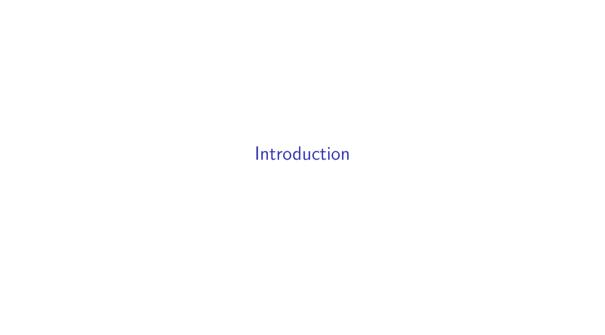
# **Data Compression**

**Edward Higgins** 

2018-07-06

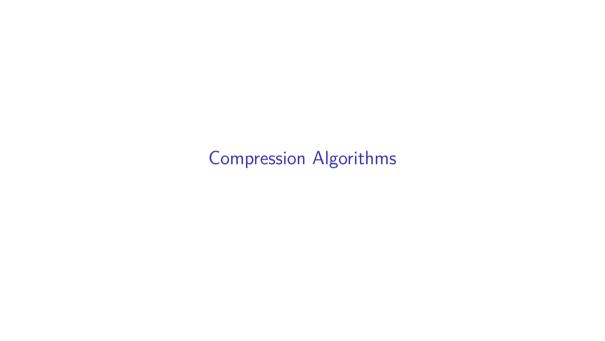


#### What is data compression?

- ▶ Used to reduce the number of bytes used to express the same information
- Can be:
  - ▶ lossy (eg. mp3, jpeg) → approximate representation of information
  - lacktriangle lossless (eg. .zip, .gz, .flac) ightarrow exact representation of information
- Example: aaaaabbcccccc => 5a,3b,6c

# Why do we care?

- Files can be big!
  - Experimental/Simulation data
  - Raw A/V
    - ▶ 24bit 96KHz audio: ~2GB per hour
    - ▶ 24bit 1080p 60Hz video: ~1.3TB per hour
  - Even plain text
    - ▶ Wikipedia, without revisions or multimedia: ~60GB
- ► Finite storage limits how much can be kept
- Finite bandwidth limits how quickly it can be transferred



# Run length encoding

- 'Runs' of data (repeated sequences) are represented as:
  - One copy of the data
  - Count of how many times it is repeated
- ightharpoonup Example: "0.00000000" ightharpoonup "101.80" (one 0, one ., eight 0s)
- Originally designed for bitmap images, where large areas of white compress well
- ▶ Good for certain data types, but can increase file size

# Dictionary encoding (eg. Lempel-Ziv variants, Snappy)

- Repeated strings are stored once, and referenced later
- Example:
  - ▶ "O Romeo, Romeo, wherefore art thou Romeo?"
  - ▶ "O Romeo, \2, wherefore art thou \3?"
- Many different variants:
  - ► How far back do you look for matches?
  - How do you perform the search?
  - ► How do you encode matches?
- Trade-off between compression ratio and compression time

```
"abracadabrarray"
<-- Search buffer ---> <-- Lookahead -->
8 7 6 5 4 3 2 1 | 1 2 3 4 5 6 | Output (off, len, next)
                   |abraca|(0,0,a)
```

```
"abracadabrarray"
<-- Search buffer ---> <-- Lookahead -->
8 7 6 5 4 3 2 1 | 1 2 3 4 5 6 | Output (off, len, next)
                   |abraca|(0,0,a)
                a | b r a c a d | (0,0,b)
```

```
"abracadabrarray"
<-- Search buffer ---> <-- Lookahead -->
8 7 6 5 4 3 2 1 | 1 2 3 4 5 6 | Output (off, len, next)
                   |abraca|(0,0,a)
                a | b r a c a d | (0,0,b)
              a b | r a c a d a | (0,0,r)
```

```
"abracadabrarray"
<-- Search buffer ---> <-- Lookahead -->
8 7 6 5 4 3 2 1 | 1 2 3 4 5 6 | Output (off, len, next)
                   |a b r a c a| (0, 0, a)
                 a | b r a c a d | (0,0,b)
              a b | r a c a d a | (0,0,r)
            a b r | a c a d a b | (3, 1, c)
```

```
"abracadabrarray"
<-- Search buffer ---> <-- Lookahead -->
8 7 6 5 4 3 2 1 | 1 2 3 4 5 6 | Output (off, len, next)
                   |a b r a c a| (0, 0, a)
                 a | b r a c a d | (0,0,b)
              a b | r a c a d a | (0,0,r)
            a b r | a c a d a b | (3, 1, c)
       a b r a c | a d a b r a | (2, 1, d)
```

```
"abracadabrarray"
<-- Search buffer ---> <-- Lookahead -->
8 7 6 5 4 3 2 1 | 1 2 3 4 5 6 | Output (off, len, next)
                  |a b r a c a| (0, 0, a)
                a | b r a c a d | (0,0,b)
             a b | r a c a d a | (0, 0, r)
           a b r | a c a d a b | (3, 1, c)
      a b r a c | a d a b r a | (2, 1, d)
  abracad labrarrl (7, 4, r)
```

\_\_\_\_\_\_

```
"abracadabrarray"
<-- Search buffer ---> <-- Lookahead -->
8 7 6 5 4 3 2 1 | 1 2 3 4 5 6 | Output (off, len, next)
                  |a b r a c a| (0, 0, a)
                a | b r a c a d | (0, 0, b)
             a b | r a c a d a | (0, 0, r)
           a b r | a c a d a b | (3, 1, c)
      a b r a c | a d a b r a | (2, 1, d)
  a b r a c a d | a b r a r r | (7, 4, r)
cadabrar|ray | (3, 2, y)
```

# Symbol reordering

- ▶ Doesn't actually compress the data
- ▶ Improves effectiveness of algorithms (eg. run length encoding, dictionary encoding)
- ightharpoonup Example: "banana" ightharpoonup "bnnaaa"

# Symbol reordering: Burrows-Wheeler Transform

1. Input 2. Make all 3. Sort 4. Take the rotations columns last column ^banana\$ anana\$^b \$^banana ana\$^ban a\$^banan a\$^banan ^banana\$ na\$^bana banana\$^ bnn^aa\$a ana\$^ban nana\$^ba nana\$^ba na\$^bana anana\$^b ^banana\$ banana\$^ \$^banana

# Entropy encoding

- Each unique symbol is given its own variable-length code
- More frequently used symbols are given shorter codes
- Eg. Morse code:
  - e is .
  - ▶ z is --.. (8 times the lengh of e)

"this is an example of a huffman tree" 36 characters = 288 bits											
	•		•	ASCII code	Letter	•		•	ASCII code		
1 1	İ	7	Ċ	00100000	's'	İ	2	Ċ	01110011		
'a'		4		01100001	't'	1	2	1	01110100		
'e'		4	١	01100101	'1'	١	1	1	01101100		
'f'	1	3	١	01100110	'0'	1	1	1	01101111		
'h'	1	2	١	01101000	'p'	1	1	1	01110000		
'i'		2	١	01101001	'r'	١	1	1	01110010		
'm'		2	١	01101101	'u'	١	1	١	01110101		
'n'	1	2	١	01101110	'x'	١	1	١	01111000		

-----

Join the 2 least frequent entries into a subtree...

-----

Repeat for all pairs with frequency of 2...

-----

Join the 2 least frequent entries into a subtree...

-----

```
7 4 4 3 2 2 2 2 2 2 2 4
''a e f h i m n s t * *

/\ __/\__
1 o * *

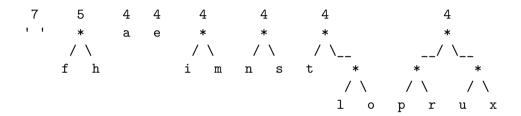
p r u x
```

Repeat for all pairs with frequency of 4...

Re-sort the list...

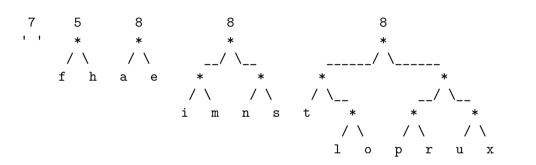
Join the 2 least frequent entries into a subtree...

Re-sort the list...



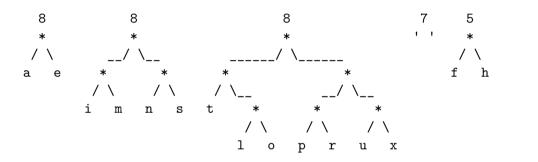
Join the 2 least frequent entries into a subtree, and repeat for all pairs with frequency 8...

-----



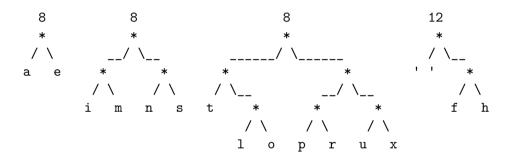
Re-sort the list...

\_\_\_\_\_

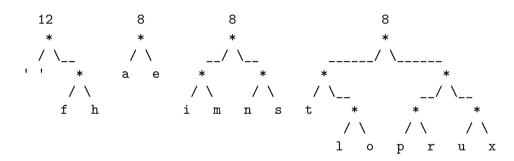


Join the 2 least frequent entries into a subtree...

-----



Re-sort the list



Join the 12-8 and 8-8 subtrees

20 16 Join the final 20-16 subtrees...

```
The Huffman Tree
```

"this is an example of a huffman tree"							36 characters = 145 bits					
	•		•	Huff. code		•		•	Huff. code			
	İ			000	's'	i	2	Ċ	1011			
'a'		4		010	't'		2	1	1100			
'e'	1	4	١	011	'1'		1	1	11010			
'f'		3	١	0010	'0'	١	1	1	11011			
'h'		2	١	0011	'p'	١	1	1	11100			
'i'	1	2	١	1000	'r'	١	1	1	11101			
'm'		2	١	1001	'u'	١	1	1	11110			
'n'		2	I	1010	'x'	١	1		11111			



#### Deflate

- ▶ Used in .zip, zlib (.gzip, .png, ssh, ...), Intel® QuickAssist Technology
- Combination of LZ77 and Huffman coding
- Good compromise between compression ratio and compression speed

### BZip2

- Uses Burrows-Wheeler and Move-To-Front transforms to make data more compressible
- Run length encoding and Huffman encoding then used to compress the data
- ► Compared to Deflate:
  - Higher compression ratio
  - Similar decompression speeds
  - Much slower compression speeds

#### LZMA

- ▶ Used in 7z (windows) and xz (unix) formats, and many package distributions (deb, rpm, . . . )
- ► Uses a modified LZ77 algorithm with range encoding (an entropy encoding algorithm)
- ▶ Higher compression ratios than bzip2, with better decompression times

#### LZ4

- ▶ High speed compression with reasonable compression ratio
- LZ77-esque dictionary encoding with no entropy encoding
- Stores data in 'blocks'

```
|----|----|-----|------|-----|
| t1 | t2 | L1 | literal string | offset | L2 |
|----|----|----|----|
| t1 + L1 = length of literal string
| t2 + L2 = length of match
```

▶ Implemented in many ZFS filesystem implementations

# How do they compare?

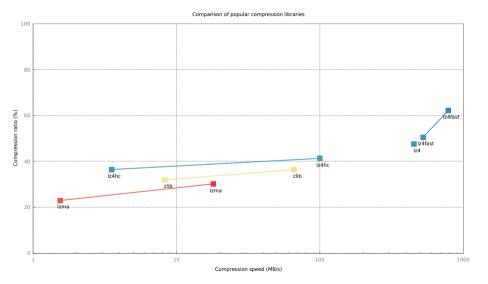


Figure 1: Compression speeds

# How do they compare?

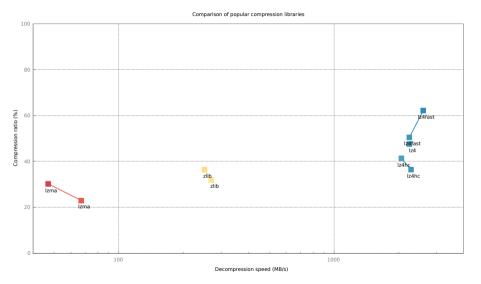


Figure 2: Decompression speeds