Compute services

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Goals

- 1. Bare Metals
- 2. Virtualization
- 3. Containers
- 4. Unikernels/MicroVM

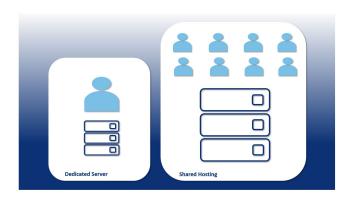
- An infrastructure is proposed to the User
 - Number of CPUs
 - Amount of Memory
 - Network Bandwidth
 - Storage capacity



- Several way to provide an infrastructure
 - BareMetal
 - Fully dedicated server
 - Virtualization
 - Abstraction of a hardware and software isolation
 - Containers
 - Lightweight virtualization, less isolated
 - Unikernel
 - The future, Strong isolation with good performance

Bare Metal (Bare Metal As A Service)

- Fully dedicated servers allocated to the clients
- The client has access fully to a server
- Can act on the hardware, e.g access hardware features
- Useful for sensitive workload with lot of computation
- But, very expensive
- Provided by AWS, GCP etc.

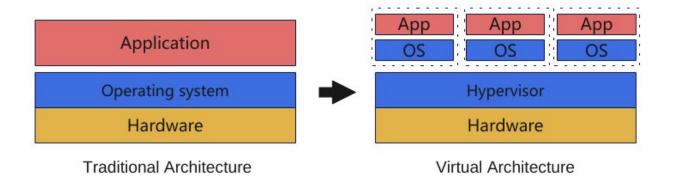


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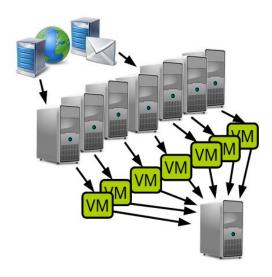
Virtualization

 Set of techniques, hardware and software, which allow managing simultaneously on a single machine several operating systems (called virtual machines (VMs)). Examples: Xen, VMWare, KVM, HyperV, etc;



Consolidation

- Creating X virtual machines from X physical ones and running them on Y physical hosts
 - With Y < X</p>
 - Historical motivation for developing virtualization technologies
 - Gives (most of) the benefits of multi-computer systems without the \$/management costs:
 - Software dependencies
 - Reliability
 - Security



- Use Cases: Software Development
 - Flexible OS diversity: different OS on the same machine
 - e.g VirtualBox with linux for kernel development
 - Rapid and cost-efficient provisioning
 - Way faster that Physical machine, we will see during lab works
 - VMs are self-contained
 - Practical way to "pack" an application with all its software dependencies
 - Model and version of the OS, libraries, etc.
 - Useful for development, automated testing, and even deployment









- Migration and Checkpoint/Restart
 - VMs are self-contained and can be migrated between hosts
 - **Live migration**: moving VMs from one host to another
 - Freeing resources for maintenance or when a fault is expected
 - Increased performance
 - Distributed resources scheduling: load balancing, consolidation for power savings, etc.
 - Checkpoint/restart for long-running jobs
 - Dump the VM state to disk in order to be able to resume it later
 - Both techniques are straightforward for VM as opposed to processes

- Use Cases: Cloud computing
 - Virtualization central technology in cloud computing
 - Allows the providers to securely share their computing infrastructure between clients (tenants)
 - Offloading local tasks to remote computing resources
 - Rent a VM to put a web server (laaS)
 - Offload mail server online to Gmail/Outlook (SaaS)
 - To save on management, infrastructure, development maintenance costs

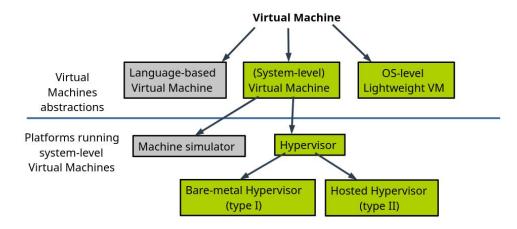


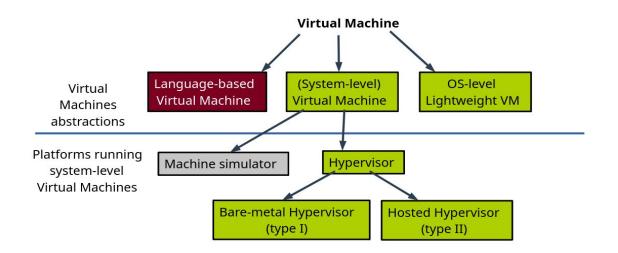




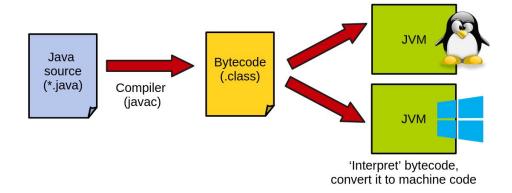
- Security
 - Virtualization provides very strong isolation between guests
 - Sandboxing
 - Cloud, virus/malware analysis, honeypots, process/task level isolation through virtualization
 - VM introspection
 - Analysis of the guest behavior from a privileged level higher than the OS's Guest
 OS cannot be trusted
 - E.g. Cloudvisor-D

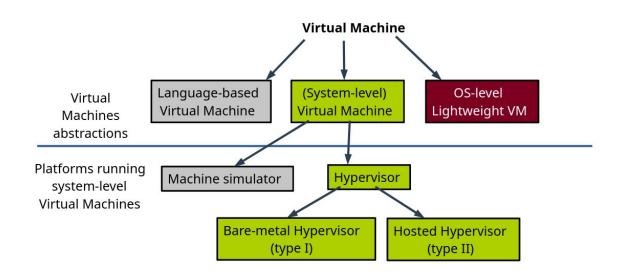
 A complete compute environment with its own isolated processing capabilities, memory and communication channels.



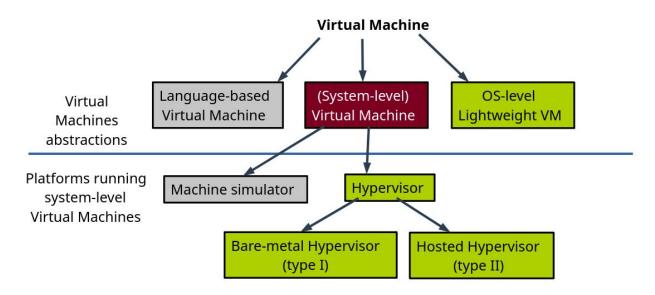


- Java Virtual Machine, JavaScript engines, Microsoft Common Language Runtime
 - Designed to run single applications

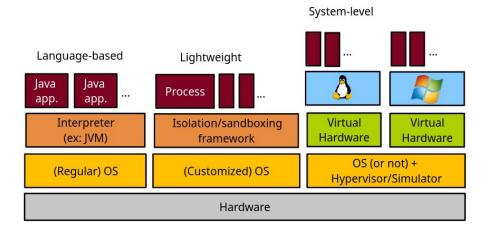




- OS-level Lightweight VMs (discussed in detail later)
 - Sandboxing: Isolation of native code running directly on the CPU though hardware/software mechanisms. Through OS mechanisms offering more isolation guarantees than regular process-based isolation
 - No attempt is mode to virtualize the hardware: Isolation enforced by the OS/framework level (breaks backward compatibility in some cases)
 - Cannot run unmodified OS
 - E.g. Containers, Google Native Client, LibOSes



- System-level Virtual Machines
 - Creates a model of the hardware for a (mostly) unmodified operating system to run on top of it
 - Each VM running on the computer has its own copy of the virtualized hardware



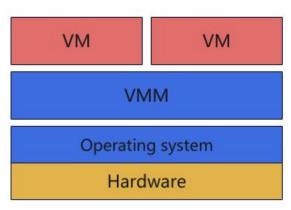
- System-level Virtual Machines
 - Full virtualization
 - Para-virtualization
 - Hardware assisted virtualization

Full virtualization:

An OS executes applications at user level. The VMM is one such application.
 Every instruction from VM is emulated by the VMM. The OS executed on a VM is not modified and can be of any type (Linux, Windows, etc.). Ex: virtualBox

Concretely:

- It's hardware simulation (slow)
- Supports several OS types



Para-virtualization (PV):

The VMM replaces the host OS and behaves as a proxy for accessing the hardware. The host OS is considered as a VM (privileged) and it is used by the VMM to perform particular tasks. Constraint: VMs' OSes have to be modified (to invoke the VMM for privileged operations). Ex: Xen, VMware, etc.

VM Priv. VM

VMM

Hardware

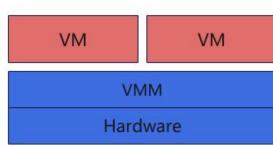
Concretely:

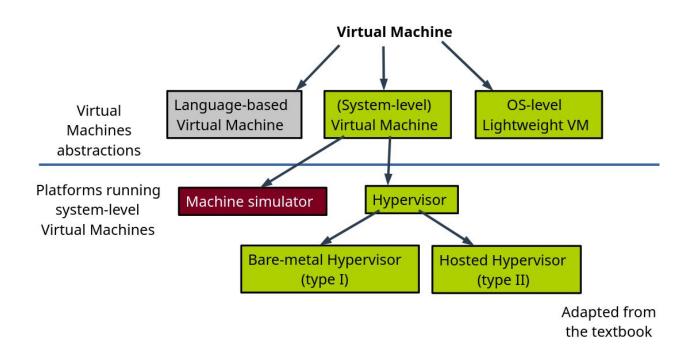
- A mix between simulation and native. Several OS types
- Each VM is allocated hardware resources and the VMM controls access to hardware
- Hypercalls (a modified OS invokes the VMM)

- Hardware assisted virtualization (HVM):
 - The hardware is aware of virtualization. VMs' OSes don't have to be modified. Ex: Xen, KVM,
 VMWare etc.



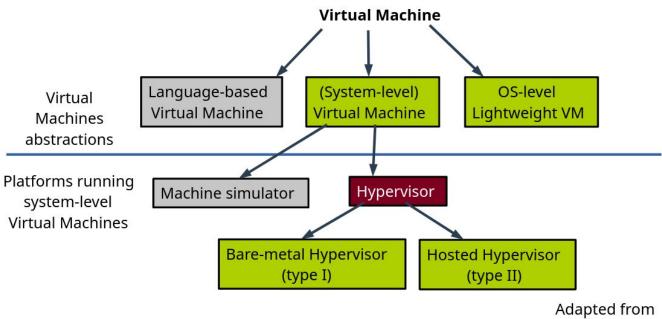
- Native, but without OS modification
- The hardware is able to manage several VMs (MMU, network device ...)
- Sometimes, para-virtualization is faster or more flexible.





Virtualization: Machine simulation

- Full virtualization
- Cross ISA emulation: for usage as substitute
 - Compatible with legacy applications
 - e.g: Qemu
- Architecture emulation
 - Help simulate a hardware architecture for study purpose
 - e.g: Android systems
- Each guest instruction is interpreted in software



Adapted from the textbook

- Hypervisor or Virtual machine monitor (VMM)
 - Relies on direct execution for performance reasons (close to native)
 - VM code executes directly on the physical CPU, at a lower privilege level than the hypervisor
 - Privileged instructions trap to the hypervisor
 - Switch to the hypervisor which determines what to do with that instruction: trap-and-emulate model
 - Then back to the VM execution
 - E.g Xen, Linux KVM, VMware ESXi, MS Hyper-V etc.

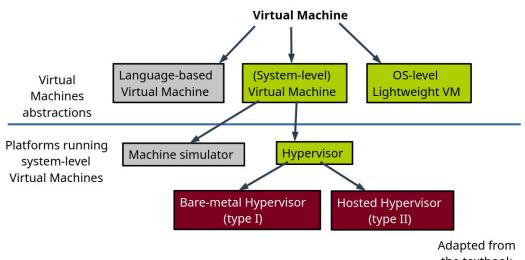






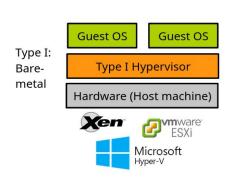
- Hypervisor or Virtual machine monitor (VMM)
 - Multiplexes physical resources between VMs-Run VMs while minimizing virtualization overheads
 - Tries to get as close as possible to native performance
 - Ensure isolation between VMs and between VMs and the hypervisor
 - Isolates physical resources: for example memory/address spaces
 - Isolate performance

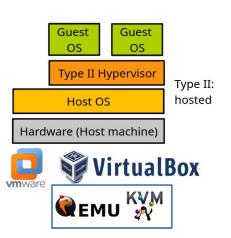
Hypervisor or Virtual machine monitor (VMM)



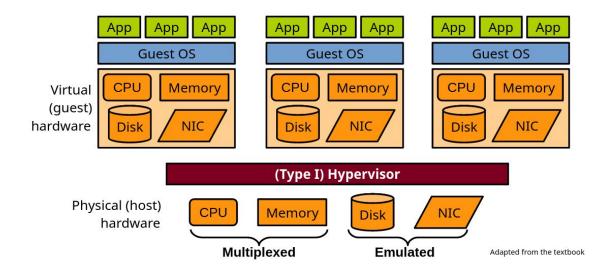
the textbook

- Hypervisor or Virtual machine monitor (VMM)
 - Resources allocation and scheduling
 - Type I: done by the hypervisor
 - Type II: more involvement from the Host OS



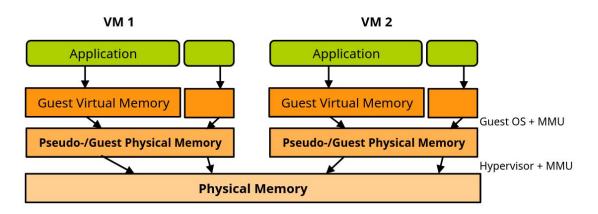


- Hypervisor or Virtual machine monitor (VMM)
 - Simplified Hypervisor Sketch



- Hypervisor or Virtual machine monitor (VMM)
 - Multiplexing CPU and Memory
 - Virtual CPU runs with reduced privileges, cannot execute privileged instructions (that could allow to escape isolation)
 - Traps to the hypervisor on such instructions (virtualization overhead)
 - Run the hypervisor in kernel mode and the VM in user mode.

- Hypervisor or Virtual machine monitor (VMM)
 - Multiplexing CPU and Memory
 - Another level of translation added, taken care of by the hypervisor: (Guest) virtual memory -> (Guest) pseudo-physical memory -> (host) physical memory.



- Hypervisor or Virtual machine monitor (VMM)
 - Emulating I/Os
 - I/O mostly emulated for compatibility reasons
 - I/O devices have well defined interfaces, for example: send a set of network packets, read 128K from disk from secto X, etc.
 - The hypervisor emulate simple virtual device (disk/NIC) that can be accessed with commonly implemented drivers (ex IDE/SCSI): **front-ent**
 - Hypervisor redirects I/O to actual devices or other abstraction, ex: disk or file: **back-end**.

