

10.1

22

Team 4 beats Team 3

Team 1, 2, 5, 6 beat Team 4

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Statement S_1, S_2, S_3, S_4 must be excuted before S_6 is excuted

10.2

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Assume the vertex below is v

The degree of v represents the number of people v knows

The neighborhood of v is the people v knows

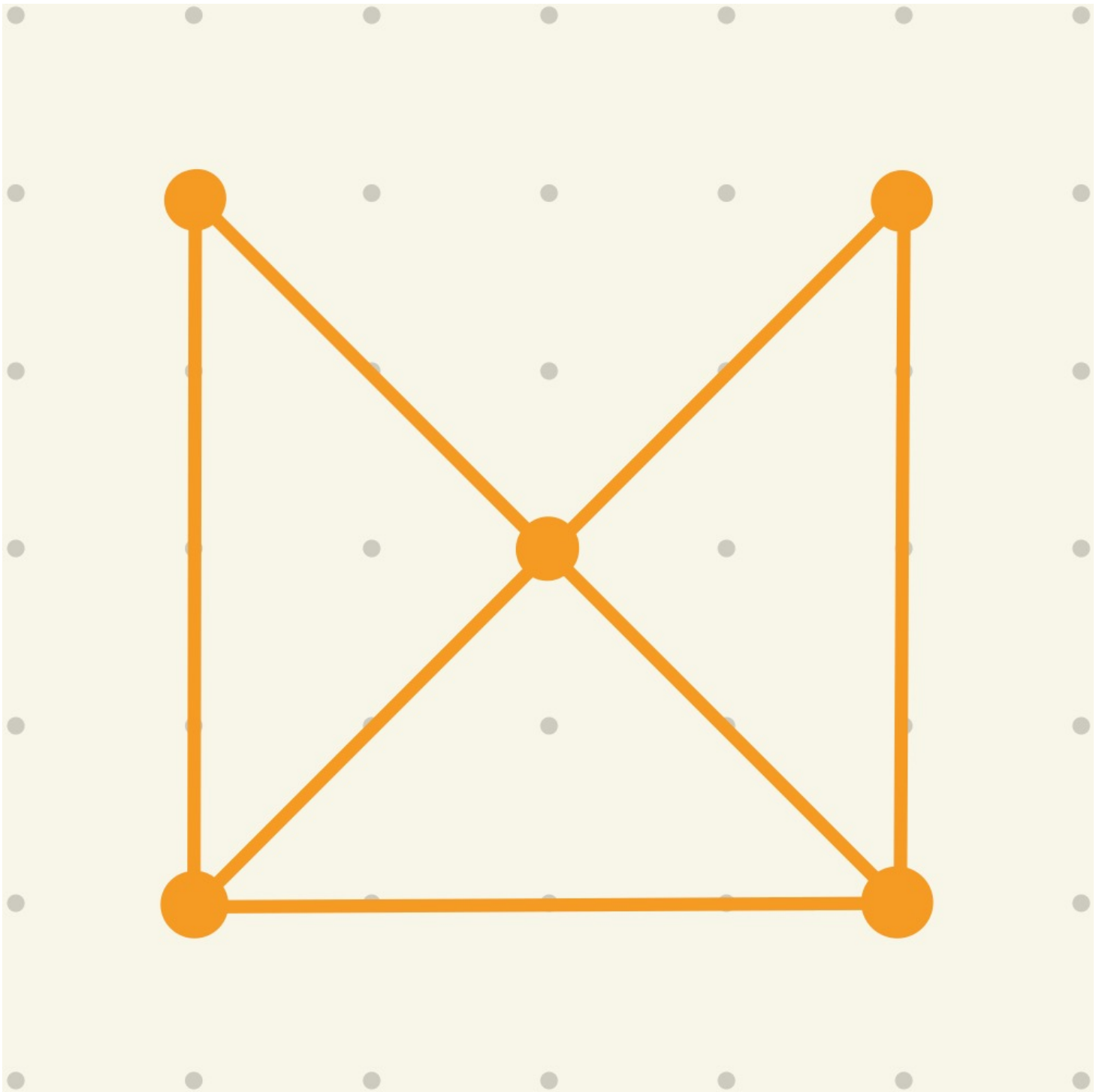
The isolated vertices represent they know nobody

The pendant vertices represent they know exactly one people

It means a person knows 1000 people on average

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Since every edge contributes 2 degrees, the number of edges is $(4 + 3 + 3 + 2 + 2)/2 = 7$



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Since $G \cap \overline{G}$ is a complete graph, the new number of edges of each vertex should be 4 - the former one, that is, 0, 1, 1, 2, 2

Hence the answer is 2, 2, 1, 1, 0

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Since G is n -regular and (V_1, V_2) is a bipartition of V , it is obvious that from the perspective of V_1 , the degree of the graph is $n|V_1|$ since $\forall v \in V_1, \deg(v) = n$, and we can get the similar

result from $V_2, n|V_2|$. Hence we have $n|V_1| = n|V_2|, |V_1| = |V_2|$

Hence they contain the same number of vertices

10.3

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1. Sparse. Since a street intersection has less than 5 edges normally, which is much less than the number of vertices
2. Dense. Since nowadays buildings are usually close to each others
3. Sparse. Most people do not have brothers or sisters
4. Neither. It is hard to say because it depends on the actual situation、

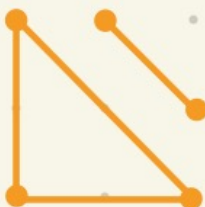
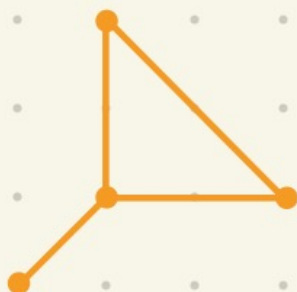
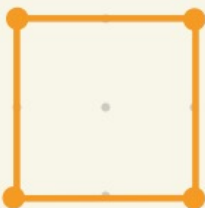
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Not isomorphic. Since the two vertices with 4 degree in the left one are adjacent, but the right one is not.

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The row of a isolated vertex should only contain 0

60



It is not isomorphic since there is no edge from out-degree = 2 (only one vertex in the graph) to in-degree = 2 (only one vertex in the graph) in the left graph, but the right one does.