Notebook UNTreeCiclo

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| 1 (| · · · · · · · · · · · · · · · · · · · |
| 1.1 | C++ plantilla |
| #ir | clude <bits stdc++.h=""></bits> |
| usi #de #de #de | <pre>ing namespace std; efine all(v) v.begin(), v.end() efine sz(arr) ((int) arr.size()) efine rep(i, a, b) for(int i = a; i < (b); ++i) edef pair<int, int=""> ii;</int,></pre> |
| tvr | <pre>pedef vector<int> vi; pedef long long ll; ust char ln = '\n';</int></pre> |
| #de | fine watch(x) cout<<#x<<"="< <x<<'\n'< td=""></x<<'\n'<> |

```
#define print(arr) for(auto& x:arr)cout<<x<<" ";cout<<"\n</pre>
typedef long double ld;
typedef vector<ii> vii;
typedef vector<long long> v1;
typedef pair<ll, ll> pll;
typedef vector<pll> vll;
const int INF = 1e9;
const ll INFL = 1e18;
const int MOD = 1e9+7;
const double EPS = 1e-9;
const ld PI = acosl(-1);
int dirx[4] = \{0, -1, 1, 0\};
int diry[4] = \{-1, 0, 0, 1\};
int dr[] = \{1, 1, 0, -1, -1, -1, 0, 1\};
int dc[] = \{0, 1, 1, 1, 0, -1, -1, -1\};
const string ABC = "abcdefghijklmnopgrstuvwxyz";
void main2(){
int main() {
        ios::sync with stdio(false);
        cin.tie(0);
        cout << setprecision(20) << fixed;</pre>
    // freopen("file.in", "r", stdin);
// freopen("file.out", "w", stdout);
        clock t start = clock();
        main2();
        cerr<<double(clock()-start)/CLOCKS PER SEC<<" s\n
        return 0;
```

1.2 Librerias

```
// En caso de que no sirva #include <bits/stdc++.h>
#include <algorithm>
#include <iostream>
#include <iterator>
#include <sstream>
#include <fstream>
#include <cassert>
#include <climits>
#include <cstdlib>
#include <cstring>
#include <string>
#include <cstdio>
#include <vector>
#include <cmath>
#include <queue>
#include <deque>
#include <stack>
```

```
#include <list>
#include <map>
#include <set>
#include <bitset>
#include <iomanip>
#include <unordered map>
#include <tuple>
#include <random>
#include <chrono>
```

1.3 Create

```
import os
def folder(problem):
        os.makedirs(problem, exist_ok=True)
        with open(os.path.join(problem, "main.cpp"), "w")
            as f:
                f.write("")
        with open (os.path.join (problem, "in.txt"), "w")
           as f:
                f.write("")
with open ("plantilla.cpp", "w") as f:
        f.write("")
with open("out.txt", "w") as f:
        f.write("")
for i in range (ord('A'), ord('P') + 1):
        folder(chr(i))
```

1.4 Bitmask

```
los corrimientos.
                 -> Verifica si x es impar
x & 1
x & (1<<i)
                -> Verifica si el i-esimo bit esta
   encendido
x = x \mid (1 << i) \rightarrow Enciende el i-esimo bit
x = x \& (1 << i) -> Apaga el i-esimo bit
x = x ^ (1 << i) -> Invierte el i-esimo bit
                -> Invierte todos los bits
x = x
                -> Devuelve el bit encendido mas a la
x & -x
   derecha (potencia de 2, no el indice)
                -> Devuelve el bit apagado mas a la
^{\sim} x & (x+1)
   derecha (potencia de 2, no el indice)
x = x \mid (x+1) -> Enciende el bit apagado mas a la
   derecha
x = x & (x-1)
                -> Apaga el bit encendido mas a la
   derecha
```

* Operaciones a nivel de bits. Si n es ll usar 111<< en

```
-> Apaga en x los bits encendidos de y
x = x \& v
* Funciones del compilador qcc. Si n es ll agregar el
   sufijo ll, por ej: __builtin_clzll(n).
__builtin_clz(x)
                      -> Cantidad de bits apagados por la
    izguierda
__builtin_ctz(x)
                      -> Cantidad de bits apagados por la
    derecha. Indice del bit encendido mas a la derecha
__builtin_popcount(x) -> Cantida de bits encendidos
builtin ffs(x)
                          -> Posicion del primer bit
   prendido (lsb+1)
* Logaritmo en base 2 (entero). Indice del bit encendido
   mas a la izquierda. Si x es ll usar 63 y clzll(x).
// 0(1)
int lg2(const int &x) { return 31-__builtin_clz(x); }
* Itera, con indices, los bits encendidos de una mascara.
// O(#bits encendidos)
for (int x = mask; x; x &= x-1) {
        int i = __builtin_ctz(x);
* Itera todas las submascaras de una mascara. (Iterar
   todas las submascaras de todas las mascaras es O(3^n))
// O(2^(#bits encendidos))
for (int sub = mask; ; sub = (sub-1)&mask) {
        // ...
        if (sub == 0) break;
// Ascendente
for(int sub = 0; ; sub = (sub-mask) &mask) {
        // ...
        if (sub == mask) break;
* retorna la siguiente mask con la misma cantidad
   encendida
ll nextMask(ll x){
        11 c = x \& -x;
        11 r = x + c;
        return (((r ^ x) >> 2) / c) | r;
// optimiza el .count de los bitsets y el popcount
#pragma GCC target("popent")
// Formulas
a \mid b = a \cdot b + a \cdot b
a (a \& b) = (a | b) b
b^{(a \& b)} = (a | b)^{a}
(a \& b) \hat{} (a | b) = a \hat{} b
a + b = a | b + a \& b
```

 $a + b = a \cdot b + 2 * (a \& b)$

```
a - b = (a ^ (a & b)) - ((a | b) ^ a)
a - b = ((a | b) ^ b) - ((a | b) ^ a)
a - b = (a ^ (a & b)) - (b ^ (a & b))
a - b = ((a | b) ^ b) - (b ^ (a & b))
a ^ b = ^ (a & b) & (a | b)
si (x < y < z) entonces min(x^y, y^z) < (x^z)
```

1.5 Custom Hashing

```
struct custom_hash {
        static long long splitmix64(long long x) {
                x += 0x9e3779b97f4a7c15;
                x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
                x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
                return x ^ (x >> 31);
        size t operator()(long long x) const {
                static const long long FIXED_RANDOM =
                   chrono::steady clock::now().
                   time since epoch().count();
                return splitmix64(x + FIXED RANDOM);
        size_t operator()(const pair<int,int>& x) const {
                return (size t) x.first * 37U + (size t)
                   x.second;
        size_t operator()(const vector<int>& v) const {
                size t s = 0;
                for(auto &e : v)
                        s^=hash<int>()(e)+0x9e3779b9+(s
                           <<6)+(s>>2);
                return s;
};
unordered_map<long long, int, custom_hash> safe_map; //
   unordered map or qp hash table
safe map.max load factor(0.25);
safe_map.reserve(1024); // potencia de 2 mas cercana
multitest - no usar reserve (por el clear, es pesado)
```

1.6 Random

```
typedef unsigned long long u64;
mt19937_64 rng (chrono::steady_clock::now().
    time_since_epoch().count());
u64 xor_hash=rng();
// return random numbers in the range [1,r]
```

```
mt19937 rng (chrono::steady_clock::now().time_since_epoch
    ().count());
double rand(double l, double r) {return
    uniform_real_distribution<double>(l, r)(rng);}
int rand(int l, int r) {return uniform_int_distribution<
    int>(l, r)(rng);}
shuffle(all(vector), rng);
```

1.7 Cosas de strings

```
// si el caracter que separa el texto es distinto al
   espacio
// entonces descomentar el segundo parametro y cambiar el
    while por el otro
vector<string> split(const string &s/*, char c */) {
        vector<string> v;
        stringstream ss(s);
        string sub:
        while(ss>>sub)v.push back(sub);
        // while (getline (ss, sub, c)) if (sz(sub)) v.push back
            (sub);
        return v;
string s;
for (char& c:s) c=toupper(c);
for (char& c:s) c=tolower(c);
int n=stoi(s); // string -> int
int n=stoi(s, nullptr, 2); // bin string -> int
double d=stod(s); // string -> double
string s=to_string(n); // int -> string
cout << "\U0001F600"; // emojis
Quitar repetidos (lo pongo aca porque no se donde mas
   ponerlo)
sort(all(bs));
bs.resize(unique(all(bs)) - bs.begin());
```

2 Arboles

2.1 Centroid Decomposition

```
// O(n*log(n))
// 1) init(adj,n);
struct CentroidDecomposition{
    vector<vi> adj;
    vi dad,sz,proc;

    int operator[](int i){return dad[i];}
    void init(vector<vi>& adj2, int n){
        proc.assign(n,false);
```

```
dad.resize(n);
                sz.resize(n);
                adj=adj2;
                build();
        void build(int v=0, int p=-1) {
                int n=dfsSz(v, p);
                int centroid=dfsCentroid(v, p, n);
                dad[centroid]=p;
                // anadir dfs para el conteo de caminos
                proc[centroid]=true;
                for(int u:adj[centroid]) {
                         if (u==p || proc[u]) continue;
                         build(u,centroid);
        int dfsSz(int v,int p) {
                sz[v]=1;
                for(int u:adj[v]){
                         if (u==p || proc[u]) continue;
                         sz[v] += dfsSz(u, v);
                return sz[v];
        int dfsCentroid(int v, int p, int n) {
                for(int u:adj[v]){
                         if(u==p || proc[u])continue;
                         if (sz[u]>n/2) return dfsCentroid(u
                            , v, n);
                return v;
};
// para el arbol de centroides
// for (int b=a;b!=-1;b=cd[b])
```

2.2 Hash Tree

```
const int MOD=1e9+97;
const int P[2]={998244353,1000000007};
const int Q[2]={1000000033,1000000021};
const int R[2]={123456789,987654321};

int add(int a, int b) {return a+b>=MOD?a+b-MOD:a+b;}
int mul(int a, int b) {return ll(a)*b%MOD;}
int binpow(int a, int b) {
    int res=1;a%=MOD;
    while(b>0) {
        if(b&1) res=mul(res,a);
        a=mul(a,a);
        b>>=1;
```

```
return res%MOD:
// O(n), 1-indexed
struct Tree{
        vector<vi> q;
        int n;
        Tree(int _n):n(_n) {q.resize(n+1);}
        void add edge(int u, int v) {
                q[u].push back(v);
                q[v].push_back(u);
        ii hash(int u, int pre=0){
                vector<vi> nw(2,vi());
                for(int v:q[u])
                         if(v!=pre){
                                 ii tmp=hash(v,u);
                                 nw[0].push_back(tmp.first
                                 nw[1].push back(tmp.
                                     second);
                ii ans=\{0,0\};
                for(int i=0;i<2;++i){</pre>
                         int& tmp=(i?ans.second:ans.first)
                         for(int x:nw[i])tmp=add(tmp,
                            binpow(P[i], x));
                         tmp=add(mul(tmp,Q[i]),R[i]);
                return ans;
        // Isomorphism
        bool iso(Tree& t) {
                vi a=get centers();
                vi b=t.get centers();
                for (int x:a) for (int y:b) if (hash (x) ==t.
                    hash(y))return 1;
                return 0;
        vi get centers(){
                auto du=bfs(1);
                int v=max element(all(du))-du.begin();
                auto dv=bfs(v);
                int u=max_element(all(dv))-dv.begin();
                du=bfs(u);
                vi ans;
                for (int i=1; i<=n; ++i) {</pre>
                         if(du[i]+dv[i]==du[v] && du[i]>=
                            du[v]/2 \&\& dv[i] >= du[v]/2) {
                                 ans.push back(i);
```

```
return ans;
        vi bfs(int s){
                 queue<int> q;
                 vi d(n+1, n+2);
                 d[0] = -1;
                 q.push(s);
                 d[s]=0;
                 while(!q.empty()){
                         int u=q.front();
                         q.pop();
                         for(int v:q[u])
                                  if(d[u]+1<d[v]){
                                          d[v]=d[u]+1;
                                          q.push(v);
                 return d:
};
```

2.3 Heavy Light Decomposition

```
typedef long long T;
T oper(T a, T b) {return max(a,b);}
T null=-1e18;
struct SegTree{}; // Add Segment tree
// O(nlog(n)) build
// O(log(n)^2) (query - update) path
// O(log(n)) (query - update) subtree, node
// 1) call build(adj,n,root)
struct HLD{
        SeaTree st;
        vector<vi> adj;
        vi dad, root, dep, sz, pos;
        int time;
        bool edges=false; // if the values are on edges
            instead of nodes
        void build(vector<vi>& adj2, int n, int v=0) { //
            v is the root
                adj=adj2;
                dad.resize(n);
                root.resize(n);
                dep.resize(n);
                sz.resize(n);
                pos.resize(n);
                root[v]=dad[v]=v;
                dep[v]=time=0;
                dfsSz(v);
```

```
dfsHld(v);
        // vector<T> palst(n);
        // for(int i=0;i<n;++i)palst[pos[i]]=vals
        // st.build(palst);
        st.build(n);
void dfsSz(int x) {
        sz[x]=0;
        for(int& y:adj[x]){
                 if (y==dad[x]) continue;
                 dad[y]=x; dep[y]=dep[x]+1;
                 dfsSz(v);
                 sz[x] += sz[y] +1;
                 if(sz[y]>sz[adj[x][0]]) swap(y, adj
                     [x][0];
void dfsHld(int x) {
        pos[x]=time++;
        for(int v:adi[x]){
                 if (y==dad[x]) continue;
                 root[y] = (y = adj[x][0]?root[x]:y);
                 dfsHld(v);
// O(log(n)^2)
template <class Oper>
void processPath(int x, int y, Oper op) {
        for (; root [x]!=root [y]; y=dad[root [y]]) {
                 if (dep[root[x]]>dep[root[y]]) swap
                     (x,y);
                 op(pos[root[y]],pos[y]);
        if(dep[x]>dep[y])swap(x,y);
        op(pos[x]+edges,pos[y]);
void modifyPath(int x, int y, int v) {
        processPath(x,y,[this,&v](int 1, int r){
                 st.upd(l,r,v);
        });
T queryPath(int x, int y) {
        T res=null;
        processPath(x,y,[this,&res](int 1, int r)
                 res=oper(res, st.get(l,r));
        });
        return res;
// O(\log(n))
```

```
void modifySubtree(int x, int v) {
                st.upd(pos[x]+edges,pos[x]+sz[x],v);
        T quervSubtree(int x) {
                return st.get(pos[x]+edges,pos[x]+sz[x]);
        void modify(int x, int v) {st.set(pos[x],v);}
        void modifyEdge(int x, int y, int v) {
                if (dep[x] < dep[y]) swap(x, y);
                modifv(x,v);
};
```

2.4 Kruskal Reconstruction Tree

```
// Kruskal Reconstruction Tree (KRT)
// the main idea is to build a tree to efficiently answer
// about the minimum or maximum edge weight between two
   nodes.
// each edge will be represented as a node in the tree.
// query (a,b) = lca(a,b)
// Add LCA
const int maxn = 1e5+5;
const int maxm = 2e5+5;
vector<vi> adi;
// sometimes it is useful
int ver[2*(maxn+maxm)]; // node at position i in euler
   tour
int st[maxn+maxm]; // start time of v
int ft[maxn+maxm]; // finish time of v
struct DSU{
        vi p, size;
        vector<bool> roots; // if the graph is a forest
        DSU(int n) {
                p.assign(n,0);
                size.assign(n,1);
                roots.assign(n,true);
                for(int i=0; i<n; ++i)p[i]=i;
        int get(int a) {return (a==p[a]?a:p[a]=get(p[a]))
        // unite node a and node b with the edge m =>
           node m
        void unite(int a, int b, int m) {
                a=get(a);b=get(b);
                if (a==b) return;
                size[m]=size[a]+size[b];
                p[a]=p[b]=m;
                roots[a]=false;
```

```
roots[b]=false;
                adj[m].push_back(a);
                adj[m].push back(b);
};
```

2.5 LCA Binary Lifting

```
// O(n*log(n)) build
// O(\log(n)) kth, lca, dist
struct LCA{
        vector<vi> up;
        vi dep;
        int n, maxlog;
        void build(vector<vi>& adj, int root) {
                 n=sz(adj);
                 \max \log = ceil(\log 2(n)) + 3;
                 up.assign(n, vi(maxlog, -1));
                 dep.assign(n,0);
                 dfs(adj,root);
                 calc(n);
        void dfs(vector<vi>& adj, int v=0, int p=-1) {
                 up[v][0]=p;
                 for(int u:adj[v]){
                         if (u==p) continue;
                         dep[u]=dep[v]+1;
                         dfs(adj, u, v);
        void calc(int n) {
                 for (int l=1; l<maxlog; ++1) {</pre>
                         for(int i=0; i<n; ++i) {
                                  if (up[i][l-1]!=-1) {
                                          up[i][l]=up[up[i
                                              ][1-1]][1-1];
        // kth ancestor, return -1 if it doesnt exits
        int kth(int u, int k){
                 for (int l=maxlog-1; l>=0; --1) {
                         if(u!=-1 && k&(1<<1)){
                                  u=up[u][1];
                 return u:
        int lca(int a, int b) {
```

2.6 LCA RMQ

```
// Add RMO - Min
typedef int T;
struct Table{
        void build(vector<T>& a);
        int get(int 1, int r);
};
// O(n*log(n)) build
// O(1) lca
struct LCA{
        Table rmq;
        vi time, path, tmp;
        int n,ti;
        void build(vector<vi>& adj, int root) {
                path.clear(); tmp.clear();
                n=sz(adj);ti=0;
                time.resize(n);
                dfs(adj, root);
                rmq.build(tmp);
        void dfs(vector<vi>& adj, int u, int p=-1) {
                time[u]=ti++;
                for(int v:adj[u]){
                        if (v==p) continue;
                        path.push_back(u);
                        tmp.push_back(time[u]);
                         dfs(adj, v, u);
        int lca(int a, int b) { // check forest
                if (a==b) return a;
```

```
a=time[a],b=time[b];
if(a>b)swap(a,b);
return path[rmq.get(a,b-1)];
};
```

2.7 Sack

```
const int maxn = 1e5+5;
vi adj[maxn];
int ver[2*maxn]; // nodo en la posicion i del euler tour
int len[maxn]; // tamano del subarbol de v
int st[maxn]; // tiempo inicial de v
int ft[maxn]; // tiempo final de v
int pos=0;
// O(n*log(n))
// 1) dfs0(root);
// 2) dfs1(root);
void dfs0(int v=0, int p=-1){
        len[v]=1;
        ver[pos]=v;
        st[v]=pos++;
        for(int u:adj[v]){
                if (u==p) continue;
                dfs0(u,v);
                len[v] +=len[u];
        ver[pos]=v;
        ft[v]=pos++;
bool vis[maxn];
void ask(int v, bool add) {
        if(vis[v] && !add) {
                vis[v]=false;
                // eliminar nodo v
                // ...
        }else if(!vis[v] && add){
                vis[v]=true;
                // anadir nodo v
                // ...
void dfs1(int v=0, int p=-1, bool keep=true) {
        int mx=0,id=-1;
        for(int u:adj[v]){
                if (u==p) continue;
                if(len[u]>mx){
                        mx=len[u];
                        id=u;
```

2.8 Virtual Tree

```
// O(k*log(k))
// 1) build(n, root, adj);
// 2) query(nodes);
LCA q; // Add LCA
int lca(int a, int b) {return q.lca(a,b);};
struct VirtualTree{
        vector<vi> adj,adjVT;
        vector<int> st,ft;
        vector<bool> important;
        int pos=0;
        void build(vector<vi>& adj2, int n, int root) {
                important.assign(n, false);
                adjVT.assign(n,vi());
                st.resize(n);
                ft.resize(n);
                adj=adj2;pos=0;
                dfs(root);
        void dfs(int v, int p=-1){
                st[v]=pos++;
                for(int u:adj[v]){
                        if (u==p) continue;
                         dfs(u, v);
                ft[v]=pos++;
        bool upper (int v, int u) {return st[v] <= st[u] &&
            ft[v]>=ft[u];}
        int getRootVirtualTree(vi nodes) {
```

```
sort(all(nodes), [&](int v, int u) {
                    return st[v] < st[u]; });</pre>
                 int m=sz(nodes);
                 for(int i=0;i<m-1;++i){</pre>
                         int v=lca(nodes[i], nodes[i+1]);
                         nodes.push back(v);
                 sort(all(nodes), [&](int v, int u){
                    return st[v] < st[u]; });</pre>
                 nodes.erase(unique(all(nodes)), nodes.end
                 for(int u:nodes)adjVT[u].clear();
                 vi s:
                 s.push back(nodes[0]);
                 m=sz (nodes);
                 for (int i=1; i < m; ++i) {</pre>
                         int v=nodes[i];
                         while (sz(s) \ge 2 \&\& !upper(s.back())
                             , v)){
                                  adiVT[s[sz(s)-2]].
                                      push_back(s.back());
                                  s.pop_back();
                         s.push_back(v);
                 while (sz(s) >= 2) {
                         adjVT[s[sz(s)-2]].push_back(s.
                             back());
                         s.pop back();
                 return s[0];
        void dfs2(int u, int p=-1){
                 if(important[u]){
                          // pass
                 }else{
                          // pass
                 for(int v:adjVT[u]){
                         if (v==p) continue;
                         dfs2(v,u);
        void query(vi& nodes){
                 for(int u:nodes)important[u]=true;
                 int root=getRootVirtualTree(nodes);
                 dfs2(root);
                 // cout ans
                 for(int u:nodes)important[u]=false;
};
```

3 Estructuras de Datos

3.1 Bit

```
// O(n) build
// O(log(n)) get, upd
typedef long long T;
struct BIT{
        vector<T> t;
        int n;
        BIT(int _n) {
                 \overline{n} = n;
                 t.assign(n+1,0);
        void upd(int i, T v) { // add v to ith element
                 for(int j=i+1; j<=n; j+=j&-j)t[j]+=v;
        T get(int i) { // get sum of range [0,i0)
                 T ans=0;
                 for(int j=i; j; j-=j&-j) ans+=t[j];
                 return ans;
        T get(int 1, int r) { // get sum of range [1,r]
                 return get(r+1)-get(l);
};
```

3.2 Bit 2D

```
// O(n*m) build
// O(\log(n) * \log(m))  get, upd
typedef long long T;
struct BIT2D{
        vector<vector<T>> bit;
        int n,m;
        BIT2D(int _n, int _m) {
                 n=n; m=m;
                 bit.assign(n+1, vector<T>(m+1,0));
        T get(int x, int y) {
                 if(x<0 || v<0) return 0;
                 T v=0;
                 for(int i=x+1;i;i-=i&-i)
                         for(int j=y+1; j; j-=j&-j) v+=bit[i
                             ][j];
                 return v:
        T get(int x, int y, int x2, int y2) {
                 return get (x_2, y_2) -get (x_1, y_2) -get (x_2, y_1)
                    +qet(x-1,y-1);
        void upd(int x, int y, T dt){
```

3.3 Cartesian Tree

```
// O(n) build
typedef long long T;
struct CartesianTree{ // 1-indexed
        vector<int> 1,r;
        int root,n;
        CartesianTree(vector<T>& a) {
                 reverse(all(a));
                 a.push back(0);
                 reverse (all(a));
                 int tot=0; n=sz(a)-1;
                 1.assign(n+1,0);
                 r.assign(n+1,0);
                 vector<int> s(n+1,0);
                 vector<bool> vis(n+1, false);
                 for (int i=1; i<=n; ++i) {</pre>
                          int k=tot;
                          while(k>0 && a[s[k-1]]>a[i])k--;
                              // < max heap
                          if(k)r[s[k-1]]=i;
                          if(k<tot)l[i]=s[k];
                          s[k++]=i;
                          tot=k;
                 for (int i=1; i<=n; ++i) vis[l[i]]=vis[r[i</pre>
                    ] ] =1;
                 root=0;
                 for (int i=1; i<=n; ++i) {</pre>
                          if(!vis[i])root=i;
};
```

3.4 Disjoint Set Union

```
struct dsu{
   vi p,size;
   int sets,maxSize;

   dsu(int n) {
       p.assign(n,0);
       size.assign(n,1);
       sets = n;
```

3.5 Disjoint Sparse Table

```
// lo mismo que sparse table, pero para st opers
// O(n*log(n)) build
// O(1)  get
typedef int T;
T null = 0;
T op (T a, T b) {return a^b;}
struct DST {
        vector<vector<T>> pre, suf;
        int k, n;
        DST(vector<T>& a) {
                 n = sz(a);
                 k = log2(n) + 2;
                 pre.assign(k + 1, vector<T > (n));
                 suf.assign(k + 1, vector < T > (n));
                 for (int \dot{j} = 0; (1 << \dot{j}) <= n; ++\dot{j}) {
                          int mask = (1 << j) - 1;
                          T nw = null;
                          for (int i = 0; i < n; ++i) {
                                  nw = op(nw, a[i]);
                                  pre[j][i] = nw;
                                  if((i \& mask) == mask) nw
                                       = null:
                         nw = null;
                          for (int i = n - 1; i >= 0; --i) {
                                  nw = op(a[i], nw);
                                  suf[j][i] = nw;
                                  if((i \& mask) == 0) nw =
                                      null;
```

```
}
T get(int 1, int r) {
    if(1 == r) return pre[0][1];
    int i = 31 - __builtin_clz(l ^ r);
    return op(suf[i][1], pre[i][r]);
}
};
```

3.6 Dynamic Connectivity Offline

```
typedef pair<int, int> ii;
struct DSU {
        vector<int> p, size, h;
        int sets;
        void build(int n) {
                 sets=n;
                 p.assign(n,0);
                 size.assign(n,1);
                 for (int i=0; i < n; ++i) p[i] = i;</pre>
        int get(int a) {return (a==p[a]?a:get(p[a]));}
        void unite(int a, int b) {
                 a=get(a); b=get(b);
                 if (a==b) return;
                 if(size[a]>size[b])swap(a,b);
                 h.push_back(a);
                 size[b]+=size[a];
                 p[a]=b; sets--;
        void rollback(int s) {
                 while (sz(h)>s) {
                         int a=h.back();
                         h.pop_back();
                         size[p[a]]-=size[a];
                         p[a]=a; sets++;
};
// O(q*log(q)*log(n))
enum { ADD, DEL, QUERY };
struct Query { int type, u, v; };
struct DynCon {
        map<ii, int> edges; DSU uf;
        vector<Query> q;
        vector<int> t;
        void add(int u, int v) {
                 if(u>v) swap(u,v);
                 edges[\{u,v\}]=sz(q);
                 q.push back({ADD, u, v});
                 t.push back(-1);
        void del(int u, int v){
```

```
if(u>v) swap(u,v);
        int i=edges[{u,v}];
        t[i]=sz(q);
        q.push_back({DEL, u, v});
        t.push_back(i);
void querv() {
        q.push_back({QUERY, -1, -1});
        t.push_back(-1);
void dnc(int 1, int r) {
        if(r-l==1){
                 if (q[1].type==QUERY)
                         cout<<uf.sets<<"\n";</pre>
                 return;
        int m=1+(r-1)/2, k=sz(uf.h);
        for(int i=m; i<r; ++i)
                 if(q[i].type==DEL && t[i]<1)
                         uf.unite(q[i].u, q[i].v);
        dnc(1, m);
        uf.rollback(k);
        for(int i=1;i<m;++i)</pre>
                 if(q[i].type==ADD && t[i]>=r)
                         uf.unite(q[i].u, q[i].v);
        dnc(m, r);
        uf.rollback(k);
void init(int n){
        uf.build(n);
        if(!sz(q))return;
        for (int & ti:t) if (ti==-1) ti=sz(q);
        dnc(0, sz(q));
```

3.7 DSU Bipartite

};

```
// Bipartite graph
// get return the leader and the parity of the distance
   to the leader
typedef pair<int, int> ii;
struct DSU{
    vector<int> p, size, len;
    DSU(int n) {
        p.assign(n,0);
        len.assign(n,0);
        size.assign(n,1);
        for(int i=0;i<n;++i)p[i]=i;
    }
    ii get(int a) {
        if(a==p[a]) return {a, 0};
        ii va=get(p[a]);</pre>
```

```
p[a]=va.first;
len[a]=(len[a]+va.second)%2;
return {p[a], len[a]};
}
void unite(int a, int b) {
    ii va=get(a);
    ii vb=get(b);
    if(va.first==vb.first)return;
    if(size[va.first]>size[vb.first])swap(va, vb);
    p[va.first]=vb.first;
    len[va.first]=(va.second+vb.second+1)%2;
    size[vb.first]+=size[va.first];
}
};
```

3.8 Dynamic Segment Tree

```
// O(q*log(n)), q => queries
typedef long long T;
T null=0, noVal=0;
T oper(T a, T b) {return a+b;}
struct Node {
        T val, lz;
        int 1,r;
        Node *pl, *pr;
        Node(int ll, int rr) {
                 val=null; lz=noVal;
                 pl=pr=nullptr;
                 l=11; r=rr;
        void update() {
                 if (r-l==1) return;
                 val=oper(pl->val, pr->val);
        void update(T v) {
                 val += ((T)(r-1)) *v;
                 1z+=v;
        void extends() {
                 if(r-l!=1 && !pl) {
                         int m = (r+1)/2;
                         pl=new Node(1, m);
                         pr=new Node(m, r);
        void propagate() {
                 if (r-l==1) return;
                 if (lz==noVal) return;
                 pl->update(lz);
                 pr->update(lz);
                 lz=noVal;
};
```

```
typedef Node* PNode;
struct SegTree{
        PNode root;
        SeqTree(int 1, int r) {root=new Node(1, r+1);}
        void upd(PNode x, int 1, int r, T v){
                 int 1x=x->1, rx=x->r;
                 if (lx>=r | | l>=rx) return;
                if(lx>=1 && rx<=r){
                         x->update(v);
                         return;
                x->extends():
                x->propagate();
                upd(x->pl,l,r,v);
                upd(x->pr, l, r, v);
                x->update();
        T get(PNode x, int 1, int r) {
                 int lx=x->1, rx=x->r;
                if(lx>=r || l>=rx)return null;
                if(lx>=1 && rx<=r) return x->val;
                x->extends();
                x->propagate();
                T v1=qet(x->pl,l,r);
                T v2=qet(x->pr,l,r);
                 return oper (v1, v2);
        T get(int 1, int r) {return get(root, 1, r+1);}
        void upd(int 1, int r, T v) {upd(root, 1, r+1, v);}
};
```

3.9 Implicit Treap

```
// Treap => Binary Search Tree + Binary Heap
// 1. create a empty root (PTreap root=nullptr;)
// 2. Append the nodes in order (left -> right)
// PTreap tmp=new Treap(x);
// root=merge(root, tmp);
typedef long long T;
typedef unsigned long long u64;
mt19937_64 rng (chrono::steady_clock::now().
   time since_epoch().count());
T \text{ null} = 0;
struct Treap{
        Treap *1, *r; // left child, right child
        u64 prior; // random
        T val, sum, lz; // value, sum subtree, lazy
        int sz; // size subtree
        Treap(T v) {
```

```
l=r=nullptr;
                 prior=rnq();
                 val=sum=v;
                 1z=0; sz=1;
         Treap(){
                 delete 1;
                 delete r;
};
typedef Treap* PTreap;
int cnt(PTreap x) {return (!x?0:x->sz);}
T sum(PTreap x) {return (!x?0:x->sum);}
void update_helper(PTreap x, T v) {
        // 1z + v
        // val += v
        // sum += v
// propagate the lazy
void push(PTreap x) {
        if (x \& \& x \rightarrow lz) \{ // check x \rightarrow lz \}
                 if (x->1) update helper (x->1, 1);
                 if (x->r) update_helper (x->r, 1);
                 x -> 1z = 0;
// updates node with its children information
void pull(PTreap x) {
        push (x->1);
        push (x->r);
        x->sz=cnt(x->1)+cnt(x->r)+1;
        x->sum=sum(x->1)+sum(x->r)+x->val;
// Updates node value += v
void upd(PTreap x, T v) {
        if(!x)return;
        pull(x);
        update_helper(x, v);
// O(log(n)) divide the treap in two parts
// [count nodes == left], [the rest of nodes]
pair<PTreap, PTreap> split(PTreap x, int left) {
        if(!x)return {nullptr, nullptr};
        push(x);
        if(cnt(x->1)>=left)
                 auto got=split(x->1, left);
                 x->l=qot.second;
                 pull(x);
                 return {got.first, x};
        }else{
                 auto got=split(x->r, left-cnt(x->1)-1);
```

```
x->r=qot.first;
                 pull(x);
                 return {x, got.second};
// O(log(n)) merge two treap
// [nodes treap x ... nodes treap y]
PTreap merge (PTreap x, PTreap y) {
        if(!x)return y;
        if(!y)return x;
        push(x); push(y);
        if (x->prior<=y->prior) {
                 x \rightarrow r = merge(x \rightarrow r, y);
                  pull(x):
                  return x;
        }else{
                  y->l=merge(x, y->l);
                 pull(y);
                  return v;
// O(n) print the treap
void dfs(PTreap x) {
        if(!x)return;
        push(x);
        dfs(x->1);
        cout << x -> val << " ";
        dfs(x->r);
```

3.10 Implicit Treap Father

```
// Treap => Binary Search Tree + Binary Heap
// 1. create a empty root (PTreap root=nullptr;)
// 2. Append the nodes in order (left -> right)
// PTreap tmp=new Treap(x);
// root=merge(root, tmp);
// si se edita un treap, se tiene que hacer un pullAll
   hasta la raiz
// si no se hace esto, el treap queda con informacion
   pasada
// si se va a modificar un treap, hacer un pushAll para
   bajar los lazy
typedef long long T;
typedef unsigned long long u64;
mt19937_64 rng (chrono::steady_clock::now().
   time since epoch().count());
T null = 0:
struct Treap{
        Treap *1,*r,*dad; // left child, right child
```

```
u64 prior; // random
        T val, sum; // value, sum subtree
        int sz; // size subtree
        Treap(T v) {
                l=r=dad=nullptr;
                prior=rng();
                val=sum=v;
                sz=1;
        <sup>2</sup>Treap(){
                delete 1:
                delete r:
};
typedef Treap* PTreap;
int cnt(PTreap x) {return (!x?0:x->sz);}
T sum(PTreap x) {return (!x?0:x->sum);}
// updates node with its children information
void pull(PTreap x) {
        x->sz=cnt(x->1)+cnt(x->r)+1;
        x->sum=sum(x->1)+sum(x->r)+x->val;
        if (x->1) x->1->dad=x; //
        if (x->r) x->r->dad=x; //
// O(log(n)) divide the treap in two parts
// [count nodes == left], [the rest of nodes]
pair<PTreap, PTreap> split(PTreap x, int left) {
        if(!x)return {nullptr, nullptr};
        if(cnt(x->1)>=left){
                auto got=split(x->1, left);
                if (got.first) got.first->dad=nullptr; //
                x->l=qot.second;
                x->dad=nullptr; //
                pull(x);
                return {got.first, x};
        }else{
                auto got=split(x->r, left-cnt(x->l)-1);
                if (got.second) got.second->dad=nullptr; //
                x->r=qot.first;
                x->dad=nullptr; //
                pull(x);
                return {x, got.second};
// O(log(n)) merge two treap
// [nodes treap x ... nodes treap y]
PTreap merge (PTreap x, PTreap y) {
        if(!x)return y;
        if(!y)return x;
        if(x->prior<=y->prior) {
                x->r=merge(x->r, y);
                pull(x);
                return x;
```

```
}else{
                y->l=merge(x, y->l);
                pull(y);
                return y;
// O(log(n)) propagate the lazy [root->x]
void pushAll(PTreap x) {
        if(!x)return;
        pushAll(x->dad);
        push(x);
// O(log(n)) update the treap [root->x]
void pullAll(PTreap x) {
        if(!x)return;
        pull(x);
        pullAll(x->dad);
// O(log(n)) return the root and the position of x (1-
   indexed)
pair<PTreap, int> findRoot(PTreap x){
        pushAll(x);
        int pos=cnt(x->1);
        while(x->dad) {
                PTreap f=x->dad;
                if (x==f->r) pos+=cnt (f->1)+1;
                x=f:
        return {x,pos+1};
```

3.11 Li Chao

```
// inf max abs value that the function may take
typedef long long ty;
struct Line {
         ty m, b;
        Line(){}
        Line(ty m, ty b): m(m), b(b){}
        ty eval(ty x){return m * x + b;}
};
struct nLiChao{
        // see coments for min
        nLiChao *left = nullptr, *right = nullptr;
        ty l, r;
        Line line;
        nLiChao(ty l, ty r): l(l), r(r){
              line = {0, -inf}; // change to {0, inf};
```

```
// T(Log(Rango)) M(Log(rango))
        void addLine(Line nline) {
                ty m = (1 + r) >> 1;
                bool lef = nline.eval(1) > line.eval(1);
                    // change > to <
                bool mid = nline.eval(m) > line.eval(m);
                    // change > to <
                if (mid) swap(nline, line);
                if (r == 1) return;
                if (lef != mid) {
                         if (!left){
                                 left = new nLiChao(l, m);
                                 left -> line = nline;
                         else left -> addLine(nline);
                else{
                         if (!right) {
                                 right = new nLiChao(m +
                                    1, r);
                                 right -> line = nline;
                         else right -> addLine(nline);
        // T(Log(Rango))
        ty get(ty x) {
                \bar{t}v m = (l + r) >> 1;
                ty op1 = -\inf, op2 = -\inf; // change to
                    inf
                if(l == r) return line.eval(x);
                else if (x < m) {
                         if (left) op1 = left -> get(x);
                         return max(line.eval(x), op1); //
                             change max to min
                else{
                         if (right) op2 = right \rightarrow get(x);
                         return max(line.eval(x), op2); //
                             change max to min
};
int main() {
        // (rango superior) * (pendiente maxima) puede
            desbordarse
        // usar double o long double en el eval para
```

```
estos casos
// (puede dar problemas de precision)
nLiChao liChao(0, 1e18);
```

3.12 Link Cut Tree

```
// 1-indexed
// All operations are O(log(n))
typedef long long T;
struct SplayTree{
        struct Node{
                 int ch[2] = \{0, 0\}, p=0;
                T val=0, path=0; // values for path
                 T sub=0, vir=0; // values for subtree
                bool flip=0; // values for lazy
        vector<Node> ns;
        SplavTree(int n):ns(n+1){}
        T path(int u) {return (u?ns[u].path:0);}
        T subsum(int u) {return (u?ns[u].sub:0);}
        void push(int x) {
                 if(!x)return;
                int l=ns[x].ch[0], r=ns[x].ch[1];
                if(ns[x].flip){
                         ns[l].flip^=1,ns[r].flip^=1;
                         swap(ns[x].ch[0], ns[x].ch[1]);
                         // if the operation is like a
                            segment tree
                         // check swap the values
                         ns[x].flip=0;
        void pull(int x) {
                 int l=ns[x].ch[0], r=ns[x].ch[1];
                push(1); push(r);
                ns[x].path=max({path(1), path(r), ns[x].}
                 ns[x].sub=ns[x].vir+subsum(1)+subsum(r)+
                    ns[x].val;
        void set(int x, int d, int y) {ns[x].ch[d]=y;ns[y
            ].p=x;pull(x);
        void splay(int x) {
                 auto dir=[&](int x){
                         int p=ns[x].p;if(!p)return -1;
                         return ns[p].ch[0] == x?0:ns[p].ch
                            [1] == x?1:-1;
                 auto rotate=[&](int x){
                         int y=ns[x].p, z=ns[y].p, dx=dir(x)
                            , dy = dir(y);
                         set (y, dx, ns[x].ch[!dx]);
                         set (x, !dx, y);
```

```
if(^{\circ}dy) set (z, dy, x);
                         ns[x].p=z;
                 };
                 for(push(x); ~dir(x);) {
                         int y=ns[x].p, z=ns[y].p;
                         push(z);push(y);push(x);
                         int dx=dir(x), dy=dir(y);
                         if(^{\circ}dy) rotate (dx!=dy?x:y);
                         rotate(x);
} ;
struct LinkCut:SplayTree{
        LinkCut(int n):SplayTree(n){}
        // return the root of us tree
        int root(int u){
                access(u); splay(u); push(u);
                 while (ns[u].ch[0]) {u=ns[u].ch[0]; push(u)
                return splay(u),u;
        // return the parent of u
        int parent(int u) {
                access(u); splay(u); push(u);
                u=ns[u].ch[0];push(u);
                while (ns[u].ch[1]) {u=ns[u].ch[1]; push(u)
                return splay(u),u;
        int access(int x){
                int u=x, v=0;
                 for(;u;v=u,u=ns[u].p){
                         splay(u);
                         int& ov=ns[u].ch[1];
                         ns[u].vir+=ns[ov].sub;
                         ns[u].vir-=ns[v].sub;
                         ov=v; pull(u);
                return splay(x), v;
        // reroot the tree with x as root
        void reroot(int x){
                access(x); ns[x].flip^=1; push(x);
        // create a edge u->v, u is the child of v
        void link(int u, int v){
                reroot (u); access (v);
                ns[v].vir+=ns[u].sub;
                ns[u].p=v;pull(v);
        // delete the edge u->v, u is the child of v
        void cut(int u, int v){
                int r=root(u);
                reroot(u); access(v);
```

};

```
ns[v].ch[0]=ns[u].p=0;pull(v);
        reroot(r);
// delete the edge u->parent(u)
void cut(int u){
        access(u);
        ns[ns[u].ch[0]].p=0;
        ns[u].ch[0]=0;pull(u);
int lca(int u, int v) {
        if (root (u) !=root (v)) return -1;
        access(u); return access(v);
// return sum of the subtree of u with v as
   father
T subtree(int u, int v) {
        int r=root(u);
        reroot (v); access (u);
        T ans=ns[u].vir+ns[u].val;
        return reroot(r), ans;
T path(int u, int v) {
        int r=root(u);
        reroot (u); access (v); pull (v);
        T ans=ns[v].path;
        return reroot (r), ans;
void set(int u, T val){
        access(u);
        ns[u].val=val;
        pull(u);
```

3.13 Link Cut Tree Lazy

```
T subsum(int u) {return (u?ns[u].sub:0);}
void push(int x){
        if(!x)return;
        int l=ns[x].ch[0], r=ns[x].ch[1];
        if(ns[x].flip){
                 ns[1].flip^=1,ns[r].flip^=1;
                 swap (ns[x].ch[0], ns[x].ch[1]);
                 // if the operation is like a
                    segment tree
                 // check swap the values
                 ns[x].flip=0;
        if (ns[x].lz) { // check the lazy
                 // propagate the lazy
                 ns[x].sub+=ns[x].lz*ns[x].ssz;
                 ns[x].vir+=ns[x].lz*ns[x].vsz;
void pull(int x) {
        int l=ns[x].ch[0],r=ns[x].ch[1];
        push(1);push(r);
        ns[x].sz=size(1)+size(r)+1;
        ns[x].path=max({path(1), path(r), ns[x].}
        ns[x].sub=ns[x].vir+subsum(1)+subsum(r)+
            ns[x].val;
        ns[x].ssz=ns[x].vsz+subsize(l)+subsize(r)
            +1;
void set(int x, int d, int y) {ns[x].ch[d]=y;ns[y
    ].p=x;pull(x);}
void splay(int x) {
        auto dir=[&](int x){
                 int p=ns[x].p;if(!p)return -1;
                 return ns[p].ch[0] == x?0:ns[p].ch
                    [1] == x?1:-1;
        };
        auto rotate=[&](int x){
                 int y=ns[x].p, z=ns[y].p, dx=dir(x)
                    \overline{dy} = dir(y);
                 set (y, dx, ns[x].ch[!dx]);
                 set (x, !dx, y);
                 if(^{\circ}dy) set (z, dy, x);
                 ns[x].p=z;
        for (push (x); ~dir(x);) {
                 int y=ns[x].p, z=ns[y].p;
                 push(z);push(y);push(x);
                 int dx=dir(x), dy=dir(y);
                 if(~dy) rotate(dx!=dy?x:y);
                 rotate(x);
```

};

```
struct LinkCut:SplayTree{
        LinkCut(int n):SplayTree(n){}
        // return the root of us tree
        int root(int u){
                access(u); splay(u); push(u);
                while (ns[u].ch[0]) {u=ns[u].ch[0]; push(u)
                return splay(u),u;
        // return the parent of u
        int parent(int u){
                access(u); splay(u); push(u);
                u=ns[u].ch[0];push(u);
                while (ns[u].ch[1]) {u=ns[u].ch[1]; push(u)
                return splay(u),u;
        int access(int x){
                int u=x, v=0;
                for(;u;v=u,u=ns[u].p){
                        splay(u);
                        int& ov=ns[u].ch[1];
                        ns[u].vir+=ns[ov].sub;
                        ns[u].vsz+=ns[ov].ssz;
                        ns[u].vir-=ns[v].sub;
                        ns[u].vsz-=ns[v].ssz;
                        ov=v; pull(u);
                return splay(x), v;
        // reroot the tree with x as root
        void reroot(int x){
                access(x); ns[x].flip^=1; push(x);
        // create a edge u->v, u is the child of v
        void link(int u, int v){
                reroot (u);
                access(v);
                ns[v].vir+=ns[u].sub;
                ns[v].vsz+=ns[u].ssz;
                ns[u].p=v;pull(v);
        // delete the edge u->v, u is the child of v
        void cut(int u, int v){
                int r=root(u);
                reroot (u);
                access(v);
                ns[v].ch[0]=ns[u].p=0;pull(v);
                reroot(r);
```

```
// delete the edge u->parent(u)
void cut(int u){
        access(u);
        ns[ns[u].ch[0]].p=0;
        ns[u].ch[0]=0;pull(u);
int lca(int u, int v) {
        if (root (u) !=root (v)) return -1;
        access(u); return access(v);
int depth(int u){
        int r=root(u);
        reroot(r);
        access(u); splay(u); push(u);
        return ns[u].sz-1;
T path(int u, int v) {
        int r=root(u);
        reroot(u); access(v); pull(v);
        T ans=ns[v].path;
        return reroot (r), ans;
void set(int u, T val){
        access(u);
        ns[u].val=val;
        pull(u);
// update the value of the nodes in the path u->v
    with += val
void upd(int u, int v, T val){
        int r=root(u);
        reroot (u); access (v); splay (v);
        // change only the lazy
        // ns[v].val+=val;
        reroot(r);
T comp size(int u) {return ns[root(u)].ssz;}
T subtree size(int u) {
        int p=parent(u);
        if(!p)return comp_size(u);
        cut(u); int ans=comp size(u);
        link(u,p); return ans;
T subtree size(int u, int v){ // subtree of u
   with v as father
        int r=root(u);
        reroot (v); access (u);
        T ans=ns[u].vsz+1;
        return reroot(r), ans;
```

```
T comp sum(int u) {return ns[root(u)].sub;}
        T subtree_sum(int u) {
                 int p=parent(u);
                if(!p)return comp_sum(u);
                 cut(u); T ans=comp_sum(u);
                link(u,p); return ans;
        T subtree_sum(int u, int v) { // subtree of u with
             v as father
                 int r=root(u);
                 reroot (v); access (u);
                T ans=ns[u].vir+ns[u].val; // por el
                    reroot
                 return reroot(r), ans;
};
```

3.14 Merge Sort Tree

```
// O(n*log(n)) build
// O(log(n)^2) get
typedef long long T;
struct SegTree{
        int size;
        vector<vector<T>> vals;
        void oper(int x) {
                merge(all(vals[2*x+1]), all(vals[2*x+2]),
                     back inserter(vals[x]));
        SegTree(vector<T>& a) {
                 size=1;
                while (size<sz(a))size*=2;</pre>
                vals.resize(2*size);
                 build(a, 0, 0, size);
        void build(vector<T>& a, int x, int lx, int rx) {
                 if(rx-lx==1) {
                         if(lx<sz(a))vals[x]={a[lx]};
                         return;
                 int m = (1x+rx)/2;
                build(a, 2*x+1, 1x, m);
                 build(a, 2*x+2, m, rx);
                 oper(x);
        int get (int 1, int r, int val, int x, int lx, int
           rx) {
                 if(lx>=r | | l>=rx) return 0;
                 if(lx>=1 && rx<=r){
                         return upper_bound(all(vals[x]),
                            val) -vals[x].begin();
```

```
int m = (1x+rx)/2;
                 int v1=get(l,r,val,2*x+1,lx,m);
                int v2=get(1,r,val,2*x+2,m,rx);
                return v1+v2;
        int get(int 1, int r,int val) {return get(1,r+1,
            val, 0, 0, size);}
};
```

3.15 MOs Algorithm

```
// O((n+q)*sq), sq=n^{(1/2)}
// 1. fill queries[]
// 2. solve(n);
// 3. print ans[]
int sq;
struct query {int l,r,idx;};
bool cmp(query& a, query& b) {
        int x=a.1/sq;
        if (a.1/sq!=b.1/sq) return a.1/sq<b.1/sq;</pre>
        return (x&1?a.r<b.r:a.r>b.r);
vector<query> queries;
vector<11> ans;
ll act();
void add(int i); // add a[i]
void remove(int i); // remove a[i]
void solve(int n) {
        sq=ceil(sqrt(n));
        sort(all(queries), cmp);
        ans.assign(sz(queries),0);
        int 1=0, r=-1;
        for(auto [li,ri,i]:queries) {
                 while (r<ri) add (++r);</pre>
                 while (1>1i) add (--1);
                 while (r>ri) remove (r--);
                 while (1<li) remove (1++);</pre>
                 ans[i]=act();
```

3.16 MOs Tree

```
// add LCA
struct LCA{};
vector<vector<int>> adj;
const int maxn=1e5+5;
```

```
int ver[2*maxn]; // node at position i in euler tour
int st[maxn]; // start time of v
int ft[maxn]; // finish time of v
int pos=0;
LCA tree;
// O((n+q)*sq), sq=n^{(1/2)}
// 1. build euler tour and lca
// 2. add queries[]
// if (st[a]>st[b]) swap (a, b);
// queries.push_back({st[a]+1,st[b],i});
// 3. solve(n);
// 4. print ans[]
int sq;
void dfs(int u=0, int p=-1) {
        ver[pos]=u;
        st[u]=pos++;
        for(int v:adj[u]){
                 if (v==p) continue;
                 dfs(v,u);
        ver[pos]=u;
        ft[u]=pos++;
struct query {int l,r,idx;};
bool cmp(query& a, query& b) {
        int x=a.l/sq;
        if (a.l/sq!=b.l/sq) return a.l/sq<b.l/sq;</pre>
        return (x&1?a.r<b.r:a.r>b.r);
vector<query> queries;
vector<ll> ans;
bool vis[maxn];
11 act();
void add(int u); // add node u
void remove(int u); // remove node u
void ask(int u) {
        if(!vis[u])add(u);
        else remove(u);
        vis[u]=!vis[u];
void solve(int n) {
         sa=ceil(sart(n));
        sort(all(queries), cmp);
    ans.resize(sz(queries));
         int 1=0, r=-1;
    for(auto [li,ri,i]:queries) {
                 while(r<ri)ask(ver[++r]);</pre>
                 while(1>li) ask (ver[--1]);
                 while(r>ri)ask(ver[r--]);
                 while(l<li) ask(ver[l++]);</pre>
                 int a=ver[l-1],b=ver[r];
```

```
int c=tree.lca(a,b);
ask(c);
ans[i]=act();
ask(c);
```

3.17 MOs Updates

```
// O(q*(s+(n/s)^2)) \Rightarrow O(q*(n^2(2/3))), s=(2*(n^2))^2(1/3) -
    s=n^{(2/3)}
// 1. fill queries[] and upds[]
// dont confuse index in queries with updates, they are
// the struct upd saves the old value and the new value
// 2. solve(n);
// 3. print ans[]
int sq;
struct upd{int i,old,cur;};
struct query {int l,r,t,idx;};
bool cmp(query& a, query& b) {
        int x=a.l/sq;
        if (a.l/sq!=b.l/sq) return a.l/sq<b.l/sq;</pre>
        if (a.r/sq!=b.r/sq) return (x&1?a.r<b.r:a.r>b.r);
        return a.t<b.t;</pre>
vector<query> queries;
vector<upd> upds;
vector<11> ans:
ll act();
void add(int i); // add a[i]
void remove(int i); // remove a[i]
void update(int i, int v, int l, int r){
        // check if the update is with an active element
        if(l<=i && i<=r){
                remove(i);
                 // a[i]=v;
                // ...
                 add(i);
        // a[i]=v;
        // ...
void solve(int n) {
        sq=ceil(pow(n, 2.0/3.0));
        sort(all(queries), cmp);
    ans.resize(sz(queries));
        int l=0, r=-1, t=0;
        for(auto [li,ri,ti,i]:queries){
                 while(t<ti) update(upds[t].i, upds[t].cur, l</pre>
                    r),++t;
```

3.18 Ordered set

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template<typename T> using ordered set = tree<T,
   null_type,less<T>, rb_tree_tag,
   tree order statistics node update>;
template<typename T> using ordered_multiset = tree<T,</pre>
   null_type,less_equal<T>, rb_tree_tag,
   tree order statistics node update>;
// ----- CONSTRUCTOR ----- //
// 1. Para ordenar por MAX cambiar less<int> por greater<
   int>
// 2. Para multiset cambiar less<int> por less_equal<int>
       Para borrar siendo multiset:
       int idx = st.order of key(value);
       st.erase(st.find_by_order(idx));
       st.swap(st2);
// ----- METHODS ----- //
st.find_by_order(k) // returns pointer to the k-th
   smallest element
st.order_of_key(x) // returns how many elements are
   smaller than x
st.find by order(k) == st.end() // true, if element does
   not exist
```

3.19 Persistent Segment Tree

```
// O(n*log(n)) build
// O(log(n)) get, set
// O((n+q)*log(n)) memory
typedef long long T;
struct Node{
        T val;
        int l,r; // saves the range of the node [l,r]
};
struct SegTree{
        vector<Node> ns;
        vector<int> roots; // roots of the differents
        versions
        T null=0;
```

```
int act=0,size; // act: number of nodes
T oper(T a, T b) {return a+b;}
SegTree(vector<T>& a, int n) {
        size=n;
        roots.push_back(build(a, 0, size));
void update(int x) {
        ns[x].val=oper(ns[ns[x].l].val, ns[ns[x].
            rl.val);
int newNode(T x){
        Node tmp=\{x, -1, -1\};
        ns.push back(tmp);
        return act++;
int newNode(int 1, int r){
        Node tmp={null,1,r};
        ns.push_back(tmp);
        update(act);
        return act++;
int build(vector<T>& a, int l, int r) {
        if(r-l==1) {return newNode(a[1]);}
        int m = (1+r)/2;
        return newNode (build (a, l, m), build (a, m,
             r));
int set(int x, int i, T v, int l, int r){
        if (r-l==1) return newNode(v);
        int m = (1+r)/2;
        if (i < m) return newNode (set (ns[x].l, i, v,</pre>
            1, m), ns[x].r);
        else return newNode(ns[x].l, set(ns[x].r,
             i, v, m, r));
T get(int x, int lx, int rx, int l, int r){
        if(lx>=r || l>=rx) return null;
        if(lx>=1 && rx<=r)return ns[x].val;</pre>
        int m = (1x+rx)/2;
        T v1=qet(ns[x].l, lx, m, l, r);
        T v2 = qet(ns[x].r, m, rx, l, r);
        return oper (v1, v2);
T get(int 1, int r, int time) {return get(roots[
   time], 0, size, 1, r+1);}
void set(int i, T v, int time) {roots.push_back(
   set(roots[time], i, v, 0, size));}
```

};

3.20 Persistent Segment Tree Lazy

```
// O(n*log(n)) build
// O(\log(n)) get, upd
// O((n+q)*log(n)) memory
typedef long long T;
struct Node {
        Node* left = nullptr;
        Node* right = nullptr;
        T val = 0, prop = 0;
typedef Node* PNode;
struct PerSegTree {
        vector<PNode> roots{};
        vector<T> vec{};
        int n = 0;
        T op (T a, T b) {
                return a+b;
        PNode newKid(PNode& curr) {
                 PNode newNode = new Node();
                newNode->left = curr->left;
                newNode->right = curr->right;
                newNode->prop = curr->prop;
                newNode->val = curr->val;
                return newNode;
        void lazy(int i, int j, PNode& curr) {
                if (!curr->prop) return;
                 curr \rightarrow val += ((T)(j - i + 1)) * curr \rightarrow
                if (i != j) {
                         curr->left = newKid(curr->left);
                         curr->right = newKid(curr->right)
                         curr->left->prop += curr->prop;
                         curr->right->prop += curr->prop;
                 curr->prop = 0;
        PNode build(int i, int j) {
                 PNode newNode = new Node();
                if (i == j) {
                         newNode->val = vec[i];
                 } else {
                         int mid = i + (j - i) / 2;
                         PNode leftt = build(i, mid);
                         PNode right = build(mid + 1, \dot{j});
                         newNode->val = op(leftt->val,
                            right->val);
                         newNode->left = leftt;
                         newNode->right = right;
```

```
return newNode;
PNode upd(int i, int j, int l, int r, T value,
   PNode& curr) {
        lazy(i, j, curr);
        if (i >= 1 && i <= r) {
                PNode newNode = newKid(curr);
                newNode->prop += value;
                lazy(i, j, newNode);
                return newNode:
        if (i > r || i < l) {
                return curr;
        PNode newNode = new Node();
        int mid = i + (j - i) / 2;
        newNode->left = upd(i, mid, l, r, value,
           curr->left);
        newNode - > right = upd (mid + 1, j, l, r,
           value, curr->right);
        newNode->val = op(newNode->left->val,
           newNode->right->val);
        return newNode;
T get(int i, int j, int l, int r, PNode& curr) {
        lazy(i, j, curr);
        if (j < l || r < i) {
                return 0;
        if (i >= 1 && i <= r) {
                return curr->val;
        int mid = i + (j - i) / 2;
        return op (get (i, mid, l, r, curr->left),
           get (mid + 1, j, l, r, curr->right));
// public methods
void build(vector<T>& vec) {
        if (vec.emptv()) return;
        n = vec.size();
        this->vec = vec;
        auto root = build(0, n - 1);
        roots.push back(root);
void upd(int 1, int r, T value, int time) {
       roots.push_back(upd(0, n - 1, 1, r, value
           , roots[time]));
T get(int 1, int r, int time) {
        return get (0, n - 1, 1, r, roots[time]);
```

```
int size() { return roots.size(); }
};
```

3.21 Polynomial Updates

```
11 gauss(ll x) {return (x*(x+111))/211;}
struct Node{
        11 sum=0; // the nodes value
        11 acum=0; // count completed levels
        ll cnt=0; // count of updates +1, +2, +3, ...
        void build(ll v) {
                acum=cnt=0;
                sum=v;
        void oper(Node& a, Node& b) {
                sum=a.sum+b.sum;
                acum=cnt=0;
        void lazy(ll len, ll acum, ll cnt){
                sum+= acum*len+gauss(len)* cnt;
                acum+=_acum;
                cnt+=_cnt;
};
struct SegTree{
        vector<Node> vals;
        Node null;
        int size;
        SegTree(vector<ll>& a) {
                size=1;
                while (size<sz(a))size*=2;</pre>
                vals.resize(2*size);
                build(a, 0, 0, size);
        void build(vector<ll>& a, int x, int lx, int rx){
                if(rx-1x==1){
                         if (lx<sz(a)) vals[x].build(a[lx]);</pre>
                         return;
                int m = (1x+rx)/2;
                build(a, 2*x+1, 1x, m);
                build(a, 2*x+2, m, rx);
                vals[x].oper(vals[2*x+1], vals[2*x+2]);
        void propagate(int x, int lx, int rx){
                if (rx-lx==1) return;
                if (vals[x].cnt==0) return;
                int m = (rx+lx)/2;
                vals[2*x+1].lazy(m-lx, vals[x].acum, vals
                    [x].cnt);
                vals[2*x+2].lazy(rx-m, vals[x].acum+ll(m-
                    lx) *vals[x].cnt, vals[x].cnt);
```

```
vals[x].acum=vals[x].cnt=0;
        void upd(int 1, int r, 11 v, int x, int 1x, int
            rx) {
                 if (rx<=l || r<=lx) return;</pre>
                 if(1<=1x && rx<=r){
                          vals[x].lazv(rx-lx,v*(lx-l),v);
                          return;
                 propagate(x,lx,rx);
                 int m = (1x+rx)/2;
                 upd(1, r, v, 2 \times x + 1, 1x, m);
                 upd(1, r, v, 2 \times x + 2, m, rx);
                 vals[x].oper(vals[2*x+1], vals[2*x+2]);
        ll get(int l, int r, int x, int lx, int rx){
                 if(rx<=l || r<=lx)return null.sum;</pre>
                 if(l<=lx && rx<=r)return vals[x].sum;</pre>
                 propagate(x,lx,rx);
                 int m = (lx + rx)/2;
                 ll v1=get(l,r,2*x+1,lx,m);
                 11 v2=qet(1,r,2*x+2,m,rx);
                 return v1+v2;
        11 get(int 1, int r) {return get(1,r+1,0,0,size);}
        void upd(int 1, int r, 11 v) {upd(1,r+1,v,0,0,size
            );}
        // v es la cantidad de veces que se aplica la
            operacion +1, +2, +3
} ;
```

3.22 Segment Tree Iterativo

```
struct segtree{
        int n; vl v; ll nulo = 0;
        11 op(ll a, ll b) {return a + b;}
         segtree(int n) : n(n) \{v = vl(2*n, nulo);\}
         segtree (vl &a): n(sz(a)), v(2*n) {
                 for(int i = 0; i<n; i++) v[n + i] = a[i];</pre>
                 for (int i = n-1; i > = 1; --i) v[i] = op(v[
                     i<<1], v[i<<1|1]);
        void upd(int k, ll nv) {
                 for (v[k += n] = nv; k > 1; k >>= 1) v[k
                     >>1] = op(v[k], v[k^1]);
        ll get(int l, int r){
                 11 \text{ vl} = \text{nulo}, \text{ vr} = \text{nulo};
```

```
for (1 += n, r += n+1; 1 < r; 1 >>= 1, r
>>= 1) {
        if (1&1) v1 = op(v1, v[1++]);
        if (r&1) vr = op(v[--r], vr);
}
return op(v1, vr);
}
};
```

3.23 Segment Tree Recursivo

```
typedef long long T;
struct SeqTree{
        vector<T> vals, lazy;
        T null=0, nolz=0;
        int size;
        T op (T a, T b) {return a+b;}
        SegTree(vector<T>& a) {
                 size=1:
                 while (size<sz(a)) size*=2;</pre>
                 vals.resize(2*size);
                 lazy.assign(2*size, nolz);
                 build(a, 0, 0, size);
        void build(vector<T>& a, int x, int lx, int rx) {
                 if(rx-lx==1) {
                         if(lx<sz(a))vals[x]=a[lx];
                         return;
                 int m = (1x+rx)/2;
                 build(a, 2*x+1, 1x, m);
                 build(a, 2*x+2, m, rx);
                 vals[x]=op(vals[2*x+1], vals[2*x+2]);
        void propagate(int x, int lx, int rx){
                 if (rx-lx==1) return;
                 if (lazy[x] == nolz) return;
                 int m = (1x+rx)/2;
                 lazy[2*x+1] += lazy[x];
                 vals[2*x+1]+=lazy[x]*((T)(m-lx));
                 lazy[2*x+2]+=lazy[x];
                 vals[2*x+2]+=lazy[x]*((T)(rx-m));
                 lazy[x]=nolz;
        void upd(int 1, int r, T v, int x, int lx, int rx)
                 if (rx<=l | | r<=lx) return;</pre>
                 if(1<=1x && rx<=r){
                         lazv[x]+=v;
                         vals[x] += v*((T)(rx-lx));
                         return;
```

```
propagate(x,lx,rx);
                 int m = (1x+rx)/2;
                 upd(1, r, v, 2 \times x + 1, 1x, m);
                 upd(1, r, v, 2 \times x + 2, m, rx);
                 vals[x] = op(vals[2*x+1], vals[2*x+2]);
        void set(int i, T v, int x, int lx, int rx){
                 if(rx-lx==1){
                          vals[x]=v;
                          return;
                 propagate(x,lx,rx);
                 int m = (1x+rx)/2;
                 if(i<m) set(i, v, 2*x+1, lx, m);
                 else set(i, v, 2*x+2, m, rx);
                 vals[x]=op(vals[2*x+1], vals[2*x+2]);
        T get(int 1, int r, int x, int lx, int rx) {
                 if(rx<=l || r<=lx)return null;</pre>
                 if(l<=lx && rx<=r)return vals[x];</pre>
                 propagate(x,lx,rx);
                 int m = (1x+rx)/2;
                 T v1=qet(1,r,2*x+1,1x,m);
                 T v2=qet(1,r,2*x+2,m,rx);
                 return op (v1, v2);
        T get(int 1, int r) {return get(1,r+1,0,0,size);}
        void upd(int 1, int r, T v) {upd(1,r+1,v,0,0,size)
        void set(int i, T val) {set(i,val,0,0,size);}
};
```

3.24 Segment Tree 2D

};

```
for (int i=0; i< n; ++i) for (int j=0; j< m; ++j)
             st[i+n][j+m]=a[i][j];
         for (int i=0; i< n; ++i) for (int j=m-1; j>=1; --
             j) st[i+n][j] = op(st[i+n][j << 1], st[i+n]
             ][j<<1|1]);
         for (int i=n-1; i>=1; --i) for (int j=0; j<2*m
             ; ++j) st[i][j]=op(st[i<<1][j], st[i]
             <<1|1|[i]);
void set(int x, int y, T v){
         st[x+n][y+m]=v;
         for(int ¬=y+m; ¬>1; ¬>>=1) st [x+n] [¬>>1] =op (
             st[x+n][j], st[x+n][j^1]);
         for (int i=x+n; i>1; i>>=1) for (int j=y+m; j; j
            >>=1) st[i>>1][j]=op(st[i][j], st[i^1][
             j]);
T get (int x0, int y0, int x1, int y1) {
         T r=neutro;
         for (int i0=x0+n, i1=x1+n+1; i0<i1; i0>>=1, i1
            >>=1) {
                  int t[4], q=0;
                  if (i0&1) t [q++]=i0++;
                  if (i1&1) t [q++]=--i1;
                  for (int k=0; k < q; ++k) for (int j0=y0
                      +m, i1=v1+m+1; i0<i1; i0>>=1, i1
                     >>=1){
                           if(j0&1)r=op(r,st[t[k]][
                               j0++]);
                           if(j1&1) r = op(r, st[t]k)
                               ] [ -- j1] );
         return r:
```

3.25 Segment Tree Beats

```
// O(n*log(n)) build
// O(log(n)) get, upd
// updMax[l,r] -> ai = max(ai, v)
// updMin[l,r] -> ai = min(ai, v)
// updAdd[l,r] -> ai = ai + v
// get[l,r] -> return sum of the range [l,r]
typedef long long T;
T null=0,noval=0;
T INF=1e18;
struct Node{
    T sum,lazy;
    T max1,max2,maxc;
    T min1,min2,minc;
```

```
void build(T x) {
                 sum=max1=min1=x;
                 maxc=minc=1;
                 lazv=noVal;
                 max2 = -INF;
                 min2=INF;
        void oper(Node& a, Node& b) {
                 sum=a.sum+b.sum;
                 if(a.max1>b.max1) {
                          max1=a.max1:
                          maxc=a.maxc;
                          \max 2 = \max (a.\max 2, b.\max 1);
                 }else if(a.max1<b.max1){</pre>
                          \max 1=b.\max 1;
                          maxc=b.maxc;
                          max2=max(b.max2, a.max1);
                 }else{
                          \max 1 = a. \max 1;
                          maxc=a.maxc+b.maxc;
                          max2=max(a.max2, b.max2);
                 if(a.min1<b.min1) {</pre>
                          min1=a.min1;
                          minc=a.minc;
                          min2=min(a.min2, b.min1);
                 }else if(a.min1>b.min1){
                          min1=b.min1;
                          minc=b.minc;
                          min2=min(b.min2, a.min1);
                 }else{
                          min1=a.min1;
                          minc=a.minc+b.minc;
                          min2=min(a.min2, b.min2);
struct SegTree{
        vector<Node> vals:
        int size;
        SegTree(vector<T>& a) {
                 size=1;
                 while (size<sz(a)) size *= 2;</pre>
                 vals.resize(2*size);
                 build(a, 0, 0, size);
        void build(vector<T>& a, int x, int lx, int rx) {
                 if(rx-lx==1){
                          if(lx<sz(a))vals[x].build(a[lx]);</pre>
                          return;
```

```
int m = (1x+rx)/2;
        build(a, 2*x+1, 1x, m);
        build(a, 2*x+2, m, rx);
        vals[x].oper(vals[2*x+1], vals[2*x+2]);
void propagateMax(T v, int x, int lx, int rx){
        if(vals[x].min1>=v) return;
        vals[x].sum-=vals[x].min1*vals[x].minc;
        vals[x].min1=v;
        vals[x].sum+=vals[x].min1*vals[x].minc;
        if(rx-lx==1) {
                vals[x].max1=v;
        }else{
                if(v>=vals[x].max1){
                         vals[x].max1=v;
                 }else if(v>vals[x].max2){
                         vals[x].max2=v;
void propagateMin(T v, int x, int lx, int rx){
        if (vals[x].max1<=v) return;</pre>
        vals[x].sum-=vals[x].max1*vals[x].maxc;
        vals[x].max1=v;
        vals[x].sum+=vals[x].max1*vals[x].maxc;
        if(rx-lx==1){
                vals[x].min1=v;
        }else{
                if(v<=vals[x].min1){</pre>
                         vals[x].min1=v;
                 }else if(v<vals[x].min2){</pre>
                         vals[x].min2=v;
void propagateAdd(T v, int x, int lx, int rx){
        vals[x].sum+=v*((T)(rx-lx));
        vals[x].lazy+=v;
        vals[x].max1+=v;
        vals[x].min1+=v;
        if(vals[x].max2!=-INF)vals[x].max2+=v;
        if (vals[x].min2!=INF) vals[x].min2+=v;
void propagate(int x, int lx, int rx){
        if (rx-lx==1) return;
        int m = (lx + rx)/2;
        if(vals[x].lazv!=noVal){
                propagateAdd(vals[x].lazy, 2*x+1,
                     lx, m);
                propagateAdd(vals[x].lazy, 2*x+2,
                     m, rx);
                vals[x].lazy=noVal;
```

```
propagateMin(vals[x].max1, 2*x+1, lx, m);
        propagateMin(vals[x].max1, 2*x+2, m, rx);
        propagateMax(vals[x].min1, 2*x+1, lx, m);
        propagateMax(vals[x].min1, 2*x+2, m, rx);
void updAdd(int 1, int r, T v, int x, int lx, int
   rx) {
        if(lx>=r || l>=rx)return;
        if(lx>=1 && rx<=r){
                propagateAdd(v, x, lx, rx);
                return;
        propagate(x,lx,rx);
        int m = (1x+rx)/2;
        updAdd(l,r,v,2*x+1,lx,m);
        updAdd(1,r,v,2*x+2,m,rx);
        vals[x].oper(vals[2*x+1], vals[2*x+2]);
void updMax(int 1, int r, T v, int x, int lx, int
   rx) {
        if(lx>=r || l>=rx || vals[x].min1>v)
           return;
        if(lx>=1 && rx<=r && vals[x].min2>v) {
                propagateMax(v, x, lx, rx);
                return;
        propagate(x,lx,rx);
        int m = (1x+rx)/2;
        updMax(1, r, v, 2 * x + 1, 1x, m);
        updMax(1,r,v,2*x+2,m,rx);
        vals[x].oper(vals[2*x+1], vals[2*x+2]);
void updMin(int 1, int r, T v, int x, int lx, int
   rx) {
        if(lx>=r || l>=rx || vals[x].max1<v)
           return:
        if(lx>=1 && rx<=r && vals[x].max2<v){
                propagateMin(v, x, lx, rx);
                return:
        propagate(x,lx,rx);
        int m = (1x+rx)/2;
        updMin(l,r,v,2*x+1,lx,m);
        updMin(l,r,v,2*x+2,m,rx);
        vals[x].oper(vals[2*x+1], vals[2*x+2]);
T get(int 1, int r, int x, int lx, int rx) {
        if(lx>=r || l>=rx)return null;
        if(lx>=1 && rx<=r)return vals[x].sum;</pre>
        propagate(x,lx,rx);
```

```
int m = (1x+rx)/2;
                 T v1=qet (1, r, 2*x+1, 1x, m);
                 T v2=qet(1,r,2*x+2,m,rx);
                 return v1+v2;
        T get(int 1, int r) {return get(1,r+1,0,0,size);}
        void updAdd(int 1, int r, T v) {updAdd(1,r+1,v)
            ,0,0,size);}
        void updMin(int 1, int r, T v) {updMin(1,r+1,v)
            ,0,0,size);}
        void updMax(int 1, int r, T v) {updMax(1,r+1,v)
            ,0,0,size);}
};
```

3.26 Sparse Table

```
// O(n*log(n)) build
// O(1) get
typedef long long T;
Top(Ta, Tb); // max, min, gcd ...
struct Table{
        vector<vector<T>> st;
        Table(vector<T>& v) {
                 st.clear();
                 int n=v.size();
                 st.push back(v);
                 for(int j=1; (1<<j)<=n;++j) {
                          st.push_back(vector<T>(n));
                          for (int i=0; i+(1<<(i-1))< n; ++i) {
                                  st[j][i] = op(st[j-1][i], st
                                      [\dot{1}-1][\dot{1}+(1<<(\dot{1}-1))];
        T get(int 1, int r) {
                 int j=31- builtin clz(r-l+1);
                 return op(st[j][l], st[j][r-(1<<j)+1]);
};
```

3.27 Sparse Table 2D

```
// O(n*m*log(n)*log(m)) build
// O(1) get
typedef int T;
const int maxn = 1000, logn = 10;
T st[logn][maxn][logn][maxn];
int lg2[maxn+1];
T op (T a, T b); // min, max, qcd...
void build(int n, int m, vector<vector<T>>& a) {
```

```
for (int i=2; i \le \max(n,m); ++i) lg2[i] = lg2[i/2] +1;
        for(int i=0;i<n;++i){</pre>
                 for(int j=0; j<m;++j)
                          st[0][i][0][j]=a[i][j];
                 for(int k2=1; k2<logn; ++k2)
                          for (int j=\bar{0}; j+(1<<(k2-1))< m;++j)
                                   st[0][i][k2][j]=op(st[0][
                                       i][k2-1][j], st[0][i][
                                       k2-1][\dot{1}+(1<<(k2-1))]);
        for(int k1=1; k1<logn; ++k1)
                 for(int i=0; i<n;++i)
                          for(int k2=0; k2<logn; ++k2)
                                   for(int j=0; j<m; ++j)
                                            st[k1][i][k2][j]=
                                                op(st[k1-1][i
                                                ][k2][j], st[
                                                k1-1 | i+(1<<(
                                                k1-1))][k2][j
                                                ]);
T get(int x1, int y1, int x2, int y2){
        x2++; y2++;
        int a=lg2[x2-x1];
        int b=lq2[y2-y1];
        return op (
                 op(st[a][x1][b][y1],
                          st[a][x2-(1<<a)][b][y1]),
                 op (st [a] [x1] [b] [y2-(1<b)],
                          st[a][x2-(1<<a)][b][y2-(1<<b)])
        );
```

Sqrt Descomposition 3.28

```
// O(n) build
// O(n/b+b) get, set
typedef long long T;
struct SQRT{
        int b; // check b
        vector<T> a,bls;
        SQRT(vector<T>& arr, int n) {
                b=ceil(sqrt(n));a=arr;
                bls.assign(b, 0);
                for(int i=0; i<n; ++i) {
                         bls[i/b] += a[i];
        void set(int x, int v){
                bls[x/b] -= a[x];
                a[x]=v;
                bls[x/b]+=a[x];
```

3.29 Treap

```
// Treap => Binary Search Tree + Binary Heap
// 1. create a empty root (PTreap root=nullptr;)
// 2. Append the nodes in asc order
// PTreap tmp=new Treap(x);
// root=merge(root, tmp);
typedef long long T;
typedef unsigned long long u64;
mt19937_64 rng (chrono::steady_clock::now().
   time since epoch().count());
T \text{ null} = 0;
struct Treap{
        Treap *1,*r,*dad; // left child, right child
        u64 prior; // random
        T val; // value
        int sz; // size subtree
        Treap(T v) {
                l=r=nullptr;
                prior=rng();
                val=v;sz=1;
        Treap(){
                delete 1;
                delete r;
};
typedef Treap* PTreap;
int cnt(PTreap x) {return (!x?0:x->sz);}
// updates node with its children information
void pull(PTreap x) {
        x->sz=cnt(x->1)+cnt(x->r)+1;
        if (x->1) x->1->dad=x;
        if (x->r) x->r->dad=x;
// O(log(n)) divide the treap in two parts
// [nodes value <= key], [nodes value > key]
pair<PTreap, PTreap> split(PTreap x, T key) {
        if(!x)return {nullptr, nullptr};
```

```
if (x->val>kev) {
                 auto got=split(x->1, key);
                 x->l=qot.second;
                 pull(x);
                 return {got.first, x};
        }else{
                 auto got=split(x->r, key);
                 x->r=qot.first;
                 pull(x);
                 return {x, got.second};
// O(log(n)) merge two treap
// if all values in treap x < all values in treap y
PTreap merge (PTreap x, PTreap y) {
        if(!x)return v;
        if(!y)return x;
        if (x->prior<=y->prior) {
                 x \rightarrow r = merge(x \rightarrow r, y);
                 pull(x);
                 return x;
        }else{
                 y->l=merge(x, y->l);
                 pull(y);
                 return v;
// O(n*log(n))
// Combine two treap into one
PTreap combine (PTreap x, PTreap y) {
        if(!x)return v;
        if(!y)return x;
        if (x->prior<y->prior) swap (x, y);
        auto z=split(y, x->val);
        x \rightarrow r = combine(x \rightarrow r, z.second);
        x->l=combine(z.first, x->l);
        return x;
// O(log(n))
// return kth element - indexed 0
T kth(PTreap& x, int k){
        if(!x)return null;
        if (k==cnt (x->1)) return x->val;
        if (k < cnt(x->1)) return kth(x->1, k);
        return kth(x->r, k-cnt(x->1)-1);
// O(log(n))
// return {index, val}
pair<int, T> lower_bound(PTreap x, T key) {
        if(!x)return {0, null};
        if(x->val<key){</pre>
                 auto y=lower bound(x->r, key);
```

```
y.first+=cnt(x->1)+1;
return y;
}
auto y=lower_bound(x->1, key);
if(y.first==cnt(x->1))y.second=x->val;
return y;
}
// O(n) print the treap
void dfs(PTreap x) {
    if(!x) return;
    dfs(x->1);
    cout<<x->val<<" ";
    dfs(x->r);
}
```

3.30 Trie Bit

```
struct node{
int childs[2]{-1, -1};
};
struct TrieBit{
        vector<node> nds;
        vi passNums;
        TrieBit(){
                nds.pb(node());
                passNums.pb(0);
        void insert(int num) {
                int cur = 0;
                for(int i = 30; i >= 0; i--){
                        bool bit = (num >> i) & 1;
                        if (nds[cur].childs[bit] == -1) {
                                 nds[cur].childs[bit] =
                                    nds.size();
                                 nds.pb(node());
                                 passNums.pb(0);
                         passNums[cur]++;
                         cur = nds[cur].childs[bit];
                passNums[cur]++;
        void remove(int num) {
                int cur = 0;
                for(int i = 30; i >= 0; i--) {
                        bool bit = (num >> i) & 1;
                        passNums[cur]--;
                         cur = nds[cur].childs[bit];
```

```
passNums[cur]--;
}
int maxXor(int num) {
    int ans = 0;
    int cur = 0;
    for(int i = 30; i >= 0; i--) {
        bool bit = (num >> i) & 1;

        int n1 = nds[cur].childs[!bit];
        if (n1 != -1 && passNums[n1]) {
            ans += (1 << i);
            bit = !bit;
        }
        cur = nds[cur].childs[bit];
}
return ans;
}
</pre>
```

3.31 Two Stacks

```
// O(1) push, pop, get
typedef long long T;
struct Node{T val,acum;};
struct TwoStacks{
        stack<Node> s1,s2;
        void push(T x){
                Node tmp=\{x, x\};
                if(!s2.empty()){
                         // tmp.acum + s2.top().acum
                s2.push(tmp);
        void pop(){
                if(s1.empty()){
                         while(!s2.empty()){
                                 Node tmp=s2.top();
                                 if(s1.empty()) {
                                          // tmp.acum = tmp
                                              . va 1
                                 }else{
                                          // tmp.acum + s1.
                                             top().acum
                                 s1.push(tmp);
                                 s2.pop();
                 s1.pop();
```

3.32 Wavelet Tree

};

```
const int maxn = 1e5+5;
const int maxv = 1e9;
const int minv = -1e9;
// O(n*log(n)) build
// O(\log(n)) kth, lte, cnt, sum
// 1. int a[maxn];
// 2. WaveletTree wt;
// 3. fill a[1;n]
// 4. wt.build(a+1, a+n+1, minv, maxv);
struct WaveletTree { // indexed 1
        int lo, hi:
        WaveletTree *1, *r;
        int *b, bsz, csz;
        ll *c;
        WaveletTree() {
                hi=bsz=csz=0;
                l=r=NULL;
                10=1;
        void build(int *from, int *to, int x, int y) {
                 lo=x, hi=v;
                if (from>=to) return;
                 int mid=lo+(hi-lo)/2;
                 auto f=[mid](int x){return x<=mid;};</pre>
                 b=(int*)malloc((to-from+2)*sizeof(int));
                bsz=0;
                b[bsz++]=0;
                 c=(ll*)malloc((to-from+2)*sizeof(ll));
                 csz=0;
                 c[csz++]=0;
                 for(auto it=from;it!=to;++it) {
                         b[bsz] = (b[bsz-1] + f(*it));
                         c[csz] = (c[csz-1] + (*it));
                         bsz++;csz++;
                 if (hi==lo) return;
```

```
auto pivot=stable partition(from, to, f);
        l=new WaveletTree();
        l->build(from, pivot, lo, mid);
        r=new WaveletTree();
        r->build(pivot, to, mid+1, hi);
//kth smallest element in [1, r]
int kth(int 1, int r, int k){
        if(l>r)return 0;
        if(lo==hi)return lo;
        int inLeft=b[r]-b[l-1], lb=b[l-1], rb=b[r
        if (k<=inLeft) return this->l->kth(lb+1, rb
        return this->r->kth(l-lb, r-rb, k-inLeft)
//count of numbers in [1, r] Less than or equal
   to k
int lte(int 1, int r, int k){
        if(1>r || k<10) return 0;
        if (hi<=k) return r-l+1;</pre>
        int lb=b[1-1], rb=b[r];
        return this->l->lte(lb+1, rb, k)+this->r
           ->lte(l-lb, r-rb, k);
//count of numbers in [l, r] equal to k
int count(int 1, int r, int k){
        if(l>r || k<lo || k>hi)return 0;
        if(lo==hi)return r-l+1;
        int lb=b[l-1], rb=b[r];
        int mid=(lo+hi)>>1;
        if (k<=mid) return this->l->count (lb+1, rb,
        return this->r->count(l-lb, r-rb, k);
//sum of numbers in [l ,r] less than or equal to
11 sum(int 1, int r, int k){
        if(1>r || k<10) return 0;
        if (hi<=k) return c[r]-c[l-1];</pre>
        int lb=b[l-1], rb=b[r];
        return this->l->sum(lb+1, rb, k)+this->r
           \rightarrowsum(l-lb, r-rb, k);
~WaveletTree(){
        delete 1;
        delete r;
```

} **;**

FLUJOS

4 Flujos

4.1 Blossom

```
// O(|E||V|^2)
struct network {
  struct struct edge { int v; struct edge * n; };
 typedef struct edge* edge;
  int n;
  struct_edge pool[MAXE]; ///2*n*n;
  edge top;
  vector<edge> adj;
  queue<int> q;
 vector<int> f, base, ing, inb, inp, match;
  vector<vector<int>> ed;
 network(int n) : n(n), match(n, -1), adj(n), top(pool),
      f(n), base(n),
                   inq(n), inb(n), inp(n), ed(n, vector<
                      int>(n)) {}
 void add edge(int u, int v) {
    if(ed[u][v]) return;
    ed[u][v] = 1;
    top->v = v, top->n = adj[u], adj[u] = top++;
    top->v = u, top->n = adj[v], adj[v] = top++;
  int get lca(int root, int u, int v) {
    fill(inp.begin(), inp.end(), 0);
    while(1) {
      inp[u = base[u]] = 1;
      if(u == root) break;
      u = f[ match[u] ];
    while(1) {
      if(inp[v = base[v]]) return v;
      else \bar{v} = f[ match[v] ];
 void mark(int lca, int u) {
    while(base[u] != lca) {
      int v = match[u];
      inb[base[u]] = 1;
      inb[base[v]] = 1;
      u = f[v];
      if(base[u] != lca) f[u] = v;
 void blossom contraction(int s, int u, int v) {
    int lca = get_lca(s, u, v);
    fill(inb.begin(), inb.end(), 0);
   mark(lca, u); mark(lca, v);
    if(base[u] != lca) f[u] = v;
    if(base[v] != lca) f[v] = u;
   for (int u = 0; u < n; u++)
      if(inb[base[u]]) {
```

```
base[u] = lca;
        if(!inq[u]) {
         inq[u] = 1;
          q.push(u);
  int bfs(int s) {
    fill(inq.begin(), inq.end(), 0);
    fill(f.begin(), f.end(), -1);
    for(int i = 0; i < n; i++) base[i] = i;</pre>
    q = queue<int>();
    q.push(s);
    inq[s] = 1;
    while(q.size()) {
      int u = q.front(); q.pop();
      for (edge e = adj[u]; e; e = e -> n) {
        int v = e -> v;
        if(base[u] != base[v] && match[u] != v) {
          if ((v == s) | | (match[v] != -1 && f[match[v]]
            blossom_contraction(s, u, v);
          else if (f[v] == -1) {
            f[v] = u;
            if (match[v] == -1) return v;
            else if(!ing[match[v]]) {
              inq[match[v]] = 1;
              q.push(match[v]);
    return -1;
  int doit(int u) {
    if(u == -1) return 0;
    int v = f[u];
    doit(match[v]);
    match[v] = u; match[u] = v;
    return u != -1;
  /// (i < net.match[i]) => means match
  int maximum matching() {
    int ans = 0;
    for(int u = 0; u < n; u++)
      ans += (match[u] == -1) && doit(bfs(u));
    return ans;
};
```

```
4.2 Dinic
```

```
// O(|E| * |V|^2)
struct edge { ll v, cap, inv, flow, ori; };
struct network {
        ll n, s, t;
        vector<ll> lvl;
        vector<vector<edge>> q;
        network(ll n) : \tilde{n}(n), lvl(n), g(n) {}
        void add_edge(int u, int v, ll c) {
                q[u].push_back(\{v, c, sz(q[v]), 0, 1\});
                q[v].push_back({u, 0, sz(q[u])-1, c, 0});
        bool bfs() {
                fill(lvl.begin(), lvl.end(), -1);
                queue<11> q;
                [vl[s] = 0;
                for(q.push(s); q.size(); q.pop()) {
                        ll u = q.front();
                        for(auto &e : q[u]) {
                                 if(e.cap > 0 && lvl[e.v]
                                    == -1) {
                                         lvl[e.v] = lvl[u]
                                             ]+1;
                                         q.push(e.v);
                return lvl[t] != -1;
        11 dfs(ll u, ll nf) {
                if(u == t) return nf;
                ll res = 0;
                for(auto &e : q[u]) {
                        if(e.cap > 0 && lvl[e.v] == lvl[u
                            ]+1) {
                                 ll tf = dfs(e.v, min(nf,
                                    e.cap));
                                 res += tf; nf -= tf; e.
                                    cap -= tf;
                                 g[e.v][e.inv].cap += tf;
                                 q[e.v][e.inv].flow -= tf;
                                 e.flow += tf;
                                 if(nf == 0) return res;
                if(!res) lvl[u] = -1;
                return res;
        ll \max flow(ll so, ll si, ll res = 0) {
                s = so; t = si;
                while(bfs()) res += dfs(s, LONG LONG MAX)
                return res;
        void min cut(){
                queue<11> q;
                vector<bool> vis(n, 0);
```

```
vis[s] = 1;
        for(q.push(s); q.size(); q.pop()) {
                ll u = q.front();
                for(auto &e : q[u]) {
                         if(e.cap > 0 && !vis[e.v
                            ]) {
                                 a.push(e.v);
                                 vis[e.v] = 1;
        vii ans;
        for (int i = 0; i<n; i++) {</pre>
                for (auto &e : q[i]) {
                         if (vis[i] && !vis[e.v]
                            && e.ori) {
                                 ans.push back({i
                                    +1, e.v+1);
        for (auto [x, y] : ans) cout << x << ' '</pre>
           << y << ln;
bool dfs2(vi &path, vector<bool> &vis, int u) {
        vis[u] = 1;
        for (auto &e : g[u]) {
                if (e.flow > 0 && e.ori && !vis[e
                    .v]){
                         if (e.v == t || dfs2(path
                            , vis, e.v)){
                                 path.push back (e.
                                    v);
                                 e.flow = 0;
                                 return 1;
        return 0;
void disjoint paths() {
        vi path;
        vector<bool> vis(n, 0);
        while (dfs2(path, vis, s)){
                path.push_back(s);
                reverse (all (path));
                cout << sz(path) << ln;
                for (int v : path) cout << v+1 <<
                cout << ln;
                path.clear(); vis.assign(n, 0);
```

} ;

4.3 Edmonds Karp

```
// O(V * E^2)
ll bfs(vector<vi> &adj, vector<vl> &capacity, int s, int
   t, vi& parent) {
        fill(parent.begin(), parent.end(), -1);
        parent[s] = -2;
        queue<pll> q;
        q.push({s, INFL});
        while (!q.empty()) {
                int cur = q.front().first;
                ll flow = q.front().second;
                q.pop();
                for (int next : adj[cur]) {
                        if (parent[next] == -1LL &&
                            capacity[cur][next]) {
                                 parent[next] = cur;
                                 ll new_flow = min(flow,
                                    capacity[cur][next]);
                                 if (next == t)
                                         return new flow;
                                 q.push({next, new_flow});
        return 0;
11 maxflow(vector<vi> &adj, vector<vl> &capacity, int s,
   int t, int n) {
        11 \text{ flow} = 0;
        vi parent(n);
        ll new flow;
        while ((new_flow = bfs(adj, capacity, s, t,
           parent))) {
                flow += new flow;
                int cur = t;
                while (cur != s) {
                        int prev = parent[cur];
                        capacity[prev][cur] -= new_flow;
                        capacity[cur][prev] += new flow;
                        cur = prev;
        return flow;
```

4.4 Hopcroft Karp

```
// O(|E|*sqrt(|V|))
struct mbm {
  vector<vector<int>> q;
  vector<int> d, match;
  int nil, l, r;
  /// u \rightarrow 0 to 1, v \rightarrow 0 to r
  mbm(int l, int r) : q(l+r), d(l+l+r, INF), match(l+r, l)
     +r),
                       nil(l+r), l(l), r(r) {}
  void add_edge(int a, int b) {
    q[a].push back(1+b);
    q[1+b].push back(a);
 bool bfs() {
    queue<int> q;
    for(int u = 0; u < 1; u++) {
      if (match[u] == nil) {
        d[u] = 0;
        q.push(u);
      } else d[u] = INF;
    d[nil] = INF;
    while(q.size()) {
      int u = q.front(); q.pop();
      if(u == nil) continue;
      for(auto v : q[u]) {
        if(d[match[v]] == INF) {
          d[match[v]] = d[u]+1;
          q.push(match[v]);
    return d[nil] != INF;
  bool dfs(int u) {
    if(u == nil) return true;
    for(int v : q[u]) {
      if (d[match[v]] == d[u]+1 && dfs(match[v])) {
        match[v] = u; match[u] = v;
        return true;
    d[u] = INF;
    return false:
  int max matching() {
    int ans = 0;
    while(bfs()) {
      for(int u = 0; u < 1; u++) {</pre>
        ans += (match[u] == nil && dfs(u));
    return ans;
  void matchs() {
```

```
for (int i = 0; i<1; i++) {
    if (match[i] == l+r) continue;
    cout << i+1 << ' ' << match[i]+1-1 << ln;
    }
};</pre>
```

4.5 Hungarian

```
#define rep(i, a, b) for(int i = a; i < (b); ++i)
typedef double type;
const type INF_TYPE = LLONG_MAX;
pair<type, vi> hungarian(const vector<vector<type>> &a) {
        if (a.empty()) return {0, {}};
        int n = sz(a) + 1, m = sz(a[0]) + 1;
        vector<type> u(n), v(m); vi p(m), ans(n-1);
        rep(i,1,n) {
                i = [0]q
                int j0 = 0; // add "dummy" worker 0
                vector<type> dist(m, INF TYPE); vi pre(m,
                vector<bool> done(m + 1);
                do { // dijkstra
                        done[j0] = true;
                        int i\bar{0} = p[i0], i1; type delta =
                            INF_TYPE;
                        rep(j,1,m) if (!done[j]) {
                                 auto cur = a[i0 - 1][j -
                                   1] - u[i0] - v[j];
                                 if (cur < dist[j]) dist[j</pre>
                                    ] = cur, pre[j] = j0;
                                 if (dist[j] < delta)</pre>
                                    delta = dist[i], i1 =
                                    j;
                        rep(j,0,m) {
                                 if (done[j]) u[p[j]] +=
                                    delta, v[j] -= delta;
                                 else dist[j] -= delta;
                        j0 = j1;
                } while (p[j0]);
                while (j0) { // update alternating path
                        int j1 = pre[j0];
                        p[j0] = p[j1], j0 = j1;
        rep(j,1,m) if (p[j]) ans[p[j] - 1] = j - 1;
        return {-v[0], ans}; // min cost
```

4.6 Maximum Bipartite Matching

```
// O(|E| * |V|)
struct mbm {
        int 1, r;
        vector<vector<int>> q;
        vector<int> match, seen;
        mbm(int 1, int r) : 1(1), r(r), g(1), match(r),
            seen(r){}
        void add_edge(int 1, int r) { g[1].push_back(r);
        bool dfs(int u) {
                for(auto v : q[u]) {
                        if(seen[v]++) continue;
                        if (match[v] == -1 || dfs(match[v])
                            ])) {
                                match[v] = u;
                                 return true;
                return false;
        int max_matching() {
                int ans = 0;
                fill(match.begin(), match.end(), -1);
                for(int u = 0; u < 1; ++u) {
                        fill(seen.begin(), seen.end(), 0)
                        ans += dfs(u);
                return ans;
        void matchs() {
                for (int i = 0; i<r; i++) {
                        if (match[i] == -1) continue;
                        cout << match[i]+1 << ' ' << i+1
                             << ln;
} ;
```

4.7 Minimum Cost Maximum Flow

```
// O(|V|*|E|^2*log(|E|))
template <class type>
struct mcmf {
    struct edge { int u, v, cap, flow; type cost; };
    int n;
    vector<edge> ed;
    vector<vector<int>> g;
    vector<iint>> p;
    vector<type> d, phi;
    mcmf(int n) : n(n), g(n), p(n), d(n), phi(n) {}
```

```
void add edge(int u, int v, int cap, type cost) {
        g[u].push_back(ed.size());
        ed.push_back({u, v, cap, 0, cost});
        q[v].push back(ed.size());
        ed.push_back({v, u, 0, 0, -cost});
bool dijkstra(int s, int t) {
        fill(d.begin(), d.end(), INF_TYPE);
        fill(p.begin(), p.end(), -1);
        set<pair<type, int>> q;
        d[s] = 0;
        for(q.insert({d[s], s}); q.size();) {
                int u = (*q.begin()).second; q.
                    erase(q.begin());
                for(auto v : q[u]) {
                        auto &e = ed[v];
                        type nd = d[e.u] + e.cost +
                            phi[e.u]-phi[e.v];
                        if(0 < (e.cap-e.flow) &&
                            nd < d[e.v]) {
                                 q.erase({d[e.v],
                                    e.v});
                                 d[e.v] = nd; p[e.
                                    v = v;
                                 q.insert({d[e.v],
                                     e.v});
        for(int i = 0; i < n; i++) phi[i] = min(</pre>
           INF_TYPE, phi[i]+d[i]);
        return d[t] != INF TYPE;
pair<int, type> max flow(int s, int t) {
        type mc = 0;
        int mf = 0;
        fill(phi.begin(), phi.end(), 0);
        while(dijkstra(s, t)) {
                int flow = INF;
                for(int v = p[t]; v != -1; v = p[
                     ed[v].u ])
                        flow = min(flow, ed[v].
                            cap-ed[v].flow);
                for (int v = p[t]; v != -1; v = p[
                     ed[v].u ]) {
                        edge &e1 = ed[v];
                        edge &e2 = ed[v^1];
                        mc += e1.cost*flow;
                        e1.flow += flow;
                        e2.flow -= flow;
                mf += flow;
        return {mf, mc};
```

} **;**

4.8 MCMF Vasito

```
// O(|E| * |F| * log(|V|))
typedef int tf;
typedef int tc;
const tf INFFLOW=1e9;
const tc INFCOST=1e9;
struct MCF {
  int n;
  vector<tc> prio, pot; vector<tf> curflow; vector<int>
     prevedge, prevnode;
  priority queue<pair<tc, int>, vector<pair<tc, int>>,
     greater<pair<tc, int>>> q;
  struct edge{int to, rev; tf f, cap; tc cost;};
  vector<vector<edge>> q;
 MCF(int n):n(n),prio(n),curflow(n),prevedge(n),prevnode
      (n), pot(n), q(n) {}
  void add edge(int s, int t, tf cap, tc cost) {
    q[s].push\_back((edge)\{t,sz(g[t]),0,cap,cost\});
    q[t].push_back((edge) {s,sz(q[s])-1,0,0,-cost});
  pair<tf,tc> get flow(int s, int t) {
    tf flow=0; tc flowcost=0;
    while(1){
      q.push({0, s});
      fill(all(prio), INFCOST);
      prio(s)=0; curflow(s)=INFFLOW;
      while(!a.emptv()) {
        auto cur=q.top();
        tc d=cur.first;
        int u=cur.second;
        q.pop();
        if(d!=prio[u]) continue;
        for (int i=0; i < sz(g[u]); ++i) {</pre>
          edge \&e=g[u][i];
          int v=e.to;
          if(e.cap<=e.f) continue;</pre>
          tc nprio=prio[u]+e.cost+pot[u]-pot[v];
          if(prio[v]>nprio) {
            prio[v]=nprio;
            q.push({nprio, v});
            prevnode[v]=u; prevedge[v]=i;
            curflow[v]=min(curflow[u], e.cap-e.f);
      if(prio[t] == INFCOST) break;
      for(int i=0;i<n;i++) pot[i]+=prio[i];</pre>
      tf df=min(curflow[t], INFFLOW-flow);
      flow+=df;
      for(int v=t; v!=s; v=prevnode[v]) {
```

```
edge &e=g[prevnode[v]][prevedge[v]];
    e.f+=df; g[v][e.rev].f-=df;
    flowcost+=df*e.cost;
}

return {flow,flowcost};
};
```

4.9 Scaling Algorithm

```
// O(|E|^2 * log(C)) C = maximum edge weight of the graph
struct MaxFlow {
    static const 11 INF = 1e18;
    struct Edge {int u,v;ll w;};
    int n, s, t;
    vector<vector<int>> q;
    vector<Edge> ed;
    vector<bool> vis;
    11 \text{ flow} = 0;
   MaxFlow(int n, int s, int t) : n(n), s(s), t(t), q(n)
    int add edge(int u, int v, ll forward, ll backward =
        const int id = (int)ed.size();
        q[u].emplace back(id);
        ed.push_back({u, v, forward});
        q[v].emplace back(id + 1);
        ed.push_back({v, u, backward});
        return id;
    bool dfs(int node, ll lim) {
        if (node == t) return true;
        if (vis[node]) return false;
        vis[node] = true;
        for (int i : q[node]) {
            auto &e = ed[i];
            auto &back = ed[i ^ 1];
            if (e.w >= lim) {
                if (dfs(e.v, lim)) {
                    e.w -= lim;
                    back.w += lim;
                    return true;
        return false:
    ll max flow() {
        for (ll bit = 1ll << 62; bit > 0; bit /= 2) {
            bool found = false;
            do {
                vis.assign(n, false);
```

```
found = dfs(s, bit);
    flow += bit * found;
} while (found);
}
return flow;
}
};
```

4.10 Weighted Matching

```
// O(|V|^3)
typedef int type;
struct matching weighted {
  int 1, r;
  vector<vector<type>> c;
  matching weighted (int 1, int r) : 1(1), r(r), c(1,
     vector<type>(r)) {
    assert (1 < = r);
  void add_edge(int a, int b, type cost) { c[a][b] = cost
     ; }
  type matching() {
    vector<type> v(r), d(r); // v: potential
    vector<int> ml(l, -1), mr(r, -1); // matching pairs
    vector<int> idx(r), prev(r);
    iota(idx.begin(), idx.end(), 0);
    auto residue = [&](int i, int j) { return c[i][j]-v[j
       ]; };
    for(int f = 0; f < 1; ++f) {
      for (int j = 0; j < r; ++j) {
        d[i] = residue(f, i);
        prev[j] = f;
      type w;
      int j, 1;
      for (int s = 0, t = 0;;) {
        if(s == t) {
          1 = s;
          w = d[idx[t++]];
          for(int k = t; k < r; ++k) {
            j = idx[k];
            type h = d[i];
            if (h <= w) {
              if (h < w) t = s, w = h;
              idx[k] = idx[t];
              idx[t++] = j;
          for (int k = s; k < t; ++k) {
            j = idx[k];
            if (mr[j] < 0) goto aug;
```

```
5 GEOMETRIA
```

```
int q = idx[s++], i = mr[q];
        for (int k = t; k < r; ++k) {
          i = idx[k];
          type h = residue(i, j) - residue(i, q) + w;
          if (h < d[i]) {
            d[i] = h;
            prev[j] = i;
            if(h == w) {
              if(mr[j] < 0) goto aug;</pre>
              idx[k] = idx[t];
              idx[t++] = j;
          }
      aug: for (int k = 0; k < 1; ++k)
       v[idx[k]] += d[idx[k]] - w;
      int i;
      do {
        mr[j] = i = prev[j];
        swap(i, ml[i]);
      } while (i != f);
    type opt = 0;
    for (int i = 0; i < 1; ++i)
      opt += c[i][ml[i]]; // (i, ml[i]) is a solution
    return opt;
};
```

5 Geometria

5.1 2D Tree

```
// given a set of points, answer queries of nearest point
    in O(\log(n))
bool onx(pt a, pt b) {return a.x < b.x;}</pre>
bool ony(pt a, pt b) {return a.y < b.y;}</pre>
struct Node {
        pt pp;
        If x0 = \inf, x1 = -\inf, y0 = \inf, y1 = -\inf;
        Node *first = 0, *second = 0;
        ll distance(pt p) {
                 11 x = min(max(x0, p.x), x1);
                 ll y = min(max(y0, p.y), y1);
                 return norm2 (pt (x, y) - p);
        Node(vector<pt>&& vp) : pp(vp[0]) {
                 for(pt p : vp) {
                         x0 = min(x0, p.x);
             x1 = max(x1, p.x);
```

```
y0 = min(y0, p.y);
            v1 = max(y1, p.y);
                if(vp.size() > 1) {
                         sort(all(vp), x1 - x0 >= y1 - y0
                             ? onx : ony);
                         int m = vp.size() / 2;
                         first = new Node({vp.begin(), vp.
                             begin() + m});
                         second = new Node({vp.begin() + m
                             , vp.end() });
};
struct KDTree {
        Node* root:
        KDTree(const vector<pt>& vp): root(new Node({all(
            {}(({qv
        pair<11, pt> search(pt p, Node *node) {
                if(!node->first){
                         // avoid query point as answer
                         // if(p.x == node->pp.x && p.y ==
                              node->pp.y) return {inf, pt()
                         return {norm2 (p-node->pp), node->
                            pp};
                Node *f = node \rightarrow first, *s = node \rightarrow second;
                ll bf = f->distance(p), bs = s ->
                    distance(p);
                if(bf > bs) swap(bf, bs), swap(f, s);
                auto best = search(p, f);
                if(bs < best.ff) best = min(best, search(</pre>
                    p, s));
                return best;
        pair<ll, pt> nearest(pt p) { return search(p, root
           ); }
};
```

5.2 3D

```
typedef double lf;
struct p3 {
    lf x, y, z;
        p3(){}
        p3(lf x, lf y, lf z): x(x), y(y), z(z){}
    p3 operator + (p3 p){ return {x + p.x, y + p.y, z + p.z}; }
    p3 operator - (p3 p){ return {x - p.x, y - p.y, z - p.z}; }
    p3 operator * (lf d){ return {x * d, y * d, z * d}; }
    p3 operator / (lf d){ return {x * d, y * d, z * d}; }
```

```
// only for floating point
    // Some comparators
    bool operator == (p3 p) { return tie(x, y, z) == tie(p
       .x, p.y, p.z); }
   bool operator != (p3 p) { return !operator == (p); }
        void print() { cout << x << " " << y << " " << z</pre>
        // scale: (newnorm / norm) * p3
};
lf dot(p3 v, p3 w) { return v.x * w.x + v.y * w.y + v.z *
   w.z; }
p3 cross(p3 v, p3 w) {
    return { v.y * w.z - v.z * w.y, v.z * w.x - v.x * w.z
       , v.x * w.y - v.y * w.x };
lf norm2(p3 v) { return dot(v, v); }
lf norm(p3 v) { return sqrt(norm2(v)); }
p3 unit(p3 v) { return v / norm(v); }
// ang(RAD)
double angle(p3 v, p3 w) {
    double cos theta = dot(v, w) / norm(v) / norm(w);
    return acos(max(-1.0, min(1.0, cos_theta)));
// orient s, pgr form a triangle pos: 'up', zero = on,
   nea = 'dow'
lf orient(p3 p, p3 q, p3 r, p3 s){
        return dot(cross((q - p), (r - p)), (s - p));
// same as 2D but in n-normal direction
lf orient_by_normal(p3 p, p3 q, p3 r, p3 n) {
        return dot(cross((q - p), (r - p)), n);
struct plane {
    p3 n; lf d; // n: normal d: dist to zero
    // From normal n and offset d
    plane(p3 n, lf d): n(n), d(d) {}
    // From normal n and point P
    plane(p3 n, p3 p): n(n), d(dot(n, p)) {}
    // From three non-collinear points P,Q,R
    plane(p3 p, p3 q, p3 r): plane(cross((q - p), (r - p)
       ), p){}
    // - these work with lf = int
    lf side(p3 p) { return dot(n, p) - d; }
    double dist(p3 p) { return abs(side(p)) / norm(n); }
    plane translate(p3 t) {return \{n, d + dot(n, t)\}; }
    /// - these require lf = double
    plane shift up(double dist) { return {n, d + dist *
       norm(n) }; }
    p3 proj(p3 p) { return p - n * side(p) / norm2(n); }
    p3 refl(p3 p) { return p - n * 2 * side(p) / norm2(n);
```

```
} ;
struct line3d {
        p3 d, o; // d: dir o: point on line
        // From two points P, Q
        line3d(p3 p, p3 q): d(q - p), o(p){}
        // From two planes p1, p2 (requires lf = double)
        line3d(plane p1, plane p2) {
                d = cross(p1.n, p2.n);
                o = cross((p2.n * p1.d - p1.n * p2.d), d)
                    / norm2(d);
        // - these work with lf = int
        double dist2(p3 p) { return norm2(cross(d, (p - o)
           )) / norm2(d); }
        double dist(p3 p) { return sqrt(dist2(p)); }
        bool cmp_proj(p3 p, p3 q) { return dot(d, p) < dot
            (d, q); }
        // - these require If = double
        p3 proj(p3 p) { return o + d * dot(d, (p - o)) /
           norm2(d);
        p3 refl(p3 p) { return proj(p) * 2 - p; }
        p3 inter(plane p) { return o - d * p.side(o) / dot
            (p.n, d); }
        // get other point: pl.o + pl.d * t;
} ;
double dist(line3d 11, line3d 12) {
        p3 n = cross(11.d, 12.d);
        if(n == p3(0, 0, 0)) return 11.dist(12.o); //
           parallel
        return abs(dot((12.o - 11.o), n)) / norm(n);
// closest point on 11 to 12
p3 closest on line1(line3d l1, line3d l2) {
        p3 n2 = cross(12.d, cross(11.d, 12.d));
        return 11.0 + 11.d * (dot((12.0 - 11.0), n2)) /
           dot(11.d, n2);
double small angle(p3 v, p3 w) { return acos(min(abs(dot(v
   (w) / norm(v) / norm(w), 1.0); } // 0.90
double angle(plane p1, plane p2) { return small_angle(p1.n
    p2.n); }
bool is_parallel(plane p1, plane p2) { return cross(p1.n,
   p2.n) == p3(0, 0, 0);
bool is perpendicular (plane p1, plane p2) { return dot (p1.
   n, p2.n) == 0;
double angle(line3d 11, line3d 12) { return small_angle(l1
   .d, 12.d); }
bool is parallel(line3d l1, line3d l2) { return cross(l1.d
   , 12.d) == p3(0, 0, 0); }
bool is perpendicular(line3d 11, line3d 12) { return dot(
   11.d, 12.d) == 0;
double angle(plane p, line3d l) { return M_PI / 2 -
```

```
small_angle(p.n, l.d); }
bool is_parallel(plane p, line3d l) { return dot(p.n, l.d)
    == 0;
bool is perpendicular (plane p, line3d 1) { return cross(p.
   n, 1.d) == p3(0, 0, 0);
line3d perp_through(plane p, p3 o) { return line3d(o, o +
plane perp through (line3d 1, p3 o) { return plane (l.d, o);
pair<p3, lf> smallest enclosing sphere(vector<p3> p) {
    int n = p.size();
    p3 c(0, 0, 0);
    for (int i = 0; i < n; i++) c = c + p[i];
    c = c / n;
    double ratio = 0.1;
    int pos = 0;
    int it = 100000;
    while (it--) {
        pos = 0;
        for (int i = 1; i < n; i++) {
            if(norm2(c - p[i]) > norm2(c - p[pos])) pos =
                i;
        c = c + (p[pos] - c) * ratio;
        ratio *= 0.998;
    return {c, sqrt(norm2(c - p[pos]))};
```

5.3 Circulos

```
// add Lines Points
enum {OUT, IN, ON};
struct circle {
       pt center; lf r;
        // (x - xo)^2 + (y - yo)^2 = r^2
        circle(pt c, lf r): center(c), r(r){};
        // circle that passes through abc
        circle(pt a, pt b, pt c) {
                b = b - a, c = c - a;
                assert (cross(b, c) != 0); // no
                    circumcircle if A, B, C aligned
                pt cen = a + rot 90 (b * norm2 (c) - c *
                   norm2(b)) / cross(b, c) / 2;
                center = cen;
                r = norm(a - cen);
        // diameter = segment pg
        circle(pt p, pt q) {
```

```
center = (p + q) * 0.5L;
                r = dis(p, q) * 0.5L;
        int contains(pt &p) {
                lf det = r * r - dis2(center, p);
                if(fabsl(det) <= EPS) return ON;</pre>
                return (det > EPS ? IN : OUT);
        bool in(circle c) { return norm(center - c.center)
            + r <= c.r + EPS; } // non strict
};
// centers of the circles that pass through ab and has
vector<pt> centers(pt a, pt b, lf r) {
        if (norm(a - b) > 2 * r + EPS) return {};
        pt m = (a + b) / 2;
        double f = sqrt(r * r / norm2(a - m) - 1);
        pt c = rot 90 (a - m) * f;
        return {m - c, m + c};
vector<pt> inter cl(circle c, line l){
        vector<pt> s;
        pt p = l.proj(c.center);
        If d = norm(p - c.center);
        if(d - EPS > c.r) return s;
        if (abs(d - c.r) <= EPS) { s.push back(p); return s</pre>
        d=sqrt(c.r * c.r - d * d);
        s.push back(p + normalize(l.v) * d);
        s.push back (p - normalize(l.v) * d);
        return s;
vector<pt> inter_cc(circle c1, circle c2) {
        pt dir = c2.center - c1.center;
        1f d2 = dis2(c1.center, c2.center);
        if(d2 \le E0) {
                //assert(fabsl(c1.r - c2.r) > E0);
                return {};
        1f td = 0.5L * (d2 + c1.r * c1.r - c2.r * c2.r)
        1f h2 = c1.r * c1.r - td / d2 * td;
        pt p = c1.center + dir * (td / d2);
        if(fabsl( h2 ) < EPS) return {p};</pre>
        if(h2 < 0.0L) return {};
        pt dir h = rot 90 (dir) * sqrtl(h2 / d2);
        return {p + dir_h, p - dir_h};
```

```
5.3 Circulos
```

```
//compute intersection of line through points a and b
//circle centered at c with radius r > 0
vector<pt> circle_line_intersection(pt c, lf r, pt a, pt
   vector<pt> ret;
    b = b - a; a = a - c;
   lf A = dot(b, b), B = dot(a, b);
   lf C = dot(a, a) - r * r, D = B * B - A * C;
    if (D < -EPS) return ret;</pre>
    ret.push back(c + a + b * (-B + sqrt(D + EPS)) / A);
    if (D > EPS) ret.push back(c + a + b * (-B - sqrt(D))
        / A);
    return ret:
// circle-line inter = 1, inner: 1 = 0x0 \ 0 = 0=0
vector<pair<pt, pt>> tangents(circle c1, circle c2, bool
   inner) {
        vector<pair<pt, pt>> out;
        if (inner) c2.r = -c2.r; // inner tangent
        pt d = c2.center - c1.center;
        double dr = c1.r - c2.r, d2 = norm2(d), h2 = d2 - c2.r
            dr * dr;
        if (d2 == 0 || h2 < 0) { assert(h2 != 0); return
           {}; } // (identical)
        for (double s : {-1, 1}) {
                pt v = (d * dr + rot 90(d) * sqrt(h2) * s)
                    / d2;
                out.push_back({c1.center + v * c1.r, c2.
                   center + v * c2.r);
        return out; // if size 1: circle are tangent
// circle targent passing through pt p
pair<pt, pt> tangent_through_pt(circle c, pt p){
        pair<pt, pt> out;
        double d = norm2(p - c.center);
        if (d < c.r) return {};
        pt base = c.center - p;
        double w = sqrt(norm2(base) - c.r * c.r);
        pt a = \{w, c.r\}, b = \{w, -c.r\};
        pt s = p + base * a / norm2(base) * w;
        pt t = p + base * b / norm2(base) * w;
        out = \{s, t\};
        return out;
lf safeAcos(lf x) {
        if (x < -1.0) x = -1.0;
        if (x > 1.0) x = 1.0;
        return acos(x);
lf areaOfIntersectionOfTwoCircles(circle c1, circle c2){
```

```
1f r1 = c1.r, r2 = c2.r, d = dis(c1.center, c2.
                          center);
                  if(d \ge r1 + r2) return 0.0L;
                  if(d \le fabs1(r2 - r1)) return PI * (r1 < r2 ? r1
                             * r1 : r2 * r2);
                  lf alpha = safeAcos((r1 * r1 - r2 * r2 + d * d) /
                             (2.0L * d * r1));
                  lf betha = safeAcos((r2 * r2 - r1 * r1 + d * d) / range = ra
                             (2.0L * d * r2));
                  lf a1 = r1 * r1 * (alpha - sinl(alpha) * cosl(
                          alpha));
                  lf a2 = r2 * r2 * (betha - sinl(betha) * cosl(
                          betha));
                  return a1 + a2;
} ;
lf intertriangle(circle& c, pt a, pt b){ // area of
        intersection with oab
                  if(abs(cross((c.center - a), (c.center - b))) <=</pre>
                          EPS) return 0.;
                  vector<pt> q = \{a\}, w = inter_cl(c, line(a, b));
                  if(w.size() == 2) for(auto p: w) if(dot((a - p),
                           (b - p)) < -EPS) q.push back(p);
                  q.push back(b);
                  if(q.size() == 4 \&\& dot((q[0] - q[1]), (q[2] - q[1]))
                           [1]) > EPS) swap(q[1], q[2]);
                  lf s = 0;
                  for(int i = 0; i < q.size() - 1; ++i){}
                                    if(!c.contains(q[i]) || !c.contains(q[i +
                                              1])) s += c.r * c.r * min angle((g[i])
                                              - c.center), q[i+1] - c.center) / 2;
                                    else s += abs(cross((q[i] - c.center), (q
                                            [i + 1] - c.center) / 2);
                  return s;
bool circumcircle contains(vector<pt> tr, pt D) { //
        triange CCW
    pt A = tr[0] - D, B = tr[1] - D, C = tr[2] - D;
    lf norm_a = norm2(tr[0]) - norm2(D);
    lf norm b = norm2(tr[1]) - norm2(D);
    If norm c = norm2(tr[2]) - norm2(D);
    lf det1 = A.x * (B.y * norm c - norm b * C.y);
    lf det2 = B.x * (C.y * norm_a - norm_c * A.y);
    lf det3 = C.x * (A.y * norm_b - norm_a * B.y);
    return det1 + det2 + det3 > E0;
// r[k]: area covered by at least k circles
// O(n^2 \log n) (high constant)
vector<lf> intercircles(vector<circle> c){
                  vector<lf> r(c.size() + 1);
                  for(int i = 0; i < c.size(); ++i){</pre>
```

```
int k = 1; pt 0 = c[i].center;
        vector<pair<pt, int>> p = {
                 \{c[i].center + pt(1,0) * c[i].r,
                 \{c[i].center - pt(1,0) * c[i].r,
                    0 } };
        for(int j = 0; j < c.size(); ++j) if(j !=
             i) {
                bool b0 = c[i].in(c[j]), b1 = c[j]
                    ].in(c[i]);
                if (b0 && (!b1 || i < j)) ++k;
                else if(!b0 && !b1){
                         auto v = inter_cc(c[i], c
                             [ j ] );
                         if(v.size() == 2){
                                 swap(v[0], v[1]);
                                 p.push_back({v
                                     [0], 1});
                                 p.push back({v
                                     [1], -1\});
                                 if (polar_cmp(v[1])
                                      - 0, v[0] - 0
                                     )) ++k;
        sort(all(p), [&](auto& a, auto& b){
            return polar cmp(a.first - 0, b.first
            - O); } j;
        for(int j = 0; j < p.size(); ++j){</pre>
                pt p0 = p[j ? j - 1 : p.size()
                    -1].first, p1 = p[j].first;
                lf a = min\_angle((p0 - c[i]).
                    center), (p1 - c[i].center));
                r[k] += (p0.x - p1.x) * (p0.y + p1.x)
                    p1.y) / 2 + c[i].r * c[i].r *
                    (a - \sin(a)) / 2;
                k += p[j].second;
return r:
```

5.4 Closest Points

```
// O(nlogn)
pair<pt, pt> closest_points(vector<pt> v) {
    sort(v.begin(), v.end());
    pair<pt, pt> ans;
    lf d2 = INF;

function<void( int, int )> solve = [&](int l, int r) {
```

```
if(l == r) return;
        int mid = (1 + r) / 2;
        lf x \text{ mid} = v[\text{mid}].x;
        solve(1, mid);
        solve (mid + 1, r);
        vector<pt> aux;
        int p1 = 1, p2 = mid + 1;
        while (p1 <= mid && p2 <= r) {
                 if (v[p1].y < v[p2].y) aux.
                    push back (v[p1++]);
                 else aux.push back(v[p2++]);
        while(p1 <= mid) aux.push back(v[p1++]);</pre>
        while (p2 \le r) aux.push back (v[p2++]);
        vector<pt> nb;
        for(int i = 1; i <= r; ++i) {
        v[i] = aux[i - 1];
        lf dx = (x_mid - v[i].x);
        if(dx * dx < d2)
                 nb.push_back(v[i]);
        for(int i = 0; i < (int) nb.size(); ++i){</pre>
        for(int k = i + 1; k < (int) nb.size();</pre>
            ++k) {
                 lf dy = (nb[k].y - nb[i].y);
                 if(dy * dy > d2) break;
                 lf nd2 = dis2(nb[i], nb[k]);
                 if(nd2 < d2) d2 = nd2, ans = {nb[}
                    il, nb[k]};
solve(0, v.size() -1);
return ans;
```

5.5 Convex Hull

5.6 Delaunay

```
// Returns planar graph representing Delaunay's
   triangulation.
// Edges for each vertex are in ccw order.
// Voronoi vertices = the circumcenters of the Delaunay
   triangles.
// O(nlogn)
typedef struct OuadEdge* O;
struct QuadEdge {
        int id, used;
        pt o;
        O rot, nxt;
        QuadEdge(int id=-1, pt o=pt(INF,INF)):id(id),used
            (0), o(o), rot(0), nxt(0) {}
        Q rev() {return rot->rot; }
        Q next() {return nxt;}
        Q prev() {return rot->next()->rot; }
        pt dest() {return rev() ->o; }
};
Q edge(pt a, pt b, int ida, int idb) {
        Q e1=new QuadEdge(ida,a);
        Q e2=new QuadEdge(idb,b);
        O e3=new OuadEdge;
        O e4=new OuadEdge;
        tie(e1->rot,e2->rot,e3->rot,e4->rot)={e3,e4,e2,e1
        tie(e1->nxt,e2->nxt,e3->nxt,e4->nxt) = \{e1,e2,e4,e3
        return e1;
void splice(Q a, Q b){
        swap(a->nxt->rot->nxt,b->nxt->rot->nxt);
        swap (a->nxt,b->nxt);
void del_edge(Q& e, Q ne) {
```

```
splice(e,e->prev()); splice(e->rev(),e->rev()->
           prev());
        delete e->rev()->rot; delete e->rev();
        delete e->rot; delete e;
        e=ne;
Q conn(Q a, Q b) {
        Q = e = e d = (a - d = st(), b - so, a - s = v() - s = d, b - s = d);
        splice(e, a->rev()->prev());
        splice(e->rev(),b);
        return e;
auto area(pt p, pt q, pt r) { return cross((q-p), (r-q)); }
bool circumcircle_contains(vector<pt> tr, pt D) {
        if (orient(tr[0], tr[1], tr[2]) < 0) reverse(all(
    pt A = tr[0] - D, B = tr[1] - D, C = tr[2] - D;
    lf norm a = norm2(tr[0]) - norm2(D);
    lf norm b = norm2(tr[1]) - norm2(D);
    lf norm c = norm2(tr[2]) - norm2(D);
    lf det1 = A.x * (B.y * norm c - norm b * C.y);
    lf det2 = B.x * (C.y * norm_a - norm_c * A.y);
    lf det3 = C.x * (A.y * norm_b - norm_a * B.y);
    return det1 + det2 + det3 > 0;
pair<Q,Q> build tr(vector<pt>& p, int 1, int r){
        if(r-1+1<=3){
                Q a=edge(p[1], p[1+1], 1, 1+1), b=edge(p[1
                    +1],p[r],l+1,r);
                if(r-1+1==2) return {a,a->rev()};
                splice(a->rev(),b);
                auto ar=area(p[1],p[1+1],p[r]);
                Q c=abs(ar)>EPS?conn(b,a):0;
                if(ar>=-EPS) return {a,b->rev()};
                return {c->rev(),c};
        int m = (1+r)/2;
        auto [la,ra]=build tr(p,l,m);
        auto [lb,rb]=build_tr(p,m+1,r);
        while (1) {
                if(orient(lb->o,ra->o, ra->dest()) > 0)
                    ra=ra->rev()->prev();
                else if(orient(lb->o,ra->o,lb->dest()) >
                    0) lb=lb->rev()->next();
                else break;
        0 b=conn(lb->rev(),ra);
        auto valid=[&](Q e) {return orient(e->dest(),b->
            dest(), b->0) > 0; ;
        if(ra->o==la->o) la=b->rev();
```

```
if(lb->o==rb->o) rb=b;
        while(1){
                 0 L=b->rev()->next();
                if(valid(L)) while(circumcircle contains
                    ({b->dest(),b->o,L->dest()},L->next()
                    ->dest())) del edge(L,L->next());
                 0 R=b->prev();
                if(valid(R)) while(circumcircle contains
                    ({b->dest(),b->o,R->dest()},R->prev()
                    ->dest())) del edge(R,R->prev());
                if(!valid(L)&&!valid(R)) break;
                 if(!valid(L) | | (valid(R) &&
                    circumcircle contains({L->dest(),L->o,
                    R->o, R->dest()))) b=conn(R,b->rev());
                 else b=conn(b->rev(),L->rev());
        return {la,rb};
vector<vector<int>> delaunay(vector<pt> v) {
        int n=v.size(); auto tmp=v;
        vector<int> id(n); iota(all(id),0);
        sort(all(id),[&](int l, int r){return v[l]<v[r</pre>
        for(int i = 0; i < n; ++i) v[i]=tmp[id[i]];</pre>
        assert (unique (all (v)) == v.end());
        vector<vector<int>> q(n);
        int col=1;
        for(int i = 2; i < n; ++i) col &= abs(area(v[i], v</pre>
            [i-1], v[i-2])) <= EPS;
        if(col){
                 for (int i = 1; i < n; i++) q[id[i-1]].pb(
                    id[i]),q[id[i]].pb(id[i-1]);
        else{
                 Q e=build tr(v, 0, n-1).first;
                vector<0> edg={e};
                 for(int i=0;i<edq.size();e=edg[i++]){</pre>
                         for(Q at=e;!at->used;at=at->next
                             ()){
                                 at->used=1;
                                 g[id[at->id]].pb(id[at->
                                     rev()->id]);
                                 edq.pb(at->rev());
        return q;
```

5.7 Halfplanes

```
const lf INF = 1e100;
struct Halfplane {
```

```
pt p, pq; // p: point on line, pq: dir, take left
        lf angle;
        Halfplane(){}
        Halfplane(pt& a, pt& b): p(a), pq(b - a) {
                angle = atan21(pq.y, pq.x);
        bool out(const pt& r) { return cross(pq, r - p) <</pre>
            -EPS; } // checks if p is inside the half plane
        bool operator < (const Halfplane& e) const {</pre>
            return angle < e.angle; }</pre>
};
// intersection pt of the lines of 2 halfplanes
pt inter(const Halfplane& s, const Halfplane& t) {
        if (abs(cross(s.pq, t.pq)) <= EPS) return {INF,</pre>
        If alpha = cross((t.p - s.p), t.pq) / cross(s.pq,
            t.pa);
        return s.p + (s.pq * alpha);
// O(nlogn) return CCW polygon
vector<pt> hp intersect(vector<Halfplane>& H) {
        pt box[4] = \{pt(INF, INF), pt(-INF, INF), pt(-INF)\}
            , -INF), pt(INF, -INF)};
        for(int i = 0; i < 4; ++i ) {</pre>
                Halfplane aux(box[i], box[(i + 1) % 4]);
                H.push back(aux);
        sort(H.begin(), H.end());
        deque < Halfplane > dq;
        int len = 0;
        for(int i = 0; i < int(H.size()); ++i){</pre>
                while (len > 1 && H[i].out(inter(dq[len -
                     1], dq[len - 2]))){
                         dq.pop_back();
                         --len:
                while (len > 1 && H[i].out(inter(dq[0],
                    dq[1]))){
                         dq.pop_front();
                         --len;
                if (len > 0 && fabsl(cross(H[i].pq, dq[
                    len - 1].pq)) < EPS){
                         if (dot(H[i].pq, dq[len - 1].pq)
                             < 0.0) return vector<pt>();
                         if (H[i].out(dq[len - 1].p)) {
                                 dq.pop back();
                                 --len:
                         } else continue;
```

5

```
dq.push back(H[i]);
                                                                           ++len;
                                     while (len > 2 && dq[0].out(inter(dq[len - 1], dq
                                                      [len - 2]))
                                                                           dq.pop_back();
                                                                           --len;
                                     while (len > 2 && dq[len - 1].out(inter(dq[0], dq
                                                      [1]))){
                                                                           dq.pop front();
                                                                           --len;
                                     if (len < 3) return vector<pt>();
                                     vector<pt> ret(len);
                                      for(int i = 0; i + 1 < len; ++i) ret[i] = inter(</pre>
                                                     dq[i], dq[i + 1]);
                                     ret.back() = inter(dq[len - 1], dq[0]);
                                      // remove repeated points if needed
                                     return ret;
// intersection of halfplanes
vector<pt> hp_intersect(vector<halfplane>& b) {
                                    vector<pt> box = \{\{\inf, \inf\}, \{-\inf\}, \{-
                                                          -inf}, {inf, -inf}};
                                     for (int i = 0; i < 4; i++) {
                                                                          b.push back(\{box[i], box[(i + 1) % 4]\});
                                     sort(b.begin(), b.end());
                                     int n = b.size(), q = 1, h = 0;
                                     vector<halfplane> c(n + 10);
                                     for(int i = 0; i < n; i++) {</pre>
                                                                          while (q < h \& \& b[i] . out (inter(c[h], c[h -
                                                                                              1]))) h--;
                                                                           while (q < h \&\& b[i].out(inter(c[q], c[q +
                                                                                              11))) g++;
                                                                           c[++h] = b[i];
                                                                           if(q < h \&\& abs(cross(c[h].pq, c[h-1].pq))
                                                                                                                if (dot (c[h].pq, c[h - 1].pq) <=
                                                                                                                                0) return {};
                                                                                                                if (b[i].out(c[h].p)) c[h] = b[i];
                                     while (q < h - 1 \&\& c[q].out(inter(c[h], c[h - 1]))
                                                    )) h--;
```

```
while (q < h - 1 \&\& c[h].out(inter(c[q], c[q + 1]))
              )) q++;
          if(h - q <= 1) return {};
          c[h + 1] = c[a];
          vector<pt> s;
          for(int i = q; i < h + 1; i++) s.pb(inter(c[i], c
              [i + 1]);
          return s:
5.8 KD Tree
  const ll INF = 2e18;
  const int D = 2; // dimension
  struct ptd{
          int p[D];
          bool operator !=(const ptd &a) const {
                  bool ok = 1;
                   for (int i = 0; i < D; i++) ok &= (p[i] ==
                       a.p[i]);
                  return !ok;
  } ;
  struct kd node{
          ptd p;
          int axis;
          kd node *left, *right;
```

} **;**

};

struct cmp points {

int axis;

11 dis2(ptd a, ptd b) {

struct KDTree{

ll ans = 0;

return ans;

vector<ptd> arr;

kd node* root;

cmp points(){}

const {

cmp points(int x): axis(x){}

bool operator () (const ptd &a, const ptd &b)

return a.p[axis] < b.p[axis];</pre>

]) * 1ll * (a.p[i] - b.p[i]);

KDTree(vector<ptd> &vptd): arr(vptd) {

build(root, 0, sz(vptd) - 1);

for (int i = 0; i < D; i++) ans += (a.p[i] - b.p[i]

```
// O(nlogn)
   void build(kd_node* &node, int 1, int r) {
            if(1 > r) {
                    node = nullptr;
                    return;
            node = new kd node();
            if(1 == r) {
                    node -> p = arr[l];
                    node->left = nullptr;
                    node->right = nullptr;
                    return;
            11 \text{ bAxis} = 0;
            11 \text{ mRange} = 0;
            for (int axis = 0; axis < D; ++axis) {
                    11 minVal = INF, maxVal = -INF;
                    for (int i = 1; i <= r; ++i) {
                             minVal = min(minVal, (ll))
                                arr[i].p[axis]);
                             maxVal = max(maxVal, (11)
                                arr[i].p[axis]);
                    if (maxVal - minVal > mRange) {
                             mRange = maxVal - minVal;
                             bAxis = axis;
            int mid = (1 + r) / 2;
            nth_element(arr.begin() + l, arr.begin()
               + mid, arr.begin() + r + 1, cmp points
                (bAxis));
            node->p = arr[mid];
            node->axis = bAxis;
            build(node->left, l, mid);
            build(node->right, mid + 1, r);
   void nearest(kd node* node, ptd q, pair<11, ptd>
       &ans){
            if(node == NULL) return;
            if(node->left == NULL && node->right ==
               NULL) {
                    if(!(q != node->p)) return; //
                        avoid query point as answer
                    if (ans.first > dis2(node->p, q))
                         ans = \{dis2(node->p, q), node\}
                        ->p};
                    return;
```

```
int axis = node->axis;
                int value = node->p.p[axis];
                if(q.p[axis] <= value) {</pre>
                         nearest(node->left, q, ans);
                         ll diff = value - q.p[axis];
                         if(diff * diff <= ans.ff) nearest</pre>
                             (node->right, q, ans);
                }else{
                         nearest(node->right, q, ans);
                         11 diff = q.p[axis] - value;
                         if(diff * diff <= ans.ff) nearest</pre>
                             (node->left, q, ans);
    // O(logn) Returns {squared distance, nearest point}
        pair<ll, ptd> nearest(ptd q){
                pair<ll, ptd> ans = {INF, ptd()};
                nearest(root, q, ans);
                return ans;
};
```

5.9 Lineas

```
// add points operators
struct line {
        pt v; lf c; // v: dir, c: mov y
        line(pt v, lf c) : v(v), c(c) {}
        line(lf a, lf b, lf c) : v({b, -a}), c(c) {} //
           ax + by = c
        line(pt p, pt q): v(q - p), c(cross(v, p)) {}
        bool operator < (line 1) { return cross(v, 1.v) >
           0; }
        bool operator == (line 1) { return (abs(cross(v, 1)))
           (v) <= E0) && c == 1.c; } // abs(c) == abs(1.
        lf side(pt p) { return cross(v, p) - c; }
        lf dist(pt p) { return abs(side(p)) / norm(v); }
        lf dist2(pt p) { return side(p) * side(p) / (lf)
           norm2(v); }
        line perp through (pt p) { return {p, p + rot90(v)
           }; } // line perp to v passing through p
        bool cmp_proj(pt p, pt q) { return dot(v, p) < dot</pre>
            (v, q); } // order for points over the line
        // use: auto fsort = [&l1] (const pt &a, const pt
           &b) { return 11.cmp_proj(a, b); };
        line translate(pt t) { return {v, c + cross(v, t)
           }; }
        line shift_left(lf d) { return {v, c + d*norm(v)};
```

```
pt proj(pt p) { return p - rot90(v) * side(p) /
           norm2(v); } // pt provected on the line
        pt refl(pt p) { return p - rot 90(v) * 2 * side(p)
           / norm2(v); } // pt reflected on the other
           side of the line
        bool has (pt p) { return abs (cross (v, p) - c) <= E0
           ; }; // pt on line
        lf evalx(lf x) {
                assert (fabsl(v.x) > EPS);
                return (c + v.v * x) / v.x;
};
pt inter_ll(line l1, line l2) {
        if (abs(cross(11.v, 12.v)) <= EPS) return {INF,
           INF \; // parallel
        return (12.v * 11.c - 11.v * 12.c) / cross(11.v,
           12.v); // floating points
// bisector divides the angle in 2 equal angles
// interior line goes on the same direction as 11 and 12
line bisector(line 11, line 12, bool interior) {
        // assert (cross(11.v, 12.v) != 0); // 11 and 12
           cannot be parallel
        lf sign = interior ? 1 : -1;
        return {12.v / norm(12.v) + 11.v / norm(11.v) *
           sian,
                        12.c / norm(12.v) + 11.c / norm(
                           11.v) * sign};
```

5.10 Manhattan

```
});
    map<11, 11, greater<11>> active; // (xs, id)
    for(auto i : ids){
                     for (auto it = active.lower_bound(
                        ps[i].x); it != active.end();
                        active.erase(it++)){
            11 j = it->second;
            if (ps[i].x - ps[i].y > ps[j].x - ps[j].y
                ) break;
            assert(ps[i].x \rightarrow= ps[j].x && ps[i].y \rightarrow=
                ps[j].y);
            edges.push_back(\{(ps[i].x - ps[j].x) + (
                ps[i].y - ps[j].y), i, j);
        active[ps[i].x] = i;
    for (auto &p : ps) { // rotate
        if (rot & 1) p.x *=-1;
        else swap(p.x, p.y);
return edges;
```

5.11 Min Circle

```
// minimo circulo que encierra todos los puntos
// Promedio: O(n), Peor: O(n^2)
Circle min circle(vector<pt> v) {
        random shuffle(v.begin(), v.end()); // shuffle(
           all(vec), rng);
        auto f2 = [\&] (int a, int b) {
                Circle ans(v[a], v[b]);
                for(int i = 0; i < a; ++ i)
                if (ans.contains(v[i]) == OUT) ans =
                    Circle(v[i], v[a], v[b]);
                return ans;
        };
        auto f1 = [&] ( int a ) {
                Circle ans (v[a], 0.0L);
                for(int i = 0; i < a; ++i)
                if(ans.contains(v[i]) == OUT) ans = f2(i)
                    , a );
                return ans:
        };
        Circle ans (v[0], 0.0L);
        for(int i = 1; i < (int) v.size(); ++i)</pre>
                if(ans.contains(v[i]) == OUT ) ans = f1(i
                   );
```

```
return ans;
}
```

5.12 Puntos

```
typedef long double lf;
const lf EPS = 1e-9;
const 1f E0 = 0.0L; //Keep = 0 for integer coordinates,
   otherwise = EPS
const lf PI = acos(-1);
struct pt {
        lf x, y;
        pt(){}
        pt(lf a, lf b): x(a), y(b) {}
        pt(lf ang): x(cos(ang)), y(sin(ang)){} // Polar
           unit point: ang(RAD)
        pt operator - (const pt &q) const { return {x - q
            .x , y - q.y \}; }
        pt operator + (const pt &q) const { return {x + q
            .x , y + q.y \}; }
        pt operator * (pt p) { return {x * p.x - y * p.y,
           x * p.y + y * p.x;
        pt operator * (const lf &t) const { return {x * t
            , y * t }; }
        pt operator / (const lf &t) const { return {x / t
             , y / t }; }
        bool operator == (pt p) { return abs(x - p.x) <=
           EPS && abs(y - p.y) <= EPS; }
        bool operator != (pt p) { return !operator==(p); }
        bool operator < (const pt & q) const { // set /
           sort
                if (fabsl(x - q.x) > E0) return x < q.x;
                return y < q.y;</pre>
        void print() { cout << x << " " << y << "\n"; }</pre>
};
pt normalize(pt p) {
        lf norm = hypotl(p.x, p.y);
        if(fabsl(norm) > EPS) return {p.x /= norm, p.y /=
             norm };
        else return p;
int cmp(lf a, lf b) { return (a + EPS < b ? -1 : (b + EPS <</pre>
    a? 1:0)); } // float comparator
// rota ccw
pt rot90(pt p) { return {-p.y, p.x}; }
// w(RAD)
pt rot(pt p, lf w) { return {cosl(w) * p.x - sinl(w) * p.y
   *, sinl(w) * p.x + cosl(w) * p.y); }
lf norm2(pt p) { return p.x * p.x + p.y * p.y; }
```

```
lf norm(pt p) { return hypotl(p.x, p.y); }
lf dis2(pt p, pt q) { return norm2(p - q); }
lf dis(pt p, pt q) { return norm(p - q); }
If arg(pt a) {return atan2(a.y, a.x); } // ang(RAD) a x-
If dot(pt a, pt b) { return a.x * b.x + a.y * b.y; } // x
       = 90 -> cos = 0
lf cross(pt a, pt b) { return a.x * b.y - a.y * b.x; } //
       x = 180 -> \sin = 0
lf orient(pt a, pt b, pt c) { return cross(b - a, c - a);
       } // AB clockwise = -
int sign(lf x) { return (EPS < x) - (x < -EPS); }
// p inside angle abc (center in a)
bool in_angle(pt a, pt b, pt c, pt p) {
                  //assert(fabsl(orient(a, b, c)) > E0);
                  if(orient(a, b, c) < -E0)
                                    return orient(a, b, p) >= -E0 || orient(a
                                            , c, p) <= E0;
                  return orient(a, b, p) >= -E0 && orient(a, c, p)
                          \leq E0;
lf min_angle(pt a, pt b) { return acos(max((lf)-1.0, min())
       lf) 1.0, dot(a, b) /norm(a) /norm(b)))); } // ang(RAD)
lf angle(pt a, pt b) { return atan2(cross(a, b), dot(a, b)
       ); } // ang(RAD)
lf angle(pt a, pt b, pt c){ // ang(RAD) AB AC ccw
                  lf ang = angle(b - a, c - a);
                  if (ang < 0) ang += 2 * PI;
                  return ang;
bool half(pt p) { // true if is in (0, 180] (line is x
       axis)
                  // assert(p.x != 0 || p.y != 0); // the argument
                          of (0, 0) is undefined
                  return p.y > 0 | | (p.y == 0 && p.x < 0);
bool half_from(pt p, pt v = \{1, 0\}) {
                  return cross(v,p) < 0 \mid \mid (cross(v,p) == 0 && dot(v,p) == 0 & dot(v,p) & 
                         v, p) < 0;
// polar sort
bool polar cmp(const pt &a, const pt &b) {
                  return make tuple(half(a), 0) < make tuple(half(b)
                         ), cross(a,b));
void polar_sort(vector<pt> &v, pt o) { // sort points in
       counterclockwise with respect to point o
                  sort(v.begin(), v.end(), [&](pt a,pt b) {
                                    return make_tuple(half(a - o), 0.0, norm2
```

```
((a - o))) < make tuple(half(b - o),
                    cross(a - o, b - o), norm2((b - o)));
        });
int cuad(pt p) { // REVISAR
        if(p.x > 0 \&\& p.y >= 0) return 0;
        if(p.x <= 0 && p.y > 0) return 1;
        if(p.x < 0 && p.y <= 0) return 2;
        if (p.x >= 0 \&\& p.v < 0) return 3;
        return -1; //x == 0 \&\& v == 0
bool cmp(pt p1, pt p2) {
        int c1 = cuad(p1), c2 = cuad(p2);
        return c1 == c2 ? p1.y * p2.x < p1.x * p2.y : c1
           < c2;
// O(n*2^d*d)
// Return the max manhattan distance between points with
   d-dimension.
ll max_distance_manhattan(vector<vi> p, int d) {
        long long ans = 0;
        for (int msk = 0; msk < (1 << d); msk++) {
                long long mx = LLONG MIN, mn = LLONG MAX;
                for (int i = 0; i < n; i++) {</pre>
                        long long cur = 0;
                        for (int j = 0; j < d; j++) {
                                 if (msk & (1 << j)) cur
                                    += p[i][j];
                                 else cur -= p[i][j];
                        mx = max(mx, cur);
                        mn = min(mn, cur);
                ans = max(ans, mx - mn);
        return ans:
ll sd_to_ll(string num, int canDec = 6) {
        string nnum = "";
        bool ok = 0;
        for (int i = 0; i < sz(num); i++) {
                if (num[i] == '.') {
                        ok = 1:
                        continue;
                if (ok) canDec--;
                nnum.pb(num[i]);
        while(canDec--) nnum.pb('0');
        return stoll(nnum);
```

5.13 Poligonos

```
// add Points Lines Segments Circles
// points in polygon(vector<pt>) ccw or cw
enum {OUT, IN, ON};
lf area(vector<pt>& p) {
        lf r = 0.;
        for(int i = 0, n = p.size(); i < n; ++i){</pre>
                r += cross(p[i], p[(i + 1) % n]);
        return r / 2; // negative if CW, positive if CCW
lf perimeter(vector<pt>& p) {
        lf per = 0;
        for (int i = 0, n = p.size(); i < n; ++i) {
                 per += norm(p[i] - p[(i + 1) % n]);
        return per;
bool is convex(vector<pt>& p) {
        bool pos = 0, neg = 0;
        for (int i = 0, n = p.size(); i < n; i++) {</pre>
                 int o = orient(p[i], p[(i + 1) % n], p[(i + 1) % n]
                     + 2) % n]);
                 if (o > 0) pos = 1;
                 if (o < 0) neg = 1;
        return ! (pos && neg);
int point_in_polygon(vector<pt>& pol, pt& p) {
        int wn = 0;
        for(int i = 0, n = pol.size(); i < n; ++i) {
                 If c = orient(p, pol[i], pol[(i + 1) % n
                 if(fabsl(c) <= E0 && dot(pol[i] - p, pol
                    [(i + 1) % n] - p) <= E0) return ON;
                    // on segment
                 if(c > 0 && pol[i].y <= p.y + E0 && pol[(
                i + 1) % n].y - p.y > E0) ++wn;
if(c < 0 && pol[(i + 1) % n].y <= p.y +
                    E0 && pol[i].y - p.y > E0) --wn;
        return wn ? IN : OUT;
// O(logn) polygon CCW, remove collinear
int point_in_convex_polygon(const vector<pt> &pol, const
   pt &p){
```

```
int low = 1, high = pol.size() - 1;
        while (high - low > 1) {
                int mid = (low + high) / 2;
                if(orient(pol[0], pol[mid], p) >= -E0)
                    low = mid;
                else high = mid;
        if(orient(pol[0], pol[low], p) < -E0) return OUT;</pre>
        if (orient (pol[low], pol[high], p) < -E0) return</pre>
        if(orient(pol[high], pol[0], p) < -E0) return OUT</pre>
        if(low == 1 \&\& orient(pol[0], pol[low], p) <= E0)
        if(orient(pol[low], pol[high], p) <= E0) return</pre>
        if(high == (int) pol.size() -1 && orient(pol[high
            ], pol[0], p) <= E0) return ON;
        return IN;
// convex polygons in some order (CCW, CW)
vector<pt> minkowski(vector<pt> P, vector<pt> Q) {
        rotate(P.begin(), min element(P.begin(), P.end())
            , P.end());
        rotate(Q.begin(), min_element(Q.begin(), Q.end())
           , Q.end());
        P.push back (P[0]), P.push back (P[1]);
        Q.push_back(Q[0]), Q.push_back(Q[1]);
        vector<pt> ans;
        size t i = 0, j = 0;
        while(i < P.size() - 2 || j < Q.size() - 2) {
                ans.push_back(P[i] + Q[j]);
                lf dt = cross(P[i + 1] - P[i], Q[j + 1] -
                if(dt \ge E0 \&\& i < P.size() - 2) ++i;
                if(dt \le E0 \&\& j < Q.size() - 2) ++j;
        return ans;
pt centroid(vector<pt>& p) {
        pt c{0, 0};
        If scale = 6. * area(p);
        for (int i = 0, n = p.size(); i < n; ++i){</pre>
                c = c + (p[i] + p[(i + 1) % n]) * cross(p)
                    [i], p[(i + 1) % n]);
        return c / scale;
void normalize(vector<pt>& p) { // polygon CCW
        int bottom = min element(p.begin(), p.end()) - p.
           begin();
```

```
vector<pt> tmp(p.begin() + bottom, p.end());
        tmp.insert(tmp.end(), p.begin(), p.begin()+bottom
           );
        p.swap(tmp);
        bottom = 0;
void remove col(vector<pt>& p) {
        vector<pt> s;
        for(int i = 0, n = p.size(); i < n; i++){</pre>
                if(!on\_segment(p[(i - 1 + n) % n], p[(i + n) % n]))
                    1) % n], p[i])) s.push_back(p[i]);
        p.swap(s);
void delete repetead(vector<pt>& p) {
        vector<pt> aux;
        sort(p.begin(), p.end());
        for (pt &pi : p) {
                if (aux.empty() || aux.back() != pi) aux.
                   push back(pi);
        p.swap(aux);
pt farthest (vector<pt>& p, pt v) { // O(log(n)) only
   CONVEX, v: dir
        int n = p.size();
        if(n < 10){
                int k = 0;
                for (int i = 1; i < n; i++) if (dot (v, (p[i
                   | - p[k])) > EPS) k = i;
                return p[k];
        pt a = p[1] - p[0];
        int s = 0, e = n, ua = dot(v, a) > EPS;
        if(!ua && dot(v, (p[n-1] - p[0])) <= EPS)
           return p[0];
        while (1) {
                int m = (s + e) / 2;
                pt c = p[(m + 1) % n] - p[m];
                int uc = dot(v, c) > EPS;
                if(!uc && dot(v, (p[(m - 1 + n) % n] - p[
                    m])) <= EPS) return p[m];
                if(ua && (!uc || dot(v, (p[s] - p[m])) >
                    EPS)) e = m;
                else if(ua || uc || dot(v, (p[s] - p[m]))
                    >= -EPS) s = m, a = c, ua = uc;
                else e = m;
                assert (e > s + 1);
vector<pt> cut(vector<pt>& p, line l) {
        // cut CONVEX polygon by line 1
```

```
// returns part at left of l.pg
        vector<pt> q;
        for(int i = 0, n = p.size(); i < n; i++) {</pre>
                int d0 = sign(l.side(p[i]));
                int d1 = sign(1.side(p[(i + 1) % n]));
                if(d0 >= 0) q.push back(p[i]);
                line m(p[i], p[(i + 1) % n]);
                if(d0 * d1 < 0 && !(abs(cross(l.v, m.v)))
                    \leq EPS)){
                         q.push back((inter ll(l, m)));
        return q;
// O(n)
vector<pair<int, int>> antipodal(vector<pt>& p) {
        vector<pair<int, int>> ans;
        int n = p.size();
        if (n == 2) ans.push_back({0, 1});
        if (n < 3) return ans;</pre>
        auto nxt = [\&] (int x) \{ return (x + 1 == n ? 0 : x = n ) \}
             + 1); };
        auto area2 = [&] (pt a, pt b, pt c) { return cross(
           b - a, c - a); };
        int b0 = 0;
        while (abs (area2 (p[n - 1], p[0], p[nxt (b0)])) >
            abs (area2(p[n - 1], p[0], p[b0]))) ++b0;
        for (int b = b0, a = 0; b != 0 && a <= b0; ++a) {
                ans.push_back({a, b});
                while (abs(area2(p[a], p[nxt(a)], p[nxt(b
                    )])) > abs(area2(p[a], p[nxt(a)], p[b
                    ]))){
                         b = nxt(b);
                         if (a != b0 || b != 0) ans.
                            push back({a, b});
                         else return ans;
                if (abs(area2(p[a], p[nxt(a)], p[nxt(b)])
                    ) == abs(area2(p[a], p[nxt(a)], p[b]))
                    ) {
                         if (a != b0 \mid | b \mid = n - 1) ans.
                            push back({a, nxt(b)});
                         else ans.push back({nxt(a), b});
        return ans;
// O(n)
// square distance of most distant points, prereq: convex
   , ccw, NO COLLINEAR POINTS
lf callipers(vector<pt>& p) {
        int n = p.size();
        lf r = 0;
```

```
for (int i = 0, j = n < 2 ? 0 : 1; <math>i < j; ++i) {
                 for (;;\dot{j} = (\dot{j} + 1) % n) {
                         r = max(r, norm2(p[i] - p[j]));
                         if(cross((p[(i + 1) % n] - p[i]),
                              (p[(j + 1) % n] - p[j])) <=
                             EPS) break;
        return r;
// O(n + m) max_dist between 2 points (pa, pb) of 2
   Convex polygons (a, b)
lf rotating callipers(vector<pt>& a, vector<pt>& b) { //
        if (a.size() > b.size()) swap(a, b); // <- del or
        pair<ll, int > start = \{-1, -1\};
        if(a.size() == 1) swap(a, b);
        for(int i = 0; i < a.size(); i++) start = max(</pre>
            start, \{norm2(b[0] - a[i]), i\});
        if(b.size() == 1) return start.first;
        lf r = 0;
        for(int i = 0, j = start.second; i < b.size(); ++</pre>
            i){
                 for(;; j = (j + 1) % a.size()){
                         r = max(r, norm2(b[i] - a[j]));
                         if(cross((b[(i + 1) % b.size()] -
                              b[i]), (a[(j + 1) % a.size()]
                              - a[j])) <= EPS) break;</pre>
        return r;
lf intercircle(vector<pt>& p, circle c){ // area of
   intersection with circle
        lf r=0.;
        for(int i = 0, n = p.size(); i < n; i++){</pre>
                 int j = (i + 1) % n;
                 If w = intertriangle(c, p[i], p[i]);
                 if(cross((p[j] - c.center), (p[i] - c.
                    center)) > 0) r += w;
                 else r -= w;
        return abs(r);
ll pick(vector<pt>& p) {
        11 boundary = 0;
        for (int i = 0, n = p.size(); i < n; i++) {</pre>
                 int j = (i + 1 == n ? 0 : i + 1);
                 boundary += \gcd((11) \operatorname{abs}(p[i].x - p[j].x)
                    ), (ll) abs(p[i].y - p[j].y));
```

```
5.13 Poligonos
```

```
oligonos
```

```
return abs(area(p)) + 1 - boundary / 2;
// minimum distance between two parallel lines (non
   necessarily axis parallel)
// such that the polygon can be put between the lines
// O(n) CCW polygon
lf width(vector<pt> &p) {
    int n = (int)p.size();
    if (n <= 2) return 0;
    lf ans = inf;
    int i = 0, j = 1;
    while (i < n) {
        while (cross(p[(i + 1) % n] - p[i], p[(j + 1) % n])
          ] - p[j]) >= 0) j = (j + 1) % n;
        line l1(p[i], p[(i + 1) \% n]);
        ans = min(ans, l1.dist(p[j]));
        i++;
    return ans;
// O(n) {minimum perimeter, minimum area} CCW polygon
pair<ld, ld> minimum enclosing rectangle(vector<pt> &p) {
        int n = p.size();
        if (n <= 2) return {perimeter(p), 0};</pre>
        int mndot = 0;
    lf tmp = dot(p[1] - p[0], p[0]);
        for (int i = 1; i < n; i++) {
                if (dot(p[1] - p[0], p[i]) <= tmp) {</pre>
                        tmp = dot(p[1] - p[0], p[i]);
                        mndot = i;
        ld ansP = inf;
        ld ansA = inf;
        int i = 0, j = 1, mxdot = 1;
        while (i < n) {
                pt cur = p[(i + 1) % n] - p[i];
        while (cross(cur, p[(j + 1) % n] - p[j]) >= 0) j
           = (i + 1) % n;
        while (dot(p[(mxdot + 1) % n], cur) >= dot(p[
           mxdot], cur)) mxdot = (mxdot + 1) % n;
        while (dot(p[(mndot + 1) % n], cur) <= dot(p[
           mndot], cur)) mndot = (mndot + 1) % n;
        line l1(p[i], p[(i + 1) % n]);
        // minimum perimeter
        ansP = min(ansP, 2.0 * ((dot(p[mxdot], cur) / 
           norm(cur) - dot(p[mndot], cur) / norm(cur)) +
           11.dist(p[i])));
        // minimum area
        ansA = min(ansA, (dot(p[mxdot], cur) / norm(cur)
           - dot(p[mndot], cur) / norm(cur)) * l1.dist(p[
           j]));
```

```
i++;
    return {ansP, ansA};
// maximum distance from a convex polygon to another
   convex polygon
lf maximum dist_from_polygon_to_polygon(vector<pt> &u,
   vector<pt> &v) \{ \frac{1}{0} (n) \}
    int n = (int)u.size(), m = (int)v.size();
    lf ans = 0;
    if (n < 3 | | m < 3) {
        for (int i = 0; i < n; i++) {</pre>
            for (int j = 0; j < m; j++) ans = max(ans,
                dis2(u[i], v[i]));
        return sqrt(ans);
    if (u[0].x > v[0].x) swap(n, m), swap(u, v);
    int i = 0, j = 0, step = n + m + 10;
    while (j + 1 < m \&\& v[j].x < v[j + 1].x) j++;
    while (step--) {
        if (cross(u[(i + 1) % n] - u[i], v[(j + 1) % m] -
            v[\dot{j}]) >= 0) \dot{j} = (\dot{j} + 1) \% m;
        else i = (i + 1) % n;
        ans = \max(ans, dis2(u[i], v[j]));
    return sgrt (ans);
pt project from point to seq(pt a, pt b, pt c) {
    double r = dis2(a, b);
    if (sign(r) == 0) return a;
   r = dot(c - a, b - a) / r;
    if (r < 0) return a;</pre>
    if (r > 1) return b;
    return a + (b - a) * r;
// minimum distance from point c to segment ab
lf pt_to_seg(pt a, pt b, pt c) {
    return dis(c, project from point to seq(a, b, c));
pair<pt, int> point_poly_tangent(vector<pt> &p, pt Q, int
    dir, int 1, int r) {
    while (r - 1 > 1) {
        int mid = (1 + r) >> 1;
        bool pvs = sign(orient(Q, p[mid], p[mid - 1])) !=
        bool nxt = sign(orient(0, p[mid], p[mid + 1])) !=
            -dir:
        if (pvs && nxt) return {p[mid], mid};
```

```
if (!(pvs || nxt)) {
                              auto p1 = point_poly_tangent(p, Q, dir, mid +
                              auto p2 = point poly tangent(p, Q, dir, 1,
                                      mid - 1);
                              return sign(orient(Q, p1.first, p2.first)) ==
                                        dir ? p1 : p2;
                   if (!pvs) {
                              if (sign(orient(0, p[mid], p[l])) == dir) r
                                      = mid - 1;
                              else if (sign(orient(Q, p[1], p[r])) == dir)
                                      r = mid - 1;
                              else l = mid + 1;
                   if (!nxt) {
                              if (sign(orient(Q, p[mid], p[l])) == dir) 1
                                      = mid + 1;
                              else if (sign(orient(Q, p[l], p[r])) == dir)
                                      r = mid - 1;
                              else l = mid + 1;
          pair<pt, int> ret = {p[1], 1};
          for (int i = 1 + 1; i <= r; i++) ret = sign(orient(Q,
                    ret.first, p[i])) != dir ? make pair(p[i], i) :
         return ret;
// (ccw, cw) tangents from a point that is outside this
        convex polygon
// returns indexes of the points
// ccw means the tangent from Q to that point is in the
        same direction as the polygon ccw direction
pair<int, int> tangents from point to polygon(vector<pt>
        &p, pt Q) {
         int ccw = point_poly_tangent(p, Q, 1, 0, (int)p.size
                   () - 1).second;
         int cw = point_poly_tangent(p, Q, -1, 0, (int)p.size
                   () - 1).second;
         return make_pair(ccw, cw);
// minimum distance from a point to a convex polygon
// it assumes point lie strictly outside the polygon
lf dist_from_point_to_polygon(vector<pt> &p, pt z) {
          lf ans = inf;
          int n = p.size();
          if (n <= 3) {
                   for (int i = 0; i < n; i++) ans = min(ans,</pre>
                            pt_{t_0} = pt_{t_0} 
                   return ans;
          pair<int, int> dum = tangents_from_point_to_polygon(p
                  , z);
```

```
int r = dum.first;
    int 1 = dum.second;
    if(1 > r) r += n;
    while (1 < r) {
        int mid = (1 + r) >> 1;
        lf left = dis2(p[mid % n], z), right= dis2(p[(mid
            + 1) % n], z);
        ans = min({ans, left, right});
        if(left < right) r = mid;</pre>
        else l = mid + 1;
    ans = sqrt(ans);
    ans = min(ans, pt_to_seq(p[1 % n], p[(1 + 1) % n], z)
    ans = min(ans, pt_to_seq(p[1 % n], p[(1 - 1 + n) % n])
       ], z));
    return ans;
// minimum distance from a convex polygon to another
   convex polygon
// the polygon doesnot overlap or touch
lf dist_from_polygon_to_polygon(vector<pt> &p1, vector<pt</pre>
   > &p2) { // O(n log n)}
    lf ans = inf;
    for (int i = 0; i < p1.size(); i++) {</pre>
        ans = min(ans, dist_from_point_to_polygon(p2, p1[
    for (int i = 0; i < p2.size(); i++) {</pre>
        ans = min(ans, dist_from_point_to_polygon(p1, p2[
           il));
    return ans;
// it returns a point such that the sum of distances
// from that point to all points in p is minimum
// O(n log^2 MX)
PT geometric_median(vector<PT> p) {
        auto tot dist = [&](PT z) {
            double res = 0;
            for (int i = 0; i < p.size(); i++) res +=</pre>
                dist(p[i], z);
            return res;
        auto findY = [&](double x) {
            double yl = -1e5, yr = 1e5;
            for (int i = 0; i < 60; i++) {
                double ym1 = yl + (yr - yl) / 3;
                double ym2 = yr - (yr - y1) / 3;
                double d1 = tot_dist(PT(x, ym1));
                double d2 = tot dist(PT(x, vm2));
                if (d1 < d2) yr = ym2;
                else yl = ym1;
```

```
return pair<double, double> (yl, tot_dist(PT(
                x, y1));
    double x1 = -1e5, xr = 1e5;
    for (int i = 0; i < 60; i++) {</pre>
        double xm1 = x1 + (xr - x1) / 3;
        double xm2 = xr - (xr - x1) / 3;
        double y1, d1, y2, d2;
        auto z = findY(\bar{x}m1); y1 = z.first; d1 = z.second;
        z = findY(xm2); y2 = \overline{z}.first; d2 = z.second;
        if (d1 < d2) xr = xm2;
        else x1 = xm1;
    return {xl, findY(xl).first };
// ear decomposition, O(n^3) but faster
vector<vector<pt>> triangulate(vector<pt> p) {
        vector<vector<pt>> v;
        while (p.size() >= 3) {
                for (int i = 0, n = p.size(); i < n; i++)
                         int pre = i == 0 ? n - 1 : i -
                            1;;
                         int nxt = i == n - 1 ? 0 : i +
                            1;;
                         lf ori = orient(p[i], p[pre], p[
                            nxtl);
                         if (ori < 0) {
                                 int ok = 1;
                                 for (int j = 0; j < n; j
                                    ++) {
                                         if (j == i || j
                                             == pre || j ==
                                              nxt) continue;
                                          vector < pt > tr = {
                                             p[i], p[pre],
                                             p[nxt]};
                                          if (
                                             point_in_polygon
                                             (tr , p[j]) !=
                                              OUT) {
                                                  ok = 0;
                                                  break:
                                 if (ok) {
                                          v.push_back({p[
                                             pre], p[i], p[
                                             p.erase(p.begin()
                                              + i);
```

```
break;
}

return v;
}
```

5.14 Segmentos

```
// add Lines Points
bool in disk(pt a, pt b, pt p) { // pt p inside ab disk
        return dot(a - p, b - p) \le E0;
bool on segment (pt a, pt b, pt p) { // p on ab
        return orient(a, b, p) == 0 \&\& in disk(a, b, p);
// ab crossing cd
bool proper_inter(pt a, pt b, pt c, pt d, pt& out) {
        lf oa = orient(c, d, a);
        lf ob = orient(c, d, b);
        lf oc = orient(a, b, c);
        lf od = orient(a, b, d);
        // Proper intersection exists iff opposite signs
        if (oa * ob < 0 && oc * od < 0)
                out = (a * ob - b * oa) / (ob - oa);
                return true;
        return false:
// intersection bwn segments
set<pt> inter_ss(pt a, pt b, pt c, pt d) {
        pt out;
        if (proper_inter(a, b, c, d, out)) return {out};
        set<pt> s;
        if (on_segment(c, d, a)) s.insert(a); // a in cd
        if (on_segment(c, d, b)) s.insert(b); // b in cd
        if (on_segment(a, b, c)) s.insert(c); // c in ab
        if (on segment(a, b, d)) s.insert(d); // d in ab
        return s;
If pt to seg(pt a, pt b, pt p) { // p to ab
        if (a != b) {
                line l(a, b);
                if (l.cmp_proj(a, p) && l.cmp_proj(p, b))
                    // if closest to projection = (a, p,
                   b)
                        return l.dist(p); // output
                            distance to line
```

```
return min(norm(p - a), norm(p - b)); //
           otherwise distance to A or B
lf seg_to_seg(pt a, pt b, pt c, pt d) {
        pt dummy;
        if (proper inter(a, b, c, d, dummy)) return 0; //
            ab intersects cd
        return min({pt_to_seg(a, b, c), pt_to_seg(a, b, d
           ), pt_to_seg(c, d, a), pt_to_seg(c, d, b)});
           // try the 4 pts
int length union(vector<pt>& a) { // REVISAR
        int n = a.size();
        vector<pair<int, bool>> x(n * 2);
        for (int i = 0; i < n; i++) {
                x[i * 2] = \{a[i].x, false\};
                x[i * 2 + 1] = \{a[i].y, true\};
        sort(x.begin(), x.end());
        int result = 0;
        int c = 0;
        for (int i = 0; i < n * 2; i++) {
                if (i > 0 && x[i].first > x[i - 1].first
                   && c > 0) result += x[i].first - x[i -
                    1].first;
                if (x[i].second) c--;
                else c++;
        return result;
```

5.15 Triangle Union

```
// Area of the union of a set of n triangles
// T(n^2 logn) M(n)

typedef double dbl;

const dbl eps = 1e-9;

inline bool eq(dbl x, dbl y) {
    return fabs(x - y) < eps;
}

inline bool lt(dbl x, dbl y) {
    return x < y - eps;
}

inline bool gt(dbl x, dbl y) {
    return x > y + eps;
}

inline bool le(dbl x, dbl y) {
    return x < y + eps;
}</pre>
```

```
inline bool ge(dbl x, dbl v){
    return x > y - eps;
struct ptT{
    dbl x, y;
    ptT(){}
    ptT(dbl x, dbl y): x(x), y(y) {}
    inline ptT operator - (const ptT & p)const{
        return ptT{x - p.x, y - p.y};
    inline ptT operator + (const ptT & p)const{
        return ptT\{x + p.x, y + p.y\};
    inline ptT operator * (dbl a)const{
        return ptT\{x * a, y * a\};
    inline dbl cross(const ptT & p)const{
        return x * p.y - y * p.x;
    inline dbl dot(const ptT & p)const{
        return x * p.x + y * p.y;
    inline bool operator == (const ptT & p)const{
        return eq(x, p.x) && eq(y, p.y);
};
struct LineT{
    ptT p[2];
    LineT(){}
    LineT(ptT a, ptT b):p\{a, b\}\{\}
    ptT vec()const{
        return p[1] - p[0];
    ptT& operator [](size_t i){
        return p[i];
};
inline bool lexComp(const ptT & 1, const ptT & r) {
    if(fabs(l.x - r.x) > eps){
        return 1.x < r.x;</pre>
    else return l.y < r.y;</pre>
vector<ptT> interSegSeg(LineT 11, LineT 12) {
    if(eq(11.vec().cross(12.vec()), 0)){
        if(!eq(l1.vec().cross(l2[0] - l1[0]), 0))
            return {};
        if(!lexComp(l1[0], l1[1]))
            swap(\bar{11}[0], 11[1]);
        if(!lexComp(12[0], 12[1]))
            swap(12[0], 12[1]);
```

```
Triangle Union
```

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GEOMETRIA
```

```
ptT l = lexComp(l1[0], l2[0]) ? l2[0] : l1[0];
        ptT r = lexComp(11[1], 12[1]) ? 11[1] : 12[1];
        if(1 == r)
            return {1};
        else return lexComp(l, r) ? vector<ptT>{l, r} :
           vector<ptT>();
    else{
        dbl s = (12[0] - 11[0]).cross(12.vec()) / 11.vec
           ().cross(12.vec());
        ptT inter = 11[0] + 11.vec() * s;
        if(qe(s, 0) \&\& le(s, 1) \&\& le((12[0] - inter).dot
            (12[1] - inter), 0))
            return {inter};
        else
            return {};
inline char get_segtype(LineT segment, ptT other_point){
    if(eq(seament[0].x, seament[1].x))
        return 0;
    if(!lexComp(segment[0], segment[1]))
        swap(segment[0], segment[1]);
    return (segment[1] - segment[0]).cross(other_point -
       segment[0]) > 0 ? 1 : -1;
dbl union area(vector<tuple<ptT, ptT, ptT> > triangles){
    vector<LineT> segments(3 * triangles.size());
    vector<char> segtype(segments.size());
    for (size t i = 0; i < triangles.size(); i++) {
        ptT a, b, c;
        tie(a, b, c) = triangles[i];
        segments [3 * i] = lexComp(a, b) ? LineT(a, b) :
           LineT(b, a);
        seqtype[3 * i] = qet_seqtype(segments[3 * i], c);
        segments [3 * i + 1] = lexComp(b, c)? LineT(b, c)
            : LineT(c, b);
        seqtype[3 * i + 1] = qet_seqtype(segments[3 * i +
            11, a);
        segments [3 * i + 2] = lexComp(c, a)? LineT(c, a)
            : LineT(a, c);
        seqtype[3 * i + 2] = qet seqtype(segments[3 * i +
            21, b);
    vector<dbl> k(segments.size()), b(segments.size());
    for (size t i = 0; i < segments.size(); i++) {
        if(seatype[i]){
            k[i] = (segments[i][1].v - segments[i][0].v)
                / (segments[i][1].x - segments[i][0].x);
            b[i] = segments[i][0].y - k[i] * segments[i]
               ][0].x;
    dbl ans = 0;
```

```
for(size t i = 0; i < segments.size(); i++){</pre>
    if(!seatvpe[i])
        continue;
    dbl l = segments[i][0].x, r = segments[i][1].x;
    vector<pair<dbl, int> > evts;
    for(size_t j = 0; j < segments.size(); j++) {</pre>
        if(!seatype[i] || i == i)
            continue;
        dbl l1 = segments[\dot{j}][0].x, r1 = segments[\dot{j}
           ][1].x;
        if(ge(l1, r) || ge(l, r1))
            continue;
        dbl common l = max(l, ll), common r = min(r, ll)
        auto pts = interSeqSeq(segments[i], segments[
           j]);
        if(pts.emptv()){
            dbl yll = k[j] * common_l + b[j];
            dbl yl = k[i] * common l + b[i];
            if(lt(yl1, yl) == (seqtype[i] == 1)){
                int evt type = -seqtype[i] * seqtype[
                evts.emplace back(common l, evt type)
                evts.emplace back(common r, -evt type
        else if(pts.size() == 1u){
            dbl vl = k[i] * common_l + b[i], yll = k[
                j] * common_l + b[j];
            int evt type = -seqtype[i] * seqtype[j];
            if(lt(vl1, vl) == (seqtype[i] == 1)){
                evts.emplace_back(common_l, evt_type)
                evts.emplace_back(pts[0].x, -evt_type
                   );
            yl = k[i] * common r + b[i], yll = k[i] *
                 common_r + b[j];
            if(lt(yl1, yl) == (seqtype[i] == 1)){
                evts.emplace back(pts[0].x, evt type)
                evts.emplace back(common r, -evt type
                    );
        else{
            if(segtype[j] != segtype[i] || j > i){
                evts.emplace_back(common_1, -2);
                evts.emplace back(common r, 2);
    evts.emplace back(1, 0);
```

```
6 GRAFOS
```

```
sort(evts.begin(), evts.end());
    size_t j = 0;
    int balance = 0;
    while(j < evts.size()){</pre>
        size t ptr = j;
        while(ptr < evts.size() && eq(evts[j].first,</pre>
            evts[ptr].first)){
            balance += evts[ptr].second;
            ++ptr;
        if(!balance && !eq(evts[j].first, r)){
            dbl next_x = ptr == evts.size() ? r :
                evts[ptr].first;
            ans -= seqtype[i] \star (k[i] \star (next_x +
                evts[j].first) + 2 * b[i]) * (next x -
                 evts[i].first);
        j = ptr;
return ans/2;
```

6 Grafos

6.1 2sat

```
// O(n+m)
// (x1 or y1) and (x2 or y2) and ... and (xn or yn)
struct sat2{
        vector<vector<vi>>> q;
        vector<bool> vis, val;
        stack<int> st;
        vi comp;
        int n;
        sat2(int n):n(n),q(2, vector<vi>(2*n)),vis(2*n),
           val(2*n), comp(2*n) {}
        int neg(int x) {return 2*n-x-1;} // get not x
        void make_true(int u) {add_edge(neg(u), u);}
        void make false(int u) {make true(neg(u));}
        void add or(int u, int v) {implication(neg(u), v);}
            // (u or v)
        void diff(int u, int v) \{eq(u, neg(v));\} // u != v
        void eq(int u, int v) {
                implication(u, v);
                implication(v, u);
        void implication(int u,int v) {
                add_edge(u, v);
```

```
add edge (neg (v), neg (u));
        void add edge(int u, int v){
                 q[0][u].push back(v);
                 q[1][v].push_back(u);
        void dfs(int id, int u, int t=0) {
                 vis[u]=true;
                 for(auto &v:g[id][u])
                          if(!vis[v])dfs(id, v, t);
                 if (id) comp[u]=t;
                 else st.push(u);
        void kosaraju() {
                 for(int u=0; u<n; ++u) {
                          if(!vis[u])dfs(0, u);
                          if(!vis[neq(u)])dfs(0, neq(u));
                 vis.assign(2*n, false);
                 int t=0;
                 while(!st.empty()){
                          int u=st.top();st.pop();
                          if(!vis[u])dfs(1, u, t++);
        // return true if satisfiable, fills val[]
        bool check(){
                 kosaraju();
                 for (int i=0; i < n; ++i) {</pre>
                          if (comp[i] == comp[neg(i)]) return
                              false;
                          val[i]=comp[i]>comp[neg(i)];
                 return true;
};
int m,n;cin>>m>>n;
sat2 s(n);
char c1, c2;
for (int a, b, i=0; i < m; ++i) {</pre>
        cin>>c1>>a>>c2>>b;
        a--;b--;
        if(c1=='-')a=s.neg(a);
        if(c2=='-')b=s.neg(b);
        s.add or(a,b);
if(s.check()){
        for (int i=0; i < n; ++i) cout << (s.val[i]?'+':'-') << " "</pre>
        cout << "\n";
}else cout<<"IMPOSSIBLE\n";</pre>
```

6.2 Bellman Ford

```
// O(V*E)
vi bellman ford(vector<vii> &adj, int s, int n) {
        vi dist(n, INF); dist[s] = 0;
        for (int i = 0; i < n-1; i++) {
                 bool modified = false;
                 for (int u = 0; u < n; u + +)
                         if (dist[u] != INF)
                                  for (auto &[v, w] : adj[u
                                     ]){
                                          if (dist[v] <=</pre>
                                              dist[u] + w)
                                              continue;
                                          dist[v] = dist[u]
                                          modified = true;
                 if (!modified) break;
        bool negativeCicle = false;
        for (int u = 0; u < n; u + +)
                 if (dist[u] != INF)
                         for (auto &[v, w] : adj[u]){
                                  if (dist[v] > dist[u] + w
                                       negativeCicle = true
        return dist;
```

6.3 Block Cut Tree

```
// O(n) build
// bi connected save the edges
const int maxn = 1e5+5;
int lowLink[maxn] , dfn[maxn];
vector<vector<ii>>> bi connected;
stack<ii> comps;
int ndfn;
void tarjan(vector<vi>& adj, int u=0, int par=-1) {
        dfn[u] = lowLink[u] = ndfn++;
        for(auto &v : adj[u]){
                if (v != par && dfn[v] < dfn[u])
                        comps.push({u, v});
                if (dfn[v] == -1) {
                        tarjan(adj, v, u);
                        lowLink[u] = min(lowLink[u] ,
                            lowLink[v]);
                        if (lowLink[v] >= dfn[u]){
```

```
bi connected.emplace back
                                     (vector<ii>());
                                 ii edge;
                                 do{
                                          edge = comps.top
                                              ();
                                          comps.pop();
                                          bi connected.back
                                              emplace back (
                                              edge);
                                  }while (edge.first != u ||
                                      edge.second != v);
                                  reverse (all (bi connected.
                                     back()));
                 }else if(v != par) {
                         lowLink[u] = min(lowLink[u], dfn
                             [V]);
void init(vector<vi>& adj){
        for(int i=0;i<sz(adj);++i)</pre>
                dfn[i]=-1;
        bi_connected.clear();
        comps=stack<ii>();
        ndfn=0;
        tarjan(adj);
```

6.4 Bridges Online

```
vector<int> par, dsu 2ecc, dsu cc, dsu cc size;
int bridges;
int lca iteration;
vector<int> last visit;
void init(int n) {
        par.resize(n);
        dsu_2ecc.resize(n);
        dsu_cc.resize(n);
        dsu cc size.resize(n);
        lca_iteration = 0;
        last visit.assign(n, 0);
        for (int i=0; i<n; ++i) {</pre>
                dsu \ 2ecc[i] = i;
                dsu cc[i] = i;
                dsu_cc_size[i] = 1;
                par[i] = -1;
        bridges = 0;
```

```
int find 2ecc(int v) {
        if (v == -1)
                return -1:
        return dsu 2ecc[v] == v ? v : dsu 2ecc[v] =
           find 2ecc(dsu 2ecc[v]);
int find cc(int v) {
        v = find 2ecc(v);
        return dsu cc[v] == v ? v : dsu cc[v] = find cc(
           dsu cc[v]);
void make root(int v) {
        int root = v;
        int child = -1;
        while (v != -1)
                int p = find_2ecc(par[v]);
                par[v] = child;
                dsu\_cc[v] = root;
                child = v;
                v = p;
        dsu cc size[root] = dsu cc size[child];
void merge path (int a, int b) {
        ++lca iteration;
        vector<int> path a, path b;
        int lca = -1;
        while (lca == -1) {
                if (a !=-1) {
                        a = find_2ecc(a);
                        path a.push back(a);
                        if (last visit[a] ==
                            lca iteration) {
                                lca = a;
                                break;
                        last_visit[a] = lca_iteration;
                        a = par[a];
                if (b !=-1) {
                        b = find 2ecc(b);
                        path_b.push_back(b);
                        if (last visit[b] ==
                            lca_iteration) {
                                lca = b;
                                break:
                        last_visit[b] = lca_iteration;
                        b = par[b];
```

```
for (int v : path_a) {
                dsu_2ecc[v] = lca;
                if (v == lca)
                        break;
                --bridges;
        for (int v : path b) {
                dsu_2ecc[v] = lca;
                if (v == lca)
                        break:
                --bridges;
void add edge(int a, int b) {
        a = find 2ecc(a);
        b = find 2ecc(b);
        if (a == b)
                return;
        int ca = find cc(a);
        int cb = find cc(b);
        if (ca != cb) {
                ++bridges;
                if (dsu_cc_size[ca] > dsu_cc_size[cb]) {
                        swap(a, b);
                        swap(ca, cb);
                make root(a);
                par[a] = dsu cc[a] = b;
                dsu_cc_size[cb] += dsu_cc_size[a];
        } else {
                merge path(a, b);
```

6.5 Camino Mas Corto De Longitud Fija

```
6.6 Clique
```

```
6 GRAFOS
```

```
}

return ans;

}

int main() {

    int n, m, k; cin >> n >> m >> k;

    vector<vl> adj(n, vl(n, INFL));

    for (int i = 0; i<m; i++) {
            ll a, b, c; cin >> a >> b >> c; a--; b--;
            adj[a][b] = min(adj[a][b], c);

}

matrix graph(n, n, adj);
graph = pow(graph, k-1);

cout << (graph.m[0][n-1]==INFL ? -1 : graph.m[0][n-1]) << "\n";

return 0;

}
</pre>
```

6.6 Clique

9

```
* Credit: kactl
 * Given a graph as a symmetric bitset matrix (without
    any self edges)
 * Finds the maximum clique
 * Can be used to find the maximum independent set by
    finding a clique of the complement graph.
 * Runs in about 1s for n=155, and faster for sparse
    graphs
 * 0 indexed
const int N = 40;
typedef vector<bitset<N>> graph;
struct Maxclique {
  double limit = 0.025, pk = 0;
  struct Vertex {
    int i, d = 0;
  typedef vector<Vertex> vv;
  graph e;
  vv V;
  vector<vector<int>> C;
  vector<int> qmax, q, S, old;
  void init(vv& r) {
    for (auto \& v : r) v.d = 0;
    for (auto& v : r) for (auto j : r) v.d += e[v.i][j.i
       ];
    sort(r.begin(), r.end(), [](auto a, auto b) {
```

```
return a.d > b.d;
    int mxD = r[0].d;
    for (int i = 0; i < sz(r); i++) r[i].d = min(i, mxD)
       + 1;
 void expand(vv& R, int lev = 1) {
    S[lev] += S[lev - 1] - old[lev];
    old[lev] = S[lev - 1];
    while (sz(R)) {
      if (sz(q) + R.back().d <= sz(qmax)) return;</pre>
      g.push back(R.back().i);
      for(auto v : R) if (e[R.back().i][v.i]) T.push_back
          (\{v.i\});
      if (sz(T)) {
        if (S[lev]++ / ++pk < limit) init(T);
        int j = 0, mxk = 1, mnk = max(sz(gmax) - sz(g) +
           1, 1);
        C[1].clear(), C[2].clear();
        for (auto v : T) {
          int k = 1;
          auto f = [&](int i) {
            return e[v.i][i];
          while (any_of(C[k].begin(), C[k].end(), f)) k
          if (k > mxk) mxk = k, C[mxk + 1].clear();
          if (k < mnk) T[j++].i = v.i;
          C[k].push_back(v.i);
        if (j > 0) T[j - 1].d = 0;
        for (int k = mnk; k \le mxk; k++) for (int i : C[k]
          T[j].i = i, T[j++].d = k;
        expand(T, lev + 1);
      } else if (sz(q) > sz(qmax)) qmax = q;
      q.pop_back(), R.pop_back();
 Maxclique(graph g) : e(g), C(sz(e) + 1), S(sz(C)), old(
    for (int i = 0; i < sz(e); i++) V.push back(\{i\});
  vector<int> solve() { // returns the clique
    init(V), expand(V);
    return qmax;
} ;
```

6.7 Cycle Directed

vector<vi> adj;
vi parent,color;

```
int cy0, cy1;
bool dfs(int v) {
        color[v]=1;
        for(int u:adj[v]){
                 if(color[u] == 0) {
                          parent[u]=v;
                         if (dfs(u)) return true;
                 }else if(color[u]==1){
                          cv1=v;
                          c\bar{y}0=u;
                          return true;
        color[v]=2;
        return false;
//O(m)
void find cycle(int n) {
        color.assign(n, 0);
        parent.assign(n, -1);
        cv0 = -1;
        for(int v=0; v<n; ++v) {
                 if(color[v]==0){
                         if (dfs(v))break;
        if(cv0==-1){
                 cout << "IMPOSSIBLE \n";
                 return;
        vi cycle;
        cycle.push back(cy0);
        for (int v=cy1; v!=cy0; v=parent[v]) cycle.push_back(
            v);
        cycle.push_back(cy0);
        reverse(cycle.begin(),cycle.end());
        print(cycle);
```

6.8 Cycle Undirected

```
vector<vi> adj;
vector<bool> visited;
int cy0,cy1;
vi parent;
bool dfs(int v, int par) {
    visited[v]=true;
    for(int u:adj[v]) {
        if(u==par) continue;
        if(visited[u]) {
            cy1=v;
            cy0=u;
```

```
return true;
                 parent[u]=v;
                 if(dfs(u,parent[u]))return true;
        return false;
//O(m)
void find_cycle(int n) {
        visited.assign(n, false);
        parent.assign(n, -1);
        cv0 = -1;
        for (int v=0; v<n; ++v) {</pre>
                 if(!visited[v]){
                 if(dfs(v, parent[v]))break;
        if(cy0==-1){
                 cout << "IMPOSSIBLE\n";</pre>
                 return;
        vi cycle;
        cycle.push back(cy0);
        for (int v=cy1; v!=cy0; v=parent[v]) cycle.push_back(
        cycle.push_back(cy0);
        print(cycle);
```

6.9 Dial Algorithm

```
const int maxn = 2e5+5;
vector<ii> adj[maxn];
// O(n*k+m)
// bfs for edge weights in the range [0, k]
void bfs(int s, int n, int k) {
        vector<queue<int>> qs(k+1, queue<int>());
        vector<bool> vis(n, false);
        vector<int> dist(n, 1e9);
        qs[0].push(s);
        dist[s]=0;
        int pos=0, num=1;
        while (num) {
                while (qs[pos%(k+1)].empty())pos++;
                int u=qs[pos%(k+1)].front();
                qs[pos%(k+1)].pop();
                num --;
                if(vis[u])continue;
                vis[u]=true;
                for(auto [w,v]:adj[u]){
                         if (dist[v]>dist[u]+w) {
                                 dist[v]=dist[u]+w;
```

```
qs[dist[v]%(k+1)].push(v)
num++;
}
}
```

6.10 Dijkstra

6.11 Dijkstra Sparse Graphs

```
// O(E*log(V))
vl dijkstra(vector<vector<pll>>> &adj, int s, int n) {
        vl dist(n, INFL); dist[s] = 0;
        set<pll> pq;
        pq.insert(pll(0, s));
        while(!pq.empty()){
                pll front = *pq.begin(); pq.erase(pq.
                   begin());
                11 d = front.first, u = front.second;
                for (auto &[v, w] : adj[u]){
                        if (dist[u] + w < dist[v]) {
                                 pq.erase(pll(dist[v], v))
                                 dist[v] = dist[u] + w;
                                 pq.insert(pll(dist[v], v)
                                    );
        return dist;
```

6.12 Eulerian Path Directed

```
const int maxn = 1e5+5;
vector<int> adj[maxn],path;
int out[maxn], in[maxn]; // remember
void dfs(int v) {
        while(!adj[v].empty()){
                int u=adi[v].back();
                adj[v].pop_back();
                dfs(u);
        path.push back(v);
// n -> nodes, m -> edges, s -> start, e -> end
void eulerian_path(int n, int m, int s, int e){
        for (int i=0;i<n;++i) {</pre>
                if(i==s || i==e)continue;
                if(in[i]!=out[i]){
                         cout << "IMPOSSIBLE \n";
                         return;
        if (out[s]-in[s]!=1 || in[e]-out[e]!=1) {
                cout << "IMPOSSIBLE\n";
                return;
        dfs(s);
        reverse(path.begin(), path.end());
        if (sz (path) !=m+1 || path.back()!=e) cout <<"
            IMPOSSIBLE\n";
        else print(path);
```

6.13 Eulerian Path Undirected

6.14 Floyd Warshall

6.15 Kosaraju

```
const int maxn = 1e5+5;
// construir el grafo inverso
// remember adj[a]->b, adj_rev[b]->a
vi adj_rev[maxn], adj[maxn];
bool used[maxn];
int idx[maxn]; // componente de cada nodo
vi order, comp;
// O(n+m)
void dfs1(int v) {
    used[v]=true;
    for(int u:adj[v])
        if(!used[u])dfs1(u);
        order.push_back(v);
```

```
void dfs2(int v){
        used[v]=true;
        comp.push_back(v);
        for(int u:adj_rev[v])
                 if(!used[u])dfs2(u);
// returna el numero de componentes
int init(int n){
        for (int i=0; i<n; ++i) if (!used[i]) dfs1(i);</pre>
        for (int i=0;i<n;++i)used[i]=false;</pre>
        reverse (all (order));
        int i=0;
        for(int v:order) {
                 if(!used[v]){
                          dfs2(v);
                          for(int u:comp)idx[u]=j;
                          comp.clear();
                          j++;
        return j;
```

6.16 kruskal

6.17 Prim

```
// O(E * log V)
// check: primer parametro de prim
// check: cuando no hay mst
vector<vii> adj;
vi tomado;
```

```
priority queue<ii> pq;
void process(int u) {
        tomado[u] = 1;
        for (auto &[v, w] : adj[u]){
                 if (!tomado[v]) pq.emplace(-w, -v);
int prim(int v, int n){
        tomado.assign(n, 0);
        process(0);
        int mst costo = 0, tomados = 0;
        while (!pq.empty()) {
                 auto [w, u] = pq.top(); pq.pop();
w = -w; u = -u;
                 if (tomado[u]) continue;
                 mst_costo += w;
                 process(u);
                 tomados++;
                 if (tomados == n-1) break;
        return mst costo;
```

6.18 Puentes y Puntos

```
const int maxn = 1e5+5;
vector<bool> vis;
vi adj[maxn]; // undirected
vi tin, low;
int timer;
void dfs(int u,int p=-1) {
        vis[u]=true;
        tin[u]=low[u]=timer++;
        int children=0;
        for(int v:adi[u]){
                if (v==p) continue;
                if (vis[v])low[u]=min(low[u],tin[v]);
                else{
                         dfs(v,u);
                         low[u] = min(low[u], low[v]);
                         if(low[v]>tin[u]); // u-v puente
                         if (low[v]>=tin[u] && p!=-1); // u
                             punto de articulacion
                         ++children;
        if (p==-1 && children>1); // u punto de
           articulacion
// O(n+m)
void init(int n) {
        timer=0;
```

```
vis.assign(n, false);
tin.assign(n,-1); low.assign(n,-1);
for(int i=0;i<n;++i) {</pre>
        if(!vis[i])dfs(i);
```

6.19 Shortest Path Faster Algorithm

```
//Algoritmo mas rapido de ruta minima
//O(V*E) peor caso, O(E) en promedio.
bool spfa(vector<vii> &adj, vector<int> &d, int s, int n)
        d.assign(n, INF);
        vector<int> cnt(n, 0);
        vector<bool> inqueue(n, false);
        queue<int> q;
        d[s] = 0;
        q.push(s);
        inqueue[s] = true;
        while (!q.empty()) {
                int v = q.front();
                q.pop();
                inqueue[v] = false;
                for (auto& [to, len] : adj[v]) {
                        if (d[v] + len < d[to]) {
                                d[to] = d[v] + len;
                                if (!inqueue[to]) {
                                         q.push(to);
                                         inqueue[to] =
                                            true;
                                         cnt[to]++;
                                         if (cnt[to] > n)
                                                 return
                                                     false:
                                                     //
                                                     ciclo
                                                     negativo
        return true;
```

6.20 Tarjan

```
// O(n+m) build graph in g[] and callt()
const int maxn = 2e5 + 5;
vi low, num, comp, q[maxn];
```

```
int scc, timer;
stack<int> st;
void tjn(int u) {
        low[u] = num[u] = timer++; st.push(u); int v;
        for(int v: q[u]) {
                 if(num[v] == -1) tjn(v);
                 if(comp[v]==-1) low[u] = min(low[u], low[u])
                     v]);
        if(low[u]==num[u]) {
                 do\{ v = st.top(); st.pop(); comp[v]=scc;
                  } while (u != v);
                 ++scc;
void callt(int n) {
        timer = scc= 0;
        num = low = comp = vector\langle int \rangle (n, -1);
        for(int i = 0; i<n; i++) if(num[i]==-1) tjn(i);</pre>
```

7 Matematicas

7.1 Bruijn sequences

```
// Given alphabet [0, k) constructs a cyclic string
// of length k n that contains every length n string as
   substr.
vi deBruijnSeq(int k, int n, int lim) {
        if (k == 1) return {0};
        vi seq, aux(n + 1);
        int cont = 0;
        function<void(int,int)> gen = [&](int t, int p) {
                if (t > n) {
                        if (n % p == 0) for (int i = 1; i)

                                if (cont >= lim) return;
                                seq.push_back(aux[i]);
                                cont++;
                } else {
                        aux[t] = aux[t - p];
                        gen(t + 1, p);
                        while (++aux[t] < k)
                                if (cont >= lim) return;
                                gen(t + 1, t);
        };
        gen(1, 1);
    // for (int i = 0; i < n-1; i++) seg.push back(0);
        return seq;
```

7.2 Convoluciones

```
// c[k] = sumatoria (i&j = k, += a[i]*b[j]) AND
   convolution
//c[k] = sumatoria (i|j = k, += a[i]*b[j]) OR
   convolution
// c[k] = sumatoria (i^i = k, += a[i]*b[i]) XOR
   convolution
// c[k] = sumatoria (qcd(i, j) = k, += a[i]*b[j]) GCD
   convolution
// c[k] = sumatoria (lcm(i, j) = k, += a[i]*b[j]) LCM
   convolution
// todas las funciones tienen operaciones con modulo
// si es indexando en 1 entonces se pone un cero al
   principio y listo
template<int MOD> struct mint {
        static const int mod = MOD;
        explicit operator int() const { return v; }
        mint():v(0) {}
        mint(ll v):v(int(v%MOD)) \{ v += (v<0)*MOD; \}
        void build(ll _v) { v=int(_v%MOD), v+=(v<0)*MOD;
        mint& operator+=(mint o) {
                if ((v += o.v) >= MOD) v -= MOD;
                return *this; }
        mint& operator-=(mint o) {
                if ((v -= o.v) < 0) v += MOD;
                return *this; }
        mint& operator *= (mint o) {
                v = int((ll)v*o.v%MOD); return *this; }
        friend mint pow(mint a, ll p) { assert(p \ge 0);
                return p==0?1:pow(a*a,p/2)*(p&1?a:1); }
        friend mint inv(mint a) { assert(a.v != 0);
           return pow(a, MOD-2); }
        friend mint operator+(mint a, mint b) { return a
        friend mint operator-(mint a, mint b) { return a
        friend mint operator*(mint a, mint b) { return a
           \star = b;
using mi = mint<998244353>;
template<typename T>
void SubsetZetaTransform(vector<T>& v) {
        const int n = v.size(); // n must be a power of 2
        for (int j = 1; j < n; j <<= 1) {
                for (int i = 0; i < n; i++)
                        if (i & j) v[i] += v[i ^ j];
```

```
template<typename T>
void SubsetMobiusTransform(vector<T>& v) {
        const int n = v.size(); // n must be a power of 2
        for (int j = 1; j < n; j <<= 1) {</pre>
                for (int i = 0; i < n; i++)
                         if (i & j) v[i] -= v[i ^ j];
template<typename T>
void SupersetZetaTransform(vector<T>& v) {
        const int n = v.size(); // n must be a power of 2
        for (int j = 1; j < n; j <<= 1) {
    for (int i = 0; i < n; i++)</pre>
                         if (i & j) v[i ^ j] += v[i];
template<typename T>
void SupersetMobiusTransform(vector<T>& v) {
        const int n = v.size(); // n must be a power of 2
        for (int j = 1; j < n; j <<= 1) {</pre>
                for (int i = 0; i < n; i++)</pre>
                         if (i & j) v[i ^ j] -= v[i];
vector<int> PrimeEnumerate(int n) {
        vector<int> P; vector<bool> B(n + 1, 1);
        for (int i = 2; i <= n; i++) {
                if (B[i]) P.push_back(i);
                for (int j: P) { if (i * j > n) break; B
                    [i * j] = 0; if (i % j == 0) break; }
        return P;
template<typename T>
void DivisorZetaTransform(vector<T>& v) {
        const int n = sz(v) - 1;
        for (int p : PrimeEnumerate(n)) {
                for (int i = 1; i * p <= n; i++)
                         v[i * p] += v[i];
template<typename T>
void DivisorMobiusTransform(vector<T>& v) {
        const int n = sz(v) - 1;
        for (int p : PrimeEnumerate(n)) {
                for (int i = n / p; i; i--)
                         v[i * p] -= v[i];
template<tvpename T>
void MultipleZetaTransform(vector<T>& v) {
```

```
const int n = sz(v) - 1;
        for (int p : PrimeEnumerate(n)) {
                for (int i = n / p; i; i--)
                        v[i] += v[i * p];
template<tvpename T>
void MultipleMobiusTransform(vector<T>& v) {
        const int n = sz(v) - 1;
        for (int p : PrimeEnumerate(n)) {
                for (int i = 1; i * p <= n; i++)
                        v[i] = v[i * p];
template<typename T>
vector<T> AndConvolution(vector<T> A, vector<T> B) {
        SupersetZetaTransform(A);
        SupersetZetaTransform(B);
        for (int i = 0; i < sz(A); i++) A[i] *= B[i];
        SupersetMobiusTransform(A);
        return A;
template<typename T>
vector<T> OrConvolution(vector<T> A, vector<T> B) {
        SubsetZetaTransform(A);
        SubsetZetaTransform(B);
        for (int i = 0; i < sz(A); i++) A[i] *= B[i];
        SubsetMobiusTransform(A);
        return A;
template<typename T>
vector<T> GCDConvolution(vector<T> A, vector<T> B) {
        MultipleZetaTransform(A);
        MultipleZetaTransform(B);
        for (int i = 0; i < sz(A); i++) A[i] *= B[i];</pre>
        MultipleMobiusTransform(A);
        return A;
template<typename T>
vector<T> LCMConvolution(vector<T> A, vector<T> B) {
        DivisorZetaTransform(A);
        DivisorZetaTransform(B);
        for (int i = 0; i < sz(A); i++) A[i] *= B[i];
        DivisorMobiusTransform(A);
        return A;
template<typename T>
vector<T> XORConvolution(vector<T> A, vector<T> B) {
        const int n = sz(A);
        auto FWT = [&](vector<T>& v) {
                for (int len = 1; len < n; len <<= 1) {</pre>
```

```
for (int i = 0; i < n; i += (len
                               << 1)) {
                                    for (int j = 0; j < len;
                                        j++) {
                                             T u(v[i + j]);
                                             T w(v[i + j + len
                                                 ]);
                                             v[i + j] = u + w;
                                                  v[i + j + len
                                                 1 = u - w;
         FWT(A); FWT(B);
         for (int i = 0; i < n; i++) A[i] *= B[i];</pre>
         FWT(A);
         T \text{ inv } n(\text{inv}(T(n)));
         for (int i = 0; i < n; i++) A[i] *= inv_n;</pre>
         return A;
void main2(){
         int n;
         cin>>n;
         vector<mi> a(1 << n), b(1 << n);
         for (int x, i=0; i < sz(a); ++i) {cin>>x; a[i].build(x);}
         for(int x, i=0; i < sz(b); ++i) { cin >> x; b[i].build(x); }
         vector<mi> ans=XORConvolution(a,b);
         for (int i=0; i < sz (ans); ++i) cout < < ans[i].v << "";</pre>
```

7.3 Criba

```
// O((R-L+1)\log(\log(R)) + \operatorname{sgrt}(R)\log(\log(\operatorname{sgrt}(R)))
// R-L+1 <= 1e7, R <= 1e14
void segmentedSieve(long long L, long long R) {
    // generate all primes up to sgrt(R)
    long long lim = sqrt(R) + 3;
    vector<bool> mark(lim + 1, false);
    vector<long long> primes;
    for (long long i = 2; i <= lim; ++i) {
        if (!mark[i]) {
             primes.emplace_back(i);
             for (long long j = i * i; j <= \lim; j += i)
                 mark[i] = true;
    vector<bool> isPrime(R - L + 1, true);
    for (long long i : primes)
        for (long long j = \max(i * i, (L + i - 1) / i * i)
            ); j <= R; j += i)
            isPrime[j - L] = false;
    if (L == 1)
        isPrime[0] = false;
```

7.4 Chinese Remainder Theorem

```
/// Complexity: |N|*log(|N|)
/// Tested: Not vet.
/// finds a suitable x that meets: x is congruent to a i
   mod n i
/** Works for non-coprime moduli.
Returns \{-1,-1\} if solution does not exist or input is
    invalid.
Otherwise, returns \{x,L\}, where x is the solution unique
   to mod L = LCM \ of \ mods
pll crt(vl A, vl M) {
        ll n = A.size(), a1 = A[0], m1 = M[0];
        for(ll i = 1; i < n; i++) {
                 11 \ a2 = A[i], \ m2 = M[i];
                 11 q = \underline{\hspace{0.2cm}} gcd(m1, m2);
                 if(a1 % q!= a2 % q) return {-1,-1};
                 ll p, q;
                 extended_euclid(m1/g, m2/g, p, g);
                 11 \mod = m1 / q * m2;
                 q %= mod; p %= mod;
                11 \times = ((111*(a1*mod)*(m2/g))*mod*q + (1)
                    11*(a2*mod)*(m1/q))*mod*p) % mod; //
                    if WA there is overflow
                 a1 = x;
                if (a1 < 0) a1 += mod;
                m1 = mod;
```

```
return {a1, m1};
}
```

7.5 Divisors

```
// d(n) = (a1+1)*(a2+1)*...*(ak+1)
11 numDiv(map<11, 11>& f) {
        ll ans=1;
        for(auto [ ,pot]:f)ans=mul(ans, (pot+111));
        return ans;
// sigma(n) = (p1^(a1+1)-1)/(p1-1) * (p2^(a2+1)-1)/(p2-1)
    * ... * (p\bar{k}^{(ak+1)-1})/(p\bar{k}-1)
// suma divisores a la xth potencia
ll sumDiv(map<11, 11>& f){
        11 ans=1,potencia=1;
        for(auto [num, pot]:f){
                11 p=binpow(num, (pot+111) *potencia) -111;
                ans=mul(ans, mul(p, inv(num-111)));
        return ans;
ll productDiv(map<ll, ll>& f) {
        ll pi=1, res=1;
        for(auto [num, pot]:f){
                11 p=binpow(num, pot*(pot+111)/211);
                res=mul(binpow(res, pot+111),binpow(p, pi
                pi=mul(pi, pot+111, MOD-111);
        return res;
// si a y b son coprimos, entonces:
// sigma(a*b) = sigma(a)*sigma(b)
// d(a*b) = d(a)*d(b)
```

7.6 Ecuaciones Diofanticas

```
// O(log(n))
ll extended_euclid(ll a, ll b, ll &x, ll &y) {
    ll xx = y = 0;
    ll yy = x = 1;
    while (b) {
        ll q = a / b;
        ll t = b; b = a % b; a = t;
        t = xx; xx = x - q * xx; x = t;
        t = yy; yy = y - q * yy; y = t;
    }
    return a;
}
```

```
// a*x+b*y=c. returns valid x and y if possible.
// all solutions are of the form (x0 + k * b / q, y0 - k
   *b/q
bool find any solution (ll a, ll b, ll c, ll &x0, ll &y0,
    ll &a) {
        if (a == 0 and b == 0) {
                if (c) return false;
                x0 = y0 = q = 0;
                return true;
        q = extended_euclid (abs(a), abs(b), x0, y0);
        if (c % q != 0) return false;
        x0 *= c / q;
        v0 \star = c / q;
        if (a < 0) x0 *= -1;
        if (b < 0) v0 *= -1;
        return true;
void shift solution(ll &x, ll &y, ll a, ll b, ll cnt) {
        x += cnt * b;
        v = cnt * a;
// returns the number of solutions where x is in the
   range[minx, maxx] and y is in the range[miny, maxy]
11 find_all_solutions(ll a, ll b, ll c, ll minx, ll maxx,
    11 miny, 11 maxy) {
        11 x, v, a;
        if (find any solution (a, b, c, x, y, g) == 0)
           return 0;
        if (a == 0 and b == 0) {
                assert(c == 0);
                return 1LL * (maxx - minx + 1) * (maxy -
                   miny + 1);
        if (a == 0) {
                return (maxx - minx + 1) * (miny <= c / b
                     and c / b <= maxy);
        if (b == 0) {
                return (maxy - miny + 1) * (minx <= c / a</pre>
                     and c / a <= maxx);
        a /= q, b /= q;
        11 \text{ sign}_a = a > 0 ? +1 : -1;
        ll sign b = b > 0 ? +1 : -1;
        shift solution (x, y, a, b, (minx - x) / b);
        if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
        if (x > maxx) return 0;
        11 \ 1x1 = x;
        shift_solution(x, y, a, b, (maxx - x) / b);
        if (x > maxx) shift_solution (x, y, a, b, -sign_b
           );
        11 rx1 = x;
        shift_solution(x, y, a, b, -(miny - y) / a);
        if (y < miny) shift_solution (x, y, a, b, -sign_a</pre>
```

```
if (y > maxy) return 0;
ll lx2 = x;
shift_solution(x, y, a, b, -(maxy - y) / a);
if (y > maxy) shift_solution(x, y, a, b, sign_a);
ll rx2 = x;
if (lx2 > rx2) swap (lx2, rx2);
ll lx = max(lx1, lx2);
ll rx = min(rx1, rx2);
if (lx > rx) return 0;
return (rx - lx) / abs(b) + 1;
}

///finds the first k / x + b * k / gcd(a, b) >= val
ll greater_or_equal_than(ll a, ll b, ll x, ll val, ll g)
{
    ld got = 1.0 * (val - x) * g / b;
    return b > 0 ? ceil(got) : floor(got);
}
```

7.7 Exponenciacion binaria

7.8 Exponenciacion matricial

```
ans.m[i][j] +=
                                              mod(m[i][k],
                                              MOD) * mod(b.m)
                                              [k][j], MOD);
                                           ans.m[i][j] = mod
                                              (ans.m[i][i],
                                              MOD);
                 return ans:
};
matrix pow(matrix &b, ll p) {
        matrix ans(b.r, b.c, vector<vl>(b.r, vl(b.c, 0)))
        for (int i = 0; i < b.r; i++) ans.m[i][i] = 1;</pre>
        while (p) {
                 if (p&1) {
                         ans = ans*b;
                 b = b*b;
                 p >>= 1;
        return ans;
```

7.9 Fast Fourier Transform

```
///Complexity: O(N log N)
///tested: https://codeforces.com/gym/104373/problem/E
#define rep(i, a, b) for(int i = a; i < (b); ++i)
#define sz (v) ((int)v.size())
#define trav(a, x) for(auto& a : x)
#define all(v) v.begin(), v.end()
typedef vector<ll> vl;
typedef vector<int> vi;
typedef complex<double> C;
typedef vector<double> vd;
void fft(vector<C>& a) {
        int n = sz(a), L = 31 - builtin clz(n);
        static vector<complex<long double>> R(2, 1);
        static vector<C> rt(2, 1); // (^ 10% faster if
           double)
        for (static int k = 2; k < n; k *= 2) {
                R.resize(n); rt.resize(n);
                auto x = polar(1.0L, acos(-1.0L) / k);
                rep(i,k,2*k) rt[i] = R[i] = i&1 ? R[i/2]
                   * x : R[i/2];
        vi rev(n);
        rep(i, 0, n) rev[i] = (rev[i / 2] | (i & 1) << L) /
            2;
```

```
rep(i,0,n) if (i < rev[i]) swap(a[i], a[rev[i]]);
        for (int k = 1; k < n; k *= 2)
                for (int i = 0; i < n; i += 2 * k) rep(j
                   ,0,k) {
                        // C z = rt[j+k] * a[i+j+k]; //
                            (25% faster if hand-rolled)
                           /// include-line
                        auto x = (double *) & rt[j+k], y =
                            (double *) &a[i+j+k];
                            / exclude-line
                        C z(x[0]*y[0] - x[1]*y[1], x[0]*y
                           [1] + x[1] * y[0]);
                            / exclude-line
                        a[i + j + k] = a[i + j] - z;
                        a[i + j] += z;
vl conv(const vl& a, const vl& b) {
        if (a.empty() || b.empty()) return {};
        vd res(sz(a) + sz(b) - 1);
        int L = 32 - \underline{\text{builtin\_clz}(\text{sz}(\text{res}))}, n = 1 << L;
        vector<C> in(n), out(n);
        copy(all(a), begin(in));
        rep(i, 0, sz(b)) in[i].imag(b[i]);
        fft(in);
        trav(x, in) x *= x;
        1);
        fft (out);
        vector<ll> resp(sz(res));
        rep(i, 0, sz(res)) resp[i] = round(imag(out[i]) /
           (4.0 * n));
        return resp;
```

7.10 Fibonacci Fast Doubling

7.11 Fraction

```
typedef __int128 T;
struct Fraction{
        T num, den;
        Fraction():num(0),den(1){}
        Fraction (T n): num(n), den(1)  }
        Fraction(T n,T d):num(n),den(d) {reduce();}
        void reduce(){
                 // assert (den!=0);
                T gcd= gcd(num, den); // <-
                num/=qcd, den/=qcd;
                if (den<0) num=-num, den=-den;</pre>
        Fraction fractional part() const{ // x - floor(x)
                Fraction fp=Fraction(num%den,den);
                if (fp<Fraction(0))fp+=Fraction(1);</pre>
                return fp;
        T compare (Fraction f) const { return num*f.den-den*f
        Fraction operator + (const Fraction& f) {return
            Fraction(num*f.den+den*f.num,den*f.den);}
        Fraction operator - (const Fraction& f) {return
            Fraction(num*f.den-den*f.num,den*f.den);}
        Fraction operator * (const Fraction& f) {
                Fraction a=Fraction(num, f.den);
                Fraction b=Fraction(f.num,den);
                return Fraction(a.num*b.num,a.den*b.den);
        Fraction operator / (const Fraction& f) {return *
           this*Fraction(f.den,f.num);}
        Fraction operator += (const Fraction& f) {return *
            this=*this+f;}
        Fraction operator -= (const Fraction& f) {return *
            this=*this-f;}
        Fraction operator *= (const Fraction& f) {return *
            this=*this*f;}
        Fraction operator /= (const Fraction& f) {return *
            this=*this/f;}
        bool operator == (const Fraction& f) const{return
            compare (f) == 0;
        bool operator != (const Fraction& f) const{return
            compare (f) !=0;
        bool operator >= (const Fraction& f) const{return
            compare (f) \ge 0;
        bool operator <= (const Fraction& f)const{return</pre>
            compare (f) \le 0;
        bool operator > (const Fraction& f)const{return
            compare (f) > 0;
        bool operator < (const Fraction& f)const{return</pre>
            compare (f) < 0;
Fraction operator - (const Fraction& f) {return Fraction(-
   f.num, f.den);}
ostream& operator << (ostream& os, const Fraction& f) {
```

```
return os<<"("<<(11) f.num<<"/"<<(11) f.den<<")";}</pre>
```

7.12 Freivalds algorithm

7.13 Gauss Jordan

```
// O(min(n, m)*n*m)
const double EPS = 1e-9;
const int INF = 2; // it doesn't actually have to be
   infinity or a big number
int gauss (vector < vector<double> > a, vector<double> &
   ans) {
        int n = (int) a.size();
        int m = (int) a[0].size() - 1;
        vector<int> where (m, -1);
        for (int col=0, row=0; col<m && row<n; ++col) {</pre>
                 int sel = row;
                 for (int i=row; i<n; ++i)</pre>
                         if (abs (a[i][col]) > abs (a[sel
                             1[col]))
                                  sel = i;
                 if (abs (a[sel][col]) < EPS)</pre>
                         continue:
                 for (int i=col; i<=m; ++i)</pre>
                         swap (a[sel][i], a[row][i]);
                 where [col] = row;
                 for (int i=0; i<n; ++i)
                         if (i != row) {
                                  double c = a[i][col] / a[
                                     rowl[col];
                                  for (int j=col; j<=m; ++j
```

```
a[i][j] = a[row]
                                     ][j] * c;
        ++row;
ans.assign (m, 0);
for (int i=0; i<m; ++i)
        if (where[i] != -1)
                 ans[i] = a[where[i]][m] / a[where
                    [i]][i];
for (int i=0; i<n; ++i) {</pre>
        double sum = 0;
        for (int j=0; j<m; ++j)
                 sum += ans[j] * a[i][j];
        if (abs (sum - a[i][m]) > EPS)
                 return 0;
for (int i=0; i<m; ++i)</pre>
        if (where [i] == -1)
                 return INF;
return 1:
```

7.14 Gauss Jordan mod 2

```
// O(min(n, m)*n*m)
int gauss (vector < bitset<N> > &a, int n, int m, bitset<</pre>
   N> & ans) {
        vector<int> where (m, -1);
        for (int col=0, row=0; col<m && row<n; ++col) {</pre>
                 for (int i=row; i<n; ++i)</pre>
                         if (a[i][col]) {
                                  swap (a[i], a[row]);
                                  break;
                 if (! a[row][col])
                         continue;
                 where [col] = row;
                 for (int i=0; i<n; ++i)</pre>
                         if (i != row && a[i][col])
                                  a[i] ^= a[row];
                 ++row:
        for (int i=0; i<m; ++i)
                 if (where[i] != -1)
                         ans[i] = a[where[i]][m] / a[where
                             [i]][i];
        for (int i=0; i<n; ++i) {
                double sum = 0;
                 for (int j=0; j<m; ++j)
                         sum += ans[j] * a[i][j];
```

```
if (abs (sum - a[i][m]) > EPS)
                return 0;
for (int i=0; i<m; ++i)
        if (where [i] == -1)
                return INF:
return 1;
```

7.15 GCD y LCM

```
//0(\log 10 \, n) \, n == \max(a, b)
int gcd(int a, int b) { return b == 0 ? a : gcd(b, a%b);
int lcm(int a, int b) { return a / gcd(a, b) * b; }
//gcd(a, b, c) = gcd(a, gcd(b, c))
//gcd(a, b) = gcd(a, b-a)
// O(\log(\min(a, b)) - a*x+b*y=\gcd(a, b)
ll gcd(ll a, ll b, ll &x, ll &y) {
        x=1, y=0;
        11 \times 1=0, y1=1, a1=a, b1=b;
        while(b1) {
                 ll q=a1/b1;
                 tie(x, x1) = make_tuple(x1, x-q*x1);
                 tie(y, y1)=make_tuple(y1, y-q*y1);
                 tie(a1, b1) = make tuple(b1, a1-q*b1);
        }return a1;
```

7.16 Integral Definida

```
const int steps = 1e6; // %2==0
double f(double x);
double simpson(double a, double b) {
        double h=(b-a)/steps;
        double s=f(a)+f(b);
        for (int i=1; i <= steps-1; i++) {</pre>
                 double x=a+h*i;
                 s+=f(x)*((i&1)?4:2);
        s*=h/3;
        return s;
```

7.17 Inverso modular

```
11 mod(ll a, ll m) {
        return ((a%m) + m) % m;
```

```
11 modInverse(ll b, ll m) {
        11 x, y;
        ll d = extEuclid(b, m, x, y); //obtiene b*x + m*
           v == d
        if (d != 1) return -1;
                                        //indica error
        // b*x + m*y == 1, ahora aplicamos (mod m) para
           obtener\ b*x == 1 \pmod{m}
        return mod(x, m);
// Otra forma
// O(log MOD)
ll inv (ll a) {
        return binpow(a, MOD-2, MOD);
//Modulo constante
inv[1] = 1;
for(int i = 2; i < p; ++i)
        inv[i] = (p - (p / i) * inv[p % i] % p) % p;
```

7.18 Lagrange

```
const int N = 1e6;
int f[N], fr[N];
void initC() {
  f[0] = 1;
  for (int i=1; i<N; i++) f[i] = mul(f[i-1], i);</pre>
  fr[N-1] = inv(f[N-1]);
  for(int i=N-1; i>=1; --i) fr[i-1] = mul(fr[i], i);
// mint C(int n, int k) { return k<0 | | k>n ? 0 : f[n] *
   fr[k] * fr[n-k]; }
struct LagrangePol {
  int n;
  vi y, den, l, r;
  LagrangePol(vector<int> f): n(sz(f)), y(f), den(n), l(n
     ), r(n) \{ / / f[i] := f(i) \}
    // Calcula interpol. pol P in O(n) := deg(P) = sz(v)
    initC();
    for (int i = 0; i<n; i++) {</pre>
      den[i] = mul(fr[n-1-i], fr[i]);
      if((n-1-i) \& 1) den[i] = mod(-den[i]);
  int eval(int x) { // Evaluate LagrangePoly P(x) in O(n)
   1[0] = r[n-1] = 1;
    for (int i = 1; i < n; i++) l[i] = mul(l[i-1], mod(x - 1))
       i + 1));
    for (int i=n-2; i>=0; --i) r[i] = mul(r[i+1], mod(x -
       i - 1));
```

```
int ans = 0;
    for (int i = 0; i < n; i++) ans = add(ans, mul(mul(l[i
       ], r[i]), mul(y[i], den[i])));
    return ans;
};
// Para Xs que no sean de [0, N]
/// Complexity: 0(|N|^2)
/// Tested: https://tinyurl.com/y23sh38k
vector<lf> X, F;
lf f(lf x) {
 lf answer = 0;
  for (int i = 0; i < sz(X); i++) {
    lf prod = F[i];
    for (int j = 0; j < sz(X); j++) {
      if(i == j) continue;
      prod = mul(prod, divide(sbt(x, X[i]), sbt(X[i], X[i]))
         1)));
    answer = add(answer, prod);
  return answer;
//given y=f(x) for x [0, degree]
vi interpolation( vi &y ) {
  int n = sz(y);
  vi u = v, ans(n), sum(n);
  ans[0] = u[0], sum[0] = 1;
  for( int i = 1; i < n; ++i )</pre>
    int inv = binpow(i, MOD - 2);
    for( int j = n - 1; j >= i; --j)
      u[\dot{j}] = 1LL * (u[\dot{j}] - u[\dot{j} - 1] + MOD) * inv % MOD;
    for ( int j = i; j > 0; --j )
      sum[j] = (sum[j - 1] - 1LL * (i - 1) * sum[j] % MOD
          + MOD) % MOD;
      ans[j] = (ans[j] + 1LL * sum[j] * u[i]) % MOD;
    sum[0] = 1LL * (i - 1) * (MOD - sum[0]) % MOD;
    ans[0] = (ans[0] + 1LL * sum[0] * u[i]) % MOD;
  return ans;
```

7.19 Logaritmo Discreto

```
// O(sqrt(m))
// Returns minimum x for which a ^ x % m = b % m.
int solve(int a, int b, int m) {
```

```
// if (a == 0) return b == 0 ? 1 : -1; Casos 0^x
   = b
a %= m, b %= m;
int k = 1, add = 0, q;
while ((g = gcd(a, m)) > 1) {
        if (b == k)
                return add:
        if (b % a)
                return -1;
        b /= q, m /= q, ++add;
        k = (\tilde{k} * 111 * a / q) % m;
int n = sqrt(m) + 1;
int an = \bar{1};
for (int i = 0; i < n; ++i)
        an = (an * 111 * a) % m;
unordered_map<int, int> vals;
for (int q = 0, cur = b; q <= n; ++q) {</pre>
        vals[cur] = q;
        cur = (cur * 111 * a) % m;
for (int p = 1, cur = k; p \le n; ++p) {
        cur = (cur * 111 * an) % m;
        if (vals.count(cur)) {
                int ans = n * p - vals[cur] + add
                 return ans;
return -1;
```

7.20 Miller Rabin

```
bool witness (ll a, ll s, ll d, ll n) {
        ll x = fpow(a, d, n);
        if (x == 1 \mid | x == n - 1) return false;
        for (int i = 0; i < s - 1; i++) {</pre>
                 x = mul(x, x, n);
                 if (x == 1) return true;
                 if (x == n - 1) return false;
        return true;
11 \text{ test}[] = \{2, 3, 5, 7, 11, 13, 17, 19, 23, 0\};
bool is prime (ll n) {
        if (n < 2) return false;</pre>
        if (n == 2) return true;
        if (n % 2 == 0) return false;
        11 d = n - 1, s = 0;
        while (d \% 2 == 0) ++s, d /= 2;
        for (int i = 0; test[i] && test[i] < n; ++i)</pre>
                 if (witness(test[i], s, d, n))
                         return false:
        return true;
```

7.21 Miller Rabin Probabilistico

```
using u64 = uint64_t;
using u128 = uint128 t;
u64 binpower(u64 base, u64 e, u64 mod) {
        u64 \text{ result} = 1;
        base %= mod;
        while (e) {
                if (e & 1)
                        result = (u128) result * base %
                            mod:
                base = (u128)base * base % mod;
                e >>= 1;
        return result:
bool check composite(u64 n, u64 a, u64 d, int s) {
        u64 x = binpower(a, d, n);
        if (x == 1 | | x == n - 1)
                return false:
        for (int r = 1; r < s; r++) {
                x = (u128)x * x % n;
                if (x == n - 1)
                        return false;
        return true;
};
bool MillerRabin (u64 n, int iter=5) { // returns true if
   n is probably prime, else returns false.
```

7.22 Mobius

```
// 1 if n is 1
// 0 if n has a squared prime factor
// (-1)^k if n is a product of k distinct prime factors
const int N = 1e6+1;
int mob[N];
void mobius() {
        mob[1] = 1;
        for (int i = 2; i < N; i++) {</pre>
                mob[i]--;
                for (int j = i + i; j < N; j += i) {
                        mob[i] -= mob[i];
// to count coprime pairs
// ans=n*(n-1)/2
// for(int x:a) {
                for(int y:divisors[a])cnt[y]++;
// ans+= (mobius[v] *cnt[v] * (cnt[v]-1))/2
```

7.23 Number Theoretic Transform

```
const int N = 1 << 20;
const int mod = 469762049; //998244353
const int root = 3;
int lim, rev[N], w[N], wn[N], inv_lim;
void reduce(int &x) { x = (x + mod) % mod; }
int POW(int x, int y, int ans = 1) {
    for (; y; y >>= 1, x = (long long) x * x % mod)
        if (y & 1) ans = (long long) ans * x % mod;
    return ans;
```

```
void precompute(int len) {
        \lim = wn[0] = 1; int s = -1;
        while (lim < len) lim <<= 1, ++s;
        for (int i = 0; i < lim; ++i) rev[i] = rev[i >>
           1) >> 1 | (i & 1) << s;
        const int g = POW(root, (mod - 1) / lim);
        inv \lim = POW(\lim, mod - 2);
        for (int i = 1; i < lim; ++i) wn[i] = (long long)
            wn[i - 1] * a % mod;
void ntt(vector<int> &a, int typ) {
        for (int i = 0; i < lim; ++i) if (i < rev[i])</pre>
           swap(a[i], a[rev[i]]);
        for (int i = 1; i < lim; i <<= 1) {</pre>
                for (int j = 0, t = \lim / i / 2; j < i;
                    ++j) w[j] = wn[j * t];
                for (int j = 0; j < lim; j += i << 1) {
                         for (int k = 0; k < i; ++k) {
                                 const int x = a[k + j], y
                                     = (long long) a[k + j]
                                     + i] * w[k] % mod;
                                 reduce(a[k + j] += y -
                                    mod), reduce(a[k + j +
                                     i = x - y);
        if (!typ) {
                reverse(a.begin() + 1, a.begin() + lim);
                for (int i = 0; i < lim; ++i) a[i] = (</pre>
                    long long) a[i] * inv lim % mod;
vector<int> multiply(vector<int> &f, vector<int> &q) {
        int n=(int)f.size() + (int)q.size() - 1;
        precompute(n);
        vector<int> a = f, b = q;
        a.resize(lim); b.resize(lim);
        ntt(a, 1), ntt(b, 1);
        for (int i = 0; i < lim; ++i) a[i] = (long long)</pre>
           a[i] * b[i] % mod;
        ntt(a, 0);
        a.resize(n + 1);
        return a;
```

7.24 Pollard Rho

```
//O(n^{(1/4)}) (?)
ll pollard rho(ll n, ll c) {
        11 x = 2, y = 2, i = 1, k = 2, d;
        while (true) {
```

```
x = (mul(x, x, n) + c);
                 if (x >= n) x -= n;
                 d = \underline{\hspace{0.2cm}} gcd(x - y, n);
                 if (d > 1) return d;
                 if (++i == k) y = x, k <<= 1;
        return n;
void factorize(ll n, vector<ll> &f) {
        if (n == 1) return;
        if (is prime(n)) {
                 f.push back(n);
                 return;
        11 d = n;
        for (int i = 2; d == n; i++)
                 d = pollard rho(n, i);
        factorize(d, f);
        factorize(n/d, f);
```

7.25 Simplex

```
// Maximizar c1*x1 + c2*x2 + c3*x3 ...
// Restricciones a11*x1 + a12*x2 <= b1, a22*x2 + a23*x3
// Retorna valor optimo y valores de las variables
// O(c^2*b), O(c*b) - variables c, restricciones b
typedef double lf;
const lf EPS = 1e-9;
struct Simplex{
        vector<vector<lf>> A;
        vector<lf> B.C:
        vector<int> X,Y;
        lf z;
        int n,m;
        Simplex(vector<vector<lf>> _a, vector<lf> _b,
            vector<lf> _c) {
                 A= a; B= b; C= c;
                 n=B.size(); m=C.size(); z=0.;
                 X=vector<int>(m); Y=vector<int>(n);
                 for (int i=0; i<m; ++i) X[i]=i;</pre>
                 for (int i=0; i<n; ++i) Y[i]=i+m;</pre>
        void pivot(int x, int y) {
                 swap(X[y],Y[x]);
                 B[x]/=A[x][y];
                 for(int i=0;i<m;++i)if(i!=y)A[x][i]/=A[x</pre>
                     ][y];
                 A[x][y]=1/A[x][y];
                 for (int i=0; i < n; ++i) if (i!=x&&abs(A[i][y])</pre>
                     >EPS) {
```

```
B[i] -= A[i][y] *B[x];
                   for (int j=0; j<m; ++j) if (j!=y) A[i] [</pre>
                       j] -= A[i][y] *A[x][j];
                   A[i][y] = -A[i][y] * A[x][y];
         z+=C[v]*B[x];
         for (int i=0; i < m; ++i) if (i!=y) C[i] -= C[y] *A[</pre>
             x][i];
         C[y] = -C[y] *A[x][y];
pair<lf, vector<lf>> maximize() {
         while (1) {
                   int x=-1, y=-1;
                   lf mn=-EPS:
                   for (int i=0;i<n;++i) if (B[i] <mn) mn</pre>
                       =B[i], x=i;
                   if (x<0) break;</pre>
                   for (int i=0; i<m; ++i) if (A[x][i]<-</pre>
                       EPS) {y=i;break;}
                   // assert (y>=0) \rightarrow y<0, no
                       solution to Ax<=B
                   pivot(x, v);
         while(1){
                   if mx=EPS;
                   int x=-1, y=-1;
                   for (int i=0; i<m; ++i) if (C[i]>mx) mx
                       =C[i], v=i;
                   if (v<0)break;</pre>
                   lf mn=1e200;
                   for (int i=0; i<n; ++i) if (A[i][y]>
                       EPS\&\&B[i]/A[i][y]<mn)mn=B[i]/A
                       [i][y],x=i;
                   // assert (x>=0) -> x<0, unbounded
                   pivot(x,y);
         vector<lf> r(m);
         for(int i=0;i<n;++i)if(Y[i]<m)r[Y[i]]=B[i</pre>
             ];
         return {z,r};
```

7.26Simplex Int

};

```
// Maximizar c1*x1 + c2*x2 + c3*x3 ...
// Restricciones a11*x1 + a12*x2 <= b1, a22*x2 + a23*x3
   <= h2 . . .
// Retorna valor optimo y valores de las variables
// O(c^2*b), O(c*b) - variables c, restricciones b (tle)
struct Fraction{};
typedef Fraction lf;
```

```
const lf ZERO(0), INF(1e18);
struct Simplex{
         vector<vector<lf>> A;
         vector<lf> B.C:
         vector<int> X,Y;
         lf z;
         int n,m;
         Simplex (vector<vector<lf>> _a, vector<lf> _b,
             vector<lf> c) {
                  A = a; B = b; C = c;
                  n=B.size(); m=C.size(); z=ZERO;
                  X=vector<int>(m); Y=vector<int>(n);
                  for (int i=0; i<m; ++i) X[i]=i;</pre>
                  for (int i=0; i < n; ++i) Y[i] = i + m;</pre>
         void pivot(int x, int y) {
                  swap (X[y], Y[x]);
                  B[x]/=A[x][y];
                  for (int i=0; i<m; ++i) if (i!=v) A[x][i]/=A[x
                      ] [y];
                  A[x][y] = Fraction(1)/A[x][y];
                  for (int i=0; i<n; ++i) if (i!=x && A[i][y]!=</pre>
                      ZERO) {
                            B[i] -= A[i][y] *B[x];
                            for (int j=0; j < m; ++ j) if (j!=y) A[i] [
                                j] -= A[i][y] *A[x][j];
                            A[i][y] = -A[i][y] *A[x][y];
                  z+=C[y]*B[x];
                  for (int i=0; i < m; ++i) if (i!=y) C[i] -= C[y] *A[</pre>
                      x][i];
                  C[y] = -C[y] *A[x][y];
         pair<lf, vector<lf>> maximize() {
                  while(1){
                            int x=-1, y=-1;
                            lf mn=ZERO;
                            for (int i=0;i<n;++i) if (B[i]<mn) mn</pre>
                                =B[i], x=i;
                            if(x<0)break;</pre>
                            for (int i=0; i<m; ++i) if (A[x][i] <</pre>
                                ZERO) {y=i;break;}
                            // assert (y>=0) -> y<0, no
                                solution to Ax<=B
                            pivot(x, y);
                  while (1) {
                            lf mx=ZERO:
                            int x=-1, y=-1;
                            for (int i=0; i<m; ++i) if (C[i]>mx) mx
                                =C[i], v=i;
                            if (v<0) break;</pre>
                            lf mn=INF;
```

```
for (int i=0; i<n; ++i) if (A[i][y]>
                     ZERO && B[i]/A[i][y] < mn) mn = B[i]
                     ]/A[i][y], x=i;
                  // assert (x \ge 0) \rightarrow x < 0, unbounded
                 pivot(x,y);
         vector<lf> r(m);
         for (int i=0; i<n; ++i) if (Y[i] <m) r[Y[i]] =B[i</pre>
         return {z,r};
pair<Fraction, vector<Fraction>> maximize_int() {
         while (1) {
                  auto sol=maximize();
                 bool all int=true;
                  for(auto &x:sol.second)all_int&=x
                      .fractional part() == ZERO;
                  if(all_int)return sol;
                  Fraction nw b=ZERO;
                  int id=-1;
                  for (int i=0; i < n; ++i) {</pre>
                          Fraction fp=B[i].
                              fractional part();
                          if (fp>=nw b) nw b=fp, id=i;
                  vector<Fraction> nw a;
                  for (auto &x:A[id]) nw a.push back
                      (-x.fractional part());
                  A.push_back(nw_a);
                  B.push_back(-nw_b);
                  Y.push_back(n+m);n++;
```

7.27 Totient y Divisores

};

```
vector<int> count_divisors_sieve() {
   bitset<mx> is_prime; is_prime.set();
   vector<int> cnt(mx, 1);
   is_prime[0] = is_prime[1] = 0;
   for(int i = 2; i < mx; i++) {
        if(!is_prime[i]) continue;
        cnt[i]++;
        for(int j = i+i; j < mx; j += i) {
            int n = j, c = 1;
            while( n%i == 0 ) n /= i, c++;
            cnt[j] *= c;
            is_prime[j] = 0;
      }
    }
   return cnt;</pre>
```

```
vector<int> euler phi sieve() {
        bitset<mx> is_prime; is_prime.set();
        vector<int> phi(mx);
        iota(phi.begin(), phi.end(), 0);
        is_prime[0] = is_prime[1] = 0;
        for(int i = 2; i < mx; i++) {</pre>
                if(!is prime[i]) continue;
                for(int j = i; j < mx; j += i) {
                         phi[j] -= phi[j]/i;
                         is prime[j] = 0;
        return phi;
ll euler phi(ll n) {
        ll ans = n;
        for(ll i = 2; i * i <= n; ++i) {</pre>
                if(n % i == 0) {
                         ans -= ans / i;
                         while (n % i == 0) n /= i;
        if(n > 1) ans -= ans / n;
        return ans:
```

7.28 Xor Basis

```
template<typename T = int, int B = 31>
struct Basis {
        T a[B];
        Basis() {
                memset(a, 0, sizeof a);
        void insert(T x){
                for (int i = B - 1; i >= 0; i--) {
                        if (x >> i & 1) {
                                if (a[i]) x ^= a[i];
                                else {
                                         a[i] = x;
                                         break:
        bool can(T x) {
                for (int i = B - 1; i >= 0; i--) {
                        x = min(x, x ^a a[i]);
                return x == 0;
        T \max xor(T ans = 0) {
                for(int i = B - 1; i >= 0; i--) {
                        ans = max(ans, ans ^a[i]);
```

```
}
return ans;
}
};

// Basis<long long, 63> B;
// Cantidad de xor diferentes es 2^sz(base)
// Cantidad de subsets xor = 0 es 2^(n-sz(base))
```

8 Programacion dinamica

8.1 Bin Packing

```
int main() {
        ll n, capacidad;
        cin >> n >> capacidad;
        vl pesos(n, 0);
        forx(i, n) cin >> pesos[i];
        vector<pll> dp((1 << n));
        dp[0] = \{1, 0\};
        // dp[X] = \{\#numero de paquetes, peso de min
            paquete}
        // La idea es probar todos los subset y en cada
            uno preguntarnos
        // quien es mejor para subirse de ultimo buscando
             minimizar
        // primero el numero de paquetes
        for (int subset = 1; subset < (1 << n); subset++)</pre>
                 dp[subset] = \{21, 0\};
                 for (int iPer = 0; iPer < n; iPer++) {</pre>
                         if ((subset >> iPer) & 1) {
                                 pll ant = dp[subset ^ (1
                                     << iPer) 1:
                                 ll k = ant.ff;
                                 ll w = ant.ss;
                                 if (w + pesos[iPer] >
                                     capacidad) {
                                          k++;
                                          w = min(pesos[
                                             iPer], w);
                                 } else {
                                          w += pesos[iPer];
                                 dp[subset] = min(dp[
                                     subset], \{k, w\});
```

```
cout << dp[(1 << n) - 1].ff << ln;
}</pre>
```

8.2 Convex Hull Trick

```
// - Me dan las pendientes ordenadas
// Caso 1: Me hacen las querys ordenadas
// O(N + O)
// Caso 2: Me hacen querys arbitrarias
// O(N + QlogN)
struct CHT {
    // funciona tanto para min como para max, depende del
        orden en que pasamos las lineas
    struct Line {
        int slope, yIntercept;
        Line (int slope, int yIntercept) : slope (slope),
           yIntercept (yIntercept) { }
        int val(int x) { return slope * x + yIntercept; }
        int intersect(Line y) {
            return (y.yIntercept - yIntercept + slope - y
                .slope - 1) / (slope - y.slope);
    };
    deque<pair<Line, int>> dq;
    void insert(int slope, int yIntercept){
        // lower hull \sin m1 < m2 < m3
        // upper hull si si m1 > m2 > m3
        Line newLine(slope, yIntercept);
        while (!dq.empty() && dq.back().second >= dq.back
            ().first.intersect(newLine)) dq.pop_back();
        if (dq.empty()) {
            dq.emplace_back(newLine, 0);
            return;
        dq.emplace back(newLine, dq.back().first.
           intersect(newLine));
    int query(int x) { // cuando las consultas son
       crecientes
        while (dq.size() > 1) {
            if (dq[1].second <= x) dq.pop_front();</pre>
            else break;
        return dq[0].first.val(x);
    int query2(int x) { // cuando son arbitrarias
        auto qry = *lower_bound(dq.rbegin(), dq.rend(),
                                make_pair(Line(0, 0), x),
```

8.3 CHT Dynamic

```
// O((N+Q) \log N) < -usando set para add y bs para q
// lineas de la forma mx + b
#pragma once
struct Line {
        mutable 11 m, b, p;
        bool operator<(const Line& o) const { return m <</pre>
        bool operator<(ll x) const { return p < x; }</pre>
};
struct CHT : multiset<Line, less<>> {
        // (for doubles, use inf = 1/.0, div(a,b) = a/b)
        static const ll inf = LLONG MAX;
        static const bool mini = 0; // <---- 1 FOR MIN
        ll div(ll a, ll b) { // floored division
                return a / b - ((a ^ b) < 0 && a % b); }
        bool isect(iterator x, iterator y) {
                if (y == end()) return x \rightarrow p = inf, 0;
                if (x->m == y->m) x->p = x->b > y->b?
                    inf : -inf;
                 else x->p = div(y->b - x->b, x->m - y->m)
                 return x->p >= y->p;
        void add(ll m, ll b) {
                if (mini) { m \star= -1, b \star= -1; }
                 auto z = insert(\{m, b, 0\}), y = z++, x =
                 while (isect(v, z)) z = erase(z);
                 if (x != begin() \&\& isect(--x, y)) isect(
                    x, y = erase(y);
                 while ((y = x) != begin() \&\& (--x)->p >=
                    y->p)
                         isect(x, erase(y));
        11 query(ll x) {
                 assert(!empty());
                 auto l = *lower bound(x);
                if (mini) return -1.m * x + -1.b;
                 else return 1.m * x + 1.b;
};
```

8.4 Digit DP

```
// dp[pos][count of d][limit]
11 dp[20][20][2];
int k,d;
// count numbers <= c with k occurrences of d
ll dfs(string& c, int x=0, int y=0, bool z=0){
        if (dp[x][y][z]!=-1) return dp[x][y][z];
        dp[x][y][z] = (y = = k);
        if(x==(int)c.size())
                return dp[x][y][z];
        int limit=9;
        if(!z){
                limit=c[x]-'0';
        dp[x][y][z]=0;
        for(int i=0;i<=limit;++i){</pre>
                if(z)dp[x][y][z]+=dfs(c, x+1, y+(i==d), z
                else dp[x][y][z] += dfs(c, x+1, y+(i==d), i
                    imit);
        return dp[x][y][z];
// count(0,m) - count(0,n-1) = count(n,m)
ll query(ll n, ll m){
        string s1=to string(m);
        string s2=to string(n-111);
        memset(dp, -1, sizeof(dp));
        ll ans=dfs(s1);
        if (n<=011) return ans; // check</pre>
        memset(dp, -1, sizeof(dp));
        return ans-dfs(s2);
```

8.5 Divide Conquer

```
PROGRAMACION DINAMICA
```

```
pair<long long, int> best = {LLONG MAX, -1};
        for (int k = optl; k \le min(mid, optr); k++) {
                best = min(best, \{(k ? dp before[k - 1] :
                     0) + C(k, mid), k);
        dp cur[mid] = best.first;
        int opt = best.second;
        compute(l, mid - 1, optl, opt);
        compute(mid + 1, r, opt, optr);
int solve() {
        for (int i = 0; i < n; i++)</pre>
                dp before[i] = C(0, i);
        for (int i = 1; i < m; i++) {
                compute (0, n - 1, 0, n - 1);
                dp_before = dp_cur;
        return dp before[n - 1];
```

8.6 Edit Distances

```
int editDistances(string& wor1, string& wor2) {
         // O(tam1*tam2)
         // minimo de letras que debemos insertar, elminar
              o reemplazar
         // de worl para obtener wor2
         11 tam1=wor1.size();
         11 tam2=wor2.size();
         vector\langle vl \rangle dp(tam2+1, vl(tam1+1,0));
         for (int i=0; i<=tam1; i++) dp [0] [i]=i;</pre>
         for (int i=0; i<=tam2; i++) dp[i][0]=i;</pre>
         dp[0][0]=0;
         for(int i=1;i<=tam2;i++) {</pre>
                  for(int j=1; j<=tam1; j++) {
                          [1] op1 = min(dp[i-1][j], dp[i][j]
                              -11)+1;
                           // el minimo entre eliminar o
                              insertar
                           11 \text{ op2} = dp[i-1][j-1]; //
                              reemplazarlo
                           if (wor1[j-1]!=wor2[i-1]) op2++;
                           // si el reemplazo tiene efecto o
                                quedo iqual
                           dp[i][j]=min(op1,op2);
         return dp[tam2][tam1];
```

8.7 Kadane 2D

```
int main() {
        11 fil,col;cin>>fil>>col;
        vector<vl> grid(fil,vl(col,0));
// Algoritmo de Kadane/DP para suma maxima de una matriz
   2D en o(n^3)
        for(int i=0;i<fil;i++) {</pre>
                 for (int e=0; e < col; e++) {</pre>
                          11 num; cin>>num;
                          if (e>0) grid[i][e]=num+grid[i][e
                              -11;
                           else grid[i][e]=num;
        ll maxGlobal = LONG_LONG_MIN;
        for (int l=0; l<col; l++) {</pre>
                 for(int r=1;r<col;r++) {</pre>
                          11 maxLoc=0;
                          for(int row=0;row<fil;row++) {</pre>
                                   if (1>0) maxLoc+=grid[row
                                       ][r]-grid[row][l-1];
                                   else maxLoc+=grid[row][r
                                       ];
                                   if (maxLoc<0) maxLoc=0;</pre>
                                   maxGlobal= max(maxGlobal,
                                       maxLoc);
```

8.8 Knuth

```
// C[b][c] <= C[a][d]
// C[a][c] + C[b][d] <= C[a][d] + C[b][c] where a < b < c
int solve() {
        int N;
        ... // read N and input
        int dp[N][N], opt[N][N];
        auto C = [\&] (int i, int j) {
                ... // Implement cost function C.
        };
        for (int i = 0; i < N; i++) {</pre>
                opt[i][i] = i;
                ... // Initialize dp[i][i] according to
                    the problem
```

```
8.9 LIS
```

```
0
```

```
9 STRINGS
```

8.9 LIS

```
// O(n*log(n))
// retorna los indices de un lis
// cambiar el tipo y revisar si permite iguales
typedef int T;
vi lis(vector<T>& a, bool equal) {
        vi prev(sz(a));
        typedef pair<T, int> p;
        vector res;
        for (int i=0; i < sz(a); ++i) {</pre>
                 auto it=lower_bound(all(res), p{a[i],(
                    equal?i:0)});
                 if (it==res.end()) res.emplace_back(), it=
                    res.end()-1;
                 *it={a[i],i};
                 prev[i] = (it == res.begin())?0:(it-1)->
                    second;
        int l=sz(res), act=res.back().second;
        vi ans(1);
        while(l--) ans[l] = act, act = prev[act];
        return ans;
```

8.10 SOS

```
const int bits = 23;
int dp[1<<bits];
// O(n*2^n)
void SOS() {</pre>
```

```
for (int i = 0; i < (1 << bits); ++i) dp[i] = A[i]
        // top - down (informacion de las submascaras)
        for(int i = 0; i < bits; ++i){</pre>
                 for(int s = 0; s < (1 << bits); ++s){</pre>
                          if(s & (1 << i)){
                                   dp[s] += dp[s ^ (1 << i)
        // bottom - up (informacion de las supermascaras)
        for(int i = 0; i < bits; ++i){</pre>
                 for (int s = (1 << bits) - 1; s >= 0; --s)
                          if(s & (1 << i)){
                                   dp[s ^ (1 << i)] += dp[s
                                      1;
int dp2[1<<bits][bits+1];</pre>
// O(2^n*n^2)
void cnt(){
        vector<int> a;
        for (int x:a) dp2[x][0]++;
        // dp[s][c] = number of s^ai with c bits
        for(int i=0;i<bits;++i) {</pre>
                 for(int c=i;c>=0;--c){
                          for (int s=0; s<(1<<bits); ++s) {</pre>
                                   dp2[s^{(1<< i)}][c+1] += dp2[s
                                      ][c];
```

9 Strings

9.1 Aho Corasick

```
// 1) init() trie and add() strings
// 2) build() aho-corasick
// 3) process the text
// 4) dfs to calculate dp

// suf: longest proper suffix that's also in the trie
// dad: closest suffix link that is terminal
```

```
// cnt: number of strings that end exactly at node v
const int maxn = 2e5+5;
const int alpha = 26;
vector<int> adj[maxn];
int to[maxn][alpha], cnt[maxn], dad[maxn], suf[maxn], act; //
    not to change
int conv(char ch) {return ((ch>='a' && ch<='z')?ch-'a':ch-
   'A'+26);}
void init(){
        for(int i=0;i<=act;++i){</pre>
                suf[i]=cnt[i]=dad[i]=0;
                adi[i].clear();
                memset(to[i], 0, sizeof(to[i]));
        act=0;
int add(string& s){
        int u=0:
        for(char ch:s){
                int c=conv(ch);
                if(!to[u][c])to[u][c]=++act;
                u=to[u][c];
        cnt[u]++;
        return u;
// O(sum(|s|)*alpha)
void build() {
        queue<int> q{{0}};
        while(!q.empty()){
                int u=q.front();q.pop();
                for(int i=0;i<alpha;++i){</pre>
                         int v=to[u][i];
                         if(!v)to[u][i]=to[suf[u]][i];
                         else q.push(v);
                         if(!u || !v)continue;
                         suf[v]=to[suf[u]][i];
                         dad[v]=cnt[suf[v]]?suf[v]:dad[suf
                             [V]];
        for(int i=1;i<=act;++i) {</pre>
                adj[i].push_back(dad[i]);
                adj[dad[i]].push back(i);
```

9.2 Hashing

```
// O(n) build - O(1) get // 1. prepare() in the main
```

```
// 2. hashing<string> hs("hello");
// 3. hs.get(1,r);
// Chars are in [1, BASE]
// BASE is prime or random(lim, MOD-lim)
// If chars are in [0, BASE) then compare the hashes for
   lenath
// 1000234999, 1000567999, 1000111997, 1000777121,
   1001265673, 1001864327, 999727999, 1070777777
// if hash(multiset 1) == hash(multiset 2) then (r+a1)*(r+a1)
   a2)...(r+an) == (r+b1) * (r+b2)...(r+bn) // (Collision n/a)
   MOD)
const ii BASE (257, 367);
const int MOD[2] = { 1001864327, 1001265673 };
int add(int a, int b, int m) {return a+b>=m?a+b-m:a+b;}
int sbt(int a, int b, int m){return a-b<0?a-b+m:a-b;}</pre>
int mul(int a, int b, int m) {return ll(a) *b%m;}
11 operator ! (const ii a) { return (11(a.first) << 32) |</pre>
    a.second; }
ii operator + (const ii& a, const ii& b) {return {add(a.
   first, b.first, MOD[0]), add(a.second, b.second, MOD
   [1])};}
ii operator - (const ii& a, const ii& b) {return {sbt(a.
   first, b.first, MOD[0]), sbt(a.second, b.second, MOD
ii operator * (const ii& a, const ii& b) {return {mul(a.
   first, b.first, MOD[0]), mul(a.second, b.second, MOD
   [1])};}
const int maxn = 1e6+5;
ii pot[maxn];
void prepare() { // remember!!!
        pot[0] = ii\{1,1\};
        rep(i,1,maxn) pot[i] = pot[i-1] * BASE;
template <class type>
struct Hashing{
        vector<ii> h;
        Hashing(type& t) {
                h.assign(sz(t)+1, ii\{0,0\});
                rep(i, 1, sz(h)) h[i] = h[i-1] * BASE + ii{
                    t[i-1], t[i-1]};
        ii get(int l, int r){
                return h[r+1] - h[1] * pot[r-1+1];
};
ii combine(ii a, ii b, int lenb){
        return a * pot[lenb] + b;
```

9.3 Hashing 2D

```
// Revisar primero los comentarios en hashing!!!
// 1-indexed
const ii BX(3731, 3731), BY(2999, 2999);
const int MOD[2] = { 998244353, 1001265673 };
int add(int a, int b, int m) {return a+b>=m?a+b-m:a+b;}
int sbt(int a, int b, int m) {return a-b<0?a-b+m:a-b;}</pre>
int mul(int a, int b, int m) {return ll(a) *b%m;}
11 operator ! (const ii a) { return (ll(a.first) << 32) |</pre>
    a.second; }
ii operator + (const ii& a, const ii& b) {return {add(a.
   first, b.first, MOD[0]), add(a.second, b.second, MOD
ii operator - (const ii& a, const ii& b) {return {sbt(a.
   first, b.first, MOD[0]), sbt(a.second, b.second, MOD
   [1])};}
ii operator * (const ii& a, const ii& b) {return {mul(a.
   first, b.first, MOD[0]), mul(a.second, b.second, MOD
   [1])};}
const int maxn = 1e6+5;
ii PX[maxn], PY[maxn];
void prepare() { // remember!!!
        PX[0] = PY[0] = ii\{1,1\};
        rep(i,1,maxn) {
                 PX[i] = PX[i-1] * BX;
                 PY[i] = PY[i-1] * BY;
template <class type>
struct Hashing2D { // 1-indexed
        vector<vector<ii>>> hs;
        int n, m;
        Hashing2D(vector<type>& s) {
                 n = sz(s), m = sz(s[0]);
                hs.assign(n + 1, vector\langle ii \rangle(m + 1, \{0,0\})
                 rep(i, 0, n) rep(j, 0, m)
                         hs[i + 1][j + 1] = {s[i][j], s[i]}
                            ][†]};
                rep(i, 0, n+1) rep(j, 0, m)
                         hs[i][j+1] = hs[i][j+1] + hs[
                            i][j] * BY;
                 rep(i, 0, n) rep(j, 0, m+1)
                         hs[i + 1][j] = hs[i + 1][j] + hs[
                            i][j] * BX;
        ii get(int x1, int y1, int x2, int y2) {
                 assert (1 \leq x1 && x1 \leq x2 && x2 \leq n);
                 assert (1 \leq y1 && y1 \leq y2 && y2 \leq m);
                x1--; y1--;
                 int dx = x^2 - x^1, dy = y^2 - y^1;
```

9.4 KMP

```
// O(n)
vector<int> phi(string& s){
        int n=sz(s):
        vector<int> tmp(n);
        for(int i=1, j=0; i<n; ++i) {
                 while(j>0 && s[j]!=s[i])j=tmp[j-1];
                 if(s[i]==s[j])j++;
                 tmp[i]=i;
        return tmp;
// O(n+m)
int kmp(string& s, string& p){
        int n=sz(s), m=sz(p), cnt=0;
        vector<int> pi=phi(p);
        for (int i=0, j=0; i < n; ++i) {</pre>
                 while ( j && s[i]!=p[j]) j=pi[j-1];
                 if(s[i]==p[j])j++;
                 if ( j==m) {
                          cnt++;
                          j=pi[j-1];
        return cnt;
```

9.5 KMP Automaton

```
p=to[p][conv(s[i])];
```

9.6 Lyndon Factorization

```
// A string is called simple if it is strictly smaller
   than all its nontrivial cyclic shifts.
// The Lyndon factorization of the string is s = w1 \ w2
// where all strings wi are simple, and they are in non-
   increasing order
// w1 >= w2 >= ... >= wk
// this factorization exists and it is unique
// O(n)
vector<string> duval(string& s){
        vector<string> factorization;
        int n=sz(s), i=0;
        while(i<n) {</pre>
                 int j=i+1, k=i;
                 while(j < n \& \& s[k] <= s[j]) {
                         if(s[k]<s[j])k=i;
                         else k++;
                         j++;
                 while (i \le k) {
                         factorization.push_back(s.substr(
                             i, j-k));
                         i += j - k;
        return factorization;
```

9.7 Manacher

```
// O(n), par (raiz, izg, der) 1 - impar 0
vi manacher(string& s, int par) {
         int l=0, r=-1, n=sz(s);
         vi m(n, 0);
         for (int i=0; i < n; ++i) {</pre>
                  int k=(i>r?(1-par):min(m[1+r-i+par], r-i
                      +par))+par;
                  while (i+k-par < n \&\& i-k>=0 \&\& s[i+k-par]==
                      s[i-k])++k;
                  m[i]=k-par; --k;
                  if (i+k-par>r) l=i-k, r=i+k-par;
         for (int i=0; i < n; ++i) m[i] = (m[i]-1+par) *2+1-par;</pre>
         return m;
```

9.8 Minimum Expression

```
// O(n)
int minimum_expression(string s){
        s=s+s;int n=sz(s), i=0, j=1, k=0;
        while (i+k< n \&\& j+k< n) {
                 if (s[i+k] == s[j+k])k++;
                 else if (s[i+k]>s[j+k])i=i+k+1, k=0; // <
                 else j=j+k+1, k=0;
                 if(i==j)j++;
        return min(i, j);
```

9.9 Next Permutation

```
// O(n)
// 1) find the last i such that ai <ai+1
// 2) find the last j such that ai<aj
// 3) swap i and j, then reverse the segment [i+1, n-1]
string nextPermutation(string& s){
        string ans(s);
        int n=sz(s);
        int i=n-2;
        while(i>=0 && ans[i]>=ans[i+1])i--;
        if(i<0)return "no permutation";</pre>
        int j=n-1;
        while (ans[i]>=ans[j])j--;
        swap(ans[i], ans[j]);
        int l=i+1, r=n-1;
        while (r>1) swap (ans[r--], ans[l++]);
        return ans;
```

9.10 Palindromic Tree

```
const int alpha = 26;
const char mini = 'a';
// tree.suf: the longest suffix-palindrome link
// tree.dad - tree.to: the parent palindrome by removing
   the first and last character
// node 0 = root with len -1 for odd
// node 1 = root with len 0 for even
struct Node {
    int to[alpha], suf, len, cnt, dad;
   Node (int x, int l = 0, int c = 1): len(x), suf(l),
       cnt(c) {
```

```
memset(to, 0, sizeof(to));
    int& operator [] (int i) { return to[i]; }
} ;
struct PalindromicTree {
    vector<Node> tree:
        vector<int> palo; // longest suffix-palindrome in
            the position i
    string s;
    int n, last; // max suffix palindrome
    PalindromicTree(string t = "") {
        n = last = 0;
        tree.push back(Node(-1));
        tree.push back(Node(0));
                for(char& c:t) add_char(c);
                // Propagate counts up the suffix links
                for(int i=sz(tree)-1;i>=2;i--){
                        tree[tree[i].suf].cnt+=tree[i].
    int getsuf(int p) {
        while (n - tree[p].len - 1 < 0 || s[n - tree[p].
           len - 1] != s[n])
                        p = tree[p].suf;
        return p;
    void add char(char ch) {
        s.push back(ch);
        int p = getsuf(last), c = ch - mini;
        if (!tree[p][c]) {
            int suf = getsuf(tree[p].suf);
            suf = max(1, tree[suf][c]);
            tree[p][c] = sz(tree);
            tree.push back(Node(tree[p].len + 2, suf, 0))
        last = tree[p][c];
        tree[last].dad = p;
        tree[last].cnt++;n++;
                palo.push back(tree[last].len);
} ;
```

9.11 Suffix Array

```
// O(n*log(n)) - char in [1, lim)
// sa: is the starting position of the i-th lex smallest
    suffix
// rnk: is the rank (position in SA) of the suffix
    starting at i
```

```
// lcp: is the longest common prefix between sa[i] and sa
auto SuffixArray(string s, int lim=256) {
        s.push\_back(0); int n = sz(s), k = 0, a, b;
        vi sa, lcp, rnk(all(s)), y(n), ws(max(n, lim));
        sa = lcp = y, iota(all(sa), 0);
        for (int j = 0, p = 0; p < n; j = max(1, j * 2),
           lim = p) {
                p = j, iota(all(y), n - j);
                rep(i, 0, n) if (sa[i] >= j) y[p++] = sa[i]
                    - j;
                fill(all(ws), 0);
                rep(i,0,n) ws[rnk[i]]++;
                rep(i,1,lim) ws[i] += ws[i-1];
                for (int i = n; i--;) sa[--ws[rnk[y[i]]]]
                    = y[i];
                swap(rnk, y), p = 1, rnk[sa[0]] = 0;
                rep(i, 1, n) a = sa[i - 1], b = sa[i], rnk[
                        (y[a] == y[b] && y[a + j] == y[b]
                           + j]) ? p - 1 : p++;
        for (int i = 0, i; i < n - 1; lcp[rnk[i++]] = k)
                for (k \& \& k--, j = sa[rnk[i] - 1]; s[i +
                   k] == s[j + k]; k++);
        reverse(all(lcp));lcp.pop_back();reverse(all(lcp)
        return tuple{sa, rnk, lcp};
```

9.12 Suffix Automaton

```
// O(n*log(alpha))
// suf: suffix link (like aho if not match)
// len: length of the longest string in this state
// minlen: smallest string of node v = (v.suf == -1.20:v.suf
   .len) + 1
// end: if this state is terminal
// count different strings [i.suf.len+1, i.len]
// para saber cuantos substrings itnee a en b, ir
   procesando los
// prefijos y al marcarlos procesar la cantidad visitando
    los sufijos de los nodos
// contribucion es u.len - u.suf.len, tener en cuenta con
    que len se llego
// puede ser un len2 para manejar eso min(u.len, u.len2) -
   u.suf.len
// a->b->c->b->c
//b \rightarrow c
template<int alpha = 26>
struct SuffixAutomaton {
        struct Node {
```

```
// array<int, alpha> to; TLE, add -> int
            conv(char ch)
        map<char, int> to;
        int len = 0, suf = 0;
        bool end = false;
};
vector<Node> sa;
int last = 0;
ll substrs = 0;
SuffixAutomaton(string &s) {
        sa.reserve(sz(s) *2);
        last = add_node();
        sa[0].suf = -1;
        for (char &c : s) add_char(c);
        for (int p = last; p; p = sa[p].suf) sa[p
            l.end = 1;
int add_node() { sa.push_back({}); return sz(sa)
void add char(char c) {
        int u = add_node(), p = last;
sa[u].len = sa[last].len + 1;
        while (p != -1 \&\& !sa[p].to.count(c)) {
                 sa[p].to[c] = u;
                substrs += p != 0 ? sa[p].len -
                    sa[sa[p].suf].len : 1;
                 p = sa[p].suf;
        if (p != -1) {
                 int q = sa[p].to[c];
                if (sa[p].len + 1 != sa[q].len) {
                         int clone = add node();
                         sa[clone] = sa[q];
                         sa[clone].len = sa[p].len
                              + 1;
                         sa[q].suf = sa[u].suf =
                             clone;
                         while (p !=-1 \&\& sa[p].
                             to[\bar{c}] == q) {
                                  sa[p].to[c] =
                                     clone;
                                  p = sa[p].suf;
                 } else sa[u].suf = q;
        last = u;
// Aplicaciones
int dfs(int u) { // count
        if (sa[u].cnt!=-1) return sa[u].cnt;
        sa[u].cnt=sa[u].end;
```

```
for(auto [ ,v]:sa[u].to){
                sa[u].cnt+=dfs(v);
        return sa[u].cnt;
void dfs2(int u) { // grade primero
        sa[u].pre--;
        if (sa[u].pre>0) return;
        for (auto [_, v]:sa[u].to) {
                sa[v].cnt2+=sa[u].cnt2;
                dfs2(v);
void dfs2(){
        vector<int> order(sz(sa)-1);
        for (int i=1; i < sz (sa); ++i) order[i-1]=i;</pre>
        sort(order.begin(), order.end(), [&](int
           a, int b) { return sa[a].len > sa[b].
           len; });
        for(auto &i : order) {
                // suf.cnt += i.cnt
int lcs(string& t){
        int u=0, l=0, ans=0;
        for(char c:t) {
                while(u && !sa[u].to.count(c)){
                         u=sa[u].suf;
                         l=sa[u].len;
                if(sa[u].to.count(c)){
                         u=sa[u].to[c];
                         1++;
                ans=max(ans, 1);
        return ans:
bool query(string& t){
        int u=0;
        for(char c:t){
                if(!sa[u].to.count(c))return
                    false;
                u=sa[u].to[c];
        return true;
void cyclic(string& t) { // dfs(0) primero
        int u=0,1=0;
        int m=sz(t);
        t+=t;
        unordered set<int> s; // vector<bool>
```

```
9.13 Suffix Tree
```

STRINGS

```
\frac{\infty}{7}
```

```
9.13 Suffix Tree
```

};

```
// O(n)
// pos: start of the edge
// len: edge length
// link: suffix link
struct SuffixTree{
        vector<map<char,int>> to;
        vector<int> pos,len,link;
        int size=0,inf=1e9;
        string s;
        int make(int _pos, int _len) {
                to.push back(map<char,int>());
                pos.push back (pos);
                len.push back (len);
                link.push_back(-1);
                return size++;
        void add(int& p, int& lef, char c) {
                s+=c;++lef;int lst=0;
                for(;lef;p?p=link[p]:lef--){
                        while (lef>1 && lef>len[to[p][s[sz
                            (s) - lef[]]) {
                                 p=to[p][s[sz(s)-lef]], lef
                                    -=len[p];
```

for(char ch:t){

11 ans=0;

cout << ans << "\n";

int c=conv(ch);

if(sa[u].to[c]){

if(l==m){

for(int u:s)ans+=sa[u].cnt;

1++;

}else{

while(u && !sa[u].to[c]){

u=sa[u].suf;

l=sa[u].len;

u=sa[u].to[c];

s.insert(u);

if(sa[u].minlen==m) {

1--;

u=sa[u].suf;

l=sa[u].len;

```
char e=s[sz(s)-lef];
                         int& q=to[p][e];
                         if(!q){
                                  q=make(sz(s)-lef,inf),
                                     link[lst]=p,lst=0;
                         }else{
                                  char t=s[pos[q]+lef-1];
                                 if (t==c) {link[lst]=p;
                                     return; }
                                 int u=make(pos[q],lef-1);
                                 to[u][c]=make(sz(s)-1, inf
                                     );
                                 to[u][t]=q;
                                 pos[a] += lef -1;
                                 if(len[q]!=inf)len[q]=
                                     lef-1;
                                 q=u,link[lst]=u,lst=u;
        SuffixTree(string& s){
                make (-1, 0); int p=0, lef=0;
                for(char c:_s)add(p,lef,c);
                add(p,lef,'$'); // smallest char
                 s.pop back();
        int query(string& p) {
                 for(int i=0, u=0, n=sz(p);;) {
                         if(i==n || !to[u].count(p[i]))
                             return i;
                         u=to[u][p[i]];
                         for (int j=0; j<len[u];++j) {</pre>
                                 if(i==n || s[pos[u]+j]!=p
                                     [i])return i;
                                 i++;
        vector<int> sa;
        void genSA(int x=0, int Len=0){
                if(!sz(to[x]))sa.push back(pos[x]-Len);
                else for (auto t:to[x]) genSA (t.second, Len+
                    len[x]);
};
```

9.14 Trie

```
const int maxn = 2e6+5;
const int alpha = 26;
```

```
int to[maxn][alpha]; // to[ul[c]: node u edge with the
   letter c
int cnt[maxn]; // count of word ending in this node
int act: // trie node cound
int conv(char ch) {return ((ch>='a' && ch<='z')?ch-'a':ch-</pre>
   'A'+26);}
void init(){
        for(int i=0;i<=act;++i){
                memset(to[i],0,sizeof(to[i]));
                 cnt[i]=0;
        act=0;
void add(string& s){
        int u=\bar{0};
        for(char ch:s){
                int c=conv(ch);
                if(!to[u][c])to[u][c]=++act;
                u=to[u][c];
        cnt[u]++;
```

9.15 Trie Bit

```
const int maxn = 5e5+5;
const int bits = 30;
const int alpha = 2;
int to[maxn*bits][alpha]; // to[u][c]: node u edge with
   the letter c
int cnt[maxn*bits]; // count of word ending in this node
int act; // trie node cound
int conv(int x, int i) {return ((x&(1<<i))?1:0);}</pre>
void init(){
        for (int i=0; i <= act; ++i) {</pre>
                 memset(to[i],0,sizeof(to[i]));
                 cnt[i]=0;
        act=0;
void add(int x){
        int u=0;
        for(int i=bits;i>=0;--i){
                 int c=conv(x,i);
                 if(!to[u][c])to[u][c]=++act;
                 cnt[u]++;
                 u=to[u][c];
        cnt[u]++;
```

9.16 Z Algorithm

9.17 El especial

```
#include<bits/stdc++.h>
using namespace std;
const int N = 1e5 + 9;
struct ST {
  \#define lc (n << 1)
  \#define rc ((n \ll 1) | 1)
  long long t[4 * N], lazy[4 * N];
  ST() {
    memset(t, 0, sizeof t);
    memset(lazy, 0, sizeof lazy);
  inline void push(int n, int b, int e) {
    if (lazy[n] == 0) return;
    t[n] = t[n] + lazy[n] * (e - b + 1);
    if (b != e) {
      lazy[lc] = lazy[lc] + lazy[n];
      lazy[rc] = lazy[rc] + lazy[n];
```

```
9.17 El especial
```

```
lazv[n] = 0;
  inline long long combine(long long a, long long b) {
    return a + b;
  inline void pull(int n) {
    t[n] = t[lc] + t[rc];
  void upd(int n, int b, int e, int i, int j, int v) {
    push(n, b, e);
    if (j < b || e < i) return;
    if (i <= b && e <= j) {
      lazy[n] = v; //set lazy
      push(n, b, e);
      return;
    int mid = (b + e) >> 1;
    upd(lc, b, mid, i, j, v);
    upd(rc, mid + 1, e, i, j, v);
    pull(n);
  long long query(int n, int b, int e, int i, int j) {
    push(n, b, e);
    if (i > e | | b > j) return 0; //return null
    if (i <= b && e <= j) return t[n];</pre>
    int mid = (b + e) >> 1;
    return combine (query (lc, b, mid, i, j), query (rc, mid
        + 1, e, i, j));
}st;
struct node
  int len, link, firstpos;
 map<char, int> nxt;
};
vector<node> t;
struct SuffixAutomaton {
  int sz, last;
  vector<int> terminal;
  vector<int> dp;
  vector<vector<int>> q;
  SuffixAutomaton() {}
  SuffixAutomaton(int n) {
    t.clear(); t.resize(2 * n);
    terminal.resize(2 * n, 0);
    dp.resize(2 * n, -1); sz = 1; last = 0;
    q.resize(2 * n);
    t[0].len = 0; t[0].link = -1; t[0].firstpos = 0;
  void extend(char c) {
    int p = last;
    int cur = sz++;
    t[cur].len = t[last].len + 1;
    t[cur].firstpos = t[cur].len;
```

```
p = last;
    while (p != -1 \&\& !t[p].nxt.count(c)) {
      t[p].nxt[c] = cur;
      p = t[p].link;
    if (p == -1) t[cur].link = 0;
      int q = t[p].nxt[c];
      if (t[p].len + 1 == t[q].len) t[cur].link = q;
      else {
        int clone = sz++;
        t[clone] = t[q];
        t[clone].len = t[p].len + 1;
        while (p != -1 \&\& t[p].nxt[c] == q) {
          t[p].nxt[c] = clone;
          p = t[p].link;
        t[q].link = t[cur].link = clone;
    last = cur;
};
pair<int, int> modifies[N * 2];
int cnt;
namespace lct {
  int par[N * 2], lazy[N * 2], last[N * 2], c[N * 2][2];
  void mark(int x, int v) {
    lazy[x] = last[x] = v;
 void push(int x) {
    if (lazy[x]) {
      if (c[x][0]) {
        mark(c[x][0], lazy[x]);
      if (c[x][1]) {
        mark(c[x][1], lazy[x]);
      lazv[x] = 0;
 bool is root(int x) {
    return c[par[x]][0] != x && c[par[x]][1] != x;
  void rotate(int x) {
    int y = par[x], z = par[y], k = c[y][1] == x;
    if (!is_root(y)) {
      c[z][c[z][1] == y] = x;
    par[c[y][k] = c[x][!k]] = y;
    par[par[c[x][!k] = y] = x] = z;
 void splay(int x) {
    static int st[N];
```

```
90
```

```
int dateToInt(int y, int m, int d){
         return 1461*(y+4800+(m-14)/12)/4+367*(m-2-(m-14)
            /12 * 12) / 12 -
                  3*((v+4900+(m-14)/12)/100)/4+d-32075;
void intToDate(int jd, int& y, int& m, int& d) {
         int x, n, i, j; x = jd + 68569;
         n=4*x/146097; x=(146097*n+3)/4;
         i = (4000 * (x+1)) / 1461001; x = 1461 * i / 4 - 31;
```

```
int top = 0;
    st[++top] = x;
    for (int i = x; !is_root(i); i = par[i]) {
      st[++top] = par[i];
    while (top) {
      push(st[top--]);
    while (!is_root(x)) {
      int y = par[x], z = par[y];
      if (!is root(y)) {
        rotate((c[y][1] == x) == (c[z][1] == y) ? y : x);
      rotate(x);
  void access(int x, int v) {
    int z = 0;
    cnt = 0;
    while (x) {
      splav(x);
      modifies[++cnt] = make pair(t[x - 1].len, last[x]);
      c[x][1] = z;
      mark(x, v);
      z = x;
      x = par[x];
int pos[N];
vector<pair<int, int>> Q[N];
long long ans[N];
int32 t main() {
  ios_base::sync_with_stdio(0);
  cin.tie(0);
  int n, q; cin >> n >> q;
  string s; cin >> s;
  SuffixAutomaton sa(n);
  for (int i = 1; i <= q; i++) {
    int 1, r; cin >> 1 >> r;
    ++1; ++r;
    Q[r].push back({l, i});
  s = "." + s;
  pos[0] = 1;
  for (int i = 1; i <= n; ++i) {
    sa.extend(s[i]);
    pos[i] = sa.last + 1;
  for (int i = 1; i <= sa.sz; ++i) {</pre>
    lct::par[i] = t[i - 1].link + 1;
  for (int i = 1; i <= n; ++i) {</pre>
    st.upd(1, 1, n, 1, i, 1);
    lct::access(pos[i], i);
```

```
int last = 0;
  for (int j = cnt; j > 1; --j) {
    pair<int, int> p = modifies[j];
    if (p.first) {
      if (p.second) {
        st.upd(1, 1, n, p.second - p.first + 1, p.
           second - last, -1);
      last = p.first;
  // st.query(l, l) = number of distinct substrings
     which lastly occured in starting position 1 for
     prefix [1, i]
  for (auto [1, id]: Q[i]) {
    ans[id] = st.query(1, 1, n, 1, i);
for (int i = 1; i <= q; i++) {
  cout << ans[i] << '\n';
return 0;
```

10 Misc

10.1 Counting Sort

```
// O(n+k)
void counting sort(vi& a) {
        int n=sz(a);
        int maxi=*max element(all(a));
        int mini=*min element(all(a));
        int k=maxi-mini+1;
        vi cnt(k,0);
        for (int i=0;i<n;++i)++cnt[a[i]-mini];</pre>
        for (int i=0, j=0; i < k; ++i)
                 while (cnt[i]--) a [j++]=i+mini;
```

10.2 Dates

```
10
```

10.3 Expression Parsing

```
// O(n) - eval() de python
bool delim(char c) {return c==' ';}
bool is op(char c){return c=='+' || c=='-' || c=='*' || c
   ==' /'; }
bool is unary (char c) {return c=='+' | | c=='-';}
int priority(char op){
        if(op<0) return 3;</pre>
        if(op=='+' || op=='-') return 1;
        if(op=='*' || op=='/') return 2;
        return -1;
void process op(stack<int>& st, char op){
        if(op<0){
                 int l=st.top();st.pop();
                 switch(-op) {
                         case '+':st.push(1);break;
                         case '-':st.push(-1);break;
        }else{
                 int r=st.top();st.pop();
                 int l=st.top();st.pop();
                 switch(op) {
                         case '+':st.push(l+r);break;
                         case '-':st.push(l-r);break;
                         case '*':st.push(l*r);break;
                         case '/':st.push(l/r);break;
int evaluate(string& s){
        stack<int> st;
        stack<char> op;
        bool may_be_unary=true;
        for (int i=0; i < sz(s); ++i) {
                if (delim(s[i])) continue;
                if(s[i] == '('){
                         op.push('(');
                         may_be_unary=true;
```

```
}else if(s[i]==')'){
                while (op.top()!='('){
                         process_op(st, op.top());
                         op.pop();
                op.pop();
                may be unary=false;
        }else if(is_op(s[i])){
                char cur op=s[i];
                if (may_be_unary && is_unary(
                    cur_op))cur_op=-cur_op;
                while(!op.empty() && ((cur op >=
                    0 && priority(op.top()) >=
                    priority(cur_op)) || (cur_op <</pre>
                     0 && priority(op.top()) >
                    priority(cur_op)))){
                         process_op(st, op.top());
                         op.pop();
                op.push(cur op);
                may be unary=true;
        }else{
                int number=0;
                while(i<sz(s) && isalnum(s[i]))</pre>
                    number=number *10+s[i++]-'0';
                st.push(number);
                may_be_unary=false;
while(!op.empty()){
        process_op(st, op.top());
        op.pop();
return st.top();
```

10.4 Hanoi

```
// hanoi(n) = 2 * hanoi(n-1) + 1
// hanoi(n, 1, 3)
vector<int> ans;
void hanoi(int x, int start, int end){
    if(!x)return;
    hanoi(x-1, start, 6-start-end);
    ans.push_back({start, end});
    hanoi(x-1, 6-start-end, end);
}
```

10.5 K mas frecuentes

// los k numeros mas frecuentes

```
// el cero es un valor neutral dentro del vector
// no usarlo en el array original (a[i] > 0, i \in [0, n-1])
// el vector guarda {valor, contador}
// pero contador es para el algo, no es la cantidad real
// algoritmo de misra-gries O(k^2)
vector<ii> null(k, {0,0});
vector<ii> init(int v){
        vector<ii> a=null;
        a[0] = \{v, 1\};
        return a;
vector<ii> oper(vector<ii> a, vector<ii> b, int k) {
        for (int i = 0; i < k; ++i) if (b[i].first) {
                int p = -1, q = -1;
                for (int j = 0; j < k; ++j) {
                        if (b[i].first == a[j].first) p =
                        if (!a[j].first) q = j;
                if (p !=-1) {
                        a[p].second += b[i].second;
                \} else if (q != -1) {
                        a[q] = b[i];
                } else {
                        int mn = b[i].second;
                        for (int j = 0; j < k; ++j) mn =
                            min(mn, a[i].second);
                        for (int j = 0; j < k; ++j) a[j].
                            second -= mn;
                        b[i].second -= mn;
                        for (int j = 0; j < k; ++j) if (!
                            a[i].second) {
                                if (b[i].second > 0) {
                                         a[j] = b[i], b[i]
                                            ].second = 0;
                                 } else {
                                         a[j].first = 0;
        return a:
```

10.6 Prefix3D

```
const int N = 100;
int A[N][N][N];
int preffix[N + 1][N + 1][N + 1];
void build(int n) {
    for (int x = 1; x <= n; x++) {
        for (int y = 1; y <= n; y++) {</pre>
```

```
for (int z = 1; z \le n; z++) {
                                preffix[x][y][z] = A[x -
                                    1][y - 1][z - 1]
                                        + preffix[x - 1][
                                            y][z] +
                                            preffix[x][v -
                                             1|[z] +
                                            preffix[x][y][
                                            z - 1]
                                         - preffix[x - 1][
                                            y - 1][z] -
                                            preffix[x -
                                            1][v][z - 1] -
                                             preffix[x][y
                                            - 1][z - 1]
                                        + preffix[x - 1][
                                            y - 1 [z - 1];
ll query(int lx, int rx, int ly, int ry, int lz, int rz){
        ll ans = preffix[rx][ry][rz]
                - preffix[lx - 1][ry][rz] - preffix[rx][
                   ly - 1][rz] - preffix[rx][ry][lz - 1]
                + preffix[lx - 1][ly - 1][rz] + preffix[
                   lx - 1 [ry] [lz - 1] + preffix[rx] [ly -
                    11[lz - 1]
                - preffix[lx - 1][ly - 1][lz - 1];
        return ans;
```

10.7 Ternary Search

```
int f(int x);
int ternary() {
    int l,r;
    while(r-l>6) {
        int m1=l+(r-l)/3;
        int m2=r-(r-l)/3;
        if(f(m1)<f(m2))l=m1; // revisar desempate
        else r=m2;
    }
    int ans=l,val=f(l);</pre>
```

11 Teoría y miscelánea

11.1 Sumatorias

•
$$\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$$

•
$$\sum_{i=1}^{n} i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

•
$$\sum_{i=1}^{n} i^5 = \frac{(n(n+1))^2 (2n^2 + 2n - 1)}{12}$$

$$\bullet \sum_{i=1}^{n} i^3 = \left(\frac{n(n+1)}{2}\right)^2$$

•
$$\sum_{i=0}^{n} x^i = \frac{x^{n+1}-1}{x-1}$$
 para $x \neq 1$

11.2 Teoría de Grafos

11.2.1 Teorema de Euler

En un grafo conectado planar, se cumple que V-E+F=2, donde V es el número de vértices, E es el número de aristas y F es el número de caras. Para varios componentes la formula es: V-E+F=1+C, siendo C el número de componentes.

11.2.2 Planaridad de Grafos

Un grafo es planar si y solo si no contiene un subgrafo homeomorfo a K_5 (grafo completo con 5 vértices) ni a $K_{3,3}$ (grafo bipartito completo con 3 vértices en cada conjunto).

11.2.3 Truco del Cow Game

Dadas restricciones de la forma:

$$x_a - x_b \le d$$

podemos transformar cada desigualdad en una arista dirigida:

$$b \to a \quad \text{con peso } d$$

```
for(int i=l+1;i<=r;++i) {
    int val2=f(i);
    if(val2>val) {
        val=val2;
        ans=i;
    }
} return val;
}
```

Luego, ejecutando un algoritmo de camino más corto desde un nodo inicial s, obtenemos:

$$\operatorname{dist}[i] = \max(x_i - x_s)$$

Nota: Pueden aparecer pesos negativos, por lo que se debe usar Bellman-Ford o SPFA, no Dijkstra.

11.3 Teoría de Números

11.3.1 Ecuaciones Diofánticas Lineales

Una ecuación diofántica lineal es una ecuación en la que se buscan soluciones enteras x e y que satisfagan la relación lineal ax+by=c, donde a, b y c son constantes dadas.

Para encontrar soluciones enteras positivas en una ecuación diofántica lineal, podemos seguir el siguiente proceso:

- 1. Encontrar una solución particular: Encuentra una solución particular (x_0, y_0) de la ecuación. Esto puede hacerse utilizando el algoritmo de Euclides extendido.
- 2. Encontrar la solución general: Una vez que tengas una solución particular, puedes obtener la solución general utilizando la fórmula:

$$x = x_0 + \frac{b}{\operatorname{mcd}(a, b)} \cdot t$$

$$y = y_0 - \frac{a}{\operatorname{mcd}(a, b)} \cdot t$$

donde t es un parámetro entero.

3. Restringir a soluciones positivas: Si deseas soluciones positivas, asegúrate de que las soluciones generales satisfagan $x \ge 0$ y $y \ge 0$. Puedes ajustar el valor de t para cumplir con estas restricciones.

11.3.2 Pequeño Teorema de Fermat

Si p es un número primo y a es un entero no divisible por p, entonces $a^{p-1} \equiv 1 \pmod{p}$.

11.3.3 Teorema de Euler

Para cualquier número entero positivo n y un entero a coprimo con n, se cumple que $a^{\phi(n)} \equiv 1 \pmod{n}$, donde $\phi(n)$ es la función phi de Euler, que representa la cantidad de enteros positivos menores que n y coprimos con n.

11.4 Geometría

11.4.1 Teorema de Pick

Sea un poligono simple cuyos vertices tienen coordenadas enteras. Si B es el numero de puntos enteros en el borde, I el numero de puntos enteros en el interior del poligono, entonces el area A del poligono se puede calcular con la formula:

$$A = I + \frac{B}{2} - 1$$

11.4.2 Fórmula de Herón

Si los lados del triángulo tienen longitudes a, b y c, y s es el semiperímetro (es decir, $s = \frac{a+b+c}{2}$), entonces el área A del triángulo está dada por:

$$A = \sqrt{s(s-a)(s-b)(s-c)}$$

11.4.3 Relación de Existencia Triangular

Para un triángulo con lados de longitud $a,\,b,\,{\bf y}\,c,$ la relación de existencia triangular se expresa como:

$$b-c < a < b+c, \quad a-c < b < a+c, \quad a-b < c < a+b$$

11.5 Combinatoria

11.5.1 Permutaciones

El número de permutaciones de n objetos distintos tomados de a r a la vez (sin repetición) se denota como P(n,r) y se calcula mediante:

$$P(n,r) = \frac{n!}{(n-r)!}$$

11.5.2 Combinaciones

El número de combinaciones de n objetos distintos tomados de a r a la vez (sin repetición) se denota como C(n,r) o $\binom{n}{r}$ y se calcula mediante:

$$C(n,r) = \binom{n}{r} = \frac{n!}{r!(n-r)!}$$

11.5.3 Permutaciones con Repetición

El número de permutaciones de n objetos tomando en cuenta repeticiones se denota como $P_{\text{rep}}(n; n_1, n_2, \dots, n_k)$ y se calcula mediante:

$$P_{\text{rep}}(n; n_1, n_2, \dots, n_k) = \frac{n!}{n_1! n_2! \cdots n_k!}$$

11.5.4 Combinaciones con Repetición

El número de combinaciones de n objetos tomando en cuenta repeticiones se denota como $C_{\text{rep}}(n; n_1, n_2, \dots, n_k)$ y se calcula mediante:

$$C_{\text{rep}}(n; n_1, n_2, \dots, n_k) = \binom{n+k-1}{n} = \binom{n+k-1}{k-1}$$

11.5.5 Números de Catalan

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

Los números de Catalan también pueden calcularse utilizando la siguiente fórmula recursiva:

$$C_0 = 1$$

$$C_{n+1} = \frac{4n+2}{n+2}C_n$$

Usos:

- Cat(n) cuenta el número de árboles binarios distintos con n vértices.
- Cat(n) cuenta el número de expresiones que contienen n pares de paréntesis correctamente emparejados.
- Cat(n) cuenta el número de formas diferentes en que se pueden colocar n+1 factores entre paréntesis, por ejemplo, para n=3 y 3+1=4 factores: a,b,c,d, tenemos: (ab)(cd),a(b(cd)),((ab)c)d y a((bc)d).

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- Los números de Catalan cuentan la cantidad de caminos no cruzados en una rejilla $n \times n$ que se pueden trazar desde una esquina de un cuadrado o rectángulo a la esquina opuesta, moviéndose solo hacia arriba y hacia la derecha.
- Los números de Catalan representan el número de árboles binarios completos con n+1 hojas.
- Cat(n) cuenta el número de formas en que se puede triangular un poligono convexo de n+2 lados. Otra forma de decirlo es como la cantidad de formas de dividir un polígono convexo en triángulos utilizando diagonales no cruzadas.

11.5.6 Estrellas y barras

Número de soluciones de la ecuación $x_1 + x_2 + \cdots + x_k = n$.

• Con $x_i \ge 0$: $\binom{n+k-1}{n}$

• Con $x_i \ge 1$: $\binom{n-1}{k-1}$

Número de sumas de enteros con límite inferior:

Esto se puede extender fácilmente a sumas de enteros con diferentes límites inferiores. Es decir, queremos contar el número de soluciones para la ecuación:

$$x_1 + x_2 + \cdots + x_k = n$$

 $con x_i \geq a_i$.

Después de sustituir $x'_i := x_i - a_i$ recibimos la ecuación modificada:

$$(x'_1 + a_i) + (x'_2 + a_i) + \dots + (x'_k + a_k) = n$$

$$\Leftrightarrow x'_1 + x'_2 + \dots + x'_k = n - a_1 - a_2 - \dots - a_k$$

con $x_i' \ge 0$. Así que hemos reducido el problema al caso más simple con $x_i' \ge 0$ y nuevamente podemos aplicar el teorema de estrellas y barras.

11.6 DP Optimization Theory

| Name | Original Recurrence | Sufficient Condition | From | То |
|-------|--|-----------------------------|-----------|---------------|
| CH 1 | $dp[i] = min_{j < i} \{dp[j] + b[j] *$ | $b[j] \ge b[j+1]$ Option- | $O(n^2)$ | O(n) |
| | a[i] | ally $a[i] \le a[i+1]$ | | |
| CH 2 | $dp[i][j] = min_{k < j} \{dp[i -]$ | $b[k] \ge b[k+1]$ Option- | $O(kn^2)$ | O(kn) |
| | [1][k] + b[k] * a[j] | ally $a[j] \le a[j+1]$ | | |
| D&Q | $dp[i][j] = min_{k < j} \{ dp[i -]$ | $A[i][j] \le A[i][j+1]$ | $O(kn^2)$ | $O(kn\log n)$ |
| | 1][k] + C[k][j] | | | |
| Knuth | dp[i][j] = | $A[i, j-1] \le A[i, j] \le$ | $O(n^3)$ | $O(n^2)$ |
| | $min_{i < k < j} \{dp[i][k] +$ | A[i+1,j] | | |
| | $dp[k][j]\} + C[i][j]$ | | | |

Notes:

- A[i][j] the smallest k that gives the optimal answer, for example in dp[i][j] = dp[i-1][k] + C[k][j]
- C[i][j] some given cost function
- We can generalize a bit in the following way $dp[i] = \min_{j < i} \{F[j] + b[j] * a[i]\}$, where F[j] is computed from dp[j] in constant time