

University of Illinois at Urbana-Champaign  
Dept. of Electrical and Computer Engineering

## ECE 120: Introduction to Computing

### Boolean Expression Terminology

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## Let's Review and Define Some New Terms

**literal** a variable or its complement

examples:  $A$ ,  $A'$ ,  $B$ ,  $B'$ ,  $C$ ,  $C'$

**sum** several terms ORed together

examples:  $A + B$ ,  $AB + B(C + D) + A'C$ ,  
 $A'B' + D(A \oplus B)(C + A')$

**product** several terms ANDed together

examples:  $AB$ ,  $(A + B)(B + CD)(A' + C)$ ,  
 $(A' + B')(D + (A \oplus B) + CA')$

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## Minterms Were Useful for Proving Logical Completeness

### minterm on N inputs

a product in which each variable or  
its complement appears exactly  
once (no other factors)

examples:  $AB'$ ,  $A'B$ ,  $AB$  (on inputs  $A$ ,  $B$ )  
 $AB'C$ ,  $AB'C'$ ,  $A'BC'$   
(on inputs  $A$ ,  $B$ ,  $C$ )

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## A Maxterm Produces a Function with One Zero Row

### maxterm on N inputs

a sum in which each variable or  
its complement appears exactly  
once (no other terms)

examples:  $(A + B')$ ,  $(A' + B)$ ,  $(A + B)$   
(on inputs  $A$ ,  $B$ )  
 $(A + B' + C)$ ,  $(A + B' + C')$ ,  
 $(A + 'B + C')$   
(on inputs  $A$ ,  $B$ ,  $C$ )

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## Sum-of-Products (SOP) Form is Quite Common

### sum-of-products (SOP)

a sum (OR)  
of products (AND)  
of literals

examples:  $AB + BC$ ,  
 $AB' + C + A'C'D'$ ,  
but NOT  $A(B + C) + D$

## Product-of-Sums (POS) Form is Also Common

### product-of-sums (POS)

a product (AND)  
of sums (OR)  
of literals

examples:  $(A + B)(B + C)$ ,  
 $(A + B')C(A' + C' + D')$ ,  
but NOT  $(A + BC)D$

## Canonical Forms Allow Easy Comparison, But Are Too Big

### canonical SOP

a sum of minterms; the expression  
produced by the logical  
completeness construction

### canonical POS

a sum of maxterms

### What does canonical mean?

**Unique** (if we assume an ordering on variables).

**Too many terms to be of practical value.**

## Do You Know Mathematical Implication?

### What does $A \rightarrow B$ mean?

**$A$  implies  $B$ .**

In other words: **if  $A$  is true,  $B$  is also true.**

What if  **$A$  is false?**

In that case, **is  $A \rightarrow B$  true or false?**

**If  $A$  is false,  $A \rightarrow B$  is true.**

## So the Following Odd Statements are True

All **purple elephants** can fly.  
(X is a **purple elephant**  $\rightarrow$  X can fly.)

Students who score **above 125%**  
in ECE120 fail the class.

(X scored **above 125%**  $\rightarrow$  X fails.)

In both, **the premise is false for any X**, so  
the **implications are true**.

## One Function Can Imply Another

A function **G is an implicant of** a second  
function **F iff G** operates on the same  
variables as **F** and  **$G \rightarrow F$** .

In other words, every row

- with an output of 1 in **G's** truth table
- also has an output of 1 in **F's** truth table.

0 rows in **G's** truth table do not matter.

## For Our Purposes, Implicants are Products of Literals

In digital design, we only refer to  
products of literals as implicants.

So we will **assume that an implicant  
can be written as a product of literals**.

## We Can Use Implicants to Simplify Functions

As a first step towards simplifying  
a function **F**, we can ask:

**Given an implicant G of F, can we  
remove any of its literals and obtain  
another implicant of F?**

For example, take  **$F = AB'C + ABC' + ABC$** .

The first term ( **$AB'C$** ) is an implicant.

**Can we remove any literals?**

### Try to Remove Each Literal to Find Only AC Implies F

Start from <b>AB'C</b> and try to remove each literal.	<b>A</b>	<b>B</b>	<b>C</b>	<b>F</b>	<b>B'C</b>	<b>AC</b>	<b>AB'</b>
	0	0	0	0	0	0	0
	0	0	1	0	1	0	0
B'C is not an implicant.	0	1	0	0	0	0	0
	0	1	1	0	0	0	0
AC is an implicant.	1	0	0	0	0	0	1
	1	0	1	1	1	1	1
AB' is not an implicant.	1	1	0	1	0	0	0
	1	1	1	1	0	1	0

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### We Remove as Many Literals as We Can

So we can simplify **F** by replacing **AB'C** with **AC**:

$$F = AC + ABC' + ABC$$

Checking the second term (**ABC'**), we find that we can eliminate **C'** to obtain:

$$F = AC + AB + ABC$$

In the third term (**ABC**), we can eliminate **B** or **C**, but not both. Let's pick **B**.

$$F = AC + AB + AC$$

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### Prime Implicants Have a Minimal Number of Literals

$$F = AC + AB + AC$$

But now we have a duplicate term, which we can eliminate to arrive at a simple form for **F**:

$$F = AC + AB$$

We can remove no more literals.

One more definition: An implicant **G** of **F** is a **prime implicant of F** iff **none of the literals in G can be removed** to produce other implicants of **F**.

**AB and AC are prime implicants of F.**

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### To Simplify, Write Function as a Sum of Prime Implicants

The conclusion is obvious:

**To simplify a function F, write it as a sum of prime implicants.**

Enjoy the algebra.

Good luck!

(Next time, we'll develop a graphical tool that lets us skip the algebra.)

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