Syntaxe abstraite:

$$e := true \mid false \mid n \mid x \mid e + e \mid e < e \mid e \text{ and } e$$

 $S := x := e \mid skip \mid S$, $S \mid if e then S else S fi \mid while e do S od$

 $\Gamma \vdash \text{true} : \text{Bool}$ $\Gamma \vdash \text{false} : \text{Bool}$ $\Gamma \; \vdash n \, : \, \mathrm{Int}$

$$\frac{\Gamma(x) = t}{\Gamma \vdash x : t}$$

$$\frac{\Gamma \vdash e_1 : Int}{\Gamma \vdash e_1 \ opa \ e_2 : Int} = \frac{\Gamma \vdash e_2 : Int}{\Gamma \vdash e_1 \ opa \ e_2 : Int}$$

$$\frac{\Gamma(x) = t}{\Gamma \vdash x : t} \qquad \frac{\Gamma \vdash e_1 : Int}{\Gamma \vdash e_1 \ opa \ e_2 : Int} \qquad \frac{\Gamma \vdash e_1 : Bool}{\Gamma \vdash e_1 \ opb \ e_2 : Bool} \qquad \frac{\Gamma \vdash e_1 : t}{\Gamma \vdash e_1 \ oprel \ e_2 : Bool} \qquad \frac{\Gamma \vdash e_1 : t}{\Gamma \vdash e_1 \ oprel \ e_2 : Bool}$$

$$\frac{\Gamma \vdash e_1 : t \quad \Gamma \vdash e_2 : t}{\Gamma \vdash e_1 \ oprel \ e_2 : Bool}$$

$$\frac{\Gamma \vdash e : t \quad \Gamma \vdash x : t}{\Gamma \vdash x := e} \qquad \qquad \frac{\Gamma \vdash skip}{\Gamma \vdash skip}$$

$$\Gamma \vdash ski$$

$$\frac{\Gamma \vdash S_1 \quad \Gamma \vdash S}{\Gamma \vdash S_1 \; ; \; S_2}$$

$$\frac{\Gamma \vdash S_1 \quad \Gamma \vdash S_2}{\Gamma \vdash S_1 \ ; \ S_2} \qquad \qquad \frac{\Gamma \vdash e : Bool \quad \Gamma \vdash S}{\Gamma \vdash while \ e \ do \ S \ od}$$

Extention du système de type:

Règle d'inference pour les blocs:

$$\frac{\Gamma \vdash D_V \mid \Gamma_1 \quad \Gamma_1 \vdash S}{\Gamma \vdash begin \ D_V \ ; \ S \ end}$$

Règle d'inférence pour les déclarations:

$$\frac{\Gamma \vdash e : t \quad \Gamma[x \mapsto t] \vdash D_V \mid \Gamma_1 \quad x \not\in DV(D_V)}{\Gamma \vdash var \ x := e \ ; \ D_V \mid \Gamma_1}$$

Règle d'évaluation collatéral:

$$\frac{\Gamma \vdash e : t \quad \Gamma \vdash D_V \mid \Gamma_1 \quad x \not\in DV(D_V)}{\Gamma \vdash var \ x := e \ ; D_V \mid \Gamma_1[x \mapsto t]}$$

Liaison statique entre procédure et variable:

Bloc:

$$\frac{\Gamma_{V} \vdash D_{V} \mid \Gamma_{V}^{'} \quad (\Gamma_{V}^{'}, \Gamma_{p}) \vdash D_{p} \quad (\Gamma_{V}^{'}, \Gamma_{p}^{'}) \vdash S}{(\Gamma_{V}, \Gamma_{p}) \vdash begin \ D_{V} \ ; \ D_{p} \ ; \ S \ end}$$

 D_p :

$$\frac{(\Gamma_V, \Gamma_p) \, \vdash \, S \quad (\Gamma_V, \Gamma_p[p \mapsto proc]) \, \vdash \, D_p \quad \ p \not\in (D_p)}{(\Gamma_V, \Gamma_p) \, \vdash \, proc \, \, p \, \, is \, \, S \, \, ; \, \, D_p}$$

Appel:

$$\frac{\Gamma_p(p) = proc}{(\Gamma_V, \Gamma_p) \vdash call \ p}$$

Liaison dynamique entre procédure et variable:

Bloc:

$$\frac{\Gamma_{V} \, \vdash D_{V} \, \mid \, \Gamma_{V}^{'} \quad (\Gamma_{V}^{'}, \Gamma_{p}^{'}) \, \vdash S \quad udef(D_{p})}{(\Gamma_{V}, \Gamma_{p}) \, \vdash begin \, D_{V} \, \, ; \, D_{p} \, \, ; \, S \, \, end}$$

Appel:

$$\tfrac{(\Gamma_V,\Gamma_p)\;\vdash\;S}{(\Gamma_V,\Gamma_p)\;\vdash\;call\;\;p}\quad\Gamma_p(p)\;=\;S$$