

University of Illinois at Urbana-Champaign
Dept. of Electrical and Computer Engineering

ECE 120: Introduction to Computing

Bit-Sliced Comparator

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How Do You Compare Unsigned Numbers?

Let's develop a bit-sliced design to
compare two unsigned numbers.

Which 8-bit unsigned number is bigger?

```
0 1 1 0 1 0 0 0
0 1 0 1 0 1 1 1
```

How did you know?

Did you start on the left or the right?

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Humans Go from Left to Right

Usually, humans start on the left. Why?

As soon as we notice a difference, we're done!

humans compare this way

0 1 1 $a_4 a_3 a_2 a_1 a_0$

0 1 0 $b_4 b_3 b_2 b_1 b_0$

Bit-sliced hardware cannot stop in the middle.

The information flows from one end to the other.

Output wires produce the answer (in bits).

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Our Design Compares from Right to Left

Our comparator design will start on the right.

humans compare this way

$a_7 a_6 a_5 a_4 a_3 a_2 a_1 a_0$

$b_7 b_6 b_5 b_4 b_3 b_2 b_1 b_0$

our design will compare this way

From least significant to most significant bit.

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Three Possible Answers for Comparison of A and B

When comparing two numbers, **A** and **B**, we have **three possible outcomes**:

$$A < B$$

$$A = B$$

$$A > B$$

To decide the answer for **N+1** bits, we need:

- the answer for **N** (less significant) bits,
- one bit of **A**, and
- one bit of **B**.

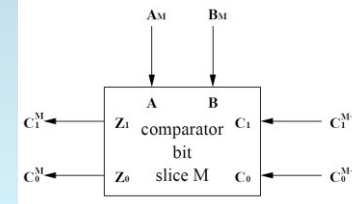
An Abstract Model of the Comparator Bit Slice

A question for you:

How many bits must pass between slices?

Two!

This figure shows an abstract model of our bit slice.



We Need a Representation for Answers

Another question for you:

How do we represent the three possible answers?

Any way we want!

Our choice of representation will affect the amount of logic we need.

Here's a good one...

C_1	C_0	meaning
0	0	A = B
0	1	A < B
1	0	A > B
1	1	not used

A Single Bit Requires Two Minterms on A, B

Let's start by solving a single bit.

In this case, there are no less significant bits.

So we consider only **A** and **B**.

Fill in the meanings, then the bits.

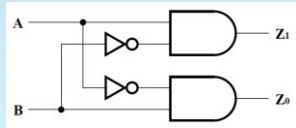
Note that **Z₁** and **Z₀** are minterms.

A	B	Z₁	Z₀	meaning
0	0	0	0	A = B
0	1	0	1	A < B
1	0	1	0	A > B
1	1	0	0	A = B

Comparing Two Bits is Fairly Easy

An implementation for a single bit appears below.

This structure forms the core of our bit slice, since it compares one bit of **A** with one bit of **B**.



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When A and B are Equal, Pass Along the Answer

Now for the full problem.

We'll start with the case of **A = 0** and **B = 0**.

A	B	C ₁	C ₀	meaning	Z ₁	Z ₀	meaning
0	0	0	0	A = B	0	0	A = B
0	0	0	1	A < B	0	1	A < B
0	0	1	0	A > B	1	0	A > B
0	0	1	1	???	x	x	don't care

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When A and B are Equal, Pass Along the Answer

Is there any difference when **A = 1** and **B = 1**?

No, outputs are the same as the last case.

A	B	C ₁	C ₀	meaning	Z ₁	Z ₀	meaning
1	1	0	0	A = B	0	0	A = B
1	1	0	1	A < B	0	1	A < B
1	1	1	0	A > B	1	0	A > B
1	1	1	1	???	x	x	don't care

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When A and B Differ, Override the Previous Answer

What about case of **A = 0** and **B = 1**?

Always output A < B (for valid inputs).

A	B	C ₁	C ₀	meaning	Z ₁	Z ₀	meaning
0	1	0	0	A = B	0	1	A < B
0	1	0	1	A < B	0	1	A < B
0	1	1	0	A > B	0	1	A < B
0	1	1	1	???	x	x	don't care

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When A and B Differ, Override the Previous Answer

And the case of **A = 1** and **B = 0**?

Always output A > B (for valid inputs).

A	B	C ₁	C ₀	Meaning	Z ₁	Z ₀	meaning
1	0	0	0	A = B	1	0	A > B
1	0	0	1	A < B	1	0	A > B
1	0	1	0	A > B	1	0	A > B
1	0	1	1	???	x	x	don't care

Z₁ is a Majority Function

Let's use a K-map to solve **Z₁**.

What are the loops?

AB'

AC₁

B'C₁

S₀

$$Z_1 = AB' + AC_1 + B'C_1$$

Z_1		AB			
		00	01	11	10
C_1C_0	00	0	0	0	1
	01	0	0	0	1
	11	x	x	x	x
	10	1	0	1	1

Z₀ is Also a Majority Function

Now let's use a K-map to solve **Z₀**.

What are the loops?

A'B

A'C₀

BC₀

S₀

$$Z_0 = A'B + A'C_0 + BC_0$$

Z_0		AB			
		00	01	11	10
C_1C_0	00	0	1	0	0
	01	1	1	1	0
	11	x	x	x	x
C_0	10	0	1	0	0

Full Implementation as SOP Expressions

$$Z_1 = AB' + AC_1 + B'C_1$$

$$Z_0 = A'B + A'C_0 + BC_0$$

Notice the symmetry.

And the single-bit core.

