§1.1 Error Analysis
- **Normalization**: 32 bit - Sign bit: 1, Sign exp: 1, Exp: 7, Normalized mantissa: 23.
- Absolute Error $= |p - p^*|$, Relative Error $= \frac{|p - p^*|}{|p|}$
- Significant Digits: RE $< 5 \times 10^{-t}$
- **f(x)**: Machine representation
- **Cancellation Error**: Subtracting nearly equal numbers
- **Example**: $p = 0.54617$, $q = 0.54601$, true $r = p - q = 0.00016$
  4-digit: $p^* = 0.5462$, $q^* = 0.5460$, $r^* = 0.002$ (RE=25%)
- **Nested Multiplication**: Reduces error
  $f(z) = 1.01z^4 - 4.62z^3 - 3.11z^2 + 12.2z - 1.99 = (((1.01z - 4.62)z - 3.11)z + 12.2)z - 1.99$
- $P_n(x) = \sum_{k=0}^{n} \frac{f^{(k)}(a)}{k!}(x - a)^k$
- Remainder: $R_n(x) = \frac{f^{(n+1)}(\xi)}{(n+1)!}(x - a)^{n+1}$, $\xi \in (a, x)$
- Linear Approx: $f(x_0 + h) \approx f(x_0) + hf'(x_0)$
- **Example**: $\xi \in (0, \pi/2)$, $\sin \xi \leq 1 \Rightarrow R_n \leq \frac{(x)^n}{n!}$

§1.3 Convergence
- $\alpha = \lim_{n \to \infty}(\alpha_n)$
- Rate: $\alpha_n = \alpha + \mathcal{O}(\beta_n)$ if $|\alpha_n - \alpha| \leq K|\beta_n|$
- Find largest p where $\alpha_n - \alpha = \mathcal{O}(1/n^p)$

§1.4 Matrix Operations
- $A = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \end{bmatrix}$, $B = \begin{bmatrix} b_{11} & b_{12} \\ b_{21} & b_{22} \\ b_{31} & b_{32} \end{bmatrix}$
- $AB = \begin{bmatrix} a_{11}b_{11} + a_{12}b_{21} + a_{13}b_{31} & a_{11}b_{12} + a_{12}b_{22} + a_{13}b_{32} \\ a_{21}b_{11} + a_{22}b_{21} + a_{23}b_{31} & a_{21}b_{12} + a_{22}b_{22} + a_{23}b_{32} \end{bmatrix}$

§6.1 Gaussian Elimination
- $\mathcal{O}(n^3)$ complexity
- **Pivoting**:
  - Partial (PP): Max element in column
  - Scaled PP: $s_i = \max_j |a_{ij}|$, pivot $\max(a_{ik}/s_i)$
  - Complete (CP): Full matrix search $(\mathcal{O}(n^3))$
- **LU Decomposition**: $PA = LU$ through GE steps and $LUx = Pb$.
- **LU Algorithm**: $L = E_{n-a, m-b}^{-1} E_{n-a+1, m-b+1}^{-1} \cdots$ $U = $ GE.
- **Choleski Factorization**: If a matrix is symmetric and positive definite, it may be factored to the form LDL^T

§6.2 Special Matrices
- **Inverse Matrix**: An inverse matrix of A is A^{-1} such that $AA^{-1} = I$
- **Properties**: $(AB)^{-1} = B^{-1}A^{-1}$, $(A^{-1})^T = (A^T)^{-1}$
- **Singular**: A matrix is singular iff its det is 0.
- **Diagonal**: $d_{ij} = 0$ for $i \neq j$: All non-diagonal entries are 0.
- **Symmetric**: $A = A^T$, $(AB)^T = B^T A^T$
- **Permutation**: Row swaps of I_n, PA reorders rows: $P^T = P^{-1}$
- **Diagonally Dominant**: $|a_{ii}| > \sum_{j \neq i} |a_{ij}|$ (nonsingular)
- **Positive Definite**: $x^T Ax > 0 \Rightarrow A = LDL^T$, $a_{ii} > 0$, $a_{ij}^2 < a_{ii}a_{jj}$
- **Minor**: M_{ij} is a submatrix of A with the row i deleted and column j deleted.
- **Band**: an $n \times n$ matrix is a band matrix if $p, q \in \mathbb{Z}: 1 \leq p, q \leq n$ exist with $a_{i_j} = 0$ for $i + p \leq j$ or $j + q \geq i$ The bandwidth is defined as $w = p + q - 1$. For a diagonal matrix, $p = 1, q = 1, w = 1$
- **Tridiagonal**: Band with $p = 2, q = 2$. It exhibits the following properties
  - $a_{ii} = l_{ii}$
  - $a_{i,i-1} = l_{i,i-1} : i = 2, 3, \cdots, n$
  - $a_{i,i+1} = l_{ii}u_{i,i+1} : i = 1 \cdots n - 1$
  - $a_{ii} = l_{i,i-1}u_{i-1,i} + l_{ii} : i = 2 \cdots n$
- Crout Factorization: This factorization happens in $O(n)$ time

§Strategies
- **RoC With inf limit**: set $h = 1/n$ and solve accordingly.
- $D \cdot (L+U)$: given D has ONLY diagonal entries and L+U has NO diagonal entries, the resulting matrix A is composed of entries $a_{ij} = d_{ii} \cdot (l + u)_{ij}$
- **Verification of Bisection**: To verify bisection can be applied, make sure that f(a) and f(b) are of different signs.
- **Error of Bisection**: To compute the accuracy of bisection to an ε, we use $\frac{b-a}{2^n} \leq \varepsilon$
- **Failure of Newton's Method**: NM Fails if $f'(x) = 0$ for some x.
- **Triangle Inequality**: $|x + y| \leq |x| + |y|$

§Key Definitions & Identities
- **Continuity**: $f \in C^n[a, b]$ reads: the nth derivative of f on [a,b] is continuous.
- **Series Expansions**
  - $e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \cdots$ $(\forall x \in \mathbb{R})$
  - $\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \cdots$
  - $\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \cdots$
  - $\arctan x = x - \frac{x^3}{3} + \frac{x^5}{5} - \frac{x^7}{7} + \cdots$ $(|x| \leq 1)$
  - $\ln(1 + x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \cdots$ $(|x| < 1)$
  - $\frac{1}{1-x} = 1 + x + x^2 + x^3 + x^4 + \cdots : |x| < 1$
  - $\frac{1}{1+x} = 1 - x + x^2 - x^3 + x^4 - \cdots : |x| < 1$

Core Identities
- $\sin^2 \theta + \cos^2 \theta = 1$
- $\sin(a \pm b) = \sin a \cos b \pm \cos a \sin b$
- $\cos(a \pm b) = \cos a \cos b \mp \sin a \sin b$

Angle Transformations
- $\sin 2\theta = 2 \sin \theta \cos \theta$
- $\cos 2\theta = 2\cos^2 \theta - 1 = 1 - 2\sin^2 \theta$
- $\sin^2 \theta = \frac{1 - \cos 2\theta}{2}$, $\cos^2 \theta = \frac{1 + \cos 2\theta}{2}$

- **Vector Norms**
  - $||x|| > 0$
  - $||x|| = 0 \Leftrightarrow x = 0$
  - $||\alpha x|| = |\alpha| ||x||$
  - $||x + y|| \leq ||x|| + ||y||$

§Useful Examples
- Suppose $g(x) = \frac{5}{x^2} + 2$ Show $p_n = g(p_{n-1})$ will converge to g for $\forall p_0 [2.5, 3]$. Since this is a decreasing function, the max of g(x) is g(2.5) and the min of g(x) is g(3).
  - First compute the max, $g(2.5) = 14/5 < 3$
  - Second compute the min, $g(3) = 5/9 + 2 > 2.5$
  - Last compute $|g'(x)| = -10/x^3 \leq max_{x \in [2.5, 3]}|g'(x)| = 16/25 < 1$
- Given $||A||$ is a natural matrix norm of matrix A. show $|\lambda| \leq ||A||$ for any nonsingular A and any λ of A. $||A|| = \max_{||x||=1} ||Ax|| \geq ||Ax||$: x is an e-vec s.t ——x——=1 $= ||\lambda x|| = |\lambda| ||x|| = |\lambda| \square$
- When performing Jacobi or GS, when computing L+U, flip the signs of all entries.
- To determine convergence for fixed point, compute $g'(p_0) \leq 1$, which gives a,b. Prove g(x) cts on [a,b], $g(x) \in [a, b]$, $g'(x)$ exists on (a,b), $|g'(x)| \leq k : \forall x \in (a, b), 0 < k < 1$

§LA Determinants
- 2x2: $|A| = ad - bc$
- nxn: $|A| = \sum_{j=1}^{n} a_{ij} A_{ij}$ via cofactors $A_{ij} = (-1)^{i+j} M_{ij}$
  Properties
  - Swap rows: $|\tilde{A}| = -|A|$
  - Scale row: $|\tilde{A}| = \lambda |A|$
  - Identical rows: $|A| = 0$
  - $|AB| = |A||B|$, $|A^T| = |A|$
  - $|A^{-1}| = \frac{1}{|A|}$

§7 Norms & Eigen
- $||x||_2 = \sqrt{\sum x_i^2}$, $||x||_\infty = \max |x_i|$
- $||A||_2 = \sqrt{\rho(A)}$
- $||A||_\infty = \max_i \sum_j |a_{ij}|$ Basically sum all rows together and determine the largest one.
- $||A|| = \max_{||x||=1} ||Ax||$
- $\forall x \in \mathbb{R}^n : ||x||_\infty \leq ||x||_2 \leq \sqrt{n} ||x||_\infty$
- A distance between matrices A and B wrt a matrix norm $|| \cdot ||$ is $||A - B||$
- Theorem: For any vector $x \neq 0$, matrix Am and abt natural norm $|| \cdot ||$ we have $||Ax|| \leq ||A|| \cdot ||x||$
- Cauchy-Schwarz: $||x + y||_2 \leq ||x||_2 ||y||_2$
- Eigen: λ is an eigenvalue if $A\lambda = v \cdot \lambda$
- Finding Eigenthings: $\det(A - \lambda I) = 0 : \forall \lambda$.
- Spectral Radius: $\rho(A) = \max |\lambda_i|$, $\rho(A) \leq ||A||$
- Theorem: If A is $n \times n$:
  - $||A||_2 = [\rho(A^t A)]^{1/2}$
  - $\rho(A) \leq ||A|| : \forall || \cdot ||$
- Convergent: $\lim_{k \to \infty} A^k = 0 \Leftrightarrow \rho(A) < 1$
- Matrix Norms: matrix norms have the following properties
  - $||A|| \geq 0$ with $||A|| = 0 \leftrightarrow A = 0$
  - $||\alpha A|| = |\alpha| \cdot ||A||$
  - $||A + B|| \leq ||A|| + ||B||$

§Iterative Methods
- **General Iteration**: $x^{k+1} = Tx^k + c$
- **Jacobi**: $x^{k+1} = D^{-1}(L + U)x^k + D^{-1}b$
- **Gauss-Seidel**: $x^{k+1} = (D - L)^{-1}Ux^k + (D - L)^{-1}b$
- **Stein-Rosenberg**: For matrices with positive diagonals: $\rho_{GS} \leq \rho_J < 1$
- **Speed of Convergence**: given matrices $T_{GS} = (D - L)^{-1}U$ and $T_J = D^{-1}(L + U)$, compare ρ. The bigger the ρ, the faster the convergence.
- Error: $||x - x^k|| \leq \frac{||T||^k}{1 - ||T||} ||x^1 - x^0||$
- Stopping: $\frac{||x^k - x^{k-1}||}{||x^k||} < \varepsilon$

§2 Nonlinear Equations
- **Bisection**:
  - While $f(p_n) \neq 0$ or $< T$: $p_n = \frac{a_1 + b_1}{2}$
  - Error: $\frac{b_n - a_n}{2} < T$, $p = a + \frac{b-a}{2}$
- **Fixed-Point**:
  - $p_n = g(p_{n-1})$, converges if $|g'(x)| \leq K < 1$
  - Algorithm: For $i < N_0$: $p = g(p_0)$, check $|p - p_0| < T$
  - A fixed point is defined as a point in which $p = f(p)$
- **Newton**:
  - $p_n = p_{n-1} - \frac{f(p_{n-1})}{f'(p_{n-1})}$
  - Quadratic convergence if $f'(p) \neq 0$
- **Secant**:
  - $p_{n+1} = p_n - \frac{f(p_n)(p_n - p_{n-1})}{f(p_n) - f(p_{n-1})}$
  - Approx derivative: $\frac{f(p_{n-1}) - f(p_{n-2})}{p_{n-1} - p_{n-2}}$
  - Algorithm: Store $q_0 = f(p_0)$, $q_1 = f(p_1)$, SET $p = p_1 - \frac{q_1(p_1 - p_0)}{q_1 - q_0}$ IF STOPPING CONDITION: RETURN p; i++, $p_0 = p_1, q_0 = q_1, p_1 = p, q_1 = f(p)$ ENDWHILE OUTPUT FAILURE.

§Theorems
- **Bisection**: Suppose $f \in C[a, b] : f(a) \cdot f(b) < 0$. Bisection generates $\{p_n\}$ approximating a zero p with $|p_n - p| \leq \frac{b-a}{2^n} : n \geq 1$
- **Fixed Point**: If $g \in C[a, b]$, $g([a, b]) \subseteq [a, b]$ g has a fixed point in $[a, b]$, additionally if $|g'| \leq K < 1$, then the fixed point is unique.
- **Fixed Point Theorem**: Let $g \in C[a, b]$ and $g(x) \in [a, b] : \forall x \in [a, b]$. Suppose as well that g' exists on (a,b) and positive $K < 1$ exists with $|g'(x)| \leq K : \forall x \in (a, b)$. Then for any number $p_0 \in [a, b]$h the sequence defined by $p_n = g(p_{n-1}) : n \geq 1$ converges to the unique point $p \in [a, b]$
- **Corollary**: If g satisfies the hypothesis of the above theorem, $|p_n - p| \leq k^n \max(p_0 - a, b - p_0)$ and $|p_n - p| \leq \frac{k^n}{1-k}|p_1 - p_0| : \forall n > 1$
- **Newton**: For $f \in C^2[a, b]$ with simple root, $\exists \delta > 0 : p_0 \in [p - \delta, p + \delta]$ converges.
- **Matrix Invertibility**: $|A| \neq 0 \Leftrightarrow$ unique solution $Ax = b \Leftrightarrow A^{-1}$ exists
- **Taylor**: With $R_n(x) \Rightarrow f(x) = P_n(x) + R_n(x)$, $R_n(x) = \frac{f^{n+1}(\xi)}{(n+1)!}(x - x_0)^{n+1} : \xi \epsilon(x, x_0)$
- **Existence of Inverse**: if A is square, $\det A \neq 0 \leftrightarrow Ax = 0$ has soln $x = 0 \leftrightarrow Ax = b$ has a unique soln for any n-vector b. $\leftrightarrow A^{-1}$ exists.
- **Diagonally Dominant Matrices**: dd matrices are nonsingular. A being dd means LU can be performed without P. A matrix is positive definite if $x^t Ax > 0$. PD matrices are nonsingular, $\forall i = 1, \cdots, n : a_{ii} > 0$, $max|a_{kj}| > max|a_{ii}|$, $(a_{ij})^2 < a_{ii}a_{jj} : \forall i \neq j$.
- **Convergence and Spectral Radii**: equivalent statements: A is convergent, $\rho < 1$, $\lim A^n x = 0 : \forall x$, $\lim ||A^n|| = 0 : \forall norms$.
- **Convergence of DD**: If A is strictly DD, Jacobi and GS converge to the unique soln. Ax=b.
- **Positive Definitive Check**: A matrix is PD if the leading principle submatrix determinants are positive.
- **Positive Definitive Check**: A matrix is PD iff it may be factored into LL^T
- **Non-singularity Check**: A matrix A has an inverse iff det $A \neq 0$.
- **Determinant of Triangular Matrices**: The determinant of a triangular matrix is Πa_{ii}.

§Proofs
- **Bisection (THM1)**: $\forall n \geq 1$: $b_n - a_n = (b - a) \cdot \frac{1}{2^{n-1}}$: $p\epsilon(a_n, b_n)$. Since $p_n = \frac{1}{2}(a_n + b_n) : \forall n \geq 1$, $|p_n - p| \leq \frac{1}{2}(b_n - a_n) = \frac{b-a}{2^n}$. $\square$
- **Fixed Point**: Part i: If $g \in [a, b], g(x) \in [a, b] : \forall x \in [a, b]$ then g(x) has a fixed point in [a,b]:
  If $g(a) = a$ or $g(b) = b$, g has a fixed point at an endpoint. Suppose for contradiction that it does not. $g(a) > a$ and $g(b) < b$. Define $h(x) = g(x) - x$. Then h is cts on [a,b] and $h(a) = g(a) - a > 0$ and $h(b) = g(b) - b < 0$ IVT states that $\exists p \in (a, b) : h(p) = 0$ Thus $g(p) - p = 0 \Rightarrow$p is a fixed point of g.
  Part ii: Suppose as well $|g'(x)| \leq k < 1 : \forall x \in (a, b)$ and that $p, q \in [a, b] : p \neq q$. By MVT, $\exists \zeta : \frac{g(p) - g(q)}{p - q} = g'(\zeta)$. $|p - q| = |g(p) - g(q)| = |g'(\zeta)||p - q| \leq k|p - q| < |p - q|$ contradiction.

§2: Error Analysis and Acelerating Convergence

Basic Methods
- **Newton's Method**: Quadratic convergence if $f'(p) \neq 0$. Iteration: $x_{n+1} = x_n - f(x_n)/f'(x_n)$.
- **Secant Method**: Superlinear convergence (order ≈ 1.618). Uses two previous points.
- **False Position**: Bracketing method combining bisection and secant approaches.

Convergence Analysis
- Order of convergence α: $\lim \frac{|p_{n+1}-p|}{|p_n-p|^{\alpha}} = \lambda$
- Linear ($\alpha = 1$), Quadratic ($\alpha = 2$)
- Fixed-point: Linear if $g'(p) \neq 0$, quadratic if $g'(p) = 0$ and g'' bounded.

Special Cases
- Multiple roots: Modify Newton's using $\mu(x) = f(x)/f'(x)$
- **Aitken's Δ^2**: Accelerates linear sequences. Is given by $\hat{p}_n = p_n - \frac{(p_{n+1}-p_n)^2}{(p_{n+2}-p_{n+1})-(p_{n+1}-p_n)}$
- Steffensen's: Combines Aitken's with fixed-point, achieves quadratic convergence
$$p_{n+1} = p_n - \frac{[g(p_n)-p_n]^2}{g(g(p_n))-2g(p_n)+p_n}$$

Polynomial Methods
- Horner's method: Efficient evaluation (n mults/adds)
- Deflation: Find roots sequentially via $P(x) \approx Q(x)(x - x_0)$
- Müller's method: Quadratic interpolation using 3 points
- Fundamental thm of alg: If $P(x)$ has a degree $n \geq 1$, $P(x)$ has at least one root.
- Cor: there also exists unique constants $x_1, ...x_k$ such that $\sum_{i=1}^{k} m_i = n, P(x) = a_n \times \prod_{i=1}^{k} (x - x_i)^{m_i}$
- Cor: these functions are unique.

Weierstrass Approximation Theorem For f continuous on $[a, b]$, $\forall \epsilon > 0$, $\exists$ polynomial $p(x)$ such that $|f(x) - p(x)| < \epsilon$ for all $x \in [a, b]$.

Lagrange Interpolation Given (x_k, f_k) for $k = 0, \ldots, N$, the unique degree $\leq N$ polynomial satisfying $P(x_k) = f_k$ is:

$$P(x) = \sum_{m=0}^{N} f_m L_m(x), \quad L_m(x) = \prod_{\substack{k=0 \\ k \neq m}}^{N} \frac{x - x_k}{x_m - x_k}$$

Interpolation Error If $f \in C^{n+1}[a, b]$, the error is:

$$f(x) - P(x) = \frac{f^{(n+1)}(\xi(x))}{(n + 1)!} \prod_{k=0}^{n} (x - x_k)$$

Newton's Divided Differences Recursive definition:
$$f[x_i] = f(x_i)$$
$$f[x_i, \ldots, x_{i+k}] = \frac{f[x_{i+1}, \ldots, x_{i+k}] - f[x_i, \ldots, x_{i+k-1}]}{x_{i+k} - x_i}$$

Interpolating polynomial:

$$P_n(x) = \sum_{k=0}^{n} f[x_0, \ldots, x_k] \prod_{j=0}^{k-1} (x - x_j)$$

Finite Differences For evenly spaced x_i with step h:
- Forward: $\Delta f_i = f_{i+1} - f_i$, $\Delta^k f_i = \Delta^{k-1} f_{i+1} - \Delta^{k-1} f_i$
- Backward: $\nabla f_i = f_i - f_{i-1}$, $\nabla^k f_i = \nabla^{k-1} f_i - \nabla^{k-1} f_{i-1}$

Newton's Formulas
- Forward: $P(x) = \sum_{k=0}^{n} \binom{s}{k} \Delta^k f_0$ where $x = x_0 + sh$
- Backward: $P(x) = \sum_{k=0}^{n} (-1)^k \binom{-s}{k} \nabla^k f_n$ where $x = x_n + sh$