$  time if a candidate element  $x$  is indeed a majority element.
3. Put these observation together to design a divide and conquer algorithm whose running time obeys the recurrence  $T(n) = 2T(n/2) + O(n)$
4. Solve the recurrence by unrolling it.

#### Problem 6

Suppose we are given an array  $A$  with  $n$  distinct numbers. We say an index  $i$  is locally optimal if  $A[i] < A[i-1]$  and  $A[i] < A[i+1]$  for  $0 < i < n-1$ , or  $A[i] < A[i+1]$  for if  $i = 0$ , or  $A[i] < A[i-1]$  for  $i = n-1$ .

Design an algorithm for finding a locally optimal index using divide and conquer. Your algorithm should run in  $O(\log n)$  time.