

University of Illinois at Urbana-Champaign  
Dept. of Electrical and Computer Engineering

## ECE 120: Introduction to Computing

### Bit-Sliced Designs

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## What's the Theory Behind a Ripple Carry Adder?

Think for a moment about addition.

**Can you add 2-digit numbers?**

**What about 5-digit numbers?**

**What about 5,000-digit numbers?**

**Does it matter if I add more digits?**

**Have you ever seen a proof  
that you're correct?**

**What kind of proof would you need?**

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## Multi-Digit Addition is Correct by Induction

Probably a **proof by induction**...

1. You know how to add 1-digit numbers. Verifying an addition table suffices.
2. **GIVEN that you can add N-digit numbers**, show (based, for example, on place value) that **you can add (N+1)-digit numbers**.

But you didn't know about proof by induction

- when you learned how to add,
- so you've probably never seen a proof.

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## The Ripple Carry Adder is Also Correct by Induction

When we designed a ripple carry adder, we also **assumed proof by induction**.

1. We know how to add one bit. We made a truth table (a binary addition table).
2. **GIVEN that we can build an N-bit adder**, show that we can build an **(N+1)-bit** adder by **attaching a full (1-bit) adder to an (N-bit) adder**.

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## Build an Addition Device Based on Human Addition

In ECE220, you will write **recursive functions**. These functions call themselves.

And you will use the same idea...

1. The answer for some base case (one or more **stopping conditions**) is known.
2. **GIVEN that we can write a function that works for input of size  $N$** , show that we can write a function that works for size  **$(N+1)$**  by **handling the extra “1”** and **calling the function recursively for the “ $N$ ”**.

## The Three Contexts are the Same Mathematically

The approach is the same.

The part that sometimes confuses people (particularly for software/recursion, but sometimes also for hardware/bit slicing) is the ASSUMPTION in the inductive step.

You **must assume that the design works for  $N$  pieces** (bits, input size, or whatever).

## All Three Approaches Require a “Leap of Faith”

You don't need to design the system all at once for  $N$  (other than some base case).

In other words,

- you must make a **“leap of faith”** and
- **assume that your answer works**
- before you actually design it!

People sometimes have trouble making such an assumption, but it's just a **standard part of an inductive proof**.

## Bit Slicing Requires Problem Decomposition

Bit slicing works for problems that

- allow us to **break off a small part** of the problem,
- say **1 bit (or a few bits)**,
- and be able to **solve the full problem using the solution for the remaining part and the 1 bit**.

(That's the inductive step.)

## Signals Between Bit Slices Must be Fixed (and Few)

For hardware, we also need

- to be able to **express the “answer”** for the remaining part
- **in a (small!) fixed number of bits.**

Otherwise, the number of inputs and outputs to the bit slice changes from slice to slice!

## Examples of Problems that Allow Bit Slicing

- Addition / subtraction
- Comparison
- Check for power of 2
- Check for multiples (of 3, 7, and so forth)
- Division by constants
- Pattern matcher
- Bitwise logic operation

## When Can't We Use Bit Slicing?

One example: **when the answer depends on ALL of the other bits** (can't summarize an answer for N bits).

For example, can you create a bit-sliced prime number identifier?

$A_{N-1} A_{N-2} \dots A_5 \leftarrow (\text{summary})$  0 1 0 0 1

What information do you pass to bit 5?

All 5 bits? 01001? I have no idea!