

University of Illinois at Urbana-Champaign
Dept. of Electrical and Computer Engineering

ECE 120: Introduction to Computing

2's Complement Overflow and Boolean Logic

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slide 1

Example: Addition of Unsigned Bit Patterns

A question for you:

What is the overflow condition for addition of two N-bit 2's complement bit patterns?

(That is, when is the sum incorrect?)

Remember that addition works exactly the same way as with **N-bit unsigned** bit patterns, so we can do some base 2 addition to find the answer.

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slide 2

Adding Two Non-Negative Patterns Can Overflow

Let's start with our first example from before:

$$\begin{array}{r}
 11 \\
 01110 \text{ (14)} \\
 + 00100 \text{ (4)} \\
 \hline
 10010 \text{ (-14)}
 \end{array}$$

Oops! We had no carry out, but the answer is wrong (an overflow occurred).

So overflow is different than for **unsigned**...

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slide 3

Carry Out Does Not Indicate 2's Complement Overflow

This example overflowed when the bits were interpreted with an **unsigned** representation.

We have no space for that bit!

$$\begin{array}{r}
 \textcircled{2}11 \\
 01110 \text{ (14)} \\
 + 10101 \text{ (-11; 21 unsigned)} \\
 \hline
 00011 \text{ (3)}
 \end{array}$$

But here the answer is still correct!

Carry out \neq overflow for 2's complement.

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slide 4

Adding Non-Negative to Negative Can Never Overflow

Claim:

Addition of two **N-bit 2's complement** bit patterns can not overflow if one pattern is **negative** (starts with 1) and the other pattern is **non-negative** (starts with 0).

Proof: **You do it!**

And THEN you can read the proof in the notes.

Long Definition for Overflow of 2's Complement Addition

Add two **N-bit 2's complement** patterns.

$$\begin{array}{r} \mathbf{A} \ a_{N-2} \ \dots \ a_0 \text{ (sign bit is A)} \\ + \ \mathbf{B} \ b_{N-2} \ \dots \ b_0 \text{ (sign bit is B)} \\ \hline \mathbf{S} \ s_{N-2} \ \dots \ s_0 \text{ (sign bit is S)} \end{array}$$

Claim: The addition overflows iff one of the following holds:

1. The two addends are non-negative, and the sum is negative.
2. The two addends are negative, and the sum is non-negative.

Boolean Algebra Gives a More Concise Expression

That's a lot of words!

Boolean algebra gives a more concise form:

$$\begin{aligned} \mathbf{OVERFLOW} = & \\ & [(\mathbf{NOT} \ \mathbf{A}) \ \mathbf{AND} \ (\mathbf{NOT} \ \mathbf{B}) \ \mathbf{AND} \ \mathbf{C}] \ \mathbf{OR} \\ & [\mathbf{A} \ \mathbf{AND} \ \mathbf{B} \ \mathbf{AND} \ (\mathbf{NOT} \ \mathbf{C})] \end{aligned}$$

(Remember: **A**, **B**, and **C** were the sign bits.)

But what do these operators (AND, OR, and NOT) mean?

Boolean Operators Were Invented in the mid-19th Century

Boolean operators were invented (by George Boole) to reason about logical propositions.

They originally operated on true/false values.

We use them with ... that's right, bits!

$$\mathbf{0} = \mathbf{false} \text{ and } \mathbf{1} = \mathbf{true}$$

Be careful not to confuse Boolean operators with English words. **The meanings are not identical.**

We Use Only a Few Boolean Functions

AND: the ALL function
returns 1 iff **ALL inputs are 1** (otherwise 0)

OR: the ANY function
returns 1 iff **ANY input is 1** (otherwise 0)

NOT: logical complement
(NOT 0) is 1; (NOT 1) is 0

XOR: the ODD function
returns 1 iff **an ODD number of inputs are 1** (otherwise 0)

A Truth Table Fully Defines a Boolean Function

The drawing to the right is
a **truth table**.

A truth table allows us to

- define a Boolean function **C**
- by listing the output value
- for all combinations of inputs
(here **A** and **B**, in base 2 order).

A	B	C
0	0	
0	1	
1	0	
1	1	

Let's write truth tables for our
four Boolean functions.

AND: The ALL Function

Let's start with AND.

AND can be written in
several ways:

- AB | We usually use these.
- $A \cdot B$
- $A \times B$
- $A \wedge B$ (math. conjunction)

A	B	A AND B
0	0	0
0	1	0
1	0	0
1	1	1

Note flat input,
rounded output.



OR: The ANY Function

And now OR.

OR can also be written
in other ways:

- $A + B$ | We usually use this one.
- $A \vee B$ (math. disjunction)

A	B	A OR B
0	0	0
0	1	1
1	0	1
1	1	1

Note rounded input,
pointed output.



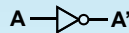
NOT: Logical Complement

And now NOT.

NOT can also be written in other ways:

- \overline{A} | We usually use these.
- $\neg A$ (math. complement)

Note triangle and inversion bubble.



A	NOT A
0	1
1	0

XOR: The ODD Function

And, finally, XOR.

XOR is usually written this way: $A \oplus B$

A	B	A XOR B
0	0	0
0	1	1
1	0	1
1	1	0

Note: like OR, but double line for inputs.



Use Definitions to Generalize to More than Two Operands

Generalize to more operands using the definitions given:

- **AND: ALL**
- **OR: ANY**
- **XOR: ODD**

As an example, fill the truth table for a 3-input XOR.

A	B	C	$A \oplus B \oplus C$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

Generalize to Sets of Bits by Pairing Bits

We can also generalize to sets of bits.

For example, if we have two N -bit patterns,

$$A = a_{N-1} \dots a_0 \text{ and } B = b_{N-1} \dots b_0,$$

we can write

$$C = A \text{ AND } B$$

To mean that

$$\text{if } C = c_{N-1} \dots c_0, c_i = a_i b_i \text{ for } 0 \leq i < N.$$

Don't Mix Algebras: Use AND/OR/NOT for Bitwise Logic

If **A** is a **2's complement** bit pattern, we might also write **$-A = (\text{NOT } A) + 1$**

Be careful about mixing

- algebraic notation for Boolean functions
- with arithmetic operations.

The “+” in the equation above means base 2 addition (and discarding any carry out), not OR.