# Notebook UNTreeCiclo

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9 S 9.		49 49	1 C++	
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9. 9.			1.1 C++ plantilla	
9.			<pre>#include <bits stdc++.h=""></bits></pre>	
9.	T		<pre>using namespace std;</pre>	
9.			<pre>#define all(v) v.begin(), v.end()</pre>	
9.			<pre>#define sz(arr) ((int) arr.size()) typedef vector<int> vi;</int></pre>	
9.			typedef long long 11;	
9.	9 Suffix Tree	52	<pre>typedef pair<int, int=""> ii; const char ln = '\n';</int,></pre>	
9.	10 Trie	53	<pre>const char ln = '\n';</pre>	
9.	11 Z Algorithm	54	<b>#define</b> watch(x) cout<< <b>#</b> x<< <b>"="</b> < <x<<b>'\n'</x<<b>	
	- -		typedef long double ld;	
10 N		54	<pre>typedef vector<ii> vii;</ii></pre>	
	0.1 Counting Sort		<pre>typedef vector<long long=""> v1; typedef pair<ll, l1=""> pl1;</ll,></long></pre>	
10	0.2 Dates	54	typeder pair<11, 11> pi1; typedef vector <pll> vll;</pll>	
	0.3 Expression Parsing		const int INF = 1e9;	

```
const 11 INFL = 1e18;
const int MOD = 1e9+7;
const double EPS = 1e-9;
const ld PI = acosl(-1);
int dirx[4] = {0,-1,1,0};
int diry[4] = {-1,0,0,1};
int dr[] = {1, 1, 0, -1, -1, -1, 0, 1};
int dc[] = {0, 1, 1, 1, 0, -1, -1, -1};
const string ABC = "abcdefghijklmnopqrstuvwxyz";
int main() {
    ios::sync_with_stdio(false);
    cin.tie(0);
    cout << setprecision(20) << fixed;
    // freopen("file.in", "r", stdin);
    // freopen("file.out", "w", stdout);
    return 0;
}</pre>
```

#### 1.2 Librerias

```
// En caso de que no sirva #include <bits/stdc++.h>
#include <algorithm>
#include <iostream>
#include <iterator>
#include <sstream>
#include <fstream>
#include <cassert>
#include <climit.s>
#include <cstdlib>
#include <cstring>
#include <string>
#include <cstdio>
#include <vector>
#include <cmath>
#include <queue>
#include <deque>
#include <stack>
#include <list>
#include <map>
#include <set>
#include <bitset>
#include <iomanip>
#include <unordered map>
////
#include <tuple>
#include <random>
#include <chrono>
```

#### 1.3 Bitmask

```
* Operaciones a nivel de bits. Si n es 11 usar 111<< en
   los corrimientos.
                -> Verifica si x es impar
x & (1<<i)
                -> Verifica si el i-esimo bit esta
   encendido
x = x \mid (1 << i) -> Enciende el i-esimo bit
x = x & (1 << i) -> Apaga el i-esimo bit
x = x^{(1 << i)} -> Invierte el i-esimo bit
                -> Invierte todos los bits
x = x
                -> Devuelve el bit encendido mas a la
x & -x
   derecha (potencia de 2, no el indice)
               -> Devuelve el bit apagado mas a la
~x & (x+1)
   derecha (potencia de 2, no el indice)
x = x \mid (x+1) -> Enciende el bit apagado mas a la
   derecha
x = x & (x-1)
               -> Apaga el bit encendido mas a la
   derecha
x = x & ~v
                -> Apaga en x los bits encendidos de y
* Funciones del compilador qcc. Si n es ll agregar el
   sufijo ll, por ej: __builtin_clzll(n).
__builtin_clz(x)
                      -> Cantidad de bits apagados por la
    izquierda
builtin ctz(x)
                      -> Cantidad de bits apagados por la
    derecha. Indice del bit encendido mas a la derecha
builtin popcount(x) -> Cantida de bits encendidos
* Logaritmo en base 2 (entero). Indice del bit encendido
   mas a la izquierda. Si x es ll usar 63 y clzll(x).
int lg2(const int &x) { return 31-__builtin_clz(x); }
* Itera, con indices, los bits encendidos de una mascara.
// O(#bits encendidos)
for (int x = mask; x; x &= x-1) {
    int i = builtin ctz(x);
* Itera todas las submascaras de una mascara. (Iterar
   todas las submascaras de todas las mascaras es O(3^n))
// O(2^(#bits encendidos))
for (int sub = mask; sub; sub = (sub-1)&mask) {}
* retorna la siguiente mask con la misma cantidad
   encendida
ll nextMask(ll x){
   11 c = x & -x;
    11 r = x + c;
    return (((r ^ x) >> 2) / c) | r;
```

## 1.4 Cosas de strings

```
int conv(char ch) {return ((ch>='a' && ch<='z')?ch-'a':ch-
'A'+26);}
vector<string> split(string& s, char c=' ') {
    vector<string> res;
    stringstream ss(s);
    string sub;
    while(getline(ss, sub, c))res.push_back(sub);
    return res;
}

for(char& c:s)c=toupper(c);
for(char& c:s)c=tolower(c);
int n=stoi(s); // de string a entero
int n=stoi(s, nullptr, 2); // base 2
double d=stod(s); // de string a double
string s=to_string(n); // de entero a string
```

# 1.5 Custom Hashing

```
struct custom hash {
        static long long splitmix64(long long x) {
                x += 0x9e3779b97f4a7c15;
                x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
                x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
                return x ^ (x >> 31);
        size t operator()(long long x) const {
                static const long long FIXED_RANDOM =
                   chrono::steady_clock::now().
                   time since epoch().count();
                return splitmix64(x + FIXED_RANDOM);
        size t operator()(const pair<int,int>& x) const {
                return (size_t) x.first * 37U + (size t)
                   x.second;
        size_t operator()(const vector<int>& v) const {
                size t s = 0;
                for(auto &e : v)
                        s^=hash<int>()(e)+0x9e3779b9+(s)
                           <<6)+(s>>2);
                return s;
unordered_map<long long, int, custom_hash> safe_map;
gp hash table<int, int, custom hash> table;
```

#### 1.6 Random

```
typedef unsigned long long u64;
mt19937_64 rng (chrono::steady_clock::now().
    time_since_epoch().count());
u64 hash=rng();
mt19937 rng (chrono::steady_clock::now().time_since_epoch
    ().count());
int rand(int a, int b){return uniform_int_distribution<
    int>(a, b)(rng);} // uniform_real_distribution
```

## 2 Arboles

# 2.1 Centroid Decomposition

```
// O(nlog(n))
struct CentroidDecomposition{
        int dad[maxn],sz[maxn];
        set<int> adj[maxn]; // check, proc
        int operator[](int i) {return dad[i];}
        void addEdge(int x,int y) {adj[x].insert(y);adj[y
           l.insert(x);}
        void build(int v=0, int p=-1) {
                int n=dfsSz(v, p);
                int centroid=dfsCentroid(v, p, n);
                dad[centroid]=p;
                // add dfs for paths
                for(int u:adj[centroid]) {
                        adj[u].erase(centroid);
                        build (u, centroid);
                adi[centroid].clear();
        int dfsSz(int v,int p) {
                sz[v]=1;
                for(int u:adi[v]){
                        if (u==p) continue;
                        sz[v] + = dfsSz(u, v);
                return sz[v];
        int dfsCentroid(int v, int p, int n) {
                for(int u:adj[v]){
                        if (u==p) continue;
                        if(sz[u]>n/2)return dfsCentroid(u
                            , v, n);
                return v;
```

```
};
// for(int b=a;b!=-1;b=cd[b])
```

# 2.2 Heavy Light Decomposition

```
typedef long long T;
T null;
T oper(T a, T b);
// Segment tree
const int maxn=1e5+1; // >= 2e5, remove struct
bool edges=false; // arista padre
struct HLD{
        int par[maxn], root[maxn], dep[maxn];
        int sz[maxn], pos[maxn], ti;
        vector<int> adj[maxn];
        SegTree st;
        void addEdge(int x, int y) {adj[x].push_back(y);
            adj[y].push_back(x);}
        void dfsSz(int x) {
                 sz[x]=0;
                 for(int& y:adj[x]){
                         if (y==par[x]) continue;
                         par[y]=x; dep[y]=dep[x]+1;
                         dfsSz(v);
                         sz[x] += sz[y] +1;
                         if(sz[y]>sz[adj[x][0]]) swap(y, adj
                             [x][0];
        void dfsHld(int x) {
                 pos[x]=ti++;
                 for(int y:adj[x]){
                         if (y==par[x]) continue;
                         root[y] = (y = adj[x][0]?root[x]:y);
                         dfsHld(v);
        void build(int n,int v=0) {
                 root[v]=par[v]=v;
                 dep[v]=ti=0;
                 dfsSz(v);
                 dfsHld(v);
                 // vl palst(n);
                 // for (int i=0; i < n; ++i) palst [pos[i]] = a[i
                 // st.build(palst, n);
                 st.build(n);
        // O(\log^2(n))
        template <class Oper>
        void processPath(int x, int y, Oper op){
                 for(; root[x]!=root[y]; y=par[root[y]]) {
                         if (dep[root[x]]>dep[root[y]]) swap
                             (x,y);
```

```
op(pos[root[y]],pos[y]);
                if (dep[x]>dep[y]) swap(x,y);
                op(pos[x]+edges,pos[y]);
        void modifyPath(int x, int y, int v) {
                processPath(x,y,[this,&v](int 1, int r){
                         st.upd(l,r,v);
        T queryPath(int x, int y) {
                T res=null;
                processPath(x, y, [this, &res] (int 1, int r)
                         res=oper(res, st.get(l,r));
                return res;
        void modifySubtree(int x, int v) {st.upd(pos[x]+
            edges, pos[x]+sz[x], v);
        int querySubtree(int x) {return st.get(pos[x]+
            edges, pos[x]+sz[x]);
        void modify(int x, int v) {st.set(pos[x],v);}
        void modifyEdge(int x, int y, int v) {
                if (dep[x] < dep[y]) swap (x, y);
                modifv(x,v);
} ;
```

## 2.3 LCA

```
2.4 Sack
```

```
2 ARBOLES
```

```
int kth(int node, int k) {
    for(int l=maxlog-1;l>=0;--l) {
        if(node!=-1 && k& (1<<l)) {
            node=up[node][l];
        }
    }
    return node;
}

int lca(int a, int b) {
    a=kth(a, dep[a]-min(dep[a], dep[b]));
    b=kth(b, dep[b]-min(dep[a], dep[b]));
    if(a==b)return a;
    for(int l=maxlog-1;l>=0;--l) {
        if(up[a][l]!=up[b][l]) {
            a=up[a][l];
            b=up[b][l];
        }
    return up[a][0];
}
```

#### 2.4 Sack

```
const int maxn = 1e5+5;
int st[maxn], ft[maxn], ver[2*maxn];
int len[maxn], n, q, pos=0;
vi adj[maxn];
bool vis[maxn];
void ask(int v, bool add) {
        if(vis[v] && !add) {
                vis[v]=false;
                // delete node
        }else if(!vis[v] && add){
                vis[v]=true;
                // add node
// O(nlogn)
void dfs(int v=0, int p=-1, bool keep=true) {
        int mx=0, id=-1;
        for(int u:adj[v]){
                if (u==p) continue;
                if(len[u]>mx){
                        mx=len[u];
                        id=u;
        for(int u:adj[v]){
                if(u!=p && u!=id)
                         dfs(u,v,0);
```

## 2.5 Virtual Tree

```
const int maxn = 2e5+5;
vi adjVT[maxn], adj[maxn];
int st[maxn], ft[maxn], n, pos=0;
bool important[maxn];
bool upper(int v, int u){return st[v]<=st[u] && ft[v]>=ft
   [u];}
bool cmp(int v, int u) {return st[v] < st[u]; }</pre>
// O(klogk)
int virtualTree(vi nodes) {
        sort(all(nodes), cmp);
        int m=sz(nodes);
        for (int i=0; i<m-1; ++i) {</pre>
                 int v=lca(nodes[i], nodes[i+1]);
                 nodes.push back(v);
        sort(all(nodes), cmp);
        nodes.erase(unique(all(nodes)), nodes.end());
        for(int u:nodes)adjVT[u].clear();
        s.push back(nodes[0]);
        m=sz(nodes);
        for (int i=1; i < m; ++i) {</pre>
                 int v=nodes[i];
                 while (sz(s) \ge 2 \&\& !upper(s.back(), v))
                         adjVT[s[sz(s)-2]].push_back(s.
                             back());
                         s.pop_back();
                 s.push_back(v);
        while (sz(s) >= 2) {
                 adjVT[s[sz(s)-2]].push_back(s.back());
                 s.pop_back();
        return s[0];
```

```
// vi nodes(k);
// for(int& x:nodes)important[x]=true;
// int root=virtualTree(nodes);
// dp(root) - output answer - reset
```

## 3 Estructuras de Datos

# 3.1 Disjoint Set Union

```
struct dsu{
        vi p, size;
        int sets,maxSize;
        dsu(int n) {
                p.assign(n,0);
                size.assign(n,1);
                sets = n;
                for (int i = 0; i < n; i++) p[i] = i;
        int find_set(int i) {return (p[i] == i) ? i : (p[
           i = find set(p[i]);
        bool is_same_set(int i, int j) {return find_set(i
           ) == find set(j);
        void unionSet(int i, int j){
                if (!is_same_set(i, j)){
                        int a = find_set(i), b = find_set
                        if (size[a] < size[b]) swap(a, b)
                        p[b] = a;
                        size[a] += size[b];
                        maxSize = max(size[a], maxSize);
                        sets--;
};
```

# 3.2 Dynamic Connectivity Offline

```
struct dsu{
    vi p,rank,h;
    int sets;
    dsu(int n) {
        sets=n;
        p.assign(n,0);
        rank.assign(n,1);
        for(int i=0;i<n;++i)p[i]=i;
}</pre>
```

```
int get(int a) {return (a==p[a]?a:get(p[a]));}
        void unite(int a, int b) {
                a=qet(a); b=qet(b);
                if (a==b) return;
                if(rank[a]>rank[b])swap(a,b);
                rank[b]+=rank[a];
                h.push back(a);
                p[a]=b; sets--;
        void rollback(int x) {
                int len=h.size();
                while(len>x) {
                         int a=h.back();
                         h.pop back();
                         rank[p[a]] -= rank[a];
                         p[a]=a; sets++; len--;
};
enum { ADD, DEL, QUERY };
struct Query{int type, u, v;};
struct DvnCon{
        vector<Query> q;
        dsu uf; vi mt;
        map<pair<int,int>, int> prv;
        DynCon(int n): uf(n){}
        void add(int i, int j) {
                if(i>j)swap(i, j);
                q.push_back({ADD, i, j});
                mt.push back(-1);
                prv[{i,j}]=sz(q)-1;
        void remove(int i, int j){
                if(i > j) swap(i, j);
                q.push back({DEL, i, j});
                int pr=prv[{i, j}];
                mt[pr]=sz(q)-1;
                mt.push_back(pr);
        void query(){q.push_back({QUERY, -1, -1});mt.
            push back (-1);
        void process() { // answers all queries in order
                if(!sz(q))return;
                for(int i=0; i<sz(q);++i){
                         if(q[i].type==ADD && mt[i]<0)mt[i
                             l=sz(q);
                }qo(0, sz(q));
        void go(int s, int e){
                if(s+1==e){
                if (q[s].type == QUERY) cout << uf.sets << "\n"</pre>
                return:
                } int k=sz(uf.h), m=(s+e)/2;
                for(int i=e-1; i>=m; --i) {
```

# 3.3 Dynamic Segment Tree

```
T null=0, nolz=0;
T oper(T a, T b);
struct Node{
         T val, lz;
         int 1, r;
         Node *pl,*pr;
         Node(int ll, int rr) {
                  val=null;lz=nolz;
                  pl=pr=nullptr;
                  l=11; r=rr;
};
typedef Node* PNode;
void update(PNode x) {
         if (x->r-x->l==1) return;
         x-val=oper(x-pl-val,x-pr-val);
void extends(PNode x){
         if (x->r-x->1!=1 \&\& !x->p1) {
                  int m = (x->r+x->1)/2;
                  x \rightarrow pl = new Node(x \rightarrow l, m);
                  x \rightarrow pr = new Node(m, x \rightarrow r);
void propagate(PNode x) {
         if (x->r-x->l==1) return;
         if(x->lz==nolz) return;
         int m = (x->r+x->1)/2;
         // pl, pr
         x \rightarrow lz = nolz;
struct SegTree{
         PNode root:
         void upd(PNode x, int 1, int r, T v){
                  int 1x=x->1, rx=x->r;
                  if(lx>=r || l>=rx)return;
                  if(lx>=1 && rx<=r){
                           // val, lz
```

```
return;
                 extends(x):
                 propagate(x);
                upd(x->pl,l,r,v);
                upd(x->pr, l, r, v);
                update(x);
        T get(PNode x, int 1, int r){
                int 1x=x->1, rx=x->r;
                if(lx>=r || l>=rx)return null;
                if(lx>=1 && rx<=r) return x->val;
                extends (x);
                propagate(x);
                T v1=qet(x->pl,l,r);
                T v2=qet(x->pr,l,r);
                return oper (v1, v2);
        T get(int 1, int r) {return get(root, 1, r+1);}
        void upd(int 1, int r, T v) {upd(root, 1, r+1, v);}
        void build(int 1, int r){root=new Node(1, r+1);}
} ;
```

### 3.4 Fenwick Tree

```
typedef long long T;
struct FwTree{
        int n;
        vector<T> bit;
        FwTree(int n): n(n),bit(n+1){}
        T get(int r) {
                 T sum=0;
                 for(++r;r;r-=r&-r)sum+=bit[r];
                 return sum;
        T get (int 1, int r) {return get (r) - (1==0.00) : get (1
            -1));}
        void upd(int r, T v) {
                 for (++r; r<=n; r+=r&-r) bit [r] +=v;
};
struct FwTree2d{
        int n, m;
        vector<vector<T>> bit;
        FwTree2d() {}
        FwTree2d(int n, int m):n(n),m(m),bit(n+1, vector<
            T > (m+1, 0) \}
        T get(int x, int y) {
                 T v=0;
                 for(int i=x+1; i; i-=i&-i)
                 for(int j=y+1; j; j-=j&-j) v+=bit[i][j];
                 return v;
```

#### 3.5 Li Chao

```
// inf max abs value that the function may take
typedef long long ty;
struct Line {
  ty m, b;
 Line(){}
 Line(ty m, ty b): m(m), b(b) {}
 ty eval(ty x) { return m * x + b; }
};
struct nLiChao{
        // see coments for min
        nLiChao *left = nullptr, *right = nullptr;
        ty 1, r;
        Line line;
        nLiChao(ty l, ty r): l(l), r(r){
                line = \{0, -inf\}; // change to \{0, inf\};
        // T(Log(Rango)) M(Log(rango))
        void addLine(Line nline) {
                ty m = (1 + r) >> 1;
                bool lef = nline.eval(1) > line.eval(1);
                   // change > to <
        bool mid = nline.eval(m) > line.eval(m); //
           change > to <
                if (mid) swap(nline, line);
                if (r == 1) return;
        if (lef != mid) {
                        if (!left){
                                 left = new nLiChao(1, m);
                                 left -> line = nline;
                        else left -> addLine(nline);
        else{
                        if (!right) {
```

```
right = new nLiChao(m +
                                     1, r);
                                 right -> line = nline;
                         else right -> addLine(nline);
        // T(Log(Rango))
        tv get(tv x) {
                tv m = (l + r) >> 1;
                ty op1 = -\inf, op2 = -\inf; // change to
                    inf
                if(l == r) return line.eval(x);
                else if (x < m) {
                         if (left) op1 = left \rightarrow get(x);
                         return max(line.eval(x), op1); //
                              change max to min
                else{
                         if (right) op2 = right \rightarrow get(x);
                         return max(line.eval(x), op2); //
                              change max to min
};
int main() {
        // (rango superior) * (pendiente maxima) puede
            desbordarse
        // usar double o long double en el eval para
            estos casos
        // (puede dar problemas de precision)
        nLiChao liChao(0, 1e18);
```

#### 3.6 Link Cut Tree

};

```
T subsum(int u) {return (u?ns[u].sub:0);}
void push(int x){
        if(!x)return;
        int l=ns[x].ch[0], r=ns[x].ch[1];
        if(ns[x].flip){
                ns[l].flip^=1,ns[r].flip^=1;
                 swap(ns[x].ch[0], ns[x].ch[1]);
                    // check with st oper
                ns[x].flip=0;
        if(ns[x].lz){
                 // ...
                ns[x].sub+=ns[x].lz*ns[x].ssz;
                 ns[x].vir+=ns[x].lz*ns[x].vsz;
                // ...
void pull(int x){
        int l=ns[x].ch[0], r=ns[x].ch[1];
        push(1);push(r);
        ns[x].sz=size(1)+size(r)+1;
        ns[x].path=max({path(1), path(r), ns[x].}
            val });
        ns[x].sub=ns[x].vir+subsum(1)+subsum(r)+
            ns[x].val;
        ns[x].ssz=ns[x].vsz+subsize(l)+subsize(r)
void set(int x, int d, int y) {ns[x].ch[d]=y;ns[y
   ].p=x;pull(x);
void splay(int x) {
        auto dir=[&](int x) {
                 int p=ns[x].p;if(!p)return -1;
                 return ns[p].ch[0] == x?0:ns[p].ch
                    [1] == x?1:-1;
        auto rotate=[&](int x){
                 int y=ns[x].p, z=ns[y].p, dx=dir(x)
                    , dy = dir(y);
                 set (y, dx, ns[x].ch[!dx]);
                 set (x, !dx, y);
                if(^{\circ}dy) set(z,dy,x);
                ns[x].p=z;
        for(push(x); ~dir(x);) {
                 int y=ns[x].p, z=ns[y].p;
                 push(z);push(y);push(x);
                int dx=dir(x), dy=dir(y);
                if(^{\circ}dv) rotate (dx!=dy?x:v);
                 rotate(x);
```

```
struct LinkCut:SplayTree{ // 1-indexed
        LinkCut(int n):SplayTree(n){}
        int root(int u){
                access(u); splay(u); push(u);
                while (ns[u].ch[0]) {u=ns[u].ch[0]; push(u)
                return splay(u),u;
        int parent(int u){
                access(u); splay(u); push(u);
                u=ns[u].ch[0];push(u);
                while (ns[u].ch[1]) {u=ns[u].ch[1]; push(u)
                return splay(u),u;
        int access(int x){
                int u=x, v=0:
                for(;u;v=u,u=ns[u].p){
                         splav(u);
                         int& ov=ns[u].ch[1];
                         ns[u].vir+=ns[ov].sub;
                         ns[u].vsz+=ns[ov].ssz;
                         ns[u].vir-=ns[v].sub;
                         ns[u].vsz-=ns[v].ssz;
                         ov=v; pull(u);
                return splay(x), v;
        void reroot(int x){
                access (x); ns [x].flip^=1; push (x);
        void link(int u, int v) { // u \rightarrow v
                reroot (u);
                access(v);
                ns[v].vir+=ns[u].sub;
                ns[v].vsz+=ns[u].ssz;
                ns[u].p=v;pull(v);
        void cut(int u, int v){
                int r=root(u);
                reroot(u);
                access(v):
                ns[v].ch[0]=ns[u].p=0;pull(v);
                reroot(r);
        void cut(int u){ // cut parent
                access(u);
                ns[ns[u].ch[0]].p=0;
                ns[u].ch[0]=0;pull(u);
```

};

```
int lca(int u, int v) {
        if (root (u) !=root (v)) return -1;
        access(u); return access(v);
int depth(int u) {
        access(u); splay(u); push(u);
        return ns[u].sz;
T path(int u, int v) {
        int r=root(u);
        reroot (u); access (v); pull (v);
        T ans=ns[v].path;
        return reroot (r), ans;
void set(int u, T val){access(u);ns[u].val=val;
   pull(u);}
void upd(int u, int v, T val) {
        int r=root(u);
        reroot (u); access (v); splay (v);
        // lazv
        reroot(r);
T comp_size(int u) {return ns[root(u)].ssz;}
T subtree size(int u) {
        int p=parent(u);
        if(!p)return comp size(u);
        cut(u);int ans=comp_size(u);
        link(u,p); return ans;
T subtree_size(int u, int v) {
        int r=root(u);
        reroot (v); access (u);
        T ans=ns[u].vsz+1;
        return reroot(r), ans;
T comp sum(int u) {return ns[root(u)].sub;}
T subtree_sum(int u) {
        int p=parent(u);
        if(!p)return comp_sum(u);
        cut(u); T ans=comp_sum(u);
        link(u,p); return ans;
T subtree_sum(int u, int v) { // subtree of u, v
   father
        int r=root(u);
        reroot (v); access (u);
        T ans=ns[u].vir+ns[u].val; // por el
            reroot
        return reroot(r), ans;
```

# 3.7 Mos Algorithm

```
// O((n+q)*s), s=n^{(1/2)}
// O(q*(s+(n/s)^2)) \Rightarrow O(q*(n^2(3))), s=(2*(n^2))^2(1/3) -
     s=n^{(2/3)}
int s,n;
struct upd{int i,old,cur;};
struct query {int l,r,t,idx;};
bool cmp(query& a, query& b) {
        int x=a.1/s;
        if (a.1/s!=b.1/s) return a.1/s<b.1/s;
        if (a.r/s!=b.r/s) return (x&1?a.r<b.r:a.r>b.r);
        return a.t<b.t;</pre>
vector<int> ans;
vector<query> qu;
vector<upd> up;
int act();
void add(int i);
void remove(int i);
void update(int i,int v,int l,int r){
        if(l<=i && i<=r); // add, remove
void solve(){
        s=(int)ceil(sqrt(n));
        sort(all(qu), cmp);
        int l=0, r=-1, t=0;
         for(int i=0;i<sz(qu);++i){</pre>
                 while (t < qu[i].t) update (up[t].i, up[t].cur,</pre>
                     l,r),++t;
                 while(t>qu[i].t)--t, update(up[t].i, up[t].
                     old, l, r);
                 while (r < qu[i].r) add (++r);</pre>
                 while (1>qu[i].1) add (--1);
                 while (r>qu[i].r) remove (r--);
                 while (1<qu[i].1) remove (1++);
                 ans [qu[i].idx] = act();
int st[maxn],ft[maxn],ver[maxn*2];
bool vis[maxn];
void ask(int v) {
        vis[v]=!vis[v];
        if (vis[v]) add(v);
        else remove(v);
// \text{ query[i]} = \{st[a]+1, st[b], i\} + lca
```

#### 3.8 Ordered set

```
#include <ext/pb ds/assoc container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template<typename T> using ordered_set = tree<T,</pre>
   null_type,less<T>, rb_tree_tag,
   tree order statistics node update>;
template<typename T> using ordered_multiset = tree<T,</pre>
   null type, less equal <T>, rb tree tag,
   tree order statistics node update>;
// ---- CONSTRUCTOR ---- //
// 1. Para ordenar por MAX cambiar less<int> por greater<
// 2. Para multiset cambiar less<int> por less equal<int>
       Para borrar siendo multiset:
       int idx = st.order of key(value);
       st.erase(st.find_by_order(idx));
// ----- METHODS ----- //
st.find_by_order(k) // returns pointer to the k-th
   smallest element
st.order_of_key(x) // returns how many elements are
   smaller than x
st.find by order(k) == st.end() // true, if element does
   not exist
```

# 3.9 Persistent Segment Tree

```
typedef long long T;
struct Node{T val;int l,r;};
struct SegTree{
        vector<Node> ns;
        int act=0,size;
        vi roots;
        T null=0;
        T oper(T a, T b);
        void update(int x) {
                ns[x].val=oper(ns[ns[x].l].val, ns[ns[x].
                    rl.val);
        int newNode(T x){
                Node tmp=\{x, -1, -1\};
                ns.push back(tmp);
                return act++;
        int newNode(int 1, int r) {
                Node tmp={null,1,r};
                ns.push back(tmp);
                update (act);
                return act++;
```

```
int build(vector<T>& a, int 1, int r){
                 if (r-l==1) {return newNode(a[1]);}
                 int m = (1+r)/2;
                 return newNode (build (a, l, m), build (a, m,
                      r));
        int set(int x, int i, T v, int l, int r){
                 if (r-l==1) return newNode(v);
                 int m = (1+r)/2;
                 if (i<m) return newNode (set (ns[x].1, i, v,</pre>
                    l, m), ns[x].r);
                 else return newNode(ns[x].l, set(ns[x].r,
                     i, v, m, r));
        T get(int x, int lx, int rx, int l, int r){
                 if(lx>=r || l>=rx)return null;
                 if(lx>=l && rx<=r) return ns[x].val;</pre>
                 int m = (1x+rx)/2;
                 T v1=qet(ns[x].l, lx, m, l, r);
                 T v2 = qet(ns[x].r, m, rx, l, r);
                 return oper (v1, v2);
        T get(int 1, int r, int time) {return get(roots[
            time], 0, size, 1, r+1);}
        void set(int i, T v, int time){roots.push_back(
            set(roots[time], i, v, 0, size));}
        void build(vector<T>& a, int n) {size=n; roots.
            push_back(build(a, 0, size));}
} ;
```

# 3.10 RMQ

# 3.11 Segment Tree Iterativo

```
struct seatree{
    int n; vl v; ll nulo = 0;
    11 op(ll a, ll b) {return a + b;}
    seqtree(int n) : n(n), v(2*n, nulo){}
    segtree(vl &a): n(sz(a)), v(2*n){
        for (int i = 0; i < n; i++) v[n + i] = a[i];
        for (int i = n-1; i > = 1; --i) v[i] = op(v[i < 1], v
           [i<<1|1]);
    void upd(int k, ll nv){
        for (v[k += n] = nv; k > 1; k >>= 1) v[k>>1] = op
            (v[k], v[k^1]);
    11 get(int 1, int r){
        11 vl = nulo, vr = nulo;
        for (1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1)
            if (1&1) vl = op(vl, v[1++]);
            if (r\&1) vr = op(v[--r], vr);
        return op (vl, vr);
};
```

# 3.12 Segment Tree Recursivo

```
void build(vector<T>& a, int n) {
        size=1;
        while (size<n) size *= 2;</pre>
        vals.resize(2*size);
        lazv.assign(2*size, nolz);
        build(a, 0, 0, size);
void propagate(int x, int lx, int rx){
        if (rx-lx==1) return;
        if(lazy[x]==nolz)return;
        int m = (1x+rx)/2;
        // 2*x+1, 2*x+2 (lazy, vals)
        lazy[x]=nolz;
void upd(int 1, int r, T v,int x, int lx, int rx)
        if(lx>=r || l>=rx) return;
        if(lx>=1 && rx<=r){
                 // lazy, vals
                 return;
        propagate(x,lx,rx);
        int m = (1x+rx)/2;
        upd(1, r, v, 2*x+1, 1x, m);
        upd(1, r, v, 2 \times x + 2, m, rx);
        vals[x]=oper(vals[2*x+1], vals[2*x+2]);
void set(int i, T v, int x, int lx, int rx){
        if(rx-lx==1){
                 vals[x]=v;
                 return;
        propagate(x,lx,rx);
        int m = (1x+rx)/2;
        if (i<m) set (i, v, 2*x+1, lx, m);
        else set(i, v, 2*x+2, m, rx);
        vals[x] = oper(vals[2*x+1], vals[2*x+2]);
T get (int 1, int r, int x, int lx, int rx) {
        if(lx>=r || l>=rx)return null;
        if(lx>=1 && rx<=r)return vals[x];</pre>
        propagate(x,lx,rx);
        int m = (1x+rx)/2;
        T v1=get (1, r, 2*x+1, 1x, m);
        T v2=qet(1,r,2*x+2,m,rx);
        return oper (v1, v2);
T get(int 1, int r) {return get(1,r+1,0,0,size);}
void upd(int 1, int r, T v) {upd(1,r+1,v,0,0,size)
   ; }
```

```
void set(int i, T val) {set(i,val,0,0,size);}
};
```

# 3.13 Segment Tree 2D

```
const int N=1000+1;
ll st[2*N][2*N];
struct SeaTree{
        int n,m,neutro=0;
        inline ll op(ll a, ll b) {return a+b;}
        SegTree(int n, int m): n(n), m(m) {
                 for (int i=0; i<2*n; ++i) for (int j=0; j<2*m
                     ;++j)st[i][j]=neutro;
        SegTree (vector\langle vi \rangle \& a): n(sz(a)), m(n ? sz(a[0])
            : 0) { build(a); }
        void build(vector<vi>& a) {
                 for (int i=0; i< n; ++i) for (int j=0; j< m; ++j)
                     st[i+n][j+m]=a[i][j];
                 for (int i=0; i<n; ++i) for (int j=m-1; j>=1; --
                     j) st[i+n][j] = op(st[i+n][j << 1], st[i+n]
                     ][i<<1|1]);
                 for (int i=n-1; i>=1; --i) for (int j=0; j<2*m
                     ;++j) st[i][j]=op(st[i<<1][j], st[i
                     <<1|1][†]);
        void upd(int x, int y, ll v){
                 st[x+n][y+m]=v;
                 for (int j=y+m; j>1; j>>=1) st [x+n] [j>>1] = op (
                     st[x+n][j], st[x+n][j^1];
                 for(int i=x+n;i>1;i>>=1) for(int j=y+m;j;j
                     >>=1) st[i>>1][j]=op(st[i][j], st[i^1][
                     j]);
        11 get(int x0, int y0, int x1, int y1){
                 11 r=neutro;
                 for (int i0=x0+n, i1=x1+n+1; i0<i1; i0>>=1, i1
                     >>=1) {
                          int t[4], q=0;
                          if (i0&1) t [q++]=i0++;
                          if (i1&1) t [q++]=--i1;
                          for (int k=0; k < q; ++k) for (int j0=y0
                              +m, j1=y1+m+1; j0<j1; j0>>=1, j1
                              >>=1) {
                                   if(j0&1)r=op(r,st[t[k]][
                                       j0++1);
                                   if(j1&1) r = op(r, st[t]k
                                       ]][-- | 1]);
```

```
return r;
};
```

# 3.14 Segment Tree Beats

```
typedef long long T;
T null=0, noVal=0;
T INF=1e18;
struct Node {
        T sum, lazv;
        T max1, max2, maxc;
        T min1, min2, minc;
struct SeqTree{
        vector<Node> vals;int size;
        void oper(int a, int b, int c); // node c, left a
            , right b;
        Node single(T x) {
                Node tmp;
                tmp.sum=tmp.max1=tmp.min1=x;
                tmp.maxc=tmp.minc=1;
                tmp.lazy=noVal;
                tmp.max2=-INF;
                tmp.min2=INF;
                return tmp;
        void build(vector<T>& a, int n);
        void propagateMin(T v, int x, int lx, int rx){
                 if (vals[x].max1<=v) return;</pre>
                vals[x].sum-=vals[x].max1*vals[x].maxc;
                vals[x].max1=v;
                vals[x].sum+=vals[x].max1*vals[x].maxc;
                if(rx-lx==1){
                         vals[x].min1=v;
                }else{
                         if(v<=vals[x].min1) {</pre>
                                 vals[x].min1=v;
                         }else if(v<vals[x].min2){</pre>
                                 vals[x].min2=v;
        void propagateAdd(T v, int x, int lx, int rx) {
                vals[x].sum+=v*((T)(rx-lx));
                vals[x].lazy+=v;
                vals[x].max1+=v;
                vals[x].min1+=v;
                if (vals[x].max2!=-INF) vals[x].max2+=v;
                if (vals[x].min2!=INF) vals[x].min2+=v;
        void propagate(int x, int lx, int rx){
```

```
if (rx-lx==1) return;
        int m = (1x+rx)/2;
        if(vals[x].lazy!=noVal){
                propagateAdd(vals[x].lazy, 2*x+1,
                     lx, m);
                propagateAdd(vals[x].lazy, 2*x+2,
                     m, rx);
                vals[x].lazy=noVal;
        propagateMin(vals[x].max1, 2*x+1, lx, m);
        propagateMin(vals[x].max1, 2*x+2, m, rx);
void updAdd(int 1, int r, T v, int x, int lx, int
   rx) {
        if(lx>=r || l>=rx) return;
        if(lx>=l && rx<=r){
                propagateAdd(v, x, lx, rx);
                return:
        propagate(x,lx,rx);
        int m = (1x+rx)/2;
        updAdd(1, r, v, 2*x+1, 1x, m);
        updAdd(1, r, v, 2*x+2, m, rx);
        oper (2*x+1, 2*x+2, x);
void updMin(int 1, int r, T v, int x, int lx, int
   rx) {
        if(lx>=r || l>=rx || vals[x].max1<v)
            return;
        if(lx>=1 && rx<=r && vals[x].max2<v){
                propagateMin(v, x, lx, rx);
                return;
        propagate(x,lx,rx);
        int m=(1x+rx)/2;
        updMin(l,r,v,2*x+1,lx,m);
        updMin(l,r,v,2*x+2,m,rx);
        oper (2*x+1, 2*x+2, x);
void updAdd(int 1, int r, T v) {updAdd(1,r+1,v)
   ,0,0,size);}
void updMin(int 1, int r, T v) {updMin(1,r+1,v)
   ,0,0,size);}
```

# 3.15 Sparse Table 2D

};

```
const int MAX_N = 100;
const int MAX_M = 100;
const int KN = log2 (MAX_N) +1;
const int KM = log2 (MAX_M) +1;
```

```
int table[KN][MAX N][KM][MAX M];
int _log2N[MAX_N+1];
int log2M[MAX M+1];
int MAT[MAX_N][MAX_M];
int n, m, ic, ir, jc, jr;
void calc log2() {
    log2N[1] = 0;
    [log2M[1] = 0;
    for (int i = 2; i <= MAX N; i++) log2N[i] = log2N[i</pre>
       /21 + 1;
    for (int i = 2; i <= MAX M; i++) log2M[i] = log2M[i</pre>
       /21 + 1;
void build() {
    for (ir = 0; ir < n; ir++) {</pre>
        for (ic = 0; ic < m; ic++)
            table[0][ir][0][ic] = MAT[ir][ic];
        for (jc = 1; jc < KM; jc++)
            for (ic = 0; ic + (1 << (jc-1)) < m; ic++)
                table[0][ir][jc][ic] = min(table[0][ir][
                    jc-1[ic], table[0][ir][jc-1][ic + (1
                    << (ic-1)));
    for (jr = 1; jr < KN; jr++)
        for (ir = 0; ir < n; ir++)
            for (jc = 0; jc < KM; jc++)
                for (ic = 0; ic < m; ic++)
                    table[jr][ir][jc][ic] = min(table[jr
                        -1][ir][jc][ic], table[jr-1][ir
                        +(1<<(jr-1))][jc][ic]);
int rmq(int x1, int y1, int x2, int y2) {
    int lenx = x2-x1+1;
    int kx = _log2N[lenx];
    int leny = y2-y1+1;
    int ky = log2M[lenv];
    int min R1 = min(table[kx][x1][ky][y1], table[kx][x1
       [ky][y2 + 1 - (1 << ky)]);
    int min_R2 = min(table[kx][x2+1-(1<<kx)][ky][y1],
       table[kx][x2+1- (1 << kx)][ky][y2 + 1 - (1 << ky)]);
    return min(min_R1, min_R2);
```

# 3.16 Sqrt Descomposition

```
typedef long long T;
struct Sqrt { // O(n/b+b)
         int b; // check b
        vector<T> nums,blocks;
```

```
void build(vector<T>& arr, int n) {
                  b=(int)ceil(sgrt(n));nums=arr;
                 blocks.assign(b, 0);
                  for (int i=0; i < n; ++i) {</pre>
                          blocks[i/b]+=nums[i];
        void set(int x, int v){
                  blocks [x/b] -= nums [x];
                  nums [x]=v;
                  blocks[x/b]+=nums[x];
        T get(int r) {
                  T res=0;
                  for(int i=0;i<r/b;++i){res+=blocks[i];}</pre>
                  for(int i=(r/b)*b;i<r;++i) {res+=nums[i];}</pre>
                  return res;
        T get(int 1, int r) {return get(r+1)-get(1);}
};
```

## 3.17 Treap

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```
// treap => order asc, implicit treap => order array
typedef long long T;
struct Treap{
        Treap *1,*r,*dad;
        u64 prior;
        T sz, value, sum, lz;
        Treap(T v) {
                 l=r=nullptr;
                 1z=0; sz=1;
                 prior=rng();
                 value=sum=v;
         ~Treap() {delete l; delete r; }
};
typedef Treap* PTreap;
T cnt (PTreap x) {return (!x?0:x->sz);}
T sum(PTreap x) {return (!x?0:x->sum);}
void propagate(PTreap x) {
        if(x && x->lz) {
                 if(x->1); // lz, value, sum ...
                 if (x->r); // lz, value, sum ...
                 x - > 1z = 0;
void update(PTreap x) {
        propagate (x->1);
```

```
propagate (x->r):
         x->sz=cnt(x->1)+cnt(x->r)+1;
         x\rightarrow sum=sum(x\rightarrow 1)+sum(x\rightarrow r)+x\rightarrow value;
         if (x->1) x->1->dad=x;
         if (x->r) x->r-> dad=x;
void upd(PTreap x, T v) {
         if(!x)return;
         update(x);
         // lz, value, sum ...
// pair<PTreap, PTreap> split(PTreap x, T kev) { // f <=</pre>
pair<PTreap, PTreap> split(PTreap x, int left){ // cnt(f)
     == left
         if(!x)return {nullptr, nullptr};
         propagate(x);
         if(cnt(x->1)>=left) { // if(x->value>key) {}
                  auto got=split(x->1, left); // , key);
                  x->l=qot.second;
                  update(x);
                  return {got.first, x};
         }else{
                  auto got=split(x->r, left-cnt(x->1)-1);
                      // , key);
                  x->r=qot.first;
                  update(x);
                  return {x, got.second};
PTreap merge (PTreap x, PTreap y) {
         if(!x)return v;
         if(!v)return x;
         propagate(x);
         propagate(y);
         if (x->prior<=y->prior) {
                  x \rightarrow r = merge(x \rightarrow r, y);
                  update(x);
                  return x;
         }else{
                  v \rightarrow l = merge(x, v \rightarrow l);
                  update(v);
                  return y;
PTreap combine (PTreap x, PTreap y) {
         if(!x)return v;
         if(!y)return x;
         if (x->prior<y->prior) swap (x, y);
         auto z=split(y, x->value);
         x \rightarrow r = combine(\bar{x} \rightarrow r, z.second);
         x->l=combine(z.first, x->l);
         return x;
```

```
3.18 Two Stacks
```

```
T kth(PTreap& x, int k){ // indexed 0
        if(!x)return null;
        if (k==cnt (x->1)) return x->value;
        if (k < cnt(x->1)) return kth(x->1, k);
        return kth(x->r, k-cnt(x->1)-1);
pair<int, T> lower_bound(PTreap x, T key) { // index,
   value
        if(!x)return {0, null};
        if(x->value<kev) {</pre>
                 auto y=lower_bound(x->r, key);
                v.first+=cnt(x->1)+1;
                return v;
        auto y=lower_bound(x->1, key);
        if (y.first==cnt(x->1))y.second=x->value;
        return v;
void dfs(PTreap x) {
        if(!x)return;
        propagate(x);
        dfs(x->1); cout<<x->value<<" "; dfs(x->r);
// PTreap root=nullptr;
// PTreap act=new Treap(c);
// root=merge(root, act);
```

#### 3.18 Two Stacks

```
typedef long long T;
struct Node{T val,acum;};
struct TwoStacks{
        stack<Node> s1,s2;
        void add(T x){
                Node tmp=\{x, x\};
                if(!s2.empty()){
                // tmp.acum + s2.top().acum
                s2.push(tmp);
        void remove(){
                if(s1.empty()){
                         while(!s2.empty()){
                                 Node tmp=s2.top();
                                 if(s1.emptv()){
                                 // tmp.acum = tmp.val
                                 // tmp.acum + s1.top().
                                    acum
```

## 3.19 Wavelet Tree

```
const int maxn = 1e5+5, maxv = 1e9, minv = -1e9;
struct WaveletTree{ // indexed 1 - O(nlogn)
        int lo, hi;
        WaveletTree *1, *r;
        int *b, bsz, csz;
        11 *c;
        WaveletTree() {
                hi=bsz=csz=0;
                l=r=NULL;
                10=1;
        void build(int *from, int *to, int x, int y) {
                lo=x, hi=y;
                if (from>=to) return;
                int mid=lo+(hi-lo)/2;
                auto f=[mid] (int x) {return x<=mid; };</pre>
                b=(int*)malloc((to-from+2)*sizeof(int));
                bsz=0:
                b[bsz++]=0;
                c=(l1*)malloc((to-from+2)*sizeof(l1));
                csz=0;
                c[csz++]=0;
                for(auto it=from;it!=to;++it) {
                         b[bsz] = (b[bsz-1] + f(*it));
                         c[csz] = (c[csz-1] + (*it));
                         bsz++; csz++;
                if (hi==lo) return;
                auto pivot=stable partition(from, to, f);
                l=new WaveletTree();
```

```
1->build(from, pivot, lo, mid);
                r=new WaveletTree();
                r->build(pivot, to, mid+1, hi);
        //kth smallest element in [1, r]
        int kth(int 1, int r, int k){
                if(l>r) return 0;
                if(lo==hi)return lo;
                 int inLeft=b[r]-b[l-1], lb=b[l-1], rb=b[r
                 if (k<=inLeft) return this->l->kth(lb+1, rb
                 return this->r->kth(l-lb, r-rb, k-inLeft)
        //count of numbers in [1, r] Less than or equal
            to k
        int lte(int 1, int r, int k) {
                if(1>r || k<10) return 0;
                 if (hi<=k) return r-l+1;</pre>
                int lb=b[l-1], rb=b[r];
                return this->l->lte(lb+1, rb, k)+this->r
                    ->lte(l-lb, r-rb, k);
        //count of numbers in [l, r] equal to k
        int count(int 1, int r, int k){
                if(1>r || k<10 || k>hi) return 0;
                if (lo==hi) return r-l+1;
                 int lb=b[l-1], rb=b[r];
                 int mid=(lo+hi)>>1;
                if (k<=mid) return this->l->count(lb+1, rb,
                return this->r->count(l-lb, r-rb, k);
        //sum of numbers in [l ,r] less than or equal to
           k
        11 sum(int 1, int r, int k){
                if(l>r || k<lo)return 0;</pre>
                if (hi<=k) return c[r]-c[l-1];</pre>
                 int lb=b[l-1], rb=b[r];
                 return this->l->sum(lb+1, rb, k)+this->r
                    \rightarrowsum(l-lb, r-rb, k);
        ~WaveletTree(){
                delete 1:
                delete r;
};
// int a[maxn];
// WaveletTree wt;
// for(int i=1;i<=n;++i)cin>>a[i];
```

```
// wt.build(a+1, a+n+1, minv, maxv);
```

### 3.20 Trie Bit

```
struct node{
 int childs[2]{-1, -1};
struct TrieBit{
    vector<node> nds;
        vi passNums;
        TrieBit(){
        nds.pb(node());
        passNums.pb(0);
    void insert(int num) {
        int cur = 0;
        for (int i = 30; i >= 0; i--) {
            bool bit = (num >> i) & 1;
            if (nds[cur].childs[bit] == -1) {
                nds[cur].childs[bit] = nds.size();
                nds.pb(node());
                passNums.pb(0);
            passNums[cur]++;
            cur = nds[cur].childs[bit];
        passNums[cur]++;
    void remove(int num) {
        int cur = 0;
        for (int i = 30; i >= 0; i--) {
            bool bit = (num >> i) & 1;
            passNums[cur]--;
            cur = nds[cur].childs[bit];
        passNums[cur]--;
    int maxXor(int num) {
        int ans = 0;
        int cur = 0;
        for(int i = 30; i >= 0; i--) {
            bool bit = (num >> i) & 1;
            int n1 = nds[cur].childs[!bit];
            if (n1 != -1 && passNums[n1]) {
                ans += (1 << i);
                bit = !bit;
```

```
cur = nds[cur].childs[bit];
}
return ans;
};
```

# 4 Flujos

### 4.1 Blossom

```
/// Complexity: O(|E||V|^2)
/// Tested: https://tinyurl.com/oe5rnpk
struct network {
  struct struct edge { int v; struct edge * n; };
  typedef struct edge* edge;
  int n;
  struct edge pool[MAXE]; ///2*n*n;
  edge top;
  vector<edge> adi;
  queue<int> q;
  vector<int> f, base, inq, inb, inp, match;
  vector<vector<int>> ed;
  network(int n) : n(n), match(n, -1), adj(n), top(pool),
      f(n), base(n),
                   ing(n), inb(n), inp(n), ed(n, vector<
                      int>(n)) {}
  void add edge(int u, int v) {
    if(ed[u][v]) return;
    ed[u][v] = 1;
    top->v = v, top->n = adj[u], adj[u] = top++;
    top->v = u, top->n = adj[v], adj[v] = top++;
  int get lca(int root, int u, int v) {
    fill(inp.begin(), inp.end(), 0);
    while (1)
      inp[u = base[u]] = 1;
      if(u == root) break;
      u = f[match[u]];
    while(1) {
      if(inp[v = base[v]]) return v;
      else v = f[ match[v] ];
  void mark(int lca, int u) {
    while(base[u] != lca) {
      int v = match[u];
      inb[base[u]] = 1;
      inb[base[v]] = 1;
      u = f[v];
      if(base[u] != lca) f[u] = v;
```

```
void blossom contraction(int s, int u, int v) {
  int lca = get lca(s, u, v);
  fill(inb.begin(), inb.end(), 0);
  mark(lca, u); mark(lca, v);
  if(base[u] != lca) f[u] = v;
  if(base[v] != lca) f[v] = u;
  for(int u = 0; u < n; u++)
    if(inb[base[u]]) {
      base[u] = lca;
      if(!inq[u]) {
          inq[u] = 1;
          q.push(u);
int bfs(int s) {
  fill(ing.begin(), ing.end(), 0);
  fill(f.begin(), f.end(), -1);
  for(int i = 0; i < n; i++) base[i] = i;</pre>
  q = queue<int>();
  q.push(s);
  inq[s] = 1;
 while(q.size()) {
    int u = q.front(); q.pop();
    for (edge e = adj[u]; e; e = e->n) {
      int v = e -> v;
      if(base[u] != base[v] && match[u] != v) {
        if((v == s) \mid | (match[v] != -1 \&\& f[match[v]])
          blossom contraction(s, u, v);
        else if (f[v] == -1) {
          f[v] = u;
          if (match[v] == -1) return v;
          else if(!ing[match[v]]) {
            ing[match[v]] = 1;
            q.push(match[v]);
  return -1;
int doit(int u) {
  if(u == -1) return 0;
  int v = f[u];
  doit(match[v]);
  match[v] = u; match[u] = v;
  return u != -1;
/// (i < net.match[i]) => means match
int maximum matching() {
  int ans = 0;
```

```
for(int u = 0; u < n; u++)
    ans += (match[u] == -1) && doit(bfs(u));
return ans;
};</pre>
```

#### 4.2 Dinic

```
// O(|E| * |V|^2)
struct edge { ll v, cap, inv, flow, ori; };
struct network {
  ll n, s, t;
  vector<ll> lvl;
  vector<vector<edge>> q;
 network(ll n) : n(n), lvl(n), g(n) {}
  void add_edge(int u, int v, ll c) {
    q[u].push back({v, c, sz(q[v]), 0, 1});
    g[v].push_back({u, 0, sz(g[u])-1, c, 0});
 bool bfs() {
    fill(lvl.begin(), lvl.end(), -1);
    queue<11> q;
    lvl[s] = 0;
    for(q.push(s); q.size(); q.pop()) {
      ll u = q.front();
      for(auto &e : q[u]) {
        if(e.cap > 0 && lvl[e.v] == -1) {
          lvl[e.v] = lvl[u]+1;
          q.push(e.v);
    return lvl[t] != -1;
  11 dfs(ll u, ll nf) {
    if(u == t) return nf;
    11 \text{ res} = 0:
    for(auto &e : q[u]) {
      if(e.cap > 0 \&\& lvl[e.v] == lvl[u]+1) {
        ll tf = dfs(e.v, min(nf, e.cap));
        res += tf; nf -= tf; e.cap -= tf;
        q[e.v][e.inv].cap += tf;
        q[e.v][e.inv].flow -= tf;
        e.flow += tf;
        if(nf == 0) return res;
    if(!res) lvl[u] = -1;
    return res;
  ll \max flow(ll so, ll si, ll res = 0) {
    s = so; t = si;
    while(bfs()) res += dfs(s, LONG_LONG_MAX);
    return res;
```

```
void min_cut() {
    queue<11> q;
    vector<bool> vis(n, 0);
    vis[s] = 1;
    for(q.push(s); q.size(); q.pop()) {
      ll u = q.front();
      for(auto &e : g[u]) {
        if(e.cap > 0 && !vis[e.v]) {
          q.push(e.v);
          vis[e.v] = 1;
    vii ans;
    for (int i = 0; i<n; i++) {
        for (auto &e : q[i]) {
            if (vis[i] && !vis[e.v] && e.ori) {
                ans.push back(\{i+1, e.v+1\});
    for (auto [x, y] : ans) cout << x << ' ' << y << ln;</pre>
  bool dfs2(vi &path, vector<bool> &vis, int u) {
    vis[u] = 1;
    for (auto &e : q[u]) {
      if (e.flow > 0 && e.ori && !vis[e.v]) {
        if (e.v == t || dfs2(path, vis, e.v)){
          path.push back(e.v);
          e.flow = 0;
          return 1;
    return 0;
  void disjoint paths() {
    vi path;
    vector<bool> vis(n, 0);
    while (dfs2(path, vis, s)){
      path.push back(s);
      reverse (all (path));
      cout << sz(path) << ln;
      for (int v : path) cout << v+1 << ' ';</pre>
      cout << ln;
      path.clear(); vis.assign(n, 0);
};
```

## 4.3 Edmonds Karp

```
4.4 Hopcroft Karp
```

```
//O(V * E^2)
ll bfs(vector<vi> &adj, vector<vl> &capacity, int s, int
   t, vi& parent) {
    fill(parent.begin(), parent.end(), -1);
    parent[s] = -2;
    queue<pll> q;
    q.push({s, INFL});
    while (!q.empty()) {
        int cur = q.front().first;
        11 flow = q.front().second;
        q.pop();
        for (int next : adj[cur]) {
            if (parent[next] == -1LL && capacity[cur][
               nextl)
                parent[next] = cur;
                11 new flow = min(flow, capacity[cur][
                    next]);
                if (next == t)
                    return new_flow;
                q.push({next, new_flow});
    return 0;
11 maxflow(vector<vi> &adj, vector<vl> &capacity, int s,
   int t, int n) {
   11 \text{ flow} = 0;
    vi parent(n);
    ll new flow;
    while ((new_flow = bfs(adj, capacity, s, t, parent)))
        flow += new flow;
        int cur = t;
        while (cur != s) {
            int prev = parent[cur];
            capacitv[prev][cur] -= new flow;
            capacity[cur][prev] += new flow;
            cur = prev;
    return flow;
```

# 4.4 Hopcroft Karp

```
// Complexity: O(|E|*sqrt(|V|))
struct mbm {
  vector<vector<int>> g;
  vector<int>> d, match;
```

```
int nil, l, r;
/// u \rightarrow 0 to 1, v \rightarrow 0 to r
mbm(int l, int r) : q(l+r), d(l+l+r, INF), match(l+r, l)
   +r),
                     nil(l+r), l(l), r(r) {}
void add_edge(int a, int b) {
  q[a].push_back(l+b);
  q[l+b].push back(a);
bool bfs() {
  queue<int> q;
  for(int u = 0; u < 1; u++) {
    if (match[u] == nil) {
      d[u] = 0;
      q.push(u);
    } else d[u] = INF;
  d[nil] = INF;
  while(q.size()) {
    int u = q.front(); q.pop();
    if(u == nil) continue;
    for(auto v : q[u]) {
      if(d[match[v]] == INF) {
        d[match[v]] = d[u]+1;
        q.push(match[v]);
  return d[nil] != INF;
bool dfs(int u) {
  if(u == nil) return true;
  for(int v : q[u]) {
    if(d[match[v]] == d[u]+1 \&\& dfs(match[v])) {
      match[v] = u; match[u] = v;
      return true;
  d[u] = INF;
  return false;
int max matching() {
  int ans = 0;
  while(bfs()) {
    for(int u = 0; u < 1; u++) {
      ans += (match[u] == nil && dfs(u));
  return ans;
void matchs() {
  for (int i = 0; i<1; i++) {
      if (match[i] == l+r) continue;
      cout << i+1 << ' ' << match[i]+1-l << ln;
```

4.5 Maximum Bipartite Matching

};

```
// O(|E| * |V|)
struct mbm {
  int 1, r;
  vector<vector<int>> a;
  vector<int> match, seen;
  mbm(int l, int r) : l(l), r(r), q(l), match(r), seen(r)
  void add edge(int 1, int r) { g[1].push back(r); }
  bool dfs(int u)
    for(auto v : q[u]) {
      if(seen[v]++) continue;
      if(match[v] == -1 \mid | dfs(match[v]))  {
        match[v] = u;
        return true;
    return false;
  int max matching() {
    int ans = 0;
    fill(match.begin(), match.end(), -1);
    for(int u = 0; u < 1; ++u) {
      fill(seen.begin(), seen.end(), 0);
      ans += dfs(u):
    return ans;
  void matchs() {
    for (int i = 0; i<r; i++) {
        if (match[i] == -1) continue;
        cout << match[i]+1 << ' ' << i+1 << ln;</pre>
};
```

### 4.6 Minimum Cost Maximum Flow

```
/// Complexity: O(|V|*|E|^2*log(|E|))
template <class type>
struct mcmf {
    struct edge { int u, v, cap, flow; type cost; };
    int n;
    vector<edge> ed;
    vector<vector<int>> g;
    vector<int> p;
    vector<type> d, phi;
```

```
mcmf(int n) : n(n), g(n), p(n), d(n), phi(n) {}
 void add_edge(int u, int v, int cap, type cost) {
    q[u].push back(ed.size());
    ed.push_back({u, v, cap, 0, cost});
    g[v].push_back(ed.size());
    ed.push back(\{v, u, 0, 0, -cost\});
 bool dijkstra(int s, int t) {
    fill(d.begin(), d.end(), INF TYPE);
    fill(p.begin(), p.end(), -1);
    set<pair<type, int>> q;
    d[s] = 0;
    for(q.insert({d[s], s}); q.size();) {
      int u = (*q.begin()).second; q.erase(q.begin());
      for(auto v : q[u]) {
        auto &e = ed[v];
        type nd = d[e.u]+e.cost+phi[e.u]-phi[e.v];
        if(0 < (e.cap-e.flow) && nd < d[e.v]) {
          q.erase({d[e.v], e.v});
          d[e.v] = nd; p[e.v] = v;
          q.insert({d[e.v], e.v});
    for(int i = 0; i < n; i++) phi[i] = min(INF TYPE, phi</pre>
       [i]+d[i]);
    return d[t] != INF TYPE;
 pair<int, type> max_flow(int s, int t) {
    type mc = 0;
    int mf = 0;
    fill(phi.begin(), phi.end(), 0);
    while(dijkstra(s, t)) {
      int flow = INF;
      for (int v = p[t]; v != -1; v = p[ed[v].u])
        flow = min(flow, ed[v].cap-ed[v].flow);
      for(int v = p[t]; v != -1; v = p[ed[v].u]) {
        edge &e1 = ed[v];
        edge &e2 = ed[v^1];
        mc += e1.cost*flow;
        e1.flow += flow;
        e2.flow -= flow;
      mf += flow;
    return {mf, mc};
};
```

# 4.7 Weighted Matching

```
/// Complexity: O(|V|^3)
typedef int type;
```

```
struct matching weighted {
  int 1, r;
 vector<vector<type>> c;
 matching_weighted(int 1, int r) : 1(1), r(r), c(1,
     vector<type>(r)) {
    assert(1 \le r);
 void add_edge(int a, int b, type cost) { c[a][b] = cost
  type matching() {
   vector<type> v(r), d(r); // v: potential
    vector\langle int \rangle ml(1, -1), mr(r, -1); // matching pairs
    vector<int> idx(r), prev(r);
    iota(idx.begin(), idx.end(), 0);
    auto residue = [&] (int i, int j) { return c[i][j]-v[j
    for(int f = 0; f < 1; ++f) {
      for (int j = 0; j < r; ++j) {
        d[j] = residue(f, j);
        prev[j] = f;
      type w;
      int j, 1;
      for (int s = 0, t = 0;;) {
        if(s == t) {
          l = s;
          w = d[ idx[t++] ];
          for (int k = t; k < r; ++k) {
            j = idx[k];
            type h = d[j];
            if (h <= w) {
              if (h < w) t = s, w = h;
              idx[k] = idx[t];
              idx[t++] = i;
          for (int k = s; k < t; ++k) {
            j = idx[k];
            if (mr[j] < 0) goto aug;
        int q = idx[s++], i = mr[q];
        for (int k = t; k < r; ++k) {
          j = idx[k];
          type h = residue(i, j) - residue(i, q) + w;
          if (h < d[i]) {
            d[j] = h;
            prev[j] = i;
            if (h == w) {
              if(mr[j] < 0) goto aug;</pre>
              idx[k] = idx[t];
              idx[t++] = j;
```

```
aug: for (int k = 0; k < 1; ++k)
      v[ idx[k] ] += d[ idx[k] ] - w;
int i;
do {
    mr[j] = i = prev[j];
    swap(j, ml[i]);
    } while (i != f);
}
type opt = 0;
for (int i = 0; i < 1; ++i)
    opt += c[i][ml[i]]; // (i, ml[i]) is a solution
return opt;
}
};</pre>
```

## 4.8 Hungarian

```
const int N = 509;
/* Complexity: O(n^3) but optimized
It finds minimum cost maximum matching.
For finding maximum cost maximum matching
add -cost and return -matching()
1-indexed */
struct Hungarian {
  long long c[N][N], fx[N], fy[N], d[N];
  int 1[N], r[N], arg[N], trace[N];
  queue<int> q;
  int start, finish, n;
  const long long inf = 1e18;
  Hungarian() {}
  Hungarian (int n1, int n2): n(max(n1, n2)) {
    for (int i = 1; i <= n; ++i) {
      fy[i] = l[i] = r[i] = 0;
      for (int j = 1; j <= n; ++j) c[i][j] = inf; // make</pre>
          it 0 for maximum cost matching (not necessarily
          with max count of matching)
 void add_edge(int u, int v, long long cost) {
    c[u][v] = min(c[u][v], cost);
  inline long long getC(int u, int v) {
    return c[u][v] - fx[u] - fy[v];
  void initBFS() {
    while (!q.empty()) q.pop();
    q.push(start);
    for (int i = 0; i <= n; ++i) trace[i] = 0;</pre>
    for (int v = 1; v \le n; ++v) {
      d[v] = getC(start, v);
      arg[v] = start;
```

```
finish = 0:
void findAugPath() {
  while (!q.empty()) {
    int u = q.front();
    q.pop();
    for (int v = 1; v <= n; ++v) if (!trace[v]) {</pre>
        long long w = getC(u, v);
        if (!w) {
          trace[v] = u;
          if (!r[v]) {
            finish = v:
            return;
          q.push(r[v]);
        if (d[v] > w) {
          d[v] = w:
          arg[v] = u;
void subX_addY() {
  long long delta = inf;
  for (int v = 1; v \le n; ++v) if (trace[v] == 0 && d[v]
     | < delta) {</pre>
      delta = d[v];
  // Rotate
  fx[start] += delta;
  for (int v = 1; v <= n; ++v) if(trace[v]) {</pre>
      int u = r[v];
      fy[v] -= delta;
      fx[u] += delta;
    } else d[v] -= delta;
  for (int v = 1; v \le n; ++v) if (!trace[v] && !d[v])
      trace[v] = arg[v];
      if (!r[v]) {
        finish = v;
        return;
      q.push(r[v]);
void Enlarge() {
  do {
    int u = trace[finish];
    int nxt = l[u];
    l[u] = finish;
    r[finish] = u;
    finish = nxt;
  } while (finish);
```

```
long long maximum matching() {
    for (int u = 1; u <= n; ++u) {
      fx[u] = c[u][1];
      for (int v = 1; v <= n; ++v) {
        fx[u] = min(fx[u], c[u][v]);
    for (int v = 1; v \le n; ++v) {
      fy[v] = c[1][v] - fx[1];
      for (int u = 1; u <= n; ++u) {
        fy[v] = min(fy[v], c[u][v] - fx[u]);
    for (int u = 1; u <= n; ++u) {
      start = u;
      initBFS();
      while (!finish) {
        findAugPath();
        if (!finish) subX_addY();
      Enlarge();
    long long ans = 0;
    for (int i = 1; i <= n; ++i) {</pre>
      if (c[i][l[i]] != inf) ans += c[i][l[i]];
      else 1[i] = 0;
    return ans;
};
```

# 5 Geometria

### 5.1 Puntos

```
typedef long double lf;
const lf EPS = 1e-9;
const lf E0 = 0.0L; //Keep = 0 for integer coordinates,
   otherwise = EPS
const lf PI = acos(-1);
struct pt {
   lf x, y;
    pt(){}
    pt(lf a, lf b): x(a), y(b) {}
   pt(lf ang): x(cos(ang)), y(sin(ang)){} // Polar unit
        point: ang(RAD)
    pt operator - (const pt &q) const { return {x - q.x,
        y - q.y }; }
    pt operator + (const pt &q) const { return {x + q.x ,
        y + q.y \}; 
    pt operator * (pt p) { return {x * p.x - y * p.y, x *
       p.y + y * p.x;
```

```
pt operator * (const lf &t) const { return {x * t , y
    pt operator / (const lf &t) const { return {x / t , y
        / t }; }
    bool operator == (pt p) { return abs(x - p.x) <= EPS</pre>
       && abs(y - p.y) <= EPS; }
    bool operator != (pt p) { return !operator==(p); }
    bool operator < (const pt & q) const { // set / sort</pre>
        if(fabsl(x - q.x) > E0) return x < q.x;
        return y < q.\bar{y};
    void print() { cout << x << " " << y << "\n"; }</pre>
};
pt normalize(pt p) {
    lf norm = hypotl(p.x, p.y);
    if(fabsl(norm) > EPS) return {p.x /= norm, p.y /=
    else return p;
int cmp(lf a, lf b) { return (a + EPS < b ? -1 : (b + EPS <</pre>
    a ? 1 : 0)); } // float comparator
// rota ccw
pt rot90(pt p) { return {-p.v, p.x}; }
// w(RAD)
pt rot(pt p, lf w) { return {cosl(w) * p.x - sinl(w) * p.y
   *, sinl(w) * p.x + cosl(w) * p.y); }
lf norm2(pt p) { return p.x * p.x + p.y * p.y; }
lf norm(pt p) { return hypotl(p.x, p.y); }
lf dis2(pt p, pt q) { return norm2(p - q); }
lf dis(pt p, pt q) { return norm(p - q); }
If arg(pt a) {return atan2(a.y, a.x); } // ang(RAD) a x-
If dot(pt a, pt b) { return a.x * b.x + a.y * b.y; } //x
   = 90 -> cos = 0
If cross(pt a, pt b) { return a.x * b.y - a.y * b.x;  } //
   x = 180 -> \sin = 0
lf orient(pt a, pt b, pt c) { return cross(b - a, c - a);
   } // AB clockwise = -
int sign(lf x) { return (EPS < x) - (x < -EPS); }
// p inside angle abc (center in a)
bool in_angle(pt a, pt b, pt c, pt p) {
    //assert(fabsl(orient(a, b, c)) > E0);
    if(orient(a, b, c) < -E0)
        return orient(a, b, p) \geq= -E0 || orient(a, c, p)
    return orient(a, b, p) \geq -E0 && orient(a, c, p) \leq
       E0;
lf min_angle(pt a, pt b) { return acos(max((lf)-1.0, min((
   lf) 1.0, dot(a, b)/norm(a)/norm(b))); } // ang(RAD)
```

```
lf angle(pt a, pt b) { return atan2(cross(a, b), dot(a, b)
   ); \} // ang (RAD)
If angle(pt a, pt b, pt c) { // ang(RAD) AB AC ccw
    lf anq = angle(b - a, c - a);
    if (anq < 0) ang += 2 * PI;
    return ang;
bool half(pt p) { // true if is in (0, 180] (line is x
    // assert (p.x != 0 || p.y != 0); // the argument of
        (0, 0) is undefined
    return p.y > 0 || (p.y == 0 && p.x < 0);
bool half_from(pt p, pt v = \{1, 0\}) {
  return cross(v,p) < 0 \mid \mid (cross(v,p) == 0 \&\& dot(v,p) <
      0);
// polar sort
bool polar cmp(const pt &a, const pt &b) {
  return make tuple(half(a), 0) < make tuple(half(b),
     cross(a,b));
void polar sort(vector<pt> &v, pt o) { // sort points in
   counterclockwise with respect to point o
    sort(v.begin(), v.end(), [&](pt a,pt b) {
        return make tuple (half (a - o), 0.0, norm2 ((a - o)
           )) < make tuple(half(b - o), cross(a - o, b -
           o), norm2((b - o));
    });
int cuad(pt p) { // REVISAR
    if(p.x > 0 && p.y >= 0) return 0;
    if(p.x <= 0 && p.y > 0) return 1;
    if(p.x < 0 && p.y <= 0) return 2;
    if(p.x >= 0 && p.y < 0) return 3;
    return -1; // x == 0 \&\& v == 0
bool cmp (pt p1, pt p2) {
  int c1 = cuad(p1), c2 = cuad(p2);
  return c1 == c2 ? p1.y * p2.x < p1.x * p2.y : <math>c1 < c2;
```

#### 5.2 Lineas

```
// add points operators
struct line {
   pt v; lf c; // v: dir, c: mov y
   line(pt v, lf c) : v(v), c(c) {}
   line(lf a, lf b, lf c) : v({b, -a}), c(c) {} // ax +
```

```
bv = c
    line (pt p, pt q) : v(q - p), c(cross(v, p)) {}
    bool operator < (line 1) { return cross(v, 1.v) > 0; }
    bool operator == (line l) { return (abs(cross(v, l.v))
        = E0) && c == 1.c; } // abs(c) == abs(1.c)
    lf side(pt p) { return cross(v, p) - c; }
    lf dist(pt p) { return abs(side(p)) / norm(v); }
    lf dist2(pt p) { return side(p) * side(p) / (lf)norm2(
    line perp_through(pt p) { return {p, p + rot90(v)}; }
       // line perp to v passing through p
    bool cmp_proj(pt p, pt q) { return dot(v, p) < dot(v,</pre>
       q); } // order for points over the line
    // use: auto fsort = [&11] (const pt &a, const pt &b) {
        return 11.cmp proj(a, b); };
    line translate(pt t) { return {v, c + cross(v, t)}; }
    line shift_left(lf d) { return {v, c + d*norm(v)}; }
    pt proj(pt p) { return p - rot90(v) * side(p) / norm2(
       v); } // pt provected on the line
    pt refl(pt p) { return p - rot90(v) * 2 * side(p) /
       norm2(v); } // pt reflected on the other side of
       the line
    bool has (pt p) { return abs (cross (v, p) - c) <= E0; };
        // pt on line
    lf evalx(lf x){
        assert (fabsl (v.x) > EPS);
        return (c + v.v * x) / v.x;
};
pt inter ll(line l1, line l2) {
    if (abs(cross(11.v, 12.v)) <= EPS) return {INF, INF};</pre>
        // parallel
    return (12.v * 11.c - 11.v * 12.c) / cross(11.v, 12.v
       ); // floating points
// bisector divides the angle in 2 equal angles
// interior line goes on the same direction as 11 and 12
line bisector(line 11, line 12, bool interior) {
    // assert (cross(11.v, 12.v) != 0); // 11 and 12
       cannot be parallel
    lf sign = interior ? 1 : -1;
    return {12.v / norm(12.v) + 11.v / norm(11.v) * sign,
            12.c / norm(12.v) + 11.c / norm(11.v) * sign
                };
```

# 5.3 Poligonos

```
// add Points Lines Segments Circles
```

```
// points in polygon(vector<pt>) ccw or cw
enum {OUT, IN, ON};
lf area(vector<pt>& p) {
    lf r = 0.;
    for (int i = 0, n = p.size(); i < n; ++i) {
        r += cross(p[i], p[(i + 1) % n]);
    return r / 2; // negative if CW, positive if CCW
lf perimeter(vector<pt>& p) {
    lf per = 0;
    for (int i = 0, n = p.size(); i < n; ++i) {
        per += norm(p[i] - p[(i + 1) % n]);
    return per;
bool is convex(vector<pt>& p) {
    bool pos = 0, neq = 0;
    for (int i = 0, n = p.size(); i < n; i++) {
        int o = orient(p[i], p[(i + 1) % n], p[(i + 2) %
           n]);
        if (o > 0) pos = 1;
        if (o < 0) neq = 1;
    return ! (pos && neg);
int point in polygon(vector<pt>& pol, pt& p) {
    int wn = 0:
    for(int i = 0, n = pol.size(); i < n; ++i) {</pre>
        If c = orient(p, pol[i], pol[(i + 1) % n]);
        if(fabsl(c) <= E0 && dot(pol[i] - p, pol[(i + 1)</pre>
            % n] - p) <= E0) return ON; // on segment
        if(c > 0 && pol[i].y <= p.y + E0 && pol[(i + 1) %
            n].y - p.y > E0) ++wn;
        if(c < 0 \&\& pol[(i + 1) % n].y \le p.y + E0 \&\& pol
            [i].y - p.y > E0) --wn;
    return wn ? IN : OUT;
// O(logn) polygon CCW, remove collinear
int point in convex polygon(const vector<pt> &pol, const
   pt &p) {
        int low = 1, high = pol.size() -1;
        while(high - low > 1) {
                int mid = (low + high) / 2;
                if (orient (pol[0], pol[mid], p) \geq -E0)
                    low = mid;
                else high = mid;
        if (orient(pol[0], pol[low], p) < -E0) return OUT;</pre>
```

```
5.3 Poligonos
```

```
if (orient (pol[low], pol[high], p) < -E0) return</pre>
        if(orient(pol[high], pol[0], p) < -E0) return OUT</pre>
        if(low == 1 \&\& orient(pol[0], pol[low], p) <= E0)
             return ON:
        if (orient (pol[low], pol[high], p) <= E0) return</pre>
        if(high == (int) pol.size() -1 && orient(pol[high
            ], pol[0], p) <= E0) return ON;
        return IN;
// convex polygons in some order (CCW, CW)
vector<pt> minkowski(vector<pt> P, vector<pt> Q) {
        rotate(P.begin(), min element(P.begin(), P.end())
            , P.end());
        rotate(Q.begin(), min element(Q.begin(), Q.end())
            , 0.end());
        P.push\_back(P[0]), P.push\_back(P[1]);
        Q.push\_back(Q[0]), Q.push\_back(Q[1]);
        vector<pt> ans;
        size t i = 0, j = 0;
        while(i < P.size() - 2 || j < Q.size() - 2){</pre>
                 ans.push_back(P[i] + Q[j]);
                 lf d\bar{t} = cross(P[i + 1] - P[i], Q[j + 1] -
                     Q[j]);
                 if(dt \ge E0 \&\& i < P.size() - 2) ++i;
                 if (dt \leq E0 && i \leq O.size() - 2) ++ i;
        return ans;
pt centroid(vector<pt>& p) {
    pt c{0, 0};
    If scale = 6. * area(p);
    for (int i = 0, n = p.size(); i < n; ++i) {
        c = c + (p[i] + p[(i + 1) % n]) * cross(p[i], p[(i + 1) % n])
           i + 1) % n]);
    return c / scale;
void normalize(vector<pt>& p) { // polygon CCW
    int bottom = min_element(p.begin(), p.end()) - p.
       begin();
    vector<pt> tmp(p.begin() + bottom, p.end());
    tmp.insert(tmp.end(), p.begin(), p.begin()+bottom);
    p.swap(tmp);
    bottom = 0;
void remove col(vector<pt>& p) {
    vector<pt> s;
    for (int i = 0, n = p.size(); i < n; i++) {
```

```
if (!on_segment(p[(i-1+n) % n], p[(i+1) % n]
           ], p[i])) s.push_back(p[i]);
    p.swap(s);
void delete repetead(vector<pt>& p) {
    vector<pt> aux;
    sort(p.begin(), p.end());
    for (pt &pi : p) {
        if (aux.empty() || aux.back() != pi) aux.
           push back(pi);
    p.swap(aux);
pt farthest (vector<pt>& p, pt v) { // O(log(n)) only
   CONVEX, v: dir
    int n = p.size();
    if(n < 10) {
        int k = 0;
        for (int i = 1; i < n; i++) if (dot (v, (p[i] - p[k
           |)) > EPS) k = i;
        return p[k];
    pt a = p[1] - p[0];
    int s = 0, e = n, ua = dot(v, a) > EPS;
    if(!ua && dot(v, (p[n-1] - p[0])) <= EPS) return p
       [0];
    while (1) {
        int m = (s + e) / 2;
        pt c = p[(m + 1) % n] - p[m];
        int uc = dot(v, c) > EPS;
        if(!uc && dot(v, (p[(m-1+n) % n] - p[m])) <=
           EPS) return p[m];
        if (ua && (!uc || dot(v, (p[s] - p[m])) > EPS)) e
        else if (ua | | uc | | dot (v, (p[s] - p[m])) >= -EPS
           ) s = m, a = c, ua = uc;
        else e = m;
        assert (e > s + 1);
vector<pt> cut(vector<pt>& p, line l) {
    // cut CONVEX polygon by line 1
    // returns part at left of l.pg
    vector<pt> q;
    for(int i = 0, n = p.size(); i < n; i++) {</pre>
        int d0 = sign(l.side(p[i]));
        int d1 = sign(l.side(p[(i + 1) % n]));
        if(d0 >= 0) q.push back(p[i]);
        line m(p[i], p[(i + 1) % n]);
        if(d0 * d1 < \overline{0} \&\& !(abs(cross(l.v, m.v)) <= EPS))
```

```
5.4 Circulos
```

```
q.push back((inter ll(l, m)));
           return q;
// O(n)
vector<pair<int, int>> antipodal(vector<pt>& p) {
           vector<pair<int, int>> ans;
           int n = p.size();
           if (n == 2) ans.push back(\{0, 1\});
           if (n < 3) return ans;</pre>
           auto nxt = [\&] (int x) \{ return (x + 1 == n ? 0 : x +
                   1); };
           auto area2 = [&](pt a, pt b, pt c){ return cross(b -
                   a, c - a); };
           int b0 = 0;
           while (abs(area2(p[n - 1], p[0], p[nxt(b0)])) > abs(
                    area2(p[n - 1], p[0], p[b0]))) ++b0;
           for (int b = b0, a = 0; b != 0 && a <= b0; ++a) {
                     ans.push_back({a, b});
                     while (abs(area2(p[a], p[nxt(a)], p[nxt(b)])) >
                              abs(area2(p[a], p[nxt(a)], p[b]))){
                               b = nxt(b);
                                if (a != b0 || b != 0) ans.push_back({a, b});
                                else return ans:
                     if (abs(area2(p[a], p[nxt(a)], p[nxt(b)])) == abs
                               (area2(p[a], p[nxt(a)], p[b]))){
                                if (a != b0 || b != n - 1) ans.push_back({a,
                                         nxt(b) });
                                else ans.push_back({nxt(a), b});
           return ans;
// O(n)
// square distance of most distant points, prereq: convex
         , ccw, NO COLLINEAR POINTS
lf callipers(vector<pt>& p) {
           int n = p.size();
           lf r = 0;
           for (int i = 0, j = n < 2 ? 0 : 1; <math>i < j; ++i) {
                     for(;;\dot{j} = (\dot{j} + 1) \% n) \{
                                r = \max(r, \text{ norm2}(p[i] - p[j]));
                                if (cross((p[(i + 1) % n] - p[i]), (p[(j + 1)
                                         % n] - p[j])) <= EPS) break;
           return r;
// O(n + m) max dist between 2 points (pa, pb) of 2
         Convex polygons (a, b)
```

```
lf rotating_callipers(vector<pt>& a, vector<pt>& b){ //
   REVISAR
    if (a.size() > b.size()) swap(a, b); // <- del or add
    pair<11, int> start = \{-1, -1\};
    if(a.size() == 1) swap(a, b);
    for (int i = 0; i < a.size(); i++) start = max(start,
        \{norm2(b[0] - a[i]), i\});
    if(b.size() == 1) return start.first;
    lf r = 0:
    for(int i = 0, j = start.second; i < b.size(); ++i){</pre>
        for(;; j = (j + 1) % a.size()){
            r = max(r, norm2(b[i] - a[j]));
            if(cross((b[(i + 1) % b.size()) - b[i]), (a[(
                j + 1) % a.size()] - a[j])) <= EPS) break;
    return r;
lf intercircle(vector<pt>& p, circle c) { // area of
   intersection with circle
    lf r=0.;
    for (int i = 0, n = p.size(); i < n; i++) {
        int j = (i + 1) % n;
        If w = intertriangle(c, p[i], p[j]);
        if(cross((p[j] - c.center), (p[i] - c.center)) >
           0) r += w;
        else r -= w;
    return abs(r);
ll pick(vector<pt>& p) {
    ll boundary = 0;
    for (int i = 0, n = p.size(); i < n; i++) {</pre>
        int j = (i + 1 == n ? 0 : i + 1);
        boundary += gcd((ll)abs(p[i].x - p[j].x), (ll)
           abs(p[i].y - p[j].y));
    return abs(area(p)) + 1 - boundary / 2;
```

## 5.4 Circulos

```
// add Lines Points
enum {OUT, IN, ON};
struct circle {
   pt center; lf r;
   // (x - xo)^2 + (y - yo)^2 = r^2
   circle(pt c, lf r): center(c), r(r){};
   // circle that passes through abc
```

```
5.4 Circulos
```

```
circle(pt a, pt b, pt c) {
        b = b - a, c = c - a;
        assert(cross(b, c) != 0); // no circumcircle if A
           , B, C aligned
        pt cen = a + rot 90 (b * norm2 (c) - c * norm2 (b)) /
            cross(b, c) / 2;
        center = cen;
        r = norm(a - cen);
    // diameter = segment pg
    circle(pt p, pt q) {
        center = (p + q) * 0.5L;
        r = dis(p, q) * 0.5L;
    int contains(pt &p) {
        lf det = r * r - dis2(center, p);
        if(fabsl(det) <= EPS) return ON;</pre>
        return (det > EPS ? IN : OUT);
    bool in(circle c) { return norm(center - c.center) + r
        <= c.r + EPS; } // non strict
};
// centers of the circles that pass through ab and has
   radius r
vector<pt> centers(pt a, pt b, lf r) {
    if (norm(a - b) > 2 * r + EPS) return {};
    pt m = (a + b) / 2;
    double f = sqrt(r * r / norm2(a - m) - 1);
    pt c = rot 90 (a - m) * f;
    return {m - c, m + c};
vector<pt> inter_cl(circle c, line l){
        vector<pt> s;
        pt p = l.proj(c.center);
        If \bar{d} = norm(\bar{p} - c.center);
        if(d - EPS > c.r) return s;
        if(abs(d - c.r) <= EPS) { s.push_back(p); return s</pre>
        d=sqrt(c.r * c.r - d * d);
        s.push back(p + normalize(l.v) * d);
        s.push_back(p - normalize(l.v) \star d);
        return s;
vector<pt> inter cc(circle c1, circle c2) {
    pt dir = c2.center - c1.center;
    1f d2 = dis2(c1.center, c2.center);
    if(d2 <= E0) {
        //assert(fabsl(c1.r - c2.r) > E0);
        return {};
```

```
lf td = 0.5L * (d2 + c1.r * c1.r - c2.r * c2.r);
    1f h2 = c1.r * c1.r - td / d2 * td;
    pt p = c1.center + dir * (td / d2);
    if(fabsl( h2 ) < EPS) return {p};</pre>
    if(h2 < 0.0L) return {};
    pt dir h = rot 90 (dir) * sqrtl(h2 / d2);
    return {p + dir h, p - dir h};
// circle-line inter = 1, inner: 1 = oxo 0 = o=o
vector<pair<pt, pt>> tangents(circle c1, circle c2, bool
   inner) {
    vector<pair<pt, pt>> out;
    if (inner) c2.r = -c2.r; // inner tangent
    pt d = c2.center - c1.center;
    double dr = c1.r - c2.r, d2 = norm2(d), h2 = d2 - dr
    if (d2 == 0 || h2 < 0) { assert(h2 != 0); return {};
       } // (identical)
    for (double s : {-1, 1}) {
        pt v = (d * dr + rot 90(d) * sqrt(h2) * s) / d2;
        out.push back({c1.center + v * c1.r, c2.center +
           \overline{v} * c2.r);
    return out; // if size 1: circle are tangent
// circle targent passing through pt p
pair<pt, pt> tangent through pt(circle c, pt p){
    pair<pt, pt> out;
    double d = norm2(p - c.center);
    if (d < c.r) return {};
    pt base = c.center - p;
    double w = sgrt(norm2(base) - c.r * c.r);
    pt a = \{w, c.r\}, b = \{w, -c.r\};
   pt s = p + base * a / norm2(base) * w;
   pt t = p + base * b / norm2(base) * w;
    out = \{s, t\};
    return out;
lf safeAcos(lf x) {
    if (x < -1.0) x = -1.0;
    if (x > 1.0) x = 1.0;
    return acos(x);
lf areaOfIntersectionOfTwoCircles(circle c1, circle c2){
    lf r1 = c1.r, r2 = c2.r, d = dis(c1.center, c2.center)
    if(d >= r1 + r2) return 0.0L;
    if(d <= fabsl(r2 - r1)) return PI * (r1 < r2 ? r1 *</pre>
       r1 : r2 * r2);
```

```
lf alpha = safeAcos((r1 * r1 - r2 * r2 + d * d) /
        (2.0L * d * r1));
    lf betha = safeAcos((r2 * r2 - r1 * r1 + d * d) /
        (2.0L * d * r2));
    lf a1 = r1 * r1 * (alpha - sinl(alpha) * cosl(alpha))
    1f a2 = r2 * r2 * (betha - sinl(betha) * cosl(betha))
    return a1 + a2;
};
lf intertriangle(circle& c, pt a, pt b){ // area of
   intersection with oab
    if(abs(cross((c.center - a), (c.center - b))) <= EPS)</pre>
        return 0.;
    vector<pt> q = \{a\}, w = inter cl(c, line(a, b));
    if(w.size() == 2) for(auto p: w) if(dot((a - p), (b -
        p)) < -EPS) q.push back(p);
    q.push back(b);
    if(q.size() == 4 \&\& dot((q[0] - q[1]), (q[2] - q[1]))
        > EPS) swap(q[1], q[2]);
    lf s = 0;
    for(int i = 0; i < q.size() - 1; ++i){}
        if(!c.contains(q[i]) \mid | !c.contains(q[i + 1])) s
           += c.r * c.r * min_angle((q[i] - c.center), q[
           i+11 - c.center) / 2;
        else s += abs(cross((q[i] - c.center), (q[i + 1]
           - c.center)) / 2);
    return s:
bool circumcircle contains(vector<pt> tr, pt D) { //
   triange CCW
  pt A = tr[0] - D, B = tr[1] - D, C = tr[2] - D;
 lf norm a = norm2(tr[0]) - norm2(D);
  lf norm b = norm2(tr[1]) - norm2(D);
 lf norm_c = norm2(tr[2]) - norm2(D);
  lf det1 = A.x * (B.y * norm_c - norm_b * C.y);
 lf det2 = B.x * (C.\bar{y} * norm_a - norm_c * A.y);
  If det3 = C.x * (A.y * norm b - norm a * B.y);
  return det1 + det2 + det3 > E0;
// r[k]: area covered by at least k circles
// O(n^2 \log n) (high constant)
vector<lf> intercircles(vector<circle> c) {
        vector<lf> r(c.size() + 1);
        for(int i = 0; i < c.size(); ++i){</pre>
                int k = 1; pt 0 = c[i].center;
                vector<pair<pt, int>> p = {
                         \{c[i].center + pt(1,0) * c[i].r,
                         \{c[i].center - pt(1,0) * c[i].r,
                            0 } };
```

```
for(int j = 0; j < c.size(); ++j) if(j !=
            i){
                bool b0 = c[i].in(c[j]), b1 = c[j]
                    ].in(c[i]);
                if(b0 && (!b1 || i < j)) ++k;
                else if(!b0 && !b1){
                         auto v = inter_cc(c[i], c
                            [j]);
                        if(v.size() == 2){
            swap(v[0], v[1]);
                                 p.push back({v
                                    [0], 1});
            p.push back(\{v[1], -1\});
                                 if(polar cmp(v[1]
                                      -0, v[0] - 0
                                    )) ++k;
        sort(all(p), [&](auto& a, auto& b) {
           return polar_cmp(a.first - 0, b.first
           - 0); });
        for(int j = 0; j < p.size(); ++j){
                pt p0 = p[j ? j - 1 : p.size()
                    -1].first, p1 = p[j].first;
                If a = \min \text{ angle}((p0 - c[i]).
                    center), (p1 - c[i].center));
                r[k] += (p0.x - p1.x) * (p0.y +
                    p1.y) / 2 + c[i].r * c[i].r *
                    (a - \sin(a)) / 2;
                k += p[j].second;
return r;
```

# 5.5 Semiplanos

```
const lf INF = 1e100;
struct Halfplane {
   pt p, pq; // p: point on line, pq: dir, take left
   lf angle;
   Halfplane() {}
   Halfplane(pt& a, pt& b): p(a), pq(b - a) {
        angle = atan2l(pq.y, pq.x);
   }

  bool out(const pt& r) { return cross(pq, r - p) < -EPS
      ;} // checks if p is inside the half plane
  bool operator < (const Halfplane& e) const { return
      angle < e.angle; }</pre>
```

```
};
// intersection pt of the lines of 2 halfplanes
pt inter(const Halfplane& s, const Halfplane& t) {
    if (abs(cross(s.pq, t.pq)) <= EPS) return {INF, INF};</pre>
    lf alpha = cross((t.p - s.p), t.pq) / cross(s.pq, t.
       pq);
    return s.p + (s.pq * alpha);
// O(nlogn) return CCW polygon
vector<pt> hp intersect(vector<Halfplane>& H) {
    pt box[4] = \{pt(INF, INF), pt(-INF, INF), pt(-INF, -
       INF), pt(INF, -INF)};
    for (int i = 0; i < 4; ++i) {
        Halfplane aux(box[i], box[(i + 1) % 4]);
        H.push back(aux);
    sort(H.begin(), H.end());
    deque < Halfplane > dq;
    int len = 0;
    for(int i = 0; i < int(H.size()); ++i){</pre>
        while (len > 1 && H[i].out(inter(dq[len - 1], dq[
           len - 2]))){}
            dq.pop_back();
            --len;
        while (len > 1 && H[i].out(inter(dq[0], dq[1]))){
            dq.pop_front();
            --len;
        if (len > 0 && fabsl(cross(H[i].pq, dq[len - 1].
           pq)) < EPS){
            if (dot(H[i].pq, dq[len - 1].pq) < 0.0)
                return vector<pt>();
            if (H[i].out(dq[len - 1].p)){
                dq.pop_back();
                --len:
            } else continue;
        dq.push_back(H[i]);
        ++len;
    while (len > 2 && dq[0].out(inter(dq[len - 1], dq[len
        - 2]))){
        dq.pop_back();
        --len;
    while (len > 2 && dq[len - 1].out(inter(dq[0], dq[1])
       )){
```

```
dq.pop front();
        --len;
    if (len < 3) return vector<pt>();
    vector<pt> ret(len);
    for(int i = 0; i + 1 < len; ++i) ret[i] = inter(dq[i</pre>
        ], da[i + 1]);
    ret.back() = inter(dg[len - 1], dg[0]);
    // remove repeated points if needed
    return ret;
// intersection of halfplanes
vector<pt> hp intersect(vector<halfplane>& b) {
    vector<pt> box = {\{\inf, \inf\}, \{-\inf, \inf\}, \{-\inf, -1\}\}
       inf}, {inf, -inf}};
    for(int i = 0; i < 4; i++) {</pre>
        b.push back(\{box[i], box[(i + 1) % 4]\});
    sort(b.begin(), b.end());
    int n = b.size(), q = 1, h = 0;
    vector<halfplane> \bar{c}(n + 10);
    for(int i = 0; i < n; i++) {</pre>
        while (q < h \& \& b[i].out(inter(c[h], c[h - 1]))) h
        while (q < h \& \& b[i].out(inter(c[q], c[q + 1]))) q
        c[++h] = b[i];
        if(q < h \&\& abs(cross(c[h].pq, c[h-1].pq)) < EPS)
            if(dot(c[h].pq, c[h - 1].pq) <= 0) return {};
            if(b[i].out(c[h].p)) c[h] = b[i];
    while (q < h - 1 \& \& c[q].out(inter(c[h], c[h - 1]))) h
    while (q < h - 1 \& c[h].out(inter(c[q], c[q + 1]))) q
    if(h - q <= 1) return {};
    c[h + 1] = c[q];
    vector<pt> s;
    for(int i = q; i < h + 1; i++) s.pb(inter(c[i], c[i +
    return s;
```

# 5.6 Segmentos

// add Lines Points

```
bool in_disk(pt a, pt b, pt p) { // pt p inside ab disk
    return dot (a - p, b - p) <= E0;
bool on_segment(pt a, pt b, pt p) { // p on ab
    return orient(a, b, p) == 0 && in_disk(a, b, p);
// ab crossing cd
bool proper_inter(pt a, pt b, pt c, pt d, pt& out) {
    lf oa = orient(c, d, a);
   lf ob = orient(c, d, b);
   lf oc = orient(a, b, c);
    lf od = orient(a, b, d);
    // Proper intersection exists iff opposite signs
    if (oa * ob < 0 && oc * od < 0) {
        out = (a * ob - b * oa) / (ob - oa);
        return true;
    return false;
// intersection bwn segments
set<pt> inter_ss(pt a, pt b, pt c, pt d) {
    pt out;
    if (proper_inter(a, b, c, d, out)) return {out}; //
       if cross -> 1
    set<pt> s;
    if (on segment(c, d, a)) s.insert(a); // a in cd
    if (on_segment(c, d, b)) s.insert(b); // b in cd
    if (on segment(a, b, c)) s.insert(c); // c in ab
    if (on_segment(a, b, d)) s.insert(d); // d in ab
    return s; // 0, 2
lf pt_to_seg(pt a, pt b, pt p) { // p to ab
    if (a != b) {
        line l(a, b);
        if (l.cmp_proj(a, p) && l.cmp_proj(p, b)) // if
           closest to projection = (a, p, b)
            return 1.dist(p); // output distance to line
    return min(norm(p - a), norm(p - b)); // otherwise
       distance to A or B
lf seg_to_seg(pt a, pt b, pt c, pt d) {
    pt dummy;
    if (proper_inter(a, b, c, d, dummy)) return 0; // ab
       intersects cd
    return min({pt_to_seg(a, b, c), pt_to_seg(a, b, d),
       pt_to_seg(c, d, a), pt_to_seg(c, d, b)}); // try
       the 4 pts
int length union(vector<pt>& a) { // REVISAR
    int n = a.size();
```

```
vector<pair<int, bool>> x(n * 2);
for (int i = 0; i < n; i++) {
    x[i * 2] = {a[i].x, false};
    x[i * 2 + 1] = {a[i].y, true};
}
sort(x.begin(), x.end());
int result = 0;
int c = 0;
for (int i = 0; i < n * 2; i++) {
    if (i > 0 && x[i].first > x[i - 1].first && c >
        0) result += x[i].first - x[i - 1].first;
    if (x[i].second) c--;
    else c++;
}
return result;
}
```

### 5.7 Convex Hull

```
// CCW order
// if colineal are needed, use > in orient and remove
   repeated points
vector<pt> chull(vector<pt>& p) {
        if(p.size() < 3) return p;</pre>
        vector<pt> r; //r.reserve(p.size());
        sort(p.begin(), p.end()); // first x, then v
        for(int i = 0; i < p.size(); i++){ // lower hull</pre>
                while (r.size() >= 2 \&\& orient(r[r.size()
                    -2], p[i], r.back()) >= 0) r.pop back
                    ();
                r.pb(p[i]);
        r.pop back();
        int k = r.size();
        for (int i = p.size() - 1; i >= 0; --i) { // upper }
           hu11
                while (r.size() >= k + 2 \&\& orient(r[r.
                    size() - 2], p[i], r.back()) >= 0) r.
                    pop back();
                r.pb(p[i]);
        r.pop back();
        return r;
```

#### 5.8 Closest Points

```
// O(nlogn)
pair<pt, pt> closest_points(vector<pt> v) {
```

```
sort(v.begin(), v.end());
pair<pt, pt> ans;
1f d2 = INF;
function<void( int, int )> solve = [&](int l, int r)
    if(l == r) return;
    int mid = (1 + r) / 2;
    lf x_mid = v[mid].x;
    solve(l, mid);
    solve(mid + 1, r);
    vector<pt> aux;
    int p1 = 1, p2 = mid + 1;
    while (p1 <= mid && p2 <= r) {</pre>
        if (v[p1].y < v[p2].y) aux.push_back (v[p1++]);
        else aux.push_back(v[p2++]);
    while(p1 <= mid) aux.push back(v[p1++]);</pre>
    while(p2 <= r) aux.push_back(v[p2++]);</pre>
    vector<pt> nb;
    for(int i = 1; i <= r; ++i) {
    v[i] = aux[i - 1];
    lf dx = (x_mid - v[i].x);
    if(dx * dx < d2)
        nb.push_back(v[i]);
    for(int i = 0; i < (int) nb.size(); ++i){</pre>
    for (int k = i + 1; k < (int) nb.size(); ++k){}
        lf dy = (nb[k].y - nb[i].y);
        if (dy * dy > d2) break;
        lf nd2 = dis2(nb[i], nb[k]);
        if(nd2 < d2) d2 = nd2, ans = {nb[i], nb[k]};
};
solve(0, v.size() -1);
return ans;
```

## 5.9 Min Circle

```
return ans;
};

auto f1 = [&]( int a ){
    Circle ans(v[a], 0.0L);
    for(int i = 0; i < a; ++i)
        if(ans.contains(v[i]) == OUT) ans = f2( i, a );
    return ans;
};

Circle ans( v[0], 0.0L );
for(int i = 1; i < (int) v.size(); ++i)
        if(ans.contains(v[i]) == OUT ) ans = f1(i);

return ans;
}</pre>
```

#### 5.10 3D

```
typedef double lf;
struct p3 {
    lf x, y, z;
        p3(){}
        p3(1f x, 1f y, 1f z): x(x), y(y), z(z) {}
    p3 operator + (p3 p) { return \{x + p.x, y + p.y, z + p\}
    p3 operator - (p3 p) { return \{x - p.x, y - p.y, z - p\}
        .z}; }
    p3 operator * (lf d) { return {x * d, y * d, z * d}; }
    p3 operator / (lf d) { return {x / d, y / d, z / d}; }
        // only for floating point
    // Some comparators
    bool operator == (p3 p) { return tie(x, y, z) == tie(p
        .x, p.y, p.z); }
    bool operator != (p3 p) { return !operator == (p); }
        void print() { cout << x << " " << y << " " << z</pre>
            << "\n"; }
        // scale: (newnorm / norm) * p3
};
lf dot(p3 v, p3 w) { return v.x * w.x + v.y * w.y + v.z *
   w.z; }
p3 cross(p3 v, p3 w) {
    return { v.y * w.z - v.z * w.y, v.z * w.x - v.x * w.z
       , v.x * w.y - v.y * w.x };
lf norm2(p3 v) { return dot(v, v); }
lf norm(p3 v) { return sqrt(norm2(v)); }
p3 unit(p3 v) { return v / norm(v); }
// ang(RAD)
double angle(p3 v, p3 w) {
    double cos_theta = dot(v, w) / norm(v) / norm(w);
    return acos (max (-1.0, min (1.0, cos_theta)));
```

```
// orient s, pqr form a triangle pos: 'up', zero = on,
   neg = 'dow'
lf orient(p3 p, p3 q, p3 r, p3 s){
        return dot(cross((q - p), (r - p)), (s - p));
// same as 2D but in n-normal direction
lf orient_by_normal(p3 p, p3 q, p3 r, p3 n) {
        return dot(cross((q - p), (r - p)), n);
struct plane {
    p3 n; lf d; // n: normal d: dist to zero
    // From normal n and offset d
    plane(p3 n, lf d): n(n), d(d) {}
    // From normal n and point P
    plane(p3 n, p3 p): n(n), d(dot(n, p)) {}
    // From three non-collinear points P,Q,R
    plane (p3 p, p3 q, p3 r): plane (cross ((q - p), (r - p)
       ), p){}
    // - these work with lf = int
    lf side(p3 p) { return dot(n, p) - d; }
    double dist(p3 p) { return abs(side(p)) / norm(n); }
    plane translate(p3 t) {return {n, d + dot(n, t)}; }
    /// - these require If = double
    plane shift up(double dist) { return {n, d + dist *
       norm(n) }; }
    p3 proj(p3 p) { return p - n * side(p) / norm2(n); }
    p3 refl(p3 p) \{ return p - n * 2 * side(p) / norm2(n); 
};
struct line3d {
        p3 d, o; // d: dir o: point on line
        // From two points P, Q
        line3d(p3 p, p3 q): d(q - p), o(p){}
        // From two planes p1, p2 (requires lf = double)
        line3d(plane p1, plane p2) {
                d = cross(p1.n, p2.n);
                o = cross((p2.n * p1.d - p1.n * p2.d), d)
                     / norm2(d);
        // - these work with lf = int
        double dist2(p3 p) { return norm2(cross(d, (p - o)
           )) / norm2(d); }
        double dist(p3 p) { return sqrt(dist2(p)); }
        bool cmp_proj(p3 p, p3 q) { return dot(d, p) < dot
            (d, q); }
        // - these require lf = double
        p3 proj(p3 p) { return o + d * dot(d, (p - o)) /
           norm2(d); }
        p3 refl(p3 p) { return proj(p) * 2 - p; }
        p3 inter(plane p) { return o - d * p.side(o) / dot
            (p.n, d); }
        // get other point: pl.o + pl.d * t;
```

```
} ;
double dist(line3d 11, line3d 12) {
        p3 n = cross(11.d, 12.d);
        if(n == p3(0, 0, 0)) return 11.dist(12.o); //
        return abs(dot((12.o - 11.o), n)) / norm(n);
// closest point on 11 to 12
p3 closest_on_line1(line3d l1, line3d l2) {
        p3 n2 = cross(12.d, cross(11.d, 12.d));
        return 11.0 + 11.d * (dot((12.0 - 11.0), n2)) /
           dot(11.d, n2);
double small_angle(p3 v, p3 w) { return acos(min(abs(dot(v
   (w) / norm(v) / norm(w), 1.0); } // 0.90
double angle(plane p1, plane p2) { return small_angle(p1.n
   , p2.n); }
bool is_parallel(plane p1, plane p2) { return cross(p1.n,
   p2.n == p3(0, 0, 0); }
bool is perpendicular (plane p1, plane p2) { return dot(p1.
   n, p2.n) == 0;}
double angle(line3d 11, line3d 12) { return small_angle(l1
   .d, 12.d); }
bool is_parallel(line3d 11, line3d 12) { return cross(11.d
   , 12.d) == p3(0, 0, 0); 
bool is perpendicular(line3d 11, line3d 12) { return dot(
   11.d, 12.d) == 0; }
double angle(plane p, line3d l) { return M_PI / 2 -
   small angle(p.n, l.d); }
bool is_parallel(plane p, line3d l) { return dot(p.n, l.d)
bool is_perpendicular(plane p, line3d l) { return cross(p.
   n, 1.d) == p3(0, 0, 0);
line3d perp_through(plane p, p3 o) { return line3d(o, o +
   p.n); }
plane perp through (line3d 1, p3 o) { return plane (l.d, o);
```

#### 5.11 KD Tree

```
// given a set of points, answer queries of nearest point
    in O(log(n))
bool onx(pt a, pt b) {return a.x < b.x;}
bool ony(pt a, pt b) {return a.y < b.y;}</pre>
struct Node {
        pt pp;
        If x0 = \inf, x1 = -\inf, y0 = \inf, y1 = -\inf;
        Node *first = 0, *second = 0;
        11 distance(pt p) {
                11 x = min(max(x0, p.x), x1);
                11 y = min(max(y0, p.y), y1);
```

```
return norm2 (pt (x, y) - p);
        Node(vector<pt>&& vp) : pp(vp[0]) {
                 for(pt p : vp) {
                         x0 = \min(x0, p.x);
            x1 = max(x1, p.x);
                         y0 = min(y0, p.y);
            y1 = max(y1, p.y);
                 if(vp.size() > 1) {
                         sort(all(vp), x1 - x0 >= y1 - y0
                             ? onx : ony);
                         int m = vp.size() / 2;
                         first = new Node({vp.begin(), vp.
                             begin() + m});
                         second = new Node({vp.begin() + m
                             , vp.end()});
};
struct KDTree {
        Node* root;
        KDTree(const vector<pt>& vp): root(new Node({all(
            ({ (qv
        pair<11, pt> search(pt p, Node *node) {
                 if(!node->first){
                         // avoid query point as answer
                         // if(p.x == node -> pp.x && p.y ==
                              node->pp.y) return {inf, pt()
                         return {norm2 (p-node->pp), node->
                             ; { qq
                 Node *f = node \rightarrow first, *s = node \rightarrow second;
                 ll bf = f->distance(p), bs = s ->
                    distance(p);
                 if(bf > bs) swap(bf, bs), swap(f, s);
                 auto best = search(p, f);
                 if(bs < best.ff) best = min(best, search(</pre>
                    p, s));
                 return best;
        pair<11, pt> nearest(pt p) { return search(p, root
           ); }
};
```

# 6 Grafos

### 6.1 Puentes

```
// O(n+m)
vector<bool> visited;
```

```
vi tin, low;
int timer;
void IS BRIDGE(int u, int v, vii &puentes) {
    puentes.push_back({min(u, v), max(u, v)});
void dfs(vector<vi> &adj, vii &puentes, int v, int p =
   -1) {
    visited[v] = true;
    tin[v] = low[v] = timer++;
    for (int to : adj[v]) {
        if (to == p) continue;
        if (visited[to]) {
            low[v] = min(low[v], tin[to]);
        } else ·
            dfs(adj, puentes, to, v);
            low[v] = min(low[v], low[to]);
            if (low[to] > tin[v])
                IS_BRIDGE(v, to, puentes);
void find_bridges(vector<vi> &adj, vii &puentes, int n) {
    timer = 0;
    visited.assign(n, false);
    tin.assign(n, -1);
    low.assign(n, -1);
    for (int i = 0; i < n; ++i) {
        if (!visited[i])
            dfs(adj, puentes, i);
```

### 6.2 Puntos de Articulación

```
// O(n+m)
int n;
vector<vector<int>> adj;

vector<bool> visited;
vector<int> tin, low;
int timer;

void dfs(int v, int p = -1) {
    visited[v] = true;
    tin[v] = low[v] = timer++;
    int children=0;
    for (int to: adj[v]) {
        if (to == p) continue;
        if (visited[to]) {
            low[v] = min(low[v], tin[to]);
        } else {
```

# 6.3 Kosajaru

```
//Encontrar las componentes fuertemente conexas en un
   grafo dirigido
//Componente fuertemente conexa: es un grupo de nodos en
   el que hay
//un camino dirigido desde cualquier nodo hasta cualquier
    otro nodo dentro del grupo.
const int maxn = 1e5+5;
vi adj rev[maxn],adj[maxn];
bool used[maxn];
vi order, comp;
// O(n+m)
void dfs1(int v) {
        used[v]=true;
        for(int u:adj[v])
                 if(!used[u])dfs1(u);
        order.push back(v);
void dfs2(int v){
        used[v]=true;
        comp.push back(v);
        for(int u:adj_rev[v])
                 if(!used[u])dfs2(u);
void init(int n){
        for (int i=0; i < n; ++i) if (!used[i]) dfs1(i);</pre>
        for (int i=0; i<n; ++i) used[i]=false;</pre>
        reverse (order.begin(), order.end());
        for(int v:order){
```

## 6.4 Tarjan

```
// O(n+m) (?)
vi low, num, comp, q[nax];
int scc, timer;
stack<int> st;
void t jn (int u) {
  low[u] = num[u] = timer++; st.push(u); int v;
  for(int v: q[u]) {
    if (num[v] == -1) tin(v);
    if (comp[v] == -1) low[u] = min(low[u], low[v]);
  if(low[u] == num[u]) {
    do\{ v = st.top(); st.pop(); comp[v]=scc;
    \} while (u != \overline{v});
    ++scc;
void callt(int n) {
 timer = scc = 0;
  num = low = comp = vector\langle int \rangle (n, -1);
  for (int i = 0; i < n; i++) if (num[i] = -1) t jn(i);
```

# 6.5 Dijkstra

```
return dist;
}
```

### 6.6 Bellman Ford

```
// O(V*E)
vi bellman_ford(vector<vii> &adj, int s, int n) {
    vi dist(n, INF); dist[s] = 0;
    for (int i = 0; i<n-1; i++) {
        bool modified = false;
        for (int u = 0; u < n; u++)
            if (dist[u] != INF)
                for (auto &[v, w] : adj[u]){
                     if (dist[v] <= dist[u] + w) continue;</pre>
                     dist[v] = dist[u] + w;
                     modified = true;
        if (!modified) break;
    bool negativeCicle = false;
    for (int u = 0; u < n; u + +)
        if (dist[u] != INF)
            for (auto &[v, w] : adj[u]) {
                if (dist[v] > dist[u] + w) negativeCicle
                    = true;
    return dist;
```

# 6.7 Floyd Warshall

# 6.8 MST Kruskal

### 6.9 MST Prim

```
// O(E * log V)
vector<vii> adi;
vi tomado;
priority_queue<ii> pq;
void process(int u) {
    tomado[u] = 1;
    for (auto &[v, w] : adj[u]) {
        if (!tomado[v]) pq.emplace(-w, -v);
int prim(int v, int n) {
    tomado.assign(n, 0);
    process(0);
    int mst costo = 0, tomados = 0;
    while (!pq.empty()) {
        auto [w, u] = pq.top(); pq.pop();
w = -w; u = -u;
        if (tomado[u]) continue;
        mst_costo += w;
        process(u);
        tomados++;
        if (tomados == n-1) break;
    return mst_costo;
```

# 6.10 Shortest Path Faster Algorithm

```
//Algoritmo mas rapido de ruta minima
//O(V*E) peor caso, O(E) en promedio.
bool spfa(vector<vii> &adj, vector<int> &d, int s, int n)
{
   d.assign(n, INF);
```

```
vector<int> cnt(n, 0);
vector<bool> inqueue(n, false);
queue<int> q;
d[s] = 0;
q.push(s);
inqueue[s] = true;
while (!q.empty())
    int v = q.front();
    q.pop();
    inqueue[v] = false;
    for (auto& [to, len] : adj[v]) {
        if (d[v] + len < d[to]) {
            d[to] = d[v] + len;
            if (!inqueue[to]) {
                q.push(to);
                inqueue[to] = true;
                cnt[to]++;
                if (cnt[to] > n)
                    return false;//ciclo negativo
return true;
```

### 6.11 Camino mas corto de longitud fija

```
for (int i = 0; i <m; i++) {
            ll a, b, c; cin >> a >> b >> c; a--; b--;
            adj[a][b] = min(adj[a][b], c);
}

matrix graph(n, n, adj);
graph = pow(graph, k-1);

cout << (graph.m[0][n-1]==INFL ? -1 : graph.m[0][n -1]) << "\n";

return 0;
}</pre>
```

#### 6.12 2sat

```
// l=(x1 \text{ or } y1) and (x2 \text{ or } y2) and ... and (xn \text{ or } yn)
struct sat2 {
        int n;
        vector<vector<vi>>> q;
        vector<bool> vis, val;
        vi comp;
        stack<int> st:
        sat2(int n):n(n),q(2, vector < vi > (2*n)),vis(2*n),
            val(2*n), comp(2*n) {}
        int neq(int x) {return 2*n-x-1;}
        void make_true(int u) {add_edge(neg(u), u);}
        void make false(int u) {make_true(neg(u));}
        void add_or(int u, int v) {implication(neg(u), v);}
        void diff(int u, int v) {eq(u, neq(v));}
        void eq(int u, int v) {
                 implication(u, v);
                 implication(v, u);
        void implication(int u,int v) {
                 add edge(u, v);
                 add edge (neg(v), neg(u));
        void add_edge(int u, int v) {
                 q[0][u].PB(v);
                 q[1][v].PB(u);
        void dfs(int id, int u, int t=0) {
                 vis[u]=true;
                 for(auto &v:q[id][u])
                         if(!vis[v])dfs(id, v, t);
                 if (id) comp[u]=t;
                 else st.push(u);
        void kosaraju() {
```

```
for(int u=0; u<n; ++u) {
                           if(!vis[u])dfs(0, u);
                           if(!vis[neq(u)])dfs(0, neq(u));
                  vis.assign(2*n, false);
                  int t=0:
                  while(!st.empty()){
                           int u=st.top();st.pop();
                           if(!vis[u])dfs(1, u, t++);
         bool check() {
                  kosaraju();
                  for(int i=0;i<n;++i) {</pre>
                           if (comp[i] == comp[neg(i)]) return
                               false;
                           val[i]=comp[i]>comp[neg(i)];
                  return true;
};
int m,n;
sat2 s(n);
char c1, c2;
for (int a, b, i=0; i < m; ++i) {</pre>
         cin>>c1>>a>>c2>>b;
         a--;b--;
         if (c1=='-') a=s.neg(a);
         if (c2=='-')b=s.neq(b);
         s.add or(a,b);
if(s.check()){
         for (int i=0; i < n; ++i) cout << (s.val[i]?'+':'-') << " "</pre>
         cout << "\n";
}else cout<<"IMPOSSIBLE\n";</pre>
```

# 7 Matematicas

# 7.1 De Bruijn sequences

```
// Given alphabet [0, k) constructs a cyclic string
// of length k^n that contains every length n string as
    substr.
vi deBruijnSeq(int k, int n, int lim) {
    if (k == 1) return {0};
    vi seq, aux(n + 1);
    int cont = 0;
    function<void(int,int)> gen = [&](int t, int p) {
        if (t > n) {
```

#### 7.2 Chinese Remainder Theorem

```
/// Complexity: |N|*log(|N|)
/// Tested: Not yet.
/// finds a suitable x that meets: x is congruent to a i
/** Works for non-coprime moduli.
Returns {-1,-1} if solution does not exist or input is
 Otherwise, returns \{x, L\}, where x is the solution unique
     to mod L = LCM \ of \ mods
pll crt(vl A, vl M) {
  ll n = A.size(), a1 = A[0], m1 = M[0];
  for(ll i = 1; i < n; i++) {</pre>
    11 \ a2 = A[i], \ m2 = M[i];
    ll g = \underline{\hspace{0.2cm}} gcd(m1, m2);
    if( a1 % g != a2 % g ) return {-1,-1};
    11 p, q;
    extended_euclid(m1/g, m2/g, p, q);
    11 \mod = m1 / q * m2;
    q %= mod; p %= mod;
    11 x = ((111*(a1*mod)*(m2/g))*mod*q + (111*(a2*mod)*(
       m1/q))%mod*p) % mod; // if WA there is overflow
    a1 = x;
    if (a1 < 0) a1 += mod;
    m1 = mod;
  return {a1, m1};
```

### 7.3 Totient y Divisores

```
vector<int> count divisors sieve() {
  bitset<mx> is_prime; is_prime.set();
  vector<int> cnt(mx, 1);
  is\_prime[0] = is\_prime[1] = 0;
  for(int i = 2; i < mx; i++) {</pre>
    if(!is prime[i]) continue;
    cnt[i]++;
    for(int j = i+i; j < mx; j += i) {
      int n = 1, c = 1;
      while ( n\%i == 0 ) n /= i, c++;
      cnt[j] *= c;
      is prime[j] = 0;
  return cnt;
vector<int> euler phi sieve() {
  bitset<mx> is_prime; is_prime.set();
  vector<int> phi(mx);
  iota(phi.begin(), phi.end(), 0);
  is prime[0] = is prime[1] = 0;
  for(int i = 2; i < mx; i++) {</pre>
    if(!is_prime[i]) continue;
    for(int j = i; j < mx; j += i) {</pre>
      phi[j] -= phi[j]/i;
      is prime[j] = 0;
  return phi;
ll euler phi(ll n) {
  ll ans = n;
  for(ll i = 2; i * i <= n; ++i) {</pre>
    if(n % i == 0) {
      ans -= ans / i;
      while (n % i == 0) n /= i;
  if(n > 1) ans -= ans / n;
  return ans;
```

#### 7.4 Ecuaciones Diofanticas

```
// O(log(n))
ll extended_euclid(ll a, ll b, ll &x, ll &y) {
    ll xx = y = 0;
    ll yy = x = 1;
    while (b) {
        ll q = a / b;
        ll t = b; b = a % b; a = t;
        t = xx; xx = x - q * xx; x = t;
        t = yy; yy = y - q * yy; y = t;
}
```

```
return a:
// a*x+b*y=c. returns valid x and y if possible.
// all solutions are of the form (x0 + k * b / q, y0 - k
   *b/a
bool find any solution (ll a, ll b, ll c, ll &x0, ll &y0,
    ll &a) {
  if (a == 0 and b == 0) {
    if (c) return false;
    x0 = v0 = q = 0;
    return true;
 q = \text{extended euclid (abs(a), abs(b), x0, y0)};
 if (c % q != 0) return false;
 x0 *= c / a;
 y0 \star = c / q;
  if (a < 0) \times 0 *= -1:
  if (b < 0) v0 *= -1;
  return true;
void shift solution(ll &x, ll &y, ll a, ll b, ll cnt) {
 x += cnt * b;
  v = cnt * a;
// returns the number of solutions where x is in the
   range[minx, maxx] and y is in the range[miny, maxy]
ll find all solutions (ll a, ll b, ll c, ll minx, ll maxx,
    11 miny, 11 maxy) {
 ll x, y, g;
  if (find_any_solution(a, b, c, x, y, g) == 0) return 0;
 if (a == 0 and b == 0) {
    assert(c == 0);
    return 1LL * (maxx - minx + 1) * (maxy - miny + 1);
    return (maxx - minx + 1) * (miny <= c / b and c / b
       \leq maxy);
  if (b == 0) {
    return (maxy - miny + 1) * (minx <= c / a and c / a
       \leq maxx);
  a /= q, b /= q;
 ll sign a = a > 0 ? +1 : -1;
 ll sign b = b > 0 ? +1 : -1;
  shift_solution(x, y, a, b, (minx - x) / b);
 if (x < minx) shift solution(x, y, a, b, sign b);
  if (x > maxx) return 0;
 11 \ 1x1 = x;
  shift solution (x, y, a, b, (maxx - x) / b);
 if (x > maxx) shift_solution (x, y, a, b, -sign_b);
  shift solution (x, y, a, b, -(miny - y) / a);
  if (y < miny) shift_solution (x, y, a, b, -sign_a);</pre>
  if (y > maxy) return 0;
```

```
ll lx2 = x;
shift_solution(x, y, a, b, -(maxy - y) / a);
if (y > maxy) shift_solution(x, y, a, b, sign_a);
ll rx2 = x;
if (lx2 > rx2) swap (lx2, rx2);
ll lx = max(lx1, lx2);
ll rx = min(rx1, rx2);
if (lx > rx) return 0;
return (rx - lx) / abs(b) + 1;
}

///finds the first k / x + b * k / gcd(a, b) >= val
ll greater_or_equal_than(ll a, ll b, ll x, ll val, ll g)
{
ld got = 1.0 * (val - x) * g / b;
return b > 0 ? ceil(got) : floor(got);
}
```

### 7.5 Exponenciacion binaria

```
1l binpow(ll b, ll n, ll m) {
    b %= m;
    ll res = 1;
    while (n > 0) {
        if (n & 1)
            res = res * b % m;
        b = b * b % m;
        n >>= 1;
    }
    return res % m;
}
```

### 7.6 Exponenciacion matricial

```
return ans;
}

};

matrix pow(matrix &b, ll p) {
    matrix ans(b.r, b.c, vector<vl>(b.r, vl(b.c, 0)));
    for (int i = 0; i<b.r; i++) ans.m[i][i] = 1;
    while (p) {
        if (p&1) {
            ans = ans*b;
        }
        b = b*b;
        p >>= 1;
    }
    return ans;
}
```

### 7.7 Fibonacci Fast Doubling

```
// O(log n) muy rapido
pair<int, int> fib (int n) {
    if (n == 0)
        return {0, 1};

    auto p = fib(n >> 1);
    int c = p.first * (2 * p.second - p.first);
    int d = p.first * p.first + p.second * p.second;
    if (n & 1)
        return {d, c + d};
    else
        return {c, d};
}
```

### 7.8 Freivalds algorithm

```
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
    ().count());
// check if two n*n matrix a*b=c within complexity (
    iteration*n^2)
// probability of error 2^(-iteration)
// 0(iter*n^2)
int Freivalds(matrix &a, matrix &b, matrix &c) {
    int n = a.r, iteration = 20;
    matrix zero(n, 1), r(n, 1);
    while (iteration--) {
        for(int i = 0; i < n; i++) r.m[i][0] = rnd() % 2;
        matrix ans = (a * (b * r)) - (c * r);
        if(ans.m != zero.m) return 0;
    }
    return 1;
}</pre>
```

#### 7.9 Gauss Jordan

```
// O(min(n, m) *n*m)
const double EPS = 1e-9;
const int INF = 2; // it doesn't actually have to be
   infinity or a big number
int gauss (vector < vector<double> > a, vector<double> &
   ans) {
    int n = (int) a.size();
    int m = (int) a[0].size() - 1;
    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {</pre>
        int sel = row;
        for (int i=row; i<n; ++i)
            if (abs (a[i][col]) > abs (a[sel][col]))
                 sel = i;
        if (abs (a[sel][col]) < EPS)</pre>
            continue:
        for (int i=col; i<=m; ++i)
             swap (a[sel][i], a[row][i]);
        where [col] = row;
        for (int i=0; i<n; ++i)</pre>
            if (i != row) {
                 double c = a[i][col] / a[row][col];
                 for (int j=col; j<=m; ++j)</pre>
                     a[i][j] -= a[row][j] * c;
        ++row;
    ans.assign (m, 0);
    for (int i=0; i<m; ++i)
        if (where[i] !=-1)
             ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i=0; i<n; ++i) {</pre>
        double sum = 0;
        for (int j=0; j<m; ++j)
            sum += ans[j] * a[i][j];
        if (abs (sum - a[i][m]) > EPS)
            return 0;
    for (int i=0; i<m; ++i)
        if (where [i] == -1)
            return INF;
    return 1;
```

### 7.10 Gauss Jordan mod 2

```
// O(min(n, m)*n*m)
```

```
int gauss (vector < bitset<N> > &a, int n, int m, bitset<</pre>
   N > \& ans) {
    vector\langle int \rangle where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {</pre>
        for (int i=row; i<n; ++i)</pre>
             if (a[i][col]) {
                 swap (a[i], a[row]);
                 break;
        if (! a[row][col])
             continue;
        where [col] = row;
        for (int i=0; i<n; ++i)
             if (i != row && a[i][col])
                 a[i] ^= a[row];
        ++row;
    for (int i=0; i<m; ++i)
        if (where[i] != -1)
             ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i=0; i<n; ++i) {
        double sum = 0;
        for (int j=0; j<m; ++j)
             sum += ans[i] * a[i][i];
        if (abs (sum - a[i][m]) > EPS)
             return 0;
    for (int i=0; i<m; ++i)</pre>
        if (where [i] == -1)
             return INF;
    return 1;
```

# 7.11 GCD y LCM

```
//0(log10 n) n == max(a, b)
int gcd(int a, int b) { return b == 0 ? a : gcd(b, a%b);
}
int lcm(int a, int b) { return a / gcd(a, b) * b; }
//gcd(a, b, c) = gcd(a, gcd(b, c))
```

# 7.12 Integral Definida

```
const int steps = 1e6; // %2==0
double f(double x);
double simpson(double a, double b) {
         double h=(b-a)/steps;
         double s=f(a)+f(b);
         for(int i=1;i<=steps-1;i++) {</pre>
```

#### 7.13 Inverso modular

```
11 mod(ll a, ll m) {
    return ((a%m) + m) % m;
ll modInverse(ll b, ll m) {
    11 x, y;
    ll d = extEuclid(b, m, x, y); //obtiene b*x + m*y ==
    if (d != 1) return -1;
                                     //indica error
    // b*x + m*y == 1, ahora aplicamos (mod m) para
       obtener\ b*x == 1 \pmod{m}
    return mod(x, m);
// Otra forma
// O(log MOD)
ll inv (ll a) {
    return binpow(a, MOD-2, MOD);
//Modulo constante
inv[1] = 1;
for (int i = 2; i < p; ++i)
        inv[i] = (p - (p / i) * inv[p % i] % p) % p;
```

# 7.14 Logaritmo Discreto

```
// O(sqrt(m))
// Returns minimum x for which a ^ x % m = b % m.
int solve(int a, int b, int m) {
    a %= m, b %= m;
    int k = 1, add = 0, g;
    while ((g = gcd(a, m)) > 1) {
        if (b == k)
            return add;
        if (b % g)
            return -1;
        b /= g, m /= g, ++add;
        k = (k * 111 * a / g) % m;
}
int n = sqrt(m) + 1;
int an = 1;
for (int i = 0; i < n; ++i)</pre>
```

```
an = (an * 111 * a) % m;
unordered_map<int, int> vals;
for (int q = 0, cur = b; q <= n; ++q) {
    vals[cur] = q;
    cur = (cur * 111 * a) % m;
}

for (int p = 1, cur = k; p <= n; ++p) {
    cur = (cur * 111 * an) % m;
    if (vals.count(cur)) {
        int ans = n * p - vals[cur] + add;
        return ans;
    }
}
return -1;
}</pre>
```

#### 7.15 Miller Rabin

```
ll mul (ll a, ll b, ll mod) {
  11 \text{ ret} = 0;
  for (a %= mod, b %= mod; b != 0;
    b >>= 1, a <<= 1, a = a >= mod ? <math>a - mod : a) {
    if (b & 1) {
      ret += a;
      if (ret >= mod) ret -= mod;
  return ret;
ll fpow (ll a, ll b, ll mod) {
  11 \text{ ans} = 1;
  for (; b; b >>= 1, a = mul(a, a, mod))
    if (b & 1)
      ans = mul(ans, a, mod);
  return ans;
bool witness (ll a, ll s, ll d, ll n) {
  ll x = fpow(a, d, n);
  if (x == 1 \mid | x == n - 1) return false;
  for (int i = 0; i < s - 1; i++) {
    x = mul(x, x, n);
    if (x == 1) return true;
    if (x == n - 1) return false;
  return true;
11 \text{ test}[] = \{2, 3, 5, 7, 11, 13, 17, 19, 23, 0\};
bool is prime (ll n) {
  if (n < 2) return false;</pre>
  if (n == 2) return true;
  if (n % 2 == 0) return false;
  11 d = n - 1, s = 0;
  while (d \% 2 == 0) ++s, d /= 2;
```

```
for (int i = 0; test[i] && test[i] < n; ++i)
   if (witness(test[i], s, d, n))
     return false;
return true;
}</pre>
```

### 7.16 Miller Rabin Probabilistico

```
using u64 = uint64 t;
using u128 = uint128 t;
u64 binpower(u64 base, u64 e, u64 mod) {
    u64 result = 1:
    base %= mod;
    while (e) {
        if (e & 1)
           result = (u128) result * base % mod;
        base = (u128)base * base % mod;
        e >>= 1;
    return result;
bool check composite (u64 n, u64 a, u64 d, int s) {
    u64 x = binpower(a, d, n);
    if (x == 1 | | x == n - 1)
        return false:
    for (int r = 1; r < s; r++) {
        x = (u128)x * x % n;
        if (x == n - 1)
            return false:
    return true;
bool MillerRabin(u64 n, int iter=5) { // returns true if
   n is probably prime, else returns false.
    if (n < 4)
        return n == 2 || n == 3;
    int s = 0;
    u64 d = n - 1;
    while ((d \& 1) == 0) {
        d >>= 1;
        s++;
    for (int i = 0; i < iter; i++) {</pre>
        int a = 2 + rand() % (n - 3);
        if (check composite(n, a, d, s))
            return false;
    return true;
```

### 7.17 Mobius

```
const int N = 1e6+1;
int mob[N];
void mobius() {
   mob[1] = 1;
   for (int i = 2; i < N; i++) {
       mob[i]--;
      for (int j = i + i; j < N; j += i) {
       mob[j] -= mob[i];
    }
}</pre>
```

#### 7.18 Pollard Rho

```
//O(n^{(1/4)}) (?)
ll pollard_rho(ll n, ll c) {
 11 x = 2, y = 2, i = 1, k = 2, d;
 while (true) {
   x = (mul(x, x, n) + c);
    if (x >= n) x -= n;
    d = \underline{gcd(x - y, n)};
    if (d > 1) return d;
    if (++i == k) v = x, k <<= 1;
 return n;
void factorize(ll n, vector<ll> &f) {
 if (n == 1) return;
 if (is prime(n)) {
    f.push back(n);
    return;
 11 d = n;
  for (int i = 2; d == n; i++)
    d = pollard rho(n, i);
  factorize(d, f);
  factorize (n/d, f);
```

# 7.19 Simplex

```
vector<int> X,Y;
double z;
int n,m;
Simplex(vector<vector<double>> _a, vector<double>
     b, vector<double> c) {
         A=_a;B=_b;C=_c;
         n=\overline{B}.size(); m=C.size(); z=0.;
         X=vector<int>(m); Y=vector<int>(n);
         for (int i=0; i<m; ++i) X[i]=i;</pre>
         for (int i=0; i<n; ++i) Y[i]=i+m;</pre>
void pivot(int x,int y) {
         swap(X[y],Y[x]);
         B[x]/=A[x][y];
         for (int i=0; i<m; ++i) if (i!=y) A[x][i] /=A[x</pre>
             ][y];
         A[x][y]=1/A[x][y];
         for(int i=0;i<n;++i)if(i!=x&&abs(A[i][y])</pre>
             >EPS) {
                  B[i] -= A[i][y] *B[x];
                  for(int j=0; j<m;++j)if(j!=y)A[i][
                      j] -= A[i][y] *A[x][j];
                  A[i][y] = -A[i][y] * A[x][y];
         z+=C[y]*B[x];
         for (int i=0; i < m; ++i) if (i!=y) C[i] -= C[y] *A[</pre>
             x][i];
         C[y] = -C[y] *A[x][y];
pair<double, vector<double>> maximize() {
         while (1) {
                  int x=-1, y=-1;
                  double mn=-EPS;
                  for (int i=0; i<n; ++i) if (B[i] <mn) mn</pre>
                      =B[i], x=i;
                  if (x<0) break;</pre>
                  for (int i=0; i<m; ++i) if (A[x][i]<-</pre>
                      EPS) {y=i;break;}
                  // y<0, no solution to Ax<=B
                  pivot(x,y);
         while(1){
                  double mx=EPS;
                  int x=-1, y=-1;
                  for (int i=0; i<m; ++i) if (C[i]>mx) mx
                      =C[i], v=i;
                  if (y<0) break;
                  double mn=1e200;
                  for (int i=0; i<n; ++i) if (A[i][y]>
                      EPS\&\&B[i]/A[i][y]<mn)mn=B[i]/A
                      [i][y],x=i;
                  // x<0, unbounded
                  pivot(x,y);
```

#### 7.20 Fast Fourier Transform

```
// O(N log N)
const double PI = acos(-1);
struct base {
  double a, b;
  base (double a = 0, double b = 0) : a(a), b(b) {}
  const base operator + (const base &c) const
    { return base(a + c.a, b + c.b); }
  const base operator - (const base &c) const
    { return base(a - c.a, b - c.b); }
  const base operator * (const base &c) const
    { return base(a * c.a - b * c.b, a * c.b + b * c.a);
void fft(vector<base> &p, bool inv = 0) {
  int n = p.size(), i = 0;
  for (int j = 1; j < n - 1; ++j) {
    for (int k = n >> 1; k > (i^= k); k >>= 1);
    if(j < i) swap(p[i], p[j]);
  for (int 1 = 1, m; (m = 1 << 1) <= n; 1 <<= 1) {
    double ang = 2 * PI / m;
    base wn = base(cos(ang), (inv ? 1. : -1.) * sin(ang))
    for (int i = 0, j, k; i < n; i += m) {
      for(w = base(1, 0), j = i, k = i + 1; j < k; ++j, w
          = w * wn) {
        base t = w * p[j + 1];
        p[j + 1] = p[j] - t;
        p[j] = p[j] + t;
  if(inv) for(int i = 0; i < n; ++i) p[i].a /= n, p[i].b
     /= n;
vector<long long> multiply(vector<int> &a, vector<int> &b
  int n = a.size(), m = b.size(), t = n + m - 1, sz = 1;
  while(sz < t) sz <<= 1;
  vector<br/><br/>base> x(sz), y(sz), z(sz);
  for (int i = 0; i < sz; ++i) {
    x[i] = i < (int)a.size() ? base(a[i], 0) : base(0, 0)
        = i < (int)b.size() ? base(b[i], 0) : base(0, 0)
```

```
}
fft(x), fft(y);
for(int i = 0; i < sz; ++i) z[i] = x[i] * y[i];
fft(z, 1);
vector<long long> ret(sz);
for(int i = 0; i < sz; ++i) ret[i] = (long long) round(
    z[i].a);

// while((int) ret. size() > 1 && ret.back() == 0) ret.
    pop_back();
return ret;
}
```

#### 7.21 Number Theoretic Transform

```
const int N = 1 \ll 20;
const int mod = 469762049; //998244353
const int root = 3;
int lim, rev[N], w[N], wn[N], inv_lim;
void reduce(int &x) { x = (x + mod) % mod; }
int POW(int x, int y, int ans = 1) {
  for (; y; y >>= 1, x = (long long) x * x % mod) if <math>(y \& x)
      1) ans = (long long) ans * x % mod;
  return ans;
void precompute(int len) {
 \lim_{n \to \infty} = wn[0] = 1; int s = -1;
  while (lim < len) lim <<= 1, ++s;
  for (int i = 0; i < lim; ++i) rev[i] = rev[i >> 1] >> 1
      | (i & 1) << s;
  const int g = POW(root, (mod - 1) / lim);
  inv_lim = POW(lim, mod - 2);
  for (int i = 1; i < lim; ++i) wn[i] = (long long) wn[i</pre>
     - 1] * q % mod;
void ntt(vector<int> &a, int typ) {
  for (int i = 0; i < lim; ++i) if (i < rev[i]) swap(a[i</pre>
     ], a[rev[i]]);
  for (int i = 1; i < lim; i <<= 1) {</pre>
    for (int j = 0, t = \lim / i / 2; j < i; ++j) w[j] =
       wn[j * t];
    for (int j = 0; j < lim; j += i << 1) {
      for (int k = 0; k < i; ++k) {
        const int x = a[k + j], y = (long long) a[k + j +
             i] * w[k] % mod;
        reduce (a[k + j] += y - mod), reduce (a[k + j + i]
           = x - y);
  if (!tvp) {
    reverse(a.begin() + 1, a.begin() + lim);
    for (int i = 0; i < lim; ++i) a[i] = (long long) a[i]
         * inv_lim % mod;
```

```
}
}
vector<int> multiply(vector<int> &f, vector<int> &g) {
    int n=(int) f.size() + (int) g.size() - 1;
    precompute(n);
    vector<int> a = f, b = g;
    a.resize(lim); b.resize(lim);
    ntt(a, 1), ntt(b, 1);
    for (int i = 0; i < lim; ++i) a[i] = (long long) a[i] *
        b[i] % mod;
    ntt(a, 0);
    a.resize(n + 1);
    return a;
}</pre>
```

# 8 Programacion dinamica

# 8.1 Bin Packing

```
int main() {
    ll n, capacidad;
    cin >> n >> capacidad;
    vl pesos(n, 0);
    forx(i, n) cin >> pesos[i];
    vector < pll > dp((1 << n));
    dp[0] = \{1, 0\};
    // dp[X] = \{ #numero de paquetes, peso de min paquete \}
    // La idea es probar todos los subset y en cada uno
       preguntarnos
    // quien es mejor para subirse de ultimo buscando
       minimizar
    // primero el numero de paquetes
    for (int subset = 1; subset < (1 << n); subset++) {</pre>
        dp[subset] = \{21, 0\};
        for (int iPer = 0; iPer < n; iPer++) {</pre>
            if ((subset >> iPer) & 1) {
                pll ant = dp[subset ^ (1 << iPer)];</pre>
                ll k = ant.ff;
                ll w = ant.ss;
                if (w + pesos[iPer] > capacidad) {
                     k++;
                     w = min(pesos[iPer], w);
                 } else {
                     w += pesos[iPer];
                dp[subset] = min(dp[subset], {k, w});
```

```
cout << dp[(1 << n) - 1].ff << ln;
```

### 8.2 CHT

```
// - Me dan las pendientes ordenadas
// Caso 1: Me hacen las querys ordenadas
// O(N + O)
// Caso 2: Me hacen querys arbitrarias
// O(N + QlogN)
struct CHT {
    // funciona tanto para min como para max, depende del
        orden en que pasamos las lineas
    struct Line {
        int slope, yIntercept;
        Line (int slope, int yIntercept) : slope (slope),
           yIntercept(yIntercept){}
        int val(int x) { return slope * x + yIntercept; }
        int intersect(Line y) {
            return (y.yIntercept - yIntercept + slope - y
                .slope - 1) / (slope - y.slope);
    };
    deque<pair<Line, int>> dq;
    void insert(int slope, int yIntercept){
                // lower hull si m1 < m2 < m3
                // upper hull si si m1 > m2 > m3
        Line newLine(slope, yIntercept);
        while (!dq.empty() && dq.back().second >= dq.back
            ().first.intersect(newLine)) dq.pop_back();
        if (dq.empty()) {
            dq.emplace_back(newLine, 0);
            return;
        dq.emplace back(newLine, dq.back().first.
           intersect(newLine));
    int query(int x) { // cuando las consultas son
       crecientes
        while (dq.size() > 1) {
            if (dq[1].second <= x) dq.pop_front();</pre>
            else break;
        return dq[0].first.val(x);
    int query2(int x) { // cuando son arbitrarias
        auto gry = *lower bound(dg.rbegin(), dg.rend(),
                                make_pair(Line(0, 0), x),
```

### 8.3 CHT Dynamic

```
// O((N+Q) \log N) < -usando set para add y bs para q
// lineas de la forma mx + b
#pragma once
struct Line {
        mutable 11 m, b, p;
        bool operator<(const Line& o) const { return m <</pre>
        bool operator<(ll x) const { return p < x; }</pre>
};
struct CHT : multiset<Line, less<>>> {
        // (for doubles, use inf = 1/.0, div(a,b) = a/b)
        static const ll inf = LLONG MAX;
        static const bool mini = 0; // <---- 1 FOR MIN</pre>
        ll div(ll a, ll b) { // floored division
                return a / b - ((a ^ b) < 0 && a % b); }
        bool isect(iterator x, iterator y) {
                if (v == end()) return x \rightarrow p = inf, 0;
                if (x->m == y->m) x->p = x->b > y->b?
                    inf : -inf;
                 else x->p = div(y->b - x->b, x->m - y->m)
                return x->p >= y->p;
        void add(ll m, ll b) {
        if (mini) { m \star= -1, b \star= -1; }
                auto z = insert(\{m, b, 0\}), y = z++, x =
                while (isect(y, z)) z = erase(z);
                 if (x != begin() && isect(--x, y)) isect(
                    x, y = erase(y);
                while ((y = x) != begin() \&\& (--x)->p >=
                    y->p)
                         isect(x, erase(y));
        11 query(11 x) {
                assert(!empty());
                 auto 1 = *lower_bound(x);
        if (mini) return -l.m * x + -l.b;
                else return l.m * x + l.b;
```

# 8.4 Divide Conquer

};

```
// C[a][c] + C[b][d] <= C[a][d] + C[b][c] where a < b < c
    < d.
int m, n;
vector<long long> dp before(n), dp cur(n);
long long C(int i, int j);
// compute dp cur[l], ... dp cur[r] (inclusive)
void compute(int 1, int r, int opt1, int optr) {
    if (l > r)
        return;
    int mid = (1 + r) >> 1;
    pair<long long, int> best = {LLONG MAX, -1};
    for (int k = optl; k <= min(mid, optr); k++) {
        best = min(best, \{(k ? dp before[k - 1] : 0) + C(
           k, mid), k);
    dp cur[mid] = best.first;
    int opt = best.second;
    compute(1, mid - 1, optl, opt);
    compute(mid + 1, r, opt, optr);
int solve() {
    for (int i = 0; i < n; i++)</pre>
        dp before[i] = C(0, i);
    for (int i = 1; i < m; i++) {
        compute (0, n - 1, 0, n - 1);
        dp_before = dp_cur;
    return dp_before[n - 1];
```

#### 8.5 Edit Distances

#### 8.6 Kadane 2D

```
int main() {
    ll fil,col;cin>>fil>>col;
    vector<vl> grid(fil, vl(col, 0));
// Algoritmo de Kadane/DP para suma maxima de una matriz
    2D en o(n^3)
    for(int i=0;i<fil;i++) {</pre>
        for(int e=0;e<col;e++){</pre>
             11 num; cin>>num;
             if (e>0) grid[i][e]=num+grid[i][e-1];
             else grid[i][e]=num;
    11 maxGlobal = LONG LONG MIN;
    for(int l=0;1<col;1++){</pre>
        for (int r=1; r < col; r++) {</pre>
             11 maxLoc=0;
             for(int row=0;row<fil;row++) {</pre>
                 if (1>0) maxLoc+=grid[row][r]-grid[row][1
                     -1];
                 else maxLoc+=grid[row][r];
                 if (maxLoc<0) maxLoc=0;</pre>
                 maxGlobal= max(maxGlobal, maxLoc);
```

### 8.7 Knuth

```
8.8 LIS
```

```
4
```

```
9 STRINGS
```

```
... // read N and input
int dp[N][N], opt[N][N];
auto C = [\&] (int i, int j) {
    ... // Implement cost function C.
for (int i = 0; i < N; i++) {</pre>
    opt[i][i] = i;
    ... // Initialize dp[i][i] according to the
       problem
for (int i = N-2; i >= 0; i--) {
    for (int j = i+1; j < N; j++) {
        int mn = INT_MAX;
        int cost = C(i, j);
        for (int k = opt[i][j-1]; k \le min(j-1, opt[i]
           +1][\dot{\gamma}]); k++) {
            if (mn >= dp[i][k] + dp[k+1][j] + cost) {
                opt[i][j] = k;
                mn = dp[i][k] + dp[k+1][j] + cost;
        dp[i][j] = mn;
cout << dp[0][N-1] << endl;
```

### 8.8 LIS

```
// 0(nlogn)
int lis(vi& a){
         int n=sz(a), last=0;
         vi dp (n+1, INT\_MAX), cnt (n, 0);
         dp[0] = INT_MIN;
         for (int i=0; i < n; ++i) {</pre>
                  int j=lower bound(all(dp), a[i])-dp.begin
                      (); // upper_bound
                  if (dp[j-1] < a[i] && a[i] < dp[j]) { // dp[j]
                      -11 <= a[i]
                            dp[j]=a[i];
                           last=max(last,j);
                  cnt[i]=j;
         int ans=0;
         for (int i=0; i<=n; i++) {</pre>
                  if (dp[i] < INT MAX) ans=i;</pre>
         vi LIS(ans);
         int act=ans;
         for(int i=n-1; i>=0; --i) {
```

#### 8.9 SOS

```
const int bits = 23;
int dp[1<<bits];</pre>
// O(n*2^n)
void SOS(){
        for (int i = 0; i < (1 << bits); ++i) dp[i] = A[i]
        // top - down
        for(int i = 0; i < bits; ++i){</pre>
                 for(int s = 0; s < (1 << bits); ++s) {
                         if(s & (1 << i)){
                                  dp[s] += dp[s ^ (1 << i)
                                      ];
        // bottom - up
        for(int i = 0; i < bits; ++i) {</pre>
                 for (int s = (1 << bits) - 1; s >= 0; --s)
                         if(s & (1 << i)){
                                  dp[s ^ (1 << i)] += dp[s]
                                      ];
```

# 9 Strings

# 9.1 Hashing

```
1000234999, 1000567999, 1000111997, 1000777121,
        1001265673, 1001864327, 999727999, 1070777777
const int mod[2] = { 1001864327, 1001265673 };
const ii base(257, 367), zero(0, 0), one(1, 1);
const int maxn = 1e6;
inline int add(int a, int b, int m) {return a+b>=m?a+b-m:a
        +b;}
```

```
inline int sbt(int a, int b, int m){return a-b<0?a-b+m:a-
inline int mul(int a, int b, int m) {return ll(a)*b%m;}
inline ll operator ! (const ii a) {return (ll(a.first)
   <<32) | a.second; }
inline ii operator + (const ii& a, const ii& b) {return {
   add(a.first, b.first, mod[0]), add(a.second, b.second,
    mod[1])};}
inline ii operator - (const ii& a, const ii& b) {return {
   sbt(a.first, b.first, mod[0]), sbt(a.second, b.second,
    mod[1])};}
inline ii operator * (const ii& a, const ii& b) {return {
   mul(a.first, b.first, mod[0]), mul(a.second, b.second,
    mod[1])};}
ii p[maxn+1];
void prepare() { // Acordate del prepare()!!
        p[0]=one;
        for (int i=1; i <= maxn; i++) p[i] = p[i-1] *base;</pre>
template <class type>
struct hashing{
        vector<ii> h:
        hashing(type& t) {
                h.resize((int)t.size()+1);
                h[0]=zero;
                for(int i=1; i<(int)h.size();++i)
                         h[i]=h[i-1]*base + ii{t[i-1], t[i]}
        ii get(int 1, int r) {return h[r+1]-h[1]*p[r-1
ii combine(ii a, ii b, int lenb) {return a*p[lenb]+b;}
```

### 9.2 KMP

```
// O(n)
vi phi(string& s) {
    int n=sz(s);
    vi tmp(n);
    for(int i=1, j=0; i < n; ++i) {
        while (j > 0 && s[j]!=s[i]) j=tmp[j-1];
        if (s[i]==s[j]) j++;
        tmp[i]=j;
    }
    return tmp;
}

// O(n+m)
int kmp(string& s, string& p) {
    int n=sz(s), m=sz(p), cnt=0;
    vi pi=phi(p);
```

### 9.3 KMP Automaton

### 9.4 Manacher

```
// O(n), par (raiz, izq, der) 1 - impar 0
vi manacher(string& s, int par) {
    int l=0, r=-1, n=sz(s); vi m(n,0);
    for(int i=0;i<n;++i) {
        int k=(i>r?(1-par):min(m[l+r-i+ par], r-i +par))+par;
        while(i+k-par<n && i-k>=0 && s[i+k-par]== s[i-k])++k;
        m[i]=k-par;--k;
        if(i+k-par>r) l=i-k, r=i+k-par;
    }
    for(int i=0;i<n;++i)m[i]=(m[i]-1+par)*2+1-par;
    return m;
}</pre>
```

# 9.5 Minimum Expression

```
// O(n)
int minimum_expression(string s) {
    s=s+s;int n=sz(s),i=0,j=1,k=0;
```

```
while(i+k<n && j+k<n) {
    if(s[i+k]==s[j+k])k++;
    else if(s[i+k]>s[j+k])i=i+k+1,k=0; //
        cambiar por < para max
    else j=j+k+1,k=0;
    if(i==j)j++;
}
return min(i, j);
}</pre>
```

### 9.6 Palindromic Tree

```
const int alpha = 26;
const char fc = 'a';
// tree suf is the longest suffix palindrome
// tree dad is the palindrome add c to the right and left
struct Node{
        int next[alpha];
        int len, suf, dep, cnt, dad;
};
// O(nlogn)
struct PalindromicTree{
        vector<Node> tree;
        string s:
        int len,n;
        int size; // node 1 - root with len -1, node 2 -
            root with len 0
        int last; // max suffix palindrome
        bool addLetter(int pos) {
                int cur=last, curlen=0;
                int let=s[pos]-fc;
                while(true) {
                         curlen=tree[cur].len;
                         if(pos-1-curlen>=0 && s[pos-1-
                            curlen] == s[pos]) break;
                         cur=tree[cur].suf;
                if(tree[cur].next[let]){
                         last=tree[cur].next[let];
                        tree[last].cnt++;
                        return false;
                size++;
                last=size;
                tree[size].len=tree[cur].len+2;
                tree[cur].next[let]=size;
                tree[size].cnt=1;
                tree[size].dad=cur;
                if (tree[size].len==1) {
```

```
tree[size].suf=2;
                         tree[size].dep=1;
                         return true;
                 while(true) {
                         cur=tree[curl.suf;
                         curlen=tree[cur].len;
                         if(pos-1-curlen>=0 && s[pos-1-
                             curlen] == s[pos]) {
                                  tree[size].suf=tree[cur].
                                     next[let];
                                  break;
                 tree[size].dep=1+tree[tree[size].suf].dep
                 return true;
        PalindromicTree(string& s2, int n) {
                 tree.assign(n+4.Node());
                 tree[1].len=-1; tree[1].suf=1;
                 tree[2].len=0;tree[2].suf=1;
                 size=2; last=2; s=s2;
                 for (int i=0; i<n; i++) {</pre>
                         addLetter(i);
                 for(int i=size; i>=3; i--) {
                         tree[tree[i].suf].cnt+=tree[i].
                             cnt;
};
```

# 9.7 Suffix Array

```
// O(nlogn)
struct SuffixArray{
    const int alpha = 256;
    string s; int n;
    vi sa,rnk,lcp;

    SuffixArray(string& _s) {
        s=_s;s.push_back('$'); // check
        n=sz(s);
        sa.assign(n, 0);
        rnk.assign(n, 0);
        lcp.assign(n-1, 0);
        buildSA();
}

void buildSA() {
```

```
vi cnt(max(alpha, n),0);
for (int i=0; i < n; ++i) cnt [s[i]] ++;</pre>
for (int i=1; i < max (alpha, n); ++i) cnt[i] +=</pre>
    cnt[i-1];
for(int i=n-1;i>=0;--i)sa[--cnt[s[i]]]=i;
for(int i=1;i<n;++i)rnk[sa[i]]=rnk[sa[i</pre>
    -1] + (s[sa[i]]!=s[sa[i-1]]);
for (int k=1; k < n; k *=2) {
         vi nsa(n),nrnk(n),ncnt(n);
         for (int i=0; i < n; ++i) sa[i] = (sa[i] -</pre>
         for (int i=0; i < n; ++i) ncnt[rnk[i</pre>
             ]]++;
         for (int i=1; i < n; ++i) ncnt[i] += ncnt</pre>
         for(int i=n-1;i>=0;--i)nsa[--ncnt
             [rnk[sa[i]]]=sa[i];
         for (int i=1; i < n; ++i) {</pre>
                  ii op1={rnk[nsa[i]], rnk
                       [(nsa[i]+k)%n]};
                  ii op2={rnk[nsa[i-1]],}
                      rnk[(nsa[i-1]+k)%n];
                  nrnk[nsa[i]]=nrnk[nsa[i
                      -1]]+(op1!=op2);
         swap(sa, nsa);swap(rnk, nrnk);
for (int i=0, k=0; i < n-1; ++i) {</pre>
         while (s[i+k] == s[sa[rnk[i]-1]+k])k
         lcp[rnk[i]-1]=k;
         if(k)k--;
```

### 9.8 Suffix Automaton

};

```
// O(n*log(alpha))
struct SuffixAutomaton{
    vector<map<char,int>> to;
    vector<bool> end;
    vi suf,len; // len, longest string
    int last;

SuffixAutomaton(string& s) {
        to.push_back(map<char,int>());
        suf.push_back(-1);
        len.push_back(0);
        last=0;

    for(int i=0;i<sz(s);i++) {
            to.push_back(map<char,int>());
        }
}
```

```
suf.push back(0);
                         len.push back(i+1);
                         int r=sz(to)-1;
                         int p=last;
                         while (p>=0 && to[p].find(s[i])==
                            to[p].end()){
                                 to[p][s[i]]=r;
                                 p=suf[p];
                         if(p!=-1){
                                 int q=to[p][s[i]];
                                 if(len[p]+1==len[q]){
                                          suf[r]=q;
                                 }else{
                                          to.push_back(to[q
                                             ]);
                                          suf.push back(suf
                                              [a]);
                                          len.push_back(len
                                              [p]+1);
                                          int qq=sz(to)-1;
                                          suf[q]=qq;
                                          suf[r]=qq;
                                          while(p>=0 && to[
                                             p[s[i]] == q) {
                                                  to[p][s[i
                                                      ] =qq;
                                                  p=suf[p];
                         last=r;
                end.assign(sz(to), false);
                int p=last;
                while (p) {
                         end[p]=true;
                         p=suf[p];
};
```

### 9.9 Suffix Tree

```
// O(n)
struct SuffixTree{
    vector<map<char,int>> to;
    vector<int> pos,len,link;
    const int inf = 1e9;
    int size=0;
    string s;
    int make(int _pos, int _len){
```

```
9.10 Trie
```

```
to.push back(map<char,int>());
        pos.push_back(_pos);
        len.push_back(_len);
        link.push back(-1);
        return size++;
void add(int& p, int& lef, char c){
        s+=c;++lef;int lst=0;
        for(;lef;p?p=link[p]:lef--){
                 while (lef>1 && lef>len[to[p][s[sz
                     (s)-lef]]]){
                         p=to[p][s[sz(s)-lef]], lef
                             -=len[p];
                 char e=s[sz(s)-lef];
                 int& q=to[p][e];
                 if(!a){
                          q=make(sz(s)-lef,inf),
                             link[lst]=p,lst=0;
                 }else{
                          char t=s[pos[q]+lef-1];
                         if(t==c){link[lst]=p;
                             return; }
                          int u=make(pos[q],lef-1);
                          to[u][c]=make(sz(s)-1,inf
                             );
                          to[u][t]=q;
                          pos[q] += lef -1;
                          if(len[q]!=inf)len[q]-=
                             lef-1;
                          q=u,link[lst]=u,lst=u;
void build(string& _s) {
        make (-1, 0); int p=0, lef=0;
        for (char c:_s) add (p, lef, c);
add (p, lef, '$');
        s.pop_back();
int query(string& p){
        for (int i=0, u=0, n=sz(p);;) {
                 if(i==n || !to[u].count(p[i]))
                     return i;
                 u=to[u][p[i]];
                 for (int j=0; j<len[u];++j) {</pre>
                         if(i==n || s[pos[u]+j]!=p
                             [i])return i;
                         i++;
```

```
vector<int> sa;
           void genSA(int x=0, int Len=0) {
                   if(!sz(to[x]))sa.push back(pos[x]-Len);
                   else for (auto t:to[x]) genSA (t.second, Len+
                       len[x]);
  };
9.10 Trie
  const int maxn = 2e6+5, alpha = 26, bits = 30;
  int to[maxn][alpha], cnt[maxn], act;
  void init(){
           for(int i=0;i<=act;++i){</pre>
                   cnt[i]=0;
                   // suf[i]=dad[i]=0;
                   // adj[i].clear();
                   memset(to[i], 0, sizeof(to[i]));
           act=0;
  int add(string& s) {
           int u=0;
           for(char ch:s){
                   int c=conv(ch);
                   if(!to[u][c])to[u][c]=++act;
                   u=to[u][c];
           cnt[u]++;
           return u;
  // Aho-Corasick
  vector<int> adj[maxn]; // dad or suf
  int dad[maxn], suf[maxn];
  // O(sum(n) *alpha)
  void build() {
           queue<int> q{{0}};
           while(!q.empty()){
                   int u=q.front();q.pop();
                   for (int i=0; i < alpha; ++i) {</pre>
                            int v=to[u][i];
                            if(!v)to[u][i]=to[suf[u]][i];
                            else q.push(v);
                            if(!u || !v)continue;
                            suf[v]=to[suf[u]][i];
                            dad[v]=cnt[suf[v]]?suf[v]:dad[suf
                                [V]];
           for(int i=1;i<=act;++i){</pre>
                   adj[i].push_back(dad[i]);
```

```
adj[dad[i]].push_back(i);
```

### 9.11 Z Algorithm

```
// O(n)
vi z function(string& s){
         int n=sz(s), l=0, r=0; vi z(n);
         for (int i=1; i < n; i++) {</pre>
                   if (i < r) z[i] = min(r-i, z[i-l]);</pre>
                   while (i+z[i] < n \& \& s[z[i]] == s[i+z[i]]) z[i]
                       ]++;
                   if(i+z[i]>r){
                             l=i:
                             r=i+z[i];
         return z;
```

### 10 Misc

### 10.1 Counting Sort

```
// O(n+k)
void counting sort(vi& a){
        int maxi=*max element(all(a));
        int mini=*min_element(all(a));
        int k=maxi-mini+1, n=sz(a);
        vi cnt(k,0);
        for (int i=0; i < n; ++i) ++cnt[a[i]-mini];</pre>
        for(int i=0, j=0; i<k; ++i)
                 while (cnt[i]--) a [j++]=i+mini;
```

# 10.2 Dates

```
int dateToInt(int y, int m, int d) {
         return 1461* (y+4800+(m-14)/12)/4+367* (m-2-(m-14)
             /12*12)/12-
                  3 \star ((y+4900+(m-14)/12)/100)/4+d-32075;
void intToDate(int jd, int& y, int& m, int& d) {
         int x,n,i,j;x=jd+68569;
         n=4*x/146097; x=(146097*n+3)/4;
         i = (4000 * (x+1)) / 1461001; x = 1461 * i / 4 - 31;
         j=80*x/2447; d=x-2447*j/80;
         x=\frac{1}{11}; m=\frac{1}{12}+2-12*x; y=100* (n-49) +i+x;
```

```
Sunday
     static int ttt[]={0, 3, 2, 5, 0, 3, 5, 1, 4, 6,
        2, 4};
     v = m < 3:
     return (y+y/4-y/100+y/400+ttt[m-1]+d)%7;
```

# 10.3 Expression Parsing

```
// O(n) - En python es eval()
bool delim(char c) {return c==' ';}
bool is op(char c){return c=='+' || c=='-' || c=='*' || c
   ==' /' ; }
bool is unary(char c){return c=='+' || c=='-';}
int priority(char op) {
        if(op<0)return 3;</pre>
        if (op=='+' | op=='-') return 1;
        if(op=='*' || op=='/')return 2;
        return -1;
void process op(stack<int>& st, char op){
        if(op<0){
                 int l=st.top();st.pop();
                 switch (-op) {
                         case '+':st.push(1);break;
                         case '-':st.push(-1);break;
        }else{
                 int r=st.top();st.pop();
                 int l=st.top();st.pop();
                 switch(op){
                         case '+':st.push(l+r);break;
                         case '-':st.push(l-r);break;
                         case '*':st.push(l*r);break;
                         case '/':st.push(l/r);break;
int evaluate(string& s) {
        stack<int> st;
        stack<char> op;
        bool may_be_unary=true;
        for (int i=0; i < sz(s); ++i) {</pre>
                 if (delim(s[i])) continue;
                 if(s[i] == '('){
                         op.push('(');
                         may be unary=true;
                 }else if(s[i]==')'){
                         while (op.top()!='('){
```

```
process op(st, op.top());
                         op.pop();
                op.pop();
                may_be_unary=false;
        }else if(is_op(s[i])){
                char cur op=s[i];
                if(may be unary && is unary(
                    cur_op))cur_op=-cur_op;
                while(!op.empty() && ((cur_op >=
                    0 && priority(op.top()) >=
                    priority(cur op)) || (cur op <</pre>
                     0 && priority(op.top()) >
                    priority(cur op)))){
                         process_op(st, op.top());
                         op.pop();
                op.push(cur_op);
                may be unary=true;
        }else{
                int number=0;
                while(i<sz(s) && isalnum(s[i]))</pre>
                    number=number \star 10+s[i++]-'0';
                 st.push(number);
                may be unary=false;
while(!op.empty()){
        process_op(st, op.top());
        op.pop();
return st.top();
```

### 10.4 Ternary Search

```
// O(log((r-l)/eps))
double ternary(){
          double l, r;
          while(r-l>eps) {
                double m1=l+(r-l)/3.0;
                double m2=r-(r-l)/3.0;
                 if(f(m1)<f(m2))l=m1;
                 else r=m2;
        }
        return max(f(l),f(r));
}</pre>
```

### 10.5 Prefix3D

```
const int N = 100;
int A[N][N][N];
int preffix [N + 1][N + 1][N + 1];
void build(int n) {
        for (int x = 1; x \le n; x++) {
        for (int y = 1; y \le n; y++) {
            for (int z = 1; z <= n; z++) {
                 preffix[x][y][z] = A[x - 1][y - 1][z - 1]
                     + preffix[x - 1][y][z] + preffix[x][y
                         - 1][z] + preffix[x][y][z - 1]
                     - preffix[x - 1][y - 1][z] - preffix[x - 1][y - 1][z]
                        x - 1|[y][z - 1] - preffix[x][y -
                        11[z - \bar{1}]
                     + preffix[x - 1][y - 1][z - 1];
11 query(int lx, int rx, int ly, int ry, int lz, int rz){
        ll ans = preffix[rx][ry][rz]
                 - preffix[lx - 1][ry][rz] - preffix[rx][
                    ly - 1][rz] - preffix[rx][ry][lz - 1]
                 + preffix[lx - 1][ly - 1][rz] + preffix[
                    lx - 1][ry][lz - 1] + preffix[rx][ly -
                     1][lz - 1]
                - preffix[lx - 1][ly - 1][lz - 1];
        return ans;
```

#### 10.6 Hanoi

```
// hanoi(n) = 2 * hanoi(n-1) + 1
// hanoi(n, 1, 3)
vi ans;
void hanoi(int x, int start, int end) {
    if(!x)return;
    hanoi(x-1, start, 6-start-end);
    ans.push_back({start, end});
    hanoi(x-1, 6-start-end, end);
}
```

# 11 Teoría y miscelánea

#### 11.1 Sumatorias

$$\bullet \ \sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$$

• 
$$\sum_{i=1}^{n} i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

$$\bullet \sum_{i=1}^{n} i^3 = \left(\frac{n(n+1)}{2}\right)^2$$

• 
$$\sum_{i=0}^{n} x^i = \frac{x^{n+1}-1}{x-1} \text{ para } x \neq 1$$

### 11.2 Teoría de Grafos

#### 11.2.1 Teorema de Euler

En un grafo conectado planar, se cumple que V-E+F=2, donde V es el número de vértices, E es el número de aristas y F es el número de caras. Para varios componentes la formula es: V-E+F=1+C, siendo C el número de componentes.

#### 11.2.2 Planaridad de Grafos

Un grafo es planar si y solo si no contiene un subgrafo homeomorfo a  $K_5$  (grafo completo con 5 vértices) ni a  $K_{3,3}$  (grafo bipartito completo con 3 vértices en cada conjunto).

#### 11.3 Teoría de Números

#### 11.3.1 Ecuaciones Diofánticas Lineales

Una ecuación diofántica lineal es una ecuación en la que se buscan soluciones enteras x e y que satisfagan la relación lineal ax+by=c, donde a, b y c son constantes dadas.

Para encontrar soluciones enteras positivas en una ecuación diofántica lineal, podemos seguir el siguiente proceso:

- 1. Encontrar una solución particular: Encuentra una solución particular  $(x_0, y_0)$  de la ecuación. Esto puede hacerse utilizando el algoritmo de Euclides extendido.
- 2. Encontrar la solución general: Una vez que tengas una solución particular, puedes obtener la solución general utilizando la fórmula:

$$x = x_0 + \frac{b}{\operatorname{mcd}(a, b)} \cdot t$$

$$y = y_0 - \frac{a}{\operatorname{mcd}(a, b)} \cdot t$$

donde t es un parámetro entero.

3. Restringir a soluciones positivas: Si deseas soluciones positivas, asegúrate de que las soluciones generales satisfagan  $x \ge 0$  y  $y \ge 0$ . Puedes ajustar el valor de t para cumplir con estas restricciones.

#### 11.3.2 Pequeño Teorema de Fermat

Si p es un número primo y a es un entero no divisible por p, entonces  $a^{p-1} \equiv 1 \pmod{p}$ .

#### 11.3.3 Teorema de Euler

Para cualquier número entero positivo n y un entero a coprimo con n, se cumple que  $a^{\phi(n)} \equiv 1 \pmod{n}$ , donde  $\phi(n)$  es la función phi de Euler, que representa la cantidad de enteros positivos menores que n y coprimos con n.

#### 11.4 Geometría

#### 11.4.1 Teorema de Pick

Sea un poligono simple cuyos vertices tienen coordenadas enteras. Si B es el numero de puntos enteros en el borde, I el numero de puntos enteros en el interior del poligono, entonces el area A del poligono se puede calcular con la formula:

$$A = I + \frac{B}{2} - 1$$

#### 11.4.2 Fórmula de Herón

Si los lados del triángulo tienen longitudes a, b y c, y s es el semiperímetro (es decir,  $s=\frac{a+b+c}{2}$ ), entonces el área A del triángulo está dada por:

$$A = \sqrt{s(s-a)(s-b)(s-c)}$$

# 11.4.3 Relación de Existencia Triangular

Para un triángulo con lados de longitud  $a,\,b,\,{\bf y}\,c,$  la relación de existencia triangular se expresa como:

$$b - c < a < b + c$$
,  $a - c < b < a + c$ ,  $a - b < c < a + b$ 

#### 11.5 Combinatoria

#### 11.5.1 Permutaciones

El número de permutaciones de n objetos distintos tomados de a r a la vez (sin repetición) se denota como P(n,r) y se calcula mediante:

$$P(n,r) = \frac{n!}{(n-r)!}$$

#### 11.5.2 Combinaciones

El número de combinaciones de n objetos distintos tomados de a r a la vez (sin repetición) se denota como C(n,r) o  $\binom{n}{r}$  y se calcula mediante:

$$C(n,r) = \binom{n}{r} = \frac{n!}{r!(n-r)!}$$

#### 11.5.3 Permutaciones con Repetición

El número de permutaciones de n objetos tomando en cuenta repeticiones se denota como  $P_{\text{rep}}(n; n_1, n_2, \dots, n_k)$  y se calcula mediante:

$$P_{\text{rep}}(n; n_1, n_2, \dots, n_k) = \frac{n!}{n_1! n_2! \cdots n_k!}$$

### 11.5.4 Combinaciones con Repetición

El número de combinaciones de n objetos tomando en cuenta repeticiones se denota como  $C_{\text{rep}}(n; n_1, n_2, \dots, n_k)$  y se calcula mediante:

$$C_{\text{rep}}(n; n_1, n_2, \dots, n_k) = \binom{n+k-1}{n} = \binom{n+k-1}{k-1}$$

#### 11.5.5 Números de Catalan

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

Los números de Catalan también pueden calcularse utilizando la siguiente fórmula recursiva:

$$C_0 = 1$$

$$C_{n+1} = \frac{4n+2}{n+2}C_n$$

Usos:

• Cat(n) cuenta el número de árboles binarios distintos con n vértices.

- Cat(n) cuenta el número de expresiones que contienen n pares de paréntesis correctamente emparejados.
- Cat(n) cuenta el número de formas diferentes en que se pueden colocar n+1 factores entre paréntesis, por ejemplo, para n=3 y 3+1=4 factores: a,b,c,d, tenemos: (ab)(cd),a(b(cd)),((ab)c)d y a((bc)d).
- Los números de Catalan cuentan la cantidad de caminos no cruzados en una rejilla  $n \times n$  que se pueden trazar desde una esquina de un cuadrado o rectángulo a la esquina opuesta, moviéndose solo hacia arriba y hacia la derecha.
- Los números de Catalan representan el número de árboles binarios completos con n+1 hojas.
- $\operatorname{Cat}(n)$  cuenta el número de formas en que se puede triangular un poligono convexo de n+2 lados. Otra forma de decirlo es como la cantidad de formas de dividir un polígono convexo en triángulos utilizando diagonales no cruzadas.

#### 11.5.6 Estrellas y barras

Número de soluciones de la ecuación  $x_1 + x_2 + \cdots + x_k = n$ .

- Con  $x_i \ge 0$ :  $\binom{n+k-1}{n}$
- Con  $x_i \ge 1$ :  $\binom{n-1}{k-1}$

Número de sumas de enteros con límite inferior:

Esto se puede extender fácilmente a sumas de enteros con diferentes límites inferiores. Es decir, queremos contar el número de soluciones para la ecuación:

$$x_1 + x_2 + \cdots + x_k = n$$

con  $x_i \geq a_i$ .

Después de sustituir  $x_i' := x_i - a_i$  recibimos la ecuación modificada:

$$(x'_1 + a_i) + (x'_2 + a_i) + \dots + (x'_k + a_k) = n$$

$$\Leftrightarrow x_1' + x_2' + \dots + x_k' = n - a_1 - a_2 - \dots - a_k$$

con  $x_i' \ge 0$ . Así que hemos reducido el problema al caso más simple con  $x_i' \ge 0$  y nuevamente podemos aplicar el teorema de estrellas y barras.

# 11.6 DP Optimization Theory

Name	Original Recurrence	Sufficient Condition	From	To
CH 1	$dp[i] = min_{j < i} \{dp[j] + b[j] *$	$b[j] \ge b[j+1]$ Option-	$O(n^2)$	O(n)
	$a[i]\}$	ally $a[i] \le a[i+1]$		
CH 2	$dp[i][j] = min_{k < j} \{ dp[i - ]$	$b[k] \ge b[k+1]$ Option-	$O(kn^2)$	O(kn)
	1][k] + b[k] * a[j]	ally $a[j] \le a[j+1]$		
D&Q	$dp[i][j] = min_{k < j} \{ dp[i - ]$	$A[i][j] \le A[i][j+1]$	$O(kn^2)$	$O(kn\log n$
	$1][k] + C[k][j]\}$			
Knuth	dp[i][j] =	$A[i, j-1] \le A[i, j] \le$	$O(n^3)$	$O(n^2)$
	$min_{i < k < j} \{dp[i][k] +$	A[i+1,j]		
	$dp[k][j]\} + C[i][j]$			

Notes:

- A[i][j] the smallest k that gives the optimal answer, for example in dp[i][j] = dp[i-1][k] + C[k][j]
- C[i][j] some given cost function
- We can generalize a bit in the following way  $dp[i] = \min_{j < i} \{F[j] + b[j] * a[i]\},$  where F[j] is computed from dp[j] in constant time