

Quantum Error Correction

Quantum Capita Selecta

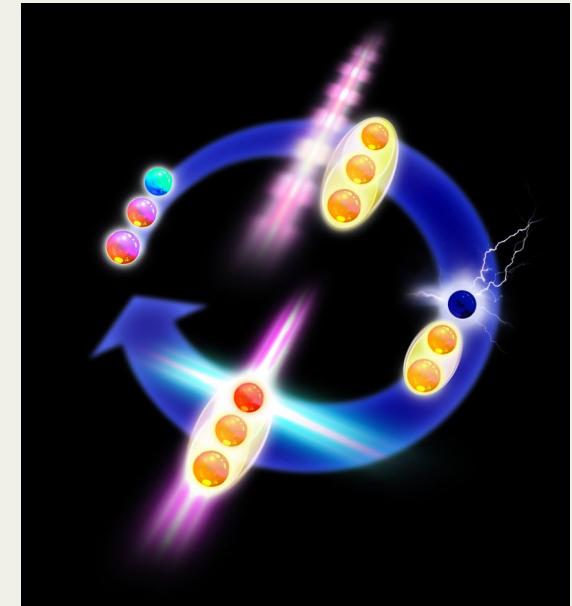
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Why Quantum Error Correction (QEC)?



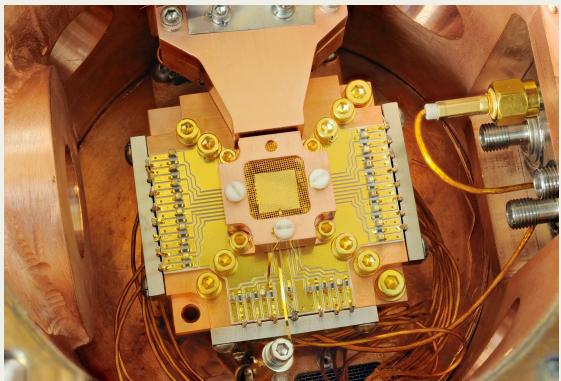
- Quantum Error Correction Codes (QECC) are algorithms too!
 - In general with Clifford gates only
- Sophisticated quantum–classic interaction
- Step towards quantum programming model



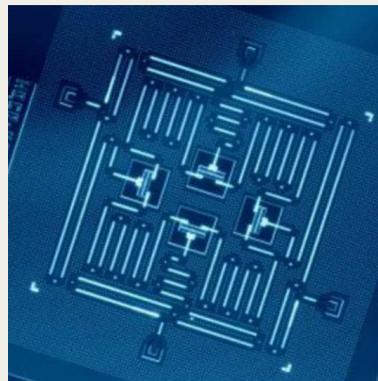
Quantum hardware



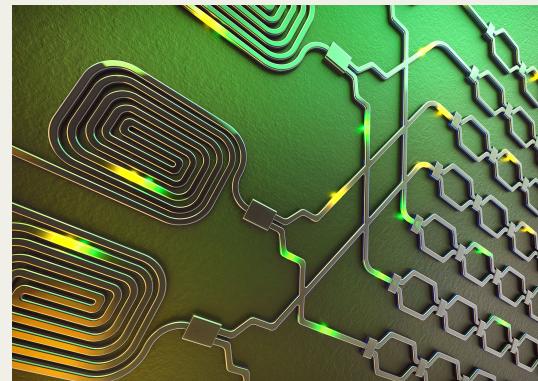
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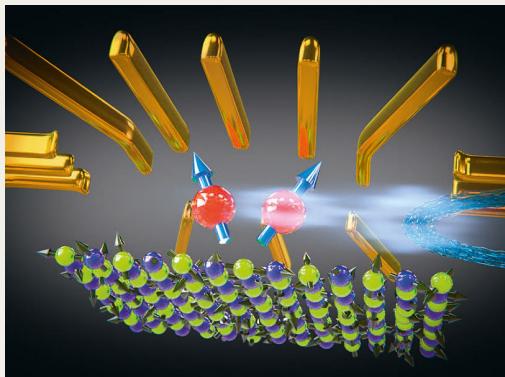
Ion traps



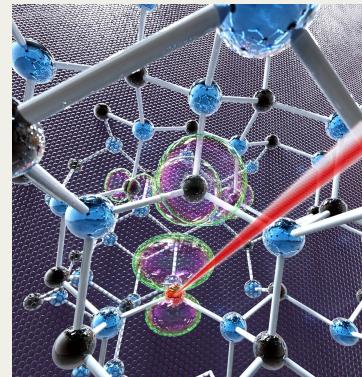
Superconductors



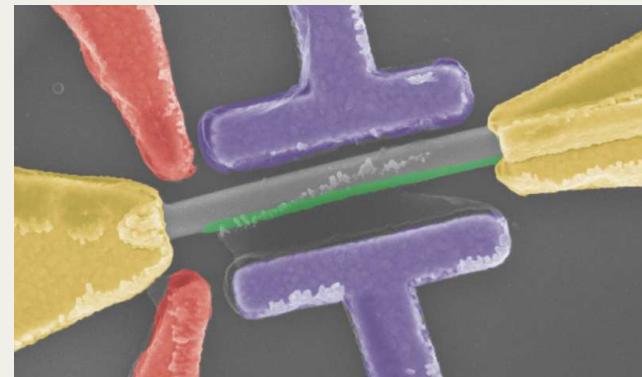
Photons



Quantum dots



NV centers

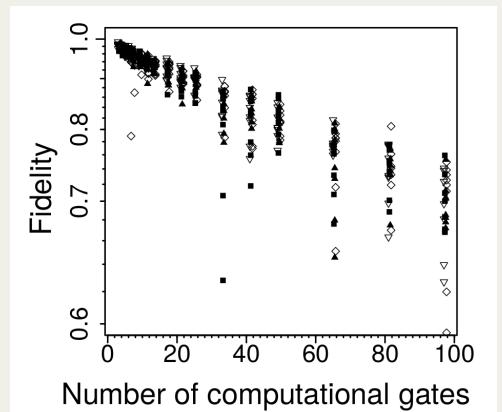
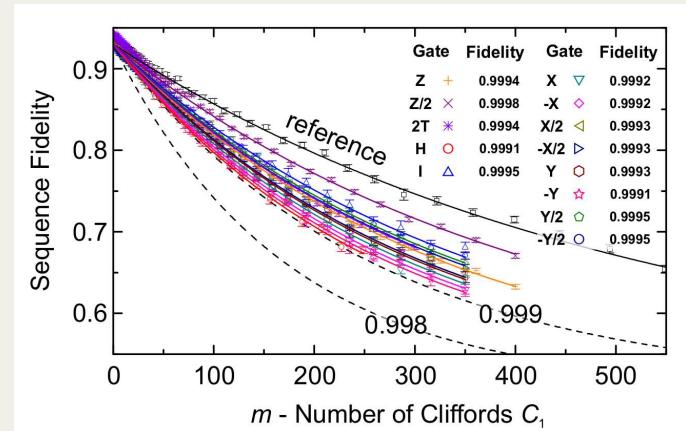


Majoranas

Qubits are fragile

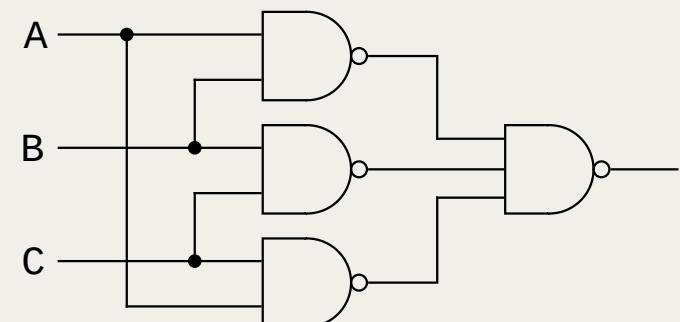
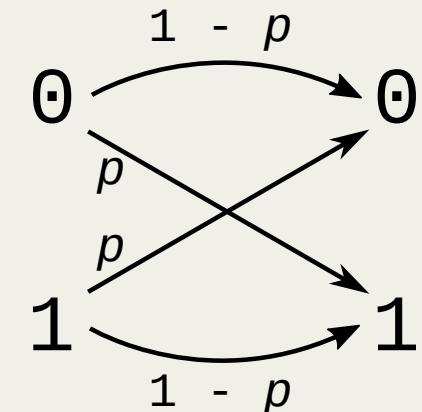


- Qubits lose their state in a short period
 - Tens of microseconds for superconducting qubits
- Quantum gates have high error rates
 - Around 0.1% for superconducting qubits
 - Factoring a 2000-bit number requires error rates around 4×10^{-13}

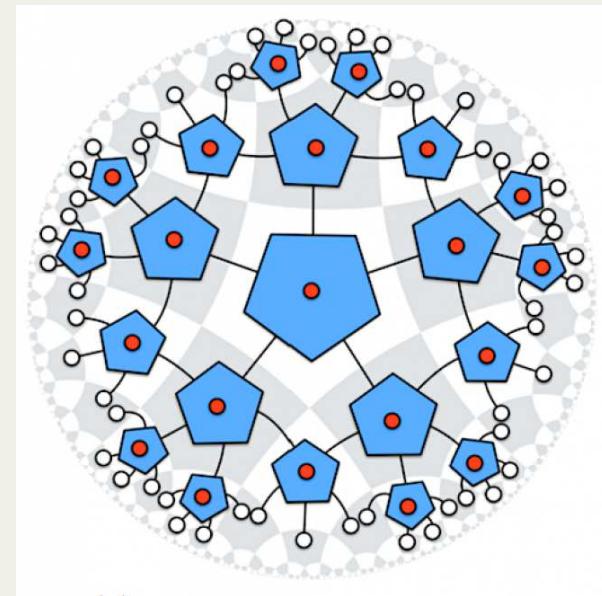


Classic Error Correction

- Error/Noise:
 - Independent single-bit errors
 - Probability of error: p
- Parity bit/Checksum
 - Detection only
- Repetition code/majority voting
- More redundancy, larger code space, better error resistance



- Encode a quantum state redundantly into a larger Hilbert space (logical qubits)
- Costs:
 - Space
 - Time
- Benefits of logical qubits:
 - Longer coherence times
 - Lower error rates





- Measurements are destructive

$$\alpha |0\rangle + \beta |1\rangle = \begin{cases} M(|\psi\rangle) = |0\rangle \\ M(|\psi\rangle) = |1\rangle \end{cases}$$

- No cloning

$$\alpha |0\rangle + \beta |1\rangle \not\Rightarrow \begin{cases} \alpha |0\rangle + \beta |1\rangle \\ \alpha |0\rangle + \beta |1\rangle \\ \alpha |0\rangle + \beta |1\rangle \end{cases}$$

- Errors are continuous:

$$|\psi\rangle |e\rangle \rightarrow \sqrt{1-p}I |\psi\rangle |e\rangle + \sqrt{p}X |\psi\rangle |\hat{e}\rangle$$

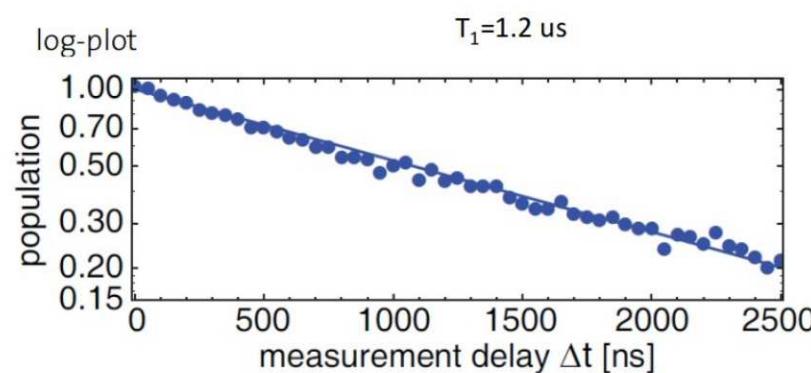
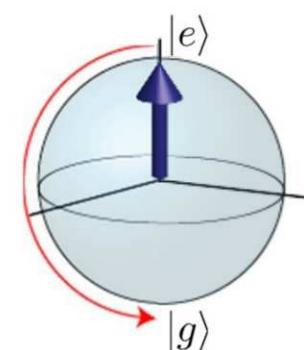
- Bit–flip channel (independent X errors)

$$|\psi\rangle |e\rangle \rightarrow X |\psi\rangle |\hat{e}\rangle$$

- Phase–flip channel (independent Z errors)

$$|\psi\rangle |e\rangle \rightarrow Z |\psi\rangle |\hat{e}\rangle$$

- Amplitude damping/relaxation



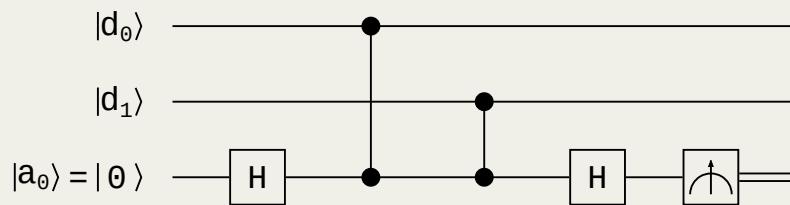


- Indirect measurements (stabilizers)
 - Learn about the state without fully collapsing it
 - Forces discrete errors
- All errors can be decomposed in X and Z errors
- Quantum memory

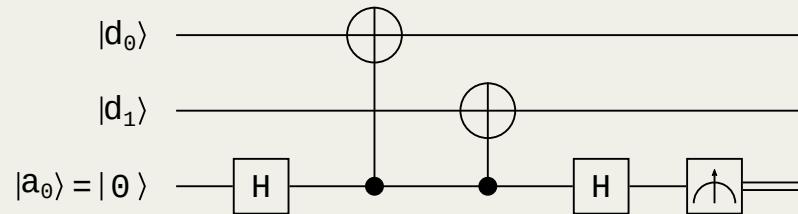
Stabilizer measurements

- Notation formalism, most important is that we can measure “stabilizers”

$Z_0 Z_1$ stabilizer



$X_0 X_1$ stabilizer





Specifications

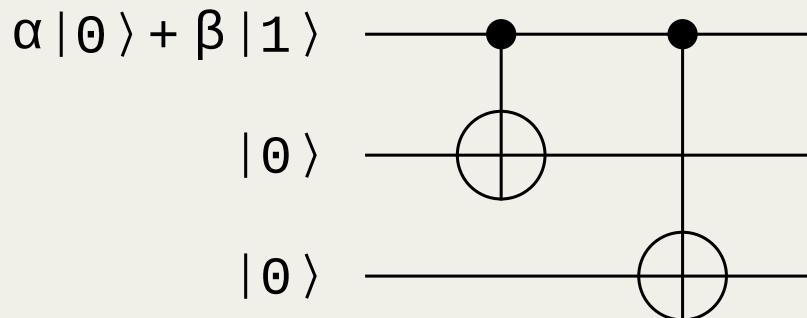
- 3 data qubits
- 2 ancilla qubits
- Detects up to two X errors
- Corrects up to one X error
- Not resistant to Z errors

Steps

- Encode state
- Extract error syndrome
- Decode error syndrome
- Correct
- Decode state

Encode state

- Encodes a single qubit state in a larger Hilbert space (redundancy)



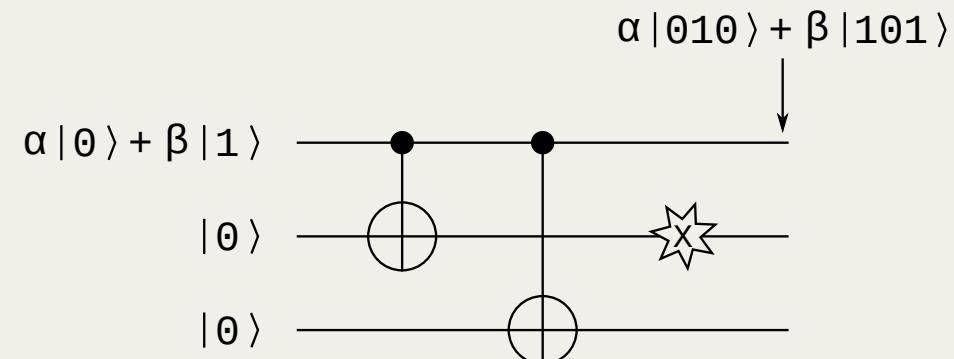
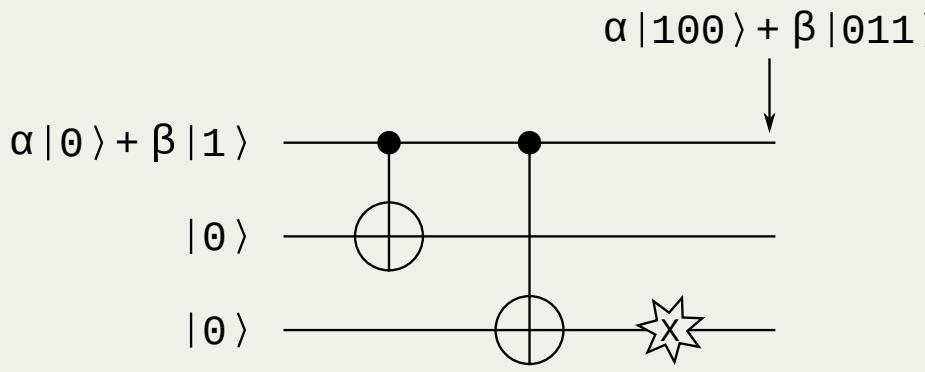
$$\begin{aligned} &\rightarrow |00\rangle \otimes |\psi\rangle \\ &\rightarrow |00\rangle \otimes (\alpha |0\rangle + \beta |1\rangle) \\ &\rightarrow \alpha |000\rangle + \beta |001\rangle \\ CNOT_{0,1} &\rightarrow \alpha |000\rangle + \beta |011\rangle \\ CNOT_{0,2} &\rightarrow \alpha |000\rangle + \beta |111\rangle \end{aligned}$$

$$|\psi_L\rangle = \alpha |000\rangle + \beta |111\rangle$$

Errors



- Independent X errors can occur with probability p (bit–flip channel)
- Single error changes parity of the encoded state



$$\rightarrow \alpha|000\rangle + \beta|001\rangle$$

$$CNOT_{0,1} \rightarrow \alpha|000\rangle + \beta|011\rangle$$

$$CNOT_{0,2} \rightarrow \alpha|000\rangle + \beta|111\rangle$$

$$X_2 \rightarrow \alpha|100\rangle + \beta|011\rangle$$

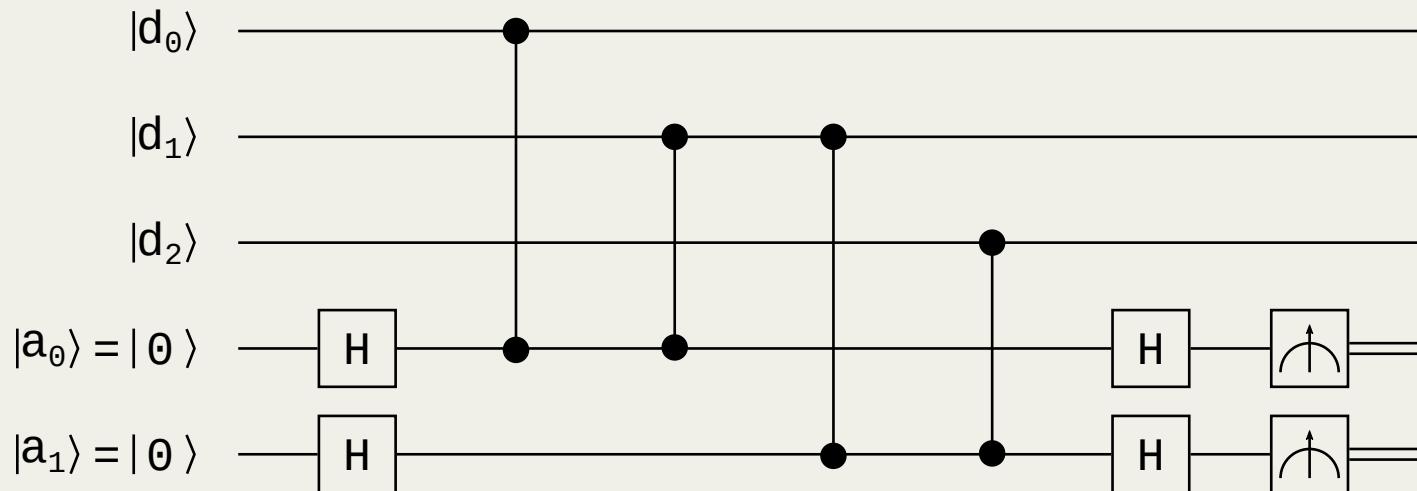
$$\rightarrow \alpha|000\rangle + \beta|001\rangle$$

$$CNOT_{0,1} \rightarrow \alpha|000\rangle + \beta|011\rangle$$

$$CNOT_{0,2} \rightarrow \alpha|000\rangle + \beta|111\rangle$$

$$X_1 \rightarrow \alpha|010\rangle + \beta|101\rangle$$

Extract error syndrome

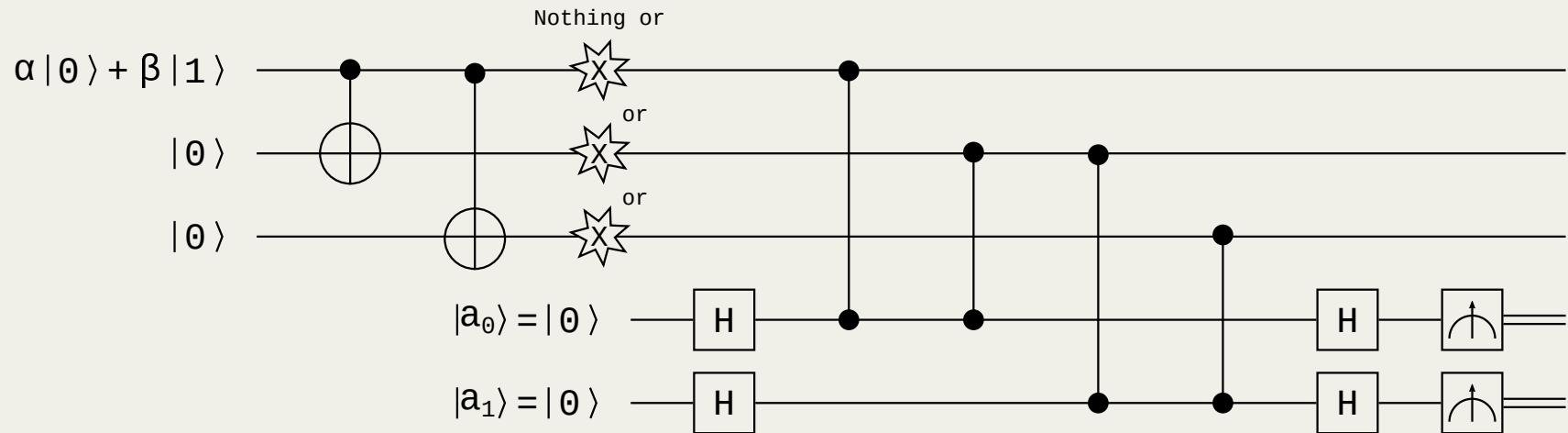


- Measuring stabilizers: Z_0Z_1 , Z_1Z_2
 - Error Syndrome Measurement
- For an encoded state in the code space:

$$M(a_0) = M(a_1) = 0$$

$$|\psi_L\rangle = \alpha |000\rangle + \beta |111\rangle$$

Error syndrome measurement



$$\begin{aligned}
 |\psi_L\rangle &= \alpha|000\rangle + \beta|111\rangle \\
 X_{d_0} \rightarrow &\alpha|001\rangle + \beta|110\rangle \\
 X_{d_1} \rightarrow &\alpha|010\rangle + \beta|101\rangle \\
 X_{d_2} \rightarrow &\alpha|100\rangle + \beta|011\rangle
 \end{aligned}$$

Error	$M(a_1)$	$M(a_0)$
None	0	0
X_{d_0}	0	1
X_{d_1}	1	1
X_{d_2}	1	0

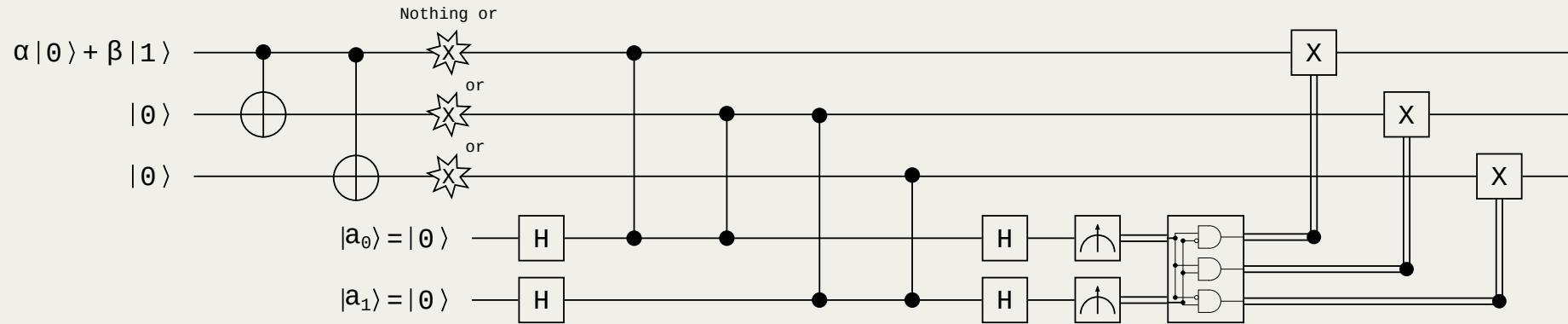


- Decoder reads error syndrome and outputs most likely error
 - Decoding is a fully classic calculation!
 - Find the least amount of errors that explains the error syndrome
 - Corrections can be wrong and generate an unrecoverable error

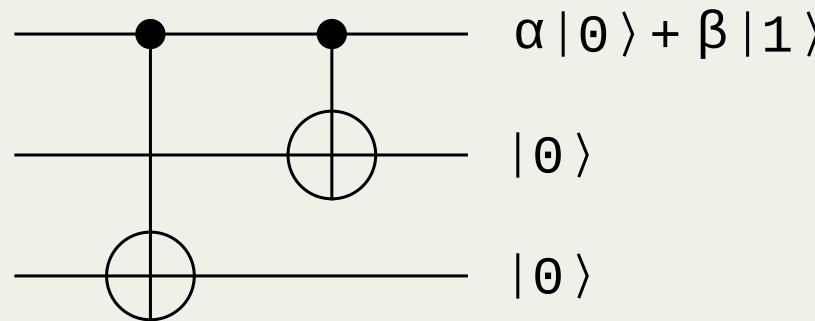
Error	$M(a_1)$	$M(a_0)$	$\Pr \{ \}$
X_{d_0}	0	1	p
$X_{d_1 d_2}$	0	1	p^2

$$\begin{aligned} |\psi_L\rangle &= \alpha |000\rangle + \beta |111\rangle \\ X_{d_0} \rightarrow &\alpha |001\rangle + \beta |110\rangle \\ X_{d_1} X_{d_2} \rightarrow &\alpha |110\rangle + \beta |001\rangle \end{aligned}$$

- Apply correction for most likely error



- Reverse encoding circuit (uncompute)
- We can also measure any qubit directly in the encoded state



$$\rightarrow \alpha |000\rangle + \beta |111\rangle$$

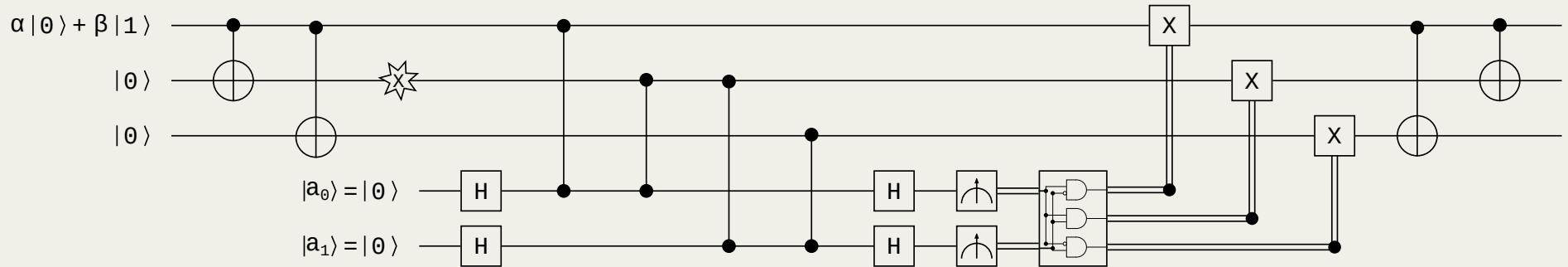
$$CNOT_{0,2} \rightarrow \alpha |000\rangle + \beta |011\rangle$$

$$CNOT_{0,1} \rightarrow \alpha |000\rangle + \beta |001\rangle$$

$$\rightarrow |00\rangle \otimes (\alpha |0\rangle + \beta |1\rangle)$$

$$\rightarrow |00\rangle \otimes |\psi\rangle$$

Example



$$\rightarrow |a_1 a_0\rangle \otimes |00\rangle \otimes |\psi\rangle$$

$$\rightarrow |00\rangle \otimes |00\rangle \otimes (\alpha|0\rangle + \beta|1\rangle)$$

$$\rightarrow \alpha|00000\rangle + \beta|00001\rangle$$

$$CNOT_{0,1} \rightarrow \alpha|00000\rangle + \beta|00011\rangle$$

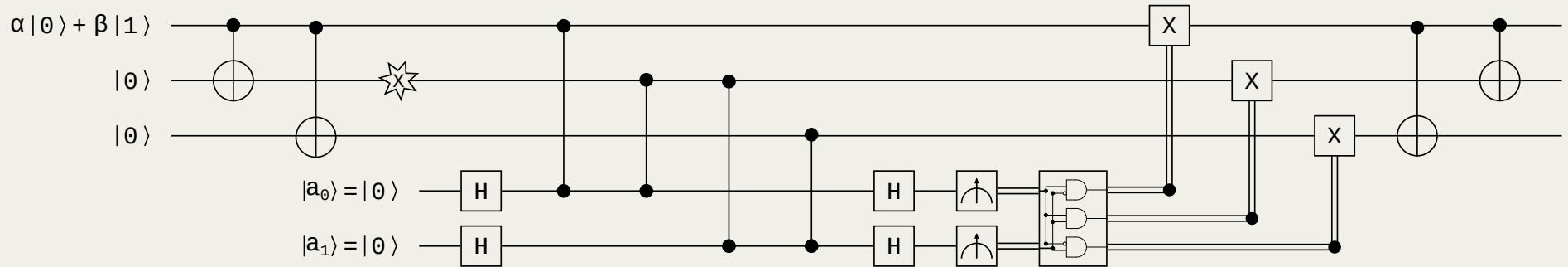
$$CNOT_{0,2} \rightarrow \alpha|00000\rangle + \beta|00111\rangle$$

$$X_1 \rightarrow \alpha|00010\rangle + \beta|00101\rangle$$

$$H_{a_0} \rightarrow \frac{1}{\sqrt{2}}(\alpha|00010\rangle + \alpha|01010\rangle + \beta|00101\rangle + \beta|01101\rangle)$$

$$H_{a_1} \rightarrow \frac{1}{2}(\alpha|00010\rangle + \alpha|10010\rangle + \alpha|01010\rangle + \alpha|11010\rangle + \beta|00101\rangle + \beta|10101\rangle + \beta|01101\rangle + \beta|11101\rangle)$$

Example



$$\rightarrow \frac{1}{2} (\alpha |00010\rangle + \alpha |10010\rangle + \alpha |01010\rangle + \alpha |11010\rangle + \beta |00101\rangle + \beta |10101\rangle + \beta |01101\rangle + \beta |11101\rangle)$$

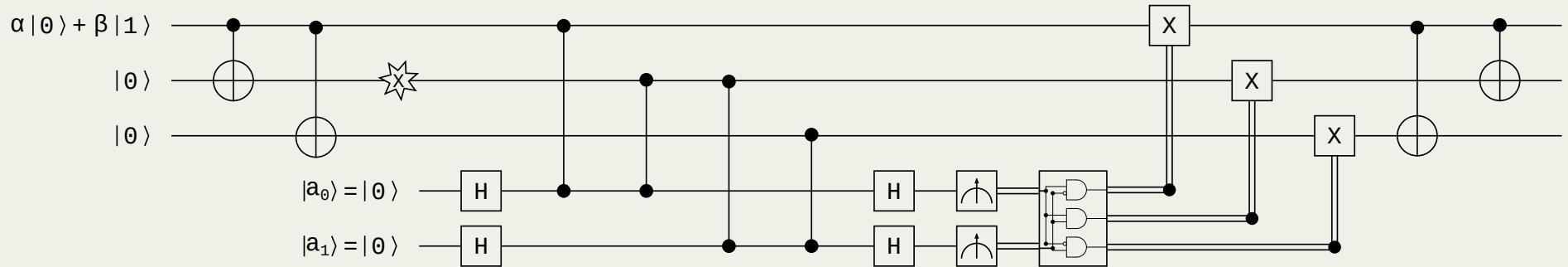
$$CZ_{0, a_0} \rightarrow \frac{1}{2} (\alpha |00010\rangle + \alpha |10010\rangle + \alpha |01010\rangle + \alpha |11010\rangle + \beta |00101\rangle + \beta |10101\rangle - \beta |01101\rangle - \beta |11101\rangle)$$

$$CZ_{1, a_0} \rightarrow \frac{1}{2} (\alpha |00010\rangle + \alpha |10010\rangle - \alpha |01010\rangle - \alpha |11010\rangle + \beta |00101\rangle + \beta |10101\rangle - \beta |01101\rangle - \beta |11101\rangle)$$

$$CZ_{1, a_1} \rightarrow \frac{1}{2} (\alpha |00010\rangle - \alpha |10010\rangle - \alpha |01010\rangle + \alpha |11010\rangle + \beta |00101\rangle + \beta |10101\rangle - \beta |01101\rangle - \beta |11101\rangle)$$

$$CZ_{2, a_1} \rightarrow \frac{1}{2} (\alpha |00010\rangle - \alpha |10010\rangle - \alpha |01010\rangle + \alpha |11010\rangle + \beta |00101\rangle - \beta |10101\rangle - \beta |01101\rangle + \beta |11101\rangle)$$

Example

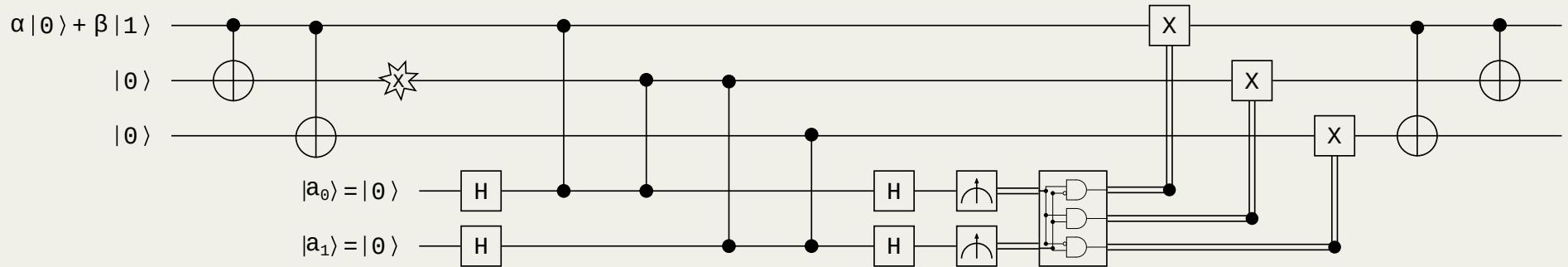


$$\rightarrow \frac{1}{2} (\alpha |00010\rangle - \alpha |10010\rangle - \alpha |01010\rangle + \alpha |11010\rangle + \beta |00101\rangle - \beta |10101\rangle - \beta |01101\rangle + \beta |11101\rangle)$$

$$\begin{aligned}
 H_{a_0} &\rightarrow \frac{1}{2\sqrt{2}} (\alpha |00010\rangle + \alpha |01010\rangle - \alpha |10010\rangle - \alpha |11010\rangle - \alpha |00010\rangle + \alpha |01010\rangle + \alpha |10010\rangle - \alpha |11010\rangle + \\
 &+ \beta |00101\rangle + \beta |01101\rangle - \beta |10101\rangle - \beta |11101\rangle - \beta |00101\rangle + \beta |01101\rangle + \beta |10101\rangle - \beta |11101\rangle) \\
 &\rightarrow \frac{1}{2\sqrt{2}} (2\alpha |01010\rangle - 2\alpha |11010\rangle + 2\beta |01101\rangle - 2\beta |11101\rangle) \\
 &\rightarrow \frac{1}{\sqrt{2}} (\alpha |01010\rangle - \alpha |11010\rangle + \beta |01101\rangle - \beta |11101\rangle)
 \end{aligned}$$

$$\begin{aligned}
 H_{a_1} &\rightarrow \frac{1}{2} (\alpha |01010\rangle + \alpha |11010\rangle - \alpha |01010\rangle + \alpha |11010\rangle + \beta |01101\rangle + \beta |11101\rangle - \beta |01101\rangle + \beta |11101\rangle) \\
 &\rightarrow \frac{1}{2} (2\alpha |11010\rangle + 2\beta |11101\rangle) \\
 &\rightarrow \alpha |11010\rangle + \beta |11101\rangle
 \end{aligned}$$

Example



$$\rightarrow \alpha |11010\rangle + \beta |11101\rangle$$

$$\rightarrow M(a_1) = M(a_0) = 1$$

CLAS $\rightarrow 010$

$$X_1 \rightarrow \alpha |11000\rangle + \beta |11111\rangle$$

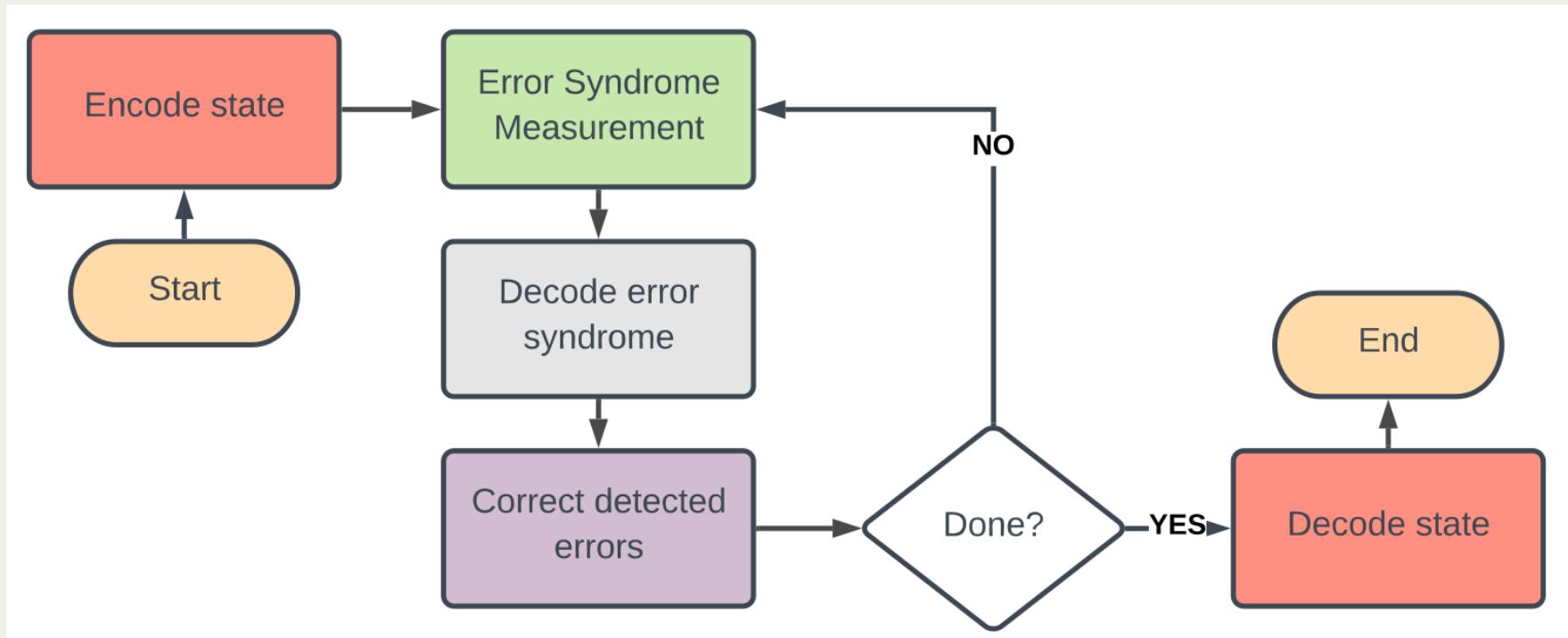
$$CNOT_{0,2} \rightarrow \alpha |11000\rangle + \beta |11011\rangle$$

$$CNOT_{0,1} \rightarrow \alpha |11000\rangle + \beta |11001\rangle$$

$$\rightarrow |11\rangle \otimes |00\rangle \otimes (\alpha |0\rangle + \beta |1\rangle)$$

$$\rightarrow |a_1 a_0\rangle \otimes |00\rangle \otimes |\psi\rangle$$

Bit-flip code summary

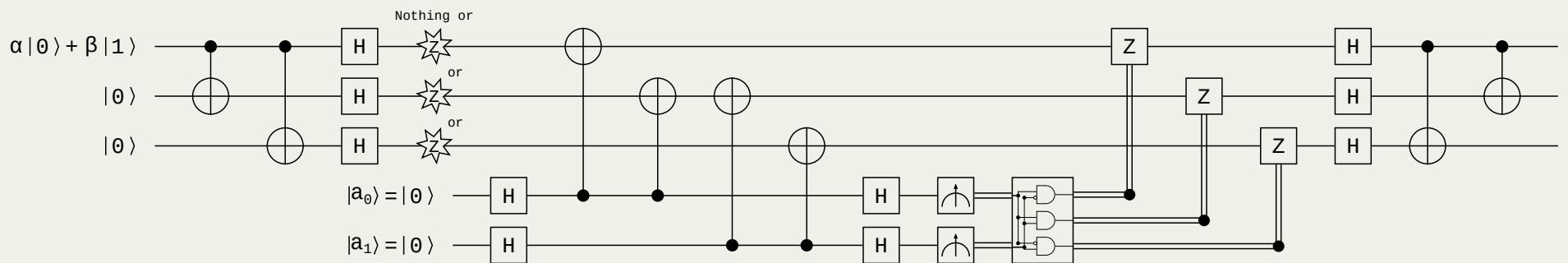


Phase–flip code

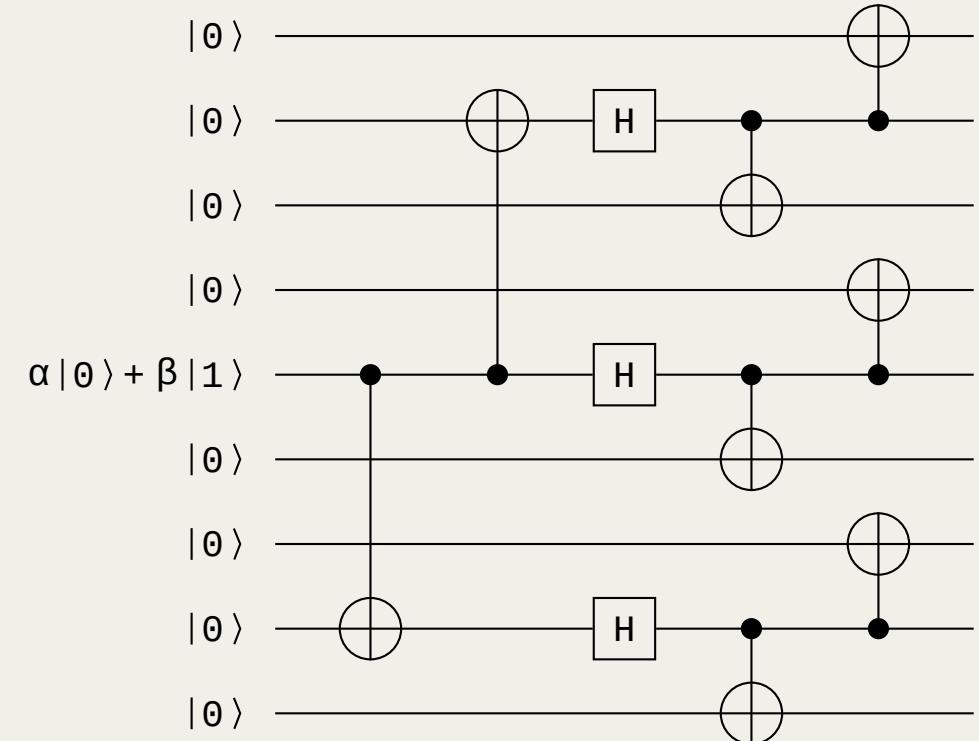


- Same principle in a different basis, now detecting Z errors

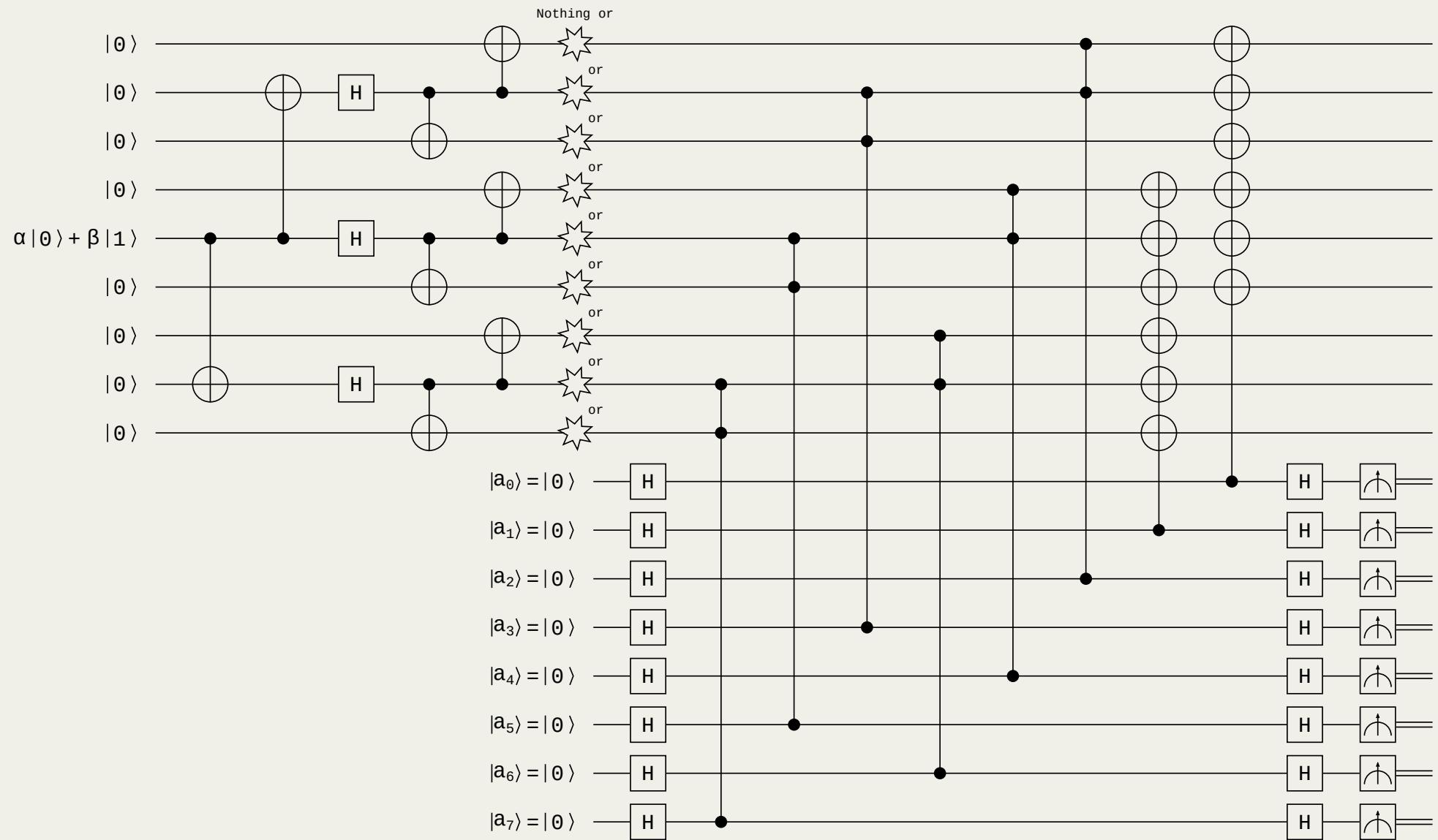
$$|\psi_L\rangle = \alpha |+++ \rangle + \beta |--- \rangle$$



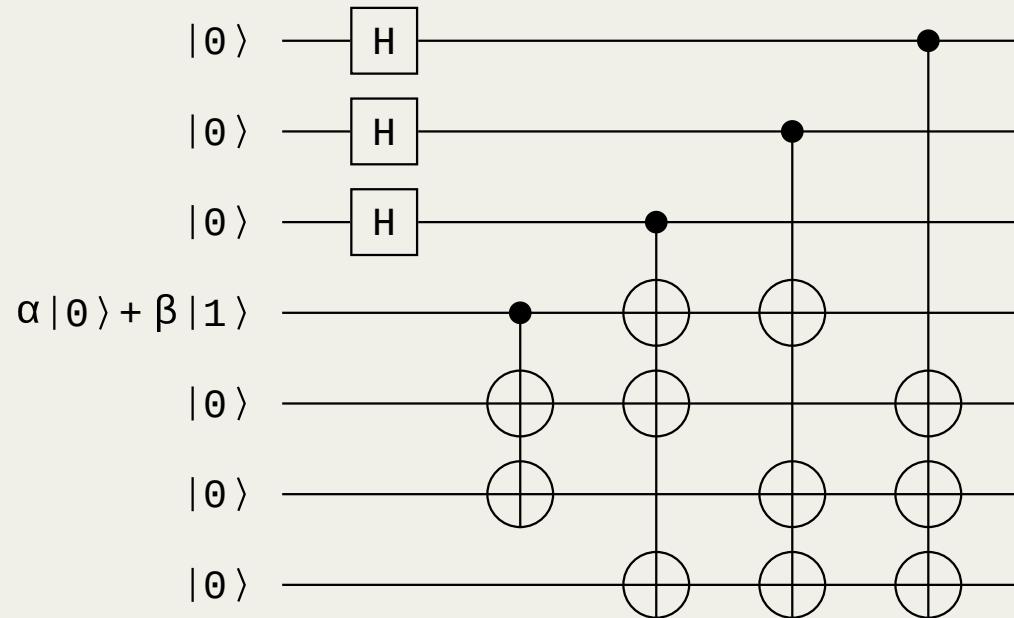
- Concatenated code
 - Phase–flip code (detects Z errors)
 - Three times bit–flip code (detects X errors)
- First complete QECC



Shor code ESM

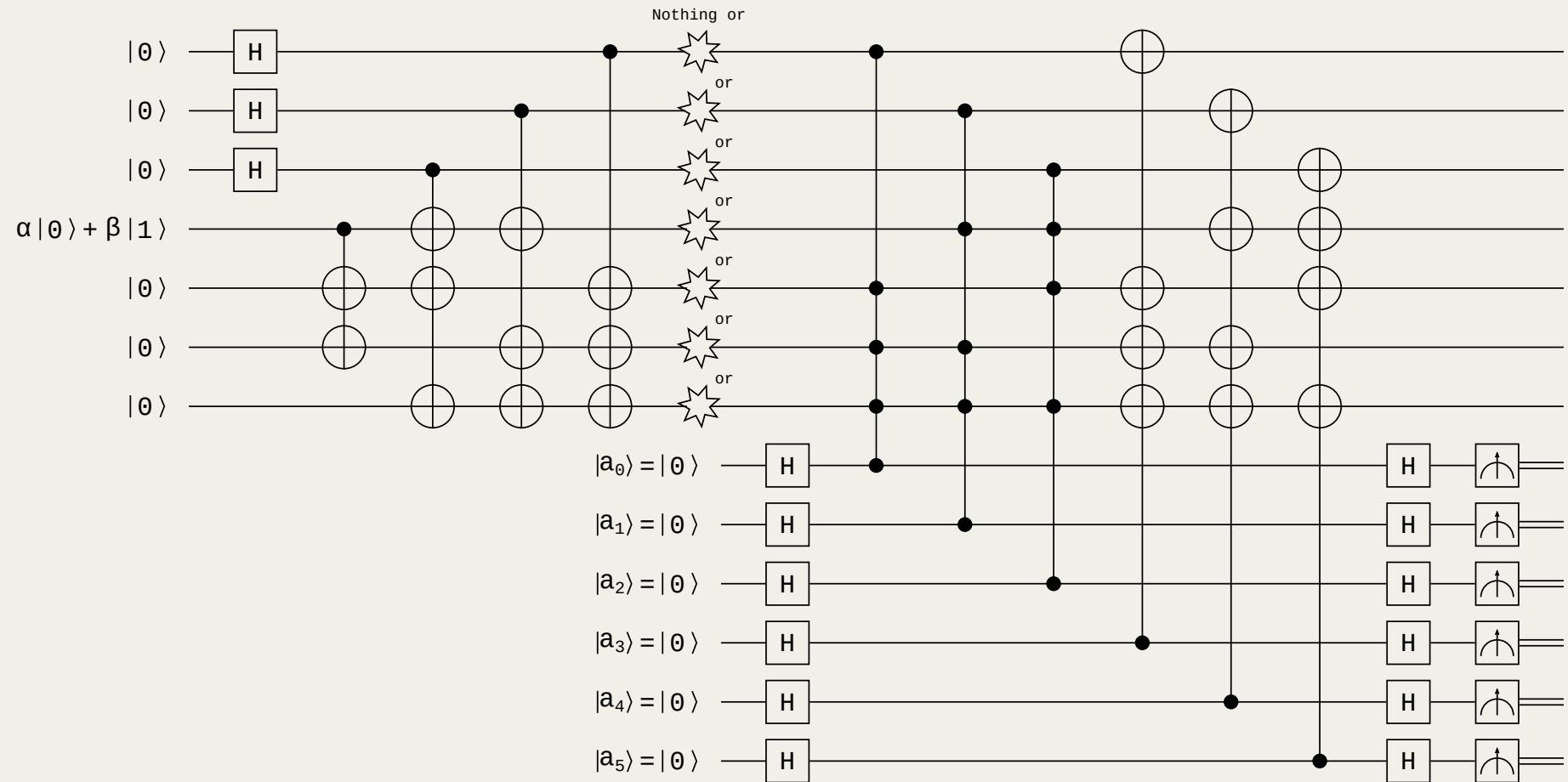


- Mathematically a very convenient code
- Corrects up to one X and one Z error



- $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle \rightarrow |\Psi\rangle = \alpha|0_{Steane}\rangle + \beta|1_{Steane}\rangle$

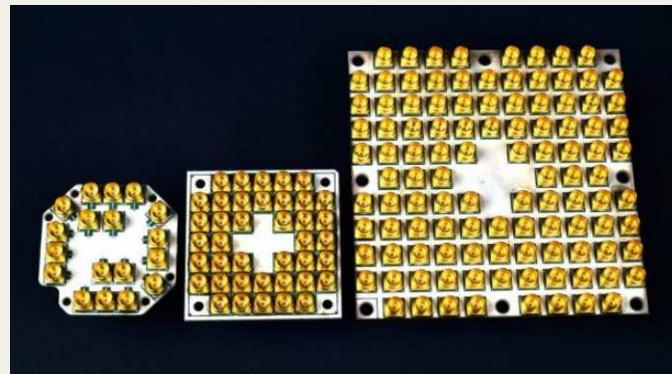
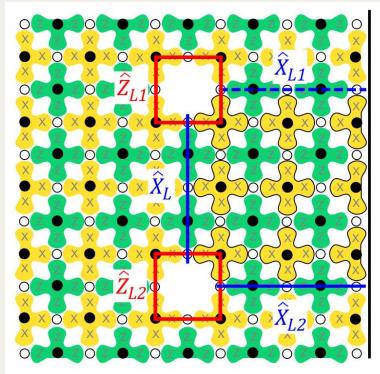
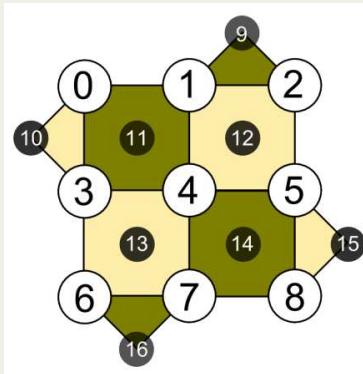
Steane code ESM



Other well-known codes



- Planar surface code (SC7, SC17, SC49) (target of QuTech)
- Double-defect surface code



Towards fault-tolerant Universal QC



- Resistance to ancilla circuit errors
- Multiple logical qubits
- Logical operations (no decoding required)
 - Universal (e.g. X , Z , H , S , $CNOT$, T)

