

# ECE 3140 / CS 3420

# EMBEDDED SYSTEMS

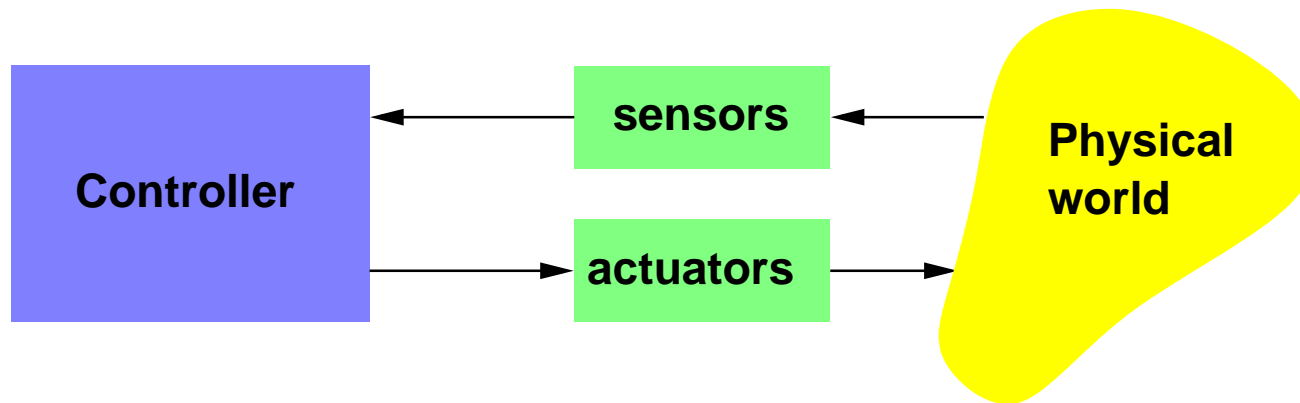
## LECTURE 14

**Prof. José F. Martínez**  
TR 1:25-2:40pm in 150 Olin



# INTRODUCTION TO REAL TIME

Time-critical systems are the norm.



Tight interaction between sensing and actuation  
⇒ need predictable timing of operations

We won't cover the details of *how* a system is controlled.



# INTRODUCTION TO REAL TIME

A computing system that is able to respond to events within *precise* timing constraints is a real-time system.

- Correct operation depends on
  - Usual properties (producing the correct output, etc)
  - Also on the *time* at which the output is produced

Some interesting observations:

- Time between different entities must be *synchronized*  
Note: time synchronization is not a simple problem
- Real time is *not* the same as fast!



# PERFORMANCE v/s REAL TIME

- Real time: have to *guarantee* timing properties
- Performance: minimize *average* response time

These are not the same!

Source of unpredictability:

- Architecture: cache, pipelining, ...
- Run-time system: scheduling, other tasks, ...
- Environment: Bursty information flow, extreme conditions, ...
- Input: no explicit notion of time in most languages



# REAL TIME SYSTEMS

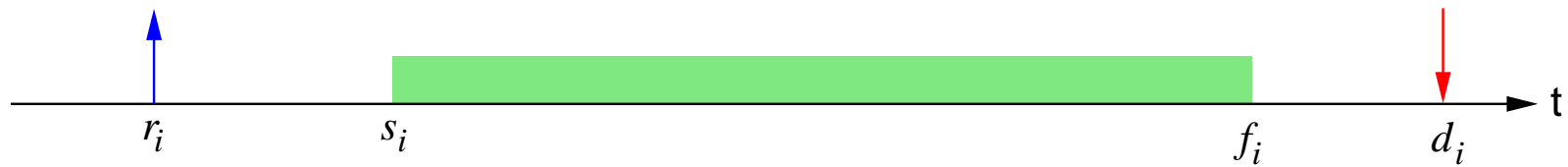
## Terminology:

- A *job* is a sequence of operations that, in the absence of any other activities, is executed by the processor
- A *task* is a sequence (possibly infinite) of jobs
- Jobs have:
  - A request time  $r_i$  (arrival time)
  - A start time  $s_i$
  - A finishing time  $f_i$
  - An absolute deadline  $d_i$



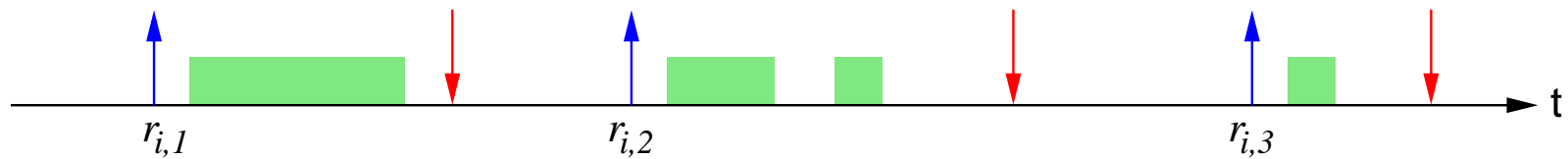
# TASKS AND JOBS

A single job:



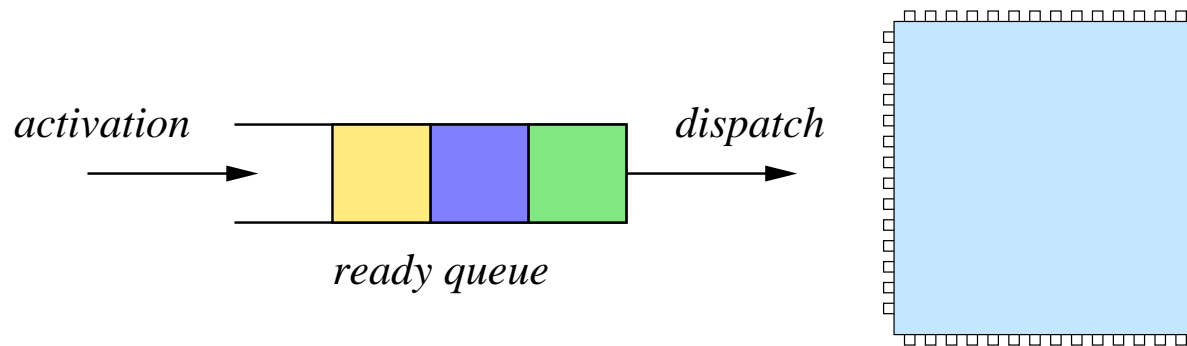
$f_i - s_i$  is the *worst case execution time* (wcet)

A task:



# SCHEDULING

**Scheduling algorithm:** the strategy used to pick a ready task for execution.



Two categories:

- **Preemptive:** The running task can be temporarily suspended to execute another task
- **Non-preemptive:** The running task cannot be suspended until completion or until it is blocked



# SCHEDULING

A **schedule** is a particular assignment of tasks to the processor.

Given a set of tasks  $\Gamma = \{\tau_1, \tau_2, \dots, \tau_n\}$ , a schedule is a mapping  $\sigma : \mathbb{R}_{>0} \rightarrow \{0, 1, \dots, n\}$  such that:

$$\sigma(t) = \begin{cases} k > 0 & \text{if } \tau_k \text{ is running} \\ 0 & \text{if the processor is idle} \end{cases}$$

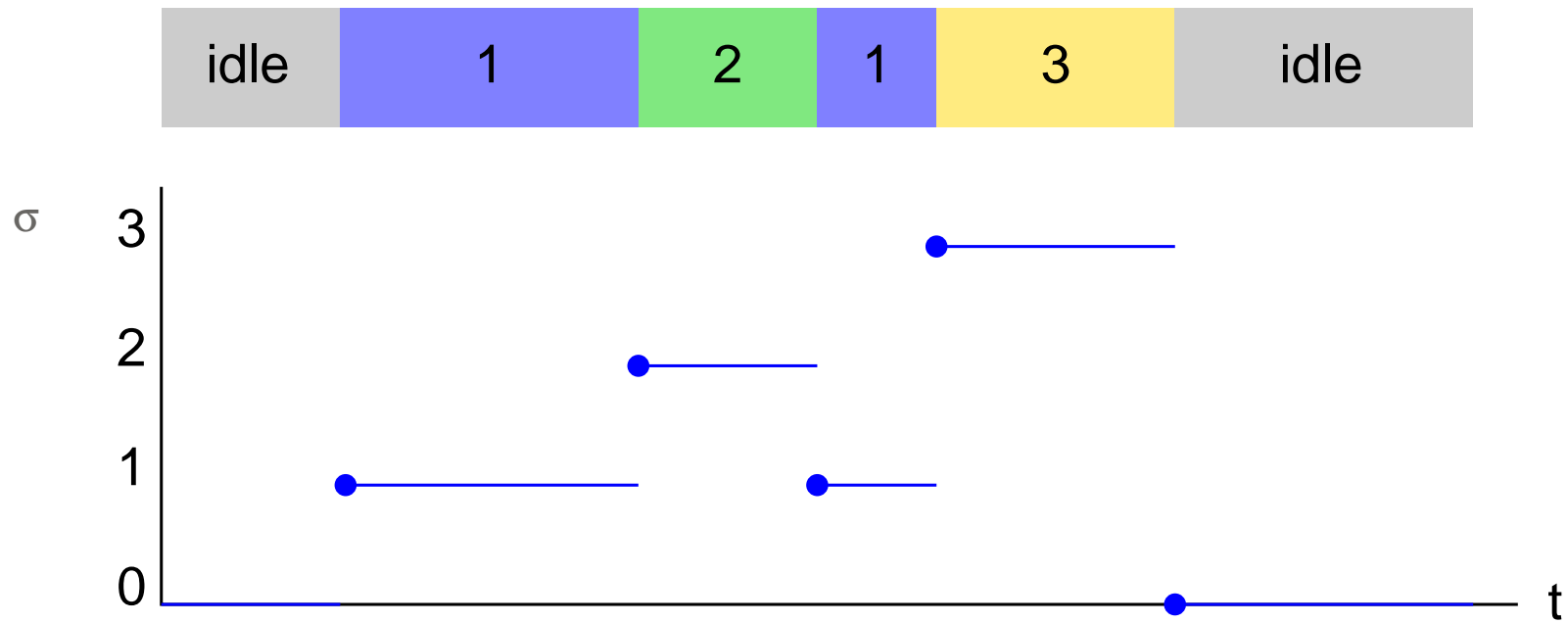
and in any interval  $[t_1, t_2) \in \mathbb{R}_{>0}$   $\sigma(t)$  can only change value a finite number of times.





# SCHEDULING

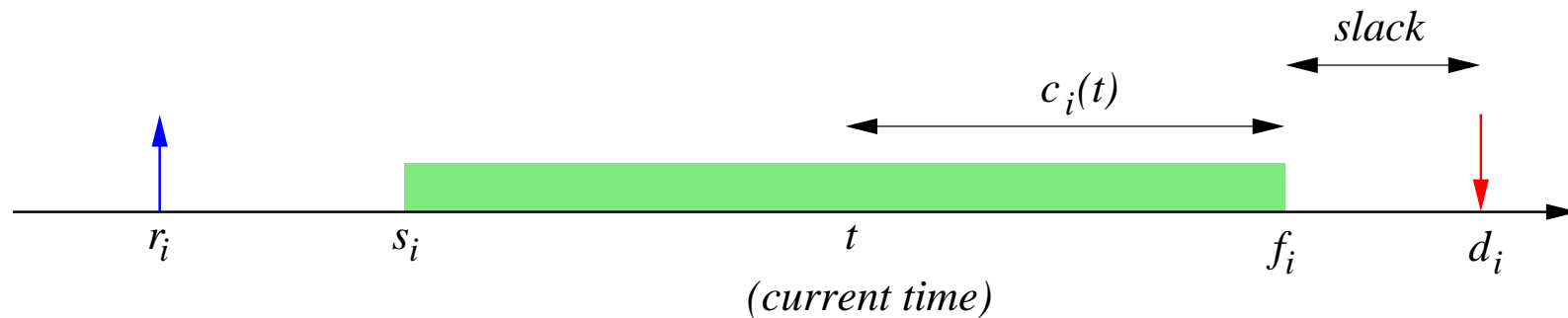
Example:



The points at which  $\sigma$  changes value is where a context switch occurs. Each interval  $[t_i, t_{i+1})$  is a time slice.



# DERIVED PARAMETERS



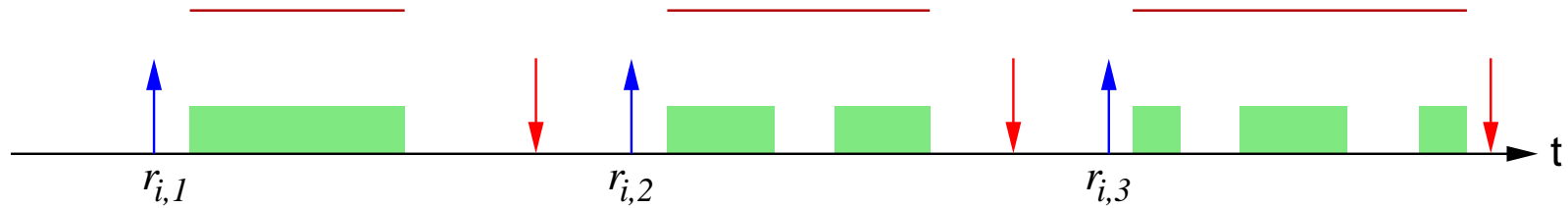
- Lateness:  $L_i = f_i - d_i$
- Tardiness:  $\max(0, L_i)$
- Residual wcet:  $c_i(t)$
- Slack:  $d_i - t - c_i(t)$



# JITTER

Jitter is the time variation of a periodic event.

Example: completion-time jitter



$$\max_k (f_{i,k} - s_{i,k}) - \min_k (f_{i,k} - s_{i,k})$$



# IMPORTANCE OF TASKS

- **Hard tasks:** All jobs must meet their deadlines. Missing a deadline has a catastrophic effect.
  - low-level control
  - sensor-actuator interactions for critical functions
- **Soft tasks:** Missing deadlines is undesirable, but only causes performance degradation
  - reading keyboard input
  - displaying a message
  - updating graphics

Tasks can be assigned *priorities*.

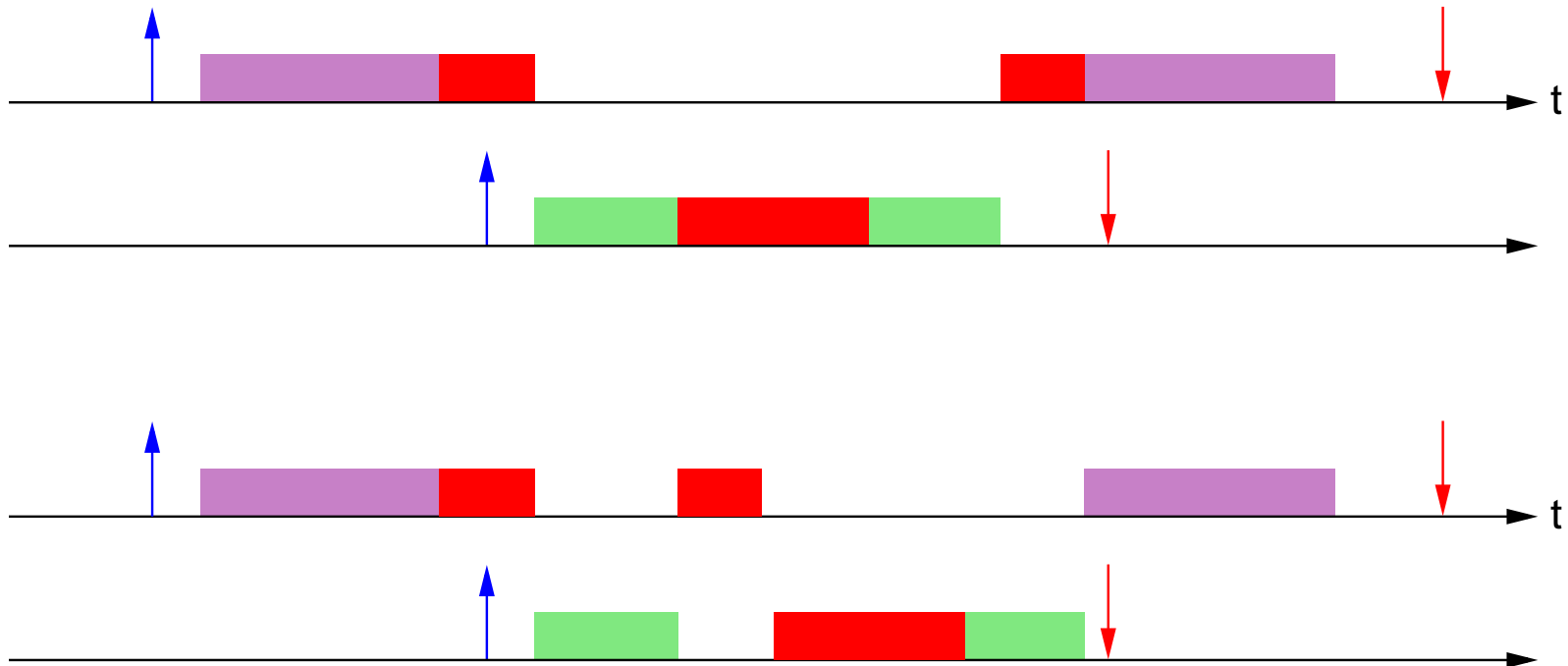
Tasks can be time-driven (periodic) or event-driven (aperiodic).



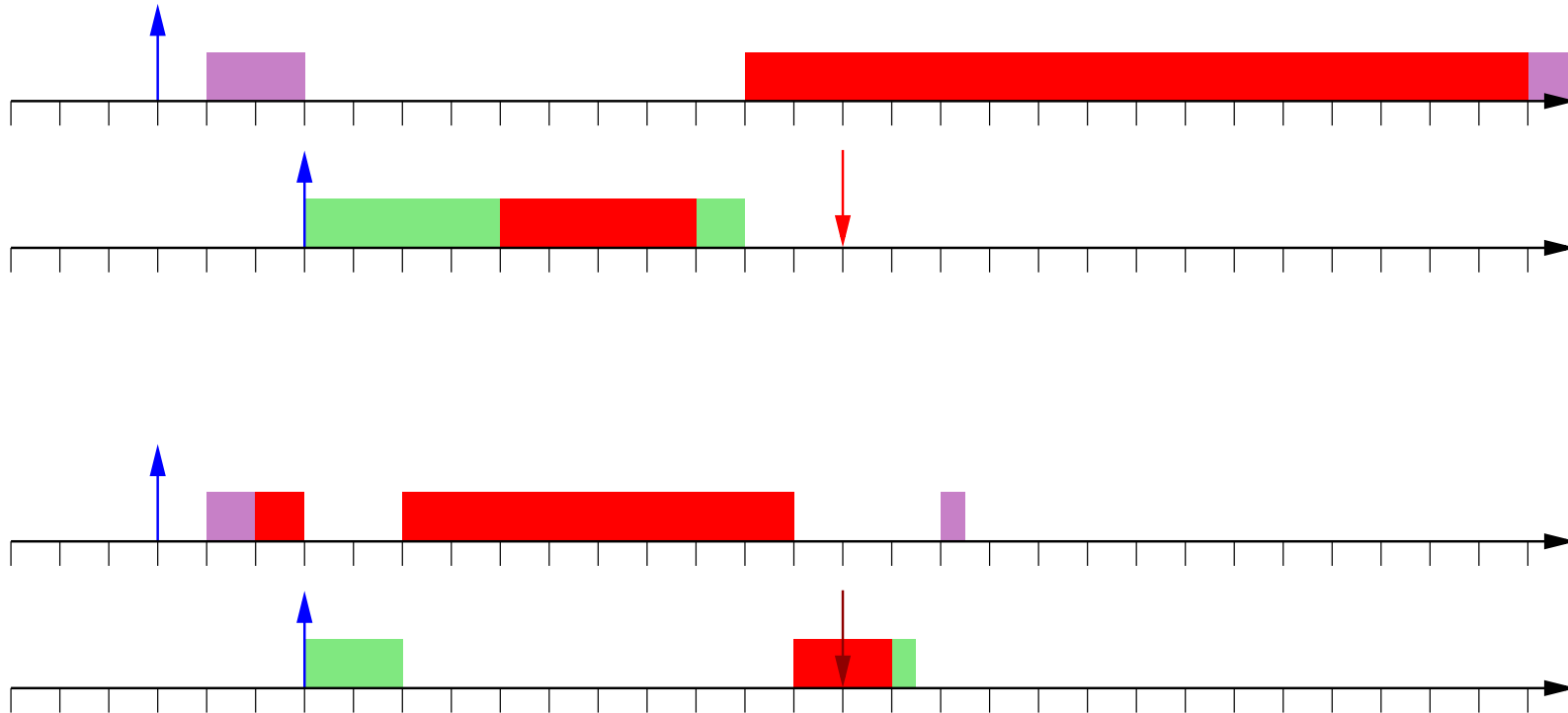
# RESOURCE CONSTRAINTS

- Resources may be limited or even unavailable
- Shared resources may require mutual exclusion

⇒ all these introduce delays



# FASTER PROCESSOR

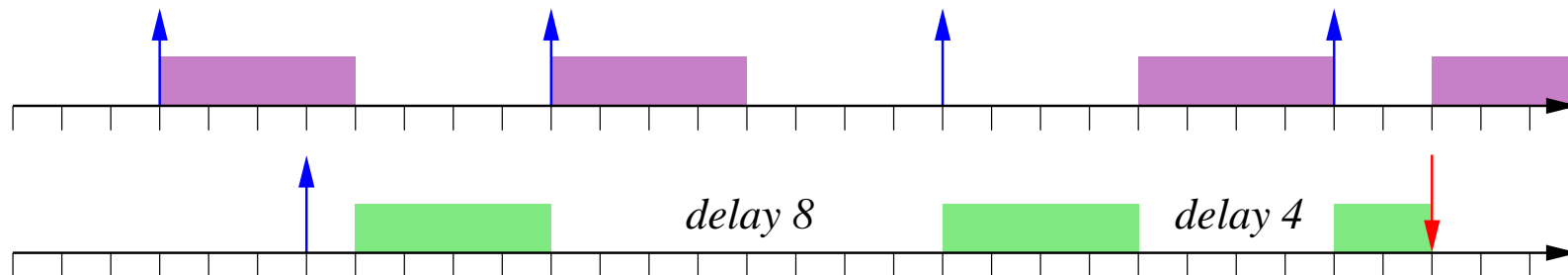
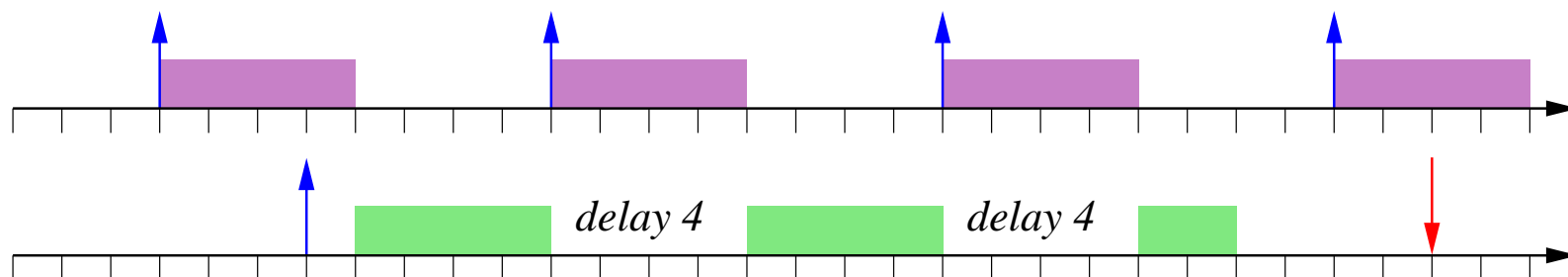


Having a faster processor doesn't automatically mean it is easier to meet deadlines.



# DELAYS

Delays can cause other tasks to take longer!



# PREDICTABILITY

- Run-time system is responsible for ensuring predictable behavior
- Scheduling algorithm matters
- Concurrency control and resource sharing matters
- Interrupt handling must also be predictable

All sources of uncertainty must be minimized.





# SCHEDULING

A schedule  $\sigma$  is *feasible* if all tasks are able to complete with their set of constraints.

A set of tasks  $\Gamma$  is set to be *schedulable* if a feasible schedule exists.

General problem: given  $\Gamma$ , a set of processors  $P$ , and a set of resources  $R$ , find an assignment of  $P$  and  $R$  that produces a feasible schedule.



# SCHEDULING ALGORITHMS

- Preemptive or nonpreemptive
- Static or dynamic: are the scheduling decisions based on parameters that change with time?
- Online or offline: are the decisions made apriori with knowledge of task activations, or are they taken at run time based on the set of active tasks?
- Optimal or heuristic: can you prove that the algorithm optimizes a certain criteria or not?

