University of Illinois at Urbana-Champaign Dept. of Electrical and Computer Engineering

ECE 120: Introduction to Computing

A Comparator for 2's Complement

ECE 120: Introduction to Computing

© 2016 Steven S. Lumetta. All rights reserved.

slide 3

Comparing 2's Complement Is Different from Unsigned

Let's design a comparator for **2's complement** numbers.

Is the function the same as with **unsigned** (like addition)?

For unsigned, 1001 > 0101.

Is the same true with 2's complement?

No.

Should we just start over?

ECE 120: Introduction to Computing

© 2016 Steven S. Lumetta. All rights reserved.

Start with the Sign Bits

Let's try a little harder first...

If we compare two non-negative numbers,

- $\circ\, the \ approach \ IS \ the \ same.$
- Right?

Maybe we can just use some extra logic to handle the sign bits?

ECE 120: Introduction to Computing

© 2016 Steven S. Lumetta. All rights reserved.

Consider All Possible Combinations of Sign Bits

Let's make a table based on the sign bits:

$A_{\rm s}$	$\mathbf{B_s}$	interpretation	solution
0	0	$A \ge 0 \text{ AND } B \ge 0$	use unsigned
			comparator
0	1	$A \ge 0$ AND $B < 0$	A > B
1	0	$A < 0 \text{ AND } B \ge 0$	A < B
1	1	A < 0 AND B < 0	unknown

ECE 120: Introduction to Computing

© 2016 Steven S. Lumetta. All rights reserved.

slide 4

Interpret 2's Complement as Unsigned

Remember our "simple" rule for translating **2's complement** bit patterns to decimal?

The pattern $A = a_{N-1}a_{N-2} \dots a_1a_0$

has value $V_A = -a_{N-1}2^{N-1} + a_{N-2}2^{N-2} + ... + a_02^0$

Let **A** be negative $(\mathbf{a}_{N-1} = \mathbf{1})$.

Interpreted as unsigned, the same bits have value $V_A + 2^N$.*

*The statement is true by definition of 2's complement, actually.

ECE 120: Introduction to Computing

© 2016 Steven S. Lumetta. All rights reserved

Negative Numbers Can be Compared Directly

What happens if we feed two negative 2's complement numbers into our unsigned comparator?

We compare $V_A + 2^N$ with $V_B + 2^N$.

And we get an answer: <, =, or >.

Let's say that we find $V_A + 2^N < V_B + 2^N$.

In that case, $V_A \le V_B$, so we have the right answer for 2's complement.

The same result holds for other answers.

ECE 120: Introduction to Computing

© 2016 Steven S. Lumetta. All rights reserved

We Need Special Logic for the Sign Bits

Now we can complete our table:

$A_{\rm s}$	$\mathbf{B_s}$	interpretation	solution
0	0	$A \ge 0 \text{ AND } B \ge 0$	use unsigned
			comparator
0	1	$A \ge 0 \text{ AND } B < 0$	A > B
1	0	$A < 0 \text{ AND } B \ge 0$	A < B
1	1	A < 0 AND B < 0	use unsigned
			comparator

ECE 120: Introduction to Computing

© 2016 Steven S. Lumetta. All rights reserved

slide 7

Simply Flip the Wires on the Most Significant Bit

Can we just flip the wires on the sign bits?

For $A_s = 0$ and $B_s = 1$,

• we feed in $A_{N-1} = 1$ and $B_{N-1} = 0$, and

• the unsigned comparator produces A > B.

For $A_s = 1$ and $B_s = 0$,

• we feed in $A_{N-1} = 0$ and $B_{N-1} = 1$, and • the unsigned comparator produces A < B.

What about when $A_s = B_s$?

Flipping the bits then has no effect!

Answers are also correct in those cases.

ECE 120: Introduction to Computing

© 2016 Steven S. Lumetta. All rights reserved.

slide 8

One Comparator with a Control Signal can Do Both

Can we use a single comparator to perform both kinds of comparisons?

Yes, if we

- add a control signal S
- o to tell the comparator whether to do **unsigned** (S=0) or **2's complement** (S=1) comparison.

Simply XOR'ing the most significant bits of A and B with S suffices. This approach leverages flexibility in the problem to reduce the logic needed.

- Analyze the design to understand how it works.

ECE 120: Introduction to Computing

© 2016 Steven S. Lumetta. All rights reserved.