Nama: Prayudha A.L NPM : (40810180008 kelas: B (4) Algoritma penjumlahah 2 matrilis 1) Tcn) = 2+4+6+8+16+ ... + 2" Deset : a(1 -1) = 2(2 -1) = 2 +1 - 2 for iel to n do for jæl ton do Notasi big 0 -> 0 (1") Mij = aij +bis T(n) & C 2" 2- = 5 & C end for 2"+1-2 < C.2" __ C ≥ I 2=3 5 0 T (n) = 1 3 (n2) 1= 1 1DZ/M D(n2) nezcn2 3) Bultillian bahwa P.q.r positif TCn) = pn+q+r adalah O (n2) 22 (n2) (n2) (n2) Karena O(n2)= ~(n2) * pembulifian Big () * pembultian Bigsl (12G7) ten) s (rCn) T(n) > C(n) (3) Algorithma Mencuri Llarily - 3 TCn) = nacio Pritintr = (.n2 りかよりかとかしかしか m'ttttc<Cn misal not p.g.rol end for select a family of the P+9+5 1 1 1 1 2 C = (8 m 8) 0 ~ 8 . 1+1+1.3C O(n) (2) (n) (5) 0 - 1 C > 3 misal n=1 n & Chical to > Charles de de como de misalkan p=q=r=1 C31 C51 1-1-15 :. O(n) = a(n) = + (0 (n) = == 190) C>3 (6) a) Operasi perbandingan T(n) = (n-1) + (n-2) + (n-3) + - + 1 - 1 Big 🕘 = n (n-1) = n2-n Karena (9 Cm²) dan sich²) ferbulli dan sama malia O(n2) pun benor. b) max pertuliaran terjadi ji duni (16 letilia n.Cn-1) 3 Tentrian (completisitas walte Wi; Wik dan Wij benjang sebanyah Am,n T(n) = n3 c) kompletisitas wallfu Big st-acn') · Best case Big () → ()(n³) n3 > Cn3 n3 4 C.n3 perbandingan -> n(h-1) káli? C >1 C < 1 T(n) = n(n-1) = n2-n Big D -> O (n3) (Larena O Cn3) = S(n3) · Mouzt case Perbandingan - h(n-1) maka O(n3) asignment = n Cn-1).

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