

Lab06 Design Consideration

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Agenda

- Soft IP
- Top Design

Soft IP

Algorithm - Extended Euclidean Algorithm

The screenshot shows a presentation slide titled "rsa*" in an ActivInspire Studio window. The slide contains mathematical equations and text related to the RSA algorithm.

Equations shown:

$$e \cdot d \bmod \varphi(n) = 1$$
$$7 \cdot d \bmod 40 = 1$$

Step 1: Euclidean algorithm

$$40x + 7y = 1$$
$$40 = 5(7) + 5$$
$$7 = 1(5) + 2$$
$$5 = 2(2) + 1$$

Step 2: Back substitution

A yellow sticky note on the right lists parameters:

- $p = 11$
- $q = 5$
- $n = 55$
- $\varphi(n) = 40$
- $e = 7$
- $d =$

At the bottom of the slide, the text "Choosing public & private RSA keys" is visible.

Ref[1]: <https://youtu.be/kYasb426Yjk>

Ref[2]: https://en.wikipedia.org/wiki/Extended_Euclidean_algorithm

Soft IP

Algorithm - Extended Euclidean Algorithm

A B r_i s_i A t_i B

- $\gcd(240, 46) = 2, -9 \times 240 + 47 \times 46 = 2$

index i	quotient q_{i-1}	Remainder r_i	s_i	t_i
0		240	1	0
1		46	0	1
2	$240 \div 46 = 5$	$240 - 5 \times 46 = 10$	$1 - 5 \times 0 = 1$	$0 - 5 \times 1 = -5$
3	$46 \div 10 = 4$	$46 - 4 \times 10 = 6$	$0 - 4 \times 1 = -4$	$1 - 4 \times -5 = 21$
4	$10 \div 6 = 1$	$10 - 1 \times 6 = 4$	$1 - 1 \times -4 = 5$	$-5 - 1 \times 21 = -26$
5	$6 \div 4 = 1$	$6 - 1 \times 4 = 2$	$-4 - 1 \times 5 = -9$	$21 - 1 \times -26 = 47$
6	$4 \div 2 = 2$	$4 - 2 \times 2 = 0$	$5 - 2 \times -9 = 23$	$-26 - 2 \times 47 = -120$

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Soft IP

Algorithm - Extended Euclidean Algorithm

- Hence, there will be two parts, the computation and the decision.

index i	quotient q_{i-1}	Remainder r_i	s_i	t_i
0		240	1	0
1		46	0	1
2	$240 \div 46 = 5$	$240 - 5 \times 46 = 10$	$1 - 5 \times 0 = 1$	$0 - 5 \times 1 = -5$
3	$46 \div 10 = 4$	$46 - 4 \times 10 = 6$	$0 - 4 \times 1 = -4$	$1 - 4 \times -5 = 21$
4	$10 \div 6 = 1$	$10 - 1 \times 6 = 4$	$1 - 1 \times -4 = 5$	$-5 - 1 \times 21 = -26$
5	$6 \div 4 = 1$	$6 - 1 \times 4 = 2$	$-4 - 1 \times 5 = -9$	$21 - 1 \times -26 = 47$
6	$4 \div 2 = 2$	$4 - 2 \times 2 = 0$	$5 - 2 \times -9 = 23$	$-26 - 2 \times 47 = -120$

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Soft IP

Algorithm - Extended Euclidean Algorithm

- Computation
 - 1. calculation of $(p-1)$, $(q-1)$, and ϕ
 - 2. Iteratively find the remainder and the quotient for six times
- Decision
 - 1. When the remainder is 1, the t is the result

Soft IP

Algorithm - Extended Euclidean Algorithm

- The detailed information can be found in the reference.

```
else if (level_idx == 3) begin: if_lv_3
    wire [WIDTH + 1: 0] r, q;
    assign q = gen_level[2].if_lv_2.phi / gen_level[1].if_lv_1.e;
    assign r = gen_level[2].if_lv_2.phi - q * gen_level[1].if_lv_1.e;

    wire signed [WIDTH + 2: 0] t;
    assign t = -q;
end
else if (level_idx == 4) begin: if_lv_4
    wire [WIDTH + 1: 0] r, q;
    assign q = gen_level[1].if_lv_1.e / gen_level[3].if_lv_3.r;
    assign r = gen_level[1].if_lv_1.e - q * gen_level[3].if_lv_3.r;

    wire signed [WIDTH + 2: 0] t;
    assign t = 1 - gen_level[3].if_lv_3.t * q;
end
else if (level_idx == 5) begin: if_lv_5
    wire [WIDTH + 1: 0] r, q;
    assign q = gen_level[3].if_lv_3.r / gen_level[4].if_lv_4.r;
    assign r = gen_level[3].if_lv_3.r - q * gen_level[4].if_lv_4.r;

    wire signed [WIDTH + 2: 0] t;
    assign t = gen_level[3].if_lv_3.t - gen_level[4].if_lv_4.t * q;
end
```

```
generate
    for (level_idx2 = WIDTH; level_idx2 != WIDTH + 1; level_idx2 = level_idx2 + 1) begin
        if (level_idx2 == 3) begin
            end
        else if (level_idx2 == 4) begin
            always @(*) begin
                if (gen_level[3].if_lv_3.r == 1)
                    tempD = gen_level[3].if_lv_3.t;
                else if (gen_level[4].if_lv_4.r == 1)
                    tempD = gen_level[4].if_lv_4.t;
                else if (gen_level[5].if_lv_5.r == 1)
                    tempD = gen_level[5].if_lv_5.t;
                else if (gen_level[6].if_lv_6.r == 1)
                    tempD = gen_level[6].if_lv_6.t;
                else if (gen_level[7].if_lv_7.r == 1)
                    tempD = gen_level[7].if_lv_7.t;
                else if (gen_level[8].if_lv_8.r == 1)
                    tempD = gen_level[8].if_lv_8.t;
                else
                    tempD = 0;
            end
        end
    end
endgenerate
```

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Top Design

Algorithm - Modular Exponentiation

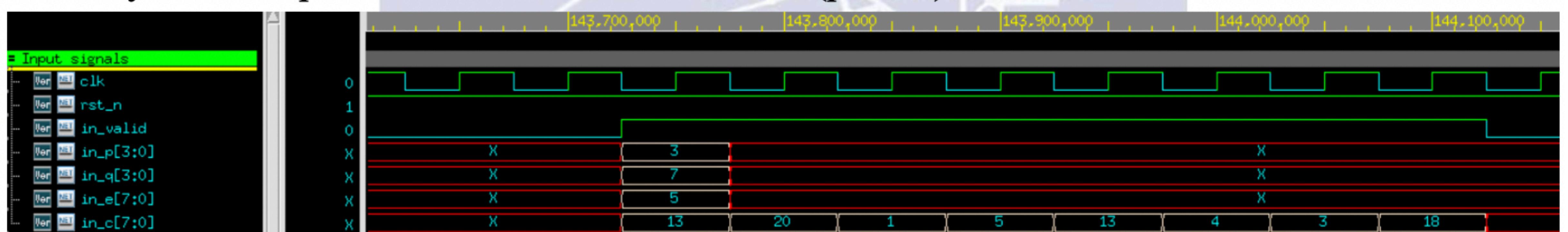
```
function modular_pow(base, exponent, modulus) is
    if modulus = 1 then
        return 0
    Assert :: (modulus - 1) * (modulus - 1) does not overflow base
    result := 1
    base := base mod modulus
    while exponent > 0 do
        if (exponent mod 2 == 1) then
            result := (result * base) mod modulus
        exponent := exponent >> 1
        base := (base * base) mod modulus
    return result
```

Top Design

Algorithm - Modular Exponentiation

- Find the “e” when it starts to input.
- Start doing the modular exponentiation by pipelining, storing them is unnecessary. -> smaller latency

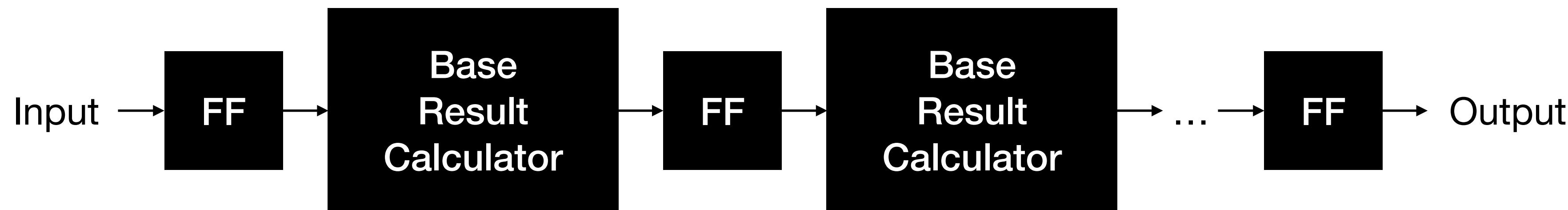
2. 8 cycles for input the information of each round(pattern)



Top Design

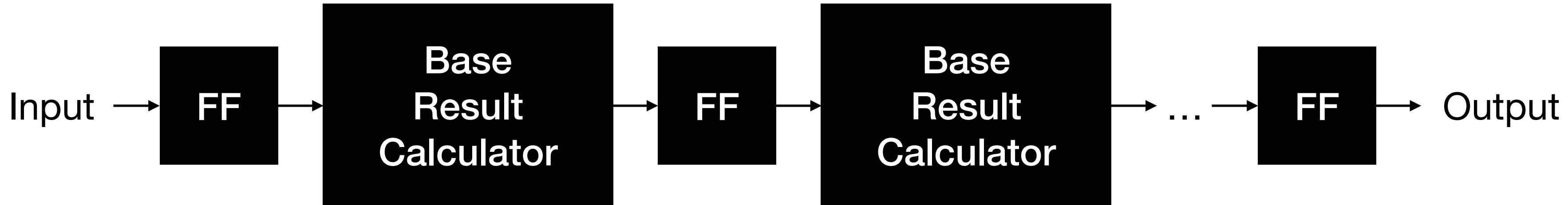
Algorithm - Modular Exponentiation

- Find the “e” when it starts to input.
- Start doing the modular exponentiation by pipelining, storing them is unnecessary. -> smaller latency
- The stages needed are the same as the amount of the input.



Top Design

Algorithm - Modular Exponentiation



```
//=====
assign base[0] = (parameter_d
                  ? (in_c_ff * in_c_ff) % parameter_n
                  : (in_c_ff) % parameter_n;
assign result[0] = (parameter_d & parameter_d[0])
                  ? (in_c_ff) % parameter_n
                  : 1;
//=====
assign base[1] = (parameter_d[6: 1])
                  ? (base_ff[0] * base_ff[0]) % parameter_n
                  : base_ff[0];
assign result[1] = (parameter_d[6: 1] & parameter_d[1])
                  ? (result_ff[0] * base_ff[0]) % parameter_n
                  : result_ff[0];
//=====
assign base[2] = (parameter_d[6: 2])
                  ? (base_ff[1] * base_ff[1]) % parameter_n
                  : base_ff[1];
assign result[2] = (parameter_d[6: 2] & parameter_d[2])
                  ? (result_ff[1] * base_ff[1]) % parameter_n
                  : result_ff[1];
```

```
always @ (posedge clk or negedge rst_n) begin
    if (~rst_n)
        for (i = 0; i != 7; i = i + 1)
            result_ff[i] <= 8'b0;
    else
        for (i = 0; i != 7; i = i + 1)
            result_ff[i] <= result[i];
end
always @ (posedge clk or negedge rst_n) begin
    if (~rst_n)
        for (i = 0; i != 7; i = i + 1)
            base_ff[i] <= 8'b0;
    else
        for (i = 0; i != 7; i = i + 1)
            base_ff[i] <= base[i];
end
```

```
always @ (posedge clk or negedge rst_n) begin
    if (~rst_n) begin
        out_valid <= 1'b0; out_m <= 8'b0;
    end
    else if (NS == OUTPUT) begin
        out_valid <= 1'b1; out_m <= result_ff[6];
    end
    else begin
        out_valid <= 1'b0; out_m <= 8'b0;
    end
end
```

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