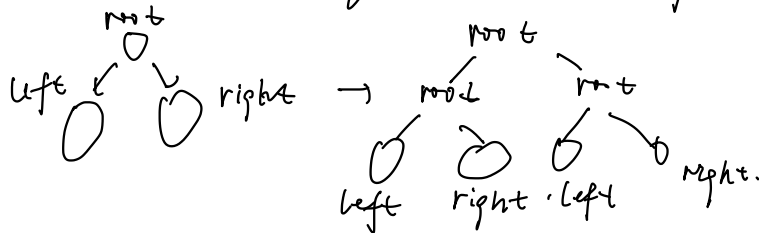


Q1 a)

1) pre + in : Yes

By pre-order sequence, we can know the root (1st node), then in-order sequence shows the left and right sub-tree.



In this way, we can finally reconstruct the whole tree

2) post + In : Yes

The same as pre + in, we can use post-order sequence to know the root, then use in-order sequence to divide left and right sub-tree.

3) pre + post : No

E.g. pre-order : AB
post-order : BA



Note: pre + post is ok if tree is full tree

b) maxDist (Node) {

```
distAndDepth (node) {
```

```
    if empty (node):
```

```
        return 0, 0
```

```
    L_dep, L_dist = distAndDepth (node->left)
```

```
    R_dep, R_dist = distAndDepth (node->right)
```

```
    dist = L_dep + R_dep
```

```
    max_dist = { L_dist, R_dist, dist }
```

```
    max_depth = { L_depth, R_depth }
```

```
    return max_depth+1, max_dist
```

```
return max_dist
```

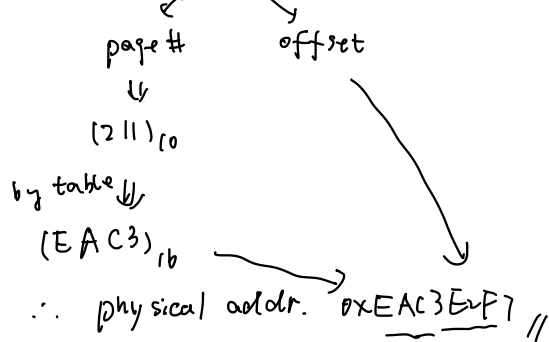
Q2 (a.1) $2^{32} = 4GB$

(a.2) 2^{16} frames, 2^{32} address \Rightarrow frame size $2^{16} \Rightarrow 16$ -bit offset

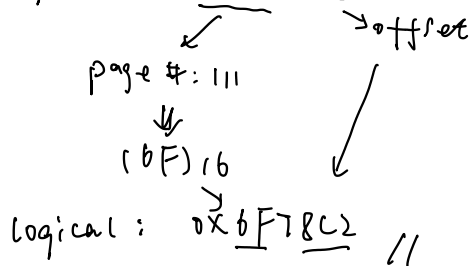
256 entries \Rightarrow 8-bit page numbers

logic addr. = page numb + offset = 16+8 = 24-bit //

(a.3) Logical: D3E2F7



(a.4) physical: A89C78C2



(b) Note: fork() function in Unix.

1. Copy itself, start a subprocess

2. parent process return subprocess ID
subprocess return 0;

SO $\left. \begin{array}{l} \text{f} \\ \swarrow \quad \searrow \\ \text{r} \quad \text{2} \\ \swarrow \quad \searrow \\ \text{2} \quad \text{2} \end{array} \right\} 10$ 2^{10} Hello //

(c)

	1	3	2	3	4	1	5	2	5	6
FIFO:	1	1	1	1	4	4	4	2	2	2
		3	3	3	3	1	1	1	1	6
			2	2	2	2	5	5	5	5
	F	F	F		F	F	F	F		F

8 page faults.

	1	3	2	3	4	1	5	2	5	6
LRU	1	1	1	1	4	4	4	2	2	2
		3	3	3	3	3	5	5	5	5
			2	2	2	1	1	1	1	6
	F	F	F		F	F	F	F		F

8 page faults.

Q3 (a) R: $10B \times 3 \times 18000 = 540000 B$ $\frac{540000}{3000} = 180 \text{ pages}$
 S: $10B \times 3 \times 48000 = 1440000 B$ $\frac{1440000}{3000} = 480 \text{ pages}$

(b) Height of B⁺ tree: $\log_{\frac{n}{2}} (1K) = \log_{\frac{20}{2}} (18000) = 5$

① if B is a primary index

of block trans: $h + b = 5 + 10 = 15$
↖
height target
 records

of seeks: $h + 1 = 5 + 1 = 6$

② if B is a secondary index:

of block trans: $h + b = 5 + 10 = 15$

of seeks: $h + 1 = 5 + 1 = 6$

(c) unsolved!

(d) unsolved!

Q4

(a) Before splitting:

+	→	<table border="1"><tr><td>C₀</td><td>5</td></tr><tr><td>C₁</td><td>5</td></tr></table>	C ₀	5	C ₁	5
C ₀	5					
C ₁	5					
-	→					

$C_{ini} = 1 - \frac{1}{4} - \frac{1}{4} = \frac{1}{2}$

A
 Yes \swarrow No \searrow
 $C_0: 4$ $C_0: 1$
 $C_1: 0$ $C_1: 5$

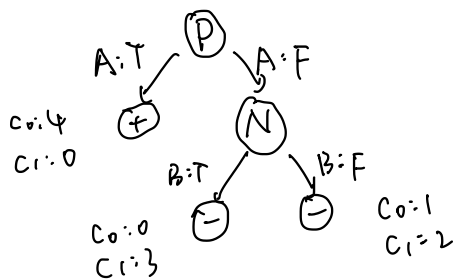
$$\text{Gain} = \frac{1}{2} - \frac{1}{6} = \frac{1}{3}$$

yes \swarrow B \searrow No

$$\begin{array}{ll} C_0 : 2 & C_0 : 3 \\ C_1 : 3 & C_1 : 2 \end{array}$$

$$\text{Gain} = \frac{1}{2} - \frac{12}{25} = \frac{1}{50}$$

(b)



(C)	predict	label
	TT +	+
	TT +	+
	FT -	-
	TF +	+
	TF +	+
	FF -	-
	FF -	-
	FF -	+
	FF -	-
	FF -	-
	FF -	-

	Preval	
	+	-
actual	+	4
	-	0
		1
		5

Confusion
matrix

$$\text{Accuracy} = \frac{9}{10}$$

$$\text{recall} = \frac{6}{5}$$

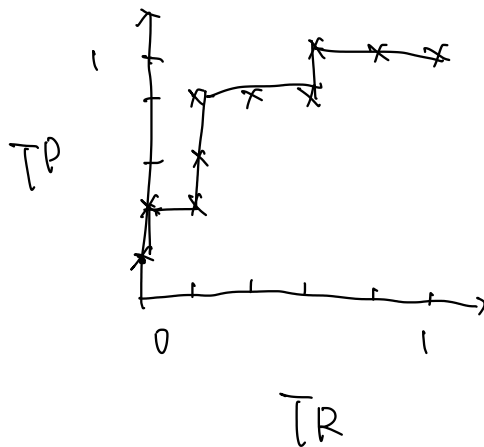
$$\begin{aligned} \text{Accuracy} &= \frac{9}{10} & \text{recall} &= \frac{4}{5} \\ \text{precision} &= 1 & F1 &= \frac{2P \cdot R}{P+R} = \frac{\frac{8}{5}}{\frac{9}{5}} = \frac{8}{9} \end{aligned}$$

 $\alpha)$

Threshold

[illegible]

TP	$\frac{1}{5}$	$\frac{2}{5}$	$\frac{2}{5}$	$\frac{3}{5}$	$\frac{4}{5}$	$\frac{4}{5}$	$\frac{6}{5}$	1	1	1
FP	0	0	$\frac{1}{5}$	$\frac{1}{5}$	$\frac{1}{5}$	$\frac{2}{5}$	$\frac{3}{5}$	$\frac{3}{5}$	$\frac{4}{5}$	1

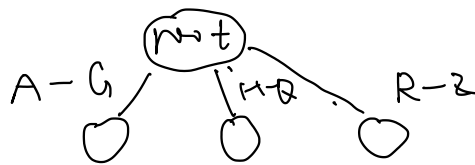


$$AUC = 0.8$$

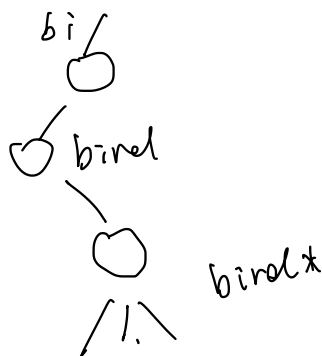
Q5 a) $1 < 3 > 3 < 2 >$

(b) (See book)

d) 1. build a B-tree like:



where the leaf or inner nodes, complete words, splitted by sub



Then to find bird*, we can traverse the sub-tree with "bird" root.

a) Reverse B-tree

c) Train Rocchio C, D

for each $c_j \in C$

do $D_j \leftarrow \{d_i \in D, c_j > GD\}$ # Samples of class j

$$\mu_j \leftarrow \frac{1}{|D_j|} \sum_{d \in D_j} \vec{v}(d) \quad \# \text{ center of samples}$$

return $\{\mu_1, \dots, \mu_k\}$

Test Rocchio ($\mu_1, \dots, \mu_k | d$)

return $\arg \min_j \cos(\mu_j, \vec{v}(d))$ # nearest center

