

# Lab10 Solution

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# Lab10.A Shopping

- ▶ Lanran likes shopping! There are  $n$  items in the shop, where each one has beauty  $w_i$  and costs  $c_i$  coins. Lanran has  $m$  coins, and he wants to get the largest sum of beauty on items he can buy. Note that, Lanran can buy at most one per item.

Sample input:

3	6
5	3
3	2
3	4

items	$W_i$	$C_i$
1	5	3
2	3	2
3	3	4



items	$sum(C_i)$	$sum(W_i)$
1	3	5
2	2	3
3	4	3
1+2	$3+2 \leq 6$	$5+3 = 8$
2+3	$3+3 \leq 6$	$2+4 = 6$
1+3	$3+4 > 6$	--
1+2+3	$3+2+4 > 6$	--



Sample Output:

8
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# Review:



$W = 11$

#	value	weight
1	1	1
2	6	2
3	18	5
4	22	6
5	28	7

$n + 1$

	$W + 1$											
	0	1	2	3	4	5	6	7	8	9	10	11
$\phi$	0	0	0	0	0	0	0	0	0	0	0	0
{ 1 }	0	1	1	1	1	1	1	1	1	1	1	1
{ 1, 2 }	0	1	6	7	7	7	7	7	7	7	7	7
{ 1, 2, 3 }	0	1	6	7	7	18	19	24	25	25	25	25
{ 1, 2, 3, 4 }	0	1	6	7	7	18	22	24	28	29	29	40
{ 1, 2, 3, 4, 5 }	0	1	6	7	7	18	22	28	29	34	34	40

# Pseudo-code

```
Input:  $n, W, w_1, \dots, w_N, v_1, \dots, v_N$ 

for  $w = 0$  to  $W$ 
     $M[0, w] = 0$ 

for  $i = 1$  to  $n$ 
    for  $w = 1$  to  $W$ 
        if  $(w_i > w)$ 
             $M[i, w] = M[i-1, w]$ 
        else
             $M[i, w] = \max \{M[i-1, w], v_i + M[i-1, w-w_i]\}$ 

return  $M[n, W]$ 
```

# Space complexity optimization of 0-1 knapsack

$n + 1$

$W + 1$

		0	1	2	3	4	5	6	7	8	9	10	11
$\phi$		0	0	0	0	0	0	0	0	0	0	0	0
{ 1 }		0	1	1	1	1	1	1	1	1	1	1	1
{ 1, 2 }		0	1	6	7	7	7	7	7	7	7	7	7
{ 1, 2, 3 }		0	1	6	7	7	18	19	24	25	25	25	25
{ 1, 2, 3, 4 }		0	1	6	7	7	18	22	24	28	29	29	40
{ 1, 2, 3, 4, 5 }		0	1	6	7	7	18	22	28	29	34	34	40

To calculate the new values of row  $i$ , only rows  $i - 1$  needed. The data from rows 0 to  $i - 2$  do not need storage space.

For  $W$  to  $w_i$   
 $M[w] = \max \{M[w], v_i + M[w - w_i]\}$

For  $W$  to  $w_i$   
 $M[i, w] = \max \{M[i, w], v_i + M[i, w - w_i]\}$

**Input:**  $n, W, w_1, \dots, w_N, v_1, \dots, v_N$

**for**  $w = 0$  to  $W$   
     $M[0, w] = 0$

**for**  $i = 1$  to  $n$

**for**  $w = 1$  to  $W$   
        **if**  $(w_i > w)$   
             $M[i, w] = M[i-1, w]$   
        **else**  
             $M[i, w] = \max \{M[i-1, w], v_i + M[i-1, w - w_i]\}$

**return**  $M[n, W]$

# Lab10.B Game

- ▶ Bob is very angry because Alice broke his wonderful TV, so he will battle with Alice in a game to avenge.
- ▶ The game is a very simple game. Initially, they are given a non-empty string  $s$ , consisting of lowercase letters. The length of the string is even. Each player also has its own empty string. In one move, a player takes either the first or the last letter of the string  $s$ , removes it from  $s$  and **appends** it to their own string (put it at the end of the string).
- ▶ The game ends when  $s$  is empty. And then Alice and Bob will compare their strings, the one that owns a lexicographically smaller string will be the winner.
- ▶ Bob is very confident with the game, so he will let **Alice move first**. Also, Alice and Bob will take their moves optimally.
- ▶ Your task is to tell who is the winner or they draw.



Sample Input: → n: the number of testcases

3

ilikealgorithm → the non-empty string in a game

ggggggggggggggg

oooooooooooooooohhhhhhhhhhhhhh

ilikealgorithm

Alice move first, if Alice get character **i** first, Bob must lose the game

ggggggggggggggg

All characters are the same, will be draw.

oooooooooooooooohhhhhhhhhhhhhh

Alice move first, if Alice get **i** first, Bob must lose the game

Sample Output:

**Alice**

**Draw**

**Alice**

only 2 characters:

ij       $S_1 < S_2$    Alice is clever    $\longrightarrow$    Alice get  $S_1$     $\longrightarrow$    Alice win

ji       $S_1 > S_2$    Alice is clever    $\longrightarrow$    Alice get  $S_2$     $\longrightarrow$    Alice win

ii       $S_1 = S_2$    Draw

4 characters:

ijkl  $S_1 < S_4$   $S_1 < S_2$  Alice get  $S_1 \longrightarrow$  Bob get  $S_2$  or  $S_4 \longrightarrow$  Alice win

jikl  $S_1 < S_4$   $S_1 > S_2$  Alice get  $S_1 \longrightarrow$  Bob get  $S_2 \longrightarrow$  Bob win

Alice get  $S_4 \longrightarrow$  Bob get  $S_1 \longrightarrow$  Bob win

jlki  $S_1 > S_4$   $S_3 > S_4$  Alice get  $S_4 \longrightarrow$  Bob get  $S_1$  or  $S_3 \longrightarrow$  Alice win

jikl  $S_1 > S_4$   $S_3 < S_4$  Alice get  $S_4 \longrightarrow$  Bob get  $S_3 \longrightarrow$  Bob win

Alice get  $S_1 \longrightarrow$  Bob get  $S_4 \longrightarrow$  Bob win

ijki  $S_1 = S_4 < S_2, S_3$  Alice get  $S_1 \longrightarrow$  Bob get  $S_2$  or  $S_4 \longrightarrow$  Alice win

Alice get  $S_4 \longrightarrow$  Bob get  $S_1$  or  $S_3 \longrightarrow$  Alice win

jikj  $S_2 < S_1 = S_4 < S_3$  Alice get  $S_4 \longrightarrow$  Bob get  $S_1 \longrightarrow$  only 2 characters  $\longrightarrow$  Alice win

jki j  $S_3 < S_1 = S_4 < S_2$  Alice get  $S_1 \longrightarrow$  Bob get  $S_4 \longrightarrow$  only 2 characters  $\longrightarrow$  Alice win

kijk  $S_2 \leq S_3 < S_1 = S_4$  Alice get  $S_4 \rightarrow$  Bob get  $S_3 \rightarrow$  Bob win

Alice get  $S_1 \rightarrow$  Bob get  $S_2 \rightarrow$  Bob win

jji  $S_3 < S_1 = S_4 = S_2$  ✓ Alice get  $S_1 \rightarrow$  Bob get  $S_4$  or  $S_2 \rightarrow$  only 2 characters  $\rightarrow$  Alice win

Alice get  $S_4 \rightarrow$  Bob get  $S_3 \rightarrow$  Bob win

jjkj  $S_1 = S_4 = S_2 < S_3$  Alice get  $S_1 \rightarrow$  Bob get  $S_4$  or  $S_2 \rightarrow$  only 2 characters  $\rightarrow$  Alice win

Alice get  $S_4 \rightarrow$  Bob get  $S_2 \rightarrow$  only 2 characters  $\rightarrow$  Alice win

jjjj  $S_1 = S_2 = S_3 = S_4$  Alice get  $S_1 \rightarrow$  Bob get  $S_4$  or  $S_2 \rightarrow$  only 2 characters  $\rightarrow$  Draw



More characters?