

Exam2: Part 3

A) a: 6 5 4 3 2 1 0 Decreasing integers

- Find the midpoint of array $\frac{n}{2}$.
- Divide array a into a_{left} or a_{right} .
- If value of midpoint is greater than i then look in a_{left} since all the numbers less than midpoint are in a_{left} , else look in a_{right} . Also if $i = \text{midpoint}$ then return -1.
- Recursively look in either a_{left} or a_{right} and if it equals midpoint return -1.

$$T(n) = T\left(\frac{n}{2}\right) + 1$$

$$f(n) = 1 \quad n^{\log_2 1} = n^1 = 1$$

$$f(n) = \Theta(n^{\log_2 1})$$

$$T(n) = \Theta(\log n) \leftarrow \because \text{By Case 2 of MT}$$

Explanation for Algorithm: This algorithm checks each midpoint and divides the problem into smaller problems to check each a_{left} or a_{right} . Since we are cutting the array in half and only doing work on just 1 subarray, then we are only checking $\log(n)$ of the array. Since we are only checking midpoint, that takes constant time.

B)

$i \backslash j$	λ	a	b	c	a	a	c	a
λ	0	1	2	3	4	5	6	7
b	1	1	1	2	3	4	5	6
c	2	2	2	1	2	3	4	5
a	3	2	3	2	1	2	3	4
b	4	3	2	3	2	2	3	4
c	5	4	3	2	3	3	2	3
a	6	5	4	3	2	3	3	2
b	7	6	5	4	3	3	4	3 [↑]

Total number of edits = 3

One example of a transformation -b

which will make $u = bcabcbab$ into
 $bcabca$