## TD 2-

# **AST, Attributions, Types**

## 2.1 Derivation trees and attributions

#### EXERCISE #1 $\triangleright$ Arithmetic expressions

Let us consider the following grammar (the end of an expression is a semicolon):

$$Z \rightarrow E;$$

$$E \rightarrow E + T$$

$$E \rightarrow T$$

$$T \rightarrow T * F$$

$$T \rightarrow F$$

$$F \rightarrow (E)$$

$$F \rightarrow i$$

• What are the derivation trees for 1 + 2 + 3;, 1 + 2 \* 3;, (1 + 2) \* 3?

## **EXERCISE** #2 **▶ Declarations of variables**

Write a grammar that accepts declarations of variables like:

```
int x=1;
float y,z;
int t;
float u,v=0;
and rejects:
int x, int y;
```

Then write an attribution that prints individual declarations (of the first case) like:

```
int x=1; float y; float z; int t; float u; float v=0;
```

## **EXERCISE** #3 ► **Prefixed expressions**

Consider prefixed expressions like \* + \* 3 4 5 7 (or \* + 1 2 \* 3 4) and assignments of such expressions to variables:

a=\* + \* 3 4 5 7. Identifiers are allowed in expressions.

- Give a grammar that recognizes lists of such assigments.
- Write derivations trees.
- Write grammar rules to compute the values of the expressions.
- Write grammar rules to construct infix assignments during parsing: the former assignment will be transformed into the *string* a=(3 \* 4 + 5)\*7. Be careful to avoid useless parentheses.
- Modify the attribution to verify that the use of each identifier is done after his first definition.

## 2.2 The MiniC language

The objective here is to be familiar with the grammar of the language we will compile.

## EXERCISE #4 ► MiniC-grammar

Here is the (simplified) grammar for the MiniC language (expr are numerical or boolean expressions):

```
grammar MiniC;
prog: function* EOF #progRule;
function: INTTYPE ID OPAR CPAR OBRACE vardecl_l block RETURN INT SCOL CBRACE #funcDecl;
vardecl_1: vardecl* #varDeclList;
vardecl: typee id_l SCOL #varDecl;
id_l
   : ID
                      #idListBase
    | ID COM id_l
                      #idList
block: stat* #statList;
    : assignment SCOL
    | if_stat
    | while_stat
    | print_stat
assignment: ID ASSIGN expr #assignStat;
if_stat: IF condition_block (ELSE IF condition_block)* (ELSE stat_block)? #ifStat;
condition_block: OPAR expr CPAR stat_block #condBlock;
stat_block
    : OBRACE block CBRACE
    stat
while_stat: WHILE OPAR expr CPAR stat_block #whileStat;
print_stat
    : PRINTINT OPAR expr CPAR SCOL
                                            #printintStat
     PRINTFLOAT OPAR expr CPAR SCOL
                                            #printfloatStat
    | PRINTSTRING OPAR expr CPAR SCOL
                                            #printstringStat
expr_l
    : /* Nothing */
                       #exprListEmpty
                        #exprListBase
    expr
    expr COM expr_l #exprList
expr
    : MINUS expr
                                            #unaryMinusExpr
     NOT expr
                                            #notExpr
     expr myop=(MULT|DIV|MOD) expr
                                            #multiplicativeExpr
     expr myop=(PLUS|MINUS) expr
                                            #additiveExpr
     expr myop=(GT|LT|GTEQ|LTEQ) expr
                                            #relationalExpr
     expr myop=(EQ|NEQ) expr
                                            #equalityExpr
     expr AND expr
expr OR expr
                                            #andExpr
                                            #orExpr
     atom
                                            #atomExpr
atom
    : OPAR expr CPAR #parExpr
     (INT | FLOAT) #numberAtom
(TRUE | FALSE) #booleanAtom
                     #idAtom
    | STRING
                     #stringAtom
```

Write a valid program for this grammar.

## 2.3 Typing

We recall the rules of the course for Typing:

$$\frac{c \in \mathbb{Z}}{\Gamma \vdash c : \mathrm{int}} \qquad \frac{\Gamma(x) = t \quad t \in \{\mathrm{int}, \mathrm{bool}\}}{\Gamma \vdash x : t} \qquad \frac{\Gamma \vdash S_1 \ \Gamma \vdash S_2}{\Gamma \vdash S_1; S_2}$$

$$\frac{\Gamma \vdash e : t \quad \Gamma \vdash x : t \quad t \in \{\mathrm{int}, \mathrm{bool}\}}{\Gamma \vdash x = e : \mathrm{void}} \qquad \frac{\Gamma \vdash b : \mathrm{bool} \quad \Gamma \vdash S : \mathrm{void}}{\Gamma \vdash \mathrm{while}(b) \{S\} : \mathrm{void}}$$

$$\frac{\Gamma \vdash b : \mathrm{bool} \quad \Gamma \vdash S_1 : \mathrm{void} \quad \Gamma \vdash S_2 : \mathrm{void}}{\Gamma \vdash \mathrm{int} \quad b : \mathrm{then} \ S_1 \ \mathrm{else} \ S_2 : \mathrm{void}}$$

$$\frac{\Gamma \vdash e_1 : \mathrm{int} \quad \Gamma \vdash e_2 : \mathrm{int}}{\Gamma \vdash e_1 + e_2 : \mathrm{int}} \qquad \frac{\Gamma \vdash e_1 : \mathrm{int} \quad \Gamma \vdash e_2 : \mathrm{int}}{\Gamma \vdash e_1 < e_2 : \mathrm{bool}}$$

## **EXERCISE** $#5 \triangleright$ **Typing under given environment**

Considering  $\Gamma = \{x_1 \mapsto \mathtt{int}\}\$ , prove that the given sequence of instructions is well typed:

```
x1 = 3;
while (x1 < 15) {
  x1 = x1 + 1;
}
```

## EXERCISE #6 ► Construction of Gamma with variable declarations

Using the following rules:

$$\frac{t \ vlist; \rightarrow_d [v \mapsto t \text{ for } v \text{ in } vlist]}{D_1 \rightarrow_d \Gamma_1 \quad D_2 \rightarrow_d \Gamma_2 \quad Dom(\Gamma_1) \cap Dom(\Gamma_2) = \emptyset}$$

$$\frac{D_1; D_2 \rightarrow_d \Gamma_1 \cup \Gamma_2}{D_1; D_2 \rightarrow_d \Gamma_1 \cup \Gamma_2}$$

Compute  $\Gamma$  the typing environnement for the given list of declarations:

```
int x;
bool b1, b2;
```

## **EXERCISE** #7 **▶ Boolean expressions**

Write all rules to type boolean expressions of MiniC.