

Lecture Notes 8: Information Measure Entropy

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Scribe:

8 Information Measure and Entropy

8.1 Discrete Cases

Definition 8.1 Given discrete random variable X , the entropy $\mathbb{H}(X)$ of X is defined by $\mathbb{H}(X) = - \sum_{x \in \mathcal{X}} p(x) \log p(x)$

- $\log e = 1$,
- $0 \log 0 = \lim_{a \rightarrow 0^+} a \log a = 0$.

Lemma 8.1 For any discrete random variable X , $\mathbb{H}(X) \geq 0$.

Proof : Since $0 \leq p(x) \leq 1$, we have $p(x) \log p(x) \leq 0$. So $\mathbb{H}(x) \geq 0$ holds.

Example 8.1 Given the random variable X with p.m.f that $p(x) = \begin{cases} 1 & \text{with probability } p \\ 0 & \text{with probability } 1-p \end{cases}$
Then, $\mathbb{H}(x) = -p \log p - (1-p) \log(1-p)$.

8.2 Joint Entropy and Conditional Entropy

Definition 8.2 The entropy $\mathbb{H}(X, Y)$ of (X, Y) is defined by

$$\mathbb{H}(X, Y) = - \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} p(x, y) \log p(x, y) = -\mathbb{E} [\log p(x, y)].$$

Definition 8.3 If $(X, Y) \sim p(x, y)$, then the conditional entropy

$$\begin{aligned} \mathbb{H}(Y|X) &= \sum_{x \in \mathcal{X}} p(x) \mathbb{H}(Y|X=x) \\ &= - \sum_{x \in \mathcal{X}} p(x) \sum_{y \in \mathcal{Y}} p(y|x) \log p(y|x) \\ &= - \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} p(x, y) \log p(y|x) \\ &= -\mathbb{E} [\log p(y|x)] \end{aligned}$$

Theorem 8.1 (The Chain Rule) $\mathbb{H}(X, Y) = \mathbb{H}(X) + \mathbb{H}(Y|X)$

Proof: Using $\log p(x, y) = \log p(x) + \log p(y|x)$, we compute the expectation in both sides about (X, Y) .

Corollary 8.1 $\mathbb{H}(X, Y|Z) = \mathbb{H}(X|Z) + \mathbb{H}(Y|X, Z)$.

8.3 Relative Entropy and Mutual Information

Definition 8.4 The relative entropy or KullbackLeibler Divergence(KLD) between p.m.f $p(x)$ and $q(x)$ is defined as follows:

$$D(p||q) = \sum_{x \in \mathcal{X}} p(x) \log \frac{p(x)}{q(x)} = \mathbb{E}_p \left[\log \frac{p(x)}{q(x)} \right]$$

- $0 \log \frac{0}{q} = 0$,
- $a \log \frac{a}{0} = \infty$,
- $0 \log \frac{0}{0} = 0$.

Definition 8.5 Given two variable X and Y with p.m.f $p(x, y)$ and the marginal p.m.f are $p(x)$ and $p(y)$. The mutual information $\mathbb{I}(X, Y)$ is

$$\begin{aligned} \mathbb{I}(X, Y) &= \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} \log \frac{p(x, y)}{p(x)p(y)} p(x, y) \\ &= D(p(x, y)||p(x)p(y)) \end{aligned}$$

Generally the KullbackLeibler Divergence is not symmetric. But we can build $D'(p||q) = \frac{1}{2}D(p||q) + \frac{1}{2}D(q||p)$ to make the KLD symmetric.

Example 8.2 Let $\mathcal{X} = \{0, 1\}$, $p(x)$ and $q(x)$ are p.m.f. let $p(X = 0) = 1 - r$, $p(X = 1) = r$, $q(X = 0) = 1 - s$, $q(X = 1) = s$, Then

$$\begin{aligned} D(p||q) &= (1 - r) \log \frac{1 - r}{1 - s} + r \log \frac{r}{s} \\ D(q||p) &= (1 - s) \log \frac{1 - s}{1 - r} + s \log \frac{s}{r} \end{aligned}$$

Theorem 8.2 (Mutual Information and Entropy)

$$\begin{aligned} \mathbb{I}(X, Y) &= \mathbb{I}(Y, X) \\ \mathbb{I}(X, X) &= \mathbb{H}(X) \\ \mathbb{I}(X, Y) &= \mathbb{H}(X) - \mathbb{H}(X|Y) \\ &= \mathbb{H}(Y) - \mathbb{H}(Y|X) \\ &= \mathbb{H}(X) + \mathbb{H}(Y) - \mathbb{H}(X, Y) \end{aligned}$$

Definition 8.6 The conditional mutual information of random variable X and Y given Z is

$$\mathbb{I}(X, Y|Z) = \mathbb{H}(X|Z) - \mathbb{H}(X|Y, Z)$$

Definition 8.7 The conditional relative entropy $D(p(y|x)||q(y|x))$ is

$$D(p(y|x)||q(y|x)) = \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} p(y|x) \log \frac{p(y|x)}{q(y|x)}$$

Theorem 8.3 Let $p(x)$ and $q(x)$ with $x \in \mathcal{X}$ be two p.m.f. Then $D(p||q) \geq 0$ with the equality if and only if $p(x) = q(x)$, for all $x \in \mathcal{X}$

Lemma 8.2 Let $\sum a_i$ and $\sum b_i$ be convergent sequence of non-negative numbers. Then the following hold:

- $\sum a_i \log \frac{b_i}{a_i} + \sum (a_i - b_i) \leq 0$ or $\sum a_i \log \frac{a_i}{b_i} + \sum (b_i - a_i) \geq 0$.
- If $\sum a_i \geq \sum b_i$, then $\sum a_i \log \frac{b_i}{a_i} \leq 0$ with equality iff $a_i = b_i$.
- Further more, if $a_i \leq 1$ and $b_i \leq 1$ for all i , then $2 \sum a_i \log \frac{a_i}{b_i} \geq \sum a_i (a_i - b_i)^2$

Proof : Considering the taylor expansion of $\log x$ at $x = 1$, we have $\log x = (x - 1) - \frac{(x-1)^2}{2} \frac{1}{\theta^2}$, where θ is between 1 and x . Hence, $\log \frac{b_i}{a_i} = (\frac{b_i}{a_i} - 1) - \frac{1}{2\theta_i^2} (\frac{b_i}{a_i} - 1)^2$, then $a_i \log \frac{b_i}{a_i} = (b_i - a_i) - \frac{a_i^3}{2a_i^2\theta_i^2} (\frac{b_i}{a_i} - 1)^2$. So $\sum a_i \log \frac{b_i}{a_i} = \sum (b_i - a_i) - \sum \frac{a_i^3}{2a_i^2\theta_i^2} (\frac{b_i}{a_i} - 1)^2$. Notice that $\theta_i \in [1, \frac{b_i}{a_i}]$, we have $a_i\theta_i \in [a_i, b_i]$, hence $\sum \frac{a_i^3}{2a_i^2\theta_i^2} (\frac{b_i}{a_i} - 1)^2 \geq 0$. So $\sum a_i \log \frac{b_i}{a_i} + \sum (a_i - b_i) \leq 0$. And the equality holds when $\frac{a_i}{b_i} = 1$. Further more, $\sum \frac{a_i^3}{2a_i^2\theta_i^2} (\frac{b_i}{a_i} - 1)^2 \leq \sum \frac{a_i}{2} (a_i - b_i)^2$, accordingly we obtain $2 \sum a_i \log \frac{a_i}{b_i} \geq \sum a_i (a_i - b_i)^2$.

Lemma 8.3 Let $\sum a_i$ and $\sum b_i$ be convergent sequences. Then

$$\sum a_i \log \frac{a_i}{b_i} \geq (\sum a_i) \log \frac{\sum a_i}{\sum b_i}$$

Proof : $\frac{\sum a_i \log \frac{b_i}{a_i}}{\sum a_i} \leq \log \sum \frac{a_i}{\sum a_i} \frac{b_i}{a_i} = \log \sum \frac{b_i}{\sum a_i}$. Both sides multiplies -1 , we have $\sum a_i \log \frac{a_i}{b_i} \geq (\sum a_i) \log \frac{\sum a_i}{\sum b_i}$. The equality holds when $\frac{a_i}{b_i}$ are the same for all i , that is $\frac{a_i}{b_i}$ are constant.

Theorem 8.4 $\mathbb{H}(X) \leq \log |\mathcal{X}|$, where $|\mathcal{X}|$ denotes the number of elements in the range of X with equality iff X has a uniform distribution over \mathcal{X} .

Proof Suppose $p(x)$ and $q(x)$ are p.m.f of random variable X , the KullbackLeibler Divergence(KLD) between p and q are

$$D(p||q) = \sum_{x \in \mathcal{X}} p(x) \log \frac{p(x)}{q(x)} \geq 0$$

Alternatively, $-\sum_{x \in \mathcal{X}} p(x) \log p(x) + \sum_{x \in \mathcal{X}} p(x) \log q(x) \leq 0$, that is $\mathbb{H}(X) \leq -\sum_{x \in \mathcal{X}} p(x) \log q(x)$.

Let $q(x) = \frac{1}{|\mathcal{X}|}$, we have $\mathbb{H}(X) \leq \log |\mathcal{X}|$, which complete the proof.

Theorem 8.5 (Condition Reduces Entropy)

$$\mathbb{H}(X|Y) \leq \mathbb{H}(X)$$

with the equality iff X and Y are independent.

Definition 8.8 (Differential Entropy)

$$\mathbb{H}(X) = - \int_{\mathcal{S}} f(x) \log f(x) dx$$

, where \mathcal{S} is the support set of random variable X (if $f(x)$ exists) and $f(x)$ is p.d.f of X .

Example 8.3 Suppose random variable X is uniformly distributed on $(0, a)$. Then $\mathbb{H}(X) = - \int_0^a \frac{1}{a} \log \frac{1}{a} dx = \log a$. $\mathbb{H}(X) \leq 0$, when $0 \leq a \leq 1$.

Similarly, suppose the p.d.f of X_1, X_2, \dots, X_n is $f(x_1, x_2, \dots, x_n)$, then

$$\mathbb{H}(x_1, x_2, \dots, x_n) = - \int f(x_1, x_2, \dots, x_n) \log f(x_1, x_2, \dots, x_n) dx_1 dx_2 \dots dx_n$$

$$\mathbb{H}(X|Y) = - \int f(x, y) \log f(x|y) dx dy$$

Example 8.4 Let $X = (X_1, \dots, X_n)$ is gaussian distribution, that is, $X \sim N(\mu, \Sigma)$.

$$\begin{aligned} \mathbb{H}(X) &= - \int f_X(x) \left(-\frac{n}{2} \log 2\pi - \frac{1}{2} \log |\Sigma| - \frac{1}{2} (X - \mu)^T \Sigma^{-1} (X - \mu) \right) dX \\ &= \frac{1}{2} (n \log 2\pi + \log |\Sigma| + \int f_X(x) \text{tr}((X - \mu)^T \Sigma^{-1} (X - \mu)) dX) \\ &= \frac{1}{2} (n \log 2\pi + \log |\Sigma| + \int f_X(x) \text{tr}(\Sigma^{-1} (X - \mu)(X - \mu)^T) dX) \\ &= \frac{1}{2} (n \log 2\pi + \log |\Sigma| + \text{tr}(\Sigma^{-1} \int f_X(x) (X - \mu)(X - \mu)^T dX)) \\ &= \frac{1}{2} (n \log 2\pi + \log |\Sigma| + \text{tr}(\Sigma^{-1} \Sigma)) \\ &= \frac{1}{2} (n \log 2\pi + \log |\Sigma| + n) \end{aligned}$$

Definition 8.9 Suppose $X \sim f(X)$ and $Y \sim g(X)$, then

$$D(f||g) = \int f(X) \log \frac{f(X)}{g(X)} dX$$

Note that $D(f||g)$ is finite only if the support of f is contained in the support of g .

Definition 8.10 (Mutual Information)

$$\mathbb{I}(X, Y) = \int f(x, y) \log \frac{f(x, y)}{f(x)f(y)} dx dy = D(f(x, y)||f(x)f(y))$$

$$\mathbb{I}(X, Y) = \mathbb{H}(X) - \mathbb{H}(X|Y) = \mathbb{H}(Y) - \mathbb{H}(Y|X)$$

Theorem 8.6

$$D(f||g) \geq 0$$

with the equality iff $f = g$ at almost everywhere.

Proof : $\int f(x) \log \frac{g(x)}{f(x)} dx \leq \log \int f(x) \frac{g(x)}{f(x)} dx = \log \int g(x) dx = 0$, So $\int f(x) \log \frac{f(x)}{g(x)} dx \geq 0$, which complete the proof.

Corollary 8.2

- $\mathbb{I}(X, Y) \geq 0$, with the equality iff X and Y are independent.
- $\mathbb{H}(X|Y) \leq \mathbb{H}(X)$, with the equality iff X and Y are independent.

Theorem 8.7 *The chain rule for differential entropy*

$$\mathbb{H}(X_1, \dots, X_n) = \sum_{i=1}^n \mathbb{H}(X_i | X_1, \dots, X_{i-1})$$

Corollary 8.3

$$\mathbb{H}(X_1, \dots, X_n) \leq \sum_{i=1}^n \mathbb{H}(X_i)$$

Example 8.5 Suppose $\Sigma \in \mathbf{S}_{++}^n$, where $\Sigma = [\sigma_{ij}]$ then

$$\det \Sigma \leq \prod_{i=1}^n \sigma_{ii}$$

Proof : Suppose $X = (X_1, \dots, X_n) \sim N(0, \Sigma)$, so $X_i \sim N(0, \sigma_{ii})$. $\mathbb{H}(X_1, \dots, X_n) = \frac{1}{2}(n \log 2\pi + \log \det \Sigma + n)$, $\mathbb{H}(X_i) = \frac{1}{2}(\log 2\pi + \log \sigma_{ii} + 1)$. Since $\mathbb{H}(X_1, \dots, X_n) \leq \sum_{i=1}^n \mathbb{H}(X_i)$, we have $\frac{1}{2}(n \log 2\pi + \log \det \Sigma + n) \leq \frac{n}{2} \log 2\pi + \frac{1}{2} \sum_{i=1}^n \log \sigma_{ii} + \frac{n}{2}$, thus $\log \det \Sigma \leq \sum_{i=1}^n \log \sigma_{ii}$. So $\det \Sigma \leq \prod_{i=1}^n \sigma_{ii}$ holds.

Theorem 8.8

$$\mathbb{H}(\alpha X + c) = \mathbb{H}(X) + \log |\alpha|$$

, where $\alpha \geq 0$

Proof : Let $Y = \alpha X + c$, then $f_Y(y) = \frac{1}{|\alpha|} f_X(\frac{Y-c}{\alpha})$

$$\begin{aligned} \mathbb{H}(\alpha X + c) &= - \int f_Y(y) \log f_Y(y) dy \\ &= - \int \frac{1}{|\alpha|} f_X\left(\frac{Y-c}{\alpha}\right) \left(\log \frac{1}{|\alpha|} + \log f_X\left(\frac{Y-c}{\alpha}\right)\right) dy \\ &= - \int f_X(X) \left(\log \frac{1}{|\alpha|} + \log f_X(X)\right) dx \\ &= \mathbb{H}(X) + \log |\alpha| \end{aligned}$$

Corollary 8.4 Suppose \mathbf{A} is nonsingular, then $\mathbb{H}(\mathbf{A}X) = \mathbb{H}(X) + \log |\mathbf{A}|$.

Theorem 8.9 Let $X \in \mathbb{R}^m$ have zero mean and covariance $\Sigma = \mathbb{E}[XX^T]$, then

$$\mathbb{H}(X) \leq \frac{1}{2} \log((2\pi)^n |\Sigma|) + \frac{n}{2}$$

Proof : Suppose $g(X)$ is p.d.f of X , we also let $f(X) = N \sim (0, \Sigma)$.

$$\begin{aligned} 0 &\leq D(g||f) \\ &= \int g \log \frac{g}{f} \\ &= \int g \log g - \int g \log f \\ &= -\mathbb{H}(X) - \int g \log f \end{aligned}$$

Hence

$$\begin{aligned} \mathbb{H}(X) &\leq - \int g \log f \\ &= - \int g(x) \left(-\frac{n}{2} \log 2\pi - \frac{1}{2} \log |\Sigma| - \frac{1}{2} (X - \mu)^T \Sigma^{-1} (X - \mu) \right) dX \\ &= \frac{1}{2} (n \log 2\pi + \log |\Sigma| + \text{tr}(\Sigma^{-1} \int g(x) (X - \mu)(X - \mu)^T dX)) \\ &= \frac{1}{2} \log((2\pi)^n |\Sigma|) + \frac{n}{2} \end{aligned}$$

8.4 The Exponential Family

Consider the p.d.f $p(x)$ which satisfies the k (independent) constraints,

$$\int_{\mathcal{X}} h_i(x) p(x) dx = m_i < \infty, \quad i = 1, \dots, k$$

, where m_1, \dots, m_k are specified constants. We want to find certain p.d.f $p(x)$ that is closest to $f(x)$. That is,

$$\min_p \int p(x) \log \frac{p(x)}{f(x)} dx \quad \text{s.t.} \quad \int_{\mathcal{X}} h_i(x) p(x) dx = m_i < \infty, \quad i = 1, \dots, k, \text{ and } \int p(x) dx = 1.$$

This is an optimization problem with the object function

$$F(p) = \int p(x) \log \frac{p(x)}{f(x)} dx + \sum_{i=1}^k \theta_i \left(\int_{\mathcal{X}} h_i(x) p(x) dx - m_i \right) + c \left(\int_{\mathcal{X}} p(x) dx - 1 \right)$$

, where $\theta_i, i = 1, \dots, k$ and c are lagrange multipliers. Besides, $f(x)$ is known.

Theorem 8.10 *The function defined above is minimized by*

$$\begin{aligned} p(x) &= E_{f_k}(X|f, g, \vec{h}, \vec{\phi}, \vec{\theta}, \vec{c}) \\ &= \frac{1}{g(\theta)} f(x) \exp\left(\sum_{i=1}^k \theta_i h_i(x)\right), \end{aligned}$$

where $c_i = 1$, and $\vec{\phi} = \vec{\theta} = (\theta_1, \dots, \theta_k)$.

Proof :

$$\begin{aligned} dF(p) &= \lim_{\alpha \rightarrow 0} \int p(x + \alpha \tau(x)) \log \frac{p(x + \alpha \tau(x))}{f(x)} dx - \lim_{\alpha \rightarrow 0} \int p(x) \log \frac{p(x)}{f(x)} \\ &\quad + \sum_{i=1}^k \theta_i \left(\int h_i(x) (p(x) + \alpha \tau(x) - p(x)) dx \right) + c \left(\int p(x) + \alpha \tau(x) - p(x) dx \right) \\ &= \lim_{\alpha \rightarrow 0} \left(\int p(x) \log \left(1 + \alpha \frac{\tau(x)}{p(x)} \right) dx + \alpha \sum_{i=1}^k \int \theta_i h_i(x) \tau(x) dx + \alpha c \int \tau(x) dx \right) \end{aligned}$$

So

$$\begin{aligned} \frac{dF(p)}{dp} &= \int p(x) \lim_{\alpha \rightarrow 0} \frac{\log(1 + \alpha \frac{\tau(x)}{p(x)})}{\alpha} dx + \int \tau(x) \log \frac{p(x)}{f(x)} dx + \sum_{i=1}^k \int \theta_i h_i(x) \tau(x) dx + c \int \tau(x) dx \\ &= (c + 1) \left(\int \tau(x) dx \right) + \int \tau(x) \log \frac{p(x)}{f(x)} dx + \sum_{i=1}^k \int \theta_i h_i(x) \tau(x) dx \end{aligned}$$

For any small $\tau(x)$, $\frac{dF(p)}{dp} = 0$. Thus

$$c + 1 + \log \frac{p(x)}{f(x)} + \sum_{i=1}^k \theta_i h_i(x) = 0$$

, which means $p(x) = \frac{1}{g(\theta)} f(x) \exp\left(\sum_{i=1}^k \theta_i h_i(x)\right)$, where $g(\theta) = \int_{x \in \mathcal{X}} f(x) \exp\left(\sum_{i=1}^k \theta_i h_i(x)\right) dx$.