

Constructing Set Constraints for ReScript

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1 Definition of set expressions

se	$::=$	\emptyset	empty set
		$-$	maximum set
		$()$	unit
		n	integer
		b	boolean
		$\lambda x.e$	function
		V_e	set variable corresponding to the possible values of e
		P_e	set variable corresponding to the possible exn packets of e
		$\text{body}_V(se)$	values that se can spit out when se is applied to something
		$\text{body}_P(se)$	exn packets that se can spit out when se is applied to something
		$\text{par}(se)$	values that can be a parameter for se
		κ	constructor
		l	field of a record
		$\text{con}(\kappa, se)$	construct
		$\text{exn}(\kappa, se)$	exception
		$\text{fld}(se, l)$	contents of the field l of a record se
		$\text{cnt}(se)$	contents of a reference
		$\text{bop}(se, se)$	binary operators, where $\text{bop} \in \{+, -, \times, \div, =, <, >\}$
		$f_{(i)}^{-1}(se)$	projection onto the i -th argument of f
		$se \cup se$	union
		$se \cap se$	intersection

The definition of the conditional set expression needs clarification.

$$se_1 \Rightarrow se_2 := \begin{cases} \emptyset & (se_1 = \emptyset) \\ se_2 & (o.w) \end{cases}$$

The conditional set expression is a naive approximation for pattern matching. Consider the case when we want to match an expression against the record pattern with fields x and y . The constraints describing the record $r = \{x = 1, y = 2\}$ are $1 \subseteq \text{fld}(V_r, x) \wedge 2 \subseteq \text{fld}(V_r, y)$. To pattern-match r against $\{x, y\}$, we want the fields x and y of V_r to be nonempty. Thus, the value of “match r with $\{x, y\} \rightarrow e$ ” is $\text{fld}(V_r, x) \Rightarrow (\text{fld}(V_r, y) \Rightarrow V_e)$.

The conditional set expression can also be used to define conditional set constraints [1].

$$se \Rightarrow \bigwedge_{i=1}^n X_i \subseteq Y_i := \bigwedge_{i=1}^n (se \Rightarrow X_i) \subseteq Y_i$$

2 Constructing set constraints

Now we are in a position to define constraint construction rules for our ReScript-like language. Hopefully this would be reasonably fast when implemented and be accurate enough...

$$\begin{array}{c}
 \text{UNIT, INT, BOOL} \quad \frac{}{\triangleright c : V_e \supseteq c} \quad c = (), n, b \\
 \text{APP} \quad \frac{\triangleright e_1 : C_1 \quad \triangleright e_2 : C_2}{\triangleright e_1 e_2 : (V_e \supseteq \text{body}_V(V_{e_1})) \wedge (P_e \supseteq (\text{body}_P(P_{e_1}) \cup P_{e_2})) \wedge (\text{par}(V_{e_1}) \supseteq V_{e_2}) \wedge C_1 \wedge C_2} \\
 \text{FN} \quad \frac{\triangleright e' : C'}{\triangleright \lambda x.e' : (V_e \supseteq \lambda x.e') \wedge (\text{body}_V(V_e) \supseteq V_{e'}) \wedge (\text{body}_P(V_e) \supseteq P_{e'}) \wedge (\text{par}(V_e) \subseteq V_x) \wedge C'} \\
 \text{LET} \quad \frac{\triangleright e_1 : C_1 \quad \triangleright e_2 : C_2}{\triangleright \text{let } x = e_1 \text{ in } e_2 : (V_x \supseteq V_{e_1}) \wedge (P_x \supseteq P_{e_2}) \wedge}
 \end{array}$$

References

- [1] Alexander Aiken. “Introduction to set constraint-based program analysis”. In: *Science of Computer Programming* 35.2 (1999), pp. 79–111. issn: 0167-6423. doi: [https://doi.org/10.1016/S0167-6423\(99\)00007-6](https://doi.org/10.1016/S0167-6423(99)00007-6).