Questions

Q1. Consider a data base with the following schema

Students(<u>ssn</u>, name, address) Course(<u>code</u>, title) Registered(ssn, code)

- 1. List the codes of courses for which no student is registered
- 2. The titles of courses for which no student is registered.
- 3. Names of the students and the titles of the courses they registered to
- 4. SSNs of students who are registered for 'Database Systems' or 'Analysis of Algorithms'.
- 5. SSNs of students who are registered for both 'Database Systems' and 'Analysis of Algorithms'.
- 6. List of courses in which all students are registered.

Answers

- 1. π code (Course) π code (Registered)
- 2. π name ((π code (Course) π code (Registered)) \bowtie Course)
- 3. **π**name,title (Student ⋈ Registered ⋈ Course)
- πssn (Student ⋈ Registered ⋈ (σ title='Database Systems' Course)) U πssn (Student
 καρίστατα και (σ title='Analysis of Algorithms' Course))
- Tssn (Student ⋈ Registered ⋈ (σ title='Database Systems' Course)) ∩ Tssn (Student
 Registered ⋈ (σ title='Analysis of Algorithms' Course))
- 6. π code, ssn (Registered) / π ssn (Student)

Give the following queries in the relational algebra using the relational schema student(id, name) enrolledIn(id, code) subject(code, lecturer)

- 1. What are the names of students enrolled in cs3020?
- 2. Which subjects is Hector taking?
- 3. Who teaches cs1500?
- 4. Who teaches cs1500 or cs3020?
- 5. Who teaches at least two different subjects?
- 6. What are the names of students in cs1500 or cs3010?
- 7. What are the names of students in both cs1500 and cs1200?

- 8. What are the names of students in at least two different subjects?
- 9. What are the codes of all the subjects taught?
- 10. What are the names of all the students?
- 11. What are the names of all the students in cs1500?
- 12. What are the names of students taking a subject taught by Roger.
- 13. What are the names of students who are taking a subject not taught by Roger?

Answers

- 1. π name(σ cs3020=code(student \bowtie enrolledIn))
- 2. π code(σ name=Hector(student \bowtie enrolledIn))
- 3. π lecturer(σ code=cs1500(subject))
- 4. π lecturer(σ code=cs1500 \vee code=cs3020(subject))
- 5. Solution: For this query we have to relate subject to itself. To disambiguate the relation, we will call the subject relation R or S.

lecturer(**O** R.lecturer = S.lecturer AND R.code < > S.code(R ⋈ S))

- 6. π name(σ code=cs1500(student \bowtie enrolledIn)) \cup π name(σ code=cs3010(student \bowtie enrolledIn))
- 7. π name(σ code=cs1500(student \bowtie enrolledIn)) \cap

 $\pi_{\text{name}}(\sigma_{\text{code}=\text{cs3010(student}}\bowtie_{\text{enrolledIn)}})$

8. For this query we have to relate *enrolledIn* to itself. To disambiguate the relation, we will call the enrolledIn relation R or S.

 π name(student \bowtie (σ R.id = S.id AND R.code < > S.code(R \bowtie S)))

- 9. $\pi_{\text{code(subject)}}$
- 10. π name(student)
- 11. π name(σ code=cs1500(student \bowtie enrolledin))
- 12. π name(σ lecturer=Roger(student \bowtie enrolledIn \bowtie subject))
- 13. π name(σ ecturer < > Roger(student \bowtie enrolledIn \bowtie subject))