

CS 152: *Programming Language Paradigms*



Modules, Structs, Hashes, and Operational Semantics

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Modules



Review Modules from HW 1

(in-class)

How do we organize code in Java?

- **Packages** provide units of code
- **Keywords** specify method access

In Java, keywords specify access

- **public**
- **protected**
- no keyword (package access)
- **private**

Organizing Code in Racket

- *Exports* specify public values:
(provide big-add)
- *Imports* specify code dependencies:
(require "big-num.rkt")

Structs and Hashes



Structs

- *Structures* allow us to create more sophisticated data structures:
`(struct name (field1 field2 ...))`
- Once we have a structure, we can *destructure* it with the **match** keyword to get at the contents.
- <Example in class>

Hashes

- Hashes are maps of key/value pairs.
- *Unlike* Java, hashes are immutable.
- <example in class>

Formal Semantics

Why do we need
formal semantics?


Everyone knows what an if statement does, right?

```
if true then
```

```
    x = 1
```

```
else
```

```
    x = 0
```



At the end of this code snippet, the value of x will be 1

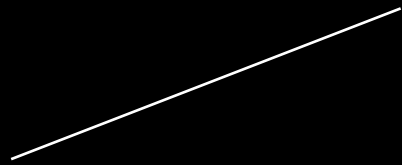
Everyone knows what an if statement does, right?

```
if false then
```

```
    x = 1
```

```
else
```

```
    x = 0
```



At the end of this code
snippet, the value of x
will be 0

Everyone knows what an if statement does, right?

```
if 0 then
```

```
    x = 1
```

```
else
```

```
    x = 0
```

Will x be set to 0,
like in C/C++?

Will x be set to 1,
like in Ruby?

Or will it be an
error, like in Java?

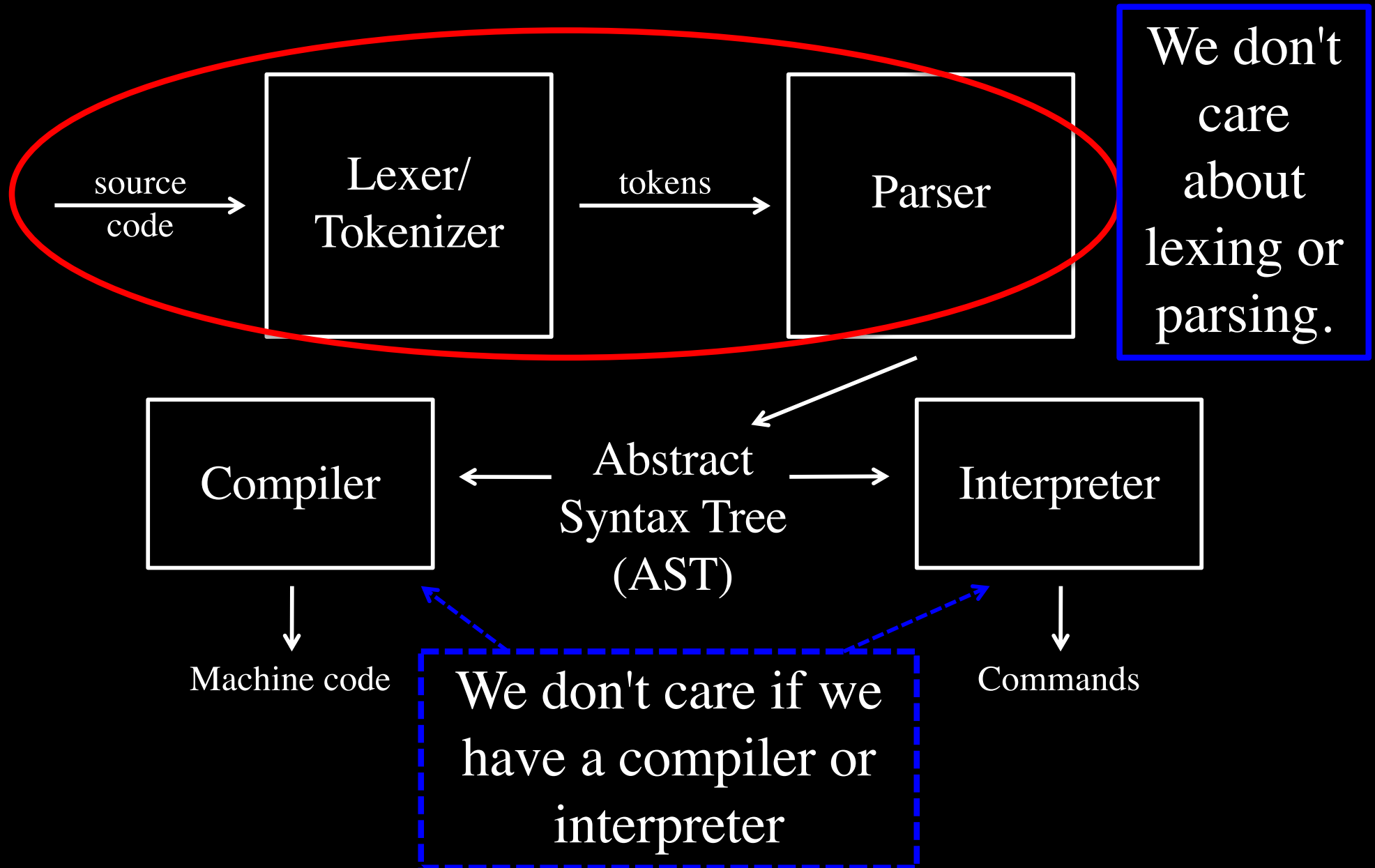
Everyone knows what an
if statement does, right?

```
x = if true  
    then 1  
    else 0
```

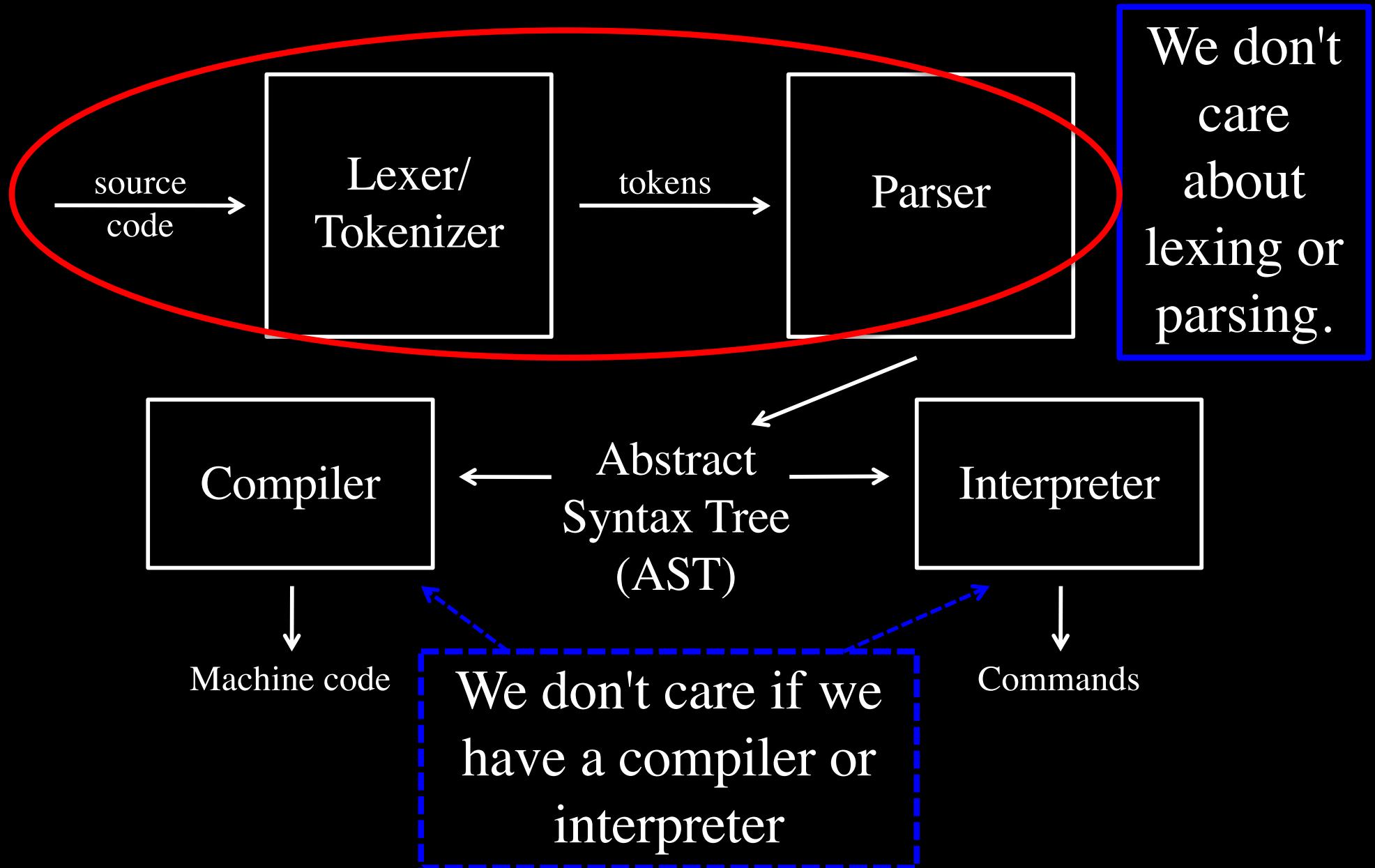
Is assignment
valid or an error?

Formal semantics define how
a language works *concisely*
and with minimal ambiguity.

A Review of Compilers



A Review of Compilers



Abstract Syntax Tree (AST)

ASTs are the
key to
understanding
a language

Bool* Language

$e ::=$

true

| false

| if e

then e

else e

expressions:

constant true

constant false

conditional

Despite appearances,
these are really ASTs

Values in Bool*

$v ::=$	<i>values:</i>
true	constant true
false	constant false

Formal Semantic Styles

- Operational semantics
 - Big-step (or "natural")
 - Small-step (or "structural")
- Axiomatic semantics
- Denotational semantics

Formal Semantic Styles

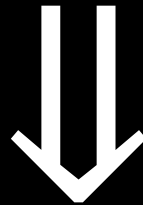
- **Operational semantics**
 - **Big-step (or "natural")**
 - **Small-step (or "structural")**
- **Axiomatic semantics**
- **Denotational semantics**

Big-Step Evaluation Relation

An expression e ...

... evaluates to ...

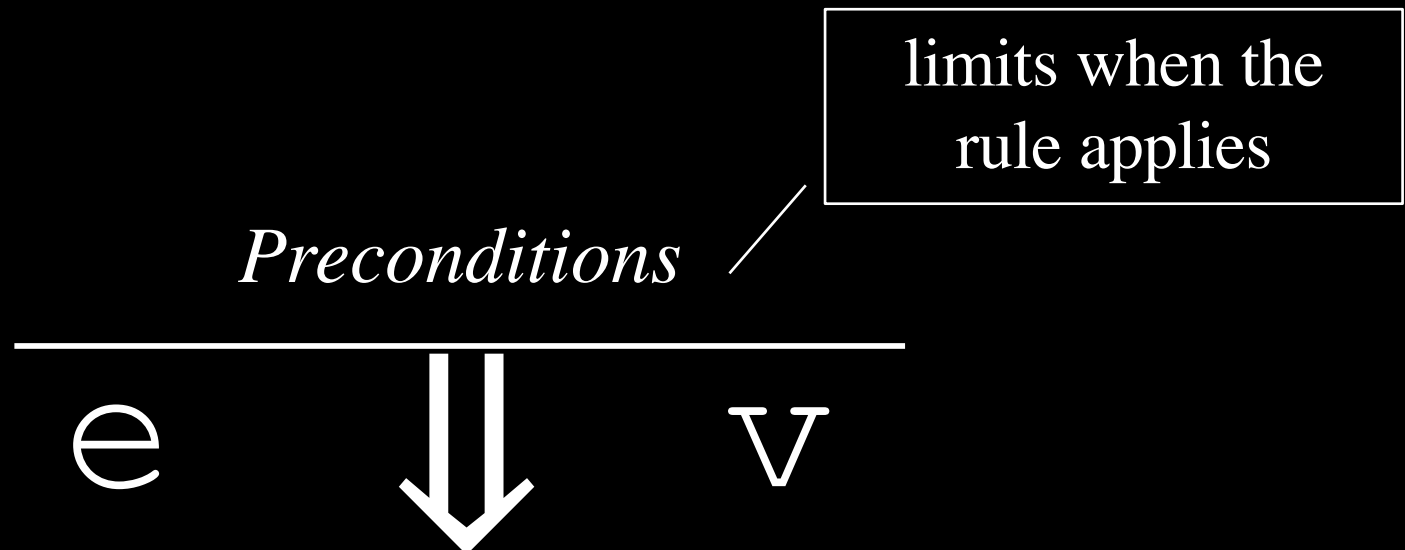
e



v

... a value v .

Big-Step Evaluation Relation



Big-step semantics for Bool*

B-IfTrue

$$\frac{e1 \Downarrow \text{true} \quad e2 \Downarrow v}{\text{if } e1 \text{ then } e2 \text{ else } e3 \Downarrow v}$$

B-IfFalse

$$\frac{e1 \Downarrow \text{false} \quad e3 \Downarrow v}{\text{if } e1 \text{ then } e2 \text{ else } e3 \Downarrow v}$$

B-Value

$$\frac{}{v \Downarrow v}$$

Bool* big-step example

true \Downarrow true false \Downarrow false

if true

 then false \Downarrow false

 else true

 false \Downarrow false

if (if true then false
 else true)

 then true

 else false

\Downarrow false

Converting our rules into code

(in-class)

Bool* extension: numbers

Users demand a new feature – numbers!

We will add 3 new features:

- Numbers, represented by `n`
- `succ`, which takes a number and returns the next highest number.
- `pred`, which takes a number and returns the next lowest number.

Extended Bool* Language

$e ::=$ true
| false
| if e then e else e
| n
| succ e
| pred e

Let's extend our semantics to handle these new language constructs

Lab: Write a Bool* Interpreter

- Starter code is available on the course website
- Extend Bool* with numbers, succ, and pred

Adding State to Semantics

SpartanLang

$e ::= !x$	dereferencing
v	values
$x := e$	assignment
$e ; e$	sequence
$e \text{ op } e$	binary operations
$\text{if } e \text{ then } e$ $\text{else } e \text{ end}$	conditionals
$\text{while } e \text{ do } e \text{ end}$	while loops

SpartanLang (continued)

$v ::= i$
 | b

integers

booleans

$op ::= + \quad | \quad -$
 | $\backslash \quad | \quad *$
 | $< \quad | \quad >$
 | $<= \quad | \quad >=$

binary operators

Bool* vs. SpartanLang evaluation

Bool* relation:

$$e \Downarrow v$$

SpartanLang relation:

$$e, \sigma \Downarrow v, \sigma'$$

A "store", represented by
the Greek letter sigma

The Store

- A mapping of references to values
- Roughly analogous to the heap in Java

Key store operations

- $\sigma(x)$
 - get the value for reference x .
- $\sigma[x := v]$
 - create a copy of store σ , except ...
 - reference x has value v .

HW 2: Write an Interpreter for SpartanLang.

Details in Canvas