ECE 3140 / CS 3420 EMBEDDED SYSTEMS

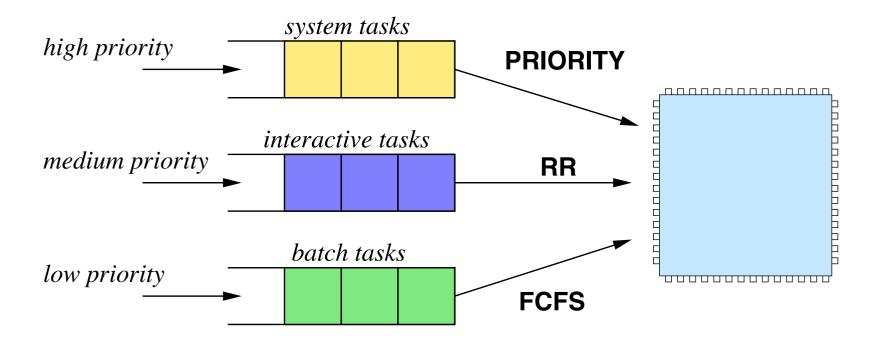
LECTURE 16

Prof. José F. Martínez

TR 1:25-2:40pm in 150 Olin

MULTI-LEVEL SCHEDULING

Mix and match different scheduling algorithms

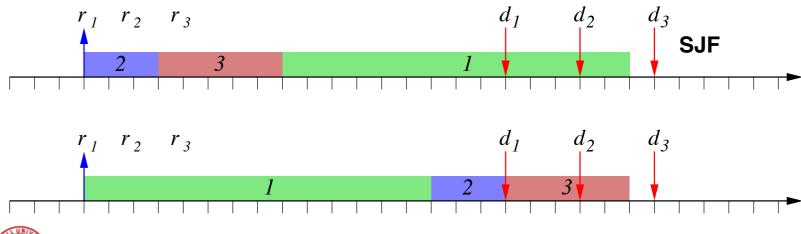


REAL-TIME ALGORITHMS

Tasks can be scheduled by:

- \blacksquare relative deadlines D_i —static
- \blacksquare absolute deadlines d_i —dynamic

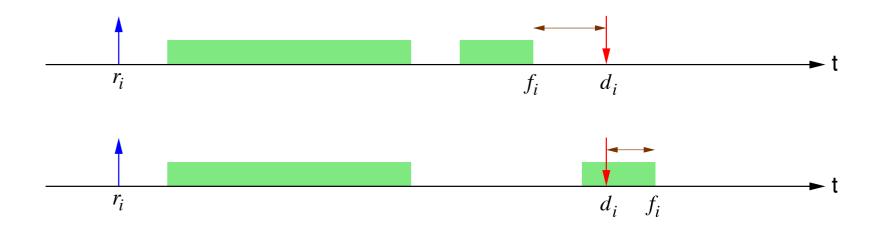
Since the goal is to meet deadlines, we should be using knowledge of deadlines to determine the schedule.



Earliest Due Date (EDD): select the task with the earliest *relative* deadline.

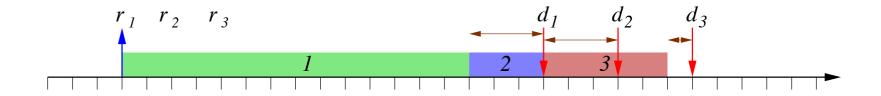
- All tasks arrive simultaneously
- Fixed priority (D_i is known)
- Preemption is not an issue
- It minimizes the maximum lateness L_{max}

$$L_i = f_i - d_i$$



MAXIMUM LATENESS

$$L_{max} = \max_{i}(L_i)$$



■ L_{max} < 0 \Rightarrow no task misses its deadline

Jackson's Rule: Given a set of n independent tasks, any algorithm that executes the tasks in order of nondecreasing deadlines is optimal with respect to maximum latenesss.

Why does it minimize the maximum lateness? Proof?

■ In particular, if $L_{max}(\sigma)$ is the maximum lateness of a schedule, then:

$$\forall \sigma: L_{max}(\sigma_{EDD}) \leq L_{max}(\sigma)$$

How can we guarantee that a task set Γ is feasible?

• We can compute L_{max} !

A task set is feasible iff $\forall i : f_i \leq d_i$

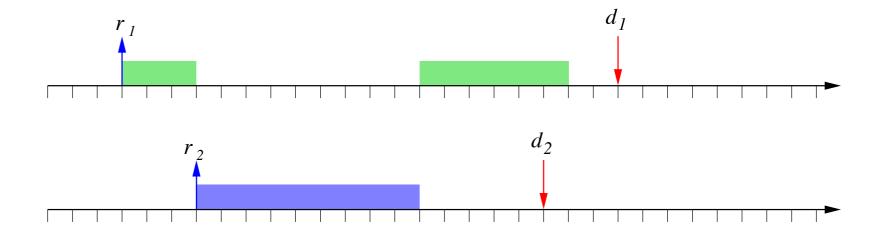
If we sort the tasks using EDD and all tasks arrive simultaneously, then

$$f_i = \sum_{k=1}^i C_i$$

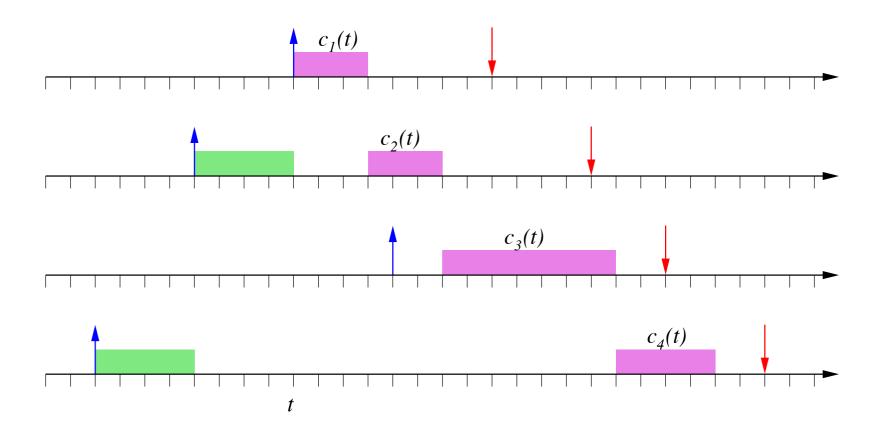
Earliest Deadline First (EDF): select the task with the earliest *absolute* deadline.

- Tasks may arrive at any time
- Dynamic priority (d_i depends on when the tasks arrive)
- Preemptive tasks
- It minimizes the maximum lateness L_{max}

Tasks that arrive with earlier deadlines pre-empt tasks with later deadlines.

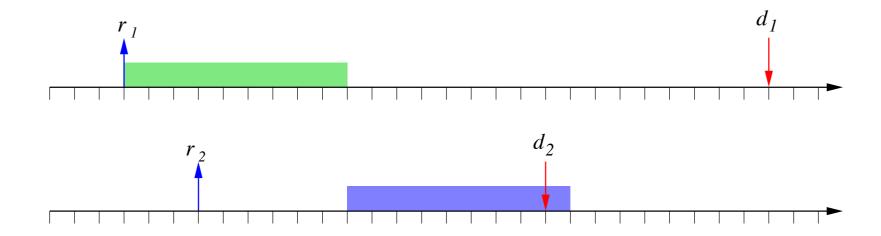


EARLIEST DEADLINE FIRST: FEASIBILITY



$$\forall i: \sum_{k=1}^{i} c_k(t) \le d_i - t$$

Under non-preemptive scheduling, EDF is not optimal



... unless the algorithm has knowledge of the future!

