## A TOUR OF AN INSERT

How we built our Postgres database in Rust





https://neon.tech https://conradludgate.com

At Neon, we have just officially launched our Serverless Postgres platform. At the heart is the same familiar C postgres implementation, but much of our special sauce is in our Rust infrastructure surrounding Postgres. Let's take a tour of an INSERT query and see how it all fits together.

```
psql "postgres://conrad:password@my-endpoint.ap-southeast-2.aws.neon.tech/main"
  -c "INSERT INTO rust_meetups (name, date) VALUES ('RustAU Sydney', '2024-04-17')"
```

Let's take this simple query.

# proxy / handshake

First things first, we accept a new TCP connection and perform the postgres handshake. This sets up the TLS connection and also receives any startup params like user and database name.

# proxy / authenticate

```
1 let res = user_info.authenticate(
2    ctx,
3    &mut stream,
4    &config.authentication_config
5 ).await;
6 let user_info = match res {
7    Ok(auth_result) => auth_result,
8    Err(e) => {
9        return stream.throw_error(e).await?;
10    }
11 };
```

# Speaker notes Now, we need to confirm you are who you say you are. We perform the SCRAM authentication protocol to securely verify your password is correct, with included MitM protection

# proxy / wake

```
1 loop {
       match handle_try_wake(api.wake_compute(ctx).await) {
           // do not retry error
           Err(e) => {
               report_error(&e, false);
 6
               return Err(e);
 8
           // retry error
           Ok(Continue(e)) => report_error(&e, true),
10
           // success
           Ok(Break(n)) => return Ok(n),
11
12
13 }
```

Neon is a serverless postgres platform. That means the servers are our problem. If your postgres server is not running, we start it now, making sure to handle any spurious or temporary failure modes.

# proxy / passthrough

```
1 let mut node = connect(ctx, &params, &user_info).await?;
2 let session = cancellation_handler.get_session();
3 prepare_client(&node, &session, &mut stream).await?;
4 proxy_pass(stream, node, session).await?;
```

You are authenticated, your database is running. Connect to the database and switch to the transparent passthrough mode.

# pageserver / GetPage@LSN

Proxy has streamed the insert query to postgres, now postgres needs to load the table schema and any index data for BTree operations. Neon has separated the storage layer out of postgres. This is what allows us to have 100ms cold starts as your database VM does not need to download the potentially several GB base backup to startup. Instead, any non-cached pages will skip the disk and go straight to our custom storage engine, called the Pageserver.

# pageserver / GetPage@LSN / lookup\_cached\_page

```
1 let lookup = self.lookup_cached_page(&key, lsn, ctx).await;
   let cached_page_img = match lookup {
       Some((cached_lsn, cached_img)) => {
           match cached_lsn.cmp(&lsn) {
               // there might be WAL between cached_lsn
 6
               // and lsn, we need to check
               Less => Some((cached_lsn, cached_img)),
 8
               // exact LSN match, return the image
 9
                Equal => return Ok(cached_img),
10
               Greater => unreachable!()
11
12
13
       None => None,
14 };
```

If we're lucky, this page is already cached at this LSN. This is backed by a hashmap and fixed size memory buffer of 128MB.

# pageserver / GetPage@LSN / reconstruct\_value

```
1 let mut reconstruct_state = ValueReconstructState {
2    records: Vec::new(),
3    img: cached_page_img,
4 };
5
6 self.get_reconstruct_data(key, lsn, &mut reconstruct_state, ctx).await?;
7 let page = self.reconstruct_value(key, lsn, reconstruct_state).await?;
```

Either we have a cached page, but not fully up to date, or we have no page at all. either way, we need to reconstruct the new page. Neon's KV page storage is built on top of a custom LSM tree, with 'image' and 'delta' layers. These are either cached in memory, saved to disk, or available in S3. Image layers contain full pages, delta layers contain compressed page diffs

# safekeeper / START\_WAL\_PUSH

```
1 let (msg_tx, msg_rx) = channel(MSG_QUEUE_SIZE);
2 let (reply_tx, reply_rx) = channel(REPLY_QUEUE_SIZE);
3 let handle = Some(WalAcceptor::spawn(tli, msg_rx, reply_tx));
4
5 select! {
6    // Forward all messages to WalAcceptor
7    r = read_network_loop(self.pgb_reader, msg_tx, next_msg).await => r,
8    // Forward all messages to Postgres
9    r = network_write(pgb, reply_rx, pageserver_feedback_rx).await => r,
10 };
11 handle.await;
```

Postgres now has all the data it needs to perform the write. It updates the page in memory, now we need to save it. This is where the 'safekeepers' come in. Postgres streams out WAL pages to many safekeepers, for... safekeeping...

# safekeeper / handle\_append\_request

```
1 self.wal_store.write_wal(msg.h.begin_lsn, &msg.wal_data).await?;
2
3 // flush wal to the disk, if required
4 if require_flush {
5     self.wal_store.flush_wal().await?;
6 }
7
8 // report that the WAL was written
9 let resp = self.append_response();
10 Ok(Some(AcceptorProposerMessage::AppendResponse(resp)))
```

The safekeepers will process all 'append' requests by saving to disk. sending a message back to Postgres that the data was written.

# safekeeper / WalBackupTask

```
1 let mut backup_lsn = Lsn(0);
   loop {
       self.commit_lsn_watch_rx.changed().await;
       let commit_lsn = *self.commit_lsn_watch_rx.borrow();
       backup_lsn = self.timeline.get_wal_backup_lsn().await;
 6
 8
       backup_lsn_range(
           &self.timeline,
           &mut backup_lsn,
10
           commit_lsn,
11
12
           self.wal_seg_size,
           &self.timeline_dir,
13
14
           &self.workspace_dir,
           self narallel inhs
15
```

Once data is flushed to several safekeeper disks, one safekeeper will occasionally back up the WAL to S3 or some other durable object store.

# pageservers / START\_REPLICATION

```
let query = format!("START_REPLICATION PHYSICAL {startpoint}");

let copy_stream = replication_client.copy_both_simple(&query).await?;

let mut physical_stream = pin!(ReplicationStream::new(copy_stream));

let mut waldecoder = WalStreamDecoder::new(startpoint, timeline.pg_version);

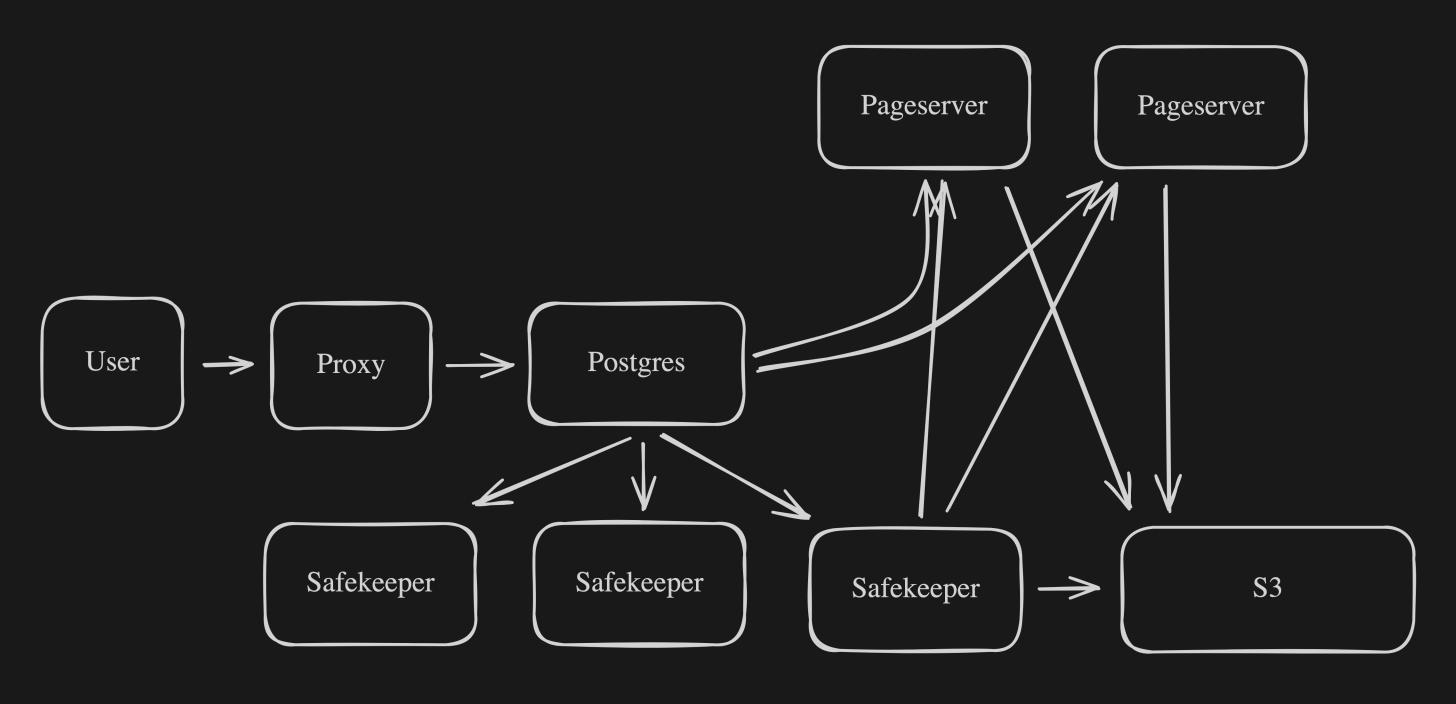
let mut walingest = WalIngest::new(timeline.as_ref(), startpoint, &ctx).await?;
```

At this point, postgres is happy that the data has been written, but the work is not yet done. Our safekeepers have saved the WAL, but pageservers do not know about this data. Our pageservers will open a 'physical replication slot' with one of the safekeepers in order to receive the WAL from safekeepers once it is considered committed.

# pageservers / START\_REPLICATION

```
1 let data = xlog_data.data();
2 let startlsn = Lsn::from(xlog_data.wal_start());
3
4 waldecoder.feed_bytes(data);
5
6 let mut decoded = DecodedWALRecord::default();
7 let mut modification = timeline.begin_modification(startlsn);
8 while let Some((lsn, recdata)) = waldecoder.poll_decode()? {
9     // Ingest the records.
10     walingest.ingest_record(recdata, lsn, &mut modification, &mut decoded, &ctx).await?
11 }
```

We read the WAL data from safekeepers, send them to the separate 'walingest' process to rebuild any page data.



# Speaker notes And 10 minutes later, we have built a database in Rust. Easy peasy. The other 120,000 lines of Rust and however much C and Go code is left as an exercise to the reader.

### THANK YOU

https://github.com/neondatabase/neon https://neon.tech/blog/category/engineering





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