# Notebook UNTreeCiclo

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# 1

# 1.1 C++ plantilla

```
#include <bits/stdc++.h>
using namespace std;
#define watch(x) cout<<#x<<"="<<x<<'\n'
#define sz(arr) ((int) arr.size())
#define all(v) v.begin(), v.end()
typedef long long li;
typedef long double ld;
typedef pair<int, int> ii;
typedef vector<ii> vii;
typedef vector<int> vi;
typedef vector<long long> v1;
typedef pair<11, 11> pll;
typedef vector<pll> vll;
const int INF = 1e9;
const ll INFL = 1e18;
const int MOD = 1e9+7;
const double EPS = 1e-9;
const ld PI = acosl(-1);
int dirx[4] = \{0, -1, 1, 0\};
int diry[4] = \{-1, 0, 0, 1\};
int dr[] = \{1, 1, 0, -1, -1, -1, 0, 1\};
int dc[] = \{0, 1, 1, 1, 0, -1, -1, -1\};
const string ABC = "abcdefghijklmnopgrstuvwxyz";
const char \bar{l}n = ' \backslash n';
int main() {
         ios::sync_with_stdio(false);
         cin.tie(0);
         cout << setprecision(20) << fixed;</pre>
    // freopen("file.in", "r", stdin);
// freopen("file.out", "w", stdout);
         return 0;
```

#### 1.2 Librerias

```
// En caso de que no sirva #include <bits/stdc++.h>
#include <algorithm>
#include <iostream>
#include <iterator>
#include <sstream>
#include <fstream>
#include <cassert>
#include <climits>
#include <cstdlib>
#include <cstring>
#include <string>
#include <cstdio>
#include <vector>
#include <cmath>
#include <queue>
#include <deque>
#include <stack>
#include <list>
#include <map>
#include <set>
#include <bitset>
#include <iomanip>
#include <unordered map>
////
#include <tuple>
#include <random>
#include <chrono>
```

#### 1.3 Bitmask

```
// Todas son O(1) Representacion
int a = 5; // Representacion binaria: 0101
int b = 3; // Representacion binaria: 0011
// Operaciones Principales
int resultado_and = a & b; // 0001 (1 en decimal)
int resultado or = a | b; // 0111 (7 en decimal)
int resultado xor = a ^ b; // 0110 (6 en decimal)
int num = 42; // Representacion binaria: 00101010
bitset<8> bits(num); // Crear un objeto bitset a partir
   del numero
cout << "Secuencia de bits: " << bits << "\n";</pre>
bits.count(); // Cantidad de bits activados
bits.set(3, true); // Establecer el cuarto bit en 1
bits.reset(6); // Establecer el septimo bit en 0
11 S,T;
// Operaciones con bits (/*) por 2 (redondea de forma
   automatica)
S=34; // == 100010
S = S << 1; // == S * 2 == 68 == 1000100
```

```
S = S >> 2; // == S/4 == 17 == 10001
S = S >> 1; // == S/2 == 8 == 1000
// Encender un bit
S = 34;
S = S | (1 << 3); // S = 42 (101010)
// Limpiar o apagar un bit
// ~: Not operacion
S = 42;
S &= (1 << 1); // S = 40 (101000)
// Comprobar si un bit esta encendido
S = 42;
T = S&(1<<3); // (!= 0): el tercer bit esta encendido
// Invertir el estado de un bit
S = 40;
S = (1 << 2); // 44 (101100)
// LSB (Primero de la derecha)
S = 40;
T = ((S) & -(S)); // 8 (001000)
__builtin_ctz(T); // nos entrega el indice del LSB
// Encender todos los bits
11 n = 3; // el tamanio del set de bits
S = 0;
S = (1 << n) - 1; // 7 (111)
// n es el tamanio de la mask (Alternativa)
// 11 n = 64:
// for (11 subset = 0; subset < (1<<n); ++subset) {
// Enumerar todos los posibles subsets de un bitmask
int mask = 18:
for (int subset = mask; subset; subset = (mask & (subset
   -1)))
    cout << subset << "\n";</pre>
// otras funciones de c++
__builtin_popcount(32); // 100000 (base 2), only 1 bit is
__builtin_popcount(30);// 11110 (base 2), 4 bits are on
__builtin_popcountl((11<<62)-11); // 2^62-1 has 62 bits
   on (near limit)
__builtin_ctz(32); // 100000 (base 2), 5 trailing zeroes
__builtin_ctz(30); // 11110 (base 2), 1 trailing zero
__builtin_ctzl(11<<62); // 2^62 has 62 trailing zeroes
```

### 1.4 Cosas de strings

```
int conv(char ch) {return ((ch>='a' && ch<='z')?ch-'a':ch-
'A'+26);}
vector<string> split(string& s, char c=' ') {
```

```
vector<string> res;
stringstream ss(s);
string sub;
while(getline(ss, sub, c))res.push_back(sub);
return res;
}

for(char& c:s)c=toupper(c);
for(char& c:s)c=tolower(c);
int n=stoi(s); // de string a entero
int n=stoi(s, nullptr, 2); // base 2
double d=stod(s); // de string a double
string s=to_string(n); // de entero a string
```

# 1.5 Custom Hashing

```
struct custom hash {
        static long long splitmix64(long long x) {
                 x + = 0x9e3779b97f4a7c15;
                 x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;

x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
                 return x ^ (x >> 31);
        size_t operator()(long long x) const {
                 static const long long FIXED RANDOM =
                     chrono::steady clock::now().
                     time_since_epoch().count();
                 return splitmix64(x + FIXED RANDOM);
        size_t operator()(const pair<int,int>& x) const {
                 return (size t) x.first * 37U + (size t)
                    x.second;
        size_t operator()(const vector<int>& v) const {
                 size t s = 0;
                 for(auto &e : v)
                          s^=hash<int>()(e)+0x9e3779b9+(s
                             <<6)+(s>>2);
                 return s;
};
unordered map<long long, int, custom hash> safe map;
gp_hash_table<int, int, custom_hash> table;
```

### 1.6 Random

```
typedef unsigned long long u64;
mt19937_64 rng (chrono::steady_clock::now().
    time_since_epoch().count());
u64 hash=rng();
```

```
mt19937 rng (chrono::steady_clock::now().time_since_epoch
    ().count());
int rand(int a, int b) {return uniform_int_distribution <
    int>(a, b) (rng);} // uniform_real_distribution
```

### 2 Arboles

# 2.1 Centroid Decomposition

```
// 0(nlogn)
struct CentroidDecomposition{
        int dad[maxn],sz[maxn];
        set < int > adi[maxn]; // check, proc
        int operator[](int i){return dad[i];}
        void addEdge(int x,int y) {adj[x].insert(y);adj[y
           l.insert(x);}
        void build(int v=0, int p=-1) {
                int n=dfsSz(v, p);
                int centroid=dfsCentroid(v, p, n);
                dad[centroid]=p;
                for(int u:adj[centroid]) {
                         if (p==u) continue;
                         // count paths
                for(int u:adj[centroid]) {
                         adj[u].erase(centroid);
                         build(u,centroid);
                adj[centroid].clear();
        int dfsSz(int v,int p) {
                sz[v]=1;
                for(int u:adj[v]){
                         if (u==p) continue;
                         sz[v] += dfsSz(u, v);
                return sz[v];
        int dfsCentroid(int v, int p, int n) {
                for(int u:adj[v]){
                         if (u==p) continue;
                         if(sz[u]>n/2)return dfsCentroid(u
                            , v, n);
                return v;
// for (int b=a;b!=-1;b=cd[b])
```

### 2.2 Heavy Light Decomposition

```
ll null=LLONG MIN;
11 oper(ll a, ll b) {return max(a,b);}
// segtree build, set, upd, get
const int maxn=1e5+1;
bool VALS IN_EDGES=false; // arista padre
struct HLD{
        int par[maxn], root[maxn], dep[maxn];
        int sz[maxn], pos[maxn], ti;
        vi adi[maxn];
        SegTree st;
        void addEdge(int x, int y) {adj[x].push_back(y);
            adj[y].push_back(x);}
        void dfsSz(int x){
                sz[x] = 0;
                 for(int& y:adj[x]){
                         if (y==par[x]) continue;
                         par[y]=x; dep[y]=dep[x]+1;
                         dfsSz(y);
                         sz[x] + = sz[y] + 1;
                         if(sz[y]>sz[adj[x][0]])swap(y,adj
                             [x][0];
        void dfsHld(int x) {
                 pos[x]=ti++;
                 for(int y:adj[x]){
                 if (y==par[x]) continue;
                 root[y] = (y = adj[x][0]?root[x]:y);
                 dfsHld(y);
        void build(int n,int v=0) {
                 root[v]=par[v]=v;
                 dep[v]=ti=0;
                 dfsSz(v);
                 dfsHld(v);
                 // vl palst(n);
                 // for (int i=0; i < n; ++i) palst [pos[i]] = a[i
                 // st.build(palst, n);
                 st.build(n);
        template <class Oper>
        void processPath(int x, int y, Oper op) {
                 for(; root[x]!=root[y]; y=par[root[y]]){
                         if (dep[root[x]]>dep[root[v]]) swap
                             (x,y);
                         op(pos[root[y]],pos[y]);
                 if (dep[x]>dep[y]) swap(x,y);
                 op(pos[x]+VALS_IN_EDGES, pos[y]);
```

```
void modifyPath(int x, int y, int v) {
                processPath(x, y, [this, &v] (int 1, int r) {
                         st.upd(l,r,v);
                });
        ll queryPath(int x, int y) {
                ll res=null;
                processPath(x,y,[this,&res](int 1, int r)
                         res=oper(res, st.get(l,r));
                });
return res;
        void modifySubtree(int x, int v) {st.upd(pos[x]+
            VALS_IN_EDGES, pos[x]+sz[x], v);}
        int querySubtree(int x) {return st.get(pos[x]+
            VALS_IN_EDGES, pos[x]+sz[x]);
        void modify(int x, int v) {st.set(pos[x],v);}
        void modifyEdge(int x, int y, int v) {
                if (dep[x] < dep[y]) swap (x, y);
                modifv(x,v);
} ;
```

### 2.3 LCA

```
const int maxn = 2e5+5, maxlog = 20+5;
int up[maxn][maxlog], dep[maxn]; // memset -1, 0
vi adj[maxn];
int n; // <-
void dfs(int v, int p=-1) {
         up[v][0]=p;
        for(int u:adj[v]){
                  if(u!=p){
                           dep[u]=dep[v]+1;
                           dfs(u, v);
void build() {
         for (int l=1; l < maxlog; ++1) {</pre>
                  for (int i=0; i < n; ++i) {</pre>
                           if (up [i] [1-1] ==-1) continue;
                           up[i][l]=up[up[i][l-1]][l-1];
int kth(int node, int k){
         for (int l=maxlog-1; l>=0; --1) {
                  if (node!=-1 \&\& k\& (1 << 1)) {
                           node=up[node][1];
```

```
2.4 Sack
```

```
2 ARBOLES
```

### 2.4 Sack

```
const int maxn = 1e5+5;
int ver[2*maxn], st[maxn], ft[maxn];
int len[maxn], dep[maxn];
vi adj[maxn];
int n,pos=0;
void prec(int v=0, int p=-1) {
        len[v]=1;ver[pos]=v;
        st[v]=pos++;
        for(int u:adi[v]){
                if (u==p) continue;
                dep[u]=dep[v]+1;
                prec(u,v);
                len[v]+=len[u];
        ver[pos]=v;
        ft[v]=pos++;
bool vis[maxn];
void ask(int v, bool add) {
        if(vis[v] && !add) // delete node
        else if(!vis[v] && add)// add node
void dfs(int v=0, int p=-1, bool keep=true) {
        int mx=0,id=-1;
        for(int u:adj[v]){
                if (u==p) continue;
                if(len[u]>mx){
                        mx=len[u];
                        id=u;
        for(int u:adj[v]){
```

#### 2.5 Virtual Tree

```
const int maxn = 2e5+5;
vi adjVT[maxn],adj[maxn];
int st[maxn],ft[maxn];
bool important[maxn];
int n,q,pos=0;
void dfs(int v, int p=-1) {
        st[v]=++pos;
        for(int u:adj[v]){
                 if (u==p) continue;
                 dfs(u, v);
        ft[v]=pos;
bool upper(int v, int u){return st[v]<=st[u] && ft[v]>=ft
   [u]; }
bool cmp(int v, int u) {return st[v] < st[u]; }</pre>
// 0(klogk)
int virtualTree(vi nodes){
        sort(all(nodes), cmp);
        int m=sz(nodes);
        for (int i=0; i < m-1; ++i) {</pre>
                 int v=lca(nodes[i], nodes[i+1]);
                 nodes.push back(v);
        sort(all(nodes), cmp);
        nodes.erase(unique(all(nodes)), nodes.end());
        for(int u:nodes)adjVT[u].clear();
        vector<int> s;
        s.push_back(nodes[0]);
        m=sz(nodes);
        for (int i=1; i < m; ++i) {</pre>
                 int v=nodes[i];
                 while (sz(s) \ge 2 \&\& !upper(s.back(), v)) {
```

### 3 Estructuras de Datos

### 3.1 Disjoint Set Union

```
struct dsu{
        vi p, size;
        int sets, maxSize;
        dsu(int n) {
                p.assign(n,0);
                size.assign(n,1);
                sets = n;
                for (int i = 0; i < n; i++) p[i] = i;
        int find_set(int i) {return (p[i] == i) ? i : (p[
           i] = find set(p[i]));
        bool is_same_set(int i, int j) {return find_set(i
           ) == find set(j);}
        void unionSet(int i, int j) {
                if (!is_same_set(i, j)) {
                         int a = find set(i), b = find set
                         if (size[a] < size[b]) swap(a, b)</pre>
                         p[b] = a;
                         size[a] += size[b];
                         maxSize = max(size[a], maxSize);
                         sets--;
};
```

### 3.2 Dynamic Connectivity Offline

```
struct dsu{
    vi p,rank,h;
```

```
int sets;
        dsu(int n) {
                sets=n;
                p.assign(n,0);
                rank.assign(n,1);
                for(int i=0;i<n;++i)p[i]=i;
        int get(int a) {return (a==p[a]?a:get(p[a]));}
        void unite(int a, int b) {
                a=get(a); b=get(b);
                if (a==b) return;
                if(rank[a]>rank[b])swap(a,b);
                rank[b]+=rank[a];
                h.push_back(a);
                p[a]=b; sets--;
        void rollback(int x) {
                int len=h.size();
                while(len>x) {
                         int a=h.back();
                         h.pop back();
                         rank[p[a]]-=rank[a];
                         p[a]=a; sets++; len--;
} ;
enum { ADD, DEL, QUERY };
struct Query{int type, u, v;};
struct DynCon{
        vector<Query> q;
        dsu uf; vi mt;
        map<pair<int,int>, int> prv;
        DynCon(int n): uf(n){}
        void add(int i, int j) {
                if(i>j)swap(i, j);
                q.push_back({ADD, i, j});
                mt.push\_back(-1);
                prv[{i,j}]=sz(q)-1;
        void remove(int i, int j) {
                if(i > j) swap(i, j);
                q.push_back({DEL, i, j});
                int pr=prv[{i, j}];
                mt[pr]=sz(q)-1;
                mt.push back(pr);
        void query() {q.push_back({QUERY, -1, -1}); mt.
           push\_back(-1);
        void process() { // answers all queries in order
                if(!sz(q))return;
                for(int i=0; i<sz(q);++i){
                         if(q[i].type==ADD && mt[i]<0)mt[i
                             =sz(q);
                go(0, sz(q));
```

```
void go(int s, int e){
                 if(s+1==e){
                 if(q[s].type == QUERY)cout<<uf.sets<<"\n"</pre>
                 return;
                  } int k=sz(uf.h), m=(s+e)/2;
                  for(int i=e-1;i>=m;--i){
                 if(mt[i] \ge 0 \&\& mt[i] \le 0 uf.unite(q[i].u, q)
                      [i].v);
                 }qo(s, m);
                 uf.rollback(k);
                 for (int i=m-1; i>=s; --i) {
                 if (mt[i]>=e) uf.unite(g[i].u, g[i].v);
                  }go(m, e);
                 uf.rollback(k);
};
```

### 3.3 Dynamic Segment Tree

```
T null=0, nolz=0;
T oper(T a, T b);
struct Node{
         T val, lz;
         int 1, r;
         Node *pl,*pr;
         Node(int ll, int rr) {
                  val=null; lz=nolz;
                  pl=pr=nullptr;
                  l=11; r=rr;
};
typedef Node* PNode;
void update(PNode x) {
         if (x->r-x->l==1) return;
         x-val=oper(x-pl-val,x-pr-val);
void extends(PNode x){
         if (x->r-x->1!=1 \&\& !x->p1) {
                  int m = (x - > r + x - > 1) / 2;
                  x->pl=new Node(x->l, m);
                  x \rightarrow pr = new Node(m, x \rightarrow r);
void propagate(PNode x){
         if (x->r-x->l==1) return:
         if(x->lz==nolz) return;
         int m=(x->r+x->1)/2;
         // pl, pr
         x \rightarrow lz = nolz;
```

```
struct SegTree{
        PNode root:
        void upd(PNode x, int 1, int r, T v){
                int 1x=x->1, rx=x->r;
                if(lx>=r || l>=rx)return;
                if(lx>=1 && rx<=r){
                         // val, 1z
                         return;
                extends (x);
                propagate(x);
                upd(x->pl,l,r,v);
                upd (x->pr, l, r, v);
                update(x);
        T get(PNode x, int 1, int r){
                int lx=x->l, rx=x->r;
                if(lx>=r || l>=rx) return null;
                if(lx>=1 && rx<=r) return x->val;
                extends (x);
                propagate(x);
                T v1=qet(x->pl,l,r);
                T v2=qet(x->pr,l,r);
                return oper (v1, v2);
        T get(int 1, int r) {return get(root,1,r+1);}
        void upd(int 1, int r, T v) {upd(root, 1, r+1, v);}
        void build(int 1, int r) {root=new Node(1, r+1);}
} ;
```

### 3.4 Fenwick Tree

```
typedef long long T;
struct FwTree{
        int n;
        vector<T> bit;
        FwTree(int n): n(n),bit(n+1){}
        T get(int r) {
                 T sum=0;
                 for(++r;r;r-=r&-r)sum+=bit[r];
                 return sum;
        T get(int 1, int r) {return get(r)-(l==0?0:get(l))
            -1));}
        void upd(int r, T v) {
                 for (++r; r<=n; r+=r&-r) bit [r] +=v;
};
struct FwTree2d{
        int n, m;
        vector<vector<T>> bit;
```

#### 3.5 Link Cut Tree

};

```
typedef long long T;
struct SplayTree{
        struct Node {
                int ch[2]={0, 0},p=0;
                T val=0, path=0, sz=1;
                                         // Path
                T sub=0, vir=0, ssz=0, vsz=0; // Subtree
                bool flip=0;T lz=0;
                                         // Lazy
        vector<Node> ns:
        SplayTree(int n):ns(n+1){}
        T path(int u) {return (u?ns[u].path:0);}
        T size(int u) {return (u?ns[u].sz:0);}
        T subsize(int u) {return (u?ns[u].ssz:0);}
        T subsum(int u) {return (u?ns[u].sub:0);}
        void push(int x) {
                if(!x)return;
                int l=ns[x].ch[0],r=ns[x].ch[1];
                if(ns[x].flip) {
                         ns[l].flip^=1,ns[r].flip^=1;
                         swap(ns[x].ch[0], ns[x].ch[1]);
                            // check with st oper
                         ns[x].flip=0;
                if(ns[x].lz){
                         ns[x].sub+=ns[x].lz*ns[x].ssz;
                         ns[x].vir+=ns[x].lz*ns[x].vsz;
                         // ...
        void pull(int x) {
                int l=ns[x].ch[0], r=ns[x].ch[1];
                push(1); push(r);
```

```
ns[x].sz=size(1)+size(r)+1;
                 ns[x].path=max({path(1), path(r), ns[x].}
                     val });
                 ns[x].sub=ns[x].vir+subsum(1)+subsum(r)+
                     ns[x].val;
                 ns[x].ssz=ns[x].vsz+subsize(1)+subsize(r)
        void set(int x, int d, int y) {ns[x].ch[d]=y;ns[y
            ].p=x;pull(x);}
        void splay(int x) {
                 auto dir=[&](int x) {
                          int p=ns[x].p;if(!p)return -1;
                          return ns[p].ch[0] == x?0:ns[p].ch
                              [1] == x?1:-1;
                 };
                 auto rotate=[&](int x){
                          int y=ns[x].p, z=ns[y].p, dx=dir(x)
                              , dy = dir(y);
                          set (y, dx, ns[x].ch[!dx]);
                          set (x, !dx, y);
                          if(^{\circ}dy) set (z, dy, x);
                          ns[x].p=z;
                 };
                 for (push (x); ~dir(x);) {
                          int y=ns[x].p, z=ns[y].p;
                          push(z); push(y); push(x);
                          int dx=dir(x), dy=dir(y);
                          if(^{\circ}dv) rotate (dx!=dv?x:v);
                          rotate(x);
};
struct LinkCut:SplayTree{ // 1-indexed
        LinkCut(int n):SplayTree(n){}
        int root(int u){
                 access(u); splay(u); push(u);
                 while (ns[u].ch[0]) \{u=ns[u].ch[0]; push(u)\}
                 return splay(u),u;
        int parent(int u){
                 access (u); splay (u); push (u);
                 u=ns[u].ch[0];push(u);
                 while (ns[u].ch[1]) \{u=ns[u].ch[1]; push(u)\}
                 return splay(u),u;
        int access(int x) {
                 int u=x, v=0;
                 for (; u; v=u, u=ns[u].p) {
                          splay(u);
```

```
int& ov=ns[u].ch[1];
                 ns[u].vir+=ns[ov].sub;
                 ns[u].vsz+=ns[ov].ssz;
                 ns[u].vir-=ns[v].sub;
                 ns[u].vsz-=ns[v].ssz;
                 ov=v; pull(u);
        return splay(x), v;
void reroot(int x){
        access(x); ns[x].flip^=1; push(x);
void link(int u, int v) { // u \rightarrow v
        reroot(u);
        access(v);
        ns[v].vir+=ns[u].sub;
        ns[v].vsz+=ns[u].ssz;
        ns[u].p=v;pull(v);
void cut(int u, int v){
        int r=root(u);
        reroot (u);
        access(v);
        ns[v].ch[0]=ns[u].p=0;pull(v);
        reroot(r);
void cut(int u){ // cut parent
        access(u);
        ns[ns[u].ch[0]].p=0;
        ns[u].ch[0]=0;pull(u);
int lca(int u, int v) {
        if (root (u) !=root (v)) return -1;
        access(u); return access(v);
int depth(int u){
        access (u); splay (u); push (u);
        return ns[u].sz;
T path(int u, int v) {
        int r=root(u);
        reroot (u); access (v); pull (v);
        T ans=ns[v].path;
        return reroot(r), ans;
void set(int u, T val){access(u);ns[u].val=val;
   pull(u);}
void upd(int u, int v, T val){
        int r=root(u);
        reroot (u); access (v); splay (v);
```

```
// lazv
                reroot(r);
        T comp_size(int u) {return ns[root(u)].ssz;}
        T subtree size(int u) {
                int p=parent(u);
                if(!p)return comp size(u);
                cut(u);int ans=comp_size(u);
                link(u,p); return ans;
        T subtree_size(int u, int v) {
                int r=root(u);
                reroot (v); access (u);
                T ans=ns[u].vsz+1;
                return reroot(r), ans;
        T comp_sum(int u) {return ns[root(u)].sub;}
        T subtree sum(int u) {
                int p=parent(u);
                if(!p)return comp sum(u);
                cut(u); T ans=comp sum(u);
                link(u,p); return ans;
        T subtree sum(int u, int v) { // subtree of u, v
           father
                int r=root(u);
                reroot (v); access (u);
                T ans=ns[u].vir+ns[u].val; // por el
                    reroot
                return reroot(r), ans;
};
```

# 3.6 Mos Algorithm

```
// O((n+q)*sqrt(n))
int sqrtn,n;
struct query {int 1, r, idx;};
bool cmp (query a, query b) {
    int x = a.l/sqrtn;
    if (x != b.l/sqrtn) return x < b.l/sqrtn;
    return (x&1 ? a.r < b.r : a.r > b.r);
}

vector<int> ans;
vector<query> q;
int act();
void add(int i);
void remove(int i);

void solve() {
    sqrtn=(int)ceil(sqrt(n));
```

```
sort(all(q), cmp);
int l=0,r=-1;
for(int i=0;i<sz(q);++i) {
    while(r<q[i].r)add(++r);
    while(l>q[i].l)add(--l);
    while(r>q[i].r)remove(r--);
    while(l<q[i].l)remove(l++);
    ans[q[i].idx]=act();
}</pre>
```

#### 3.7 Ordered set

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template<typename T> using ordered_set = tree<T,</pre>
   null_type,less<T>, rb_tree_tag,
   tree order statistics node update>;
template<typename T> using ordered multiset = tree<T,</pre>
   null_type, less_equal<T>, rb_tree_tag,
   tree_order_statistics_node_update>;
// ----- CONSTRUCTOR ----- //
// 1. Para ordenar por MAX cambiar less<int> por greater<
   int>
// 2. Para multiset cambiar less<int> por less equal<int>
      Para borrar siendo multiset:
       int idx = st.order_of_key(value);
      st.erase(st.find_by_order(idx));
// ---- METHODS ---- //
st.find_by_order(k) // returns pointer to the k-th
   smallest element
st.order of key(x) // returns how many elements are
   smaller than x
st.find_by_order(k) == st.end() // true, if element does
   not exist
```

# 3.8 Persistent Segment Tree

```
int newNode(T x){
                Node tmp=\{x, -1, -1\};
                ns.push_back(tmp);
                return act++;
        int newNode(int 1, int r){
                Node tmp={null, l, r};
                ns.push_back(tmp);
                update(act);
                return act++;
        int build(vector<T>& a, int 1, int r){
                if(r-l==1) {return newNode(a[1]);}
                int m = (1+r)/2;
                return newNode (build (a, l, m), build (a, m,
                     r));
        int set(int x, int i, T v, int l, int r){
                if (r-l==1) return newNode (v);
                int m = (1+r)/2;
                if (i<m) return newNode (set (ns[x].l, i, v,</pre>
                    1, m), ns[x].r);
                 else return newNode(ns[x].l, set(ns[x].r,
                     i, v, m, r));
        T get(int x, int lx, int rx, int l, int r){
                 if(lx>=r || l>=rx) return null;
                if(lx>=l && rx<=r) return ns[x].val;</pre>
                int m = (1x+rx)/2;
                T v1=qet(ns[x].l, lx, m, l, r);
                T v2 = qet(ns[x].r, m, rx, l, r);
                return oper (v1, v2);
        T get(int 1, int r, int time) {return get(roots[
            time], 0, size, 1, r+1);}
        void set(int i, T v, int time){roots.push_back(
            set(roots[time], i, v, 0, size));}
        void build(vector<T>& a, int n){size=n;roots.
            push back(build(a, 0, size));}
} ;
```

### 3.9 RMQ

```
typedef long long T;
T oper(T a, T b); // max, min, gcd ...
struct RMQ {
    vector<vector<T>> table;
    void build(vector<T>& v) {
```

### 3.10 Segment Tree Iterativo

```
struct segtree{
    int n; vl v; ll nulo = 0;
    ll op(ll a, ll b) {return a + b;}
    segtree (int n) : n(n), v(2*n, nulo) {}
    segtree(vl &a): n(sz(a)), v(2*n){
        for (int i = 0; i < n; i++) v[n + i] = a[i];
        for (int i = n-1; i > = 1; --i) v[i] = op(v[i << 1], v
            [i<<1|1]);
    void upd(int k, ll nv) {
        for (v[k += n] = nv; k > 1; k >>= 1) v[k>>1] = op
            (v[k], v[k^1]);
    ll get(int l, int r){
        11 \text{ vl} = \text{nulo}, \text{ vr} = \text{nulo};
        for (1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1)
             if (1&1) vl = op(vl, v[1++]);
             if (r\&1) vr = op(v[--r], vr);
        return op (vl, vr);
};
```

# 3.11 Segment Tree Recursivo

```
typedef long long T;
struct SegTree{
    vector<T> vals,lazy;
    T null=0,nolz=0;
```

```
int size;
T oper(T a, T b);
void build(vector<T>& a, int x, int lx, int rx) {
        if(rx-lx==1){
                 if(lx<sz(a))vals[x]=a[lx];
                 return:
        int m = (1x+rx)/2;
        build(a, 2*x+1, 1x, m);
        build(a, 2*x+2, m, rx);
        vals[x]=oper(vals[2*x+1], vals[2*x+2]);
void build(vector<T>& a, int n) {
        size=1:
        while (size<n) size *= 2;</pre>
        vals.resize(2*size);
        lazv.assign(2*size, nolz);
        build(a, 0, 0, size);
void propagate(int x, int lx, int rx){
        if (rx-lx==1) return;
        if (lazy[x]==nolz) return;
        int m = (1x+rx)/2;
        // 2*x+1, 2*x+2 (lazy, vals)
        lazy[x]=nolz;
void upd(int 1, int r, T v,int x, int lx, int rx)
        if(lx>=r || l>=rx)return;
        if(lx>=l && rx<=r){
                 // lazv, vals
                 return;
        propagate(x,lx,rx);
        int m = (1x+rx)/2;
        upd(1, r, v, 2*x+1, 1x, m);
        upd(1, r, v, 2 \times x + 2, m, rx);
        vals[x] = oper(vals[2*x+1], vals[2*x+2]);
void set(int i, T v, int x, int lx, int rx){
        if(rx-lx==1){
                 vals[x]=v;
                 return;
        propagate(x,lx,rx);
        int m = (1x+rx)/2;
        if (i<m) set (i, v, 2*x+1, lx, m);
        else set(i,v,2*x+2,m,rx);
        vals[x]=oper(vals[2*x+1], vals[2*x+2]);
T get (int 1, int r, int x, int lx, int rx) {
```

```
if(lx>=r || l>=rx)return null;
                 if(lx>=1 && rx<=r) return vals[x];</pre>
                 propagate (x, lx, rx);
                 int m = (lx+rx)/2;
                 T v1=qet (1, r, 2*x+1, 1x, m);
                 T v2=qet(1,r,2*x+2,m,rx);
                 return oper (v1, v2);
        T get(int 1, int r) {return get(1,r+1,0,0,size);}
        void upd(int 1, int r, T v) {upd(1,r+1,v,0,0,size)
        void set(int i, T val) {set(i,val,0,0,size);}
};
```

### 3.12 Segment Tree 2D

```
const int N=1000+1;
ll st[2*N][2*N];
struct SegTree{
        int n,m,neutro=0;
        inline ll op(ll a, ll b) {return a+b;}
        SegTree(int n, int m): n(n), m(m) {
                 for(int i=0;i<2*n;++i)for(int j=0;j<2*m</pre>
                     ;++i)st[i][i]=neutro;
        SegTree(vector\langle vi \rangle \& a): n(sz(a)), m(n ? sz(a[0])
            : 0) { build(a); }
        void build(vector<vi>& a) {
                 for (int i=0; i< n; ++i) for (int j=0; j< m; ++j)
                     st[i+n][j+m]=a[i][j];
                 for(int i=0;i<n;++i) for(int j=m-1;j>=1;--
                     j) st[i+n][j]=op(st[i+n][j<<1], st[i+n
                     ][ † << 1 | 1 ] );
                 for (int i=n-1; i>=1; --i) for (int j=0; j<2*m
                     ;++j)st[i][j]=op(st[i<<1][j], st[i
                     <<1|1|[i]);
        void upd(int x, int y, ll v){
                 st[x+n][y+m]=v;
                 for (int j=y+m; j>1; j>>=1) st [x+n] [j>>1] =op (
                     st[x+n][i], st[x+n][i^1];
                 for (int i=x+n; i>1; i>>=1) for (int j=y+m; j; j
                     >>=1) st[i>>1][j]=op(st[i][j], st[i^1][
                     j]);
        ll get (int x0, int y0, int x1, int y1) {
                 ll r=neutro:
                 for(int i0=x0+n,i1=x1+n+1;i0<i1;i0>>=1,i1
                     >>=1) {
```

```
int t[4], q=0;
                            if (i0&1) t[q++]=i0++;
                            if (i1&1) t [q++]=--i1;
                            for (int k=0; k < q; ++k) for (int j0=y0
                                 +m, j1=y1+m+1; j0<j1; j0>>=1, j1
                                >>=\bar{1}) {
                                      if(j0&1) r = op(r, st[t[k]][
                                          j0++1);
                                      if(j1&1) r=op(r, st[t[k]
                                          ]][-- | 1]);
                   return r;
} ;
```

### 3.13 Segment Tree Beats

```
typedef long long T;
T null=0, noVal=0;
T INF=1e18;
struct Node{
        T sum, lazv:
        T max1, max2, maxc;
        T min1, min2, minc;
struct SegTree{
        vector<Node> vals;int size;
        void oper(int a, int b, int c); // node c, left a
            , right b;
        Node single (T x) {
                 Node tmp;
                 tmp.sum=tmp.max1=tmp.min1=x;
                 tmp.maxc=tmp.minc=1;
                 tmp.lazv=noVal;
                 tmp.max2=-INF;
                 tmp.min2=INF;
                 return tmp;
        void build(vector<T>& a, int n);
        void propagateMin(T v, int x, int lx, int rx){
                 if (vals[x].max1<=v) return;</pre>
                 vals[x].sum-=vals[x].max1*vals[x].maxc;
                 vals[x].max1=v;
                 vals[x].sum+=vals[x].max1*vals[x].maxc;
                 if(rx-lx==1){
                         vals[x].min1=v;
                 }else{
                         if(v<=vals[x].min1) {</pre>
                                  vals[x].min1=v;
                         }else if(v<vals[x].min2){</pre>
                                  vals[x].min2=v;
```

```
void propagateAdd(T v, int x, int lx, int rx){
        vals[x].sum+=v*((T)(rx-lx));
        vals[x].lazy+=v;
        vals[x].max1+=v;
        vals[x].min1+=v;
        if (vals[x].max2!=-INF) vals[x].max2+=v;
        if (vals[x].min2!=INF) vals[x].min2+=v;
void propagate(int x, int lx, int rx) {
        if (rx-lx==1) return;
        int m = (lx + rx)/2;
        if(vals[x].lazy!=noVal){
                propagateAdd(vals[x].lazy, 2*x+1,
                     lx, m);
                propagateAdd(vals[x].lazy, 2*x+2,
                     m, rx);
                vals[x].lazy=noVal;
        propagateMin(vals[x].max1, 2*x+1, lx, m);
        propagateMin(vals[x].max1, 2*x+2, m, rx);
void updAdd(int 1, int r, T v, int x, int lx, int
   rx) {
        if(lx>=r || l>=rx) return;
        if(lx>=1 && rx<=r){
                propagateAdd(v, x, lx, rx);
                 return;
        propagate(x,lx,rx);
        int m = (lx+rx)/2;
        updAdd(l,r,v,2*x+1,lx,m);
        updAdd(1,r,v,2*x+2,m,rx);
        oper (2*x+1, 2*x+2, x);
void updMin(int 1, int r, T v, int x, int lx, int
   rx) {
        if(lx>=r || l>=rx || vals[x].max1<v)
            return:
        if(lx>=1 && rx<=r && vals[x].max2<v){</pre>
                propagateMin(v, x, lx, rx);
                return;
        propagate(x,lx,rx);
        int m = (1x+rx)/2;
        updMin(l,r,v,2*x+1,lx,m);
        updMin(l,r,v,2*x+2,m,rx);
        oper (2*x+1, 2*x+2, x);
void updAdd(int 1, int r, T v) {updAdd(1,r+1,v)
   ,0,0,size);}
```

### 3.14 Sqrt Descomposition

```
// O(n/sqrt(n)+sqrt(n)) query
struct Sqrt {
        int block_size;
         vl nums, blocks;
         Sgrt (vi &arr) {
                 block_size=(int)ceil(sqrt(sz(arr)));
                 blocks.assign(block_size, 0);
                 nums=arr;
                 for (int i=0; i < sz (nums); ++i) {</pre>
                          blocks[i/block size] += nums[i];
         void update(int x, int v) {
                 blocks[x/block_size] -= nums[x];
                 nums [x]=v;
                 blocks[x/block size] += nums[x];
         11 query(int r) {
                  11 \text{ res}=0;
                  for(int i=0;i<r/block size;++i) {res+=</pre>
                     blocks[i];}
                  for(int i=(r/block size)*block size;i<r</pre>
                     ; ++i) { res+=nums[i]; }
                 return res;
         ll query(int 1, int r) {return query(r)-query(l-1)
            ; }
} ;
```

# 3.15 Treap

```
// treap => order asc, implicit treap => order array
typedef long long T;
struct Treap{
    Treap *1,*r,*dad;
    u64 prior;
    T sz,value,sum,lz;
    Treap(T v) {
        l=r=nullptr;
        lz=0;sz=1;
        prior=rng();
        value=sum=v;
}
```

```
~Treap() {delete l;delete r;}
};
typedef Treap* PTreap;
T cnt (PTreap x) {return (!x?0:x->sz);}
T sum(PTreap x) {return (!x?0:x->sum);}
void propagate(PTreap x) {
        if(x && x->1z) {
                 if(x->1); // lz, value, sum ...
                 if (x->r); // lz, value, sum ...
                 x - > 1z = 0;
void update(PTreap x) {
        propagate (x->1);
        propagate(x->r);
        x->sz=cnt(x->1)+cnt(x->r)+1;
        x \rightarrow sum = sum(x \rightarrow 1) + sum(x \rightarrow r) + x \rightarrow value;
        if (x->1) x->1->dad=x;
        if (x->r)x->r->dad=x;
void upd(PTreap x, T v) {
        if(!x)return;
        update(x);
         // lz, value, sum ...
// pair<PTreap, PTreap> split(PTreap x, T key) { // f <=
    kev < s
pair<PTreap, PTreap> split(PTreap x, int left){ // cnt(f)
         if(!x)return {nullptr, nullptr};
        propagate(x);
        if(cnt(x->1)>=left) { // if(x->value>key) {}
                  auto got=split(x->1, left); //, key);
                 x->l=qot.second;
                  update(x);
                 return {got.first, x};
         }else{
                  auto got=split(x->r, left-cnt(x->1)-1);
                     // , key);
                 x->r=qot.first;
                  update(x);
                 return {x, qot.second};
PTreap merge(PTreap x, PTreap y) {
        if(!x)return v;
        if(!y)return x;
        propagate(x);
        propagate(y);
        if (x->prior<=y->prior) {
                 x \rightarrow r = merge(x \rightarrow r, y);
```

```
update(x);
                return x;
        }else{
                y->1=merge(x, y->1);
                update(y);
                return y;
PTreap combine (PTreap x, PTreap y) {
        if(!x)return y;
        if(!y)return x;
        if (x->prior<y->prior) swap(x, y);
        auto z=split(y, x->value);
        x->r=combine(x->r, z.second);
        x->l=combine(z.first, x->l);
        return x;
T kth(PTreap& x, int k){ // indexed 0
        if(!x)return null;
        if (k==cnt (x->1)) return x->value;
        if (k < cnt(x->1)) return kth(x->1, k);
        return kth(x->r, k-cnt(x->1)-1);
pair<int, T> lower bound(PTreap x, T key) { // index,
   value
        if(!x)return {0, null};
        if(x->value<kev) {</pre>
                auto y=lower_bound(x->r, key);
                y.first+=cnt(x->1)+1;
                return v;
        auto y=lower bound(x->1, key);
        if (y.first==cnt(x->1))y.second=x->value;
        return v;
void dfs(PTreap x) {
        if(!x)return;
        propagate(x);
        dfs(x->1); cout << x-> value << ""; <math>dfs(x->r);
// PTreap root=nullptr;
// PTreap act=new Treap(c);
// root=merge(root, act);
```

#### 3.16 Two Stacks

```
void add(T x){
        Node tmp=\{x, x\};
        if(!s2.empty()){
        // tmp.acum + s2.top().acum
        s2.push(tmp);
void remove(){
        if(s1.empty()){
                while(!s2.empty()){
                        Node tmp=s2.top();
                        if(s1.emptv()){
                        // tmp.acum = tmp.val
                         }else{
                         // tmp.acum + s1.top().
                        s1.push(tmp);
                        s2.pop();
        s1.pop();
bool good() {
        if(s1.empty() && s2.empty())return false;
        else if(!s1.empty() && s2.empty()){
                return true; // eval sl.top();
        }else if(s1.empty() && !s2.empty()){
                return true; // eval s2.top();
        }else{
                return true; // eval s1.top() +
                    s2.top()
```

#### 3.17 Wavelet Tree

};

```
const int maxn = 1e5+5, maxv = 1e9, minv = -1e9;
struct WaveletTree{ // indexed 1 - O(nlogn)
    int lo, hi;
    WaveletTree *1, *r;
    int *b, bsz, csz;
    ll *c;

    WaveletTree() {
        hi=bsz=csz=0;
        l=r=NULL;
        lo=1;
    }

    void build(int *from, int *to, int x, int y) {
        lo=x, hi=y;
        if(from>=to)return;
```

```
int mid=lo+(hi-lo)/2;
        auto f=[mid](int x){return x<=mid;};</pre>
        b=(int*)malloc((to-from+2)*sizeof(int));
        bsz=0;
        b[bsz++]=0;
        c=(l1*)malloc((to-from+2)*sizeof(l1));
        c[csz++]=0;
        for(auto it=from; it!=to; ++it) {
                b[bsz] = (b[bsz-1] + f(*it));
                c[csz] = (c[csz-1] + (*it));
                bsz++;csz++;
        if (hi==lo) return;
        auto pivot=stable partition(from, to, f);
        l=new WaveletTree();
        1->build(from, pivot, lo, mid);
        r=new WaveletTree();
        r->build(pivot, to, mid+1, hi);
//kth smallest element in [1, r]
int kth(int 1, int r, int k){
        if(l>r)return 0;
        if(lo==hi)return lo;
        int inLeft=b[r]-b[1-1], lb=b[1-1], rb=b[r
        if (k<=inLeft) return this->l->kth(lb+1, rb
        return this->r->kth(l-lb, r-rb, k-inLeft)
//count of numbers in [1, r] Less than or equal
   to k
int lte(int l, int r, int k){
        if(1>r || k<10) return 0;
        if (hi<=k) return r-l+1;</pre>
        int lb=b[l-1], rb=b[r];
        return this->l->lte(lb+1, rb, k)+this->r
           ->lte(l-lb, r-rb, k);
//count of numbers in [l, r] equal to k
int count(int 1, int r, int k){
        if(l>r || k<lo || k>hi)return 0;
        if (lo==hi) return r-l+1;
        int lb=b[1-1], rb=b[r];
        int mid=(lo+hi)>>1;
        if (k<=mid) return this->l->count (lb+1, rb,
        return this->r->count(l-lb, r-rb, k);
//sum of numbers in [l ,r] less than or equal to
```

```
4 FLUJO
```

# 4 Flujos

#### 4.1 Dinic

```
// O(|E| * |V|^2)
struct edge { ll v, cap, inv, flow; };
struct network {
  ll n, s, t;
  vector<ll> lvl;
  vector<vector<edge>> q;
  network(ll n) : n(n), lvl(n), g(n) {}
  void add_edge(int u, int v, ll c) {
    g[u].push_back({v, c, g[v].size(), 0});
    g[v].push_back({u, 0, g[u].size()-1, c});
 bool bfs() {
    fill(lvl.begin(), lvl.end(), -1);
    queue<11> q;
    [vl[s] = 0;
    for (q.push(s); q.size(); q.pop()) {
      ll u = q.front();
      for(auto &e : g[u]) {
        if(e.cap > 0 && lvl[e.v] == -1) {
          lvl[e.v] = lvl[u]+1;
          q.push(e.v);
    return lvl[t] != -1;
  void min cut(){
    queue<11> q;
    vector<bool> vis(n, 0);
```

```
vis[s] = 1;
    for(q.push(s); q.size(); q.pop()) {
      ll u = q.front();
      for(auto &e : g[u]) {
        if(e.cap > 0 && !vis[e.v]) {
          a.push(e.v);
          vis[e.v] = 1;
    set<ii> ans;
    for (int i = 0; i<n; i++) {</pre>
        for (auto &e : q[i]) {
            if (vis[i] && !vis[e.v]){
                ans.insert(\{i+1, e.v+1\});
    for (auto [x, y] : ans) cout << x << ' ' << y << ln;
  ll dfs(ll u, ll nf) {
    if(u == t) return nf;
    11 \text{ res} = 0;
    for(auto &e : g[u]) {
      if(e.cap > 0 \&\& lvl[e.v] == lvl[u]+1) {
        ll tf = dfs(e.v, min(nf, e.cap));
        res += tf; nf -= tf; e.cap -= tf;
        g[e.v][e.inv].cap += tf;
        q[e.v][e.inv].flow -= tf;
        e.flow += tf;
        if(nf == 0) return res;
    if(!res) lvl[u] = -1;
    return res;
 ll max_flow(ll so, ll si, ll res = 0) {
    s = so; t = si;
    while(bfs()) res += dfs(s, LONG LONG MAX);
    return res;
};
```

# 4.2 Edmonds Karp

```
//o(V * E^2)
11 bfs(vector<vi> &adj, vector<vl> &capacity, int s, int
    t, vi& parent) {
    fill(parent.begin(), parent.end(), -1);
    parent[s] = -2;
    queue<pll> q;
    q.push({s, INFL});
    while (!q.empty()) {
```

```
int cur = q.front().first;
        11 flow = q.front().second;
        q.pop();
        for (int next : adj[cur]) {
            if (parent[next] == -1LL && capacity[cur][
                next]) {
                parent[next] = cur;
                ll new_flow = min(flow, capacity[cur][
                    next]);
                if (next == t)
                    return new flow;
                q.push({next, new flow});
    return 0;
11 maxflow(vector<vi> &adj, vector<vl> &capacity, int s,
   int t, int n) {
   11 \text{ flow} = 0;
    vi parent(n);
    ll new flow;
    while ((new_flow = bfs(adj, capacity, s, t, parent)))
        flow += new_flow;
        int cur = t;
        while (cur != s) {
            int prev = parent[cur];
            capacity[prev][cur] -= new_flow;
            capacity[cur][prev] += new flow;
            cur = prev;
    return flow;
```

# 4.3 Hopcroft Karp

```
bool bfs() {
    queue<int> q;
    for(int u = 0; u < 1; u++) {
      if (match[u] == nil) {
        d[u] = 0;
        a.push(u);
      } else d[u] = INF;
    d[nil] = INF;
    while(q.size()) {
      int \bar{u} = q.front(); q.pop();
      if(u == nil) continue;
      for(auto v : q[u]) {
        if(d[ match[v] ] == INF) {
   d[ match[v] ] = d[u]+1;
           q.push(match[v]);
    return d[nil] != INF;
  bool dfs(int u) {
    if(u == nil) return true;
    for(int v : q[u]) {
      if(d[match[v]] == d[u]+1 && dfs(match[v])) {
        match[v] = u; match[u] = v;
        return true;
    d[u] = INF;
    return false;
  int max matching() {
    int ans = 0;
    while(bfs()) {
      for(int u = 0; u < 1; u++) {</pre>
        ans += (match[u] == nil && dfs(u));
    return ans;
  void matchs() {
    for (int i = 0; i<1; i++) {
        if (match[i] == l+r) continue;
        cout << i+1 << ' ' << match[i]+1-l << ln;</pre>
};
```

# 4.4 Maximum Bipartite Matching

```
// O(|E|*|V|)
struct mbm {
   int 1, r;
```

```
vector<vector<int>> q;
 vector<int> match, seen;
 mbm(int 1, int r) : l(1), r(r), seen(r), match(r), q(1)
 void add_edge(int 1, int r) { g[1].push_back(r); }
 bool dfs(int u)
    for (auto v : q[u]) {
      if (seen[v]++) continue;
      if (match[v] == -1 || dfs(match[v])) {
        match[v] = u;
        return true;
    return false;
  int max matching() {
    int ans = 0;
    fill (match.begin(), match.end(), -1);
    for(int u = 0; u < 1; ++u) {
      fill(seen.begin(), seen.end(), 0);
      ans += dfs(u);
    return ans;
 void matchs() {
    for (int i = 0; i<r; i++) {
        if (match[i] == -1) continue;
        cout << match[i]+1 << ' ' << i+1 << ln;
};
```

### 4.5 Minimum Cost Maximum Flow

```
/// Complexity: O(|V|*|E|^2*log(|E|))
template <class type>
struct mcmf {
  struct edge { int u, v, cap, flow; type cost; };
  int n;
  vector<edge> ed;
  vector<vector<int>> q;
  vector<int> p;
  vector<type> d, phi;
  mcmf(int n) : n(n), q(n), p(n), d(n), phi(n) {}
  void add_edge(int u, int v, int cap, type cost) {
    q[u].push back(ed.size());
    ed.push back({u, v, cap, 0, cost});
    g[v].push back(ed.size());
    ed.push_back({v, u, 0, 0, -cost});
  bool dijkstra(int s, int t) {
    fill(d.begin(), d.end(), INF TYPE);
    fill(p.begin(), p.end(), -1);
```

```
set<pair<type, int>> q;
    d[s] = 0;
    for(q.insert({d[s], s}); q.size();) {
      int u = (*q.begin()).second; q.erase(q.begin());
      for(auto v : q[u]) {
        auto &e = ed[v];
        type nd = d[e.u]+e.cost+phi[e.u]-phi[e.v];
        if(0 < (e.cap-e.flow) && nd < d[e.v]) {
          q.erase({d[e.v], e.v});
          \bar{d}[e.v] = nd; p[e.v] = v;
          q.insert({d[e.v], e.v});
    for (int i = 0; i < n; i++) phi[i] = min(INF TYPE, phi
       [i]+d[i]);
    return d[t] != INF TYPE;
  pair<int, type> max_flow(int s, int t) {
    type mc = 0;
    int mf = 0;
    fill(phi.begin(), phi.end(), 0);
    while(dijkstra(s, t)) {
      int flow = INF;
      for (int v = p[t]; v != -1; v = p[ed[v].u])
        flow = min(flow, ed[v].cap-ed[v].flow);
      for (int v = p[t]; v != -1; v = p[ed[v].u]) {
        edge &e1 = ed[v];
        edge &e2 = ed[v^1];
        mc += e1.cost*flow;
        e1.flow += flow;
        e2.flow -= flow;
      mf += flow;
    return {mf, mc};
};
```

### 4.6 Weighted Matching

```
vector<type> v(r), d(r); // v: potential
vector\langle int \rangle ml(1, -1), mr(r, -1); // matching pairs
vector<int> idx(r), prev(r);
iota(idx.begin(), idx.end(), 0);
auto residue = [&](int i, int j) { return c[i][j]-v[j
for(int f = 0; f < 1; ++f) {
  for (int j = 0; j < r; ++j) {
    d[i] = residue(f, i);
    prev[j] = f;
  type w;
  int j, 1;
  for (int s = 0, t = 0;;) {
   if(s == t) {
     1 = s;
      w = d[ idx[t++] ];
      for (int k = t; k < r; ++k) {
        j = idx[k];
        type h = d[i];
        if (h <= w) {
          if (h < w) t = s, w = h;
          idx[k] = idx[t];
          idx[t++] = i;
      for (int k = s; k < t; ++k) {
        i = idx[k];
        if (mr[j] < 0) goto aug;
    int q = idx[s++], i = mr[q];
    for (int k = t; k < r; ++k) {
      j = idx[k];
      type h = residue(i, j) - residue(i, q) + w;
      if (h < d[j]) {
        d[j] = h;
        prev[j] = i;
        if(h == w) {
          if(mr[j] < 0) goto aug;</pre>
          idx[k] = idx[t];
          idx[t++] = i;
      }
  aug: for (int k = 0; k < 1; ++k)
   v[idx[k]] += d[idx[k]] - w;
  int i:
  do {
   mr[j] = i = prev[j];
    swap(j, ml[i]);
 } while (i != f);
type opt = 0;
```

### 5 Geometria

#### 5.1 Puntos

```
typedef long double lf;
const lf EPS = 1e-9;
const 1f E0 = 0.0L; //Keep = 0 for integer coordinates,
   otherwise = EPS
const lf PI = acos(-1);
struct pt {
    If x, y;
    pt(){}
    pt(lf a, lf b): x(a), y(b) {}
    pt operator - (const pt &q ) const { return {x - q.x
        , y - q.y }; }
    pt operator + (const pt &q ) const { return {x + q.x
        , y + q.y }; }
    pt operator * (const lf &t ) const { return {x * t ,
       y * t }; }
    pt operator / (const lf &t ) const { return {x / t ,
        y / t }; }
    bool operator == (pt p) \{ return abs(x - p.x) \le EPS \}
        && abs(v - p.v) <= EPS; }
    bool operator != (pt p) { return !operator==(p); }
    bool operator < ( const pt & q ) const {</pre>
        if (fabsl(x - q.x) > E0) return x < q.x;
        return y < q.y;</pre>
    void print() { cout << x << " " << y << "\n"; }</pre>
};
pt normalize(pt p) {
    lf norm = hypotl(p.x, p.y);
    if(fabsl(norm) > EPS) return {p.x /= norm, p.v /=
        norm };
    else return p;
int cmp(lf a, lf b) { return (a + EPS < b ? -1 : (b + EPS <</pre>
    a ? 1 : 0)); } // float comparator
// rota ccw
pt rot90(pt p) { return {-p.y, p.x}; }
// w(RAD)
pt rot(pt p, lf w) { return \{\cos l(w) * p.x - \sin l(w) * p.y\}
   *, sinl(w) * p.x + cosl(w) * p.y); }
lf norm2(pt p) { return p.x * p.x + p.y * p.y; }
```

```
5.2 Lineas
```

```
lf norm(pt p) { return hypotl(p.x, p.y); }
lf dis2(pt p, pt g) { return norm2(p - g); }
lf dis(pt p, pt q) { return norm(p - q); }
If arg(pt a) \{ return atan2(a.y, a.x); \} // ang(RAD) a x-
If dot(pt a, pt b) { return a.x * b.x + a.y * b.y; } // x
   = 90 -> \cos = 0
lf cross(pt a, pt b) { return a.x * b.y - a.y * b.x; } //
   x = 180 -> \sin = 0
lf orient(pt a, pt b, pt c) { return cross(b - a, c - a);
   } // clockwise = -
int sign(lf x) { return (lf(0) < x) - (x < lf(0)); }
// x inside angle abc (center in a)
bool in_angle(pt a, pt b, pt c, pt p) {
    //assert(fabsl(orient(a, b, c)) > E0);
    if(orient(a, b, c) < -E0)
        return orient(a, b, p) \geq= -E0 || orient(a, c, p)
    return orient(a, b, p) \geq -E0 && orient(a, c, p) \leq
// If angle(pt a, pt b) { return acos(max((lf)-1.0, min((
   1f)1.0, dot(a, b)/norm(a)/norm(b)))); } // min ang(RAD
lf angle(pt a, pt b) { return atan2(cross(a, b), dot(a, b)
   ); } // ang(RAD)
If angle (pt a, pt b, pt c) { // ang (RAD) AB AC ccw
    lf ang = angle(b - a, c - a);
    if (ang < 0) ang += 2 * PI;
    return ang;
bool half (pt p) { // true if is in (0, 180) (line is x
   axis)
    // assert (p.x != 0 \mid \mid p.y \mid = 0); // the argument of
        (0, 0) is undefined
    return p.v > 0 \mid | (p.v == 0 && p.x < 0);
bool half_from(pt p, pt v = \{1, 0\}) {
  return cross(v,p) < 0 \mid \mid (cross(v,p) == 0 \&\& dot(v,p) <
// polar sort
bool polar cmp (const pt &a, const pt &b) {
  return make tuple(half(a), 0) < make tuple(half(b),</pre>
     cross(a,b));
void polar_sort(vector<pt> &v, pt o) { // sort points in
   counterclockwise with respect to point o
    sort(v.begin(), v.end(), [&](pt a,pt b) {
```

```
return make_tuple(half(a - o), 0.0, norm2((a - o)
        )) < make_tuple(half(b - o), cross(a - o, b -
        o), norm2((b - o)));
});
</pre>
```

#### 5.2 Lineas

```
struct line {
    pt v; T c; // v:direction c: pos in y axis
    line(pt v, T c) : v(v), c(c) {}
    line(T a, T b, T c) : v(\{b,-a\}), c(c)\{\} // ax + by =
    line(pt p, pt q): v(q-p), c(cross(v,p)) {}
    T side(pt p) { return cross(v,p)-c; }
   lf dist(pt p) { return abs(side(p)) / abs(v); }
   lf sq dist(pt p) { return side(p) * side(p) / (lf) norm(
    line perp_through(pt p) { return {p, p + rot90ccw(v)}
       }; } // line perp to v passing through p
   bool cmp_proj(pt p, pt q) { return dot(v,p) < dot(v,q)</pre>
       ); } // order for points over the line
    line translate(pt t) { return {v, c + cross(v,t)}; }
    line shift_left(double d) { return {v, c + d*abs(v)};
   pt proj(pt p) { return p - rot90ccw(v) *side(p) / norm(v
       ); } // pt provected on the line
    pt refl(pt p) { return p - rot90ccw(v) *2*side(p) / norm
       (v); } // pt reflected on the other side of the
       line
} ;
bool inter_ll(line 11, line 12, pt &out) {
    T d = cross(11.v, 12.v);
    if (d == 0) return false;
    out = (12.v*11.c - 11.v*12.c) / d; // floating points
    return true;
//bisector divides the angle in 2 equal angles
//interior line goes on the same direction as 11 and 12
line bisector(line 11, line 12, bool interior) {
    assert(cross(11.v, 12.v) != 0); /// 11 and 12 cannot
       be parallel!
    lf sign = interior ? 1 : -1;
    return {12.v/abs(12.v) + 11.v/abs(11.v) * sign,
            12.c/abs(12.v) + 11.c/abs(11.v) * sign};
```

# 5.3 Poligonos

```
enum {IN, OUT, ON};
struct polygon {
```

```
vector<pt> p;
polygon(int n) : p(n) {}
int top = -1, bottom = -1;
void delete_repetead() {
    vector<pt> aux;
    sort(p.begin(), p.end());
    for(pt &i : p)
        if(aux.empty() || aux.back() != i)
          aux.push back(i);
    p.swap(aux);
bool is convex() {
    bool pos = 0, neg = 0;
    for (int i = 0, n = p.size(); i < n; i++) {
        int o = orient(p[i], p[(i+1)%n], p[(i+2)%n]);
        if (o > 0) pos = 1;
        if (o < 0) neg = 1;
    return ! (pos && neg);
lf area(bool s = false) { // better on clockwise
   order
    lf ans = 0;
    for (int i = 0, n = p.size(); i < n; i++)
        ans += cross(p[i], p[(i+1)%n]);
    ans /= 2;
    return s ? ans : abs(ans);
lf perimeter() {
    lf per = 0;
    for (int i = 0, n = p.size(); i < n; i++)
       per += abs(p[i] - p[(i+1)%n]);
    return per;
bool above(pt a, pt p) { return p.y >= a.y; }
bool crosses ray(pt a, pt p, pt q) { // pq crosses
   rav from a
    return (above (a,q) -above (a,p)) *orient (a,p,q) > 0;
int in polygon(pt a) {
    int crosses = 0;
    for(int i = 0, n = p.size(); i < n; i++) {</pre>
        if (on_segment(p[i], p[(i+1)%n], a)) return ON
        crosses += crosses ray(a, p[i], p[(i+1)%n]);
    return (crosses&1 ? IN : OUT);
void normalize() { /// polygon is CCW
    bottom = min element(p.begin(), p.end()) - p.
       begin();
    vector<pt> tmp(p.begin()+bottom, p.end());
    tmp.insert(tmp.end(), p.begin(), p.begin()+bottom
    p.swap(tmp);
```

```
bottom = 0:
    top = max_element(p.begin(), p.end()) - p.begin()
int in convex(pt a) {
    assert (bottom == 0 \&\& top != -1);
    if(a < p[0] || a > p[top]) return OUT;
    T orientation = orient(p[0], p[top], a);
    if(orientation == 0) {
        if(a == p[0] || a == p[top]) return ON;
        return top == 1 || top + 1 == p.size() ? ON :
    } else if (orientation < 0) {</pre>
        auto it = lower_bound(p.begin()+1, p.begin()+
        T d = orient(*prev(it), a, *it);
        return d < 0 ? IN : (d > 0 ? OUT: ON);
        auto it = upper_bound(p.rbegin(), p.rend()-
           top-1, a);
        T d = orient(*it, a, it == p.rbegin() ? p[0]
           : *prev(it));
        return d < 0 ? IN : (d > 0 ? OUT: ON);
polygon cut(pt a, pt b) { // cuts polygon on line ab
    line l(a, b);
    polygon new_polygon(0);
    for(int i = 0, n = p.size(); i < n; ++i) {
        pt c = p[i], d = p[(i+1)%n];
        lf abc = cross(b-a, c-a), abd = cross(b-a, d-a)
        if(abc >= 0) new polygon.p.push back(c);
        if(abc*abd < 0) {
          pt out; inter ll(l, line(c, d), out);
          new polygon.p.push back(out);
    return new polygon;
void convex hull() {
    sort(p.begin(), p.end());
    vector<pt> ch;
    ch.reserve(p.size()+1);
    for(int it = 0; it < 2; it++) {
        int start = ch.size();
        for(auto &a : p) {
            /// if colineal are needed, use < and
               remove repeated points
            while(ch.size() >= start+2 && orient(ch[
               ch.size()-2], ch.back(), a) <= 0)
                ch.pop back();
            ch.push back(a);
        ch.pop_back();
```

```
reverse(p.begin(), p.end());
    if(ch.size() == 2 \&\& ch[0] == ch[1]) ch.pop back
    /// be careful with CH of size < 3
    p.swap(ch);
vector<pii> antipodal() {
    vector<pii> ans;
    int n = p.size();
    if(n == 2) ans.push back({0, 1});
    if(n < 3) return ans;</pre>
    auto nxt = [\&] (int x) \{ return (x+1 == n ? 0 : x \} \}
       +1); };
    auto area2 = [&] (pt a, pt b, pt c) { return cross
        (b-a, c-a); };
    int b0 = 0;
    while(abs(area2(p[n - 1], p[0], p[nxt(b0)])) >
       abs(area2(p[n - 1], p[0], p[b0]))) ++b0;
    for (int b = b0, a = 0; b != 0 && a <= b0; ++a) {
        ans.push_back({a, b});
        while (abs(area2(p[a], p[nxt(a)], p[nxt(b)]))
            > abs(area2(p[a], p[nxt(a)], p[b]))) {
            b = nxt(b);
            if(a != b0 || b != 0) ans.push back({ a,
            else return ans;
        if(abs(area2(p[a], p[nxt(a)], p[nxt(b)])) ==
            abs(area2(p[a], p[nxt(a)], p[b]))) {
            if (a != b0 \mid | b \mid = n-1) ans.push_back({ a
                , nxt(b) });
            else ans.push_back({ nxt(a), b });
    return ans;
pt centroid() {
    pt c{0, 0};
    lf scale = 6. * area(true);
    for(int i = 0, n = p.size(); i < n; ++i) {
        int j = (i+1 == n ? 0 : i+1);
        c = c + (p[i] + p[j]) * cross(p[i], p[j]);
    return c / scale;
ll pick() {
    11 boundary = 0;
    for(int i = 0, n = p.size(); i < n; i++) {</pre>
        int j = (i+1 == n ? 0 : i+1);
        boundary += \underline{gcd((ll)abs(p[i].x - p[j].x)}, (
           ll) abs (p[i].y - p[j].y);
    return area() + 1 - boundary/2;
```

```
pt& operator[] (int i) { return p[i]; }
};
```

### 5.4 Circulos

```
struct circle {
   pt c; T r;
// (x-xo)^2 + (y-yo)^2 = r^2
//circle that passes through abc
circle center(pt a, pt b, pt c) {
    b = b-a, c = c-a;
    assert(cross(b,c) != 0); /// no circumcircle if A,B,C
    pt cen = a + rot90ccw(b*norm(c) - c*norm(b))/cross(b,
       c) /2:
    return {cen, abs(a-cen)};
//centers of the circles that pass through ab and has
   radius r
vector<pt> centers(pt a, pt b, T r) {
    if (abs(a-b) > 2*r + eps) return {};
    pt m = (a+b)/2;
    double f = sqrt(r*r/norm(a-m) - 1);
    pt c = rot90ccw(a-m)*f;
    return {m-c, m+c};
int inter_cl(circle c, line l, pair<pt, pt> &out) {
    lf h2 = c.r*c.r - l.sq_dist(c.c);
    if(h2 >= 0) { // line touches circle
        pt p = 1.proj(c.c);
        pt h = 1.v * sqrt(h2) / abs(1.v); // vector of len h
           parallel to line
        out = \{p-h, p+h\};
    return 1 + sign(h2); // if 1 -> out.F == out.S
int inter cc(circle c1, circle c2, pair<pt, pt> &out) {
    pt d = c2.c-c1.c;
    double d2 = norm(d);
    if(d2 == 0) { assert(c1.r != c2.r); return 0; } //
       concentric circles (identical)
    double pd = (d2 + c1.r*c1.r - c2.r*c2.r)/2; // = /
       0.1PI * d
    double h2 = c1.r*c1.r - pd*pd/d2; // = h^2
    if(h2 >= 0) {
        pt p = c1.c + d*pd/d2, h = rot90ccw(d)*sqrt(h2/d2
        out = \{p-h, p+h\};
    return 1 + sign(h2);
//circle-line inter = 1
```

```
int tangents(circle c1, circle c2, bool inner, vector<</pre>
   pair<pt,pt>> &out) {
    if(inner) c2.r = -c2.r; // inner tangent
    pt d = c2.c-c1.c;
    double dr = c1.r-c2.r, d2 = norm(d), h2 = d2-dr*dr;
    if(d2 == 0 \mid | h2 < 0)  { assert(h2 != 0); return 0; }
        // (identical)
    for(double s : {-1,1}) {
        pt v = (d*dr + rot 90ccw(d)*sqrt(h2)*s)/d2;
        out.push_back({c1.c + v*c1.r, c2.c + v*c2.r});
    return 1 + (h2 > 0); // if 1: circle are tangent
//circle targent passing through pt p
int tangent through pt(pt p, circle c, pair<pt, pt> &out)
    double d = abs(p - c.c);
    if(d < c.r) return 0;</pre>
    pt base = c.c-p;
    double w = sqrt(norm(base) - c.r*c.r);
    pt a = \{w, c.r\}, b = \{w, -c.r\};
    pt s = p + base*a/norm(base)*w;
    pt t = p + base*b/norm(base) *w;
    out = \{s, t\};
    return 1 + (abs(c.c-p) == c.r);
```

### 5.5 Semiplanos

```
struct halfplane{
    double angle;
    pt p, pq;
    halfplane(){}
    halfplane(pt a, pt b): p(a), pq(b - a) {
        angle = atan2(pq.y,pq.x);
    bool operator < (halfplane b) const{return angle < b.</pre>
       angle; }
    bool out (pt q) {return cross(pq, (q-p)) < -eps; } //
       checks if p is inside the half plane
};
const lf inf = 1e100;
// intersection pt of the lines of 2 halfplanes
pt inter(halfplane& h1, halfplane& h2) {
    if(abs(cross(unit(h1.pq), unit(h2.pq))) <= eps)return</pre>
         {inf, inf};
    lf alpha = cross((h2.p - h1.p), h2.pq) / cross(h1.pq,
        h2.pa);
    return h1.p + (h1.pq * alpha);
// intersection of halfplanes
vector<pt> intersect(vector<halfplane>& b) {
```

```
vector < pt > box = { \{inf, inf\}, \{-inf, inf\}, \{-inf, -inf\}, \{-inf, -i
               inf}, {inf, -inf} };
for(int i = 0; i < 4; i++) {
                 b.push_back(\{box[i], box[(i + 1) % 4]\});
sort(b.begin(), b.end());
int n = b.size(), q = 1, h = 0;
vector<halfplane> c(n + 10);
for(int i = 0; i < n; i++) {</pre>
                 while (q < h \&\& b[i].out(inter(c[h], c[h-1]))) h
                 while (q < h \& \& b[i].out(inter(c[q], c[q+1]))) q
                 c[++h] = b[i];
                 if(q < h \&\& abs(cross(c[h].pq, c[h-1].pq)) < eps)
                                  if(dot(c[h].pq, c[h-1].pq) <= 0) return {};
                                  if(b[i].out(c[h].p)) c[h] = b[i];
while (q < h-1 \&\& c[q].out(inter(c[h], c[h-1]))) h--;
while (q < h-1 \& \& c[h].out(inter(c[q], c[q+1]))) q++;
if(h - q <= 1) return {};
c[h+1] = c[q];
vector<pt> s;
for (int i = q; i < h+1; i++) s.pb(inter(c[i], c[i+1])
return s;
```

### 5.6 Segmentos

```
bool in_disk(pt a, pt b, pt p) { // pt p inside ab disk
    return dot(a-p, b-p) <= 0;
bool on_segment(pt a, pt b, pt p) { // p on ab
    return orient (a,b,p) == 0 \&\& in disk(a,b,p);
// ab crossing cd
bool proper_inter(pt a, pt b, pt c, pt d, pt &out) {
   T oa = orient(c,d,a),
    ob = orient(c,d,b),
    oc = orient(a,b,c),
    od = orient(a,b,d);
    /// Proper intersection exists iff opposite signs
    if (oa*ob < 0 && oc*od < 0) {
        out = (a*ob - b*oa) / (ob-oa);
        return true;
    return false;
// intersection bwn segments
set<pt> inter_ss(pt a, pt b, pt c, pt d) {
```

```
25
```

```
6 Grafos
```

### 6.1 Puentes

pt dummy;

pts

pt out;

set<pt> s;

cross -> 1

**return** s; // 0, 2

line l(a,b);

intersects cd

distance to A or B

lf seq\_to\_seq(pt a, pt b, pt c, pt d) {

**if**(a != b) {

```
vector<bool> visited;
vi tin, low;
int timer;
void IS_BRIDGE(int u, int v, vii &puentes) {
    puentes.push back({min(u, v), max(u, v)});
void dfs(vector<vi> &adj, vii &puentes, int v, int p =
   -1) {
    visited[v] = true;
    tin[v] = low[v] = timer++;
    for (int to : adj[v]) {
        if (to == p) continue;
        if (visited[to]) {
            low[v] = min(low[v], tin[to]);
        } else
            dfs(adj, puentes, to, v);
            low[v] = min(low[v], low[to]);
            if (low[to] > tin[v])
```

if (proper\_inter(a,b,c,d,out)) return {out}; //if

if (l.cmp\_proj(a,p) && l.cmp\_proj(p,b)) /// if
 closest to projection = (a, p, b)

pt\_to\_seq(c,d,a), pt\_to\_seq(c,d,b)}); // try the 4

return l.dist(p); /// output distance to line

if (on\_segment(c,d,a)) s.insert(a); // a in cd
if (on\_segment(c,d,b)) s.insert(b); // b in cd

if (on\_segment(a,b,c)) s.insert(c); // c in ab
if (on segment(a,b,d)) s.insert(d); // d in ab

return min(abs(p-a), abs(p-b)); /// otherwise

if (proper\_inter(a,b,c,d,dummy)) return 0; // ab

return min({pt\_to\_seg(a,b,c), pt\_to\_seg(a,b,d),

If pt to seg(pt a, pt b, pt p) { // p to ab

### 6.2 Puntos de Articulación

```
int n;
vector<vector<int>> adj;
vector<bool> visited;
vector<int> tin, low;
int timer;
void dfs(int v, int p = -1) {
    visited[v] = true;
    tin[v] = low[v] = timer++;
    int children=0;
    for (int to : adj[v]) {
        if (to == p) continue;
        if (visited[to]) {
            low[v] = min(low[v], tin[to]);
        } else {
            dfs(to, v);
            low[v] = min(low[v], low[to]);
            if (low[to] >= tin[v] && p!=-1)
                IS_CUTPOINT(v);
            ++children;
    if(p == -1 && children > 1)
        IS_CUTPOINT(v);
void find_cutpoints() {
    timer = 0;
    visited.assign(n, false);
    tin.assign(n, -1);
    low.assign(n, -1);
    for (int i = 0; i < n; ++i) {
        if (!visited[i])
            dfs (i);
```

GRAFOS

# 6.3 Kosajaru

```
//Encontrar las componentes fuertemente conexas en un
   grafo dirigido
//Componente fuertemente conexa: es un grupo de nodos en
   el que hav
//un camino dirigido desde cualquier nodo hasta cualquier
    otro nodo dentro del grupo.
void Kosaraju(int u, int pass) {
    dfs num[u] = 1;
    vii &neighbor = (pass == 1) ? AL[u] : AL_T[u];
    for (auto &[v, w] : neighbor)
        if (dfs num[v] == UNVISITED)
            Kosaraju(v, pass);
    S.push_back(u);
int main(){
    S.clear():
    dfs_num.assign(N, UNVISITED);
    for (int u = 0; u < N; ++u)
        if (dfs num[u] == UNVISITED)
            Kosaraju(u, 1);
    numSCC = 0;
    dfs num.assign(N, UNVISITED);
    for (int i = N-1; i >= 0; --i)
        if (dfs_num[S[i]] == UNVISITED)
            ++numSCC, Kosaraju(S[i], 2);
    printf("There are %d SCCs\n", numSCC);
```

# 6.4 Tarjan

```
vi low, num, comp, g[nax];
int scc, timer;
stack<int> st;
void tjn(int u) {
  low[u] = num[u] = timer++; st.push(u); int v;
  for(int v: g[u]) {
    if(num[v]==-1) tjn(v);
    if(comp[v]==-1) low[u] = min(low[u], low[v]);
  }
  if(low[u]==num[u]) {
    do{ v = st.top(); st.pop(); comp[v]=scc;
    }while(u != v);
    +*scc;
  }
}
void callt(int n) {
  timer = scc= 0;
```

```
num = low = comp = vector<int>(n,-1);
for(int i = 0; i<n; i++) if(num[i]==-1) tjn(i);
}</pre>
```

### 6.5 Dijkstra

```
//Camino mas cortos
//NO USAR CON PESOS NEGATIVOS, usar Bellman Ford o SPFA(
   mas rapido)
// O ((V+\bar{E})*log V)
vi dijkstra(vector<vii> &adj, int s, int V) {
    vi dist(V+1, INT_MAX); dist[s] = 0;
    priority queue<ii, vii, greater<ii>> pq; pq.push(ii
        (0, s);
    while(!pa.emptv()){
        ii front = pq.top(); pq.pop();
        int d = front.first, u = front.second;
        if (d > dist[u]) continue;
        for (int j = 0; j < (int)adj[u].size(); j++){</pre>
            ii v = adj[u][j];
            if (dist[u] + v.second < dist[v.first]) {</pre>
                dist[v.first] = dist[u] + v.second;
                pq.push(ii(dist[v.first], v.first));
    return dist;
```

### 6.6 Bellman Ford

```
vi bellman_ford(vector<vii> &adj, int s, int n) {
    vi dist(n, INF); dist[s] = 0;
    for (int i = 0; i<n-1; i++) {</pre>
        bool modified = false;
        for (int u = 0; u < n; u + +)
            if (dist[u] != INF)
                 for (auto &[v, w] : adj[u]) {
                     if (dist[v] <= dist[u] + w) continue;</pre>
                     dist[v] = dist[u] + w;
                     modified = true;
        if (!modified) break;
    bool negativeCicle = false;
    for (int u = 0; u<n; u++)
        if (dist[u] != INF)
            for (auto &[v, w] : adj[u]){
                if (dist[v] > dist[u] + w) negativeCicle
                    = true;
```

```
return dist;
}
```

### 6.7 Floyd Warshall

### 6.8 MST Kruskal

```
//Arbol de minima expansion
//O(E*log V)
int main() {
    int n, m;
    cin >> n >> m;
    vector<pair<int, ii>> adj; //Los pares son: {peso, {
       vertice, vecino}}
    for (int i = 0; i<m; i++) {
        int x, y, w; cin >> x >> y >> w;
        adj.push back (make pair (w, ii(x, y)));
    sort(adj.begin(), adj.end());
    int mst costo = 0, tomados = 0;
    dsu UF(n);
    for (int i = 0; i < m && tomados < n-1; i++) {</pre>
        pair<int, ii> front = adj[i];
        if (!UF.is_same_set(front.second.first, front.
           second.second)){
            tomados++;
            mst_costo += front.first;
```

### 6.9 MST Prim

```
vector<vii> adi;
vi tomado;
priority_queue<ii>> pq;
void process(int u) {
    tomado[u] = 1;
    for (auto &[v, w] : adj[u]){
        if (!tomado[v]) pq.emplace(-w, -v);
int prim(int v, int n){
    tomado.assign(n, 0);
    process(0);
    int mst_costo = 0, tomados = 0;
    while (!pq.empty()){
        auto [w, u] = pq.top(); pq.pop();
w = -w; u = -u;
        if (tomado[u]) continue;
        mst costo += w;
        process(u);
        tomados++;
        if (tomados == n-1) break;
    return mst_costo;
```

### 6.10 Shortest Path Faster Algorithm

```
//Algoritmo mas rapido de ruta minima
//O(V*E) peor caso, O(E) en promedio.
bool spfa(vector<vii> &adj, vector<int> &d, int s, int n)
{
    d.assign(n, INF);
    vector<int> cnt(n, 0);
    vector<bool> inqueue(n, false);
    queue<int> q;

    d[s] = 0;
    q.push(s);
    inqueue[s] = true;
    while (!q.empty()) {
        int v = q.front();
        q.pop();
        inqueue[v] = false;
```

### 6.11 Camino mas corto de longitud fija

```
Modificar operacion * de matrix de esta forma:
En la exponenciacion binaria inicializar matrix ans = b
matrix operator * (const matrix &b) {
    matrix ans (this->r, b.c, vector<vl>(this->r, vl(b.c,
       INFL)));
    for (int i = 0; i<this->r; i++) {
        for (int k = 0; k<b.r; k++) {</pre>
            for (int j = 0; j<b.c; j++) {
                ans.m[i][j] = min(ans.m[i][j], m[i][k] +
                    b.m[k][j]);
    return ans;
int main() {
    int n, m, k; cin >> n >> m >> k;
    vector<vl> adj(n, vl(n, INFL));
    for (int i = 0; i<m; i++) {</pre>
        ll a, b, c; cin >> a >> b >> c; a--; b--;
        adj[a][b] = min(adj[a][b], c);
    matrix graph(n, n, adj);
    graph = pow(graph, k-1);
    cout << (graph.m[0][n-1] == INFL ? -1 : graph.m[0][n
```

```
-11) << "\n";
      return 0;
6.12 2sat
  // O(n+m)
  // l=(x1 or y1) and (x2 or y2) and ... and (xn or yn)
  struct sat2 {
          int n;
          vector<vector<vi>>> q;
          vector<bool> vis, val;
          vi comp;
          stack<int> st;
           sat2(int n):n(n),q(2, vector<vi>(2*n)),vis(2*n),
              val(2*n), comp(2*n) {}
          int neq(int x) {return 2*n-x-1;}
          void make true(int u) {add_edge(neg(u), u);}
          void make false(int u) {make true(neg(u));}
          void add_or(int u, int v) {implication(neg(u), v);}
          void diff(int u, int v) {eq(u, neq(v));}
          void eq(int u, int v){
                   implication(u, v);
                   implication(v, u);
          void implication(int u,int v) {
                   add edge(u, v);
                   add_edge(neg(v),neg(u));
          void add edge(int u, int v) {
                   q[0][u].PB(v);
                   q[1][v].PB(u);
          void dfs(int id, int u, int t=0){
                   vis[u]=true;
                   for(auto &v:g[id][u])
                           if(!vis[v])dfs(id, v, t);
                   if (id) comp[u]=t;
                   else st.push(u);
          void kosaraju() {
                   for(int u=0; u<n; ++u) {
                           if(!vis[u])dfs(0, u);
                           if(!vis[neq(u)])dfs(0, neq(u));
                   vis.assign(2*n, false);
                   int t=0;
                   while(!st.emptv()){
                           int u=st.top();st.pop();
```

**if**(!vis[u])dfs(1, u, t++);

```
7 MATEMATICAS
```

```
bool check() {
                  kosaraju();
                  for(int i=0;i<n;++i) {</pre>
                           if (comp[i] == comp[neq(i)]) return
                               false;
                           val[i]=comp[i]>comp[neg(i)];
                  return true;
};
int m,n;
sat2 s(n);
char c1, c2;
for(int a,b,i=0;i<m;++i) {</pre>
        cin>>c1>>a>>c2>>b;
         a--;b--;
         if (c1=='-') a=s.neq(a);
         if (c2=='-')b=s.neg(b);
         s.add or(a,b);
if(s.check()){
         for(int i=0;i<n;++i)cout<<(s.val[i]?'+':'-')<<" "</pre>
         cout << "\n";
}else cout<<"IMPOSSIBLE\n";</pre>
```

### 7 Matematicas

### 7.1 Coeficientes binomiales

```
const int MAX N = 100010;
                             // MOD > MAX N
// O (log MOD)
ll inv (ll a) {
    return binpow(a, MOD-2, MOD);
11 fact[MAX N];
// O(log MOD)
11 C(int n, int k) {
    if (n < k) return 0;
    return (((fact[n] * inv(fact[k])) % MOD) * inv(fact[n
       -k])) % MOD;
int main() {
    fact[0] = 1;
    for (int i = 1; i<MAX N; i++) {</pre>
        fact[i] = (fact[i-1]*i) % MOD;
    cout << C(100000, 50000) << "\n";
```

```
return 0;
```

#### 7.2 Criba Modificada

```
//Criba modificada
Si hay que determinar el numero de factores primos para
   muchos (o un rango) de enteros.
La mejor solucion es el algoritmo de criba modificada O(N
    log log N)
int numDiffPFarr[MAX N+10] = \{0\}; // e.g., MAX N = 10^7
for (int i = 2; i <= MAX N; ++i)</pre>
    if (numDiffPFarr[i] == 0) // i is a prime number
        for (int j = i; j <= MAX N; j += i)
            ++numDiffPFarr[j]; // j is a multiple of i
//Similar para EulerPhi
int EulerPhi[MAX_N+10];
for (int i = 1; i <= MAX_N; ++i) EulerPhi[i] = i;</pre>
for (int i = 2; i <= MAX N; ++i)</pre>
    if (EulerPhi[i] == i) // i is a prime number
        for (int j = i; j <= MAX N; j += i)
            EulerPhi[i] = (EulerPhi[i]/i) * (i-1);
```

#### 7.3 Ecuaciones Diofanticas

```
// O(\log(\min(a, b)))
ll extEuclid(ll a, ll b, ll &x, ll &y) {
    11 xx = y = 0;
    11 yy = \bar{x} = 1;
    while (b) {
        11 q = a/b;
        ll \hat{t} = b; b = a b; a = t;
        t = xx; xx = x-q*xx; x = t;
        t = yy; yy = y - q * yy; y = t;
    return a; //Devuelve gcd(a, b)
bool find any solution (ll a, ll b, ll c, ll &x0, ll &y0,
   ll &a) {
    q = extEuclid(abs(a), abs(b), x0, y0);
    if (c % a) {
        return false;
    x0 \star = c / q;
    y0 \star = c / q;
    if (a < 0) x0 = -x0;
    if (b < 0) v0 = -v0;
    return true;
```

#### 7.4 Funcion Totient de Euler

### 7.5 Exponenciacion binaria

```
ll binpow(ll b, ll n, ll m) {
    b %= m;
    ll res = 1;
    while (n > 0) {
        if (n & 1)
            res = res * b % m;
        b = b * b % m;
        n >>= 1;
    }
    return res % m;
}
```

### 7.6 Exponenciacion matricial

```
}
}
return ans;
}

matrix pow(matrix &b, ll p) {
    matrix ans(b.r, b.c, vector<vl>(b.r, vl(b.c, 0)));
    for (int i = 0; i<b.r; i++) ans.m[i][i] = 1;
    while (p) {
        if (p&1) {
            ans = ans*b;
        }
        b = b*b;
        p >>= 1;
    }
    return ans;
}
```

### 7.7 Fibonacci Fast Doubling

```
pair < int, int > fib (int n) {
    if (n == 0)
        return {0, 1};

auto p = fib(n >> 1);
    int c = p.first * (2 * p.second - p.first);
    int d = p.first * p.first + p.second * p.second;
    if (n & 1)
        return {d, c + d};
    else
        return {c, d};
}
```

### 7.8 Freivalds algorithm

```
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
   ().count());
// check if two n*n matrix a*b=c within complexity (
    iteration*n^2)
// probability of error 2^(-iteration)
int Freivalds(matrix &a, matrix &b, matrix &c) {
   int n = a.r, iteration = 20;
   matrix zero(n, 1), r(n, 1);
   while (iteration--) {
        for(int i = 0; i < n; i++) r.m[i][0] = rnd() % 2;
        matrix ans = (a * (b * r)) - (c * r);
        if(ans.m != zero.m) return 0;
   }
   return 1;
}</pre>
```

#### 7.9 Gauss Jordan

```
const double EPS = 1e-9;
const int INF = 2; // it doesn't actually have to be
   infinity or a big number
int gauss (vector < vector < double > > a, vector < double > &
   ans) {
    int n = (int) a.size();
    int m = (int) a[0].size() - 1;
    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {</pre>
        int sel = row;
        for (int i=row; i<n; ++i)
            if (abs (a[i][col]) > abs (a[sel][col]))
                 sel = i;
        if (abs (a[sel][col]) < EPS)</pre>
             continue:
        for (int i=col; i<=m; ++i)
             swap (a[sel][i], a[row][i]);
        where [col] = row;
        for (int i=0; i<n; ++i)</pre>
            if (i != row) {
                 double c = a[i][col] / a[row][col];
                 for (int j=col; j<=m; ++j)</pre>
                     a[i][j] -= a[row][j] * c;
        ++row;
    ans.assign (m, 0);
    for (int i=0; i<m; ++i)</pre>
        if (where [i] !=-1)
            ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i=0; i<n; ++i) {
        double sum = 0;
        for (int j=0; j<m; ++j)
            sum += ans[j] * a[i][j];
        if (abs (sum - a[i][m]) > EPS)
            return 0;
    for (int i=0; i<m; ++i)
        if (where [i] == -1)
            return INF;
    return 1;
```

# 7.10 Gauss Jordan mod 2

```
// O(min(n, m)*n*m)
```

```
int gauss (vector < bitset<N> > &a, int n, int m, bitset<</pre>
   N > \& ans) {
    vector\langle int \rangle where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {</pre>
         for (int i=row; i<n; ++i)</pre>
             if (a[i][col]) {
                 swap (a[i], a[row]);
                 break;
        if (! a[row][col])
             continue;
        where [col] = row;
        for (int i=0; i<n; ++i)
             if (i != row && a[i][col])
                 a[i] ^= a[row];
         ++row;
    for (int i=0; i<m; ++i)
        if (where [i] != -1)
             ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i=0; i<n; ++i) {
         double sum = 0;
        for (int j=0; j<m; ++j)
             sum += ans[j] * a[i][j];
        if (abs (sum - a[i][m]) > EPS)
             return 0;
    for (int i=0; i<m; ++i)</pre>
         if (where [i] == -1)
             return INF;
    return 1;
```

# 7.11 GCD y LCM

```
//o(log10 n) n == max(a, b)
int gcd(int a, int b) { return b == 0 ? a : gcd(b, a%b);
}
int lcm(int a, int b) { return a / gcd(a, b) * b; }
//gcd(a, b, c) = gcd(a, gcd(b, c))
```

### 7.12 Integral Definida

```
double f(double x) {
   return x*x;
}

const int N = 1000 * 1000; // number of steps (already
   multiplied by 2)
double simpson_integration(double a, double b) {
```

```
double h=(b-a)/N;
double s=f(a)+f(b);
for (int i=1;i<=N-1;i++) {
    double x=a+h*i;
    s+=f(x)*((i & 1)?4:2);
}
s*=h/3;
return s;
}</pre>
```

### 7.13 Inverso modular

```
11 mod(ll a, ll m) {
    return ((a%m) + m) % m;
}

11 modInverse(ll b, ll m) {
    ll x, y;
    ll d = extEuclid(b, m, x, y); //obtiene b*x + m*y ==
        d
    if (d!=1) return -1; //indica error
        // b*x + m*y == 1, ahora aplicamos (mod m) para
        obtener b*x == 1 (mod m)
    return mod(x, m);
}

// Otra forma
// O(log MOD)
ll inv (ll a) {
    return binpow(a, MOD-2, MOD);
}
```

### 7.14 Logaritmo Discreto

```
// Returns minimum x for which a \hat{x} % m = b % m.
int solve(int a, int b, int m) {
    a %= m, b %= m;
    int k = 1, add = 0, q;
    while ((q = qcd(a, m)) > 1) {
        if (b == k)
            return add;
        if (b % q)
            return -1;
        b /= g, m /= g, ++add;
        k = (k * 111 * a / q) % m;
    int n = sqrt(m) + 1;
    int an = 1;
    for (int i = 0; i < n; ++i)
        an = (an * 111 * a) % m;
    unordered_map<int, int> vals;
```

```
for (int q = 0, cur = b; q <= n; ++q) {
    vals[cur] = q;
    cur = (cur * 111 * a) % m;
}

for (int p = 1, cur = k; p <= n; ++p) {
    cur = (cur * 111 * an) % m;
    if (vals.count(cur)) {
        int ans = n * p - vals[cur] + add;
        return ans;
    }
}
return -1;
}</pre>
```

#### 7.15 Miller Rabin

```
ll mul (ll a, ll b, ll mod) {
  11 \text{ ret} = 0;
  for (a %= mod, b %= mod; b != 0;
    b >>= 1, a <<= 1, a = a >= mod ? <math>a - mod : a) {
    if (b & 1) {
      ret += a;
      if (ret >= mod) ret -= mod;
  return ret:
ll fpow (ll a, ll b, ll mod) {
 ll ans = 1;
  for (; b; b >>= 1, a = mul(a, a, mod))
    if (b & 1)
      ans = mul(ans, a, mod);
  return ans:
bool witness (ll a, ll s, ll d, ll n) {
  ll x = fpow(a, d, n);
  if (x == 1 \mid | x == n - 1) return false;
  for (int i = 0; i < s - 1; i++) {
    x = mul(x, x, n);
    if (x == 1) return true;
    if (x == n - 1) return false;
  return true;
11 \text{ test}[] = \{2, 3, 5, 7, 11, 13, 17, 19, 23, 0\};
bool is prime (ll n) {
 if (n < 2) return false;</pre>
  if (n == 2) return true;
  if (n % 2 == 0) return false;
  11 \dot{d} = n - 1, s = 0;
  while (d \% 2 == 0) ++s, d /= 2;
  for (int i = 0; test[i] && test[i] < n; ++i)</pre>
    if (witness(test[i], s, d, n))
```

```
return false;
return true;
}
```

#### 7.16 Miller Rabin Probabilistico

```
using u64 = uint64 t;
using u128 = __uint128_t;
u64 binpower(u64 base, u64 e, u64 mod) {
    u64 \text{ result} = 1;
    base %= mod;
    while (e) {
        if (e & 1)
            result = (u128) result * base % mod;
        base = (u128)base * base % mod;
        e >>= 1;
    return result;
bool check_composite(u64 n, u64 a, u64 d, int s) {
    u64 x = binpower(a, d, n);
    if (x == 1 | x == n - 1)
        return false:
    for (int r = 1; r < s; r++) {
        x = (u128) x * x % n;
        if (x == n - 1)
            return false;
    return true;
};
bool MillerRabin(u64 n, int iter=5) { // returns true if
   n is probably prime, else returns false.
    if (n < 4)
        return n == 2 || n == 3;
    int s = 0;
    u64 d = n - 1;
    while ((d & 1) == 0) {
        d >>= 1;
        s++;
    for (int i = 0; i < iter; i++) {
        int a = 2 + rand() % (n - 3);
        if (check_composite(n, a, d, s))
            return false;
    return true;
```

### 7.17 Mobius

```
const int N = 1e6+1;
int mob[N];
void mobius() {
   mob[1] = 1;
   for (int i = 2; i < N; i++) {
       mob[i]--;
      for (int j = i + i; j < N; j += i) {
       mob[j] -= mob[i];
    }
}</pre>
```

#### 7.18 Pollard Rho

```
//O(n^{(1/4)}) (?)
ll pollard rho(ll n, ll c) {
 11 x = 2, y = 2, i = 1, k = 2, d;
 while (true) {
    x = (mul(x, x, n) + c);
    if (x >= n) x -= n;
   d = \underline{gcd(x - y, n)};
    if (d > 1) return d:
    if (++i == k) v = x, k <<= 1;
 return n;
void factorize(ll n, vector<ll> &f) {
 if (n == 1) return;
  if (is prime(n)) {
    f.push_back(n);
    return;
 11 d = n;
  for (int i = 2; d == n; i++)
    d = pollard rho(n, i);
  factorize(d, f);
  factorize (n/d, f);
```

# 7.19 Simplex

```
vi X, Y;
void pivot(int x, int y) {
         swap(X[y], Y[x]);
         b[x]/=a[x][y];
         for (int i=0; i < m; ++i) {</pre>
                  if(i!=v)a[x][i]/=a[x][v];
         a[x][y]=1/a[x][y];
         for(int i=0;i<n;++i) {</pre>
                  if(i!=x && abs(a[i][y])>EPS){
                           b[i] -= a[i][y] *b[x];
                           for(int j=0; j<m; j++) {
                                    if(j != y)a[i][j
                                        ]-=a[i][v]*a[x]
                                        ][†];
                           a[i][y]=-a[i][y]*a[x][y];
         z + = c[y] * b[x];
         for(int i=0;i<m;++i){
                  if(i != y)c[i]-=c[y]*a[x][i];
         c[v] = -c[v] *a[x][v];
Simplex(vector<vector<double>> &A.vector<double>
   &B, vector<double> &C) {
         a=A; b=B; c=C;
         n=b.size(); m=c.size(); z= 0.0;
         X.resize(m); iota(X.begin(), X.end(), 0);
         Y.resize(n); iota(Y.begin(), Y.end(), m);
// \{z, \{x1, x2, x3...\}\}
pair<double, vector<double>> maximize() {
         while(true) {
                  int x=-1, y=-1;
                  double mn=-EPS;
                  for (int i=0; i < n; ++i) {</pre>
                           if(b[i]<mn)mn=b[i],x=i;
                  if (x<0) break;</pre>
                  for (int i=0; i<m; ++i) {</pre>
                           if(a[x][i]<-EPS){
                                    y=i;
                                    break;
                  assert(y>=0); // no hay solucion
                      para Ax<=b
                  pivot(x,y);
         while(true) {
                  double mx=EPS;
                  int x=-1, y=-1;
                  for (int i=0; i < m; ++i) {</pre>
```

### 7.20 Fast Fourier Transform

```
const double PI = acos(-1);
struct base {
  double a, b;
  base (double a = 0, double b = 0) : a(a), b(b) {}
  const base operator + (const base &c) const
    { return base(a + c.a, b + c.b); }
  const base operator - (const base &c) const
    { return base(a - c.a, b - c.b); }
  const base operator * (const base &c) const
    { return base(a * c.a - b * c.b, a * c.b + b * c.a);
void fft(vector<base> &p, bool inv = 0) {
  int n = p.size(), i = 0;
  for (int j = 1; j < n - 1; ++j) {
    for (int k = n >> 1; k > (i = k); k >>= 1);
    if(j < i) swap(p[i], p[j]);
  for(int 1 = 1, m; (m = 1 << 1) <= n; 1 <<= 1) {
    double ang = 2 * PI / m;
    base wn = base(cos(ang), (inv ? 1. : -1.) * sin(ang))
    for (int i = 0, j, k; i < n; i += m) {
      for(w = base(1, 0), j = i, k = i + 1; j < k; ++j, w
          = w * wn) {
        base t = w * p[j + l];
        p[j + 1] = p[j] - t;
       p[j] = p[j] + t;
  if(inv) for(int i = 0; i < n; ++i) p[i].a /= n, p[i].b
```

```
/= n;
vector<long long> multiply(vector<int> &a, vector<int> &b
  int n = a.size(), m = b.size(), t = n + m - 1, sz = 1;
  while(sz < t) sz <<= 1;
  vector<br/><br/>base> x(sz), v(sz), z(sz);
  for(int i = 0 ; i < sz; ++i) {</pre>
    x[i] = i < (int)a.size() ? base(a[i], 0) : base(0, 0)
    y[i] = i < (int)b.size() ? base(b[i], 0) : base(0, 0)
  fft(x), fft(y);
  for (int i = 0; i < sz; ++i) z[i] = x[i] * y[i];
  fft(z, 1);
  vector<long long> ret(sz);
  for (int i = 0; i < sz; ++i) ret[i] = (long long) round(
     z[i].a);
// while((int)ret.size() > 1 && ret.back() == 0) ret.
   pop back();
  return ret;
```

#### 7.21 Number Theoretic Transform

```
const int N = 1 \ll 20;
const int mod = 469762049; \frac{1}{998244353}
const int root = 3;
int lim, rev[N], w[N], wn[N], inv_lim;
void reduce(int &x) { x = (x + mod) % mod; }
int POW(int x, int y, int ans = 1) {
  for (; y; y >>= 1, x = (long long) x * x % mod) if (y &
      1) ans = (long long) ans * x % mod;
  return ans;
void precompute(int len) {
  \lim = wn[0] = 1; int s = -1;
  while (lim < len) lim <<= 1, ++s;
  for (int i = 0; i < lim; ++i) rev[i] = rev[i >> 1] >> 1
      | (i & 1) << s;
  const int g = POW(root, (mod - 1) / lim);
  inv_lim = POW(lim, mod - 2);
  for (int i = 1; i < lim; ++i) wn[i] = (long long) wn[i</pre>
     -11 * q % mod;
void ntt(vector<int> &a, int typ) {
  for (int i = 0; i < lim; ++i) if (i < rev[i]) swap(a[i</pre>
     l, a[rev[i]]);
  for (int i = 1; i < lim; i <<= 1) {</pre>
    for (int j = 0, t = \lim / i / 2; j < i; ++j) w[j] =
       wn[j * t];
    for (int j = 0; j < lim; j += i << 1) {
      for (int k = 0; k < i; ++k) {
```

```
const int x = a[k + j], y = (long long) a[k + j +
            i \mid * w[k] % mod;
        reduce(a[k + j] += y - mod), reduce(a[k + j + i]
           = x - y);
  if (!typ) {
    reverse(a.begin() + 1, a.begin() + lim);
    for (int i = 0; i < lim; ++i) a[i] = (long long) a[i]
        * inv lim % mod;
vector<int> multiply(vector<int> &f, vector<int> &q) {
  int n=(int)f.size() + (int)q.size() - 1;
  precompute(n);
  vector<int> a = f, b = q;
  a.resize(lim); b.resize(lim);
  ntt(a, 1), ntt(b, 1);
  for (int i = 0; i < lim; ++i) a[i] = (long long) a[i] *
      b[i] % mod;
 ntt(a, 0);
  a.resize(n + 1);
  return a;
```

# 8 Programacion dinamica

### 8.1 LIS

```
int main() {
    ios::sync with stdio(false);
    cin.tie(0);
    int n; cin >> n;
    vl vals(n);
    for (int i = 0; i < n; i++) cin >> vals[i];
    vl copia(vals);
    sort(copia.begin(),copia.end());
    map <11,11> dicc;
    for (int i=0;i<n;i++)if (!dicc.count(copia[i])) dicc[</pre>
        copia[i]]=i;
    vl baseSt(n,0);
    nodeSt st(baseSt, 0, n - 1);
    11 \text{ maxi} = 0;
    for (ll pVal:vals) {
        ll op =st.get(0,dicc[pVal]-1)+1;
        maxi = max(maxi, op);
        st.actl(dicc[pVal],op);
    cout << maxi << ln;
```

# 8.2 Bin Packing

```
int main() {
    ll n, capacidad;
    cin >> n >> capacidad;
    vl pesos(n, 0);
    forx(i, n) cin >> pesos[i];
    vector < pll > dp((1 << n));
    dp[0] = \{1, 0\};
    // dp[X] = \{ #numero de paquetes, peso de min paquete \}
    // La idea es probar todos los subset y en cada uno
        preguntarnos
    // quien es mejor para subirse de ultimo buscando
       minimizar
    // primero el numero de paquetes
    for (int subset = 1; subset < (1 << n); subset++) {</pre>
        dp[subset] = \{21, 0\};
        for (int iPer = 0; iPer < n; iPer++) {</pre>
            if ((subset >> iPer) & 1) {
                 pll ant = dp[subset ^ (1 << iPer)];</pre>
                 ll k = ant.ff;
                 ll w = ant.ss;
                 if (w + pesos[iPer] > capacidad) {
                     k++;
                     w = min(pesos[iPer], w);
                 } else {
                     w += pesos[iPer];
                 dp[subset] = min(dp[subset], {k, w});
    cout << dp[(1 << n) - 1].ff << ln;
```

# 8.3 Algoritmo de Kadane 2D

```
int main() {
    ll fil,col;cin>>fil>>col;
    vector<vl> grid(fil,vl(col,0));

// Algoritmo de Kadane/DP para suma maxima de una matriz
    2D en o(n^3)
    for(int i=0;i<fil;i++) {
        for(int e=0;e<col;e++) {
            ll num;cin>>num;
        }
}
```

### 8.4 Knuth Clasico

```
const int N = 1010;
const int INF = (int) 1e9;
int v[N], dp[N][N], sum[N], best[N][N];
int main() {
    ios::sync_with_stdio(0);
    cin.tie(0);
    int n;
    while(cin >> n) {
        if(n == 0) break;
        for(int i = 0; i < n; i++) cin >> v[i];
        for(int i = 0; i < n; i++) {</pre>
             sum[i+1] = sum[i] + v[i];
        for(int i = 0; i < n; i++) best[i][i] = i;</pre>
        for(int len = 2; len <= n; ++len) {</pre>
             for (int i = 0; i+len-1 < n; ++i) {
                 int j = i + len - 1;
                 int &ref = dp[i][j];
                 ref = INF;
                 for(int k = best[i][j-1]; k <= best[i+1][</pre>
                     j]; ++k) {
                     if(k < j) {
                          int cur = dp[i][k] + dp[k+1][j];
                          if(cur < ref) {</pre>
                              best[i][j] = k;
                              ref = cur;
```

```
ref += sum[j+1] - sum[i];
}
cout << dp[0][n-1] << '\n';
}
return 0;
}</pre>
```

#### 8.5 Edit Distances

```
int editDistances(string& wor1, string& wor2) {
    // O(tam1*tam2)
    // minimo de letras que debemos insertar, elminar o
        reemplazar
    // de wor1 para obtener wor2
    11 tam1=wor1.size();
    11 tam2=wor2.size();
    vector<vl> dp(tam2+1, vl(tam1+1, 0));
    for (int i=0; i<=tam1; i++) dp[0][i]=i;</pre>
    for (int i=0;i<=tam2;i++)dp[i][0]=i;</pre>
    0=[0][0]qb
    for(int i=1;i<=tam2;i++) {</pre>
        for (int j=1; j<=tam1; j++) {</pre>
             ll op1 = min(dp[i-1][j], dp[i][j-1])+1;
             // el minimo entre eliminar o insertar
             11 \text{ op2} = dp[i-1][j-1]; // reemplazarlo
             if (wor1[j-1]!=wor2[i-1]) op2++;
             // si el reemplazo tiene efecto o quedo igual
             dp[i][j]=min(op1,op2);
    return dp[tam2][tam1];
```

# 8.6 Divide Conquer

```
int m, n;
vector<long long> dp_before(n), dp_cur(n);
long long C(int i, int j);
// compute dp_cur[l], ... dp_cur[r] (inclusive)
void compute(int l, int r, int optl, int optr) {
    if (l > r)
        return;

    int mid = (l + r) >> 1;
    pair<long long, int> best = {LLONG_MAX, -1};

    for (int k = optl; k <= min(mid, optr); k++) {
        best = min(best, {(k ? dp_before[k - 1] : 0) + C(k, mid), k});
}</pre>
```

```
dp_cur[mid] = best.first;
    int opt = best.second;

    compute(1, mid - 1, opt1, opt);
    compute(mid + 1, r, opt, optr);
}
int solve() {
    for (int i = 0; i < n; i++)
        dp_before[i] = C(0, i);

    for (int i = 1; i < m; i++) {
        compute(0, n - 1, 0, n - 1);
        dp_before = dp_cur;
    }
    return dp_before[n - 1];
}</pre>
```

#### 8.7 Knuth

```
#Condiciones
\#C(b,c) <= C(a,d)
\#C(a,c)+C(b,d) \le C(a,d)+C(b,c)
int solve() {
    int N:
    ... // read N and input
    int dp[N][N], opt[N][N];
    auto C = [\&](int i, int j) {
        ... // Implement cost function C.
    for (int i = 0; i < N; i++) {
        opt[i][i] = i;
        ... // Initialize dp[i][i] according to the
            problem
    for (int i = N-2; i >= 0; i--) {
        for (int j = i+1; j < N; j++) {
            int mn = INT MAX;
            int cost = C(i, j);
            for (int k = opt[i][j-1]; k <= min(j-1, opt[i</pre>
                +1][\dot{1}]); k++) {
                if (mn \ge dp[i][k] + dp[k+1][j] + cost) {
                     opt[i][j] = k;
                     mn = dp[i][k] + dp[k+1][j] + cost;
            dp[i][j] = mn;
    cout << dp[0][N-1] << endl;
```

# 9 Strings

### 9.1 Hashing

```
// 1000234999, 1000567999, 1000111997, 1000777121,
   1001265673, 1001864327, 999727999, 1070777777
const int MOD[2] = { 1001864327, 1001265673 };
const ii BASE(257, 367), ZERO(0, 0), ONE(1, 1);
const int MAXN = 1e6;
inline int add(int a, int b, int mod) {return a+b>=mod?a+b
   -mod:a+b;}
inline int sbt(int a, int b, int mod) {return a-b<0?a-b+</pre>
   mod:a-b;}
inline int mul(int a, int b, int mod) {return ll(a) *b%mod
inline ll operator ! (const ii a) {return (ll(a.first)
   <<32) | a.second; }
inline ii operator + (const ii& a, const ii& b) {return {
   add(a.first, b.first, MOD[0]), add(a.second, b.second,
    MOD[1])};}
inline ii operator - (const ii& a, const ii& b) {return {
   sbt(a.first, b.first, MOD[0]), sbt(a.second, b.second,
    MOD[1])};}
inline ii operator * (const ii& a, const ii& b) {return {
   mul(a.first, b.first, MOD[0]), mul(a.second, b.second,
    MOD[1])};}
ii p[MAXN+1];
void prepare() { // Acordate del prepare()!!
        p[0]=ONE;
        for (int i=1; i<=MAXN; i++) p[i]=p[i-1] *BASE;
ii combine(ii a, ii b, int lenb){return a*p[lenb]+b;}
template <class type>
struct hashing{
        vector<ii> h:
        hashing(type& t) {
                h.resize(sz(t)+1);
                h[0] = ZERO;
                 for (int i=1; i < sz(h); ++i)</pre>
                         h[i]=h[i-1]*BASE+ii\{t[i-1], t[i]\}
                             -1 ] };
        ii get(int 1, int r) {return h[r+1]-h[1]*p[r-1
            +1];}
};
```

#### 9.2 KMP

```
// O(n)
vi phi(string& s){
        int n=sz(s);
        vi tmp(n);
        for(int i=1, j=0; i<n;++i) {
                 while(j>0 && s[j]!=s[i])j=tmp[j-1];
                 if(s[i]==s[j])j++;
                 tmp[i]=i;
        return tmp;
// O(n+m)
int kmp(string& s, string& p) {
        int n=sz(s), m=sz(p), cnt=0;
        vi pi=phi(p);
        for(int i=0, j=0; i<n; ++i) {
                 while(j && s[i]!=p[j])j=pi[j-1];
                 if(s[i]==p[j])j++;
                 if ( j==m) {
                         cnt++;
                         j=pi[j-1];
        return cnt;
```

### 9.3 KMP Automaton

### 9.4 Manacher

```
// O(n), par (raiz, izq, der) 1 - impar 0
vi manacher(string& s, int par) {
    int l=0,r=-1,n=sz(s);vi m(n,0);
```

# 9.5 Minimum Expression

```
// O(n)
int minimum_expression(string s) {
    s=s+s;int n=sz(s),i=0,j=1,k=0;
    while(i+k<n && j+k<n) {
        if(s[i+k]==s[j+k])k++;
        else if(s[i+k]>s[j+k])i=i+k+1,k=0; //
            cambiar por < para max
        else j=j+k+1,k=0;
        if(i==j)j++;
    }
    return min(i, j);
}</pre>
```

#### 9.6 Palindromic Tree

```
const int alpha = 26;
const char fc = 'a';
// tree suf is the longest suffix palindrome
// tree dad is the palindrome add c to the right and left
struct Node {
        int next[alpha];
        int len, suf, dep, cnt, dad;
};
// O(nlogn)
struct PalindromicTree{
        vector<Node> tree;
        string s:
        int len,n;
        int size; // node 1 - root with len -1, node 2 -
            root with len 0
        int last; // max suffix palindrome
        bool addLetter(int pos) {
                int cur=last, curlen=0;
                int let=s[pos]-fc;
                while (true) {
                         curlen=tree[cur].len;
```

```
if(pos-1-curlen>=0 && s[pos-1-
                            curlen] == s[pos]) break;
                         cur=tree[cur].suf;
                if(tree[cur].next[let]){
                         last=tree[cur].next[let];
                         tree[last].cnt++;
                         return false;
                size++;
                last=size;
                tree[size].len=tree[cur].len+2;
                tree[cur].next[let]=size;
                tree[size].cnt=1;
                tree[size].dad=cur;
                if(tree[size].len==1){
                         tree[size].suf=2;
                         tree[size].dep=1;
                         return true;
                while(true) {
                         cur=tree[cur].suf;
                         curlen=tree[cur].len;
                         if(pos-1-curlen>=0 && s[pos-1-
                            curlen] == s[pos]) {
                                 tree[size].suf=tree[cur].
                                     next[let];
                                 break;
                tree[size].dep=1+tree[tree[size].suf].dep
                return true;
        PalindromicTree(string& _s, int n) {
                tree.assign(n+4, Node());
                tree[1].len=-1; tree[1].suf=1;
                tree[2].len=0;tree[2].suf=1;
                size=2; last=2; s= s;
                for(int i=0;i<n;i++)addLetter(i);</pre>
                for (int i=size; i>=3; i--) tree[tree[i].suf
                    l.cnt+=tree[i].cnt;
};
```

# 9.7 Suffix Array

```
// O(nlogn)
struct SuffixArray{
```

```
const int alpha = 256;
string s; int n;
vi sa, rnk, lcp;
void build(string& _s) {
         s= s;s.push_back('$'); // check
         n=sz(s);
         sa.assign(n, 0);
         rnk.assign(n, 0);
         lcp.assign(n-1, 0);
         buildSA();
void buildSA() {
         vi cnt(max(alpha, n),0);
         for (int i=0; i < n; ++i) cnt[s[i]]++;</pre>
         for (int i=1; i < max (alpha, n); ++i) cnt[i] +=</pre>
             cnt[i-1];
         for (int i=n-1; i>=0; --i) sa[--cnt[s[i]]]=i;
         for(int i=1;i<n;++i)rnk[sa[i]]=rnk[sa[i</pre>
             -1] + (s[sa[i]]!=s[sa[i-1]]);
         for (int k=1; k < n; k *=2) {
                  vi nsa(n),nrnk(n),ncnt(n);
                  for (int i=0;i<n;++i)sa[i]=(sa[i]-</pre>
                      k+n)%n;
                  for (int i=0;i<n;++i)ncnt[rnk[i</pre>
                      ]]++;
                  for (int i=1; i < n; ++i) ncnt[i] += ncnt</pre>
                      [i-1];
                  for(int i=n-1; i>=0; --i) nsa[--ncnt
                      [rnk[sa[i]]]]=sa[i];
                  for (int i=1; i < n; ++i) {</pre>
                           ii op1={rnk[nsa[i]], rnk
                               [(nsa[i]+k)%n]};
                           ii op2={rnk[nsa[i-1]],}
                               rnk[(nsa[i-1]+k)%n];
                           nrnk[nsa[i]]=nrnk[nsa[i
                               -1]]+(op1!=op2);
                  swap(sa, nsa);swap(rnk, nrnk);
         for (int i=0, k=0; i < n-1; ++i) {</pre>
                  while (s[i+k] == s[sa[rnk[i]-1]+k])k
                  lcp[rnk[i]-1]=k;
                  if(k)k--;
```

#### 9.8 Suffix Automaton

};

```
// O(n*log(alpha))
```

```
struct SuffixAutomaton{
        vector<map<char,int>> to;
        vector<bool> end; vi suf, len; // len, longest
            string
        int last;
        SuffixAutomaton(string& s){
                to.push back(map<char,int>());
                suf.push back (-1);
                len.push back(0);
                last=0;
                for(int i=0;i<sz(s);i++) {</pre>
                         to.push_back(map<char,int>());
                         suf.push_back(0);
                         len.push back(i+1);
                         int r=sz(to)-1;
                         int p=last;
                         while (p>=0 \&\& to[p].find(s[i])==
                            to[p].end()){
                                 to[p][s[i]]=r;
                                 p=suf[p];
                         if(p!=-1){
                                  int q=to[p][s[i]];
                                 if(len[p]+1==len[q]){
                                          suf[r]=q;
                                  }else{
                                          to.push back(to[q
                                              1);
                                          suf.push back(suf
                                              [q]);
                                          len.push back(len
                                              [p]+1);
                                          int qq=sz(to)-1;
                                          suf[q]=qq;
                                          suf[\bar{r}]=qq;
                                          while(p>=0 && to[
                                             p][s[i]] == q){
                                                   to[p][s[i
                                                      ] =qq;
                                                   p=suf[p];
                         last=r;
                end.assign(sz(to), false);
                int p=last;
                while(p){
                         end[p]=true;
                         p=suf[p];
```

# 9.9 Suffix Tree

};

```
// O(n)
struct SuffixTree{
        vector<map<char,int>> to;
        vector<int> pos,len,link;
        const int inf = 1e9;
        int size=0;
        string s;
        int make(int pos, int len) {
                to.push back(map<char,int>());
                 pos.push_back(_pos);
                len.push back (len);
                link.push back (-1);
                return size++;
        void add(int& p, int& lef, char c) {
                 s+=c;++lef;int lst=0;
                 for(;lef;p?p=link[p]:lef--){
                         while (lef>1 && lef>len[to[p][s[sz
                             (s)-lef]]){
                                 p=to[p][s[sz(s)-lef]], lef
                                     -=len[p];
                         char e=s[sz(s)-lef];
                         int& q=to[p][e];
                         if(!q){
                                 q=make(sz(s)-lef,inf),
                                     link[lst]=p,lst=0;
                         }else{
                                 char t=s[pos[q]+lef-1];
                                 if(t==c) {link[lst]=p;
                                     return; }
                                 int u=make(pos[q],lef-1);
                                 to[u][c]=make(sz(s)-1,inf
                                 );
to[u][t]=q;
                                  pos[q] += lef-1;
                                 if(len[q]!=inf)len[q]-=
                                     lef-1;
                                 q=u,link[lst]=u,lst=u;
        void build(string& s){
                make (-1, \bar{0}); int p=0, lef=0;
                for (char c:_s) add (p, lef, c);
                 add(p,lef,'$');
                 s.pop_back();
```

#### 9.10 Trie

```
const int maxn = 2e6+5, alpha = 26, bits = 30;
int to[maxn][alpha],cnt[maxn],act;
void init(){
        for (int i=0;i<=act;++i) {</pre>
                cnt[i]=0;
                // suf[i]=dad[i]=0;
                // adi[i].clear();
                memset(to[i], 0, sizeof(to[i]));
        act=0:
int add(string& s){
        int u=0;
        for(char ch:s) {
                int c=conv(ch);
                if(!to[u][c])to[u][c]=++act;
                u=to[u][c];
        cnt[u]++;
        return u;
// Aho-Corasick
vector<int> adj[maxn]; // dad or suf
int dad[maxn], suf[maxn];
// O(sum(n)*alpha)
void build(){
        queue<int> q{{0}};
```

```
while(!q.empty()) {
    int u=q.front();q.pop();
    for(int i=0;i<alpha;++i) {
        int v=to[u][i];
        if(!v)to[u][i]=to[suf[u]][i];
        else q.push(v);
        if(!u || !v)continue;
        suf[v]=to[suf[u]][i];
        dad[v]=cnt[suf[v]]?suf[v]:dad[suf[v]];
        }
}
for(int i=1;i<=act;++i) {
        adj[i].push_back(dad[i]);
        adj[dad[i]].push_back(i);
}</pre>
```

# 9.11 Z Algorithm

### 10 Misc

# 10.1 Counting Sort

# 10.2 Expression Parsing

```
// En python es eval()
bool delim(char c) {return c==' ';}
bool is op(char c){return c=='+' || c=='-' || c=='*' || c
bool is unary (char c) {return c=='+' | | c=='-';}
int priority(char op) {
        if(op<0)return 3;</pre>
        if (op=='+' | op=='-') return 1;
        if(op=='*' || op=='/') return 2;
        return -1;
void process_op(stack<int>& st, char op){
        if(op<0){
                int l=st.top();st.pop();
                switch(-op) {
                         case '+':st.push(1);break;
                         case '-':st.push(-1);break;
        }else{
                int r=st.top();st.pop();
                int l=st.top();st.pop();
                switch(op){
                         case '+':st.push(l+r);break;
                         case '-':st.push(l-r);break;
                         case '*':st.push(l*r);break;
                         case '/':st.push(l/r);break;
int evaluate(string& s){
        stack<int> st;
        stack<char> op:
        bool may be unary=true;
        for (int i=0; i < sz(s); ++i) {</pre>
                if (delim(s[i])) continue;
                if(s[i] == '('){
                         op.push('(');
                         may be unary=true;
                }else if(s[i]==')'){
                         while (op.top()!='('){
                                 process_op(st, op.top());
                                 op.pop();
                         op.pop();
                         may_be_unary=false;
                }else if(is op(s[i])){
                         char cur op=s[i];
                         if (may_be_unary && is_unary(
```

```
cur op))cur op=-cur op;
                 while(!op.empty() && ((cur_op >=
                    0 && priority(op.top()) >=
                    priority(cur op)) || (cur op <</pre>
                     0 && priority(op.top()) >
                    priority(cur_op)))){
                         process_op(st, op.top());
                         op.pop();
                 op.push(cur op);
                 may be unary=true;
        }else{
                 int number=0;
                 while(i<sz(s) && isalnum(s[i]))</pre>
                    number=number * 10+s[i++]-'0';
                 st.push(number);
                may be unary=false;
while(!op.empty()){
        process op(st, op.top());
```

# 11 Teoría y miscelánea

#### 11.1 Sumatorias

• 
$$\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$$

• 
$$\sum_{i=1}^{n} i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

• 
$$\sum_{i=1}^{n} i^5 = \frac{(n(n+1))^2 (2n^2 + 2n - 1)}{12}$$

$$\bullet \ \sum_{i=1}^{n} i^3 = \left(\frac{n(n+1)}{2}\right)^2$$

• 
$$\sum_{i=0}^{n} x^i = \frac{x^{n+1}-1}{x-1}$$
 para  $x \neq 1$ 

### 11.2 Teoría de Grafos

#### 11.2.1 Teorema de Euler

En un grafo conectado planar, se cumple que V-E+F=2, donde V es el número de vértices, E es el número de aristas y F es el número de caras. Para varios componentes la formula es: V-E+F=1+C, siendo  ${\bf C}$  el número de componentes.

#### 11.2.2 Planaridad de Grafos

Un grafo es planar si y solo si no contiene un subgrafo homeomorfo a  $K_5$  (grafo completo con 5 vértices) ni a  $K_{3,3}$  (grafo bipartito completo con 3 vértices en cada conjunto).

```
op.pop();
}
return st.top();
}
```

## 10.3 Ternary Search

```
// O(log((r-1)/eps))
double ternary(){
          double l, r;
          while(r-1>eps) {
                double m1=l+(r-1)/3.0;
                double m2=r-(r-1)/3.0;
                if(f(m1) < f(m2)) l=m1;
                else r=m2;
        }
        return max(f(1),f(r));
}</pre>
```

### 11.3 Teoría de Números

#### 11.3.1 Ecuaciones Diofánticas Lineales

Una ecuación diofántica lineal es una ecuación en la que se buscan soluciones enteras x e y que satisfagan la relación lineal ax+by=c, donde a, b y c son constantes dadas.

Para encontrar soluciones enteras positivas en una ecuación diofántica lineal, podemos seguir el siguiente proceso:

- 1. Encontrar una solución particular: Encuentra una solución particular  $(x_0, y_0)$  de la ecuación. Esto puede hacerse utilizando el algoritmo de Euclides extendido.
- 2. Encontrar la solución general: Una vez que tengas una solución particular, puedes obtener la solución general utilizando la fórmula:

$$x = x_0 + \frac{b}{\operatorname{mcd}(a, b)} \cdot t$$

$$y = y_0 - \frac{a}{\operatorname{mcd}(a, b)} \cdot t$$

donde t es un parámetro entero.

3. Restringir a soluciones positivas: Si deseas soluciones positivas, asegúrate de que las soluciones generales satisfagan  $x \ge 0$  y  $y \ge 0$ . Puedes ajustar el valor de t para cumplir con estas restricciones.

#### 11.3.2 Pequeño Teorema de Fermat

Si p es un número primo y a es un entero no divisible por p, entonces  $a^{p-1} \equiv 1 \pmod{p}$ .

#### 11.3.3 Teorema de Euler

Para cualquier número entero positivo n y un entero a coprimo con n, se cumple que  $a^{\phi(n)} \equiv 1 \pmod{n}$ , donde  $\phi(n)$  es la función phi de Euler, que representa la cantidad de enteros positivos menores que n y coprimos con n.

#### 11.4 Geometría

#### 11.4.1 Teorema de Pick

Sea un poligono simple cuyos vertices tienen coordenadas enteras. Si B es el numero de puntos enteros en el borde, I el numero de puntos enteros en el interior del poligono, entonces el area A del poligono se puede calcular con la formula:

$$A = I + \frac{B}{2} - 1$$

#### 11.4.2 Fórmula de Herón

Si los lados del triángulo tienen longitudes a, b y c, y s es el semiperímetro (es decir,  $s = \frac{a+b+c}{2}$ ), entonces el área A del triángulo está dada por:

$$A = \sqrt{s(s-a)(s-b)(s-c)}$$

### 11.4.3 Relación de Existencia Triangular

Para un triángulo con lados de longitud  $a,\,b,\,{\bf y}\,c,$  la relación de existencia triangular se expresa como:

$$b - c < a < b + c$$
,  $a - c < b < a + c$ ,  $a - b < c < a + b$ 

### 11.5 Combinatoria

#### 11.5.1 Permutaciones

El número de permutaciones de n objetos distintos tomados de a r a la vez (sin repetición) se denota como P(n,r) y se calcula mediante:

$$P(n,r) = \frac{n!}{(n-r)!}$$

#### 11.5.2 Combinaciones

El número de combinaciones de n objetos distintos tomados de a r a la vez (sin repetición) se denota como C(n,r) o  $\binom{n}{r}$  y se calcula mediante:

$$C(n,r) = \binom{n}{r} = \frac{n!}{r!(n-r)!}$$

#### 11.5.3 Permutaciones con Repetición

El número de permutaciones de n objetos tomando en cuenta repeticiones se denota como  $P_{\text{rep}}(n; n_1, n_2, \dots, n_k)$  y se calcula mediante:

$$P_{\text{rep}}(n; n_1, n_2, \dots, n_k) = \frac{n!}{n_1! n_2! \cdots n_k!}$$

#### 11.5.4 Combinaciones con Repetición

El número de combinaciones de n objetos tomando en cuenta repeticiones se denota como  $C_{\text{rep}}(n; n_1, n_2, \dots, n_k)$  y se calcula mediante:

$$C_{\text{rep}}(n; n_1, n_2, \dots, n_k) = \binom{n+k-1}{n} = \binom{n+k-1}{k-1}$$

#### 11.5.5 Números de Catalan

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

Los números de Catalan también pueden calcularse utilizando la siguiente fórmula recursiva:

$$C_0 = 1$$

$$C_{n+1} = \frac{4n+2}{n+2}C_n$$

Usos:

- Cat(n) cuenta el número de árboles binarios distintos con n vértices.
- Cat(n) cuenta el número de expresiones que contienen n pares de paréntesis correctamente emparejados.
- Cat(n) cuenta el número de formas diferentes en que se pueden colocar n+1 factores entre paréntesis, por ejemplo, para n=3 y 3+1=4 factores: a,b,c,d, tenemos: (ab)(cd),a(b(cd)),((ab)c)d y a((bc)d).
- Los números de Catalan cuentan la cantidad de caminos no cruzados en una rejilla  $n \times n$  que se pueden trazar desde una esquina de un cuadrado o rectángulo a la esquina opuesta, moviéndose solo hacia arriba y hacia la derecha.

- $\bullet$  Los números de Catalan representan el número de árboles binarios completos con n+1 hojas.
- $\operatorname{Cat}(n)$  cuenta el número de formas en que se puede triangular un poligono convexo de n+2 lados. Otra forma de decirlo es como la cantidad de formas de dividir un polígono convexo en triángulos utilizando diagonales no cruzadas.

#### 11.5.6 Estrellas y barras

Número de soluciones de la ecuación  $x_1 + x_2 + \cdots + x_k = n$ .

• Con  $x_i \ge 0$ :  $\binom{n+k-1}{n}$ 

• Con  $x_i \ge 1$ :  $\binom{n-1}{k-1}$ 

Número de sumas de enteros con límite inferior:

Esto se puede extender fácilmente a sumas de enteros con diferentes límites inferiores. Es decir, queremos contar el número de soluciones para la ecuación:

$$x_1 + x_2 + \dots + x_k = n$$

 $con x_i \geq a_i$ .

Después de sustituir  $x_i' := x_i - a_i$  recibimos la ecuación modificada:

$$(x'_1 + a_i) + (x'_2 + a_i) + \dots + (x'_k + a_k) = n$$

$$\Leftrightarrow x_1' + x_2' + \dots + x_k' = n - a_1 - a_2 - \dots - a_k$$

con  $x_i' \ge 0$ . Así que hemos reducido el problema al caso más simple con  $x_i' \ge 0$  y nuevamente podemos aplicar el teorema de estrellas y barras.

### 11.6 DP Optimization Theory

Name	Original Recurrence	Sufficient Condition	From	То
CH 1	$dp[i] = min_{j < i} \{dp[j] + b[j] *$	$b[j] \ge b[j+1]$ Option-	$O(n^2)$	O(n)
	$a[i]\}$	ally $a[i] \le a[i+1]$		
CH 2	$dp[i][j] = min_{k < j} \{ dp[i - $	$b[k] \ge b[k+1]$ Option-	$O(kn^2)$	O(kn)
	1][k] + b[k] * a[j]	ally $a[j] \le a[j+1]$		
D&Q	$dp[i][j] = min_{k < j} \{ dp[i - ]$	$A[i][j] \le A[i][j+1]$	$O(kn^2)$	$O(kn\log n$
	$1][k] + C[k][j]\}$			
Knuth	dp[i][j] =	$A[i, j-1] \le A[i, j] \le$	$O(n^3)$	$O(n^2)$
	$min_{i < k < j} \{dp[i][k] +$	A[i+1,j]		
	$dp[k][j]\} + C[i][j]$			

Notes:

- A[i][j] the smallest k that gives the optimal answer, for example in dp[i][j] = dp[i-1][k] + C[k][j]
- C[i][j] some given cost function
- We can generalize a bit in the following way  $dp[i] = \min_{j < i} \{F[j] + b[j] * a[i]\},$  where F[j] is computed from dp[j] in constant time